

Note that in case when all sets  $B_i$  have exactly 2 elements (i.e.  $b = 2$ ), the Hitting Set problem is equivalent to the Vertex Cover problem (two-element sets  $B_i$  correspond to edges). In the chapter we saw two approximation algorithm for Vertex Cover; here we generalize the one based on linear programming to arbitrary  $b$ .

Consider the following problem for Linear Programming:

$$\begin{aligned} \text{Min} \quad & \sum_{i=1}^n w_i x_i \\ \text{s.t.} \quad & 0 \leq x_i \leq 1 \quad \text{for all } i = 1, \dots, n \\ & \sum_{i: a_i \in B_j} x_i \geq 1 \quad \text{for all } j = 1, \dots, m \text{ (all sets are hit)} \end{aligned}$$

Let  $x$  be the solution of this problem, and  $w_{LP}$  is a value of this solution (i.e.  $w_{LP} = \sum_{i=1}^n w_i x_i$ ).

Now define the set  $S$  to be all those elements where  $x_i \geq 1/b$ :

$$S = \{a_i \mid x_i \geq 1/b\}$$

(1)  $S$  is a hitting set.

*Proof.* We want to prove that any set  $B_j$  intersects with  $S$ . We know that the sum of all  $x_i$  where  $a_i \in B_j$  is at least 1. The set  $B_j$  contains at most  $b$  elements. Therefore some  $x_i \geq 1/b$ , for some  $a_i \in B_j$ . By definition of  $S$ , this element  $a_i \in S$ . So,  $B_j$  intersects with  $S$  by  $a_i$ . ■

(2) The total weight of all elements in  $S$  is at most  $b \cdot w_{LP}$ .

*Proof.* For each  $a_i \in S$  we know that  $x_i > 1/b$ , i.e.,  $1 < bx_i$ . Therefore

$$w(S) = \sum_{a_i \in S} w_i \leq \sum_{a_i \in S} w_i \cdot bx_i \leq b \sum_{i=1}^n w_i x_i = bw_{LP}$$

■

(3) Let  $S^*$  be the optimal hitting set. Then  $w_{LP} \leq w(S^*)$ .

*Proof.* Set  $x_i = 1$  if  $a_i$  is in  $S^*$ , and  $x_i = 0$  otherwise. Then the vector  $x$  satisfy constraints of our problem for Linear Programming:

$$\begin{aligned} 0 \leq x_i \leq 1 \quad & \text{for all } i = 1, \dots, n \\ \sum_{i: a_i \in B_j} x_i \geq 1 \quad & \text{for all } j = 1, \dots, m \text{ (because all sets are hit)} \end{aligned}$$

Therefore the optimal solution is not worse than this particular one. That is,

$$w_{LP} \leq \sum_{i=1}^n w_i x_i = \sum_{a_i \in S^*} w_i = w(S^*)$$

■

Therefore we have a hitting set  $S$ , such that  $w(S) \leq b \cdot w(S^*)$ .

---

<sup>1</sup>ex53.496.888