### **EasyCommit: A Non-blocking Two-phase Commit Protocol**

# Presentation for ECS 265: Distributed Database Systems

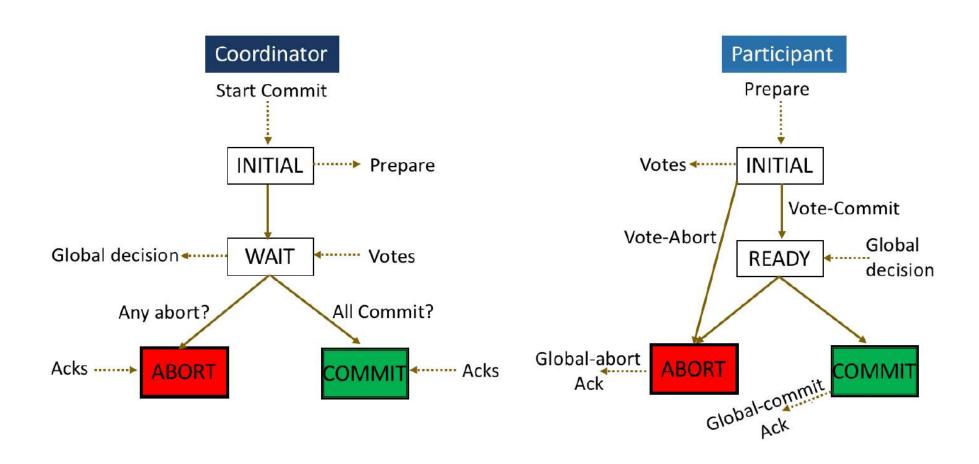
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**October 4, 2018** 

# Roadmap

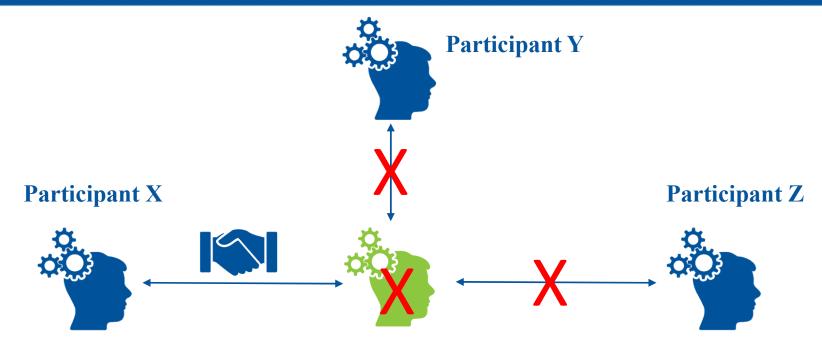
- 1. Background
- 2. Two-phase Commit (2PC) Protocol
- 3. Three-phase Commit (3PC) Protocol
- 4. EasyCommit Protocol
- 5. Comparison and Analysis

### **Two-phase Commit (2PC) Protocol**



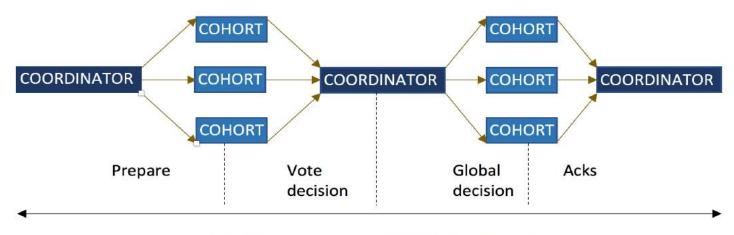
Source: EasyCommit: A Non-blocking Two-phase Commit Protocol. EDBT'18.

# **Blocking Problem in 2PC Protocol**

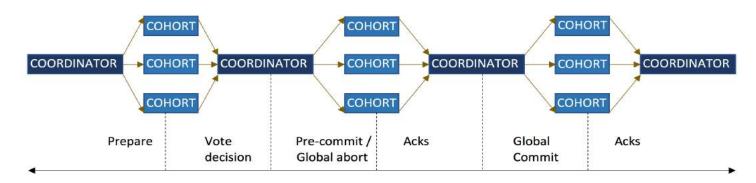


**Coordinator C** 

### **Three-phase Commit (3PC) Protocol**



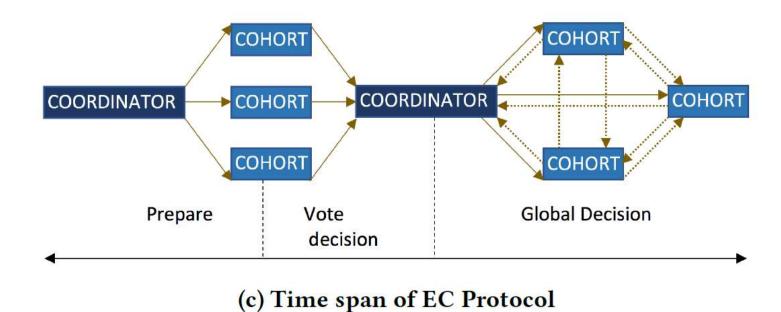
#### (a) Time span of 2PC Protocol



#### (b) Time span of 3PC Protocol

Source: EasyCommit: A Non-blocking Two-phase Commit Protocol. EDBT'18.

# **EasyCommit Protocol**



Source: EasyCommit: A Non-blocking Two-phase Commit Protocol. EDBT'18.

### **Termination Protocol**

#### **Three different situations:**

- 1. Coordinator Timeout in WAIT State
- 2. Cohort Timeout in INITIAL State
- 3. Cohort Timeout in READY State

# **Independent Recovery**

### EasyCommit Protocol can do the independent recovery when:

- 1. If a cohort fails before transmitting its vote, then on recovery, it can simply abort the transaction.
- 2. If the coordinator fails before transmitting the global decision, then it aborts the transaction on recovery.
- 3. If either coordinator or participant fail after transmitting the global decision and writing the log, then on recovery they can use this entry to reach the consistent state.

### **Conclusion**

The easy commit protocol takes the advantages of both 2PC Protocol and 3PC Protocol. It is able to:

- 1. Handle commit transactions efficiently (Two phases).
- 2. Resolve the transactions safely if the coordinator fails (Non-blocking).
- 3. Execute independent recovery in three different cases. (Independent recovery).