



# *subtractive mixture models*

## *representation, learning & inference*

**antonio vergari** (he/him)

 @tetradosse

18th Feb 2025 - Foundations of AI Seminar Warwick

*april*

april-tools.github.io

# *april*

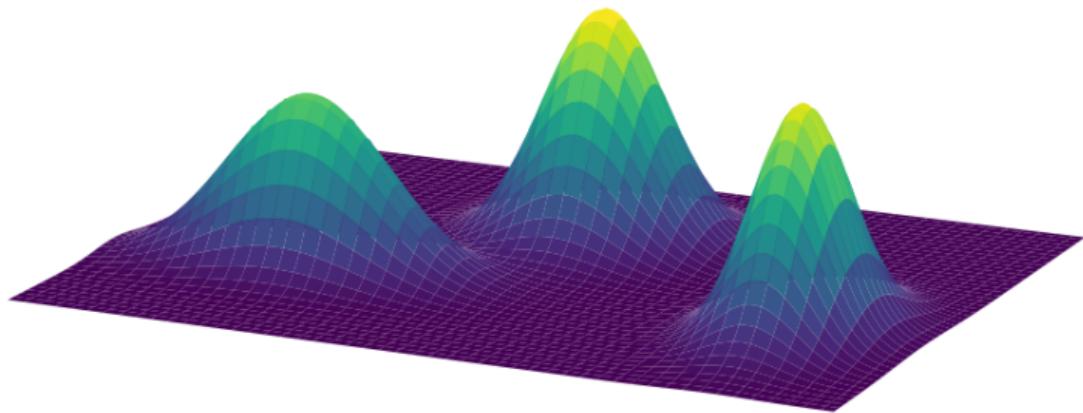
*about  
probabilities  
integrals &  
logic*

# *april*

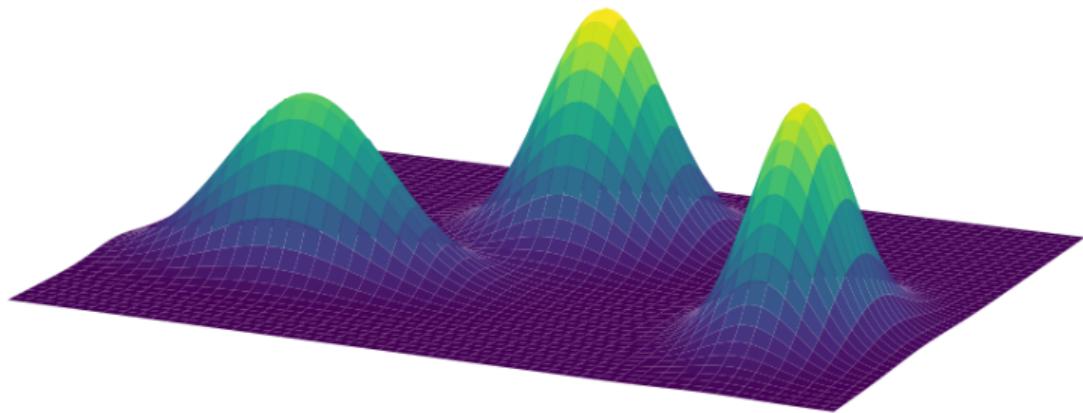
*autonomous &  
provably  
reliable  
intelligent  
learners*

# *april*

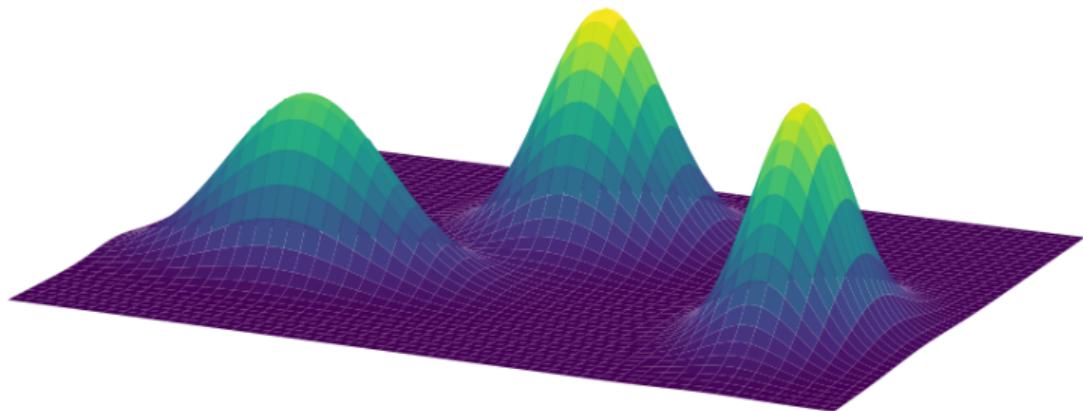
*april is  
probably a  
recursive  
identifier of a  
lab*



***who knows mixture models?***



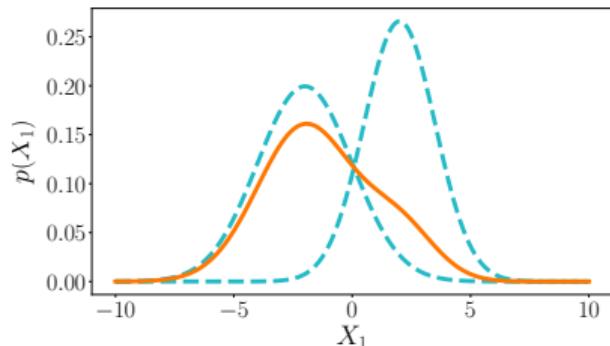
***who loves mixture models?***



$$c(\mathbf{X}) = \sum_{i=1}^K w_i c_i(\mathbf{X}), \quad \text{with} \quad w_i \geq 0, \quad \sum_{i=1}^K w_i = 1$$

# GMMS

as computational graphs

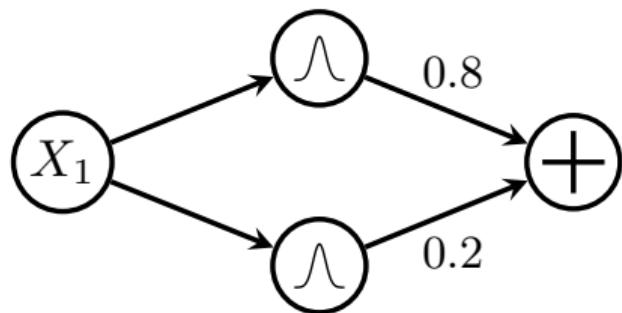


$$p(X) = w_1 \cdot p_1(X_1) + w_2 \cdot p_2(X_1)$$

⇒ translating inference to data structures...

# GMMs

as computational graphs

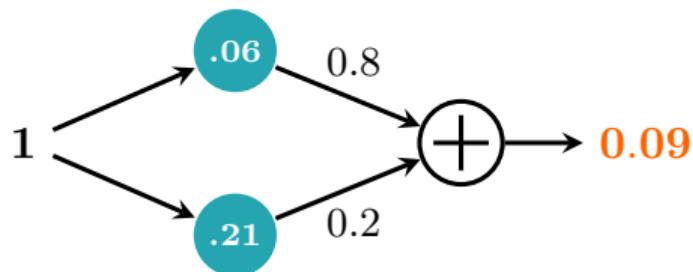


$$p(X_1) = 0.2 \cdot p_1(X_1) + 0.8 \cdot p_2(X_1)$$

⇒ ...e.g., as a weighted sum unit over Gaussian input distributions

# GMMS

as computational graphs

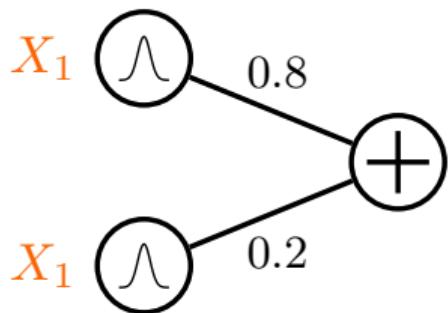


$$p(X = 1) = 0.2 \cdot p_1(X_1 = 1) + 0.8 \cdot p_2(X_1 = 1)$$

⇒ inference = feedforward evaluation

# GMMs

as computational graphs

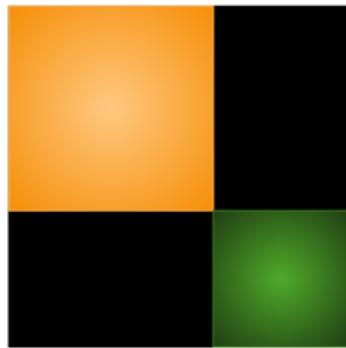
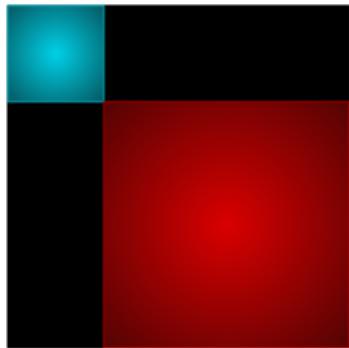


A simplified notation:

- ⇒ **scopes** attached to inputs
- ⇒ edge directions omitted

# GMMS

as computational graphs

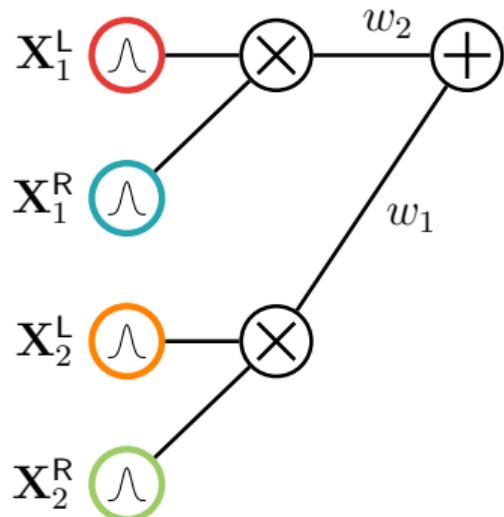


$$p(\mathbf{X}) = w_1 \cdot p_1(\mathbf{X}_1^L) \cdot p_1(\mathbf{X}_1^R) + \\ w_2 \cdot p_2(\mathbf{X}_2^L) \cdot p_2(\mathbf{X}_2^R)$$

⇒ local factorizations...

# GMMs

as computational graphs



$$p(\mathbf{X}) = w_1 \cdot p_1(\mathbf{X}_1^L) \cdot p_1(\mathbf{X}_1^R) + \\ w_2 \cdot p_2(\mathbf{X}_2^L) \cdot p_2(\mathbf{X}_2^R)$$

⇒ ...are product units

# **probabilistic circuits (PCs)**

*a grammar for tractable computational graphs*

I. A simple tractable function is a circuit

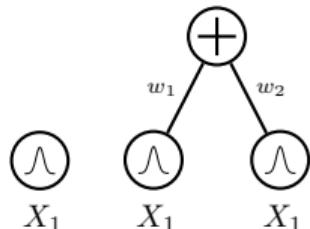
$$\bigcirc \wedge \\ X_1$$

# **probabilistic circuits (PCs)**

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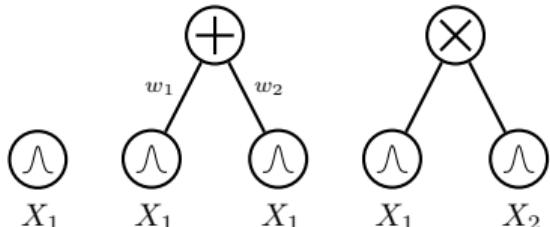
II. A weighted combination of circuits is a circuit



# **probabilistic circuits (PCs)**

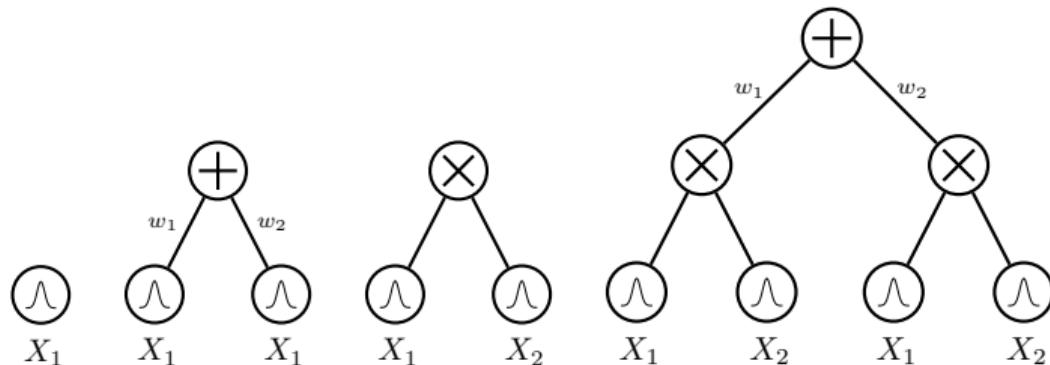
*a grammar for tractable computational graphs*

- I. A simple tractable function is a circuit
- II. A weighted combination of circuits is a circuit
- III. A product of circuits is a circuit



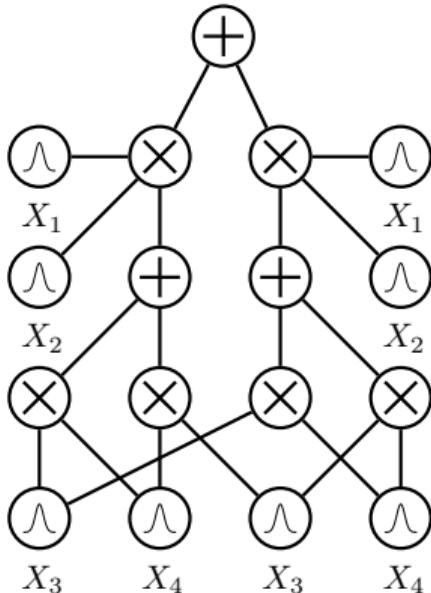
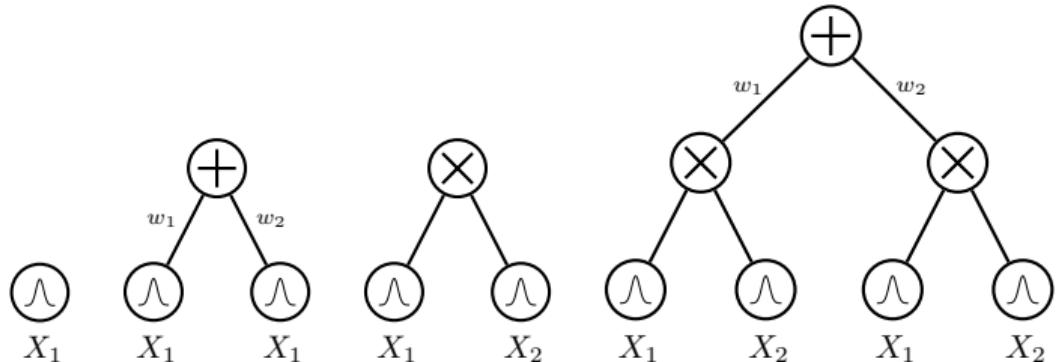
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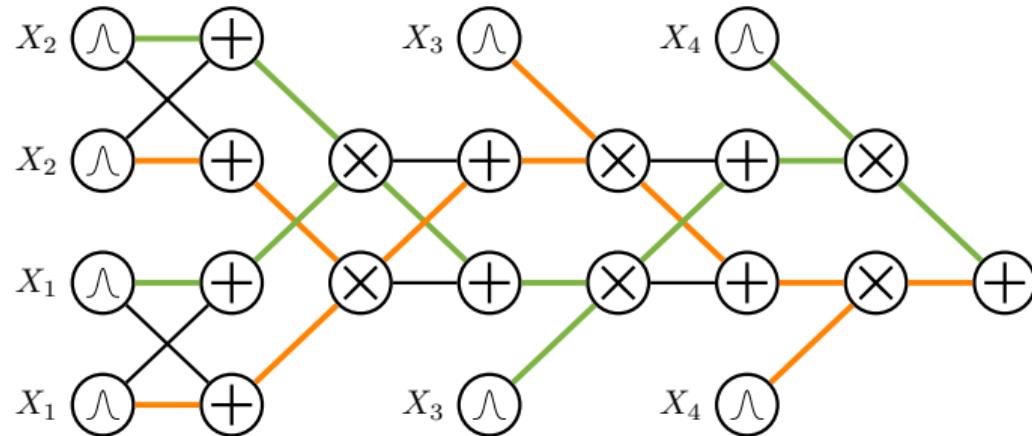


# **probabilistic circuits (PCs)**

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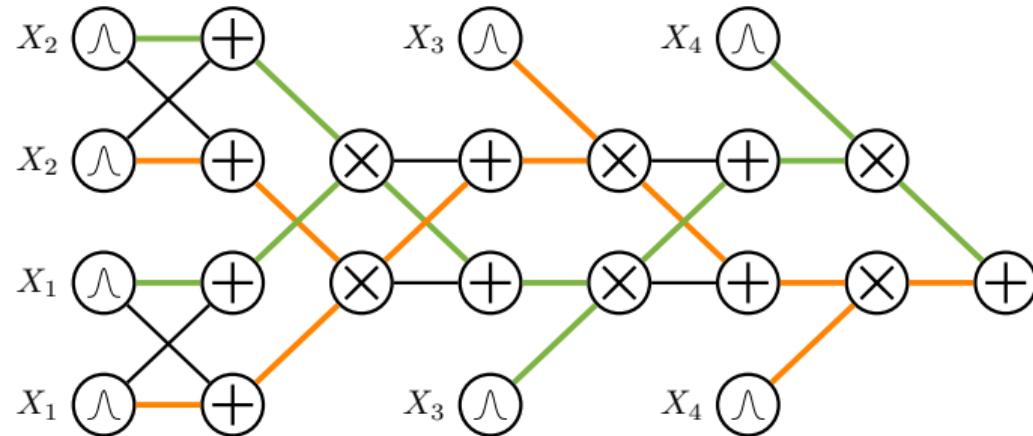


## deep mixtures



$$p(\mathbf{x}) = \sum_{\mathcal{T}} \left( \prod_{w_j \in \mathbf{w}_{\mathcal{T}}} w_j \right) \prod_{l \in \text{leaves}(\mathcal{T})} p_l(\mathbf{x})$$

## *deep mixtures*



*an exponential number of mixture components!*

*circuits*  
**(and variants)**  
*everywhere*

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# Semantic Probabilistic Layers for Neuro-Symbolic Learning

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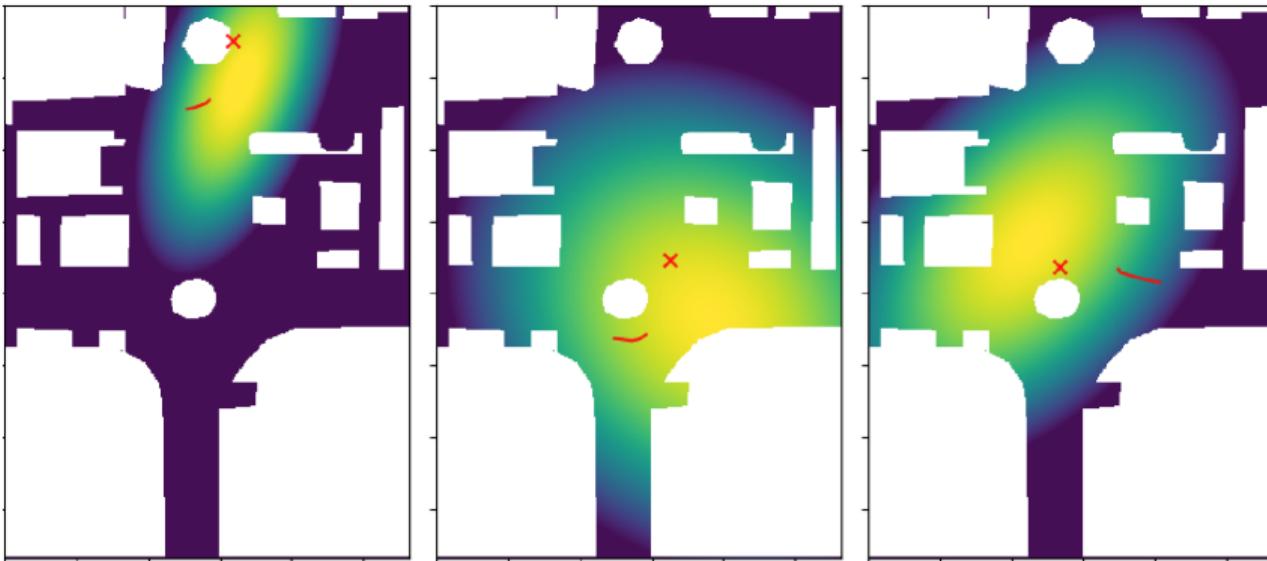
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***enforce constraints in neural networks at NeurIPS 2022***



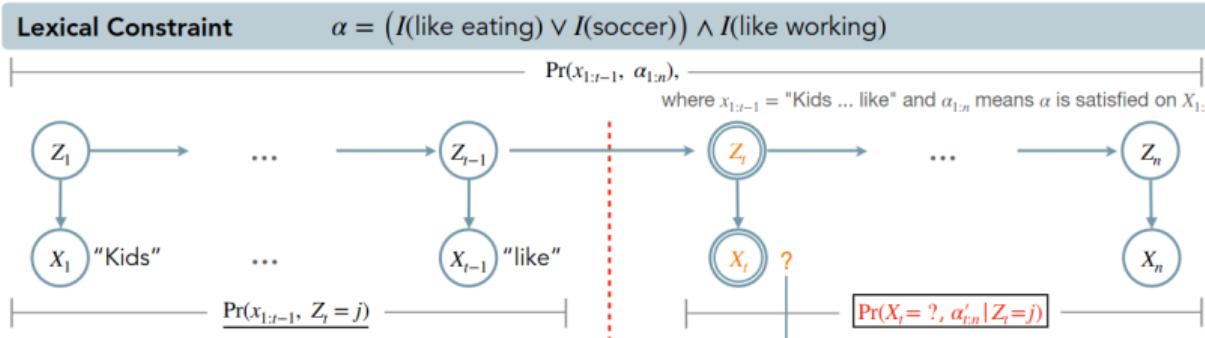
*extending it to SMT constraints*

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## Tractable Control for Autoregressive Language Generation

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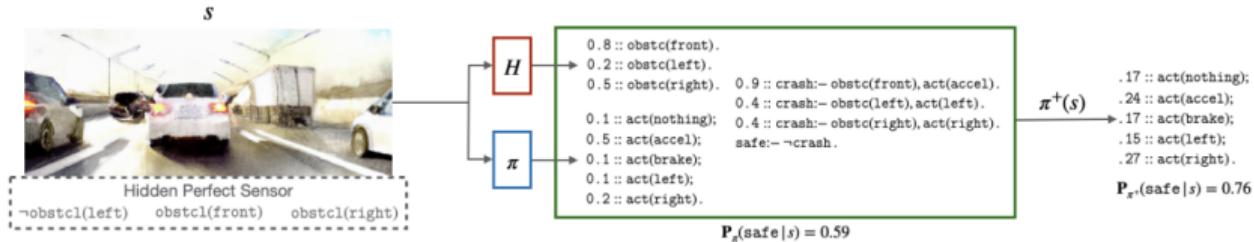
Honghua Zhang<sup>\* 1</sup> Meihua Dang<sup>\* 1</sup> Nanyun Peng<sup>1</sup> Guy Van den Broeck<sup>1</sup>



***constrained text generation with LLMs (ICML 2023)***

# Safe Reinforcement Learning via Probabilistic Logic Shields

Wen-Chi Yang<sup>1</sup>, Giuseppe Marra<sup>1</sup>, Gavin Rens and Luc De Raedt<sup>1,2</sup>



***reliable reinforcement learning (AAAI 23)***

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# How to Turn Your Knowledge Graph Embeddings into Generative Models

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*enforce constraints in knowledge graph embeddings  
oral at NeurIPS 2023*

---

# Logically Consistent Language Models via Neuro-Symbolic Integration

---

The figure displays three panels illustrating logical reasoning:

- Forward Implication**:  $A \rightarrow \neg B$ .  
A: (albatross, isA, bird)  
B: (albatross, isA, fish)  
Is an albatross a bird? (Yes) Yes.  
Is an albatross a fish? (Yes) No.
- Reverse Implication**:  $\neg B \rightarrow \neg A$ .  
B: (albatross, isNotA, organism)  
A: (albatross, isNotA, living thing)  
Is it true that an albatross is not an organism? (Yes) No.  
Is it true that an albatross is not a living thing? (Yes) No.
- Negation**:  $A \bullet \neg A$ .  
A: (computer, isA, airplane)  
A: (computer, isNotA, airplane)  
Is a computer a airplane? (Yes) No.  
Is it true that a computer is not a airplane? (Yes) No.

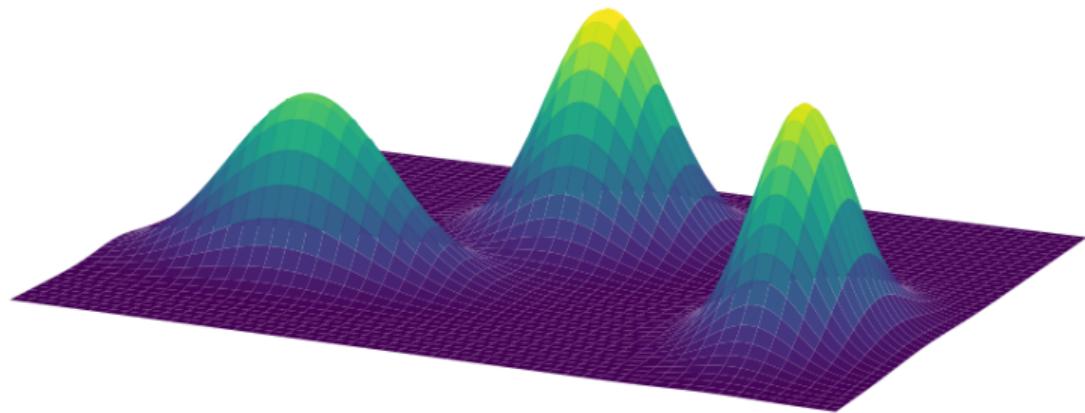
Legend:  
= LLaMA 2  
= LoCo-LLM 2

*improving logical (self-)consistency in LLMs at ICLR 2025*



***learning & reasoning with circuits in pytorch***

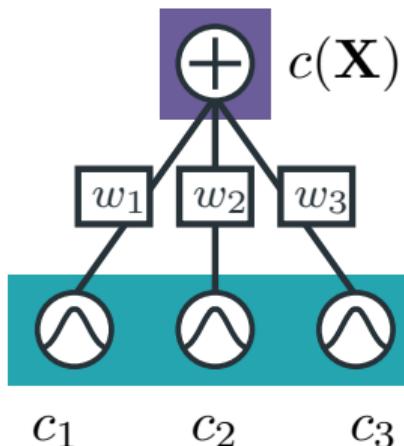
[github.com/april-tools/cirkit](https://github.com/april-tools/cirkit)



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# **additive MMs**

*are so cool!*



easily represented as shallow PCs

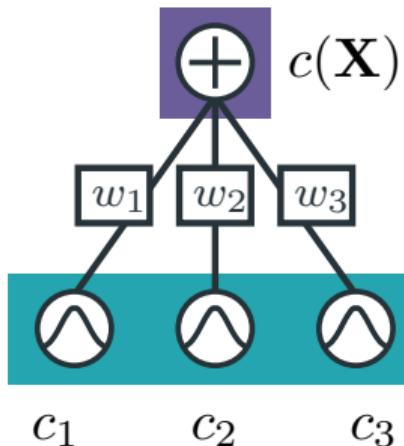
these are **monotonic** PCs

if marginals/conditionals are tractable for the components, they are tractable for the MM

universal approximators...

# **additive MMs**

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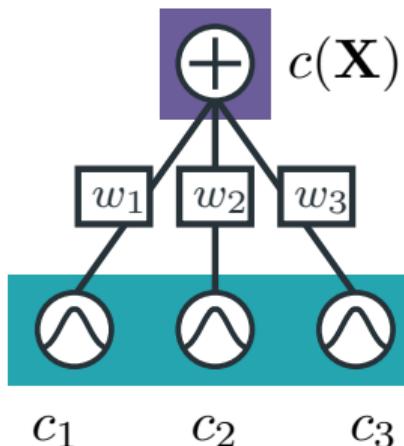
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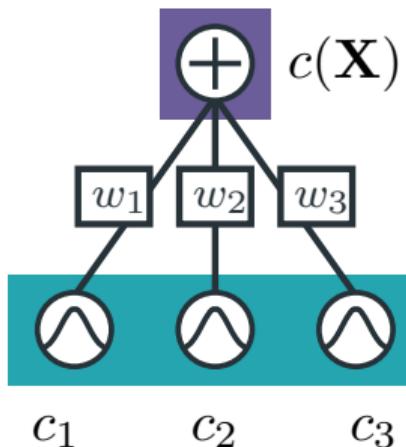
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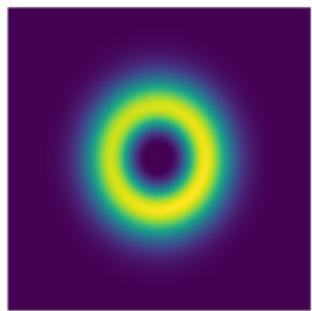
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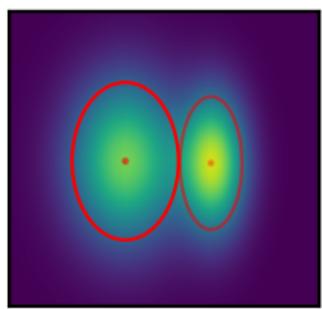
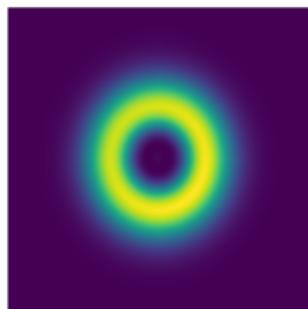
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universal approximators...

*however...*

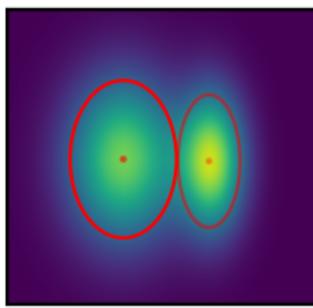
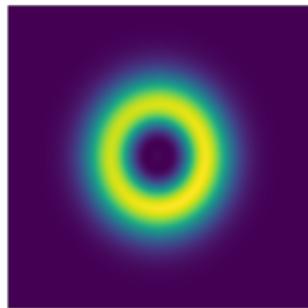


*however...*

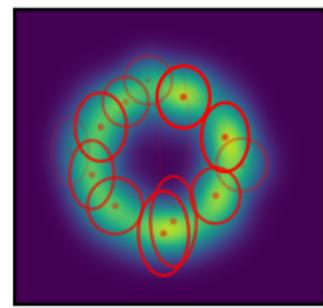


GMM ( $K = 2$ )

*however...*

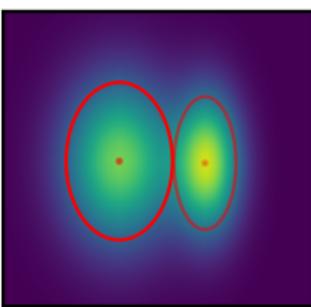
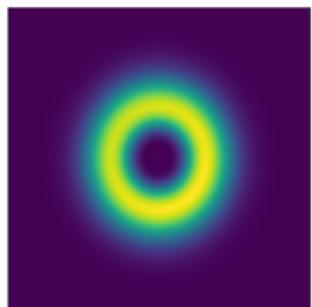


GMM ( $K = 2$ )

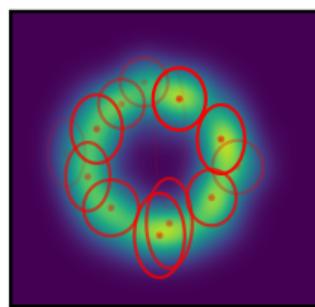


GMM ( $K = 16$ )

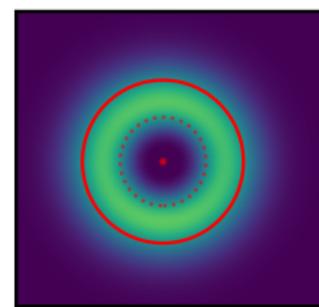
**however...**



GMM ( $K = 2$ )



GMM ( $K = 16$ )

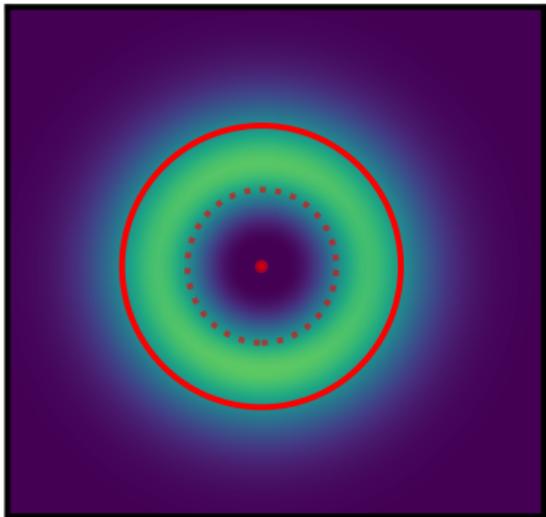


nGMM<sup>2</sup> ( $K = 2$ )

*spoiler*

**shallow mixtures  
with negative parameters  
can be *exponentially more compact* than  
deep ones with positive parameters.**

# **subtractive MMs**



also called negative/signed/**subtractive** MMs

⇒ or **non-monotonic** circuits,...

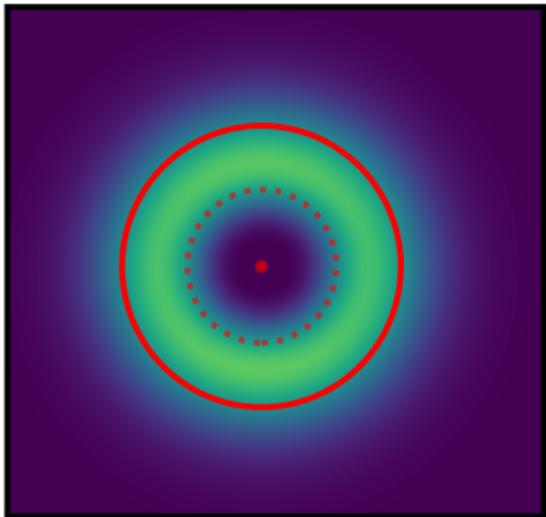
**issue:** how to preserve non-negative outputs?

well understood for simple parametric forms

e.g., Weibulls, Gaussians

⇒ constraints on variance, mean

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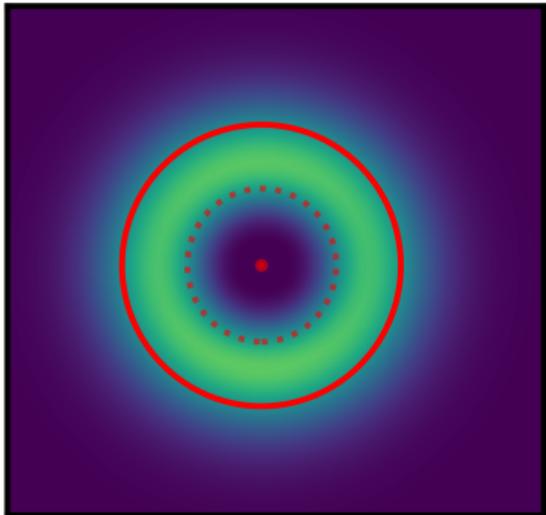
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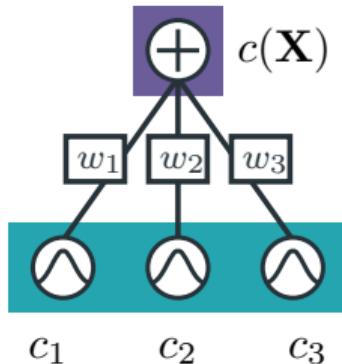
**tl;dr**

***“Understand when and how  
we can use negative parameters  
in deep subtractive mixture models”***

**tl;dr**

***“Understand when and how  
we can use negative parameters  
in deep **non-monotonic squared circuits**”***

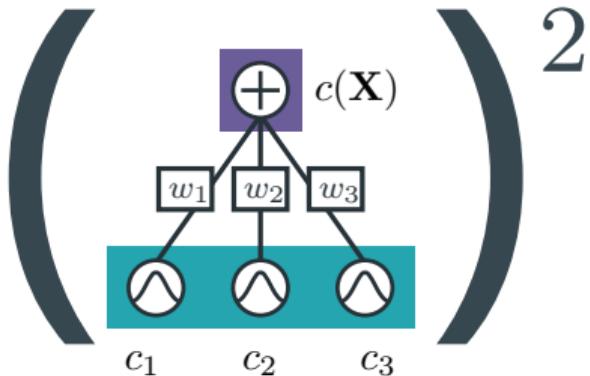
## *subtractive MMs as circuits*



a **non-monotonic** smooth and (structured) decomposable circuit  
⇒ possibly with negative outputs

$$c(\mathbf{X}) = \sum_{i=1}^K w_i c_i(\mathbf{X}), \quad w_i \in \mathbb{R},$$

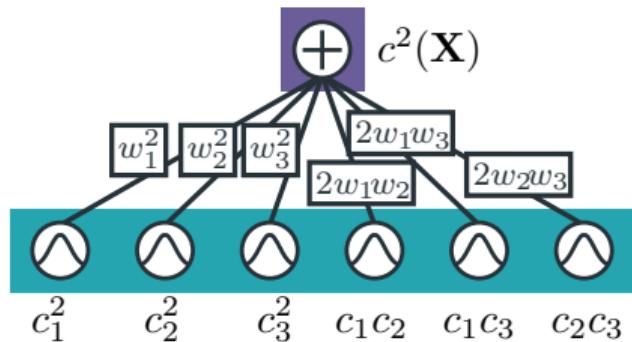
## *squaring shallow MMs*



$$c^2(\mathbf{X}) = \left( \sum_{i=1}^K w_i c_i(\mathbf{X}) \right)^2$$

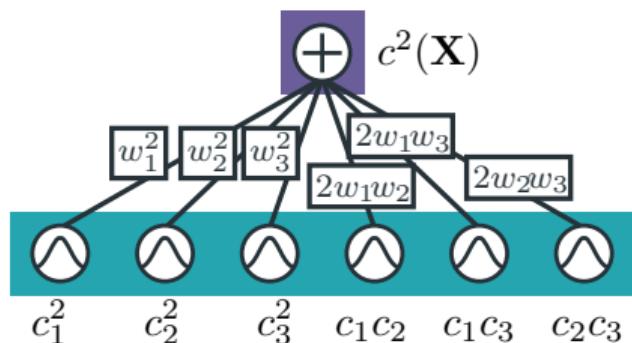
$\Rightarrow$  ensure non-negative output

# *squaring shallow MMs*



$$\begin{aligned} c^2(\mathbf{X}) &= \left( \sum_{i=1}^K w_i c_i(\mathbf{X}) \right)^2 \\ &= \sum_{i=1}^K \sum_{j=1}^K w_i w_j c_i(\mathbf{X}) c_j(\mathbf{X}) \end{aligned}$$

## *squaring shallow MMs*



$$\begin{aligned} c^2(\mathbf{X}) &= \left( \sum_{i=1}^K w_i c_i(\mathbf{X}) \right)^2 \\ &= \sum_{i=1}^K \sum_{j=1}^K w_i w_j c_i(\mathbf{X}) c_j(\mathbf{X}) \end{aligned}$$

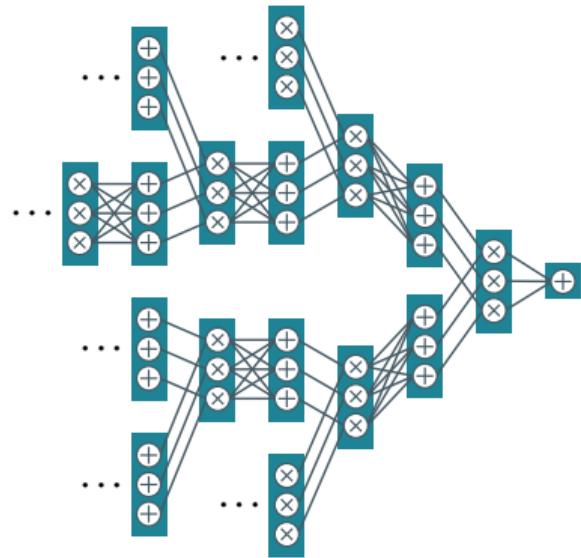
still a smooth and (str) decomposable PC with  $\mathcal{O}(K^2)$  components!

$\Rightarrow$  but still  $\mathcal{O}(K)$  parameters

*wait...*

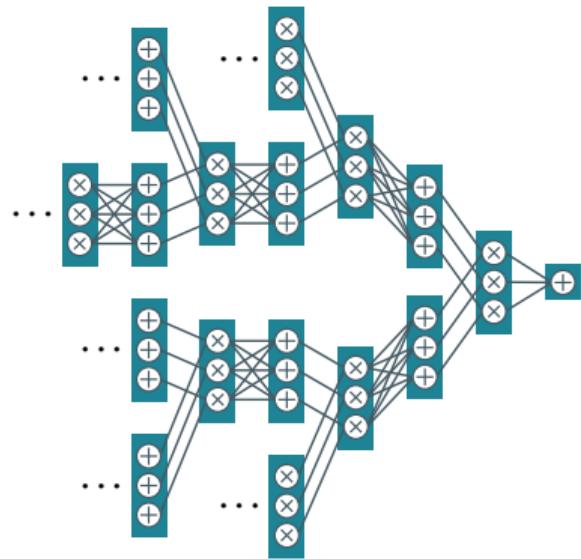
*“do negative parameters  
really boost expressiveness?  
and...always?”*

# theorem



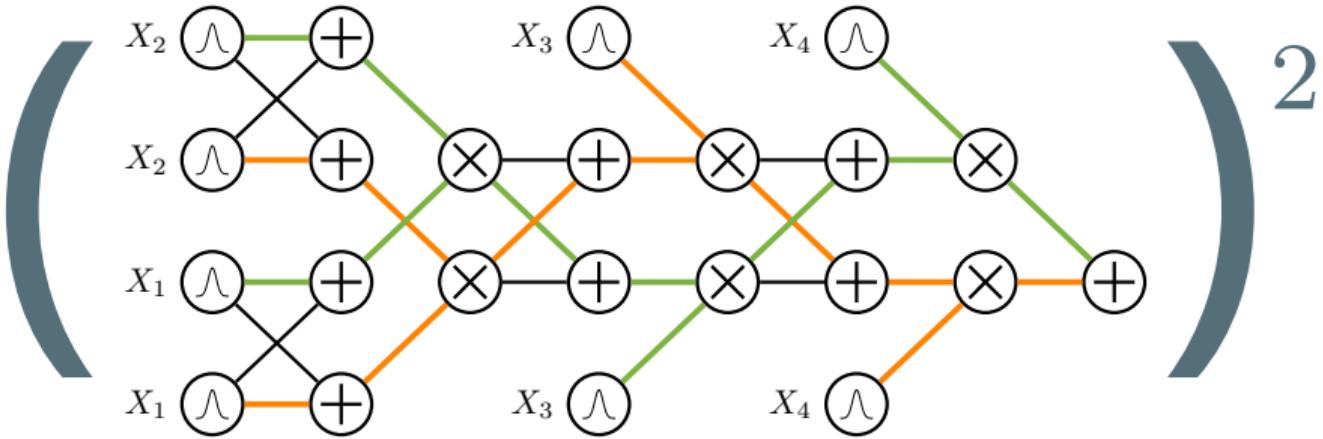
$\exists p$  requiring exponentially large  
monotonic circuits...

# *theorem*



$$\left( \begin{array}{c} \text{...} \\ \text{...} \\ \text{...} \end{array} \right)^2$$

**...but compact  
squared non-monotonic circuits**



**how to efficiently square (and *renormalize*) a deep PC?**

# *compositional inference I*



```
1 from cirkit.symbolic.functional import integrate, multiply  
2  
3 #  
4 # create a deep circuit  
5 c = build_symbolic_circuit('quad-tree-4')  
6  
7 #  
8 # compute the partition function of c^2  
9 def renormalize(c):  
10     c2 = multiply(c, c)  
11     return integrate(c2)
```

# **probabilistic circuits (PCs)**

*the unit-wise definition*

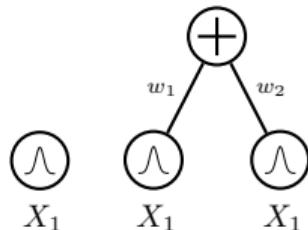
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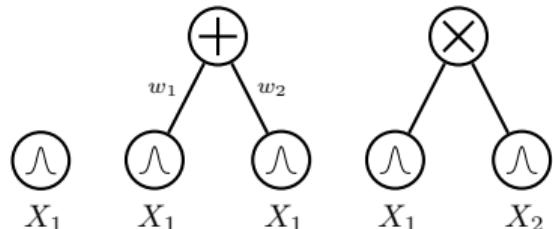
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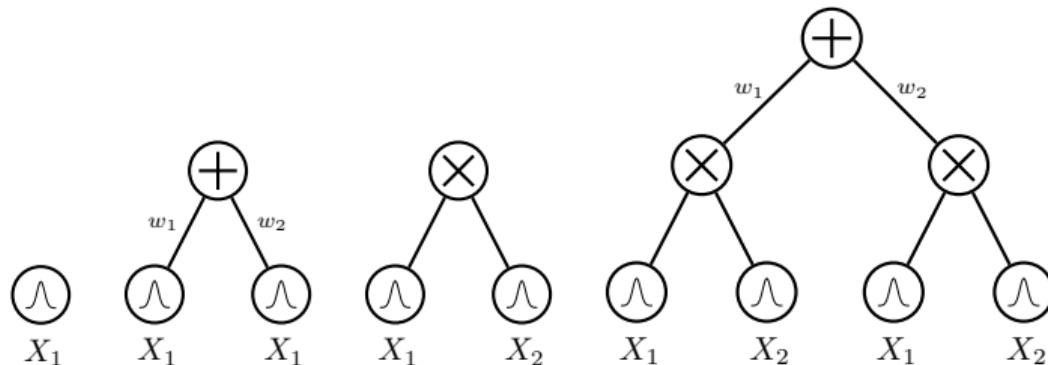
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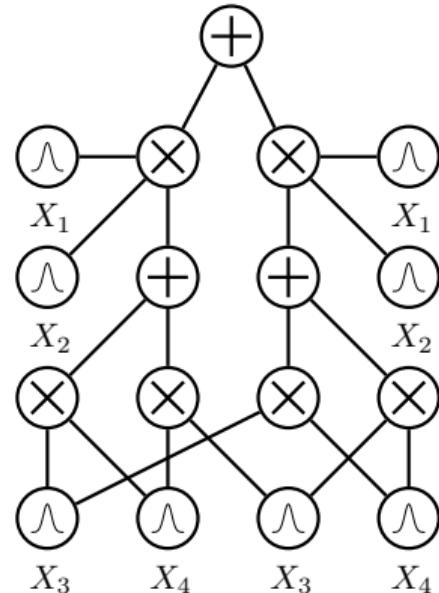
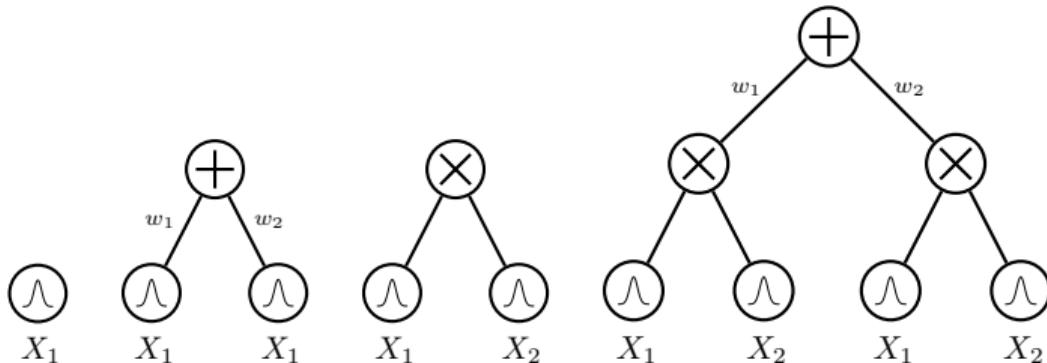
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# **probabilistic circuits (PCs)**

*a tensorized definition*

- I. A set of tractable functions is a circuit layer



# **probabilistic circuits (PCs)**

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- I. A set of tractable functions is a circuit layer
- II. A linear projection of a layer is a circuit layer

$$c(\mathbf{x}) = \mathbf{W}l(\mathbf{x})$$



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- III. The product of two layers is a circuit layer

$$c(\mathbf{x}) = \mathbf{l}(\mathbf{x}) \odot \mathbf{r}(\mathbf{x}) \quad // \text{ Hadamard}$$



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$$c(\mathbf{x}) = \mathbf{l}(\mathbf{x}) \odot \mathbf{r}(\mathbf{x}) \quad // \text{ Hadamard}$$

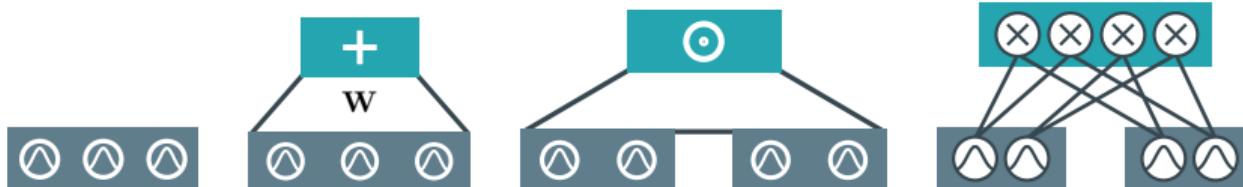


# **probabilistic circuits (PCs)**

*a tensorized definition*

- I. A set of tractable functions is a circuit layer
- II. A linear projection of a layer is a circuit layer
- III. The product of two layers is a circuit layer

$$c(\mathbf{x}) = \text{vec}(\mathbf{l}(\mathbf{x})\mathbf{r}(\mathbf{x})^\top) \quad // \text{ Kronecker}$$

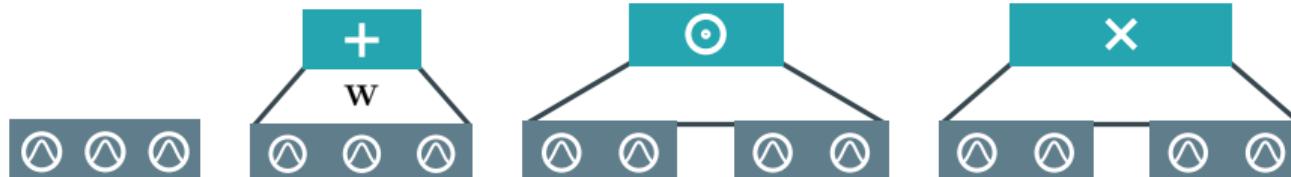


# **probabilistic circuits (PCs)**

*a tensorized definition*

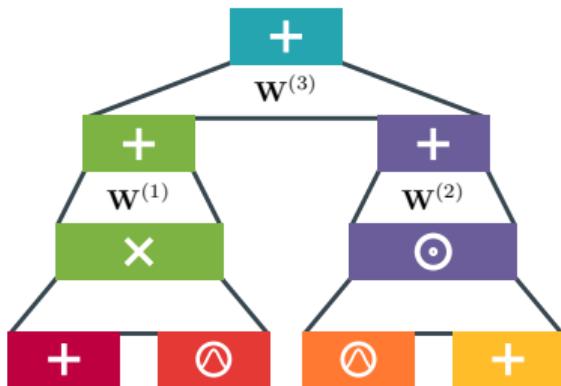
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# **probabilistic circuits (PCs)**

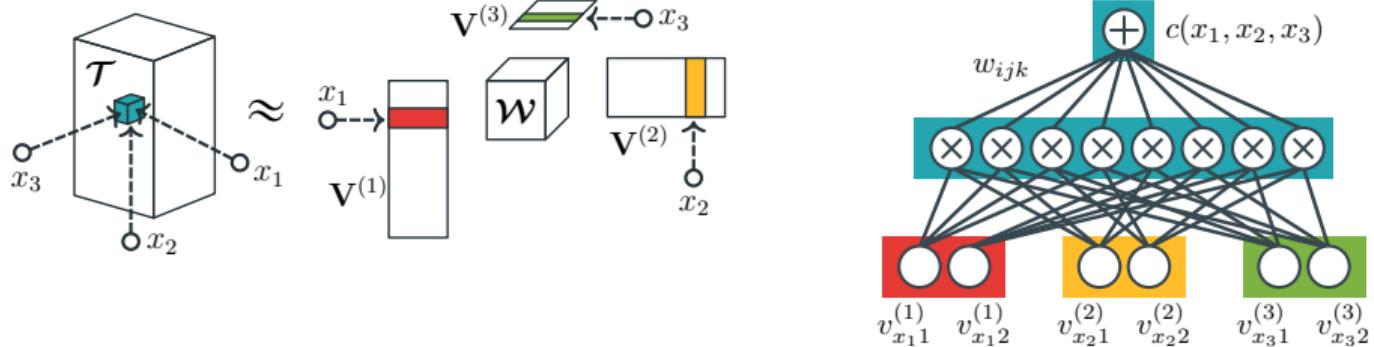
*a tensorized definition*



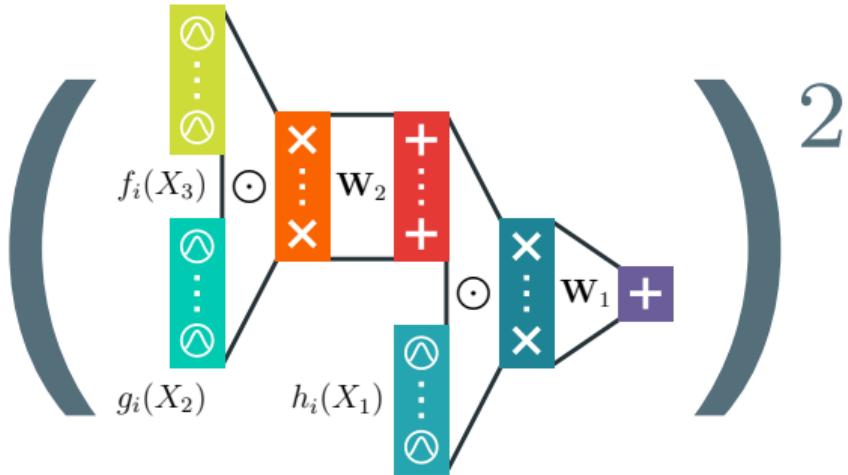
- I. A set of tractable functions is a circuit layer
  - II. A linear projection of a layer is a circuit layer
  - III. The product of two layers is a circuit layer
- stack layers to build a deep circuit!**

# *circuits layers*

*as tensor factorizations*



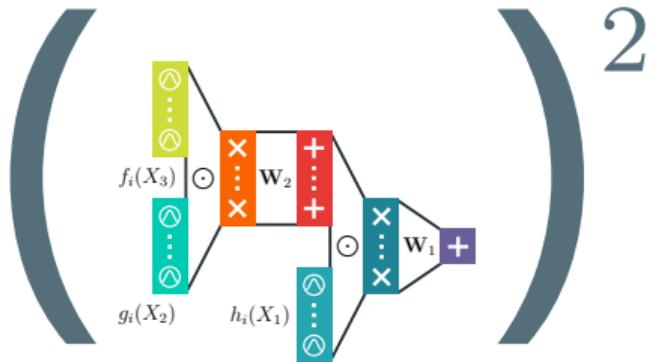
Loconte et al., "What is the Relationship between Tensor Factorizations and Circuits (and How Can We Exploit it)?", TMLR, 2025



**how to efficiently square (and *renormalize*) a deep PC?**

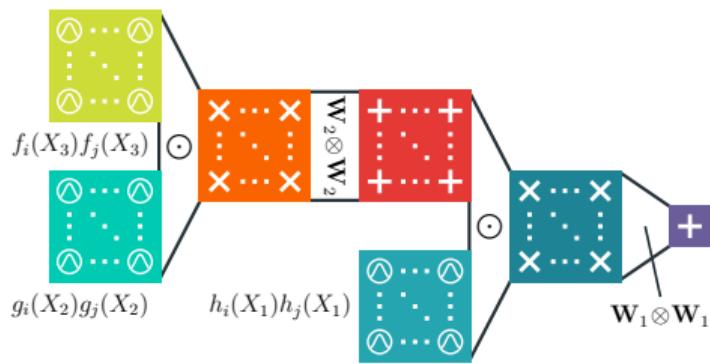
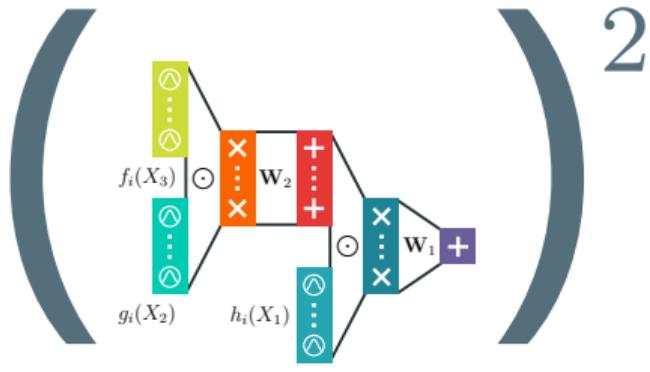
# *squaring deep PCs*

*the tensorized way*



# *squaring deep PCs*

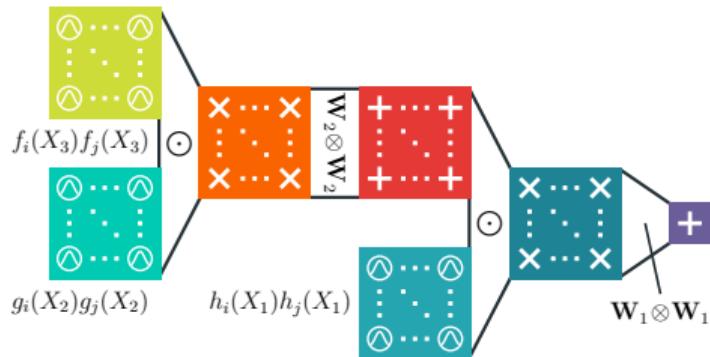
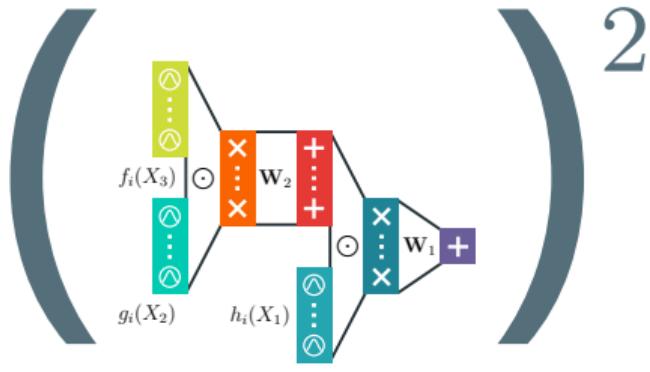
*the tensorized way*



**squaring a circuit = squaring layers**

# *squaring deep PCs*

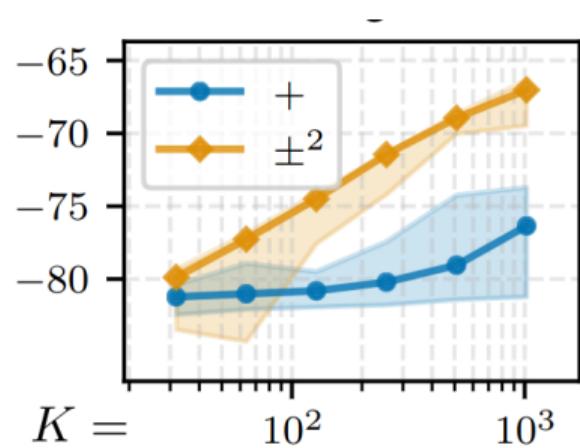
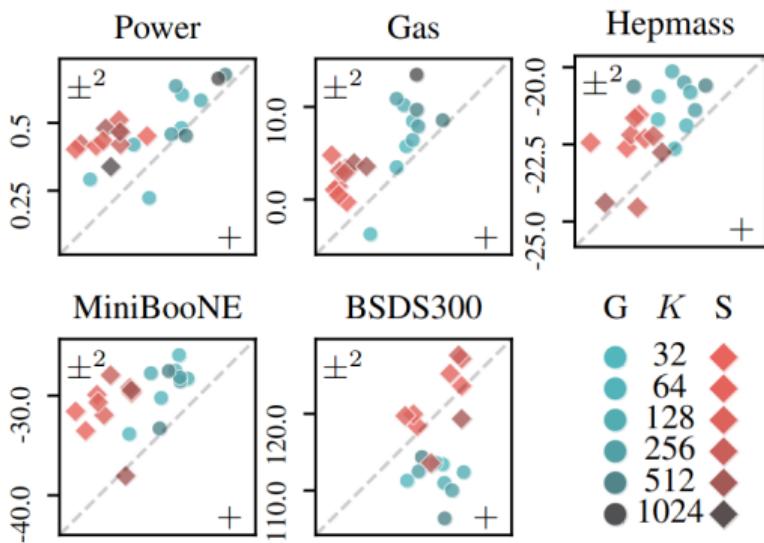
*the tensorized way*



**exactly compute**  $\int c(x)c(x)dX$  **in time**  $O(LK^2)$

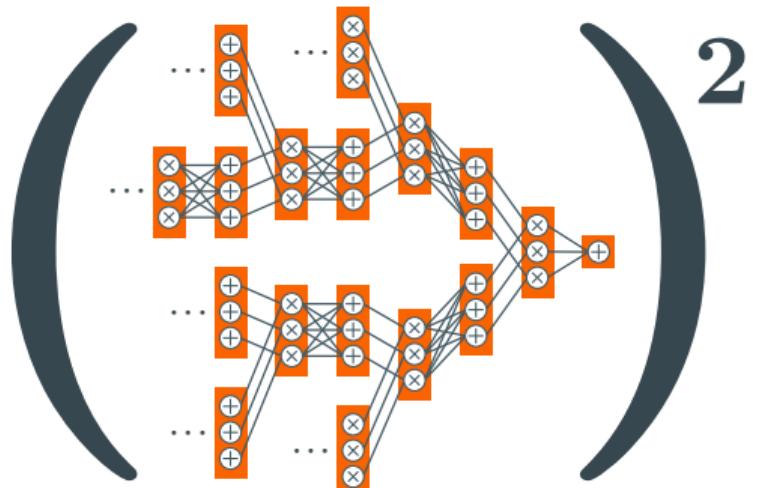
# how more expressive?

for the ML crowd



# *theorem*

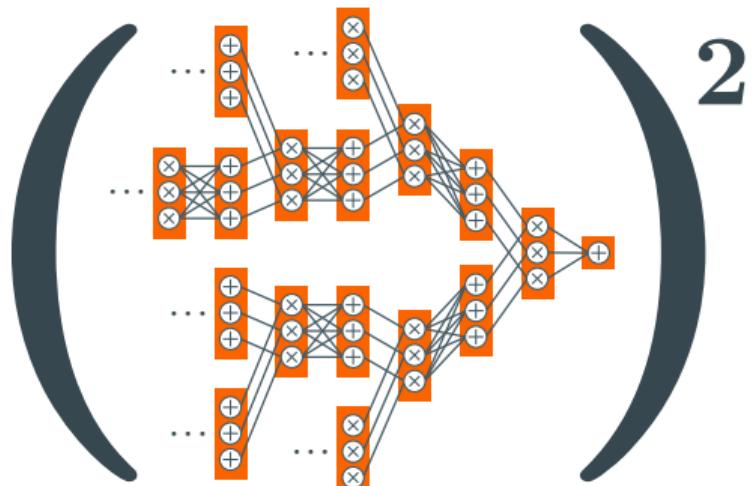
$\exists p$  requiring exponentially large  
squared non-mono circuits...

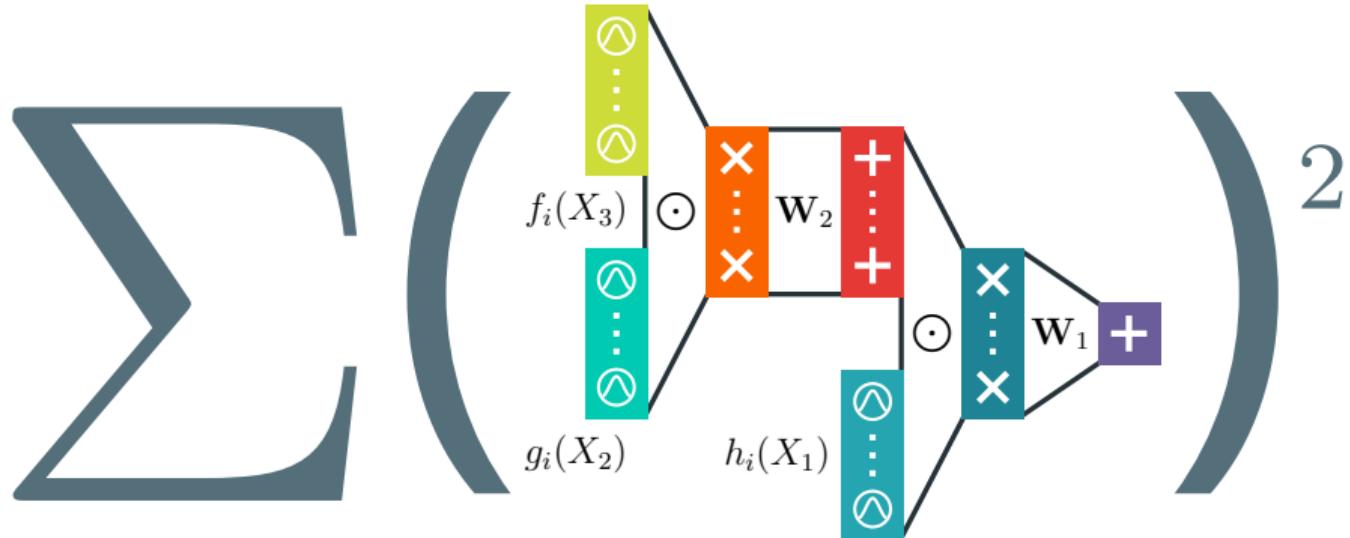


# *theorem*



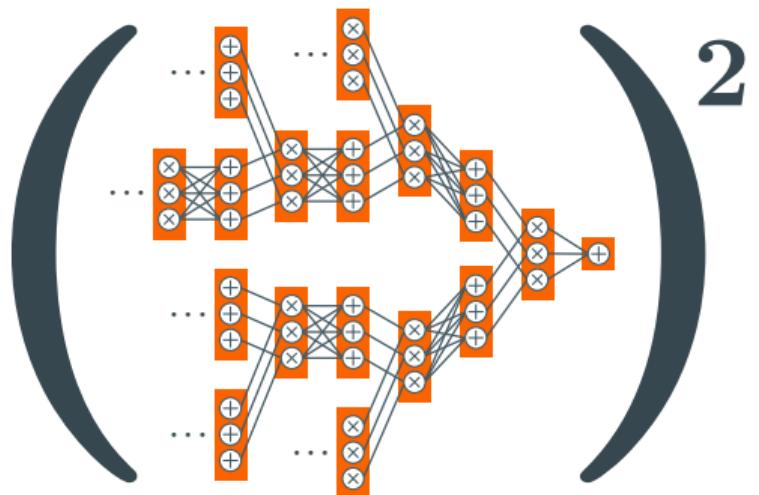
**...but compact  
monotonic circuits...!**





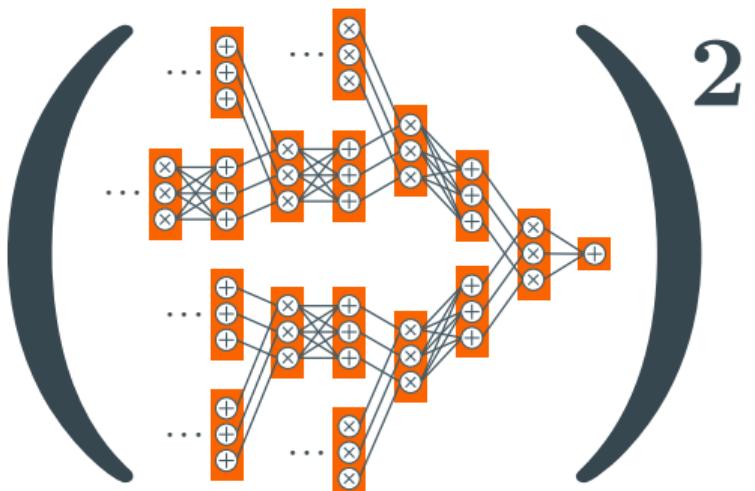
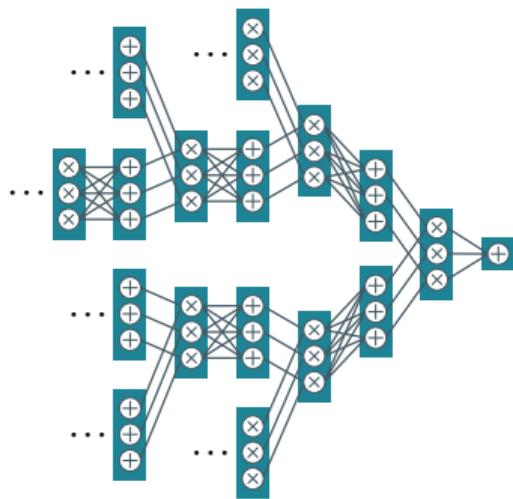
***what if we use more than one square?***

# *theorem*



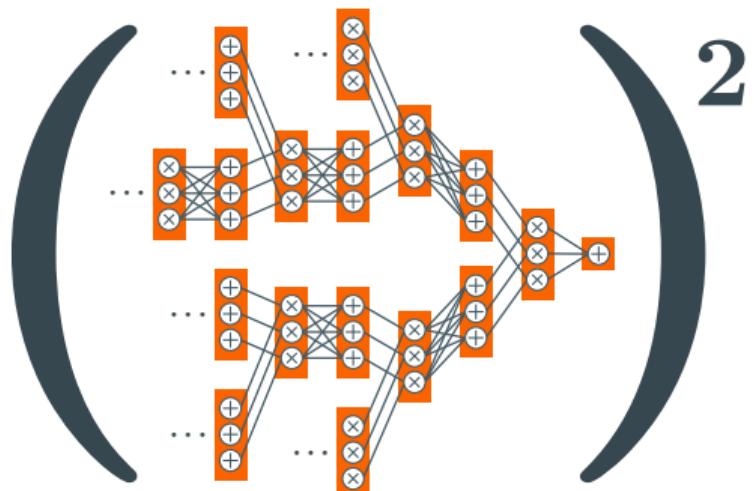
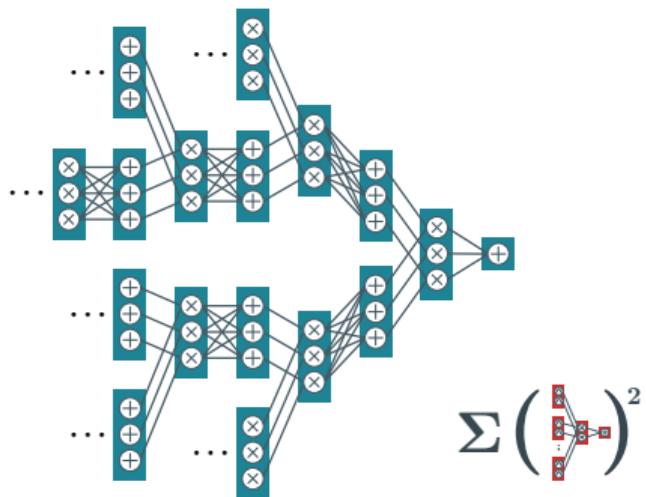
$\exists p$  requiring exponentially large **squared non-mono circuits...**

# *theorem*



**...exponentially large monotonic circuits...**

# theorem



...but compact **SOS circuits...**!

$$\pm_{\text{sd}} = \Delta \Sigma_{\text{cmp}}^2$$

(Theorem 5)

$$\Sigma_{\text{cmp}}^2 = \text{psd}$$

(Proposition 2)

$$+_{\text{sd}}$$

•  
Open Question 1

•  
Open Question 2

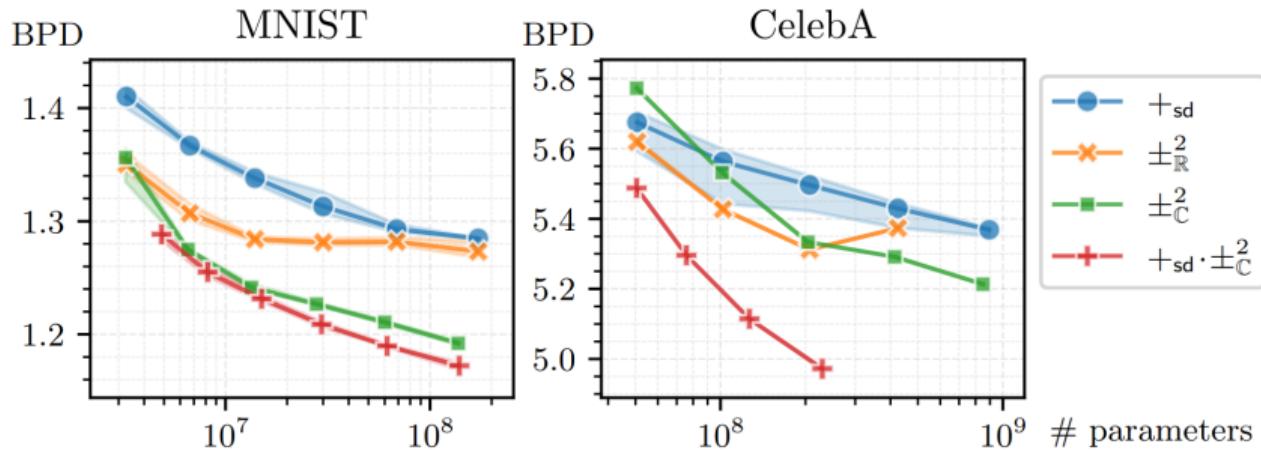
$$\pm_{\mathbb{R}}^2$$

• UDISJ  
(Theorem 0)

• UPS  
(Theorem 2)

• UTQ  
(Theorem B.3)

***a hierarchy of subtractive mixtures***



***complex circuits are SOS (and scale better!)***

# *compositional inference I*



```
1 from cirkit.symbolic.functional import integrate, multiply,
2     conjugate
3
4 # create a deep circuit with complex parameters
5 c = build_symbolic_complex_circuit('quad-tree-4')
6
7 # compute the partition function of c^2
8 def renormalize(c):
9     c1 = conjugate(c)
10    c2 = multiply(c, c1)
11    return integrate(c2)
```

# **approximate inference**

e.g., via sampling

Can we use a subtractive mixture model to approximate expectations?

$$\mathbb{E}_{\mathbf{x} \sim q(\mathbf{x})} [f(\mathbf{x})] \approx \frac{1}{S} \sum_{i=1}^S f(\mathbf{x}^{(i)}) \quad \text{with} \quad \mathbf{x}^{(i)} \sim q(\mathbf{x})$$

$\Rightarrow$  but how to sample from  $q$ ?

# approximate inference

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$\Rightarrow$  but how to sample from  $q$ ?

use **autoregressive inverse transform sampling**:

$$x_1 \sim q(x_1), \quad x_i \sim q(x_i | \mathbf{x}_{<i}) \quad \text{for } i \in \{2, \dots, d\}$$

$\Rightarrow$  can be slow for large dimensions, requires **inverting the CDF**

# **approximate inference**

*difference of expectation estimator*

**Idea:** represent  $q$  as a difference of two additive mixtures

$$q(\mathbf{x}) = Z_+ \cdot q_+(\mathbf{x}) - Z_- \cdot q_-(\mathbf{x})$$

⇒ *expectations will break down in two “parts”*

# **approximate inference**

*difference of expectation estimator*

**Idea:** represent  $q$  as a difference of two additive mixtures

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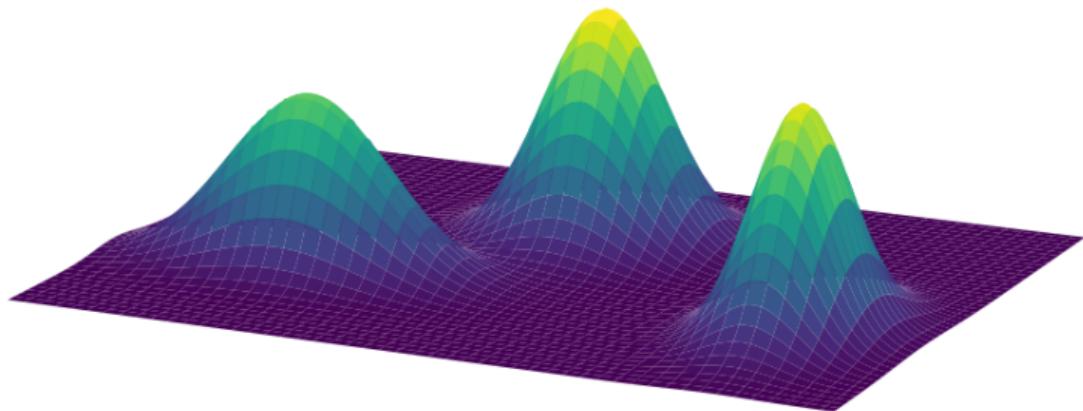
$$\frac{Z_+}{S_+} \sum_{s=1}^{S_+} f(\mathbf{x}_+^{(s)}) - \frac{Z_-}{S_-} \sum_{s=1}^{S_-} f(\mathbf{x}_-^{(s)}), \text{ where } \begin{aligned} \mathbf{x}_+^{(s)} &\sim q_+(\mathbf{x}_+) \\ \mathbf{x}_-^{(s)} &\sim q_-(\mathbf{x}_-) \end{aligned}, \quad (1)$$

# *approximate inference*

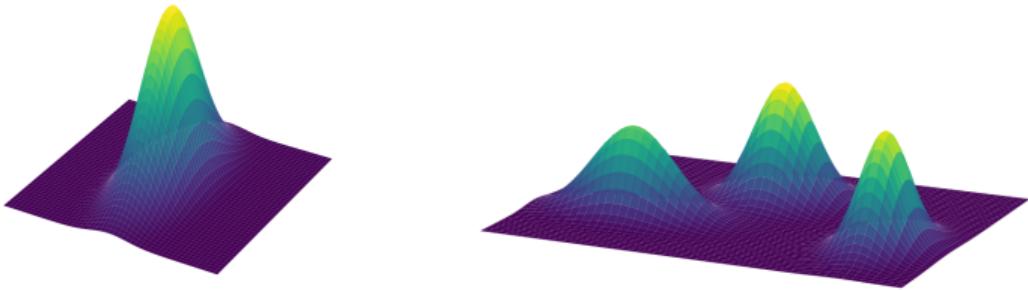
*difference of expectation estimator*

Method	$d$	Number of components ( $K$ )					
		2		4		6	
		$\log( \hat{I} - I )$	Time (s)	$\log( \hat{I} - I )$	Time (s)	$\log( \hat{I} - I )$	Time (s)
ΔExS	16	-19.507 ± 1.025	0.293 ± 0.004	-19.062 ± 0.823	1.049 ± 0.077	-19.497 ± 1.974	2.302 ± 0.159
ARITS	16	-19.111 ± 1.103	7.525 ± 0.038	-19.299 ± 1.611	7.52 ± 0.023	-18.739 ± 1.024	7.746 ± 0.032
ΔExS	32	-48.411 ± 1.265	0.325 ± 0.012	-48.046 ± 0.972	1.027 ± 0.107	-48.34 ± 0.814	2.213 ± 0.177
ARITS	32	-47.897 ± 1.165	15.196 ± 0.059	-47.349 ± 0.839	15.535 ± 0.059	-47.3 ± 0.978	17.371 ± 0.06
ΔExS	64	-108.095 ± 1.094	0.38 ± 0.034	-107.56 ± 0.616	0.9 ± 0.14	-107.653 ± 0.945	1.512 ± 0.383
ARITS	64	-107.898 ± 1.129	30.459 ± 0.098	-107.33 ± 0.929	33.892 ± 0.119	-107.374 ± 1.138	52.02 ± 0.127

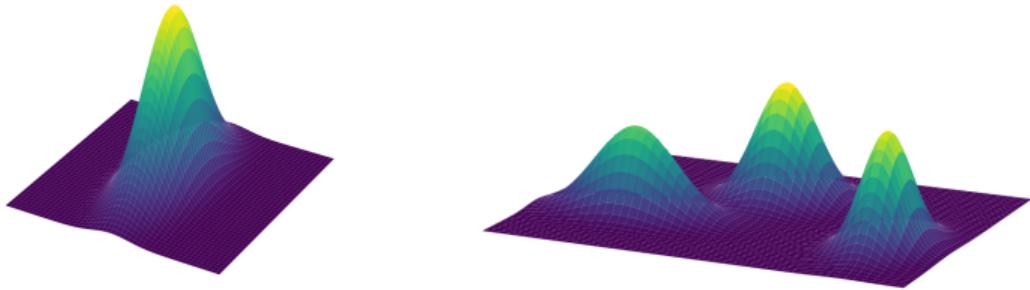
*faster than autoregressive sampling*



*oh mixtures, you're so fine you blow my mind!*



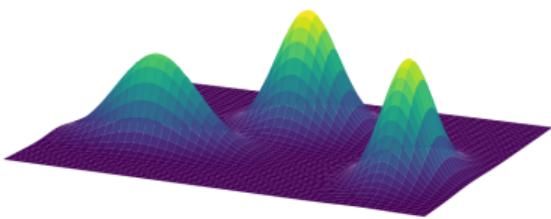
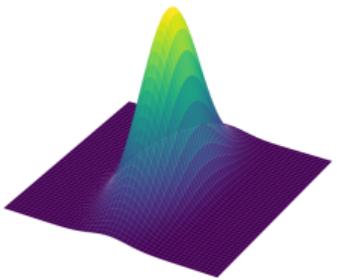
$$p(\mathbf{X}) \quad \xrightarrow{\text{orange arrow}} \quad \sum_{i=1}^K w_i p_i(\mathbf{X}) \quad w_i > 0$$



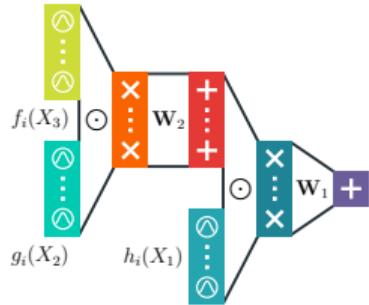
$$p(\mathbf{X}) \longrightarrow \sum_{i=1}^K w_i p_i(\mathbf{X}) \quad w_i > 0$$

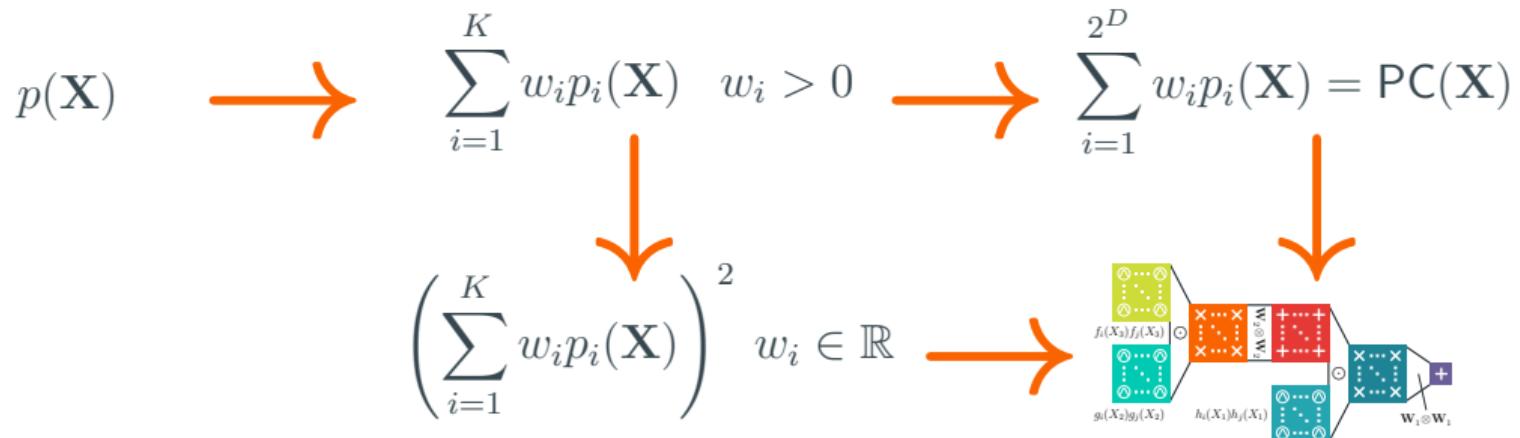
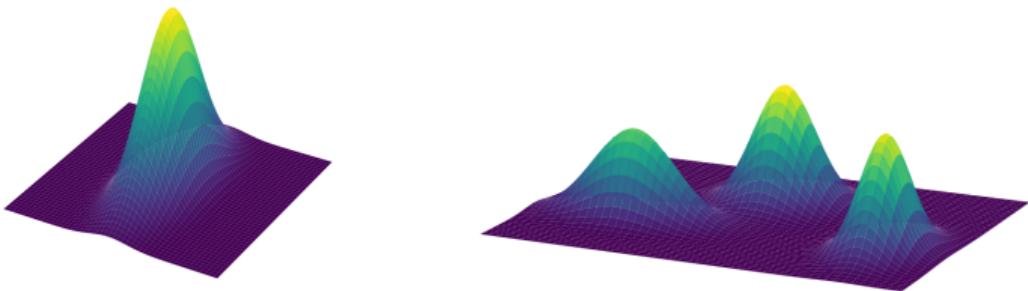
*"if someone publishes a paper on **model A**, there will be a paper about  
**mixtures of A** soon, with high probability"*

A. Vergari



$$p(\mathbf{X}) \rightarrow \sum_{i=1}^K w_i p_i(\mathbf{X}) \quad w_i > 0 \rightarrow \sum_{i=1}^{2^D} w_i p_i(\mathbf{X}) = \text{PC}(\mathbf{X})$$

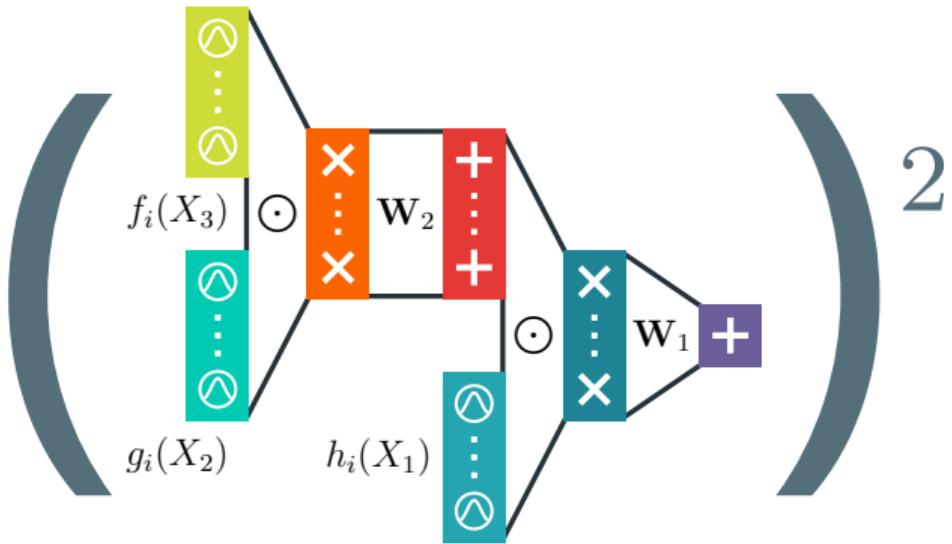






***learning & reasoning with circuits in pytorch***

[github.com/april-tools/cirkit](https://github.com/april-tools/cirkit)



**questions?**