



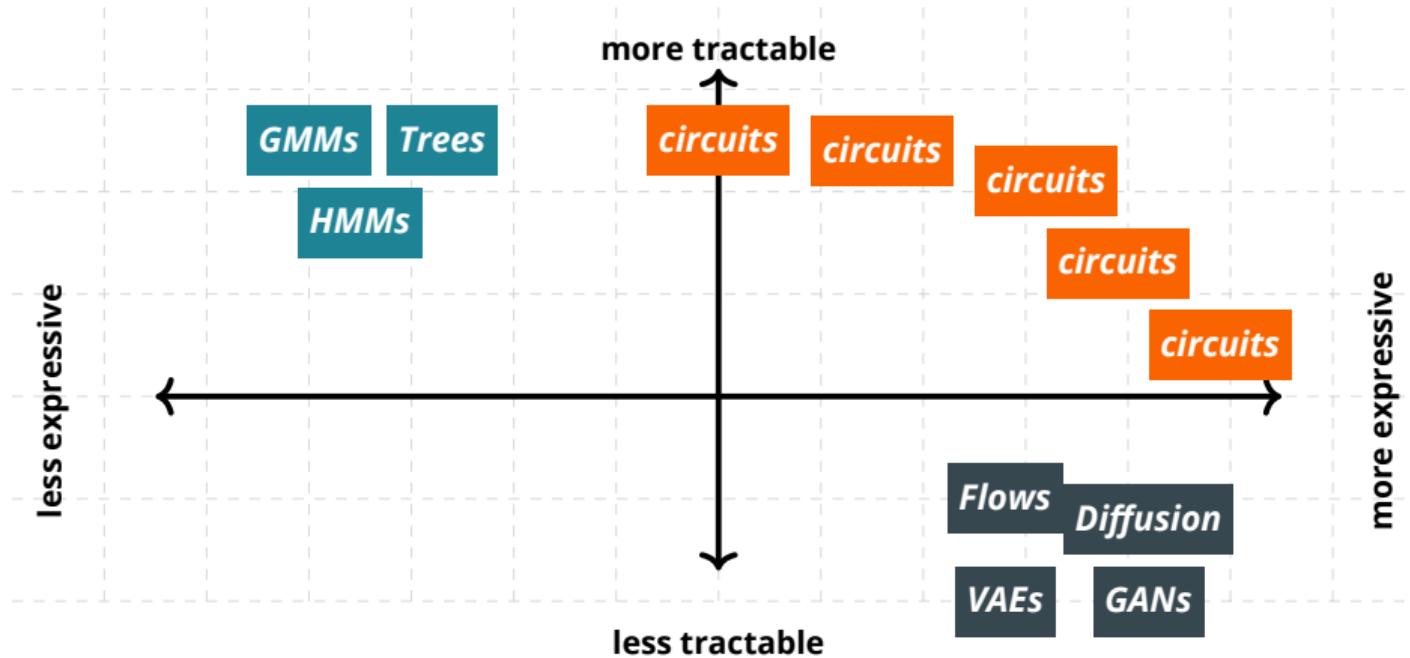
advanced

probabilistic modeling & reasoning

antonio vergari (he/him)

 @tetraduzione

18th Oct 2024 - PICS PhD School Copenhagen



trade-off tractability vs expressiveness

Reasoning about ML models



q₁

*"What is the probability of a treatment for a patient with **unavailable records**?"*



q₂

*"How **fair** is the prediction is a certain protected attribute changes?"*



q₃

*"Can we certify no **adversarial examples** exist?"*

Reasoning about ML models



q₁ $\int p(\mathbf{x}_o, \mathbf{x}_m) d\mathbf{X}_m$
(missing values)

q₂ $\mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{X}_c | X_s=0)} [f_0(\mathbf{x}_c)] - \mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{X}_c | X_s=1)} [f_1(\mathbf{x}_c)]$
(fairness)

q₃ $\mathbb{E}_{\mathbf{e} \sim \mathcal{N}(\mathbf{0}, \sigma^2 \mathbf{I}_D)} [f(\mathbf{x} + \mathbf{e})]$
(adversarial robust.)

...in the language of probabilities

Reasoning about ML models



q₁ $\int p(\mathbf{x}_o, \mathbf{x}_m) d\mathbf{X}_m$
(missing values)

q₂ $\mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{X}_c | X_s=0)} [f_0(\mathbf{x}_c)] - \mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{X}_c | X_s=1)} [f_1(\mathbf{x}_c)]$
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(adversarial robust.)

it is crucial we compute them exactly and in polytime!

Reasoning about ML models



q₁ $\int p(\mathbf{x}_o, \mathbf{x}_m) d\mathbf{X}_m$
(missing values)

q₂ $\mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{X}_c | X_s=0)} [f_0(\mathbf{x}_c)] - \mathbb{E}_{\mathbf{x}_c \sim p(\mathbf{X}_c | X_s=1)} [f_1(\mathbf{x}_c)]$
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q₃ $\mathbb{E}_{\mathbf{e} \sim \mathcal{N}(\mathbf{0}, \sigma^2 \mathbf{I}_D)} [f(\mathbf{x} + \mathbf{e})]$
(adversarial robust.)

it is crucial we compute them tractably!

Which structural properties

for complex reasoning



smooth + decomposable

Which structural properties

for complex reasoning



smooth + decomposable



???????



decomposable

Which structural properties

for complex reasoning



smooth + decomposable



compatibility



decomposable

The problem!

queries (behaviors)



	M_1	M_2	M_3	M_4	M_5	...
HMMs	✓	?	✓	✓	✗	...
decision trees	✓	✓	✓	✓	✓	...
neural nets w/ ReLUs	✗	?	?	✓	✗	...
tensor factorizations	✓	?	✓	✓	?	...
transformers	?	?	?	?	?	...
...

model classes

✓ reliable & efficient algo

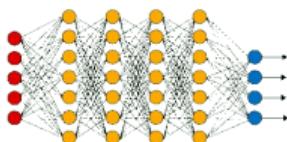
✗ intractable

? do not know

"How uncertain will the classifier be if some input is missing?"

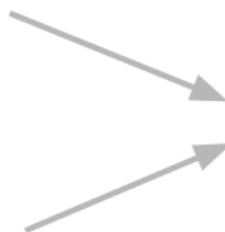
behaviors

to be inspected



ML models

(classifiers, generative models, ...)



???



```
def var(p, f):
    r = pow(f, 2)
    t = mult(r, p)
    return integrate(t)
```

efficient & reliable complex reasoning routines

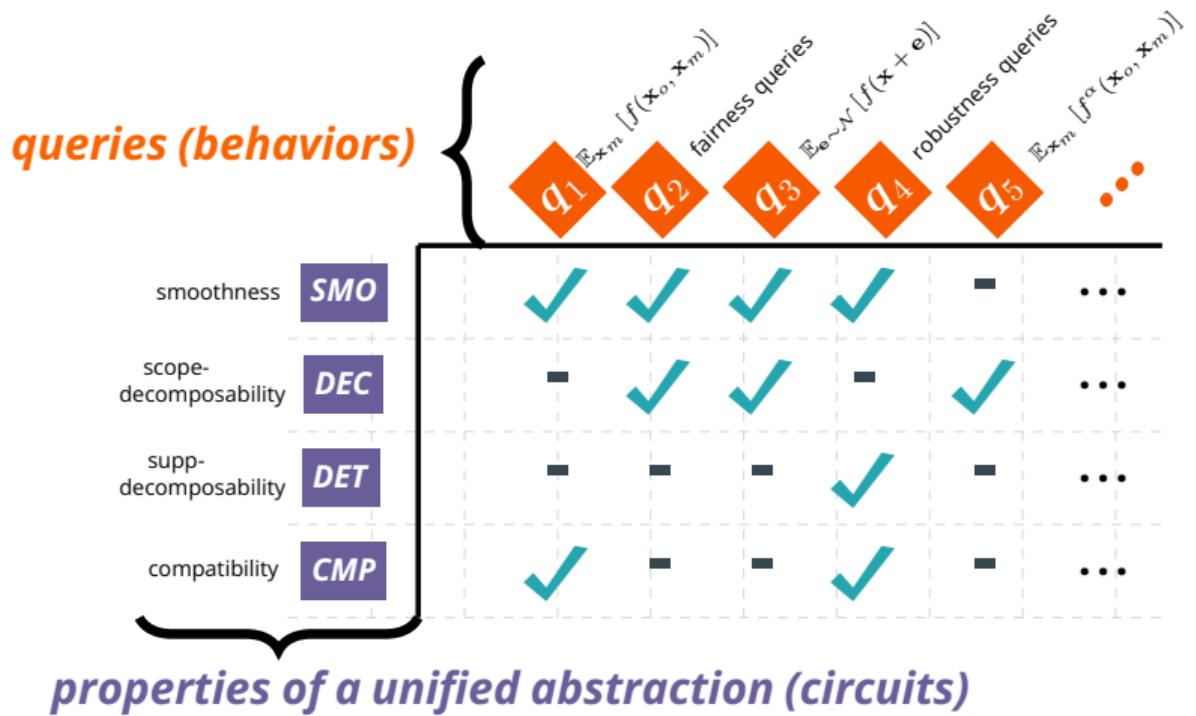
(fast routines to inspect and guarantee a ML system's behavior)

How can we solve this *engineering bottleneck*?

		q_1 $\mathbb{E}_{\mathbf{x}_m} [f(\mathbf{x}_o, \mathbf{x}_m)]$	q_2 fairness queries	q_3 $\mathbb{E}_{\mathbf{e} \sim \mathcal{N}} [f(\mathbf{x} + \mathbf{e})]$	q_4 robustness queries	q_5 $\mathbb{E}_{\mathbf{x}_m} [f^\alpha(\mathbf{x}_o, \mathbf{x}_m)]$			
		queries (behaviors)							
		M_1	M_2	M_3	M_4	M_5	\dots		
HMMs	M_1	✓	?	✓	✓	✗	...		
decision trees	M_2	✓	✓	✓	✓	✓	...		
neural nets w/ ReLUs	M_3	✗	?	?	✓	✗	...		
tensor factorizations	M_4	✓	?	✓	✓	?	...		
transformers	M_5	?	?	?	?	?	...		
		

✓ reliable & efficient algo
✗ intractable
? do not know

model classes



<i>atomic queries</i>		\times	$p \times q$	$p + q$	pow_N	pow_R	p^n	p^q	$/$	p/q	\log	$\log p$	\exp	$\exp p$
smoothness	SMO	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓
scope-decomposability	DEC	-	-	-	-	✓	-	-	-	-	-	-	-	-
supp-decomposability	DET	-	-	-	-	✓	-	✓	✓	✓	✓	-	-	-
compatibility	CMP	-	✓	✓	-	-	-	✓	-	-	-	-	-	-

✓ necess. conditions for reliability and efficiency

properties of a unified abstraction (circuits)

(De)composing queries

*"What is the **expected prediction** for a patient with unavailable records?"*

$$\int [p(\mathbf{x}_m \mid \mathbf{x}_o) \times f(\mathbf{x}_o, \mathbf{x}_m)]$$

(De)composing queries

"What is the **expected prediction** for a patient with unavailable records?"

$$\int [p(\mathbf{x}_m \mid \mathbf{x}_o) \times f(\mathbf{x}_o, \mathbf{x}_m)]$$

"What's the expected **variance**?"

$$\int [p(\mathbf{x}_m \mid \mathbf{x}_o) \times \text{pow}(f(\mathbf{x}_o, \mathbf{x}_m), 2)] - \text{pow} \left(\int [p(\mathbf{x}_m \mid \mathbf{x}_o) \times f(\mathbf{x}_o, \mathbf{x}_m)], 2 \right)$$

(De)composing queries

"What is the **expected prediction** for a patient with unavailable records?"

$$\int [p(\mathbf{x}_m \mid \mathbf{x}_o) \times f(\mathbf{x}_o, \mathbf{x}_m)]$$

"What's the expected **variance**?"

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"How **fair** is the prediction is a certain protected attribute changes?"

$$\int [p(\mathbf{x}_c \mid X_s = 0) \times f(\mathbf{x}_c, 0)] - \int [p(\mathbf{x}_c \mid X_s = 1) \times f(\mathbf{x}_c, 1)]$$

A language for queries

Integral expressions that can be formed by composing these operators

`+ , × , pow , log , exp` and `/`

⇒ *many divergences and information-theoretic queries*

A language for queries

Integral expressions that can be formed by composing these operators

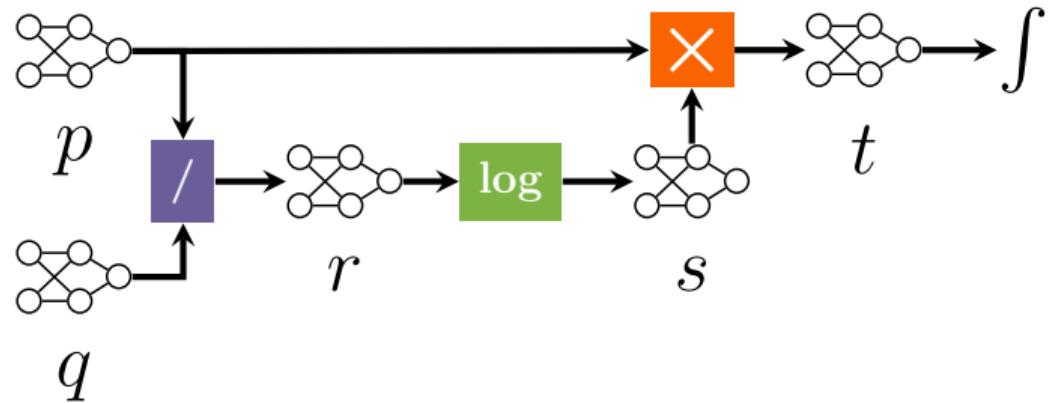
`+ , × , pow , log , exp` and `/`

⇒ *many divergences and information-theoretic queries*

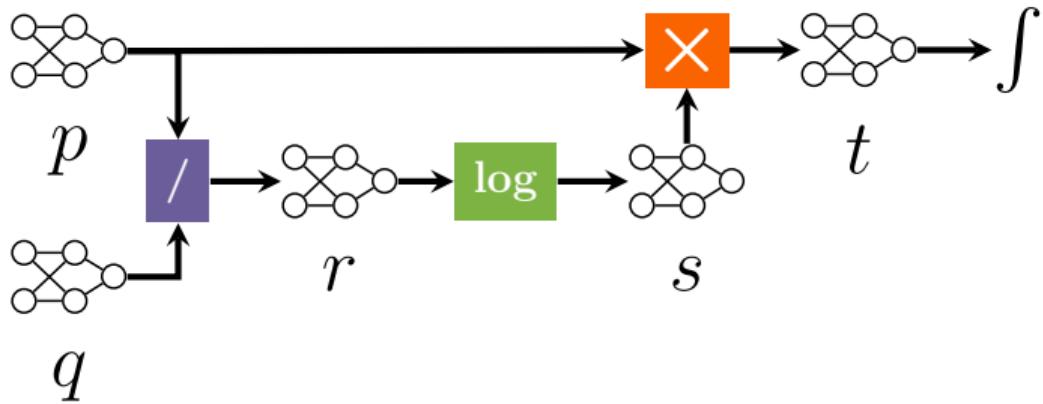
Represented as ***higher-order computational graphs***—pipelines—operating over circuits!

⇒ *re-using intermediate transformations across queries*

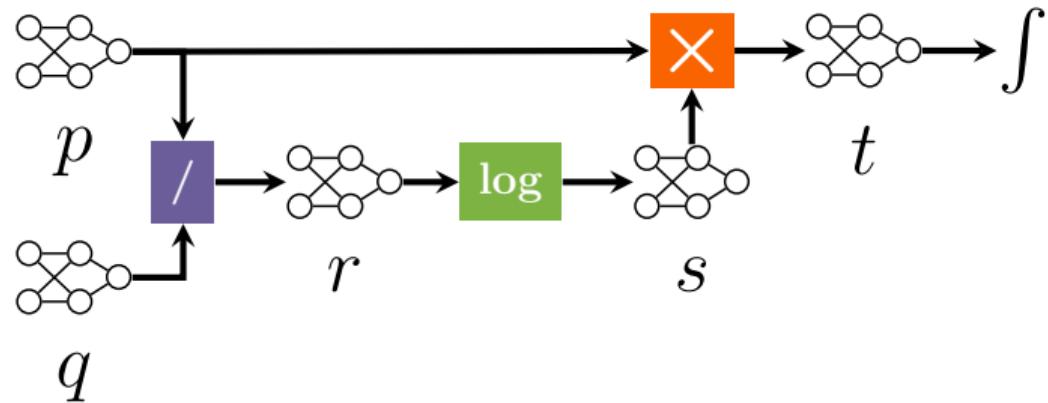
$$\mathbb{KLD}(p \parallel q) = \int_{\text{val}(\mathbf{X})} p(\mathbf{x}) \times \log(p(\mathbf{x})/q(\mathbf{x})) \, d\mathbf{X}$$



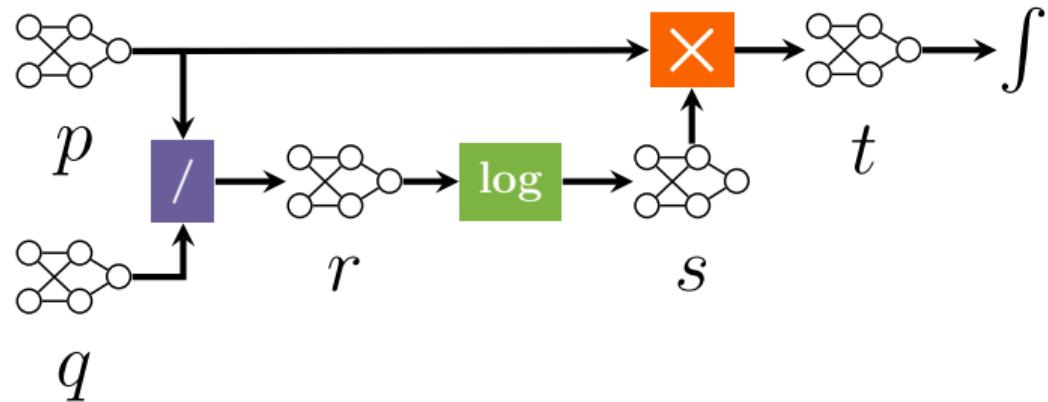
$$\text{KLD}(p \parallel q) = \int_{\text{val}(\mathbf{X})} p(\mathbf{x}) \times \log \left(p(\mathbf{x}) / q(\mathbf{x}) \right) d\mathbf{X}$$



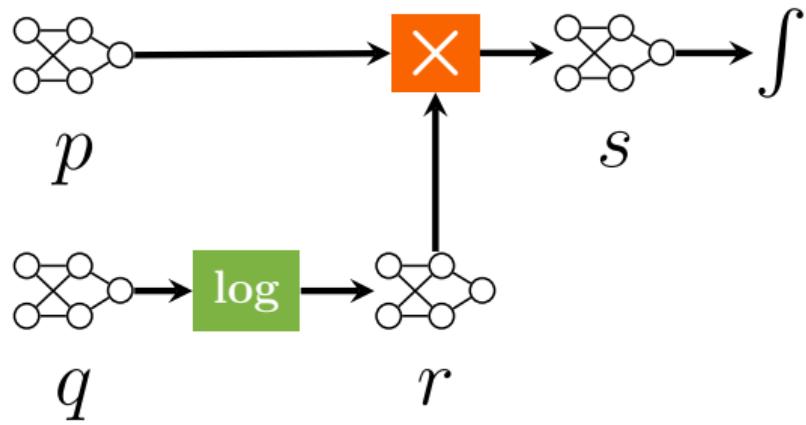
$$\mathbb{KLD}(p \parallel q) = \int_{\text{val}(\mathbf{X})} p(\mathbf{x}) \times \log(p(\mathbf{x})/q(\mathbf{x})) d\mathbf{X}$$



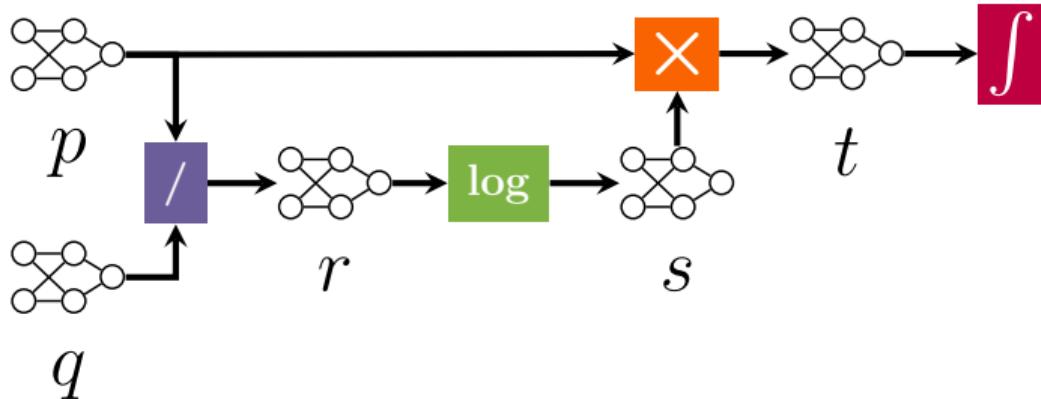
$$\mathbb{KLD}(p \parallel q) = \int_{\text{val}(\mathbf{X})} p(\mathbf{x}) \times \log(p(\mathbf{x})/q(\mathbf{x})) d\mathbf{X}$$



$$\text{XENT}(p \parallel q) = \int p(\mathbf{x}) \times \log q(\mathbf{x}) d\mathbf{X}$$

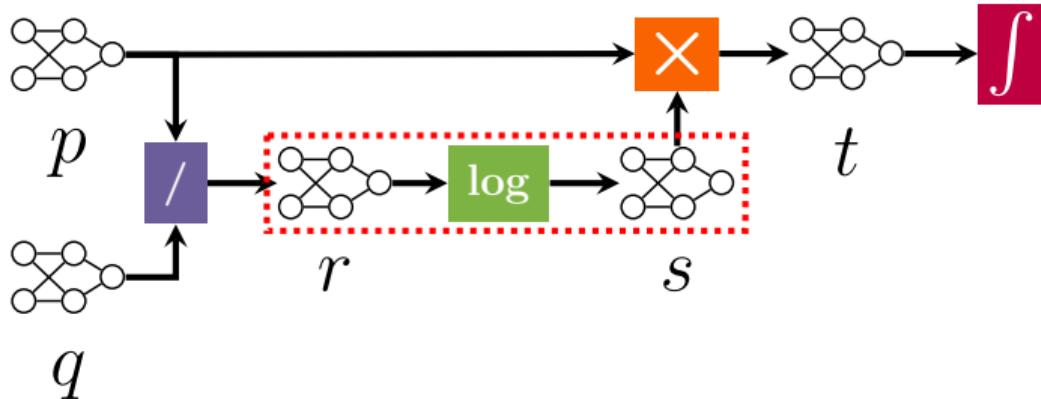


$$\int p(\mathbf{x}) \times \log \left(p(\mathbf{x}) / q(\mathbf{x}) \right) d\mathbf{X}$$



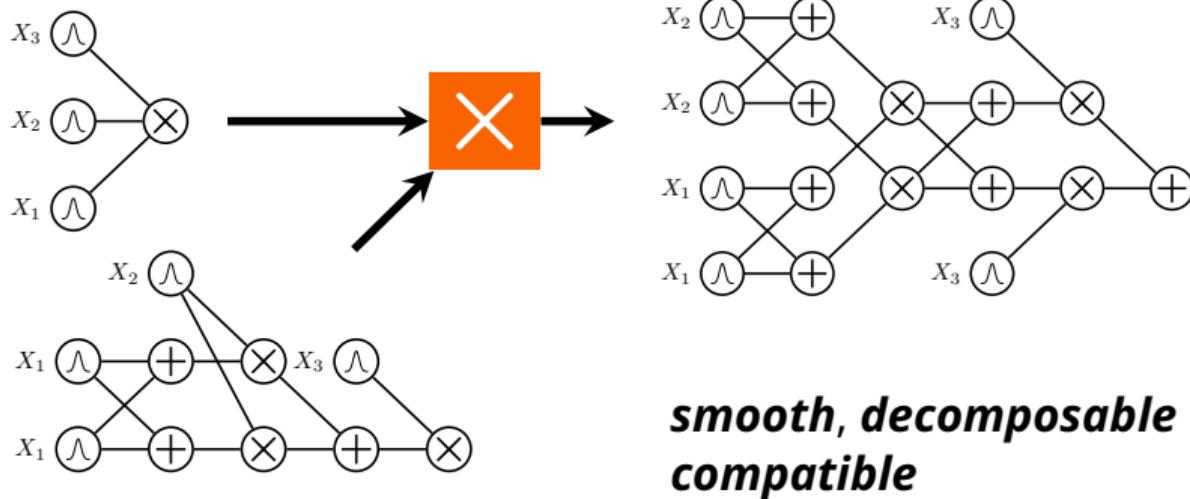
build a LEGO-like query calculus...

$$\int p(\mathbf{x}) \times \log \left(p(\mathbf{x}) / q(\mathbf{x}) \right) d\mathbf{X}$$

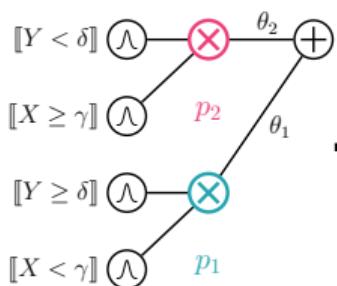


build a LEGO-like query calculus...

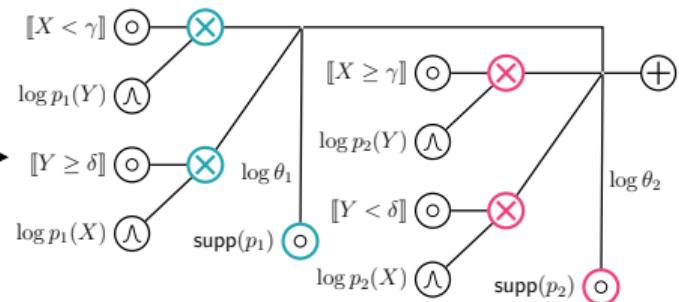
Tractable operators



Tractable operators

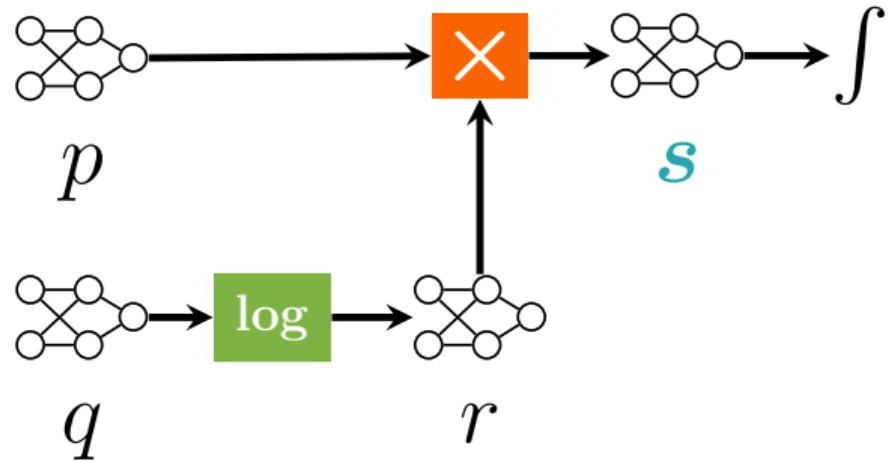


log

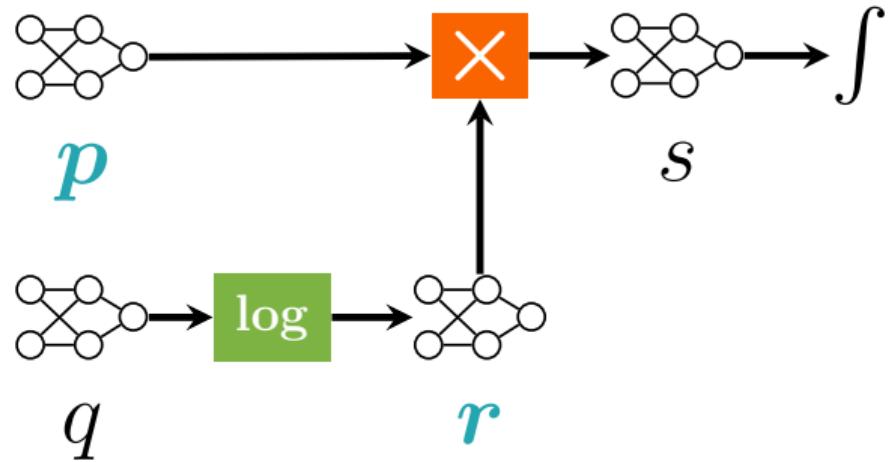


*smooth, decomposable
deterministic*

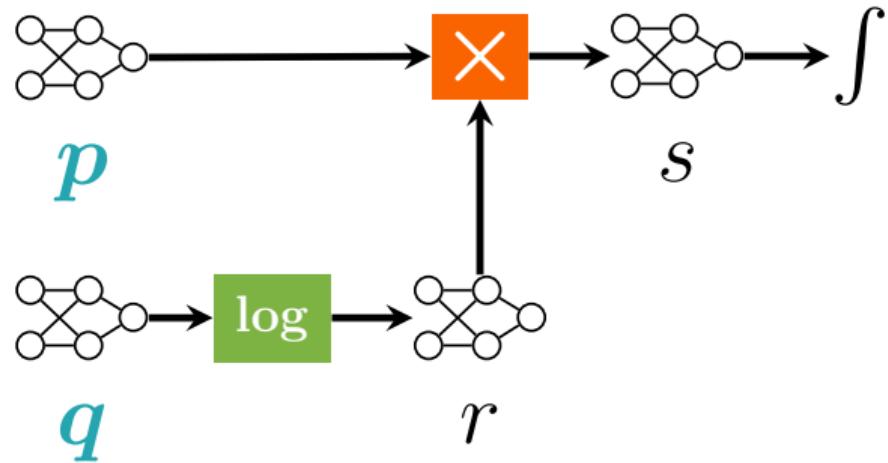
smooth, decomposable



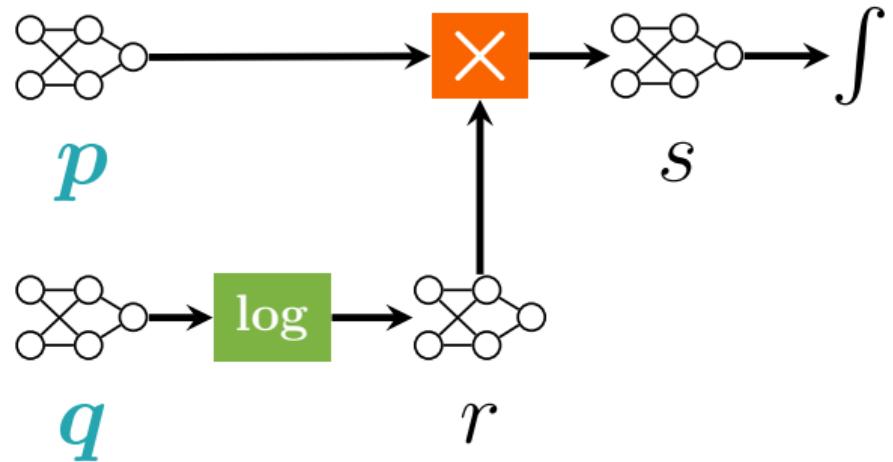
To perform tractable integration we need s to be *smooth and decomposable*...



hence we need p and r to be smooth, decomposable and **compatible**...



therefore q must be smooth, decomposable and **deterministic**...



we can compute XENT tractably if p and q are smooth, decomposable, compatible and q is deterministic...

Query	Tract. Conditions	Hardness
CROSS ENTROPY	$-\int p(\mathbf{x}) \log q(\mathbf{x}) d\mathbf{X}$	Cmp, q Det #P-hard w/o Det
SHANNON ENTROPY	$-\sum p(\mathbf{x}) \log p(\mathbf{x})$	Sm, Dec, Det coNP-hard w/o Det
RÉNYI ENTROPY	$(1 - \alpha)^{-1} \log \int p^\alpha(\mathbf{x}) d\mathbf{X}, \alpha \in \mathbb{N}$ $(1 - \alpha)^{-1} \log \int p^\alpha(\mathbf{x}) d\mathbf{X}, \alpha \in \mathbb{R}_+$	SD #P-hard w/o SD
MUTUAL INFORMATION	$\int p(\mathbf{x}, \mathbf{y}) \log(p(\mathbf{x}, \mathbf{y})/(p(\mathbf{x})p(\mathbf{y})))$	Sm, Dec, Det* #P-hard w/o Det
KULLBACK-LEIBLER DIV.	$\int p(\mathbf{x}) \log(p(\mathbf{x})/q(\mathbf{x})) d\mathbf{X}$	Sm, SD, Det* coNP-hard w/o SD
RÉNYI'S ALPHA DIV.	$(1 - \alpha)^{-1} \log \int p^\alpha(\mathbf{x}) q^{1-\alpha}(\mathbf{x}) d\mathbf{X}, \alpha \in \mathbb{N}$ $(1 - \alpha)^{-1} \log \int p^\alpha(\mathbf{x}) q^{1-\alpha}(\mathbf{x}) d\mathbf{X}, \alpha \in \mathbb{R}$	Cmp, Det #P-hard w/o Det
ITAKURA-SAITO DIV.	$\int [p(\mathbf{x})/q(\mathbf{x}) - \log(p(\mathbf{x})/q(\mathbf{x})) - 1] d\mathbf{X}$	Cmp, q Det #P-hard w/o Det
CAUCHY-SCHWARZ DIV.	$-\log \frac{\int p(\mathbf{x}) q(\mathbf{x}) d\mathbf{X}}{\sqrt{\int p^2(\mathbf{x}) d\mathbf{X} \int q^2(\mathbf{x}) d\mathbf{X}}}$	Cmp, Det #P-hard w/o Det
SQUARED LOSS	$\int (p(\mathbf{x}) - q(\mathbf{x}))^2 d\mathbf{X}$	Cmp #P-hard w/o Cmp

compositionally derive the tractability of many more queries

	Query	Tract. Conditions	Hardness
CROSS ENTROPY	$-\int p(\mathbf{x}) \log q(\mathbf{x}) d\mathbf{X}$	Cmp, q Det	#P-hard w/o Det
SHANNON ENTROPY	$-\sum p(\mathbf{x}) \log p(\mathbf{x})$	Sm, Dec, Det	coNP-hard w/o Det
RÉNYI ENTROPY	$(1 - \alpha)^{-1} \log \int p^\alpha(\mathbf{x}) d\mathbf{X}, \alpha \in \mathbb{N}$ $(1 - \alpha)^{-1} \log \int p^\alpha(\mathbf{x}) d\mathbf{X}, \alpha \in \mathbb{R}_+$	SD Sm, Dec, Det	#P-hard w/o SD #P-hard w/o Det
MUTUAL INFORMATION	$\int p(\mathbf{x}, \mathbf{y}) \log(p(\mathbf{x}, \mathbf{y})/(p(\mathbf{x})p(\mathbf{y})))$	Sm, SD, Det*	coNP-hard w/o SD
KULLBACK-LEIBLER DIV.	$\int p(\mathbf{x}) \log(p(\mathbf{x})/q(\mathbf{x})) d\mathbf{X}$	Cmp, Det	#P-hard w/o Det
RÉNYI'S ALPHA DIV.	$(1 - \alpha)^{-1} \log \int p^\alpha(\mathbf{x}) q^{1-\alpha}(\mathbf{x}) d\mathbf{X}, \alpha \in \mathbb{N}$ $(1 - \alpha)^{-1} \log \int p^\alpha(\mathbf{x}) q^{1-\alpha}(\mathbf{x}) d\mathbf{X}, \alpha \in \mathbb{R}$	Cmp, q Det Cmp, Det	#P-hard w/o Det #P-hard w/o Det
ITAKURA-SAITO DIV.	$\int [p(\mathbf{x})/q(\mathbf{x}) - \log(p(\mathbf{x})/q(\mathbf{x})) - 1] d\mathbf{X}$	Cmp, Det	#P-hard w/o Det
CAUCHY-SCHWARZ DIV.	$-\log \frac{\int p(\mathbf{x}) q(\mathbf{x}) d\mathbf{X}}{\sqrt{\int p^2(\mathbf{x}) d\mathbf{X} \int q^2(\mathbf{x}) d\mathbf{X}}}$	Cmp	#P-hard w/o Cmp
SQUARED LOSS	$\int (p(\mathbf{x}) - q(\mathbf{x}))^2 d\mathbf{X}$	Cmp	#P-hard w/o Cmp

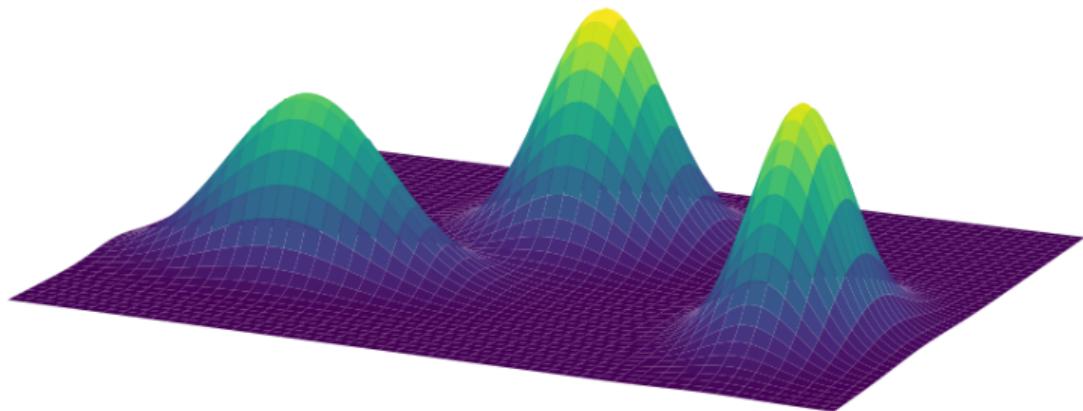
and prove hardness when some input properties are not satisfied

compositional inference I

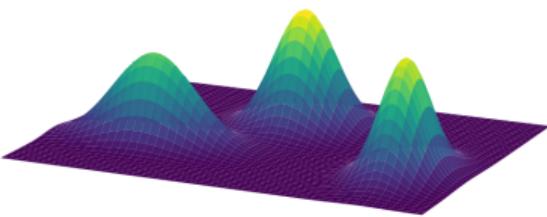
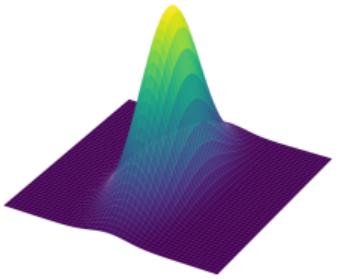


```
1 from cirkit.symbolic.functional import integrate, multiply  
2  
3 def expectation(p, f):  
4     i = multiply(p, f)  
5     return integrate(i)  
6  
7 def squared_loss(p, q): # \int (p - q)^2  
8     p2 = multiply(p, p)  
9     q2 = multiply(q, q)  
10    pq = multiply(p, q)  
11    return integrate(p2) + integrate(q2) - 2*integrate(pq)
```

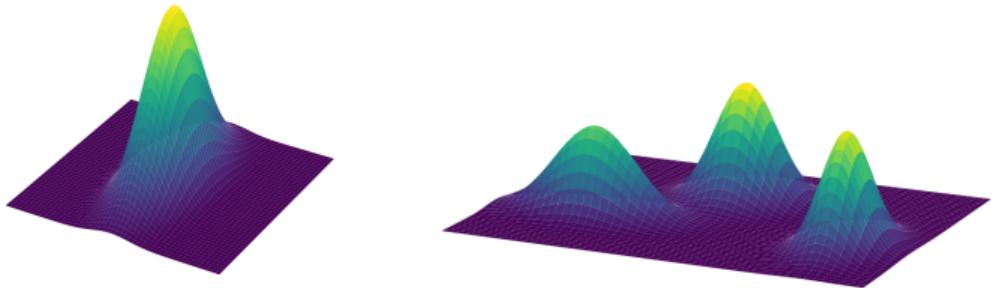
so far...



oh mixtures, you're so fine you blow my mind!



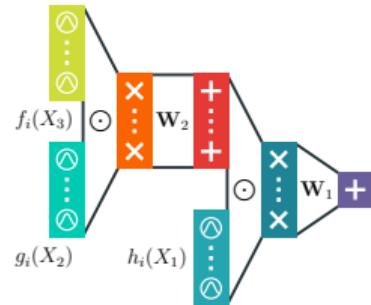
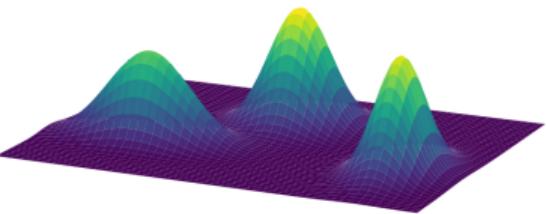
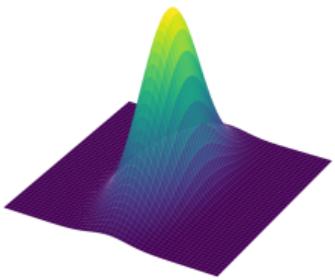
$$p(\mathbf{X}) \quad \xrightarrow{\hspace{1cm}} \quad \sum_{i=1}^K w_i p_i(\mathbf{X})$$



$$p(\mathbf{X}) \longrightarrow \sum_{i=1}^K w_i p_i(\mathbf{X})$$

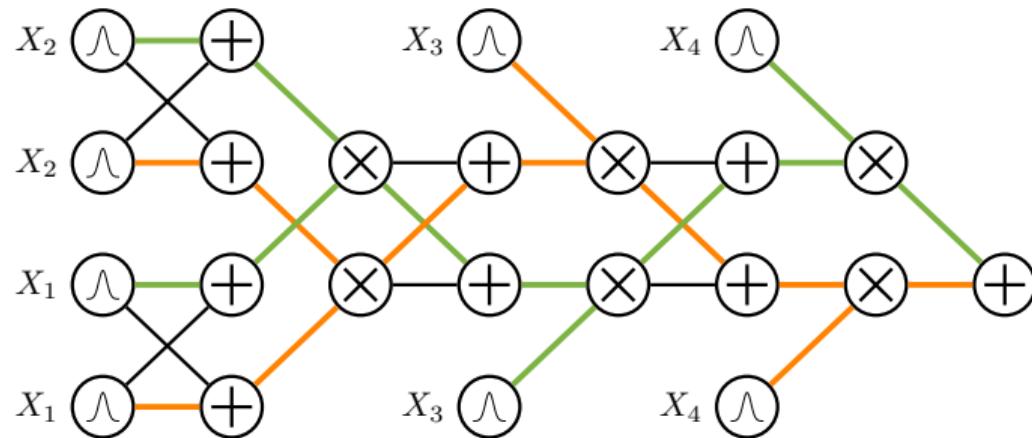
"if someone publishes a paper on model A, there will be a paper about mixtures of A soon with high probability"

A. Vergari



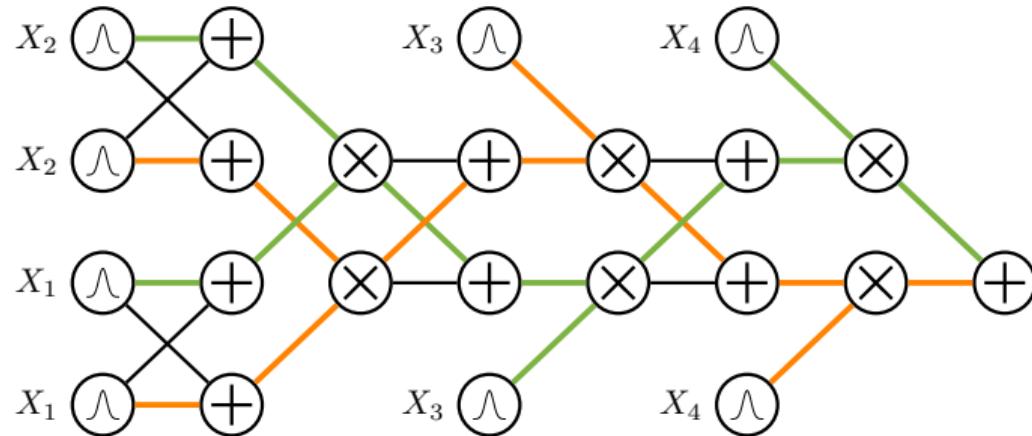
$$p(\mathbf{X}) \longrightarrow \sum_{i=1}^K w_i p_i(\mathbf{X}) \longrightarrow \sum_{i=1}^{2^D} w_i p_i(\mathbf{X}) = \text{PC}(\mathbf{X})$$

expressive efficiency

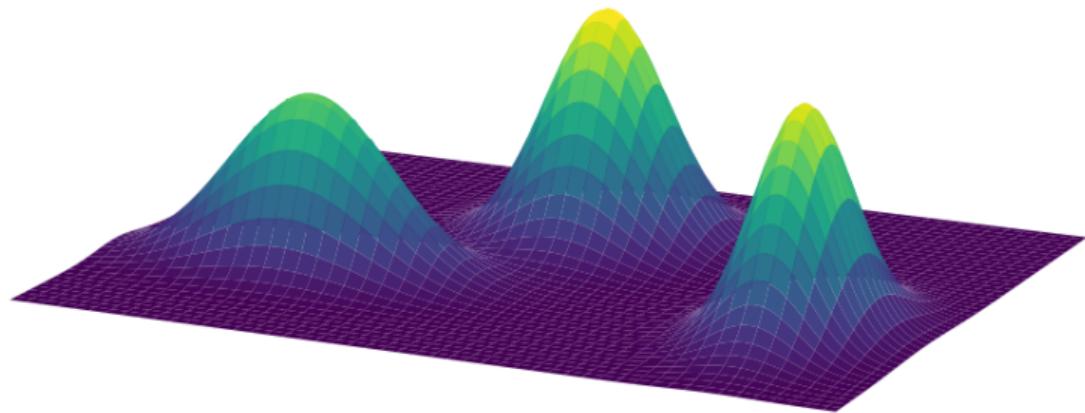


$$p(\mathbf{x}) = \sum_{\mathcal{T}} \left(\prod_{w_j \in \mathbf{w}_{\mathcal{T}}} w_j \right) \prod_{l \in \text{leaves}(\mathcal{T})} p_l(\mathbf{x})$$

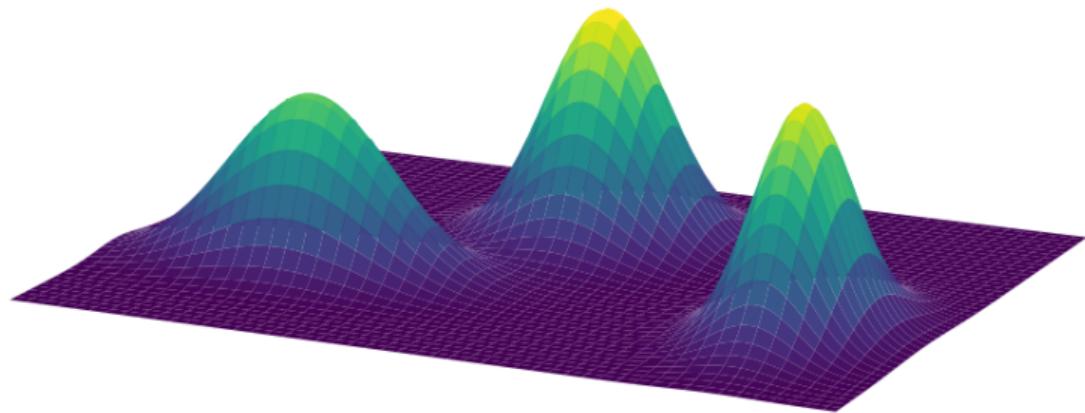
expressive efficiency



an exponential number of mixture components!

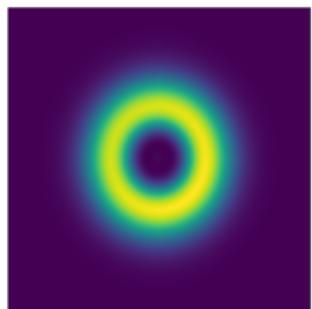


$$c(\mathbf{X}) = \sum_{i=1}^K w_i c_i(\mathbf{X}), \quad \text{with} \quad w_i \geq 0, \quad \sum_{i=1}^K w_i = 1$$

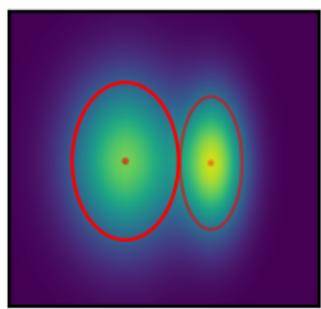
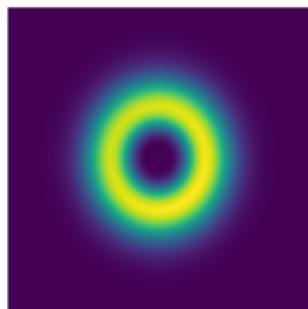


$$c(\mathbf{X}) = \sum_{i=1}^K w_i c_i(\mathbf{X}), \quad \text{with } w_i \geq 0, \quad \sum_{i=1}^K w_i = 1$$

however...

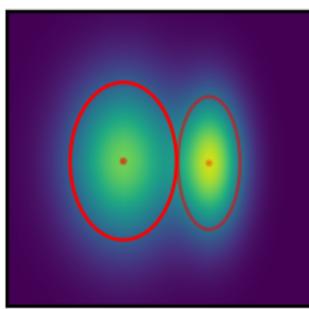
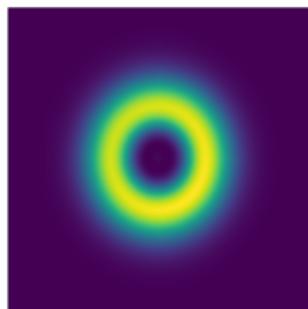


however...

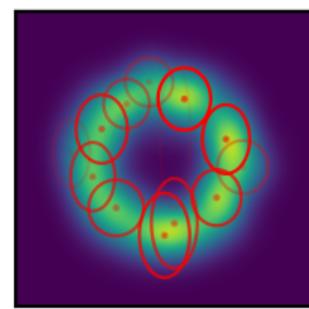


GMM ($K = 2$)

however...

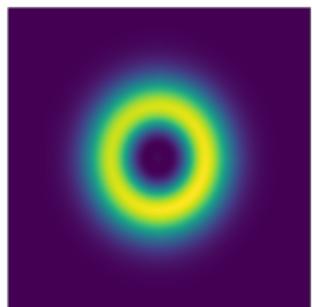


GMM ($K = 2$)

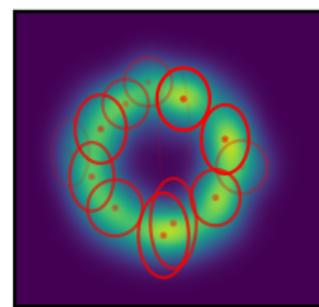


GMM ($K = 16$)

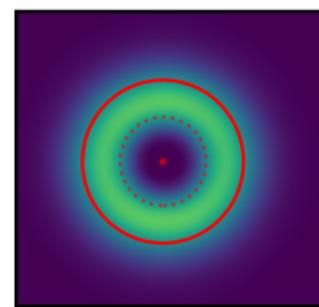
however...



GMM ($K = 2$)



GMM ($K = 16$)

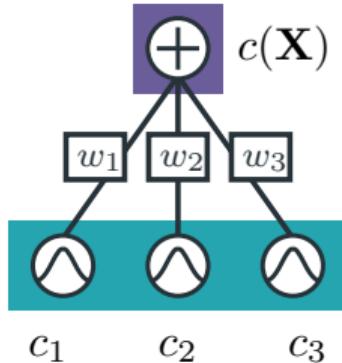


nGMM² ($K = 2$)

theorem

**shallow mixtures
with negative parameters
can be *exponentially more compact* than
deep ones with positive ones.**

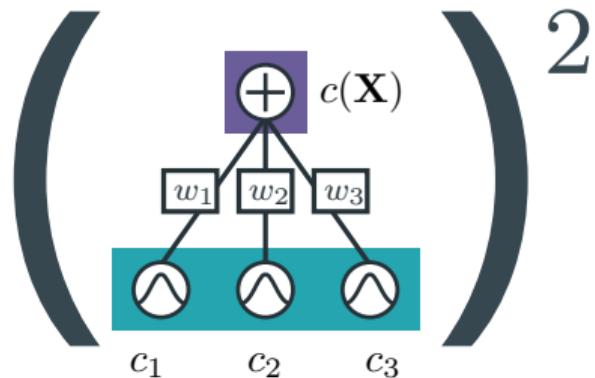
subtractive MMs as circuits



a **non-monotonic** smooth and (structured) decomposable circuit
⇒ possibly with negative outputs

$$c(\mathbf{X}) = \sum_{i=1}^K w_i c_i(\mathbf{X}), \quad w_i \in \mathbb{R},$$

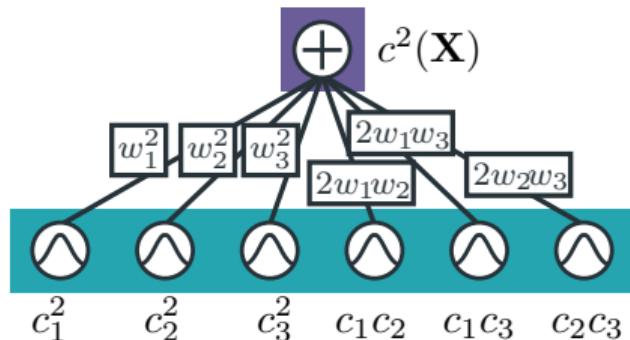
squaring shallow MMs



$$c^2(\mathbf{X}) = \left(\sum_{i=1}^K w_i c_i(\mathbf{X}) \right)^2$$

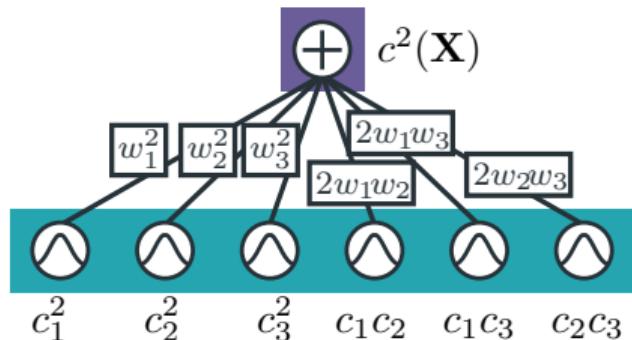
⇒ ensure non-negative output

squaring shallow MMs



$$\begin{aligned} c^2(\mathbf{X}) &= \left(\sum_{i=1}^K w_i c_i(\mathbf{X}) \right)^2 \\ &= \sum_{i=1}^K \sum_{j=1}^K w_i w_j c_i(\mathbf{X}) c_j(\mathbf{X}) \end{aligned}$$

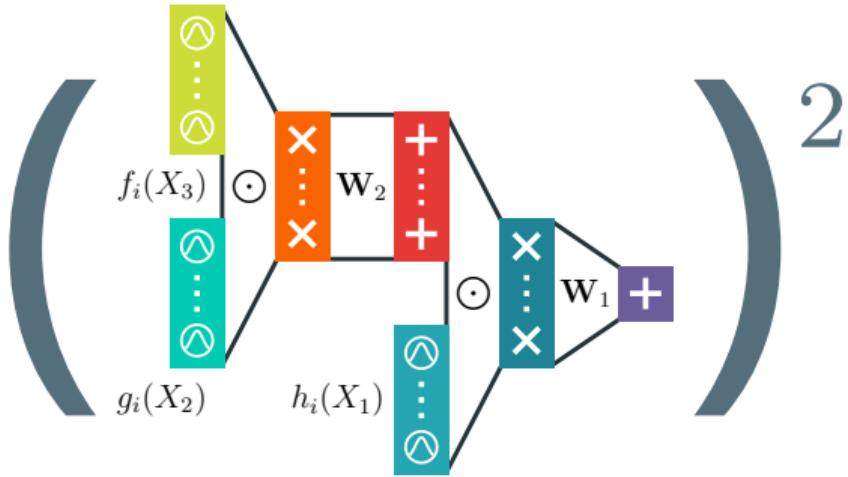
squaring shallow MMs



$$\begin{aligned} c^2(\mathbf{X}) &= \left(\sum_{i=1}^K w_i c_i(\mathbf{X}) \right)^2 \\ &= \sum_{i=1}^K \sum_{j=1}^K w_i w_j c_i(\mathbf{X}) c_j(\mathbf{X}) \end{aligned}$$

still a smooth and (str) decomposable PC with $\mathcal{O}(K^2)$ components!

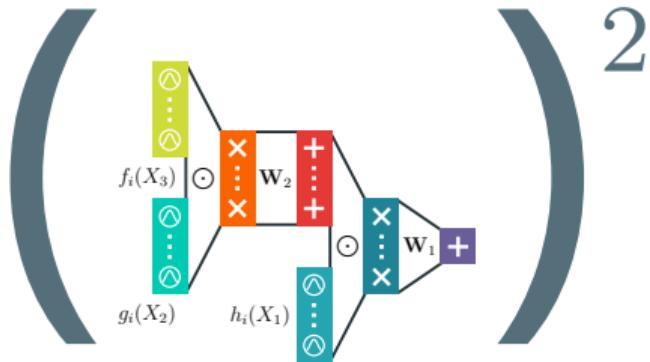
\Rightarrow but still $\mathcal{O}(K)$ parameters



how to efficiently square (and renormalize) a deep PC?

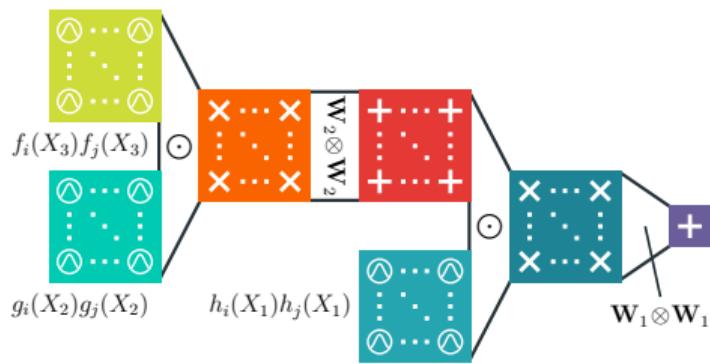
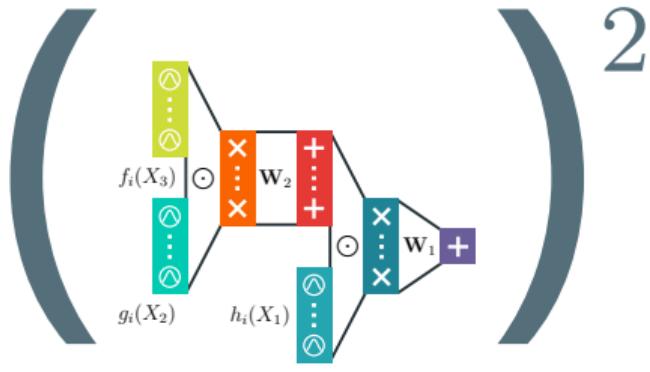
squaring deep PCs

the tensorized way



squaring deep PCs

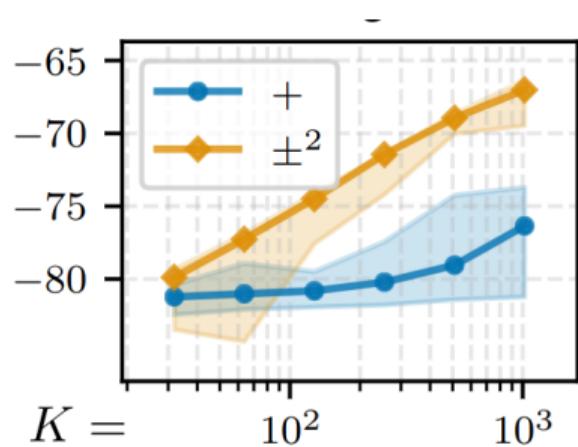
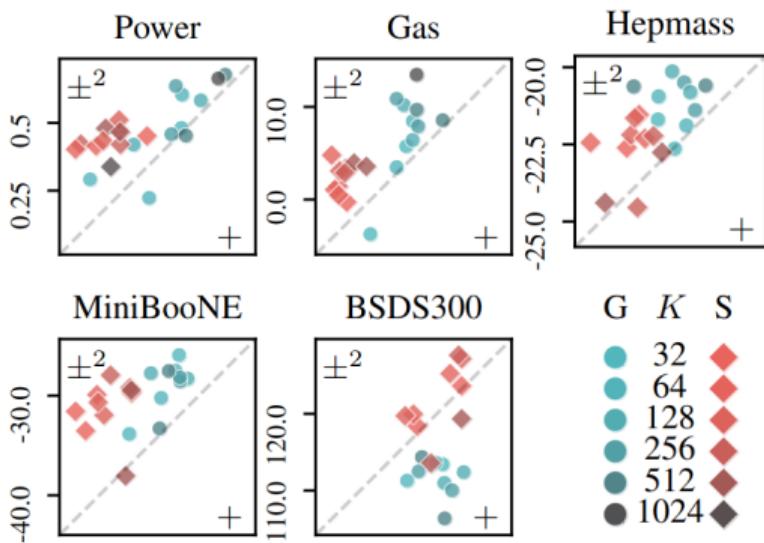
the tensorized way

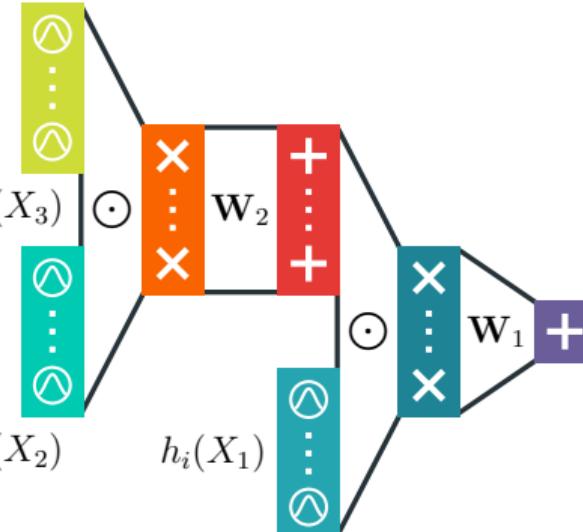


squaring a circuit = to squaring layers

how more expressive?

for the ML crowd



$$\sum \left(f_i(X_3) \odot \begin{array}{c} \times \\ \vdots \\ \times \end{array} \right. \text{W}_2 \left. \begin{array}{c} + \\ \vdots \\ + \end{array} \right) \odot \left(g_i(X_2) \odot \begin{array}{c} \times \\ \vdots \\ \times \end{array} \right. \text{W}_1 \left. \begin{array}{c} + \\ \vdots \\ + \end{array} \right) \Big)^2$$


more than a single square?

$$\pm_{\text{sd}} = \Delta \Sigma_{\text{cmp}}^2$$

(Theorem 5)

$$\Sigma_{\text{cmp}}^2 = \text{psd}$$

(Proposition 2)

$+$ _{sd}

•
Open Question 1

•
Open Question 2

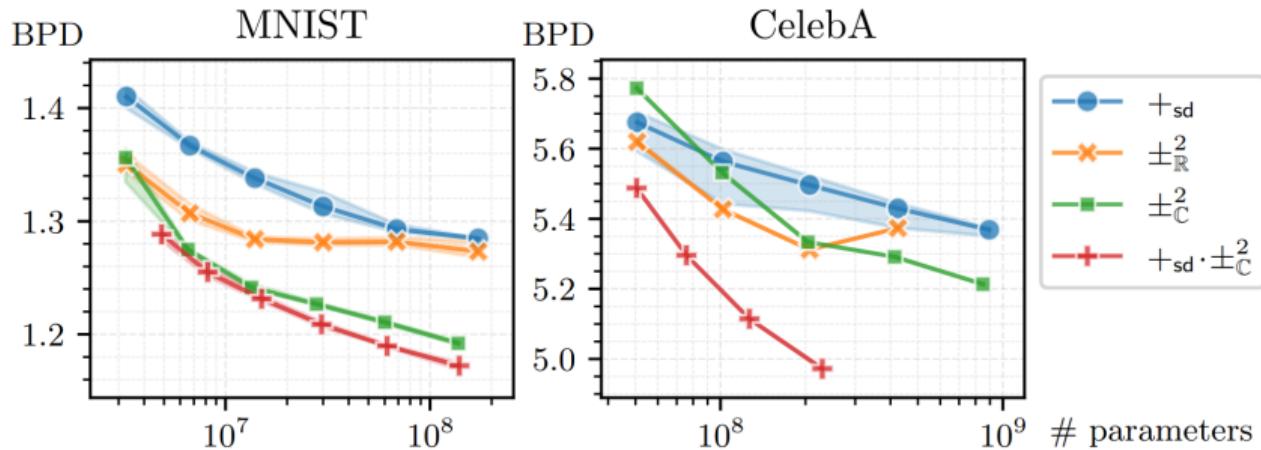
$$\pm_{\mathbb{R}}^2$$

• UDISJ
(Theorem 0)

• UPS
(Theorem 2)

• UTQ
(Theorem B.3)

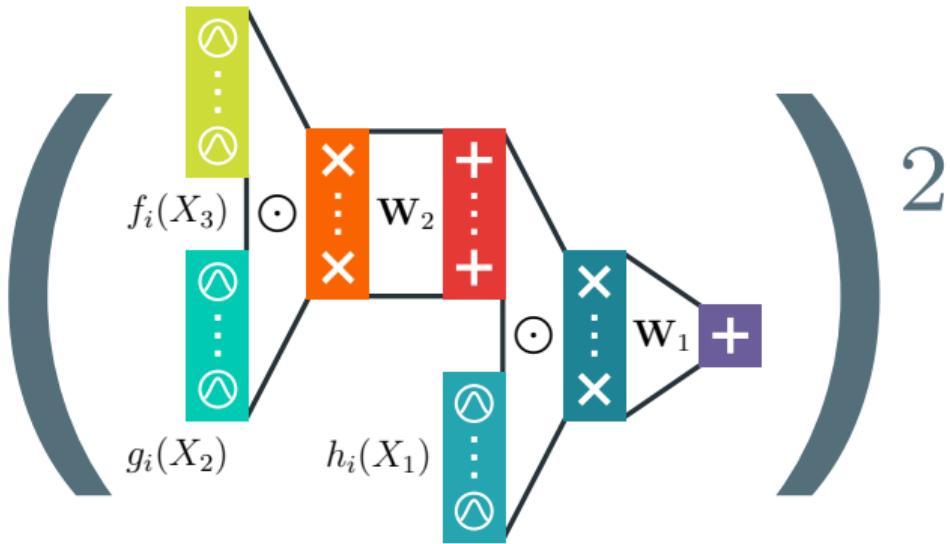
SOS circuits are more expressive



complex circuits are SOS (and scale better!)



learning & reasoning with circuits in pytorch



questions?