

# Arrays

- 2.1** Let A be a two-dimensional array declared as follows:

A : array [1 ... 10] [1 ... 15] of integer;

Assuming that each integer takes one memory location. The array is stored in row-major order and the first element of the array is stored at location 100, what is the address of the element A[i][j]?

- (A) 15i + j + 84      (B) 15j + i + 84  
 (C) 10i + j + 89      (D) 10j + i + 89  
**[1998 : 2 M]**

```
i = 0; j = 1;
while (j < n) {
    if (E) j++;
    else if (a[j] - a[i] == S) break;
    else i++;
}
```

```
if (j < n) printf("yes");
else printf ("no");
```

Choose the correct expression for E.

- (A) a[j] - a[i] > S      (B) a[j] - a[i] < S  
 (C) a[i] - a[j] < S      (D) a[i] - a[j] > S  
**[2005 : 2 M]**

- 2.6** Let a and b be two sorted arrays containing n integers each, in non-decreasing order. Let c be a sorted array containing 2n integers obtained by merging the two arrays a and b. Assuming the arrays are indexed starting from 0, consider the following four statements:

- I. a[i] ≥ b[i] => c[2i] ≥ a[i]  
 II. a[i] ≥ b[i] => c[2i] ≤ b[i]  
 III. a[i] ≥ b[i] => c[2i] ≤ a[i]  
 IV. a[i] ≥ b[i] => c[2i] ≤ b[i]

Which of the following is TRUE

- (A) Only I and II      (B) Only I and IV  
 (C) Only II and III      (D) Only III and IV  
**[2005 : 2 M]**

- 2.7** Consider the following C function in which size is the number of elements in the array E. int MyX(int \*E, unsigned int size) {

```
int Y = 0;
int Z;
int i, j, k;
for(i = 0; i < size; i++)
    Y = Y + E[i];
for(j = 1; j < size; j++)
{
    Z = 0;
    for(k = i; k <= j; k++)
        Z = Z + E[k];
    if (Z > Y)
        Y = Z;
}
return Y;
```

}

```
Z = 0;
for(k = i; k <= j; k++)
    Z = Z + E[k];
if (Z > Y)
    Y = Z;
```

}

return Y;

}

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}

**2.13** Consider the following C program:

```
#include <stdio.h>
void strcpy(char *, char *);
int main() {
    char a[30] = "@#Hello World!";
    strcpy(a, a + 2);
    printf("%s\n", a);
    return 0;
}
void strcpy(char *s, char *t) {
    while(*t)
        *s++ = *t++;
}
```

Which ONE of the following will be the output of the program?

- (a) @#Hello World!    (b) Hello World!  
 (c) ello World!    (d)>Hello World!

[2025 (Set-2) : 1 M]

**2.14** An array  $A$  of length  $n$  with distinct elements is said to be bitonic if there is an index  $1 \leq i \leq n$  such that  $A[1..i]$  is sorted in the non-decreasing order and  $A[i+1..n]$  is sorted in the non-increasing order.

Which ONE of the following represents the best possible asymptotic bound for the worst-case number of comparisons by an algorithm that searches for an element in a bitonic array  $A$ ?

- (a)  $\Theta(n)$     (b)  $\Theta(1)$   
 (c)  $\Theta(\log^2 n)$     (d)  $\Theta(\log n)$

[2025 (Set-2) : 2 M]



### Answers Arrays

- 2.1 (a) 2.2 (a) 2.3 (a) 2.4 (a) 2.5 (b) 2.6 (c) 2.7 (a) 2.8 (a) 2.9 (a)

- 2.10 (c) 2.11 (0) 2.12 (19) 2.13 (d) 2.14 (d)

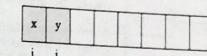
### Explanations Arrays

**2.1** (a)

Let  $r$  be number of elements in a row.

$$\begin{aligned} \text{Address of the element } A[i][j] \\ &= \text{Base address} + (i-1)*r + (j-1) \\ &= 100 + (i-1) \times 15 + (j-1) \\ &= 100 + 15i - 15 + j - 1 \\ &= 100 + 15i + j - 16 \\ &= 84 + 5i + j \end{aligned}$$

**2.5** (b)



So, if condition will check  $a[j] - a[i] < S$  then increment  $j$  so that difference between two element increase and when  $a[j] - a[i] > S$  then increment  $i$  for making difference between two element indexed at  $i$  and  $j$  respectively is decrease. So,  $a[j] - a[i] < S$  is correct condition.

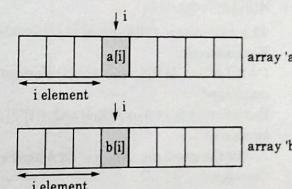
**2.2** (a)

Array is

0	-1	-2	...	1-n
1	0	-1	...	
2	1	0	...	
3	2	1	...	
:	:	:	...	
-(1-n)	-(2-n)			0

$$\text{Sum} = 0 + 1 + 2 + \dots + (-1) + (-2) + \dots = 0.$$

**2.6** (c)



**2.3** (a)

Effect of the given 3 reversals for any  $k$  is equivalent to left rotation of the array of size  $n$  by  $k$ .

Let,  $S[1 \dots 7]$

1	2	3	4	5	6	7
---	---	---	---	---	---	---

$\therefore n = 7, k = 2$

reverse ( $S, 1, 2$ ) we get  $[2, 1, 3, 4, 5, 6, 7]$

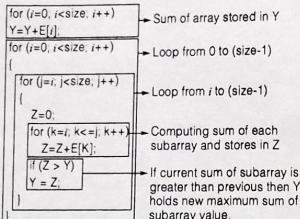
reverse ( $S, 3, 7$ ) we get  $[2, 1, 7, 6, 5, 4, 3]$

reverse ( $S, 1, 7$ ) we get  $[3, 4, 5, 6, 7, 1, 2]$

**2.4** (a)

There are 500 students the score range is 0 to 100. Print the frequency of those student whose score above 50. So frequency range contains score from 50 to 100, so an array of 50 numbers is suitable for representing the frequency.

- When  $a[i] \geq b[i]$  then  $2i$  elements smaller than equal to  $a[i]$ . Since ' $i$ ' elements to array 'b' that are less than equal to  $b[i]$  and ' $i$ ' element of array 'a' which are less than equal to  $a[i]$ . So  $a[i]$  will comes after  $c[2i]$  i.e.  $c[2i] \leq a[i]$  is true.
- When  $a[i] \geq b[i]$  then  $b[i]$  always be less than equal to  $c[2i]$  i.e. when ' $i$ ' elements of array 'a' and ' $i$ ' elements of array 'b' are same then maximum index of  $b[i]$  is  $c[2i]$ . So,  $c[2i] \geq b[i]$  is true.

**2.7 (a)**

∴ The value (Y) returned by function is the maximum possible sum of elements in any subarray of array E.

**2.8 (a)**

$$\text{Temp} = A[i][j] + C;$$

$$A[i][j] = A[j][i];$$

$$A[j][i] = \text{Temp} - C;$$

This code swaps  $A[i][j]$  and  $A[j][i]$  elements

**For example**

$i = 1, j = 2 \Rightarrow A[1][2]$  and  $A[2][1]$  elements are exchanged

But when  $i = 2, j = 1 \Rightarrow A[2][1]$  and  $A[1][2]$  again swapped

∴ For the given code, the matrix A itself will be the output.

**2.9 (a)**

0 <sup>th</sup> row	2000	2004	2008
	1	2	3
1 <sup>st</sup> row	2012	2016	2020
	4	5	6
2 <sup>nd</sup> row	2024	2028	2032
	7	8	9
3 <sup>rd</sup> row	2036	2040	2044
	10	11	12

Arithmetic of two-dimensional array addresses.

'x' is the base address of the 0<sup>th</sup> row.

'x+1' is the base address of the 1<sup>st</sup> row.

'x+i' is the base address of the i<sup>th</sup> row.

Similarly  $(x+i)$  is the i<sup>th</sup> row. So  $(x+i)$  the address of the 0<sup>th</sup> element in the i<sup>th</sup> row (by putting i = 0)

The outputs printed by printf statements are:  
 $x+3 \rightarrow 2036$  (starting address of 3<sup>rd</sup> row)

\* $(x+3) \rightarrow 2036$  (address of 0<sup>th</sup> element in 3<sup>rd</sup> row)  
\* $(x+2)+3 \rightarrow 2036$  (gives address in 2<sup>nd</sup> row but +3 makes the starting address of 2<sup>nd</sup> row (2024) incremented by  $3 \times 4 = 12$ )  
∴ Output is 2036, 2036, 2036.

**2.10 (c)**

Y is one dimensional array  
X is two dimensional array

Y	0	0	0	0	.	.	0
	0	1	n-1				
X	0	0	1	2	3	4	n-1
	1	2	3	4	...	n	
	2	3	4	5	...	n+1	
	3	4	5	6	...	n+2	
	..						
	n-1	n	n+1	n+2	...	n+n-2	

**Code segment 1:**

```

    for (i = 0; i < n; i++)
        Y[i] += X[0][i];
    
```

**Code segment 2:**

```

    for (i = 0; i < n; i++)
        Y[i] += X[i][0];
    
```

S1: Final contents of array Y will be same in both code segments.

Y	0	1	2	3	4	5	6	7	n-1
	0	1	2	3	4	5	6	7	...

is the output of both segments.

S2: X contents are stored row wise in C-program hence code segment row '0' accessed contiguous in memory.

S3: In code segment 2, content of X are not accessed in contiguous.

**2.11 (0)**

$$K = 4$$

0	1	2	3
1	0	1	1

$$a = 2, n = 8, Z = 1,$$

$$i = 0 \quad Z = 1^2 \bmod 8 = 1$$

$$\text{if } (C[0] == 1) \quad Z = (1 \times 2) \bmod 8 = 2$$

$$i = 1 \quad Z = 2^2 \bmod 8 = 4$$

$$i = 2 \quad Z = 4^2 \bmod 8 = 0$$

$$\text{if } (C[2] == 1) \quad Z = (0 \times 2) \bmod 8 = 0$$

$$i = 3 \quad Z = 0^2 \bmod 8 = 0$$

$$\text{if } (C[3] == 1) \quad Z = (0 \times 2) \bmod 8 = 0$$

Returns 0

**2.12 (19)**

'a' is a two dimensional array as shown below:

0, 0	0, 1	0, 2	0, 3	0, 4
1, 0	1, 1	1, 2	1, 3	1, 4
6	7	8	9	10
2, 0	2, 1	2, 2	2, 3	2, 4
11	12	13	14	15
3, 0	3, 1	3, 2	3, 3	3, 4
16	17	18	19	20

a = address of 0<sup>th</sup> index element of 0<sup>th</sup> index 1-D array

\*a =

\*a is value of 0<sup>th</sup> index element of 0<sup>th</sup> index 1-D array

\*a = 1

\*\*a + 2 = 1 + 2 = 3

a + 3 = address of 3<sup>rd</sup> index 1-D array

\*(a + 3) = address of 0<sup>th</sup> index element of 3<sup>rd</sup> index 1-D array

\*a + 3 = address of 3<sup>rd</sup> index element of 3<sup>rd</sup> index 1-D array

\*\*\*(a + 3) = value of 3<sup>rd</sup> index element of 3<sup>rd</sup> index 1-D array

index 1-D array = 19

Hence it will print 19.

**2.13 (d)**

a	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
	@	#	H	e	i	l	o	w	o	r	i	d	!	/	

s	100	t	102
	101		103
	102		104
	103		105
	104		106
	105		107
	106		108
	107		109
	108		110
	109		111
	110		112
	111		113
	112		114
	113		115

• Value at location 102 (H) copy into value at location 100:

• Value at location 103 (e) copy into value at location 101.

⋮

• The same process continues until it reaches the null character.

**2.14 (d)**

1	2	3	4	5	6	7	8	9	10
4	6	8	10	12	15	18	12	8	7

Using Binary Search:

```

low = 1; high = n;
while(low ≤ high)
{
    m = (low + high)/2;
    if (C[0] == 1) Z = (1 × 2) mod 8 = 2
    if (C[2] == 1) Z = (0 × 2) mod 8 = 0
    if (C[3] == 1) Z = (0 × 2) mod 8 = 0
}
    
```

```

if (a[mid - 1] < a[mid] < a[mid + 1])
    low = mid + 1;
else if (a[mid - 1] > a[mid] > a[mid + 1])
    high = mid - 1;
else if (a[mid - 1] > a[mid] > a[mid + 1])
    return(mid);
}
    
```

TC:  $\Theta(\log n)$