## CSE 505 Fall 2015

## Assignment 4: Higher-Order Functions and Advanced Control

Assigned: Mon, Nov 9, 2015 Due: Mon, Nov 23 (11:59 pm)

Note: This assignment may be done by a pair of students.

**Problem 1.** Consider the following function definition in the C programming language for carrying out the summation expressed by the operator  $\Sigma$  discussed in Lecture 16:

```
int sigma(int *k, int low, int high, int expr()) {
   int sum = 0;
   for (*k=low; *k<=high; (*k)++) {
      sum = sum + expr();
   }
   return sum;
}</pre>
```

Write a main program in the C programming language that makes repeated use of sigma in order to compute and print out the value of the following expression:

$$\sum_{i=0}^{4} \ \left( \ i \ * \sum_{j=0}^{4} \ \left( \ (i+j) \ * \sum_{k=0}^{4} (j*k-i) \right) \right)$$

Place the code for sigma and main in a file sigma.c. Use the gcc compiler for testing your code.

**Problem 2.** Lecture 14 showed an ML function flatten: 'a list list  $\rightarrow$  'a list, which takes a two-level list, say II, and returns a single-level list by appending together all the sub-lists in II. For example, flatten([[1],[2,3],[4],[]]) = [1,2,3,4].

a. Why could we **not** write in ML a function that would flatten an input list with an arbitrary number of levels: 2-level list, 3-level list, etc.? That is, in addition to the example shown above, we would also like to have an example such as:

```
flatten([[1],[2]],[[3]],[[4,5],[6]]) = [1,2,3,4,5,6]
```

b. Show how we can write a general flatten function using a **Python generator**. This generator should yield the values in a multi-level list one by one. Then, we can create a single-level list from any multi-level list, I, by executing: [x for x in flatten(1)].

Create a single file, flatten.py, containing the definition of flatten, as well as your answer to part a.

**Problem 3.** Consider an infinite list of strings of the form:

```
"Lf.Lx.(f x)"
"Lf.Lx.(f (f x))"
"Lf.Lx.(f (f (f x)))"
"Lf.Lx.(f (f (f x)))"
```

These strings represent the numbers 1, 2, 3, 4, ... in the pure lambda-calculus. Here, L stands for  $\lambda$ . Each string is called a *Church numeral* – in honor of Alonzo Church who invented the  $\lambda$ -calculus.

Referring to the **infinite list** ML type discussed in Lecture 17:

```
datatype 'a inf list = lcons of 'a * (unit -> 'a inf list)
```

Define a function church: string -> string inf\_list which generates an infinite list of Church numerals starting from 1. Test out church by executing the following "main" program:

```
fun take(0, _) = []
  | take(n, lcons(h, thk)) = h :: take(n-1, thk());
take(5,church("x"))
```

Create a file called **church.sml** with all relevant definitions.

**Problem 4.** Give a formal definition for a function LI which defines the **leftmost-innermost redex** of a lambda-term, along the lines of the substitution operation given in the notes on Lambda Calculus (pages 5-6). If there is no redex in the input term, LI should return  $\bot$ , which stands for "undefined". Examples:

```
\begin{array}{lll} LI & \textbf{x} = \bot & LI & \lambda \textbf{x}. (\textbf{a} & \textbf{x}) = \lambda \textbf{x}. (\textbf{a} & \textbf{x}) \\ LI & \lambda \textbf{x}. \textbf{y} = \bot & LI & (\lambda \textbf{x}. \textbf{y} & \textbf{a}) = (\lambda \textbf{x}. \textbf{y} & \textbf{a}) \\ LI & \lambda \textbf{x}. (\textbf{x} & \textbf{x}) = \bot & LI & (\lambda \textbf{x}. \textbf{y} & \textbf{y}) = (\lambda \textbf{y}. \textbf{y} & \textbf{x}) \\ LI & (\textbf{a} & \textbf{b}) = \bot & LI & ((\textbf{a} & \textbf{b}) & (\textbf{c} & (\lambda \textbf{z}. \textbf{z} & \textbf{b}))) = (\lambda \textbf{z}. \textbf{z} & \textbf{b}) \\ LI & ((\textbf{a} & \textbf{b}) & (\textbf{c} & \textbf{d})) = \bot & LI & ((\lambda \textbf{z}. \textbf{z} & \textbf{a}) & (\lambda \textbf{z}. \textbf{z} & \textbf{b})) = (\lambda \textbf{z}. \textbf{z} & \textbf{a}) \\ \end{array}
```

The definition of *LI* involves about 8 rules. Two of the rules to help you get started are:

```
LI \quad V = \bot

LI \quad \lambda V.T = LI \quad T, \quad \text{if } LI \quad T \neq \bot
```

**Important:** The LHS of each of these rules must be mutually-exclusive of other rules.

You may use the variable *V* to stand for any variable and *T*, *T1*, *T2* for arbitrary lambda-terms. In addition to the standard *occurs\_free\_in* test, you may also use *is\_etaredex(T)* and *is\_betaredex(T)* to test whether *T* is an eta- or a beta-redex, respectively. Place your solution in a file called **lambda.pdf**.

What to Submit: Prepare a top-level directory named A4\_UBITId1\_UBITId2 if the assignment is done by two students; otherwise, name it as A4\_UBITId if the assignment is done solo. (Order the UBITId's in alphabetic order, in the former case.) In this directory, place sigma.c, flatten.py, church.sml, and lambda.pdf. Compress the directory and submit the resulting compressed file using the submit\_cse505 command. For more details regarding online submission, see Resources → Homeworks → Online Submission 2015.pdf.