



### **Department of Information Technology**

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Academic Year: 2025-2026 Semester: V

Class / Branch: TE IT
Subject: Security Lab

Subject Incharge: Prof. Apeksha Mohite

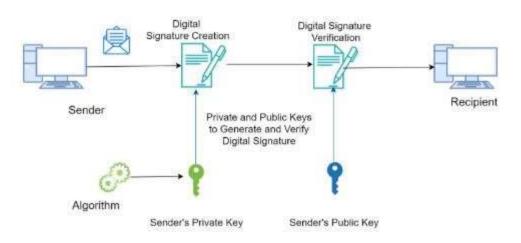
#### **Experiment No. 13(ii)**

1. Aim: To study and analyze RSA cryptosystem and digital signature scheme.

2. Software Required : CrypTool 1.4.41

3. Theory:

#### **Digital Signatures**



Digital signatures can provide the added assurances of evidence of origin, identity and status of an electronic document, transaction or message and can acknowledge informed consent by the signer.

#### How digital signatures work

Digital signatures are based on public key cryptography, also known as asymmetric



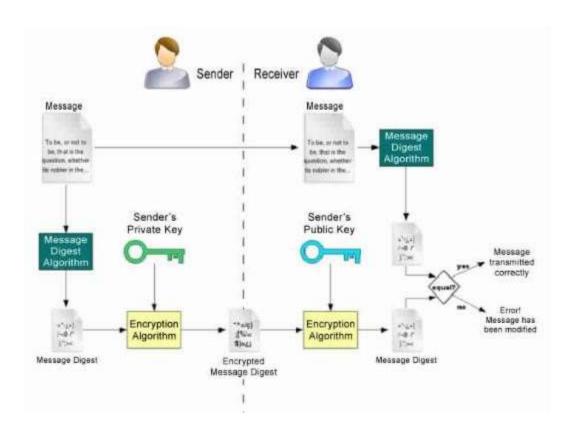


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cryptography. Using a public key algorithm, such as RSA, one can generate two keys that are mathematically linked: one private and one public.

Digital signatures work because public key cryptography depends on two mutually authenticating cryptographic keys. The individual who is creating the digital signature uses their own private key to encrypt signature-related data; the only way to decrypt that data is with the signer's public key. This is how digital signatures are authenticated.



#### How to create a digital signature

To create a digital signature, signing software such as an email program -- creates a one-way hash of the electronic data to be signed. The private key is then used to encrypt the hash. The encrypted hash along with other information, such as the hashing algorithm is the digital signature.





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The reason for encrypting the hash instead of the entire message or document is that a hash function can convert an arbitrary input into a fixed length value, which is usually much shorter. This saves time as hashing is much faster than signing.

The value of a hash is unique to the hashed data. Any change in the data, even a change in a single character, will result in a different value. This attribute enables others to validate the integrity of the data by using the signer's public key to decrypt the hash.

If the decrypted hash matches a second computed hash of the same data, it proves that the data hasn't changed since it was signed. If the two hashes don't match, the data has either been tampered with in some way -- integrity -- or the signature was created with a private key that doesn't correspond to the public key presented by the signer -- authentication.

RSA algorithm is asymmetric cryptography algorithm. Asymmetric actually means that it works on two different keys i.e. Public Key and Private Key. As the name describes that the Public Key is given to everyone and Private key is kept private.

An example of asymmetric cryptography:

- 1. A client (for example browser) sends its public key to the server and requests for some data.
- 2. The server encrypts the data using client's public key and sends the encrypted data.
- 3. Client receives this data and decrypts it.

Since this is asymmetric, nobody else except browser can decrypt the data even if a third party has public key of browser.

The idea of RSA is based on the fact that it is difficult to factorize a large integer. The public key consists of two numbers where one number is multiplication of two large prime numbers. And private key is also derived from the same two prime numbers. So if somebody can factorize the large number, the private key is compromised. Therefore encryption strength totally lies on the key size and if we double or triple the key size, the strength of encryption increases exponentially. RSA keys can be typically 1024 or 2048 bits long, but experts believe



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that 1024 bit keys could be broken in the near future. But till now it seems to be an infeasible task.

We demonstrate RSA with the help of cryptool







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	ulated such that $d = e^{-1} \pmod{phi(N)}$ .	n but must be coprime to the totient. The prival ed the public RSA parameters: the modulus N
and the public ke	y e.	sa the public from parameters, the modulus fr
rime number entry—	82	50.45
<sup>o</sup> rime number p	211	Generate prime numbers
Prime number q	233	
SA parameters		
RSA modulus N	49163	(public)
ohi(N) = (p-1)(q-1)	48720	(secret)
Public key e	2^16+1	
Private key d	44273	Update parameters
SA encryption using	e / decryption using d [alphabet size: 256]-	
nput as C text	• numbers	Alphabet and number system options
Ciphertext coded in n		
		64 # 09394 # 00622 # 13392 # 16226 # 3927
Decruption into plainte		00082 # 00089 # 00080 # 00084 # 00073 # (
	correction fints accreents of size 1: the cumb	val '#' in used as separator)
00082 # 00083 # 00 Dutput text from the d	lecryption (into segments of size 1; the symb	
00082 # 00083 # 00 Dutput text from the d	lecryption (into segments of size 1; the symb	





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A digital signature scheme typically consists of 3 algorithms;

- A key generation algorithm that selects a private key uniformly at random from a set of
  possible private keys. The algorithm outputs the private key and a corresponding public
  key.
- A signing algorithm that, given a message and a private key, produces a signature.
- A *signature verifying* algorithm that, given the message, public key and signature, either accepts or rejects the message's claim to authenticity.

Two main properties are required. First, the authenticity of a signature generated from a fixed message and fixed private key can be verified by using the corresponding public key. Secondly, it should be computationally infeasible to generate a valid signature for a party without knowing that party's private key. A digital signature is an authentication mechanism that enables the creator of the message to attach a code that acts as a signature.

One digital signature scheme (of many) is based on RSA. To create signature keys, generate a RSA key pair containing a modulus, N, that is the product of two random secret distinct large primes, along with integers, e and d, such that e  $d \equiv 1 \pmod{\phi(N)}$ , where  $\phi$  is the Euler phifunction. The signer's public key consists of N and e, and the signer's secret key contains d.

To sign a message, m, the signer computes a signature,  $\sigma$ , such that  $\sigma \equiv md \pmod{N}$ . To verify, the receiver checks that  $\sigma \equiv m \pmod{N}$ .



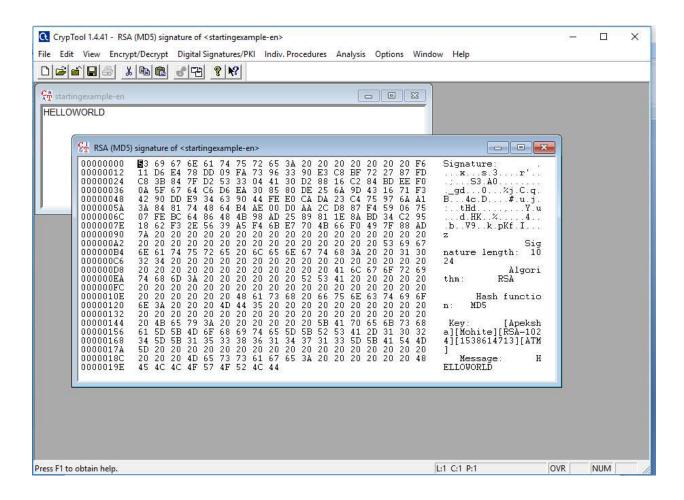


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#### **Digital Signatures using Cryptool**

MD5 digital signature example





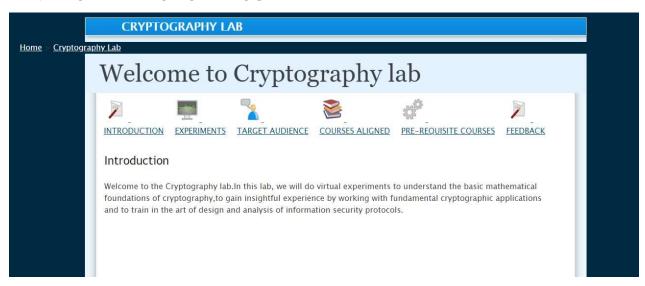
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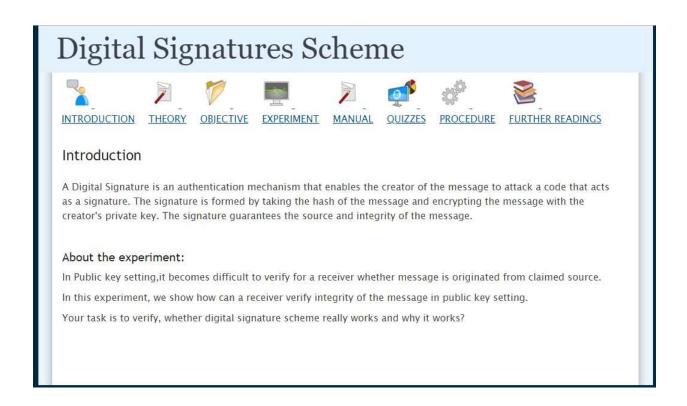
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