

# **Auction-Based Collegiate Athletic Conference Scheduling**

## **An Iterative Market Mechanism for Sports Schedule Formation**

Art Steinmetz

2026-02-06

### **Abstract**

This article presents an iterative auction-based model for forming collegiate athletic conference schedules. The mechanism treats away games as differentiated goods, categorized by opponent strength match and travel requirements. Schools bid on game types according to their preferences, and a market-clearing process determines the final schedule. The model enforces hard constraints including exact home/away game requirements and pairwise meeting limits, while prices emerge endogenously to reflect relative scarcity and desirability. The prices for game types are established early in the auction while many further iterations are required to optimize the schedule. Analysis of simulation results reveals how geographic isolation creates systematic disadvantages for remote schools, who must spend more tokens to acquire less desirable schedules. The auction framework provides a transparent, preference-respecting approach to the complex combinatorial problem of conference scheduling.

### **Introduction**

In the last few decades, collegiate athletic conferences have undergone significant realignment driven by financial incentives, media rights considerations, and competitive balance concerns. This is a particular problem for Division III conferences, where schools have limited travel budgets and fewer peer competitors. Further, schools in a conference play each other in every sport, while schools have varying levels of strength by sport. Some conferences have made agreements to bring in outside schools to some sports but these are ad-hoc arrangements.

A typical Division III conference has ten schools. An alternative structure might allow highly fluid “conference” definitions where many more schools participate. Opponents could be different for each sport and different for each year. Such a structure would encompass a larger geographic area and allow more flexibility for schools to find appropriate opponents. However, the scheduling problem becomes much more complex. Travel time and cost also grow as the number of schools increases.

This article explores an such an approach: using an iterative auction mechanism to form conference schedules. In this framework, away games are treated as differentiated goods, and schools express their preferences through bidding behavior. The market mechanism then allocates games in a way that respects both preferences and hard constraints. While the literature contains many approaches to athletic schedule optimization, I did not see any specifically suggesting an auction market.

The auction-based approach offers several advantages:

1. **Preference revelation:** Schools' bids reveal their true preferences over different types of games
2. **Price discovery:** Equilibrium prices emerge that reflect the relative scarcity of desirable game types
3. **Transparency:** The allocation process is governed by clear rules rather than opaque negotiations
4. **Flexibility:** The mechanism can accommodate heterogeneous preferences across schools

Notably, the auction does not directly allocate specific opponents. Instead, schools bid on categories of games (e.g., “evenly matched bus games”) and the clearing mechanism determines which specific opponents they receive based on geographic compatibility and mutual availability.

We implement and simulate this mechanism for a 20-school conference where each school plays 12 games per season (6 home, 6 away). The simulation demonstrates how the auction produces feasible schedules while generating meaningful price signals about game type desirability.

This paper is written as an Quarto embedded code notebook, so all the code necessary to reproduce and critique the results is visible. I believe this is the future of academic publishing. Quarto publishing, the Positron IDE, the R language were used. While I am an experienced R coder, Claude Opus 4.5 wrote most of the code. “The programming language of the future is English.”

DISCLAIMER: This is a “toy” model, created by a data science hobbyist without domain expertise. There may be errors. I share this in hopes it might inspire some discussion around the future of athletic conferences and provide ideas for student data science explorations.

## Model Description

### Agents and Season Parameters

The model considers a set of  $N = 20$  schools that form an athletic conference. Each school must play exactly  $G = 12$  games per season, split evenly between home and away:

- Home games:  $H = 6$
- Away games:  $A = 6$

Each school begins the season with a budget of  $B_0 = 100$  tokens, used to bid on away games.

### School Strengths

Each school  $i$  is assigned a discrete strength score  $s_i \in \{1, 2, 3\}$ :

- $s_i = 1$ : Weak
- $s_i = 2$ : Moderate
- $s_i = 3$ : Strong

The strength match between two schools is defined by the absolute difference in their strength scores:

$$b_{ij} = \begin{cases} 3 & \text{if } |s_i - s_j| = 0 \text{ (evenly matched)} \\ 2 & \text{if } |s_i - s_j| = 1 \text{ (close match)} \\ 1 & \text{if } |s_i - s_j| = 2 \text{ (mismatched)} \end{cases}$$

Schools prefer evenly matched opponents (band 3) over mismatched opponents (band 1).

## Geography and Travel

Each school has a geographic location characterized by latitude and longitude coordinates. For each ordered pair of schools  $(i, j)$ , travel requirements are determined by:

- **Distance:** Haversine distance in miles between school locations
- **Bus travel time:**  $T_{ij}^{bus} = \text{distance}_{ij}/60$  hours
- **Travel class  $\tau_{ij}$ :**
  - $B$  (bus) if  $T_{ij}^{bus} \leq 5$  hours
  - $P$  (plane) if  $T_{ij}^{bus} > 5$  hours

Travel costs are fixed rates:

Travel Mode	Cost
Home	\$0
Bus	\$1,500
Plane	\$7,500

Plane travel time is treated as constant at 5 hours regardless of distance.

## Game Types

Away games are categorized into 6 types based on the combination of strength match and travel class:

$$k = (\text{strength\_match}, \text{travel\_class}) \in \{1, 2, 3\} \times \{B, P\}$$

Each ordered away match  $(i, j)$  belongs to exactly one game type determined by  $b_{ij}$  and  $\tau_{ij}$ .

## Preference Specification

Schools express preferences over game types through disutilities  $v_{i,k} \leq 0$ . More negative values indicate less desirable game types. Disutility captures:

- Strength band preference (evenly matched preferred)
- Travel time burden

- Travel cost burden

To convert disutilities to willingness-to-pay (value):

$$\text{value}_{i,k} = v_{i,k} - \min_k v_{i,k}$$

This ensures the most-preferred game type has the highest value and the least-preferred has value zero.

## Feasibility Constraints

The final schedule must satisfy:

(a) **Total games per school:**

$$\sum_{j \neq i} x_{ij} + \sum_{j \neq i} x_{ji} = G \quad \forall i$$

(b) **Exact home/away split:**

$$\sum_{j \neq i} x_{ji} = H \quad \forall i \quad (\text{home games})$$

$$\sum_{j \neq i} x_{ij} = A \quad \forall i \quad (\text{away games})$$

(c) **Pairwise meeting cap:**

$$x_{ij} + x_{ji} \leq 1 \quad \forall i < j$$

The constraint above prefers single meetings between any pair of schools. However, if a feasible schedule cannot be constructed with all single matchups, the algorithm relaxes this constraint to allow teams to play each other up to two times (one home, one away).

(d) **Binary decision variables:**

$$x_{ij} \in \{0, 1\}$$

## Iterative Auction Process

The auction proceeds iteratively:

1. **Demand determination:** Each school determines integer demand for each game type given current prices and remaining budget
2. **Clearing problem:** An LP solver finds the value-maximizing allocation respecting feasibility and demand constraints
3. **Price update:** Prices increase for game types with excess demand
4. **Budget update:** Schools pay for allocated games at current prices
5. **Schedule update:** Cumulative schedule matrix is updated

The process terminates when all schools have complete schedules or a maximum iteration limit is reached.

## The Bidder's Experience

From a school's perspective, the auction unfolds as follows. At the start of each round, the school reviews the current prices for each of the six game types and its remaining token budget. The athletic director must decide: how many games of each type do we want at these prices?

Consider a typical school, "Overthinkon." Their ideal schedule would include mostly evenly-matched bus games—nearby opponents of similar strength. But if prices for these premium games have risen to 15 tokens each while mismatched plane games remain at 2 tokens, the school faces a trade-off. Spending heavily on preferred games leaves fewer tokens for later rounds when competition for remaining slots intensifies.

At the same time, a nearby evenly matched opponent which is very desirable for Overthinkon will be seen by another school as a mismatched long-distance game and will not want to pay much for it.

Schools cannot directly bid on specific opponents. Instead, they express demand for *categories* of games. If Overthinkon requests 3 evenly-matched bus games, the clearing mechanism determines which specific opponents they receive based on geographic compatibility and mutual availability. A school might request games against strong nearby teams but end up matched with whoever remains feasible given everyone else's requests.

This indirect allocation creates strategic uncertainty. Schools with many nearby peers can often satisfy their preferences, while geographically isolated schools find themselves price-takers—forced to accept whatever game types their location permits, regardless of willingness to pay.

## Auction Simulation

Now we will implement and simulate the auction mechanism described above. The code is organized into sections corresponding to the model components. Code visibility can be toggled in the interactive version of this article.

### Section 1: Initialize Season Parameters

This section establishes the fundamental parameters that govern the auction simulation, including the number of schools, games per season, and budget allocations. The code is visible in the interactive version of this article.

The conference consists of 20 schools, each playing 12 games with an even split between home and away. Each school starts with 100 tokens to bid on away games.

### Section 2: Create Schools and Geographic Locations

Schools are created with randomly assigned strength scores and random geographic locations within a region approximating the eastern United States. The names are fictional.

## School Locations

Geographic distribution of conference schools in auction

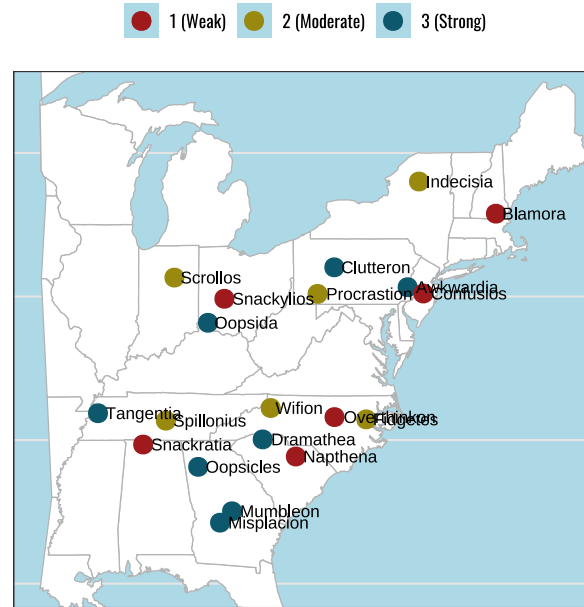


Figure 1: Map of conference school locations with state boundaries

Figure 1 shows the geographic distribution of conference schools with state boundaries for context.

### Section 3: Define Game Types and Preferences

Schools have preferences over different types of away games. Game types are defined by the combination of opponent strength match and travel requirements. In practice schools would reveal their preferences through the bidding process. Here we assign somewhat realistic preferences to simulate the auction. Each school's preferences are generated with some random variation around base disutility values.

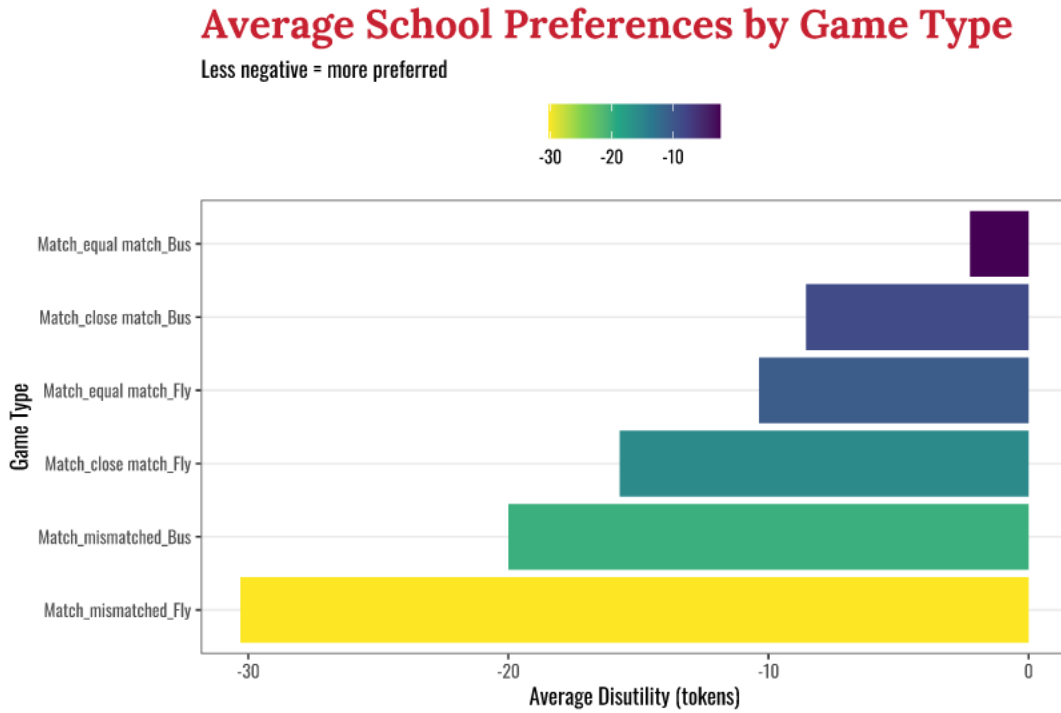


Figure 2: Average school preferences by game type

Schools strongly prefer evenly matched bus games (low disutility) and dislike mismatched plane games (high disutility).

#### Section 4: Create School Pairs with Travel Information

This section computes the travel requirements between all pairs of schools, determining which games require plane travel versus bus travel. If the bus ride is longer than 5 hours, the game is classified as a plane game with a fixed travel time of 5 hours. The cost for a bus trip is \$1,500 while the cost for a plane trip is \$7,500. This is not intended to be an accurate model of travel costs, just a simple way to differentiate travel types. Likewise, we use the straight-line distance rather than actual driving distance, which would not be realistic if we were using actual schools.

## Distribution of Travel Times Between Schools

Travel class determines transportation method

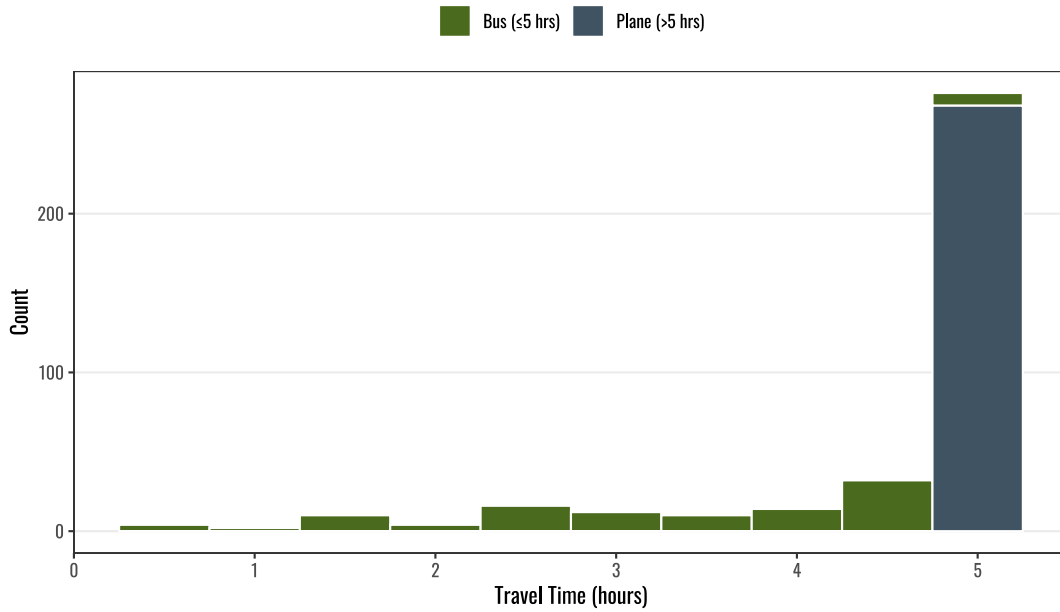


Figure 3: Distribution of travel times between schools

## Distribution of Away Game Types

By match strength and travel class

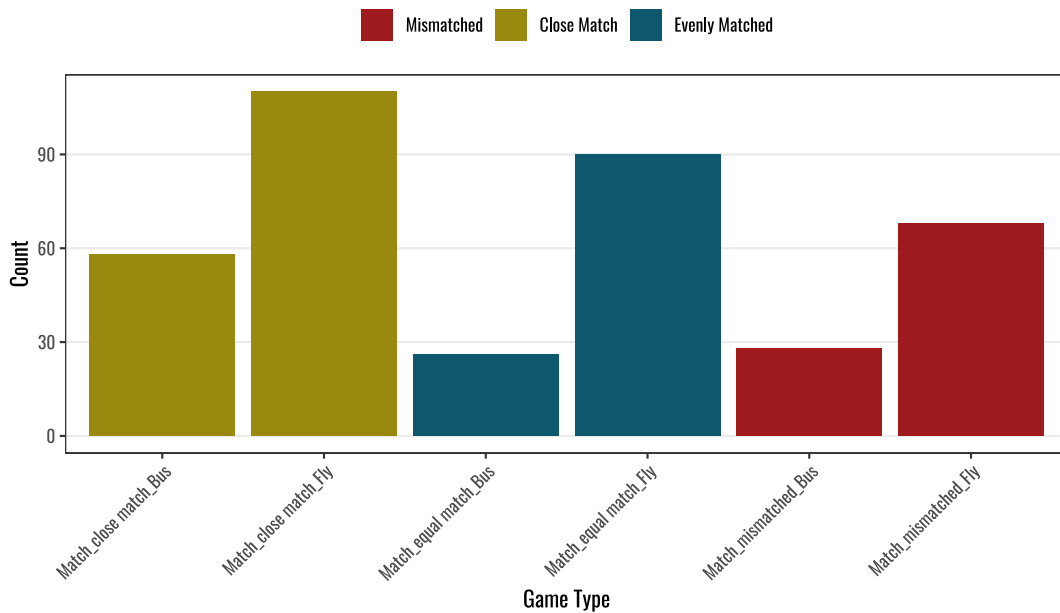


Figure 4: Distribution of away game types



Some schools are more geographically isolated than others. The table below quantifies each school's isolation score, which reflects the average distance to its opponents weighted by match quality. Inevitably, schools located far from the conference core must travel longer distances to reach well-matched opponents.

Table 1: Geographic isolation scores by school

School	Strength	Avg	Distance (mi)	Weighted	Avg (mi)	Bus	OpEqual+Bus tions	Isolation Score
Blamora	1		693		671	3	1	150
Indecisia	2		623		621	4	0	139
Tangen- tia	3		537		518	3	1	116
Snackra- tia	1		466		512	5	0	115
Awkwar- dia	3		471		504	5	1	113
Confu- sios	1		492		478	5	1	107
Scrollos	2		460		460	2	0	103
Mispla- cion	3		481		456	6	3	102
Clut- teron	3		417		437	5	1	98
Mum- bleon	3		455		431	7	3	97
Spillo- nius	2		414		423	8	1	95
Fidgetes	2		416		420	4	1	94
Snack- yllos	1		386		400	5	0	90
Napthena	1		380		399	7	1	89
Oopsi- cles	3		413		391	8	4	87
Procras- tion	2		381		379	7	1	85
Over- thinkon	1		375		378	5	1	85
Opsida	3		378		377	5	0	84
Dra- mathea	3		353		345	8	3	77
Wifion	2		334		338	10	3	76

The isolation score weights distance by opponent match quality—schools whose evenly-matched opponents are far away score higher than schools with nearby well-matched peers.

## Section 5: Define Auction Helper Functions

The auction requires several helper functions for demand calculation, feasibility checking, and schedule updates.

## Section 6: LP-Based Clearing Mechanism

The heart of the auction is the clearing problem, solved as an integer linear program to find the value-maximizing allocation.

## Section 7: Run the Auction

With all components in place, we execute the iterative auction to produce the conference schedule.

The auction completed in 100 iterations, producing a complete schedule for all 20 schools.

# Results

## Schedule Verification

We first verify that the produced schedule satisfies all feasibility constraints.

Table 2: Schedule constraint verification

Constraint	Required	Observed	Satisfied
Total games per school	<b>12</b>	12-12	TRUE
Home games per school	<b>6</b>	6-6	TRUE
Away games per school	<b>6</b>	6-6	TRUE
Max pairwise meetings	$\leq 2$	2	TRUE

## Final Prices

The auction produces prices for each game type that reflect relative demand and scarcity. Note that prices reach their terminal values quickly, often within the first 20 iterations. Beyond that, prices stabilize as schools exhaust their budgets.

## Price Evolution During Auction

How prices adjusted based on demand

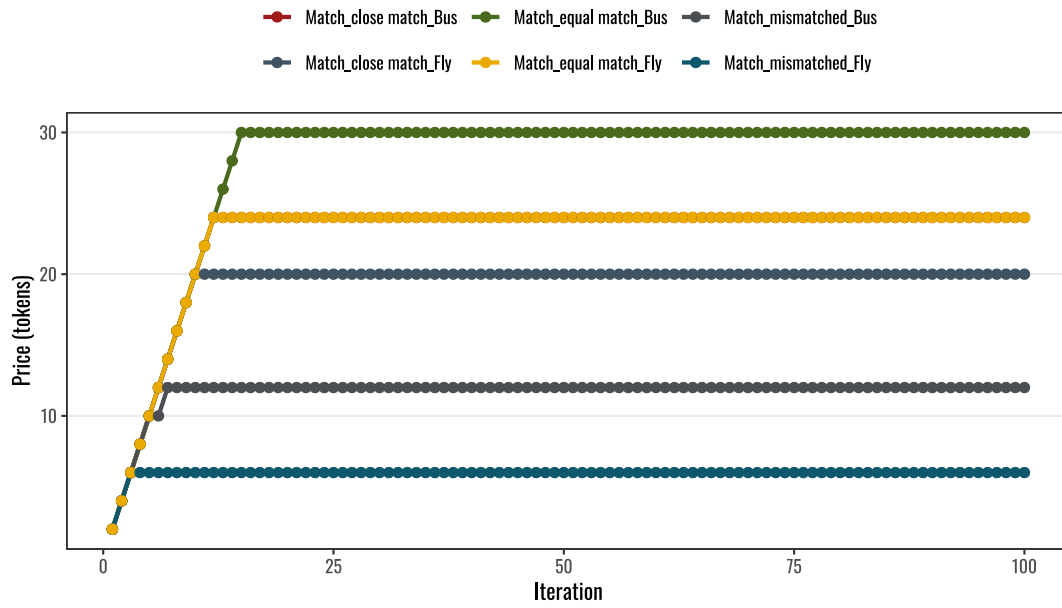


Figure 5: Price evolution during the auction

## Travel Analysis

Geographic location significantly impacts school outcomes in terms of travel burden.

## Total Travel Cost by School

Bus: \$1,500 | Plane: \$7,500 | Color indicates schedule disutility

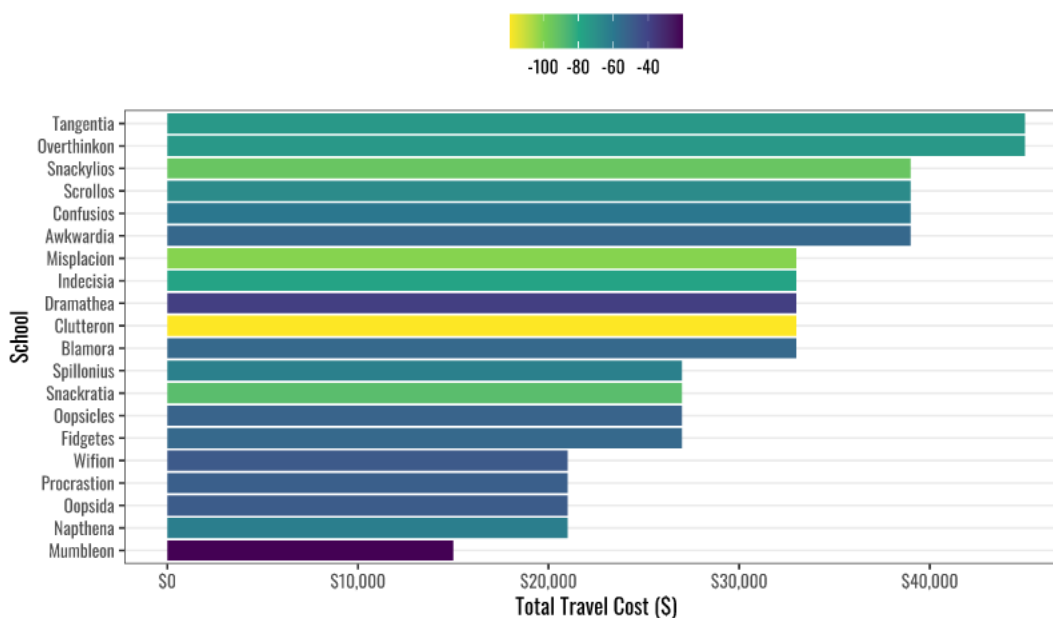


Figure 6: Total travel cost by school

## Total Travel Hours by School

Time spent traveling to away games | Color indicates schedule disutility

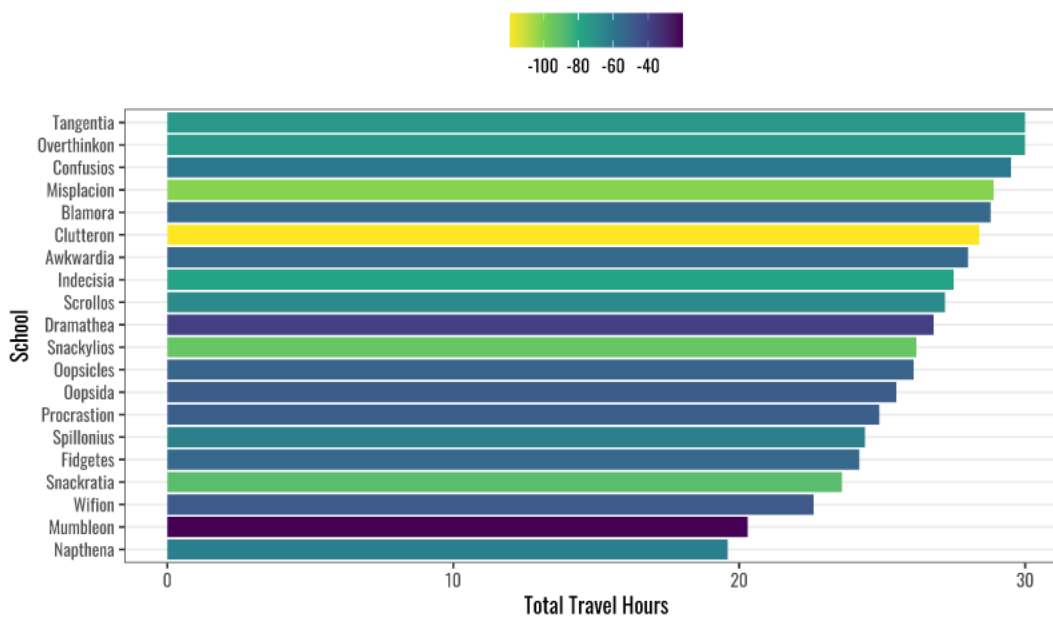


Figure 7: Total travel hours by school

### Budget and Disutility Analysis

We examine how token spending relates to schedule quality (total disutility).

### Tokens Spent vs Total Disutility Accepted

Schools spending more tokens should get better (less negative) schedules

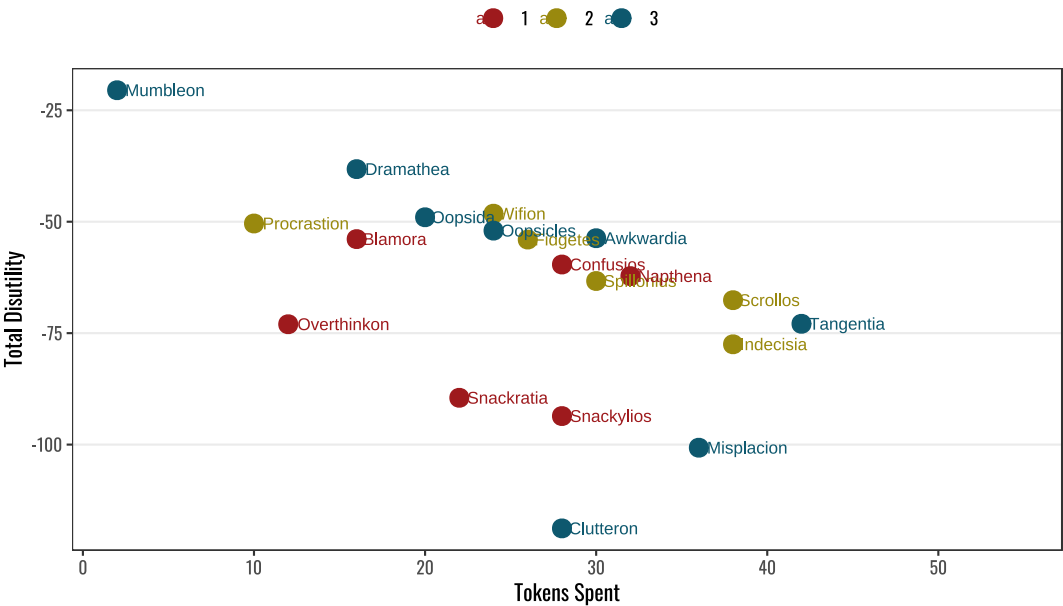


Figure 8: Tokens spent versus total disutility

### Geographic Isolation Effects

A key finding is that geographic isolation systematically disadvantages schools.

## Geographic Isolation vs Schedule Quality

Correlation: -0.12 | Schools farther from others get worse schedules

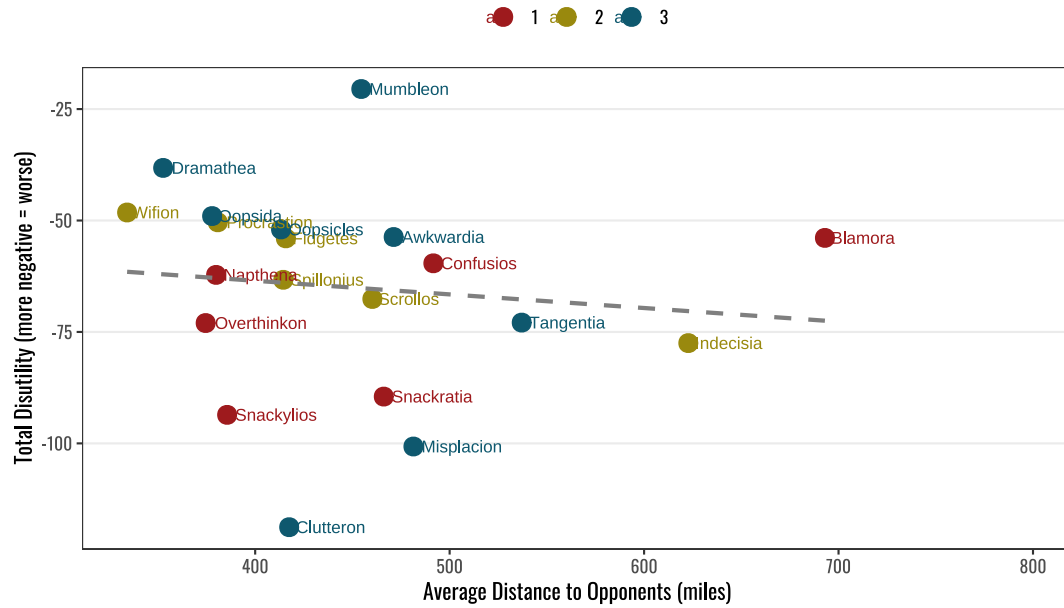


Figure 9: Geographic isolation versus schedule quality

## Conference Schedule Map

### Conference Schedule Map

Lines connect schools that play each other | 114 unique matchups

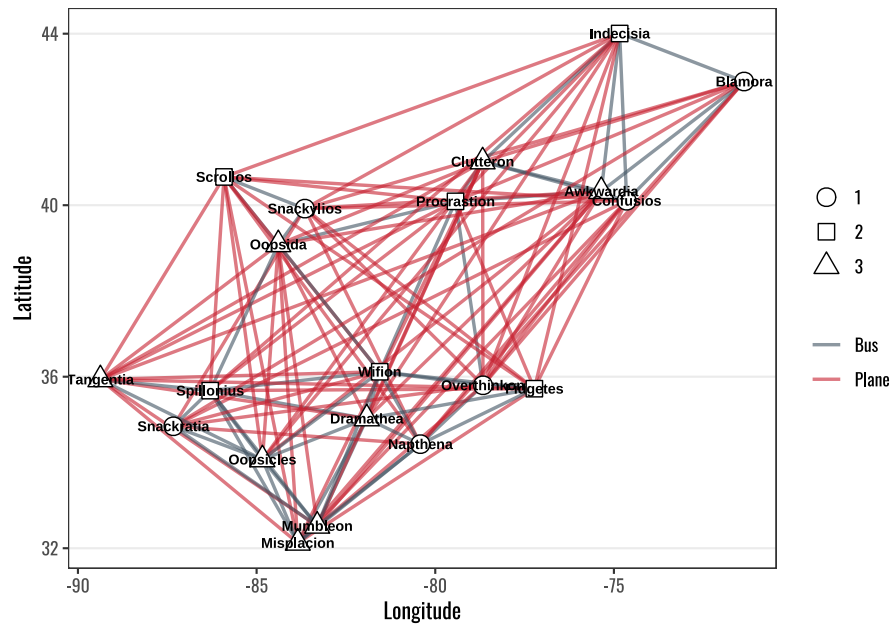


Figure 10: Conference schedule map showing all matchups

## Overthinkon Schedule

Lines connect to opponents | Dashed lines indicate air travel



Figure 11: Example schedule for a single school



Table 3: Overthinkon’s complete season schedule

Week	Opponent	Location	Opp. Strength	Travel Mode	Opp. Mode	Distance (mi)	Travel (hrs)	Time
1	Procrastion	Home	2	—	Bus	299	5	
2	Snackyllos	Away	1	Plane	—	394	5	
3	Clutteron	Home	3	—	Plane	360	5	
4	Confusios	Away	1	Plane	—	369	5	
5	Indecisia	Home	2	—	Plane	602	5	
6	Awkwardia	Away	3	Plane	—	361	5	
7	Fidgetes	Home	2	—	Bus	81	1.3	
8	Blamora	Away	1	Plane	—	626	5	
9	Napthena	Home	1	—	Bus	137	2.3	
10	Mumbleon	Away	3	Plane	—	349	5	
11	Wifion	Home	2	—	Bus	163	2.7	
12	Snackratia	Away	1	Plane	—	493	5	

In this particular example, Overthinkon has received a package where all their away games are by plane, resulting in the highest travel cost among all schools. Let’s look at how their utility budget might have resulted in this outcome.

Table 4: Overthinkon’s game type preferences compared to conference average

Game Type	Overthinkon	Conf. Avg	Difference	Interpretation
Match_equal match_Bus	−3.1	−2.3	−0.8	Similar
Match_equal match_Fly	−5.5	−10.4	4.9	Less averse
Match_close match_Bus	−11.5	−8.6	−2.9	More averse
Match_close match_Fly	−17.2	−15.7	−1.5	More averse
Match_mismatched_Bus	−22.9	−20.0	−2.9	More averse
Match_mismatched_Fly	−25.5	−30.3	4.8	Less averse

Schools with less aversion to plane travel (less negative disutility for plane game types - bigger travel budgets) would be more willing to bid on those games. The problem for Overthinkon is their well-matched opponents are all far away, forcing them to acquire plane games to get quality matchups.

## Schedule Heatmap

Schools will not meet the same opponent multiple times in a season unless the schedule cannot be filled otherwise. The schedule heatmap below shows how many times each pair of schools meet during the season. In this run Snackyllos and Figetes meet twice, while all other pairs meet either once or not at all. The likelihood of multiple meetings depends on the utility budgets and geographic distribution of schools.

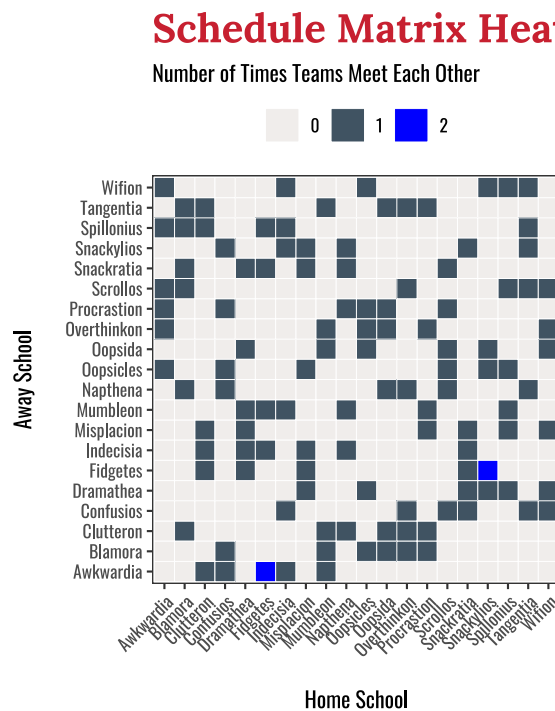


Figure 12: Schedule matrix heatmap

## Conclusion

This simulation demonstrates that an iterative auction mechanism can successfully produce feasible conference schedules while respecting school preferences and hard constraints. Key findings include:

1. **Successful schedule formation:** The auction consistently produces complete schedules satisfying all constraints (12 games per school, exactly 6 home and 6 away, at most two meetings per pair).
2. **Meaningful price discovery:** Prices for desirable game types (evenly matched, bus travel) rise during the auction, reflecting genuine scarcity and preferences.

3. **Geographic disadvantage:** Schools in geographically isolated locations face structural disadvantages. They must acquire more plane games regardless of preference, leading to:
  - Higher token expenditure
  - Worse schedule quality (higher total disutility)
  - This inverse correlation between spending and outcomes reflects the geographic constraints embedded in the game type structure.
4. **Preference heterogeneity matters:** Schools with different strength profiles have different sets of opponents available for “evenly matched” games, affecting their ability to achieve preferred schedules.

The auction framework provides a transparent, rules-based approach to conference scheduling that could be adapted for real-world applications. Future extensions could incorporate additional constraints such as mandatory rivalries and dollar budgets for travel costs.

The key feature of this model is that auction prices reflect game *types*, not individual matchups. Schools cannot escape their geographic constraints through bidding behavior alone—a school surrounded by distant opponents must acquire plane games regardless of their willingness to pay for bus games. This structural feature of the mechanism has important equity implications for conference design.