## **NAME**

get\_thread\_area, set\_thread\_area - manipulate thread-local storage information

#### **LIBRARY**

```
Standard C library (libc, -lc)
```

#### **SYNOPSIS**

```
#include <sys/syscall.h> /* Definition of SYS_* constants */
#include <unistd.h>
#if defined __i386__ || defined __x86_64__
# include <asm/ldt.h> /* Definition of struct user_desc */
int syscall(SYS_get_thread_area, struct user_desc *u_info);
int syscall(SYS_set_thread_area, struct user_desc *u_info);
#elif defined __m68k__
int syscall(SYS_get_thread_area);
int syscall(SYS_set_thread_area, unsigned long tp);
#elif defined __mips__
int syscall(SYS_set_thread_area, unsigned long addr);
#endif
```

*Note*: glibc provides no wrappers for these system calls, necessitating the use of **syscall**(2).

# DESCRIPTION

These calls provide architecture-specific support for a thread-local storage implementation. At the moment, **set\_thread\_area**() is available on m68k, MIPS, and x86 (both 32-bit and 64-bit variants); **get\_thread\_area**() is available on m68k and x86.

On m68k and MIPS, **set\_thread\_area**() allows storing an arbitrary pointer (provided in the **tp** argument on m68k and in the **addr** argument on MIPS) in the kernel data structure associated with the calling thread; this pointer can later be retrieved using **get\_thread\_area**() (see also NOTES for information regarding obtaining the thread pointer on MIPS).

On x86, Linux dedicates three global descriptor table (GDT) entries for thread-local storage. For more information about the GDT, see the Intel Software Developer's Manual or the AMD Architecture Programming Manual.

Both of these system calls take an argument that is a pointer to a structure of the following type:

```
struct user_desc {
    unsigned int entry_number;
    unsigned int base_addr;
    unsigned int limit;
    unsigned int seg_32bit:1;
    unsigned int contents:2;
    unsigned int read_exec_only:1;
    unsigned int limit_in_pages:1;
    unsigned int seg_not_present:1;
    unsigned int useable:1;
#ifdef __x86_64__
    unsigned int lm:1;
#endif
};
```

**get\_thread\_area**() reads the GDT entry indicated by  $u\_info->entry\_number$  and fills in the rest of the fields in  $u\_info$ .

set\_thread\_area() sets a TLS entry in the GDT.

The TLS array entry set by  $set\_thread\_area()$  corresponds to the value of  $u\_info->entry\_number$  passed in by the user. If this value is in bounds,  $set\_thread\_area()$  writes the TLS descriptor pointed to by  $u\_info$  into the thread's TLS array.

When  $set\_thread\_area()$  is passed an  $entry\_number$  of -1, it searches for a free TLS entry. If  $set\_thread\_area()$  finds a free TLS entry, the value of  $u\_info->entry\_number$  is set upon return to show which entry was changed.

A *user\_desc* is considered "empty" if *read\_exec\_only* and *seg\_not\_present* are set to 1 and all of the other fields are 0. If an "empty" descriptor is passed to **set\_thread\_area**(), the corresponding TLS entry will be cleared. See BUGS for additional details.

Since Linux 3.19, **set\_thread\_area**() cannot be used to write non-present segments, 16-bit segments, or code segments, although clearing a segment is still acceptable.

### **RETURN VALUE**

On x86, these system calls return 0 on success, and -1 on failure, with errno set to indicate the error.

On MIPS and m68k, **set\_thread\_area**() always returns 0. On m68k, **get\_thread\_area**() returns the thread area pointer value (previously set via **set\_thread\_area**()).

### **ERRORS**

## **EFAULT**

*u\_info* is an invalid pointer.

#### **EINVAL**

*u\_info->entry\_number* is out of bounds.

#### **ENOSYS**

get\_thread\_area() or set\_thread\_area() was invoked as a 64-bit system call.

#### **ESRCH**

(set thread area()) A free TLS entry could not be located.

### **VERSIONS**

set\_thread\_area() first appeared in Linux 2.5.29. get\_thread\_area() first appeared in Linux 2.5.32.

### **STANDARDS**

**set\_thread\_area**() and **get\_thread\_area**() are Linux-specific and should not be used in programs that are intended to be portable.

## **NOTES**

These system calls are generally intended for use only by threading libraries.

**arch\_prctl**(2) can interfere with **set\_thread\_area**() on x86. See **arch\_prctl**(2) for more details. This is not normally a problem, as **arch\_prctl**(2) is normally used only by 64-bit programs.

On MIPS, the current value of the thread area pointer can be obtained using the instruction:

```
rdhwr dest, $29
```

This instruction traps and is handled by kernel.

## **BUGS**

On 64-bit kernels before Linux 3.19, one of the padding bits in *user\_desc*, if set, would prevent the descriptor from being considered empty (see **modify\_ldt**(2)). As a result, the only reliable way to clear a TLS entry is to use **memset**(3) to zero the entire *user\_desc* structure, including padding bits, and then to set the *read\_exec\_only* and *seg\_not\_present* bits. On Linux 3.19, a *user\_desc* consisting entirely of zeros except for *entry\_number* will also be interpreted as a request to clear a TLS entry, but this behaved differently on older kernels.

Prior to Linux 3.19, the DS and ES segment registers must not reference TLS entries.

## **SEE ALSO**

 $\begin{array}{lll} \textbf{arch\_prctl}(2), & \textbf{modify\_ldt}(2), & \textbf{ptrace}(2) & (\textbf{PTRACE\_GET\_THREAD\_AREA} & \text{and} \\ \textbf{PTRACE\_SET\_THREAD\_AREA}) \\ \end{array}$