## **NAME**

fuse - Filesystem in Userspace (FUSE) device

#### **SYNOPSIS**

#include ux/fuse.h>

#### DESCRIPTION

This device is the primary interface between the FUSE filesystem driver and a user-space process wishing to provide the filesystem (referred to in the rest of this manual page as the *filesystem daemon*). This manual page is intended for those interested in understanding the kernel interface itself. Those implementing a FUSE filesystem may wish to make use of a user-space library such as *libfuse* that abstracts away the low-level interface.

At its core, FUSE is a simple client-server protocol, in which the Linux kernel is the client and the daemon is the server. After obtaining a file descriptor for this device, the daemon may **read**(2) requests from that file descriptor and is expected to **write**(2) back its replies. It is important to note that a file descriptor is associated with a unique FUSE filesystem. In particular, opening a second copy of this device, will not allow access to resources created through the first file descriptor (and vice versa).

#### The basic protocol

Every message that is read by the daemon begins with a header described by the following structure:

```
struct fuse_in_header {
   uint32_t len;
                       /* Total length of the data,
                          including this header */
   uint32_t opcode;    /* The kind of operation (see below) */
                       /* A unique identifier for this request */
   uint64 t unique;
   uint64_t nodeid;
                       /* ID of the filesystem object
                          being operated on */
   uint32 t uid;
                       /* UID of the requesting process */
   uint32_t gid;
                       /* GID of the requesting process */
   uint32_t pid;
                       /* PID of the requesting process */
   uint32_t padding;
};
```

The header is followed by a variable-length data portion (which may be empty) specific to the requested operation (the requested operation is indicated by *opcode*).

The daemon should then process the request and if applicable send a reply (almost all operations require a reply; if they do not, this is documented below), by performing a **write**(2) to the file descriptor. All replies must start with the following header:

This header is also followed by (potentially empty) variable-sized data depending on the executed request. However, if the reply is an error reply (i.e., *error* is set), then no further payload data should be sent, independent of the request.

# **Exchanged messages**

This section should contain documentation for each of the messages in the protocol. This manual page is currently incomplete, so not all messages are documented. For each message, first the struct sent by the kernel is given, followed by a description of the semantics of the message.

## **FUSE\_INIT**

```
struct fuse_init_in {
    uint32_t major;
    uint32_t minor;
    uint32_t max_readahead; /* Since protocol v7.6 */
    uint32_t flags; /* Since protocol v7.6 */
};
```

This is the first request sent by the kernel to the daemon. It is used to negotiate the protocol version and other filesystem parameters. Note that the protocol version may affect the layout of any structure in the protocol (including this structure). The daemon must thus remember the negotiated version and flags for each session. As of the writing of this man page, the highest supported kernel protocol version is 7.26.

Users should be aware that the descriptions in this manual page may be incomplete or incorrect for older or more recent protocol versions.

The reply for this request has the following format:

If the major version supported by the kernel is larger than that supported by the daemon, the reply shall consist of only *uint32\_t major* (following the usual header), indicating the largest major version supported by the daemon. The kernel will then issue a new **FUSE\_INIT** request conforming to the older version. In the reverse case, the daemon should quietly fall back to the kernel's major version.

The negotiated minor version is considered to be the minimum of the minor versions provided by the daemon and the kernel and both parties should use the protocol corresponding to said minor version.

## FUSE\_GETATTR

The requested operation is to compute the attributes to be returned by **stat**(2) and similar operations for the given filesystem object. The object for which the attributes should be computed is indicated either by *header->nodeid* or, if the **FUSE\_GETATTR\_FH** flag is set, by the file handle *fh*. The latter case of operation is analogous to **fstat**(2).

For performance reasons, these attributes may be cached in the kernel for a specified duration of time. While the cache timeout has not been exceeded, the attributes will be served from the cache and will not cause additional **FUSE\_GETATTR** requests.

The computed attributes and the requested cache timeout should then be returned in the following structure:

```
struct fuse_attr_out {
              /* Attribute cache duration (seconds + nanoseconds) */
             uint64_t attr_valid;
             uint32_t attr_valid_nsec;
             uint32_t dummy;
             struct fuse_attr {
                 uint64_t ino;
                  uint64 t size;
                  uint64_t blocks;
                  uint64_t atime;
                  uint64_t mtime;
                  uint64_t ctime;
                  uint32_t atimensec;
                  uint32_t mtimensec;
                  uint32_t ctimensec;
                  uint32_t mode;
                  uint32_t nlink;
                  uint32_t uid;
                  uint32_t gid;
                  uint32_t rdev;
                  uint32_t blksize;
                  uint32_t padding;
              } attr;
          };
FUSE_ACCESS
         struct fuse_access_in {
             uint32_t mask;
             uint32_t padding;
```

If the *default\_permissions* mount options is not used, this request may be used for permissions checking. No reply data is expected, but errors may be indicated as usual by setting the *error* field in the reply header (in particular, access denied errors may be indicated by returning **–EACCES**).

# FUSE\_OPEN and FUSE\_OPENDIR

};

The requested operation is to open the node indicated by *header->nodeid*. The exact semantics of what this means will depend on the filesystem being implemented. However, at the very least the filesystem should validate that the requested *fla gs* are valid for the indicated resource and then send a reply with the following format:

```
struct fuse_open_out {
    uint64_t fh;
    uint32_t open_flags;
    uint32_t padding;
};
```

The fh field is an opaque identifier that the kernel will use to refer to this resource The open\_flags field is a bit mask of any number of the flags that indicate properties of this file handle to the

kernel:

## FOPEN\_DIRECT\_IO

Bypass page cache for this open file.

## FOPEN\_KEEP\_CACHE

Don't invalidate the data cache on open.

### FOPEN\_NONSEEKABLE

The file is not seekable.

# FUSE\_READ and FUSE\_READDIR

```
struct fuse_read_in {
    uint64_t fh;
    uint64_t offset;
    uint32_t size;
    uint32_t read_flags;
    uint64_t lock_owner;
    uint32_t flags;
    uint32_t padding;
};
```

The requested action is to read up to *size* bytes of the file or directory, starting at *offset*. The bytes should be returned directly following the usual reply header.

#### **FUSE INTERRUPT**

```
struct fuse_interrupt_in {
    uint64_t unique;
};
```

The requested action is to cancel the pending operation indicated by *unique*. This request requires no response. However, receipt of this message does not by itself cancel the indicated operation. The kernel will still expect a reply to said operation (e.g., an *EINTR* error or a short read). At most one **FUSE\_INTERRUPT** request will be issued for a given operation. After issuing said operation, the kernel will wait uninterruptibly for completion of the indicated request.

## FUSE\_LOOKUP

Directly following the header is a filename to be looked up in the directory indicated by *header->nodeid*. The expected reply is of the form:

The combination of *nodeid* and *generation* must be unique for the filesystem's lifetime.

The interpretation of timeouts and *attr* is as for **FUSE GETATTR**.

## FUSE\_FLUSH

```
struct fuse_flush_in {
    uint64_t fh;
    uint32_t unused;
    uint32_t padding;
    uint64_t lock_owner;
};
```

The requested action is to flush any pending changes to the indicated file handle. No reply data is expected. However, an empty reply message still needs to be issued once the flush operation is complete.

### FUSE\_RELEASE and FUSE\_RELEASEDIR

```
struct fuse_release_in {
    uint64_t fh;
    uint32_t flags;
    uint32_t release_flags;
    uint64_t lock_owner;
};
```

These are the converse of **FUSE\_OPEN** and **FUSE\_OPENDIR** respectively. The daemon may now free any resources associated with the file handle *fh* as the k ernel will no longer refer to it. There is no reply data associated with this request, but a reply still needs to be issued once the request has been completely processed.

#### **FUSE STATES**

This operation implements **statfs**(2) for this filesystem. There is no input data associated with this request. The expected reply data has the following structure:

```
struct fuse_kstatfs {
    uint64_t blocks;
    uint64_t bfree;
    uint64_t files;
    uint64_t files;
    uint64_t ffree;
    uint32_t bsize;
    uint32_t namelen;
    uint32_t padding;
    uint32_t spare[6];
};

struct fuse_statfs_out {
    struct fuse_kstatfs st;
};
```

For the interpretation of these fields, see **statfs**(2).

#### **ERRORS**

**E2BIG** Returned from **read**(2) operations when the kernel's request is too large for the provided buffer and the request was **FUSE\_SETXATTR**.

## **EINVAL**

Returned from **write**(2) if validation of the reply failed. Not all mistakes in replies will be caught by this validation. However, basic mistakes, such as short replies or an incorrect *unique* value, are detected.

**EIO** Returned from **read**(2) operations when the kernel's request is too large for the provided buffer.

*Note*: There are various ways in which incorrect use of these interfaces can cause operations on the provided filesystem's files and directories to fail with **EIO**. Among the possible incorrect uses are:

- changing *mode & S\_IFMT* for an inode that has previously been reported to the kernel; or
- giving replies to the kernel that are shorter than what the kernel expected.

#### **ENODEV**

Returned from **read**(2) and **write**(2) if the FUSE filesystem was unmounted.

#### **EPERM**

Returned from operations on a /de v/fuse file descriptor that has not been mounted.

# **STANDARDS**

The FUSE filesystem is Linux-specific.

### **NOTES**

The following messages are not yet documented in this manual page:

```
FUSE_BATCH_FORGET
FUSE_BMAP
FUSE_CREATE
FUSE_DESTROY
FUSE_FALLOCATE
FUSE_FORGET
FUSE_FSYNC
FUSE_FSYNCDIR
FUSE_GETLK
FUSE_GETXATTR
FUSE_IOCTL
FUSE_LINK
FUSE_LISTXATTR
FUSE_LSEEK
FUSE_MKDIR
FUSE_MKNOD
FUSE_NOTIFY_REPLY
FUSE_POLL
FUSE_READDIRPLUS
FUSE_READLINK
FUSE_REMOVEXATTR
FUSE_RENAME
FUSE_RENAME2
FUSE_RMDIR
FUSE_SETATTR
FUSE_SETLK
FUSE_SETLKW
```

# **SEE ALSO**

fusermount(1), mount.fuse(8)

FUSE\_SYMLINK FUSE\_UNLINK FUSE\_WRITE