

# Gravel2 Kernel API

*18–342 Fundamentals of Embedded Systems*

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## 1 Kernel API

This document lists all the syscalls that the Gravel real-time kernel supports. All syscalls are to be made using the SWI instruction. The SWI number to be used is the syscall number of your syscall added to a base of 0x900000. Arguments to a syscall are put in order in r0, r1, r2, r3. The return value of a syscall is put in r0. No registers other than r0 are modified during a syscall. Control is transferred to the instruction succeeding the SWI after the syscall is complete. A syscall with an invalid syscall number results in the task being forcibly exited with an exit status of 0x0badc0de. A syscall returning an error returns the negative of the error codes mentioned below.

## 2 SWI Numbers

SWI Number	Syscall Name
0x900003	read
0x900004	write
0x900006	time
0x900007	sleep

## 3 Syscalls

### 3.1 read

**Syscall Number** 3

**Description** read attempts to read a requested number of bytes from the given file descriptor. The read values are written into the provided buffer sequentially.

**Argument 0** int fd

The file descriptor to read from — e.g. STDIN\_FILENO

**Argument 1** void \*buf

The target buffer to write any read data to.

**Argument 2** size\_t count

The number of bytes to read into buf. If this value is zero, then zero is returned and nothing else happens. If this value is greater than SSIZE\_MAX, the result is unspecified.

**Return** ssize\_t ret

On success, the number of bytes read is returned. This number can be less than count if a short read occurs due to an end-of-file or another error. The following errors may be returned.

**EBADF** fd was an invalid or unreadable file descriptor.

**EFAULT** buf points to region whose bounds lie outside valid address space.

**EIO** An internal I/O error occurred.

## 3.2 write

### Syscall Number 4

**Description** `write` attempts to write a requested number of bytes to the given file descriptor. The values to be written are read from the provided buffer sequentially.

**Argument 0** `int fd`

The file descriptor to write to — e.g. `STDOUT_FILENO`

**Argument 1** `const void *buf`

The target buffer to read data from.

**Argument 2** `size_t count`

The number of bytes to read from `buf`. If this value is zero, then zero is returned and nothing else happens. If this value is greater than `SSIZE_MAX`, the result is unspecified. Note that this behavior is subtly different from the linux behavior and is used here to remain consistent with `read`.

**Return** `ssize_t ret`

On success, the number of bytes written is returned. This number can be less than `count` if the underlying medium cannot accept any more bytes at the moment. The following errors may be returned on failure.

**EBADF** `fd` was an invalid or unwritable file descriptor.

**EFAULT** `buf` points to region whose bounds lie outside valid address space.

**EIO** An internal I/O error occurred.

## 3.3 time

### Syscall Number 6

**Description** `time` returns the time in milliseconds that have elapsed since the kernel booted up. This can be interpreted as the number of milliseconds to elapse since the kernel timer was turned on.

**Return** `unsigned long ret`

The function returns the time elapsed since the kernel booted up in milliseconds. There is no error case. Note that even though the number returned is in milliseconds, the actual resolution may be lower.

## 3.4 sleep

### Syscall Number 7

**Description** `sleep` suspends the execution of the current task for the given time. The task may continue executing after atleast the given time has elapsed.

**Argument 0** `unsigned long millis`

The number of milliseconds to sleep. Note that the kernel may round up the time provided if its timer resolution is lower than a millisecond.

**Return** `void`

There is no return value. The function always succeeds.

### 3.5 task\_create

**Syscall Number** 10

**Description** task\_create takes a set of tasks, verifies that they are schedulable, and then begins to schedule them. All currently present tasks are cancelled, and an entirely new task set is scheduled upon success.

**Argument 0** task\_t \*

An array of task descriptors that qualify the tasks to create.

**Argument 0** size\_t num\_tasks

The number of elements in the array.

**Return** int

Upon success, this function does not return. It instead begins scheduling tasks. On failure, an error code may be returned.

**EINVAL** num\_tasks is greater than the maximum number of tasks the OS supports (64).

**EFAULT** tasks points to region whose bounds lie outside valid address space.

**ESCHED** The given task set is not schedulable – some tasks may not meet their deadlines.

### 3.6 mutex\_create

**Syscall Number** 15

**Description** This function returns a unique identifier to a mutex object that the kernel has created. If the kernel is out of identifiers, it returns an error. All valid ids are positive integers. The returned ids start at zero and monotonically increase by one for every successful subsequent call. This function should return successfully for at least 32 invocation.

**Return** int

A positive integer is returned if successful. Upon failure, an error code may be returned.

**ENOMEM** The kernel cannot produce any more mutices.

### 3.7 mutex\_lock

**Syscall Number** 16

**Description** This function instructs the kernel to acquire the indicated mutex. If the mutex is acquirable, the return is zero. If the mutex is unacquirable, the kernel will put the task to sleep until it is acquirable and will then acquire it.

**Argument 0** int mutex

The unique id for a mutex object that was obtained from a successful invocation of mutex\_create.

**Return** int

A zero is returned upon the successful acquisition of the mutex. Upon failure, an error code may be returned.

**EINVAL** The provided mutex identifier is invalid.

**EDEADLOCK** The current task is already holding the lock.

### 3.8 mutex\_unlock

**Syscall Number** 17

**Description** This function instructs the kernel to release an already acquired mutex. If the mutex has other tasks waiting for its release, the first waiting task will be woken up and the mutex will be handed over to that task atomically (with respect to the scheduler).

**Argument 0** `int mutex`

The unique id for a mutex object that was obtained from a successful invocation of `mutex_create`.

**Return** `int`

A zero is returned upon the successful acquisition of the mutex. Upon failure, an error code may be returned.

**EINVAL** The provided mutex identifier is invalid.

**EPERM** The current task does not hold the mutex.

### 3.9 event\_wait

**Syscall Number** 20

**Description** Puts the calling task to sleep on a given device numbers sleep queue until the device next signals the device event. Device events are signalled regularly based on the signalling frequency of the device.

**Argument 0** `unsigned int dev`

A valid device number – 0, 1, 2 or 3.

**Return** `int`

A zero is returned upon the successful acquisition of the mutex. Upon failure, an error code may be returned.

**EINVAL** The provided device identifier is invalid.

## 4 Task Descriptor

```
typedef struct task
{
void (*lambda)(void*);
void *data;
void *stack_pos;
unsigned long C;
unsigned long T;
} task_t;
```

This is the definition of a task descriptor structure. `lambda` is a pointer to the function to be invoked when the task first starts. `data` is the argument that is blindly passed to `lambda`. `stack_pos` is the initial value of userspace `sp` when the task first starts executing. There is no restriction on this value because it is possible to write user applications that never use the stack and hence use `sp` as a general purpose register. `C` and `T` are the worst case compute time and the periodicity of the task respectively. They are both measured in milliseconds. The system is allowed to round/approximate the given values based on the internal resolution of the OS timer as long as it does not incorrectly cause a task to miss its deadline. In general, this means rounding up `C` and rounding down `T`.