

Administrativia: no discussions, no extra material consulted

Problem

Let S denote all subsets of some set of elements U .

For $A, B \in S$, define $A \Delta B = \{x \in U \mid (x \in A) \text{ XOR } (x \in B)\}$.

Note that Δ is commutative and associative since XOR is commutative and associative.

Let $n \in \mathbb{Z}^+$ and let $A_1, \dots, A_n \in S$.

Formally prove (using induction) that, for all $x \in U$, $x \in A_1 \Delta A_2 \Delta \dots \Delta A_n$ if and only if $\#\{i \in \{1, \dots, n\} \mid x \in A_i\}$ is odd.

Solution

For any $i \in \mathbb{N}$, let T_i be the set of all sets consisting of i subsets in S .

Let A_i be the i^{th} element in each $Q_i \in T_i$.

For $n \in \mathbb{N}$, let

$P(n) = \forall \tau_n \in T_n. \forall x \in U. (x \in A_1 \Delta \dots \Delta A_n \text{ IFF } \#\{i \in \{1, \dots, n\} \mid x \in A_i\} \text{ is odd})$.

(1)	$P(1)$	<u>Base Case</u> for all $x \in U$, if $x \in A_1$, then, $\#\{i \in \{1, \dots, n\} \mid x \in A_i\} = 1$, which is odd
(2)	$\forall i \in [1, k] \cap \mathbb{N}. P(i)$ for some $k \in \mathbb{N}$	Inductive Step Assumption
(3)	$P(k)$	Specialisation, (2)
(4)	$\forall Q'_k \in T_k. \forall x \in U.$ $(x \in A_1 \Delta \dots \Delta A_k \text{ IFF } \#\{i \in \{1, \dots, k\} \mid x \in A_i\} \text{ is odd})$	Definition of P , (3)
(5)	$\boxed{\text{Let } S_k \in T_k \text{ be arbitrary.}}$	
(6)	$\boxed{\text{Let } x \in U \text{ be arbitrary.}}$	
(7)	$\boxed{\text{Let } A_1, \dots, A_k \text{ be distinct elements in } S_k.}$	
(8)	$x \in A_1 \Delta \dots \Delta A_k \text{ IFF } \#\{i \in \{1, \dots, k\} \mid x \in A_i\} \text{ is odd}$	Specialisation, (4)
(9)	$\boxed{\text{Let } A_{k+1} \in S \text{ be arbitrary.}}$	
(10)	$A_k \Delta A_{k+1} \in S$	Definition of S
(11)	$\{A_1, \dots, A_{k-1}, A_k \Delta A_{k+1}\}$ is a set of k subsets in S	(5), (10)
(12)	$\boxed{\text{Let } Q_k = A_k \Delta A_{k+1} \text{ and } Q_i = A_i \text{ for } i \in [1, k-1] \cap \mathbb{N}.}$	

(13) $\left| \begin{array}{c} | \\ | \\ | \\ | \end{array} \right| \quad Q = \{Q_1, \dots, Q_k\} \text{ is a set of } k \text{ subsets in } S$ Substitution; (11), (13)

(14) $\left| \begin{array}{c} | \\ | \\ | \\ | \end{array} \right| \quad Q \in T_k$ Definition of Q , (13)

(15) $\left| \begin{array}{c} | \\ | \\ | \\ | \end{array} \right| \quad \begin{array}{l} \forall y \in U. \\ (y \in Q_1 \Delta \dots \Delta Q_k \text{ IFF} \\ \#\{i \in \{1, \dots, k\} \mid y \in Q_i\} \text{ is odd}) \end{array}$ Specialisation for Q , (14)

(16) $\left| \begin{array}{c} | \\ | \\ | \\ | \end{array} \right| \quad \begin{array}{l} (x \in Q_1 \Delta \dots \Delta Q_k \text{ IFF} \\ \#\{i \in \{1, \dots, k\} \mid x \in Q_i\} \text{ is odd}) \end{array}$ Specialisation for x , (15)

(17) $\left| \begin{array}{c} | \\ | \\ | \\ | \end{array} \right| \quad Q_1 \Delta \dots \Delta Q_k = A_1 \Delta \dots \Delta A_k \Delta A_{k+1}$ Definition of Q_i for $i \in [1, k+1] \cap \mathbb{N}$; (12)

(18) $\left| \begin{array}{c} | \\ | \\ | \\ | \end{array} \right| \quad \begin{array}{l} (x \in A_1 \Delta \dots \Delta A_k \Delta A_{k+1} \text{ IFF} \\ \#\{i \in \{1, \dots, k\} \mid x \in Q_i\} \text{ is odd}) \end{array}$ Substitution, (17)

(19) $\left| \begin{array}{c} | \\ | \\ | \\ | \end{array} \right| \quad \begin{array}{l} \#\{i \in \{1, \dots, k\} \mid x \in Q_i\} = \\ \#\{i \in \{1, \dots, k-1\} \mid x \in Q_i\} + \#\{i \in \{k\} \mid x \in Q_i\} \end{array}$ Definition of $\#\{i \in \{1, \dots, k\} \mid x \in Q_i\}$

(20) $\left| \begin{array}{c} | \\ | \\ | \\ | \end{array} \right| \quad \text{Let } R = \{A_k, A_{k+1}\}.$

(21) $\left| \begin{array}{c} | \\ | \\ | \\ | \end{array} \right| \quad R \in T_2$ Definition of T_2 and R , (20)

(22) $\left| \begin{array}{c} | \\ | \\ | \\ | \end{array} \right| \quad \boxed{P(2)}$ Specialisation, (2)

(23) $\left| \begin{array}{c} | \\ | \\ | \\ | \end{array} \right| \quad \begin{array}{l} \forall w \in U. (w \in A_k \Delta A_{k+1} \text{ IFF} \\ \#\{i \in \{k, k+1\} \mid w \in A_i\} \text{ is odd}) \end{array}$ Specialisation for R , (20), (22)

(24) $\left| \begin{array}{c} | \\ | \\ | \\ | \end{array} \right| \quad \begin{array}{l} (x \in A_k \Delta A_{k+1} \text{ IFF} \\ \#\{i \in \{k, k+1\} \mid x \in A_i\} \text{ is odd}) \end{array}$ Specialisation for x , (23)

(25) $\left| \begin{array}{c} | \\ | \\ | \\ | \end{array} \right| \quad \begin{array}{l} (x \in Q_k \text{ IFF} \\ \#\{i \in \{k, k+1\} \mid x \in A_i\} \text{ is odd}) \end{array}$ Substitution, (9), (24)

(26) $\boxed{P(1)}$ Specialisation, (2)

(27) $Q_k \in T_1$ Definition of T_1 and Q_k , (12)

(28) $(x \in Q_k \text{ IFF } \#\{i \in \{k\} \mid x \in Q_i\} \text{ is odd})$ Specialisation for Q_k and x , (26)

(29) Let A, B, C be arbitrary propositional variables.

(30) $((A \text{ IFF } B) \text{ AND } (A \text{ IFF } C)) \text{ IMPLIES } (B \text{ IFF } C)$ Tautology

(31) $\#\{i \in \{k\} \mid x \in Q_i\} \text{ is odd IFF } \#\{i \in \{k, k+1\} \mid x \in A_i\} \text{ is odd}$ Use of Tautology (Modus Ponens), (25), (28), (30)

(32) $\forall A \in \mathbb{N} \cup \{0\}. \forall B \in \mathbb{N} \cup \{0\}. \forall C \in \mathbb{N} \cup \{0\}. [(A \text{ is odd IFF } B \text{ is odd}) \text{ IFF } (A + C \text{ is odd IFF } B + C \text{ is odd})]$ Tautology: Parity Equivalence is Conserved Under Addition

(33) $(\#\{i \in \{k\} \mid x \in Q_i\} \text{ is odd IFF } \#\{i \in \{k, k+1\} \mid x \in A_i\} \text{ is odd}) \text{ IFF } (\#\{i \in \{k\} \mid x \in Q_i\} + \#\{i \in \{1, \dots, k-1\} \mid x \in Q_i\} \text{ is odd IFF } \#\{i \in \{k, k+1\} \mid x \in A_i\} + \#\{i \in \{1, \dots, k-1\} \mid x \in Q_i\} \text{ is odd})$ Specialisation, (32)

(34) $\#\{i \in \{k\} \mid x \in Q_i\} + \#\{i \in \{1, \dots, k-1\} \mid x \in Q_i\} \text{ is odd IFF } \#\{i \in \{k, k+1\} \mid x \in A_i\} + \#\{i \in \{1, \dots, k-1\} \mid x \in Q_i\} \text{ is odd}$ Modus Ponens, (31), (35)

(35) $\#\{i \in \{1, \dots, k\} \mid x \in Q_i\} \text{ is odd IFF } \#\{i \in \{k, k+1\} \mid x \in A_i\} + \#\{i \in \{1, \dots, k-1\} \mid x \in A_i\} \text{ is odd}$ Substitution, (19)

$$(36) \quad \begin{array}{|l} \\ \\ \\ \\ \\ \end{array} \quad \begin{array}{l} \# \{i \in \{k, k+1\} \mid x \in A_i\} \\ + \# \{i \in \{1, \dots, k-1\} \mid x \in A_i\} \\ = \# \{i \in \{1, \dots, k+1\} \mid x \in A_i\} \end{array} \quad \begin{array}{l} \text{Definition of} \\ \# \{i \in \{1, \dots, k+1\} \mid x \in A_i\} \end{array}$$

$$(37) \quad \begin{array}{|l} \\ \\ \\ \\ \\ \end{array} \quad \begin{array}{l} \# \{i \in \{1, \dots, k\} \mid x \in Q_i\} \text{ is odd IFF} \\ \# \{i \in \{1, \dots, k+1\} \mid x \in A_i\} \text{ is odd} \end{array} \quad \text{Substitution, (35)}$$

$$(38) \quad \begin{array}{|l} \\ \\ \\ \\ \\ \end{array} \quad \begin{array}{l} x \in A_1 \Delta \dots \Delta A_k \Delta A_{k+1} \text{ IFF} \\ \# \{i \in \{1, \dots, k+1\} \mid x \in A_i\} \text{ is odd} \end{array} \quad \begin{array}{l} \text{Use of Tautology (Modus Ponens),} \\ (18), (30), (37) \end{array}$$

$$(39) \quad \begin{array}{|l} \\ \end{array} \quad \begin{array}{l} \forall x \in U. (x \in A_1 \Delta \dots \Delta A_k \Delta A_{k+1} \text{ IFF} \\ \# \{i \in \{1, \dots, k+1\} \mid x \in A_i\} \text{ is odd}) \end{array} \quad \text{Generalisation, (6)}$$

$$(40) \quad \begin{array}{|l} \end{array} \quad \begin{array}{l} \forall S \in T_k. \forall x \in U. (x \in A_1 \Delta \dots \Delta A_k \Delta A_{k+1} \text{ IFF} \\ \# \{i \in \{1, \dots, k+1\} \mid x \in A_i\} \text{ is odd}) \end{array} \quad \text{Generalisation, (5)}$$

$$(41) \quad \begin{array}{|l} \end{array} \quad P(k+1) \quad \text{Direct Proof; (2)-(40)}$$

$$(42) \quad \forall n \in \mathbb{N}. P(n) \quad \text{Induction; (1)-(41)}$$