

DELETION IN BST

```
def delete(self, item)
```

```
    """ REMOVE ONE OCCURRENCE OF AN ITEM FROM THE TREE """
```

```
    → IF self.is_empty:
```

```
        → pass
```

```
    → elif self.root == item:
```

```
        → self.delete_root()
```

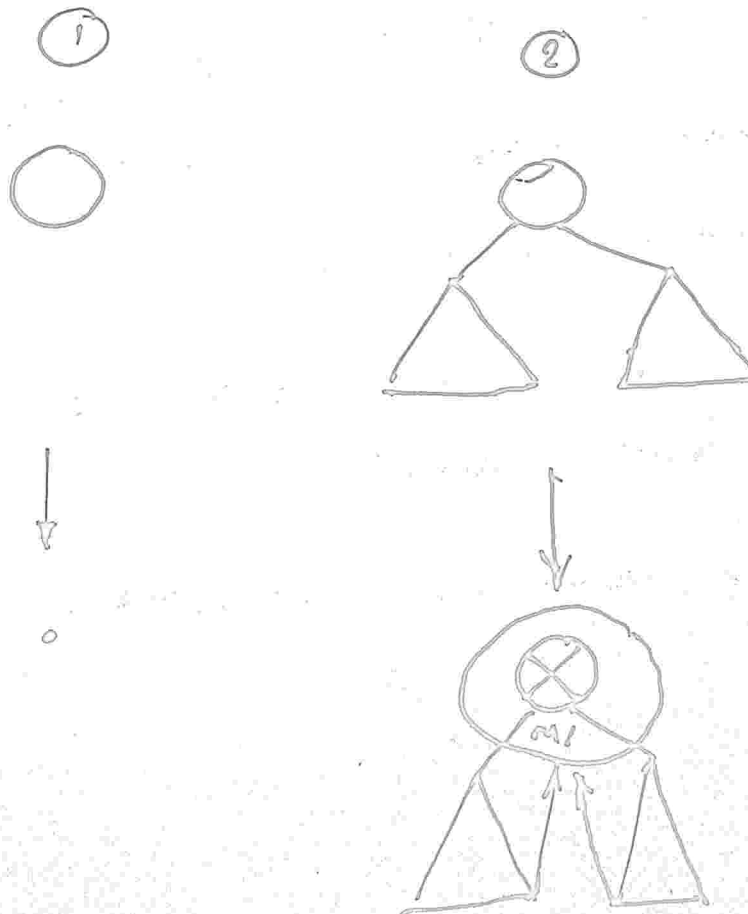
```
    → elif item < self.root:
```

```
        → self._left.delete(item)
```

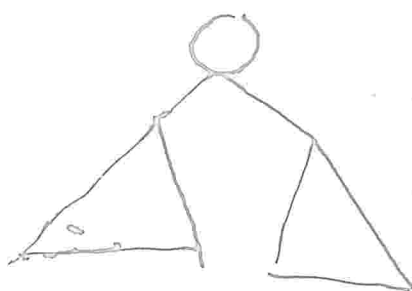
```
    → else:
```

```
        → self._right.delete(item)
```

```
def _delete_root(self):
```



BST EFFICIENCY



• EACH RECURSIVE CALL GOES DOWN BY ONE, WHICH MEANS THAT WHAT MATTERS IN DETERMINING COMPLEXITY IS THE HEIGHT OF THE TREE.

THE LAST RECURSIVE CALL IS ON THE LAST NODE IN THE LIST.

OF RECURSIVE CALLS = TREE HEIGHT (H)

27 THE RUNNING TIME OF DELETE IS PROPORTIONAL TO $O(H)$.

HOWEVER, THIS IS NOT A FULL STORY.
WHAT ABOUT SEARCH?

• H IS THE WORST CASE SCENARIO.

