

**Problem.** Prove  $q(n)$  is true for all even natural numbers.

*Solution.* Let  $p(k) = q(2k)$ .

$\forall k \in \mathbb{N}. p(k)$

means the same as

$\forall k \in \mathbb{N}. q(2k)$ , which is the same as

$\forall n \in \mathbb{N}. (n \text{ is even} \implies q(n))$ .

Base Case:

$p(0) = q(0)$

Induction Step:

$p(k) \implies p(k+1)$ ,

which is the same as

$q(2k) \implies q(2k+2)$ .

It is sufficient to prove

$q(0) \text{ AND } \forall n \in \mathbb{N} (q(n) \implies q(n+2))$  .

□

### Theorem 0.1

For all  $n \in \mathbb{Z}^+$  and all  $a_1, \dots, a_n \in \mathbb{R}^+$ ,

$$\left(\prod_{i=1}^n a_i\right)^{\frac{1}{n}} \leq \frac{\sum_{i=1}^n a_i}{n}$$

*Proof.* We prove  $\forall n \in \mathbb{Z}^+. P(n)$ .

Base Case:

$n = 2$

Let  $a_1, a_2 \in \mathbb{R}^+$  be arbitrary.

Then  $0 \leq (a_1 - a_2)^2 = a_1^2 - 2a_1a_2 + a_2^2$ .

Hence,  $a_1^2 + a_2^2 \geq 2a_1a_2$ .

Thus,

$$\left(\frac{a_1 + a_2}{2}\right)^2 = \frac{a_1^2 + a_2^2 + 2a_1a_2}{4} \geq a_1a_2$$

Hence,  $P(2)$  is true by generalisation.

Induction Step:

Let  $n \in \mathbb{Z}^+$  be arbitrary and suppose  $n \geq 2$ .

Assume  $P(n)$ .

Let  $a_1, \dots, a_{n-1} \in \mathbb{R}^+$  be arbitrary.

Let  $b_i = a_i$  for  $i = 1, \dots, n-1$ .

Let  $b_n = \frac{a_1 + a_2 + \dots + a_{n-1}}{n-1}$ .

By specialisation of  $p(n)$ ,

$$b_1 \cdots b_{n-1} b_n \leq \left( \frac{b_1 + \cdots + b_n}{n} \right)^n = \left( \frac{b_1 + \cdots + b_n}{n} \right)^n \quad (1)$$

$$= \left( \frac{a_1 + \cdots + a_{n-1} + b_n}{n} \right)^n \quad (2)$$

$$= \left( \frac{(n-1)b_n + b_n}{n} \right)^n \quad (3)$$

$$= b_n^n \quad (4)$$

Therefore,  $b_1 b_2 \cdots b_{n-1} \leq b_n^{n-1}$ .

Hence,  $P(n-1)$  is true by generalisation.

Let  $a_1, \dots, a_n \in \mathbb{R}^+$  be arbitrary.

Let  $b_1 = \frac{a_1 + \cdots + a_n}{n}$  and  $b_2 = \frac{a_{n+1} + \cdots + a_{2n}}{n}$ .

By specialisation of  $P(n)$ ,

$$\prod_{i=1}^n a_i \leq \left( \frac{1}{n} \sum_{i=1}^n a_i \right)^n$$

and

$$\prod_{i=n+1}^{2n} a_i \leq \left( \frac{1}{n} \sum_{i=n+1}^{2n} a_i \right)^n$$

and by specialisation of  $P(2)$ ,

$$b_1 b_2 \leq \left( \frac{b_1 + b_2}{2} \right)^2$$

Hence

$$\prod_{i=1}^{2n} a_i \leq \left( \frac{\sum_{i=1}^n a_i}{n} \right) \left( \frac{\sum_{i=n+1}^{2n} a_i}{n} \right)^n = (b_1 b_2)^n \leq \left( \frac{b_1 + b_2}{2} \right)^{2n}.$$

Note that  $\left( \frac{b_1 + b_2}{2} \right)^{2n} = \left( \frac{1}{2n} \sum_{i=1}^{2n} a_i \right)^{2n}$ .

By generalisation,  $P(2n)$  is true.

$\forall n \in \mathbb{N}[(n \geq 2 \text{ AND } P(n)) \text{ IMPLIES } P(2n)]$ .

Therefore, by induction,

$\forall n \in \mathbb{Z}^+. P(n)$

□

## 0.1 Induction in Finite Sets

**Problem.** Prove  $\forall i \in \{0, \dots, n\}. P(i)$ .

*Solution.* Base Case:

$p(0)$

Induction Step:

Let  $i \in \{0, \dots, n-1\}$  be arbitrary.

Assume  $p(i)$ .

$\vdots$

$p(i+1)$ .

$\forall i \in \{0, \dots, n-1\}. [p(i) \text{ IMPLIES } p(i+1)]$

$\forall i \in \{0, \dots, n\} p(i)$  by induction

□

## 0.2 Strong Induction

To prove  $\forall i \in \mathbb{N}. p(i)$  it suffices to prove that

$$\forall i \in \mathbb{N}. \forall j \in \mathbb{N}. [(j < i) \text{ IMPLIES } p(j)] \text{ IMPLIES } P(i)$$

The only difference of the strong induction from the weak induction is  $p(0), \dots, p(i-1)$ .

A template proof follows.

*Proof.* Let  $i \in \mathbb{N}$  be arbitrary.

Assume  $\forall j \in \mathbb{N}. (j < i \text{ IMPLIES } p(j))$ .

... various cases, including the base case ...

$p(i)$

$\forall i \in \mathbb{N} [\forall j \in \mathbb{N}. (j < i) \text{ IMPLIES } p(j)] \text{ IMPLIES } p(i)$  by direct proof and generalization.

$\forall i \in \mathbb{N}. p(i)$  by strong induction

□

### Theorem 0.2

For all  $n \geq 4$ , exactly a sum of  $n$  can be exchanged in coins with nomination 2 and 5\$ bills.

*Proof.* Let  $p(n) = \exists f \in \mathbb{N}. \exists g \in \mathbb{N}. (n = 2f + 5g)$  for all  $n \in \mathbb{N}$ .

Let  $n \in \mathbb{N}$  be arbitrary.

Suppose  $n \geq 4$  and  $\forall j \in \mathbb{N}. (4 \leq j < n \text{ IMPLIES } p(j))$ .

If  $n = 4$ , then  $n = 2 \cdot 2 + 0 \cdot 5$ .

If  $n = 5$ , then  $n = 0 \cdot 2 + 1 \cdot 5$ .

If  $n \leq 6$ , then  $4 \leq n-2 < n$ . Then  $P(n-2)$  is true by specialisation.

□