



**VIT**<sup>®</sup>  
**Vellore Institute of Technology**  
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Subject Code: BCSE204P

Course Title: Design and Analysis of Algorithms  
Lab

Lab Slot: L39 + L40

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Lab Assessment 3

1. Design a solution to see if a content  $C = \text{PGGA}$  is plagiarized in Text  $T = \text{SAQSPAPGPGGAS}$  using the following Algorithms.
  - a. KMP Algorithm
  - b. Rabin-Karp Algorithm

a. KMP Algorithm

Algorithm:

KMP\_String\_Matcher( $T, P$ ):

$n = \text{length}(T)$

$m = \text{length}(P)$

$\pi = \text{Compute\_LPS}(P)$

$i = 0$

$j = 0$

while  $i < n$ :

  if  $T[i] == P[j]$ :

$i = i + 1$

$j = j + 1$

  if  $j == m$ :

    print("Pattern occurs at index",  $i - j$ )

$j = \pi[j - 1]$

  else:

    if  $j != 0$ :

$j = \pi[j - 1]$

    else:

$i = i + 1$

Compute\_LPS( $P$ ):

$m = \text{length}(P)$

LPS = array of size  $m$

LPS[0] = 0

len = 0

$i = 1$

while  $i < m$ :

  if  $P[i] == P[\text{len}]$ :

    len = len + 1

    LPS[i] = len

$i = i + 1$

```
    else:
        if len != 0:
            len = LPS[len - 1]
        else:
            LPS[i] = 0
            i = i + 1

return LPS
```

Source Code:

```
#include <iostream>
#include <string>
#include <vector>
using namespace std;

void constructLps(string &pat, vector<int> &lps) {

    int len = 0;

    lps[0] = 0;

    int i = 1;
    while (i < pat.length()) {

        if (pat[i] == pat[len]) {
            len++;
            lps[i] = len;
            i++;
        }

        else {
            if (len != 0) {

                len = lps[len - 1];
            }
            else {
                lps[i] = 0;
            }
        }
    }
}
```

```

        i++;
    }
}
}
}

vector<int> search(string &pat, string &txt) {
    int n = txt.length();
    int m = pat.length();

    vector<int> lps(m);
    vector<int> res;

    constructLps(pat, lps);

    int i = 0;
    int j = 0;

    while (i < n) {

        if (txt[i] == pat[j]) {
            i++;
            j++;

            if (j == m) {
                res.push_back(i - j);

                j = lps[j - 1];
            }
        }

        else {

            if (j != 0)
                j = lps[j - 1];
            else
                i++;
        }
    }
}

```

```

    }
    return res;
}

int main() {
    string txt = "SAQSPAPGPGGAS";
    string pat = "PGGA";

    vector<int> res = search(pat, txt);
    for (int i = 0; i < res.size(); i++)
        cout << res[i] << " ";

    return 0;
}

```

Input:

```

string txt = "SAQSPAPGPGGAS";
string pat = "PGGA";

```

Output:

```

Analysis\ of\ Algorithms\ Iyappan/Lab/Assessment3/KMP ; exit;
8

```

= 8 is the index

Time Complexity Analysis:

From the <RS table

$$T(m) = T(m-1) + O(1)$$

$$a=1, b=1, k=0$$

$$\log_b a = 1, k=0$$

$$\log_b a > k$$

$$O(m^{\log_b a})$$

$$O(m^1)$$

$$= O(m)$$

From pattern matching

$$T(n) = T(n-1) + O(1)$$

$$= O(n)$$

$$T(m+n) = O(m) + O(n) = O(m+n)$$
$$= O(m+n)$$

## b. Rabin-Karp Algorithm

Algorithm:

Rabin\_Karp\_Matcher(T, P, d, q):

n = length(T)

m = length(P)

h =  $d^{(m-1)} \bmod q$

p = 0

t = 0

for i = 0 to m-1:

p =  $(d * p + P[i]) \bmod q$

$t = (d * t + T[i]) \bmod q$

for  $s = 0$  to  $n-m$ :

if  $p == t$ :

if  $P[0..m-1] == T[s..s+m-1]$ :

print "Pattern occurs at index",  $s$

if  $s < n-m$ :

$t = (d * (t - T[s] * h) + T[s + m]) \bmod q$

if  $t < 0$ :

$t = t + q$

Source Code:

```
#include <iostream>
#include <string>
using namespace std;

void search(string pat, string txt, int q)
{
    int M = pat.size();
    int N = txt.size();
    int i, j;
    int p = 0;
    int t = 0;
    int h = 1;
    int d = 256;
    for (i = 0; i < M - 1; i++)
        h = (h * d) % q;

    for (i = 0; i < M; i++) {
        p = (d * p + pat[i]) % q;
        t = (d * t + txt[i]) % q;
    }

    for (i = 0; i <= N - M; i++) {
        if (p == t) {
            for (j = 0; j < M; j++) {
                if (txt[i + j] != pat[j]) {
                    break;
                }
            }
        }
    }
}
```

```

    }

    if (j == M)
        cout << "Pattern found at index " << i
            << endl;
    }
    if (i < N - M) {
        t = (d * (t - txt[i] * h) + txt[i + M]) % q;
        if (t < 0)
            t = (t + q);
    }
}

int main()
{
    string txt = "SAQSPAPGPGGAS";
    string pat = "PGGA";

    int q = INT_MAX;
    search(pat, txt, q);
    return 0;
}

```

Input:

```

string txt = "SAQSPAPGPGGAS";
string pat = "PGGA";

```

Output:

```

\ and\ Analysis\ of\ Algorithms\
Pattern found at index 8

```

At index 8 again

Time Complexity Analysis:



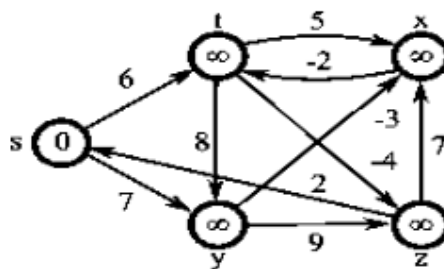
Rolling hash function updates in  $O(1)$  in each shift.

so  
Average cost =  $\Theta(n - m + 1) = \Theta(n)$

If all hash values collide, a full character comparison is done at each position.

$$\Theta(m(n - m + 1)) \\ = \Theta(mn)$$

2. Consider the given directed graph  $G = \{V, E\}$ , Determine the shortest path from Source vertex to all the remaining vertices of the Graph by using appropriate algorithm and analyze its time complexity.



Algorithm:

createGraph(V, E):

    Initialize Graph with V vertices and E edges

    Return Graph

displayDistances(distances):

    Print "Vertex   Distance from Source"

    For each (vertex, distance) in distances:

        Print vertex, distance

bellmanFord(graph, source):

    Initialize distances for all vertices as  $\infty$ , except source = 0

    Repeat (V-1) times:

        For each edge (u, v, w) in graph:

            If  $\text{distances}[u] + w < \text{distances}[v]$ :

                Update  $\text{distances}[v] = \text{distances}[u] + w$

    For each edge (u, v, w) in graph:

        If  $\text{distances}[u] + w < \text{distances}[v]$ :

            Print "Negative cycle detected"

    Return

displayDistances(distances)

Source Code:

```
#include <iostream>
#include <cstdlib>
#include <climits>
#include <map>
using namespace std;

struct Edge {
    char startVertex, endVertex;
    int edgeWeight;
};

struct Graph {
    int numVertices, numEdges;
    struct Edge* edges;
};

struct Graph* initializeGraph(int vertices, int edges) {
```

```

    struct Graph* graph = (struct Graph*)malloc(sizeof(struct Graph));
    graph->numVertices = vertices;
    graph->numEdges = edges;
    graph->edges = (struct Edge*)malloc(graph->numEdges * sizeof(struct Edge));
    return graph;
}

void displayFinalDistances(map<char, int>& distanceMap) {
    cout << "\nVertex\tMinimum Distance from Source" << endl;
    for (auto& entry : distanceMap) {
        cout << entry.first << "\t\t" << entry.second << endl;
    }
}

void bellmanFord(struct Graph* graph, char sourceVertex) {
    int vertices = graph->numVertices;
    int edges = graph->numEdges;
    map<char, int> shortestDistances;

    for (int i = 0; i < edges; i++) {
        shortestDistances[graph->edges[i].startVertex] = INT_MAX;
        shortestDistances[graph->edges[i].endVertex] = INT_MAX;
    }

    shortestDistances[sourceVertex] = 0;

    for (int i = 1; i <= vertices - 1; i++) {
        for (int j = 0; j < edges; j++) {
            char u = graph->edges[j].startVertex;
            char v = graph->edges[j].endVertex;
            int weight = graph->edges[j].edgeWeight;

            if (shortestDistances[u] != INT_MAX && shortestDistances[u] + weight < shortestDistances[v]) {
                shortestDistances[v] = shortestDistances[u] + weight;
            }
        }
    }
}

```

```

for (int i = 0; i < edges; i++) {
    char u = graph->edges[i].startVertex;
    char v = graph->edges[i].endVertex;
    int weight = graph->edges[i].edgeWeight;

    if (shortestDistances[u] != INT_MAX && shortestDistances[u] + weight < shortestDistances[v]) {
        cout << "\nWarning: The graph contains a negative weight cycle!" << endl;
        return;
    }
}

displayFinalDistances(shortestDistances);
}

int main() {
    int vertices, edges;
    char source;

    cout << "Enter the number of vertices: ";
    cin >> vertices;
    cout << "Enter the number of edges: ";
    cin >> edges;
    cout << "Enter the source vertex: ";
    cin >> source;

    struct Graph* graph = initializeGraph(vertices, edges);

    for (int i = 0; i < edges; i++) {
        cout << "\nEnter properties for edge " << i + 1 << " (Start Vertex, End Vertex, Weight): ";
        cin >> graph->edges[i].startVertex >> graph->edges[i].endVertex >> graph->edges[i].edgeWeight;
    }

    bellmanFord(graph, source);

    return 0;
}

```

Input:

Enter the number of vertices: 5

Enter the number of edges: 10

Enter the source vertex: s

Enter properties for edge 1 (Start Vertex, End Vertex, Weight): s t 6

Enter properties for edge 2 (Start Vertex, End Vertex, Weight): s y 7

Enter properties for edge 3 (Start Vertex, End Vertex, Weight): t x 5

Enter properties for edge 4 (Start Vertex, End Vertex, Weight): t y 8

Enter properties for edge 5 (Start Vertex, End Vertex, Weight): t z -4

Enter properties for edge 6 (Start Vertex, End Vertex, Weight): y x -3

Enter properties for edge 7 (Start Vertex, End Vertex, Weight): y z 9

Enter properties for edge 8 (Start Vertex, End Vertex, Weight): z x 7

Enter properties for edge 9 (Start Vertex, End Vertex, Weight): z s 2

Enter properties for edge 10 (Start Vertex, End Vertex, Weight): x t -2

Output:

```
Vertex  Minimum Distance from Source
s        0
t        2
x        4
y        7
z       -2

Saving session...
...copying shared history...
...saving history...truncating history files...
...completed.

[Process completed]
```

Time Complexity Analysis:

$$T(V, E) = T(V-1, E) + O(E)$$

where  $V$  = vertices &  $E$  = edges

if  $T(V, E) = T(V-1, E) + O(E)$  then,

$$T(V, E) = T(V-2, E) + O(E) + O(E) \text{ and}$$
$$T(V, E) = T(V-3, E) + O(E) + O(E) + O(E)$$

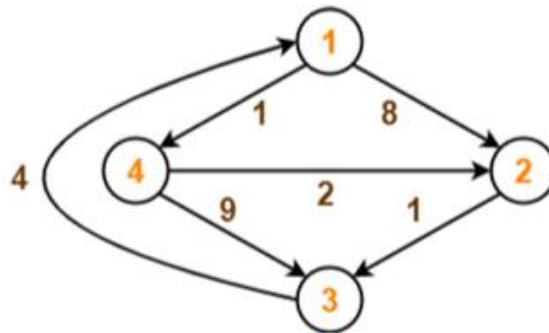
After  $V-1$  expansion

$$T(V, E) = T(1, E) + (V-1) O(E)$$

$$T(V, E) = O(VE)$$

$$O(VE)$$

3. Determine step-by-step process of finding the shortest path distance between every pair of vertices using Floyd warshall algorithm, from the directed weighted graph.



Algorithm

FloydWarshall(graph, vertices):  
    numVertices = size of vertices

```
for k from 0 to numVertices - 1:
    for i from 0 to numVertices - 1:
        for j from 0 to numVertices - 1:
            start = vertices[i]
            end = vertices[j]
            intermediate = vertices[k]
```

```
            if graph[start][intermediate] is not INF and graph[intermediate][end] is not INF:
                if graph[start][end] is INF or graph[start][end] > graph[start][intermediate] +
graph[intermediate][end]:
                    graph[start][end] = graph[start][intermediate] + graph[intermediate][end]
```

Source Code

```
#include <iostream>
#include <vector>
#include <map>
#include <limits>
```

```

using namespace std;

void floydWarshall(map<char, map<char, int> > &graph, vector<char> &vertices) {
    int numVertices = vertices.size();

    for (int k = 0; k < numVertices; k++) {
        for (int i = 0; i < numVertices; i++) {
            for (int j = 0; j < numVertices; j++) {
                char start = vertices[i], end = vertices[j], intermediate = vertices[k];

                if (graph[start][intermediate] != INT_MAX && graph[intermediate][end] != INT_MAX &&
                    (graph[start][end] == INT_MAX || graph[start][end] > graph[start][intermediate] +
graph[intermediate][end])) {
                    graph[start][end] = graph[start][intermediate] + graph[intermediate][end];
                }
            }
        }
    }
}

int main() {
    int numVertices, numEdges;

    cout << "Enter the number of vertices: ";
    cin >> numVertices;

    cout << "Enter the number of edges: ";
    cin >> numEdges;

    vector<char> vertices(numVertices);
    map<char, map<char, int> > graph;

    cout << "Enter the vertex labels: ";
    for (int i = 0; i < numVertices; i++) {
        cin >> vertices[i];

        for (int j = 0; j < numVertices; j++) {
            graph[vertices[i]][vertices[j]] = (i == j) ? 0 : INT_MAX;
        }
    }
}

```



```

}

cout << "Enter edges in format (source destination weight):" << endl;
for (int i = 0; i < numEdges; i++) {
    char source, destination;
    int weight;
    cin >> source >> destination >> weight;
    graph[source][destination] = weight;
}

floydWarshall(graph, vertices);

cout << "\nAll-Pairs Shortest Paths:\n ";
for (char vertex : vertices) {
    cout << vertex << "\t";
}
cout << "\n-----\n";

for (char start : vertices) {
    cout << start << " | ";
    for (char end : vertices) {
        if (graph[start][end] == INT_MAX)
            cout << "INF\t";
        else
            cout << graph[start][end] << "\t";
    }
    cout << endl;
}

return 0;
}

```

#### Input

```
orithms\ Iyappan/Lab/Assessment3/FM ; exit;  
Enter the number of vertices: 4  
Enter the number of edges: 6  
Enter the vertex labels: 1  
2  
3  
4  
Enter edges in format (source destination weight):  
1 4 1  
1 2 8  
4 2 2  
4 3 9  
2 3 1  
3 1 4
```

#### Output

All-Pairs Shortest Paths:

|   | 1 | 2 | 3 | 4 |
|---|---|---|---|---|
| 1 | 0 | 2 | 0 | 0 |
| 2 | 5 | 0 | 1 | 0 |
| 3 | 4 | 2 | 0 | 0 |
| 4 | 7 | 2 | 3 | 0 |

#### Time Complexity Analysis

$$T(N) = N \times T(N-1) + O(1)$$

Since each iteration needs  $O(N^3)$  for a vertex

$$T(N) = N^3 + (N-1)^3 + (N-2)^3 + \dots + (N-(N-1))^3$$

$$= \left( \frac{N \cdot (N+1)}{2} \right)^2$$

$$= \frac{N^2 (N+1)^2}{4}$$

$$= O\left(\frac{N^4}{4}\right) \text{ which is simplified to}$$

$$O(N^4) //$$