CS738: Advanced Compiler Optimizations

Interprocedural Data Flow Analysis Functional Approach

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Interprocedurally Valid Paths

- G^* ignores the special nature of call and return edges
- Not all paths in G^* are feasible
  - do not represent potentially valid execution paths
- IVP(r_1, n): set of all interprocedurally valid paths from r_1 to n
- Path $q \in \text{path}_{G^*}(r_1, n)$ is in IVP(r_1, n)
  - iff sequence of all E^1 edges in q (denoted q_1)is *proper*

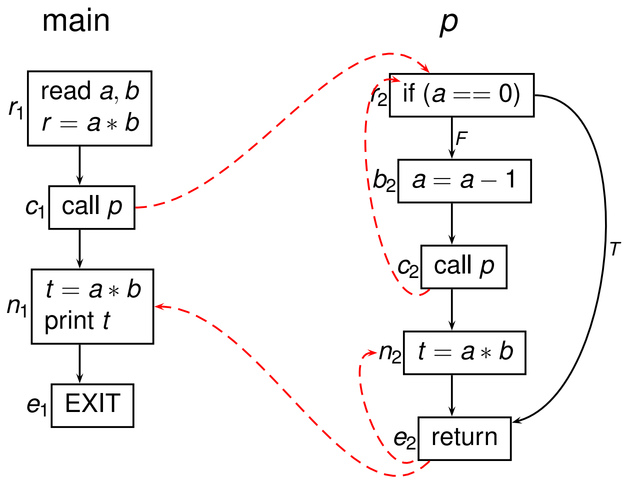
Proper sequence

- q_1 without any return edge is proper
- let $q_1[i]$ be the first return edge in q_1. q_1 is proper if
  - $i > 1$; and
  - $q_1[i-1]$ is call edge corresponding to $q_1[i]$; and
  - q_1' obtained from deleting $q_1[i-1]$ and $q_1[i]$ from q_1 is proper

Interprocedurally Valid Complete Paths

- IVP$_0(r_p, n)$ for procedure p and node $n \in N_p$
- set of all interprocedurally valid paths q in G^* from r_p to n s.t.
  - Each call edge has corresponding return edge in q restricted to E^1

IVPs



$$r_1 \to c_1 \to r_2 \to c_2 \to r_2 \to e_2 \to n_2 \to e_2 \to n_1 \to e_1$$
$$\in \text{IVP}(r_1, e_1)$$
$$r_1 \to c_1 \to r_2 \to c_2 \to r_2 \to e_2 \to n_1 \to e_1 \notin \text{IVP}(r_1, e_1)$$
$$r_2 \to c_2 \to r_2 \to e_2 \to n_2 \in \text{IVP}_0(r_2, n_2)$$
$$r_2 \to c_2 \to r_2 \to c_2 \to e_2 \to n_2 \notin \text{IVP}_0(r_2, n_2)$$

Path Decomposition

$$\begin{array}{rcl} q & \in & \text{IVP}(r_{\text{main}}, n) \\ & \Leftrightarrow & \\ q & = & q_1 \parallel (c_1, r_{p_2}) \parallel q_2 \parallel \cdots \parallel (c_{j-1}, r_{p_j}) \parallel q_j \\ & & \text{where for each } i < j, q_i \in \text{IVP}_0(r_{p_i}, c_i) and q_j \in \text{IVP}_0(r_{p_j}, n) \end{array}$$

Functional Approach

- (L, F): a *distributive* data flow framework
- Procedure p, node $n \in N_p$
- $\phi_{(r_p,n)} \in F$ describes flow of data flow information from start of r_p to start of n
  - along paths in $\text{IVP}_0(r_p, n)$

Functional Approach Constraints

$$\begin{array}{rcl} \phi_{(r_p,r_p)} & \leq & id_L \\ \phi_{(r_p,n)} & = & \displaystyle\bigwedge_{(m,n)\in E_p} (h_{(m,n)} \circ \phi_{(r_p,m)}) \qquad \text{for } n \in N_p \end{array}$$

where

$$h_{(m,n)} = \begin{cases} f_{(m,n)} & \text{if } (m,n) \in E_p^0, \\ & \quad f_{(m,n)} \in F \text{ associated flow function} \\ \phi_{(r_q,e_q)} & \text{if } (m,n) \in E_p^1 \text{ and } m \text{ calls procedure } q \end{cases}$$

Information x at r_p translated to information $\phi_{(r_p,n)}(x)$ at n

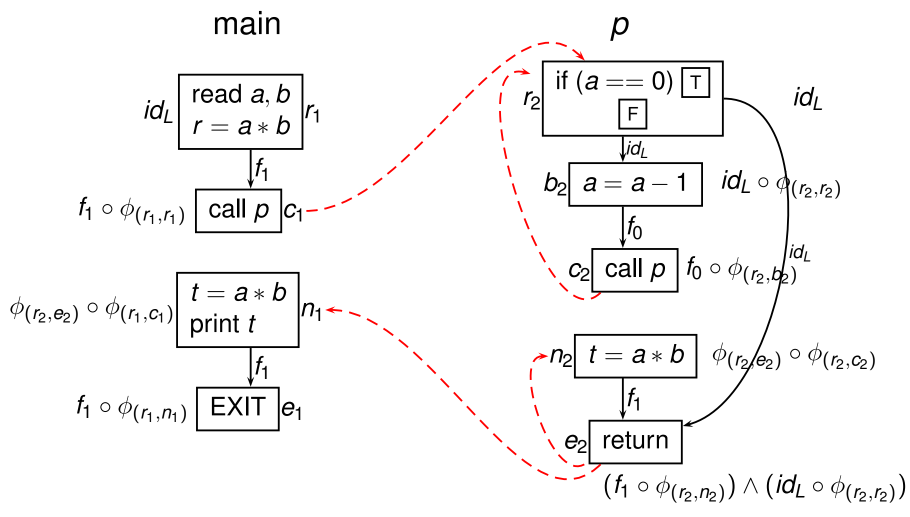
Solving ϕ Constraints

- Round-robin iterative approximations to initial values

$$
\begin{aligned}
\phi^0_{(r_p,r_p)} &\leq id_L \\
\phi^0_{(r_p,n)} &\leq f_\Omega \qquad n \in N_p - \{r_p\}
\end{aligned}
$$

- Reach maximal fixed point

Example



Iterative Solution

Function	Constraint	Iteration # Init	1^{st}	2^{nd}	3^{rd}
$\phi_{(r_1,r_1)}$	id_L	id	id	id	id
$\phi_{(r_1,c_1)}$	$f_1 \circ \phi_{(r_1,r_1)}$	f_Ω	f_1	f_1	f_1
$\phi_{(r_1,n_1)}$	$\phi_{(r_2,e_2)} \circ \phi_{(r_1,c_1)}$	f_Ω	f_Ω	f_1	f_1
$\phi_{(r_1,e_1)}$	$f_1 \circ \phi_{(r_1,n_1)}$	f_Ω	f_1	f_1	f_1
$\phi_{(r_2,r_2)}$	id_L	id	id	id	id
$\phi_{(r_2,b_2)}$	$id_L \circ \phi_{(r_2,r_2)}$	f_Ω	id	id	id
$\phi_{(r_2,c_2)}$	$f_0 \circ \phi_{(r_2,b_2)}$	f_Ω	f_0	f_0	f_0
$\phi_{(r_2,n_2)}$	$\phi_{(r_2,e_2)} \circ \phi_{(r_2,c_2)}$	f_Ω	f_Ω	f_0	f_0
$\phi_{(r_2,e_2)}$	$(f_1 \circ \phi_{(r_2,n_2)}) \wedge (id_L \circ \phi_{(r_2,r_2)})$	f_Ω	id	id	id

Solving Data Flow Problem

- The above process gives solution to ϕ functions
- Use it to compute data flow information x_n associated with start of block n

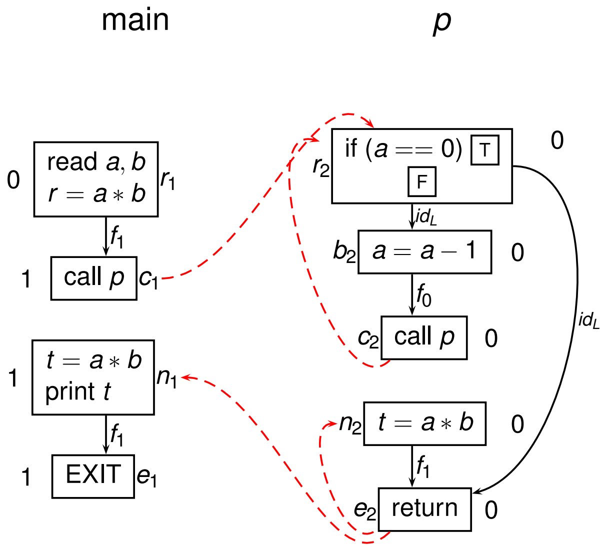
$$x_{r_{\text{main}}} = BoundaryInfo$$

for each procedure p

$$
\begin{aligned}
x_{r_p} &= \bigwedge \left\{ \phi_{(r_q,c)}(x_{r_q}) : \begin{array}{l} q \text{ is a procedure and} \\ c \text{ is a call to } p \text{ in } q \end{array} \right\} \\
x_n &= \phi_{(r_p,n)} \qquad n \in N_p - \{r_p\}
\end{aligned}
$$

- Iterative algorithm for solution, maximal fixed point solution

Example



Functions
$\phi_{(r_1,r_1)} = id_L$
$\phi_{(r_1,c_1)} = f_1$
$\phi_{(r_1,n_1)} = f_1$
$\phi_{(r_1,e_1)} = f_1$
$\phi_{(r_2,r_2)} = id_L$
$\phi_{(r_2,b_2)} = id_L$
$\phi_{(r_2,c_2)} = f_0$
$\phi_{(r_2,n_2)} = f_0$
$\phi_{(r_2,e_2)} = id_L$

Interprocedural MOP

$$\begin{aligned}\Psi_n &= \bigwedge\{f_q : q \in \mathrm{IVP}(r_{\mathrm{main}}, n)\} \in F \qquad \forall n \in N^* \\ y_n &= \Psi_n(BoundaryInfo) \qquad \forall n \in N^*\end{aligned}$$

y_n is the *meet-over-all-paths solution* (MOP).

IVP_0 Lemma

$$\phi_{(r_p,n)} = \bigwedge\{f_q : q \in \mathrm{IVP}_0(r_p, n)\} \qquad \forall n \in N_p$$

Proof: By induction (**Exercise/Reading Assignment**)

MOP

$$\begin{aligned}\Psi_n &= \bigwedge\{f_q : q \in \mathrm{IVP}(r_{\mathrm{main}}, n)\} \in F \qquad \forall n \in N^* \\ \mathcal{X}_n &= \wedge\{\phi_{(r_{p_j},n)} \circ \phi_{(r_{p_{j-1}},c_{j-1})} \circ \ldots \circ \phi_{(r_{p_1},c_1)} \mid c_i \text{ calls } p_{i+1}\}\end{aligned}$$

Then

$$\Psi_n = \mathcal{X}_n$$

Proof: IVP_0 Lemma and Path decomposition

$$y_n = \Psi_n(BoundaryInfo) = \mathcal{X}_n(BoundaryInfo)$$

MOP vs MFP

- F is distributive $\Rightarrow$ $MFP = MOP$
- F is monotone $\Rightarrow$ $MFP \leq MOP$

Practical Issues

- How to compute ϕs effectively?
  - For general frameworks
    - L not finite
    - F not bounded
  - Does the solution process converge?
- How much space is required to represent ϕ functions?
- Is it possible to avoid explicit function compositions and meets?