CS738: Advanced Compiler Optimizations SSA Continued

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Agenda

- Properties of SSA
- ► SSA to Executable
- SSA for Optimizations

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- ► Computation of SSA: *O*(*R*³)
- In practice, worst case is rare.
- Practical complexity: O(R)

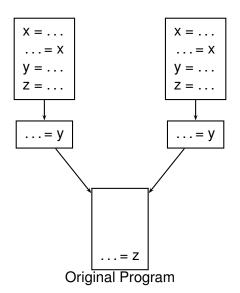
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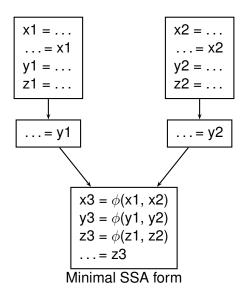
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- Uses a new data structure called DJ-graph
- Linear time is achieved by careful ordering of nodes in the DJ-graph
- DF for a node is computed only once an reused later if required.

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- Semi-Pruned SSA

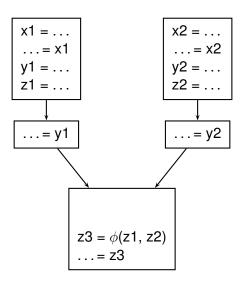
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- Requires global Live variable analysis

Variants of SSA Form: Pruned SSA Example



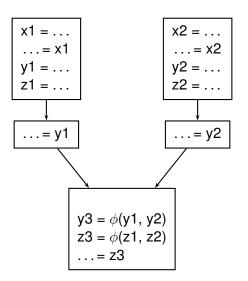
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- ▶ Total number of ϕ -functions between minimal and pruned SSA
- Needs only local Live information
- Non-locals can be computed without iteration or elimination

Variants of SSA Form: Semi-pruned SSA Example



```
foreach block B {
```

```
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  defined = {}
```

```
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```

```
foreach block B {
  defined = {}
  foreach instruction V = X op Y {
    if X not in defined
```

```
foreach block B {
    defined = {}
    foreach instruction V = X \circ p y {
        if X not in defined
            non-locals = non-locals \cup \{X\}
```

Computing Non-locals

```
foreach block B {
  defined = {}
  foreach instruction V = X op Y {
    if X not in defined
       non-locals = non-locals U {X}
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       non-locals = non-locals U {Y}
```

Computing Non-locals

```
foreach block B {
   defined = {}
   foreach instruction v = x op y {
      if x not in defined
            non-locals = non-locals U {x}
   if y not in defined
            non-locals = non-locals U {y}
   defined = defined U {v}
   }
}
```

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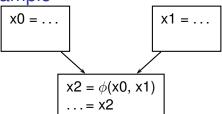
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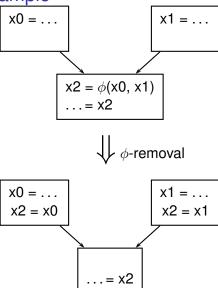
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 - ▶ Need to fix up the ϕ -function
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 - Insert copies in predecessors to mimick φ-function
 - Simple algorithm
 - Works in most cases, but not always
 - Adds lots of copies
 - Many of them will be optimized by later passes

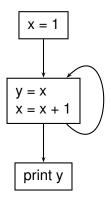
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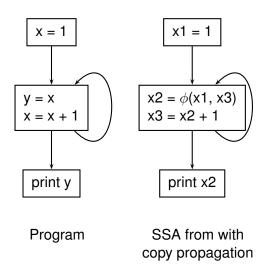


Lost Copy Problem

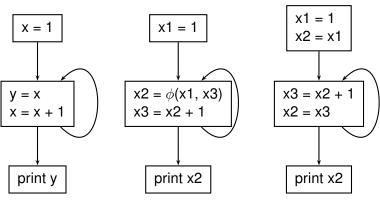


Program

Lost Copy Problem



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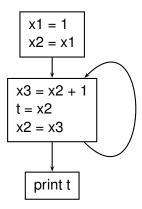


Program

SSA from with copy propagation

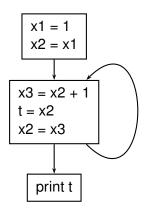
After ϕ -removel

Lost Copy Problem: Solutions

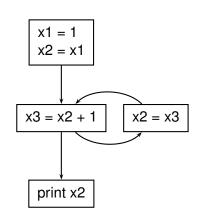


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Lost Copy Problem: Solutions

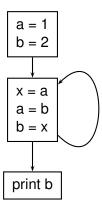


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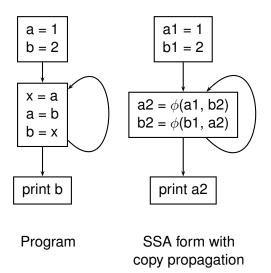
2. Critical Edge Split

Swap Problem

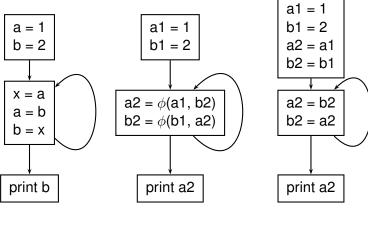


Program

Swap Problem



Swap Problem



Program

SSA form with copy propagation

After ϕ -removel

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- May require temporary if cyclic dependency exists.

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