CS738: Advanced Compiler Optimizations Overview of Optimizations

Amey Karkare

karkare@cse.iitk.ac.in

http://www.cse.iitk.ac.in/~karkare/cs738 Department of CSE, IIT Kanpur



Recap

- Optimizations
 - ► To improve efficiency of generated executable (time, space, resources, ...)
 - ► Maintain semantic equivalence
- ▶ Two levels
 - ► Machine Independent
 - ▶ Machine Dependent

Machine Independent Optimizations

Machine Independent Code Optimizations

- Scope of optimizations
 - Intraprocedural
 - Local
 - Global
 - Interprocedural

Local Optimizations Global Optimizations Restricted to a basic block Typically restricted within a procedure/function Simplifies the analysis ► Could be restricted to a smaller scope, e.g. a loop Not all optimizations can be applied locally ▶ Most compiler implement up to global optimizations ► E.g. Loop optimizations ► Well founded theory Gains are also limited Practical gains Simplify global/interprocedural optimizations Interprocedural Optimizations ► Spans multiple procedures, files A Catalog of In some cases multiple languages! **Code Optimizations** Not as popular as global optimizations ► No single theory applicable to all scenarios ▶ Time consuming

Compile-time Evaluation

- ► Move run-time actions to compile-time
- Constant Folding

Volume =
$$\frac{4}{3} \times \pi \times r \times r \times r$$

- ▶ Compute $\frac{4}{3} \times \pi$ at compile-time
- Applied frequently for linearizing indices of multidimensional arrays
- ► When should we NOT apply it?

Compile-time Evaluation

- ► Constant Propagation
 - ► Replace a variable by its "constant" value

- May result in the application of constant folding
- When should we NOT apply it?

Common Subexpression Elimination

► Reuse a computation if already "available"

- ▶ How to check if an expression is already available?
- ► When should we NOT apply it?

Copy Propagation

- ► Replace (use of) a variable by another variable
 - ► If they are guaranteed to have the "same value"

- ► May result in dead code, common subexpression
- ► When should we NOT apply it?

Code Movement

- ► Move the code around in a program
- Benefits
 - Code size reduction
 - ► Reduction in the frequency of execution
- ▶ How to find out which code to move?

Code Movement

- Code size reduction
 - ➤ Suppose the operator ⊕ results in the generation of a large number of machine instructions. Then,

$$\begin{array}{c} \text{if } (a < b) \\ u = x \oplus y \\ \text{else} \\ v = x \oplus y \end{array} \text{ can be replaced by } \begin{array}{c} t = x \oplus y \\ \text{if } (a < b) \\ u = t \\ \text{else} \\ v = t \end{array}$$

► When should we NOT apply it?

Code Movement

► Execution frequency reduction

$$\begin{array}{c} \text{if } (a < b) \\ u = \dots \\ \text{else} \\ v = x * y \\ w = x * y \end{array} \\ \text{can be replaced by} \\ \begin{array}{c} \text{if } (a < b) \\ u = \dots \\ t = x * y \\ \text{else} \\ t = x * y \\ v = t \\ w = t \end{array}$$

► When should we NOT apply it?

Loop Invariant Code Movement

► Move loop invariant code out of the loop

for
$$(...)$$
 {
...
 $u = a + b$
...
} can be replaced by
...
 $u = t$
...
}

► When should we NOT apply it?

Code Movement	Other Optimizations	
Safety of code motion Profitability of code motion	 Dead code elimination Remove unreachable and/or unused code. Can we always do it? Is there ever a need to introduce unused code? Strength Reduction Use of low strength operators in place of high strength ones. i * i instead of i * 2, pow(i, 2) i << 1 instead of i * 2 Typically performed for integers only – Why? 	
Agenda	Assumptions	
 Static analysis and compile-time optimizations For the next few lectures Intraprocedural Data Flow Analysis Classical Examples Components 	 Intraprocedural: Restricted to a single function Input in 3-address format Unless otherwise specified 	

3-address Code Format

Assignments

▶ Jump/control transfer

```
goto L
if x relop y goto L
```

► Statements can have label(s)

L: . . .

Arrays, Pointers and Functions to be added later when needed