CS738: Advanced Compiler Optimizations Welcome & Introduction

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About the Course

- Program Analysis
- ► Analysis of a Program, by a Program, for a Program¹
 - ► Of a Program User Program
 - ► By a Program Analyzer (Compiler, Runtime)
 - For a Program Optimizer, Verifier
- Transforming user program based on the results of the analysis

Expectations from You

- ► Basic Compiler Knowledge
- Write Code
- ▶ Willingness to understand and modify large code bases
- Read and present state-of-the-art reseach papers

Your Expectations

?

^{1 &}quot;Democracy is the government of the people, by the people, for the people" - Abraham Lincoln

Quick Quizzes (QQs)

- ► There will be small quizzes (10-15 min duration) during the class.
- ► These can be announced or un-announced (surprize quizzes).
- ► Always bring a pen and some loose papers to the class

QQ #1 (Ungraded)

► What are the vaious phases of a typical compiler? (5 minutes)

$$\mathsf{file.c} \to \hspace{-2em} \to \hspace{-2em} \to \cdots \to \hspace{-2em} \to \mathsf{file.asm}$$

Assignments

- ▶ Short assignments to apply the lecture material.
- Assignments will have some written and some programming tasks.
- ▶ 4–5 Assignments for the semester

Using Program Analysis

- ► Compiler Code Optimizations
- ► Why are optimizations important?
- ▶ Why not write optimized code to begin with?
- ▶ Where do optimizations fit in the compiler flow?

Code Optimization

- ► Machine Independent
 - Remove redundancy introduced by the Programmer
 - Remove redundancy not required by later phases of compiler
 - ► Take advantage of algebraic properties of operators
- Machine dependent
 - ► Take advantage of the properties of target machine
- Optimization must preserve the semantics of the original program!

Machine Independent Optimizations

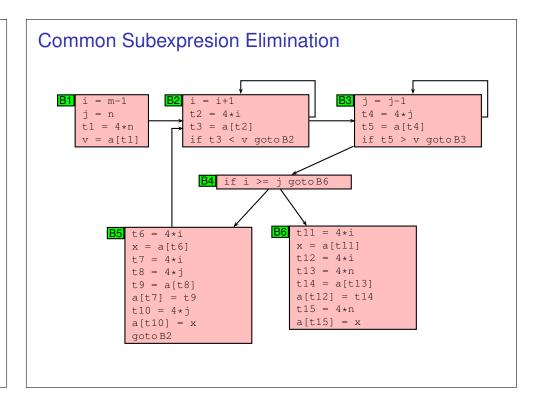
Motivational Example

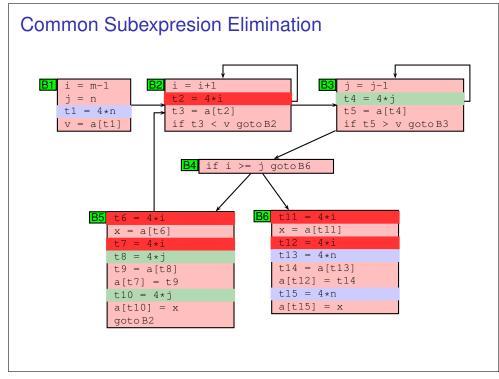
```
void quicksort(int m, int n)
/* recursively sort a[m] through a[n] */
{
    int i, j;
    int v, x;
    if(n <= m) return;
    i = m-1; j = n; v = a[n];
    while (1) {
        do i = i+1; while (a[i] < v);
        do j = j-1; while (a[j] > v);
        if (i > j) break;
        x = a[i]; a[i] = a[j]; a[j] = x;
}

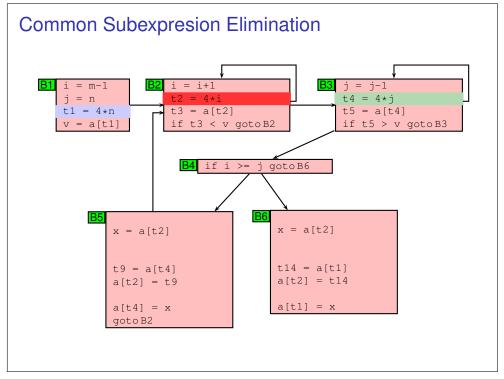
x = a[i]; a[i] = a[n]; a[n] = x;
quicksort(m,j); quicksort(i+1,n);
}
```

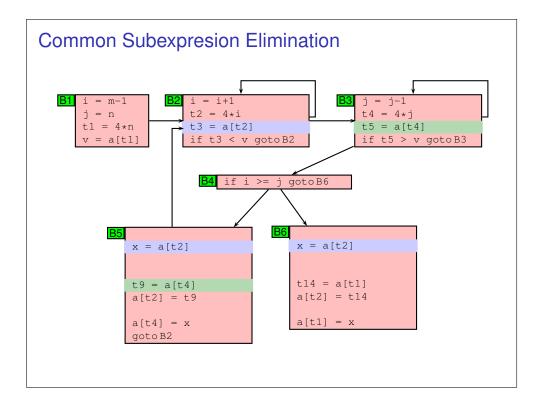
```
(14) t6 = 4*i
                            (15) x = a[t6]
(1) i = m-1
                            (16) t7 = 4 * i
(2) j = n
                            (17) t8 = 4*\dot{7}
(3) t1 = 4*n
                            (18) t9 = a[t8]
                            (19) a[t7] = t9
(4) v = a[t1]
                            (20) t10 = 4 * j
(5) i = i+1
                            (21) a[t10] = x
(6) t2 = 4*i
(7) t3 = a[t2]
                            (22) goto (5)
(8) if t3 < v goto (5)
                            (23) t11 = 4*i
(9) \dot{j} = \dot{j} - 1
                            (24) x = a[t11]
(10) t4 = 4 * j
                            (25) t12 = 4*i
(11) t5 = a[t4]
                            (26) t13 = 4*n
(12) if t5 > v goto (9)
                            (27) t14 = a[t13]
(13) if i >= j \text{ qoto } (23)
                            (28) a[t12] = t14
                            (29) t15 = 4*n
                            (30) a[t15] = x
```

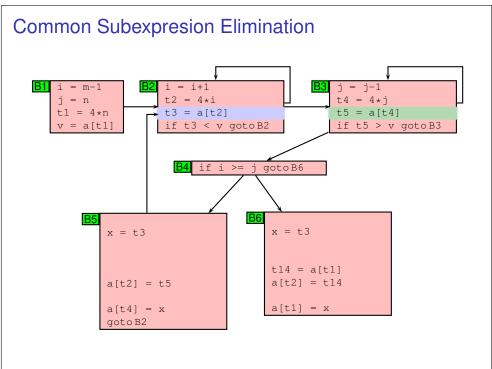
```
(14)
                                 t6 = 4 * i
                            (15) x = a[t6]
(1) i = m-1
                            (16) t7 = 4*i
(2) j = n
                            (17) t8 = 4*\dot{7}
                            (18) t9 = a[t8]
(3) t1 = 4*n
(4) v = a[t1]
                            (19) \quad a[t7] = t9
(5) i = i+1
                            (20) t10 = 4 * j
(6) t2 = 4 * i
                            (21) a[t10] = x
(7) t3 = a[t2]
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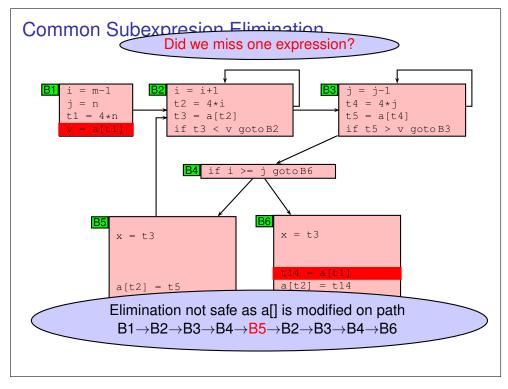


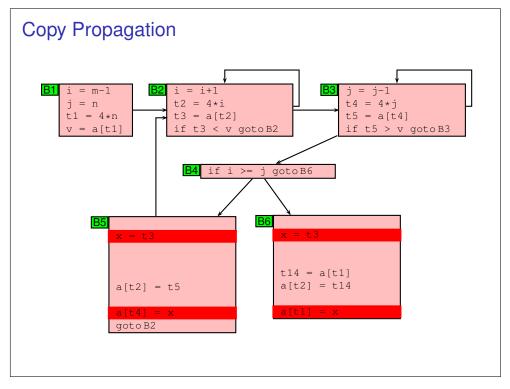


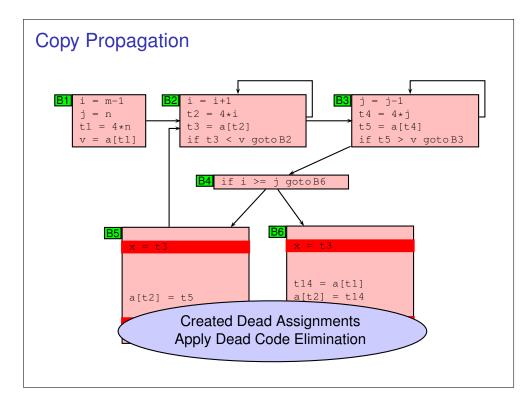


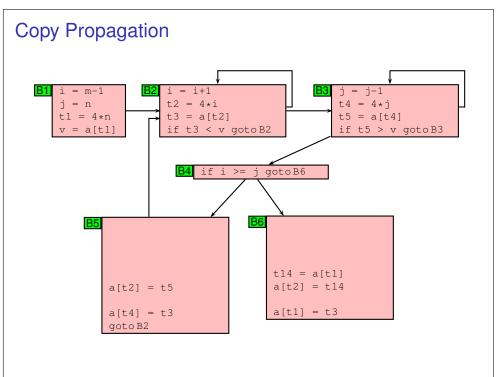


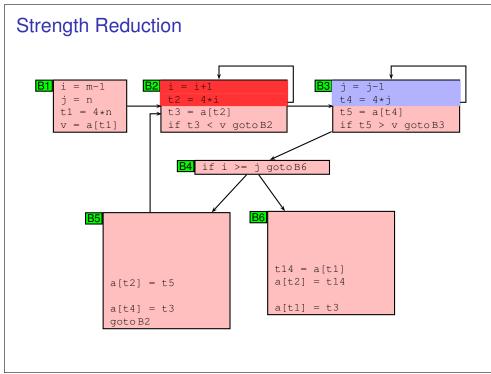


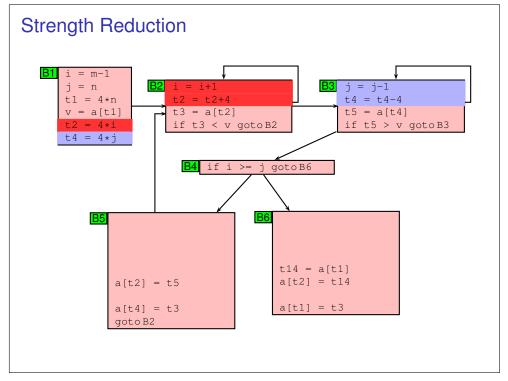


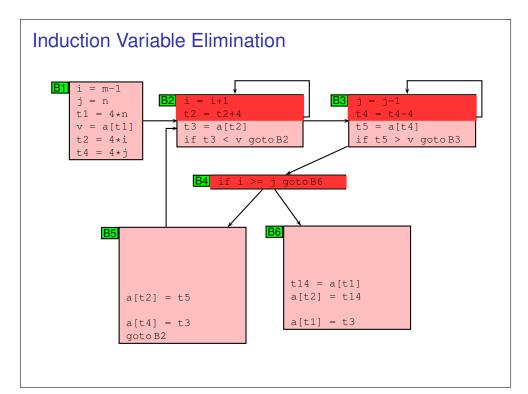


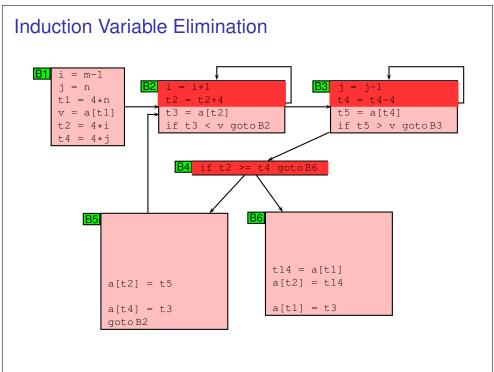


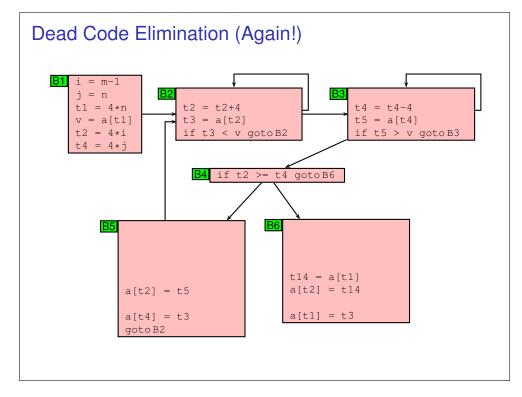












Benefits

В#	# Stmts before Opts	# Stmts after Opts
B1	4	6
B2	4	3
B3	4	3
B4	1	1
B5	9	3
B6	8	3

► Assumptions:

- ► Unit cost for each stmt
- Outer loop: 10 iterations
- ► Inner loops: 100 iterations each

Cost of Execution:

▶ Original Program:

1*4 + 100*4 + 100*4 + 10*1 + 10*9 + 1*8 = 912

Optimized Program:

1*6 + 100*3 + 100*3 + 10*1 + 10*3 + 1*3 = 649

Machine Dependent Optimizations

Peephole Optimizations

- ▶ Target code often contains redundant instructions and suboptimal constructs
- ► Examine a short sequence of target instruction (peephole) and replace by a shorter or faster sequence
- ▶ Peephole is a small moving window on the target systems

Peephole Optimizations: Examples

- Redundant loads and stores
- ► Consider the code sequence

move
$$R_0$$
, a move a , R_0

- ▶ Is instruction 2 redundant? Can we always remove it?
 - ► YES, if it does not have label

Peephole Optimizations: Unreachable code

► Consider the following code

```
int debug = 0;
if (debug) {
   print debugging info
}
```

► This may be translated as

```
int debug = 0;
if (debug == 1) goto L1
  goto L2
L1: print debugging info
L2:
```

Peephole Optimizations: Unreachable code

► Eliminate Jumps

```
int debug = 0;
if (debug != 1) goto L2
print debugging info
L2:
```

► Constant propagation

```
int debug = 0;
if (0 != 1) goto L2
print debugging info
L2:
```

Peephole Optimizations: Unreachable code

► Constant folding and simplification: Since if condition is always true, the code becomes:

```
goto L2
print debugging info
L2:
```

➤ The print statement is now unreachable. Therefore, the code becomes

L2:

Peephole Optimizations: Jump Optimizations

► Replace jump-over-jumps

```
goto L1
: can be replaced by
L1: goto L2
L1: goto L2
```

Peephole Optimizations: Simplify Algebraic Expressions

► Remove

```
x = x + 0;

x = x * 1;
```

Peephole Optimizations: Strength Reduction

- ► Replace *X*² by *X* * *X*
- ► Replace multiplication by left shift
- ► Replace divison by right shift

Peephole Optimizations: Use of Faster Instructions

► Replace
Add #1, R
by
Inc R

Course Logistics

Evaluation

- Assignments
- ► Course project
- ► Mid semester exam
- ► End semester exan
- Quizzes/Class participation
- ► Refer to course webpage for details.