

## CS738: Advanced Compiler Optimizations

### Constant Propagation

Amey Karkare  
karkare@cse.iitk.ac.in

<http://www.cse.iitk.ac.in/~karkare/cs738>  
Department of CSE, IIT Kanpur



## Agenda

- ▶ Using data flow analysis to identify “constant expressions” in a program
- ▶ Identify similarity/differences with bit-vector data flow analyses discussed earlier
- ▶ Other properties of constant propagation

## Special Values for CP

- ▶ **NAC**: not a constant
  - ▶ If variable is inferred not to be a constant
  - ▶ Multiple (different valued) defs, non-const defs, assigned an “un-interpreted” value, ...
- ▶ **Undef**: No definition of the variable is seen yet - nothing known!

## NAC vs Undef

- ▶ **NAC**  $\Rightarrow$  *too many* definitions seen for a variable  $v$  to declare  $v$  is **NOT** a constant
- ▶ **Undef**  $\Rightarrow$  *too few* definitions seen to declare anything about the variable
- ▶  $\top$  is **Undef**;  $\perp$  is **NAC**

## Constant Propagation

- ▶ CP: Replace expressions that evaluate to same constant “c” every time they are executed, by the value “c”

## DF Framework for CP

- ▶ Domain
  - ▶ For a single variable  $v$  of type  $\tau$ , all possible constants of type  $\tau$
- ▶ Semilattice
  - ▶ What is  $\wedge$ ?
  - ▶ What is  $\top$ ?
  - ▶ What is  $\perp$ ?

## CP Meet $\wedge$

- ▶ Recall the requirement

$$\top \wedge x = x$$

$$\perp \wedge x = \perp$$

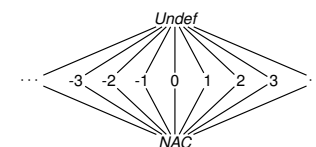
$$Undef \wedge c = c$$

$$NAC \wedge c = NAC$$

$$c_1 \wedge c_2 = NAC \text{ when } c_1 \neq c_2$$

$$c \wedge c = c$$

## CP Semilattice for an integer variable



- ▶ Infinite domain, but finite height

## CP Semilattice

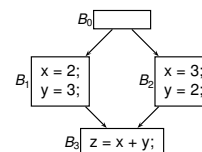
- ▶ Previous figure was semilattice for one variable of one type
- ▶ CP Semilattice = Product of such lattices for all variables (of all types)
- ▶ Each semilattice has a **finite** height

## Computing GEN

- ▶ Informal representation

Statement	GEN
$x = c // \text{const}$	$\{x \rightarrow c\}$
$x = y + z$	if $\{y \rightarrow c_1, z \rightarrow c_2\}$ in IN then $\{x \rightarrow c_1 + c_2\}$ else if $\{y \rightarrow NAC\}$ in IN then $\{x \rightarrow NAC\}$ else if $\{z \rightarrow NAC\}$ in IN then $\{x \rightarrow NAC\}$ else $\{x \rightarrow Undef\}$
$x = \text{complicated expr}$	$\{x \rightarrow NAC\}$

## Nondistributivity of CP



- ▶ All paths:
  - ▶  $B_0 \rightarrow B_1 \rightarrow B_3$
  - ▶  $B_0 \rightarrow B_2 \rightarrow B_3$
- ▶ Value of  $z$  is 5 along both the paths.
- ▶ MOP value for  $z$  is 5.
- ▶ MFP value for  $z$  is **NAC**. (Exercise)
- ▶ MFP value  $\neq$  MOP value (MFP < MOP)

## Monotonicity of CP

- ▶ Case analysis on transfer function  $f$
- ▶  $NAC \leq c \leq Undef$
- ▶  $x = c$  has constant transfer function.
- ▶  $x = \text{complicated expr}$  also has constant transfer function
- ▶ See the next slide for  $x = y + z$  (and similar statements)

## Monotonicity of CP: $x = y + z$

- ▶ Fix  $z$  to be one of **Undef**,  $c_2$ , **NAC**
- ▶ Vary  $y$  over **Undef**,  $c_1$ , **NAC**
- ▶ Confirm that  $x$  does not “increase”
- ▶ Do this for all  $z$  choices.
- ▶ Similarly, fix  $y$  and vary  $z$ .