

CS738: Advanced Compiler Optimizations

SSA Continued

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Agenda

- ▶ Properties of SSA
- ▶ SSA to Executable
- ▶ SSA for Optimizations

Complexity of Construction

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- ▶ Practical complexity: $O(R)$

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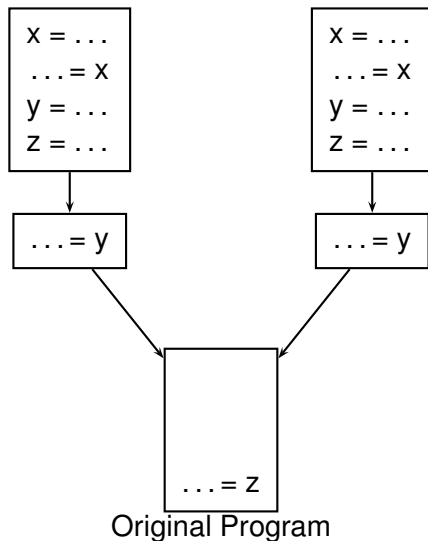
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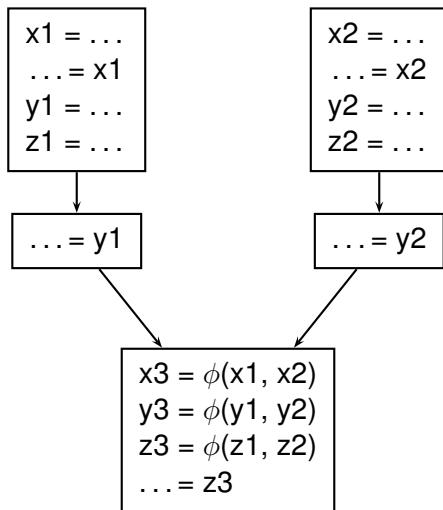
Linear Time Algorithm for ϕ -functions

- ▶ By Sreedhar and Gao, in POPL'95
- ▶ Uses a new data structure called DJ-graph
- ▶ Linear time is achieved by careful ordering of nodes in the DJ-graph
- ▶ DF for a node is computed only once and reused later if required.

Variants of SSA Form: Simple Example



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Minimal SSA form

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- ▶ Semi-Pruned SSA

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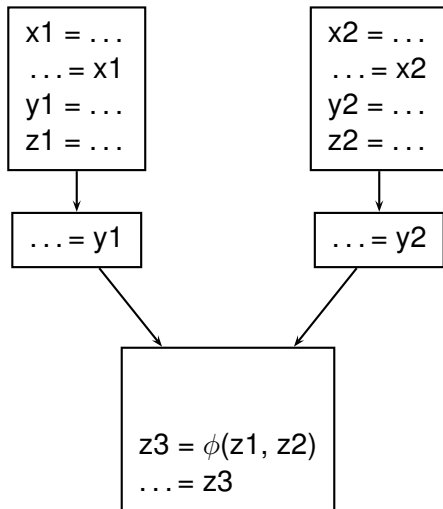
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- ▶ Requires global Live variable analysis

Variants of SSA Form: Pruned SSA Example



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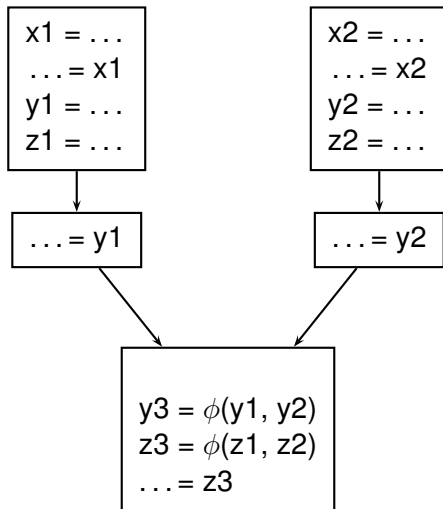
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- ▶ Total number of ϕ -functions between minimal and pruned SSA
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- ▶ Non-locals can be computed without iteration or elimination

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        defined = defined  $\cup \{v\}$   
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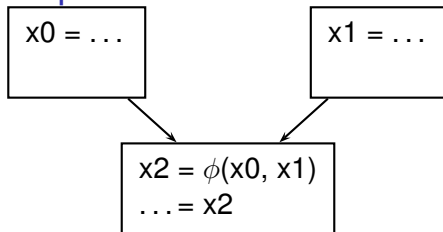
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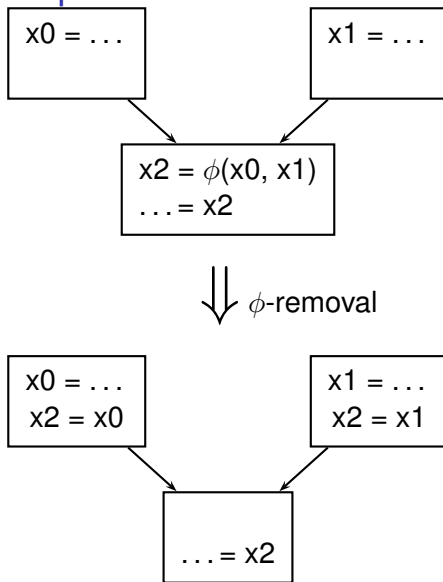
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 - ▶ Works in most cases, but **not always**
 - ▶ Adds lots of copies
 - ▶ Many of them will be optimized by later passes

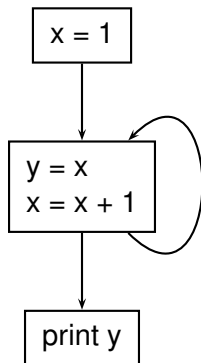
ϕ -removal: Example



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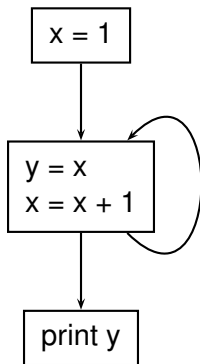


Lost Copy Problem

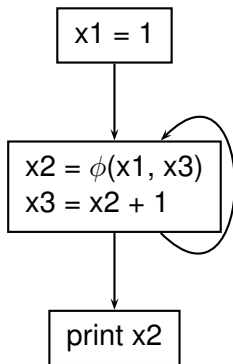


Program

Lost Copy Problem

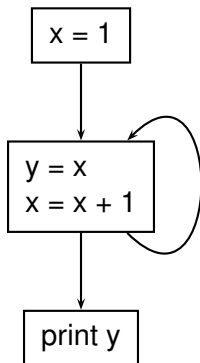


Program

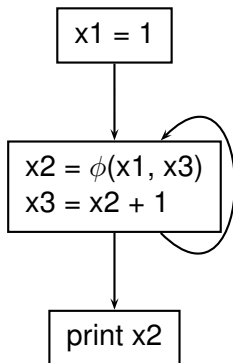


SSA from with
copy propagation

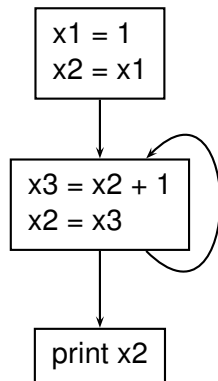
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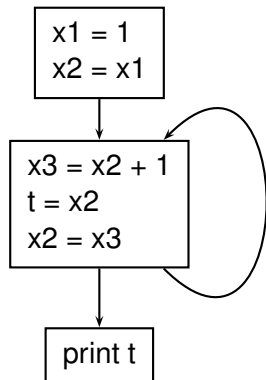


SSA form with
copy propagation



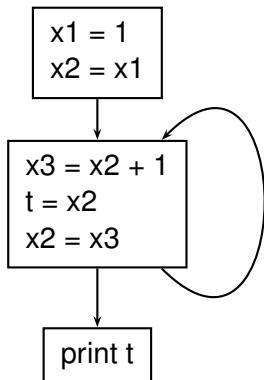
After ϕ -removal

Lost Copy Problem: Solutions

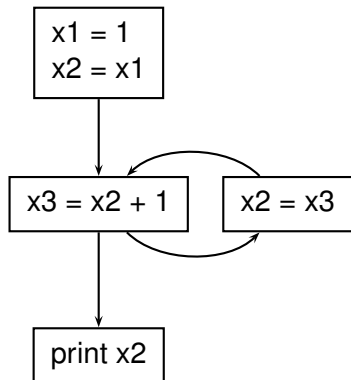


1. Use of Temporary

Lost Copy Problem: Solutions

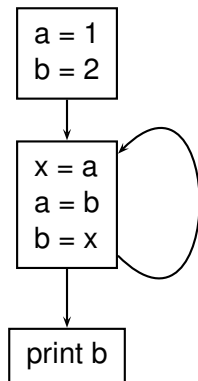


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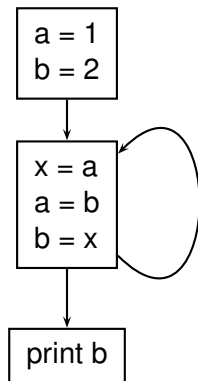
2. Critical Edge Split

Swap Problem

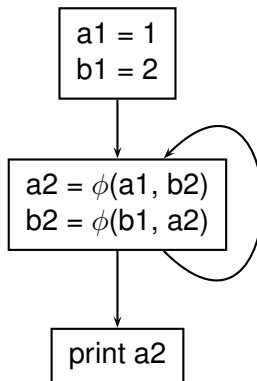


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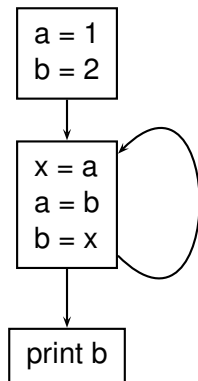


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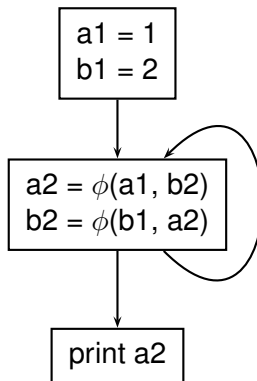


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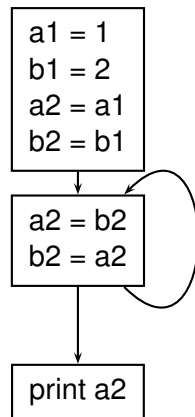
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- ▶ May require temporary if cyclic dependency exists.

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