

CS738: Advanced Compiler Optimizations

Interprocedural Data Flow Analysis

Functional Approach

Amey Karkare

karkare@cse.iitk.ac.in

<http://www.cse.iitk.ac.in/~karkare/cs738>

Department of CSE, IIT Kanpur



Interprocedurally Valid Paths

- ▶ G^* ignores the special nature of call and return edges

Interprocedurally Valid Paths

- ▶ G^* ignores the special nature of call and return edges
- ▶ Not all paths in G^* are feasible

Interprocedurally Valid Paths

- ▶ G^* ignores the special nature of call and return edges
- ▶ Not all paths in G^* are feasible
 - ▶ do not represent potentially valid execution paths

Interprocedurally Valid Paths

- ▶ G^* ignores the special nature of call and return edges
- ▶ Not all paths in G^* are feasible
 - ▶ do not represent potentially valid execution paths
- ▶ $IVP(r_1, n)$: set of all interprocedurally valid paths from r_1 to n

Interprocedurally Valid Paths

- ▶ G^* ignores the special nature of call and return edges
- ▶ Not all paths in G^* are feasible
 - ▶ do not represent potentially valid execution paths
- ▶ $IVP(r_1, n)$: set of all interprocedurally valid paths from r_1 to n
- ▶ Path $q \in \text{path}_{G^*}(r_1, n)$ is in $IVP(r_1, n)$

Interprocedurally Valid Paths

- ▶ G^* ignores the special nature of call and return edges
- ▶ Not all paths in G^* are feasible
 - ▶ do not represent potentially valid execution paths
- ▶ $IVP(r_1, n)$: set of all interprocedurally valid paths from r_1 to n
- ▶ Path $q \in \text{path}_{G^*}(r_1, n)$ is in $IVP(r_1, n)$
 - ▶ iff sequence of all E^1 edges in q (denoted q_1) is *proper*

Proper sequence

- ▶ q_1 without any return edge is proper

Proper sequence

- ▶ q_1 without any return edge is proper
- ▶ let $q_1[i]$ be the first return edge in q_1 . q_1 is proper if

Proper sequence

- ▶ q_1 without any return edge is proper
- ▶ let $q_1[i]$ be the first return edge in q_1 . q_1 is proper if
 - ▶ $i > 1$; and

Proper sequence

- ▶ q_1 without any return edge is proper
- ▶ let $q_1[i]$ be the first return edge in q_1 . q_1 is proper if
 - ▶ $i > 1$; and
 - ▶ $q_1[i - 1]$ is call edge corresponding to $q_1[i]$; and

Proper sequence

- ▶ q_1 without any return edge is proper
- ▶ let $q_1[i]$ be the first return edge in q_1 . q_1 is proper if
 - ▶ $i > 1$; and
 - ▶ $q_1[i - 1]$ is call edge corresponding to $q_1[i]$; and
 - ▶ q'_1 obtained from deleting $q_1[i - 1]$ and $q_1[i]$ from q_1 is proper

Interprocedurally Valid Complete Paths

- ▶ $IVP_0(r_p, n)$ for procedure p and node $n \in N_p$

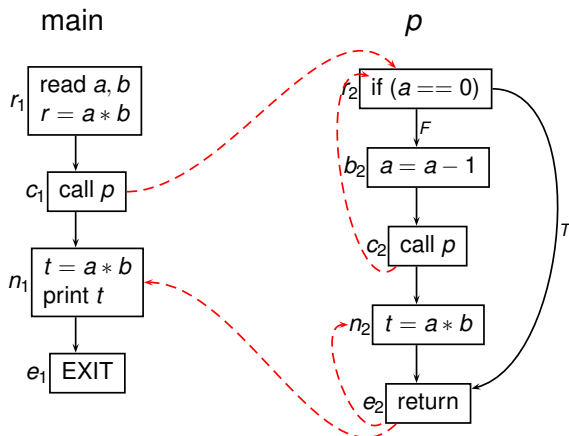
Interprocedurally Valid Complete Paths

- ▶ $IVP_0(r_p, n)$ for procedure p and node $n \in N_p$
- ▶ set of all interprocedurally valid paths q in G^* from r_p to n s.t.

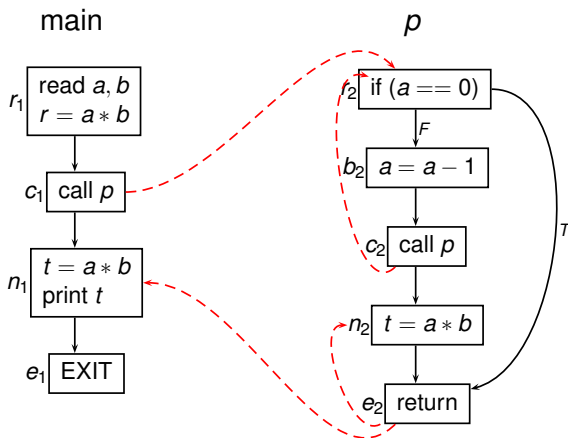
Interprocedurally Valid Complete Paths

- ▶ $IVP_0(r_p, n)$ for procedure p and node $n \in N_p$
- ▶ set of all interprocedurally valid paths q in G^* from r_p to n s.t.
 - ▶ Each call edge has corresponding return edge in q restricted to E^1

IVPs

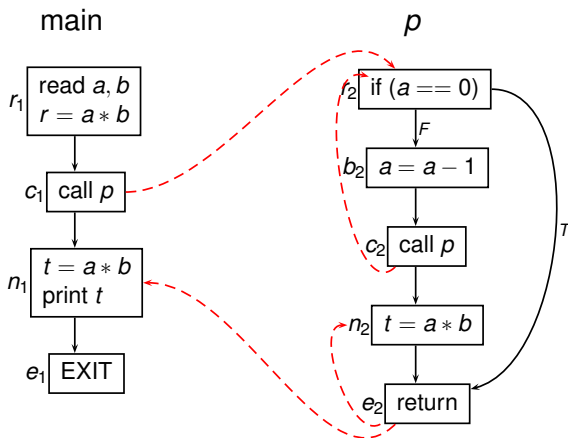


IVPs



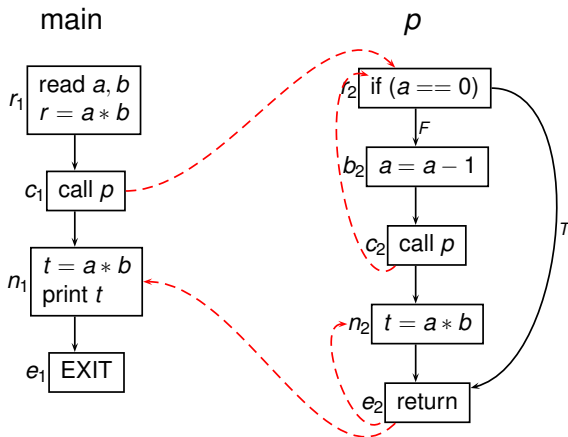
$r_1 \rightarrow c_1 \rightarrow r_2 \rightarrow c_2 \rightarrow r_2 \rightarrow e_2 \rightarrow n_2 \rightarrow e_2 \rightarrow n_1 \rightarrow e_1$

IVPs



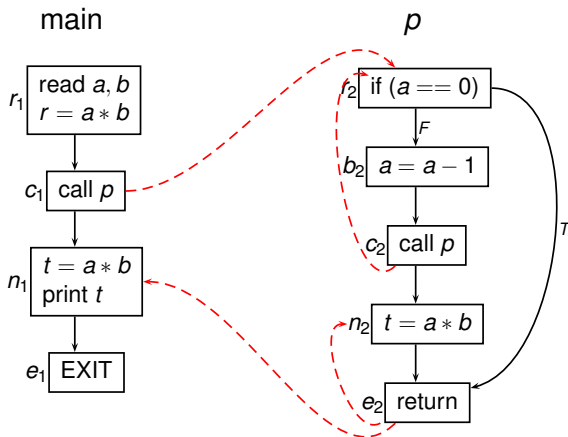
$$r_1 \rightarrow c_1 \rightarrow r_2 \rightarrow c_2 \rightarrow r_2 \rightarrow e_2 \rightarrow n_2 \rightarrow e_2 \rightarrow n_1 \rightarrow e_1 \\ \in \text{IVP}(r_1, e_1)$$

IVPs



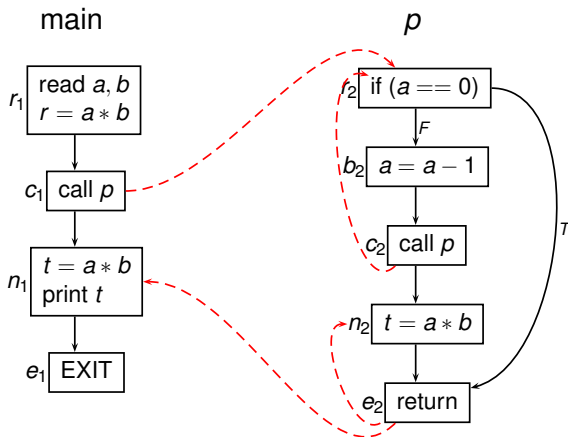
$r_1 \rightarrow c_1 \rightarrow r_2 \rightarrow c_2 \rightarrow r_2 \rightarrow e_2 \rightarrow n_1 \rightarrow e_1$

IVPs



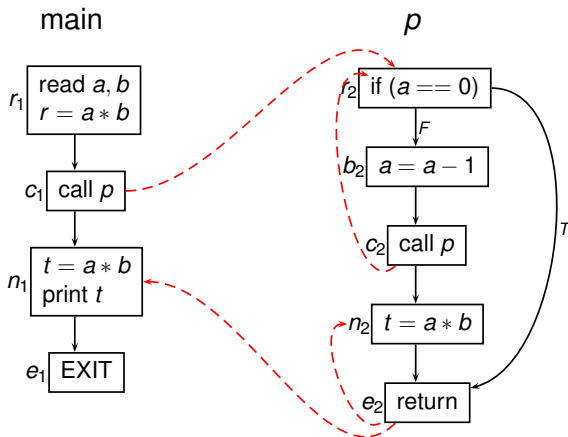
$r_1 \rightarrow c_1 \rightarrow r_2 \rightarrow c_2 \rightarrow r_2 \rightarrow e_2 \rightarrow n_1 \rightarrow e_1 \notin \text{IVP}(r_1, e_1)$

IVPs



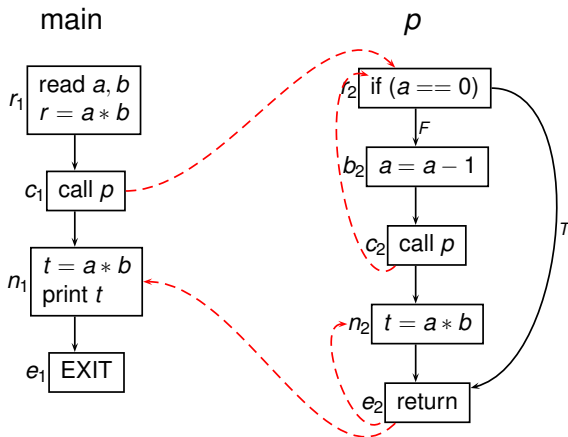
$r_2 \rightarrow c_2 \rightarrow r_2 \rightarrow e_2 \rightarrow n_2$

IVPs



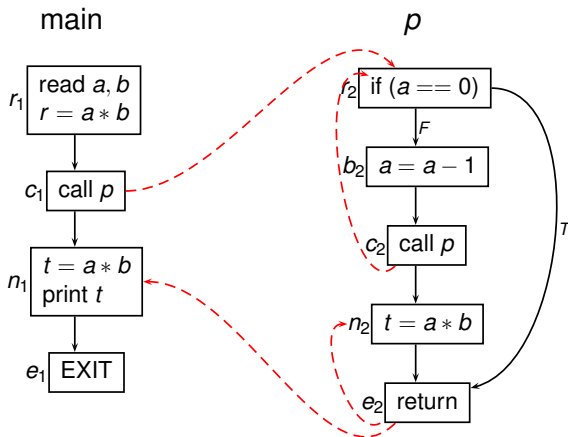
$$r_2 \rightarrow c_2 \rightarrow r_2 \rightarrow e_2 \rightarrow n_2 \in \text{IVP}_0(r_2, n_2)$$

IVPs



$r_2 \rightarrow c_2 \rightarrow r_2 \rightarrow c_2 \rightarrow e_2 \rightarrow n_2$

IVPs



$$r_2 \rightarrow c_2 \rightarrow r_2 \rightarrow c_2 \rightarrow e_2 \rightarrow n_2 \notin \text{IVP}_0(r_2, n_2)$$

Path Decomposition

$$q \in \text{IVP}(r_{\text{main}}, n)$$

\Leftrightarrow

$$q = q_1 \parallel (c_1, r_{p_2}) \parallel q_2 \parallel \cdots \parallel (c_{j-1}, r_{p_j}) \parallel q_j$$

where for each $i < j$, $q_i \in \text{IVP}_0(r_{p_i}, c_i)$ and $q_j \in \text{IVP}_0(r_{p_j}, n)$

Functional Approach

- ▶ (L, F) : a *distributive* data flow framework

Functional Approach

- ▶ (L, F) : a *distributive* data flow framework
- ▶ Procedure p , node $n \in N_p$

Functional Approach

- ▶ (L, F) : a *distributive* data flow framework
- ▶ Procedure p , node $n \in N_p$
- ▶ $\phi_{(r_p, n)} \in F$ describes flow of data flow information from start of r_p to start of n

Functional Approach

- ▶ (L, F) : a *distributive* data flow framework
- ▶ Procedure p , node $n \in N_p$
- ▶ $\phi_{(r_p, n)} \in F$ describes flow of data flow information from start of r_p to start of n
 - ▶ along paths in $IVP_0(r_p, n)$

Functional Approach Constraints

$$\phi_{(r_p, r_p)} \leq id_L$$

$$\phi_{(r_p, n)} = \bigwedge_{(m, n) \in E_p} (h_{(m, n)} \circ \phi_{(r_p, m)}) \quad \text{for } n \in N_p$$

where

$$h_{(m, n)} = \begin{cases} f_{(m, n)} & \text{if } (m, n) \in E_p^0, \\ & f_{(m, n)} \in F \text{ associated flow function} \\ \phi_{(r_q, e_q)} & \text{if } (m, n) \in E_p^1 \text{ and } m \text{ calls procedure } q \end{cases}$$

Information x at r_p translated to information $\phi_{(r_p, n)}(x)$ at n

Solving ϕ Constraints

- ▶ Round-robin iterative approximations to initial values

$$\begin{aligned}\phi_{(r_p, r_p)}^0 &\leq id_L \\ \phi_{(r_p, n)}^0 &\leq f_\Omega \quad n \in N_p - \{r_p\}\end{aligned}$$

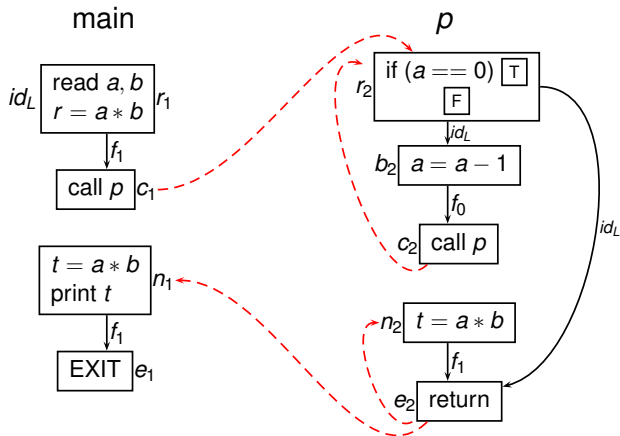
Solving ϕ Constraints

- ▶ Round-robin iterative approximations to initial values

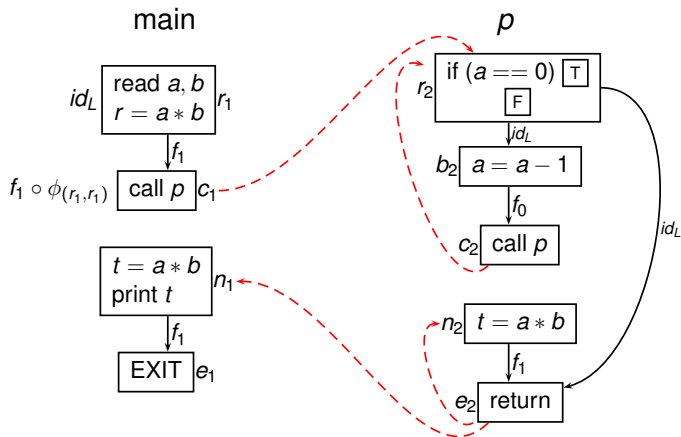
$$\begin{aligned}\phi_{(r_p, r_p)}^0 &\leq id_L \\ \phi_{(r_p, n)}^0 &\leq f_\Omega \quad n \in N_p - \{r_p\}\end{aligned}$$

- ▶ Reach maximal fixed point

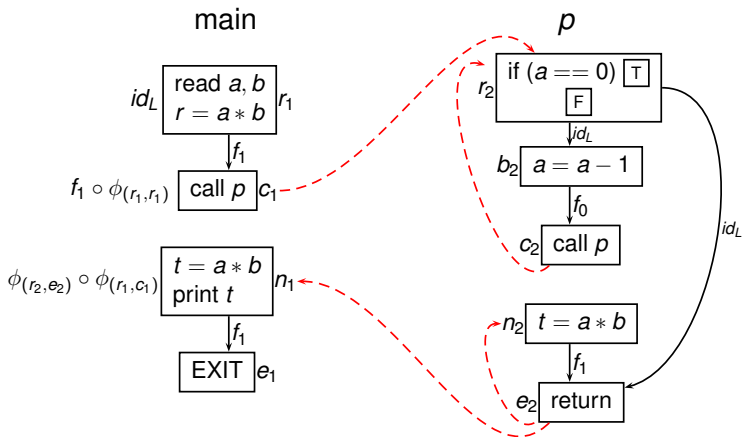
Example



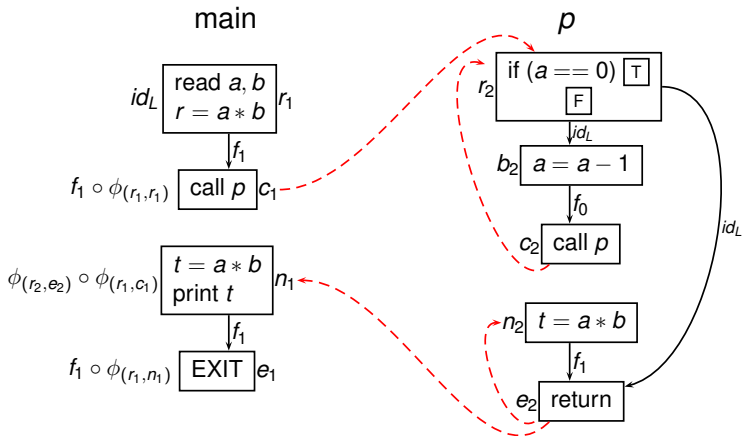
Example



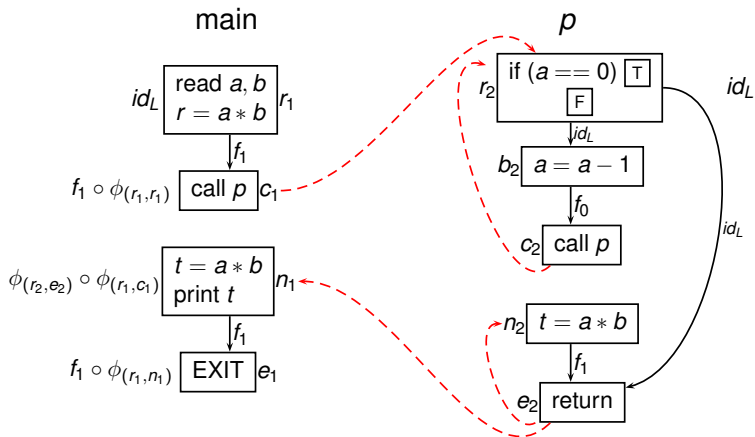
Example



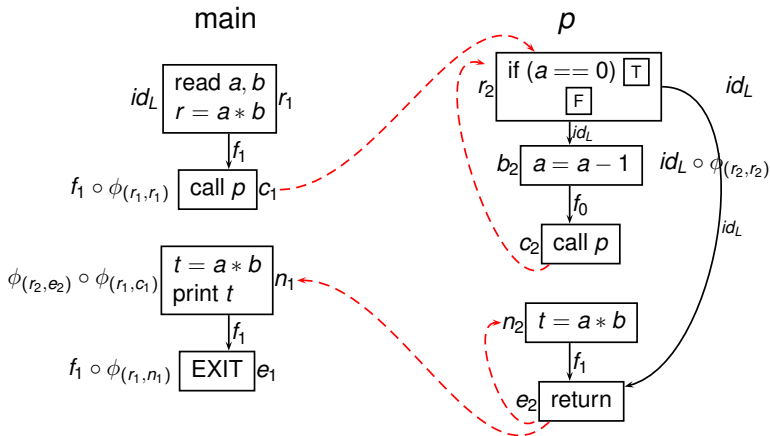
Example



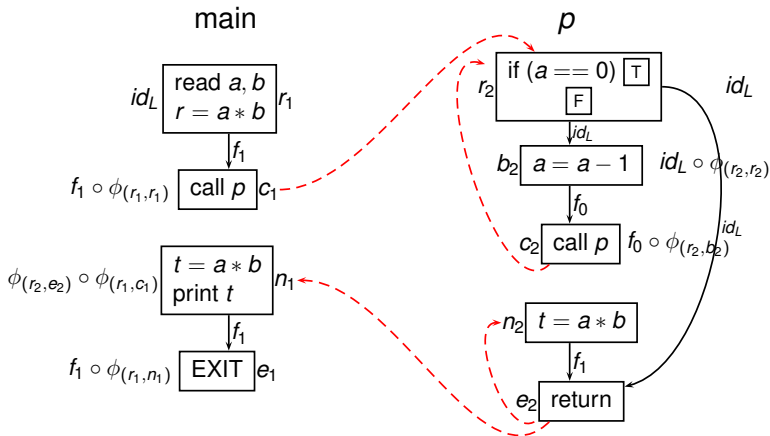
Example



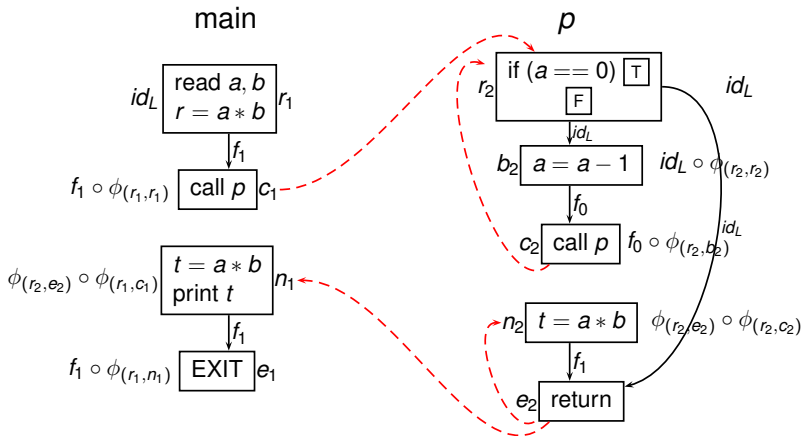
Example



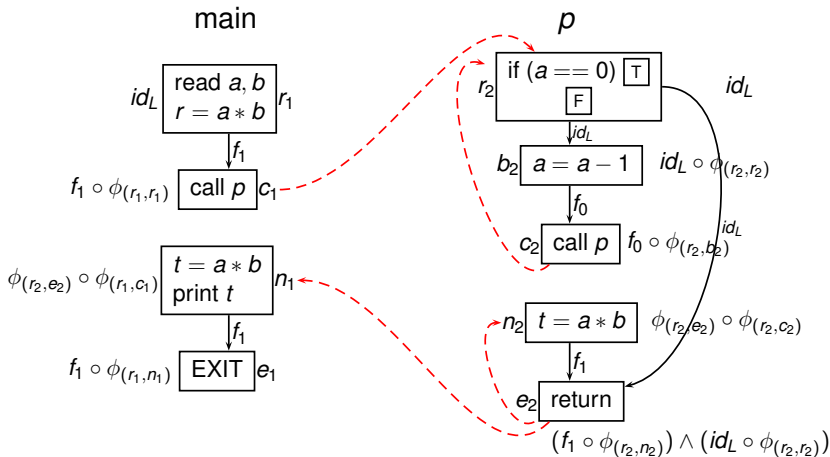
Example



Example



Example



Iterative Solution

Function	Constraint	Iteration #		
		Init	1 st	2 nd 3 rd

Iterative Solution

Function	Constraint	Iteration #			
		Init	1 st	2 nd	3 rd
$\phi(r_1, r_1)$	id_L	id	id	id	id

Iterative Solution

Function	Constraint	Iteration #			
		Init	1 st	2 nd	3 rd
$\phi_{(r_1, r_1)}$	id_L	id	id	id	id
$\phi_{(r_1, c_1)}$	$f_1 \circ \phi_{(r_1, r_1)}$	f_Ω	f_1	f_1	f_1

Iterative Solution

Function	Constraint	Iteration #			
		Init	1 st	2 nd	3 rd
$\phi(r_1, r_1)$	id_L	id	id	id	id
$\phi(r_1, c_1)$	$f_1 \circ \phi(r_1, r_1)$	f_Ω	f_1	f_1	f_1
$\phi(r_1, n_1)$	$\phi(r_2, e_2) \circ \phi(r_1, c_1)$	f_Ω	f_Ω	f_1	f_1

Iterative Solution

Function	Constraint	Iteration #			
		Init	1 st	2 nd	3 rd
$\phi(r_1, r_1)$	id_L	id	id	id	id
$\phi(r_1, c_1)$	$f_1 \circ \phi(r_1, r_1)$	f_Ω	f_1	f_1	f_1
$\phi(r_1, n_1)$	$\phi(r_2, e_2) \circ \phi(r_1, c_1)$	f_Ω	f_Ω	f_1	f_1
$\phi(r_1, e_1)$	$f_1 \circ \phi(r_1, n_1)$	f_Ω	f_1	f_1	f_1

Iterative Solution

Function	Constraint	Iteration #			
		Init	1 st	2 nd	3 rd
$\phi(r_1, r_1)$	id_L	id	id	id	id
$\phi(r_1, c_1)$	$f_1 \circ \phi(r_1, r_1)$	f_Ω	f_1	f_1	f_1
$\phi(r_1, n_1)$	$\phi(r_2, e_2) \circ \phi(r_1, c_1)$	f_Ω	f_Ω	f_1	f_1
$\phi(r_1, e_1)$	$f_1 \circ \phi(r_1, n_1)$	f_Ω	f_1	f_1	f_1
$\phi(r_2, r_2)$	id_L	id	id	id	id

Iterative Solution

Function	Constraint	Iteration #			
		Init	1 st	2 nd	3 rd
$\phi(r_1, r_1)$	id_L	id	id	id	id
$\phi(r_1, c_1)$	$f_1 \circ \phi(r_1, r_1)$	f_Ω	f_1	f_1	f_1
$\phi(r_1, n_1)$	$\phi(r_2, e_2) \circ \phi(r_1, c_1)$	f_Ω	f_Ω	f_1	f_1
$\phi(r_1, e_1)$	$f_1 \circ \phi(r_1, n_1)$	f_Ω	f_1	f_1	f_1
$\phi(r_2, r_2)$	id_L	id	id	id	id
$\phi(r_2, b_2)$	$id_L \circ \phi(r_2, r_2)$	f_Ω	id	id	id

Iterative Solution

Function	Constraint	Iteration #			
		Init	1 st	2 nd	3 rd
$\phi(r_1, r_1)$	id_L	id	id	id	id
$\phi(r_1, c_1)$	$f_1 \circ \phi(r_1, r_1)$	f_Ω	f_1	f_1	f_1
$\phi(r_1, n_1)$	$\phi(r_2, e_2) \circ \phi(r_1, c_1)$	f_Ω	f_Ω	f_1	f_1
$\phi(r_1, e_1)$	$f_1 \circ \phi(r_1, n_1)$	f_Ω	f_1	f_1	f_1
$\phi(r_2, r_2)$	id_L	id	id	id	id
$\phi(r_2, b_2)$	$id_L \circ \phi(r_2, r_2)$	f_Ω	id	id	id
$\phi(r_2, c_2)$	$f_0 \circ \phi(r_2, b_2)$	f_Ω	f_0	f_0	f_0

Iterative Solution

Function	Constraint	Iteration #			
		Init	1 st	2 nd	3 rd
$\phi(r_1, r_1)$	id_L	id	id	id	id
$\phi(r_1, c_1)$	$f_1 \circ \phi(r_1, r_1)$	f_Ω	f_1	f_1	f_1
$\phi(r_1, n_1)$	$\phi(r_2, e_2) \circ \phi(r_1, c_1)$	f_Ω	f_Ω	f_1	f_1
$\phi(r_1, e_1)$	$f_1 \circ \phi(r_1, n_1)$	f_Ω	f_1	f_1	f_1
$\phi(r_2, r_2)$	id_L	id	id	id	id
$\phi(r_2, b_2)$	$id_L \circ \phi(r_2, r_2)$	f_Ω	id	id	id
$\phi(r_2, c_2)$	$f_0 \circ \phi(r_2, b_2)$	f_Ω	f_0	f_0	f_0
$\phi(r_2, n_2)$	$\phi(r_2, e_2) \circ \phi(r_2, c_2)$	f_Ω	f_Ω	f_0	f_0

Iterative Solution

Function	Constraint	Iteration #			
		Init	1 st	2 nd	3 rd
$\phi(r_1, r_1)$	id_L	id	id	id	id
$\phi(r_1, c_1)$	$f_1 \circ \phi(r_1, r_1)$	f_Ω	f_1	f_1	f_1
$\phi(r_1, n_1)$	$\phi(r_2, e_2) \circ \phi(r_1, c_1)$	f_Ω	f_Ω	f_1	f_1
$\phi(r_1, e_1)$	$f_1 \circ \phi(r_1, n_1)$	f_Ω	f_1	f_1	f_1
$\phi(r_2, r_2)$	id_L	id	id	id	id
$\phi(r_2, b_2)$	$id_L \circ \phi(r_2, r_2)$	f_Ω	id	id	id
$\phi(r_2, c_2)$	$f_0 \circ \phi(r_2, b_2)$	f_Ω	f_0	f_0	f_0
$\phi(r_2, n_2)$	$\phi(r_2, e_2) \circ \phi(r_2, c_2)$	f_Ω	f_Ω	f_0	f_0
$\phi(r_2, e_2)$	$(f_1 \circ \phi(r_2, n_2)) \wedge (id_L \circ \phi(r_2, r_2))$	f_Ω	id	id	id

Solving Data Flow Problem

- ▶ The above process gives solution to ϕ functions

Solving Data Flow Problem

- ▶ The above process gives solution to ϕ functions
- ▶ Use it to compute data flow information x_n associated with start of block n

$$x_{r_{\text{main}}} = \textit{BoundaryInfo}$$

for each procedure p

$$x_{r_p} = \bigwedge \left\{ \phi_{(r_q, c)}(x_{r_q}) : \begin{array}{l} q \text{ is a procedure and} \\ c \text{ is a call to } p \text{ in } q \end{array} \right\}$$
$$x_n = \phi_{(r_p, n)} \quad n \in N_p - \{r_p\}$$

Solving Data Flow Problem

- ▶ The above process gives solution to ϕ functions
- ▶ Use it to compute data flow information x_n associated with start of block n

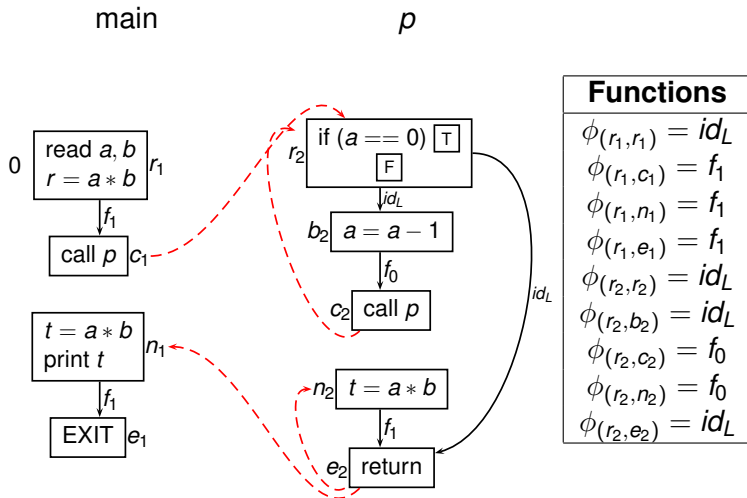
$$x_{r_{\text{main}}} = \textit{BoundaryInfo}$$

for each procedure p

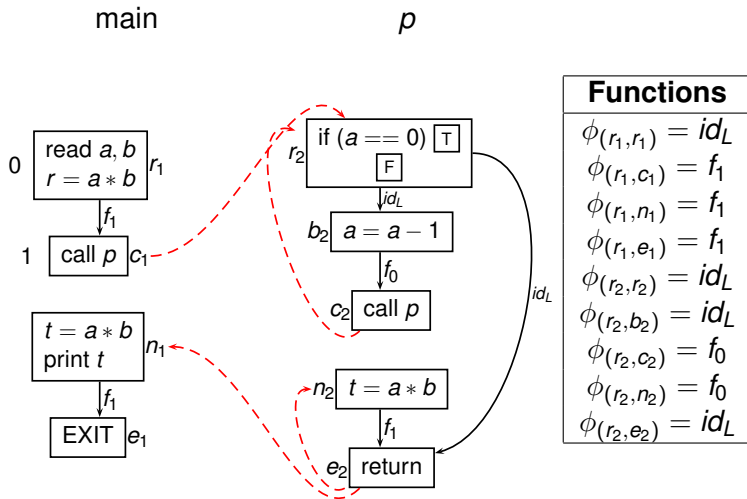
$$x_{r_p} = \bigwedge \left\{ \phi_{(r_q, c)}(x_{r_q}) : \begin{array}{l} q \text{ is a procedure and} \\ c \text{ is a call to } p \text{ in } q \end{array} \right\}$$
$$x_n = \phi_{(r_p, n)} \quad n \in N_p - \{r_p\}$$

- ▶ Iterative algorithm for solution, maximal fixed point solution

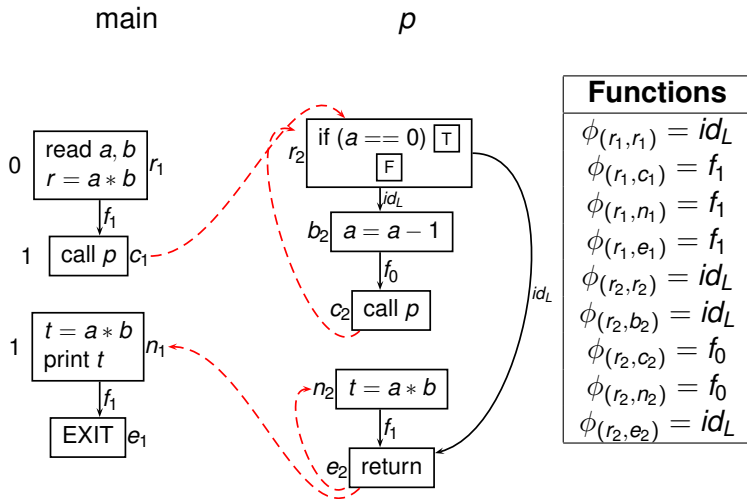
Example



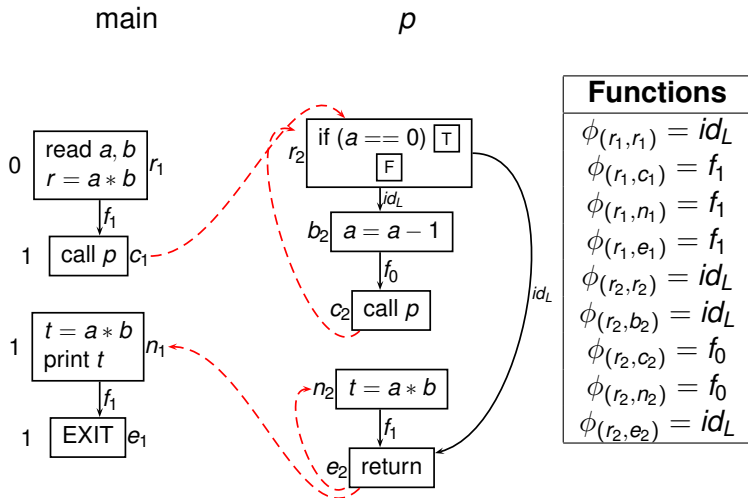
Example



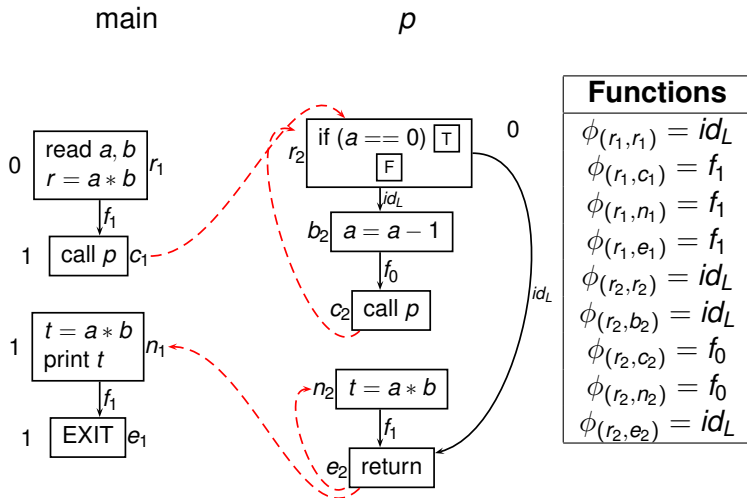
Example



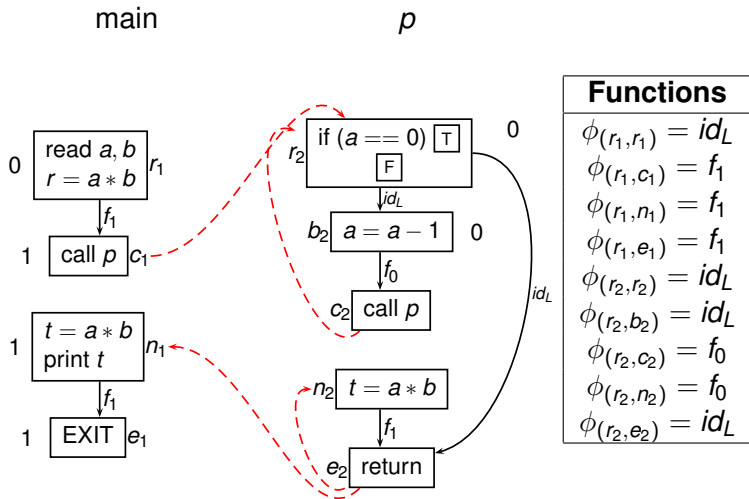
Example



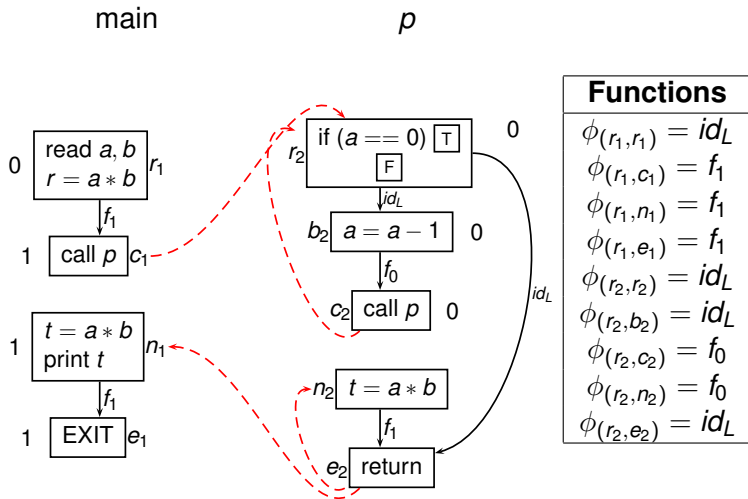
Example



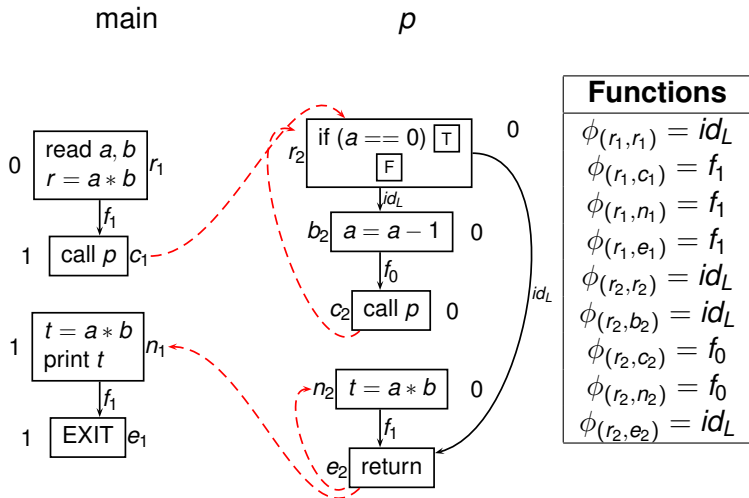
Example



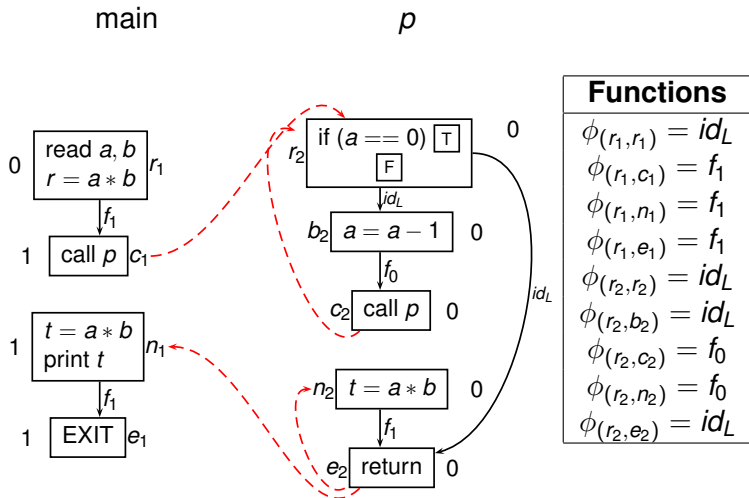
Example



Example



Example



Interprocedural MOP

$$\begin{aligned}\psi_n &= \bigwedge \{f_q : q \in \text{IVP}(r_{\text{main}}, n)\} \in F & \forall n \in N^* \\ y_n &= \psi_n(\text{BoundaryInfo}) & \forall n \in N^*\end{aligned}$$

Interprocedural MOP

$$\begin{aligned}\psi_n &= \bigwedge \{f_q : q \in \text{IVP}(r_{\text{main}}, n)\} \in F & \forall n \in N^* \\ y_n &= \psi_n(\text{BoundaryInfo}) & \forall n \in N^*\end{aligned}$$

y_n is the *meet-over-all-paths solution* (MOP).

IVP₀ Lemma

$$\phi_{(r_p, n)} = \bigwedge \{f_q : q \in \text{IVP}_0(r_p, n)\} \quad \forall n \in N_p$$

IVP₀ Lemma

$$\phi_{(r_p, n)} = \bigwedge \{f_q : q \in \text{IVP}_0(r_p, n)\} \quad \forall n \in N_p$$

Proof: By induction (**Exercise/Reading Assignment**)

$$\begin{aligned}\Psi_n &= \bigwedge \{f_q : q \in \text{IVP}(r_{\text{main}}, n)\} \in F \quad \forall n \in N^* \\ \mathcal{X}_n &= \wedge \{\phi_{(r_{p_j}, n)} \circ \phi_{(r_{p_{j-1}}, c_{j-1})} \circ \dots \circ \phi_{(r_{p_1}, c_1)} \mid c_i \text{ calls } p_{i+1}\}\end{aligned}$$

$$\begin{aligned}\psi_n &= \bigwedge \{f_q : q \in \text{IVP}(r_{\text{main}}, n)\} \in F \quad \forall n \in N^* \\ \mathcal{X}_n &= \wedge \{\phi_{(r_{p_j}, n)} \circ \phi_{(r_{p_{j-1}}, c_{j-1})} \circ \dots \circ \phi_{(r_{p_1}, c_1)} \mid c_i \text{ calls } p_{i+1}\}\end{aligned}$$

Then

$$\psi_n = \mathcal{X}_n$$

Proof: IVP₀ Lemma and Path decomposition

$$\begin{aligned}\Psi_n &= \bigwedge \{f_q : q \in \text{IVP}(r_{\text{main}}, n)\} \in F \quad \forall n \in N^* \\ \mathcal{X}_n &= \wedge \{\phi_{(r_{p_j}, n)} \circ \phi_{(r_{p_{j-1}}, c_{j-1})} \circ \dots \circ \phi_{(r_{p_1}, c_1)} \mid c_i \text{ calls } p_{i+1}\}\end{aligned}$$

Then

$$\Psi_n = \mathcal{X}_n$$

Proof: IVP₀ Lemma and Path decomposition

$$y_n = \Psi_n(\text{BoundaryInfo}) = \mathcal{X}_n(\text{BoundaryInfo})$$

MOP vs MFP

- ▶ F is distributive $\Rightarrow MFP = MOP$

MOP vs MFP

- ▶ F is distributive $\Rightarrow MFP = MOP$
- ▶ F is monotone $\Rightarrow MFP \leq MOP$

Practical Issues

- ▶ How to compute ϕ s effectively?

Practical Issues

- ▶ How to compute ϕ s effectively?
 - ▶ For general frameworks

Practical Issues

- ▶ How to compute ϕ s effectively?
 - ▶ For general frameworks
 - ▶ L not finite

Practical Issues

- ▶ How to compute ϕ s effectively?
 - ▶ For general frameworks
 - ▶ L not finite
 - ▶ F not bounded

Practical Issues

- ▶ How to compute ϕ s effectively?
 - ▶ For general frameworks
 - ▶ L not finite
 - ▶ F not bounded
 - ▶ Does the solution process converge?

Practical Issues

- ▶ How to compute ϕ s effectively?
 - ▶ For general frameworks
 - ▶ L not finite
 - ▶ F not bounded
 - ▶ Does the solution process converge?
- ▶ How much space is required to represent ϕ functions?

Practical Issues

- ▶ How to compute ϕ s effectively?
 - ▶ For general frameworks
 - ▶ L not finite
 - ▶ F not bounded
 - ▶ Does the solution process converge?
- ▶ How much space is required to represent ϕ functions?
- ▶ Is it possible to avoid explicit function compositions and meets?