

CS738: Advanced Compiler Optimizations

SSAPRE: SSA based Partial Redundancy Elimination

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PRE without SSA

- ▶ Based on well known DF analyses
 - ▶ Availability
 - ▶ Anticipability
 - ▶ Partial Availability
 - ▶ Partial Anticipability
- ▶ Identifies partially redundant computations, make them totally redundant by inserting new computations
- ▶ Remove totally redundant computations (CSE)

PRE without SSA

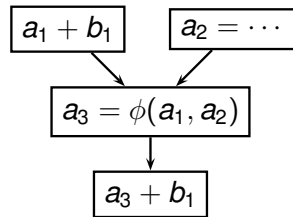
- ▶ Iterative data flow analysis
- ▶ Operates on control flow graph
- ▶ Computes global and local versions of data flow information

SSAPRE

- ▶ Information flow along SSA edges
- ▶ No distinction between global and local information

SSAPRE: Challenge

- ▶ SSA form defined for variables
- ▶ How to identify potentially redundant expressions
 - ▶ Expressions having different variable versions as operands



- ▶ Here $a_1 + b_1$ is same as $a_3 + b_1$ when control follows the left branch. Lexically different, but computationally identical

SSAPRE: Key Idea

- ▶ Redundancy Class Variables (RCVs)
 - ▶ variable (say h) to represent computation of an expression (say E)
- ▶ Computation of expression could represent either a *def* or a *use*
 - ▶ definition of $E \Rightarrow$ store into h
 - ▶ use of $E \Rightarrow$ load from h
- ▶ PRE on SSA form of RCVs (h) to remove redundancies
- ▶ Final program will be in SSA form

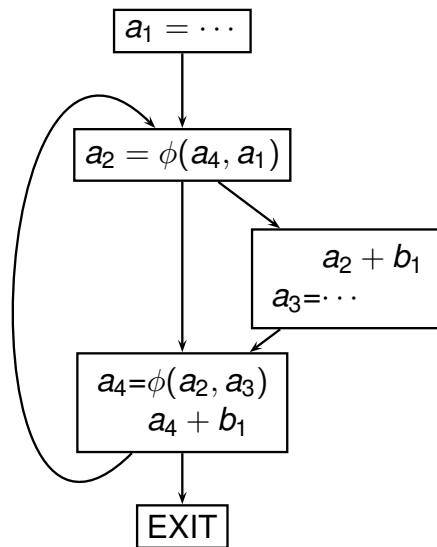
SSAPRE: Preparations

- ▶ Split all the *critical edges* in the flow graph
 - ▶ Edge from a node with more than one successor to a node with more than one predecessor
 - ▶ WHY is this important?
- ▶ Single pass to identify identical expressions
 - ▶ Ignoring the version number of the operands
 - ▶ In the earlier example, $a_3 + b_1$ and $a_1 + b_1$ could be identical

SSAPRE Steps

- ▶ Six step algorithm
 1. Φ -insertion
 2. Renaming
 3. Down-safety computation
 4. WillBeAvail computation
 5. Finalization
 6. Code Motion

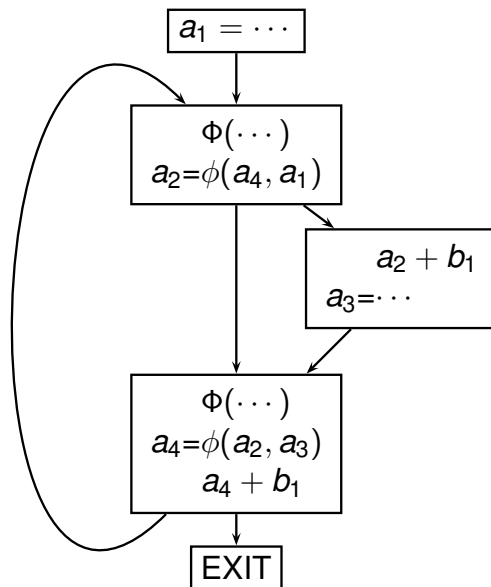
Running Example



Φ -insertion

- ▶ Φ for an expression E is required where two potentially different values of an expression merge
- ▶ At iterated dominance frontiers of occurrences of E
- ▶ At each block having a ϕ for some argument of E
 - ▶ Potential change in the expression's value

Φ -insertion



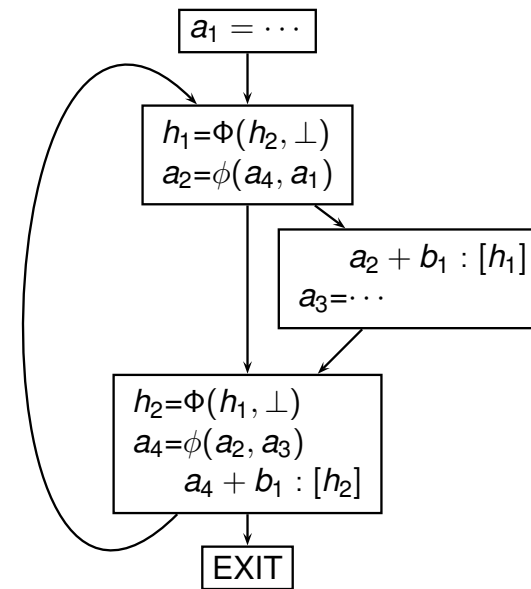
Rename

- ▶ Similar to SSA variable renaming
- ▶ Stack of every expression is maintained
- ▶ Three kinds of occurrences of E
 - ▶ Real occurrences (present in original program)
 - ▶ Results of Φ operators inserted
 - ▶ Operands of inserted Φ
- ▶ After renaming
 - ▶ Identical SSA instances of h represent identical values of E
 - ▶ A control flow path with two different instances of h has to cross either an assignment to an operand of E or a Φ of h

Rename Algorithm

- ▶ Runs with variable renaming
- ▶ When an E is encountered
 - ▶ if E is result of Φ , assign a new version to h and push it on E stack
 - ▶ if E is the real occurrence
 - ▶ for each operand, compare the version of operand with the top of the rename stack for operand
 - ▶ If all match, h gets same version as the top of E stack
 - ▶ If any mismatch, assign a new version to h and push it on E stack
- ▶ if E is operand of Φ , in the corresponding predecessor block
 - ▶ for each operand of E , compare the version of operand with the top of the rename stack for operand
 - ▶ If all match, h gets same version as the top of E stack
 - ▶ If any mismatch, replace E by \perp in the operand push it on E stack (WHY?)

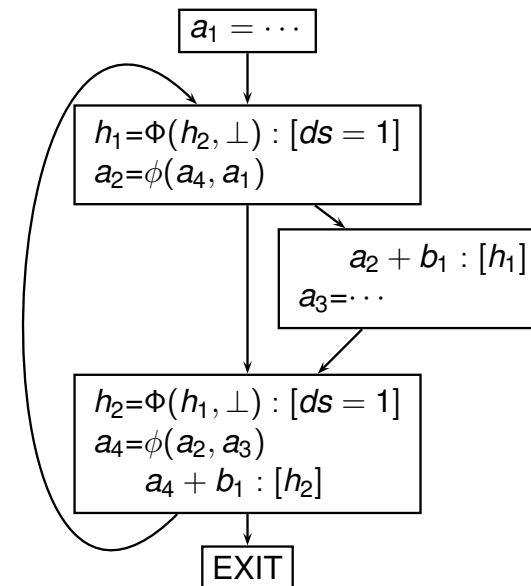
Rename



Down-safety

- ▶ Down-safety is same as very-busy (anticipability) property of expressions
 - ▶ Do not want to introduce new computation of E
- ▶ We only need to compute down-safety for inserted Φ -operators
- ▶ A Φ computation is **NOT** down-safe if
- ▶ there is a path to **EXIT** from Φ along which the result of Φ is
 - ▶ either not used
 - ▶ used only as an operand of another Φ that itself is **NOT** down-safe
- ▶ *HasRealUse*: Real occurrence of an expression

Down-safety (ds = ...)



WillBeAvail

- ▶ The set of Φ s where the expression must be available in any computationally optimal placement
 - ▶ Computation of *two forward* properties:
 - ▶ *CanBeAvail*: Φ s for which E is either available or anticipable or both
 - ▶ *Later*: Φ s beyond which insertion can not be postponed without introducing new redundancy
- $$WillBeAvail = CanBeAvail \wedge \neg Later$$

CanBeAvail

- ▶ Initialized to *true* for all Φ s
- ▶ Boundary Φ s:
 - ▶ Not Down-safe, and
 - ▶ At least one argument is \perp
- ▶ Set *false* for boundary Φ s
- ▶ Propagate *false* value along the chain of def-use to other Φ s
 - ▶ exclude edges along which *HasRealUse* is *true*

Later

- ▶ Determines latest (final) insertion points
- ▶ Initialize *Later* to *true* wherever *CanBeAvail* is *true*, otherwise *false*
- ▶ Assign *false* for Φ s with at least one operand with *HasRealUse* flag *true*
- ▶ Propagate *false* value forward to other Φ s
- ▶ *Later* \Rightarrow Φ s that are *CanBeAvail*, but do not reach any real occurrence of E

Insertion Points

- ▶ Insertions are done for Φ operands
- ▶ Along the corresponding predecessor edges
- ▶ Insertion done along i^{th} predecessor of Φ if *Insert* is *true*, i.e.
 - ▶ *WillBeAvail*(Φ) == *true*; AND
 - ▶ *Arg_i* is \perp ; OR
 - ▶ (*HasRealUse*(*Arg_i*) == *false*), AND
 - ▶ *Arg_i* is defined by Φ' with *WillBeAvail*(Φ') == *false*

Finalize

- ▶ Transforms the program with RCVs into a valid SSA form
- ▶ For every real occurrence of E , decide whether it is a *def* or a *use*
- ▶ For every ϕ with *WillBeAvail* being *true*, insert E along incoming edges with *Insert* being *true*
- ▶ For each ϕ for E
 - ▶ If *WillBeAvail* is *true*, it is replaced by SSA temporary with appropriate version (h_x)
 - ▶ If *WillBeAvail* is *false*, it is not part of SSA form, and is removed

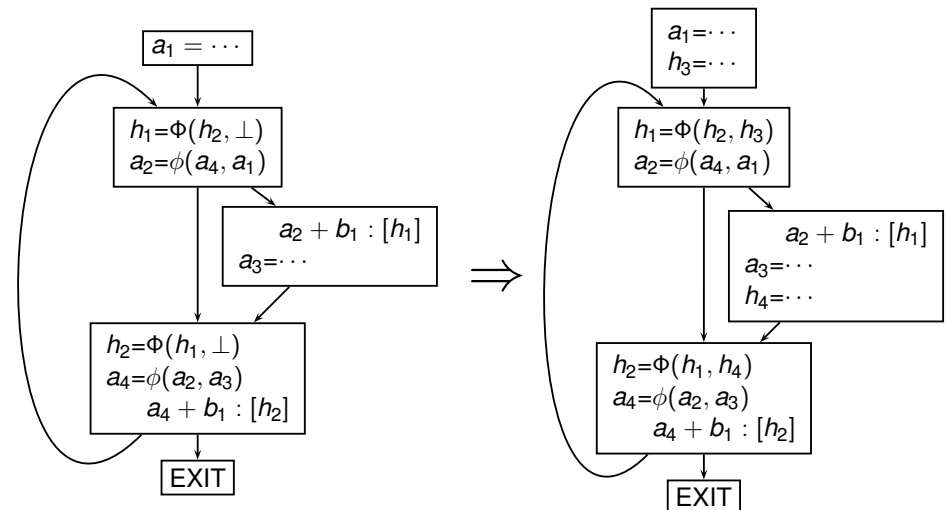
Finalize: AvailDef

- ▶ AvailDef: Table to mark def of expression occurrences
- ▶ Computed for each class (say h_x) of E
- ▶ Preorder traversal of dominator tree

AvailDef Computation

- ▶ Initialize: $\text{AvailDef}[x] = \perp \forall x$ (all classes of all expressions)
- ▶ During course of traversal, process occurrence x of E
 - ▶ ϕ occurrence:
 - ▶ If *WillBeAvail* is *false*, ignore.
 - ▶ Otherwise $\text{AvailDef}[x] = \text{this } \phi$ (we must be visiting x for first time) – WHY?
 - ▶ Real occurrence:
 - ▶ If $\text{AvailDef}[x]$ is \perp , mark this occurrence as def
 - ▶ Else, if $\text{AvailDef}[x]$ does not dominate this occurrence, mark this occurrence as def
 - ▶ Else, mark this occurrence as use of $\text{AvailDef}[x]$
 - ▶ ϕ operand (processed in predecessor block P)
 - ▶ If *WillBeAvail* of ϕ is *false*, ignore.
 - ▶ Else, if *Insert* is *true* for the operand, insert computation of E in block P , set it as a def, mark this occurrence as use of inserted.
 - ▶ Else (*Insert* is *false*), mark this occurrence as use of $\text{AvailDef}[x]$

Finalize



Code Motion

- ▶ For real *def* occurrence of E , compute E in a new version of temporary t
- ▶ For real *use* occurrence of E , replace E by current version of t
- ▶ For inserted occurrence of E , compute E in a new version of temporary t
- ▶ For a Φ occurrence, insert appropriate ϕ for t

Code Motion

