CS738: Advanced Compiler Optimizations SSA Continued

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Agenda

- Properties of SSA
- ► SSA to Executable
- SSA for Optimizations

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- In practice, worst case is rare.
- Practical complexity: O(R)

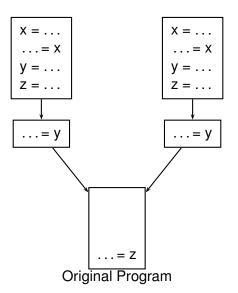
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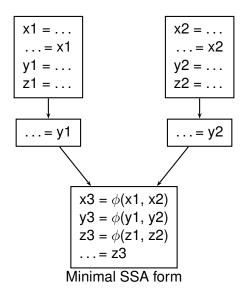
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- Linear time is achieved by careful ordering of nodes in the DJ-graph
- DF for a node is computed only once an reused later if required.

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- Semi-Pruned SSA

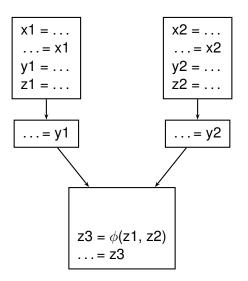
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- Requires global Live variable analysis

Variants of SSA Form: Pruned SSA Example



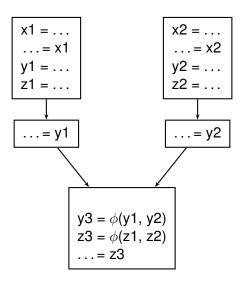
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- Non-locals can be computed without iteration or elimination

Variants of SSA Form: Semi-pruned SSA Example



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foreach block B {
```

```
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  defined = {}
```

```
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```

```
foreach block B {
  defined = {}
  foreach instruction V = X op Y {
    if X not in defined
```

```
foreach block B {
   defined = {}
   foreach instruction V = X \circ p y {
      if X not in defined
            non-locals = non-locals \cup \{X\}
```

Computing Non-locals

```
foreach block B {
  defined = {}
  foreach instruction V = X op Y {
    if X not in defined
       non-locals = non-locals U {X}
  if Y not in defined
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Computing Non-locals

```
foreach block B {
  defined = {}
  foreach instruction V = X op Y {
    if X not in defined
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  if Y not in defined
        non-locals = non-locals U {Y}
    defined = defined U {V}
  }
}
```

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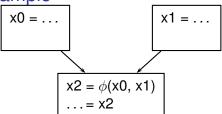
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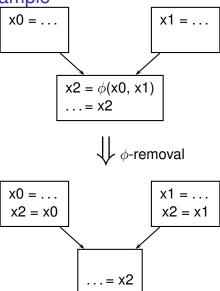
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 - ▶ Need to fix up the ϕ -function
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 - Simple algorithm
 - Works in most cases, but not always
 - Adds lots of copies
 - Many of them will be optimized by later passes

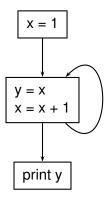






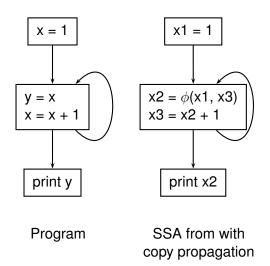


Lost Copy Problem

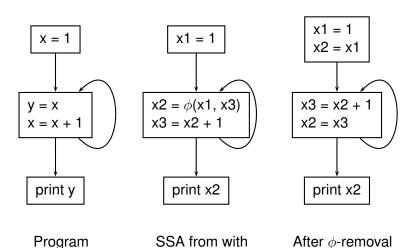


Program

Lost Copy Problem

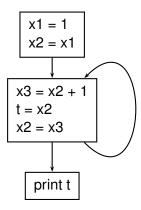


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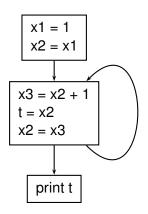
copy propagation

Lost Copy Problem: Solutions

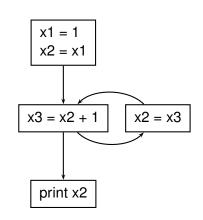


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Lost Copy Problem: Solutions

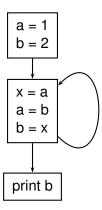


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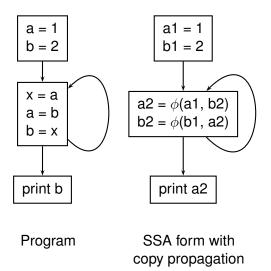
2. Critical Edge Split

Swap Problem

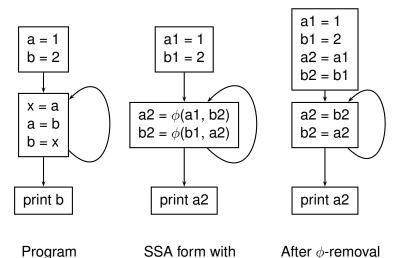


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