

Honors Single Variable Calculus

Apurva Nakade
JHU

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1 Proofs

We'll start by learning about the various kinds of proofs that we'll encounter in this class.¹

Proofs are the heart of mathematics. You must come to terms with proofs – you must be able to read, understand and write them. What is the secret? What magic do you need to know? The short answer is: there is no secret, no mystery, no magic. All that is needed is some common sense and a basic understanding of a few trusted and easy to understand techniques.

The Structure of a Proof

The basic structure of a proof is easy: it is just a series of statements, each one being either

- An assumption or
- A conclusion, clearly following from an assumption or previously proved result

And that is all. Occasionally there will be the clarifying remark, but this is just for the reader and has no logical bearing on the structure of the proof.

A well written proof will flow. That is, the reader should feel as though they are being taken on a ride that takes them directly and inevitably to the desired conclusion without any distractions about irrelevant details.

Each step should be clear or at least clearly justified. A good proof is easy to follow. When you are finished with a proof, apply the above simple test to every sentence: is it clearly

1. an assumption
2. a justified conclusion?

If the sentence fails the test, maybe it doesn't belong in the proof.

Example

In order to write proofs, you must be able to read proofs. See if you can follow the proof below. Don't worry about how you would have (or would not have) come up with the idea for the proof. Read the proof with an eye towards the criteria listed above. Is each sentence clearly an assumption or a conclusion? Does the proof flow? Was the theorem in fact proved?

Theorem 1.1. *The square root of 2 is an irrational number.*

Proof. Let's represent the square root of 2 by s . Then, by definition, s satisfies the equation

$$s^2 = 2.$$

If s were a rational number, then we could write $s = p/q$ where p and q are a pair of integers. In fact, by dividing out the common multiple if necessary, we may even assume p and q have no common multiple (other than 1). If we now substitute this into the first equation we obtain, after a little algebra, the equation

$$p^2 = 2q^2.$$

But now, 2 must appear in the prime factorization of the number p^2 (since it appears in the same number $2q^2$). Since 2 itself is a prime number, 2 must then appear in the prime factorization of the number p . But then, $2 \cdot 2$ would appear in the prime factorization of p^2 , and hence in $2q^2$. By dividing out a 2, it then appears that 2 is in the prime factorization of q^2 . Like before (with p^2) we can now conclude 2 is a prime factor of q . But now we have p and q sharing a prime factor, namely 2. This violates our assumption above (see if you can find it) that p and q have no common multiple other than 1. \square

¹This section is taken almost verbatim from <http://zimmer.csufresno.edu/~larryc/proofs/proofs.html>.

1.1 Direct Proofs

Most theorems that you want to prove are either explicitly or implicitly in the form

“If P , then Q ”.

This is the standard form of a theorem (though it can be disguised). A direct proof should be thought of as a flow of implications beginning with P and ending with Q .

$$P \implies \dots \implies Q$$

Most proofs are (and should be) direct proofs. Always try direct proof first, unless you have a good reason not to. If you find a simple proof, and you are convinced of its correctness, then don't be shy about. Many times proofs are simple and short.

Exercise. 1.2. Prove each of the following.

1. If a divides b and a divides c then a divides $b + c$, where a , b , and c are positive integers.
2. For all real numbers a and b , $a^2 + b^2 \geq 2ab$.
3. If a is a rational number and b is a rational number, then $a + b$ is a rational number.

1.2 Proof by Contradiction

In a proof by contradiction we assume, along with the hypotheses, the logical negation of the result we wish to prove, and then reach some kind of contradiction. That is, if we want to prove “If P , then Q ”, we assume P and Not Q . The contradiction we arrive at could be some conclusion contradicting one of our assumptions, or something obviously untrue like $1 = 0$. The proof of Theorem 1.1 is an example of this.

Exercise. 1.3. Use the method of Proof by Contradiction to prove each of the following.

1. The cube root of 2 is irrational.
2. If a is a rational number and b is an irrational number, then $a + b$ is an irrational number.
3. There are infinitely many prime numbers.¹

1.3 Proof by Contrapositive

Proof by contrapositive takes advantage of the logical equivalence between “ P implies Q ” and “Not Q implies Not P ”. For example, the assertion “If it is my car, then it is red” is equivalent to “If that car is not red, then it is not mine”. So, to prove “If P , then Q ” by the method of contrapositive means to prove

“If Not Q , then Not P ”.

How Is This Different From Proof by Contradiction? The difference between the Contrapositive method and the Contradiction method is subtle. Let's examine how the two methods work when trying to prove “If P , then Q ”.

Method of Contradiction: Assume P and Not Q and prove some sort of contradiction.

Method of Contrapositive: Assume Not Q and prove Not P .

The method of Contrapositive has the advantage that your goal is clear: Prove Not P . In the method of Contradiction, your goal is to prove a contradiction, but it is not always clear what the contradiction is going to be at the start.

¹There are dozens of proofs of this theorem, originally due to Euclid. Feel free to look one up online.

Exercise. 1.4. Use the method of Proof by Contrapositive to prove each of the following.

1. If the product of two integers is even, then at least one of the two must be even.
2. If the product of two integers is odd, then both must be odd.
3. If the product of two real numbers is an irrational number, then at least one of the two must be an irrational number.

1.4 Converse

The converse of an assertion in the form “If P , then Q ” is the assertion

“If Q , then P ”.

A common logical fallacy is to assume that if an assertion is true then so is its converse.

1.4.1 If and Only If

Many theorems are stated in the form “ P , if, and only if, Q ”. Another way to say the same thing is: “ Q is necessary, and sufficient for P ”. This means two things:

“If P , Then Q ” and “If Q , Then P ”.

So to prove an “If, and Only If” theorem, you must prove the theorem and also its *converse*.

Exercise. 1.5. Go back to the problems in Exercises 1.2, 1.3, 1.4 and find the ones which are of the form “If P , Then Q ”. For each of these:

- State the converse.
- Prove or disprove the converse (by providing either a proof or a counterexample).
- For the problems where the converse is also true rewrite the assertion as an “If, and Only If” statement.

1.5 Quantifiers

It is extremely important in mathematics to be able to formulate very precise statements. To prevent ambiguous statements and logical fallacies the vocabulary used is very limited and every term has a well-defined meaning. The two concepts we need to get used to are *logical operators* and *quantifiers*.

Logical Operators: Operators allow us to form complex statements by combining simpler ones. The basic logical operators are *and*, *or*, and *not*. Their usage is the same as in everyday language. The more complicated logical operator that we’ll be using a lot is *If - then -*. We’ve already encountered several examples of its usage in the previous sections.

Quantifiers: Quantifiers allow us to make abstract statements which are “universally true” without having to specify a concrete element. There are two commonly used quantifiers:

For all/every

There exists

Understanding and formulating complex statements using these requires a lot of practice. You’ll get used to these as the course progresses. Here are a few examples,

Example 1.6.

1. *For every* odd integer a , the integer $a + 1$ is even.
2. An integer a is even if and only if *there exists* an integer b such that $a = 2b$.
3. *There do not exist* integers p, q such that $\frac{p}{q} = \sqrt{2}$.
4. *For every* non-zero rational number x *there exists* a rational number y such that $x \cdot y = 1$.

1.5.1 Nesting Quantifiers

When dealing with statements involving multiple quantifiers the order really matters; changing the order changes the meaning of a statement completely. When you're using multiple quantifiers in a single sentence you should always pause and check if you're using the right quantifiers in the right order.

Exercise. 1.7. For each of the following pairs, explain how changing the order of the quantifiers changes the meaning of the statement.¹

1. (a) *For every even integer a there exists an integer b such that $a = 2b$.*
 (b) *There exists an integer b such that for every even integer a , $a = 2b$.*
2. (a) *For every positive integer n there exists a positive real number ϵ such that $\epsilon < 1/n$.*
 (b) *There exists a positive real number ϵ such that for every positive integer n , $\epsilon < 1/n$.*²
3. (a) *For every even integer a , for every odd integer b , $a + b$ is odd.*
 (b) *For every odd integer b , for every even integer a , $a + b$ is odd.*
4. Let $f : \mathbb{R} \rightarrow \mathbb{R}$ be a function and let ϵ be a positive real number.
 (a) *For every x , there exists a $\delta > 0$ such that for every y , if $|x - y| < \delta$ then $|f(x) - f(y)| < \epsilon$.*
 (b) *There exists a $\delta > 0$ such that for every x , for every y , if $|x - y| < \delta$ then $|f(x) - f(y)| < \epsilon$.*

1.5.2 Negating Nested Quantifiers

As we saw in Sections 1.2 and 1.3 we often need to negate mathematical statements.

In order to systematically negate complicated logical expressions, we observe that we can negate simple statements in the following manner. (These are called *De Morgan's laws*.)

Statement	Negation
For every x the statement P is true.	There exists an x such that the statement P is false.
There exists an x such that the statement P is true.	For every x the statement P is false.
P is true and Q is true	P is false or Q is false.
P is true or Q is true	P is false and Q is false.
If P then Q	P but not Q .

If we have multiple quantifiers in a statement then we start with the outermost quantifier and recursively move inwards.

Example 1.8. The negation of “For every x , there exists a y such that, the statement P is true” equals “There exists an x such that, for every y , the statement P is false.”

¹It's ok to a bit vague in your answers to this question.

²Such an ϵ is called an **infinitesimal**. There are number systems, for example the **hyperreal numbers**, which extend the real numbers by incorporating infinitesimals.

Exercise. 1.9. Negate each of the following.

1. *For every even integer a there exists an integer b such that $a = 2b$.*
2. *There exists an integer b such that for every even integer a , $a = 2b$.*
3. *For every positive integer n there exists a positive real number ϵ such that $\epsilon < 1/n$.*
4. *There exists a positive real number ϵ such that for every integer n , $\epsilon < 1/n$.*
5. *For every even integer a , for every odd integer b , $a + b$ is odd.*
6. *For every $\epsilon > 0$, for every x , there exists a $\delta > 0$ such that for every y , if $|x - y| < \delta$ then $|f(x) - f(y)| < \epsilon$.*
7. *For every $\epsilon > 0$, there exists a $\delta > 0$ such that for every x , for every y , if $|x - y| < \delta$ then $|f(x) - f(y)| < \epsilon$.*

Word of caution: When the word “and” is used in a list it is not considered a logical operator (this is a deficiency of the English language). For example, negating the following expression

For every even integer a and for every odd integer b , $a + b$ is odd,

does not change the *and* to an *or* (what happens if you do change the *and* to an *or*?).

2 Inequalities

In this section we'll learn how to solve inequalities and in the process figure out how to write proofs of statements involving inequalities.

The biggest different between equalities and inequalities is that, unlike equalities, there are usually infinitely many (if any) solutions to inequalities. When we're trying to solve an inequality we're not looking for the best solution, we're not even looking for a good solution, we're simply looking for **one** solution. For example, both $x = 1.001$ and $x = 10^{100}$ are equally valid solutions to the inequality $x > 1$.

Exercise. 2.1. Find a positive real solution to each of the following inequalities.

- | | |
|------------------------|-------------------------|
| 1. $x^3 + x < 1$ | 4. $x^3 + 2x < 1 + x^2$ |
| 2. $(1 + x)^2 - 1 < 1$ | 5. $x^{10} + x > 10$ |
| 3. $1 - (1 - x)^2 < 1$ | 6. $x^3 - x > 1$ |

When we try to solve equalities we try to simplify the equation until it becomes easy to solve. The same thing is true for inequalities, however, the ways to simplify an inequality are much more subtle.

Example 2.2. If we're trying to solve $x > P$, where P is some complicated expression, then it is sufficient to solve $x > P + 1$ as every solution of $x > P + 1$ is also a solution for $x > P$. Maybe the right simplification to do is to try and solve $x > P^2$ assuming that $P^2 > P$. If this does not work then we try something else, and so on. There is no *canonical way* in which an inequality can be solved. We try some which somehow simplify the inequality and hope that we get a solution.

For solving complicated inequalities, we need a systematic approach towards estimating functions. There are different tricks and identities for estimating different kinds of functions. For now, we'll focus mainly on estimating polynomials.

Exercise. 2.3. Prove the following extremely useful set of inequalities

$$\begin{array}{ll} \text{if } 1 \geq x > 0 & \text{then } 1 \geq x \geq x^2 \geq x^3 \geq \dots \\ \text{if } 1 \leq x & \text{then } 1 \leq x \leq x^2 \leq x^3 \leq \dots \end{array}$$

Word of Caution: You should be very careful when negative numbers are involved. Multiplying an inequality by a negative number changes its sign, for example, $2 < 3$ but $-2 > -3$. So,

$$\begin{array}{ll} \text{if } 1 \geq x > 0 & \text{then } -1 \leq -x \leq -x^2 \leq -x^3 \leq \dots \\ \text{if } 1 \leq x & \text{then } -1 \geq -x \geq -x^2 \geq -x^3 \geq \dots \end{array}$$

2.1 Proofs involving Inequalities

We'll encounter several statements in this class like: *For every $\epsilon > 0$ there exists a positive real number x such that $x^3 + 3x < \epsilon$.* When we want to prove such a statement, we are essentially trying to solve $x^3 + 3x < \epsilon$. (Why?) For example,

Q. 2.4. Prove that for every $\epsilon > 0$ there exists a positive real number x such that $x^3 + 3x < \epsilon$.

We're asking for a solution to the equation $x^3 + 3x < \epsilon$ for an arbitrary positive real number ϵ .

1. The terms x^3 and x have different degrees. We can try to use the inequalities in Exercise 2.3 to simplify their sum, for this we need to know if $x \geq 1$ or $x \leq 1$.
2. We're probably looking for a small x so we could assume $x \leq 1$. If this does not work we'll come back try something else. (Remember that we're trying to find **one** solution.)

3. As $x \leq 1$, $x^3 \leq x$ so that $x^3 + 3x \leq x + 3x = 4x$.
4. As $x^3 + 3x \leq 4x$ and we want to find a solution to $x^3 + 3x < \epsilon$, it suffices to solve $4x < \epsilon$. As ϵ is positive, **one** solution for this is $x = \epsilon/5$.
5. We still need to check if this solution actually works. We can do this by trying to write down a direct proof and making sure that all the implications are logically sound.
6. We required the condition $x \leq 1$ for this solution to work. To ensure this we'll set $x = \min\{1, \epsilon/5\}$.

Now we know what the solution is, but we still need to write down a proof for the original statement. A proof should always start with some assumptions and logically derive the required conclusion. For us the assumption will be $\epsilon > 0$ and $x = \min\{1, \epsilon/5\}$ and the expected conclusion is $x^3 + 3x < \epsilon$.

Proof of Q. 2.4. Let $\epsilon > 0$ and let $x = \min\{1, \epsilon/5\}$. Note that $x \leq 1$ and hence $x^3 \leq x$. Then,

$$\begin{aligned}
 x^3 + 3x &\leq x + 3x \\
 &= 4x \\
 &= 4 \min\{1, \epsilon/5\} \\
 &\leq 4\epsilon/5 \\
 &< \epsilon
 \end{aligned}
 \qquad \text{as } \epsilon \text{ is positive}$$

Hence, $x = \min\{1, \epsilon/5\}$ is a solution of $x^3 + 3x < \epsilon$, which proves the proposition. \square

First we had to solve the inequality then reverse the steps: start with the solution and write a *direct proof* for the proposition. This is how proofs are usually discovered. You make an “educated guess” and hope that it works. Sometimes it doesn’t, so you go back to finding another “educated guess”.

Remark 2.5 (A note on quantifiers). For proving the statement “*prove that for every $\epsilon > 0$ there exists a positive real number x such that $x^3 + 3x < \epsilon$* ” we had to start with the statement “*let $\epsilon > 0$ and $x = \min\{1, \epsilon/5\}$* ”. The variable x depends on the variable ϵ . This is fine because of the order of quantifiers: “*prove that for every $\epsilon > 0$ there exists a positive real number $x \dots$* ”; the quantifier for ϵ comes before the quantifier for x and hence the variable x can depend on the variable ϵ .

If instead, suppose we were trying to prove “*there exists a positive real number x such that for every $\epsilon > 0$, $x^3 + 3x < \epsilon$* ”. We’re still trying to solve the equation $x^3 + 3x < \epsilon$ but now the quantifier for x comes before the quantifier for ϵ and hence x cannot depend on the variable ϵ but instead should be a constant that universally solves the equation $x^3 + 3x < \epsilon$ for every $\epsilon > 0$. No such x exists (why?) and hence the statement “*there exists a positive real number x such that for every $\epsilon > 0$, $x^3 + 3x < \epsilon$* ” is false.

Exercise. 2.6. Prove that for every $\epsilon > 0$ there exists a positive real number x such that ...

1. $x^3 + x < \epsilon$.
2. $(1 + x)^2 - 1 < \epsilon$.
3. $1 - (1 - x)^2 < \epsilon$.
4. $x^3 + 2x < \epsilon + x^2$.
5. $x^{10} + x > \epsilon$.
6. $x^3 - x > \epsilon$.¹

Finally, we are sometimes required to find not one solution but a range of solutions. We try to come up an “educated guess” by the same method, and many a times this gives us the required range for free with a few changes in the final proof.

¹Hint: Break x^3 as $\frac{z}{x^3} + \frac{z}{x^3}$ and find the conditions on x for which $0 < \left(x - \frac{z}{x^3}\right)$.

Q. 2.7. Prove that for every $\epsilon > 0$, there exists a positive real number δ such that, for all x , if $0 < x < \delta$ then $x^3 + 3x < \epsilon$.

Proof. Let $\epsilon > 0$, let $\delta = \min\{1, \epsilon/5\}$ and let $0 < x < \delta$. Because $\delta \leq 1$, we have $x < 1$ and hence $x^3 < x$. As before,

$$\begin{aligned} x^3 + 3x &\leq x + 3x \\ &= 4x \\ &< 4\delta \\ &= 4\min\{1, \epsilon/5\} \\ &\leq 4\epsilon/5 \\ &< \epsilon \end{aligned} \quad \text{as } \epsilon \text{ is positive}$$

Hence, every $0 < x < \min\{1, \epsilon/5\}$ is a solution of $x^3 + 3x < \epsilon$, which proves the proposition. \square

Exercise. 2.8. Prove that for every $\epsilon > 0$, there exists a positive real number δ such that, for all x , if $0 < x < \delta$ then ...

- | | |
|---------------------------------|----------------------------------|
| 1. $x^3 + x < \epsilon$. | 3. $1 - (1 - x)^2 < \epsilon$. |
| 2. $(1 + x)^2 - 1 < \epsilon$. | 4. $x^3 + 2x < \epsilon + x^2$. |

Prove that for every $\epsilon > 0$, there exists a positive real number δ such that, for all x , if $x > \delta$ then

- | | |
|------------------------------|---------------------------|
| 1. $x^{10} + x > \epsilon$. | 2. $x^3 - x > \epsilon$. |
|------------------------------|---------------------------|

Exercise. 2.9. (Optional Challenge Problems)

1. Prove that for every $\epsilon > 0$, for every real number $x > 0$, there exists a $\delta > 0$ such that, for all real numbers $y > 0$, if $0 < y - x < \delta$ then $y^2 - x^2 < \epsilon$.
2. Prove that the following statement is false: For every $\epsilon > 0$, there exists a $\delta > 0$ such that, for every real number $x > 0$, for all real numbers $y > 0$, if $0 < y - x < \delta$ then $y^2 - x^2 < \epsilon$.