Rationalization

- Thomas Vikolaus

Space = sSet

Def": X -> Y rational and the managements

if the induced map $H_*(X; \mathbb{Q}) \longrightarrow H_*(X; \mathbb{Q})$ wis an isomorphism (or equivalently for H^*).

. ≥ is rational if X → ≥Y rational equivalence

⇒ [X,Z] → [Y,Z] is an equivalence.

(=> RMap (Y,Z) --> RMap(X,Z) equivalence)

· Z -> Za is called rationalization if it is an equivalence and Zo is rational.

Ex: «K(Q;n) is rational. for X > Y rational equivalence

[Y, K(Q; n] ~ H*(Y; Q)
∫ / Is
[x, K(Q; n)] ~ H*(X; Q)

· V = Q vector space

K(V; n) is rational.

. Homotopy limits of rational spaces are rational.

Z = holim Zi

Rmap (Y, Z) = holim & Rmap (Y, Zi)

1

5

Rmap (X,Z) = holim Rmap (X, Zi)

· Z simply-connected space s.t. Tx(Z) is are rational.

See this by induction on Postnikov tower:

Z = Rolim (T < Z)

 $\pi_*(\tau_{\leq n}Z) \cong \begin{cases} 0 & \text{if } x > n \\ \pi_*Z & \text{if } x < n \end{cases}$

n-even
hofib
$$gMuu_{49}[K(Q,n)]$$

is a forational homotopy equivalence.

· Ho(sSets) = this how a left adjoint (-) a

Full Subcategory
of rational spaces

(ssets) this how a left adjoint (-) a which is the rationalization functor.

This OI-connected space X is rational iff Tt X is rational.

② X -> y is between I-connected spaces is a rational equivalence > TI × X & Q -> TI × Y Q Q is an isomorphism

3 × 1-connected, 3 ×→× 2 with X-1-connected.

Proof: (2) Assume X -> Y induces iso on TT*(-) & Q.

Want to show rational eq.

homotopy fiber: Fthen the Serve SS gives $H^*(Y; H^*(F; \mathbb{R})) \Rightarrow H^*(X; \mathbb{R})$

enough to show $H_*(F;Q) = 0$, $* \ge 1$.
Use induction on Postnikov to were of F:

Hn (F; Q) = Hn (TsnF; Q)

Serve SS for

K(A;n) -> Tin = Tin = Tin-1 > F

A - torsion

=> Reduces the claim to show that

H*(K(Ajn); R) =0 *>n.

By the fibration $K(A, n-1) \rightarrow * \rightarrow K(A, n)$ Reduces to K(A, 1) and this $\stackrel{>}{=}$ can be Shown using group homology of A. lem: We construct for every 1- connected space X a map

 $X \longrightarrow X_{Q}$ such that (functionial in X)

i) Tx (xq) are rational

2) iso on TT*(-) ⊗ Q.

This proves (1), @ and (3).

Construct XQ by induction over Postnikov towers

if X=K(A'N)

XQ = K(A & Q,n)

F→X→Y for ~ suppose such a cliagram exists X→YQ (can be shown by using some cylinder object argument)

FQ = befiber of XQ -> \$YQ

Then F -> For its a satisfies the by conditions in the lemma.

. Apply this to and use $X = holim \ T \leq n \times X$ $X := holim \ (T \leq n \times)_{\mathfrak{A}}$ $T \leq n \times X$

Corollary: $\pi_{\nu}(s^n)\otimes\mathbb{R}=$ * = 2n-1 if n-even

Thm (Bousfield)

Every space admits a rationalization. More precisely there is a Quillen model structure · Cofib = mono

- · co.e. = rational equivalences
- . fib objects = rational Kan complexes.
- . IF M is a model category then the a left Bousfield localization is a model category M' with same underlying category and same cofib but more weak equivalences. (Wm & Wm')
- · Ho (M') Ho (M)

This functor is fully faithful. Rid has a right adjoint Lid this is like the the thought sublacement. · Wm, 2 Wm , Fib(M) & Fib(M)

6this determines what the rufib objects are.
But it is hard to prove that this is indeed a mudel structure.

"Sketch of broof":

- . Fibrations have RLP w. r.t acyclic cofib.
- . Can use Kan-Quillen model shucture to derive most of factorization thms.

= X -> X need to construct this factorization

Idea: · Use small object argument for X >>> Y

reational equivalence and cofib.

- · restrict to mono S= { x >>> y with > countable }
- . This gives us XIN, FIRLP with respect to S

Key lemma: Every rational trivial cofib X -> Y is a transfirulte compositions of pushouts of major in S.

· Given any map A -> B want a lift

$$\begin{array}{c} X \longrightarrow A \\ \downarrow \\ Y \longrightarrow B \end{array}$$

Replace B by bushowt P

 $\begin{array}{ccc}
\times & \longrightarrow & A \\
\downarrow & & \downarrow \\
Y & \longrightarrow & P & \longrightarrow & B
\end{array}$

- Use colimit of all such P's say Poo and respect. This is terms
 - Then repeat using bansfinite induction.
 - . By small object argument this process terminates.