

# Decision Theory

## Lecture 8

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## Organization

- decision theory only for Data Science specialization "optimization"
- evaluation

## Recap: What did we do?

- Decision under uncertainty
- Methods:
  - Minimax/Maximin
  - Minimax Regret
  - Hurwicz
  - Average
- Properties

## Decision Under Uncertainty

	$s_1$	$s_2$	$\dots$	$s_n$
$a_1$	$e_{11}$	$e_{12}$	$\dots$	$e_{1n}$
$a_2$	$e_{21}$	$e_{22}$	$\dots$	$e_{2n}$
$\vdots$	$\vdots$	$\vdots$	$\vdots$	$\vdots$
$a_m$	$e_{m1}$	$e_{m2}$	$\dots$	$e_{mn}$

- Any of the  $s_1, \dots, s_n$  might happen
- We do not know any likelihoods
- Which  $a_i$  to choose?

## Approach 1: Maximin/Minimax

- Maximisation problem: maximin
  - = maximise the smallest (worst) value
- Minimisation problem: minimax
  - = minimise the largest (worst) value
- Evaluate a row (decision) by its worst value
- Finds the best worst-case guarantee

## Approach 2: Minimax Regret

- For each column (scenario) find the best value
- Consider the difference to the best in each column
- Choose alternative that minimizes the largest difference

## Example

Maximise:

	$s_1$	$s_2$
$a_1$	3	7
$a_2$	5	2

Regret:

	$s_1$	$s_2$
$a_1$	2	0
$a_2$	0	5

- Max regret of  $a_1$  is 2
- Max regret of  $a_2$  is 5
- Choose  $a_1$

## Approach 3: Hurwicz's Pessimism-Optimism Index

- Given trade-off value  $\alpha \in [0, 1]$
- For each alternative  $i$ :
  - Determine the worst outcome  $m_i$
  - Determine the best outcome  $M_i$
  - Calculate  $\alpha m_i + (1 - \alpha)M_i$
- Compromise between best and worst
- $\alpha = 1$ : optimise worst-case
- $\alpha = 0$ : optimise best-case

## Approach 4: Laplace's Principle of Insufficient Reason (Average)

- Assumption:
  - If we don't know anything about probabilities, just assume all options are equally likely
- Optimise with respect to average value



## Example

Maximise:

	$s_1$	$s_2$	$s_3$	min	max	Hurwicz	regret	average
$a_1$	3	3	4					
$a_2$	2	3	1					
$a_3$	6	2	2					
$a_4$	5	2	3					

- Maximin:
- Hurwicz  $\alpha = 1/2$ :
- Regret:
- Average:

## Example

Maximise:

	$s_1$	$s_2$	$s_3$	min	max	Hurwicz	regret	average
$a_1$	3	3	4	3	4	$7/2$	3	$10/3$
$a_2$	2	3	1	1	3	$4/2$	4	$6/3$
$a_3$	6	2	2	2	6	$8/2$	2	$10/3$
$a_4$	5	2	3	2	5	$7/2$	1	$10/3$
max	6	3	4					

- Maximin:  $a_1$
- Hurwicz  $\alpha = 1/2$ :  $a_3$
- Regret:  $a_4$
- Average:  $a_1$ ,  $a_3$  or  $a_4$

## Axioms

- Axiom 1: Complete ranking
- Axiom 2: Independence of labelling
- Axiom 3: Independence of value scale
- Axiom 4: Strong domination
- Axiom 5: Independence of irrelevant alternatives
- Axiom 6: Independence of addition of a constant to a column
- Axiom 7: Independence of row permutation
- Axiom 8: Independence of column duplication

## Table of Results

	Minimax	Hurwicz	Regret	Average
Ax 1	✓	✓	✓	✓
Ax 2	✓	✓	✓	✓
Ax 3	✓	✓	✓	✓
Ax 4	✓	✓	✓	✓
Ax 5	✓	✓	✗	✓
Ax 6	✗	✗	✓	✓
Ax 7	✓	✓	✗	✓
Ax 8	✓	✓	✓	✗

## No Free Lunch

- Any criterion that fulfils
  - Axiom 1: Complete ranking
  - Axiom 4: Strong domination
  - Axiom 5: Independence of irrelevant alternatives
  - Axiom 6: Independence of addition of a constant to a column
  - Axiom 7: Independence of row permutationis equivalent to the average criterion
- The average criterion does not fulfil Axiom 8
- There is no criterion that fulfils all axioms

## Warm-Up

Find a counterexample to axiom 6: "adding a constant to a column does not change the solution"

- Minimax does not satisfy axiom 6
- Hurwicz does not satisfy axiom 6

## Warm-Up

Minimize:

	$s_1$	$s_2$
$a_1$	1	1
$a_2$	2	0

- Minimax: choose  $a_1$
- The same for Hurwicz  $\alpha = 1$

	$s_1$	$s_2$
$a_1$	1	3
$a_2$	2	2

- Minimax: choose  $a_2$
- The same for Hurwicz  $\alpha = 1$

## Today

- Last session: Decision under Uncertainty
- Ordered Weighted Averaging (OWA)



## OWA

- Let a weight vector  $w$  be given

$$w = (w_1, \dots, w_n)$$

with  $w_j \in [0, 1]$ ,  $\sum_{j \in [n]} w_j = 1$

- For  $A = (a_1, \dots, a_n)$ , let the vector  $B = (b_1, \dots, b_n)$  be all entries from  $a$  ordered by size (from large to small)
- A function  $F : \mathbb{R}^n \rightarrow \mathbb{R}$  is called an Ordered Weighted Averaging (OWA) operator if

$$F(a_1, \dots, a_n) = w_1 b_1 + w_2 b_2 + \dots + w_n b_n$$

- Weights  $w_i$  are bound to the **position**, not the scenario

## Example

- Let  $w = (0.2, 0.3, 0.1, 0.4)$
- Find  $F(6, 10, 3, 5)$ 
  - Ordered:  $B = (10, 6, 5, 3)$

$$\begin{aligned} F(6, 10, 3, 5) \\ &= 0.2 \cdot 10 + 0.3 \cdot 6 + 0.1 \cdot 5 + 0.4 \cdot 3 \\ &= 5.5 \end{aligned}$$

- Find  $F(0, 7, 10, 2)$ 
  - Ordered:  $B = (10, 7, 2, 0)$

$$\begin{aligned} F(0, 7, 10, 2) \\ &= 0.2 \cdot 10 + 0.3 \cdot 7 + 0.1 \cdot 2 + 0.4 \cdot 0 \\ &= 4.3 \end{aligned}$$

## Example

Maximize. Calculate OWA for each alternative with  $w = (1/2, 1/4, 0, 1/4)$

	$s_1$	$s_2$	$s_3$	$s_4$
$a_1$	6	3	5	2
$a_2$	4	8	2	0
$a_3$	3	2	5	6

- $D(a_1)$
- $D(a_2)$
- $D(a_3)$

## Example

Maximize. Calculate OWA for each alternative with  $w = (1/2, 1/4, 0, 1/4)$

	$s_1$	$s_2$	$s_3$	$s_4$
$a_1$	6	3	5	2
$a_2$	4	8	2	0
$a_3$	3	2	5	6

- $D(a_1) = F(6, 3, 5, 2) = 0.5 \cdot 6 + 0.25 \cdot 5 + 0.25 \cdot 2 = 4.75$
- $D(a_2) = F(4, 8, 2, 0) = 0.5 \cdot 8 + 0.25 \cdot 4 + 0.25 \cdot 0 = 5$
- $D(a_3) = F(3, 2, 5, 6) = 0.5 \cdot 6 + 0.25 \cdot 5 + 0.25 \cdot 2 = 4.75$

## Properties

- $0 \leq F(A) \leq 1$  if all values are in  $[0, 1]$
- Why?

$$F(A) = \sum_{i \in [n]} w_i b_i \leq \sum_{i \in [n]} w_i b_1 = b_1 \leq 1$$

$$F(A) = \sum_{i \in [n]} w_i b_i \geq \sum_{i \in [n]} w_i b_n = b_n \geq 0$$

## Properties

- Let  $A = (a_1, \dots, a_n)$
- Set  $A' = (a'_1, \dots, a'_n)$  with
$$a_j \geq a'_j \text{ for all } j \in [n]$$
- Then  $F(A) \geq F(A')$
- If each individual value improves, the result improves

## Why?

- Let  $B$  and  $B'$  be the corresponding sorted vectors
- Then  $b_j \geq b'_j$
- Thus,  $\sum_{i \in [n]} w_i b_i \geq \sum_{i \in [n]} w_i b'_i$

## Properties

- $F(a_1, \dots, a_n) = F(a_{\pi(1)}, \dots, a_{\pi(n)})$  for all permutations  $\pi$
- Thus: order of scenarios doesn't matter

## Why?

- According to the definition of OWA, values are sorted anyway

## Remark

- If we only assign weights  $w_j$  to each scenario  $j$  (not to positions), the property does not hold
- So:

$$G(a_1, \dots, a_n) = \sum_{j \in [n]} w_j a_j$$

does not satisfy independence of column ordering

- Example:  $G(a_1, a_2) = 0.7a_1 + 0.3a_2$
- $G(1, 0) \neq G(0, 1)$



## Special Cases

- Reminder:  $w_1$  is the weight for the highest outcome,  $w_n$  for the lowest outcome
- What if ... (how to maximize)
  - $w_j = 1/n$  for all  $j$ ?
  - $w_1 = 1$ ,  $w_j = 0$  for all  $j \neq 1$ ?
  - $w_n = 1$ ,  $w_j = 0$  for all  $j \neq n$ ?
  - $w_n = \alpha$ ,  $w_1 = 1 - \alpha$ ,  $w_j = 0$  for all  $j \neq 1, n$ ?

## Special Cases

- Reminder:  $w_1$  is the weight for the highest outcome,  $w_n$  for the lowest outcome
- What if ... (how to maximize)
  - $w_j = 1/n$  for all  $j$ ? Average
  - $w_1 = 1$ ,  $w_j = 0$  for all  $j \neq 1$ ? Best case
  - $w_n = 1$ ,  $w_j = 0$  for all  $j \neq n$ ? Worst case
  - $w_n = \alpha$ ,  $w_1 = 1 - \alpha$ ,  $w_j = 0$  for all  $j \neq 1, n$ ? Hurwicz
- OWA generalizes Average, Minimax, Hurwicz

## Properties

- Let  $\underline{w} = (0, 0, \dots, 0, 1)$  and  $\overline{w} = (1, 0, 0, \dots, 0)$
- Let  $F_w(a_1, \dots, a_n)$  be the OWA operator with respect to weight  $w$
- Then:
  - $F_{\underline{w}}(a_1, \dots, a_n) = \min_{j \in [n]} a_j$
  - $F_{\overline{w}}(a_1, \dots, a_n) = \max_{j \in [n]} a_j$
  - $F_{\underline{w}}(a_1, \dots, a_n) \leq F_w(a_1, \dots, a_n) \leq F_{\overline{w}}(a_1, \dots, a_n)$  for any  $w$

## How do I determine $w$ ? First Approach

- Let  $S_k = \sum_{j=1}^k w_j$
- $S_n = 1$
- $S_0 := 0$
- Assuming we have a sorted vector  $B$  such that  $b_j = 1$  for  $j \leq k$ , and  $b_j = 0$  otherwise
- Inquiry: How good do you find it if the best outcome is achieved in  $k$  out of  $n$  scenarios, and the remaining  $n - k$  are worst possible?
- It is  $F(B) = \sum_{j \in [n]} w_j b_j = \sum_{j=1}^k w_j = S_k$

## How do I determine $w$ ? First Approach

- $S_k$  corresponds to the utility when  $k$  out of  $n$  scenarios are good
- $w_j = 1/n$ : linear utility gain
- $\underline{w}$ : no satisfaction unless all scenarios are good
- $\overline{w}$ : satisfaction even if one scenario is already good
- If  $S_k$  can be determined, then set

$$w_j = S_j - S_{j-1}$$

## How do I determine $w$ ? Second Approach

- Given observations on how someone has decided
- Based on that, estimate weights
- Let  $\mathcal{B} = (B^1, \dots, B^K)$  with  $B^k = (b_1^k, \dots, b_n^k)$  be some alternatives we have observed so far
- Let  $R \subseteq \mathcal{B} \times \mathcal{B}$  be pairs of alternatives where the decision-maker prefers the first alternative to the second alternative

## How do I determine $w$ ? Second Approach

- OWA of  $B^k$  is  $\sum_{j=1}^n w_j b_j^k$
- Problem: find  $w$  such that  $F_w(B^i) \geq F_w(B^k)$  for all  $(B^i, B^k) \in R$
- Might not be possible  $\rightarrow$  Optimization problem

max  $z$

$$\text{s.t. } \sum_{j=1}^n w_j b_j^i - \sum_{j=1}^n w_j b_j^k \geq z$$

$$\forall (B^i, B^k) \in R$$

$$\sum_{j=1}^n w_j = 1$$

$$w_j \geq 0$$

$$\forall j \in [n]$$

## "Andness" and "Orness"

- OWA can also be considered as a mixture of "And" (andness) and "Or" (orness)
- $w_n = 1$ : each scenario must be good – "and"
- $w_1 = 1$ : one scenario must be good – "or"
- Define "orness" as

$$orness(w) = \frac{1}{n-1} \sum_{j \in [n]} w_j (n-j)$$

- $andness(w) = 1 - orness(w)$



## "Andness" and "Orness"

$$\text{orness}(w) = \frac{1}{n-1} \sum_{j \in [n]} w_j (n-j)$$

- $\underline{w} = (0, 0, \dots, 0, 1),$

$$\text{orness}(\underline{w}) = \frac{1}{n-1} \cdot 0 = 0$$

- $\overline{w} = (1, 0, \dots, 0)$

$$\text{orness}(\overline{w}) = \frac{1}{n-1} \cdot (n-1) = 1$$

## "Andness" and "Orness"

$$\text{orness}(w) = \frac{1}{n-1} \sum_{j \in [n]} w_j (n-j)$$

- $w = (1/n, 1/n, \dots, 1/n)$ ,

$$\begin{aligned} \text{orness}(w) &= \frac{1}{n-1} \cdot \frac{1}{n} \cdot \sum_{j \in [n]} (n-j) \\ &= \frac{1}{n-1} \cdot \frac{1}{n} \cdot \frac{n(n-1)}{2} = \frac{1}{2} \end{aligned}$$

## "Andness" and "Orness"

$$\text{orness}(w) = \frac{1}{n-1} \sum_{j \in [n]} w_j (n-j)$$

- What does this weight vector do?

$$w = (0, 0, 1, 0, 0)$$

## "Andness" and "Orness"

$$orness(w) = \frac{1}{n-1} \sum_{j \in [n]} w_j (n-j)$$

- What does this weight vector do?

$$w = (0, 0, 1, 0, 0)$$

- Chooses the median (the moderately good scenario)
- What is its *orness*?

## "Andness" and "Orness"

$$orness(w) = \frac{1}{n-1} \sum_{j \in [n]} w_j (n-j)$$

- What does this weight vector do?

$$w = (0, 0, 1, 0, 0)$$

- Chooses the median (the moderately good scenario)
- What is its *orness*?

$$orness(w) = \frac{1}{4} \cdot (5 - 3) = \frac{1}{2}$$

## OWA for Combinatorial Problems

- So far: OWA for a single alternative
- Approach possible for decision matrix and combinatorial problems
- So far only maximization problems, also possible for minimization problems
- If maximizing:  $w = (1, 0, \dots, 0)$  chooses the largest – best case
- If minimizing:  $w = (1, 0, \dots, 0)$  chooses the smallest – worst case

## OWA for Combinatorial Problems

- Calculate minimal OWA solution with  $w = (1/2, 1/2, 0, 0)$ , choose 2 out of 3 items

item	$c^1$	$c^2$	$c^3$	$c^4$
1	3	4	3	5
2	4	2	5	1
3	4	2	4	5

- $\{1, 2\}$ :
- $\{2, 3\}$ :
- $\{1, 3\}$ :

## OWA for Combinatorial Problems

- Calculate minimal OWA solution with  $w = (1/2, 1/2, 0, 0)$ , choose 2 out of 3 items

item	$c^1$	$c^2$	$c^3$	$c^4$
1	3	4	3	5
2	4	2	5	1
3	4	2	4	5

- $\{1, 2\}$ :  $A = (7, 6, 8, 6)$ ,  $OWA(\{1, 2\}) = 7.5$
- $\{2, 3\}$ :  $A = (8, 4, 9, 6)$ ,  $OWA(\{2, 3\}) = 8.5$
- $\{1, 3\}$ :  $A = (7, 6, 7, 10)$ ,  $OWA(\{1, 3\}) = 8.5$



## OWA for Combinatorial Problems

- Calculate minimal OWA solution with  $w = (2/3, 1/3, 0, 0)$ , choose 2 out of 3 items

item	$c^1$	$c^2$	$c^3$	$c^4$
1	3	7	2	4
2	2	1	8	3
3	3	4	7	1

- $\{1, 2\}$ :
- $\{2, 3\}$ :
- $\{1, 3\}$ :

## OWA for Combinatorial Problems

- Calculate minimal OWA solution with  $w = (2/3, 1/3, 0, 0)$ , choose 2 out of 3 items

item	$c^1$	$c^2$	$c^3$	$c^4$
1	3	7	2	4
2	2	1	8	3
3	3	4	7	1

- $\{1, 2\}$ :  $A = (5, 8, 10, 7)$ ,  $OWA(\{1, 2\}) = 9.3$
- $\{2, 3\}$ :  $A = (5, 5, 15, 4)$ ,  $OWA(\{2, 3\}) = 11.7$
- $\{1, 3\}$ :  $A = (6, 11, 9, 5)$ ,  $OWA(\{1, 3\}) = 10.3$

## Approximation

- If minimizing, let the weights  $w_i$  be non-decreasing:

$$w_1 \geq w_2 \geq \dots \geq w_n$$

- i.e., we are conservative / risk-averse: worse outcomes are always more important than better ones
- Worst-case with  $w_1 = 1$  is a special case
- For item  $i$ , let  $\hat{c}_{i1} \geq \hat{c}_{i2} \geq \dots \hat{c}_{in}$  be the sorted cost vector over all scenarios
- Set  $\hat{c}_i = \sum_{j \in [n]} w_j \hat{c}_{ij}$

## Approximation

- Solve with respect to scenario  $\hat{c}$
- The solution is a  $w_1 \cdot n$ -approximation for OWA
- Meaning, the objective function value is at most a factor of  $w_1 n$  larger than the optimum
- Worst-case:  $n$ -approximation

## Example

- Calculate minimal OWA solution with  $w = (2/3, 1/3, 0, 0)$ , choose 2 out of 3 items

item	$c^1$	$c^2$	$c^3$	$c^4$	$\hat{c}$
1	3	7	2	4	6.0
2	2	1	8	3	6.3
3	3	4	7	1	6.0

- $\hat{c}_1 = w_1 c_1^2 + w_2 c_1^4 + w_3 c_1^1 + w_4 c_1^3 = (2/3) \cdot 7 + (1/3) \cdot 4 = 6$
- Heuristic: include  $\{1, 3\}$  (unfortunately not optimal)
- Approximation guarantee:  $w_1 n = 2.7$

## Extension: WOWA

- We have seen so far: OWA generalizes Maximin, Hurwicz, Average
- A weighting of scenarios of the form

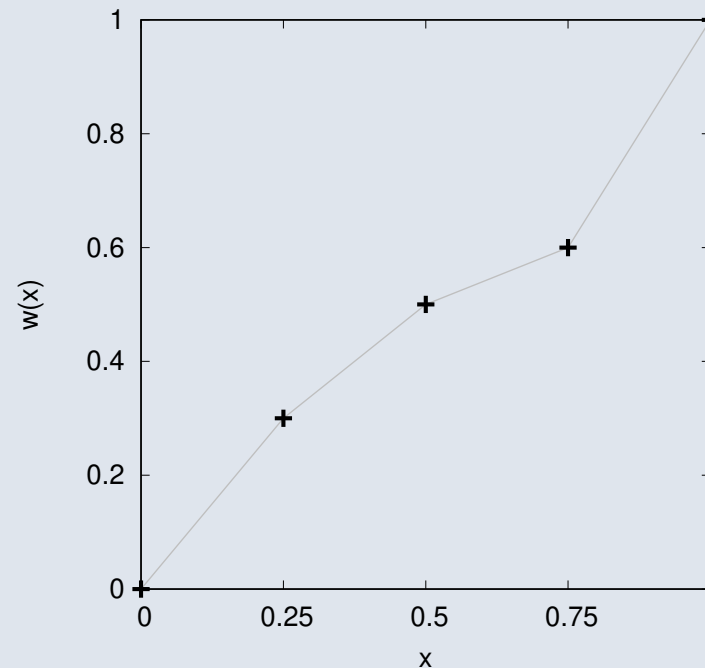
$$F(a_1, \dots, a_n) = p_1 a_1 + p_2 a_2 + \dots + p_n a_n$$

is generally not possible

- Reminder: weights are bound to position and not to scenarios
- There is another generalization of OWA, called the Weighted Ordered Weighted Averaging (WOWA) Operator

## WOWA: Definition

- Weight vector  $w = (w_1, w_2, \dots, w_n)$
- $w(x)$  is an interpolating function of the points  $(i/n, \sum_{j \leq i} w_j)$
- Example:  $w = (0.3, 0.2, 0.1, 0.4)$



## WOWA: Definition

- Given:
  - $p_i \in [0, 1]$ ,  $\sum_{i \in [n]} p_i = 1$
  - $w_i \in [0, 1]$ ,  $\sum_{i \in [n]} w_i = 1$

- Define

$$WOWA(a_1, \dots, a_n) = \sum_{i \in [n]} \omega_i a_{\pi(i)}$$

with

- $\pi$  a permutation such that  $A$  is sorted in ascending order
- $\omega_i = w(\sum_{j \leq i} p_{\pi(j)}) - w(\sum_{j < i} p_{\pi(j)})$
- $w(x)$  is the interpolating function



## WOWA: Properties

- Let  $p_i = 1/n$  for all  $i \in [n]$
- Intuitively: all scenarios are equally important
- Then:

$$\begin{aligned}\omega_i &= w \left( \sum_{j \leq i} p_{\pi(j)} \right) - w \left( \sum_{j < i} p_{\pi(j)} \right) \\ &= w(i/n) - w((i-1)/n) = w_i\end{aligned}$$

- WOWA corresponds to the normal OWA

## WOWA: Properties

- Let  $w_i = 1/n$  for all  $i \in [n]$
- Intuitively: all positions are equally important
- Then:

$$\begin{aligned}\omega_i &= w \left( \sum_{j \leq i} p_{\pi(j)} \right) - w \left( \sum_{j < i} p_{\pi(j)} \right) \\ &= \sum_{j \leq i} p_{\pi(j)} - \sum_{j < i} p_{\pi(j)} = p_{\pi(i)}\end{aligned}$$

- And WOWA corresponds to the weighting  $\sum_{i \in [n]} p_i a_i$

## Quiz

### Question 1

True or false?

OWA is a generalization of Maximin, Regret, Hurwicz, and Average.

### Question 2

For  $A = (4, 8, 2, 6, 3, 0)$ , what is OWA with  $w = (0, 0, 1/2, 1/2, 0, 0)$ ?

## Quiz

### Question 1

True or false?

OWA is a generalization of Maximin, Regret, Hurwicz, and Average.

### Solution

False: it is not a generalization of Regret.

### Question 2

For  $A = (4, 8, 2, 6, 3, 0)$ , what is OWA with  $w = (0, 0, 1/2, 1/2, 0, 0)$ ?

### Solution

It is the average of the two middle outcomes.

$$\begin{aligned} B &= (8, 6, 4, 3, 2, 0) \\ F(A) &= (1/2) \cdot 4 + (1/2) \cdot 3 \\ &= 3.5 \end{aligned}$$