

Decision Theory

Lecture 11

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Today

- Game theory

What is it about?

- So far:
 - Decision under certainty = no opponent
 - Decision under uncertainty = "against nature"
- Game theory: consciously acting opponent
- Examples:
 - Wage negotiations
 - Auctions
 - Competitive situations, e.g., in an oligopoly
 - Cold War

Example: Prisoner's Dilemma (see Lecture 1)

- You and your accomplice stole apples and got arrested
- You are in different cells
- Confess or not?
 - You confess
 - ▶ Accomplice confesses: 4 years in prison for both
 - ▶ Accomplice does not confess: 1 year in prison for you, 6 years for him
 - You do not confess
 - ▶ Accomplice confesses: 6 years in prison for you, 1 year for him
 - ▶ Accomplice does not confess: 2 years in prison for both

Example: Prisoner's Dilemma (see Lecture 1)

- Decision matrix:

	He confesses	He does not confess
I confess	(4,4)	(1,6)
I do not confess	(6,1)	(2,2)

- First value: my outcome
- Second value: accomplice's outcome
- Minimize prison years
- What should I do? Decision via survey.

Example, continued

	He confesses	He does not confess
I confess	(4,4)	(1,6)
I do not confess	(6,1)	(2,2)

- From the perspective of game theory: choose the best strategy assuming that everyone follows that strategy
- So, I do not confess?
- "Optimality" for game theory usually: "Nash Equilibria," more on that later

Distinctions

- Number of players:
 - Two-person games
 - N-person games ($N > 2$) with or without the possibility of coalition formation
- Rules of the game:
 - Finite games (number of repetitions is fixed)
 - Infinite games (number of repetitions is not fixed)
- Strategy space:
 - Discrete strategy space
 - Continuous strategy space

Distinctions, continued

- Strategy choice:
 - Pure strategy: Consistent behavior
 - Mixed strategy: Choosing a strategy using a random mechanism
- Timing of strategy choice:
 - Simultaneous, i.e., players have no knowledge (only assumptions) about the opponent's strategy choice
 - Sequential, i.e., players know the previous move or moves of the opponent

Distinctions, continued

- Coordination between players:
 - Cooperative games (agreements between players)
 - Non-cooperative games (no agreements between players)
- Outcomes:
 - Games with variable gains
 - Constant-sum games
 - Zero-sum games

Terms and Representations

- Strategy:
 - Plan that contains, for each information available to player i at the time of executing a move, a (conditional) instruction on how to execute the move.
- Representation of games in normal form (payoff matrix):

		Player 2		
		Strategy b_1	\dots	Strategy b_n
Player 1	Strategy a_1	(u_{11}^1, u_{11}^2)	\dots	(u_{1n}^1, u_{1n}^2)
	\vdots	\vdots		\vdots
	Strategy a_m	(u_{m1}^1, u_{m1}^2)	\dots	(u_{mn}^1, u_{mn}^2)

- Representation as a game tree is also possible

Examples

		Player 2		
		Strategy b_1	\dots	Strategy b_n
Player 1	Strategy a_1	(u_{11}^1, u_{11}^2)	\dots	(u_{1n}^1, u_{1n}^2)
	\vdots	\vdots		\vdots
	Strategy a_m	(u_{m1}^1, u_{m1}^2)	\dots	(u_{mn}^1, u_{mn}^2)

	b_1	b_2	b_3
a_1	(10, 10)	(4, 4)	(0, 18)
a_2	(8, 7)	(8, -6)	(7, 8)

Examples

		Player 2		
		Strategy b_1	\dots	Strategy b_n
Player 1	Strategy a_1	(u_{11}^1, u_{11}^2)	\dots	(u_{1n}^1, u_{1n}^2)
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Task!

Write a payoff matrix for Rock-Paper-Scissors

Strong and Weak Dominance

- Strong Dominance:
 - Always better outcomes regardless of the opponent's strategy.
- Weak Dominance:
 - At least as good and better in one aspect of the outcome.

Strong and Weak Dominance

	b_1	b_2	b_3
a_1	(10, 10)	(4, 4)	(0, 18)
a_2	(8, 7)	(8, -6)	(7, 8)

- Strong dominance: b_3 dominates b_1 and b_2
- Assumption: No one chooses strongly dominated strategies
- Then, player 1 will choose strategy a_2
- "Iterated Dominance": Repeatedly applying dominance checking

Example of Iterated Dominance

	b_1	b_2	b_3
a_1	(3, 6)	(19, 27)	(15, 15)
a_2	(18, 16)	(13, 28)	(4, 8)
a_3	(30, 5)	(13, 15)	(10, 26)

Nash Equilibrium: Strategy Profile

- Assuming there are N players
- Player i chooses strategy $s_i \in S_i$
- Obtain strategy profile $s = (s_1, \dots, s_N)$
- Corresponds to a cell in our payoff matrix

Nash Equilibrium

- When no one wants to change their strategy in response
- A strategy profile where, for each player, there is no better strategy if others do not change theirs
- Formal: $\hat{s} = (\hat{s}_1, \dots, \hat{s}_N)$ is a nash equilibrium if, for every player i and $s_i \in S_i$, it holds that;

$$u_i(\hat{s}) \geq u_i(\hat{s}_1, \dots, \hat{s}_{i-1}, s_i, \hat{s}_{i+1}, \dots, \hat{s}_N)$$

Examples

Games with a unique nash equilibrium:

- Example 1

	b_1	b_2	b_3
a_1	(10, 10)	(4, 4)	(0, 18)
a_2	(8, 7)	(8, -6)	(7, 8)

Examples

Games with a unique nash equilibrium:

- Example 1

	b_1	b_2	b_3
a_1	(10, 10)	(4, 4)	(0, 18)
a_2	(8, 7)	(8, -6)	(7, 8)

Examples

Games with a unique nash equilibrium:

- Example 1

	b_1	b_2	b_3
a_1	(10, 10)	(4, 4)	(0, 18)
a_2	(8, 7)	(8, -6)	(7, 8)

- Example 2

	b_1	b_2	b_3
a_1	(7, -2)	(0, 2)	(9, 1)
a_2	(4, 4)	(3, 2)	(-2, 4)
a_3	(-2, 7)	(8, 8)	(7, 8)

Examples

Games with a unique nash equilibrium:

- Example 1

	b_1	b_2	b_3
a_1	(10, 10)	(4, 4)	(0, 18)
a_2	(8, 7)	(8, -6)	(7, 8)

- Example 2

	b_1	b_2	b_3
a_1	(7, -2)	(0, 2)	(9, 1)
a_2	(4, 4)	(3, 2)	(-2, 4)
a_3	(-2, 7)	(8, 8)	(7, 8)

Examples

- Games with multiple nash equilibria:

	b_1	b_2	b_3
a_1	(1, 1)	(0, 0)	(0, 0)
a_2	(0, 0)	(0, 0)	(100, 100)
a_3	(0, 0)	(10, 10)	(0, 0)

- Game without nash equilibrium:

	b_1	b_2	b_3
a_1	(7, 7)	(4, 4)	(18, 0)
a_2	(8, 8)	(8, 5)	(7, 10)
a_3	(0, 14)	(7, 9)	(-2, 18)

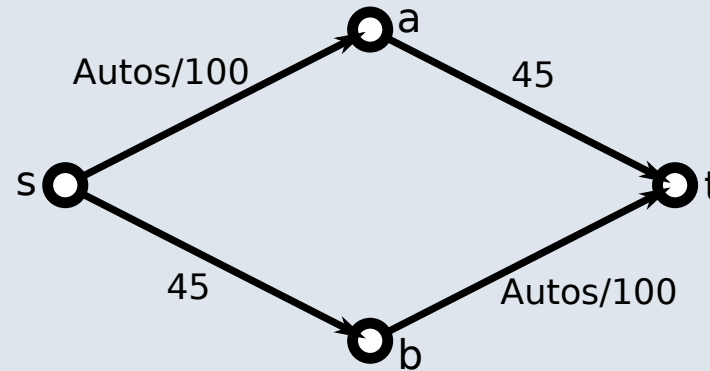
- Is there an equilibrium in the Prisoner's Dilemma?

How to Find nash equilibria?

- Approach 1: check cell by cell
- Approach 2:
 - Mark highest values on the left per column
 - Mark highest values on the right per row
 - Check for overlap

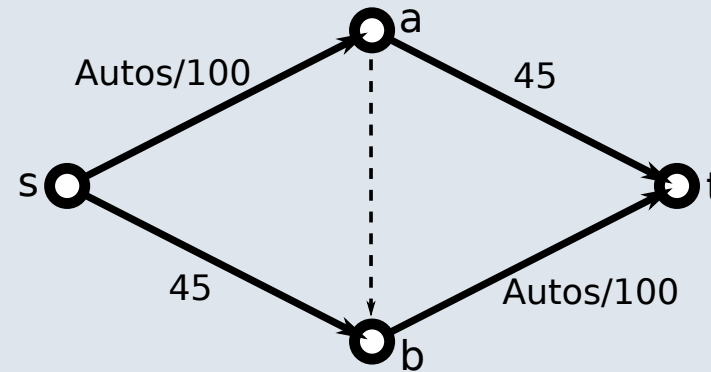
Braess's Paradox

- Introducing a new alternative can make everyone worse off



- 4000 drivers want to travel from *s* to *t*
- A*: drivers going above. Travel time *A*: $A/100 + 45$
- B*: drivers going below. Travel time *B*: $B/100 + 45$
- $A + B = 4000$
- Nash equilibrium: $A = B = 2000$, each $2000/100 + 45 = 65$ minutes

Braess's Paradox



- Build new road $a \rightarrow b$ with travel time 0
- For someone currently in A or B : travel time $s - a - b - t$ is now $2000/100 + 2001/100 = 40.01$ minutes
- Players switch, switch, switch, ...
- $s - a - b - t$ remains the best alternative but gets worse over time
- In the end: everyone uses $s - a - b - t$
- Duration: $4000/100 + 4000/100 = 80$ minutes

Games of the Prisoner's Dilemma Type

	b_1	b_2
a_1	(4, 4)	(0, 6)
a_2	(6, 0)	(2, 2)

- In matrix: payoff instead of damage
- Wide application: two companies want to set prices
- a_1/b_1 : keep the price
- a_2/b_2 : reduce the price
- Nash equilibrium at the "bad" point (a_2, b_2)

Games of the "Battle of the Sexes" Type

- A and B want to go out
- Theater or boxing match
- Independently choose either theater ticket (a_1/b_1) or ticket for boxing match (a_2/b_2)
- A prefers a joint theater evening
- B prefers a joint boxing match
- Separated worse than being together

	b_1	b_2
a_1	(2, 1)	(-1, -1)
a_2	(-1, -1)	(1, 2)

- Nash equilibria are at (2, 1) and (1, 2)

Cooperation

- Simultaneous, independent decision-making is difficult to lead to the desired outcome
- Assume both flip a coin (mixed strategy)
- A: π_A for a_1 , $(1 - \pi_A)$ for a_2
- B: π_B for b_1 , $(1 - \pi_B)$ for b_2
- Expected outcomes:

- For A:

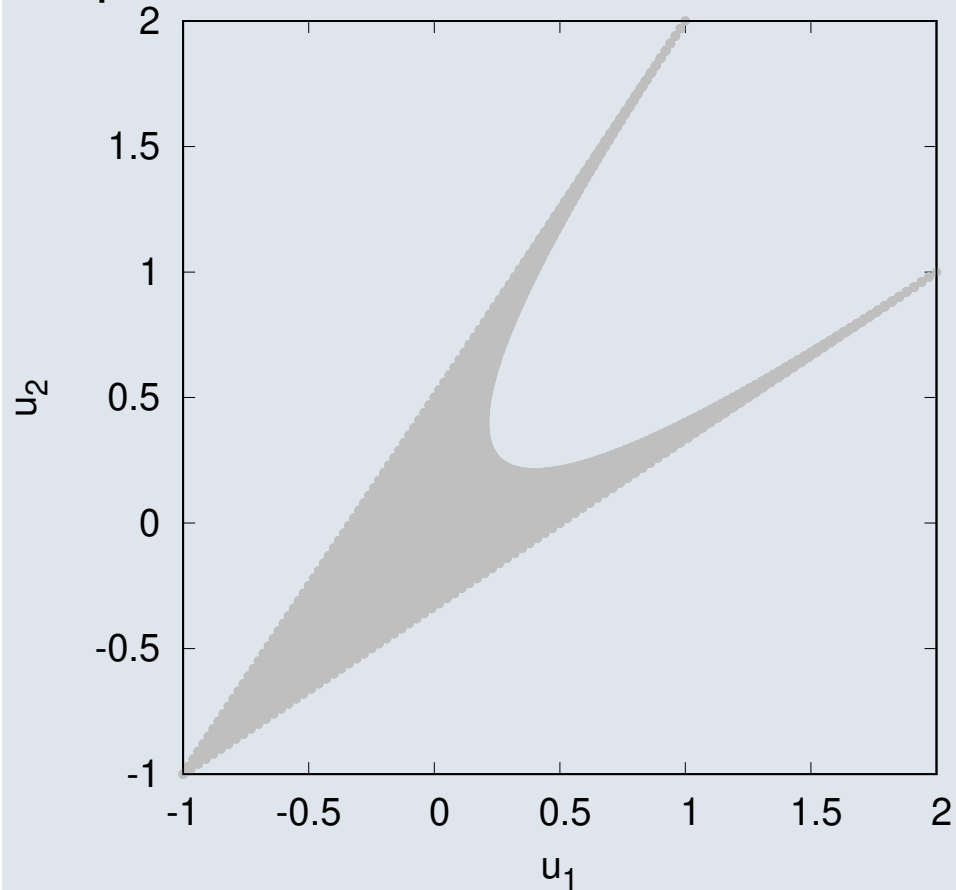
$$u_1(\pi_A, \pi_B) = 2\pi_A\pi_B - \pi_A(1 - \pi_B) - (1 - \pi_A)\pi_B + (1 - \pi_A)(1 - \pi_B)$$

- For B:

$$u_2(\pi_A, \pi_B) = 1\pi_A\pi_B - \pi_A(1 - \pi_B) - (1 - \pi_A)\pi_B + 2(1 - \pi_A)(1 - \pi_B)$$

Cooperation

All possible outcomes:



$$u_1 \left(\frac{1}{2}, \frac{1}{2} \right) = u_2 \left(\frac{1}{2}, \frac{1}{2} \right) = \frac{1}{4}$$

Cooperation

- Suppose we allow agreements instead
- Coin toss would lead to utility expectation of $(1.5, 1.5)$
- Alternating agreement on Boxing/Theater also averages utility value of $(1.5, 1.5)$
- Point $(1.5, 1.5)$ is outside the previously possible range!
- We call all cooperatively attainable outcomes K

Nash Solutions

- What are the "optimal solutions" in a cooperative game?
- Let (\bar{u}_1, \bar{u}_2) be a "special value"
- The value one would get without an agreement, e.g.,
 - No trade: $(\bar{u}_1, \bar{u}_2) = (0, 0)$
 - No cooperation: guaranteed payout in the non-cooperative game
- Goal: find point (\hat{u}_1, \hat{u}_2) in data $(K; \bar{u}_1, \bar{u}_2)$

$$F : (K; \bar{u}_1, \bar{u}_2) \mapsto (\hat{u}_1, \hat{u}_2) \in K$$

Requirements for Nash Solution

- Requirement 1: Independence from linear transformations
 - If utility zero points and utility units are changed, the solution (\hat{u}_1, \hat{u}_2) changes in the same proportion
- Requirement 2: Individual Rationality
 - Solution (\hat{u}_1, \hat{u}_2) must satisfy inequalities $\hat{u}_1 \geq \bar{u}_1$, $\hat{u}_2 \geq \bar{u}_2$
 - No rational decision-maker would accept an agreement that puts them in a worse position than a failure of negotiations

Cooperative Games

- Requirement 3: Pareto-Optimality
 - Solution (\hat{u}_1, \hat{u}_2) is undominated
- Requirement 4: Symmetry
 - If the roles of both players are completely symmetrical, the solution is also symmetrical, i.e., $\hat{u}_1 = \hat{u}_2$.
- Requirement 5: Independence of irrelevant alternatives
 - If \tilde{K} is a subset of K that includes both the point (\bar{u}_1, \bar{u}_2) and the solution (\hat{u}_1, \hat{u}_2) of $(K; \bar{u}_1, \bar{u}_2)$, then $F_0(\tilde{K}; \bar{u}_1, \bar{u}_2) = (\hat{u}_1, \hat{u}_2)$ also holds.

Summary of Requirements

$F_0(K; \bar{u}_1, \bar{u}_2) = (\hat{u}_1, \hat{u}_2)$ is the maximum of the following optimization problem:

$$\begin{aligned} \max & (u_1 - \bar{u}_1) \cdot (u_2 - \bar{u}_2) \\ \text{s.t.} & (u_1, u_2) \in K \\ & u_1 \geq \bar{u}_1 \\ & u_2 \geq \bar{u}_2 \end{aligned}$$

Theorem

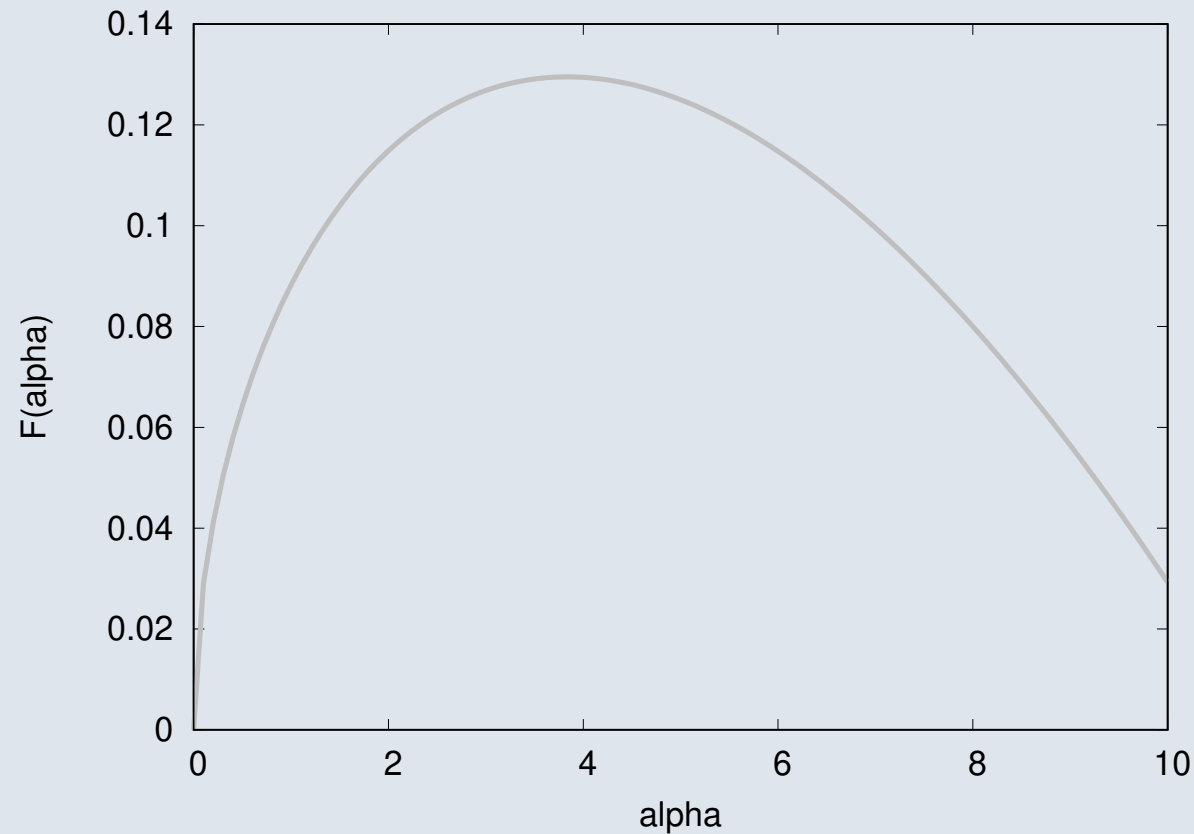
There is exactly one arbitration mechanism defined on the set of all bargaining games $(K; \bar{u}_1, \bar{u}_2)$ that simultaneously satisfies all five requirements; this is given by the F_0 described above.

Example: Resolution of a Wage Dispute

- Unions and employers negotiate a wage increase α in the range from 0 to 10 percent ($0 \leq \alpha \leq 10$)
- Unions: $u_1(\alpha) = \sqrt{\frac{\alpha}{20}}$, $\bar{u}_1 = \frac{1}{5}$
- Employers: $u_2(\alpha) = 1 - \frac{\alpha}{20}$, $\bar{u}_2 = \frac{3}{5}$
- Maximize:

$$\begin{aligned}(u_1 - \bar{u}_2) \cdot (u_2 - \bar{u}_2) &= \left(\sqrt{\frac{\alpha}{20}} - \frac{1}{5} \right) \cdot \left(1 - \frac{\alpha}{20} - \frac{3}{5} \right) \\ &= \frac{2}{5} \cdot \sqrt{\frac{\alpha}{20}} - \frac{2}{25} - \left(\frac{\alpha}{20} \right)^{\frac{3}{2}} + \frac{\alpha}{100}\end{aligned}$$

Example: Resolution of a Wage Dispute



Optimal Solution: $\alpha = 3.8$; $u_1 = 0.44$; $u_2 = 0.81$

Issues and Applications

- Corresponding data is usually not available
- In reality, negotiation games are often divided into a threat game and a demand game (two-stage game): Threats (e.g., strike and lockout) are only implemented if the negotiation solution fails
- Player behavior often does not meet the rationality postulate (problem applies to the entire analytical game theory)
- However, game theory is very suitable for analyzing conflict and decision situations

Quiz

Question 1

Is there a nash equilibrium?
(non-coop, pure strategies)

	b_1	b_2
a_1	(0,1)	(1,0)
a_2	(1,0)	(0,1)

Question 3

What is the best cooperative solution?

	b_1	b_2
a_1	(0,1)	(1,0)
a_2	(1,0)	(0,1)

Question 3

Is there a nash equilibrium?
(non-coop, pure strategies)

	b_1	b_2
a_1	(3,2)	(1,4)
a_2	(2,2)	(4,3)

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Solution

No.

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What is the best cooperative solution?

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a_1	(0,1)	(1,0)
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Solution

Everyone plays 1/2 with utility (1/2, 1/2).

Question 3

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(non-coop, pure strategies)

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(non-coop, pure strategies)

	b_1	b_2
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a_2	(2,2)	(4,3)

Solution

Yes: (a_2, b_2) .