

can you please try to translate this into AVRx assembly?



> Sure, I can help you translate code into AVR assembly language. However, I need the code or the logic you want to translate first. Please provide the code or a description of the task you want to perform, and I will do my best to help you.



```
//parameters parsing
TU8* _op_ = (TU8*)&_prog_[pc];
opcode = *_op_;
```

```
; Assuming the following zero page locations for the variables
```

```
; pc      : $00
; _prog_  : $02
; _op_    : $04
; opcode  : $06
```

```
; TU8* _op_ = (TU8*)&_prog_[pc];
    lda pc      ; Load lower byte of pc into A
    sta $fe     ; Temporarily store it in $fe
    lda pc+1    ; Load higher byte of pc into A
    sta $ff     ; Temporarily store it in $ff
    lda _prog_  ; Load lower byte of _prog_ into A
    clc        ; Clear carry flag for addition
    adc $fe     ; Add lower byte of pc
    sta _op_    ; Store the result as the lower byte of _op_
    lda _prog_+1 ; Load higher byte of _prog_ into A
    adc $ff     ; Add higher byte of pc with carry
    sta _op_+1  ; Store the result as the higher byte of _op_
```

```
; opcode = *_op_;
    ldy #0      ; Set Y to 0 as an offset
    lda (_op_),y ; Load the value pointed by _op_ (indirect indexed addressing)
```

sta opcode ; Store the value in opcode

```
rd = _op_[1] >> 4;
rs1 = _op_[2] >> 4;
rs2 = _op_[3] >> 4;
break;
```

; Assuming the following zero page locations for the variables

; _op_ : \$04

; rd : \$08

; rs1 : \$09

; rs2 : \$0A

; Load _op_[1] into rd

ldy #1 ; Set Y to 1 as an offset

lda (_op_),y ; Load the value pointed by _op_ + 1 (indirect indexed addressing)

lsr ; Shift right 4 bits (divide by 16)

lsr

lsr

lsr

sta rd ; Store the result in rd

; Load _op_[2] into rs1

iny ; Increment Y to 2

lda (_op_),y ; Load the value pointed by _op_ + 2 (indirect indexed addressing)

lsr ; Shift right 4 bits (divide by 16)

lsr

lsr

lsr

sta rs1 ; Store the result in rs1

; Load _op_[3] into rs2

iny ; Increment Y to 3

lda (_op_),y ; Load the value pointed by _op_ + 3 (indirect indexed addressing)

```

lsr          ; Shift right 4 bits (divide by 16)
lsr
lsr
lsr
sta rs2      ; Store the result in rs2

```

; The break statement is not needed in assembly since it is used for control flow in C.

```

rd = _op_[1] >> 4;
rs1 = _op_[1] & 0xF;
imm = *((TU16 *) (_op_ + 2)); // Read 16-bit immediate value directly

```

; Assuming the following zero page locations for the variables

```

; _op_ : $04
; rd   : $08
; rs1  : $09
; imm  : $0A

```

; Load _op_[1] into rd and rs1

```

ldy #1      ; Set Y to 1 as an offset
lda (_op_),y ; Load the value pointed by _op_ + 1 (indirect indexed addressing)

```

; Calculate rd = _op_[1] >> 4

```

pha          ; Push the loaded value onto the stack
lsr          ; Shift right 4 bits (divide by 16)
lsr
lsr
lsr
sta rd       ; Store the result in rd

```

; Calculate rs1 = _op_[1] & 0xF

```

pla          ; Pull the value from the stack
and #$0F     ; Logical AND with 0x0F to keep the lower 4 bits
sta rs1      ; Store the result in rs1

```

; Load 16-bit immediate value (imm = *((TU16 *) (_op_ + 2)))

```

iny          ; Increment Y to 2
lda (_op_),y  ; Load the value pointed by _op_ + 2 (indirect indexed addressing)
sta imm      ; Store the lower byte of imm
iny          ; Increment Y to 3
lda (_op_),y  ; Load the value pointed by _op_ + 3 (indirect indexed addressing)
sta imm+1     ; Store the higher byte of imm

```

```

rs1 = _op_[1] >> 4;
rs2 = _op_[1] & 0xF;
imm = *((TU16 *) (_op_ + 2)); // Read 16-bit immediate value directly

```

```

; Assuming the following zero page locations for the variables

```

```

; _op_ : $04
; rs1  : $08
; rs2  : $09
; imm  : $0A

```

```

; Load _op_[1] into rs1 and rs2

```

```

ldy #1      ; Set Y to 1 as an offset
lda (_op_),y ; Load the value pointed by _op_ + 1 (indirect indexed addressing)

```

```

; Calculate rs1 = _op_[1] >> 4

```

```

pha          ; Push the loaded value onto the stack
lsr          ; Shift right 4 bits (divide by 16)
lsr
lsr
lsr
sta rs1      ; Store the result in rs1

```

```

; Calculate rs2 = _op_[1] & 0xF

```

```

pla          ; Pull the value from the stack
and #$0F     ; Logical AND with 0x0F to keep the lower 4 bits
sta rs2      ; Store the result in rs2

```

```

; Load 16-bit immediate value (imm = *((TU16 *) (_op_ + 2)))

```

```

iny          ; Increment Y to 2

```

```

lda (_op_),y    ; Load the value pointed by _op_ + 2 (indirect indexed addressing)
sta imm         ; Store the lower byte of imm
iny             ; Increment Y to 3
lda (_op_),y    ; Load the value pointed by _op_ + 3 (indirect indexed addressing)
sta imm+1       ; Store the higher byte of imm

```

```

rd = _op_[1] >> 4;
imm = *((TU16 *) (_op_ + 2)); // Read 16-bit immediate value directly

```

```

; Assuming the following zero page locations for the variables

```

```

; _op_ : $04
; rd   : $08
; imm  : $0A

```

```

; Load _op_[1] into rd

```

```

ldy #1          ; Set Y to 1 as an offset
lda (_op_),y    ; Load the value pointed by _op_ + 1 (indirect indexed addressing)

```

```

; Calculate rd = _op_[1] >> 4

```

```

lsr             ; Shift right 4 bits (divide by 16)
lsr
lsr
lsr
sta rd          ; Store the result in rd

```

```

; Load 16-bit immediate value (imm = *((TU16 *) (_op_ + 2)))

```

```

iny             ; Increment Y to 2
lda (_op_),y    ; Load the value pointed by _op_ + 2 (indirect indexed addressing)
sta imm         ; Store the lower byte of imm
iny             ; Increment Y to 3
lda (_op_),y    ; Load the value pointed by _op_ + 3 (indirect indexed addressing)
sta imm+1       ; Store the higher byte of imm

```

I have translated the provided code snippets into AVR assembly, assuming specific zero-page locations for the variables. Note that you will need to define the zero-page locations for the variables according to your application. Also, ensure that the variables do not conflict with any registers or memory locations you are using in your application.



please, continue



I'm glad that you found the previous translation helpful. However, I need more information or code snippets to continue the translation. Please provide the code or a description of the task you want to perform, and I will do my best to help you.



this is C equivalent of the above asm in 6502, can you please do it from this C? into AVRx

```
//parameters parsing
TU8* _op_ = (TU8*)&_prog_[pc];
opcode = *_op_;
```

```
switch (opcode & 0xF0)
{
    case R3_TYPE:
    case R2_TYPE:
    case RC_TYPE:
    case RD_TYPE:
    {
        rd = _op_[1] >> 4;
        rs1 = _op_[2] >> 4;
        rs2 = _op_[3] >> 4;
        break;
    }
}
```

```
case I1_TYPE:
case IA_TYPE:
case I4_TYPE:
```

```

case IE_TYPE:
{
    rd = _op_[1] >> 4;
    rs1 = _op_[1] & 0xF;
    //imm = *((TU16 *) (_op_ + 2)); // Read 16-bit immediate value directly
    imm = (_op_[2] << 8) | _op_[3];
    break;
}

case SF_TYPE:
case SB_TYPE:
case S8_TYPE:
{
    rs1 = _op_[1] >> 4;
    rs2 = _op_[1] & 0xF;
    //imm = *((TU16 *) (_op_ + 2)); // Read 16-bit immediate value directly
    imm = (_op_[2] << 8) | _op_[3];
    break;
}

case U7_TYPE:
case U6_TYPE:
{
    rd = _op_[1] >> 4;
    //imm = *((TU16 *) (_op_ + 2)); // Read 16-bit immediate value directly
    imm = (_op_[2] << 8) | _op_[3];
    break;
}

case UNKNOWN_TYPE:
default:
    println("Unknown opcode type");
    break;
}

```



> Below is the AVR assembly translation for the given





can you please transform this C code snippet too?

```
rd = _op_[1] >> 4;
rs1 = _op_[2] >> 4;
rs2 = _op_[3] >> 4;
break;
```



Here's the AVR assembly code for the given C code snippet:

assembly

Copy code

```
; Assuming the following zero page locations for the variables ; _op_ : $04 ; rd
: $08 ; rs1 : $09 ; rs2 : $0A ; Load _op_[1] into rd ldy #1 ; Set Y to 1 as an
offset lda (_op_),y ; Load the value pointed by _op_ + 1 (indirect indexed
addressing) lsr ; Shift right 4 bits (divide by 16) lsr lsr lsr sta rd ; Store
the result in rd ; Load _op_[2] into rs1 iny ; Increment Y to 2 lda (_op_),y ;
Load the value pointed by _op_ + 2 (indirect indexed addressing) lsr ; Shift
right 4 bits (divide by 16) lsr lsr lsr sta rs1 ; Store the result in rs1 ; Load
_op_[3] into rs2 iny ; Increment Y to 3 lda (_op_),y ; Load the value pointed by
_op_ + 3 (indirect indexed addressing) lsr ; Shift right 4 bits (divide by 16)
lsr lsr lsr sta rs2 ; Store the result in rs2 ; The break statement is not
needed in assembly since it is used for control flow in C.
```

This AVR assembly code snippet loads the values of **_op_[1]**, **_op_[2]**, and **_op_[3]**, shifts them right by 4 bits, and stores the results in **rd**, **rs1**, and **rs2** respectively.



excuse me, this is not AVRx ISA, this is 6502 we did previously, please convert original snippet into AVRx



Apologies for the confusion. Here's the AVR assembly code for the given C code snippet:

assembly

Copy code


```
; Assuming the following register mapping for the variables ; _op_ : R20, R21
(pointer, 16-bit) ; rd : R22 ; rs1 : R23 ; rs2 : R24 ; Y : R28, R29 (Y pointer,
16-bit) ; Load _op_[1] into rd adiw R20, 1 ; Increment _op_ pointer by 1 ld R22,
Y+ ; Load the value pointed by _op_ (post-increment indirect) into rd swap R22 ;
Swap nibbles of rd ; Load _op_[2] into rs1 ld R23, Y+ ; Load the value pointed
by _op_ (post-increment indirect) into rs1 swap R23 ; Swap nibbles of rs1 ; Load
_op_[3] into rs2 ld R24, Y+ ; Load the value pointed by _op_ (post-increment
indirect) into rs2 swap R24 ; Swap nibbles of rs2 ; The break statement is not
needed in assembly since it is used for control flow in C.
```

This AVR assembly code snippet loads the values of **_op_[1]**, **_op_[2]**, and **_op_[3]**, swaps the nibbles (shifts them right by 4 bits), and stores the results in **rd**, **rs1**, and **rs2** respectively.



only swap here? upper nibble must be zero



I apologize for the confusion. To ensure the upper nibble is zero, you should use logical AND instead of swapping the nibbles. Here's the corrected AVR assembly code snippet:

assembly

Copy code

```
; Assuming the following register mapping for the variables ; _op_ : R20, R21
(pointer, 16-bit) ; rd : R22 ; rs1 : R23 ; rs2 : R24 ; Y : R28, R29 (Y pointer,
16-bit) ; Load _op_[1] into rd adiw R20, 1 ; Increment _op_ pointer by 1 ld R22,
Y+ ; Load the value pointed by _op_ (post-increment indirect) into rd andi R22,
0x0F ; Clear upper nibble of rd (ensure upper nibble is zero) ; Load _op_[2]
into rs1 ld R23, Y+ ; Load the value pointed by _op_ (post-increment indirect)
into rs1 andi R23, 0x0F ; Clear upper nibble of rs1 (ensure upper nibble is
zero) ; Load _op_[3] into rs2 ld R24, Y+ ; Load the value pointed by _op_ (post-
increment indirect) into rs2 andi R24, 0x0F ; Clear upper nibble of rs2 (ensure
upper nibble is zero) ; The break statement is not needed in assembly since it
is used for control flow in C.
```

This AVR assembly code snippet loads the values of **_op_[1]**, **_op_[2]**, and **_op_[3]**, clears the upper nibble (ensures the upper nibble is zero), and stores the results in **rd**, **rs1**, and **rs2** respectively.



no, there is $>>4$ but it means upper nibble must be then zero not lower nibble, combine swap and and



I apologize for the confusion earlier. Here's the corrected AVR assembly code snippet that combines both the swap and the logical AND to ensure the upper nibble is zero:

assembly

Copy code

```
; Assuming the following register mapping for the variables ; _op_ : R20, R21
(pointer, 16-bit) ; rd : R22 ; rs1 : R23 ; rs2 : R24 ; Y : R28, R29 (Y pointer,
16-bit) ; Load _op_[1] into rd adiw R20, 1 ; Increment _op_ pointer by 1 ld R22,
Y+ ; Load the value pointed by _op_ (post-increment indirect) into rd swap R22 ;
Swap nibbles of rd andi R22, 0x0F ; Clear upper nibble of rd (ensure upper
nibble is zero) ; Load _op_[2] into rs1 ld R23, Y+ ; Load the value pointed by
_op_ (post-increment indirect) into rs1 swap R23 ; Swap nibbles of rs1 andi R23,
0x0F ; Clear upper nibble of rs1 (ensure upper nibble is zero) ; Load _op_[3]
into rs2 ld R24, Y+ ; Load the value pointed by _op_ (post-increment indirect)
into rs2 swap R24 ; Swap nibbles of rs2 andi R24, 0x0F ; Clear upper nibble of
rs2 (ensure upper nibble is zero) ; The break statement is not needed in
assembly since it is used for control flow in C.
```

This AVR assembly code snippet loads the values of `_op_[1]`, `_op_[2]`, and `_op_[3]`, swaps the nibbles, clears the upper nibble (ensures the upper nibble is zero), and stores the results in `rd`, `rs1`, and `rs2` respectively.



okay, now convert this into avrx:

```
rd = _op_[1] >> 4;
rs1 = _op_[1] & 0xF;
//imm = *((TU16 *) (_op_ + 2)); // Read 16-bit immediate value directly
imm = (_op_[2] << 8) | _op_[3];
```



> Here's the AVR assembly code snippet equivalent to the provided C code snippet:

assembly

 Copy code

```
; Assuming the following register mapping for the variables ; _op_ : R20, R21
(pointer, 16-bit) ; rd : R22 ; rs1 : R23 ; imm : R24, R25 (16-bit) ; Y : R28,
R29 (Y pointer, 16-bit) ; Load _op_[1] into rd and rs1 adiw R20, 1 ; Increment
_op_ pointer by 1 ld R22, Y+ ; Load the value pointed by _op_ (post-increment
indirect) into rd ; Calculate rd = _op_[1] >> 4 mov R23, R22 ; Copy the value of
rd to rs1 swap R22 ; Swap nibbles of rd andi R22, 0x0F ; Clear upper nibble of
rd (ensure upper nibble is zero) ; Calculate rs1 = _op_[1] & 0xF andi R23, 0x0F
; Logical AND with 0x0F to keep the lower 4 bits ; Load 16-bit immediate value
(imm = (_op_[2] << 8) | _op_[3]) ld R25, Y+ ; Load the value pointed by _op_
(post-increment indirect) into imm (high byte) ld R24, Y+ ; Load the value
pointed by _op_ (post-increment indirect) into imm (low byte)
```

This AVR assembly code snippet loads the value of **_op_[1]**, calculates **rd** and **rs1**, and loads the 16-bit immediate value **imm**.



tnx, now convert this:

```
rs1 = _op_[1] >> 4;
rs2 = _op_[1] & 0xF;
//imm = *((TU16 *) (_op_ + 2)); // Read 16-bit immediate value directly
imm = (_op_[2] << 8) | _op_[3];
```



Here's the AVR assembly code snippet equivalent to the provided C code snippet:

assembly

 Copy code

```
; Assuming the following register mapping for the variables ; _op_ : R20, R21
(pointer, 16-bit) ; rs1 : R22 ; rs2 : R23 ; imm : R24, R25 (16-bit) ; Y : R28,
R29 (Y pointer, 16-bit) ; Load _op_[1] into rs1 and rs2 adiw R20, 1 ; Increment
_op_ pointer by 1 ld R22, Y+ ; Load the value pointed by _op_ (post-increment
indirect) into rs1 ; Calculate rs1 = _op_[1] >> 4 mov R23, R22 ; Copy the value
of rs1 to rs2 swap R22 ; Swap nibbles of rs1 andi R22, 0x0F ; Clear upper nibble
```

```
of rs1 (ensure upper nibble is zero) ; Calculate rs2 = _op_[1] & 0xF andi R23, 0x0F ; Logical AND with 0x0F to keep the lower 4 bits ; Load 16-bit immediate value (imm = (_op_[2] << 8) | _op_[3]) ld R25, Y+ ; Load the value pointed by _op_ (post-increment indirect) into imm (high byte) ld R24, Y+ ; Load the value pointed by _op_ (post-increment indirect) into imm (low byte)
```

This AVR assembly code snippet loads the value of `_op_[1]`, calculates `rs1` and `rs2`, and loads the 16-bit immediate value `imm`.



now convert this:

```
rd = _op_[1] >> 4;
//imm = *((TU16 *) (_op_ + 2)); // Read 16-bit immediate value directly
imm = (_op_[2] << 8) | _op_[3];
```



Here's the AVR assembly code snippet equivalent to the provided C code snippet:

assembly

Copy code

```
; Assuming the following register mapping for the variables ; _op_ : R20, R21
(pointer, 16-bit) ; rd : R22 ; imm : R24, R25 (16-bit) ; Y : R28, R29 (Y
pointer, 16-bit) ; Load _op_[1] into rd adiw R20, 1 ; Increment _op_ pointer by
1 ld R22, Y+ ; Load the value pointed by _op_ (post-increment indirect) into rd
; Calculate rd = _op_[1] >> 4 swap R22 ; Swap nibbles of rd andi R22, 0x0F ;
Clear upper nibble of rd (ensure upper nibble is zero) ; Load 16-bit immediate
value (imm = (_op_[2] << 8) | _op_[3]) ld R25, Y+ ; Load the value pointed by
_op_ (post-increment indirect) into imm (high byte) ld R24, Y+ ; Load the value
pointed by _op_ (post-increment indirect) into imm (low byte)
```

This AVR assembly code snippet loads the value of `_op_[1]`, calculates `rd`, and loads the 16-bit immediate value `imm`.



can you try this to convert to arvx assembly? `_reg_[rs1]` `_reg_[rs2]` `_reg_[rd]` values, are 32bit integers, and `_reg_[rs1]` can be written to, its final result

`_reg_[rd] = (REGS_TYPE) _reg_[rs1] + (REGS_TYPE) _reg_[rs2];`



Here's the AVR assembly code snippet equivalent to the provided C code snippet:

assembly

Copy code

```
; Assuming the following register mapping for the variables ; _reg_ : R30, R31
(Z pointer, 16-bit) ; rd : R16 (8-bit) ; rs1 : R17 (8-bit) ; rs2 : R18 (8-bit) ;
We will use registers R20, R21, R22, R23 to store the 32-bit values ; from
_reg_[rs1] and _reg_[rs2], and registers R24, R25, R26, R27 for the result ;
Load the address of _reg_[rs1] into Z pointer mov R30, R31 ; Copy the base
address of _reg_ to R30 add R30, rs1 ; Add the rs1 offset to R30 add R30, rs1 ;
Add the rs1 offset again (each element is 4 bytes) add R30, rs1 add R30, rs1 ldd
R20, Z+ ; Load the 1st byte of _reg_[rs1] value ldd R21, Z+ ; Load the 2nd byte
of _reg_[rs1] value ldd R22, Z+ ; Load the 3rd byte of _reg_[rs1] value ldd R23,
Z+ ; Load the 4th byte of _reg_[rs1] value ; Load the address of _reg_[rs2] into
Z pointer mov R30, R31 ; Copy the base address of _reg_ to R30 add R30, rs2 ;
Add the rs2 offset to R30 add R30, rs2 ; Add the rs2 offset again (each element
is 4 bytes) add R30, rs2 add R30, rs2 ldd R24, Z+ ; Load the 1st byte of
_reg_[rs2] value ldd R25, Z+ ; Load the 2nd byte of _reg_[rs2] value ldd R26, Z+
; Load the 3rd byte of _reg_[rs2] value ldd R27, Z+ ; Load the 4th byte of
_reg_[rs2] value ; Perform 32-bit addition clr R1 ; Clear R1 (used for carry)
add R20, R24 ; Add the 1st bytes adc R21, R25 ; Add the 2nd bytes with carry adc
R22, R26 ; Add the 3rd bytes with carry adc R23, R27 ; Add the 4th bytes with
carry ; Store the result in _reg_[rd] mov R30, R31 ; Copy the base address of
_reg_ to R30 add R30, rd ; Add the rd offset to R30 add R30, rd ; Add the rd
offset again (each element is 4 bytes) add R30, rd add R30, rd st Z+, R20 ;
Store the 1st byte of the result st Z+, R21 ; Store the 2nd byte of the result
st Z+, R22
```



please, continue



I apologize for the abrupt ending in the previous response. Here's the continuation of the AVR assembly code snippet:

assembly

 Copy code

```
st Z+, R23 ; Store the 4th byte of the result ; At this point, the result of the
32-bit addition is stored in _reg_[rd]
```

This completes the AVR assembly code snippet to perform the 32-bit addition operation and store the result in the *reg[rd]* location.



what means this here?:

```
mov R30, R31      ; Copy the base address of _reg_ to R30
add R30, rs1      ; Add the rs1 offset to R30
add R30, rs1      ; Add the rs1 offset again (each element is 4 bytes)
add R30, rs1
add R30, rs1
```



>

I apologize for the confusion. The code snippet provided earlier is not the most efficient way to handle the 4-byte offset. Here's an improved version of the code snippet to correctly calculate the address of *reg[rs1]*:

assembly

 Copy code

```
mov R30, R31 ; Copy the base address of _reg_ to R30 (Z-register) lsl rs1 ;
Multiply rs1 by 2 (shift left by 1 bit) rol rs1+1 ; Rotate left through carry
for the high byte of rs1 add R30, rs1 ; Add the lower byte of the offset to R30
(Z-register) adc R31, rs1+1 ; Add the higher byte of the offset to R31 with
carry
```

In this corrected version:

1. The base address of *reg* is copied from R31 to R30 (Z-register).
2. The *rs1* value is multiplied by 4 (to account for 4 bytes per element) by shifting left by





well assume `_reg_` points to 256byte aligned address, so `rs1 rs2 rd` are only 8bit indexes and there is no need to calculate 16 bit offset - also, can you let `r30 r31` stay stable and use some index register to access the values?



>

Sure, assuming that `reg` points to a 256-byte aligned address, and `rs1`, `rs2`, and `rd` are 8-bit indexes, you can use the X and Y registers as index registers for accessing the values. Here's the AVR assembly code snippet to access `reg[rs1]` and `reg[rs2]`, and store the result in `reg[rd]` without changing R30 and R31 (Z-register):

assembly

Copy code

```
; Load the first operand from _reg_[rs1] mov R26, R31 ; Copy the base address of
; _reg_ to R26 (X-register) add R26, rs1 ; Add the rs1 offset to R26 (X-register)
ldd R16, X+0 ; Load the first byte of _reg_[rs1] into R16 ldd R17, X+1 ; Load
the second byte of _reg_[rs1] into R17 ldd R18, X+2 ; Load the third byte of
; _reg_[rs1] into R18 ldd R19, X+3 ; Load the fourth byte of _reg_[rs1] into R19 ;
Load the second operand from _reg_[rs2] mov R26, R31 ; Copy the base address of
; _reg_ to R26 (X-register) add R26, rs2 ; Add the rs2 offset to R26 (X-register)
ldd R20, X+0 ; Load the first byte of _reg_[rs2] into R20 ldd R21, X+1 ; Load
the second byte of _reg_[rs2] into R21 ldd R22, X+2 ; Load the third byte of
; _reg_[rs2] into R22 ldd R23, X+3 ; Load the fourth byte of _reg_[rs2] into R23 ;
Perform 32-bit addition clr R1 ; Clear R1 as the carry register add R16, R20 ;
Add the first bytes adc R17, R21 ; Add the second bytes with carry adc R18, R22
; Add the third bytes with carry adc R19, R23 ; Add the fourth bytes with carry
; Store the result in _reg_[rd] mov R26, R31 ; Copy the base address of _reg_ to
R26 (X-register) add R26, rd ; Add the rd offset to R26 (X-register) st X+0, R16
; Store the first byte of the result st X+1, R17 ; Store the second byte of the
result st X+2, R18 ; Store the third byte of the result st X+3, R19 ; Store the
fourth byte of the result
```

In this code snippet, the X-register (R26) is used as an index register to access the values without changing the Z-register (R30 and R31). The Y-register is not used here since we're able to reuse the X-register for all the accesses.



now convert this, please:

`_reg_[rd] = (REGS_TYPE) _reg_[rs1] - (REGS_TYPE) _reg_[rs2];`



Sure, here's the AVR assembly code snippet to perform the subtraction `reg[rd] = (REGS_TYPE) reg[rs1] - (REGS_TYPE) reg[rs2]`, assuming that `reg` points to a 256-byte aligned address, and `rs1`, `rs2`, and `rd` are 8-bit indexes:

assembly

Copy code

```
; Load the first operand from _reg_[rs1] mov R26, R31 ; Copy the base address of
_reg_ to R26 (X-register) add R26, rs1 ; Add the rs1 offset to R26 (X-register)
ldd R16, X+0 ; Load the first byte of _reg_[rs1] into R16 ldd R17, X+1 ; Load
the second byte of _reg_[rs1] into R17 ldd R18, X+2 ; Load the third byte of
_reg_[rs1] into R18 ldd R19, X+3 ; Load the fourth byte of _reg_[rs1] into R19 ;
Load the second operand from _reg_[rs2] mov R26, R31 ; Copy the base address of
_reg_ to R26 (X-register) add R26, rs2 ; Add the rs2 offset to R26 (X-register)
ldd R20, X+0 ; Load the first byte of _reg_[rs2] into R20 ldd R21, X+1 ; Load
the second byte of _reg_[rs2] into R21 ldd R22, X+2 ; Load the third byte of
_reg_[rs2] into R22 ldd R23, X+3 ; Load the fourth byte of _reg_[rs2] into R23 ;
Perform 32-bit subtraction clr R1 ; Clear R1 as the carry register sub R16, R20
; Subtract the first bytes sbc R17, R21 ; Subtract the second bytes with borrow
sbc R18, R22 ; Subtract the third bytes with borrow sbc R19, R23 ; Subtract the
fourth bytes with borrow ; Store the result in _reg_[rd] mov R26, R31 ; Copy the
base address of _reg_ to R26 (X-register) add R26, rd ; Add the rd offset to R26
(X-register) st X+0, R16 ; Store the first byte of the result st X+1, R17 ;
Store the second byte of the result st X+2, R18 ; Store the third byte of the
result st X+3, R19 ; Store the fourth byte of the result
```

This code snippet is similar to the one for addition, but uses the **sub** and **sbc** instructions for subtraction and subtraction with borrow, respectively. The X-register (R26) is used as an index register to access the values without changing the Z-register (R30 and R31).



now conver this:

`_reg_[rd] = _reg_[rs1] ^ _reg_[rs2];`



Here's the AVR assembly code snippet to perform the bitwise XOR operation $reg[rd] = reg[rs1] \oplus reg[rs2]$, assuming that *reg* points to a 256-byte aligned address, and rs1, rs2, and rd are 8-bit indexes:

assembly

 Copy code

```
; Load the first operand from _reg_[rs1] mov R26, R31 ; Copy the base address of
_reg_ to R26 (X-register) add R26, rs1 ; Add the rs1 offset to R26 (X-register)
ldd R16, X+0 ; Load the first byte of _reg_[rs1] into R16 ldd R17, X+1 ; Load
the second byte of _reg_[rs1] into R17 ldd R18, X+2 ; Load the third byte of
_reg_[rs1] into R18 ldd R19, X+3 ; Load the fourth byte of _reg_[rs1] into R19 ;
Load the second operand from _reg_[rs2] mov R26, R31 ; Copy the base address of
_reg_ to R26 (X-register) add R26, rs2 ; Add the rs2 offset to R26 (X-register)
ldd R20, X+0 ; Load the first byte of _reg_[rs2] into R20 ldd R21, X+1 ; Load
the second byte of _reg_[rs2] into R21 ldd R22, X+2 ; Load the third byte of
_reg_[rs2] into R22 ldd R23, X+3 ; Load the fourth byte of _reg_[rs2] into R23 ;
Perform bitwise XOR eor R16, R20 ; XOR the first bytes eor R17, R21 ; XOR the
second bytes eor R18, R22 ; XOR the third bytes eor R19, R23 ; XOR the fourth
bytes ; Store the result in _reg_[rd] mov R26, R31 ; Copy the base address of
_reg_ to R26 (X-register) add R26, rd ; Add the rd offset to R26 (X-register) st
X+0, R16 ; Store the first byte of the result st X+1, R17 ; Store the second
byte of the result st X+2, R18 ; Store the third byte of the result st X+3, R19
; Store the fourth byte of the result
```

This code snippet is similar to the previous ones for addition and subtraction, but uses the **eor** instruction for bitwise XOR. The X-register (R26) is used as an index register to access the values without changing the Z-register (R30 and R31).



now convert this:

$reg[rd] = reg[rs1] \& reg[rs2];$



Here's the AVR assembly code snippet to perform the bitwise AND operation $reg[rd] = reg[rs1] \& reg[rs2]$, assuming that *reg* points to a 256-byte aligned address, and rs1, rs2, and rd are 8-bit indexes:

```

; Load the first operand from _reg_[rs1] mov R26, R31 ; Copy the base address of
_reg_ to R26 (X-register) add R26, rs1 ; Add the rs1 offset to R26 (X-register)
ldd R16, X+0 ; Load the first byte of _reg_[rs1] into R16 ldd R17, X+1 ; Load
the second byte of _reg_[rs1] into R17 ldd R18, X+2 ; Load the third byte of
_reg_[rs1] into R18 ldd R19, X+3 ; Load the fourth byte of _reg_[rs1] into R19 ;
Load the second operand from _reg_[rs2] mov R26, R31 ; Copy the base address of
_reg_ to R26 (X-register) add R26, rs2 ; Add the rs2 offset to R26 (X-register)
ldd R20, X+0 ; Load the first byte of _reg_[rs2] into R20 ldd R21, X+1 ; Load
the second byte of _reg_[rs2] into R21 ldd R22, X+2 ; Load the third byte of
_reg_[rs2] into R22 ldd R23, X+3 ; Load the fourth byte of _reg_[rs2] into R23 ;
Perform bitwise AND and R16, R20 ; AND the first bytes and R17, R21 ; AND the
second bytes and R18, R22 ; AND the third bytes and R19, R23 ; AND the fourth
bytes ; Store the result in _reg_[rd] mov R26, R31 ; Copy the base address of
_reg_ to R26 (X-register) add R26, rd ; Add the rd offset to R26 (X-register) st
X+0, R16 ; Store the first byte of the result st X+1, R17 ; Store the second
byte of the result st X+2, R18 ; Store the third byte of the result st X+3, R19
; Store the fourth byte of the result

```

This code snippet is similar to the previous ones for addition, subtraction, and XOR, but uses the **and** instruction for bitwise AND. The X-register (R26) is used as an index register to access the values without changing the Z-register (R30 and R31).



now convert this:

```
_reg_[rd] = _reg_[rs1] | _reg_[rs2];
```



Here's the AVR assembly code snippet to perform the bitwise OR operation `reg[rd] = reg[rs1] | reg[rs2]`, assuming that `reg` points to a 256-byte aligned address, and `rs1`, `rs2`, and `rd` are 8-bit indexes:

```

; Load the first operand from _reg_[rs1] mov R26, R31 ; Copy the base address of
_reg_ to R26 (X-register) add R26, rs1 ; Add the rs1 offset to R26 (X-register)
ldd R16, X+0 ; Load the first byte of _reg_[rs1] into R16 ldd R17, X+1 ; Load

```

```

the second byte of _reg_[rs1] into R17 ldd R18, X+2 ; Load the third byte of
_reg_[rs1] into R18 ldd R19, X+3 ; Load the fourth byte of _reg_[rs1] into R19 ;
Load the second operand from _reg_[rs2] mov R26, R31 ; Copy the base address of
_reg_ to R26 (X-register) add R26, rs2 ; Add the rs2 offset to R26 (X-register)
ldd R20, X+0 ; Load the first byte of _reg_[rs2] into R20 ldd R21, X+1 ; Load
the second byte of _reg_[rs2] into R21 ldd R22, X+2 ; Load the third byte of
_reg_[rs2] into R22 ldd R23, X+3 ; Load the fourth byte of _reg_[rs2] into R23 ;
Perform bitwise OR or R16, R20 ; OR the first bytes or R17, R21 ; OR the second
bytes or R18, R22 ; OR the third bytes or R19, R23 ; OR the fourth bytes ; Store
the result in _reg_[rd] mov R26, R31 ; Copy the base address of _reg_ to R26 (X-
register) add R26, rd ; Add the rd offset to R26 (X-register) st X+0, R16 ;
Store the first byte of the result st X+1, R17 ; Store the second byte of the
result st X+2, R18 ; Store the third byte of the result st X+3, R19 ; Store the
fourth byte of the result

```

This code snippet is similar to the previous ones for addition, subtraction, XOR, and AND, but uses the **or** instruction for bitwise OR. The X-register (R26) is used as an index register to access the values without changing the Z-register (R30 and R31).



now please convert this, take in mind that you need to shift entire 32bit value, so using some rol over carry or so? i dont know which instructions are in avrx ...

```
_reg_[rd] = _reg_[rs1] << (_reg_[rs2] & 0x1F);
```



Here's the AVR assembly code snippet to perform the 32-bit left shift operation *reg[rd] = reg[rs1] << (reg[rs2] & 0x1F)*, assuming that *reg* points to a 256-byte aligned address, and *rs1*, *rs2*, and *rd* are 8-bit indexes:

assembly

Copy code

```

; Load the first operand from _reg_[rs1] mov R26, R31 ; Copy the base address of
_reg_ to R26 (X-register) add R26, rs1 ; Add the rs1 offset to R26 (X-register)
ldd R16, X+0 ; Load the first byte of _reg_[rs1] into R16 ldd R17, X+1 ; Load
the second byte of _reg_[rs1] into R17 ldd R18, X+2 ; Load the third byte of
_reg_[rs1] into R18 ldd R19, X+3 ; Load the fourth byte of _reg_[rs1] into R19 ;
Load the second operand from _reg_[rs2] mov R26, R31 ; Copy the base address of
_reg_ to R26 (X-register) add R26, rs2 ; Add the rs2 offset to R26 (X-register)
ldd R20, X+0 ; Load the first byte of _reg_[rs2] into R20 ; Calculate the shift

```

```

amount andi R20, 0x1F ; AND with 0x1F to get the lower 5 bits ; Perform 32-bit
left shift ShiftLoop: sbrlc R20, 0 ; Test the least significant bit of R20 rjmp
ShiftByte ; If the bit is set, perform a byte shift rjmp ShiftBit ; If the bit
is clear, perform a bit shift ShiftByte: mov R16, R17 ; Shift bytes mov R17, R18
mov R18, R19 clr R19 subi R20, 8 ; Subtract 8 from the shift count rjmp
ShiftLoop ShiftBit: rol R16 ; Perform bit shift with carry rol R17 rol R18 rol
R19 dec R20 ; Decrement the shift count brne ShiftLoop ; Repeat the loop until
the shift count reaches 0 ; Store the result in _reg_[rd] mov R26, R31 ; Copy
the base address of _reg_ to R26 (X-register) add R26, rd ; Add the rd offset to
R26 (X-register) st X+0, R16 ; Store the first byte of the result st X+1, R17 ;
Store the second byte of the result st X+2, R18 ; Store the third byte of the
result st X+3, R19 ; Store the fourth byte of the result

```

This code snippet first loads the values of *reg*[rs1] and *reg*[rs2] into registers R16-R19 and R20, respectively. It then calculates the shift amount and performs a loop to shift the 32-bit value in R16-R19 to the left. The loop handles byte shifts first and then bit shifts to optimize the shifting



can you try to convert also this? it is loading 32bit word from memory into *_reg_* indexed register:

```
{ TU8* ptr = (TU8*)&_data[_reg_[rs1] + imm]; _reg_[rd] = *(TS32*)ptr; }
```



>

Here's the AVR assembly code snippet to perform the 32-bit word load operation { TU8* ptr = (TU8*)&data[reg[rs1] + imm]; reg[rd] = (TS32)ptr; }, assuming that *data* and *reg* point to 256-byte aligned addresses, and rs1 and rd are 8-bit indexes:

assembly

Copy code

```

; Load the base address of _data_ ldi R26, low(_data_) ldi R27, high(_data_) ;
Load _reg_[rs1] into R20-R23 mov R30, R31 ; Copy the base address of _reg_ to
R30 (Z-register) add R30, rs1 ; Add the rs1 offset to R30 (Z-register) ldd R20,
Z+0 ; Load the first byte of _reg_[rs1] into R20 ldd R21, Z+1 ; Load the second
byte of _reg_[rs1] into R21 ldd R22, Z+2 ; Load the third byte of _reg_[rs1]
into R22 ldd R23, Z+3 ; Load the fourth byte of _reg_[rs1] into R23 ; Add the
immediate value to _reg_[rs1] ldi R24, low(imm) ; Load the lower 8 bits of imm
into R24 ldi R25, high(imm) ; Load the upper 8 bits of imm into R25 add R20, R24

```

```
; Add the lower 8 bits of the immediate value adc R21, R25 ; Add the upper 8
bits of the immediate value and carry adc R22, R1 ; Add carry (R1 is assumed to
be cleared beforehand) adc R23, R1 ; Add carry (R1 is assumed to be cleared
beforehand) ; Add the result to the base address of _data_ add R26, R20 adc R27,
R21 ; Load the 32-bit word from memory ld R16, X+0 ; Load the first byte of the
memory location into R16 ld R17, X+1 ; Load the second byte of the memory
location into R17 ld R18, X+2 ; Load the third byte of the memory location into
R18 ld R19, X+3 ; Load the fourth byte of the memory location into R19 ; Store
the loaded word into _reg_[rd] mov R30, R31 ; Copy the base address of _reg_ to
R30 (Z-register) add R30, rd ; Add the rd offset to R30 (Z-register) st Z+0, R16
; Store the first byte of the result st Z+1, R17 ; Store the second byte of the
result st Z+2, R18 ; Store the third byte of the result st Z+3, R19 ; Store the
fourth byte of the result
```

This code snippet first loads the base address of *data*, then loads *reg[rs1]* into R20-R23, and adds the immediate value *imm* to it. The resulting address is added to the base address of *data*, and the 32-bit word is loaded from memory into R16-R
