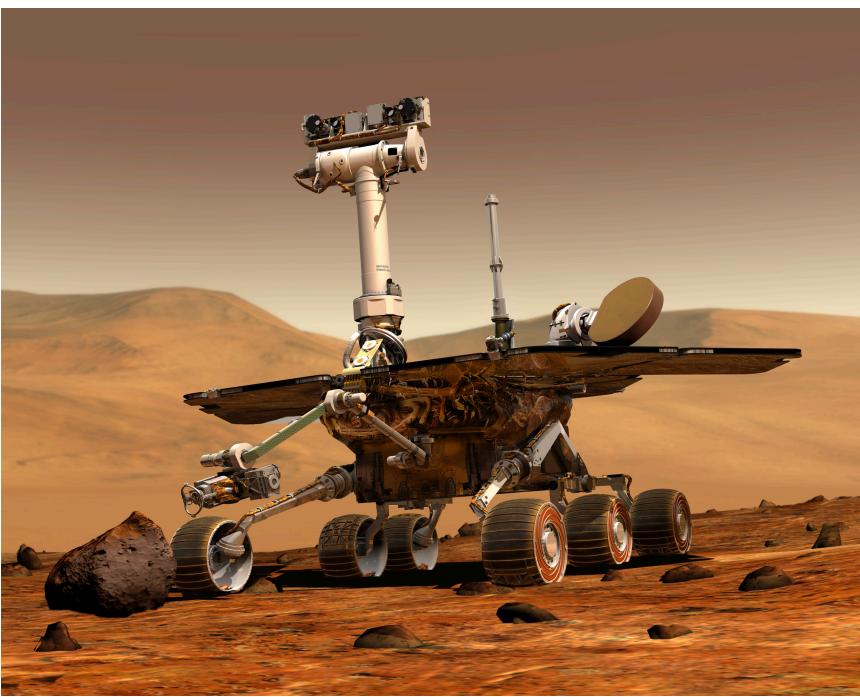
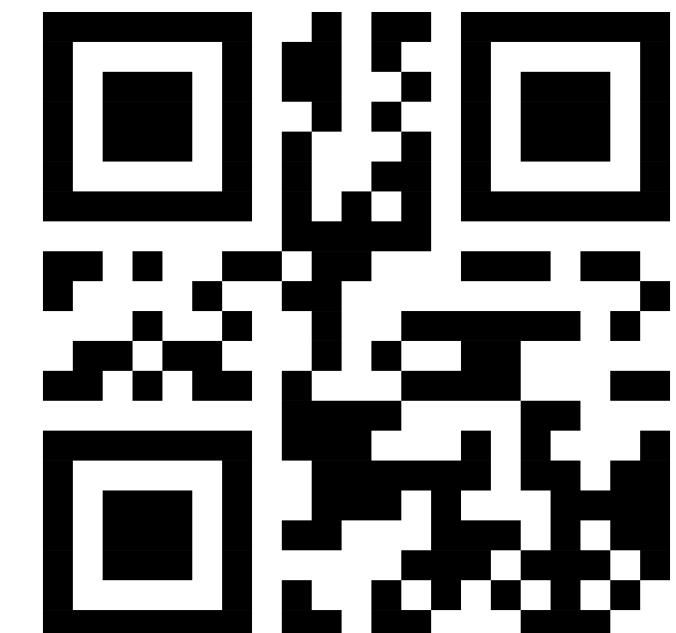
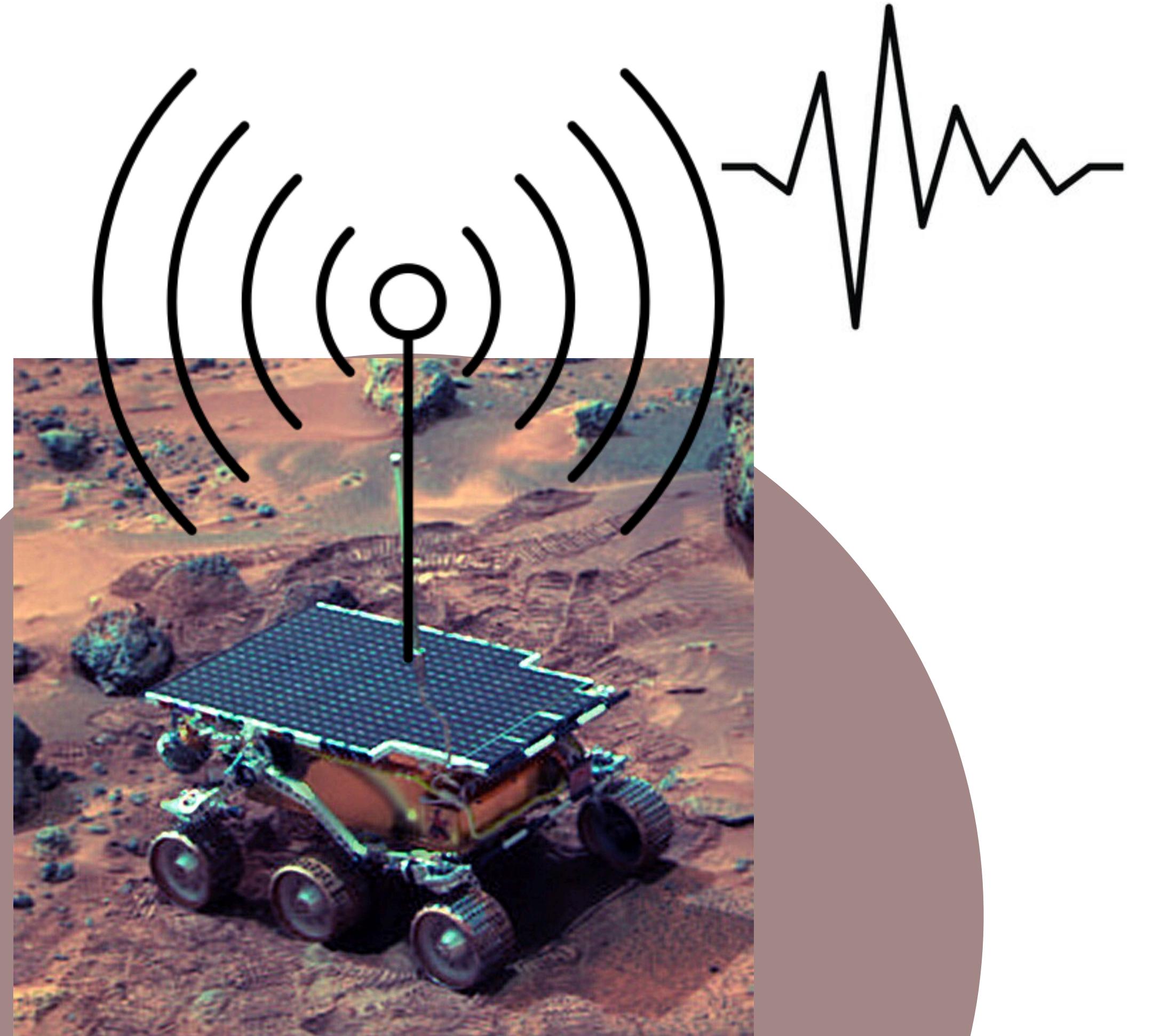


# Signal recovery from very noisy codes

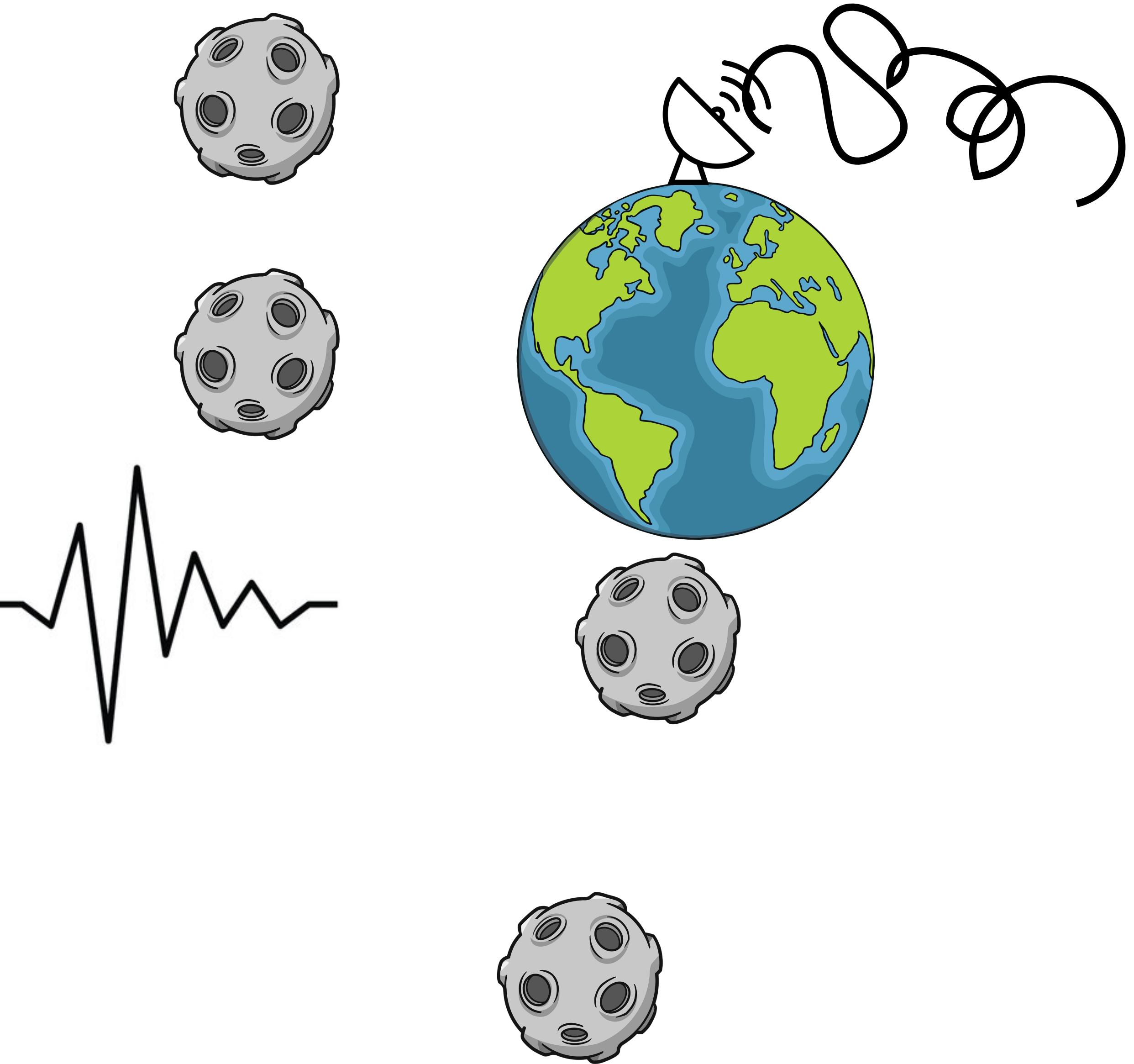


Anqi Li  
University of Cambridge



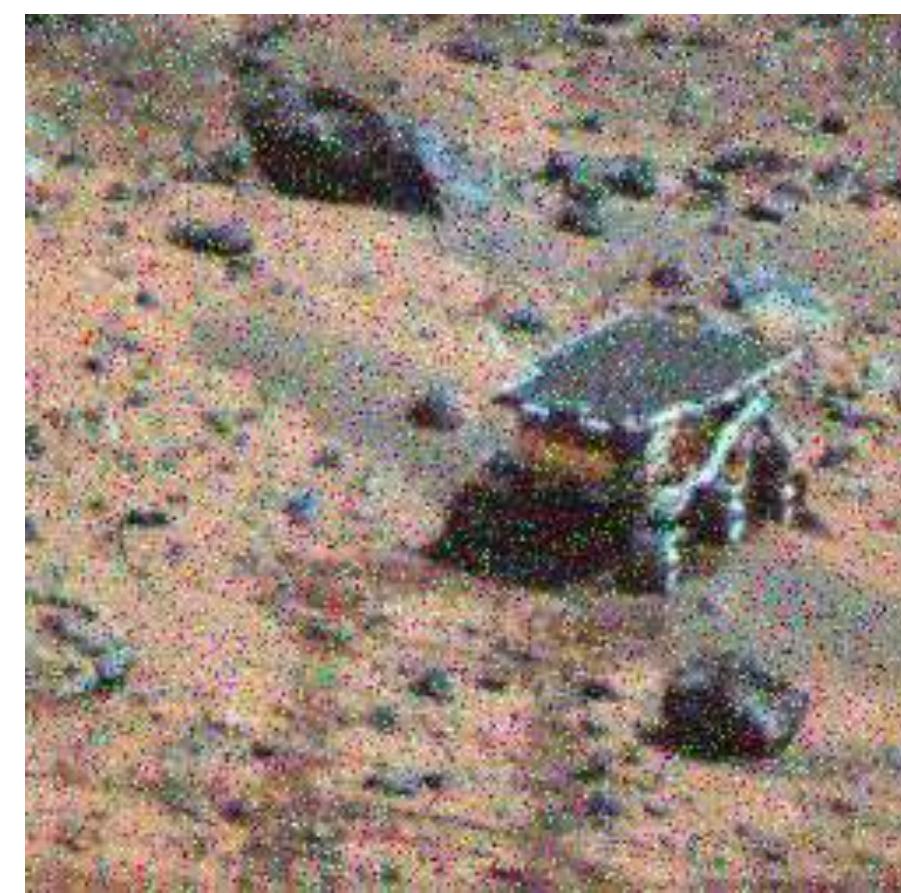


Sojourner (Mars rover)





0 %



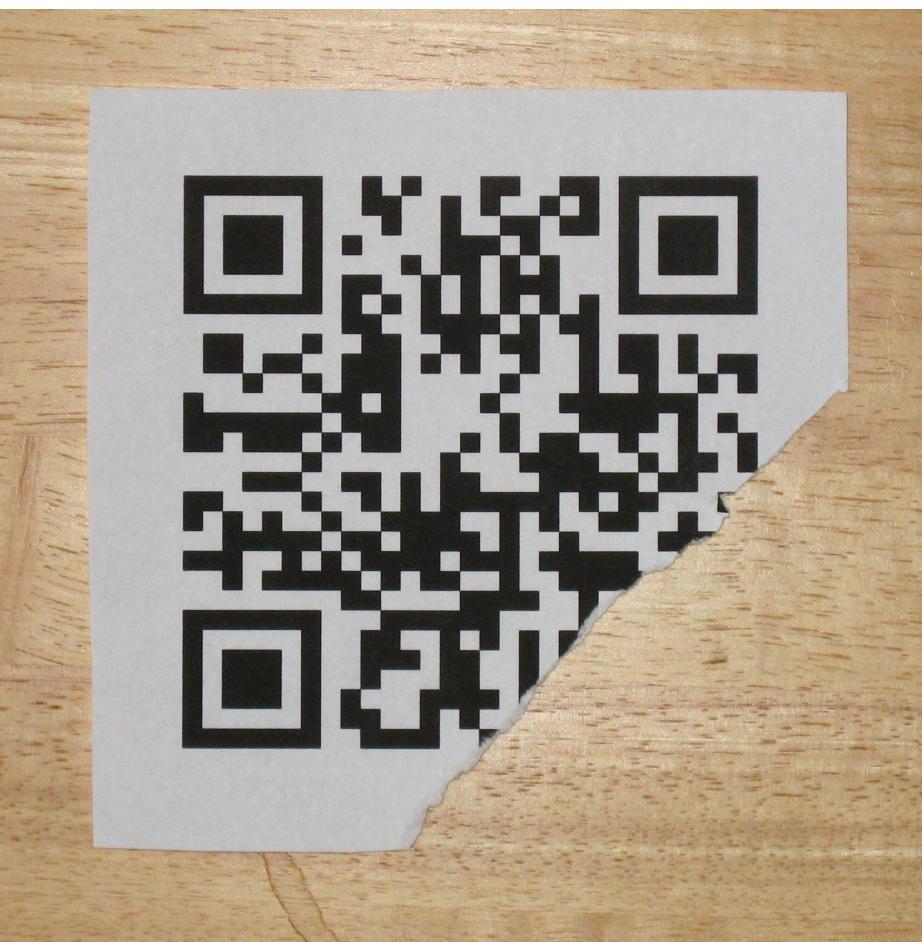
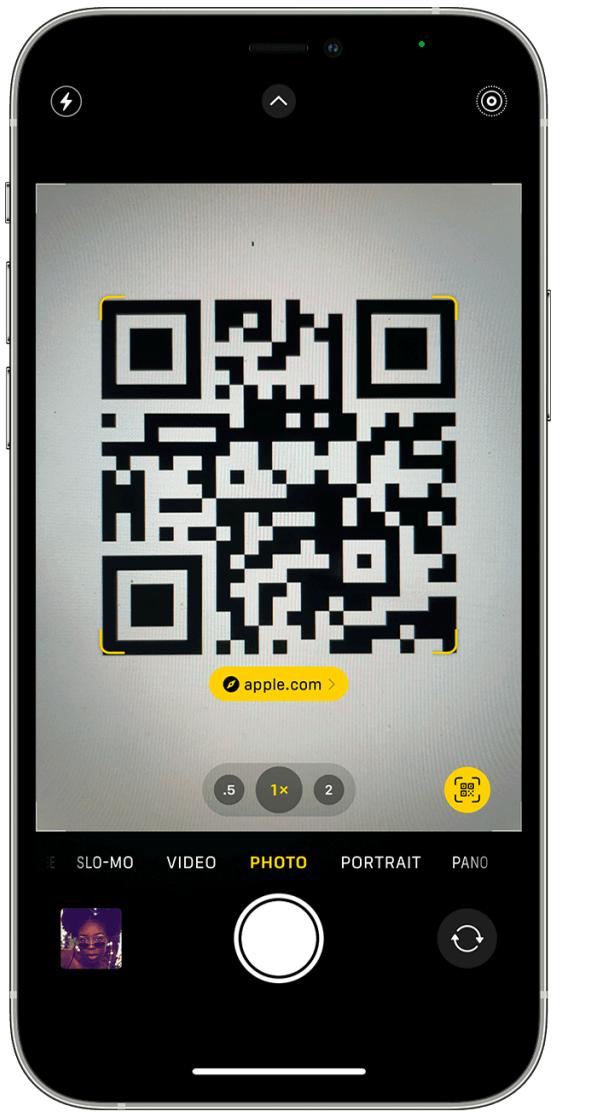
5 %



20 %

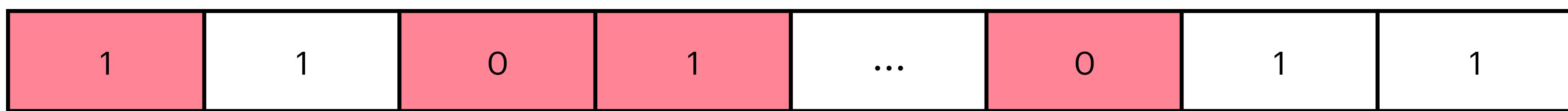
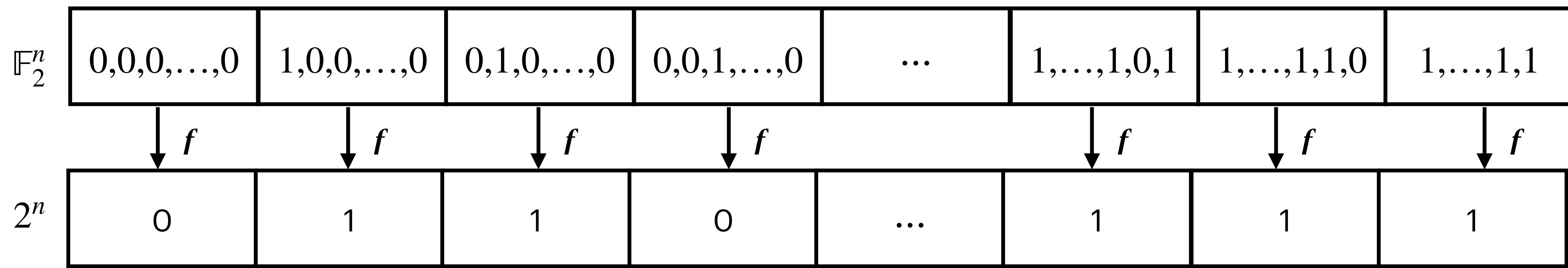


40 %



# Error-correcting Reed-Muller codes

encode polynomial  $f: \mathbb{F}_2^n \rightarrow \mathbb{F}_2$  in terms of its outputs



# Reed-Muller codes

Uncorrupted

0	1	1	0	...	1	1	1
---	---	---	---	-----	---	---	---

Corrupted

1	1	0	1	...	0	1	1
---	---	---	---	-----	---	---	---

- If 1% of the bits are corrupted, can we efficiently recover the original bits? 
- If 70% of the bits are corrupted, can we ~~efficiently~~ recover the original bits? 
- If 49% of the bits are corrupted, can we efficiently recover the original bits?  ish

# List decoding

## When there's lots of noise

- **Input:**  $(50 - \varepsilon)$  % corrupted output of degree  $d$  polynomial
- **Given:** query access to the corrupted  $2^n$  bit string
- **Goal:** Output a list of candidate polynomials (with high probability) using polynomial in  $n$  number of queries



# List decoding

## When there's lots of noise

- **Input:**  $(50 - \varepsilon)$  % corrupted output of degree  $d$  polynomial
- **Goal:** Output a list of possible starting polynomials efficiently

### History

- $d = 1$ : Goldreich-Levin (1989); applications in *cryptography, learning theory, ...*
- $d = 2$ : Quadratic Goldreich-Levin (2011 - Tulsiani, Wolf)
- $d = 3$ : **Cubic Goldreich-Levin** (2023 - Kim, L., Tidor)

# Scenario with less noise

- **Input:** 20 % corrupted output of linear function
- **Given:** query access to the corrupted  $2^n$  bit string
- **Goal:** Output the actual value of the  $i$ -th bit with high probability efficiently

$\mathbb{F}_2^4$	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
	1	0	1	1	0	1	0	0	1	0	?	1	0	1	1	0

1	0	1	1	1	0	0	0	1	0	?	1	1	1	1	0	0
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
$g$	1	0	1	1	1	0	0	0	1	0	?	1	1	1	0	0



$g(11) = g(11 + b) - g(b)$  for  $\geq 60\%$  of  $b \in \{1, \dots, 16\}$



When there is 49% noise, this no longer decodes uniquely

**Differencing** picks out linear functions  $\Rightarrow$  (discrete) Fourier analysis



Large Fourier coefficients = candidate linear functions



Parseval's theorem  $\Rightarrow$  there are few such candidate functions



For polynomials, “successive differencing”  $\Rightarrow$  higher order Fourier analysis

Codes from polynomials are  
very resistant to corruption by  
noise.

# Overview of other projects

# **extremal & additive combinatorics ∩ theory CS**

# discrete Fourier analysis $\cap$ probabilistic methods

