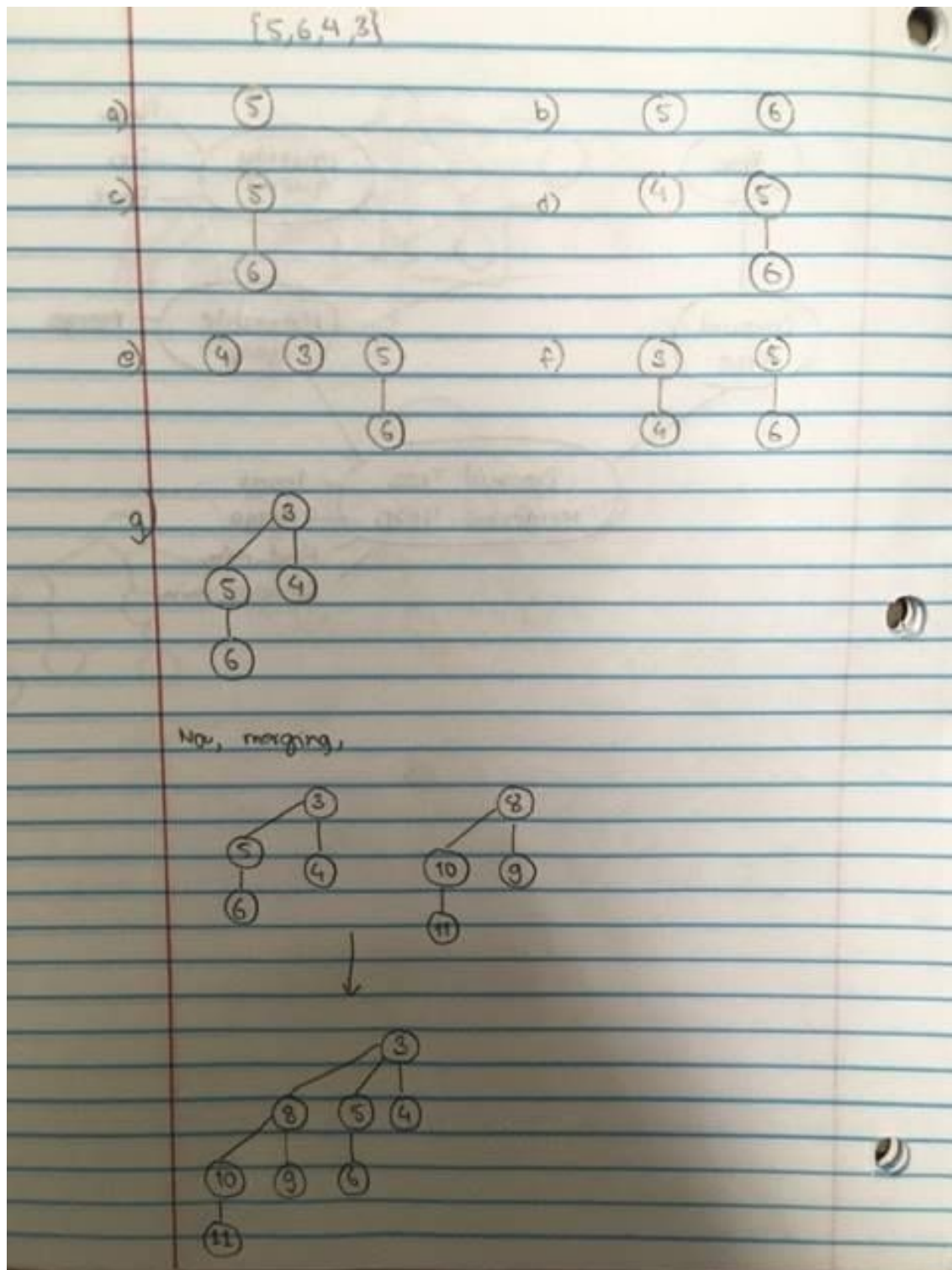
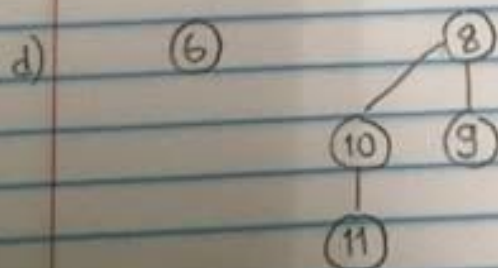
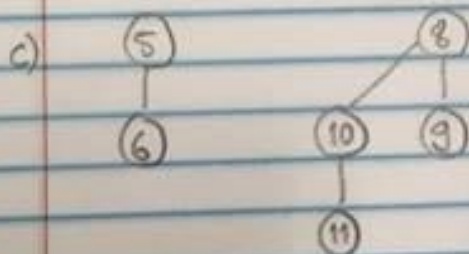
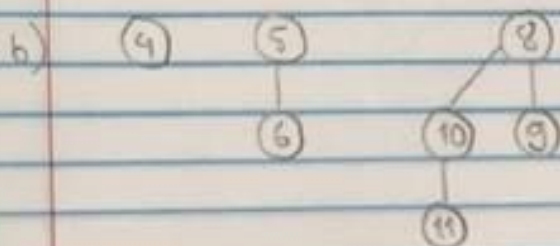
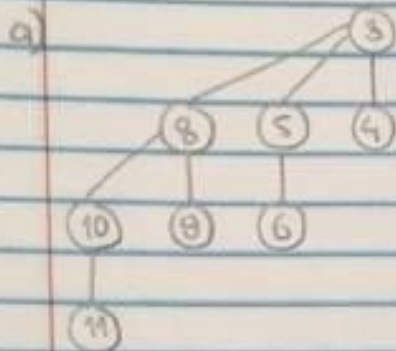


Lab 5: Merging Heaps - Tugii & Ash

Writeup A:



Not, removing 3 element



Writeup B:

The greatest advantage of a binomial heap compared to a typical binary tree heap is that binomial trees allows the heaps to be merged in $O(\log n)$ time, whereas a typical binary tree heap takes $O(n)$ time.

The greatest disadvantage of a binomial heap is that method to find the minimum takes $O(\log n)$ time in binomial heap compared to $O(1)$ for binary tree heap. Implementation for a binomial heap can also be much more complex because we need to keep track of multiple children instead of just two, as in binary heap, and multiple trees.

Writeup C:

a) Insert - $O(1)$ amortized time

To insert a new element into a heap, we can create a new heap that has only the new element and merge the two heaps. The merge will take $O(\log n)$, but the amortized time of insertion is $O(1)$.

b) Merge - $O(\log n)$

Merging two heaps with the same degree is simple since it only entails comparing the roots and adding one heap as a subtree to another. Since each tree has at most $\log n$ order, the merge operation will take $O(\log n)$.

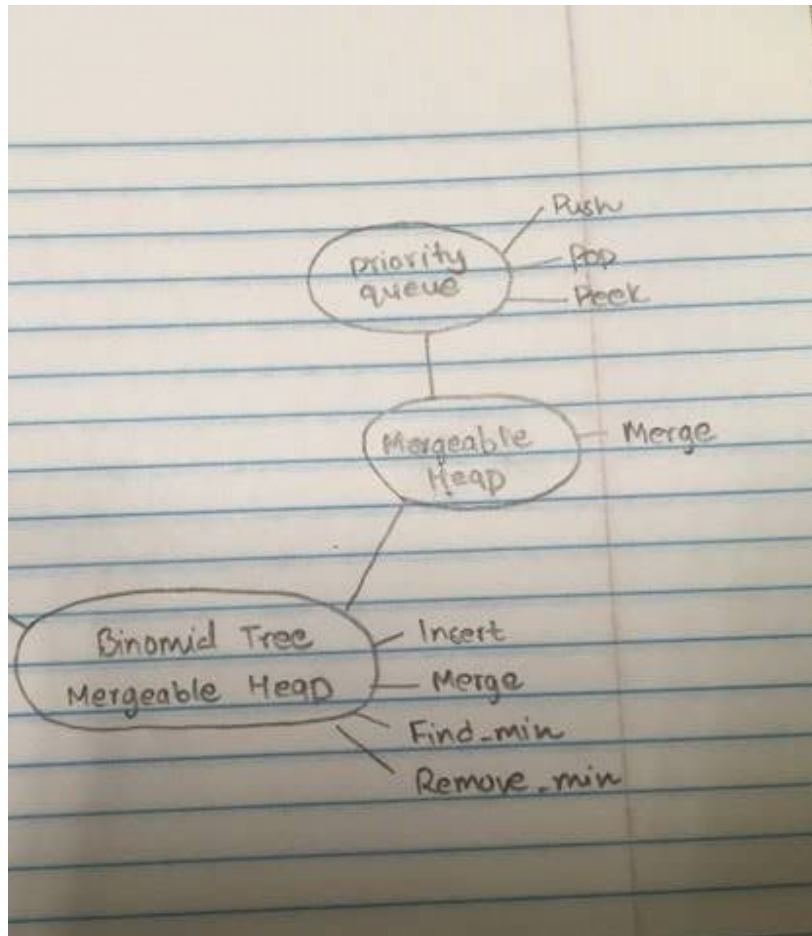
c) Find_min - $O(\log n)$

The finding minimum would take $O(\log n)$ as the number of trees in a binomial heap could be $(\log n)$ and the root nodes of all the trees need to be checked to find the minimum.

d) Remove_min $O(\log n)$

Removing the minimum item would take $O(\log n)$ as the number of trees in a binomial heap could be $(\log n)$ and the root nodes of all the trees need to be checked to find and remove the minimum. After removing the minimum element, we reverse the children into a separate binomial heap and merge it with the original.

Design:



Implementation of the tests, the binomial heap as well as the extension is in the attached python script.