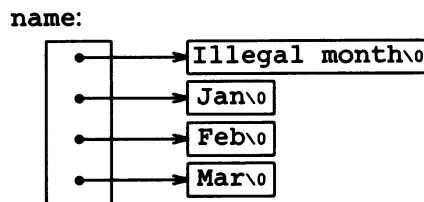


point to a twenty-element vector; some may point to two elements, some to fifty, and some to none at all.

Although we have phrased this discussion in terms of integers, by far the most frequent use of arrays of pointers is to store character strings of diverse lengths, as in the function `month_name`. Compare the declaration and picture for an array of pointers:

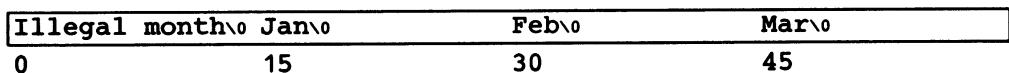
```
char *name[] = { "Illegal month", "Jan", "Feb", "Mar" };
```



with those for a two-dimensional array:

```
char aname[][15] = { "Illegal month", "Jan", "Feb", "Mar" };
```

aname:



Exercise 5-9. Rewrite the routines `day_of_year` and `month_day` with pointers instead of indexing. □

5.10 Command-line Arguments

In environments that support C, there is a way to pass command-line arguments or parameters to a program when it begins executing. When `main` is called, it is called with two arguments. The first (conventionally called `argc`, for argument count) is the number of command-line arguments the program was invoked with; the second (`argv`, for argument vector) is a pointer to an array of character strings that contain the arguments, one per string. We customarily use multiple levels of pointers to manipulate these character strings.

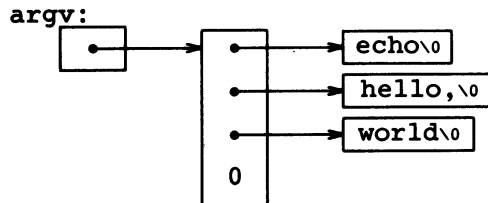
The simplest illustration is the program `echo`, which echoes its command-line arguments on a single line, separated by blanks. That is, the command

```
echo hello, world
```

prints the output

```
hello, world
```

By convention, `argv[0]` is the name by which the program was invoked, so `argc` is at least 1. If `argc` is 1, there are no command-line arguments after the program name. In the example above, `argc` is 3, and `argv[0]`, `argv[1]`, and `argv[2]` are "echo", "hello,", and "world" respectively. The first optional argument is `argv[1]` and the last is `argv[argc-1]`; additionally, the standard requires that `argv[argc]` be a null pointer.



The first version of `echo` treats `argv` as an array of character pointers:

```
#include <stdio.h>

/* echo command-line arguments; 1st version */
main(int argc, char *argv[])
{
    int i;

    for (i = 1; i < argc; i++)
        printf("%s%s", argv[i], (i < argc-1) ? " " : "");
    printf("\n");
    return 0;
}
```

Since `argv` is a pointer to an array of pointers, we can manipulate the pointer rather than index the array. This next variation is based on incrementing `argv`, which is a pointer to pointer to char, while `argc` is counted down:

```
#include <stdio.h>

/* echo command-line arguments; 2nd version */
main(int argc, char *argv[])
{
    while (--argc > 0)
        printf("%s%s", ++argv, (argc > 1) ? " " : "");
    printf("\n");
    return 0;
}
```

Since `argv` is a pointer to the beginning of the array of argument strings, incrementing it by 1 (`++argv`) makes it point at the original `argv[1]` instead of `argv[0]`. Each successive increment moves it along to the next argument; `*argv` is then the pointer to that argument. At the same time, `argc` is decremented; when it becomes zero, there are no arguments left to print.

Alternatively, we could write the `printf` statement as

```
printf((argc > 1) ? "%s " : "%s", ++argv);
```

This shows that the format argument of `printf` can be an expression too.

As a second example, let us make some enhancements to the pattern-finding program from Section 4.1. If you recall, we wired the search pattern deep into the program, an obviously unsatisfactory arrangement. Following the lead of the UNIX program `grep`, let us change the program so the pattern to be matched is specified by the first argument on the command line.

```
#include <stdio.h>
#include <string.h>
#define MAXLINE 1000

int getline(char *line, int max);

/* find: print lines that match pattern from 1st arg */
main(int argc, char *argv[])
{
    char line[MAXLINE];
    int found = 0;

    if (argc != 2)
        printf("Usage: find pattern\n");
    else
        while (getline(line, MAXLINE) > 0)
            if (strstr(line, argv[1]) != NULL) {
                printf("%s", line);
                found++;
            }
    return found;
}
```

The standard library function `strstr(s,t)` returns a pointer to the first occurrence of the string `t` in the string `s`, or `NULL` if there is none. It is declared in `<string.h>`.

The model can now be elaborated to illustrate further pointer constructions. Suppose we want to allow two optional arguments. One says “print all lines *except* those that match the pattern;” the second says “precede each printed line by its line number.”

A common convention for C programs on UNIX systems is that an argument that begins with a minus sign introduces an optional flag or parameter. If we choose `-x` (for “except”) to signal the inversion, and `-n` (“number”) to request line numbering, then the command

```
find -x -n pattern
```

will print each line that doesn’t match the pattern, preceded by its line number.

Optional arguments should be permitted in any order, and the rest of the program should be independent of the number of arguments that were present. Furthermore, it is convenient for users if option arguments can be combined, as

in

```
find -nx pattern
```

Here is the program:

```
#include <stdio.h>
#include <string.h>
#define MAXLINE 1000

int getline(char *line, int max);

/* find: print lines that match pattern from 1st arg */
main(int argc, char *argv[])
{
    char line[MAXLINE];
    long lineno = 0;
    int c, except = 0, number = 0, found = 0;

    while (--argc > 0 && (++argv)[0] == '-')
        while (c = ++argv[0])
            switch (c) {
                case 'x':
                    except = 1;
                    break;
                case 'n':
                    number = 1;
                    break;
                default:
                    printf("find: illegal option %c\n", c);
                    argc = 0;
                    found = -1;
                    break;
            }
    if (argc != 1)
        printf("Usage: find -x -n pattern\n");
    else
        while (getline(line, MAXLINE) > 0) {
            lineno++;
            if ((strstr(line, *argv) != NULL) != except) {
                if (number)
                    printf("%ld:", lineno);
                printf("%s", line);
                found++;
            }
        }
    return found;
}
```

`argc` is decremented and `argv` is incremented before each optional argument. At the end of the loop, if there are no errors, `argc` tells how many arguments remain unprocessed and `argv` points to the first of these. Thus `argc`

should be 1 and `*argv` should point at the pattern. Notice that `***argv` is a pointer to an argument string, so `(***argv)[0]` is its first character. (An alternate valid form would be `****argv`.) Because `[]` binds tighter than `*` and `++`, the parentheses are necessary; without them the expression would be taken as `***(argv[0])`. In fact, that is what we used in the inner loop, where the task is to walk along a specific argument string. In the inner loop, the expression `***argv[0]` increments the pointer `argv[0]`!

It is rare that one uses pointer expressions more complicated than these; in such cases, breaking them into two or three steps will be more intuitive.

Exercise 5-10. Write the program `expr`, which evaluates a reverse Polish expression from the command line, where each operator or operand is a separate argument. For example,

```
expr 2 3 4 + *
```

evaluates $2 \times (3+4)$. \square

Exercise 5-11. Modify the programs `entab` and `datab` (written as exercises in Chapter 1) to accept a list of tab stops as arguments. Use the default tab settings if there are no arguments. \square

Exercise 5-12. Extend `entab` and `datab` to accept the shorthand

```
entab -m +n
```

to mean tab stops every n columns, starting at column m . Choose convenient (for the user) default behavior. \square

Exercise 5-13. Write the program `tail`, which prints the last n lines of its input. By default, n is 10, let us say, but it can be changed by an optional argument, so that

```
tail -n
```

prints the last n lines. The program should behave rationally no matter how unreasonable the input or the value of n . Write the program so it makes the best use of available storage; lines should be stored as in the sorting program of Section 5.6, not in a two-dimensional array of fixed size. \square

5.11 Pointers to Functions

In C, a function itself is not a variable, but it is possible to define pointers to functions, which can be assigned, placed in arrays, passed to functions, returned by functions, and so on. We will illustrate this by modifying the sorting procedure written earlier in this chapter so that if the optional argument `-n` is given, it will sort the input lines numerically instead of lexicographically.

A sort often consists of three parts—a comparison that determines the

ordering of any pair of objects, an exchange that reverses their order, and a sorting algorithm that makes comparisons and exchanges until the objects are in order. The sorting algorithm is independent of the comparison and exchange operations, so by passing different comparison and exchange functions to it, we can arrange to sort by different criteria. This is the approach taken in our new sort.

Lexicographic comparison of two lines is done by `strcmp`, as before; we will also need a routine `numcmp` that compares two lines on the basis of numeric value and returns the same kind of condition indication as `strcmp` does. These functions are declared ahead of `main` and a pointer to the appropriate one is passed to `qsort`. We have skimped on error processing for arguments, so as to concentrate on the main issues.

```
#include <stdio.h>
#include <string.h>

#define MAXLINES 5000      /* max #lines to be sorted */
char *lineptr[MAXLINES]; /* pointers to text lines */

int readlines(char *lineptr[], int nlines);
void writelines(char *lineptr[], int nlines);

void qsort(void *lineptr[], int left, int right,
           int (*comp)(void *, void *));
int numcmp(char *, char *);

/* sort input lines */
main(int argc, char *argv[])
{
    int nlines;          /* number of input lines read */
    int numeric = 0;      /* 1 if numeric sort */

    if (argc > 1 && strcmp(argv[1], "-n") == 0)
        numeric = 1;
    if ((nlines = readlines(lineptr, MAXLINES)) >= 0) {
        qsort((void **) lineptr, 0, nlines-1,
              (int (*)(void*,void*))(numeric ? numcmp : strcmp));
        writelines(lineptr, nlines);
        return 0;
    } else {
        printf("input too big to sort\n");
        return 1;
    }
}
```

In the call to `qsort`, `strcmp` and `numcmp` are addresses of functions. Since they are known to be functions, the `&` operator is not necessary, in the same way that it is not needed before an array name.

We have written `qsort` so it can process any data type, not just character

strings. As indicated by the function prototype, `qsort` expects an array of pointers, two integers, and a function with two pointer arguments. The generic pointer type `void *` is used for the pointer arguments. Any pointer can be cast to `void *` and back again without loss of information, so we can call `qsort` by casting arguments to `void *`. The elaborate cast of the function argument casts the arguments of the comparison function. These will generally have no effect on actual representation, but assure the compiler that all is well.

```
/* qsort:  sort v[left]...v[right] into increasing order */
void qsort(void *v[], int left, int right,
           int (*comp)(void *, void *))
{
    int i, last;
    void swap(void *v[], int, int);

    if (left >= right)    /* do nothing if array contains */
        return;          /* fewer than two elements */
    swap(v, left, (left + right)/2);
    last = left;
    for (i = left+1; i <= right; i++)
        if ((*comp)(v[i], v[left]) < 0)
            swap(v, ++last, i);
    swap(v, left, last);
    qsort(v, left, last-1, comp);
    qsort(v, last+1, right, comp);
}
```

The declarations should be studied with some care. The fourth parameter of `qsort` is

```
int (*comp)(void *, void *)
```

which says that `comp` is a pointer to a function that has two `void *` arguments and returns an `int`.

The use of `comp` in the line

```
if ((*comp)(v[i], v[left]) < 0)
```

is consistent with the declaration: `comp` is a pointer to a function, `*comp` is the function, and

```
(*comp)(v[i], v[left])
```

is the call to it. The parentheses are needed so the components are correctly associated; without them,

```
int *comp(void *, void *)    /* WRONG */
```

says that `comp` is a function returning a pointer to an `int`, which is very different.

We have already shown `strcmp`, which compares two strings. Here is `numcmp`, which compares two strings on a leading numeric value, computed by

calling `atof`:

```
#include <stdlib.h>

/* numcmp: compare s1 and s2 numerically */
int numcmp(char *s1, char *s2)
{
    double v1, v2;

    v1 = atof(s1);
    v2 = atof(s2);
    if (v1 < v2)
        return -1;
    else if (v1 > v2)
        return 1;
    else
        return 0;
}
```

The `swap` function, which exchanges two pointers, is identical to what we presented earlier in the chapter, except that the declarations are changed to `void *`.

```
void swap(void *v[], int i, int j)
{
    void *temp;

    temp = v[i];
    v[i] = v[j];
    v[j] = temp;
}
```

A variety of other options can be added to the sorting program; some make challenging exercises.

Exercise 5-14. Modify the sort program to handle a `-r` flag, which indicates sorting in reverse (decreasing) order. Be sure that `-r` works with `-n`. □

Exercise 5-15. Add the option `-f` to fold upper and lower case together, so that case distinctions are not made during sorting; for example, `a` and `A` compare equal. □

Exercise 5-16. Add the `-d` (“directory order”) option, which makes comparisons only on letters, numbers and blanks. Make sure it works in conjunction with `-f`. □

Exercise 5-17. Add a field-handling capability, so sorting may be done on fields within lines, each field sorted according to an independent set of options. (The index for this book was sorted with `-df` for the index category and `-n` for the page numbers.) □

5.12 Complicated Declarations

C is sometimes castigated for the syntax of its declarations, particularly ones that involve pointers to functions. The syntax is an attempt to make the declaration and the use agree; it works well for simple cases, but it can be confusing for the harder ones, because declarations cannot be read left to right, and because parentheses are over-used. The difference between

```
int *f();          /* f: function returning pointer to int */
```

and

```
int (*pf)();      /* pf: pointer to function returning int */
```

illustrates the problem: `*` is a prefix operator and it has lower precedence than `()`, so parentheses are necessary to force the proper association.

Although truly complicated declarations rarely arise in practice, it is important to know how to understand them, and, if necessary, how to create them. One good way to synthesize declarations is in small steps with `typedef`, which is discussed in Section 6.7. As an alternative, in this section we will present a pair of programs that convert from valid C to a word description and back again. The word description reads left to right.

The first, `dcl1`, is the more complex. It converts a C declaration into a word description, as in these examples:

```
char **argv
  argv:  pointer to pointer to char
int (*daytab)[13]
  daytab: pointer to array[13] of int
int *daytab[13]
  daytab: array[13] of pointer to int
void *comp()
  comp:  function returning pointer to void
void (*comp)()
  comp:  pointer to function returning void
char ((*x())[5])()
  x:     function returning pointer to array[] of
        pointer to function returning char
char ((*x[3])())[5]
  x:     array[3] of pointer to function returning
        pointer to array[5] of char
```

`dcl1` is based on the grammar that specifies a declarator, which is spelled out precisely in Appendix A, Section 8.5; this is a simplified form:

```
dcl:          optional *'s direct-dcl
direct-dcl:   name
              (dcl)
              direct-dcl ( )
              direct-dcl [optional size]
```

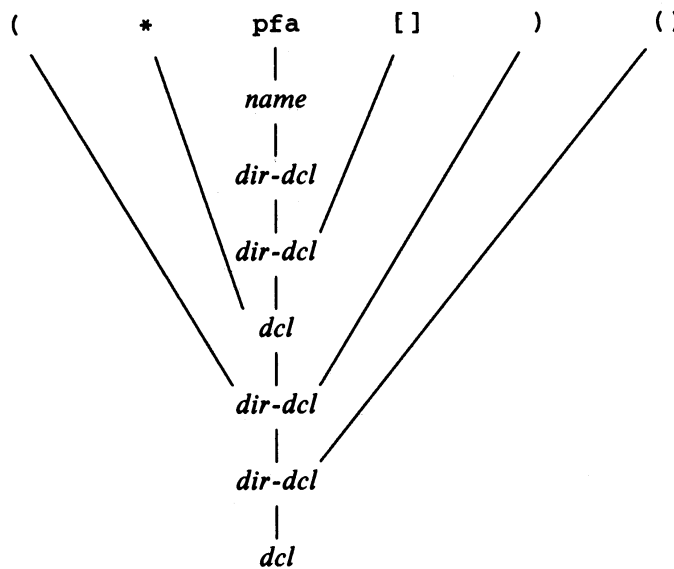
In words, a *dcl* is a *direct-dcl*, perhaps preceded by `*`'s. A *direct-dcl* is a

name, or a parenthesized *dcl*, or a *direct-dcl* followed by parentheses, or a *direct-dcl* followed by brackets with an optional size.

This grammar can be used to parse declarations. For instance, consider this declarator:

```
(*pfa[])( )
```

pfa will be identified as a *name* and thus as a *direct-dcl*. Then *pfa[]* is also a *direct-dcl*. Then **pfa[]* is recognized as a *dcl*, so *(*pfa[])* is a *direct-dcl*. Then *(*pfa[])()* is a *direct-dcl* and thus a *dcl*. We can also illustrate the parse with a parse tree like this (where *direct-dcl* has been abbreviated to *dir-dcl*):



The heart of the *dcl* program is a pair of functions, *dcl* and *dirdcl*, that parse a declaration according to this grammar. Because the grammar is recursively defined, the functions call each other recursively as they recognize pieces of a declaration; the program is called a recursive-descent parser.

```

/* dcl:  parse a declarator */
void dcl(void)
{
    int ns;

    for (ns = 0; gettoken() == '*'; )    /* count *'s */
        ns++;
    dirdcl();
    while (ns-- > 0)
        strcat(out, " pointer to");
}

```

```

/* dirdcl: parse a direct declarator */
void dirdcl(void)
{
    int type;

    if (tokentype == '(') {                /* ( dcl ) */
        dcl();
        if (tokentype != ')')
            printf("error: missing )\n");
    } else if (tokentype == NAME)          /* variable name */
        strcpy(name, token);
    else
        printf("error: expected name or (dcl)\n");
    while ((type=gettoken()) == PARENS || type == BRACKETS)
        if (type == PARENS)
            strcat(out, " function returning");
        else {
            strcat(out, " array");
            strcat(out, token);
            strcat(out, " of");
        }
}

```

Since the programs are intended to be illustrative, not bullet-proof, there are significant restrictions on `dcl`. It can only handle a simple data type like `char` or `int`. It does not handle argument types in functions, or qualifiers like `const`. Spurious blanks confuse it. It doesn't do much error recovery, so invalid declarations will also confuse it. These improvements are left as exercises.

Here are the global variables and the main routine:

```

#include <stdio.h>
#include <string.h>
#include <ctype.h>

#define MAXTOKEN 100

enum { NAME, PARENS, BRACKETS };

void dcl(void);
void dirdcl(void);

int gettoken(void);
int tokentype;          /* type of last token */
char token[MAXTOKEN];   /* last token string */
char name[MAXTOKEN];    /* identifier name */
char datatype[MAXTOKEN]; /* data type = char, int, etc. */
char out[1000];         /* output string */

```

```

main() /* convert declaration to words */
{
    while (gettoken() != EOF) { /* 1st token on line */
        strcpy(datatype, token); /* is the datatype */
        out[0] = '\0';
        dcl(); /* parse rest of line */
        if (tokentype != '\n')
            printf("syntax error\n");
        printf("%s: %s %s\n", name, out, datatype);
    }
    return 0;
}

```

The function `gettoken` skips blanks and tabs, then finds the next token in the input; a “token” is a name, a pair of parentheses, a pair of brackets perhaps including a number, or any other single character.

```

int gettoken(void) /* return next token */
{
    int c, getch(void);
    void ungetch(int);
    char *p = token;

    while ((c = getch()) == ' ' || c == '\t')
        ;
    if (c == '(') {
        if ((c = getch()) == ')') {
            strcpy(token, "()");
            return tokentype = PARENS;
        } else {
            ungetch(c);
            return tokentype = '(';
        }
    } else if (c == '[') {
        for (*p++ = c; (*p++ = getch()) != ']'; )
            ;
        *p = '\0';
        return tokentype = BRACKETS;
    } else if (isalpha(c)) {
        for (*p++ = c; isalnum(c = getch()); )
            *p++ = c;
        *p = '\0';
        ungetch(c);
        return tokentype = NAME;
    } else
        return tokentype = c;
}

```

`getch` and `ungetch` were discussed in Chapter 4.

Going in the other direction is easier, especially if we do not worry about generating redundant parentheses. The program `undcl` converts a word

description like “*x* is a function returning a pointer to an array of pointers to functions returning *char*,” which we will express as

```
x () * [] * () char
```

to

```
char ((*x())[ ]())
```

The abbreviated input syntax lets us reuse the `gettoken` function. `undcl` also uses the same external variables as `dcl` does.

```
/* undcl: convert word description to declaration */
main()
{
    int type;
    char temp[MAXTOKEN];

    while (gettoken() != EOF) {
        strcpy(out, token);
        while ((type = gettoken()) != '\n')
            if (type == PARENS || type == BRACKETS)
                strcat(out, token);
            else if (type == '*') {
                sprintf(temp, "(%s)", out);
                strcpy(out, temp);
            } else if (type == NAME) {
                sprintf(temp, "%s %s", token, out);
                strcpy(out, temp);
            } else
                printf("invalid input at %s\n", token);
        printf("%s\n", out);
    }
    return 0;
}
```

Exercise 5-18. Make `dcl` recover from input errors. □

Exercise 5-19. Modify `undcl` so that it does not add redundant parentheses to declarations. □

Exercise 5-20. Expand `dcl` to handle declarations with function argument types, qualifiers like `const`, and so on. □

CHAPTER 6: Structures

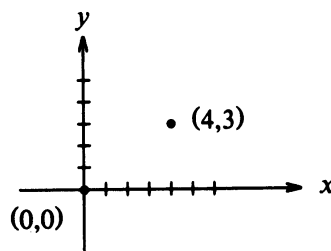
A structure is a collection of one or more variables, possibly of different types, grouped together under a single name for convenient handling. (Structures are called “records” in some languages, notably Pascal.) Structures help to organize complicated data, particularly in large programs, because they permit a group of related variables to be treated as a unit instead of as separate entities.

One traditional example of a structure is the payroll record: an employee is described by a set of attributes such as name, address, social security number, salary, etc. Some of these in turn could be structures: a name has several components, as does an address and even a salary. Another example, more typical for C, comes from graphics: a point is a pair of coordinates, a rectangle is a pair of points, and so on.

The main change made by the ANSI standard is to define structure assignment—structures may be copied and assigned to, passed to functions, and returned by functions. This has been supported by most compilers for many years, but the properties are now precisely defined. Automatic structures and arrays may now also be initialized.

6.1 Basics of Structures

Let us create a few structures suitable for graphics. The basic object is a point, which we will assume has an x coordinate and a y coordinate, both integers.



The two components can be placed in a structure declared like this:

```
struct point {
    int x;
    int y;
};
```

The keyword `struct` introduces a structure declaration, which is a list of declarations enclosed in braces. An optional name called a *structure tag* may follow the word `struct` (as with `point` here). The tag names this kind of structure, and can be used subsequently as a shorthand for the part of the declaration in braces.

The variables named in a structure are called *members*. A structure member or tag and an ordinary (i.e., non-member) variable can have the same name without conflict, since they can always be distinguished by context. Furthermore, the same member names may occur in different structures, although as a matter of style one would normally use the same names only for closely related objects.

A `struct` declaration defines a type. The right brace that terminates the list of members may be followed by a list of variables, just as for any basic type. That is,

```
struct { ... } x, y, z;
```

is syntactically analogous to

```
int x, y, z;
```

in the sense that each statement declares `x`, `y` and `z` to be variables of the named type and causes space to be set aside for them.

A structure declaration that is not followed by a list of variables reserves no storage; it merely describes a template or the shape of a structure. If the declaration is tagged, however, the tag can be used later in definitions of instances of the structure. For example, given the declaration of `point` above,

```
struct point pt;
```

defines a variable `pt` which is a structure of type `struct point`. A structure can be initialized by following its definition with a list of initializers, each a constant expression, for the members:

```
struct point maxpt = { 320, 200 };
```

An automatic structure may also be initialized by assignment or by calling a function that returns a structure of the right type.

A member of a particular structure is referred to in an expression by a construction of the form

```
structure-name . member
```

The structure member operator `“.”` connects the structure name and the member name. To print the coordinates of the point `pt`, for instance,

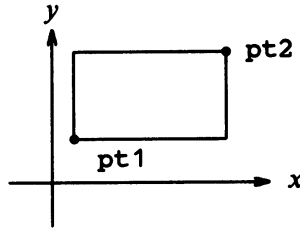
```
printf("%d,%d", pt.x, pt.y);
```

or to compute the distance from the origin (0,0) to pt,

```
double dist, sqrt(double);
```

```
dist = sqrt((double)pt.x * pt.x + (double)pt.y * pt.y);
```

Structures can be nested. One representation of a rectangle is a pair of points that denote the diagonally opposite corners:



```
struct rect {
    struct point pt1;
    struct point pt2;
};
```

The `rect` structure contains two `point` structures. If we declare `screen` as

```
struct rect screen;
```

then

```
screen.pt1.x
```

refers to the x coordinate of the `pt1` member of `screen`.

6.2 Structures and Functions

The only legal operations on a structure are copying it or assigning to it as a unit, taking its address with `&`, and accessing its members. Copy and assignment include passing arguments to functions and returning values from functions as well. Structures may not be compared. A structure may be initialized by a list of constant member values; an automatic structure may also be initialized by an assignment.

Let us investigate structures by writing some functions to manipulate points and rectangles. There are at least three possible approaches: pass components separately, pass an entire structure, or pass a pointer to it. Each has its good points and bad points.

The first function, `makepoint`, will take two integers and return a point structure:


```

/* makepoint: make a point from x and y components */
struct point makepoint(int x, int y)
{
    struct point temp;

    temp.x = x;
    temp.y = y;
    return temp;
}

```

Notice that there is no conflict between the argument name and the member with the same name; indeed the re-use of the names stresses the relationship.

`makepoint` can now be used to initialize any structure dynamically, or to provide structure arguments to a function:

```

struct rect screen;
struct point middle;
struct point makepoint(int, int);

screen.pt1 = makepoint(0, 0);
screen.pt2 = makepoint(XMAX, YMAX);
middle = makepoint((screen.pt1.x + screen.pt2.x)/2,
                  (screen.pt1.y + screen.pt2.y)/2);

```

The next step is a set of functions to do arithmetic on points. For instance,

```

/* addpoint: add two points */
struct point addpoint(struct point p1, struct point p2)
{
    p1.x += p2.x;
    p1.y += p2.y;
    return p1;
}

```

Here both the arguments and the return value are structures. We incremented the components in `p1` rather than using an explicit temporary variable to emphasize that structure parameters are passed by value like any others.

As another example, the function `ptinrect` tests whether a point is inside a rectangle, where we have adopted the convention that a rectangle includes its left and bottom sides but not its top and right sides:

```

/* ptinrect: return 1 if p in r, 0 if not */
int ptinrect(struct point p, struct rect r)
{
    return p.x >= r.pt1.x && p.x < r.pt2.x
        && p.y >= r.pt1.y && p.y < r.pt2.y;
}

```

This assumes that the rectangle is represented in a standard form where the `pt1` coordinates are less than the `pt2` coordinates. The following function returns a rectangle guaranteed to be in canonical form:

```

#define min(a, b) ((a) < (b) ? (a) : (b))
#define max(a, b) ((a) > (b) ? (a) : (b))

/* canonrect: canonicalize coordinates of rectangle */
struct rect canonrect(struct rect r)
{
    struct rect temp;

    temp.pt1.x = min(r.pt1.x, r.pt2.x);
    temp.pt1.y = min(r.pt1.y, r.pt2.y);
    temp.pt2.x = max(r.pt1.x, r.pt2.x);
    temp.pt2.y = max(r.pt1.y, r.pt2.y);
    return temp;
}

```

If a large structure is to be passed to a function, it is generally more efficient to pass a pointer than to copy the whole structure. Structure pointers are just like pointers to ordinary variables. The declaration

```
struct point *pp;
```

says that `pp` is a pointer to a structure of type `struct point`. If `pp` points to a `point` structure, `*pp` is the structure, and `(*pp).x` and `(*pp).y` are the members. To use `pp`, we might write, for example,

```

struct point origin, *pp;

pp = &origin;
printf("origin is (%d,%d)\n", (*pp).x, (*pp).y);

```

The parentheses are necessary in `(*pp).x` because the precedence of the structure member operator `.` is higher than `*`. The expression `*pp.x` means `*(pp.x)`, which is illegal here because `x` is not a pointer.

Pointers to structures are so frequently used that an alternative notation is provided as a shorthand. If `p` is a pointer to a structure, then

p->member-of-structure

refers to the particular member. (The operator `->` is a minus sign immediately followed by `>`.) So we could write instead

```
printf("origin is (%d,%d)\n", pp->x, pp->y);
```

Both `.` and `->` associate from left to right, so if we have

```
struct rect r, *rp = &r;
```

then these four expressions are equivalent:

```

r.pt1.x
rp->pt1.x
(r.pt1).x
(rp->pt1).x

```

The structure operators `.` and `->`, together with `()` for function calls and `[]` for subscripts, are at the top of the precedence hierarchy and thus bind very tightly. For example, given the declaration

```
struct {
    int len;
    char *str;
} *p;
```

then

```
++p->len
```

increments `len`, not `p`, because the implied parenthesization is `++(p->len)`. Parentheses can be used to alter the binding: `(++p)->len` increments `p` before accessing `len`, and `(p++)->len` increments `p` afterward. (This last set of parentheses is unnecessary.)

In the same way, `*p->str` fetches whatever `str` points to; `*p->str++` increments `str` after accessing whatever it points to (just like `*s++`); `(*p->str)++` increments whatever `str` points to; and `*p++->str` increments `p` after accessing whatever `str` points to.

6.3 Arrays of Structures

Consider writing a program to count the occurrences of each C keyword. We need an array of character strings to hold the names, and an array of integers for the counts. One possibility is to use two parallel arrays, `keyword` and `keycount`, as in

```
char *keyword[NKEYS];
int keycount[NKEYS];
```

But the very fact that the arrays are parallel suggests a different organization, an array of structures. Each keyword entry is a pair:

```
char *word;
int count;
```

and there is an array of pairs. The structure declaration

```
struct key {
    char *word;
    int count;
} keytab[NKEYS];
```

declares a structure type `key`, defines an array `keytab` of structures of this type, and sets aside storage for them. Each element of the array is a structure. This could also be written

```

struct key {
    char *word;
    int count;
};

struct key keytab[NKEYS];

```

Since the structure `keytab` contains a constant set of names, it is easiest to make it an external variable and initialize it once and for all when it is defined. The structure initialization is analogous to earlier ones—the definition is followed by a list of initializers enclosed in braces:

```

struct key {
    char *word;
    int count;
} keytab[] = {
    "auto", 0,
    "break", 0,
    "case", 0,
    "char", 0,
    "const", 0,
    "continue", 0,
    "default", 0,
    /* ... */
    "unsigned", 0,
    "void", 0,
    "volatile", 0,
    "while", 0
};

```

The initializers are listed in pairs corresponding to the structure members. It would be more precise to enclose initializers for each “row” or structure in braces, as in

```

{ "auto", 0 },
{ "break", 0 },
{ "case", 0 },
...

```

but the inner braces are not necessary when the initializers are simple variables or character strings, and when all are present. As usual, the number of entries in the array `keytab` will be computed if initializers are present and the `[]` is left empty.

The keyword-counting program begins with the definition of `keytab`. The main routine reads the input by repeatedly calling a function `getword` that fetches one word at a time. Each word is looked up in `keytab` with a version of the binary search function that we wrote in Chapter 3. The list of keywords must be sorted in increasing order in the table.

```

#include <stdio.h>
#include <ctype.h>
#include <string.h>

#define MAXWORD 100

int getword(char *, int);
int binsearch(char *, struct key *, int);

/* count C keywords */
main()
{
    int n;
    char word[MAXWORD];

    while (getword(word, MAXWORD) != EOF)
        if (isalpha(word[0]))
            if ((n = binsearch(word, keytab, NKEYS)) >= 0)
                keytab[n].count++;
    for (n = 0; n < NKEYS; n++)
        if (keytab[n].count > 0)
            printf("%4d %s\n",
                keytab[n].count, keytab[n].word);
    return 0;
}

/* binsearch: find word in tab[0]...tab[n-1] */
int binsearch(char *word, struct key tab[], int n)
{
    int cond;
    int low, high, mid;

    low = 0;
    high = n - 1;
    while (low <= high) {
        mid = (low+high) / 2;
        if ((cond = strcmp(word, tab[mid].word)) < 0)
            high = mid - 1;
        else if (cond > 0)
            low = mid + 1;
        else
            return mid;
    }
    return -1;
}

```

We will show the function `getword` in a moment; for now it suffices to say that each call to `getword` finds a word, which is copied into the array named as its first argument.

The quantity `NKEYS` is the number of keywords in `keytab`. Although we

could count this by hand, it's a lot easier and safer to do it by machine, especially if the list is subject to change. One possibility would be to terminate the list of initializers with a null pointer, then loop along `keytab` until the end is found.

But this is more than is needed, since the size of the array is completely determined at compile time. The size of the array is the size of one entry times the number of entries, so the number of entries is just

size of keytab / size of struct key

C provides a compile-time unary operator called `sizeof` that can be used to compute the size of any object. The expressions

`sizeof object`

and

`sizeof(type name)`

yield an integer equal to the size of the specified object or type in bytes. (Strictly, `sizeof` produces an unsigned integer value whose type, `size_t`, is defined in the header `<stddef.h>`.) An object can be a variable or array or structure. A type name can be the name of a basic type like `int` or `double`, or a derived type like a structure or a pointer.

In our case, the number of keywords is the size of the array divided by the size of one element. This computation is used in a `#define` statement to set the value of `NKEYS`:

```
#define NKEYS (sizeof keytab / sizeof(struct key))
```

Another way to write this is to divide the array size by the size of a specific element:

```
#define NKEYS (sizeof keytab / sizeof keytab[0])
```

This has the advantage that it does not need to be changed if the type changes.

A `sizeof` can not be used in a `#if` line, because the preprocessor does not parse type names. But the expression in the `#define` is not evaluated by the preprocessor, so the code here is legal.

Now for the function `getword`. We have written a more general `getword` than is necessary for this program, but it is not complicated. `getword` fetches the next "word" from the input, where a word is either a string of letters and digits beginning with a letter, or a single non-white space character. The function value is the first character of the word, or `EOF` for end of file, or the character itself if it is not alphabetic.

```

/* getword: get next word or character from input */
int getword(char *word, int lim)
{
    int c, getch(void);
    void ungetch(int);
    char *w = word;

    while (isspace(c = getch()))
        ;
    if (c != EOF)
        *w++ = c;
    if (!isalpha(c)) {
        *w = '\0';
        return c;
    }
    for ( ; --lim > 0; w++)
        if (!isalnum(*w = getch())) {
            ungetch(*w);
            break;
        }
    *w = '\0';
    return word[0];
}

```

`getword` uses the `getch` and `ungetch` that we wrote in Chapter 4. When the collection of an alphanumeric token stops, `getword` has gone one character too far. The call to `ungetch` pushes that character back on the input for the next call. `getword` also uses `isspace` to skip white space, `isalpha` to identify letters, and `isalnum` to identify letters and digits; all are from the standard header `<ctype.h>`.

Exercise 6-1. Our version of `getword` does not properly handle underscores, string constants, comments, or preprocessor control lines. Write a better version. □

6.4 Pointers to Structures

To illustrate some of the considerations involved with pointers to and arrays of structures, let us write the keyword-counting program again, this time using pointers instead of array indices.

The external declaration of `keytab` need not change, but `main` and `binsearch` do need modification.

```

#include <stdio.h>
#include <ctype.h>
#include <string.h>
#define MAXWORD 100

int getword(char *, int);
struct key *binsearch(char *, struct key *, int);

/* count C keywords; pointer version */
main()
{
    char word[MAXWORD];
    struct key *p;

    while (getword(word, MAXWORD) != EOF)
        if (isalpha(word[0]))
            if ((p=binsearch(word, keytab, NKEYS)) != NULL)
                p->count++;
    for (p = keytab; p < keytab + NKEYS; p++)
        if (p->count > 0)
            printf("%4d %s\n", p->count, p->word);
    return 0;
}

/* binsearch: find word in tab[0]...tab[n-1] */
struct key *binsearch(char *word, struct key *tab, int n)
{
    int cond;
    struct key *low = &tab[0];
    struct key *high = &tab[n];
    struct key *mid;

    while (low < high) {
        mid = low + (high-low) / 2;
        if ((cond = strcmp(word, mid->word)) < 0)
            high = mid;
        else if (cond > 0)
            low = mid + 1;
        else
            return mid;
    }
    return NULL;
}

```

There are several things worthy of note here. First, the declaration of `binsearch` must indicate that it returns a pointer to `struct key` instead of an integer; this is declared both in the function prototype and in `binsearch`. If `binsearch` finds the word, it returns a pointer to it; if it fails, it returns `NULL`.

Second, the elements of `keytab` are now accessed by pointers. This

requires significant changes in `binsearch`.

The initializers for `low` and `high` are now pointers to the beginning and just past the end of the table.

The computation of the middle element can no longer be simply

```
mid = (low+high) / 2    /* WRONG */
```

because the addition of two pointers is illegal. Subtraction is legal, however, so `high-low` is the number of elements, and thus

```
mid = low + (high-low) / 2
```

sets `mid` to point to the element halfway between `low` and `high`.

The most important change is to adjust the algorithm to make sure that it does not generate an illegal pointer or attempt to access an element outside the array. The problem is that `&tab[-1]` and `&tab[n]` are both outside the limits of the array `tab`. The former is strictly illegal, and it is illegal to dereference the latter. The language definition does guarantee, however, that pointer arithmetic that involves the first element beyond the end of an array (that is, `&tab[n]`) will work correctly.

In `main` we wrote

```
for (p = keytab; p < keytab + NKEYS; p++)
```

If `p` is a pointer to a structure, arithmetic on `p` takes into account the size of the structure, so `p++` increments `p` by the correct amount to get the next element of the array of structures, and the test stops the loop at the right time.

Don't assume, however, that the size of a structure is the sum of the sizes of its members. Because of alignment requirements for different objects, there may be unnamed "holes" in a structure. Thus, for instance, if a `char` is one byte and an `int` four bytes, the structure

```
struct {
    char c;
    int i;
};
```

might well require eight bytes, not five. The `sizeof` operator returns the proper value.

Finally, an aside on program format: when a function returns a complicated type like a structure pointer, as in

```
struct key *binsearch(char *word, struct key *tab, int n)
```

the function name can be hard to see, and to find with a text editor. Accordingly an alternate style is sometimes used:

```
struct key *
binsearch(char *word, struct key *tab, int n)
```

This is a matter of personal taste; pick the form you like and hold to it.

6.5 Self-referential Structures

Suppose we want to handle the more general problem of counting the occurrences of *all* the words in some input. Since the list of words isn't known in advance, we can't conveniently sort it and use a binary search. Yet we can't do a linear search for each word as it arrives, to see if it's already been seen; the program would take too long. (More precisely, its running time is likely to grow quadratically with the number of input words.) How can we organize the data to cope efficiently with a list of arbitrary words?

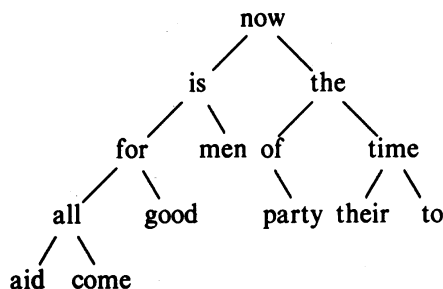
One solution is to keep the set of words seen so far sorted at all times, by placing each word into its proper position in the order as it arrives. This shouldn't be done by shifting words in a linear array, though—that also takes too long. Instead we will use a data structure called a *binary tree*.

The tree contains one “node” per distinct word; each node contains

- a pointer to the text of the word
- a count of the number of occurrences
- a pointer to the left child node
- a pointer to the right child node

No node may have more than two children; it might have only zero or one.

The nodes are maintained so that at any node the left subtree contains only words that are lexicographically less than the word at the node, and the right subtree contains only words that are greater. This is the tree for the sentence “now is the time for all good men to come to the aid of their party”, as built by inserting each word as it is encountered:



To find out whether a new word is already in the tree, start at the root and compare the new word to the word stored at that node. If they match, the question is answered affirmatively. If the new word is less than the tree word, continue searching at the left child, otherwise at the right child. If there is no child in the required direction, the new word is not in the tree, and in fact the empty slot is the proper place to add the new word. This process is recursive, since the search from any node uses a search from one of its children. Accordingly, recursive routines for insertion and printing will be most natural.

Going back to the description of a node, it is conveniently represented as a structure with four components: