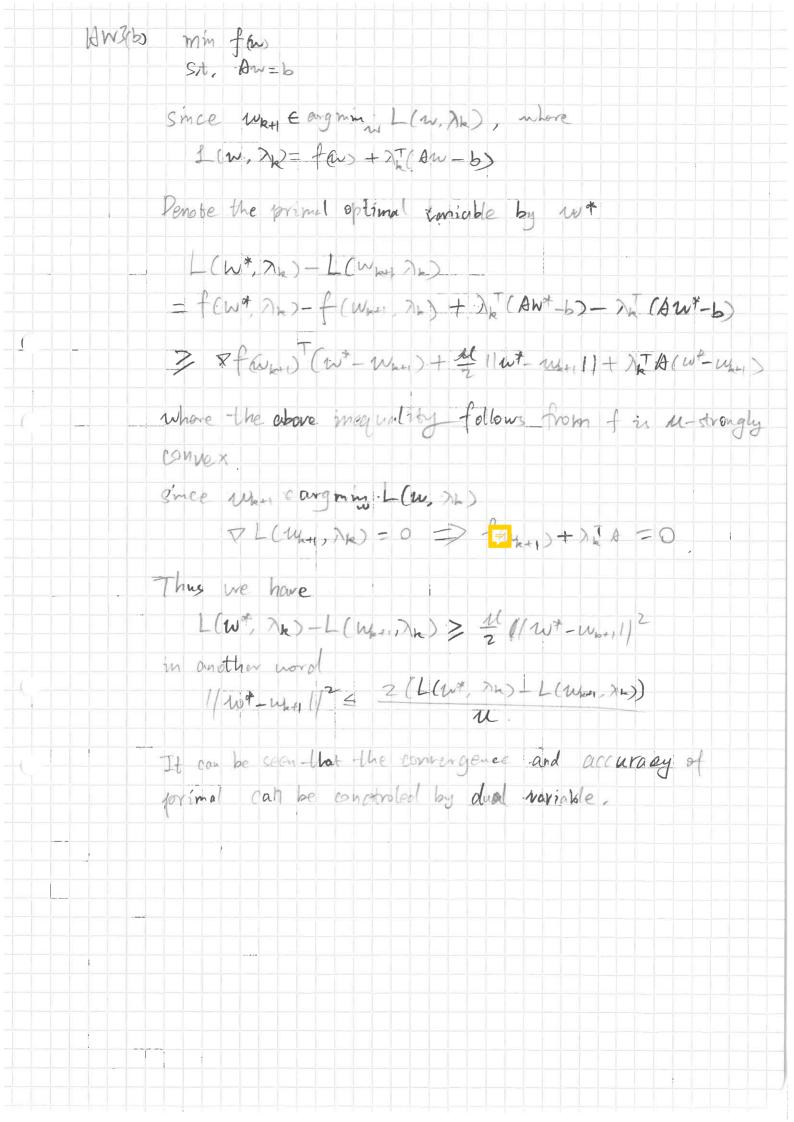
Overall passed. Hwz 10) min fix> s+ 1+x=6 For convex and cliesed function f, thereis ft - T Looks fine! x+ 2f797 + 2 f**(x) =) y + 2 fw 46 8fix) =) f (w-fix) Zy'(u-x), VM => y a - fux = g x - f(x) - 6-11. >fry) = max(gTu -fau)= yTx -fax) > x= argmin fun-yTu also let f*(1)= max (vin-fr=) Z f*(y)+ x*(v-y)=> x+ x*(y)
The duen problem. max ga)=-f*(-Aix) - Aib Since for comes and dase function of xt of type yt of(x) (=) x = aig min fun) - gty DWESTIL-ATA) & Weargmin (tw) + ATAW.

-> AW-6+ + 9CA) & Weargmin funt AW Thus Hwb & agran for on We arguin fair + STAW -



	Hw3 (6) - Convergence analysis of dual canable
	$11\lambda^{k+1} - \lambda^{\dagger} _{2}^{2} = 11\lambda^{k} + \lambda k \left(Aw_{k} - b\right) - \lambda^{\dagger} _{2}^{2}$
	= 111/2 x + 20/k (Aur - 6, x 1 - x + > + d k 2 1 + 4 mk - 6 11)
	f: u-strong L-5 mooth => 9 1 - strong comes and II - smooth =
	those 13, 11) 141- 121/2 < (1-20/2) 11/2 - 20/9 (21)-90% + 2/2 11/44/2
-(
	(By wany &-Strong convexity). Eurther
	Further Also, there is $ \lambda^{k+1} - \lambda^* ^2 \le C -\lambda_k / 2 $ $ \lambda^k - \lambda^* + ^2 - 2 \delta_k (g \circ k^*) - g \circ g$
	$+2d\kappa^{2}\left(g\alpha^{(k)}-g(\lambda^{2})\right)$
-(,	= (1-dk/L) // \\ - 2dk (f-dk) (qu')-qu') Convergence quaranteed when
-(,	oct-ak/coloned 1-dk 20 => 2k5U
	and or - de 1 Prince 20511
	Men then require 474 Apart form the given comments, looks fine.
	And the convergence vote is this I

minimize $\sum_{\omega',\ldots,\omega'}^{N} f^{i}(\omega')$ z.t. $\omega' = \omega^2 = \cdots = \omega^{N}$ where there are N numbers of nodes Let wij represents node i's estimate of nodes j's internal state. Let λij , $i,j \in \{1,2,\cdots,N\}$, be the dual variables. For instance Dij is associated with the constraint Let $\lambda = \begin{bmatrix} \lambda_{11}^{\mathsf{T}}, \cdots, \lambda_{1N}^{\mathsf{T}}, \lambda_{21}^{\mathsf{T}}, \cdots, \lambda_{2N}^{\mathsf{T}}, \cdots, \lambda_{NN} \end{bmatrix}^{\mathsf{T}}$ The Langrangian: $L(\omega', \omega^2, \dots, \omega', \lambda) = \sum_{i=1}^{N} f^i(\omega_i, \omega_2', \dots, \omega_N')$ $+\sum_{i=1}^{N}\sum_{i=1}^{N}\lambda_{ij}^{T}\left(\omega_{i}^{2}-\omega_{i}^{3}\right)$ The dual function $g(\lambda) = \inf_{\omega', \dots, \omega^N \in \mathcal{M}} L(\omega', \dots, \omega^N, \lambda)$ We can write the dual function in subproblems such as $\phi^{i}(\lambda) = \inf_{\omega \in W} f^{i}(\omega_{i}^{i}, \dots, \omega_{N}^{i}) + \sum_{j=1}^{N} \lambda_{ij}^{T} \cdot \omega_{i}^{i} - \sum_{j=1}^{N} \lambda_{ji}^{T} \cdot \omega_{j}^{i}$ Since $\phi'(\lambda)$ only depends on λ and w^i , node i can compute $\phi'(\lambda)$ locally.

The Langrangian is the sum $g(\lambda) = \sum_{i=1}^{N} \phi^{i}(\lambda)$ Finally, the dual problem is the maximization $g^* = \max_{\lambda} g(\lambda) = \max_{\lambda} \sum_{i=1}^{N} \phi^i(\lambda)$ Communication cost Suppose that there are N nodes, each part is transferred over N-1 nodes. Summing over all parts will give the the communication cost for primal method $\sum_{i=1}^{N} (N-1) \operatorname{dim}(w_{i}^{2}) = (N-1) = \operatorname{dim}(w)$ Similarly, the communication cost for dual method $\sum_{i=1}^{N} \sum_{i=1}^{N} (N-1) \dim \lambda_{ij} = (N-1) \lim_{n \to \infty} (\lambda)$ Since $\dim(\lambda)$ > $\dim(\omega)$, the communication cost of the dual method is higher. But, one can decrease this amount since a mode only needs the sum of dual variables, \$\frac{1}{2} \lambda_{ij}, for each iteration.