Module 2 - HW 1 - Extracting Parallelism

Extracting dependency from code is an almost automatic process. You need to choose a granularity. But once that is chosen, the entire analysis follows.

In the whole activity, you should express the metrics in complexity notation as a function of the parameters of the functions.

1 Coin Collection (from Midterm Spring 2018) (50 pts)

The Coin Collection problem is defined as follows:

Several coins are placed on an $n \times m$ board with at most one coin per cell of the board. A robot is initially located at the upper left cell of the board. The robot can only move to the right or down; it can not move up or left. When the robot visits a cell where a coin is located, it picks it up. At most, how many coins can the robot collect?

This problem can be solved by the following method:

```
void RobotCoin(int n, int m, //size of the board
                     int C[n][m] //Is there a coin in (i,j)
  int F[n][m]; //How many coins can be collected while on (i,j)
  F[0][0] = C[0][0];
  for (int k=1; k < m; ++k) {
     F[0][k] = F[0][k-1] + C[0][k];
  for (int i=1; i < n; ++i) {
     F[i][0] = F[i-1][0] + C[i][0];
     \quad \textbf{for} \quad (\textbf{int} \quad j = 1; \quad j < m; \quad +\!\!\!+\!\! j \;) \quad \{
        F\,[\,\,i\,\,]\,[\,\,j\,\,] \;=\; \max\;\; (F\,[\,\,i\,\,-1]\,[\,\,j\,\,]\,\,,\;\; F\,[\,\,i\,\,]\,[\,\,j\,-1]) \;+\; C\,[\,\,i\,\,]\,[\,\,j\,\,]\,;
  }
  return F[n-1][m-1];
Question: What is the complexity of this function?
Question: Extract the dependencies.
Question: What is the width?
Question: What is the work?
Question: What is the critical path? What is its length?
```

2 Knapsack (50 pts)

The Knapsack problem aims at finding the best set of objects to pack in a bag. Often the following dynamic programming algorithm is used to solve the problem.

```
void \text{ knapsack (int } n, \text{ int } W, \text{ int } value[], \text{ int } weight[], \text{ int } val[][]) 
  for (int a = 0; a < W; ++a) {
    val[0][a] = 0;
  for (int j=0; j<=W; ++j) {
       val[i][j] = val[i-1][j];
       \mathbf{if} (weight [i-1] \ll j) {
         val[i][j] = max (val[i-1][j], value[i-1]+val[i-1][j-weight[i-1]]);
    }
  }
  (You can assume weight is positive.)
Question: What is the complexity of this function?
Question: Extract the dependencies.
Question: What is the width?
Question: What is the work?
Question: What is the critical path? What is its length?
```