



Weekly Quiz 30

Discussion

Digital Logic:

Boolean Algebra, Minimization, Number System



Instructor:

Deepak Poonia

MTech, IISc Bangalore

GATE CSE AIR 53; AIR 67;
AIR 107; AIR 206; AIR 256

Digital Logic Complete Course:

<https://www.goclasses.in/courses/Digital-Logic>



Test Series

Here it Comes!!

GATE Overflow + GO Classes

2-IN-1 TEST SERIES

Most Awaited

GO Test Series
is Here

R E G I S T E R N O W

<http://tests.gatecse.in/>

100+

Number of tests

20+

Number of Full Length Mock Tests

15th APRIL 2023

+91 - 7906011243

+91- 6398661679

On
“**GATE Overflow**
Website



Join **GO+ GO Classes Combined Test Series** for BEST quality tests, matching GATE CSE Level:

Visit www.gateoverflow.in website to join Test Series.

1. **Quality Questions:** No Ambiguity in Questions, All Well-framed questions.
2. Correct, **Detailed Explanation**, Covering Variations of questions.
3. **Video Solutions.**

<https://gateoverflow.in/blog/14987/gate-overflow-and-go-classes-test-series-gate-cse-2024>



Join GO Classes **GATE CSE Complete Course** now:

<https://www.goclasses.in/s/pages/gatecompletecourse>

1. Quality Learning: No Rote-Learning. **Understand Everything**, from basics, **In-depth**, with variations.
2. Daily Homeworks, **Quality Practice Sets**, Weekly Quizzes.
3. **Summary Lectures** for Quick Revision.
4. Detailed Video Solutions of Previous ALL **GATE Questions**.
5. **Doubt Resolution**, **Revision**, Practice, a lot more.



Digital Logic

Download the GO Classes Android App:

<https://play.google.com/store/apps/details?id=com.goclasses.courses>

Search “GO Classes”
on Play Store.

Hassle-free learning
On the go!
Gain expert knowledge



www.goclasses.in



NOTE :

Complete Discrete Mathematics & Complete Engineering Mathematics Courses, by GO Classes, are **FREE** for ALL learners.

Visit here to watch : <https://www.goclasses.in/s/store/>

SignUp/Login on Goclasses website for free and start learning.



We are on Telegram. **Contact us** for any help.

Link in the Description!!

Join GO Classes **Doubt Discussion** Telegram Group :



Username:

@GATECSE_GOCLASSES



We are on Telegram. Contact us for any help.

Join GO Classes [Telegram Channel](#), Username: **@GOCLASSES_CSE**

Join GO Classes **Doubt Discussion** Telegram Group :

Username: **@GATECSE_Goclasses**

(Any doubt related to Goclasses Courses can also be asked here.)

Join [GATEOverflow Doubt Discussion](#) Telegram Group :

Username: **@GateOverflow_CSE**







Weekly Quiz 30

Discussion

Digital Logic:

Boolean Algebra, Minimization, Number System



Weekly Quiz 30: Digital Logic

In this Quiz, ALL the questions
are from GATE EC PYQs.

GATE ECE 2022 | Question: 17

asked in **Others** Feb 15, 2022 • retagged Mar 21, 2022 by **Lakshman Bhaiya**

128 views



2

Select the Boolean function(s) equivalent to $x + yz$, where x , y , and z are Boolean variables, and $+$ denotes logical OR operation.



- A. $x + z + xy$
- B. $(x + y)(x + z)$
- C. $x + xy + yz$
- D. $x + xz + xy$

gateece-2022

multiple-selects

GATE ECE 2022 | Question: 17

asked in **Others** Feb 15, 2022 • retagged Mar 21, 2022 by Lakshman Bhaiya

128 views



2



Select the Boolean function(s) equivalent to $x + yz$, where x , y , and z are Boolean variables, and $+$ denotes logical OR operation.

- A. $x + z + xy$
- B. $(x + y)(x + z)$
- C. $x + xy + yz$
- D. $x + xz + xy$

Ans.

gateece-2022

multiple-selects



(A) $x + z + xy = \cancel{x + xy} + z$
 $= x + z \neq x + yz$

GO
CLASSES

$$\underbrace{A + AB = A}_{\text{Absorption Law}}$$

B

$$\begin{aligned} (x+y)(x+z) &= \cancel{x}x + \cancel{y}z + yx + yz \\ &= \cancel{x} + \cancel{x}z + \cancel{y}x + yz \\ &\quad \swarrow \quad \searrow \\ &= \cancel{x} + \cancel{yz} \checkmark \end{aligned}$$

$$\overline{x + yz} = (x+y)(x+z)$$

Distributive Law

$$x + xy = x$$

$$x + x\beta = x$$



C

$$x + \bar{y}y + yz$$

$$\underline{x + yz}$$

$$\boxed{\overline{A + AB} = A}$$

Absorption Law

GO
CLASSES

(D)

$$x + xz + xy$$



$$x + xy = x \neq x + yz$$

≡ GO Electronics

GATE ECE 2010 | Question: 11

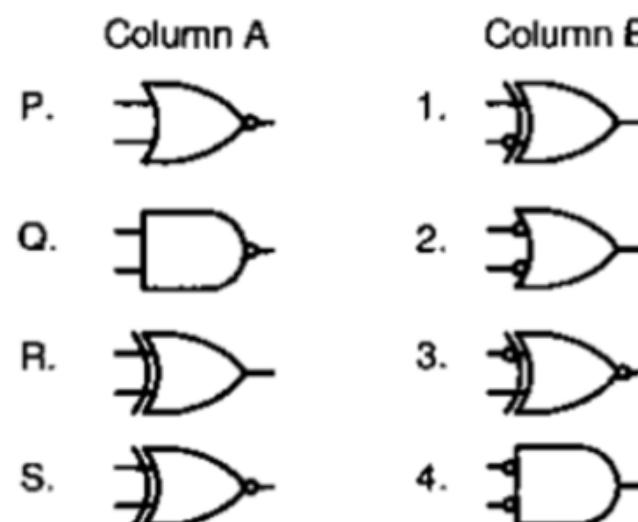
asked in Combinational Circuits Sep 15, 2022 • recategorized Jan 27 by Lakshman Bhaiya

24 views



Match the logic gates in **Column A** with their equivalents in **Column B**.

2



- A. P-2, Q-4, R-1, S-3
- B. P-4, Q-2, R-1, S-3
- C. P-2, Q-4, R-3, S-1
- D. P-4, Q-2, R-3, S-1

GATE ECE 2010 | Question: 11

P-4 Q-2
 R-3 S-1

asked in Combinational Circuits Sep 15, 2022 · recategorized Jan 27 by Lakshman Bhaiya

24 views



Match the logic gates in **Column A** with their equivalents in **Column B**.

2

$$\overline{ab} = \overline{a+b}$$

NAND

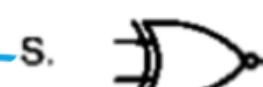
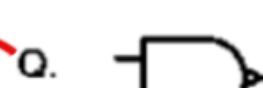
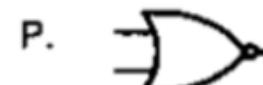
$$a \oplus b$$

Exor

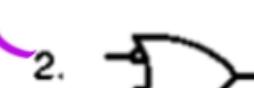
$$a \odot b$$

XNor

Column A



Column B



$$\overline{a} + \overline{b}$$

$$\overline{a} \odot b = a \oplus b$$

- A. P-2, Q-4, R-1, S-3
- B. P-4, Q-2, R-1, S-3
- C. P-2, Q-4, R-3, S-1
- D. P-4, Q-2, R-3, S-1

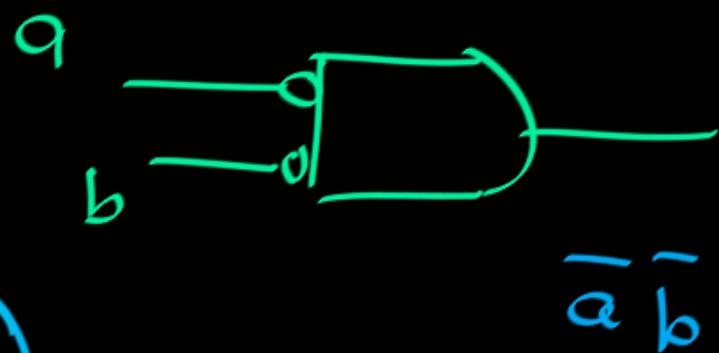


P : NOR gate



$$= \overline{a} \overline{b}$$

P - Y



Y

CLASSES



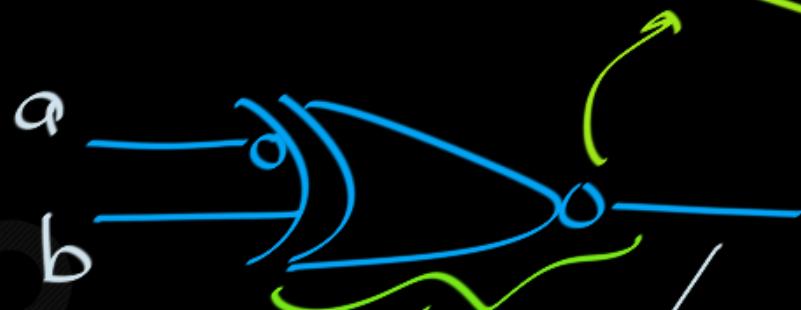
R: ExOR gate



$$\overline{a} \odot b = a \oplus b$$

prove it

③



XNOR
gate

$$\equiv \overline{a} \odot b$$

$$a \oplus b = \bar{a} \oplus \bar{b}$$

$$\bar{a} \oplus b = a \odot b$$

$$a \oplus \bar{b} = a \odot b$$

$$\overline{a \oplus b} = a \odot b$$

$$\bar{a} \odot b = a \oplus b$$

PROOF: by case method

Case 1:

$$a=0$$

$$\begin{aligned} LHS &= b \\ RHS &= b \end{aligned}$$

$$\begin{aligned} LHS &= \bar{b} \\ RHS &= \bar{b} \end{aligned}$$

Case 2:

$$a=1$$

$$\begin{aligned} LHS &= \bar{b} \\ RHS &= \bar{b} \end{aligned}$$

same



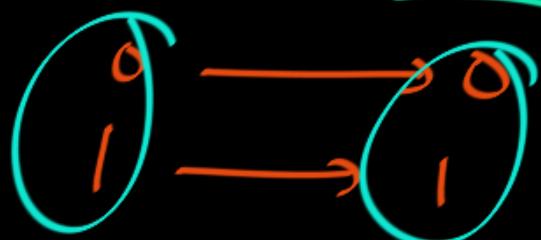
Digital Logic

$$1 \odot b = \boxed{b}$$

$$1 \odot b = b \checkmark$$



$$0 \oplus b = \boxed{b}$$

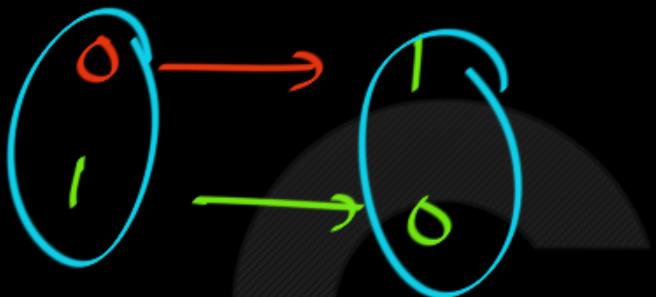


GO
CLASSES

$$0 \oplus b = b$$



$$0 \oplus b = \boxed{-} \rightarrow \bar{b}$$



GO
CLASSES

$$1 \oplus b = \boxed{-} \rightarrow \bar{b}$$

$$0 \rightarrow 1$$

$$1 \rightarrow 0$$

$$\alpha \odot b = Q + b$$
$$\bar{a}b + a\bar{b} = \bar{\alpha}b + a\bar{b}$$

$$x \oplus y = \bar{x}y + x\bar{y}$$

$$x \odot y = \bar{x}\bar{y} + xy$$

$$\alpha \odot b = \bar{\alpha} \oplus b$$



1

GO
CLASSES

$$\bar{\alpha} \oplus b$$

[←](#) [→](#) [C](#)ec.gateoverflow.in/459/gate-ece-2014-set-2-question-15[≡](#) GO Electronics

GATE ECE 2014 Set 2 | Question: 15

asked in Digital Circuits Mar 26, 2018 • recategorized Feb 26, 2021 by Lakshman Bhaiya

129 views



1



The number of bytes required to represent the decimal number 1856357 in packed BCD (Binary Coded Decimal) form is _____.

[gate2014-ec-2](#)[digital-circuits](#)[number-system](#)[number-representation](#)

[←](#) [→](#) [C](#)[ec.gateoverflow.in/459/gate-ece-2014-set-2-question-15](#)[≡ GO Electronics](#)

GATE ECE 2014 Set 2 | Question: 15

asked in Digital Circuits Mar 26, 2018 • recategorized Feb 26, 2021 by Lakshman Bhaiya

129 views



1



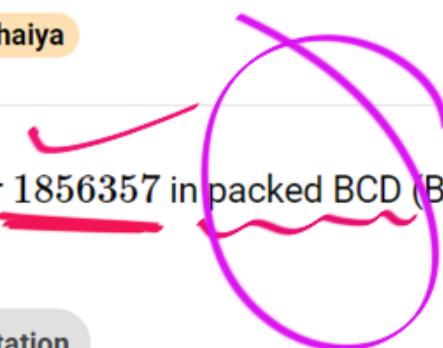
The number of bytes required to represent the decimal number 1856357 in packed BCD (Binary Coded Decimal) form is _____.

gate2014-ec-2

digital-circuits

number-system

number-representation



18 5635]

Decimal
Number

7 Digits

with 3 Bytes

= only 24 bits

can be represented

Packed BCD:

$$7 \times 4 = 28 \text{ Bits}$$

$$\left\lceil \frac{28}{8} \right\rceil = 4 \text{ Bytes}$$

18 5625]

Packed BCD

0001 10000101 0110 0011 01010111

In Bytes

4 bytes

BCD : Binary Coded Decimal

Decimal

Decimal	BCD form
0	0000
1	0001
2	0010
3	0011
4	0100
5	0101
6	0110
7	0111
8	1000
9	1001

BCD : Binary Coded Decimal

BCD

Packed BCD

Use 4 bits to
encode a Decimal
Digit.

Unpacked BCD

Use 8 bits to
encode a Decimal
Digit.



By Default;

BCD = Packed BCD

we use 4-bits to
encode a Decimal Digit.



Digital Logic

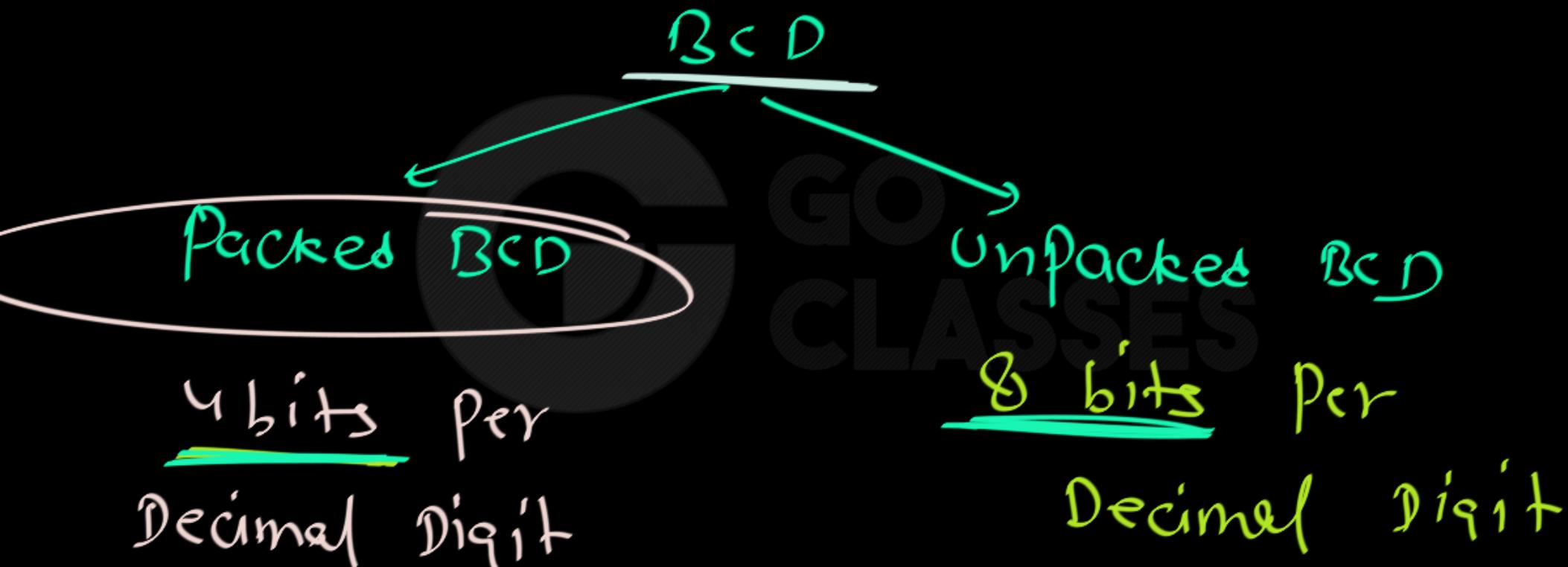
Decimal	(Packed) BCD	Unpacked BCD
0	0000	0000 0000
1	0001	0000 0001
2	0010	0000 0010
3	0011	0000 0011
4	0100	0000 0100
:		
8	1000	0000 1000
9	1001	0000 1001



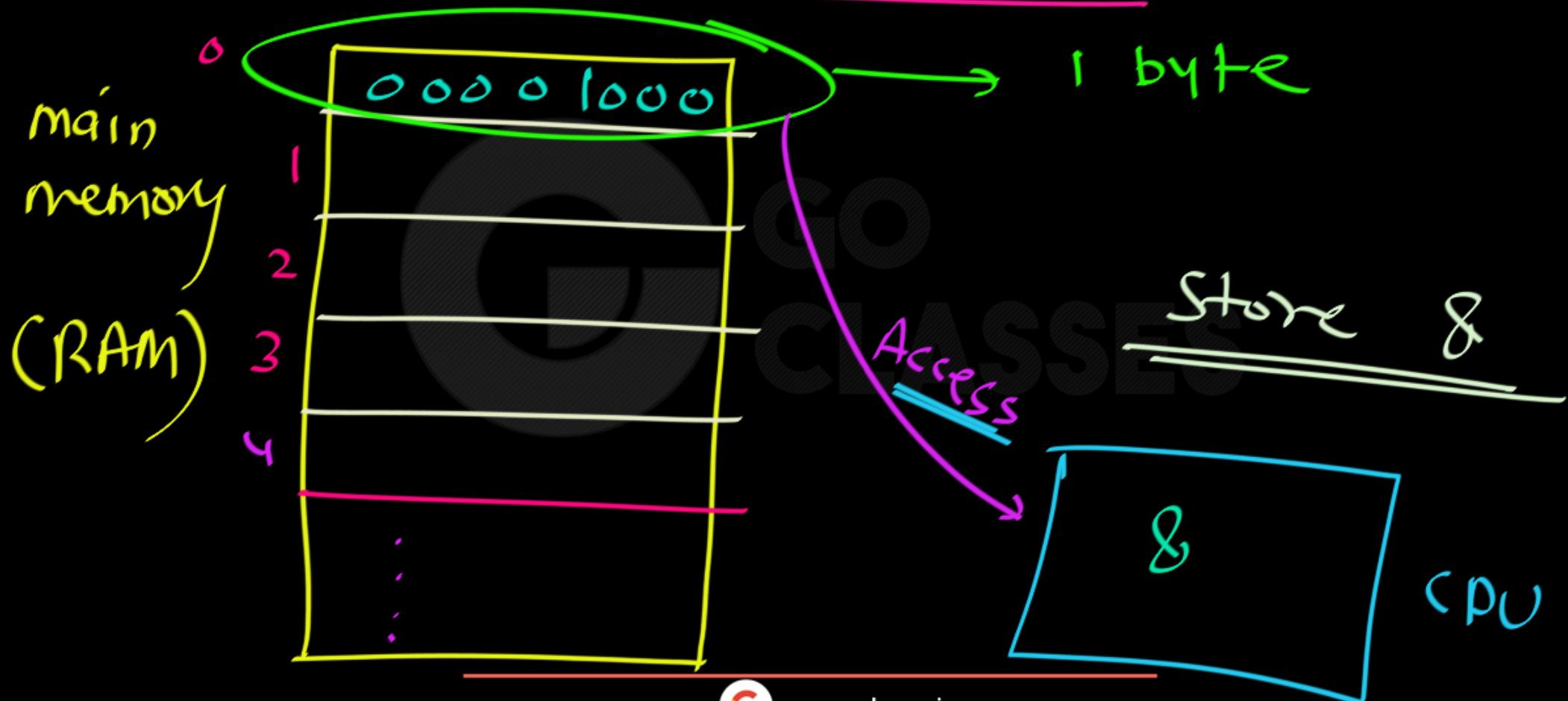
Digital Logic

Decimal	(Packed) BCD	Unpacked BCD
0	0000	0000 0000
1	0001	0000 0001
2	0010	0000 0010
3	0011	0000 0011
4	0100	0000 0100
:		
8	1000	0000 1000
9	1001	0000 1001

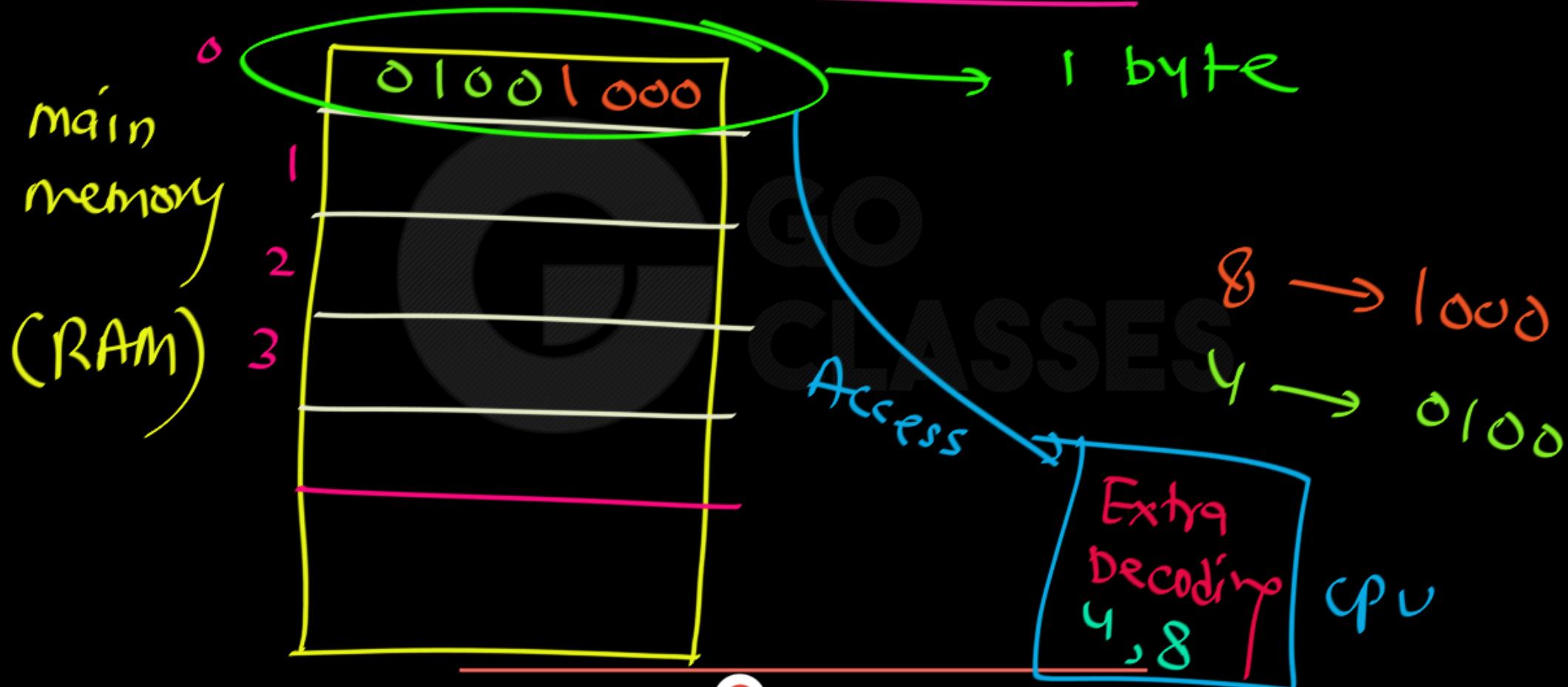
BCD : Binary Coded Decimal



Benefit of Unpacked BCD:



In Packed BCD:



GATE ECE 2014 Set 1 | Question: 15

asked in **Digital Circuits** Mar 26, 2018 · edited Nov 17, 2020 by **soujanyareddy13**

161 views



The Boolean expression $(X + Y)(X + \bar{Y}) + \overline{(X \bar{Y})} + \overline{\bar{X}}$ simplifies to

1

- A. X
- B. Y
- C. XY
- D. $X + Y$



gate2014-ec-1

digital-circuits

boolean-algebra

GATE ECE 2014 Set 1 | Question: 15

asked in [Digital Circuits](#) Mar 26, 2018 · edited Nov 17, 2020 by [soujanyareddy13](#)

161 views



The Boolean expression $(X + Y)(X + \bar{Y}) + \overline{(X \bar{Y})} + \overline{\bar{X}}$ simplifies to

1

- A. X ✓
- B. Y
- C. XY
- D. $X + Y$



[gate2014-ec-1](#)

[digital-circuits](#)

[boolean-algebra](#)



$$\underline{(x+y)(x+\bar{y}) + \overline{x\bar{y} + \bar{x}}}$$

$$\overline{\bar{x} + x\bar{y}} = \overline{\bar{x} + \bar{y}} = \textcircled{xy} \checkmark$$

$$\overline{\bar{x} + x\bar{y}} = \bar{x} + \bar{y}$$

$$A + \bar{A}B = A + \cancel{\bar{A}}B = A+B$$

$$A + \bar{A}B = \underbrace{(A + \bar{A})}_{1} (A + B) = \underline{\underline{A+B}}$$

$$\alpha + \bar{\alpha}\beta = \alpha + \beta$$

$$\bar{\alpha} + \cancel{\alpha}\beta = \bar{\alpha} + \beta$$

$$\overline{(x+y)(x+\bar{y})} + \overline{\bar{x}\bar{y} + \bar{x}} = xy$$

$$\begin{aligned}(x+y)(x+\bar{y}) &= x\cancel{x} + x\bar{y} + xy + \cancel{y\bar{y}} \\ &= x + x\bar{y} + xy \\ &= x + xy = x\end{aligned}$$

$$\overline{A} + \overline{AB} = \underline{A}$$

Absorption law

$$(x+y)(x+\bar{y})$$

+

$$\overline{x\bar{y} + \bar{x}}$$

 $x\bar{y}$ x

$$= x + x\bar{y}$$

$$= \underline{\underline{x}}$$

GATE ECE 2015 Set 1 | Question: 38

asked in [Digital Circuits](#) Mar 28, 2018 • recategorized Nov 15, 2020 by [soujanyareddy13](#)

81 views

1

A 3-input majority gate is defined by the logic function $M(a, b, c) = ab + bc + ca$. Which one of the following gates is represented by the function $M(\overline{M(a, b, c)}, M(a, b, \bar{c}), c)$?

- A. 3-input NAND gate
B. 3-input XOR gate
C. 3-input NOR gate
D. 3-input XNOR gate

gate2015-ec-1

digital-circuits

logic-gates



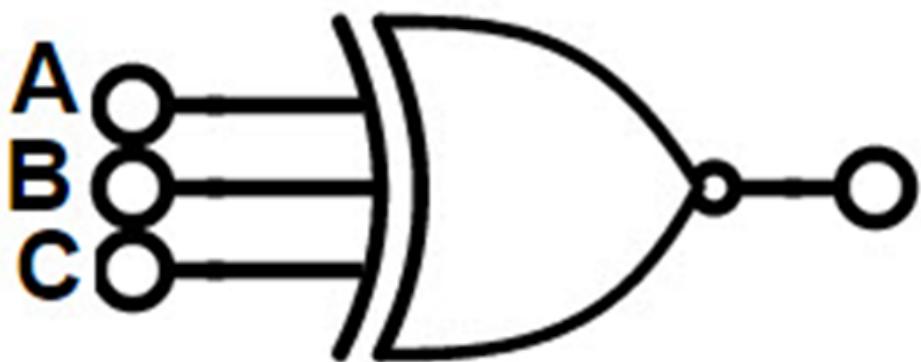
An Important NOTE:

Difference between

XNOR GATE & XNOR Function



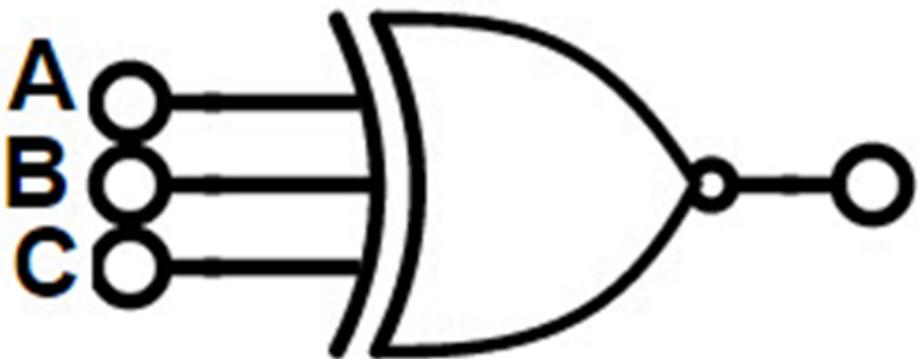
Are these SAME???



$$A \oplus B \oplus C$$



Are these SAME???



$$A \odot B \odot C$$



Quick Recap of XOR, XNOR functions:

XOR, XNOR functions over 2 input variables:

a	b	$a \oplus b$	$a \odot b$
0	0	0	1
0	1	1	0
1	0	1	0
1	1	0	1

$$a \oplus b = \overline{a \odot b}$$



$$\left\{ \begin{array}{l} a \oplus b = \overline{a \odot b} \\ a \odot b = \overline{a \oplus b} \end{array} \right.$$

for 2 input variables



XNOR function is Associative:

O is Associative.

$$(a \oplus b) \oplus c = a \oplus (b \oplus c)$$

Can Prove using Truth Table.

x	y	z	Digital Logic	$(x \oplus y) \oplus z$	$x \oplus (y \oplus z)$
0	0	0		0	0
0	0	1		1	1
0	1	0		1	0
0	1	1		0	1
1	0	0		0	1
1	0	1		0	0
1	1	0		0	0
1	1	1		1	1

XNOR is Associative.

$$(a \odot b) \odot c = a \odot (b \odot c)$$

We can write it

$$\underbrace{a \odot b \odot c}$$



XOR, XNOR function

are Associative & Commutative:

$$a \oplus b = b \oplus a$$

$$a \odot b = b \odot a$$

Commutative
Property

$$(a \odot b) \odot c = a \odot (b \odot c)$$

$$(a \oplus b) \oplus c = a \oplus (b \oplus c)$$

Associative Prop.



XOR, XNOR functions over 3 input variables:



x	y	z
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1

x ⊕ y ⊕ z
0
1
1
1
0
0
1
0



XOR, XNOR functions over 3 input variables:

$$a \oplus b \oplus c = a \odot b \odot c$$

GO
CLASSES



XOR, XNOR :

2 variables

$$a \oplus b = \overline{a \odot b}$$

3 variables

$$a \oplus b \oplus c = a \odot b \odot c$$



XOR, XNOR :

Even variables

$$a \oplus b = \overline{a \odot b}$$

Odd variables

$$a \oplus b \oplus c = a \odot b \odot c$$

ExOR = 1

iff

Odd number of 1's.

ExNOR = 1

iff

Even number of 0's

XOR is Odd 1's Detector.

XNOR is Even 0's Detector.



Digital Logic

$$| \oplus \circ \oplus | \oplus \circ \oplus \circ \oplus \circ = \circ$$

$$| \circ \circ \circ | \circ \circ \circ \circ \circ \circ = 1$$



for odd number of input Variables :

$$\oplus(a_1, a_2, \dots, a_{\text{odd}}) = \odot(a_1, a_2, \dots, a_{\text{odd}})$$

for even number of input Variables :

$$\oplus(a_1, \dots, a_{\text{even}}) = \overline{\odot(a_1, a_2, \dots, a_{\text{even}})}$$



Let A, B, C be three boolean variables. \oplus and \odot are exclusive-or(ExOr) and exclusive-nor(ExNor) operations respectively.

1

Consider the following statements:



1. $(A \oplus B) \oplus C = A \oplus (B \oplus C)$
2. $(A \odot B) \odot C = A \odot (B \odot C)$
3. $A \oplus B \oplus C = A \odot B \odot C$
4. $A \oplus B \oplus C = \overline{(A \odot B \odot C)}$



Which of the above statements is/are correct?

- A. 1 and 3 only
- B. 2 and 4 only
- C. 1, 2 and 3 only
- D. 1, 2 and 4 only

1

Let A, B, C be three boolean variables. \oplus and \odot are exclusive-or(ExOr) and exclusive-nor(ExNor) operations respectively.

Consider the following statements:

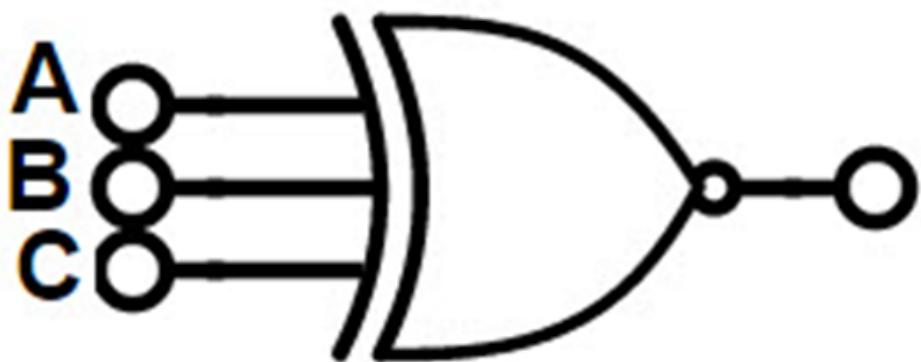
1. $(A \oplus B) \oplus C = A \oplus (B \oplus C)$
 2. $(A \odot B) \odot C = A \odot (B \odot C)$
 3. $A \oplus B \oplus C = A \odot B \odot C$
 4. $A \oplus B \oplus C = \overline{(A \odot B \odot C)}$

Which of the above statements is/are correct?

- A. 1 and 3 only
 - B. 2 and 4 only
 - C. 1, 2 and 3 only
 - D. 1, 2 and 4 only

GO CLASSES

Are these SAME???



$$A \oplus B \oplus C$$



Are these SAME???



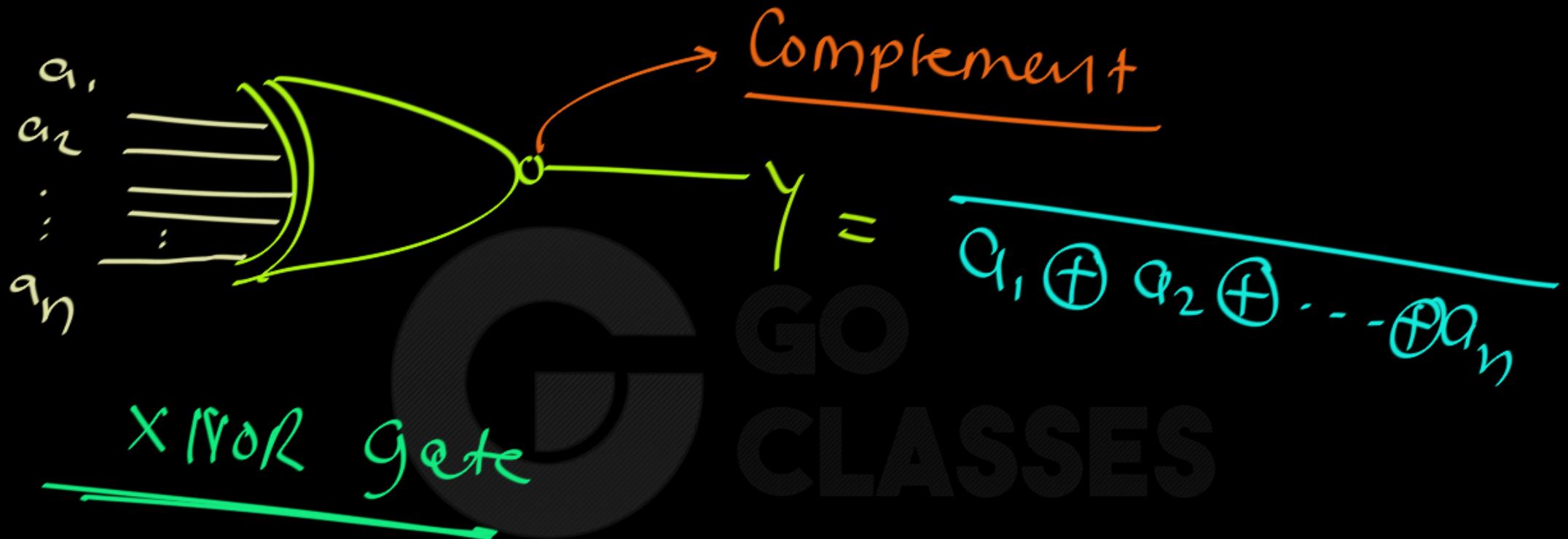
✗ $A \odot B \odot C$

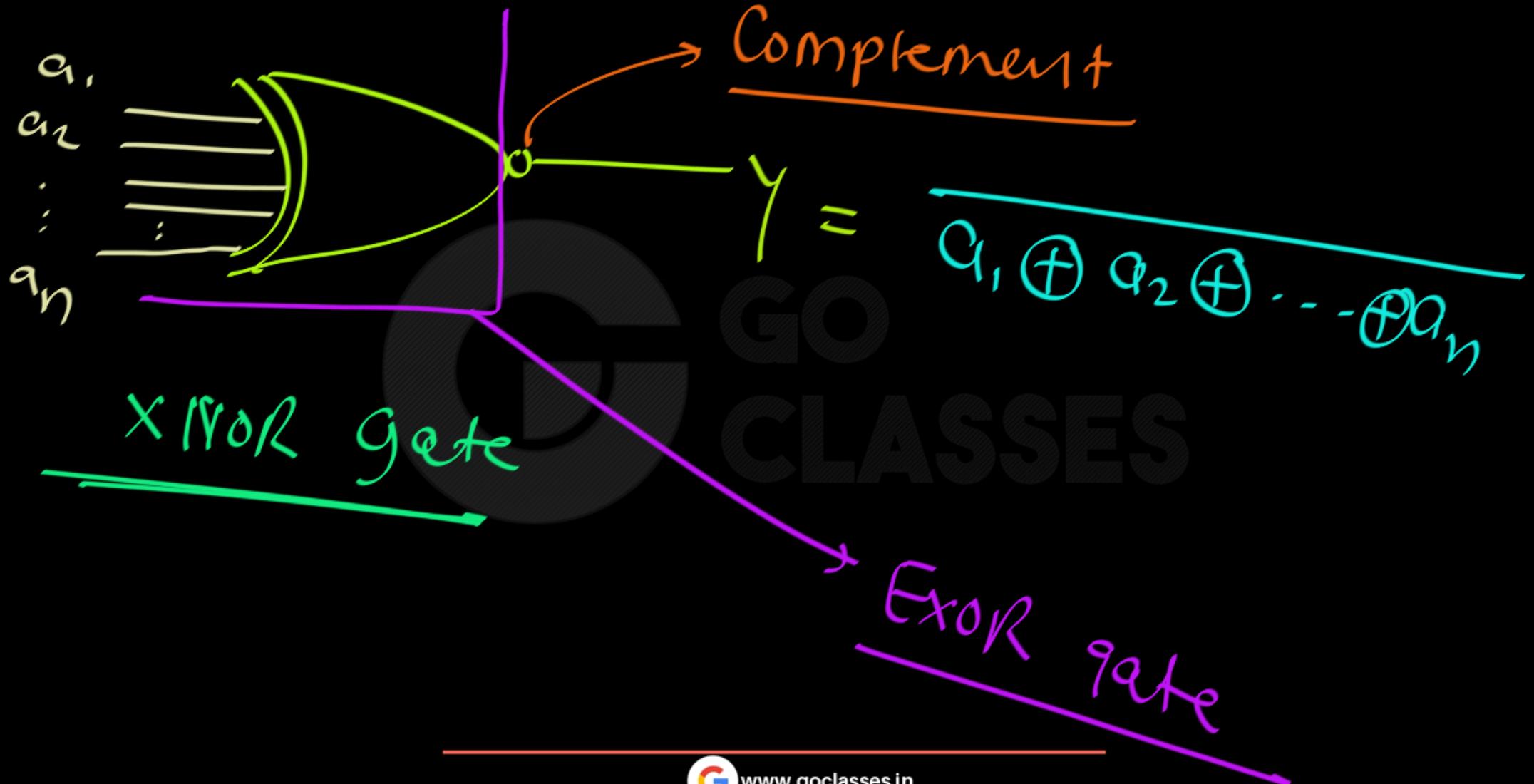


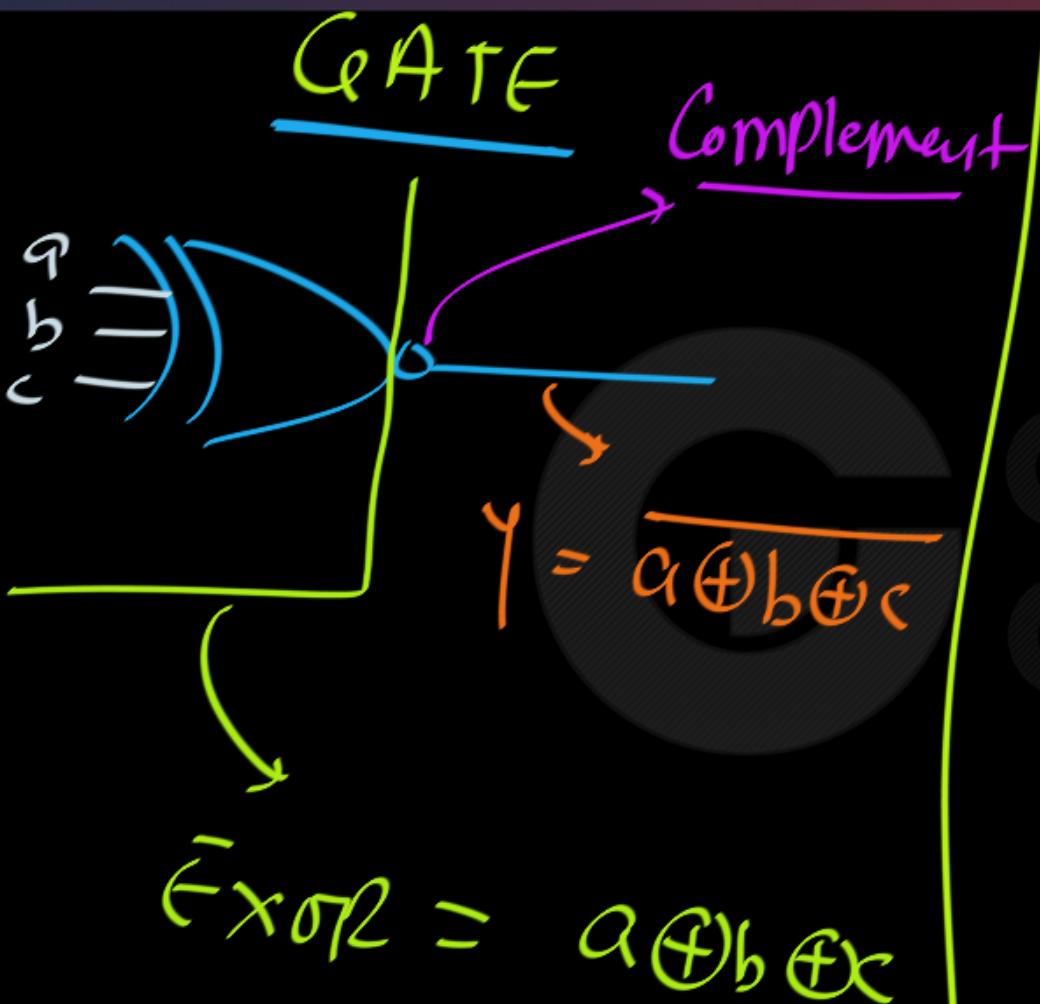
NOTE:

XNOR Gate is Defined as NOT of XOR Gate.

Always



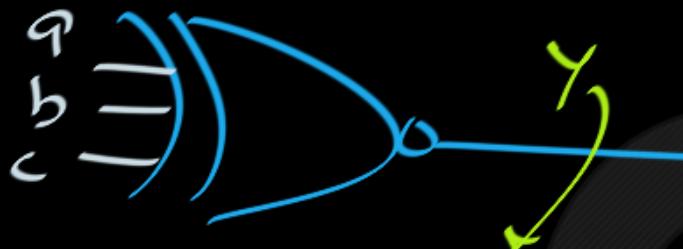




function

$$a \oplus b \oplus c =$$

$$a \oplus b \oplus c$$

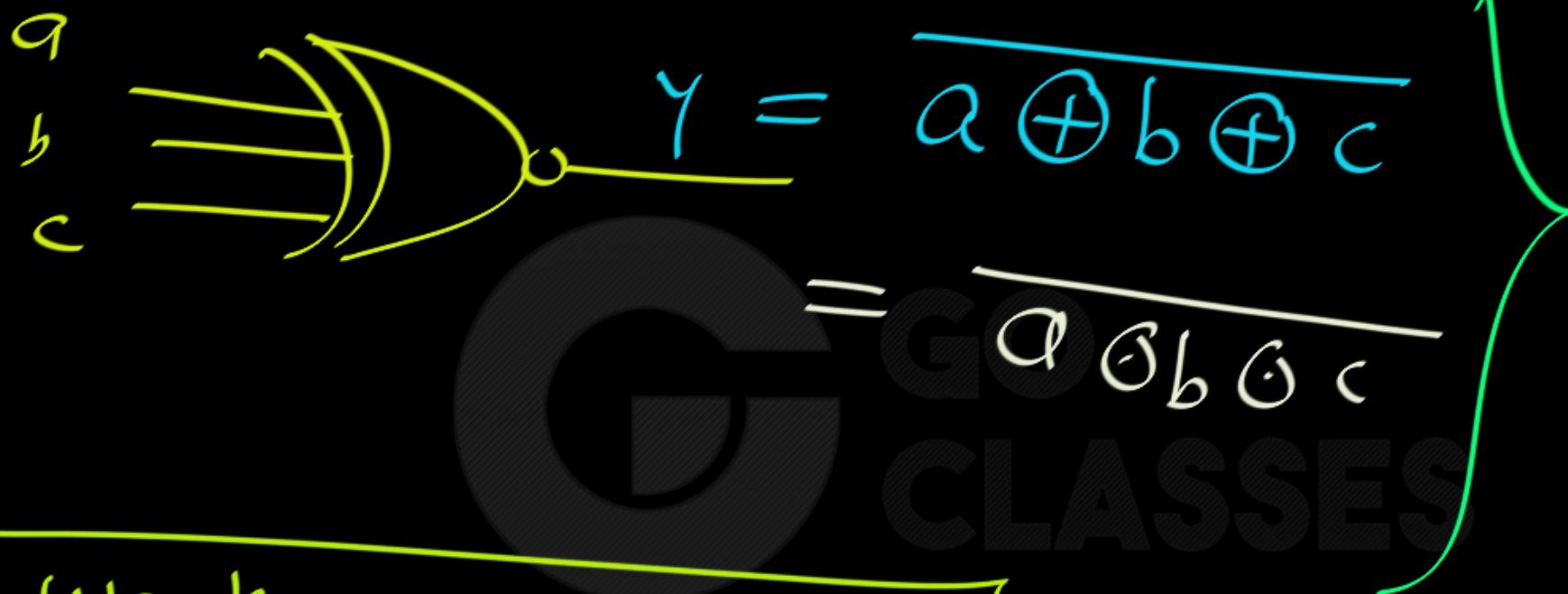
GATE

$$\overline{a \oplus b \oplus c}$$

function

$$a \oplus b \oplus c =$$

$$a \odot b \odot c$$



We know :

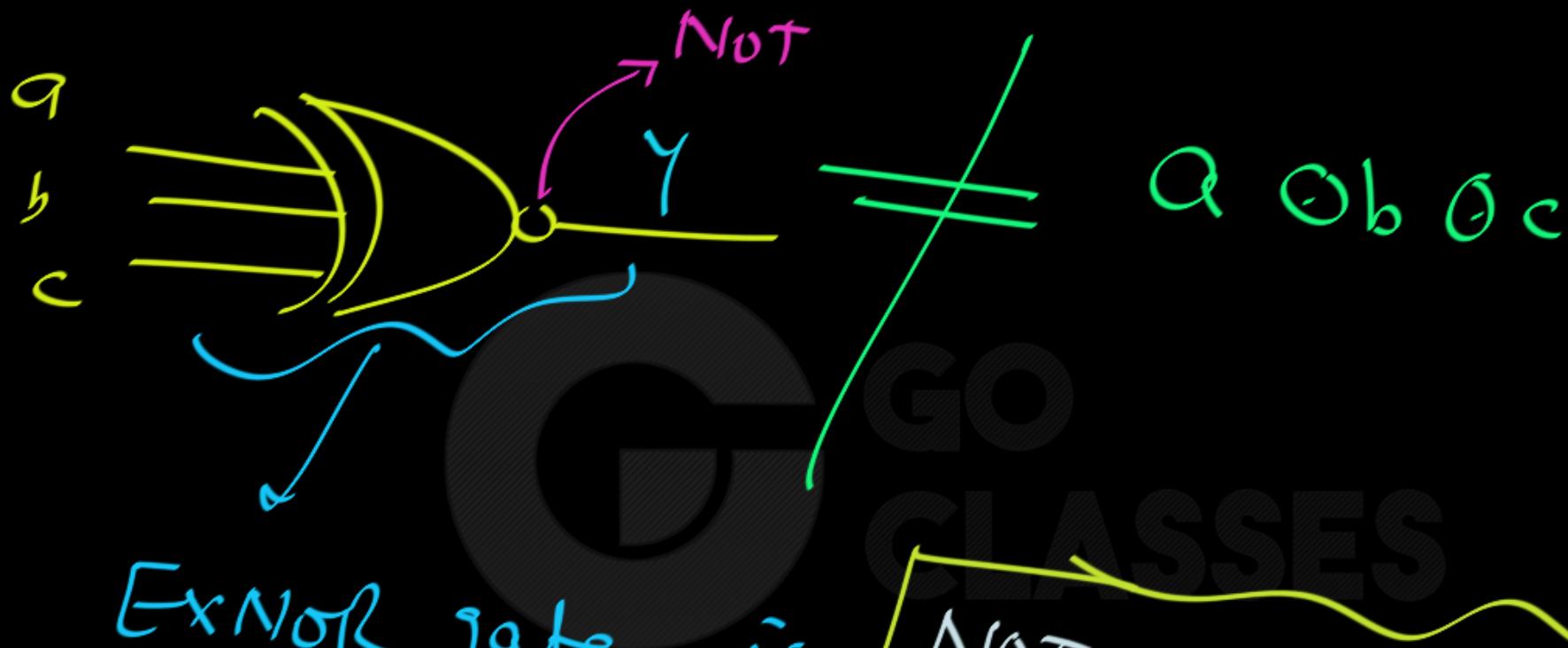
$$a \oplus b \oplus c = a \odot b \odot c$$

gatefunction

$$\begin{array}{c} a \\ b \end{array} \Rightarrow \text{Do} \quad \overline{a \oplus b} \\ = a \odot b \quad |$$

$$a \oplus b \oplus c = a \odot b \odot c$$

$$\begin{array}{c} a \\ b \\ c \end{array} \Rightarrow \text{Do} \quad \overline{a \oplus b \oplus c} \\ \neq a \odot b \odot c \quad |$$



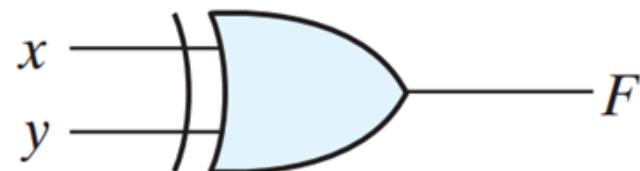
NOT of XOR gate

Note:

XNOR Gate for any
number of input variables is
Always Defined as

Not of ExOR gate

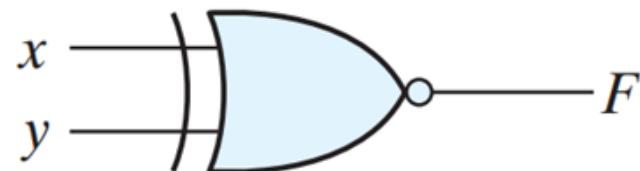
Exclusive-OR
(XOR)



$$\begin{aligned}F &= xy' + x'y \\&= x \oplus y\end{aligned}$$

x	y	F
0	0	0
0	1	1
1	0	1
1	1	0

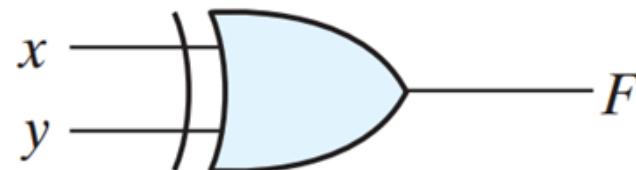
Exclusive-NOR
or
equivalence



$$\begin{aligned}F &= xy + x'y' \\&= (x \oplus y)'\end{aligned}$$

x	y	F
0	0	1
0	1	0
1	0	0
1	1	1

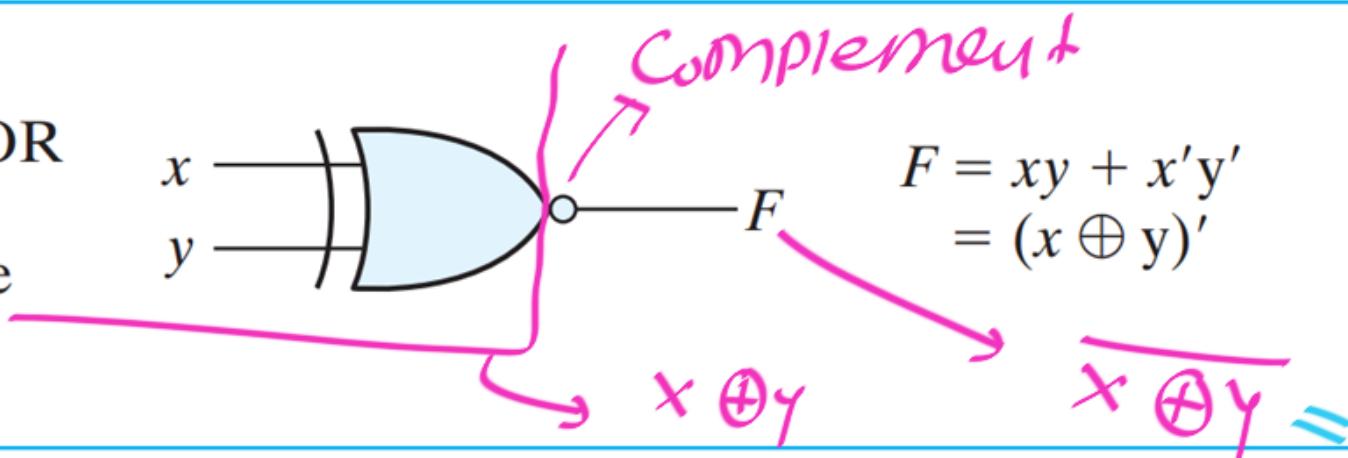
Exclusive-OR
(XOR)



$$\begin{aligned}F &= xy' + x'y \\&= x \oplus y\end{aligned}$$

x	y	F
0	0	0
0	1	1
1	0	1
1	1	0

Exclusive-NOR
or
equivalence



$$\begin{aligned}F &= xy + x'y' \\&= (x \oplus y)'\end{aligned}$$

x	y	F
0	0	1
0	1	0
1	0	0
1	1	1



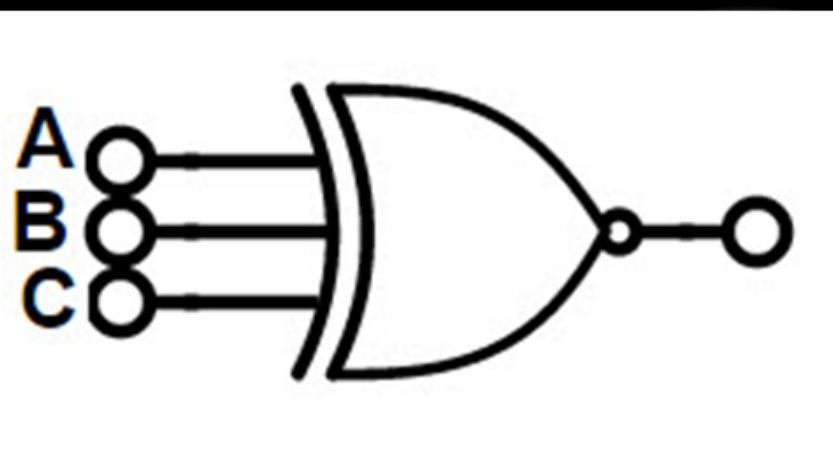
NOTE:

XNOR Gate is Defined as NOT of XOR Gate.

Always



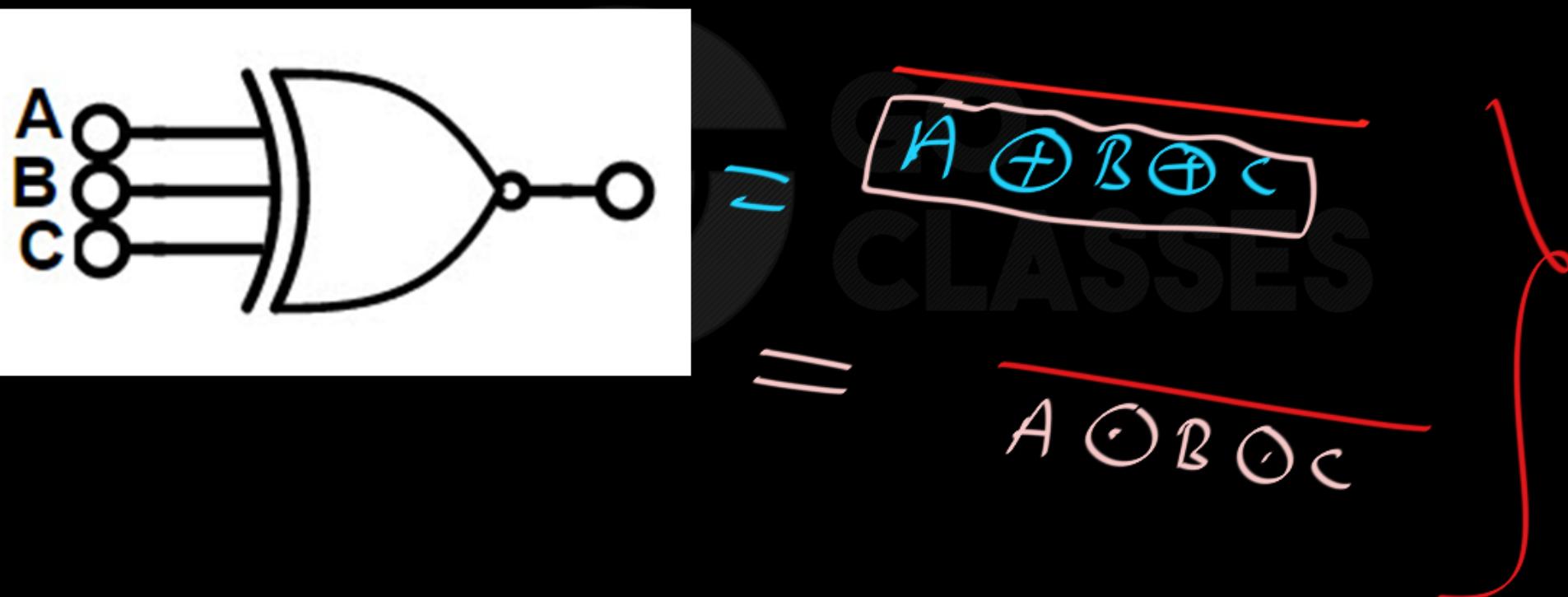
XNOR gate is defined as NOT of XOR gate:



$$\begin{aligned} &= \overline{A \oplus B \oplus C} \\ &\neq A \odot B \odot C \end{aligned}$$



XNOR gate is defined as NOT of XOR gate:





62 Chapter 2 Boolean Algebra and Logic Gates

The exclusive-OR gate has a graphic symbol similar to that of the OR gate, except for the additional curved line on the input side. The equivalence, or exclusive-NOR, gate is the complement of the exclusive-OR, as indicated by the small circle on the output side of the graphic symbol.

Source: Morris Mano

< > Celectronics.stackexchange.com/questions/285392/3-input-xnor-gate-operation

StackExchange



Search on Electrical Engineering...

Home

PUBLIC

Questions

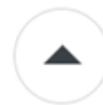
Tags

Users

Unanswered

TEAMS

4 Answers

Sorted by: Highest score (default) 

The misunderstanding is that, given XOR as a logic gate, XNOR is *defined* as being always its negation.

12



Having defined your XOR-3 as an *odd parity checker* (by accepting the minterm xyz - otherwise it would be a *one hot checker*), the correct interpretation of a XNOR-3 would then be an *even parity checker* (as pointed by Bradman175). This simply means that the expression for your algebraic XNOR-3 is not correct in this context.



In other words, $x \odot y \odot z \neq \text{XNOR-3}$.



The XNOR Gate is logically equivalent to an inverted XOR i.e., XOR gate followed by a NOT gate (inventor). Thus XNOR produces 1 (high) output when the input combination has even number of 1's.

Table 4.20 3-input XNOR gate



Conclusion:

A Very Important NOTE is:

ExNOR – 3 gate with inputs a, b, c is NOT same as $a \odot b \odot c$.

ExNOR – n gate is defined as Complement of ExOR – n gate.

NOTE that:

$$a \odot b \odot c = a \oplus b \oplus c$$

BUT $ExNOR - 3 \neq ExOR - 3$

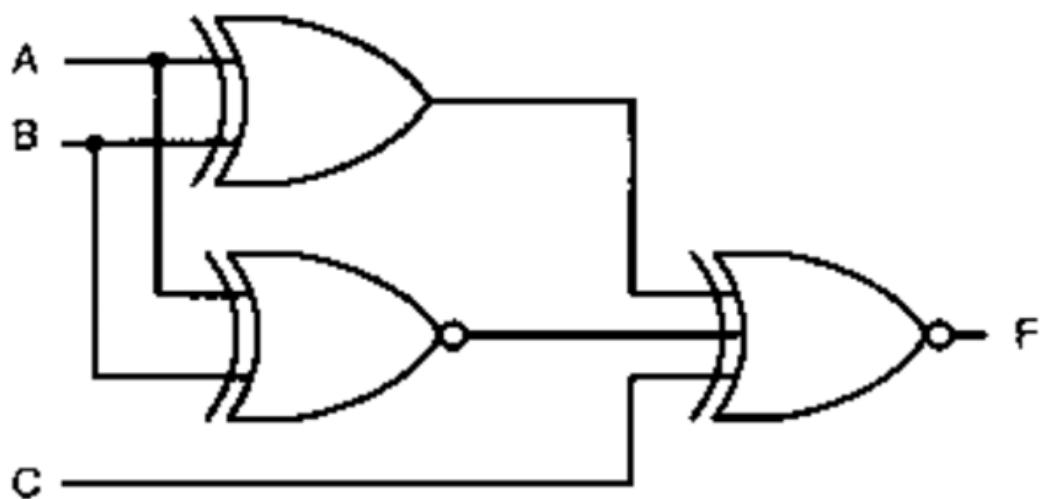
The exclusive-NOR gate is the complement of the exclusive-OR gate, as indicated by the small circle on the output side of the graphic symbol.

GATE EC 2010 Question:

EC2010.pdf

4 / 24 | - 157% + | ↻ ↺

Q.12 For the output F to be 1 in the logic circuit shown, the input combination should be



- (A) $A = 1, B = 1, C = 0$
(C) $A = 0, B = 1, C = 0$

- (B) $A = 1, B = 0, C = 0$
(D) $A = 0, B = 0, C = 1$



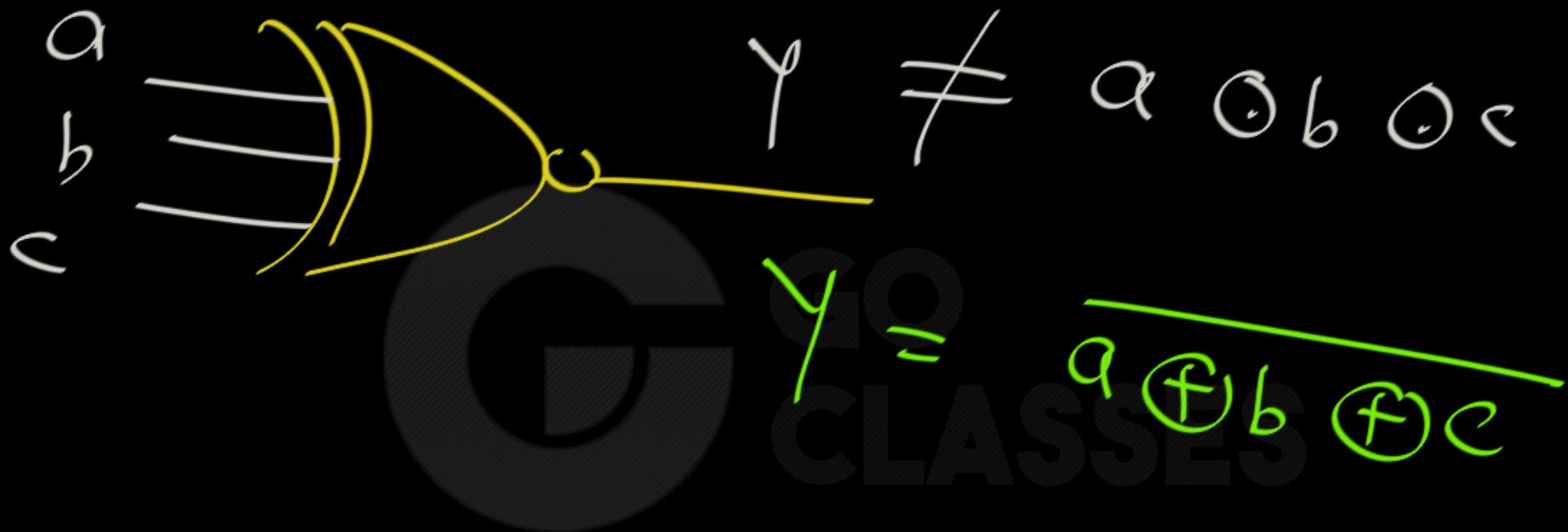
Digital Logic



$$F = x \odot \bar{x} \oplus c$$

$$\overline{x \oplus \bar{x} \oplus c}$$





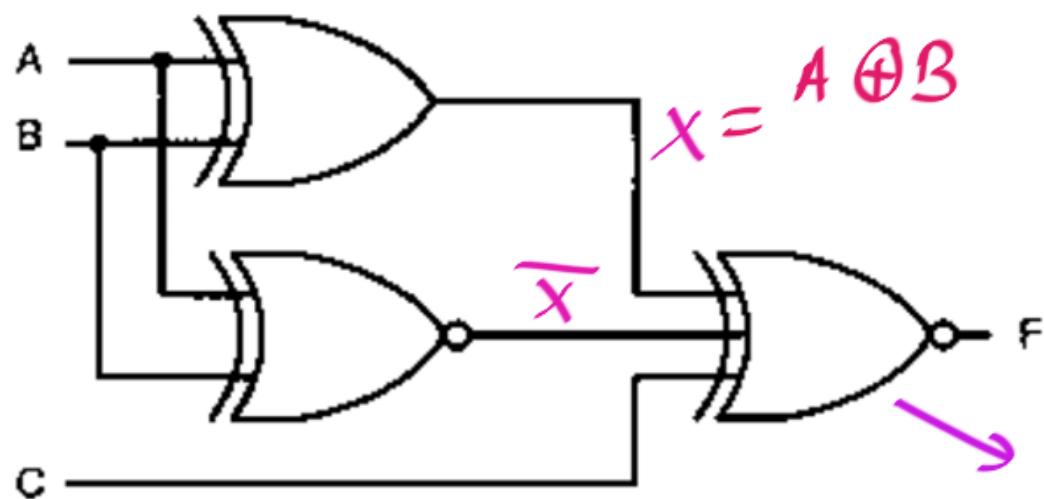
GATE EC 2010 Question:

EC2010.pdf

4 / 24 | - 157% + | ☰ ⌂

Q.12 For the output F to be 1 in the logic circuit shown, the input combination should be

$$F = C$$



- (A) A = 1, B = 1, C = 0
(C) A = 0, B = 1, C = 0

- (B) A = 1, B = 0, C = 0
(D) A = 0, B = 0, C = 1

XNOR-GATE

$$F = \overline{x \oplus \bar{x} \oplus c}$$

$$= \overline{1 \oplus c} = \overline{\bar{c}} = c$$

$$F = c$$

$$\Rightarrow F = 1 \text{ iff } c = 1$$

$$x \oplus \bar{x} = 1$$

$$a \oplus b = 1 \quad \text{iff} \quad a \neq b,$$

$$I \oplus C = ? \rightarrow \bar{C}$$

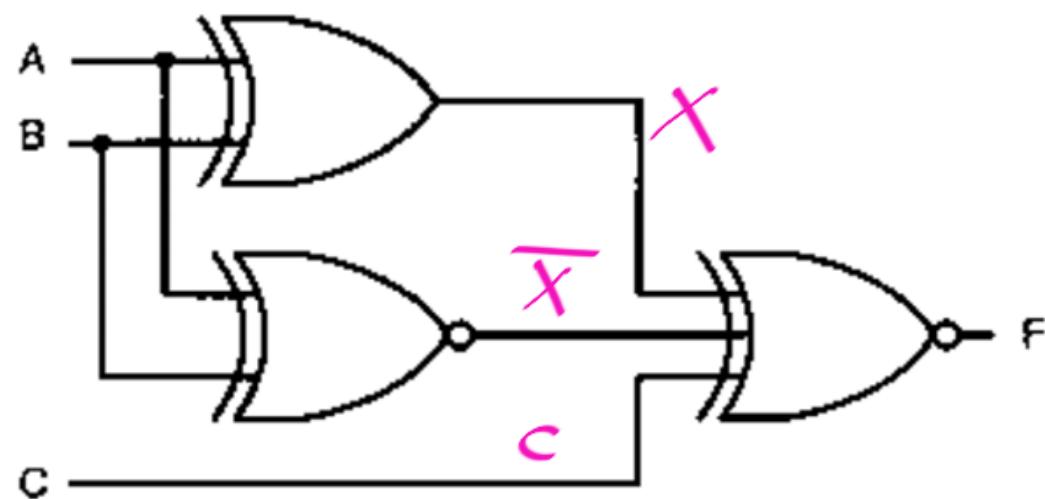
A diagram illustrating the union operation. At the bottom left, there is a set represented by an oval containing the letter 'I'. An arrow points from this set to another set at the bottom right, also represented by an oval containing the letter 'I'. Above these two sets, there is a third set represented by a blue-bordered box containing a question mark. An arrow points from the bottom right set to this box. To the right of the box, the symbol ' \rightarrow ' is followed by the expression ' \bar{C} '.

GATE EC 2010 Question:

EC2010.pdf

4 / 24 | - 157% + | ☰ ⌂

Q.12 For the output F to be 1 in the logic circuit shown, the input combination should be



- (A) $A = 1, B = 1, C = 0$
(C) $A = 0, B = 1, C = 0$

- (B) $A = 1, B = 0, C = 0$
(D) $A = 0, B = 0, C = 1$

$$F = x \odot \bar{x} \odot c$$



WRONG

$$= 0 \odot c$$

$$= \bar{c}$$

$$\overline{f = \bar{c}}$$

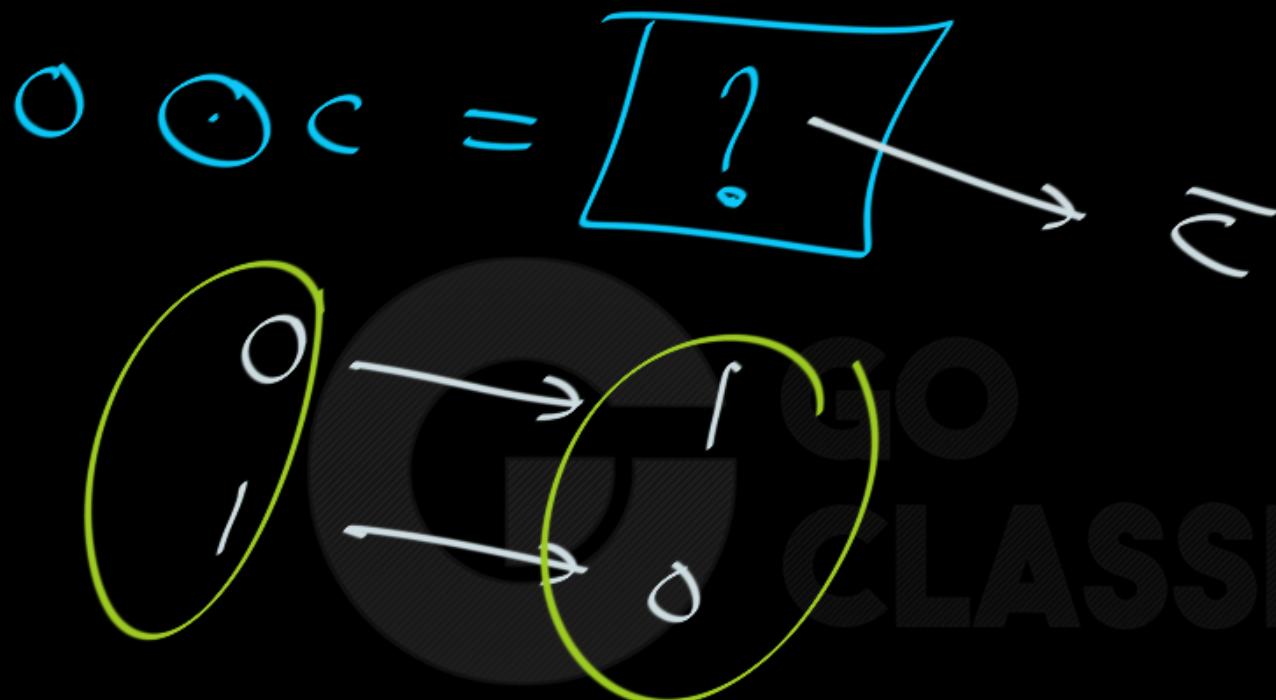
OPTIONS
A, B, C

$$F = 1 \text{ iff } \bar{c} = 0$$

$x \odot \bar{x} = 0$
 $a \odot b = 1 \text{ iff } a = b$



Digital Logic

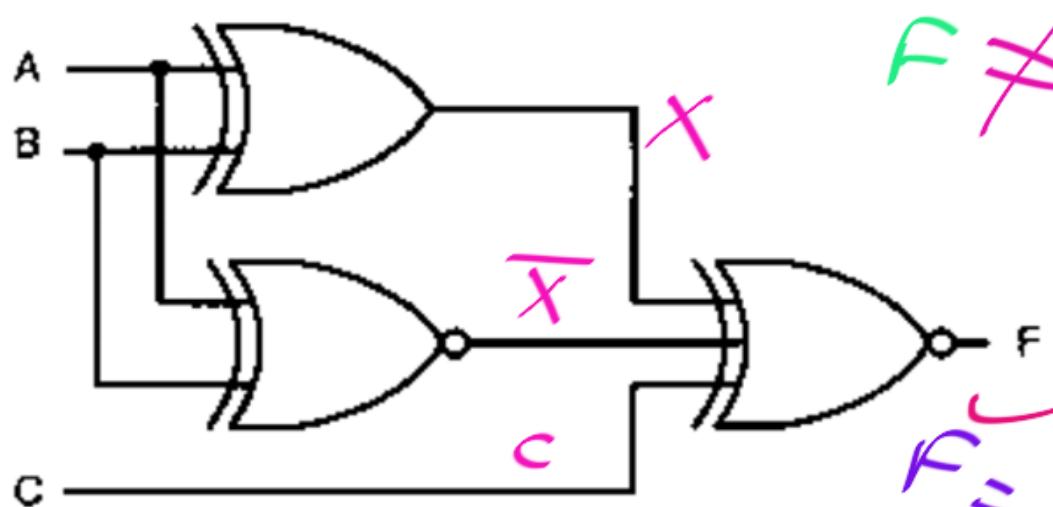


GATE EC 2010 Question:

EC2010.pdf

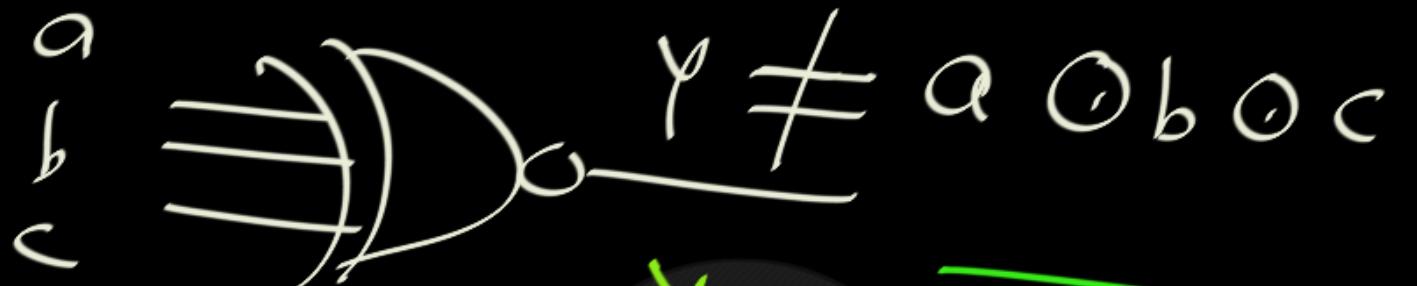
4 / 24 | - 157% + | ☰ ⌂

Q.12 For the output F to be 1 in the logic circuit shown, the input combination should be



- (A) A = 1, B = 1, C = 0
(C) A = 0, B = 1, C = 0

- (B) A = 1, B = 0, C = 0
(D) A = 0, B = 0, C = 1

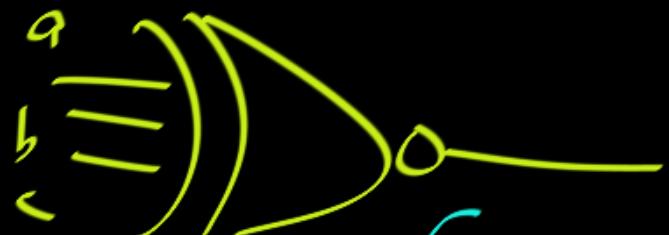


$$Y = \overline{a \oplus b \oplus c}$$

ExNOR gate is NOT of Exor gate



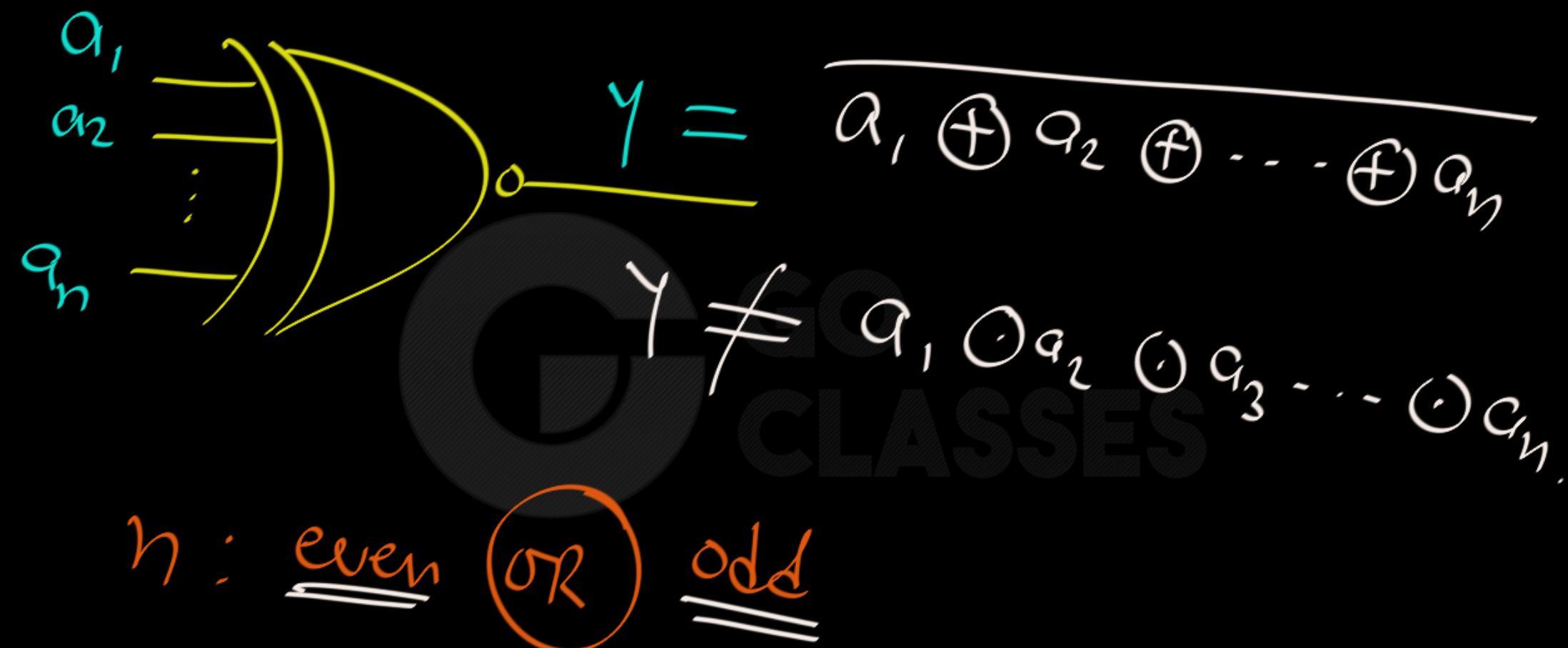
Digital Logic

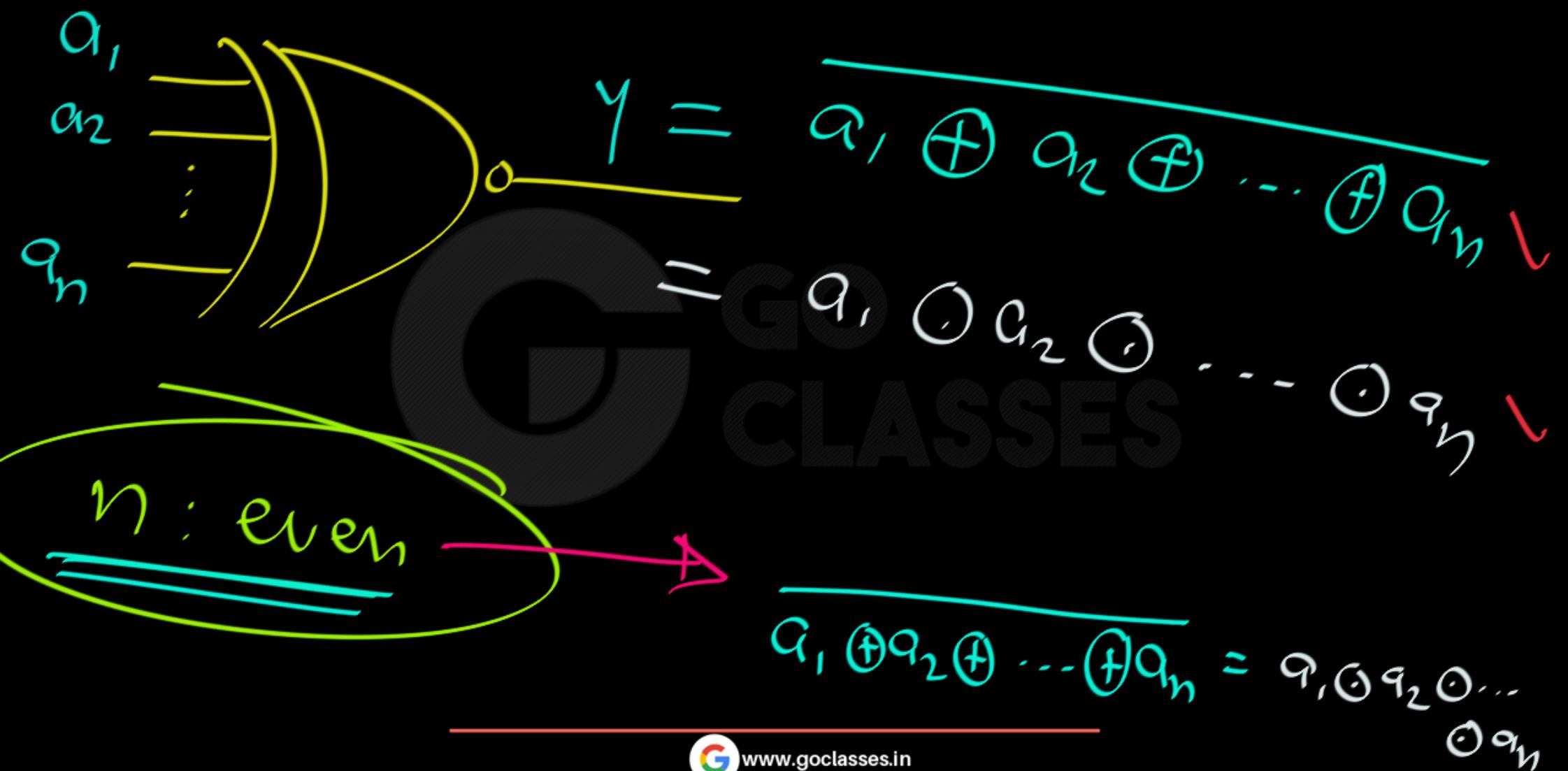


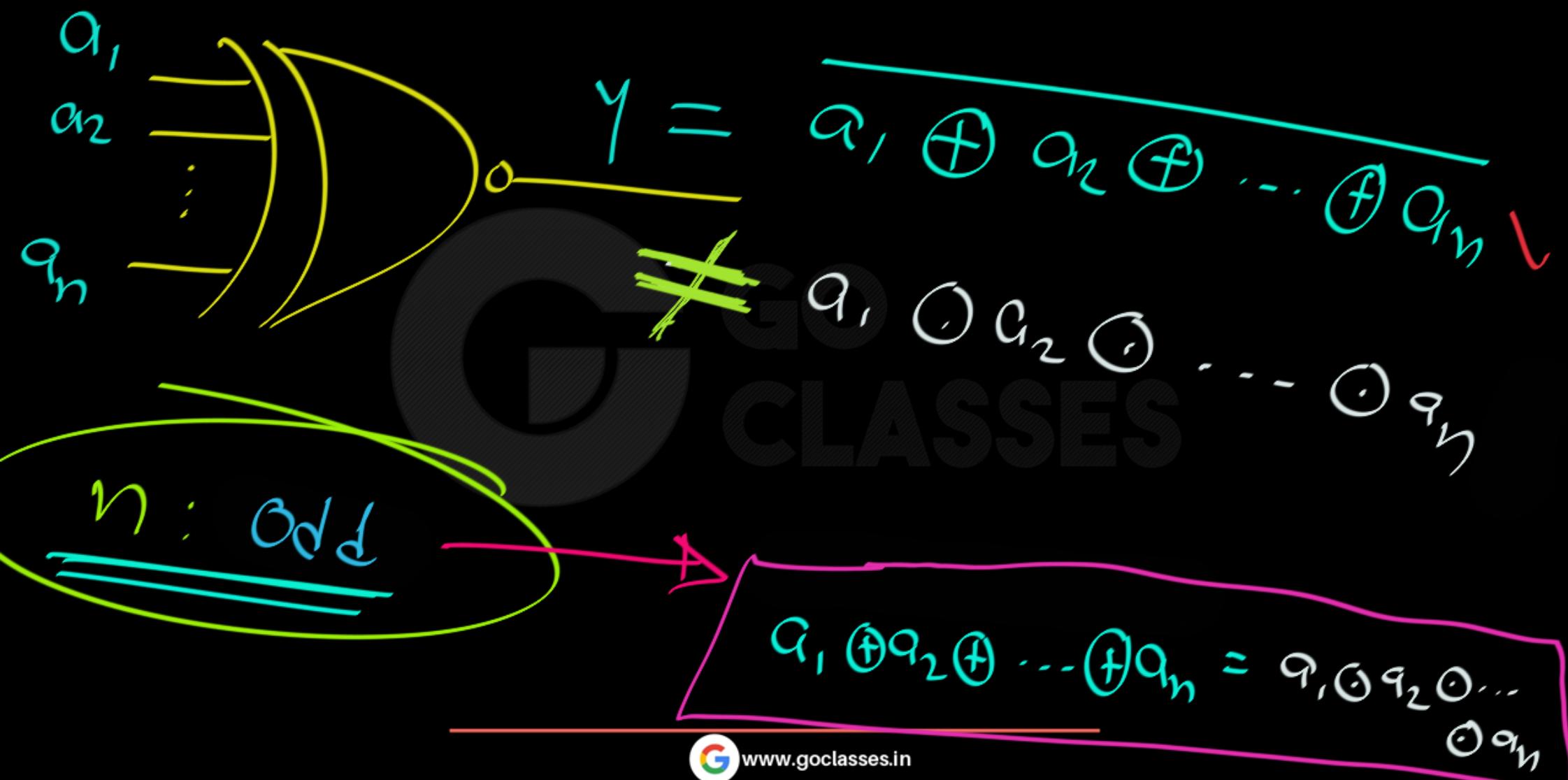
XNOR
gate

$$Y = \overline{a \oplus b \oplus c}$$

$$\begin{aligned} & a \oplus b \oplus c \\ = & a \odot b \odot c \end{aligned}$$







GATE ECE 2015 Set 1 | Question: 38

asked in [Digital Circuits](#) Mar 28, 2018 • recategorized Nov 15, 2020 by [soujanyareddy13](#)

81 views

1

A 3-input majority gate is defined by the logic function $M(a, b, c) = ab + bc + ca$. Which one of the following gates is represented by the function $M(\overline{M(a, b, c)}, M(a, b, \bar{c}), c)$?

- A. 3-input NAND gate
B. 3-input XOR gate
C. 3-input NOR gate
D. 3-input XNOR gate

gate2015-ec-1

digital-circuits

logic-gates

GATE ECE 2015 Set 1 | Question: 38

asked in Digital Circuits Mar 28, 2018 • recategorized Nov 15, 2020 by soujanyareddy13

81 views

1 A 3-input majority gate is defined by the logic function $M(a, b, c) = ab + bc + ca$. Which one of the following gates is represented by the function $M(\overline{M(a, b, c)}, M(a, b, \bar{c}), c)$?

- A. 3-input NAND gate
- B. 3-input XOR gate
- C. 3-input NOR gate
- D. 3-input XNOR gate

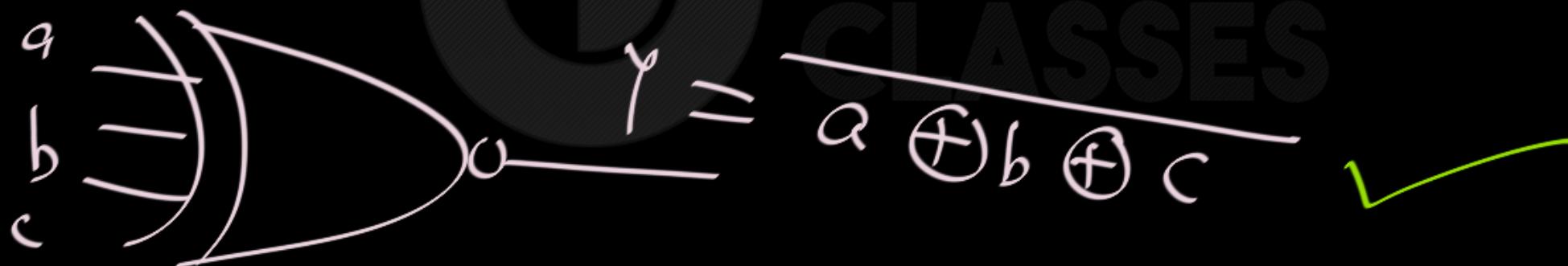
Complement of each other.

gate2015-ec-1

digital-circuits

logic-gates

XNOR gate is Always
Complement of ExOR gate.





3-input XNOR gate



3 input XOR gate

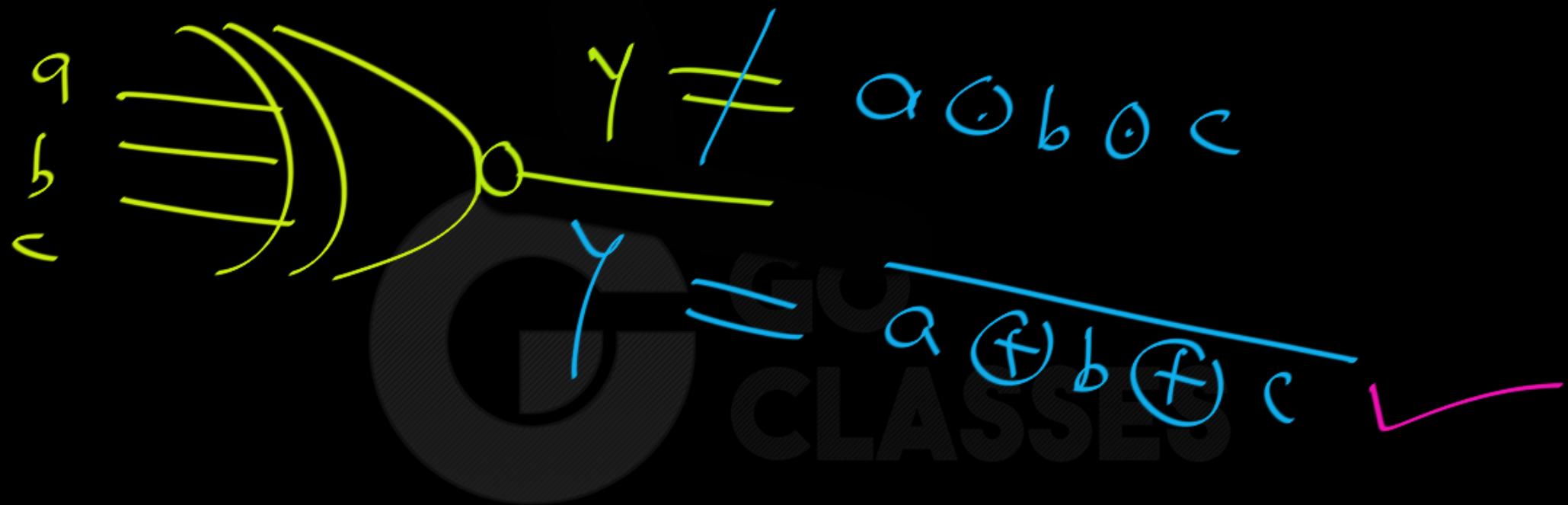


$$a \oplus b \oplus c$$

$$= abc \oplus ac \oplus bc$$



Digital Logic





$$M(a, b, c) = ab + bc + ca$$

$$M(\overline{m(a, b, c)}, M(a, b, \bar{c}), c) = ?$$

X Y

$$M(x, y, c) = xy + xc + yc$$



$$M(a, b, c) = ab + bc + \bar{c}a$$

$$X = \overline{M_{(a,b,c)}} = \overline{ab + bc + \bar{c}a}$$

$$= [(\bar{a} + \bar{b})(\bar{b} + \bar{c})] (\bar{c} + \bar{a})$$

$$= [\cancel{\bar{a} \bar{b}} + \bar{a} \bar{c} + \cancel{\bar{b} + \bar{b}\bar{c}}] (\bar{c} + \bar{a})$$

$$= (\bar{b} + \bar{a}\bar{c})(\bar{c} + \bar{a}) = \rightarrow$$

$$X = (\bar{b} + \bar{a}\bar{c})(\bar{c} + \bar{a})$$

$$= \bar{b}\bar{c} + \bar{b}\bar{a} + \bar{a}\bar{c}\bar{c} + \bar{a}\bar{c}\bar{a}$$

$$X = \boxed{\bar{a}\bar{b} + \bar{a}\bar{c} + \bar{b}\bar{c}}$$

$\bar{a}\bar{c}$

$$\alpha + \alpha\beta = \alpha$$

$$(\bar{b} + \bar{b}\bar{c}) + \bar{a}\bar{b}$$

$$\bar{b} + \bar{b}\bar{a} = \bar{b} \checkmark$$



$$M(a, b, c) = ab + bc + ca$$

$$Y = M(a, b, \bar{c}) = ab + b\bar{c} + a\bar{c}$$

$$X = \overline{M(a, b, c)} = \overline{ab} + \overline{bc} + \overline{ac}$$



$$M(a,b,c) = ab + bc + ca$$

$$M(x,y,c) = xy + xc + yc$$

GO
CLASSES

$$xy = (\bar{a}\bar{b} + \bar{a}\bar{c} + \bar{b}\bar{c}) (ab + b\bar{c} + a\bar{c})$$

$$= 0+0+0 + 0 + \bar{a}\bar{c}b + 0+0+0+$$

$$xy = \underbrace{\bar{a}\bar{c}b}_{a\bar{b}\bar{c}} + \underbrace{a\bar{b}\bar{c}}$$



$$\begin{aligned} X_C &= \left(\bar{a} \bar{b} + \bar{a} \bar{c} + \bar{b} \bar{c} \right) C \\ &= \bar{a} \bar{b} C + 0 + 0 = \bar{a} \bar{b} C \end{aligned}$$

$$\begin{aligned} Y_C &= (ab + a\bar{c} + b\bar{c})(\bar{C}) \\ &= ab\bar{C} \end{aligned}$$

$$M(a, b, c) = ab + bc + \bar{a} \bar{b} c$$

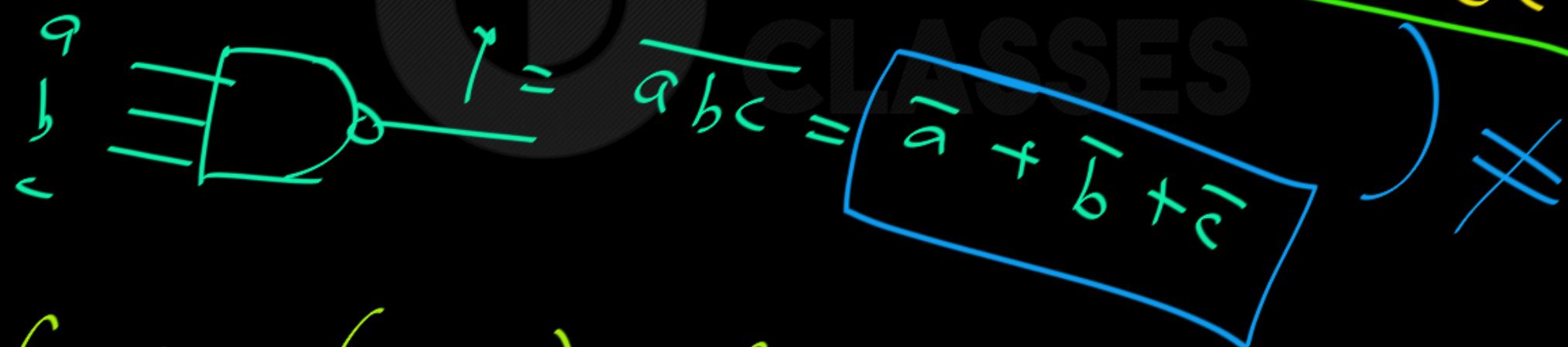
$$M(x, y, c) = xy + xc + yc$$

$\bar{a} \bar{c} b + a \bar{b} \bar{c}$

$$M(x, y, c) = \bar{a} \bar{b} c + abc + \bar{a} b \bar{c} + a \bar{b} \bar{c}$$

$$M(x,y,z) = \overbrace{abz + \bar{a}\bar{b}z + \bar{a}b\bar{z}}^{\text{Option A:}} + \bar{a}\bar{b}\bar{z}$$

Option A :


$$z = \bar{a}\bar{b}\bar{c} = \overbrace{\bar{a} + \bar{b} + \bar{c}}^{\text{abz}}$$

Consider $(a, b, c) = (0, 0, 0)$

$$M(x,y,z) = \overbrace{abz + \bar{a}\bar{b}z + \bar{a}b\bar{z} + ab\bar{z}}^{\text{option c?}}$$

option c?

$a+b+c \neq \bar{a}'\bar{b}'z$

Consider $(a, b, c) = (0, 0, 0); (1, 1, 1)$



$$M(x,y,z) = ab\bar{c} + \bar{a}\bar{b}\bar{c} + \bar{a}b\bar{c} + ab\bar{c}$$



ab̄c̄



$$\begin{aligned}M(x,y,z) &= \overbrace{abz + \bar{a}\bar{b}z + \bar{a}b\bar{z} +}^{\text{abz}} \\&\quad \overbrace{a\bar{b}\bar{z}}^{\text{abz}} \\&= \overbrace{a \oplus b \oplus z}^{\text{abz}}\end{aligned}$$



$$a \oplus b \oplus c = 1$$

iff

odd number
of variables
are 1.

$$a \oplus b \oplus c = abc + a\bar{b}\bar{c} + \bar{a}b\bar{c} + \bar{a}\bar{b}c$$



$$M(a, b, c) = ab + bc + ca$$

$$M\left(\overline{m_{(a,b,c)}}, M(a, b, \bar{c}), c\right) =$$

$$= a \oplus b \oplus c$$

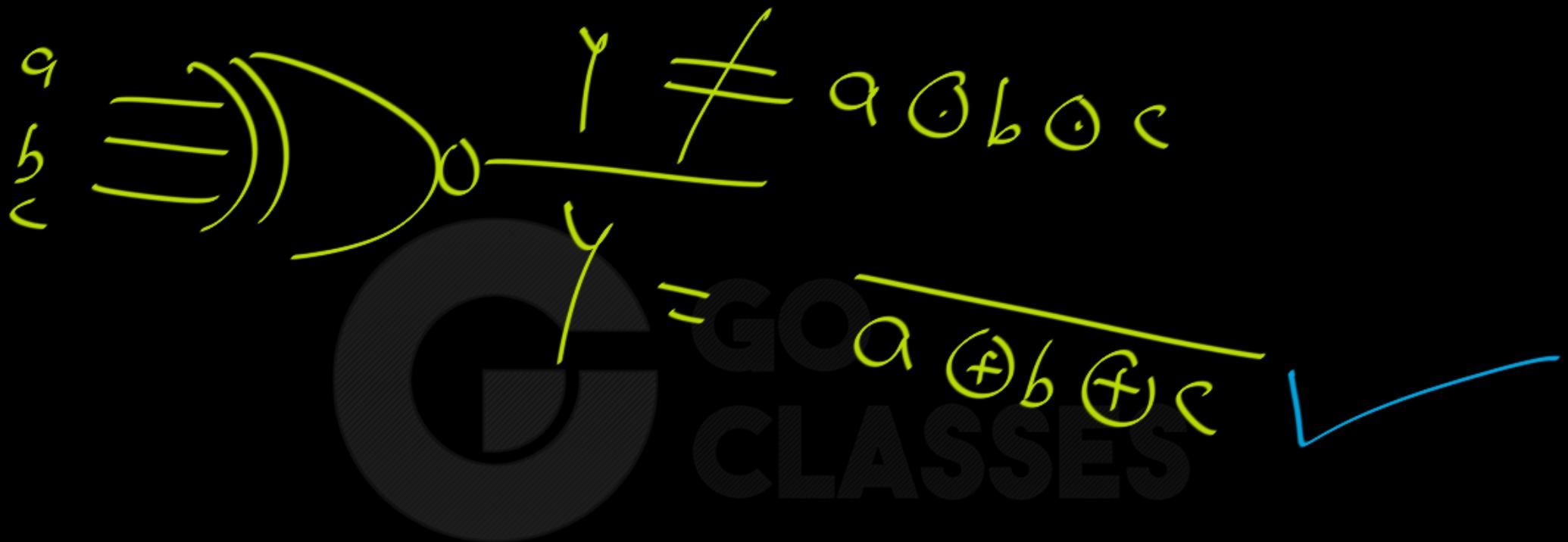
$$= a \odot b \odot c$$

Ans :





Digital Logic



GATE ECE 2015 Set 1 | Question: 38

asked in Digital Circuits Mar 28, 2018 • recategorized Nov 15, 2020 by soujanyareddy13

81 views

- 1 A 3-input majority gate is defined by the logic function $M(a, b, c) = ab + bc + ca$. Which one of the following gates is represented by the function $M(\overline{M(a, b, c)}, M(a, b, \bar{c}), c)$?
- A. 3-input NAND gate
 - B. 3-input XOR gate
 - C. 3-input NOR gate
 - D. 3-input XNOR gate

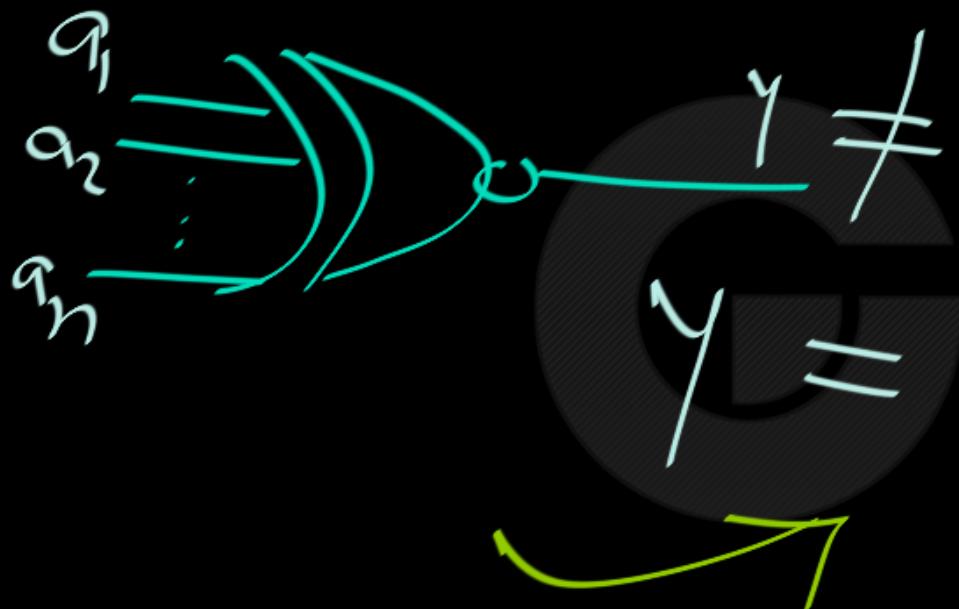
$$\begin{aligned} &= abc + \bar{a}\bar{b}c + \bar{a}b\bar{c} + ab\bar{c} \\ &= a \oplus b \oplus c \end{aligned}$$

gate2015-ec-1

digital-circuits

logic-gates

Very Important:



$$\begin{array}{c} y \neq q_1 \odot q_2 \odot \dots \odot q_n \\ \hline y = q_1 \oplus q_2 \oplus \dots \oplus q_n \end{array}$$

XNOR gate is Always Complement of XOR gate.