



First Order Logic  
Next Chapter:

# Practice

Bounded Variable, Free Variable  
Scope of a Quantifier

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GATE CSE AIR 53; AIR 67;  
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# First Order Logic

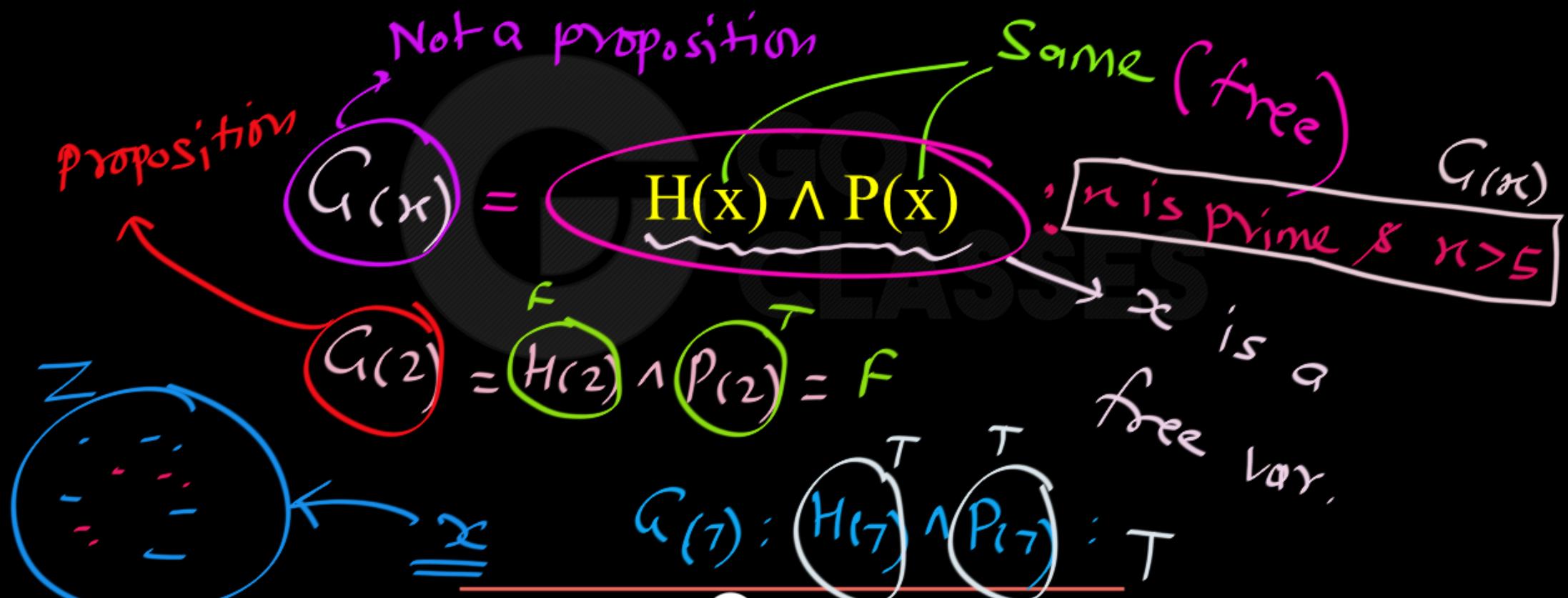
Next Topic:

G GO  
PracticeS

Bounded Variable, Free Variable

Assume: Domain: Set of Integers

$P(x)$ :  $x$  is Prime.       $H(x)$ :  $x > 5$ .



## Mathematics :

$$y = x^2 + x + 2$$

free (same  $n$ )

$$x=2$$

$$y = 2^2 + 2 + 2 \quad ; \quad y = (10)^2 + 10 + 2$$

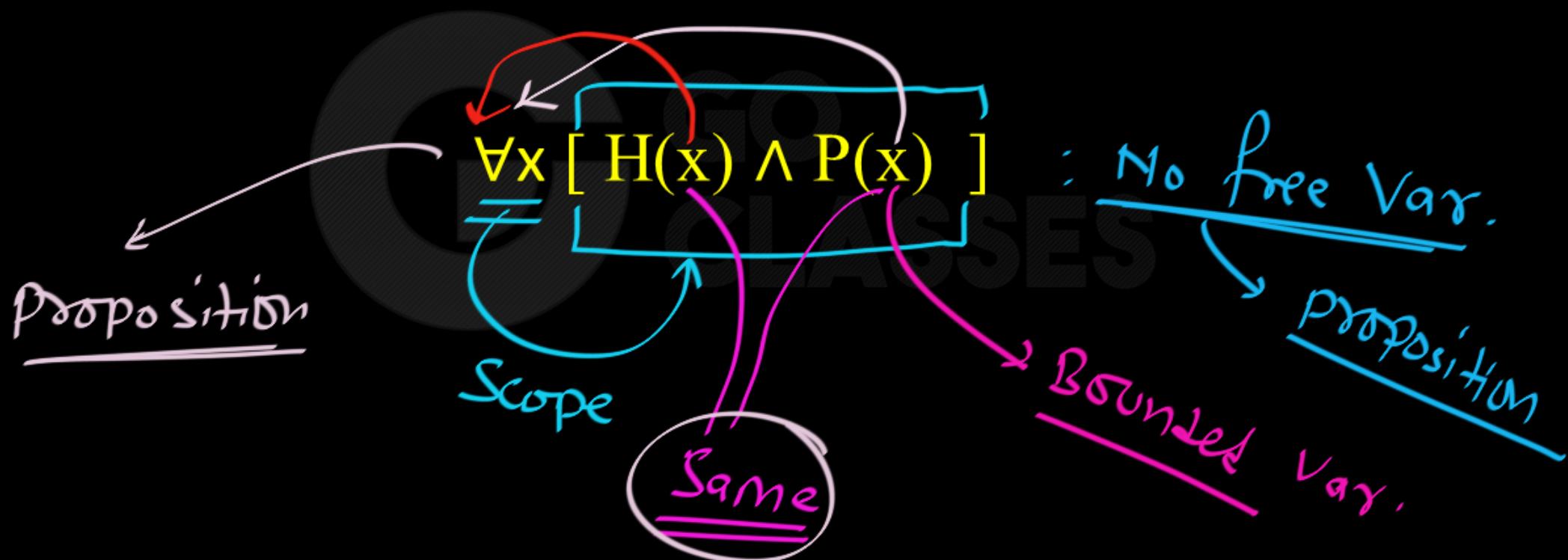
$$y = 3^2 + 4 + 2$$

WRONG



Assume: Domain: Set of Integers

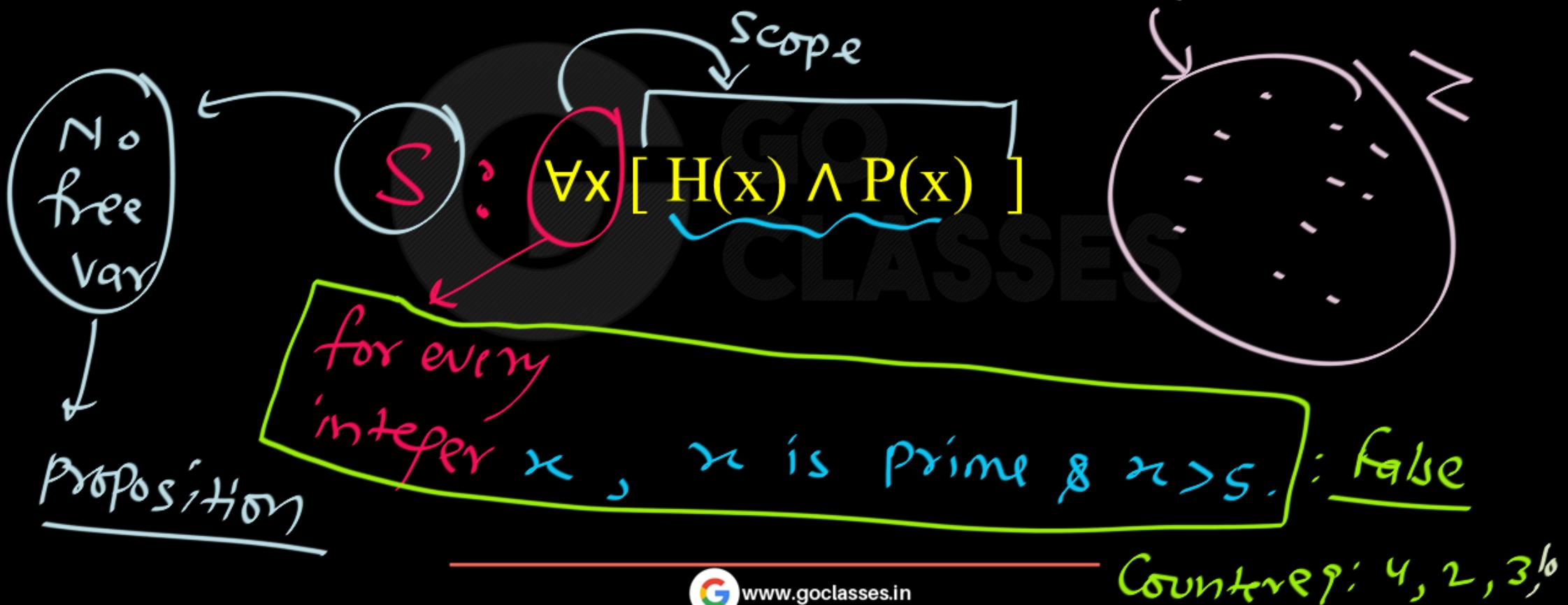
P(x): x is Prime.      H(x): x > 5.



Assume: Domain: Set of Integers

P(x): x is Prime.

H(x): x > 5.





Assume: Domain: Set of Integers

P(x): x is Prime.      H(x): x > 5.

$$\forall x [ H(x) \wedge P(x) ]$$

false

Counterex:  $x = 2$

$$x = 10$$

$$x = 8$$

$$x = 5 \checkmark$$

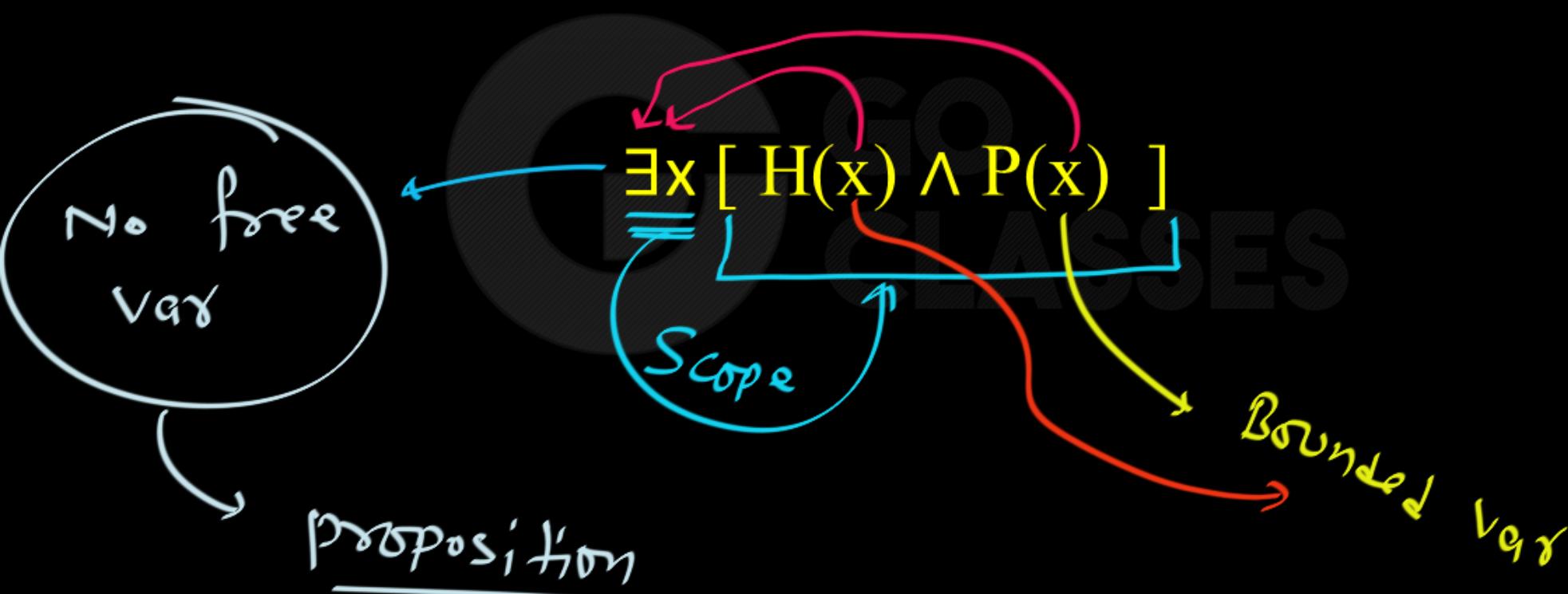
$$x = 4 \checkmark$$

$$x = -5 \checkmark$$



Assume: Domain: Set of Integers

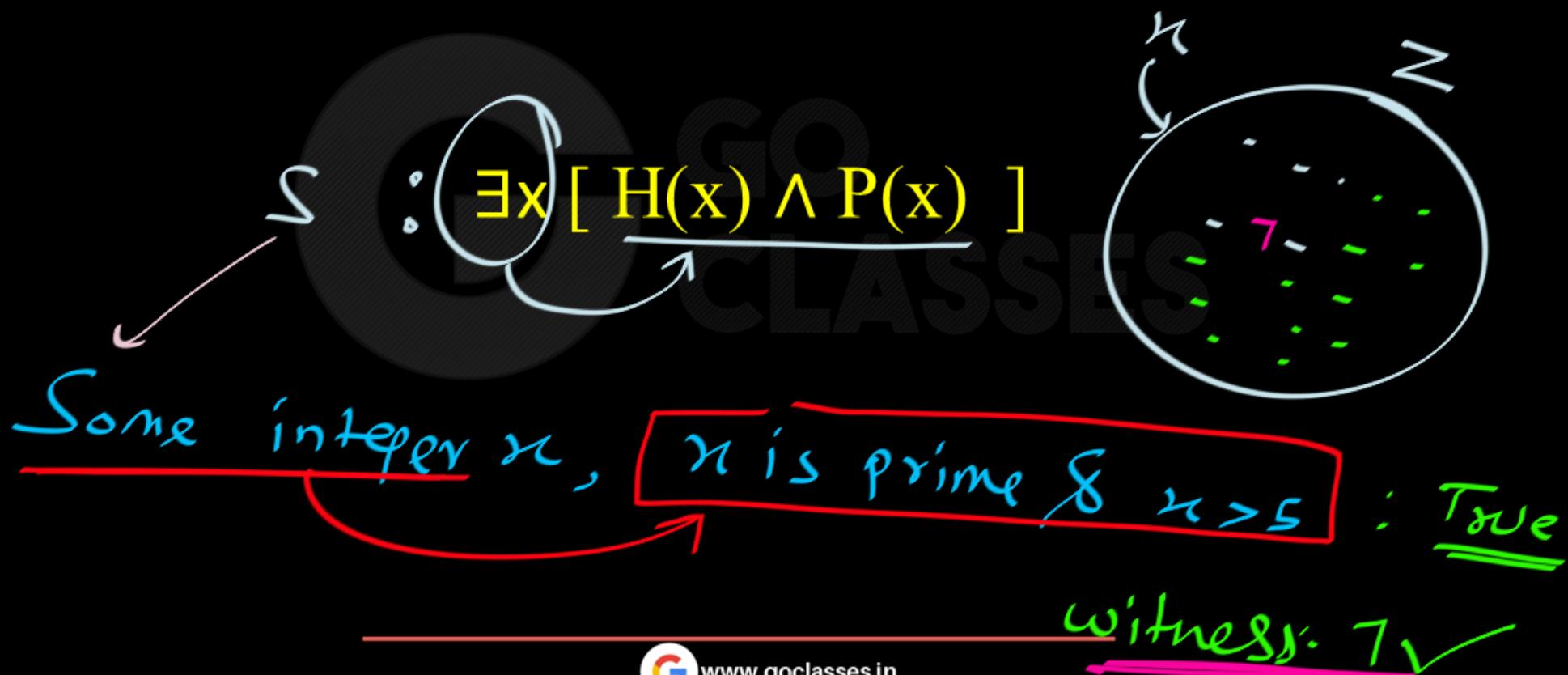
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Assume: Domain: Set of Integers

P(x): x is Prime.      H(x): x > 5.





Assume: Domain: Set of Integers

P(x): x is Prime.      H(x): x > 5.

$$\exists x [ H(x) \wedge P(x) ]$$

: True



Assume: Domain: Set of Integers

P(x): x is Prime.      H(x): x > 5.

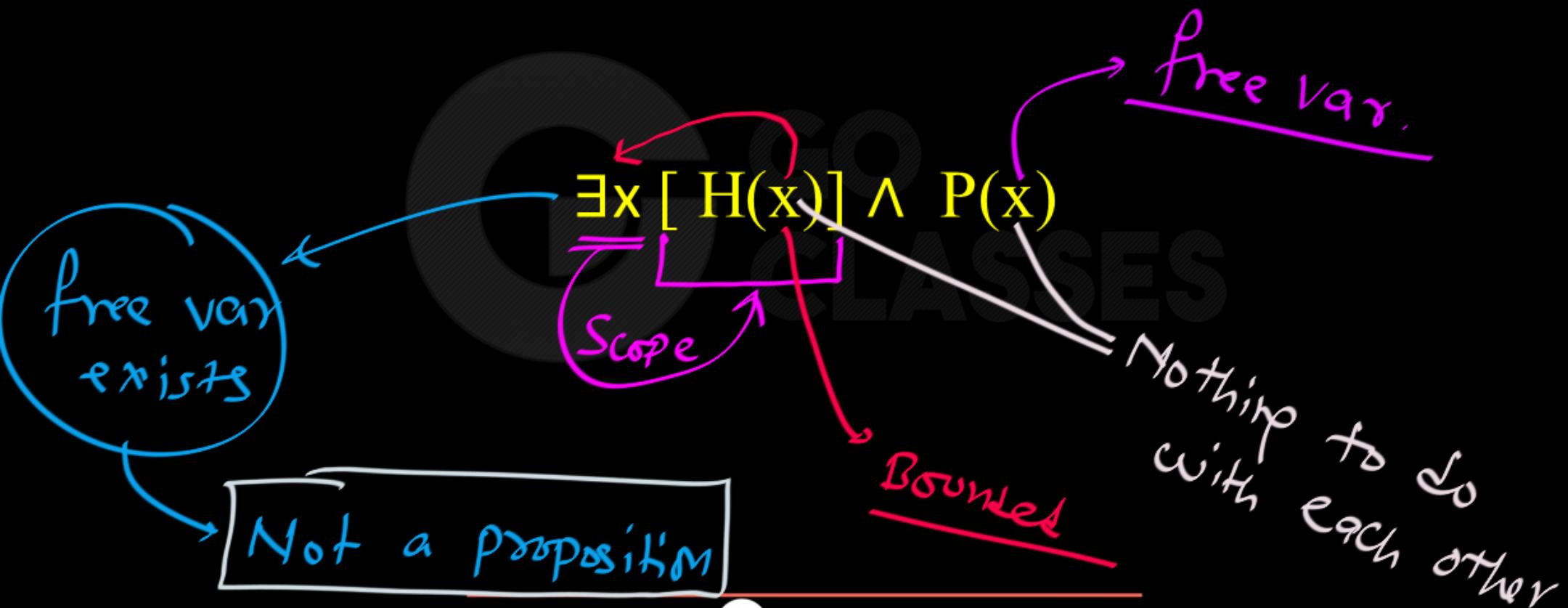
$$\exists x [ H(x) \wedge P(x) ]$$



Assume: Domain: Set of Integers

P(x): x is Prime.

H(x): x > 5.

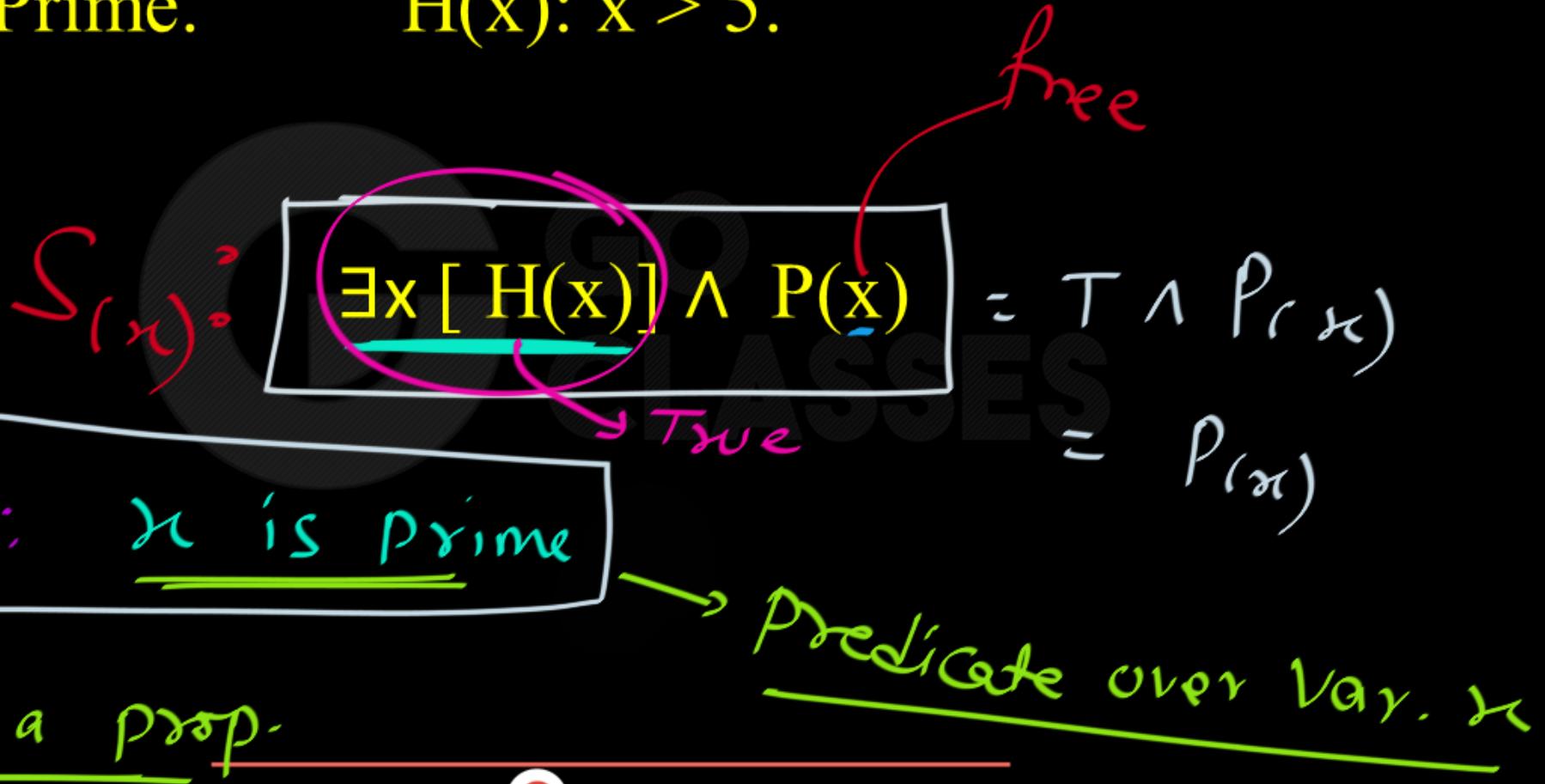




Assume: Domain: Set of Integers

$P(x)$ :  $x$  is Prime.

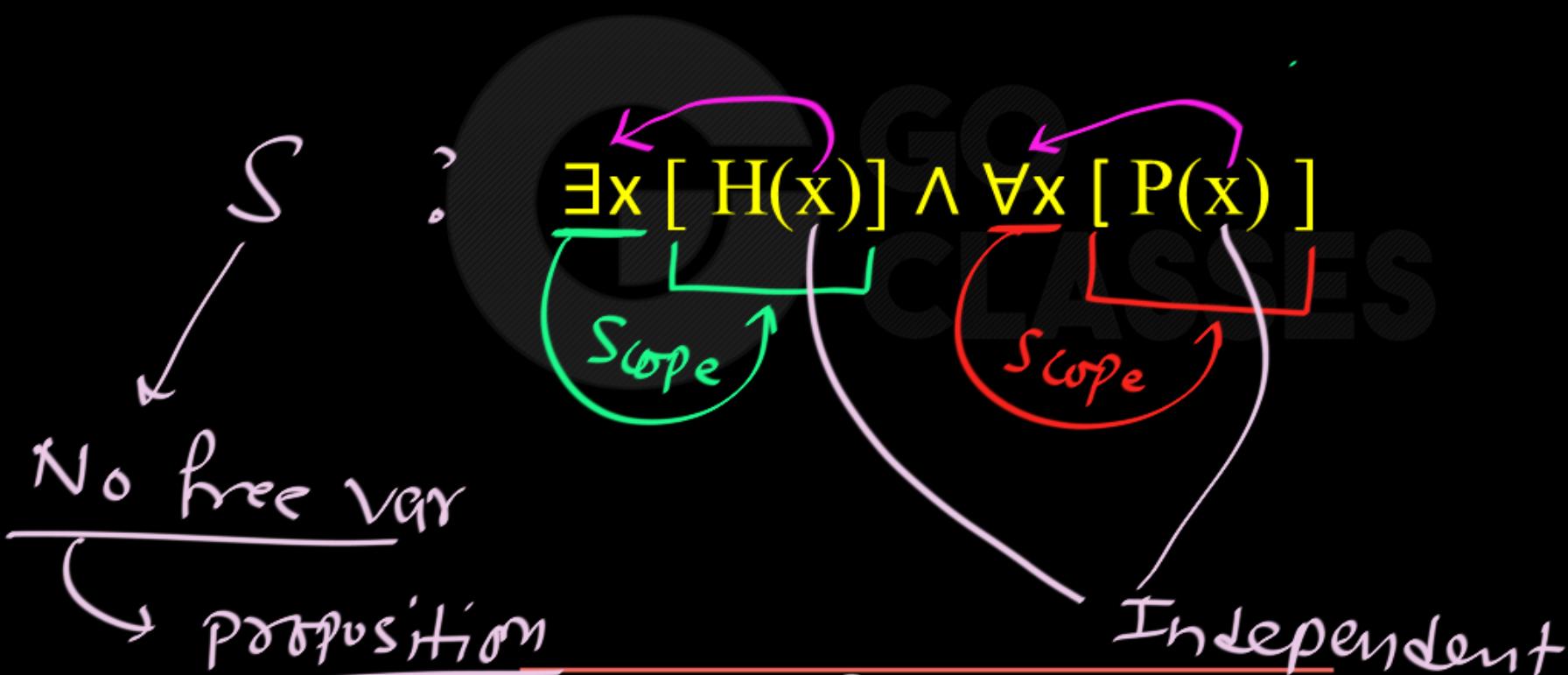
$H(x)$ :  $x > 5$ .





Assume: Domain: Set of Integers

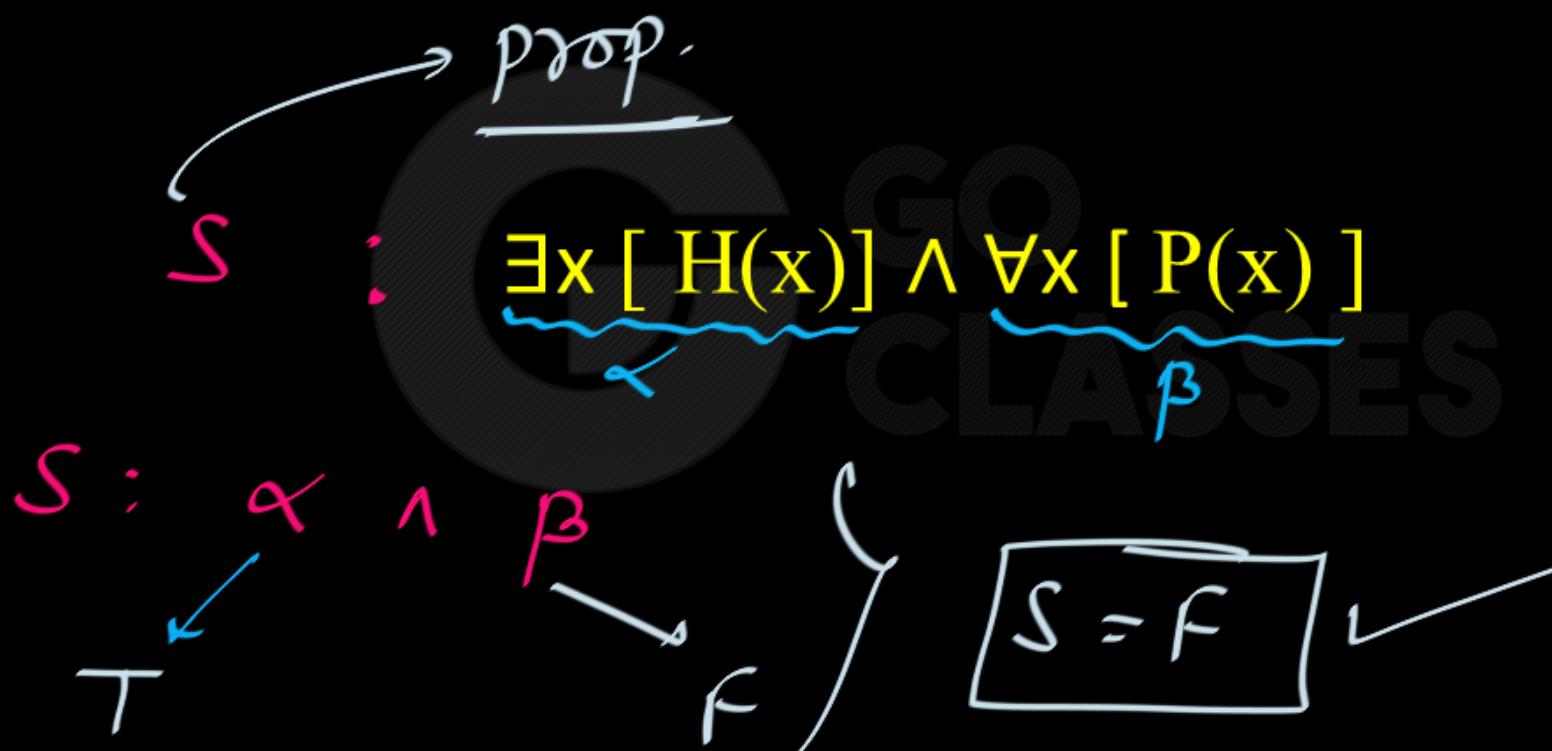
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Assume: Domain: Set of Integers

P(x): x is Prime.      H(x): x > 5.





Assume: Domain: Set of Integers

P(x): x is Prime.      H(x): x > 5.

$$\exists x [ H(x) ] \vee \forall y [ P(y) ]$$



Assume: Domain: Set of Integers

P(x): x is Prime.      H(x): x > 5.

S :

$$\exists x [ H(x) ] \vee \forall y [ P(y) ]$$

No free  
var

Proposition



Assume: Domain: Set of Integers

P(x): x is Prime.      H(x): x > 5.

$$\mathcal{S} : \exists x [ H(x) ] \vee \forall y [ P(y) ]$$

$\alpha$        $\beta$

$$\mathcal{T} : \mathcal{F}$$

$\mathcal{S} = \mathcal{F}$



$$\boxed{\exists_n P(n) \wedge \forall_n P(n)} \quad \equiv \quad \boxed{\exists_n P(n) \wedge \forall_y P(y)}$$

*same*

Q What is the difference between these three formulae?

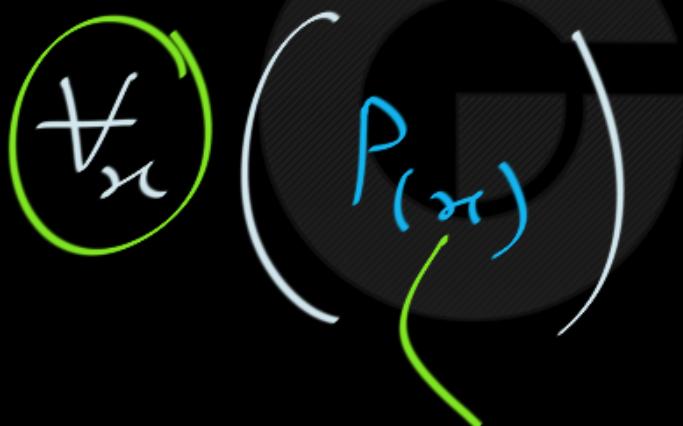
$G(x) : \text{Hx} \wedge Px$  → free var. x  
Same

$S : \forall x[\text{Hx} \wedge Px]$  : proposition (No free var)

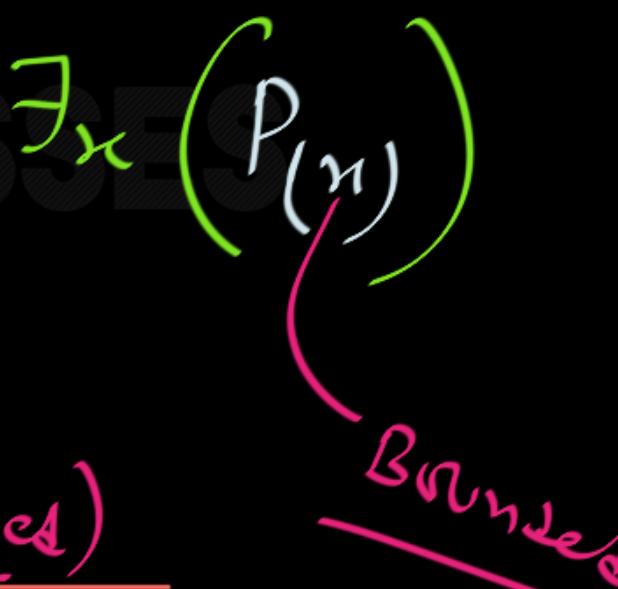
$M(x) : \forall x[\text{Hx}] \wedge Px$  → Not a prop.  
Predicate



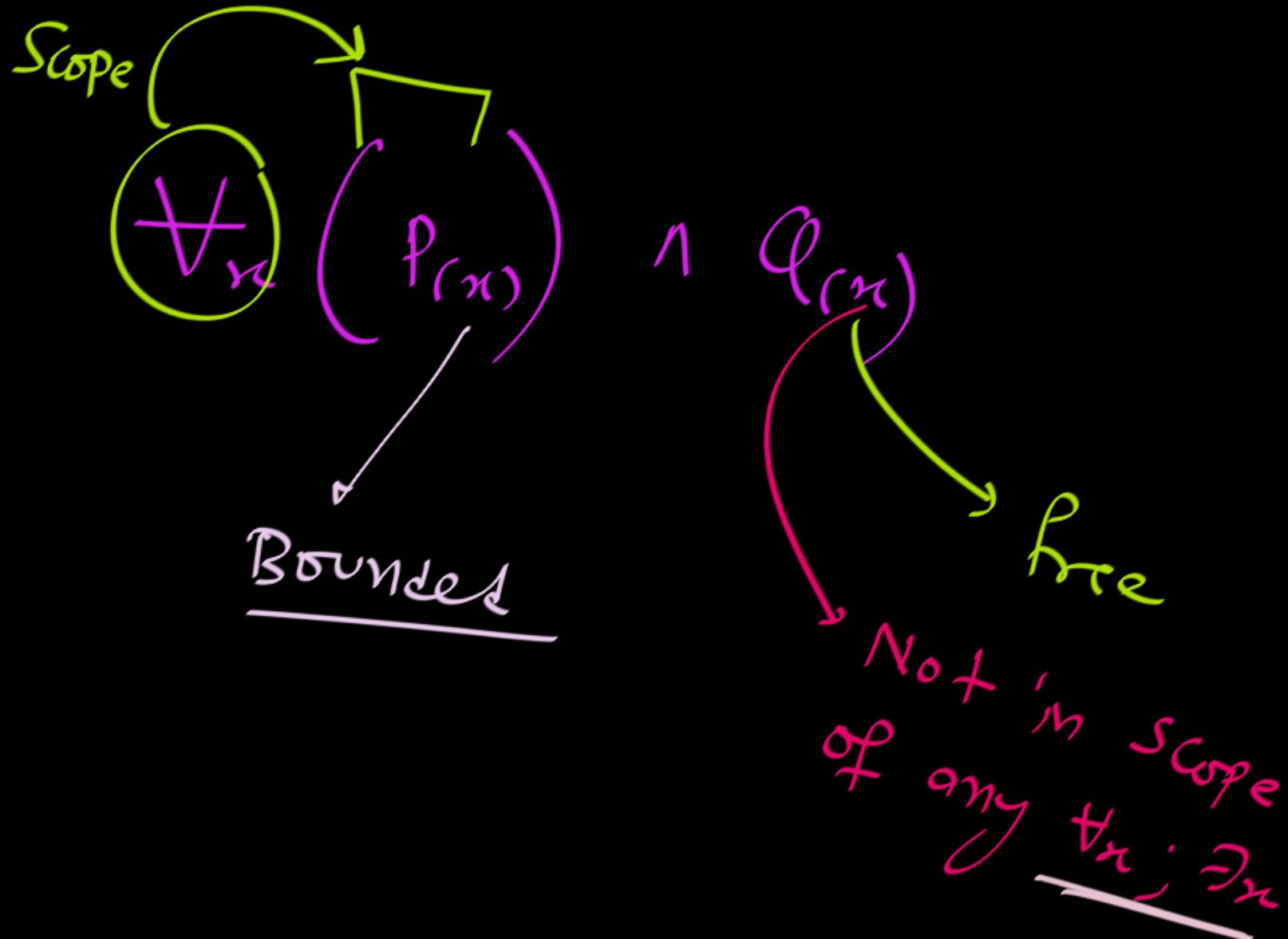
An occurrence of  $x$  within the scope of  $\forall x$  or  $\exists x$  is bound, otherwise it is free.



Bounded (Quantified)



Bounded





**Bound variables:** Occurrence of a variable is bound *iff* it is within the scope of a quantifier.

**DESIRABLE**

**Free variables:** Occurrence of a variable is free *iff* it is NOT within the scope of a quantifier.

**NOT DESIRABLE**



free var. exists  $\longleftrightarrow$  Not a Proposition

Proposition  $\longleftrightarrow$  No free var



**Q:** In the following predicate formula, which *variables* are *free* and which are *bound*?

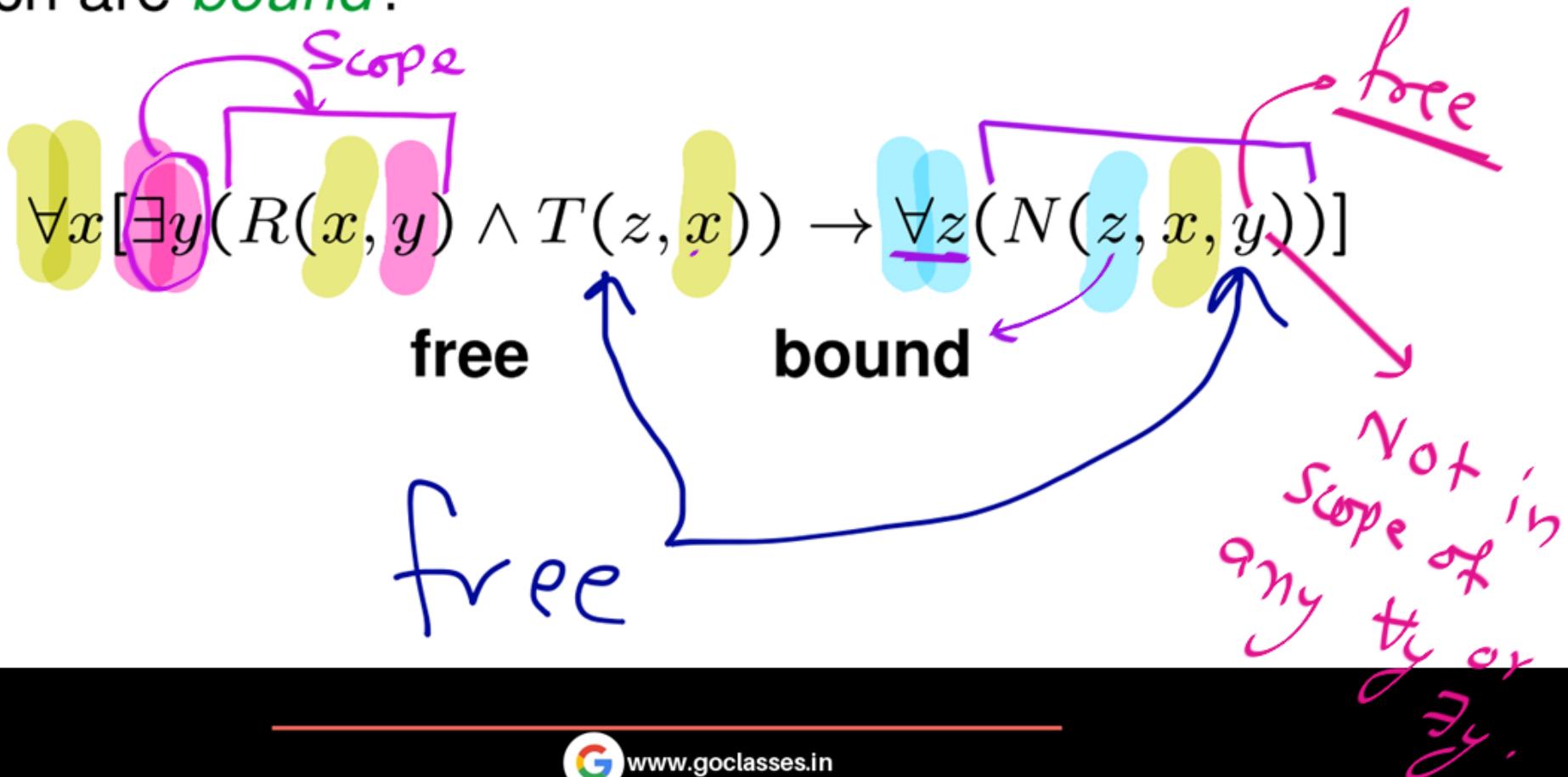
$$\forall x [\exists y (R(x, y) \wedge T(z, x)) \rightarrow \forall z (N(z, x, y))]$$

free                          bound  
free

The formula is  $\forall x [\exists y (R(x, y) \wedge T(z, x)) \rightarrow \forall z (N(z, x, y))]$ . The regions are color-coded: yellow for free  $x$ , pink for free  $y$ , light blue for bound  $z$ , and yellow for bound  $z$  again. Blue arrows point from the labels 'free' and 'bound' to their respective colored regions.

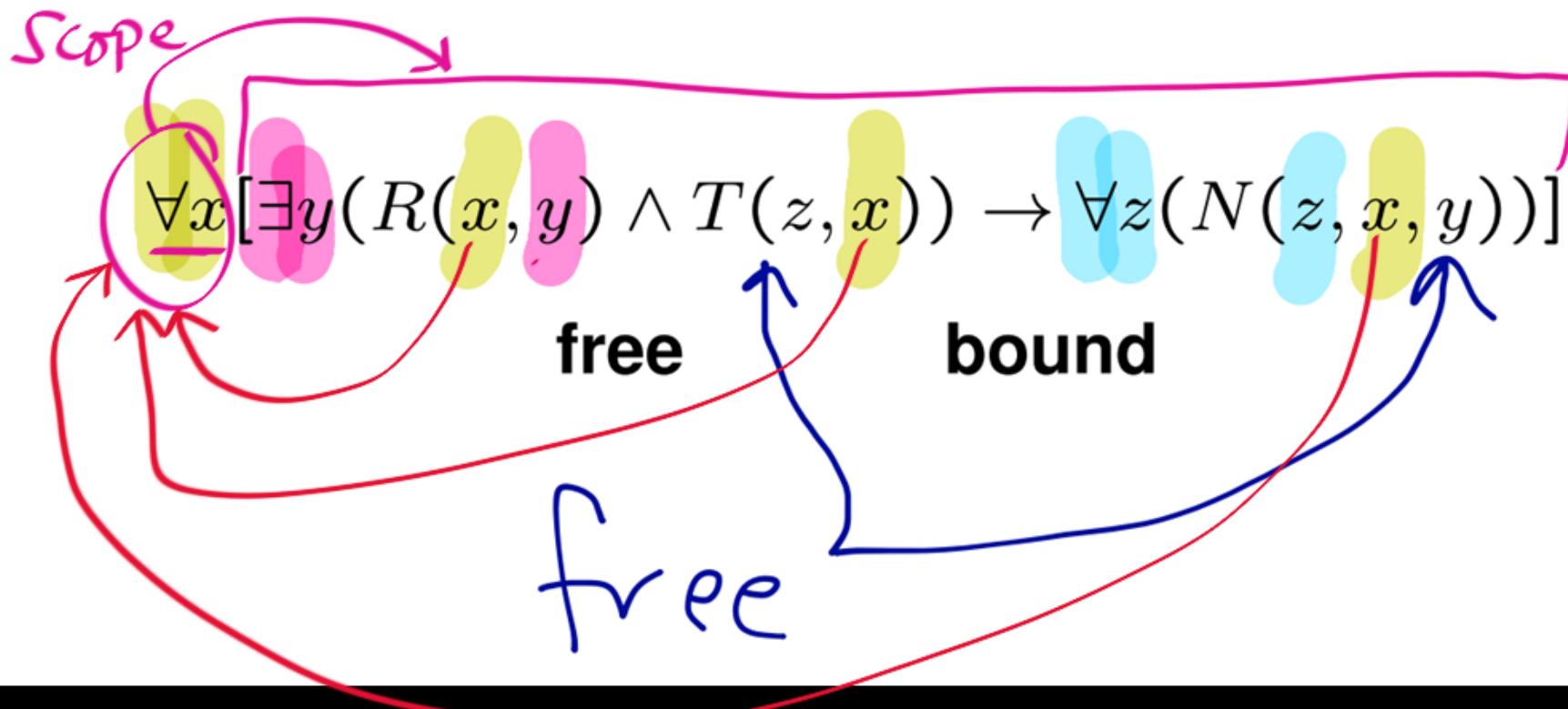


Q: In the following predicate formula, which **variables** are **free** and which are **bound**?





Q: In the following predicate formula, which *variables* are *free* and which are *bound*?

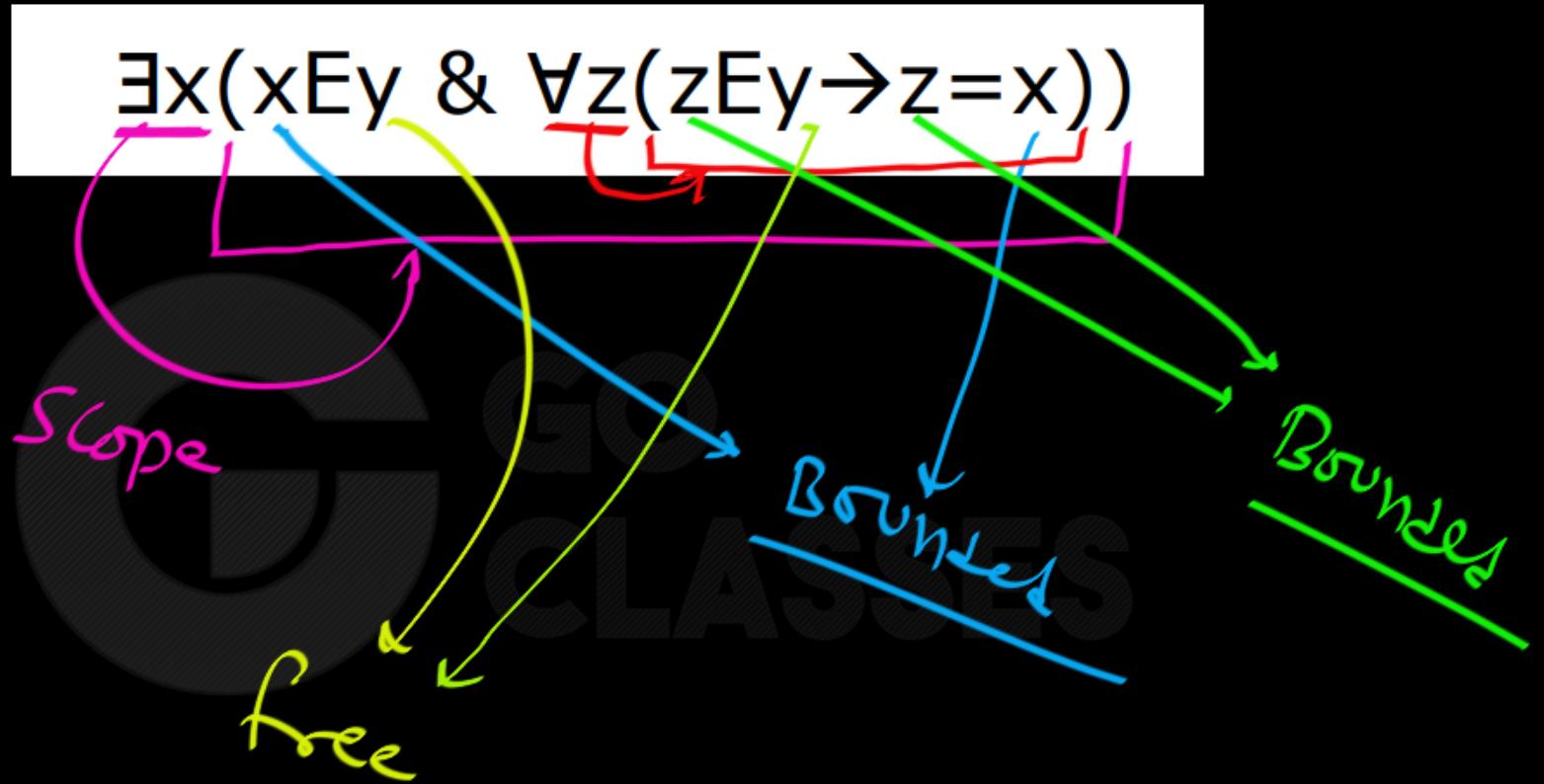




$$\exists x(xEy \ \& \ \forall z(zEy \rightarrow z=x))$$

Find free & Bounded Var.






$$\exists x(xEy \ \& \ \forall z(zEy \rightarrow z=x))$$

Bound occurrence



$$\exists x(xEy \ \& \ \exists y(\neg yEx))$$

finl Bounded & free Var.





$$\exists x(x \in y \text{ & } \exists y(\neg y \in x))$$

Scope

Bounded

Bounded

Free

$\exists x(xEy \ \& \ \exists y(\neg yEx))$  $\exists x(xE\textcolor{pink}{y} \ \& \ \exists \textcolor{pink}{y}(\neg \textcolor{pink}{y}Ex))$ 

y free

y bound

# Quantifiers: Example

Consider this formula

$$\Phi = \forall y. ((P(y) \wedge Q(x)) \vee \forall x. (P(x) \vee Q(y) \wedge P(z)))$$

How many free variables does this formula contain? **Clicker**

- ① One
- ② Two
- ③ Three
- ④ Four
- ⑤ None

# Quantifiers: Example

Consider this formula

$$\Phi = \forall y. ((P(y) \wedge Q(x)) \vee \forall x. (P(x) \vee Q(y) \wedge P(z)))$$

How many free variables does this formula contain? Clicker

- ① One
- ② Two
- ③ Three
- ④ Four
- ⑤ None

Bounded

Free

Bounded  
Bounded

Free



Q: For each formula, determine which of its two variables is free. Determine whether the formula is true when the value of the free variable is 0. If it is true then find a witness.  
(Assume domain is set of real numbers)

(a)  $\exists n(x > 2^n).$

(b)  $\exists x(x > 2^n).$



Q: For each formula, determine which of its two variables is free.  
Determine whether the formula is true when the value of the  
free variable is 0. If it is true then find a witness.  
(Assume domain is set of real numbers)

$$(a) \exists n(x > 2^n).$$

free

bounded

$$(b) \exists x(x > 2^n).$$

free

bounded



Q: For each formula, determine which of its two variables is free.  
Determine whether the formula is true when the value of the  
free variable is 0. If it is true then find a witness.  
(Assume domain is set of real numbers)

(a)  $\exists n(x > 2^n).$  :  $S_{(n)}$

(b)  $\exists x(x > 2^n).$  :  $G_{(n)}$

free

Not propositions

Q: For each formula, determine which of its two variables is free.  
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(Assume domain is set of real numbers)

(a)  $\exists n(x > 2^n). : S(x)$

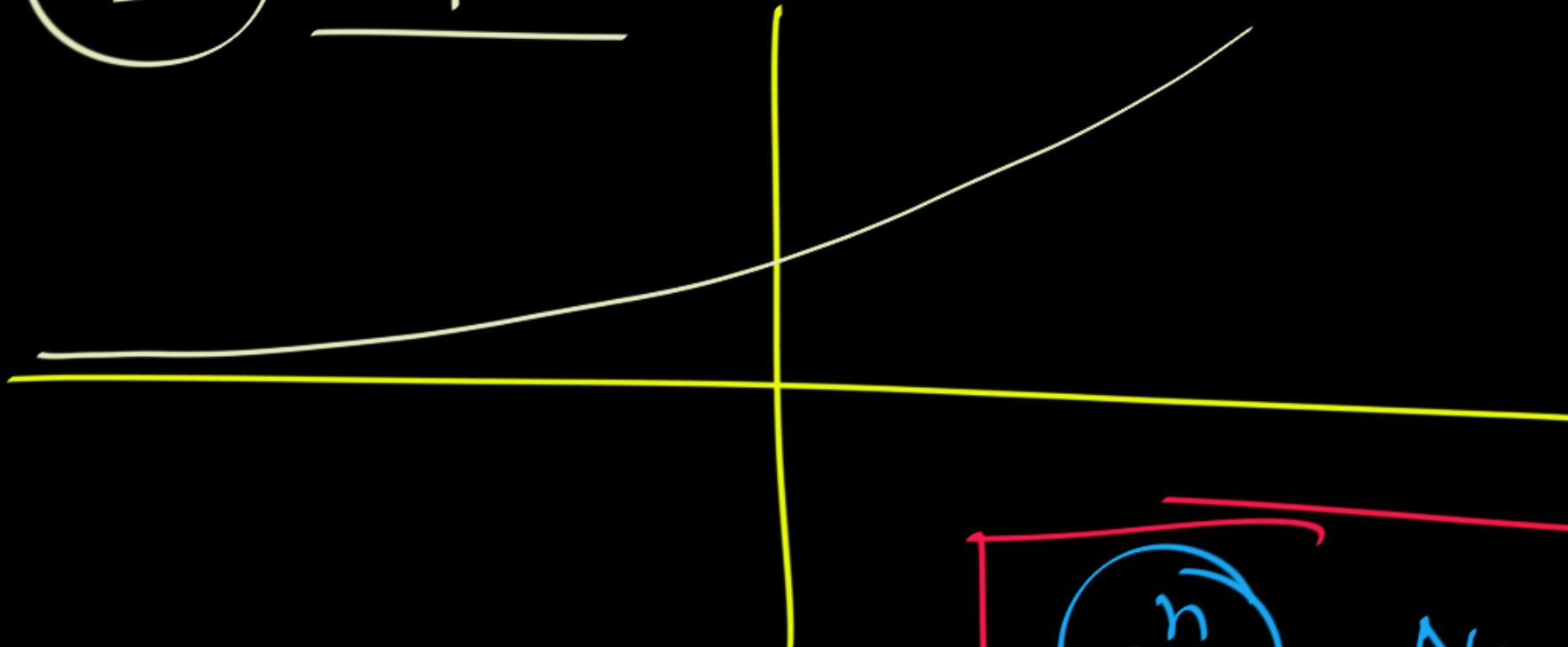
$S(0) : \exists n(0 > 2^n)$

(b)  $\exists x(x > 2^n).$



$2^n$

Graph



$2^n$  : Never negative

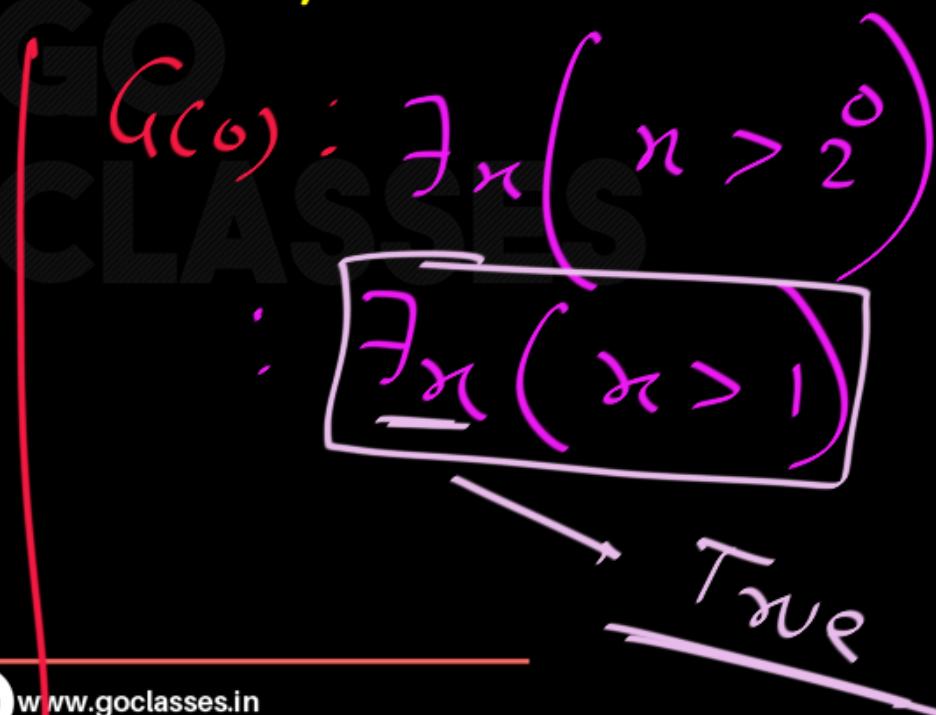


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(a)  $\exists n(x > 2^n).$

(b)  $\boxed{\exists x(x > 2^n).} : G(n)$

Free





- { ① Scope of a Quantifier
- ② free Var
- ③ Bounded Var

- We need quantifiers to formally express the meaning of the words “all” and “some”.
- The two most important quantifiers are:
  - ▶ Universal quantifier, “For all”. Symbol:  $\forall$
  - ▶ Existential quantifier, “There exists”. Symbol:  $\exists$
- $\forall x P(x)$  asserts that  $P(x)$  is true for **every**  $x$  in the domain.
- $\exists x P(x)$  asserts that  $P(x)$  is true for **some**  $x$  in the domain.
- The quantifiers are said to **bind** the variable  $x$  in these expressions.
- Variables in the scope of some quantifier are called **bound variables**. All other variables in the expression are called **free variables**.
- A propositional function that does not contain any free variables is a proposition and has a truth value.

# Free and bound variables

Two essentially different ways in which we use individual variables in first-order formulae:

1. **Free variables**: used to denote *unknown or unspecified objects*, as in  $(5 < x) \vee (x^2 + x - 2 = 0)$ .
2. **Bound variables**: used to *quantify*, as in  
 $\exists x((5 < x) \vee (x^2 + x - 2 = 0))$   
and  $\forall x((5 < x) \vee (x^2 + x - 2 = 0))$ .

Note that the same variable can be *both free and bound in a formula*, e.g.  $x$  in the formula  $x > 0 \wedge \exists x(5 < x)$ .

A formula with no bound variables is an **open formula**.

A formula with no free variables is a **closed formula**, or a **sentence**.