

- 1) Applicative
2) Monads

Eq: a

class Functor t where

fmap :: $(a \rightarrow b) \rightarrow t\ a \rightarrow t\ b$

list :: Type \rightarrow Type

[] :: [a]

[] :: Type \rightarrow Type

[] Int = [Int]

[] Bool = [Bool]

instance Functor [] where

fmap :: $(a \rightarrow b) \rightarrow [a] \rightarrow [b]$

fmap = map

~~functor~~ fmap id = id

fmap f . fmap g = fmap (f . g)

These laws should be followed by the programmer

instance Functor Maybe where

fmap :: $(a \rightarrow b) \rightarrow \text{Maybe}\ a \rightarrow \text{Maybe}\ b$

fmap f (Just a) = Just (f a)

fmap f Nothing = Nothing

(.) :: $(b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c)$

(\\$) :: $(a \rightarrow b) \rightarrow a \rightarrow b$

Just $\underbrace{(f a)}_{\text{link}}$ can be written as $\$ f a$
link will suggest to replace such code with $\$$

```
data Tree a = Empty  
            | Node (Tree a) a (Tree a)
```

instance Functor Tree where

fmap :: $(a \rightarrow b) \rightarrow (\text{Tree } a) \rightarrow (\text{Tree } b)$

$\text{fmap } f (\text{Node } l \ a \ r) = \text{Node } (\text{fmap } f l) (f a) (\text{fmap } f r)$

$\text{fmap } f \text{ Empty} = \text{Empty}$

"There is only one rule in the Fight Club, i.e.

"There is no rule in the Fight Club"

-Fight Club

"Recursion is the Big Daddly of looping; Anything that you can do with loop it can be done with recursion"

- Piyush

class Functor t where

`fmap :: (a → b) → t a → t b`

instance Functor Maybe where

$$\text{fmap } f \text{ (Maybe } a) = \text{f} \circ \text{Maybe } (f a)$$

$f \circ g \circ f \quad \text{Nothing} = \text{Nothing}$

instance Functor `[]` where `[]` is used as a

$$[c]_a - [a]$$

instance Fundov Tree where

Properties to follow

$$f_{\text{max}} \circ d = d$$

$$\text{fmap } f \text{ fmap } g = \text{fmap} (f \cdot g)$$

Io a

+0 u
type of i/o actions that result in a value
of type α .

gLine :: 10String

~~read :: "123" :: Int~~ read :: Read a \Rightarrow String \rightarrow a

(read "123" :: Int)

123 :: Int

foo :: String → Int

foo ← read → read

foo s = read s + 1

-- we write the different definition on the of the getLine
getLine :: IO Int

IO is known to be an Functor :- It has a fmap.

getLine = fmap read getLine

foo = fmap read (Read a, Functor f) ⇒

tString → t a.

data Person = Person String

dat type foo = Int -- type aliasing

data Person = Person String.

newtype Person = Person String

14 : Head - Define a function $\text{getPerson} :: \text{IO Person}$.
make use of newtype Person.

$\text{getPerson} = \text{fmap Person getLine}$.

Applicative

$\text{data Person} = \text{Person String Int}$

$\text{get} = \text{fmap Person getLine}$

$\text{get} :: \text{IO}(\text{Int} \rightarrow \text{Person})$

we can't ~~do~~ ~~get~~ get IO Person (the person).

There is no function ~~to~~ to get 'a' from 'IO a'

$\text{IO a} \rightarrow a$. There is no function of this type.

This is made by design.

$\text{class Functor } t \Rightarrow \text{Applicative } t \text{ where}$

$$(\#) : (a \rightarrow b) \rightarrow t a \rightarrow t b$$

$\text{pure} :: a \rightarrow t a$

$(\langle *\rangle) :: t(a \rightarrow b) \rightarrow t a \rightarrow t b$

$\text{get} = (\text{fmap Person getLine}) \leftarrow \text{fmap getLine}$

Q) Read 3 integers from the user and add them up.

$\text{add2} :: \text{Int} \rightarrow \text{Int} \rightarrow \text{Int} \rightarrow \text{Int}$

$\text{add3 } x \ y \ z = x + y + z$

$\text{fmap add3} :: \text{IO(Int)} \rightarrow \text{IO(Int} \rightarrow \text{Int} \rightarrow \text{Int})$

$((\text{fmap add3 getIntLine}) \leftarrow \text{getIntLine}) \leftarrow \text{getIntLine}$

$f \$ t_n = \text{fmap } f \ t_n$

$f \$ x = f x$

\therefore It can also be written as

$\text{add3} \$ \text{getIntLine} \leftarrow \text{getIntLine} \leftarrow \text{getIntLine}$

$f: a \rightarrow b \rightarrow c \rightarrow d \rightarrow e$

$t_a :: \text{IO } a$

$t_b :: \text{IO } b$

$t_c :: \text{IO } c$

$t_d :: \text{IO } d$

$f \leftarrow f \langle \$ \rangle t_a \langle + \rangle t_b \langle \Rightarrow \rangle t_c \langle * \rangle t_d \quad :: \text{IO}(e)$

Law of Applicative

$$\text{pure } f \langle * \rangle t_a \equiv \text{fmap } f t_a$$

$$\text{pure } f \langle * \rangle t_a \equiv \text{fmap } f t_a$$

Hint

data Person = Person String String Int

define getPerson for this definition of person.

getPerson = Person <\\$> getLine <*> getLine <#> getIntLine

Hint:

instance Applicative Maybe where

4/3/24

class Functor t \Rightarrow Applicative t where

pure :: $a \rightarrow t a$

($\langle * \rangle$) :: $t(a \rightarrow b) \rightarrow t a \rightarrow t b$

instance Applicative Maybe where

-- pure :: $a \rightarrow \text{Maybe } a$

pure = Just

-- (\leftrightarrow) :: Maybe (a \rightarrow b) \rightarrow Maybe a \rightarrow Maybe b

(\leftrightarrow) (Just f) (Just a) = Just \$ f a

(\leftrightarrow) _ _ = None Nothing

data Exp = C Int

| PLUS Exp Exp

| DIV Exp Exp

eval :: Exp \rightarrow Maybe (Int)

eval (Cx) = Just x

eval (DIV e₁ e₂) = -- complete yourself

eval (PLUS e₁ e₂) = case eval e₁ of

Just x \rightarrow case eval e₂ of

Just y \rightarrow Just xy

Nothing \rightarrow Nothing

Nothing \rightarrow Nothing.

but this is very large definition and not neat
and ~~plus~~ plus.

$f: a \rightarrow b \rightarrow c \rightarrow d$

$- f <\$> ta = \text{map } f \text{ to}$

$ma : \text{Maybe } a$

$f <\$> ta = \text{pure } f <\$> ma$

$mb : \text{Maybe } b$

$mc : \text{Maybe } c$

$\text{eval plus} = (+) <\$> \text{eval } (e_1) \quad (+) \text{ eval } (e_2)$

instance Applicative: $[]$ where

$\text{pure} :: a \rightarrow []$

$\text{pure } x = []$

$(\langle *\rangle) :: [a \rightarrow b] \rightarrow [a] \rightarrow [b]$

$fs \langle * \rangle xs = [fx \mid f \leftarrow fs, x \leftarrow xs]$

Notation:

$[x:y \mid x \leftarrow x_1, y \leftarrow y_1], \quad x < 10, y < 10^0]$

Step

bottom

top

at

in

12 maybe

adjective : ["cute", "age", "innovative"]

noun : ["product", "technology"]

join $x : ys = \text{unwords}(xs, ys)$

all = join $\langle\$ \rangle$ adjective $\langle*\rangle$ noun

11-3-24

class Functor t \Rightarrow Applicative t where

pure :: $a \rightarrow t a$

$\langle * \rangle :: t(b \rightarrow a) \rightarrow ta \rightarrow tb$

$f \langle \$ \rangle [x_1, \dots, x_n] \langle * \rangle [y_1, \dots, y_m]$

= $f[x_i y_j \dots]$

$[f_1, \dots, f_m] \langle * \rangle [y_1, \dots, y_n]$

$[\dots, f_i y_j \dots]$

$g :: a \rightarrow b \rightarrow c \rightarrow d$

$g \langle \$ \rangle [x_1, \dots, x_l] \langle * \rangle [y_1, \dots, y_m] \langle * \rangle [z_1, \dots, z_k]$

the earlier definition of Applicative of a list doesn't compute all possible combinations.

(Cauchy sequence)

so,

$$\forall \epsilon \exists n \quad \forall m_1, m_2 > n \quad |x_{m_1} - x_{m_2}| < \epsilon$$

newtype ZipList a = ZipList [a]

instance Functor ZipList where

$$fmap :: (a \rightarrow b) \rightarrow ZipList a \rightarrow ZipList b$$

$$fmap f (ZipList []) = ZipList []$$

$$fmap f (ZipList (x : xs)) = * ZipList ($$

$$fmap f ZipList lst = ZipList (fmap f lst)$$

instance instance Applicative ZipList where

$$pure :: a \rightarrow ZipList a$$

$$(\langle * \rangle) :: ZipList ((a \rightarrow b) \rightarrow ZipList a \rightarrow ZipList b)$$

(\Leftarrow) $(\text{ZipList } fs) (\text{ZipList } xs) = \text{ZipList} (\text{zipWith } (\lambda) \text{ } fs \text{ } xs)$

-- $fs : [a \rightarrow b]$

-- $xs : [a]$

-- to write ZipList_3 or ZipList_n

-- $f <\$> \text{ZipList} [1, 2, 3] \leftrightarrow \text{ZipList} [\dots] \leftrightarrow$

we need to satisfy the problem

pure $f <\#> xs = fmap f xs$

pure $x = \text{ZipList } xs = ((x : x) : xs) : xs$

where $xs = x : xs$

13-3

• Functor : Applying pure function "inside" the data structure
→ $fmap : (a \rightarrow b) \rightarrow [a] \rightarrow [b]$

• Applicative | Monoidal functors : Extending Functor interface
(Applicative Functors)
for "multi-argument" function
→ $(\langle \# \rangle) : t (a \rightarrow b) \rightarrow ta \rightarrow tb$
(apply)

• Monad : ($\gg=$) Bind operator

(1) Read an int, read a string, ...

(2) Read an int, if n is read them read n things.

$\langle \cdot \rangle$: map of

$\langle * \rangle$: apply of

$\gg=$: bind of

$\gg= :: ta \rightarrow (a \rightarrow tb) \rightarrow tb$

class Applicative t \Rightarrow Monad t where

return :: $a \rightarrow ta$ λ : return is same as pure
is a useless relic from the past

$(\gg=) :: ta \rightarrow (a \rightarrow tb) \rightarrow (tb)$

$(\gg) :: ta \rightarrow tb \rightarrow tb$ -- $(\gg) ta tb = (\gg=) ta (\lambda a \rightarrow tb)$

$(\langle * \rangle) tf ta = tf \gg= (\lambda f \rightarrow (ta \gg= (\lambda a \rightarrow \text{return}(ta))))$

Syntactic Sugar

The do-notation

stmt = $\text{do } \overbrace{\text{stmt}}^{\text{ta}} \text{ do } \overbrace{\text{action}}^{\text{ta}}; \text{stmt}$ (1)

| do $\lambda \leftarrow \text{action}; \text{stmt}$ (2)

| action.

(1) action $\gg=$ stmt

(2) action $\gg= \lambda x \rightarrow \text{stmt}$

$(\langle \star \rangle) \quad \text{if } ta = \text{ do } f \leftarrow t \cdot f$
 $\qquad \qquad \qquad x \leftarrow ta$
 $\qquad \qquad \qquad \text{return } (f x)$

15-324

$(\langle \$ \rangle) :: (a \rightarrow b) \rightarrow ta \rightarrow tb$ (map) Functor
 $(\langle * \rangle) :: t(a \rightarrow b) \rightarrow ta \rightarrow tb$ (apply) Applicative
 $(\gg=) :: ta \rightarrow (a \rightarrow tb) \rightarrow tb$ (bind) Monad

$\mathbb{Q}[x]$ is an ∞ dimensional vector space \mathbb{C}

$ta_1 \gg = \lambda x_1 \rightarrow (ta_2 \gg = \lambda x_2 \rightarrow ta_3 \dots)$

do $x_1 \leftarrow ta_1$

$x_2 \leftarrow ta_2$

:

$x_m \leftarrow ta_m$

\vdash

$\text{print} :: \text{Show } a \Rightarrow a \rightarrow \text{IO}()$

$\text{putStr} :: \text{String} \rightarrow \text{IO}()$

$\text{print } x = \text{putStr } (\text{show } x)$

$\text{print} = \text{putStr} . \text{show}$

class Show a where
 $\text{show} :: a \rightarrow \text{String}$

class Read a where
 $\text{read} :: \text{String} \rightarrow a$

getline :: IO String

get :: Read a \Rightarrow IO a

get = fmap read getline

get = read <\$> getline

point

do putStrLn "Give an integer".

inp \leftarrow get

print (inp + 1)

~~putStrLn "Give an integer"~~

get $\gg=$ ~~inp \rightarrow (print (inp + 1))~~

In sugr the above statement

action $\gg=$
 $\lambda x \rightarrow (x^2)$

putStrLn "Give an integer" \gg get $\gg=$ $\lambda inp \rightarrow (\text{print} (inp + 1))$

- Q) It will print prompt the user for an integer
and should print its square and the program
should be infinite.

Square :: IO()

Square = do putStrLn "Give an integer"

inp \leftarrow get

print (inp * inp)

Square.

inf_square = pushr "Enter n" >> (get >>= \ n -> (n >> inf_sq)

foo = putStrLn "Hello"

main = putStrLn "World"

In Haskell, the program has one main whose type is $\text{IO}()$. The entire computation is forced by eval of it.

foo = putStrLn "Hello"

main = putStrLn "World"

main = putStrLn "World"

foo = putStrLn "Hello"

main = do
 x ← getLine
 y ← getLine

// Takes two strings and
concatenate them.

putStrLn(x ++ y)

for = getLine

// It is also the same
as the first one.

main = do
 x ← getLine
 y ← for
 putStrLn(x ++ y)

for = getLine

// It is also the same
as the first one

main = do
 x ← for
 y ← for
 putStrLn(x ++ y)

as for is a plan
to generate a string
and the plan is used
twice in the plan of main.

foo = getLine

// Takes two strings
and concatenate

main = do
 x ← getLine
 y ← getLine
 putStrLn(x ++ y)

them
The first line is ignored
as it is never executed

getLine is a plan to generate a string using
some side effectful computation.

```
foo = do return "Hello"           -- The output will
          return "World"           -- be World in the
                                         terminal.
```

① main = do $x \leftarrow \text{foo}$

pet. Styln \approx

return "Hello" → return "World"

zeller "Hello" => _ → zeller "World"

main = do x < gelline

$y \leftarrow (+t) \langle \$ \rangle$ gelline

② main = do x ← (+) <\\$> getLine <*> getLine
return x.

- everything we can do with applicative can be done with a monad. Because monad happens to be a applicative.
 - always try to ~~try~~ do things generally whenever possible. the second definition of is more general - that it can used for both applicative and monad.

Take a number and check whether it is prime.

test function

instance Monad Maybe where

($\gg=$) $\therefore \text{Maybe} \rightarrow (\alpha \rightarrow \text{Maybe} \beta) \rightarrow \text{Maybe} \beta$

($\gg=$) $\text{Nothing} \gg= \text{Nothing}$

($\gg=$) ($\text{Just } a$) $f \gg= \text{Just}(f a)$

do
 $x_1 \leftarrow \text{foo}_1$,
 $x_2 \leftarrow \text{foo}_2$,
 :
 $x_m \leftarrow \text{foo}_m$
 something

$\text{Nothing} \gg= f = \text{Nothing}$

$\text{Just } x \gg= f = f x$

return = Just.

f -- here bracket is not

-- here bracket is not required around $\text{Just } x$

as ' $\gg=$ ' is an operator and Just is a function.

data Exp v = Var v

| Plus (Exp v) (Exp v)

| Mult (Exp v) (Exp v)

| Const Int

instance Functor Exp where

-- fmap :: ($v_1 \rightarrow v_2$) \rightarrow Exp $v_1 \rightarrow$ Exp v_2

fmap ut (Var x) = Var (ut x)

fmap ut (Plus e₁ e₂) = Plus (fmap ut e₁) (fmap ut e₂)

fmap ut (Mult e₁ e₂) = Mult (fmap ut e₁) (fmap ut e₂)

Entire point of abstraction is selective amnesia

- Piguska

do x \leftarrow e₁

e₁ :: $(\lambda x \rightarrow e_2)$

e₂

join :: Monad m \Rightarrow m(m a) \rightarrow m a

join mma = do ma \leftarrow mma | mma :: $\lambda x \rightarrow x$

No Monad can be defined in terms of other ~~named~~ ~~integers~~
of "join" or "bind".

($\gg=$) :: Monad $m \Rightarrow ma \rightarrow (a \rightarrow mb) \rightarrow mb$

($\gg=$) $mf = \text{join}(fma + ma)$

join :: $\text{Exp}(\text{Exp } u) \rightarrow \text{Exp } u$

join (var v) = v

join (Plus e₁ e₂) = Plus (join e₁) (join e₂)

join (Mul e₁ e₂) = Mul (join e₁) (join e₂)

~~join~~ (Con)

join x = x

join (Const x) = Const x

3.2x Syntactic Monad

data Exp v = Var v

| const Int

| Plus (Exp v) (Exp v)

what is the meaning of $\gg=$ in Exp (Exp substitution?)

List Monad

do $x \leftarrow \text{alist}$
~~for~~ x

$\text{alist} \text{ alist } \gg= \text{foo}$

a list :: [a], foo :: a → [b]

alist $\gg=$ foo = concat \$

alist $\gg=$ foo = concat \$ map foo alist

Distr a : generate a ~~some value from~~ value from

given a : ^{random}generate an element of type a.

Let

The idea of list monad is sampling.

gI : Gen Int

getI :: Gen (Int, Int)

n = gI

do $x \leftarrow gI$

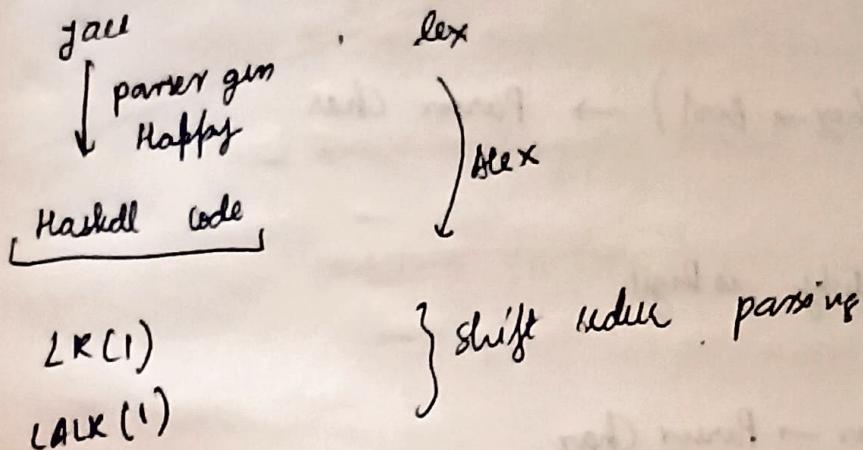
$y \leftarrow gI$

return (n, y)

- Read the Gen monad

Parser Monad

Build a library for constructing Parser.



LL - Parsing

Recursive descent parsing

(-- Ocaml - module LR)

newtype Parser a = Parser (string → a)

Parser Combinator libraries

Parsec, megaParsec, altoParsec

Parser should handle these

- (1) "Rest of the input"
- (2) "Parse error".

$\text{String} \rightarrow \text{Maybe}(a, \text{String})$

new type $\text{Parser } a = \text{Parser } (\text{String} \rightarrow \text{Maybe}(a, \text{String}))$

$\text{satisfy} : (\text{Char} \rightarrow \text{Bool}) \rightarrow \text{Parser Char}$

$\text{digit} = \text{satisfy } \text{isDigit}$

$\text{char} : \text{Char} \rightarrow \text{Parser Char}$

$\text{char } x = \text{satisfy } (\lambda c \rightarrow c == x)$

$\text{char } x = \text{satisfy } (c == x)$

$\text{satisfy } pr = \text{Parser } fn$

where

$fn : \text{String} \rightarrow \text{Maybe}(\text{char}, \text{String})$

$fn (x : xs) = \begin{cases} \text{Just}(x, xs) & \text{if } pr x \\ \text{Nothing} & \text{else} \end{cases}$

$fn (\text{nil}) = \text{Nothing}$

$(\langle 1 \rangle) : \text{Parser } a \rightarrow \text{Parser } a \rightarrow \text{Parser } a$

$p_1 \sqcap p_2 =$

type Result a = Maybe(a, String)

newtype Parser a = Parser(string → Result a)

runParser : Parser a → string → Result a

runParser (Parser fn) = fn

$p_1 \sqcup p_2 = \text{Parser } fn$

where fn input = case runParser p_1 input of
Nothing → runParser p_2 input
 $n \rightarrow n$

many : Parser a → Parser (a)

many p = Parser fn

where fn input =

newtype Parser a = Parser(string → Result a)

// this can be done
if the input doesn't
match 'a' then return
empty else try
matching the remaining
this can be done
differently.

instance Functor Parser where

fmap :: $(a \rightarrow b) \rightarrow \text{Parser } a \rightarrow \text{Parser } b$

fmap f pa = Parser fn

where fn :: String \rightarrow Result b

fn input = runParser

fn input = fmap f (runParser pa input)

data Result a = Ok a | Err

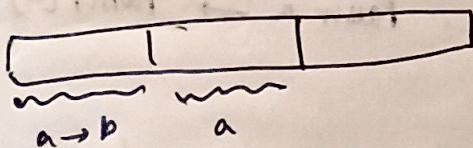
Instance Functor Result where

fmap f (Ok a) = Ok (f a)

fmap f Err = Err

Hw
Define the applicative instance
of

\leftrightarrow Parser $(a \rightarrow b) \rightarrow \text{Parser } a$
 $\rightarrow \text{Parser } b$



Tryout megaparsec

a Parser
library to
write small
parsers

Parser as Monads

data Result a = Ok a | String
| Err

newtype Parser a = Parser (String → Result a)

~~runParser~~:
runParser :: Parser a → (String → Result a)

runParser (Parser fn) = fn

instance Functor Result where

-- fmap :: (a → b) → Result a → Result b

fmap f (Ok x str) = Ok (f x) str

fmap f Err = Err

instance Functor Parser

-- fmap :: f (a → b) → Parser a → Parser b

fmap f pa = Parser (pfn :: String → Result b)

where pfn str = fmap f (runParser pa str)

instance Applicative Parser where

-- pure :: $a \rightarrow \text{Parser } a$

pure a = Parser (pf : String \rightarrow Result a)

where pf str = OK a str

($\langle *\rangle$) :: Parser($a \rightarrow b$) \rightarrow Parser(a) \rightarrow Parser(b)

pf $\langle *\rangle$ pa = Parser (pb : String \rightarrow Result b)

where

pb str = case runParser pf str of
OK f rest \rightarrow f rest
Err \rightarrow Err

pf $\langle *\rangle$ pa = Parser (pb : String \rightarrow Result b)

where

pb str = case runParser pf str of

OK f rest \rightarrow case runParser pa rest of

Good \rightarrow rest \rightarrow OK a rest \rightarrow OK (f a) rest

Err \rightarrow Err

Err \rightarrow Err

$p1 \leftrightarrow p \circ = \text{Parser } (pb :: \text{String} \rightarrow \text{Result } b)$

where

$pb \text{ str} = \text{case runParser pf str of}$
 $\text{Ok } f \text{ rest} \rightarrow \text{fmap } f (\text{runParser pa rest})$
 $\text{Err} \rightarrow \text{Err}$

$\text{digit} = \text{satisfy isDigit}$

$\text{alpha} = \text{satisfy isAlpha}$

$\text{many} :: \text{Parser a} \rightarrow \text{Parser [a]}$

$\text{many1} :: \text{Parser a} \rightarrow \text{Parser [a]}$

data Person String Int

$\text{name} = \text{many1 alpha}$

$\text{age} = \text{read } \langle \# \rangle (\text{many1 digit})$

Person \$> name \$> age

Define many in terms of applicative interface.

$\text{many1 } p = \text{map } (\lambda) \langle \$ \rangle p \langle * \rangle \text{ many } p$

$\text{many } p = \text{many1 } p \langle ! \rangle \text{ pure } []$