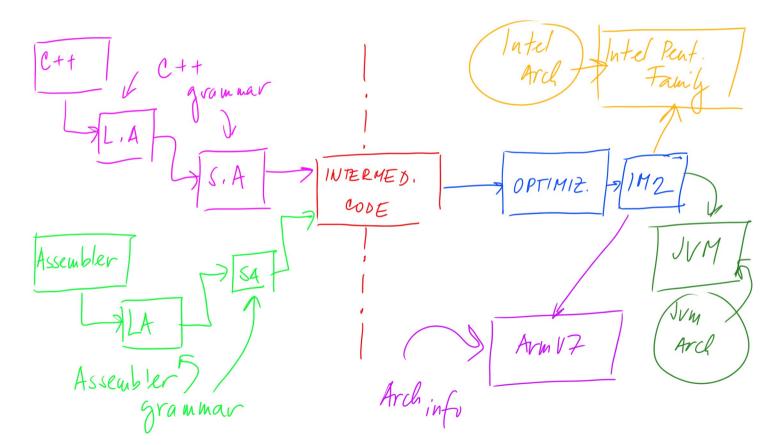
NTI/PRK Intermediate languages and gcc as a use-case

Lenka Kosková Třísková, <u>lenka.koskova.triskova@tul.cz</u>, LS 2024

Intermediate code



Type validation

- Attribute grammar with type identification
 - A special form of a grammar with type rules
- In the derivation tree, each node has associated type

https://www.cs.csub.edu/~melissa/cs350-f15/notes/notes04.html

Structured type equivalence

Naming is different.

Control flow diagramm

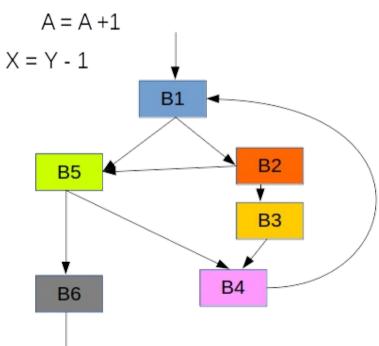
A basic block sequence of consecutive statements, without branching or loops.

A program control flow graph is an oriented graph (B, H, s), where: B is the set of blocks, H is the set of edges, and s is a mapping $BxB \rightarrow H$ such that: $s(h)=\{Bi,Bj\}$ just if Bi is a potential successor of Bj in terms of program control flow.

Příklad

Originální kód:

while (A < B) and (C < D) or (X = Y):



Tříadresní kód

1: if (A >B) goto (4)

2: if (C >D) goto (4)

3: goto 5

4: if X<>Y goto 8

5: T1 = A + 1

6: A = T1

7: goto 1

8: T2 = Y - 1

9: X = T2

Intermediate code - design goal

Preservation of structure (validity ranges, flow control, expressions, etc.) and semantic meaning.

Shading of the original syntax - conversion to a "universal language". Syntax does not have to be "human readable".

Preparation for generating machine code or other output.

There may even be several different ones.

Approach:

Tree structure vs. Three-address code

Tree structure

The information is stored in a tree structure.

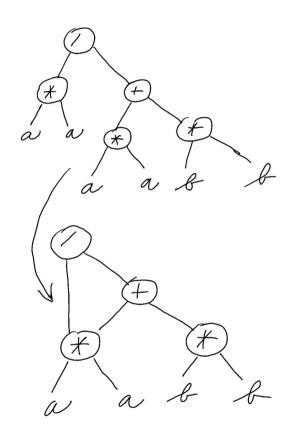
It is often used for code optimization.

Example: AST: abstract syntax tree - described standard for GCC.

Tree edits

axa/(axa+bxb)

Subtree replacement



Three Address Code (Linear Representation)

The code is very close to assembly - it can be viewed as virtual machine code.

Mathematical operations rewritten for registers => expressions only combinations of two operators.

Always longer, always more variables.

Original:

$$x = y*3 + z;$$
 $t1 = y * 3;$
 $t2 = t1 + z;$

Static single assignment - helpful for control flow analysis

Orignal:

$$x = a + b$$

$$y = x + 1$$

$$x = b + c$$

$$z = x + y$$

SSA edit:

$$x1 = a + b$$

$$y1 = x1 + 1$$

$$x2 = b + c$$

$$z = x2 + y1$$

Internal gcc structure

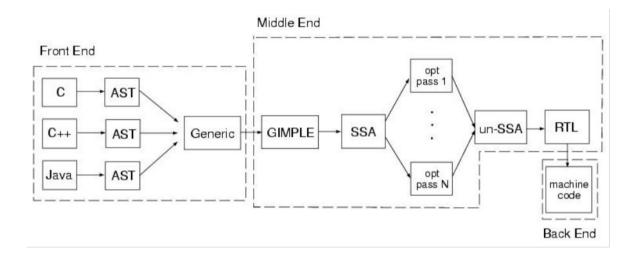
AST – abstract syntax tree

GENERIC – syntaktický strom

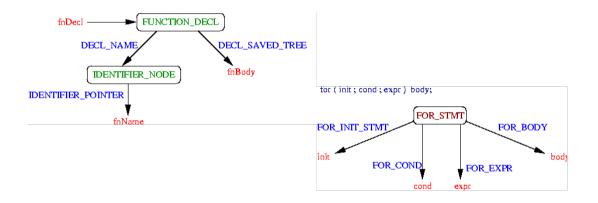
GIMPLE – zjednodušená podmnožina GENERIC užitá k optimalizaci

SSA – static simple assignment

RTL – register transfer language

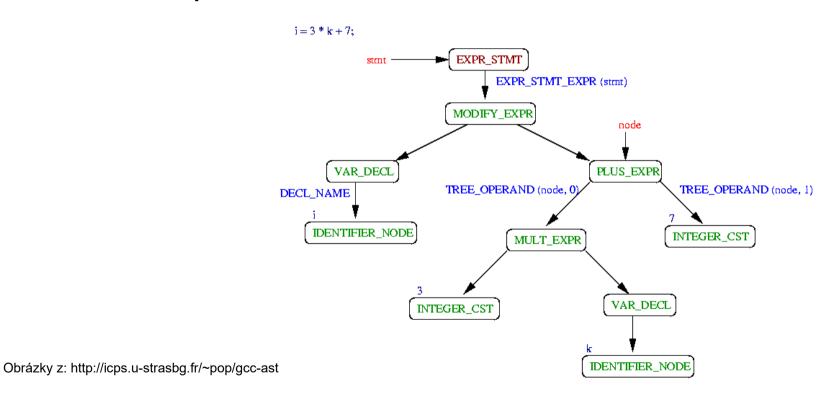


AST - Example





AST – Examples



14

Gimple - Example

```
struct A { A(); ~A(); };
   int i;
   int g();
   void f()
   {
      A a;
      int j = (--i, i ? 0 : 1);
      for (int x = 42; x > 0; --x)
        {
         i += g()*4 + 32;
      }
}
```

```
void f()
       int i.0;
       int T.1;
       int iftmp.2;
       int T.3;
       int T.4;
       int T.5;
       int T.6;
         struct A a;
         int j;
         __comp_ctor (&a);
         try
             i.0 = i;
             T.1 = i.0 - 1;
             i = T.1;
             i.0 = i;
```

RTL

```
#include <stdio.h>
int
main (void)
{
   printf ("Hello, world!\n");
   return 0;
}
```

```
(note 1 0 3 NOTE INSN DELETED)
(note 3 1 18 2 [bb 2] NOTE INSN BASIC BLOCK)
(insn/f 18 3 19 2 (set (mem:DI (pre dec:DI (reg/f:DI 7
sp)) [0 S8 A8])
       (reg/f:DI 6 bp)) hello.c:5 65 {*pushdi2 rex64}
     (nil))
(insn/f 19 18 20 2 (set (reg/f:DI 6 bp)
       (reg/f:DI 7 sp)) hello.c:5 89 {*movdi internal}
     (nil))
(note 20 19 2 2 NOTE INSN PROLOGUE END)
(note 2 20 5 2 NOTE INSN FUNCTION BEG)
(insn 5 2 6 2 (set (reg:DI 5 di)
       (symbol ref/f:DI ("*.LCO") [flags 0x2]
<var decl 0x7fc588b8db40 *.LC0>)) hello.c:6 89
{*movdi internal}
     (nil))
(call insn 6 5 7 2 (set (reg:SI 0 ax)
       (call (mem:QI (symbol ref:DI ("puts") [flags
0x41] <function decl 0x7fc588ae2360 builtin puts>)
[0 builtin puts S1 A8])
           (const int 0 [0]))) hello.c:6 669
{*call value}
```