

Lecture 8: Revisiting the quantum Fourier transform

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Recap

The previous lecture discussed an algorithm to extract the phase factor $w \in [0, 1]$ from a generic n -qubit quantum state. Writing w as $\frac{x}{2^n}$, for x an integer representable in n qubits, the estimation process was described by

$$\frac{1}{\sqrt{2^n}} \sum_{y \in 2^n} e^{2\pi i (\frac{x}{2^n}) y} |y\rangle \leadsto |\textcolor{red}{x}\rangle$$

Its inverse is [QFT](#), the quantum Fourier transform, a most useful routine in Quantum Computation.

Let us revisit its construction in a systematic way.

Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
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The quantum Fourier transform

Essentially, the QFT performs a **change-of-basis** operation which encodes information of computational basis states in **local phases**.

For 1 qubit state this is exactly what the Hadamard gate accomplishes:

$$H|0\rangle = \frac{1}{\sqrt{2}}(|0\rangle + 1|1\rangle) \quad H|1\rangle = \frac{1}{\sqrt{2}}(|0\rangle + (-1)|1\rangle)$$

Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
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QFT: 1 qubit

Thus, $QFT_1 = H$:

$$QFT_1 |0\rangle = \frac{1}{\sqrt{2}}(|0\rangle + 1|1\rangle) \quad QFT_1 |1\rangle = \frac{1}{\sqrt{2}}(|0\rangle + (-1)|1\rangle)$$

Operation H^{-1} allows to extract information encoded in local phases

$$\downarrow \\ = H$$

Exercise

Let $\omega_1 = e^{i2\pi\frac{1}{2}}$. Show that $QFT_1 |x\rangle = \frac{1}{\sqrt{2}}(|0\rangle + \omega_1^{1 \cdot x}|1\rangle)$

$$\downarrow \\ \text{angle of } \pi \text{ radians}$$

QFT: 1 qubit

The complex ω_1 represents a rotation of π radians, dividing the **unit circle** into two slices.

Actually, the two corresponding points in the circle correspond to the **2^{th} -roots of the identity**

$$\omega_1^0 = 1 \quad \text{and} \quad \omega_1^1 = e^{\frac{i2\pi}{2}} = e^{i\pi} = -1$$

Note that

$$\omega_1^{1x} = e^{\frac{i2\pi x}{2}} = e^{i2\pi \frac{x}{2}} = e^{i2\pi(0.x)}$$

as used in the previous lecture.

Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
ooo

QFT: 2 qubits

Let $\omega_2 = e^{i2\pi \frac{1}{4}}$

$$QFT_2 |00\rangle = \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{2 \cdot 0}|1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{1 \cdot 0}|1\rangle)$$

$$QFT_2 |01\rangle = \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{2 \cdot 1}|1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{1 \cdot 1}|1\rangle)$$

$$QFT_2 |10\rangle = \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{2 \cdot 2}|1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{1 \cdot 2}|1\rangle)$$

$$QFT_2 |11\rangle = \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{2 \cdot 3}|1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{1 \cdot 3}|1\rangle)$$

In general

$$QFT_2 |\textcolor{red}{x}\rangle = \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{2 \cdot \textcolor{red}{x}}|1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{1 \cdot \textcolor{red}{x}}|1\rangle)$$

Exercise

Show that, for $x = |x_1 x_2\rangle$, $QFT_2 |x\rangle$ can be written as

$$QFT_2 |x\rangle = \frac{1}{\sqrt{2}}(|0\rangle + e^{i2\pi(0 \cdot x_2)}|1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle + e^{i2\pi(0 \cdot x_1 x_2)}|1\rangle)$$

Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
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QFT: 2 qubits

The basic observation is that for every ω_2 -rotation on the second qubit there are *two* such rotations on the first qubit.

Exercise

Compute the phase coefficients in the expressions above and use the Bloch sphere to study $QFT_2 |x\rangle$.

Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
ooo

QFT: 2 qubits

Hint

$$\begin{array}{ll} \omega_2^{2,0} = 1 & \omega_2^{1,0} = 1 \\ \omega_2^{2,1} = -1 & \omega_2^{1,1} = e^{i\frac{\pi}{2}} \\ \omega_2^{2,2} = 1 & \omega_2^{1,2} = -1 \\ \omega_2^{2,3} = -1 & \omega_2^{1,3} = e^{i\frac{3}{2}\pi} \end{array}$$

Note that the information on $|x\rangle$ previously encoded by vectors pointing to the poles becomes encoded by vectors in the **xz-plane**

Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
ooo

QFT: 2 qubits

In order to derive a circuit for QFT_2 , compute

$$\begin{aligned} QFT_2 |x\rangle &= \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{2 \cdot x} |1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{1 \cdot x} |1\rangle) \\ &= \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{2(2x_1+x_2)} |1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{2x_1+x_2} |1\rangle) \\ &= \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{4x_1+2x_2} |1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{2x_1+x_2} |1\rangle) \\ &= \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{4x_1} \omega_2^{2x_2} |1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{2x_1} \omega_2^{x_2} |1\rangle) \\ &= \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{2x_2} |1\rangle) \otimes \frac{1}{\sqrt{2}}(|0\rangle + \omega_2^{2x_1} \omega_2^{x_2} |1\rangle) \\ &= \underbrace{\frac{1}{\sqrt{2}}(|0\rangle + (-1)^{x_2} |1\rangle)}_{H|x2\rangle} \otimes \underbrace{\frac{1}{\sqrt{2}}(|0\rangle + (-1)^{x_1} \omega_2^{x_2} |1\rangle)}_{\text{some controlled rot. on } H|x1\rangle} \end{aligned}$$

Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
ooo

QFT: 2 qubits

Define

$$R_2 |0\rangle = |0\rangle \quad \text{and} \quad R_2 |1\rangle = \omega_2 |1\rangle$$

which rotates a vector in the xz -plane $\frac{\pi}{2}$ radians

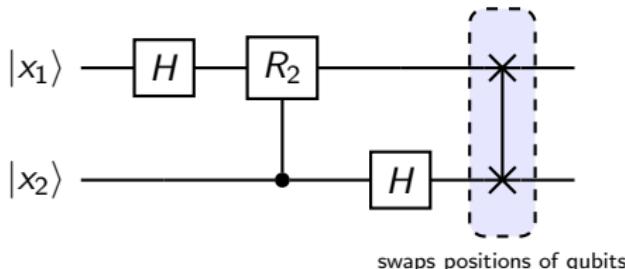
It yields a **controlled- R_2** operation

$$|x\rangle |0\rangle \mapsto |x\rangle |0\rangle \quad |x\rangle |1\rangle \mapsto R_2 |x\rangle |1\rangle$$

or, equivalently,

$$|0\rangle |y\rangle \mapsto |0\rangle |y\rangle \quad |1\rangle |y\rangle \mapsto \omega_2^y |1\rangle |y\rangle$$

Putting all the pieces together to derive the QFT circuit for 2 qubits:



Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
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QFT: 3 qubits

$$QFT_3 |x\rangle = \frac{1}{\sqrt{2}} (|0\rangle + \omega_3^{4 \cdot x} |1\rangle) \otimes (|0\rangle + \omega_3^{2 \cdot x} |1\rangle) \otimes (|0\rangle + \omega_3^{1 \cdot x} |1\rangle)$$

for $\omega_n = e^{i2\pi \cdot \frac{1}{2^n}}$.

N.B.

In the sequel the normalisation factor $\frac{1}{\sqrt{2}}$ will be dropped in each state to increase readability

Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
ooo

QFT: 3 qubits

Recalling that a binary number $x_1 \dots x_n$ represents the natural number

$$2^{n-1} \cdot x_1 + \dots + 2^0 \cdot x_n$$

and that

$$\omega_n^2 = \omega_{n-1} \text{ and } \omega_n^{2^{n-1}} = e^{i\pi} = -1$$

define QFT_3 as follows:

Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
ooo

QFT: 3 Qubits

$QFT_3 |x\rangle$

$$\begin{aligned} &= (|0\rangle + \omega_3^{4 \cdot x} |1\rangle) \otimes (|0\rangle + \omega_3^{2 \cdot x} |1\rangle) \otimes (|0\rangle + \omega_3^{1 \cdot x} |1\rangle) \\ &= (|0\rangle + (-1)^x |1\rangle) \otimes (|0\rangle + \omega_3^{2 \cdot x} |1\rangle) \otimes (|0\rangle + \omega_3^{1 \cdot x} |1\rangle) \\ &= (|0\rangle + (-1)^{x_3} |1\rangle) \otimes (|0\rangle + \omega_3^{2 \cdot x} |1\rangle) \otimes (|0\rangle + \omega_3^{1 \cdot x} |1\rangle) \\ &= H|x_3\rangle \otimes (|0\rangle + \omega_3^{2 \cdot (4x_1+2x_2+x_3)} |1\rangle) \otimes (|0\rangle + \omega_3^{1 \cdot x} |1\rangle) \\ &= H|x_3\rangle \otimes (|0\rangle + \omega_3^{2 \cdot (4x_1+2x_2)} \omega_3^{2 \cdot x_3} |1\rangle) \otimes (|0\rangle + \omega_3^{1 \cdot x} |1\rangle) \\ &= H|x_3\rangle \otimes (|0\rangle + \omega_2^{2 \cdot (2x_1+x_2)} \omega_2^{x_3} |1\rangle) \otimes (|0\rangle + \omega_3^{4x_1+2x_2+x_3} |1\rangle) \\ &= H|x_3\rangle \otimes (|0\rangle + \omega_2^{2 \cdot (2x_1+x_2)} \omega_2^{x_3} |1\rangle) \otimes (|0\rangle + \omega_3^{4x_1+2x_2} \omega_3^{x_3} |1\rangle) \\ &= H|x_3\rangle \otimes (|0\rangle + \omega_2^{2 \cdot (2x_1+x_2)} \omega_2^{x_3} |1\rangle) \otimes (|0\rangle + \omega_3^{2 \cdot (2x_1+x_2)} \omega_3^{x_3} |1\rangle) \\ &= H|x_3\rangle \otimes \underbrace{\left(|0\rangle + \omega_2^{2 \cdot (2x_1+x_2)} \omega_2^{x_3} |1\rangle \right)}_{\text{some controlled-rotations on } QFT_2|x_1x_2\rangle} \otimes (|0\rangle + \omega_2^{2x_1+x_2} \omega_3^{x_3} |1\rangle) \end{aligned}$$

Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
ooo

QFT: 3 qubits

Take $R_3 |0\rangle = |0\rangle$ and $R_3 |1\rangle = \omega_3 |1\rangle$. Intuitively, R_3 rotates a vector in the xz -plane ‘one 2^3 -th of the unit circle’.

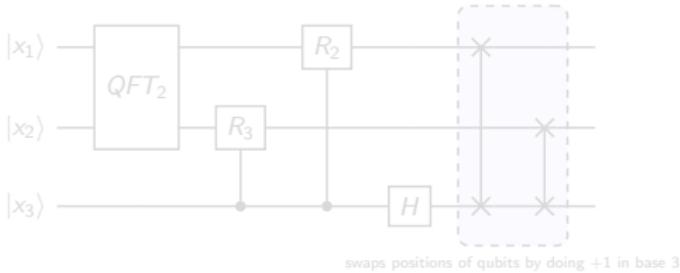
It yields a **controlled- R_3** operation defined by

$$|x\rangle |0\rangle \mapsto |x\rangle |0\rangle \text{ and } |x\rangle |1\rangle \mapsto R_3 |x\rangle |1\rangle$$

Equivalently

$$|0\rangle |y\rangle \mapsto |0\rangle |y\rangle \text{ and } |1\rangle |y\rangle \mapsto \omega_3^y |1\rangle |y\rangle$$

Putting all pieces together we derive the QFT circuit for 3 qubits



Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
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QFT: 3 qubits

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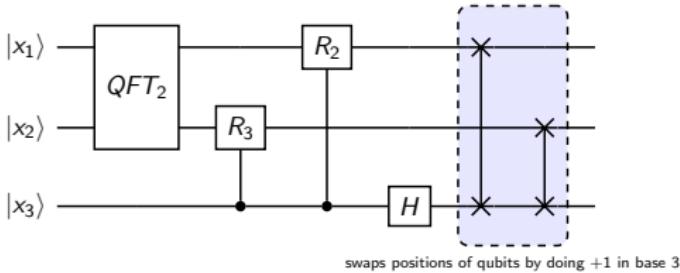
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Equivalently

$$|0\rangle |y\rangle \mapsto |0\rangle |y\rangle \text{ and } |1\rangle |y\rangle \mapsto \omega_3^y |1\rangle |y\rangle$$

Putting all pieces together we derive the QFT circuit for 3 qubits



Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
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QFT: n qubits

Calculation easily extends to QFT_n (*in lieu of QFT_3*) :

Let $\omega_n = e^{i2\pi \cdot \frac{1}{2^n}}$ (division of the **unit circle** in 2^n slices)

$$QFT_n |\textcolor{red}{x}\rangle = \frac{1}{\sqrt{2}} (|0\rangle + \omega_n^{2^{n-1} \cdot \textcolor{red}{x}} |1\rangle) \otimes \cdots \otimes (|0\rangle + \omega_n^{2^0 \cdot \textcolor{red}{x}} |1\rangle)$$

Take $R_n |0\rangle = |0\rangle$ and $R_n |1\rangle = \omega_n |1\rangle$. Intuitively, R_n rotates a vector in the xz -plane '**one 2^n -th of the unit circle**'

It yields a **controlled- R_n** operation defined by $|x\rangle |0\rangle \mapsto |x\rangle |0\rangle$ and $|x\rangle |1\rangle \mapsto R_n |x\rangle |1\rangle$. Equivalently

$$|0\rangle |y\rangle \mapsto |0\rangle |y\rangle \text{ and } |1\rangle |y\rangle \mapsto \omega_n^y |1\rangle |y\rangle$$

Recalling the basic idea
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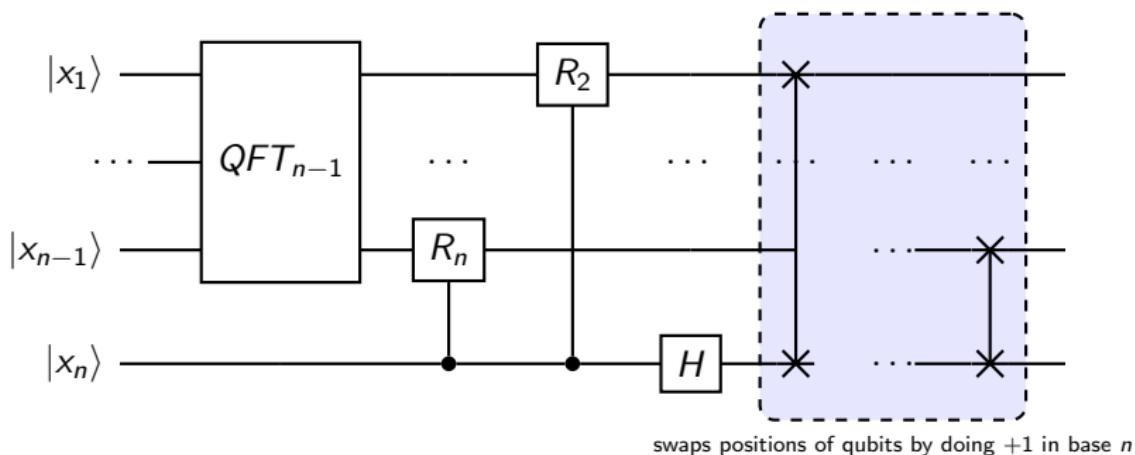
QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
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QFT: n qubits

This suggests a recursive definition for the general QFT circuit:



Recalling the basic idea
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QFT on 2 qubits
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QFT on 3 qubits
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QFT: The general case
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An equivalent formulation of QFT

Although we have been working with

$$QFT_n |x\rangle = \frac{1}{\sqrt{2}}(|0\rangle + \omega_n^{2^{n-1} \cdot x} |1\rangle) \otimes \cdots \otimes \frac{1}{\sqrt{2}}(|0\rangle + \omega_n^{1 \cdot x} |1\rangle)$$

we are already familiar with an equivalent, useful definition

$$QFT_n |x\rangle = \frac{1}{\sqrt{2^n}} \sum_{y=0}^{2^n - 1} \omega_n^{x \cdot y} |y\rangle$$

Examples with $n = 1$ and $n = 2$

$$QFT_1 |x\rangle = \frac{1}{\sqrt{2}}(|0\rangle + \omega_1^x |1\rangle)$$

$$QFT_2 |x\rangle = \frac{1}{\sqrt{2^2}}(|00\rangle + \omega_2^x |01\rangle + \omega_2^{2 \cdot x} |10\rangle + \omega_2^{3 \cdot x} |11\rangle)$$