

Machine-Level Programming: Loops, Switch statements

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Loops

- C provides several looping constructs - namely, *do-while*, *while*, and *for*
- No corresponding instructions exist in machine code
 - Instead, combinations of conditional tests and jumps are used to implement the effect of loops
- We will study the translation of loops as a progression, starting with *do-while* and then working toward ones with more complex implementations
 - Most compilers generate loop code based on the *do-while* form of a loop

“Do-While” loop example

■ Count number of 1's in argument **x**

C code

```
int pcount(unsigned int x)
{
    int result = 0;
    do{
        result += x & 0x1;
        x >>= 1;
    }while(x);
    return result;
}
```

Goto version

```
int pcount(unsigned int x)
{
    int result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if(x)
        goto loop;
    return result;
}
```

■ Use conditional branch to either continue or exit loop

“Do-While” loop in Assembly

■ Assume:

%edx	x
%ecx	result

```

    movl    $0,%ecx        # result = 0
.L2:                                # loop:
    movl    %edx,%eax
    andl    $1,%eax        # t = x & 1
    addl    %eax,%ecx      # result += t
    shrl    %edx           # x >>= 1
    jne     .L2            # if !0, goto loop

```

Goto version

```

int pcount(unsigned int x)
{
    int result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if(x)
        goto loop;
    return result;
}

```

General “Do-While” translation

C code

```
do{  
    statement1  
    ...  
    statementn  
}while(Test);
```

} Body



Goto version

```
loop:  
    statement1  
    ...  
    statementn  
    if(Test)  
        goto loop
```

} Body

■ Test returns integer

- = 0 interpreted as false
- ≠ 0 interpreted as true

“While” loop example

C code

```
int pcount(unsigned int x){  
    long result = 0;  
    while(x){  
        result += x & 0x1;  
        x >>= 1;  
    }  
    return result;  
}
```

- It differs from *do-while* in that *test-expr* is first evaluated
 - The loop is potentially terminated before the first execution of *body-statement*
- There are a number of ways to translate a *while* loop into machine code

General “While” translation #1

■ “Jump-to-middle” translation

- Used with `-Og`

While version

```
while (Test) {  
    Body  
}
```



Goto version

```
goto test;  
loop:  
    Body  
test:  
    if (Test)  
        goto loop;  
done:
```

- Avoids duplicating test code
- Unconditional jump incurs no performance penalty on modern CPUs
 - It occupies a decode unit but never makes it into the main pipeline

While loop – Jump to middle translation

C code

```
int pcount(unsigned int x){
    long result = 0;
    while(x) {
        result += x & 0x1;
        x >>= 1;
    }
    return result;
}
```

Jump to middle version

```
int pcount(unsigned int x){
    long result = 0;
    goto test;
loop:
    result += x & 0x1;
    x >>= 1;
test:
    if(x)
        goto loop;
    return result;
}
```


General “While” translation #2

While version

```
while (Test) {  
    Body  
}
```



Do-While version

```
if (!Test)  
    goto done;  
do {  
    Body  
} while (Test);  
done:
```



Goto version

```
if (!Test)  
    goto done;  
loop:  
    Body  
    if (Test)  
        goto loop;  
done:
```

■ “Do-while” conversion

- Used with -O1

“While” loop – do while translation

C code

```
int pcount(unsigned int x)
{
    int result = 0;

    while(x) {
        result += x & 0x1;
        x >>= 1;
    }

    return result;
}
```

Goto Version

```
int pcount(unsigned int x)
{
    int result = 0;
    if (!x) goto done;
loop:
    result += x & 0x1;
    x >>= 1;
    if (x)
        goto loop;
done:
    return result;
}
```

■ The compiler can often optimize the initial test

- For example, determining that the test condition will always hold

General “For” loop form

General form

```
for (Init; Test; Update) {  
    Body  
}
```

Example

```
for(i = 0; i < WSIZE; i++) {  
    unsigned mask = 1 << i;  
    result += (x & mask) != 0;  
}
```

Init

```
i = 0
```

Test

```
i < WSIZE
```

Update

```
i++
```

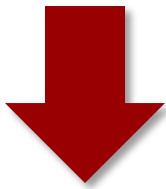
Body

```
{  
    unsigned mask = 1 << i;  
    result += (x & mask) != 0;  
}
```

“For” loop → ... → Goto

For version

```
for(Init; Test; Update) {  
  Body  
}
```



While version

```
Init;  
while(Test) {  
  Body  
  Update;  
}
```



Do-While version

```
Init;  
if(!Test)  
  goto done;  
do {  
  Body  
  Update  
} while(Test);  
done:
```



Goto Version

```
Init;  
if(!Test)  
  goto done;  
loop:  
  Body  
  Update  
  if(Test)  
    goto loop;  
done:
```

“For” loop conversion example

C code

```
#define WSIZE 8*sizeof(int)

int pcount(unsigned int x)
{
    int i;
    int result = 0;

    for(i = 0; i < WSIZE; i++){
        unsigned mask = 1 << i;
        result += (x & mask) != 0;
    }

    return result;
}
```

Goto version

```
int pcount(unsigned int x){
    int i;
    int result = 0;
    i = 0;
    if (!(i < WSIZE))
        goto done;
loop:
    {
        unsigned mask = 1 << i;
        result += (x & mask) != 0;
    }
    i++;
    if (i < WSIZE)
        goto loop;
done:
    return result;
}
```

Diagram illustrating the Goto version flow:

- Init**: `i = 0;`
- !Test**: `if (!(i < WSIZE))`
- Body**: `loop:` block containing the loop body.
- Update**: `i++;`
- Test**: `if (i < WSIZE)`

The *loop* instructions

- Use the ECX register as a counter and automatically decrease its value as the loop instruction is executed
 - Without affecting the EFLAGS register flag bits when ECX reaches zero
- Support only an 8-bit offset, so only short jumps can be performed

loopX	Condition	Description
<code>loop</code>	<code>ECX != 0</code>	Loop until the ECX register is zero
<code>loope/loopz</code>	<code>ECX != 0 or ZF</code>	Loop until either the ECX register is zero, or the ZF flag is not set
<code>loopne/loopnz</code>	<code>ECX != 0 and ~ZF</code>	Loop until either the ECX register is zero, or the ZF flag is set

The *loop* instructions example

C code

```
for (i = 100; i > 0; i--)  
{  
    ...  
}
```

Assembly *loop* version

```
movl $100,%ecx  
for_loop:  
    ...  
    loop for_loop
```

■ Be careful with code inside the loop

- If the ECX register is modified, it will affect the operation of the loop
- Function calls within the loop can easily trash the value of the ECX register without you knowing it
- If ECX is already ≤ 0 before the loop, it will eventually exit when the register overflows

Today

- Loops
- **Switch statements**

Switch statements

- **Provide a multi-way branching capability based on the value of an integer index**
 - Particularly useful when dealing with tests where there can be a large number of possible outcome
- **Large blocks are implemented using a *jump table***
 - An array where entry i is the address of a code segment implementing the action the program should take when the switch index equals i
- **The time taken to perform the switch is independent of the number of switch cases**
 - As opposed to a long sequence of if-else statements

Example

```
int switch_eg(int x, int y, int z){  
    int w = 1;  
    switch(x) {  
        case 1:  
            w = y*z;  
            break;  
        case 2:  
            w = y/z;  
            /* Fall through */  
        case 3:  
            w += z;  
            break;  
        case 5:  
        case 6:  
            w -= z;  
            break;  
        default:  
            w = 2;  
    }  
    return w;  
}
```

■ Multiple case labels

- case 5 & 6

■ Fall through cases

- case 2

■ Missing cases

- case 4

Jump table structure

Switch form

```
switch(x) {
  case val_0:
    Block 0
  case val_1:
    Block 1
    . . .
  case val_n-1:
    Block n-1
}
```

Approximate translation

```
target = JTab[x]
goto *target;
```

Jump table

JTab:	Targ0
	Targ1
	Targ2
	•
	•
	•
	Targn-1

Jump targets

Targ0:

Code Block
0

Targ1:

Code Block
1

Targ2:

Code Block
2

•
•
•

Targn-1:

Code Block
n-1

Switch statement example

```
int switch_eg(int x, int y, int z){
    int w = 1;
    switch(x) {
        . . .
    }
    return w;
}
```

Setup:

```
switch_eg:
    pushl %ebp
    movl  %esp, %ebp
    movl  8(%ebp), %eax    # %eax = x
    cmpl  $6, %eax        # Compare x:6
    ja    .L2             # if > goto default
    jmp   *.L7(, %eax, 4)   # goto *JTab[x]
```

What range of values
takes *default*?

Switch statement example

```
int switch_eg(int x, int y, int z){
    int w = 1;
    switch(x) {
        . . .
    }
    return w;
}
```

Setup:

```
switch_eg:
    pushl %ebp
    movl %esp, %ebp
    movl 8(%ebp), %eax    # %eax = x
    cmpl $6, %eax        # Compare x:6
    ja .L2               # if > goto default
    jmp *.L7(, %eax, 4)    # goto *JTab[x]
```

Jump table

```
.section .rodata
    .align 4
.L7:
    .int .L2 # x = 0
    .int .L3 # x = 1
    .int .L4 # x = 2
    .int .L5 # x = 3
    .int .L2 # x = 4
    .int .L6 # x = 5
    .int .L6 # x = 6
```

← Indirect jump

Setup explanation

■ Table structure

- Each target requires 4 bytes
- Base address at `.L7`

■ Jumping

- **Direct:** `jmp .L2`
- Jump target is denoted by label `.L2`
- **Indirect:** `jmp *.L7(,%eax,4)`
- Start of jump table: `.L7`
- Must scale by factor of 4 (addresses are 4 bytes on IA32)
- Fetch target from effective address `.L7 + %eax*4`
 - Only for $0 \leq x \leq 6$

Jump table

```
.section .rodata
.align 4
.L7:
.int     .L2  # x = 0
.int     .L3  # x = 1
.int     .L4  # x = 2
.int     .L5  # x = 3
.int     .L2  # x = 4
.int     .L6  # x = 5
.int     .L6  # x = 6
```

Jump table

```
.section .rodata
.align 4
.L7:
.int    .L2 # x = 0
.int    .L3 # x = 1
.int    .L4 # x = 2
.int    .L5 # x = 3
.int    .L2 # x = 4
.int    .L6 # x = 5
.int    .L6 # x = 6
```

```
switch(x) {
case 1:      /* .L3 */
    w = y*z;
    break;
case 2:      /* .L4 */
    w = y/z;
    /* Fall Through */
case 3:      /* .L5 */
    w += z;
    break;
case 5:
case 6:      /* .L6 */
    w -= z;
    break;
default:    /* .L2 */
    w = 2;
}
```

- Duplicates have same label
- Missing cases use label for the default case

Code blocks (x == 1, default)

```

switch(x) {
case 1:      /* .L3 */
    w = y*z;
    break;

    ...
default:    /* .L2 */
    w = 2;
}

```

```

.L3:
    movl    16(%ebp), %eax    # z
    imull   12(%ebp), %eax    # w = y*z
    jmp     .L8              # Goto done

.L2:
    movl    $2, %eax         # w = 2
    jmp     .L8              # Goto done

```

- **Jump table avoids sequencing through cases**
 - Constant time, rather than linear

Code blocks (x == 2, x == 3)

```

int w = 1;
. . .
switch(x) {
. . .
case 2:          /* .L4 */
    w = y/z;
    /* Fall Through */

case 3:          /* .L5 */
    w += z;
    break;
. . .
}

```

```

.L4:
    movl    12(%ebp), %edx    # y
    movl    %edx, %eax
    sarl    $31, %edx
    idivl   16(%ebp)          # w = y/z
    jmp     .L9

.L5:
    movl    $1, %eax         # w = 1

.L9:
    addl    16(%ebp), %eax    # w += z
    jmp     .L8              # goto done

```

- Do not initialize `w = 1` unless really need it
- Use program sequencing to handle fall-through

Handling fall-through

```
int w = 1;  
.  
.  
.  
switch(x) {  
.  
.  
.  
case 2:  
    w = y/z;  
    /* Fall through */  
case 3:  
    w += z;  
    break;  
.  
.  
.  
}
```

case 2:
 w = y/z;
 goto merge;

case 3:
 w = 1;

merge:
 w += z;

Code blocks ($x == 5$, $x == 6$)

```

switch(x) {
    . . .
    case 5:      /* .L6 */
    case 6:      /* .L6 */
        w -= z;
        break;
    . . .
}

return w;

```

```

.L6:
    movl    $1,%eax           # w = 1
    subl    16(%ebp),%eax     # w -= z

.L8:                          # done
    movl    %ebp,%esp
    popl    %ebp
    ret

```

- Use jump table to handle holes and duplicate tags

Summary

■ C control

- do-while
- while, for
- switch

■ Assembler control

- Conditional and unconditional jumps
- Indirect jump (via jump tables)

■ Standard techniques

- Loops converted to do-while form
- Large switch statements use jump tables