

$\text{len}^{\text{arr}} = \frac{1}{2} \times 5$

$$\begin{aligned} \text{inc}[i] \\ = \text{dec}[j] + 1 \\ \xleftarrow{\quad \text{dec}[i] \\ = \text{inc}[j] + 1 \quad} \end{aligned}$$

$I \rightarrow \{1, 17\}$   
 $D \rightarrow \{17\}$

$I \rightarrow \{1, 17, 5, 10\}$   
 $D \rightarrow \{1, 17, 10\}$

$I \rightarrow \{1, 17, 5, 15\}$   
 $D \rightarrow \{17, 15\}$

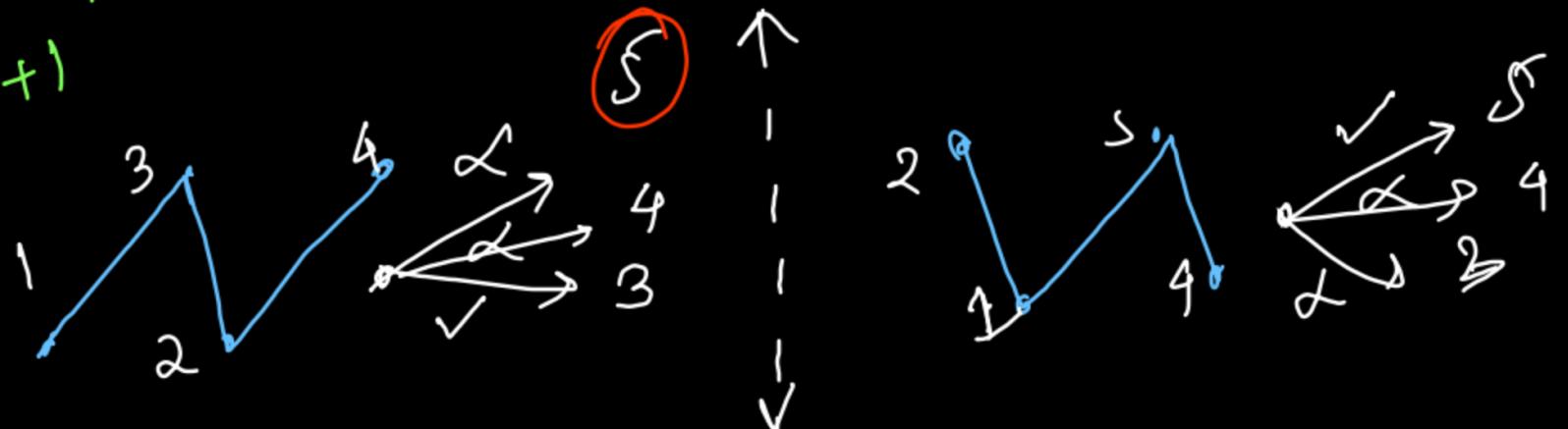
$[1, 17, 5, 10, 13, 15, 10, 5, 16, 8]$

$I \rightarrow \{1\}$   
 $D \rightarrow \{1\}$

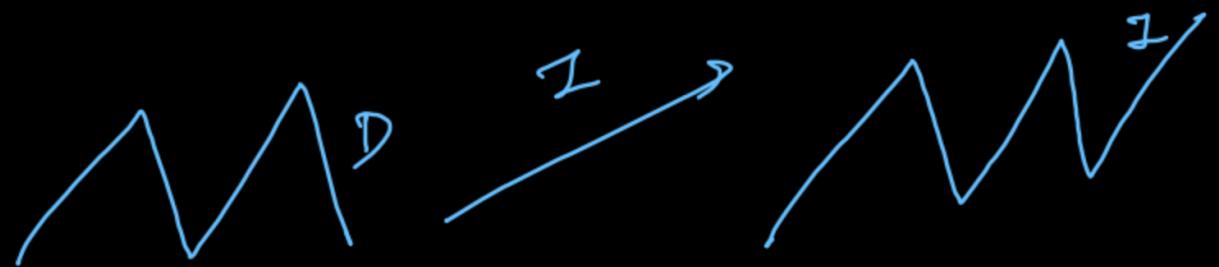
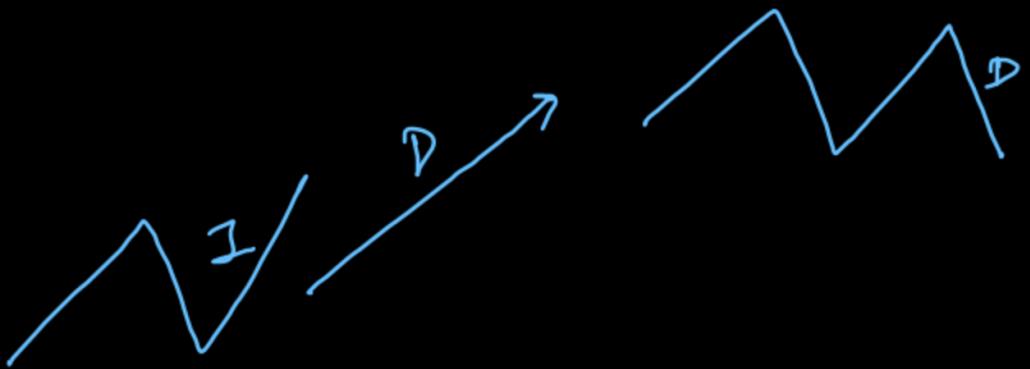
$I \rightarrow \{1, 5\}$   
 $D \rightarrow \{1, 17, 5\}$

$I \rightarrow \{1, 17, 5, 13\}$   
 $D \rightarrow \{1, 17, 13\}$

$I \rightarrow \{1, 17, 5, 10\}$   
 $D \rightarrow \{1, 17, 5, 10\}$



$$\text{dec}(i) = \text{inc}(j) + 1$$



$$\begin{aligned}\text{inc}(i) \\ = \text{dec}(j) + 1\end{aligned}$$

```

public int wiggleMaxLength(int[] nums) {
    int[] inc = new int[nums.length];
    Arrays.fill(inc, 1);

    int[] dec = new int[nums.length];
    Arrays.fill(dec, 1);

    int maxLength = 0;
    for(int i=0; i<nums.length; i++){

        for(int j=0; j<i; j++){
            if(nums[i] > nums[j]){
                // We are Increasing
                inc[i] = Math.max(inc[i], dec[j] + 1);
            } else if(nums[i] < nums[j]){
                // We are Decreasing
                dec[i] = Math.max(dec[i], inc[j] + 1);
            }
        }

        maxLength = Math.max(maxLength, Math.max(inc[i], dec[i]));
    }

    return maxLength;
}

```

Time  $\rightarrow O(2 \cdot n^2)$   
 $= O(n^2)$

Space  $\rightarrow O(2 \cdot n)$

2 rows } 1D DP

## LIS on 2D Coordinate

Overlapping Boxes  
Practice - GFG

Russian Doll

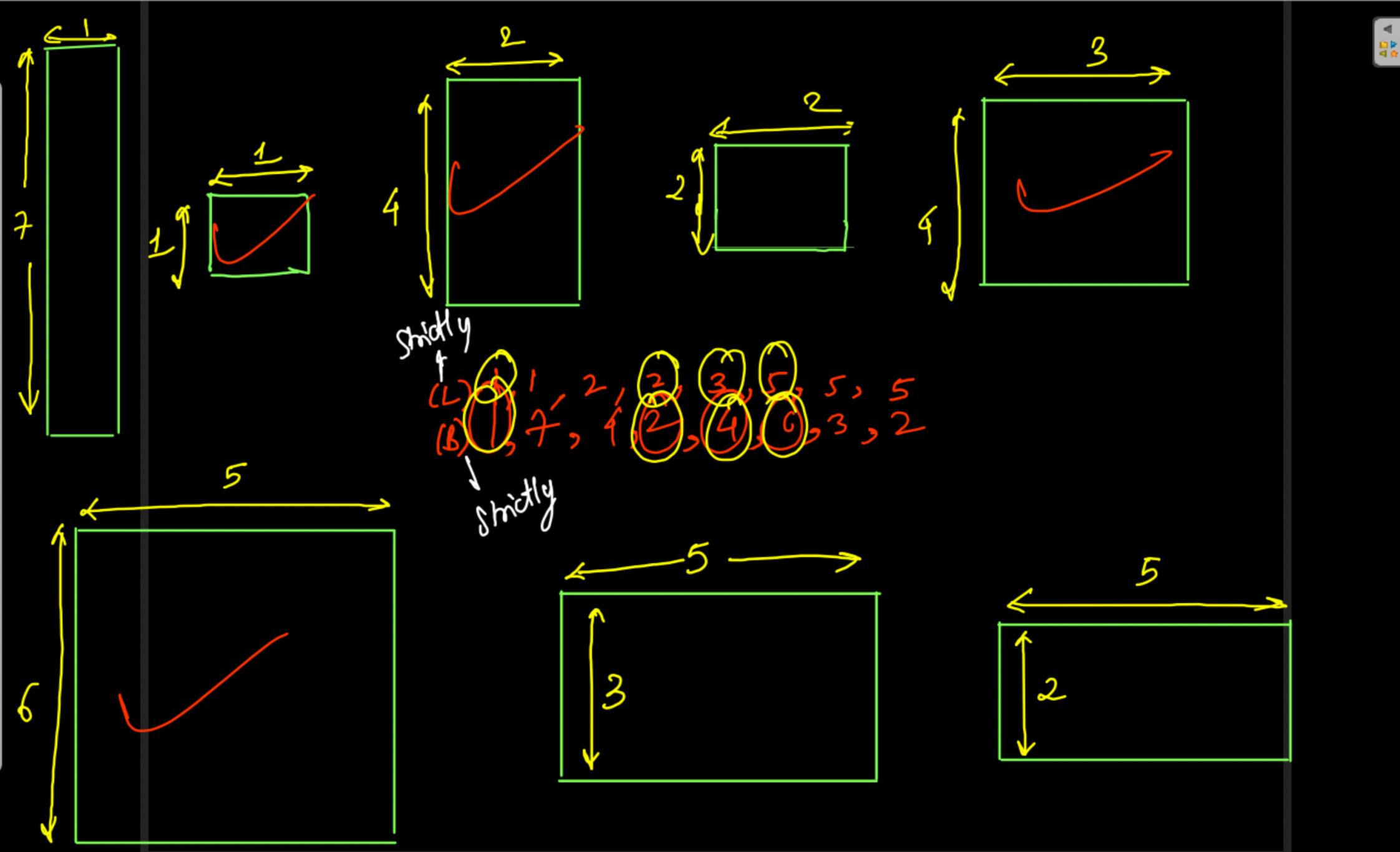


Bon Stacking

GFG  
IV-1

Leetcode  
II

One envelope can fit into another if and only if both the width and height of one envelope are greater than the other envelope's width and height.

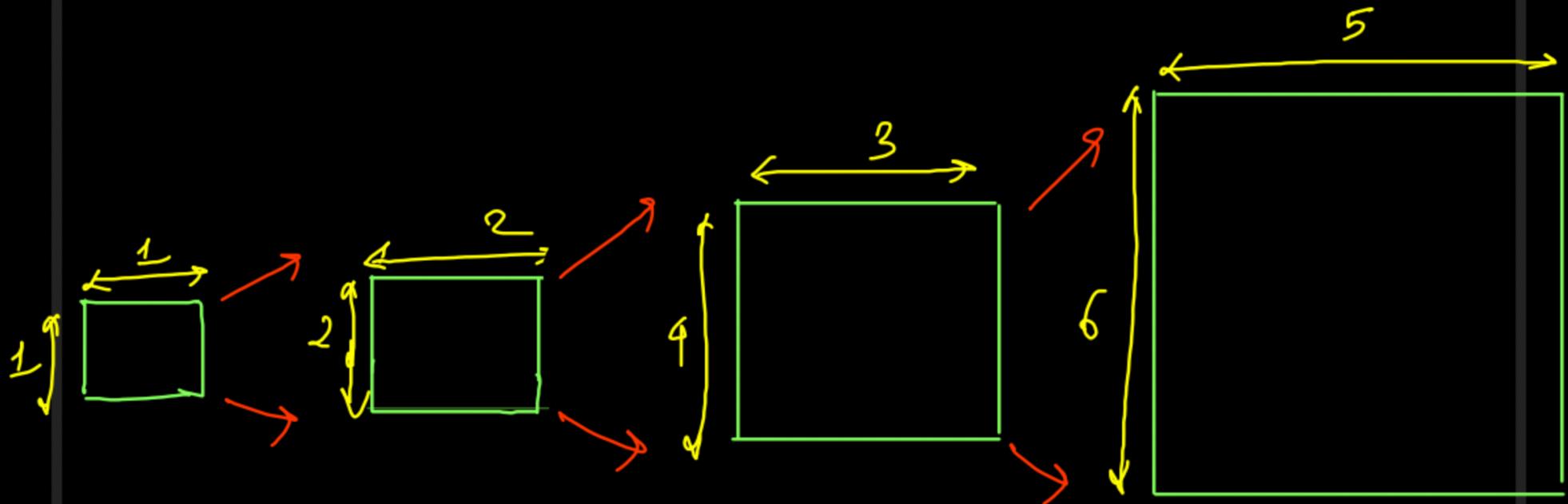


Time:  $O(N \log N + N \log N)$   
↓  
Breadth  
length

Space:  $O(N)$

```
public int lowerBound(ArrayList<Integer> nums, int target){  
    int low = 0, high = nums.size() - 1;  
    int idx = nums.size();  
  
    while(low <= high){  
        int mid = low + (high - low) / 2;  
  
        if(nums.get(mid) < target){  
            low = mid + 1;  
        } else {  
            high = mid - 1;  
            idx = mid;  
        }  
    }  
  
    return idx;  
}
```

```
public int maxEnvelopes(int[][] envelopes) {  
    Arrays.sort(envelopes, (int[] first, int[] second) -> {  
        return (first[0] != second[0]) ? first[0] - second[0]  
            : second[1] - first[1];  
    });  
  
    int n = envelopes.length;  
    ArrayList<Integer> sorted = new ArrayList<>();  
  
    for(int i=0; i<envelopes.length; i++){  
        int lb = lowerBound(sorted, envelopes[i][1]);  
        if(lb == sorted.size()){  
            sorted.add(envelopes[i][1]);  
            // Current Element larger than the largest  
            // LIS of one length more  
        } else {  
            sorted.set(lb, envelopes[i][1]);  
        }  
    }  
  
    return sorted.size(); // This Sorted Array has same size as LIS  
}
```



# LIS on length

# LIS on breadth

1, 2, 3, 5

1, 2, 4, 6

2D LIS

1<sup>st</sup> coordinate → Sorting  
2<sup>nd</sup> coordinate → LIS

# longest Common Subsequence

Lecture-1 Saturday (28 May)

9 AM to 12 PM

→ length of LCS

→ Print LCS

Any one

All

→ longest Repeated subset (LRS)

→ Palindromic  
Subsequence

length

Count

Highest Common Subsequence

is

order of characters

e.g.

"a<sup>.</sup>b<sup>.</sup>c<sup>.</sup>d<sup>.</sup>e" , "i<sup>.</sup>d<sup>.</sup>c<sup>.</sup>a"

Brute force

$O(2^N \times 2^M)$

Compare all  
Subsequences

$S[i:j] = S[j:j]$   
 $LCS(i, j)$   
 ↓  
 $\text{char}$   
 $a$   
 $\uparrow$   
 $\downarrow$   
 $LCS(i+1, j+1)$

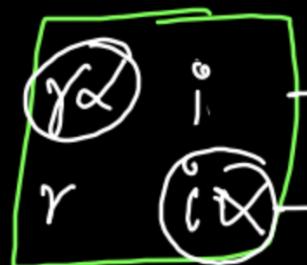
char taken (1)

	$a\nu$	$a\nu$	$\rightarrow n$
1	$a\nu$	$a\nu$	length l
2	$a\nu$	$a\nu$	
3	$a\nu$	$a\nu$	

$S_1[i:j] = S_2[i:j]$  

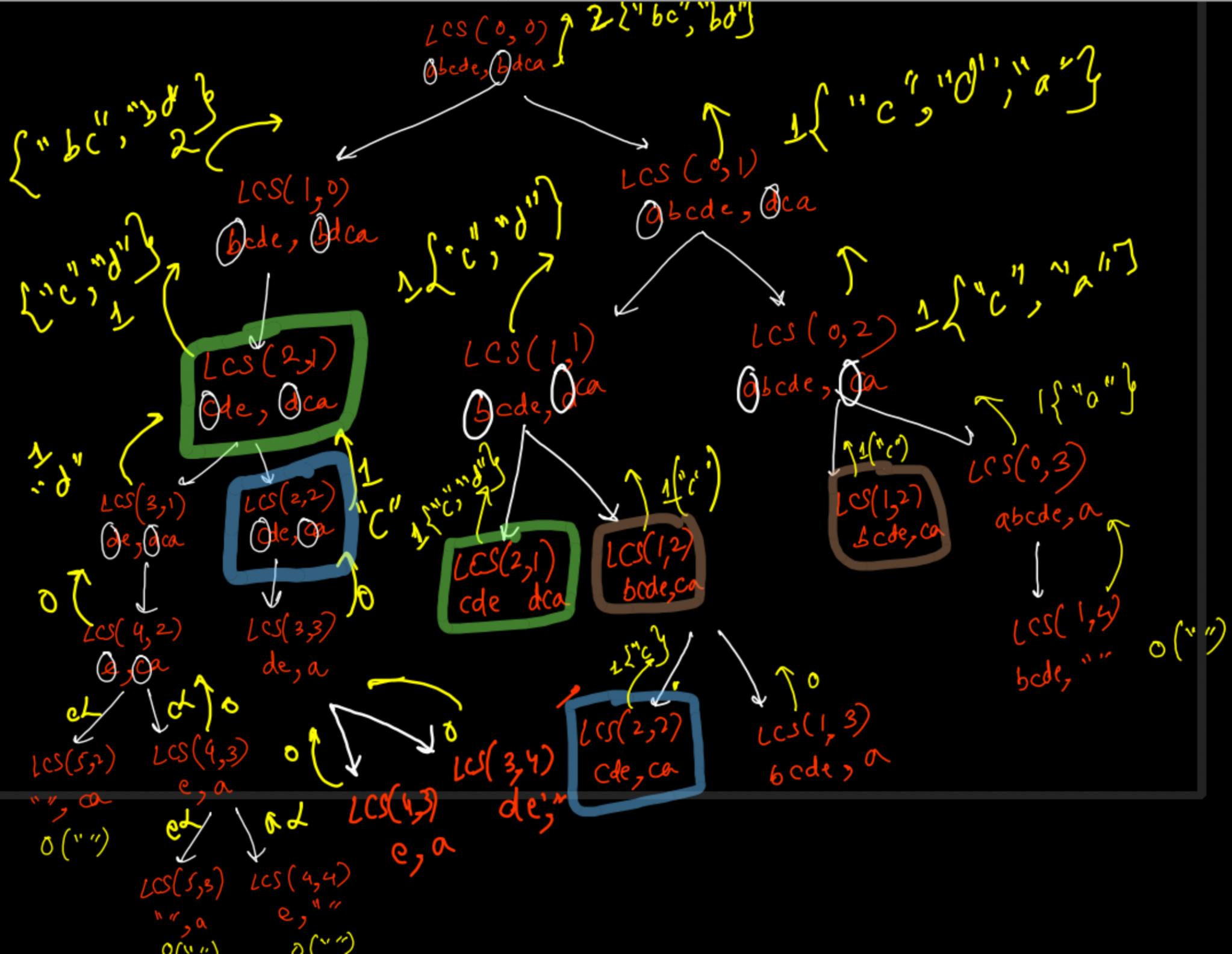
LCS(i, j) 

char not taken =  $\delta + LCS(i+1, j)$   $\delta + LCS(i, j+1)$




$r \vee i \vee c \rightarrow$  Common  $\wedge$   
 $r \wedge i \wedge c \rightarrow$  Length  $\downarrow$

$LCS(i, j)$   
 $\delta + LCS(i+1, j+1)$



```

public int LCS(int i, int j, String s1, String s2, int[][] dp){
    if(i == s1.length() || j == s2.length())
        return 0; // LCS of Empty String with Other String is Empty String only

    if(dp[i][j] != -1) return dp[i][j];

    char ch1 = s1.charAt(i);
    char ch2 = s2.charAt(j);

    if(ch1 == ch2) // If characters are same, take the common from both of them
        return dp[i][j] = 1 + LCS(i + 1, j + 1, s1, s2, dp);

    // If character is uncommon, either not take s1[i] or not take s2[j]
    int option1 = LCS(i + 1, j, s1, s2, dp);
    int option2 = LCS(i, j + 1, s1, s2, dp);
    return dp[i][j] = 0 + Math.max(option1, option2);
}

```

```

public int longestCommonSubsequence(String s1, String s2) {
    int[][] dp = new int[s1.length() + 1][s2.length() + 1];
    for(int i=0; i<dp.length; i++){
        for(int j=0; j<dp[0].length; j++){
            dp[i][j] = -1;
        }
    }

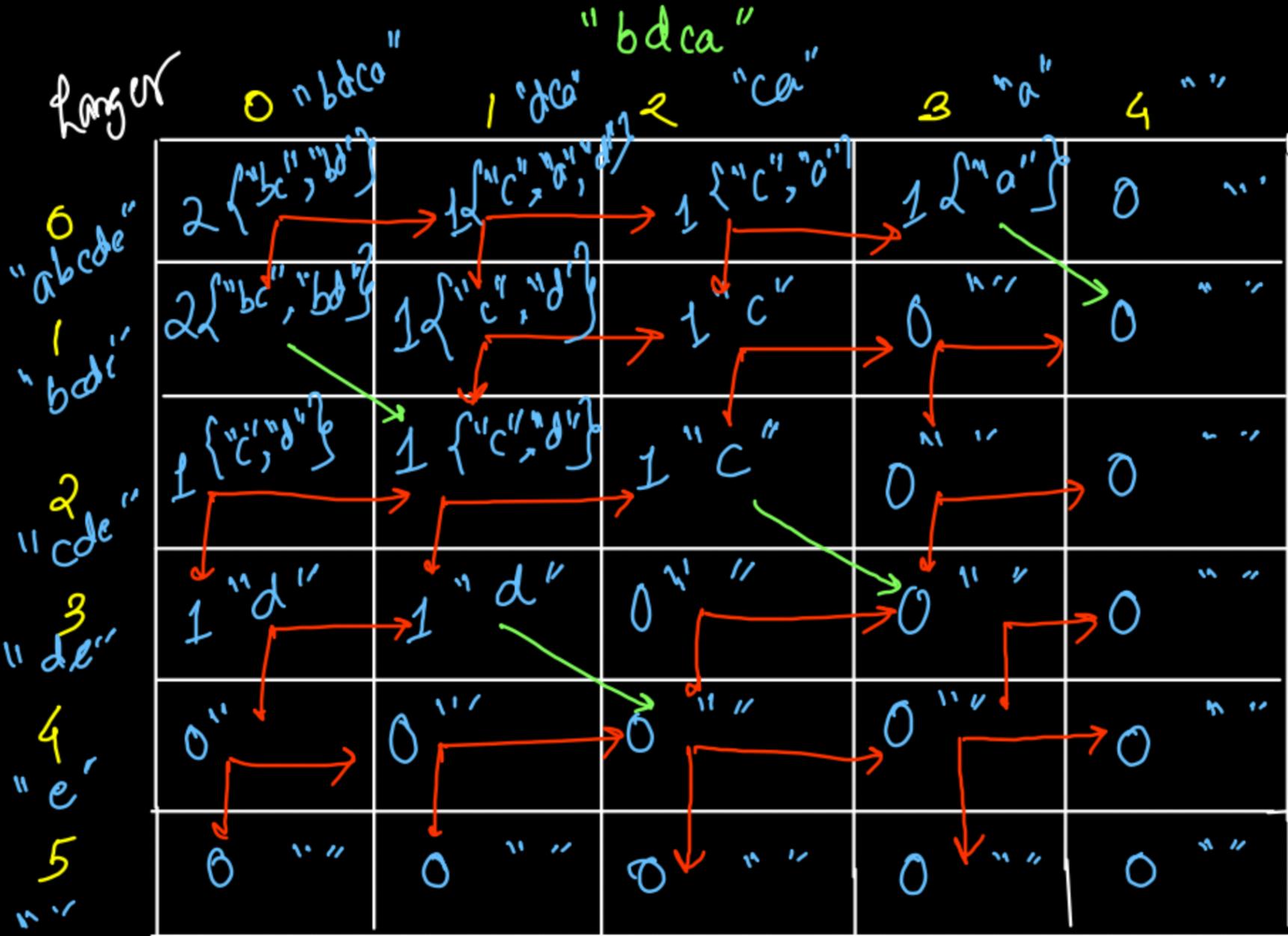
    return LCS(0, 0, s1, s2, dp);
}

```

Time  $\rightarrow O(N \times M)$

Space  $\rightarrow 2D DP$   
 $O(N \times M)$

eg "abcde", "bdca"



Smallest

```

public int longestCommonSubsequence(String s1, String s2) {
    int[][] dp = new int[s1.length() + 1][s2.length() + 1];

    for(int i=s1.length()-1; i>=0; i--){
        for(int j=s2.length()-1; j>=0; j--){
            char ch1 = s1.charAt(i);
            char ch2 = s2.charAt(j);

            if(ch1 == ch2)
                dp[i][j] = 1 + dp[i + 1][j + 1];
            else dp[i][j] = Math.max(dp[i + 1][j], dp[i][j + 1]);
        }
    }

    return dp[0][0];
}

```

Time  $\rightarrow O(N \times M)$

Space  $\rightarrow O(N \times M)$

$\boxed{2D DP}$

```

int[] next = new int[s2.length() + 1];

for(int i=s1.length()-1; i>=0; i--){
    int[] curr = new int[s2.length() + 1];

    for(int j=s2.length()-1; j>=0; j--){
        char ch1 = s1.charAt(i);
        char ch2 = s2.charAt(j);

        if(ch1 == ch2)
            curr[j] = 1 + next[j + 1];
        else curr[j] = Math.max(next[j], curr[j + 1]);
    }

    next = curr;
}

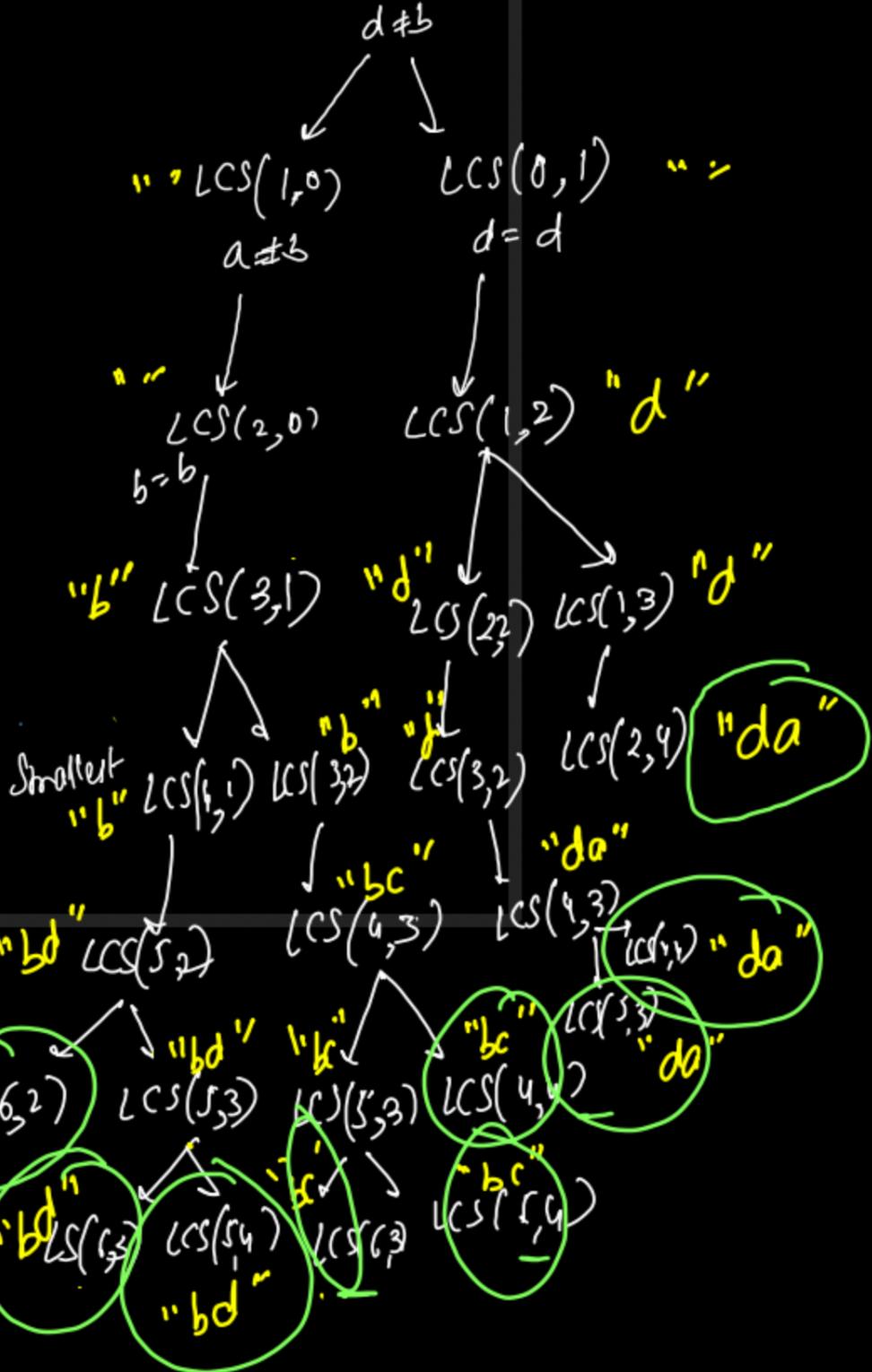
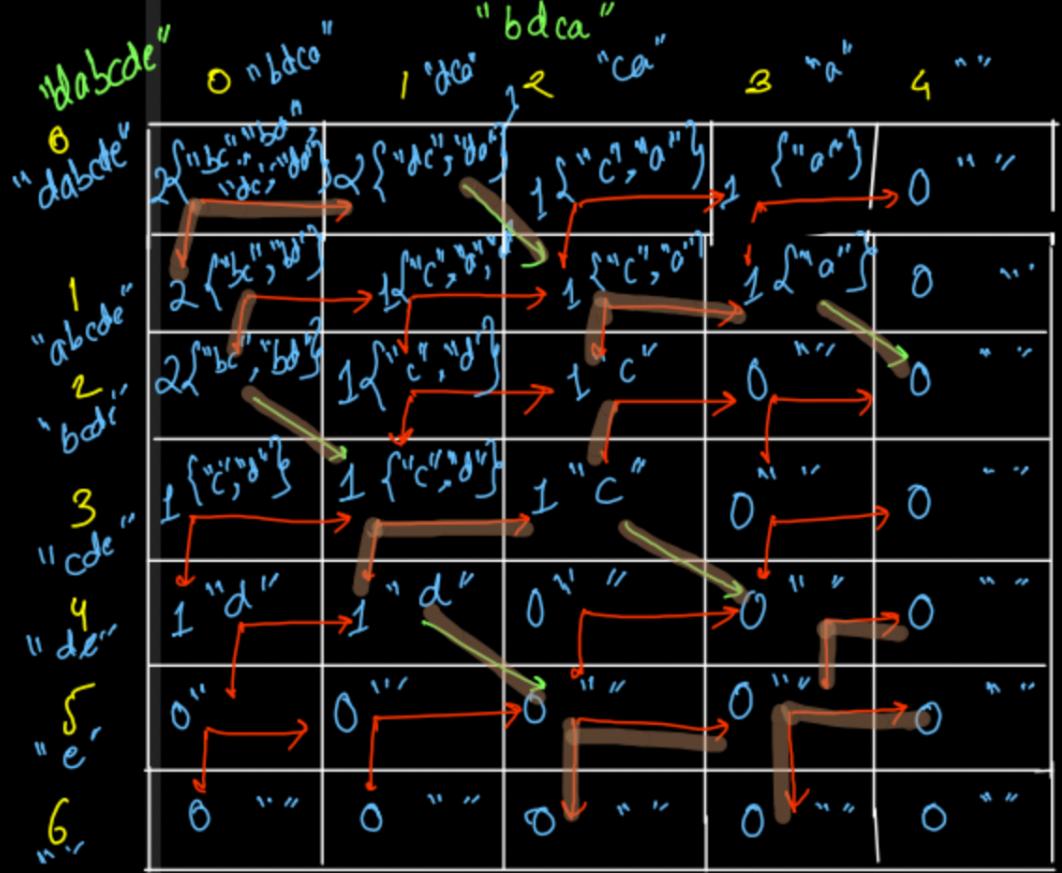
return next[0];

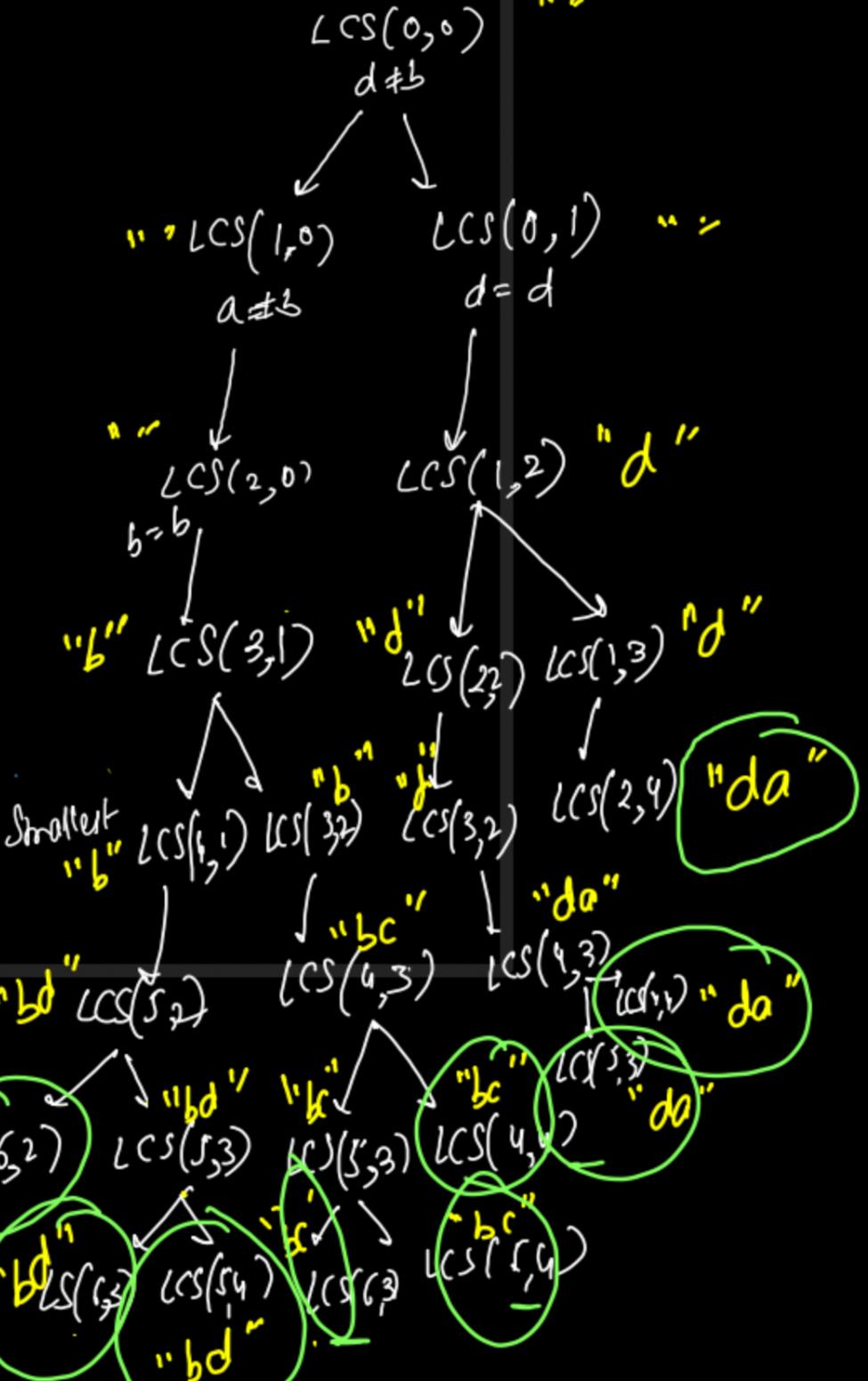
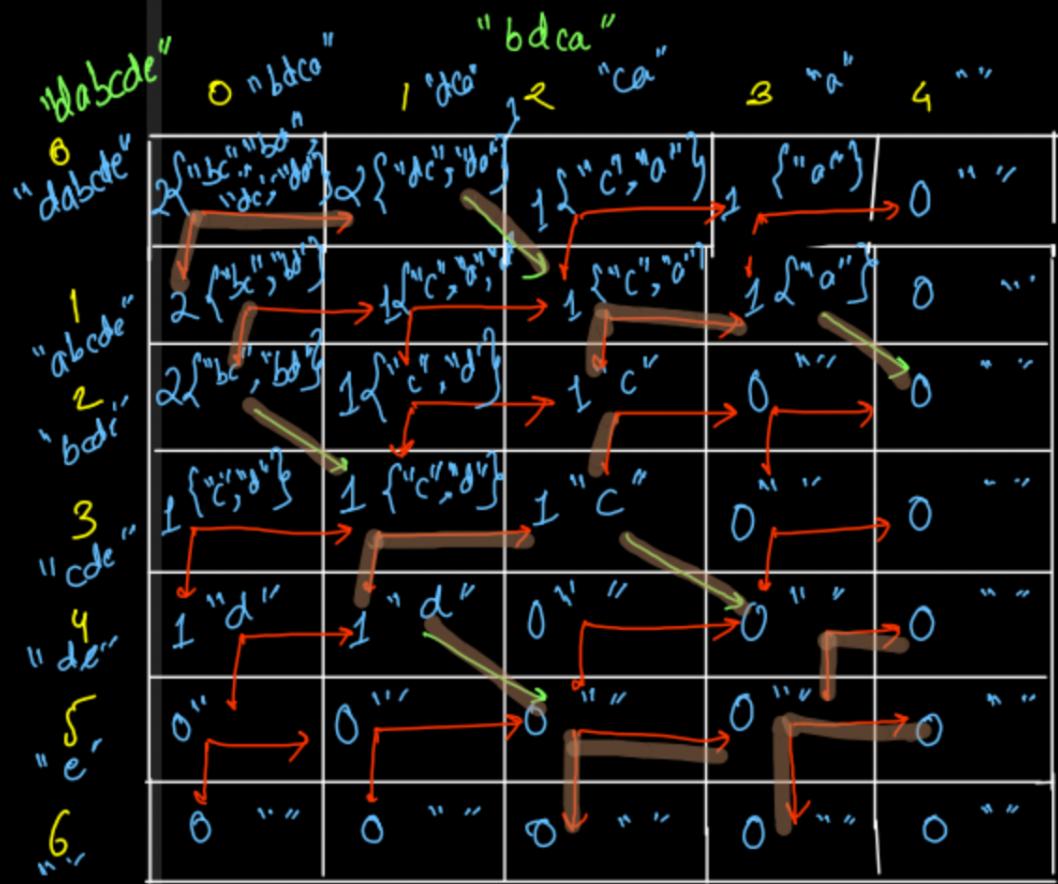
```

Time  $\rightarrow O(N \times M)$

Space  $\rightarrow O(2 \times M)$

$\approx \boxed{1D DP}$





```

TreeSet<String> answers; // Both Ordered (Lexicographical Order), Unique

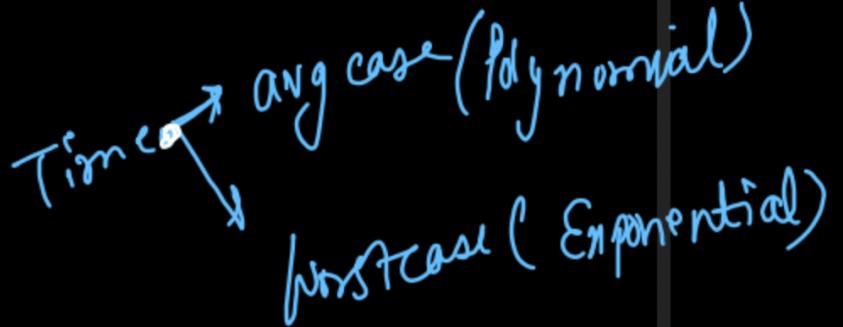
public void helper(int i, int j, String s1, String s2, int[][] dp, String lcs){
    if(i == s1.length() || j == s2.length()){
        answers.add(lcs);
        return;
    }

    char ch1 = s1.charAt(i);
    char ch2 = s2.charAt(j);

    if(ch1 == ch2){
        // Character Taken (Same)
        helper(i + 1, j + 1, s1, s2, dp, lcs + ch1);
    } else {
        // Character Not Taken
        if(dp[i][j] == dp[i + 1][j]){
            helper(i + 1, j, s1, s2, dp, lcs);
        }

        if(dp[i][j] == dp[i][j + 1]){
            helper(i, j + 1, s1, s2, dp, lcs);
        }
    }
}

```


 Time  
 Best Case (Exponential)  
 Worst Case (Polynomial)

```

public List<String> all_longest_common_subsequences(String s1, String s2)
{
    int[][] dp = new int[s1.length() + 1][s2.length() + 1];

    for(int i=s1.length()-1; i>=0; i--){
        for(int j=s2.length()-1; j>=0; j--){
            char ch1 = s1.charAt(i);
            char ch2 = s2.charAt(j);

            if(ch1 == ch2)
                dp[i][j] = 1 + dp[i + 1][j + 1];

            else dp[i][j] = Math.max(dp[i + 1][j], dp[i][j + 1]);
        }
    }

    answers = new TreeSet<>();
    helper(0, 0, s1, s2, dp, "");
    List<String> res = new ArrayList<>();
    for(String str: answers){
        res.add(str);
    }
    return res;
}

```

718 Longest Common Substring

Contiguous

~~eg~~

"dabcde" vs "bdca"

Substitution

bc → ✓ ✗

bd → ✗ ✓

dc → ✗ ✓

da → ✗ ✗

Longest common substring  
→ "d", "a", "b", "c"

~~eg~~

"abcdabc",

$$S_1[i] = S_2[j]$$

LCSS(i, j)

LCSS(i+1, j+1)

$S_1[i] \neq S_2[j]$

~~LCSS(i+1, j+1)~~

	b	c	d	b	c	a	..
a	0	0	0	0	0	1	0
b	3 "bd"	0	0	2 "bc"	0	0	0
c	0	2 "cd"	0	0	1 "c"	0	0
d	0	0	1 "d"	0	0	0	0
a	0	0	0	0	0	1 "a"	0
b	2 "b"	0	0	2 "bc"	0	0	0
c	0	1 "c"	0	0	1 "c"	0	0
..	0	0	0	0	0	0	0

```

class Solution {
    public int findLength(int[] s1, int[] s2) {
        int[][] dp = new int[s1.length + 1][s2.length + 1];

        int maxLen = 0;
        for(int i=s1.length-1; i>=0; i--){
            for(int j=s2.length-1; j>=0; j--){
                if(s1[i] == s2[j])
                    dp[i][j] = 1 + dp[i + 1][j + 1];

                // else dp[i][j] = Math.max(dp[i + 1][j], dp[i][j + 1]);
                maxLen = Math.max(maxLen, dp[i][j]);
            }
        }

        return maxLen;
    }
}

```

Time  $\rightarrow O(N \times M)$

Space  $\rightarrow O(N \times M)$



④ It can be optimized  
to 1D DP

# Longest Common Subsequence - II

29<sup>th</sup> May, Sunday, 9 AM - 12 PM

- Longest Repeating / Duplicate Subsequence (QFG)
- Longest Duplicate String
  - I (104 LC)
  - II (QFG)
- Longest Palindromic Subsequence (LC 516)
  - Insert (QFG)
- Min<sup>m</sup> Steps to Make Palindrome
  - Delete (1312 LC)
- Palindromic Substrings
  - Longest (5 LC)
  - Count (647 LC)



# Longest Repeating Subsequence

LCS

"ab<sup>.</sup>a<sup>.</sup>c<sup>.</sup>b<sup>.</sup>c<sup>.</sup>"

	aback	backc	acbc	c <sup>.</sup> bc	bc	c	" "
abacbc	3 "abc"	3 "abc"	3 "abc"	2 "bu"	2 "bu"	1 "c"	0
bacbc	3 "bk"	2 "bc"	2 "bc"	2 "bc"	2 "bc"	1 "c"	0
acbc	3 "abc"	2 "bc"	1 "c"	1 "c"	1 "c"	1 "c"	0
dbc	2 "bc"	2 "bc"	1 "c"	1 "c"	1 "c"	1 "c"	0
bc	2 "bc"	2 "bc"	1 "c"	1 "c"	0	0	0
c	1 "c"	1 "c"	1 "c"	1 "c"	0	0	0
" "	0	0	0	0	0	0	0

$$\begin{aligned} \text{diff char} \\ dp[i][j] \\ = dp[i+1][j] \\ dp[i][j+1] \end{aligned}$$

$$\begin{aligned} \text{Same char} \\ 2 & \& 1 = j \\ dp[i][j] &= dp[i+1][j+1] \\ \neq 1 \end{aligned}$$

```

public int LongestRepeatingSubsequence(String s)
{
    int[][] dp = new int[s.length() + 1][s.length() + 1];

    for(int i=s.length()-1; i>=0; i--){
        for(int j=s.length()-1; j>=0; j--){
            char ch1 = s.charAt(i);
            char ch2 = s.charAt(j);

            if(ch1 == ch2 && i != j) → longest repeating subset
                dp[i][j] = 1 + dp[i + 1][j + 1];
            else dp[i][j] = Math.max(dp[i + 1][j], dp[i][j + 1]);
        }
    }

    return dp[0][0];
}

```

with same  
↳ LCS ↳  
longest repeating subset

Longest Repeating Subsequence									
	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9
2	2	3	4	5	6	7	8	9	10
3	3	4	5	6	7	8	9	10	11
4	4	5	6	7	8	9	10	11	12
5	5	6	7	8	9	10	11	12	13
6	6	7	8	9	10	11	12	13	14
7	7	8	9	10	11	12	13	14	15
8	8	9	10	11	12	13	14	15	16
9	9	10	11	12	13	14	15	16	17

Longest Repeating Substring

Lecture 1044

Non-overlapping (GFG)

Overlapping (Leetcode)

e.g. "banana"

	b	a	n	a	n	a	..
b	0	0	0	0	0	0	0
a	0	0	0	3 "ana"	0	1 "a"	0
n	0	0	0	0	2 "na"	0	0
a	0	3 "ana"	0	0	0	1 "a"	0
n	0	0	2 "na"	0	0	0	0
a	0	1 "a"	0	1 "a"	0	0	0
..	0	0	0	0	0	0	0

longest common  
substring  
in 2 same  
strings

$\{dp(i)(j) = \text{LCS } \sigma\}$

$s1.\text{substr}(i)$

$s2.\text{substr}(j)$

meaning

1044

```
public String longestDupSubstring(String s) {  
    int[][] dp = new int[s.length() + 1][s.length() + 1];  
  
    int idx = s.length(), len = 0;  
  
    for(int i=s.length()-1; i>=0; i--){  
        for(int j=s.length()-1; j>=0; j--){  
            char ch1 = s.charAt(i);  
            char ch2 = s.charAt(j);  
  
            if(ch1 == ch2 && i != j)  
                dp[i][j] = 1 + dp[i + 1][j + 1];  
  
            if(dp[i][j] > len){  
                idx = i;  
                len = dp[i][j];  
            }  
  
            // else dp[i][j] = Math.max(dp[i + 1][j], dp[i][j + 1]);  
        }  
    }  
  
    return s.substring(idx, idx + len);  
}
```

(TLE) → Expected Rolling Hash

Time →  $O(N^2)$

Space →  $O(N^2)$  2D DP

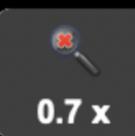


Largest Repeated / Duplicate Substrings  
with non-overlapping (QFG)

"banana" length = 2 ("an")

b	0	0	0	0	0	0	0
a'	0	0	0	3 "ana"	0	1 "a"	0
n'	0	0	0	0	2 "na"	0	0
a''	0	3 "na"	0	0	0	1 "a"	0
n''	0	0	2 "na"	0	0	0	0
a'''	0	1 "a"	0	1 "a"	0	0	0
..	0	0	0	0	0	0	0

$\star$   
 $\min(i,j) + \text{dplis}(j)$   
 $< \max(i,j)$



0.7 x

```
int[][] dp = new int[s.length() + 1][s.length() + 1];
int idx = s.length(), len = 0;

for(int i=s.length()-1; i>=0; i--){
    for(int j=s.length()-1; j>=0; j--){
        char ch1 = s.charAt(i);
        char ch2 = s.charAt(j);

        if(ch1 == ch2 && i != j)
            dp[i][j] = 1 + dp[i + 1][j + 1];

        if(Math.min(i, j) + dp[i][j] >= Math.max(i, j)){
            // Overlapping Substrings
            dp[i][j] = 0;
        }

        if(dp[i][j] >= len){
            idx = i;
            len = dp[i][j];
        }
    }
}
```

```
if(len == 0) return "-1";
return s.substring(idx, idx + len);
```

$\text{abs}(j-i) \leq d\varphi(i,j)$  (1)  
⊕ To compute only non-overlapping substring  
 Time  $\rightarrow O(n^2)$   
 Space  $\rightarrow O(n^2)$

Large Regional Depthless Anabiosis						
	1	2	3	4	5	6
1	0	0	0	0	0	0
2	0	0	0	0	0	0
3	0	0	0	0	0	0
4	0	0	0	0	0	0
5	0	0	0	0	0	0
6	0	0	0	0	0	0
7	0	0	0	0	0	0
8	0	0	0	0	0	0
9	0	0	0	0	0	0
10	0	0	0	0	0	0
11	0	0	0	0	0	0
12	0	0	0	0	0	0
13	0	0	0	0	0	0
14	0	0	0	0	0	0
15	0	0	0	0	0	0
16	0	0	0	0	0	0
17	0	0	0	0	0	0
18	0	0	0	0	0	0
19	0	0	0	0	0	0
20	0	0	0	0	0	0
21	0	0	0	0	0	0
22	0	0	0	0	0	0
23	0	0	0	0	0	0
24	0	0	0	0	0	0
25	0	0	0	0	0	0
26	0	0	0	0	0	0
27	0	0	0	0	0	0
28	0	0	0	0	0	0
29	0	0	0	0	0	0
30	0	0	0	0	0	0
31	0	0	0	0	0	0
32	0	0	0	0	0	0
33	0	0	0	0	0	0
34	0	0	0	0	0	0
35	0	0	0	0	0	0
36	0	0	0	0	0	0
37	0	0	0	0	0	0
38	0	0	0	0	0	0
39	0	0	0	0	0	0
40	0	0	0	0	0	0
41	0	0	0	0	0	0
42	0	0	0	0	0	0
43	0	0	0	0	0	0
44	0	0	0	0	0	0
45	0	0	0	0	0	0
46	0	0	0	0	0	0
47	0	0	0	0	0	0
48	0	0	0	0	0	0
49	0	0	0	0	0	0
50	0	0	0	0	0	0
51	0	0	0	0	0	0
52	0	0	0	0	0	0
53	0	0	0	0	0	0
54	0	0	0	0	0	0
55	0	0	0	0	0	0
56	0	0	0	0	0	0
57	0	0	0	0	0	0
58	0	0	0	0	0	0
59	0	0	0	0	0	0
60	0	0	0	0	0	0
61	0	0	0	0	0	0
62	0	0	0	0	0	0
63	0	0	0	0	0	0
64	0	0	0	0	0	0
65	0	0	0	0	0	0
66	0	0	0	0	0	0
67	0	0	0	0	0	0
68	0	0	0	0	0	0
69	0	0	0	0	0	0
70	0	0	0	0	0	0
71	0	0	0	0	0	0
72	0	0	0	0	0	0
73	0	0	0	0	0	0
74	0	0	0	0	0	0
75	0	0	0	0	0	0
76	0	0	0	0	0	0
77	0	0	0	0	0	0
78	0	0	0	0	0	0
79	0	0	0	0	0	0
80	0	0	0	0	0	0
81	0	0	0	0	0	0
82	0	0	0	0	0	0
83	0	0	0	0	0	0
84	0	0	0	0	0	0
85	0	0	0	0	0	0
86	0	0	0	0	0	0
87	0	0	0	0	0	0
88	0	0	0	0	0	0
89	0	0	0	0	0	0
90	0	0	0	0	0	0
91	0	0	0	0	0	0
92	0	0	0	0	0	0
93	0	0	0	0	0	0
94	0	0	0	0	0	0
95	0	0	0	0	0	0
96	0	0	0	0	0	0
97	0	0	0	0	0	0
98	0	0	0	0	0	0
99	0	0	0	0	0	0
100	0	0	0	0	0	0

"abba"      "aaa"      "ababa"      "aca"

# Longest Palindromic Subsequence

$$\text{Rev}("abca\bar{c}ba") = ab\bar{c}a\bar{c}ba$$

acaca, ababa

"abcaca.b-a

```

public int LCS(String s1, String s2){
    int[][] dp = new int[s1.length() + 1][s2.length() + 1];

    for(int i=s1.length()-1; i>=0; i--){
        for(int j=s2.length()-1; j>=0; j--){
            char ch1 = s1.charAt(i);
            char ch2 = s2.charAt(j);

            if(ch1 == ch2)
                dp[i][j] = 1 + dp[i + 1][j + 1];

            else dp[i][j] = Math.max(dp[i + 1][j], dp[i][j + 1]);
        }
    }

    return dp[0][0];
}

public int longestPalindromeSubseq(String s) {
    StringBuilder rev = new StringBuilder(s);
    rev = rev.reverse();
    String s2 = new String(rev);

    return LCS(s, s2);
}

```

LPS  
 $= LCS(s, \text{rev}(s))$

Time  $\rightarrow O(n^2)$

Space  $\approx O(n^2)$

"abcabca"  
 min deletion  
 to make palindrome  
 ↓  
 "abcabca"

min insertion  
 to make  
 palindrome

"abcabca"  
 \* \* \* \*

"a b c <sup>b</sup> a b c <sup>b</sup> a"  
 \* \* \* \*

$$2ns = \text{Str.length} - LPS$$

$$2ns = \text{Str.length} - LPS$$

```

class Solution {
    public int minInsertions(String s) {
        int n = s.length();
        int lps = lcs(s, s);
        return n - lps;
    }

    private int lcs(String s1, String s2) {
        int m = s1.length();
        int n = s2.length();
        int dp[][] = new int[m+1][n+1];
        for(int i=0; i<m+1; i++) {
            for(int j=0; j<n+1; j++) {
                if(i==0 || j==0)
                    dp[i][j] = 0;
                else if(s1.charAt(i-1) == s2.charAt(j-1))
                    dp[i][j] = dp[i-1][j-1] + 1;
                else
                    dp[i][j] = Math.max(dp[i-1][j], dp[i][j-1]);
            }
        }
        return dp[m][n];
    }
}
    
```

```

class Solution {
    public int LCS(String s1, String s2){
        int[][] dp = new int[s1.length() + 1][s2.length() + 1];

        for(int i=s1.length()-1; i>=0; i--){
            for(int j=s2.length()-1; j>=0; j--){
                char ch1 = s1.charAt(i);
                char ch2 = s2.charAt(j);

                if(ch1 == ch2)
                    dp[i][j] = 1 + dp[i + 1][j + 1];

                else dp[i][j] = Math.max(dp[i + 1][j], dp[i][j + 1]);
            }
        }

        return dp[0][0];
    }

    public int minInsertions(String s) {
        StringBuilder rev = new StringBuilder(s);
        rev = rev.reverse();
        String s2 = new String(rev);

        return s.length() - LCS(s, s2);
    }
}

```

Min Insertions  
 Deletions  
 to make  
 Palindrome



# # Longest Palindromic Substring

"abc abac" LPSL = "aba" → ③

$$LPSS = LCSS(s, \text{rev}(s))$$

Longest Palindromic Substring      ↗      Longest Common Subsequence

Time  $\sim O(N^2)$

Space  $\rightarrow O(n^2)$



```

public String longestCommonSubstring(String s1, String s2){
    int[][] dp = new int[s1.length() + 1][s2.length() + 1];

    String res = "";

    for(int i=s1.length()-1; i>=0; i--){
        for(int j=s2.length()-1; j>=0; j--){
            char ch1 = s1.charAt(i);
            char ch2 = s2.charAt(j);

            if(ch1 == ch2)
                dp[i][j] = 1 + dp[i + 1][j + 1];

            if(dp[i][j] >= res.length()){
                String curr = s1.substring(i, i + dp[i][j]);

                if(isPalindrome(curr) == true){
                    res = curr;
                }
            }
        }
    }

    return res;
}

```

```

public String longestPalindrome(String s) {
    StringBuilder rev = new StringBuilder(s);
    rev = rev.reverse();
    String s2 = new String(rev);

    return longestCommonSubstring(s, s2);
}

```

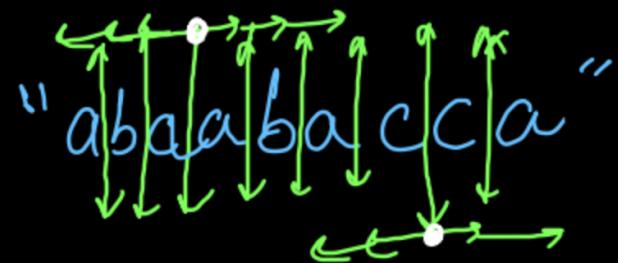
This type of  
 DP might go  
 wrong for  
 some test cases

abcaba  
 abcaba  
 Time = O(n^2)  
 Space = O(n^2)

Expand Around  
 Centers  
 Greedy  
 odd length  
 palindromic substrings  
 even length  
 palindromic  
 substrings



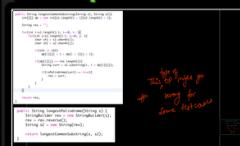
↳ Treat each character as  
 center for odd length  
 palindromes



↳ Take in  $2w^2$  characters  
 as middle for even length  
 palindromes

(odd)  
 LPSS = "a"  
 "b"  
 "c"  
 "aba"  
 "cabac"

(even)  
 LPSS = "abba"  
 "abaca"



Time →  $O(N^2)$  →  $N \times O(N)$  *expand time in worst case*, Space →  $O(1)$

```

String res = "";
for(int i=0; i<s.length(); i++){
    // Odd Length Palindromes
    int left = i - 1, right = i + 1, len = 1;
    while(left >= 0 && right < s.length()){
        if(s.charAt(left) == s.charAt(right)){
            len = len + 2;
            left--; right++;
        } else {
            break;
        }
    }

    if(len > res.length()){
        res = s.substring(left + 1, left + 1 + len);
    }
}

```

```

for(int i=0; i<s.length(); i++){
    // Even Length Palindromes
    if(i + 1 == s.length() || s.charAt(i) != s.charAt(i + 1)){
        continue;
    }

    int left = i - 1, right = i + 2, len = 2;
    while(left >= 0 && right < s.length()){
        if(s.charAt(left) == s.charAt(right)){
            len = len + 2;
            left--; right++;
        } else break;
    }

    if(len > res.length()){
        res = s.substring(left + 1, left + 1 + len);
    }
}

return res;

```

Diagram illustrating the search for palindromes in the string "abbaabacca". The string is shown with a green box highlighting the segment "abbaabacca". A blue box highlights the center of the palindrome "abbaabacca". Handwritten annotations show  $l = i - 1$ ,  $r = i + 2$ , and  $len = 7$ .

