

$\text{len}^{\text{arr}} = \frac{1}{2} \times 5$

$$\begin{aligned} \text{inc}[i] \\ = \text{dec}[j] + 1 \\ \xleftarrow{\quad \text{dec}[i] \\ = \text{inc}[j] + 1 \quad} \end{aligned}$$

$I \rightarrow \{1, 17\}$
 $D \rightarrow \{17\}$

$I \rightarrow \{1, 17, 5, 10\}$
 $D \rightarrow \{1, 17, 10\}$

$I \rightarrow \{1, 17, 5, 15\}$
 $D \rightarrow \{17, 15\}$

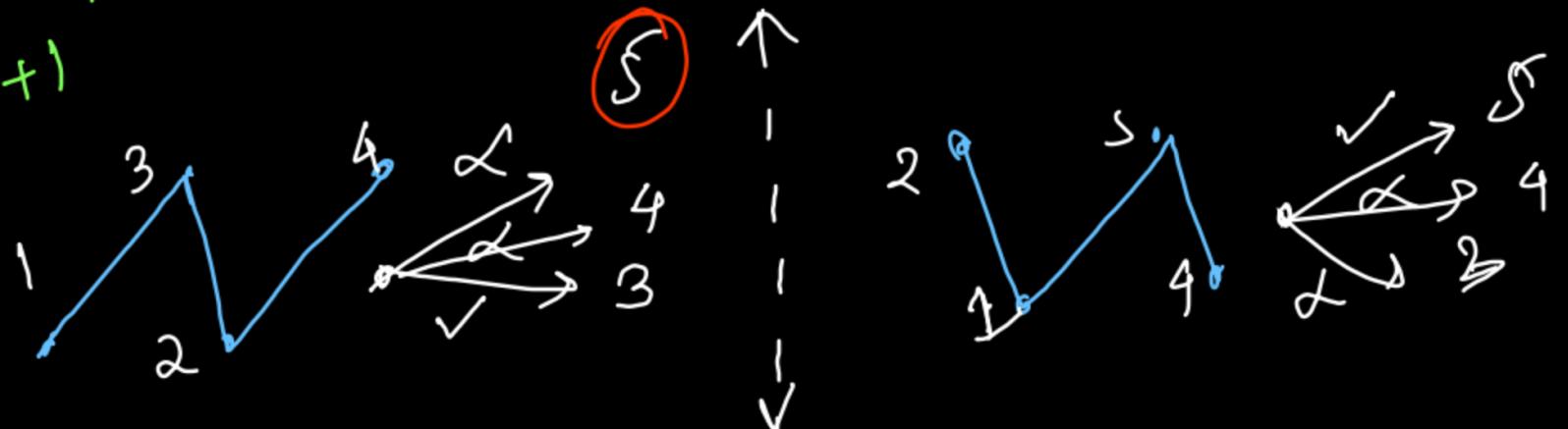
[1, 17, 5, 10, 13, 15, 10, 5, 16, 8]

$I \rightarrow \{1\}$
 $D \rightarrow \{1\}$

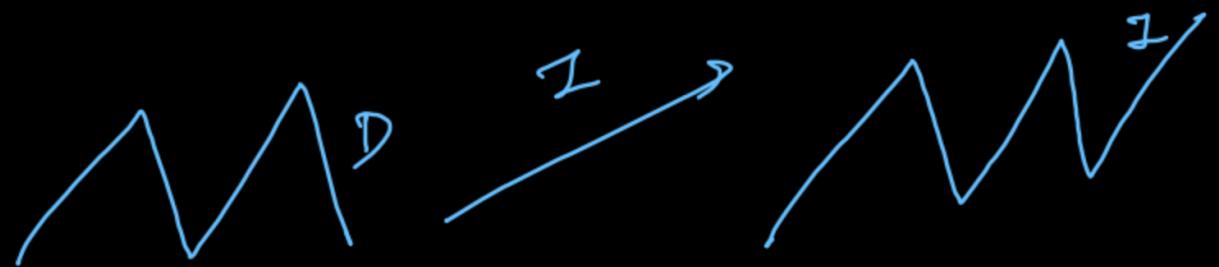
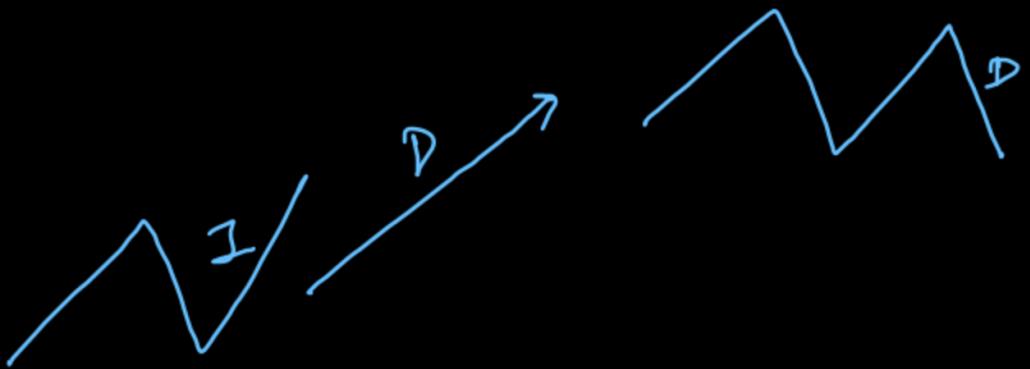
$I \rightarrow \{1, 5\}$
 $D \rightarrow \{1, 17, 5\}$

$I \rightarrow \{1, 17, 5, 13\}$
 $D \rightarrow \{1, 17, 13\}$

$I \rightarrow \{1, 17, 5, 10\}$
 $D \rightarrow \{1, 17, 5, 10\}$



$$\text{dec}(i) = \text{inc}(j) + 1$$



$$\begin{aligned}\text{inc}(i) \\ = \text{dec}(j) + 1\end{aligned}$$

```

public int wiggleMaxLength(int[] nums) {
    int[] inc = new int[nums.length];
    Arrays.fill(inc, 1);

    int[] dec = new int[nums.length];
    Arrays.fill(dec, 1);

    int maxLength = 0;
    for(int i=0; i<nums.length; i++){

        for(int j=0; j<i; j++){
            if(nums[i] > nums[j]){
                // We are Increasing
                inc[i] = Math.max(inc[i], dec[j] + 1);
            } else if(nums[i] < nums[j]){
                // We are Decreasing
                dec[i] = Math.max(dec[i], inc[j] + 1);
            }
        }

        maxLength = Math.max(maxLength, Math.max(inc[i], dec[i]));
    }

    return maxLength;
}

```

Time $\rightarrow O(2 \cdot n^2)$
 $= O(n^2)$

Space $\rightarrow O(2 \cdot n)$

2 rows } 1D DP

LIS on 2D Coordinate

Overlapping Boxes
Practice - GFG

Russian Doll

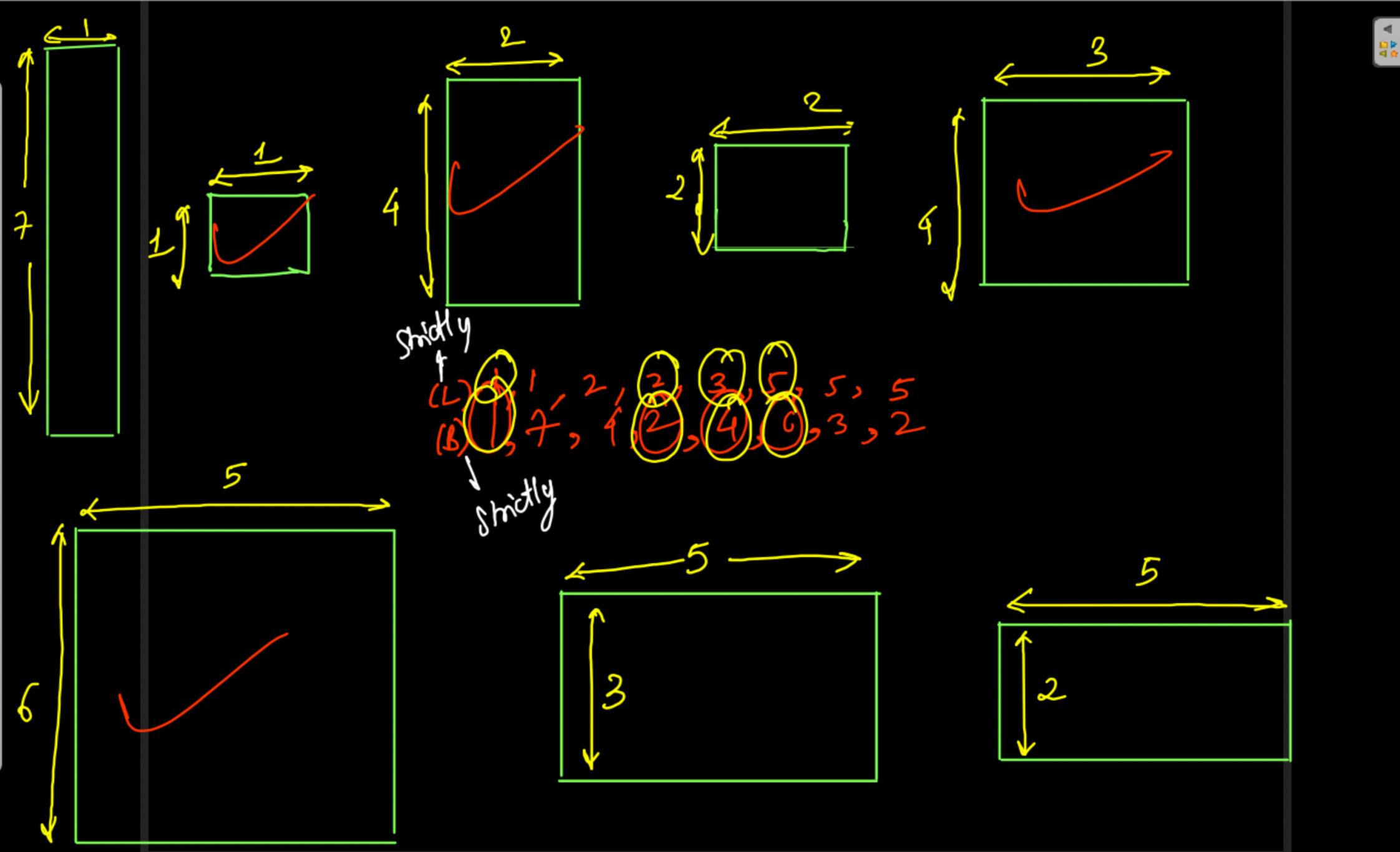


Bon Stacking

GFG
IV-1

Leetcode
II

One envelope can fit into another if and only if both the width and height of one envelope are greater than the other envelope's width and height.

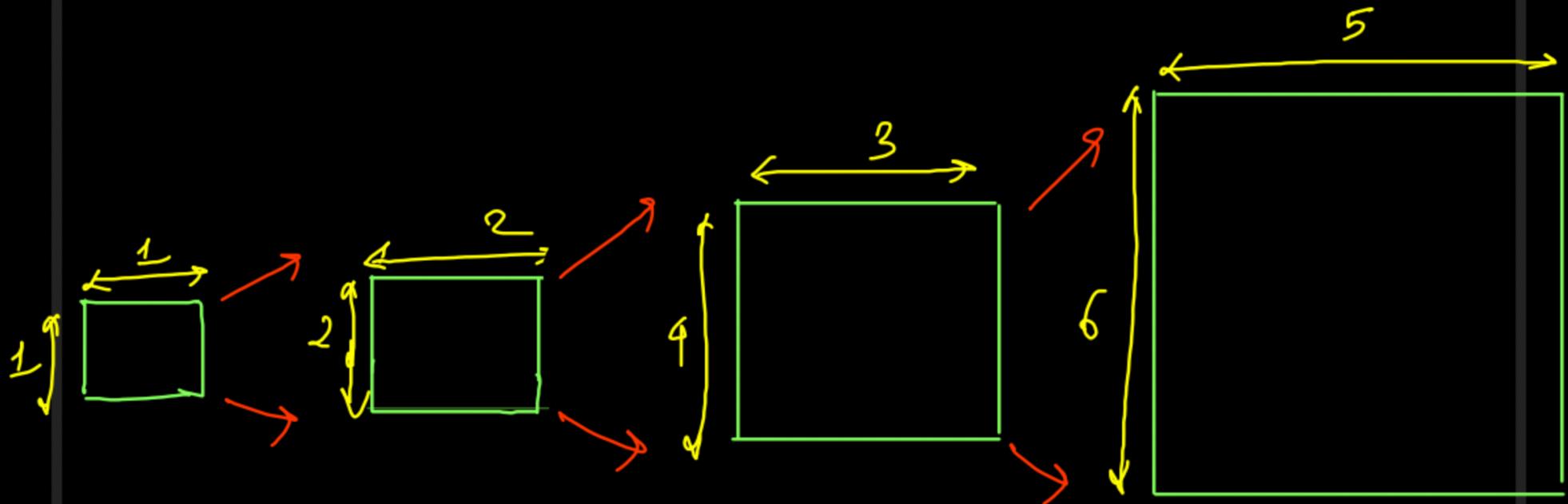


Time: $O(N \log N + N \log N)$
↓
Breadth
length

Space: $O(N)$

```
public int lowerBound(ArrayList<Integer> nums, int target){  
    int low = 0, high = nums.size() - 1;  
    int idx = nums.size();  
  
    while(low <= high){  
        int mid = low + (high - low) / 2;  
  
        if(nums.get(mid) < target){  
            low = mid + 1;  
        } else {  
            high = mid - 1;  
            idx = mid;  
        }  
    }  
  
    return idx;  
}
```

```
public int maxEnvelopes(int[][] envelopes) {  
    Arrays.sort(envelopes, (int[] first, int[] second) -> {  
        return (first[0] != second[0]) ? first[0] - second[0]  
            : second[1] - first[1];  
    });  
  
    int n = envelopes.length;  
    ArrayList<Integer> sorted = new ArrayList<>();  
  
    for(int i=0; i<envelopes.length; i++){  
        int lb = lowerBound(sorted, envelopes[i][1]);  
        if(lb == sorted.size()){  
            sorted.add(envelopes[i][1]);  
            // Current Element larger than the largest  
            // LIS of one length more  
        } else {  
            sorted.set(lb, envelopes[i][1]);  
        }  
    }  
  
    return sorted.size(); // This Sorted Array has same size as LIS  
}
```



LIS on length

LIS on breadth

1, 2, 3, 5

1, 2, 4, 6

2D LIS

1st coordinate → Sorting
2nd coordinate → LIS

longest Common Subsequence

Lecture-1 Saturday (28 May)

9 AM to 12 PM

→ length of LCS

→ Print LCS

Any one

All

→ longest Repeated subset (LRS)

→ Palindromic
subsequence

length

Count

Highest Common Subsequence

is

order of characters

e.g.

"a[.]b[.]c[.]d[.]e" , "i[.]d[.]c[.]a"

Brute force

$O(2^N \times 2^M)$

Compare all
subsequences

$S[i:j] = S[j:j]$
 $LCS(i, j)$
 ↓
 char
 a
 \uparrow
 \downarrow
 $LCS(i+1, j+1)$

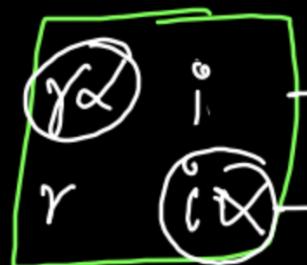
char taken (1)

	$a\nu$	$a\nu$	$\rightarrow n$
1	$a\nu$	$a\nu$	length l
2	$a\nu$	$a\nu$	
3	$a\nu$	$a\nu$	

$S_1[i:j] = S_2[i:j]$ 

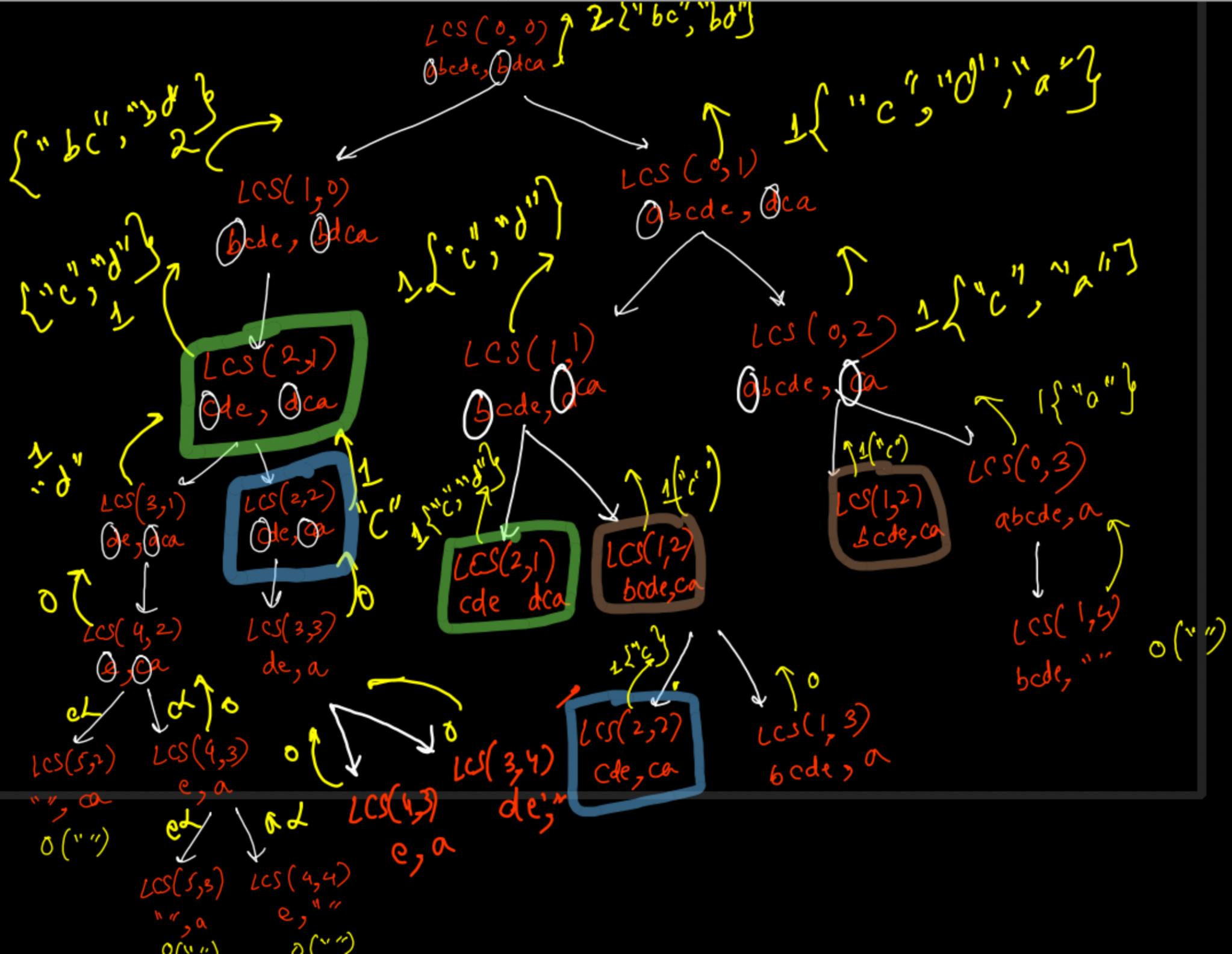
LCS(i, j) 

char not taken = $\delta + LCS(i+1, j)$ $\delta + LCS(i, j+1)$



$r \vee i \vee c \rightarrow$ Common \sqsubset
 $r c \quad i c \quad c \rightarrow$ Length \downarrow

$LCS(i, j)$
 $\delta + LCS(i+1, j+1)$



```

public int LCS(int i, int j, String s1, String s2, int[][] dp){
    if(i == s1.length() || j == s2.length())
        return 0; // LCS of Empty String with Other String is Empty String only

    if(dp[i][j] != -1) return dp[i][j];

    char ch1 = s1.charAt(i);
    char ch2 = s2.charAt(j);

    if(ch1 == ch2) // If characters are same, take the common from both of them
        return dp[i][j] = 1 + LCS(i + 1, j + 1, s1, s2, dp);

    // If character is uncommon, either not take s1[i] or not take s2[j]
    int option1 = LCS(i + 1, j, s1, s2, dp);
    int option2 = LCS(i, j + 1, s1, s2, dp);
    return dp[i][j] = 0 + Math.max(option1, option2);
}

```

```

public int longestCommonSubsequence(String s1, String s2) {
    int[][] dp = new int[s1.length() + 1][s2.length() + 1];
    for(int i=0; i<dp.length; i++){
        for(int j=0; j<dp[0].length; j++){
            dp[i][j] = -1;
        }
    }

    return LCS(0, 0, s1, s2, dp);
}

```

Time $\rightarrow O(N \times M)$

Space $\rightarrow 2D DP$
 $O(N \times M)$

eg "abcde", "bdca"

	0	1	2	3	4
0	"abcde"	"bdcba"	"dca"	"a"	" "
1	"bde"	"cda"	"ca"	" "	" "
2	"bc", "bd"	"c", "d", "a"	"c", "a"	"a"	" "
3	"bc", "bd", "d"	"c", "d", "a"	"c"	" "	" "
4	"c", "d"	"c", "d", "a"	"c"	" "	" "
5	"d"	"d", "a"	"d", "a"	"a"	" "
6	" "	" "	" "	" "	" "
7	" "	" "	" "	" "	" "
8	" "	" "	" "	" "	" "
9	" "	" "	" "	" "	" "
10	" "	" "	" "	" "	" "

Smallest

```

public int longestCommonSubsequence(String s1, String s2) {
    int[][] dp = new int[s1.length() + 1][s2.length() + 1];

    for(int i=s1.length()-1; i>=0; i--){
        for(int j=s2.length()-1; j>=0; j--){
            char ch1 = s1.charAt(i);
            char ch2 = s2.charAt(j);

            if(ch1 == ch2)
                dp[i][j] = 1 + dp[i + 1][j + 1];
            else dp[i][j] = Math.max(dp[i + 1][j], dp[i][j + 1]);
        }
    }

    return dp[0][0];
}

```

Time $\rightarrow O(N \times M)$

Space $\rightarrow O(N \times M)$

$\boxed{2D DP}$

```

int[] next = new int[s2.length() + 1];

for(int i=s1.length()-1; i>=0; i--){
    int[] curr = new int[s2.length() + 1];

    for(int j=s2.length()-1; j>=0; j--){
        char ch1 = s1.charAt(i);
        char ch2 = s2.charAt(j);

        if(ch1 == ch2)
            curr[j] = 1 + next[j + 1];
        else curr[j] = Math.max(next[j], curr[j + 1]);
    }

    next = curr;
}

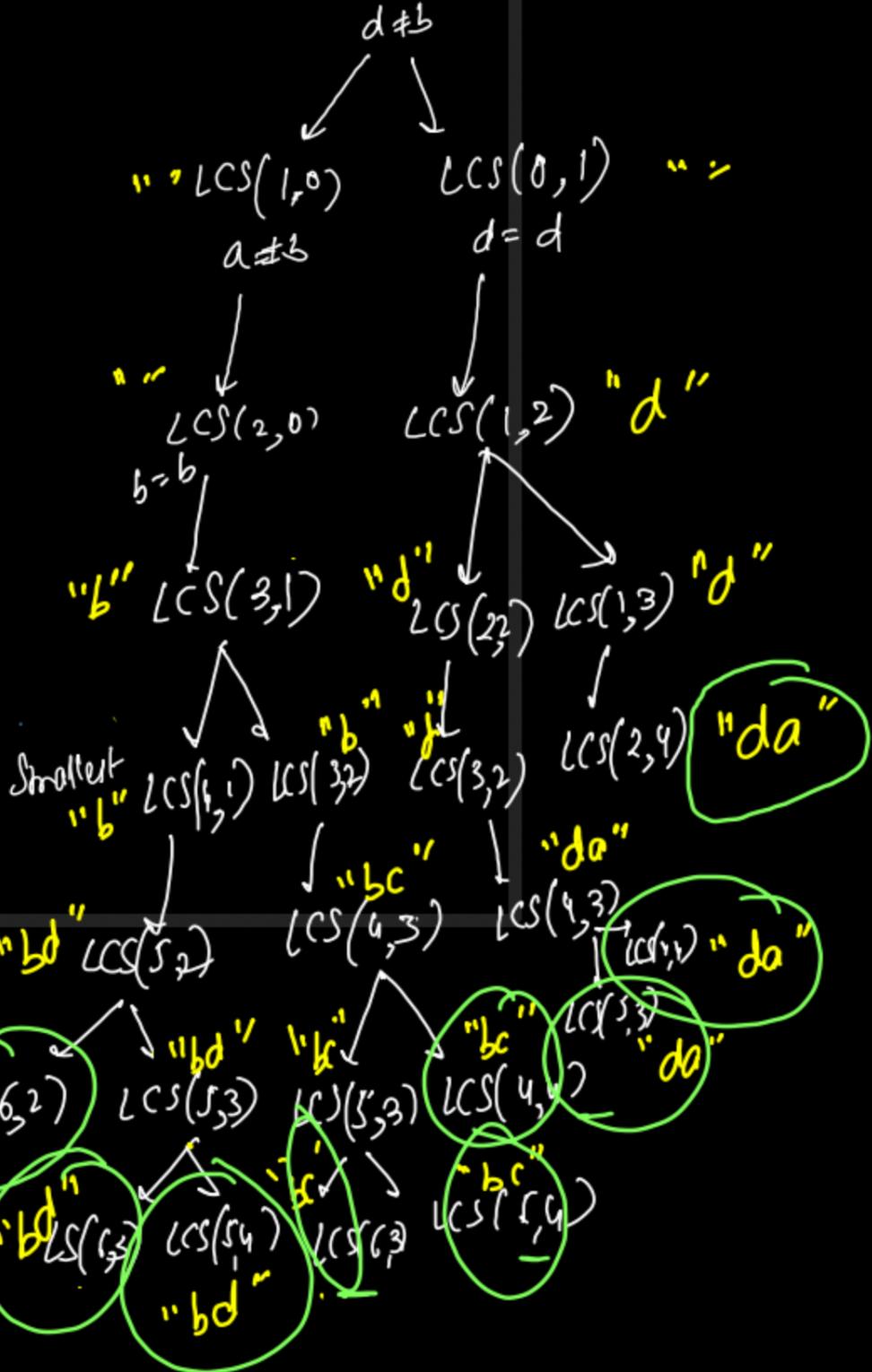
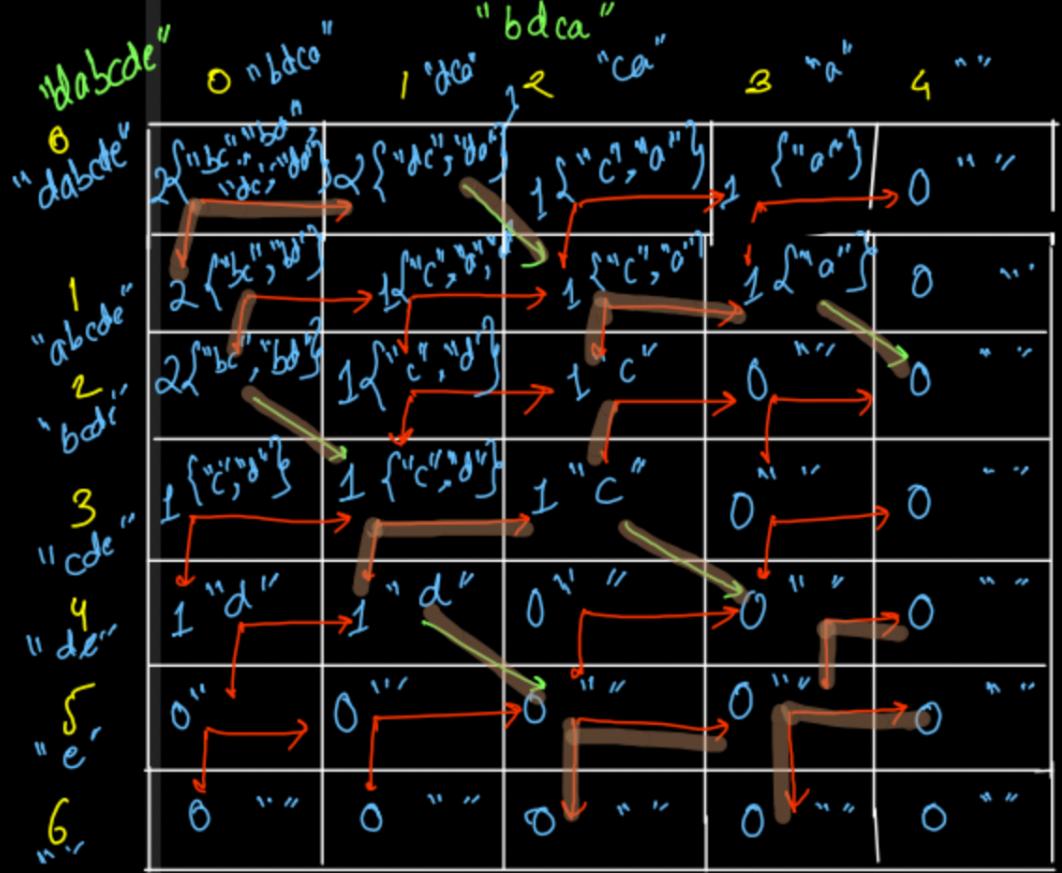
return next[0];

```

Time $\rightarrow O(N \times M)$

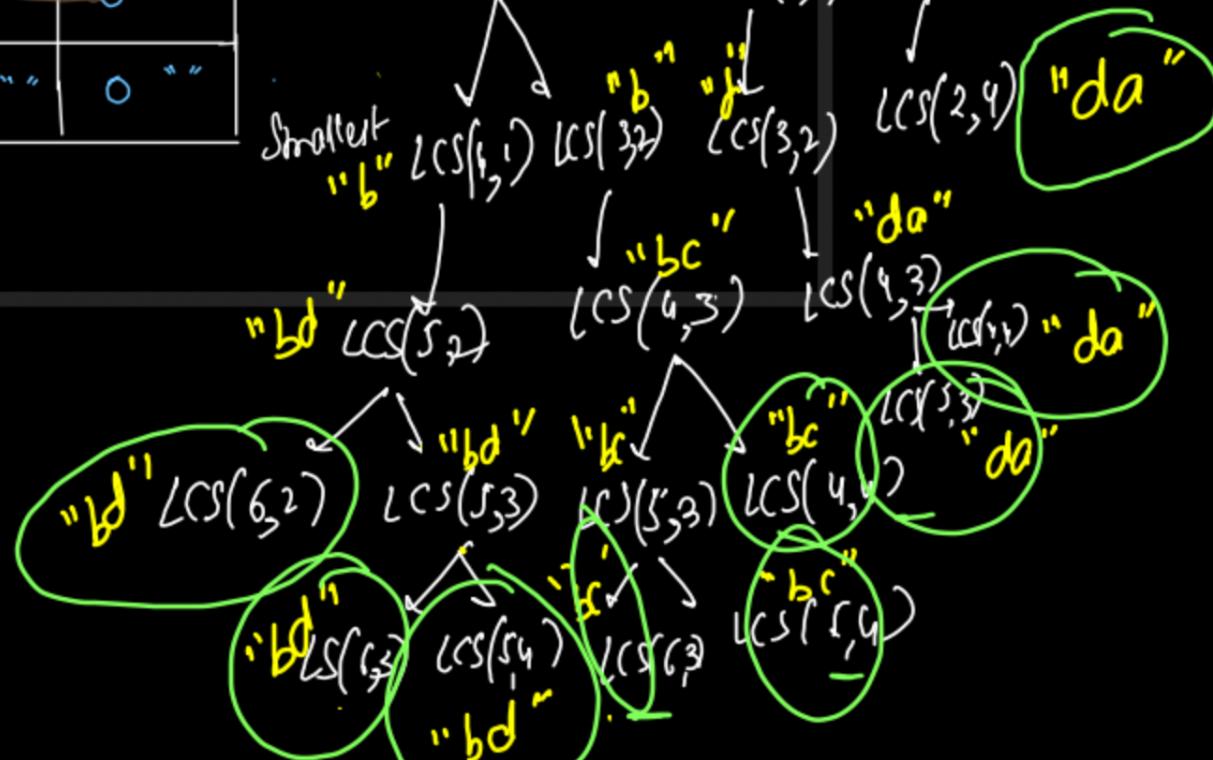
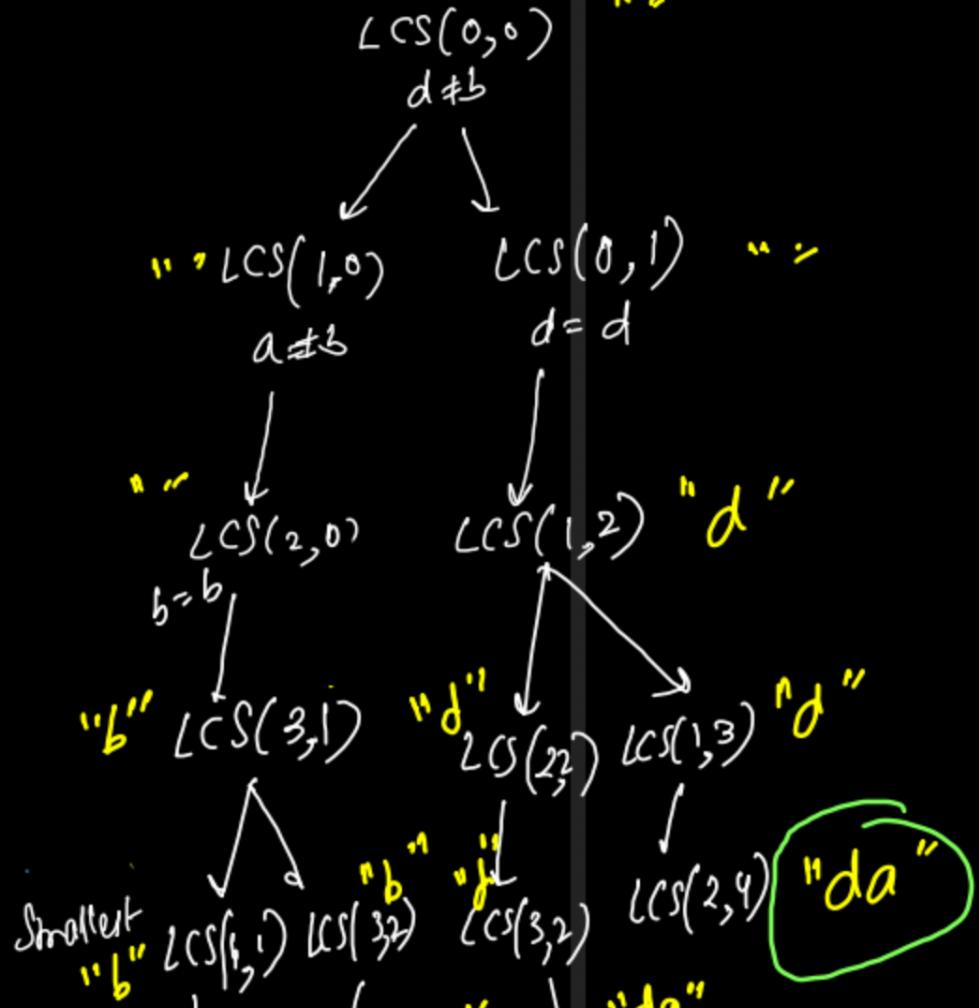
Space $\rightarrow O(2 \times M)$

$\approx \boxed{1D DP}$



	0 "bdc"	1 "dc"	2 "ca"	3 "a"	4 "
0 "abcde"	"bdc"	"dc"	"ca"	"a"	" "
1 "dabce"	"bc"	"dc", "bdc"	"c", "o"	"a"	" "
2 "abcde"	"bc", "bd"	"c", "o"	"a"	" "	" "
3 "cde"	"c", "d"	"c", "d"	"c"	" "	" "
4 "de"	"d"	"d"	" "	" "	" "
5 "e"	" "	" "	" "	" "	" "
6 "	" "	" "	" "	" "	" "

- ✗ ArrayList → Sorted ✓
- ✗ HashSet → Unique ✓
- ✓ TreeSet → Unique ✓
- Sorted ✓



```

TreeSet<String> answers; // Both Ordered (Lexicographical Order), Unique

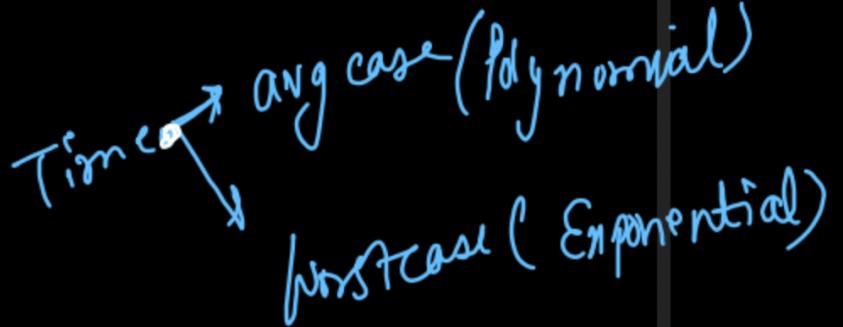
public void helper(int i, int j, String s1, String s2, int[][] dp, String lcs){
    if(i == s1.length() || j == s2.length()){
        answers.add(lcs);
        return;
    }

    char ch1 = s1.charAt(i);
    char ch2 = s2.charAt(j);

    if(ch1 == ch2){
        // Character Taken (Same)
        helper(i + 1, j + 1, s1, s2, dp, lcs + ch1);
    } else {
        // Character Not Taken
        if(dp[i][j] == dp[i + 1][j]){
            helper(i + 1, j, s1, s2, dp, lcs);
        }

        if(dp[i][j] == dp[i][j + 1]){
            helper(i, j + 1, s1, s2, dp, lcs);
        }
    }
}

```


 Time
 Best Case (Exponential)
 Worst Case (Polynomial)

```

public List<String> all_longest_common_subsequences(String s1, String s2)
{
    int[][] dp = new int[s1.length() + 1][s2.length() + 1];

    for(int i=s1.length()-1; i>=0; i--){
        for(int j=s2.length()-1; j>=0; j--){
            char ch1 = s1.charAt(i);
            char ch2 = s2.charAt(j);

            if(ch1 == ch2)
                dp[i][j] = 1 + dp[i + 1][j + 1];

            else dp[i][j] = Math.max(dp[i + 1][j], dp[i][j + 1]);
        }
    }

    answers = new TreeSet<>();
    helper(0, 0, s1, s2, dp, "");
    List<String> res = new ArrayList<>();
    for(String str: answers){
        res.add(str);
    }
    return res;
}

```

718 Longest Common Substring

Contiguous

~~eg~~

"dabcde" vs "bdca"

Substitution

bc → ✓ ✗

bd → ✗ ✓

dc → ✗ ✓

da → ✗ ✗

Longest common substring
→ "d", "a", "b", "c"

~~eg~~

"abcdabc",

$$S_1[i] = S_2[j]$$

LCSS(i, j)

LCSS(i+1, j+1)

$S_1[i] \neq S_2[j]$

~~LCSS(i+1, j+1)~~

	b	c	d	b	c	a	..
a	0	0	0	0	0	1	0
b	3 "bd"	0	0	2 "bc"	0	0	0
c	0	2 "cd"	0	0	1 "c"	0	0
d	0	0	1 "d"	0	0	0	0
a	0	0	0	0	0	1 "a"	0
b	2 "b"	0	0	2 "bc"	0	0	0
c	0	1 "c"	0	0	1 "c"	0	0
..	0	0	0	0	0	0	0

```

class Solution {
    public int findLength(int[] s1, int[] s2) {
        int[][] dp = new int[s1.length + 1][s2.length + 1];

        int maxLen = 0;
        for(int i=s1.length-1; i>=0; i--){
            for(int j=s2.length-1; j>=0; j--){
                if(s1[i] == s2[j])
                    dp[i][j] = 1 + dp[i + 1][j + 1];

                // else dp[i][j] = Math.max(dp[i + 1][j], dp[i][j + 1]);
                maxLen = Math.max(maxLen, dp[i][j]);
            }
        }

        return maxLen;
    }
}

```

Time $\rightarrow O(N \times M)$

Space $\rightarrow O(N \times M)$



④ It can be optimized
to 1D DP