

Dynamic Programming

→ Lecture - 8 9 AM to 12 PM Saturday
30 April

① Paint fence

② Ugly No - II

③ Ugly No - III

④ Egg Jump

⑤ Jump Game - All Paths

Paint fence

(n houses, k colors)

no more than 2 houses
have the same color

$$\# \overbrace{N=1}^{=} \quad k=1 \quad \boxed{\text{RRR}} \quad ①$$

$$\# \overbrace{N=2}^{=} \quad k=1 \quad \boxed{\text{RR}} \quad ①$$

$$k=2 \quad \boxed{\text{RR}}, \quad \boxed{\text{RR}} \quad ②$$

$$k=2 \quad \boxed{\text{RR}}, \quad \boxed{\text{RR}} \quad ③$$

$$k=3 \quad \boxed{\text{RR}}, \quad \boxed{\text{RR}}, \quad \boxed{\text{RR}} \quad ④$$

$$\vdots \downarrow \quad k \rightarrow \boxed{k}$$

$$\boxed{\text{RR}} \quad \boxed{\text{RR}} \quad \boxed{\text{RR}} \quad \boxed{\text{RR}} \quad \boxed{\text{RR}}$$

$N=2$
 $k=3$



⑨



$N=2$
 k

\mathbb{R}^2

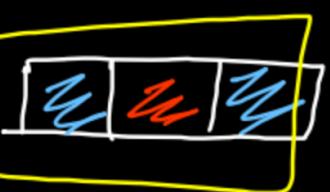
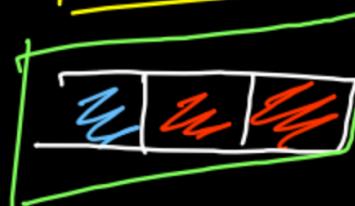
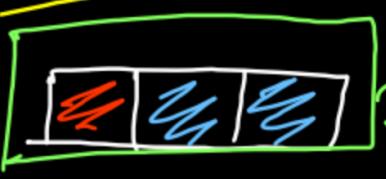
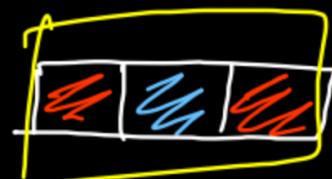
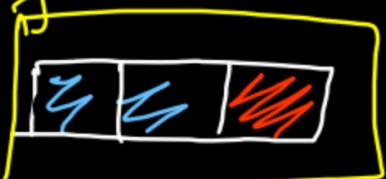
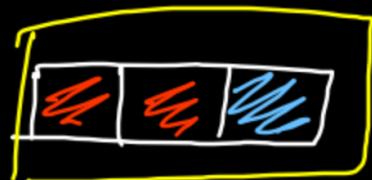
$N=3$

$k=1$



$k=2$ ⑥

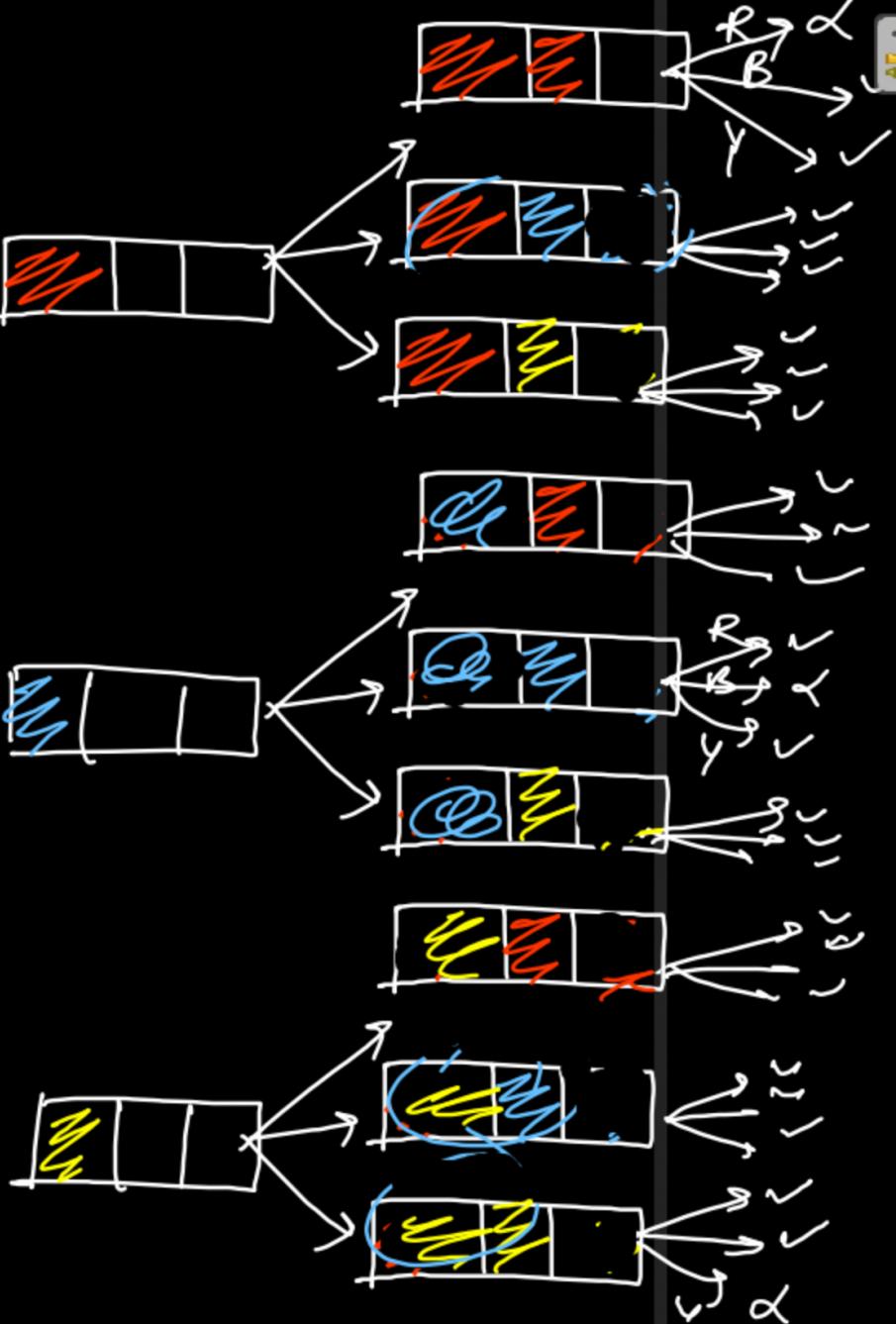
$fd(n, k)$



$k=3$

24

$fs(n, k)$



$$f_n(n, k) = \underbrace{f_{n \text{ same}}(n, k)}_{\substack{n \& n+1 \\ \text{last 2 houses}}} + \underbrace{f_{n \text{ diff}}(n, k)}_{\substack{n \& n+1 \\ \text{last 2 houses} \\ \text{have diff color}}}$$

Painting $n-1$ houses
with k colors
such that
 $(n-1)^{\text{th}}$ and $(n^{\text{th}})^{\text{th}}$
have diff color

$$\boxed{f(n-2, k) * (k-1)}$$

Painting $n-1$ houses
with $\$$ colors
 \Rightarrow we are painting
 n^{th} house with
any color other
than $(n-1)^{\text{th}}$ house

$$\boxed{f(n, k) = \left[f(n-2, k) + f(n-1, k) \right] * (k-1)}$$

Tabulation

(last 2 houses)
Same

| | $N=1$ | $N=2$ | $N=3$ | $N=4$ |
|------|---|--------------|----------------------------------|--------------------------|
| Same |  | 3 | 3×2 | 18 |
| diff |  | 3×2 | $(3 + 3 \times 2) \times 2 = 18$ | $(18 + 6) \times 2 = 48$ |



$k=3$

~~TC $\Rightarrow O(n)$~~
~~SC $\Rightarrow O(n)$~~ \rightarrow space can be optimized

```
public class Solution {
    public int memo(int n, int k, int[] dp){
        if(n == 1) return k;
        if(n == 2) return k * k;
        if(dp[n] != -1) return dp[n];

        int ans = (memo(n - 1, k, dp) + memo(n - 2, k, dp)) * (k - 1);
        return dp[n] = ans;
    }

    public int numWays(int n, int k) {
        if(n == 1) return k;
        if(n == 2) return k * k;
        if(k == 1) return 0;

        int[] dp = new int[n + 1];
        Arrays.fill(dp, -1);
        return memo(n, k, dp);
    }
}
```



0.7 x

```
public int numWays(int n, int k) {  
    if(n == 1) return k;  
    if(n == 2) return k * k;  
    if(k == 1) return 0;  
  
    int[] same = new int[n + 1];  
    int[] diff = new int[n + 1];  
    same[2] = k; diff[2] = k * (k - 1);  
  
    for(int i=3; i<=n; i++){  
        same[i] = diff[i - 1];  
        diff[i] = (same[i - 1] + diff[i - 1]) * (k - 1);  
    }  
  
    return same[n] + diff[n];  
}
```



0.7 x

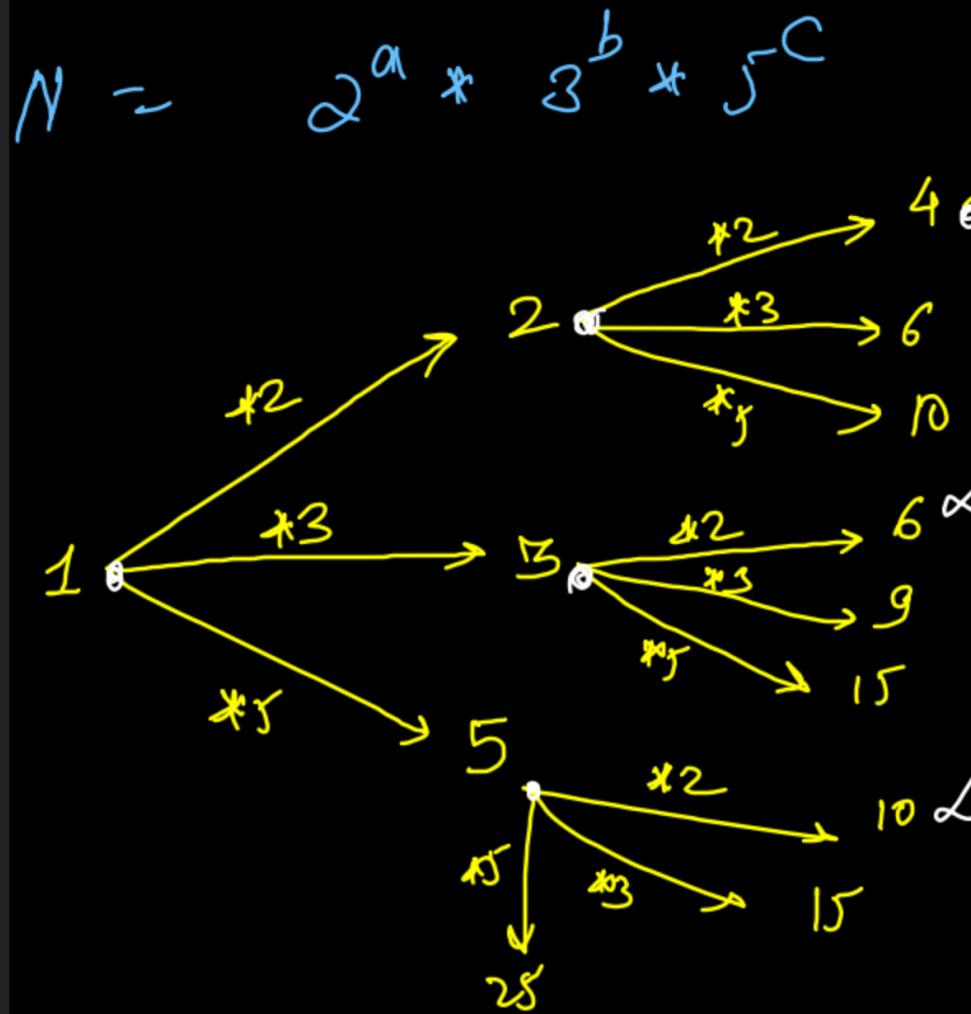
```
public int numWays(int n, int k) {  
    if(n == 1) return k;  
    if(n == 2) return k * k;  
    if(k == 1) return 0;  
  
    int same = k;  
    int diff = k * (k - 1);  
  
    for(int i=3; i<=n; i++){  
        int newSame = diff;  
        int newDiff = (same + diff) * (k - 1);  
  
        same = newSame; diff = newDiff;  
    }  
  
    return same + diff;  
}
```



0.7 x

~~# Ugly No - 11~~

1, 2, 3, 4, 5, 6, 8, 10



Recursion
DFS or BFS?

Overlapping subproblems (DP)

$2 \times 3 = 3 \times 2$
Optimal substructure (Faith)

Bigger Ugly = $\min_{i=2} \text{Ugly} * (2/3/5)$

~~ptr 2~~
~~5th - 6th~~

4th
9th 3
4th 3

~~ptr 5~~
~~2nd 3rd~~
~~1st 5~~

①, ②, ③, ④, ⑤, ⑥, ⑦, ⑧, ⑨, ⑩
1st 2nd 3rd 4th 5th 6th 7th 8th 9th 10th

```
// O(N) Time, O(N) Space
public int nthUglyNumber(int n) {
    if(n == 1) return 1;

    // Pointers Pointing to Indices not Values
    int ptr2 = 0, ptr3 = 0, ptr5 = 0;

    ArrayList<Integer> ugly = new ArrayList<>(); // visited array
    ugly.add(1); // to add the 1st ugly no at index 0

    for(int i=1; i<n; i++){
        int a = ugly.get(ptr2) * 2;
        int b = ugly.get(ptr3) * 3;
        int c = ugly.get(ptr5) * 5;

        int min = Math.min(a, Math.min(b, c));
        ugly.add(min);

        if(min == a) ptr2++;
        if(min == b) ptr3++;
        if(min == c) ptr5++;
    }

    return ugly.get(n - 1);
}
```



0.7 x

```
// O(N * Log N) Time, O(N) Space
public int nthUglyNumber(int n) {
    if(n == 1) return 1;

    PriorityQueue<Long> q = new PriorityQueue<>();
    q.add(1L);
    HashSet<Long> vis = new HashSet<>();

    int idx = 0;
    while(q.size() > 0){
        long min = q.remove();
        if(vis.contains(min) == true)
            continue;

        idx++;
        if(idx == n) return (int)min;

        vis.add(min);
        q.add(min * 2L);
        q.add(min * 3L);
        q.add(min * 5L);
    }

    return 1;
}
```

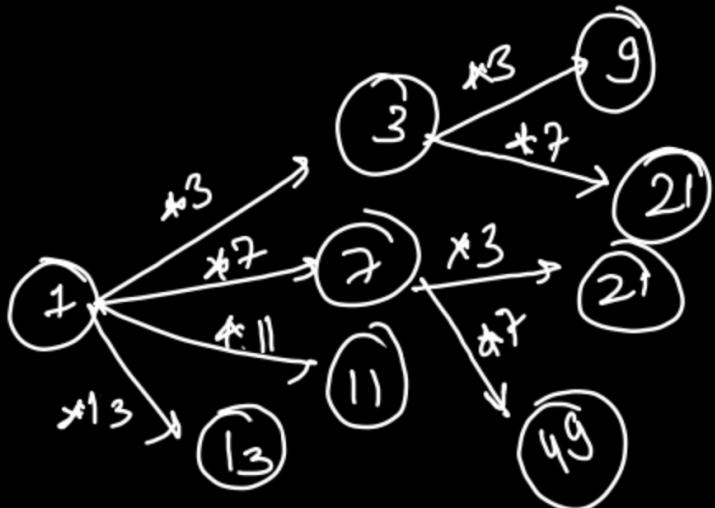


0.5 x

~~Super Ugly No~~

Q (Input) prime = { 3, 7, 11, 13 } Q ugly ?
 $N=7$

(Extra)
= O(F) prime = { 4, 3, 2, 2 }



RCS (Extra) = O(N)
1st, 2nd, 3rd, 4th, 5th, 6th, 7th

space
comp

$O(Nl+k)$
l =
k = pointers index

```
// O(N * K) Time, O(N + K) Space
public int nthSuperUglyNumber(int n, int[] primes) {
    int[] ptr = new int[primes.length];

    ArrayList<Integer> ugly = new ArrayList<>();
    ugly.add(1); // Add 1st Ugly No at Index 0

    for(int i=1; i<n; i++){
        // Finding the next Smallest Ugly No
        int min = Integer.MAX_VALUE;
        for(int j=0; j<primes.length; j++)
            min = Math.min(min, ugly.get(ptr[j]) * primes[j]);

        ugly.add(min);

        // Updating All Pointers Pointing to Min
        for(int j=0; j<primes.length; j++)
            if(ugly.get(ptr[j]) * primes[j] == min)
                ptr[j]++;
    }

    return ugly.get(n - 1);
}
```



0.5 x