

# # Recursion & Backtracking

## Today's Class { lecture - 7 }

[N Queen - Combinations](#)

[N Queen - Permutations](#)

[N-Queens - I \(IB\)](#)

[N-Queens - II](#)

[Using Backtracking](#)

[Code](#)

[Using Branch & Bound](#)

[Using Bit Manipulation](#)

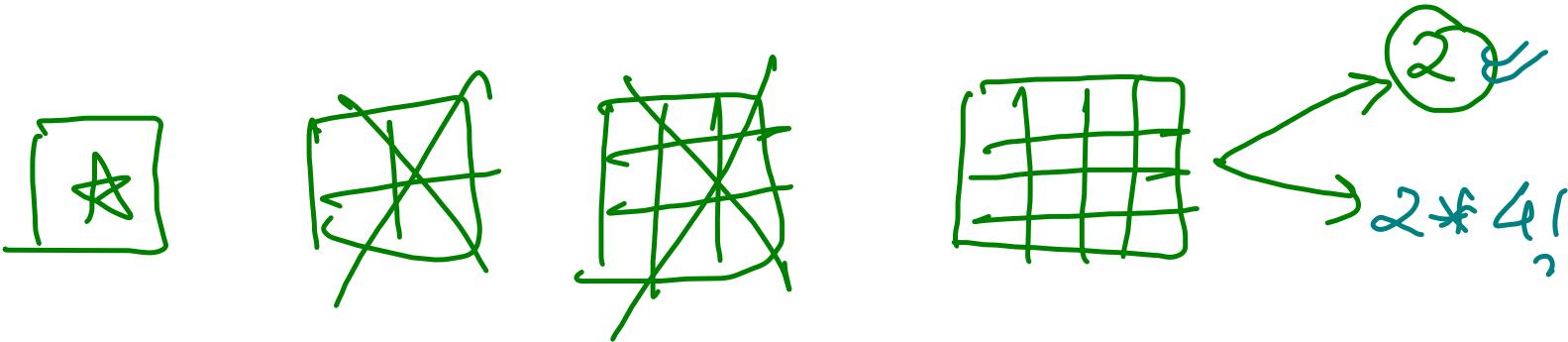
[N Knight](#)

[Knight's Tour](#)

[GfG Article](#)

[Knight's Tour](#)

[Code](#)



	$q_1$		
$q_2$			*
*			$q_3$
		$q_4$	*

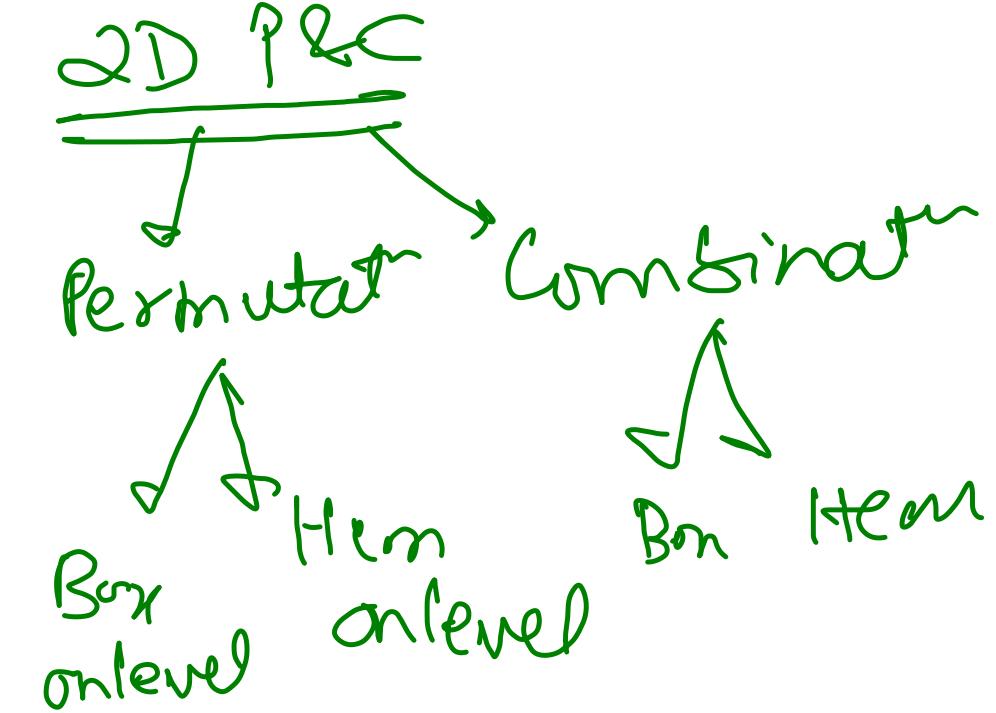
$$N^2 = 4 \times 4$$

$N = 4$  queens

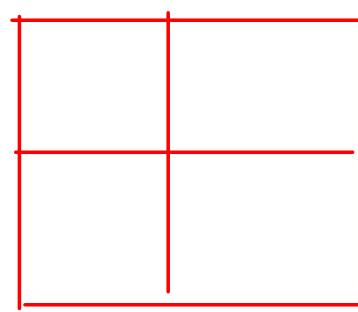
Combination

$$\{q_1, q_2, q_3, q_4\}$$

constraint  $\rightarrow$  No queen pair should  
kill each other.



# ① N Queen - Combination



$\{q_1, q_2\}$

$$2^{N^2}$$

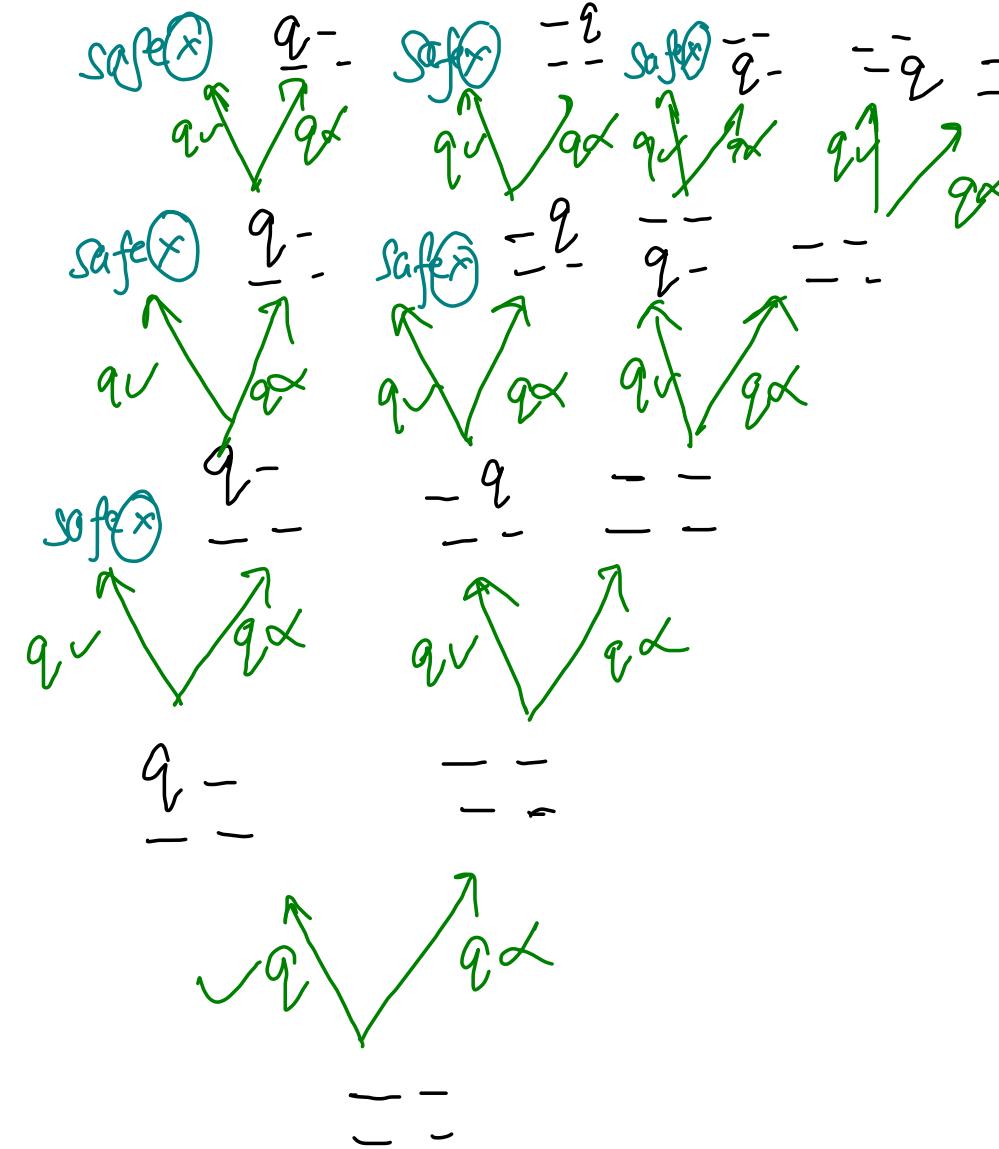
$N = 2 \Rightarrow$   
# <sup>0</sup> valid  
configurations

(1, 1)

(1, 0)

(0, 1)

(0, 0)



(call(s))<sup>height</sup>

+

(pre + post) \* height

$(2^n)$

+

$n * (n^2) = O(2^n + n^3)$

$N = 9$   
 $\frac{81}{2} + 3^3$

# #IsQueenSafe

	c0	c1	c2	c3
r0	dD	dE	dG	dH
r1	dC	dD	dF	dG
r2	dS	dC	dD	dF
r3	dA	dB	dC	dD

① Left Diagonal

	c0	c1	c2	c3
r0	dI	dII	dIII	dIV
r1	dIX	dIII	dIV	dV
r2	dIII	dIV	dV	dVI
r3	dV	dVI	dVII	dVIII

② Right Diagonal

	c0	c1	c2	c3
r0	r0	r0	r0	r0
r1	r1	r1	r1	r1
r2	r2	r2	r2	r2
r3	r3	r3	r3	r3

③ Row

	c0	c1	c2	c3
r0	c0	c1	c2	c3
r1	c0	c1	c2	c3
r2	c0	c1	c2	c3
r3	c0	c1	c2	c3

④ column

Total left or right diagonals  
 $= 2n - 1$

Total rows or columns =  $n$

$O(N)$  time,  $O(1)$  space

```
public static boolean isQueenSafe(int r, int c, boolean[][] vis){  
    // row (left)  
    for(int j=0; j<c; j++){  
        if(vis[r][j] == true)  
            return false;  
    }  
    // col (up)  
    for(int i=0; i<r; i++){  
        if(vis[i][c] == true)  
            return false;  
    }  
  
    // left diagonal (up)  
    int i = r, j = c;  
    while(i >= 0 && j >= 0){  
        if(vis[i][j] == true){  
            return false;  
        }  
        i--; j--;  
    }  
  
    // right diagonal (down)  
    i = r; j = c;  
    while(j < vis.length && i >= 0){  
        if(vis[i][j] == true){  
            return false;  
        }  
        i--; j++;  
    }  
  
    return true;  
}
```

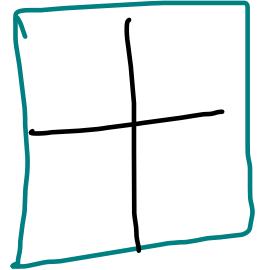
→ Recursion Space:  $\rightarrow O(n^2)$   
→ Output:  $O(n^2)$

```
public static void nqueen(int r, int c, int qpsf, boolean[][] vis) {  
    if (qpsf == vis.length) {  
        List<String> ans = new ArrayList<>();  
        for (int i = 0; i < vis.length; i++) {  
            res.add(ans);  
        }  
        return;  
    }  
  
    if(r == vis.length){  
        return;  
    }  
  
    if (isQueenSafe(r, c, vis)) {  
        vis[r][c] = true;  
  
        if (c == vis.length - 1) {  
            nqueen(r + 1, 0, qpsf + 1, vis);  
        } else {  
            nqueen(r, c + 1, qpsf + 1, vis);  
        }  
  
        vis[r][c] = false;  
    }  
  
    if (c == vis.length - 1) {  
        nqueen(r + 1, 0, qpsf, vis);  
    } else {  
        nqueen(r, c + 1, qpsf, vis);  
    }  
}
```

# N-Queen Permutation

$\{q_1, q_2\}$  Branch level

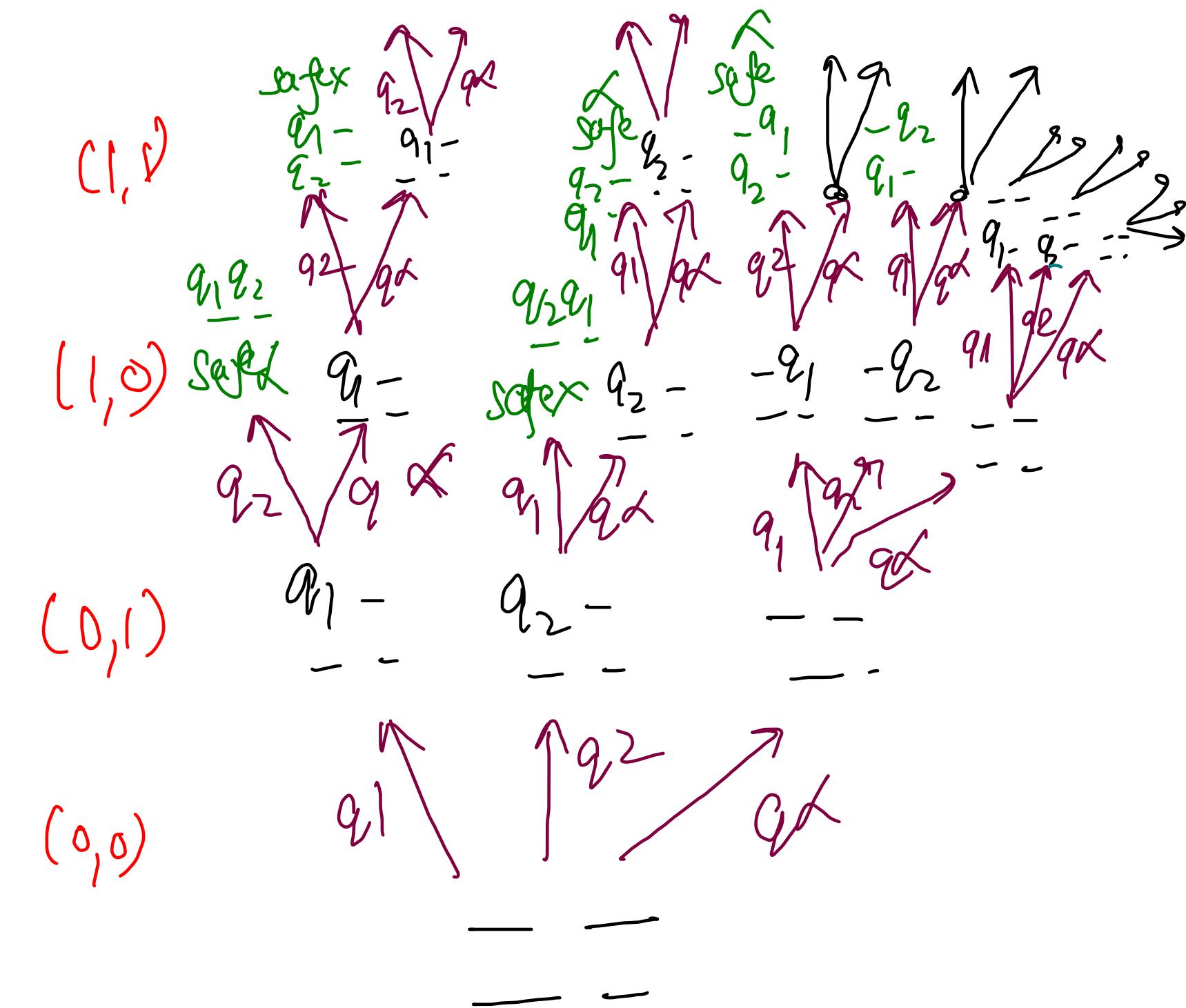
$$2 \times 2 \quad N = 2$$



20 valid configurations

$$N = 2 \Rightarrow \text{Permutations} = 2 \times 2!$$

Visited queens array ~~array~~  $= 2^2 = 4$  = 48



Ban on level approach { Order Different }

```
public static void nqueens(int qpsf, int r, int c, boolean[] queenPlaced, int[][] chess) {
    int n = queenPlaced.length;
    if(qpsf == n){
        for(int i=0; i<n; i++){}
        System.out.println();
        return;
    }

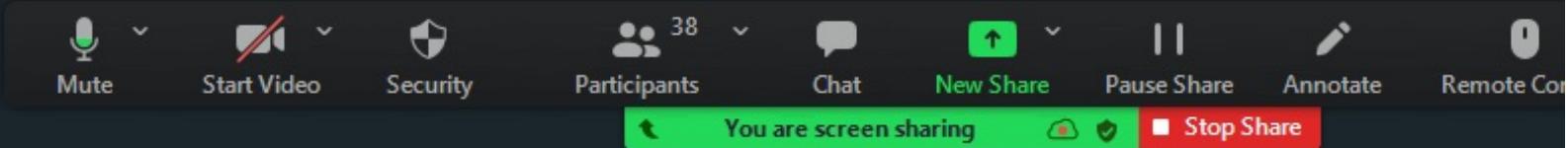
    if(r == n){ return; }

    if(c == n-1){
        nqueens(qpsf, r+1, 0, queenPlaced, chess);
    } else {
        nqueens(qpsf, r, c+1, queenPlaced, chess);
    }

    for(int i=0; i<n; i++){
        if(queenPlaced[i] == false && isQueenSafe(r, c, chess)){
            queenPlaced[i] = true;
            chess[r][c] = (i + 1);

            if(c == n-1){
                nqueens(qpsf+1, r+1, 0, queenPlaced, chess);
            } else {
                nqueens(qpsf+1, r, c+1, queenPlaced, chess);
            }

            chess[r][c] = 0;
            queenPlaced[i] = false;
        }
    }
}
```



A screenshot of a video conferencing interface at the bottom of the slide. It includes controls for Mute, Start Video, Security, Participants (38), Chat, New Share, Pause Share, Annotate, and Remote Control. A message at the bottom says "You are screen sharing".

# #IsQueenSafe

	c0	c1	c2	c3
r0	d <sub>D</sub>	d <sub>E</sub>	d <sub>G</sub>	d <sub>H</sub>
r1	d <sub>C</sub>	d <sub>D</sub>	d <sub>F</sub>	d <sub>G</sub>
r2	d <sub>S</sub>	d <sub>C</sub>	d <sub>D</sub>	d <sub>F</sub>
r3	d <sub>A</sub>	d <sub>B</sub>	d <sub>C</sub>	d <sub>D</sub>

① Left Diagonal

	c0	c1	c2	c3
r0	d <sub>I</sub>	d <sub>II</sub>	d <sub>III</sub>	d <sub>IV</sub>
r1	d <sub>II</sub>	d <sub>III</sub>	d <sub>IV</sub>	d <sub>V</sub>
r2	d <sub>III</sub>	d <sub>IV</sub>	d <sub>V</sub>	d <sub>VI</sub>
r3	d <sub>IV</sub>	d <sub>V</sub>	d <sub>VI</sub>	d <sub>VII</sub>

② Right Diagonal

	c0	c1	c2	c3
r0	r <sub>0</sub>	r <sub>0</sub>	r <sub>0</sub>	r <sub>0</sub>
r1	r <sub>1</sub>	r <sub>1</sub>	r <sub>1</sub>	r <sub>1</sub>
r2	r <sub>2</sub>	r <sub>2</sub>	r <sub>2</sub>	r <sub>2</sub>
r3	r <sub>3</sub>	r <sub>3</sub>	r <sub>3</sub>	r <sub>3</sub>

③ Row

	c0	c1	c2	c3
r0	c <sub>0</sub>	c <sub>1</sub>	c <sub>2</sub>	c <sub>3</sub>
r1	c <sub>0</sub>	c <sub>1</sub>	c <sub>2</sub>	c <sub>3</sub>
r2	c <sub>0</sub>	c <sub>1</sub>	c <sub>2</sub>	c <sub>3</sub>
r3	c <sub>0</sub>	c <sub>1</sub>	c <sub>2</sub>	c <sub>3</sub>

④ column

[T<sub>f</sub> T<sub>f</sub> T<sub>f</sub> T<sub>f</sub> T<sub>f</sub>]

Total left or right

diagonals

$$= 2n - 1$$

Total rows or

columns = n

# # NQueen {Combinat} {Backtracking}

	★	#	
#			★
★			#
	#	★	

4 rows  
4 columns

4 queens

row m  
level  
column as  
option

$$(n)^n + (n) \times n$$



isQueenSafe  
 $\downarrow$   
 $(n) \times n$

$$\Rightarrow O(n^n + n^2)$$

$$N = g \\ g^g + g^2$$

$O(n)$  time,  $O(1)$  space

```
public static boolean isQueenSafe(int r, int c, boolean[][] vis){  
    // row (left)  
    for(int j=0; j<c; j++){  
        if(vis[r][j] == true)  
            return false;  
    }  
    // col (up)  
    for(int i=0; i<r; i++){  
        if(vis[i][c] == true)  
            return false;  
    }  
    // left diagnol (up)  
    int i = r, j = c;  
    while(i >= 0 && j >= 0){  
        if(vis[i][j] == true){  
            return false;  
        }  
        i--; j--;  
    }  
    // right diagnol (down)  
    i = r; j = c;  
    while(j < vis.length && i >= 0){  
        if(vis[i][j] == true){  
            return false;  
        }  
        i--; j++;  
    }  
  
    return true;  
}
```

Recursion call stack  $\rightarrow O(n)$

output :  $O(n^2)$

```
public static void nqueen(int r, boolean[][] vis) {  
    if (r == vis.length) {  
        List<String> ans = new ArrayList<>();  
        for (int i = 0; i < vis.length; i++) {  
            res.add(ans);  
        }  
        return;  
    }  
  
    for(int c=0; c<vis.length; c++){  
        if(isQueenSafe(r, c, vis)){  
            vis[r][c] = true;  
            nqueen(r + 1, vis);  
            vis[r][c] = false;  
        }  
    }  
}
```

### ③ Branch & Bound

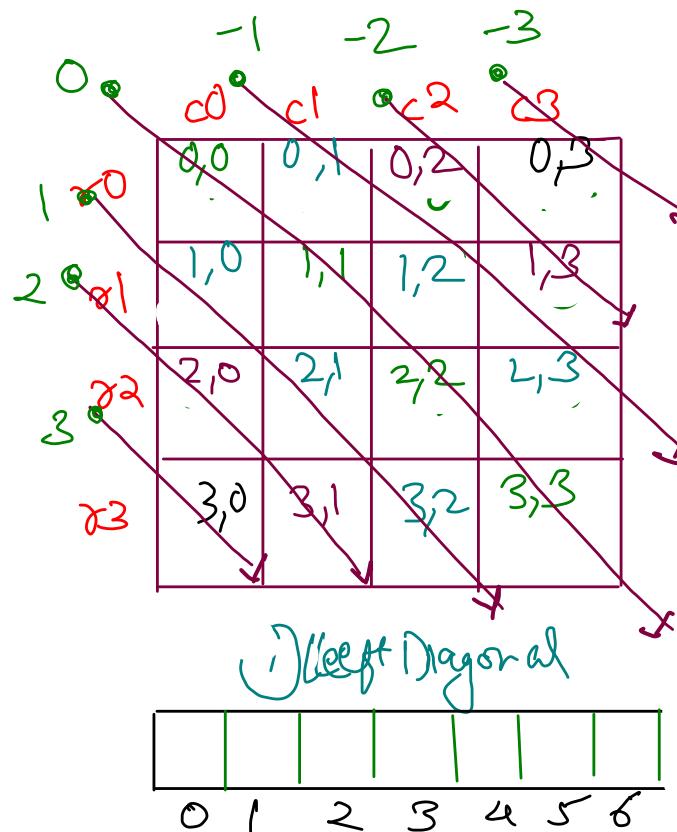
$\{ \text{optimization in } \text{isQueenSafe} \}$   
 in terms of  
 time complexity

	$c_0$	$c_1$	$c_2$	$c_3$
$r_0$	$c_0$	$c_1$	$c_2$	$c_3$
$r_1$	$c_0$	$c_1$	$c_2$	$c_3$
$r_2$	$c_0$	$c_1$	$c_2$	$c_3$
$r_3$	$c_0$	$c_1$	$c_2$	$c_3$



Column  
 $\{ n \}$

BytE

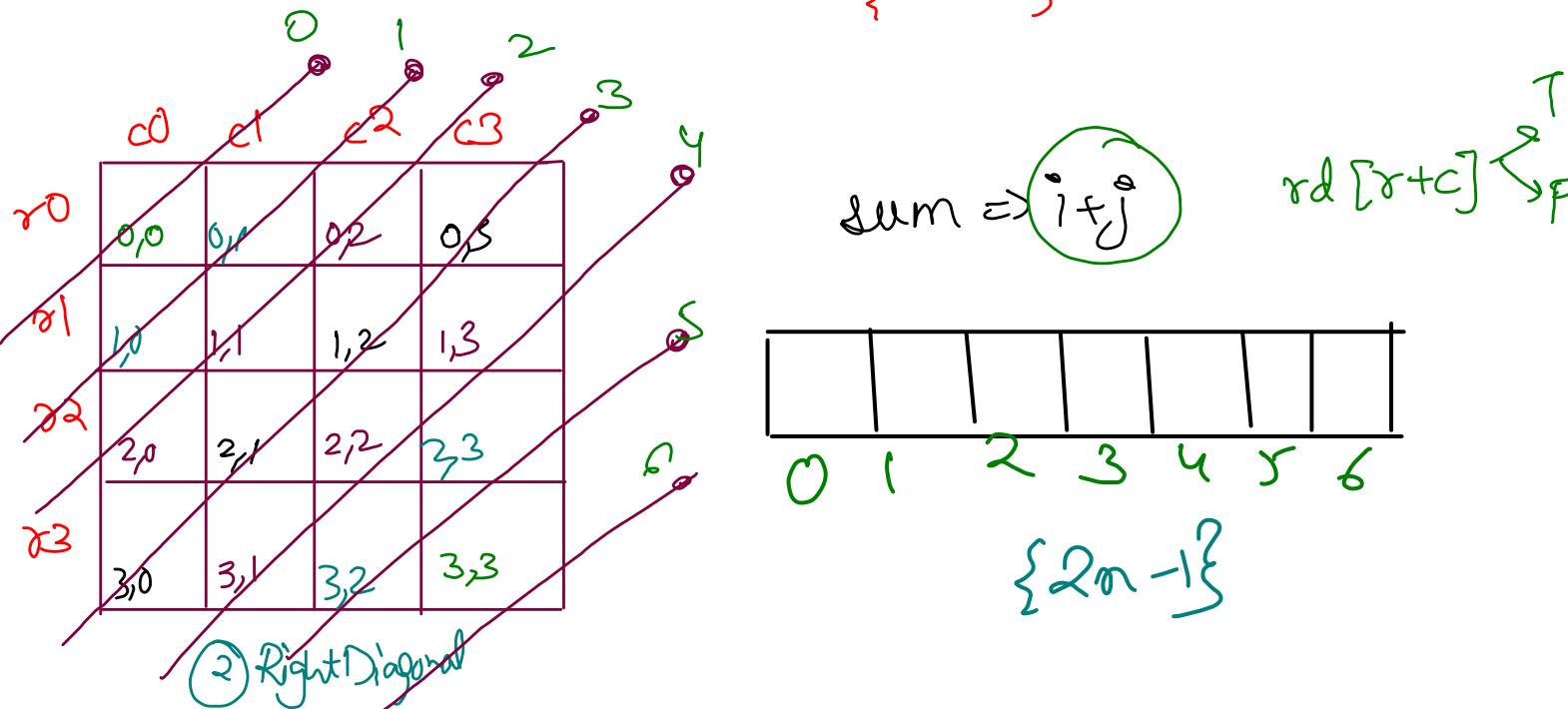


left diagonal for  $r, c$

$$= \text{ld} [ \text{rd} [ r - c + (n-1) ] ]$$

$\frac{-3}{4+3}, \frac{-2}{5+3}, \frac{-1}{6+3}, \frac{0}{7+3}, \frac{1}{8+3}, \frac{2}{9+3}, \frac{3}{10+3}$

$\{ 2n-1 \}$



$\{ 2n-1 \}$

```

public static boolean isQueenSafe(int r, int c, boolean[] ld, boolean[] rd, boolean[] col){
    int n = col.length;
    return ((ld[r - c + n - 1] == true) || (rd[r + c] == true) || (col[c] == true)) ? false : true;
}

```

→ Extra Space

$O(1)$  time  
 $O(n)$  space

```

public static void nqueen(int r, boolean[][] chess, boolean[] ld, boolean[] rd, boolean[] col) {
    int n = chess.length;
    if (r == n) {
        List<String> ans = new ArrayList<>();
        for (int i = 0; i < n; i++) {
            String curr = "";
            for (int j = 0; j < n; j++) {
                if (chess[i][j] == true) {
                    curr = curr + "Q";
                } else {
                    curr = curr + ".";
                }
            }
            ans.add(curr);
        }
        res.add(ans);
    }
    return;
}

for(int c=0; c<n; c++){
    if(isQueenSafe(r, c, ld, rd, col)){
        chess[r][c] = ld[r - c + n - 1] = rd[r + c] = col[c] = true;
        nqueen(r + 1, chess, ld, rd, col);
        chess[r][c] = ld[r - c + n - 1] = rd[r + c] = col[c] = false;
    }
}

```

$\mathcal{O}(n)$  extra space

Recursion call  
Stack :  $\rightarrow O(n)$

$\mathcal{O}(n)$  output

Time  
Comp

$$(\mathcal{G})^n + (\mathcal{R})^n$$

$\uparrow$   
isQueenSafe

$$\Rightarrow 4^n + n$$

## N Queen Combin<sup>n</sup>

Time Complexity

Brute force on level  $\Rightarrow 2^{n^2} + n^3 \approx 120\text{ ms}$

Backtracking  $\Rightarrow 4^n + n^2 \approx 4\text{ ms}$

Branch & Bound  $\Rightarrow 4^n + n \approx 3\text{ ms}$

Space Complexity

Brute force on level  $\Rightarrow O(n^2)$  Recursion call stack

Backtracking  $\Rightarrow O(n)$  Recursion = ..

Branch & Bound  $\Rightarrow O(N)$  extra Space

$O(N)$  recursion call stack

$O(n^2)$  output  
(chess)

Bit Manipulation  $\Rightarrow O(1)$  extra space

$O(N)$  Recursion call stack

## ④ Bit Manipulation

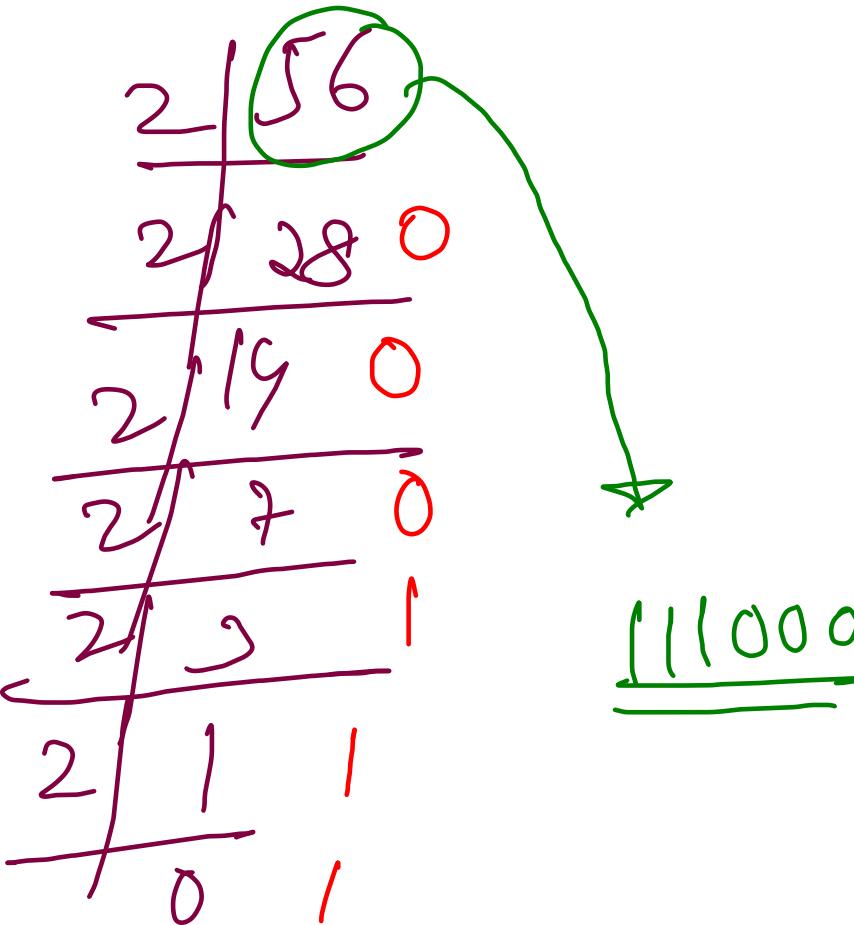
Replacing

boolean visited array by

$\downarrow$   
 $O(n)$  bytes

T/F

bitset  
 $\overbrace{\quad}$   
 $\downarrow$   
 $n$  bits  
 $\uparrow$  {0/1}  
 $\downarrow$   
 $O(1)$  space



integer: 4 bytes : 32 bits  
 $\overbrace{\quad}$   
↳ ld  
integer rd

integer<sup>e-1</sup>  
↳ cool

```
public static boolean isQueenSafe(int r, int c, int n, BitSet ld, BitSet rd, BitSet col){
    return ((ld.get(r - c + n - 1) == true) || (rd.get(r + c) == true) || (col.get(c) == true)) ? false : true;
}

public static void nqueen(int r, boolean[][] chess, BitSet ld, BitSet rd, BitSet col) {
    int n = chess.length;
    if (r == n) {
        List<String> ans = new ArrayList<>();
        for (int i = 0; i < n; i++) {ans}
        res.add(ans);
        return;
    }

    for(int c=0; c<n; c++){
        if(isQueenSafe(r, c, n, ld, rd, col)){
            chess[r][c] = true;
            ld.set(r - c + n - 1, true);
            rd.set(r + c, true);
            col.set(c, true);

            nqueen(r + 1, chess, ld, rd, col);

            ld.set(r - c + n - 1, false);
            rd.set(r + c, false);
            col.set(c, false);
            chess[r][c] = false;
        }
    }
}
```

```
bitset<1000000> col,d1,d2;
void solve(int i,int n,int &ans)
{
    if(i == n) {ans++; return;}
    for(int j=0;j<n;j++)
    {
        if(!col[j] && !d1[i-j+n-1] && !d2[i+j])
        {
            col[j] = d1[i-j+n-1] = d2[i+j] = 1;
            solve(i+1,n,ans);
            col[j] = d1[i-j+n-1] = d2[i+j] = 0;
        }
    }
}
class Solution {
public:
    int totalNQueens(int n) {
        if(n <= 1) return n;
        int ans = 0;
        solve(0,n,ans);
        return ans;
    }
};
```

{ N Queen - 11 ?  
using C++ }  
Count

# R&B - level-2 - Lecture 8

## Sudoku & Other Puzzles

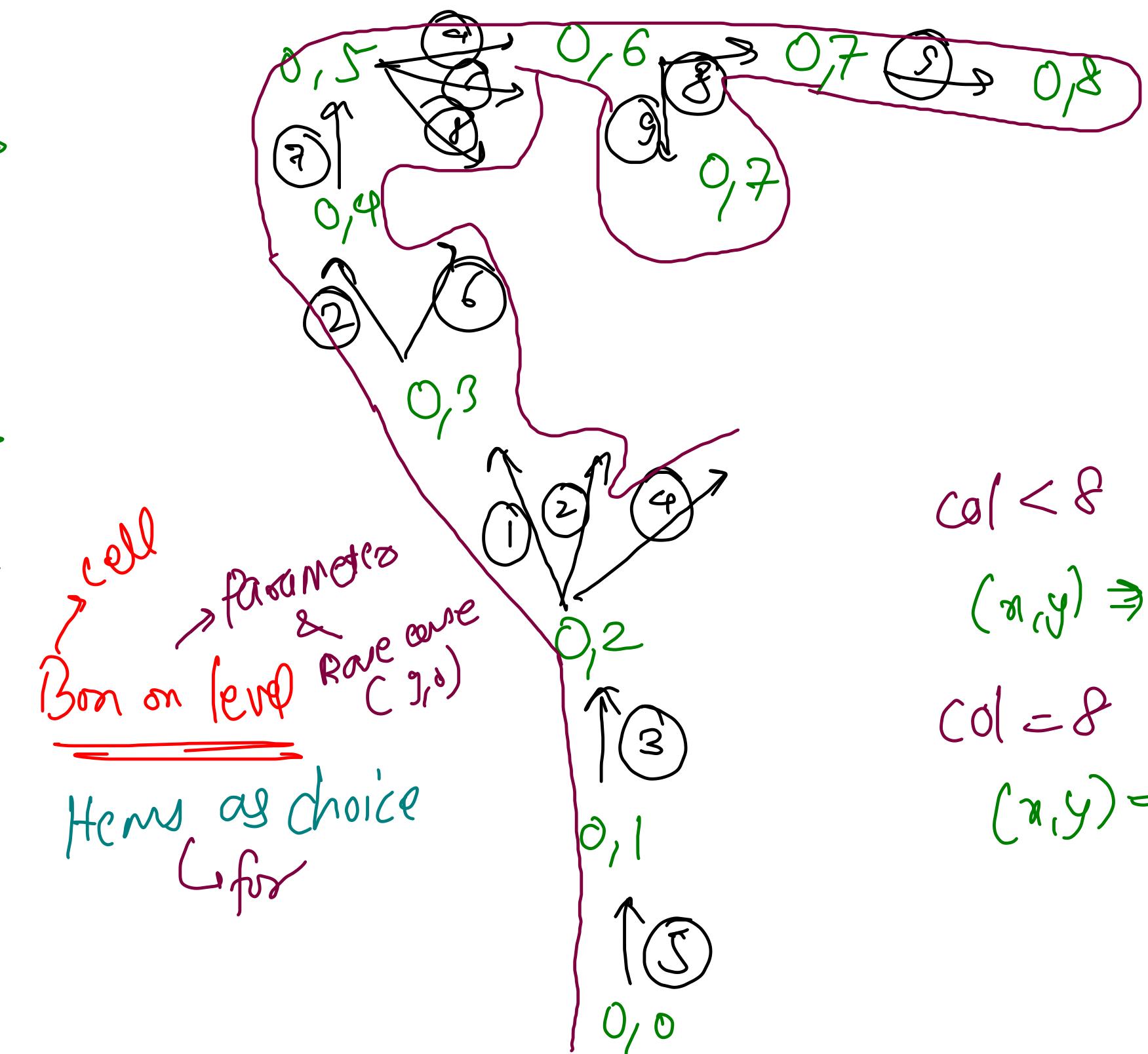
Puzzles	
1	<a href="#"><u>Valid Sudoku (IB)</u></a>
1	<a href="#"><u>Sudoku Solver (IB)</u></a>
1	<a href="#"><u>Using Backtracking</u></a>
1	<a href="#"><u>Using Bit Manipulation</u></a>
2	<a href="#"><u>Crossword Puzzle</u></a>
3	<a href="#"><u>Cryptarithmetic Puzzle</u></a>
3	<a href="#"><u>Cryptarithmetic Puzzle</u></a>
4	<a href="#"><u>Max Score</u></a>
4	<a href="#"><u>Max Score</u></a>
5	<a href="#"><u>N Knight</u></a> , <a href="#"><u>Knight's Tour</u></a>

	0	1	2	3	4	5	6	7	8
0	5	3	4	6	7	3	9	1	2
1	6		1	9	5				
2	9	8			6				
3	8		6			3			
4	4	8	3				1		
5	7		2		6				
6	6			2	8				
7		4	1	9			5		
8		8			7	9			

① In each row, we should have 1-9  
 $\hookrightarrow$  9 columns

② In each column, we should have 1-9  
 $\hookrightarrow$  9 rows

③ In each  $3 \times 3$  matrix, we should have 1-9



	0	1	2	3	4	5	6	7	8
0	5	3		7		1			
1	6			1	9	5			
2		9	8				6		
3	8			6		1		3	
4			8		3				1
5	7			2				6	
6		6		1	2		8		
7				4	1	9			5
8							7	9	

f top-left

val  
 0, 1, 2  $\Rightarrow 0 = 0 \times 3 - (\text{val}/3) * 3$   
 3, 4, 5  $\Rightarrow 3 = 1 \times 3 - (\text{val}/3) * 3$   
 6, 7, 8  $\Rightarrow 6 = 2 \times 3 - (\text{val}/3) * 3$

8, 4  $\Rightarrow$  6, 3

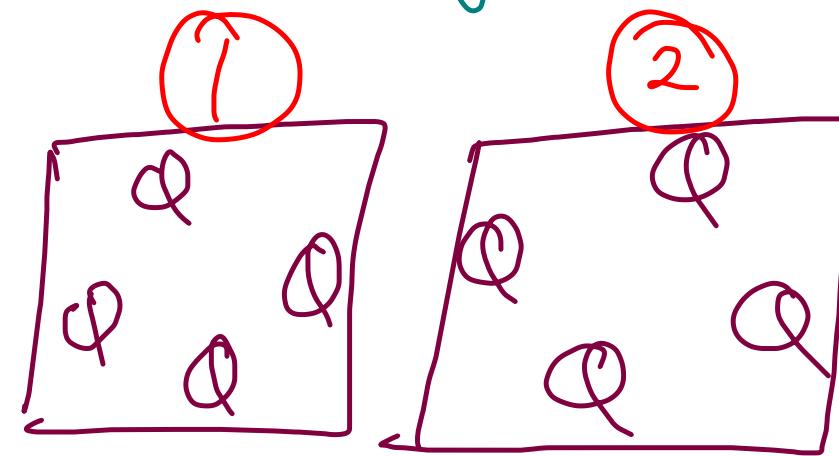
Diagram showing connections between cells:

- Cells (6,3), (6,4), (6,5) are connected to (6,3).
- Cells (7,3), (7,4), (7,5) are connected to (7,3).
- Cells (8,3), (8,4), (8,5) are connected to (8,3).
- Cell (7,3) is connected to (7,4) and (7,5).
- Cell (8,3) is connected to (8,4) and (8,5).
- Cell (6,3) is connected to (7,3) and (8,3).
- Cell (7,3) is connected to (6,3), (7,4), and (7,5).
- Cell (8,3) is connected to (6,3), (8,4), and (8,5).
- Cell (6,3) is connected to (6,4) and (6,5).
- Cell (7,3) is connected to (7,4) and (7,5).
- Cell (8,3) is connected to (8,4) and (8,5).
- Cell (6,3) is connected to (6,3).
- Cell (7,3) is connected to (7,3).
- Cell (8,3) is connected to (8,3).

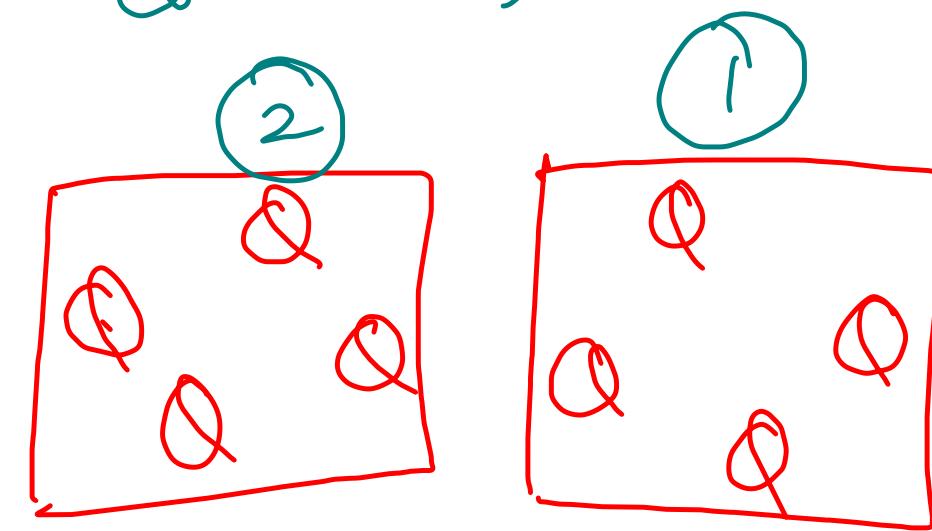
rw + 2

nqueen

row by row



column by column



## Max Score

```
Input: words = ["dog", "cat", "dad", "good"], letters =  
["a", "a", "c", "d", "d", "g", "o", "o"], score =  
[1, 0, 9, 5, 0, 0, 3, 0, 0, 0, 0, 0, 2, 0, 0, 0, 0, 0, 0, 0, 0, 0]  
Output: 23
```

a-2 / 7 / 2

c-1

d-3 / 0 / 2 / 3

g-1 / 0 / 0 /

o-2 / 0 / 1 / 1

"dog" ✓ "dad" ✓ "good" ✓ "cat" ✓ "dad" ✓ "dog" ✓

"dad"

"good"

s3, o, freq  
curr(score)

```
public int maxScore(int idx, String[] words, int[] freq, int[] score){  
    if(idx == words.length) return 0;  
  
    // yes  
    int currScore = 0;  
    boolean flag = true;  
    for(char ch: words[idx].toCharArray()){  
        freq[ch - 'a']--;  
        if(freq[ch - 'a'] < 0){  
            flag = false;  
        }  
  
        currScore += score[ch - 'a'];  
    }  
  
    if(flag == true){  
        currScore += maxScore(idx + 1, words, freq, score);  
    } else {  
        currScore = 0;  
    }  
  
    // backtrack  
    for(char ch: words[idx].toCharArray()){  
        freq[ch - 'a']++;  
    }  
  
    // no  
    return Math.max(currScore, maxScore(idx + 1, words, freq, score));  
}
```

## Move Zeros to End

{ 0, 1, 3, 12 }  
3, 12 }

→ { 1, 3, 12, 0, 0 }

$f(0,0) \rightarrow f(1,0) \rightarrow f(2,1) \rightarrow f(3,1) \rightarrow f(4,2) \rightarrow f(5,3)$

```
public static void moveZeroes(int idx, int countOfNonZero, int[] nums) {
    if(idx == nums.length){
        int countOfZeros = nums.length - countOfNonZero;
        for(int i=0; i<countOfZeros; i++){
            nums[nums.length - 1 - i] = 0;
        }
        return;
    }

    if(nums[idx] != 0){
        nums[countOfNonZero] = nums[idx];
        moveZeroes(idx + 1, countOfNonZero + 1, nums);
    }
    else {
        moveZeroes(idx + 1, countOfNonZero, nums);
    }
}
```

## Today's Questions

- ① Crossword
- ② Word Search - (I, II, III)
- ③ Cryptarithmefic
- ④ Josephus

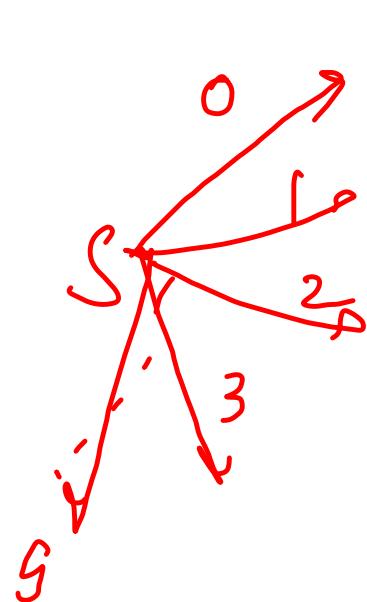
# Cryptarithm

$$\text{"SEND"} + \text{"MORE"} = \text{"MONEY"}$$

1:1 mapping  
 ↓      ↓  
 char digit  
 (A-Z) (0-9)

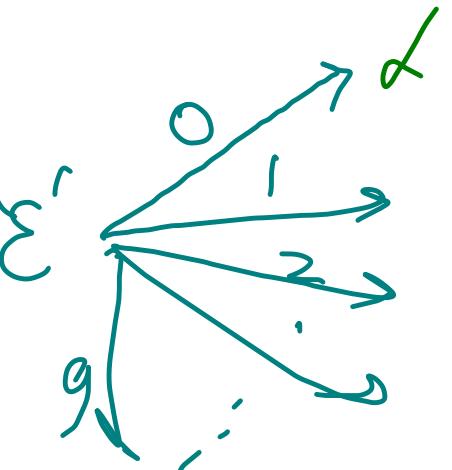
unique  
char  
on  
level

{ 'S', 'E', 'N', 'D', 'M', 'O', 'R', 'Y' }  
 base are  
 ↗ equivalence check



$$'S' \rightarrow 0$$

$$0 \rightarrow 'S'$$



$$'S' \rightarrow 0$$

$$0 \rightarrow 'E'$$

$$0 \rightarrow 'S'$$

$$1 \rightarrow 'E'$$

$$2 \rightarrow 'N'$$

$$3 \rightarrow 'D'$$

$$4 \rightarrow 'Y'$$

$$5 \rightarrow 'M'$$

$$6 \rightarrow 'O'$$

$$7 \rightarrow 'R'$$

$$8 \rightarrow 'L'$$

$$9 \rightarrow 'I'$$

$$'S' \rightarrow 0$$

$$'E' \rightarrow 1$$

$$'N' \rightarrow 2$$

$$'D' \rightarrow 3$$

$$'Y' \rightarrow 4$$

$$'M' \rightarrow 5$$

$$'O' \rightarrow 6$$

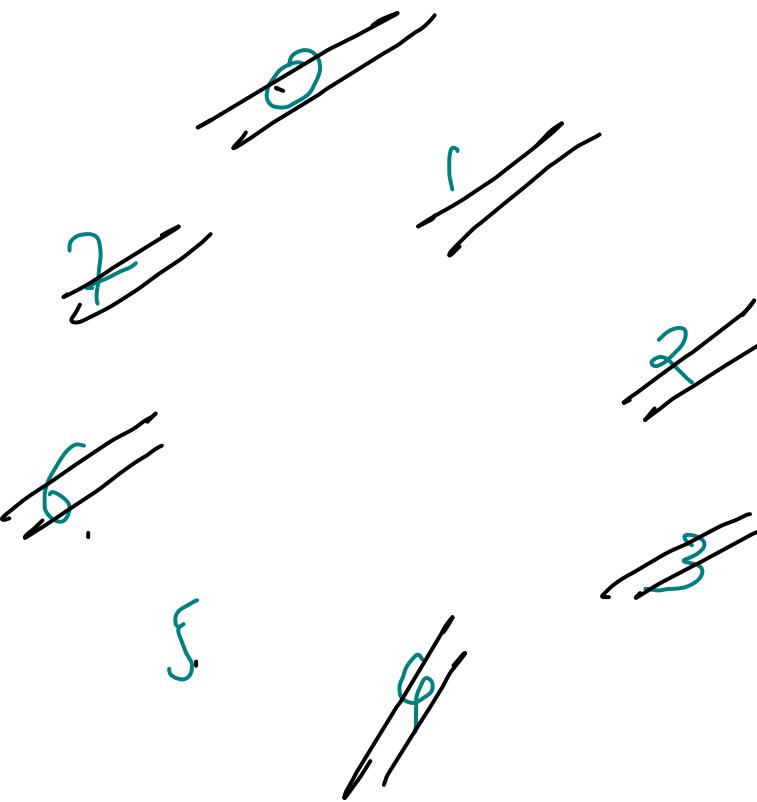
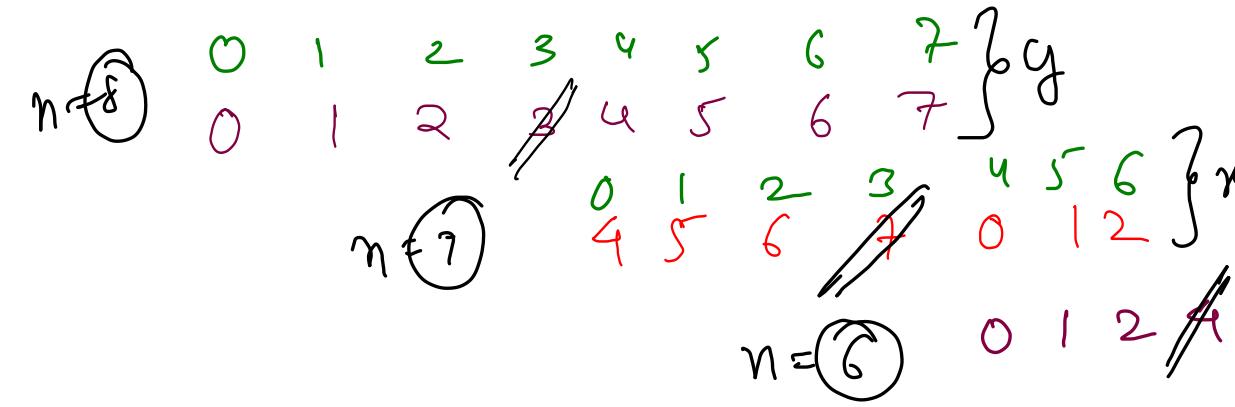
$$'R' \rightarrow 7$$

$$'L' \rightarrow 8$$

$$'I' \rightarrow 9$$

Josephus

$k=4$



$$\begin{cases} \{0, 1, 2, 3, 4, 5, 6\} x \\ \{4, 5, 6, 7, 0, 1, 2\} y \end{cases}$$

$n=8$   
 $k=4$

actual value  
 $y = (x + k) / n$

Meeky Exp

Shifting on  $\odot$   
 rotation on  $\odot$

$$(\text{josephu}(n-1, k) + k) \cdot f(n)$$

```

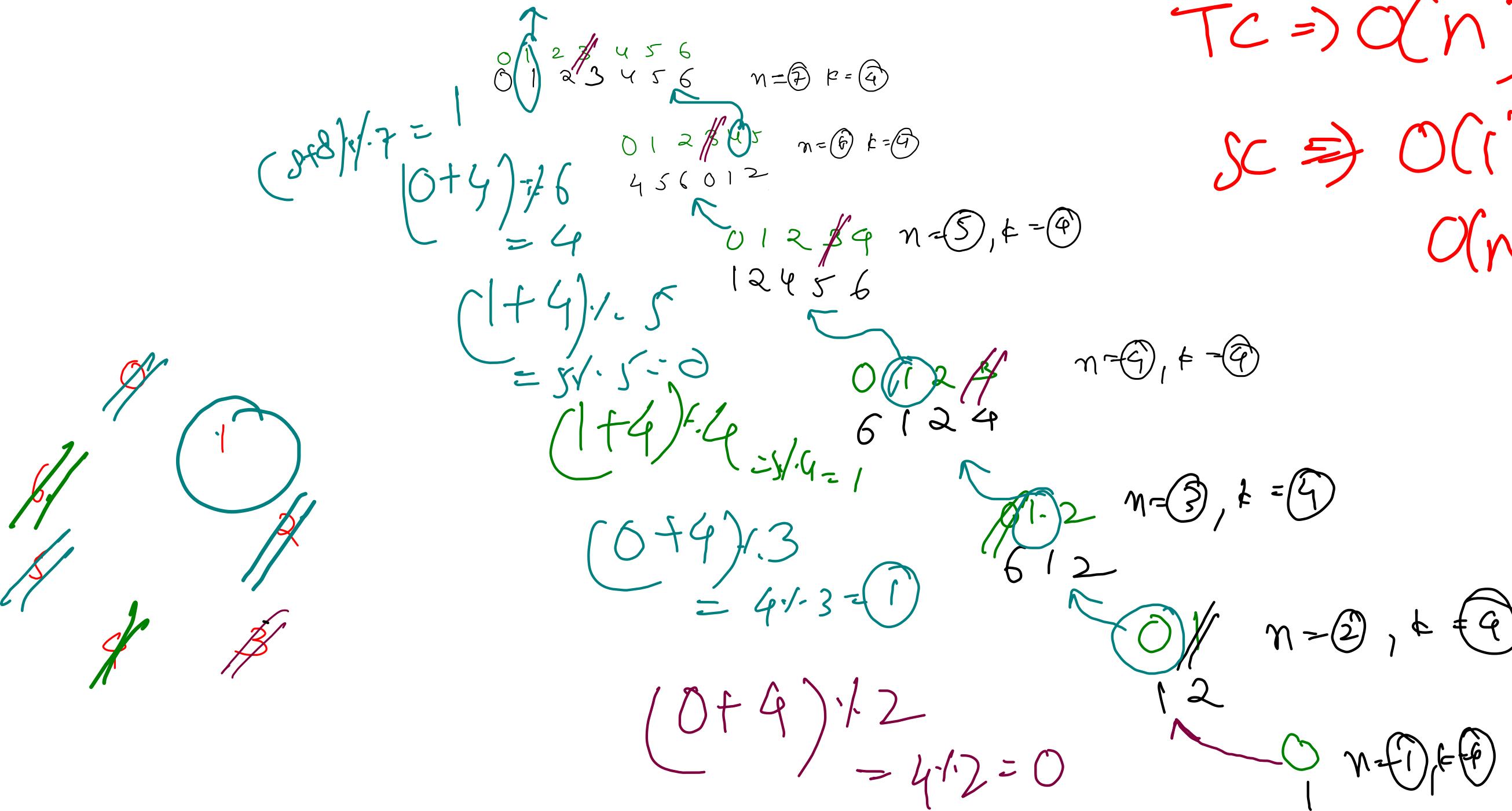
public static int solution(int n, int k){
    if(n == 1) return 0; // return index
    return (solution(n - 1, k) + k) % n;
}

```

$Tc \Rightarrow O(n)$

$Sc \Rightarrow O(1)$  extra space

$O(n)$  recursion



```

public static int solution(int n, int k){
    Queue<Integer> q = new ArrayDeque<>();
    for(int i=0; i<n; i++){
        q.add(i);
    }
    int count = 0;
    while(q.size() > 1){
        int val = q.remove();
        count++;

        if(count == k){
            count = 0;
        } else {
            q.add(val);
        }
    }
    return q.peek();
}

```

constructive  
algorithm

Queue { Brute force }

$O(n \times k)$   $\Rightarrow$  Time

$O(n)$   $\Rightarrow$  Queue  
Space

# henicographical Nos

1 ↙

10 100 101 102 103 104 ... 109

11 110 111 112 113 ... 119

12 120 121 122 ... 129

13 191 192 ... 199

$$N = 999$$

2 ↙

20 200 201 202 ... 209

21 210 211 ... 219

22 220 221 ... 229

23 230 231 ... 239

24 240 241 ... 249

25 250 251 ... 259

26 260 261 ... 269

27 270 271 ... 279

28 280 281 ... 289

29 290 291 ... 299

3

30 → 301 ... 309

31 → 310 ... 319

32 -

⋮

39 → 390 ... 399

4

40, 41, 42, ... 49

5

50, ... 59

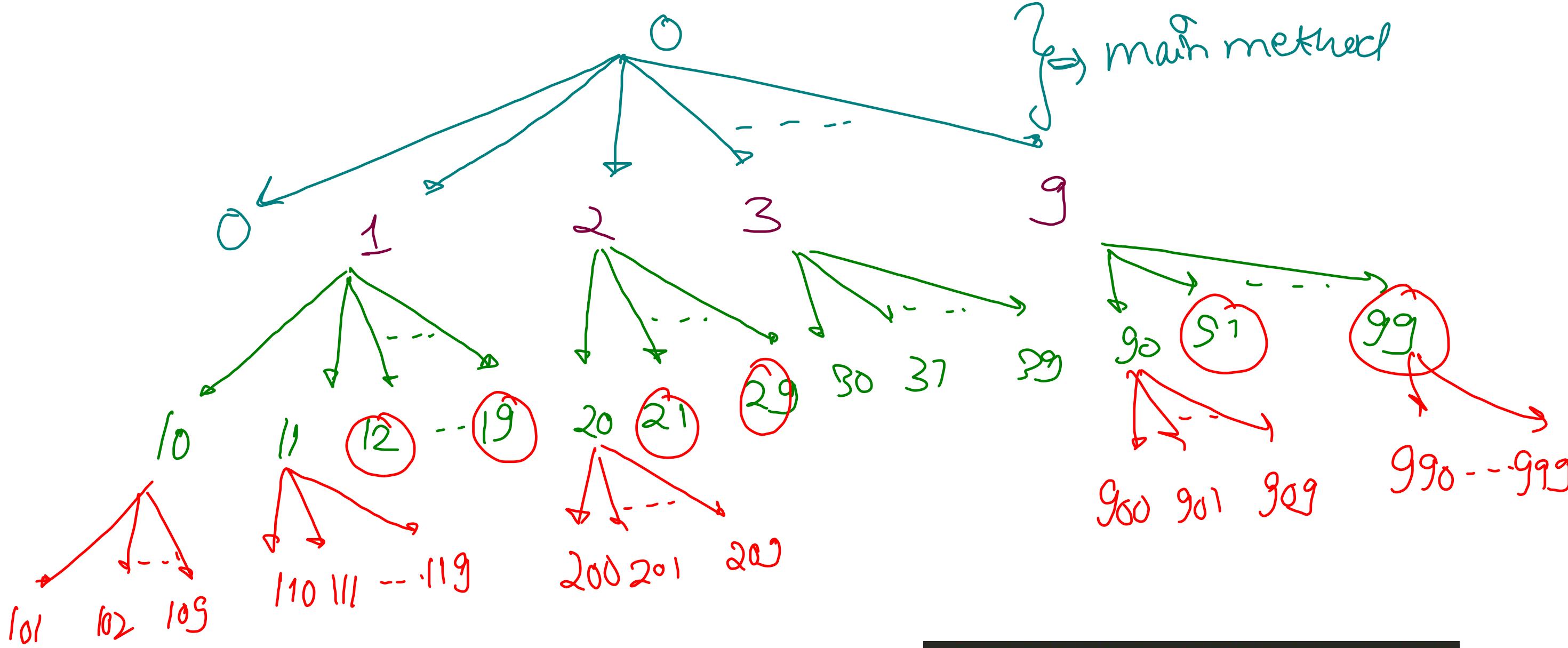
6

7

8

9

90, 91, 92, ... 99  
900, 901, ... 909  
990-999



```

public void lexicalOrder(int input, int n){
    if(input > n) return;

    res.add(input);
    for(int i=0; i<10; i++){
        lexicalOrder(input * 10 + i, n);
    }
}

```

# R&B - Level 2 - Lecture 10

- ① Crossword
- ② Word search - II
- Word Search - III { Homework }
- ③ Friends Pairings
- ④ Largest No in k swaps
- ⑤ Restore IP Address

More questions to practice

- Paranthesis Addition Ways
- Stepping Numbers
- Kth Symbol in Grammar
- Split into Fibonacci Sequence
- Distinct Subsets Permutations

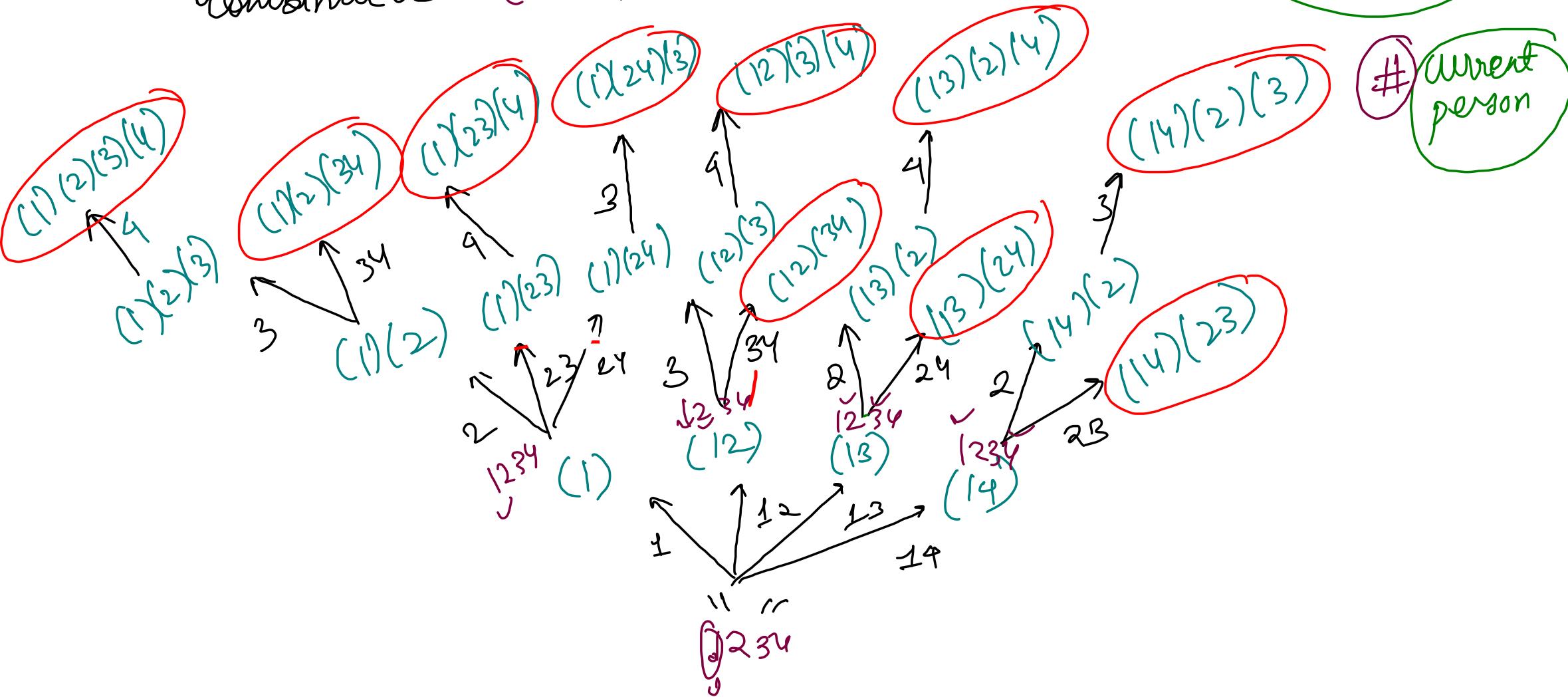
N knights

## Friends Pairing

combinations

$$n = 4$$

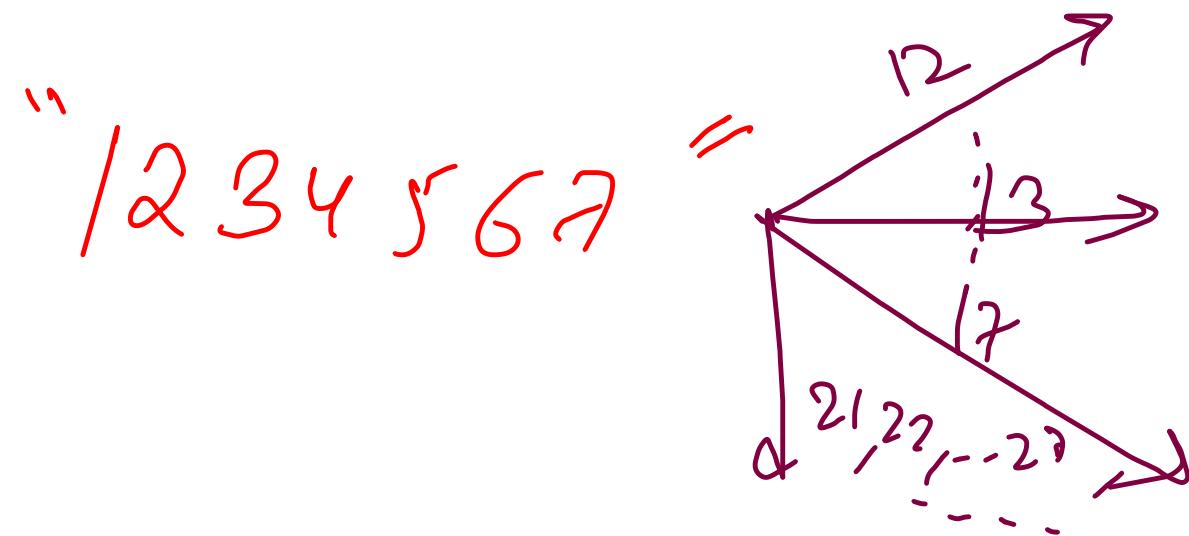
$$\{1, 2, 3, 4\}$$



# visited array

# current person

Largest No in K swaps



2134567

32145(7)

7234561

13  
12

14  
17

23  
24  
21

7634521

7654321

$k=3$

level ①  
swap = ①

level ②  
swap ②

level ③  
swap ③

# Word Search - 11

o	a	a	n
e	t	a	e
i	h	k	r
i	f	l	v

{ soan, geo, eat, rain }

```
public static void solution(int totalFriends, boolean[] friendInGroup, String groupsSofar) {
    int i = 0;
    for(int j=1; j<=totalFriends; j++){
        if(friendInGroup[j] == false){
            i = j;
            break;
        }
    }

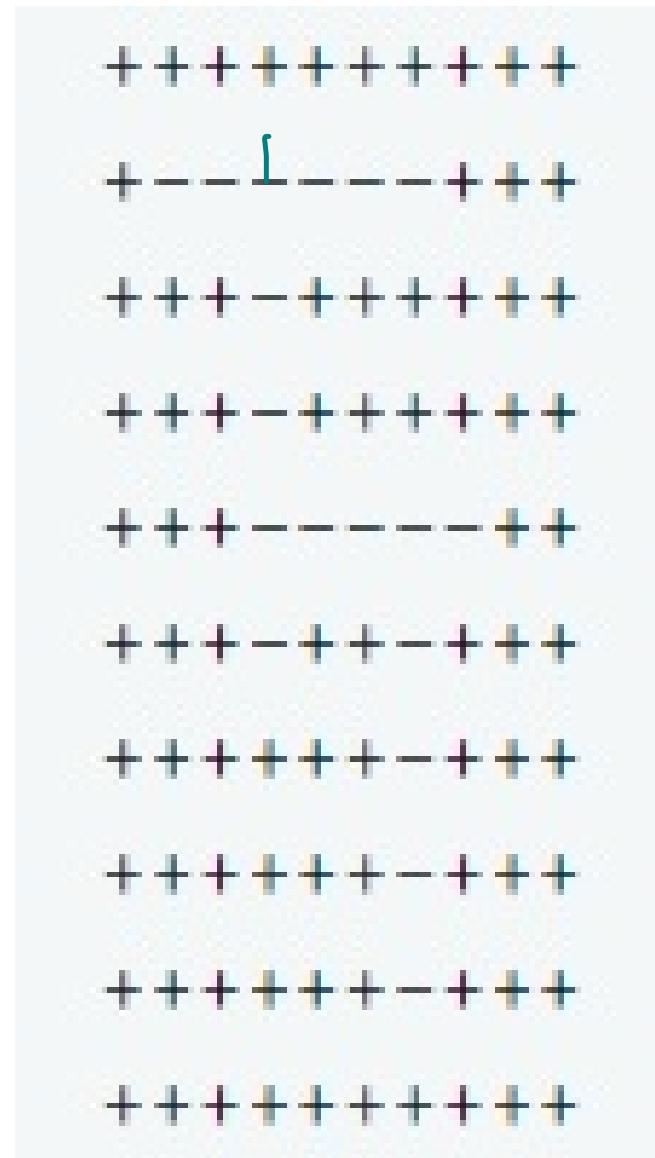
    if(i == 0){}

    // Group of 1 Member
    friendInGroup[i] = true;
    solution(totalFriends, friendInGroup, groupsSofar + "(" + i + ") ");

    for(int j=1; j<=totalFriends; j++){
        if(friendInGroup[j] == false){
            friendInGroup[j] = true;
            solution(totalFriends, friendInGroup, groupsSofar + "(" + i + "," + j + ") ");
            friendInGroup[j] = false;
        }
    }

    friendInGroup[i] = false;
}
```

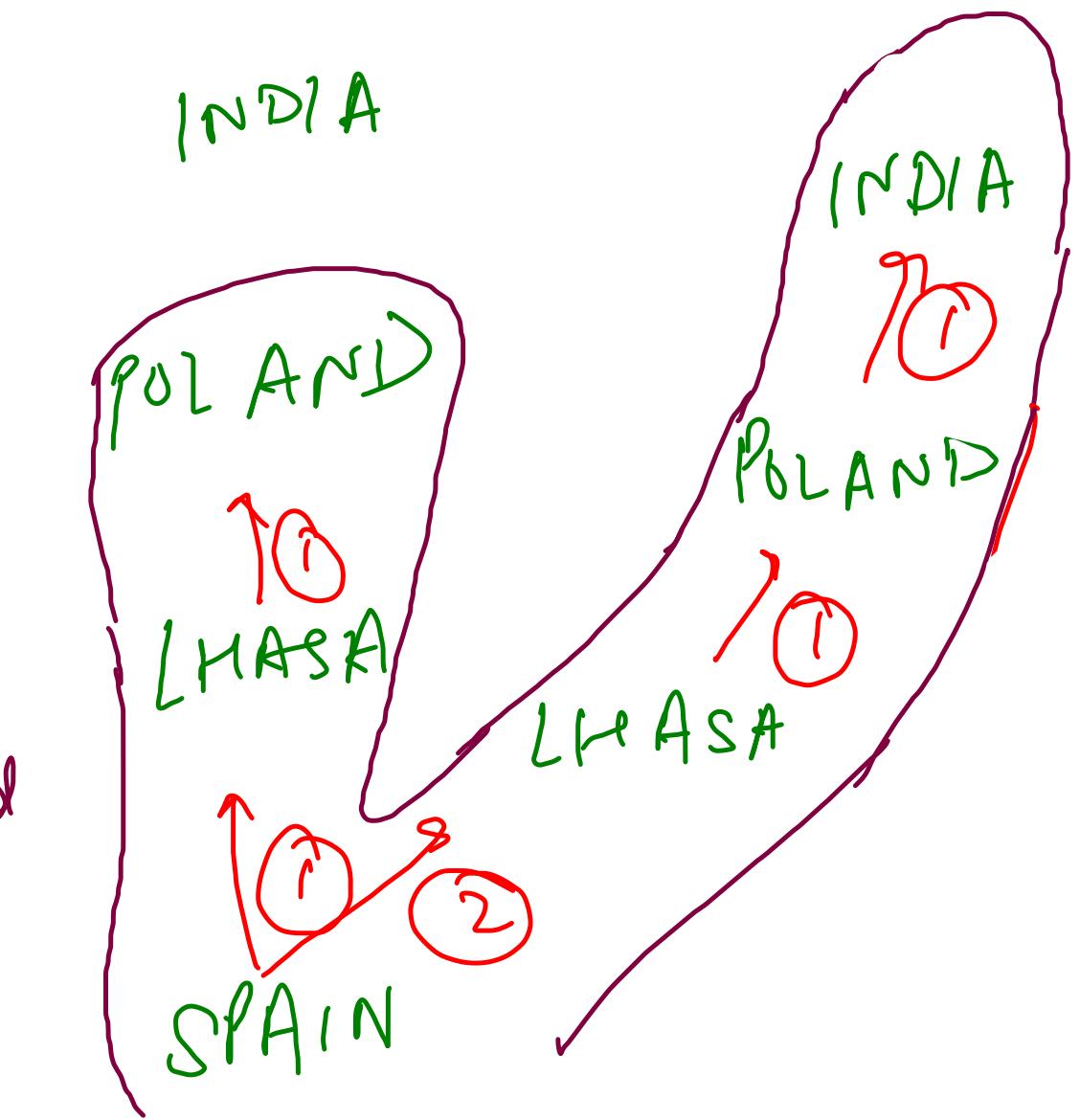
## Crossword



[POLAND, LHASA, SPAIN, INDIA]:

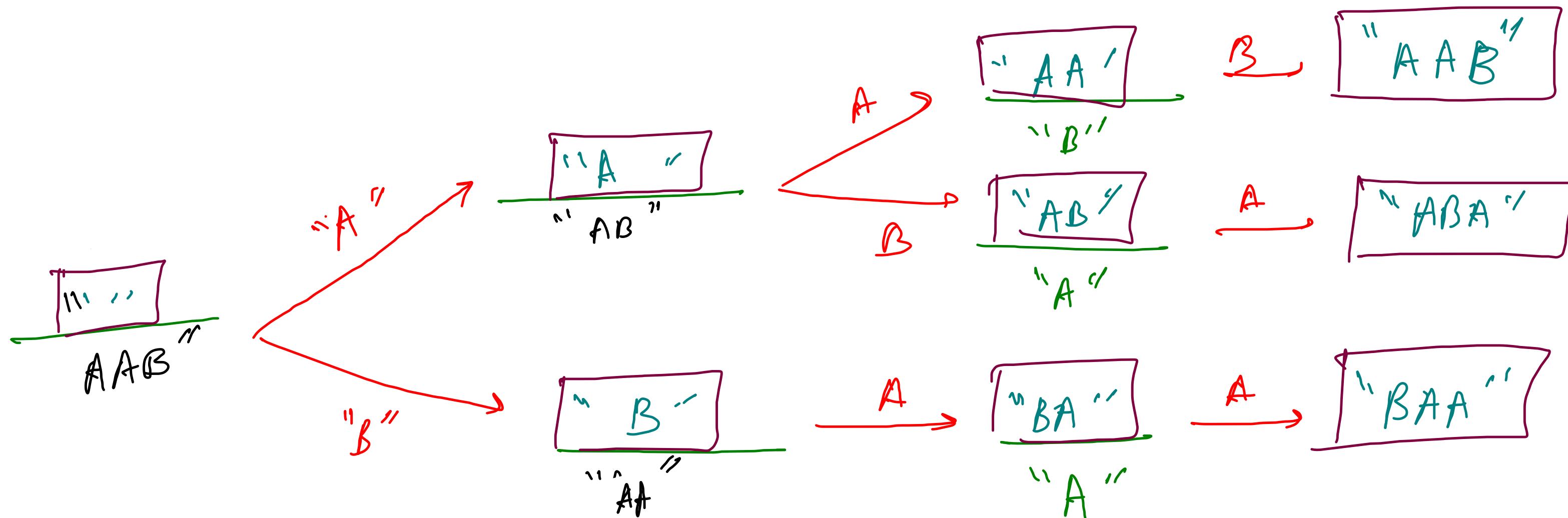
- ① Completely fill +, out of bound

- ② Backtrack only those which are filled by you! !



## Subsets + Permutations

$\{0, 0, 0, 0, 0\}$  letter tile possibilities



Box on level  
Item at  
option

## # Tentative Lecture Plan of DSA - Level 2 }

① Recursion & Backtracking → 10 lectures (Complete)

② Searching & Sorting → 10 lectures

③ Arrays & Strings

- ↳ Greedy Algo → 5 lectures
- ↳ 1D & 2D arrays → 6-7 lectures
- ↳ Strings → 2-3 lectures

④ HashMap & Heap

- ↳ Two pointer / Sliding window → 5 lectures
- ↳ Hashmap → 3-4 lectures
- ↳ Heap → 1-2 lectures

- # Binary Tree & BST → 10 lectures
- # Dynamic Programming → 10-12 lectures
- # linked list → 5 lectures
- # Stack & Queue → 5 lectures
- # Graphs → 7-8 lectures
- # Trees → 3-4 lectures
- # Bit Manipulation → 4-5 lectures

