

Greedy Algorithms

we will start by
9:10

~~WHAT~~

local optimal choice at each step helps
solving global optimum (minima) maxima

APPLICATIONS

- ① Fractional knapsack

- ② Activity Selection / Overlapping Intervals / Meeting Rooms

- ③ Minimum Spanning Tree
(Prim's (PQ), Kruskal(DSU))

- ④ Huffman Encoding & Decoding /
optimal Merge Pattern

- ⑤ Single source Shortest Path
Algo { DIJKSTRA }

- ⑥ Job Sequencing

- ⑦ Problems based on $\log n$
SORTING

Lecture 1

① Job Sequencing

② Meeting Rooms - I

③ Disjoint Intervals - I

④ Disjoint Intervals - II

⑤ Maximum Chain Length

⑥ Minimum Balloon Burst

Saturday Morning

Lecture 2

① Meeting Rooms - II

② Minimum Platforms

③ Car Pooling

④ Merge Overlapping Intervals

⑤ Insert Interval

⑥ Intervals Intersection

Saturday Evening

Lecture ③

- ① Circular Tour (Salesforce)
- ② Car Fleet (Sprinklr)
- ③ Two City Scheduling
- ④ Candy/Temple Offerings
- ⑤ Chocolate Distribution
- ⑥ Queue Reconstruction Height

Sunday Morning

Lecture ④

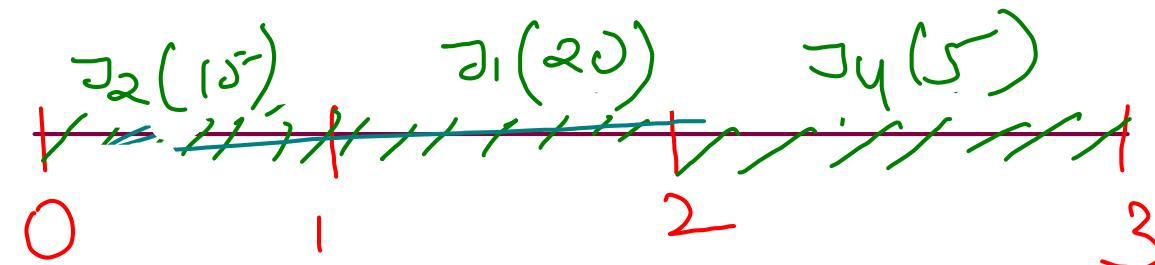
- ① Biased Standings
- ② Defense Kingdom
- ③ Georgivia
- ④ Amplifiers
- ⑤ Load Balancing
- # Huffman Coding & Decoding

Sunday Evening

Job Sequencing

1 job \rightarrow 1 unit of time
maximum prof's

$n=5$	$i \Rightarrow 0$	1	2	3	4	5
Jobs	J_1	J_2	J_3	J_4	J_5	
profits	20	15	10	5	1	
deadlines	2	2	1	3	3	



- ① Order of picking jobs
 \Rightarrow decreasing order of profit

- ② Picked job should be placed as last as possible

$J_1 \neq 3$

Benefit $\neq 3 \neq 40$

```

public static class MyComparator implements Comparator<Job>{
    public int compare(Job obj1, Job obj2){
        if(obj1.profit != obj2.profit){
            return obj2.profit - obj1.profit;
        }
        return obj2.deadline - obj1.deadline;
    }
}

//Function to find the maximum profit and the number of jobs done.
int[] Jobscheduling(Job arr[], int n)
{
    Arrays.sort(arr, new MyComparator());  $\Rightarrow$  n log n
    int maxDeadline = 0;
    for(int i=0; i<n; i++){
        maxDeadline = Math.max(arr[i].deadline, maxDeadline);
    }

    boolean[] slots = new boolean[maxDeadline];
    int maxProfit = 0;
    int jobsAllocated = 0;

    for(int i=0; i<n; i++) {
        for(int j=arr[i].deadline-1; j>=0; j--){
            if(slots[j] == false){
                slots[j] = true;
                jobsAllocated++;
                maxProfit += arr[i].profit;
                break;
            }
        }
    }

    return new int[]{jobsAllocated, maxProfit};
}

```

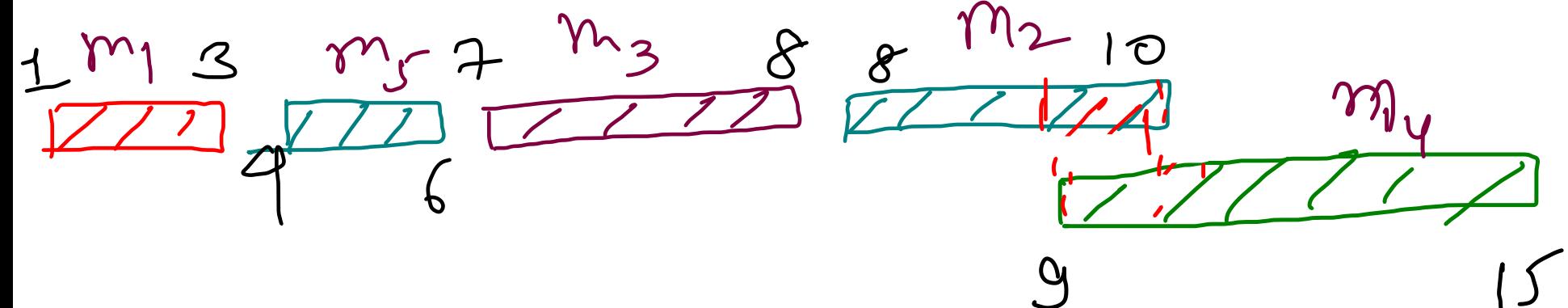
$O(n \log n)$

\downarrow
 $d \times n$
max deadline

Meeting Rooms - I

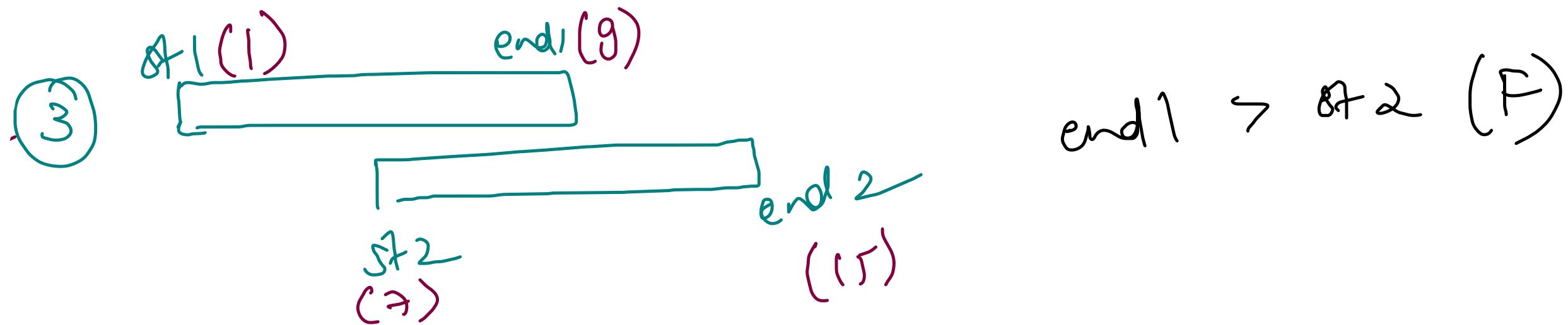
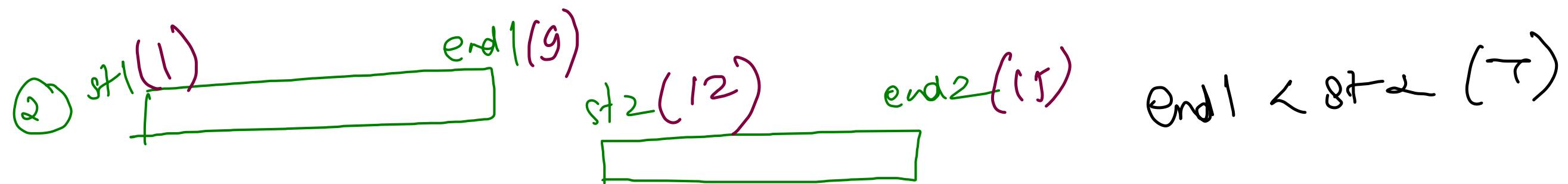
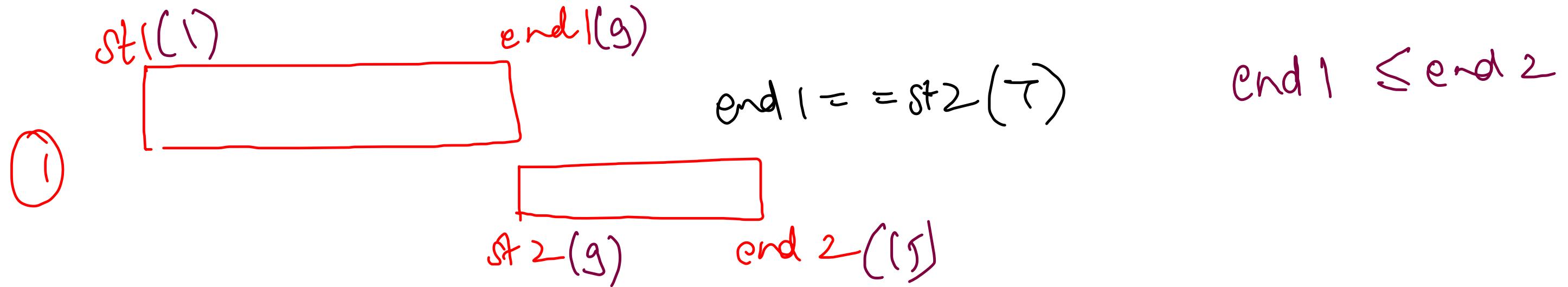
One room \Rightarrow one meeting at a time

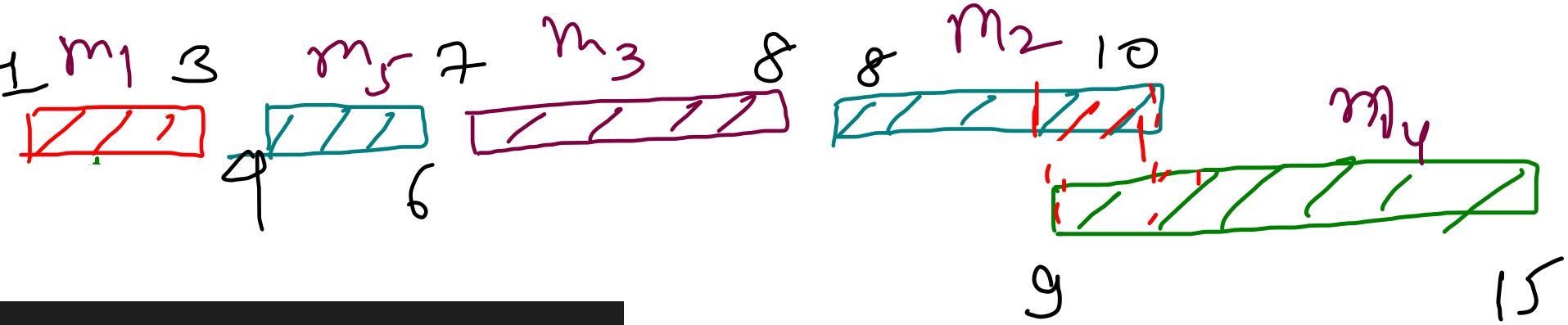
	0	1
① m_1	1	3
② m_2	8	10
③ m_3	7	8
④ m_4	9	15
⑤ m_5	4	2



- ① Sort intervals based on ending index
- ② Check overlapping with previous slot.

false





```

public class Solution {
    public static class MyComparator implements Comparator<Interval>{
        public int compare(Interval obj1, Interval obj2){
            if(obj1.end != obj2.end)
                return obj1.end - obj2.end;
            return obj1.start - obj2.start;
        }
    }

    public boolean canAttendMeetings(List<Interval> intervals) {
        Collections.sort(intervals, new MyComparator()); } O(N log N)

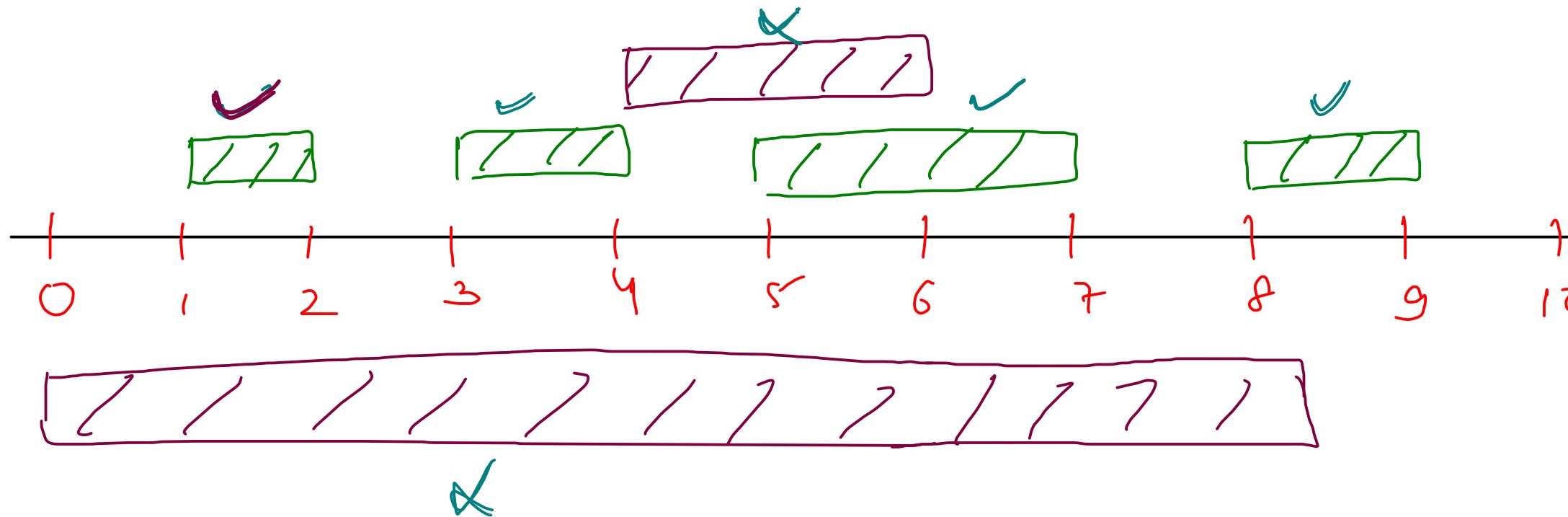
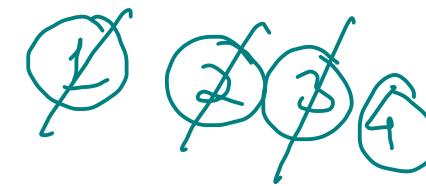
        int limit = Integer.MIN_VALUE; // last interval's ending time

        for(int i=0; i<intervals.size(); i++){
            if(limit > intervals.get(i).start){ }  $\Rightarrow O(N)$ 
                return false;
            }
            limit = intervals.get(i).end;
        }

        return true;
    }
}

```

$s[] =$	1	0	3	8	5	8	4
$f[] =$	2	18	4	9	7	16	
	1	2	3	4	5	6	



```

int limit = Integer.MIN_VALUE; // last interval's ending time
int count = 0;

for(int i=0; i<n; i++){
    if(limit < intervals[i].start){
        count++;
        limit = intervals[i].end;
    }
}

return count;

```

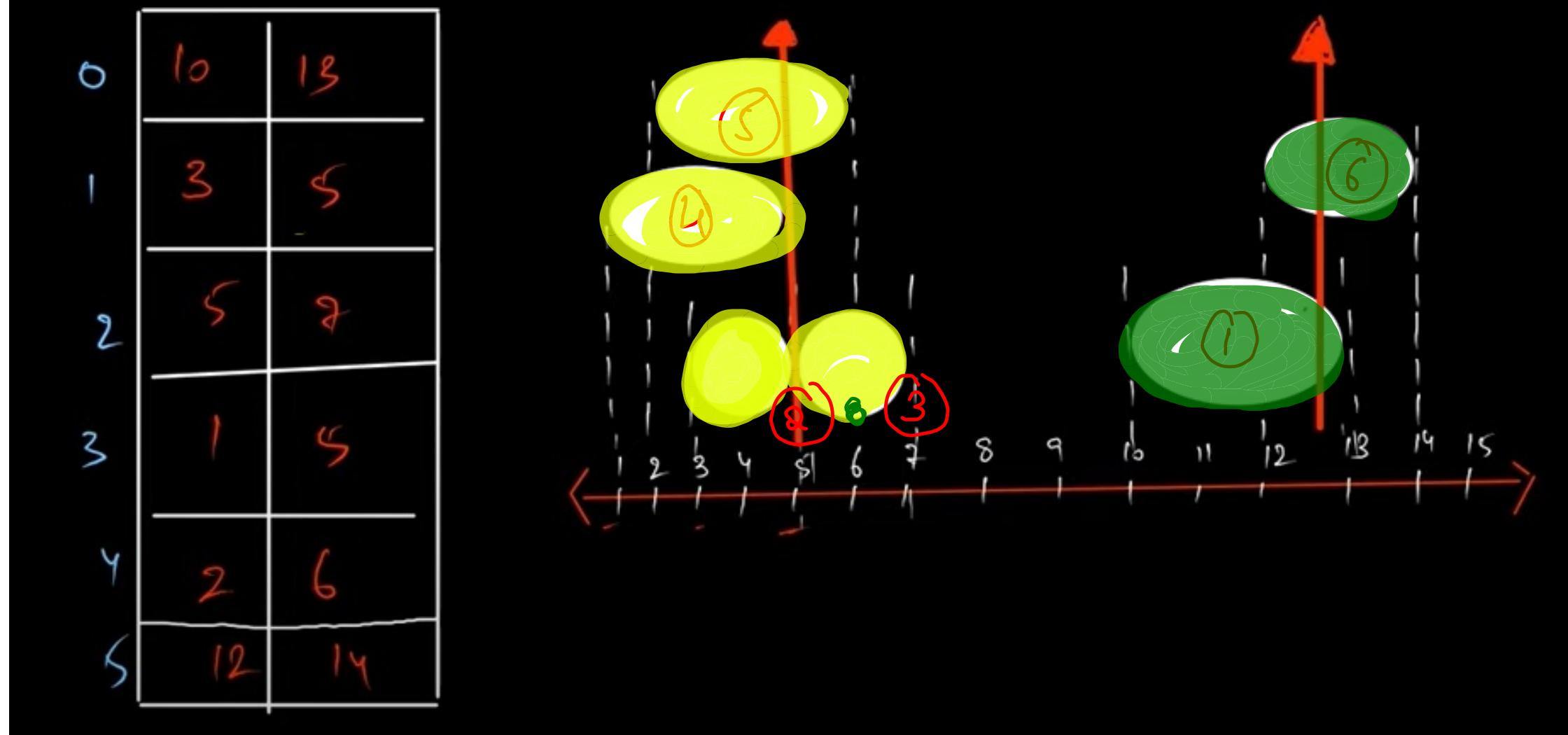
\Rightarrow $\text{end}_2 = 8 + 1$

\Rightarrow non-overlapping

\Rightarrow overlapping

Minimum Balloon Bursts

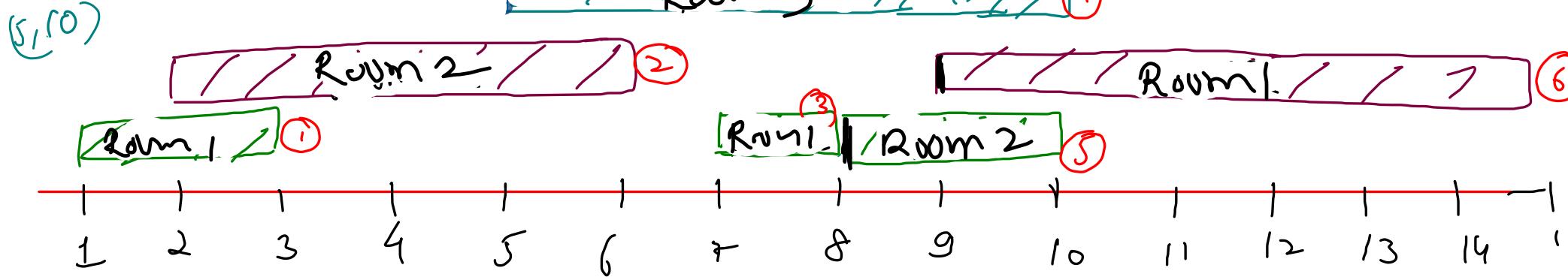
```
coordinates[] : [[10,13],[3,5],[5,7],[1,5],[2,6],[12,14]]  
output : 2
```



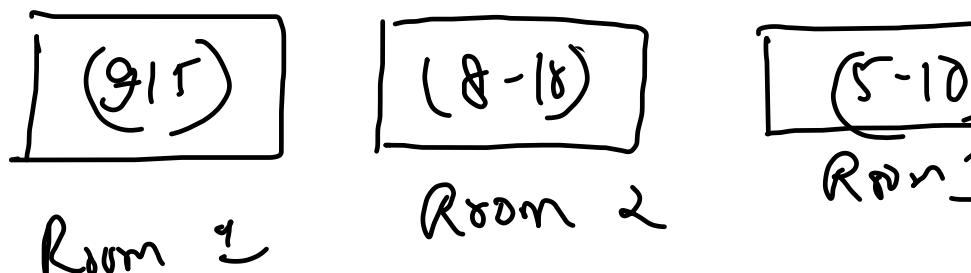
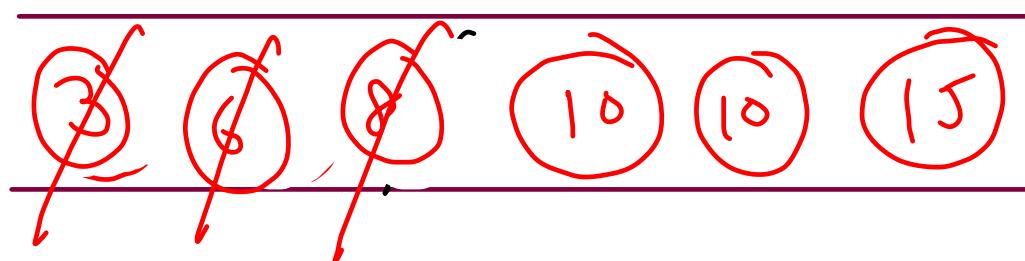
minimum
arrows
= maximum
meetings
that
can be
accommodated
in
one jump

5
1 3
8 10
7 8
9 15
2 6

minimum room



Count of Rooms = ~~0 1 2 3~~



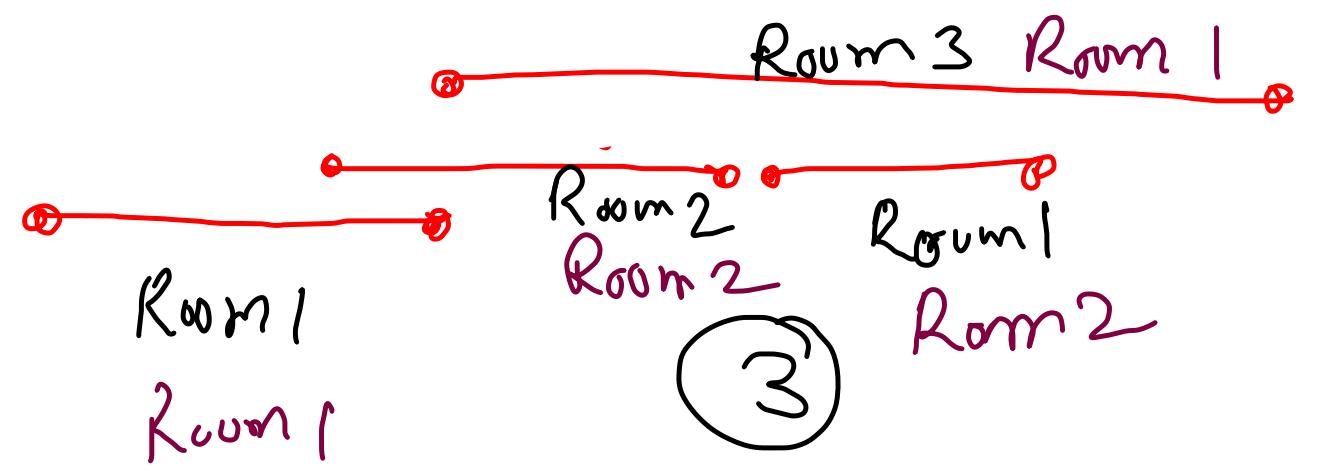
```
public static int meetingRooms(int intervals[][]){
    int maxRooms = 0;
    Arrays.sort(intervals, (a, b) -> a[1] - b[1]);

    Queue<Integer> q = new ArrayDeque<>();

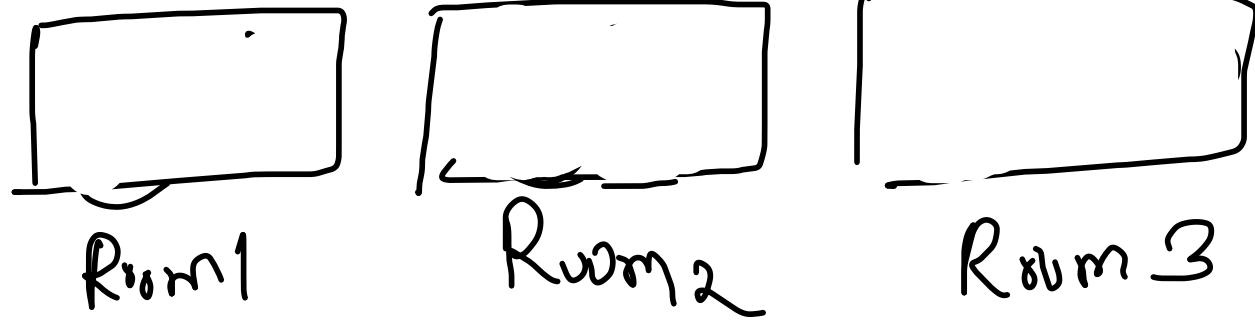
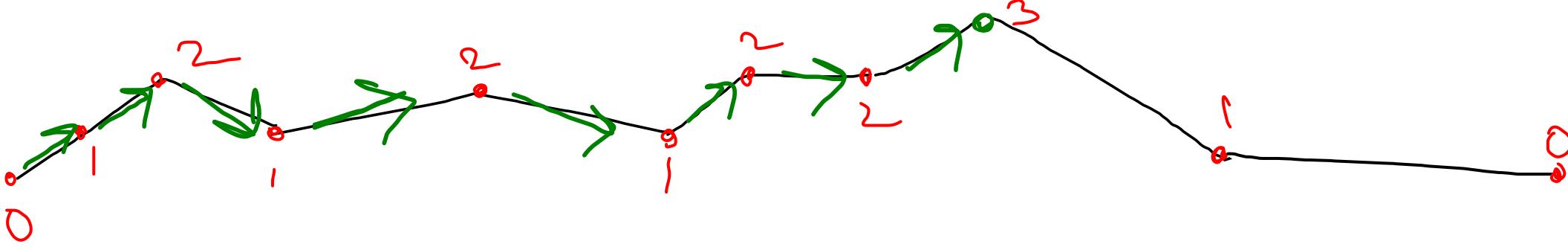
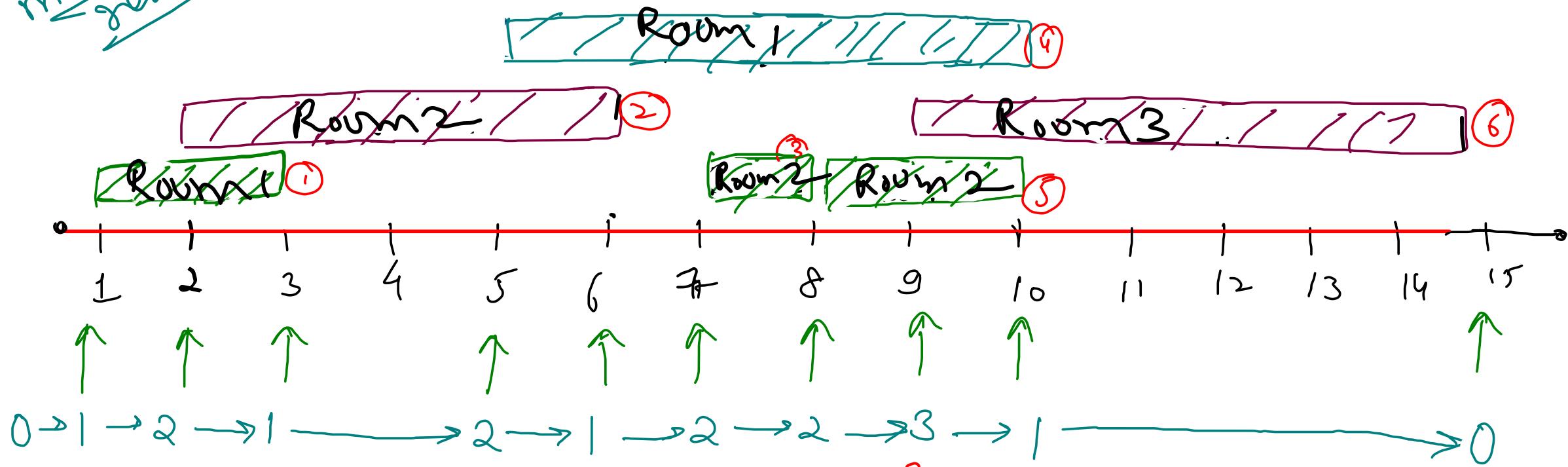
    for(int i=0; i<intervals.length; i++){
        if(q.size() == 0 || q.peek() > intervals[i][0]){
            maxRooms++;
        } else {
            q.remove();
        }
        q.add(intervals[i][1]);
    }
    return maxRooms;
}
```

Time $\rightarrow O(N \log N + N)$

Space $\rightarrow O(N)$



minimum
geoms



1 2 5 8 9
3 6 8 10 15
11

```
public int minMeetingRooms(List<Interval> intervals) {  
    ArrayList<Integer> start = new ArrayList<>();  
    ArrayList<Integer> end = new ArrayList<>();  
    for(int i=0; i<intervals.size(); i++){  
        start.add(intervals.get(i).start);  
        end.add(intervals.get(i).end);  
    }  
  
    Collections.sort(start);  
    Collections.sort(end);  
  
    int currentRooms = 0, maxRooms = 0;  
    int startIdx = 0, endIdx = 0;  
  
    while(startIdx < intervals.size()){  
        if(start.get(startIdx) < end.get(endIdx)){  
            startIdx++;  
            currentRooms++;  
        } else if(end.get(endIdx) < start.get(startIdx)){  
            endIdx++;  
            currentRooms--;  
        } else{  
            startIdx++;  
            endIdx++;  
        }  
        maxRooms = Math.max(maxRooms, currentRooms);  
    }  
    return maxRooms;  
}
```

$O(N)$

$O(N \log N)$

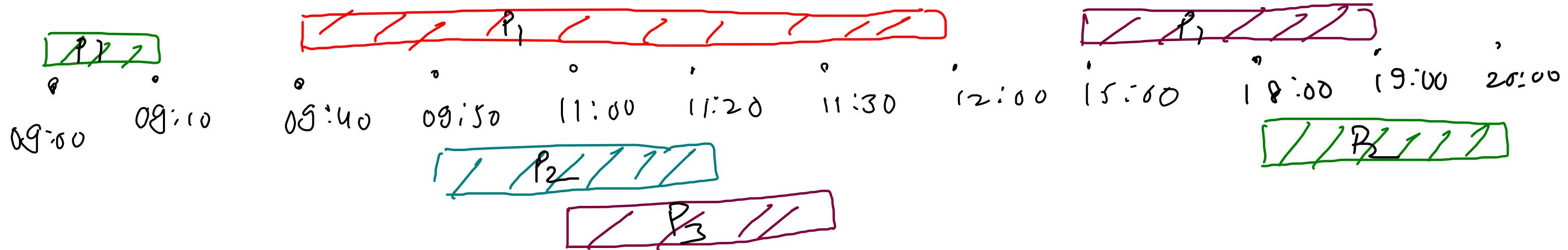
$O(2N)$

minimum Platforms

{ same as meeting Rooms - 11 }

{0900, 0940, 0950, 1100, 1500, 1800}

{0910, 1200, 1120, 1130, 1900, 2000}



Car Pooling

{2, 1, 5}

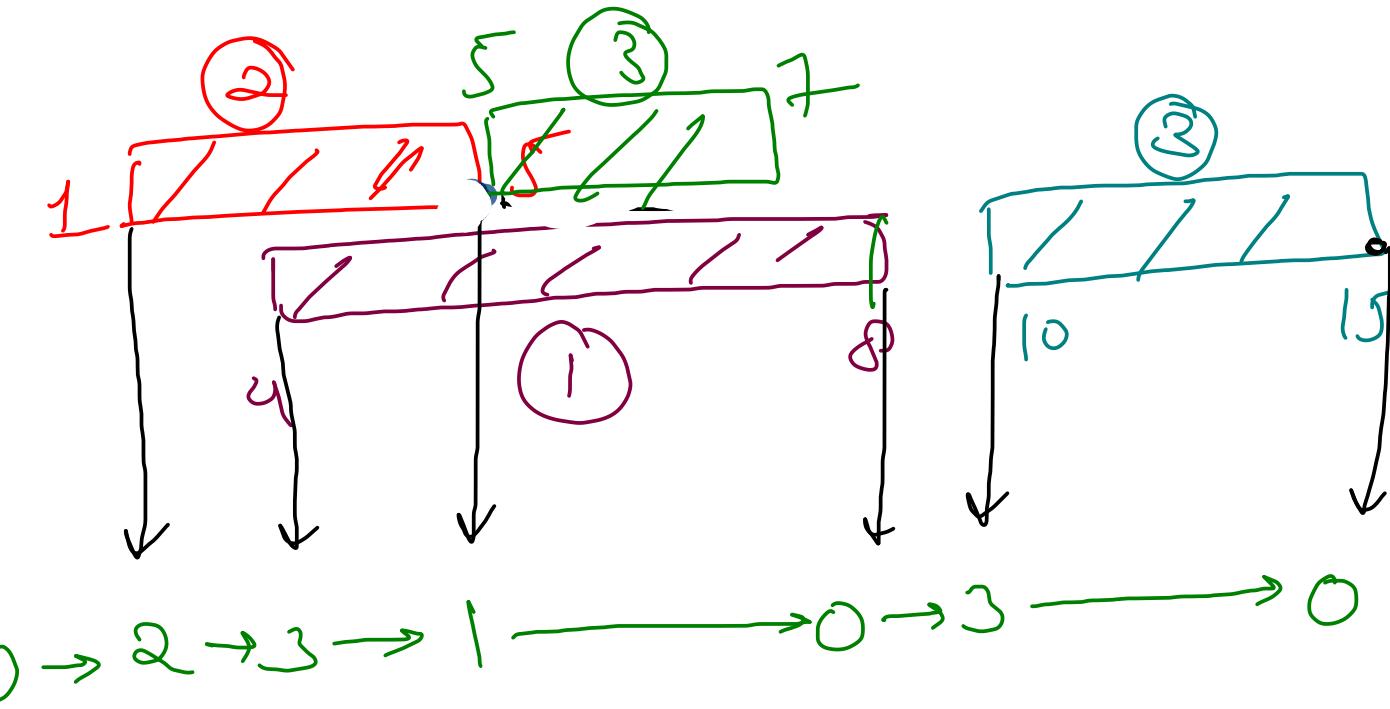
{1, 4, 8}

{3, 10, 15}

(3)

HINT
TreeMap { Red Black Tree }
→ ordered (keys)

direction → left to right , max^m capacity = ③



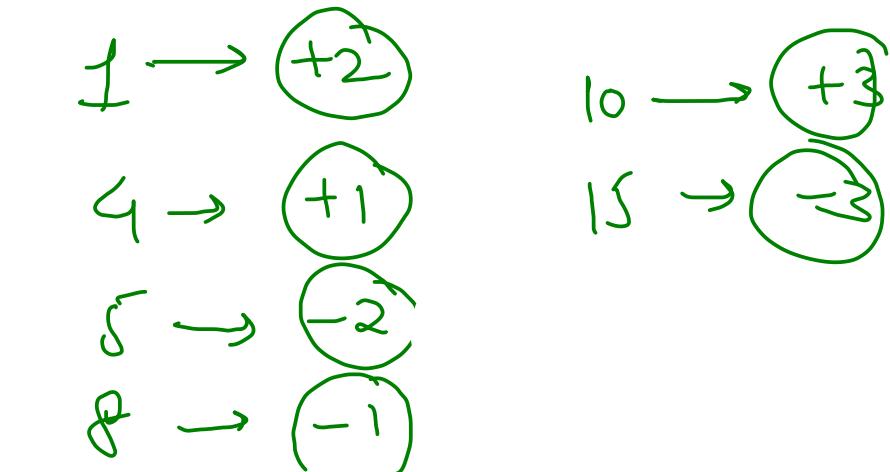
```

TreeMap<Integer, Integer> changes = new TreeMap<>();
for(int i=0; i<trips.length; i++){
    int passengers = trips[i][0];
    int start = trips[i][1];
    int end = trips[i][2];

    changes.put(start, changes.getOrDefault(start, 0) + passengers);
    changes.put(end, changes.getOrDefault(end, 0) - passengers);
}

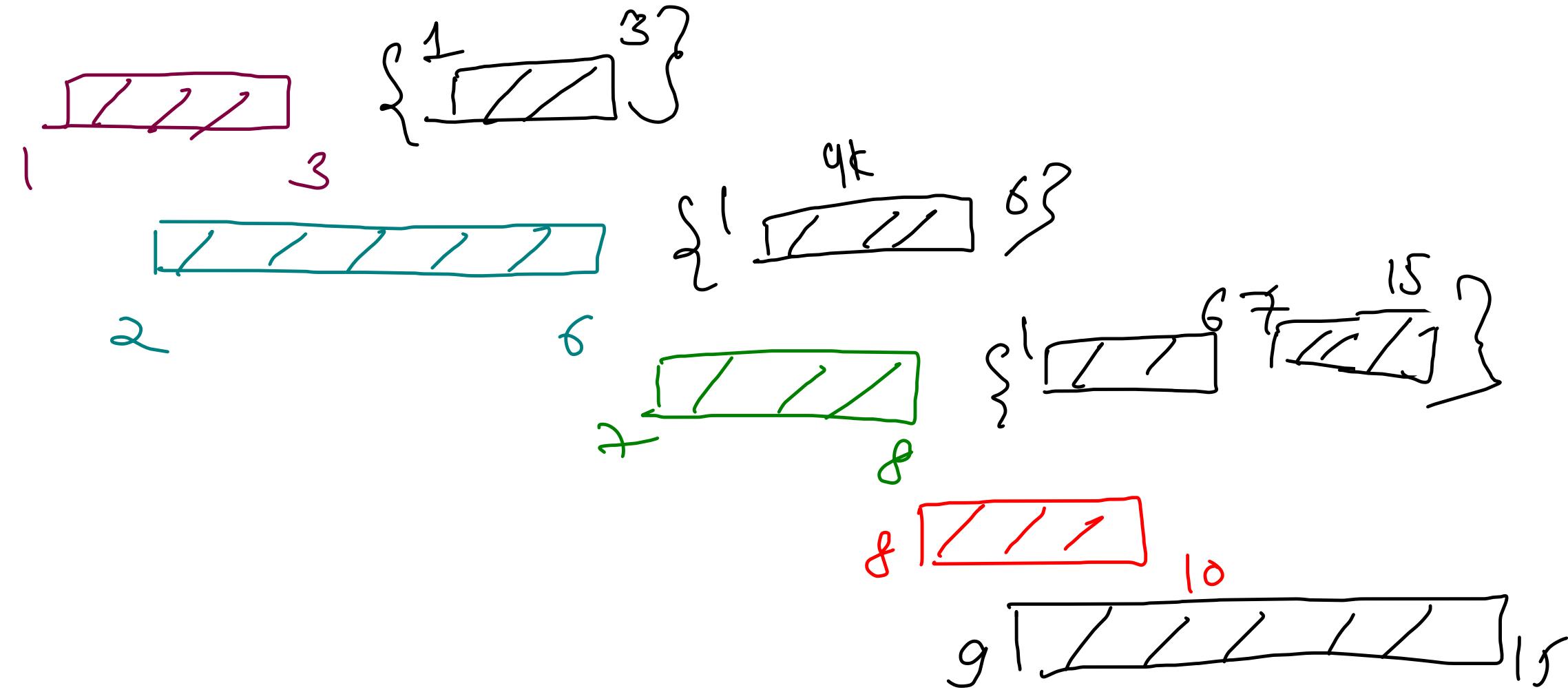
int currPassengers = 0;
for(Integer key: changes.keySet()){
    Integer val = changes.get(key);
    currPassengers += val;

    if(currPassengers > capacity) return false;
}
return true;
    
```

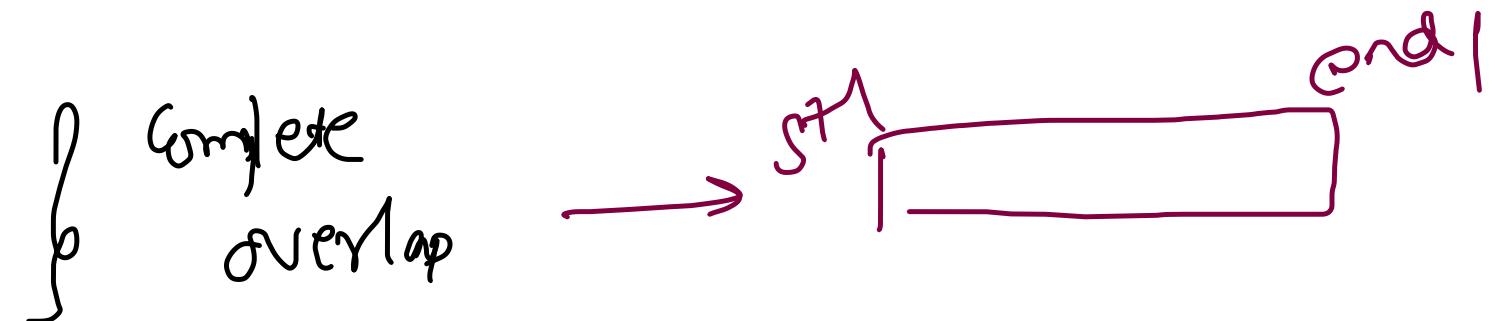


Merge Overlapping Intervals

0	1
0	1 3
1	8 10
2	7 8
3	9 15
4	2 6



Sort on basis of starting index \Rightarrow $st_1 \leq st_2$



```

public int[][] merge(int[][] intervals) {
    Arrays.sort(intervals, (a, b) -> a[0] - b[0]);
    ArrayList<int[]> merged = new ArrayList<>();
    merged.add(intervals[0]);

    for(int i=1; i<intervals.length; i++){
        int[] lastInt = merged.get(merged.size() - 1);
        int[] currInt = intervals[i];

        if(lastInt[1] >= currInt[0]){
            // merge
            lastInt[1] = Math.max(lastInt[1], currInt[1]);
        } else {
            merged.add(currInt);
        }
    }

    int[][] res = new int[merged.size()][2];
    for(int i=0; i<merged.size(); i++){
        res[i] = merged.get(i);
    }
    return res;
}

```

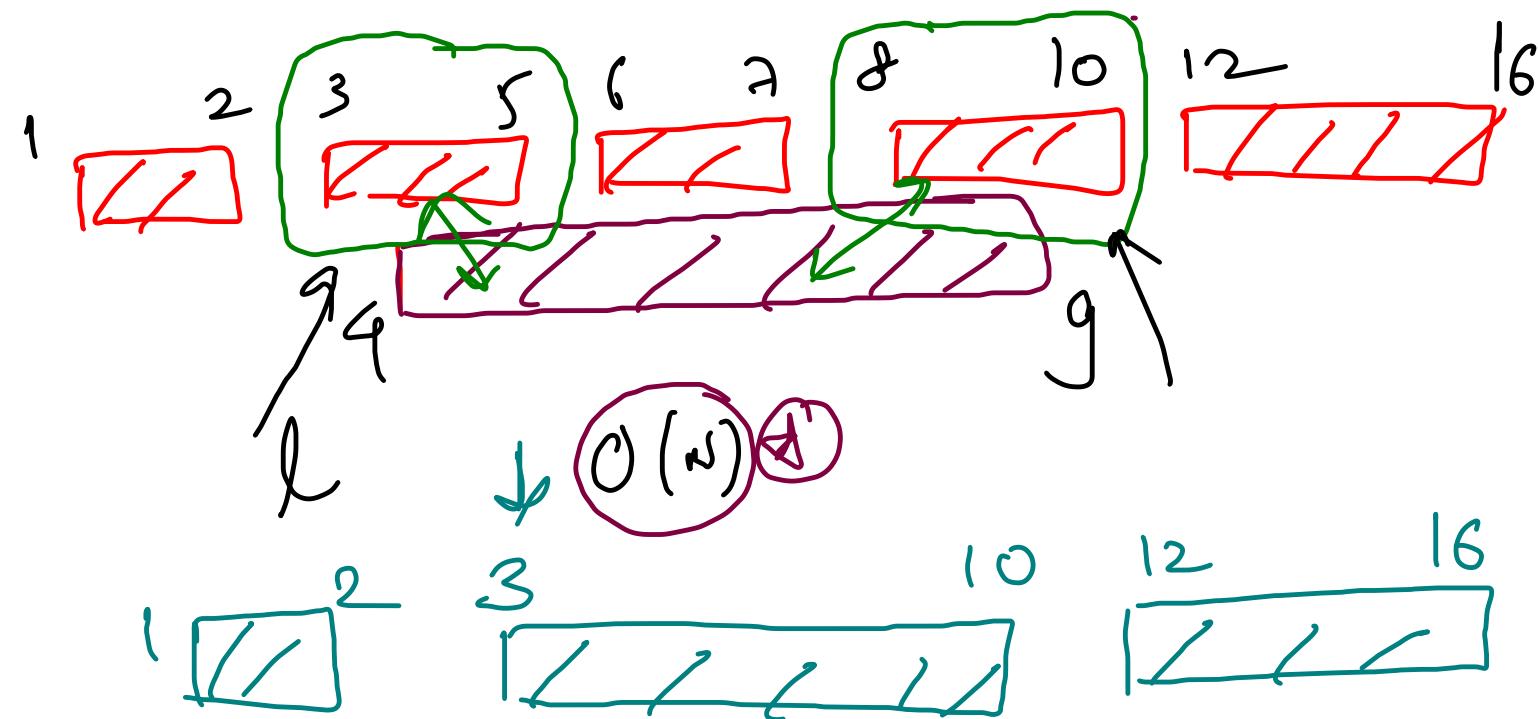
$T \in \Theta(N \log N)$

$\frac{SC}{T}$
 $O(N)$ + $O(N)$ + $O(1)$
 output Input extra space

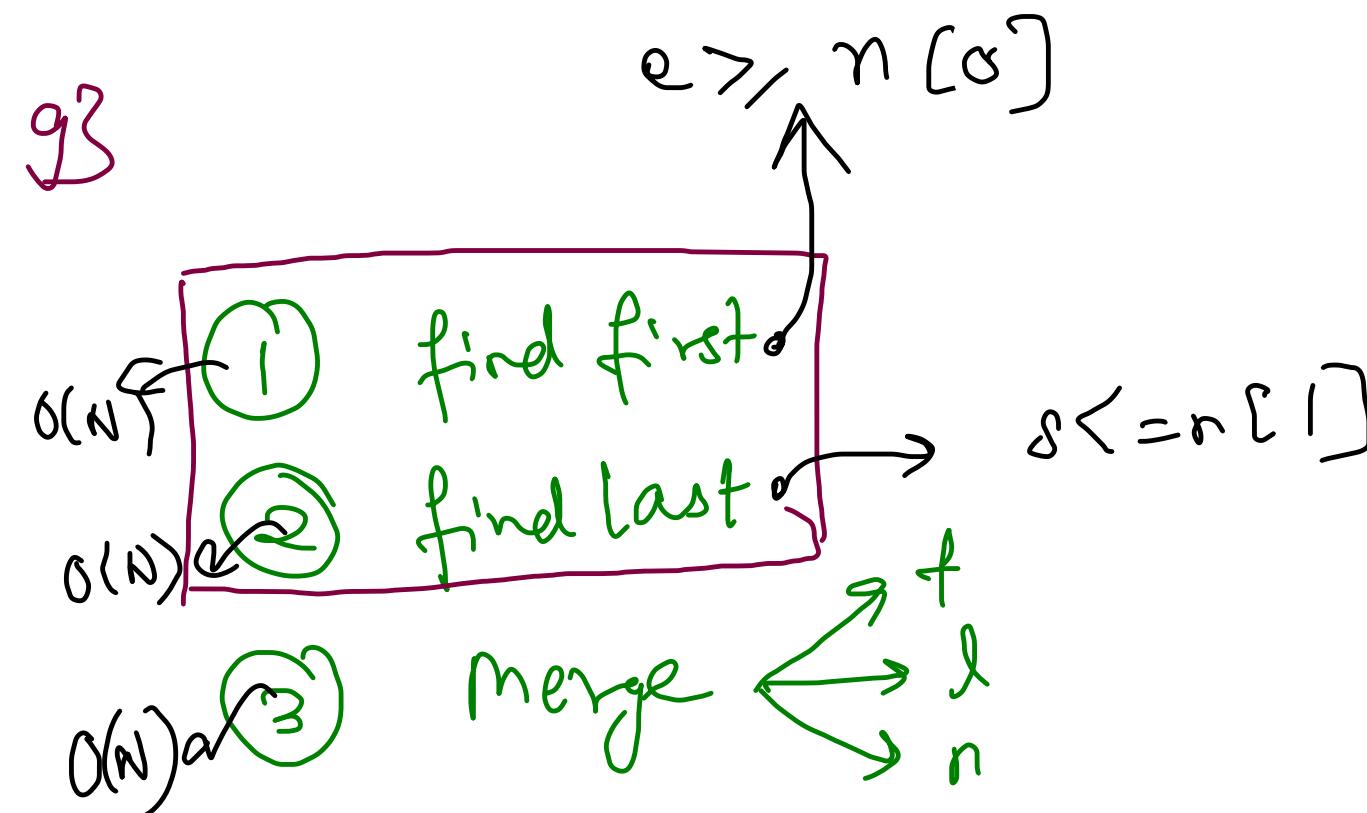
Insert Interval

`[[1, 2], [3, 5], [6, 7], [8, 10], [12, 16]]`

24, 93



```
int[] merged = new int[2];
merged[0] = Math.min(intervals[firstIdx][0], newInterval[0]);
merged[1] = Math.max(intervals[lastIdx][1], newInterval[1]);
res.add(merged);
```



$\{ \min(f[0], n[0]),$
 $\max(l[1], n[1]) \}$

```

int findFirst(int[][] intervals, int[] newInterval){
    for(int idx=0; idx<intervals.length; idx++){
        if(intervals[idx][1] >= newInterval[0]){
            return idx;
        }
    }
    return intervals.length;
}

int findLast(int[][] intervals, int[] newInterval){
    for(int idx=intervals.length-1; idx>=0; idx--){
        if(intervals[idx][0] <= newInterval[1]){
            return idx;
        }
    }
    return -1;
}

int firstIdx = findFirst(intervals, newInterval);
int lastIdx = findLast(intervals, newInterval);

ArrayList<int[]> res = new ArrayList<>();

```

```

// non merging -> firstIdx > lastIdx
if(firstIdx > lastIdx){
    for(int i=0; i<lastIdx; i++){
        res.add(intervals[i]);
    }
    res.add(newInterval);
    for(int i=firstIdx; i<intervals.length; i++){
        res.add(intervals[i]);
    }
}

```

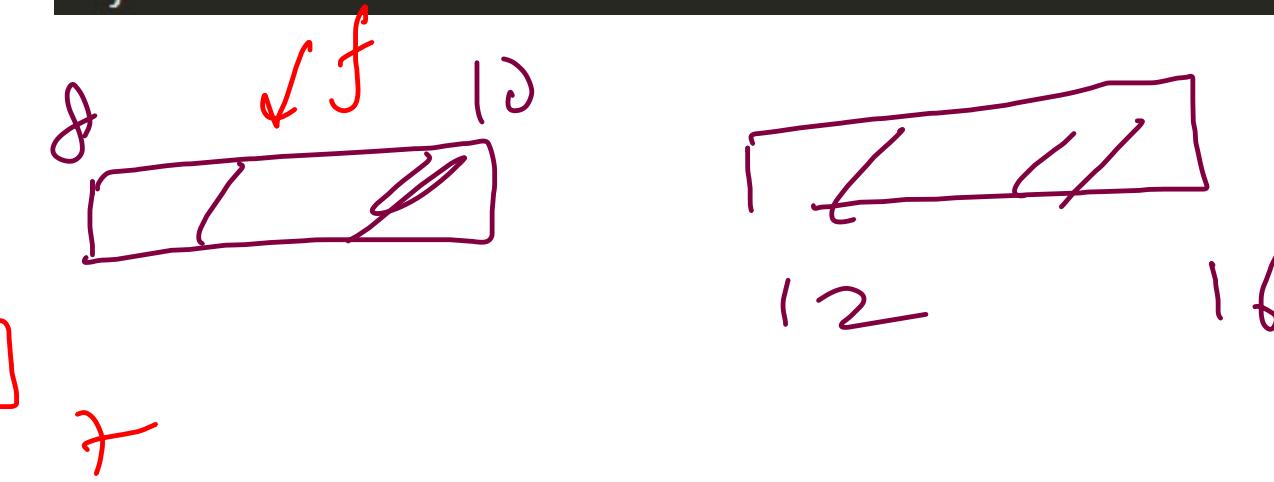
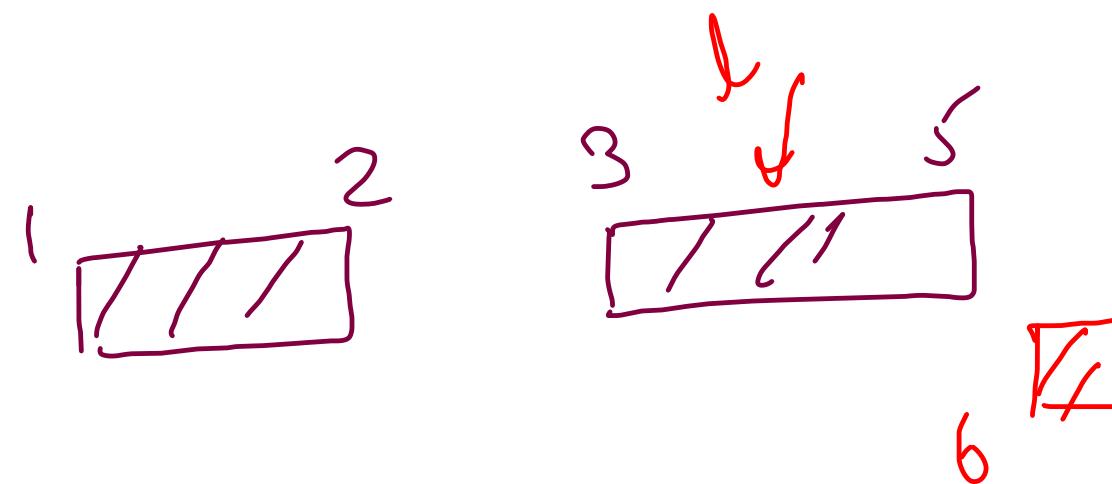
```

} else {
    for(int i=0; i<firstIdx; i++){
        res.add(intervals[i]);
    }

    int[] merged = new int[2];
    merged[0] = Math.min(intervals[firstIdx][0], newInterval[0]);
    merged[1] = Math.max(intervals[lastIdx][1], newInterval[1]);
    res.add(merged);

    for(int i=lastIdx+1; i<intervals.length; i++){
        res.add(intervals[i]);
    }
}

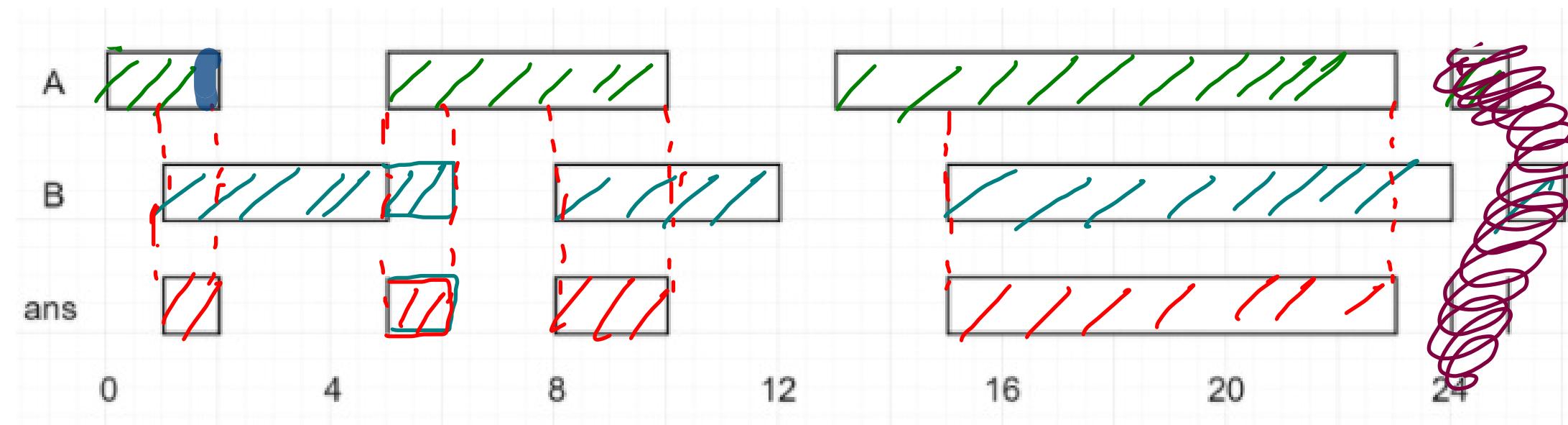
```



Intervals Intersection

non-overlapping

Each list of intervals is pairwise **disjoint** and in **sorted order**.



```
public int[][] intervalIntersection(int[][][] firstList, int[][][] secondList) {
    int firstIdx = 0, secondIdx = 0;
    ArrayList<int[]> intersect = new ArrayList<>();

    while(firstIdx < firstList.length && secondIdx < secondList.length){
        int s1 = firstList[firstIdx][0];
        int e1 = firstList[firstIdx][1];
        int s2 = secondList[secondIdx][0];
        int e2 = secondList[secondIdx][1];

        int start = Math.max(s1, s2);
        int end = Math.min(e1, e2);

        if(start <= end){
            intersect.add(new int[]{start, end});
        }

        if(e1 < e2) firstIdx++;
        else secondIdx++;
    }

    int[][] res = new int[intersect.size()][2];
    for(int i=0; i<intersect.size(); i++){
        res[i] = intersect.get(i);
    }
    return res;
}
```

- ① S&S - GitHub
- ② Notes
- ③ Recording
- ④ CB videos
- ⑤ RL R-level ②