

Recursion & Backtracking

Today's Class { lecture - 7 }

[N Queen - Combinations](#)

[N Queen - Permutations](#)

[N-Queens - I \(IB\)](#)

[N-Queens - II](#)

[Using Backtracking](#)

[Code](#)

[Using Branch & Bound](#)

[Using Bit Manipulation](#)

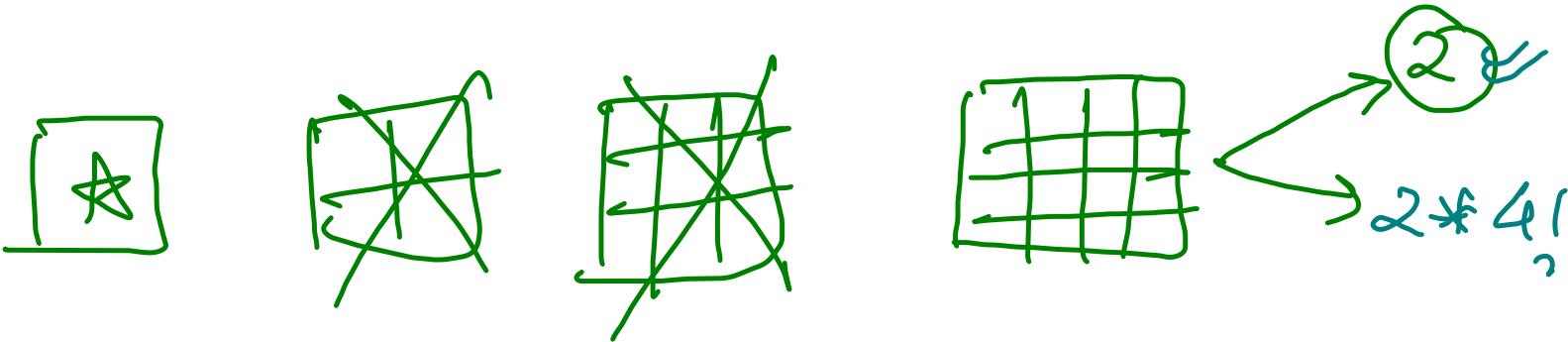
[N Knight](#)

[Knight's Tour](#)

[GfG Article](#)

[Knight's Tour](#)

[Code](#)



	q_1		
q_2			*
*			q_3
		q_4	*

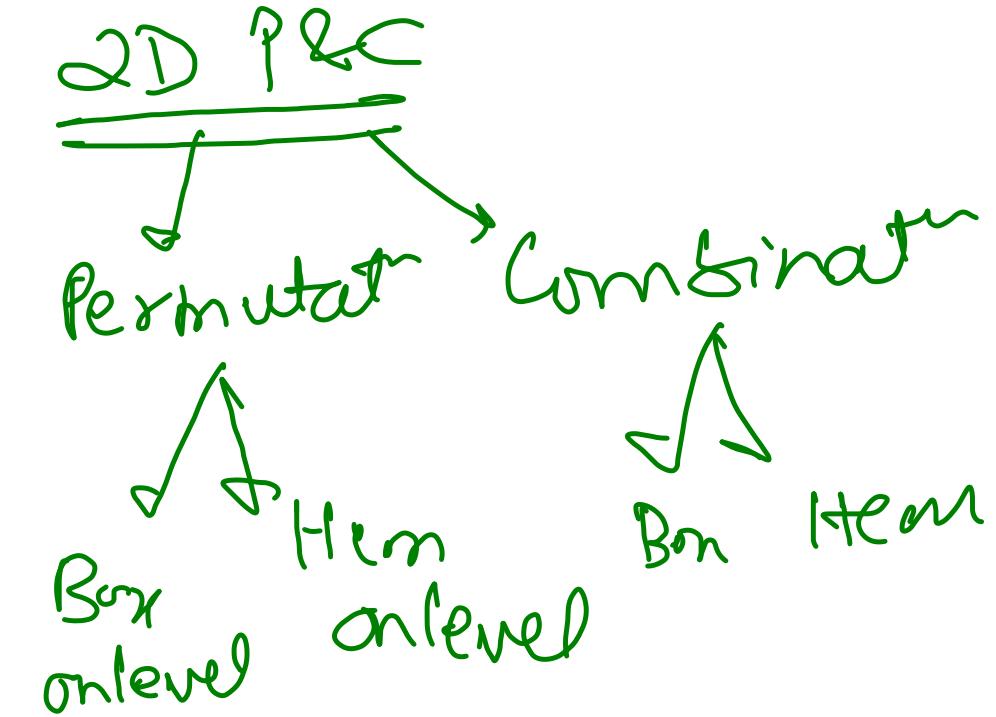
$$N^2 = 4 \times 4$$

$N = 4$ queens

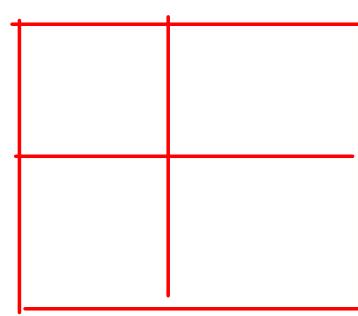
Combination

$$\{q_1, q_2, q_3, q_4\}$$

constraint \rightarrow No queen pair should
kill each other.



① N Queen - Combination



$\{q_1, q_2\}$

$$2^{N^2}$$

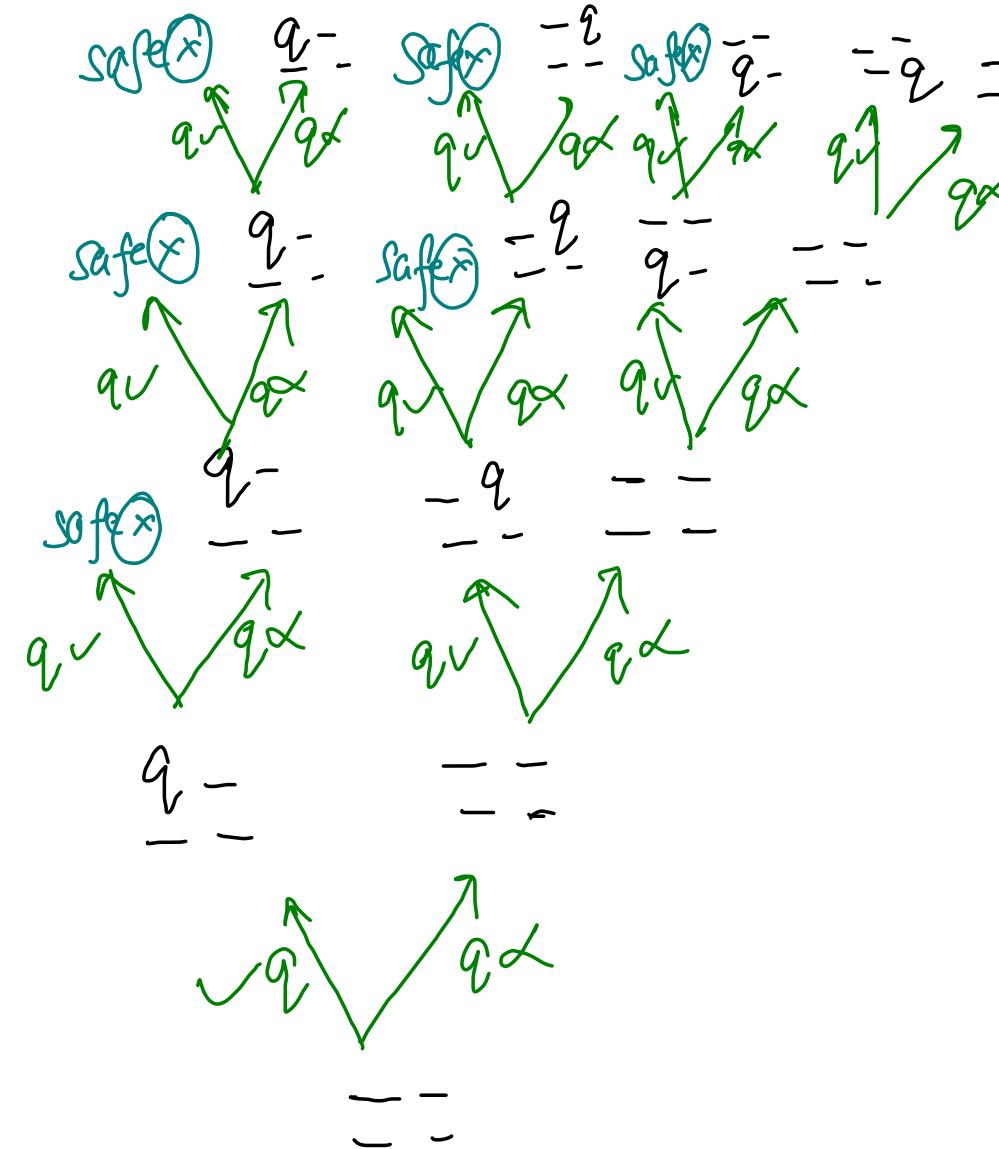
$N = 2 \Rightarrow$
⁰ valid
configurations

(1, 1)

(1, 0)

(0, 1)

(0, 0)



(call(s))^{height}

+

(pre + post) * height

(2^n)

+

$n * (n^2) = O(2^n + n^3)$

$N = 9$
 $\frac{81}{2} + 3^3$

#IsQueenSafe

	c0	c1	c2	c3
r0	dD	dE	dG	dH
r1	dC	dD	dF	dG
r2	dS	dC	dD	dF
r3	dA	dB	dC	dD

① Left Diagonal

	c0	c1	c2	c3
r0	dI	dII	dIII	dIV
r1	dIX	dIII	dIV	dV
r2	dIII	dIV	dV	dVI
r3	dV	dVI	dVII	dVIII

② Right Diagonal

	c0	c1	c2	c3
r0	r0	r0	r0	r0
r1	r1	r1	r1	r1
r2	r2	r2	r2	r2
r3	r3	r3	r3	r3

③ Row

	c0	c1	c2	c3
r0	c0	c1	c2	c3
r1	c0	c1	c2	c3
r2	c0	c1	c2	c3
r3	c0	c1	c2	c3

④ column

Total left or right diagonals
 $= 2n - 1$

Total rows or columns = n

$O(N)$ time, $O(1)$ space

```
public static boolean isQueenSafe(int r, int c, boolean[][] vis){  
    // row (left)  
    for(int j=0; j<c; j++){  
        if(vis[r][j] == true)  
            return false;  
    }  
    // col (up)  
    for(int i=0; i<r; i++){  
        if(vis[i][c] == true)  
            return false;  
    }  
  
    // left diagonal (up)  
    int i = r, j = c;  
    while(i >= 0 && j >= 0){  
        if(vis[i][j] == true){  
            return false;  
        }  
        i--; j--;  
    }  
  
    // right diagonal (down)  
    i = r; j = c;  
    while(j < vis.length && i >= 0){  
        if(vis[i][j] == true){  
            return false;  
        }  
        i--; j++;  
    }  
  
    return true;  
}
```

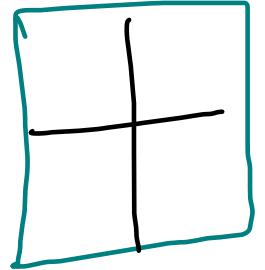
→ Recursion Space: $\rightarrow O(n^2)$
→ Output: $O(n^2)$

```
public static void nqueen(int r, int c, int qpsf, boolean[][] vis) {  
    if (qpsf == vis.length) {  
        List<String> ans = new ArrayList<>();  
        for (int i = 0; i < vis.length; i++) {  
            res.add(ans);  
        }  
        return;  
    }  
  
    if(r == vis.length){  
        return;  
    }  
  
    if (isQueenSafe(r, c, vis)) {  
        vis[r][c] = true;  
  
        if (c == vis.length - 1) {  
            nqueen(r + 1, 0, qpsf + 1, vis);  
        } else {  
            nqueen(r, c + 1, qpsf + 1, vis);  
        }  
  
        vis[r][c] = false;  
    }  
  
    if (c == vis.length - 1) {  
        nqueen(r + 1, 0, qpsf, vis);  
    } else {  
        nqueen(r, c + 1, qpsf, vis);  
    }  
}
```

N-Queen Permutation

$\{q_1, q_2\}$ Branch level

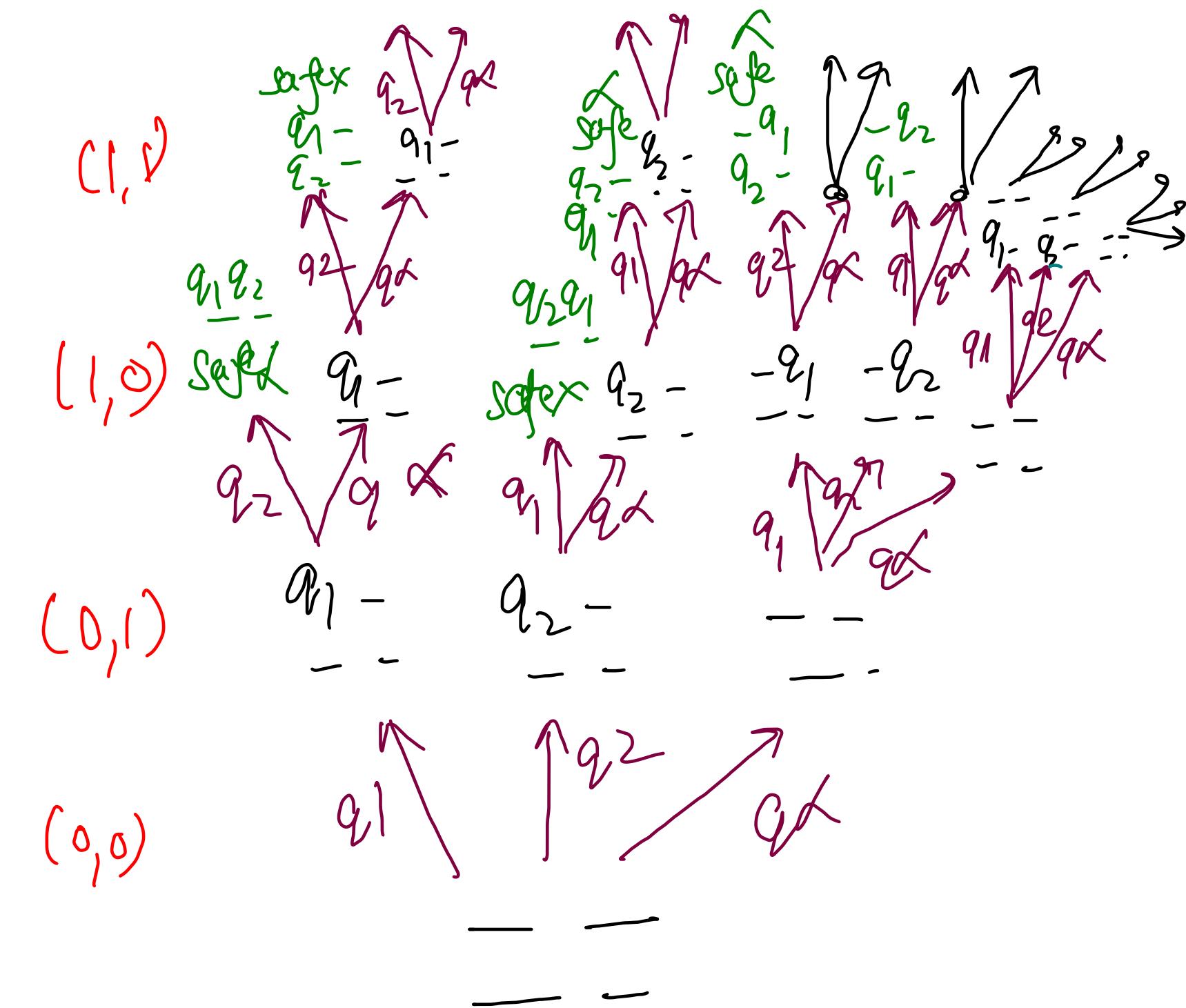
$$2 \times 2 \quad N = 2$$



20 valid configurations

$$N = 2 \Rightarrow \text{Permutations} = 2 \times 2!$$

Visited queens array ~~array~~ $= 2^2 = 4$ = 48



Ban on level approach { Order Different }

```
public static void nqueens(int qpsf, int r, int c, boolean[] queenPlaced, int[][] chess) {
    int n = queenPlaced.length;
    if(qpsf == n){
        for(int i=0; i<n; i++){}
        System.out.println();
        return;
    }

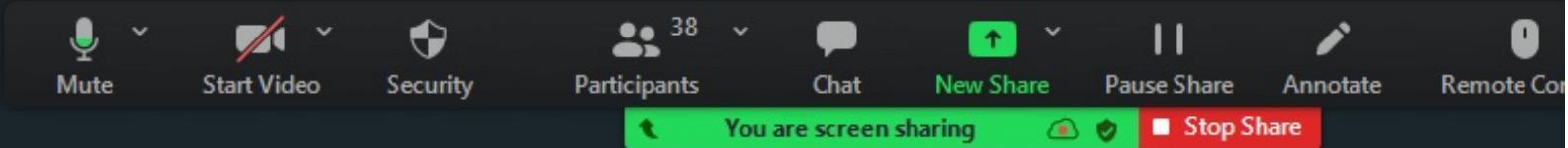
    if(r == n){ return; }

    if(c == n-1){
        nqueens(qpsf, r+1, 0, queenPlaced, chess);
    } else {
        nqueens(qpsf, r, c+1, queenPlaced, chess);
    }

    for(int i=0; i<n; i++){
        if(queenPlaced[i] == false && isQueenSafe(r, c, chess)){
            queenPlaced[i] = true;
            chess[r][c] = (i + 1);

            if(c == n-1){
                nqueens(qpsf+1, r+1, 0, queenPlaced, chess);
            } else {
                nqueens(qpsf+1, r, c+1, queenPlaced, chess);
            }

            chess[r][c] = 0;
            queenPlaced[i] = false;
        }
    }
}
```



A screenshot of a video conferencing interface at the bottom of the slide. It includes controls for Mute, Start Video, Security, Participants (38), Chat, New Share, Pause Share, Annotate, and Remote Control. A message at the bottom says "You are screen sharing".

#IsQueenSafe

	c0	c1	c2	c3
r0	d _D	d _E	d _G	d _H
r1	d _C	d _D	d _F	d _G
r2	d _S	d _C	d _D	d _F
r3	d _A	d _B	d _C	d _D

① Left Diagonal

	c0	c1	c2	c3
r0	d _I	d _{II}	d _{III}	d _{IV}
r1	d _{II}	d _{III}	d _{IV}	d _V
r2	d _{III}	d _{IV}	d _V	d _{VI}
r3	d _{IV}	d _V	d _{VI}	d _{VII}

② Right Diagonal

	c0	c1	c2	c3
r0	r ₀	r ₀	r ₀	r ₀
r1	r ₁	r ₁	r ₁	r ₁
r2	r ₂	r ₂	r ₂	r ₂
r3	r ₃	r ₃	r ₃	r ₃

③ Row

	c0	c1	c2	c3
r0	c ₀	c ₁	c ₂	c ₃
r1	c ₀	c ₁	c ₂	c ₃
r2	c ₀	c ₁	c ₂	c ₃
r3	c ₀	c ₁	c ₂	c ₃

④ column

[T_f T_f T_f T_f T_f]

Total left or right

diagonals

$$= 2n - 1$$

Total rows or

columns = n

NQueen {Combinat} {Backtracking}

	★	#	
#			★
★			#
	#	★	

4 rows
4 columns

4 queens

row m
level
column as option

$$(n)^n + (n) \times n$$



isQueenSafe
 \downarrow
 $(n) \times n$

$$\Rightarrow O(n^n + n^2)$$

$$N = g \\ g^g + g^2$$

$O(n)$ time, $O(1)$ space

```
public static boolean isQueenSafe(int r, int c, boolean[][] vis){  
    // row (left)  
    for(int j=0; j<c; j++){  
        if(vis[r][j] == true)  
            return false;  
    }  
    // col (up)  
    for(int i=0; i<r; i++){  
        if(vis[i][c] == true)  
            return false;  
    }  
    // left diagnol (up)  
    int i = r, j = c;  
    while(i >= 0 && j >= 0){  
        if(vis[i][j] == true){  
            return false;  
        }  
        i--; j--;  
    }  
    // right diagnol (down)  
    i = r; j = c;  
    while(j < vis.length && i >= 0){  
        if(vis[i][j] == true){  
            return false;  
        }  
        i--; j++;  
    }  
  
    return true;  
}
```

Recursion call stack $\rightarrow O(n)$

output : $O(n^2)$

```
public static void nqueen(int r, boolean[][] vis) {  
    if (r == vis.length) {  
        List<String> ans = new ArrayList<>();  
        for (int i = 0; i < vis.length; i++) {  
            res.add(ans);  
        }  
        return;  
    }  
  
    for(int c=0; c<vis.length; c++){  
        if(isQueenSafe(r, c, vis)){  
            vis[r][c] = true;  
            nqueen(r + 1, vis);  
            vis[r][c] = false;  
        }  
    }  
}
```

③ Branch & Bound

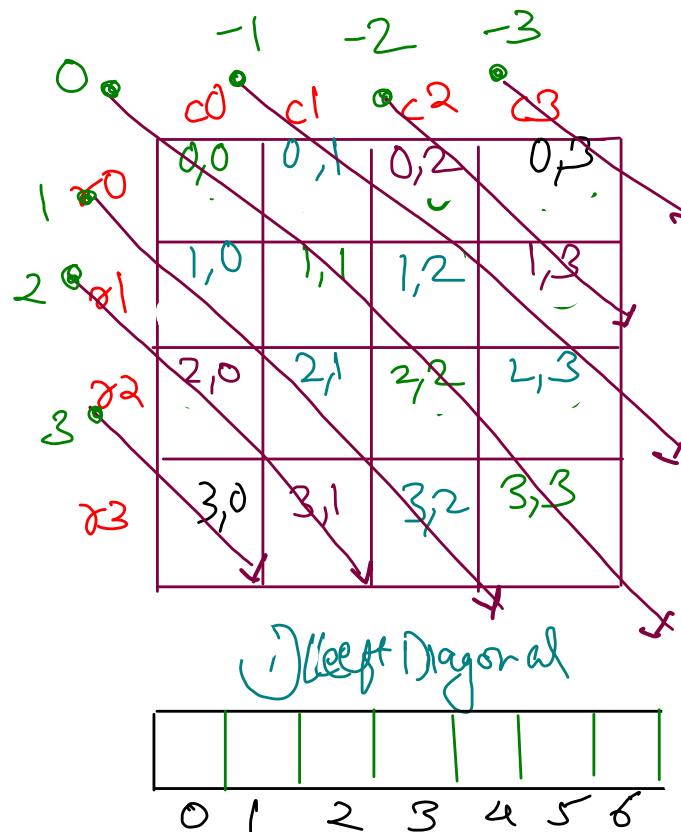
{ optimization in isQueenSafe }
 in terms of
 time complexity

	c_0	c_1	c_2	c_3
r_0	c_0	c_1	c_2	c_3
r_1	c_0	c_1	c_2	c_3
r_2	c_0	c_1	c_2	c_3
r_3	c_0	c_1	c_2	c_3

$\boxed{T_f \quad T_{1f} \quad T_{2f} \quad T_{3f}}$

Column
 $\{n\}$

BytE



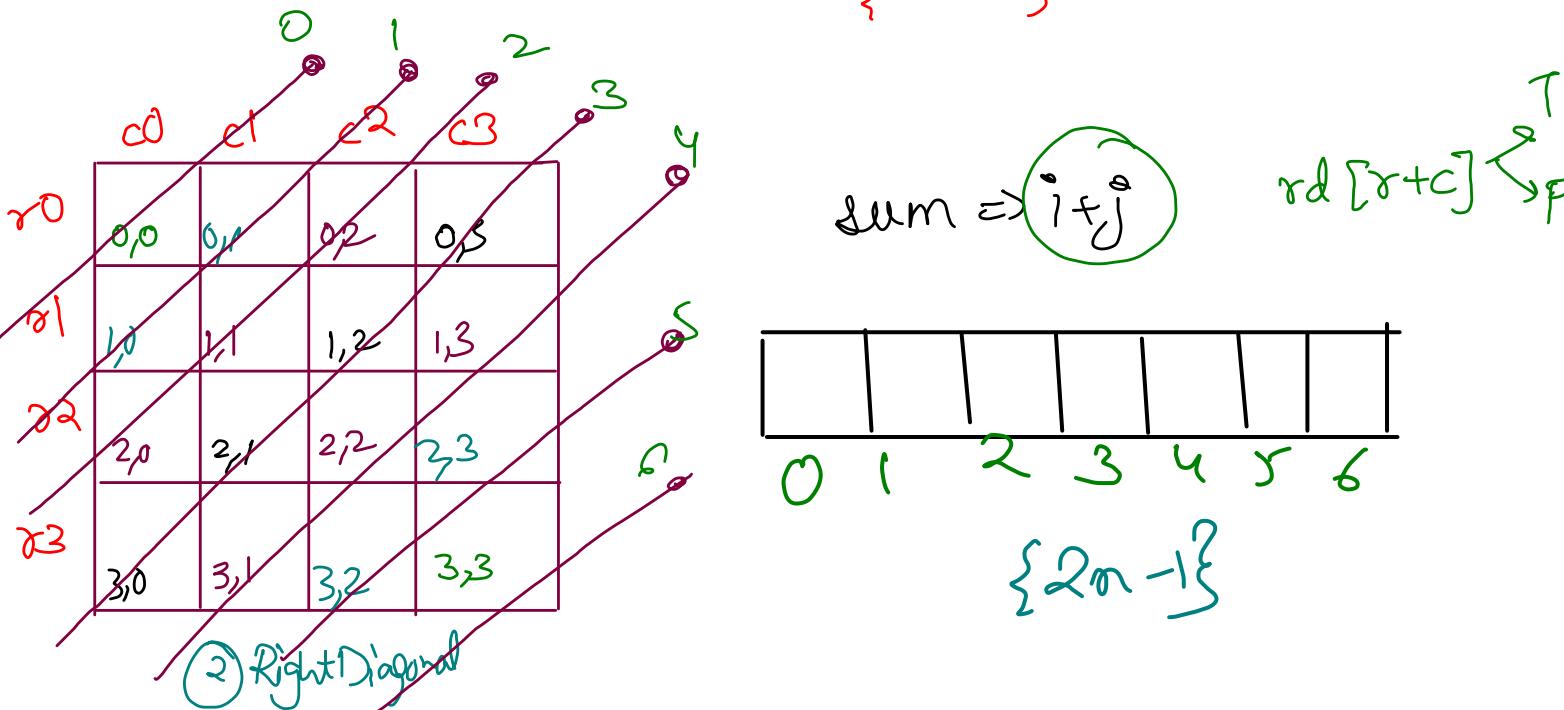
left diagonal for r, c

$$= ld [\cancel{r-c} + (n-1)]$$

$\frac{-3}{4+3}, \frac{-2}{4+3}, \frac{-1}{4+3}, \frac{0}{4+3}, \frac{1}{4+3}, \frac{2}{4+3}, \frac{3}{4+3}$

0 1 2 3 4 5 6

$\{2n-1\}$



0 1 2 3 4 5 6

$\{2n-1\}$

sum $\Rightarrow i+j$

$rd [r+c]$

T_f

```

public static boolean isQueenSafe(int r, int c, boolean[] ld, boolean[] rd, boolean[] col){
    int n = col.length;
    return ((ld[r - c + n - 1] == true) || (rd[r + c] == true) || (col[c] == true)) ? false : true;
}

```

→ Extra Space

$O(1)$ time
 $O(n)$ space

```

public static void nqueen(int r, boolean[][] chess, boolean[] ld, boolean[] rd, boolean[] col) {
    int n = chess.length;
    if (r == n) {
        List<String> ans = new ArrayList<>();
        for (int i = 0; i < n; i++) {
            String curr = "";
            for (int j = 0; j < n; j++) {
                if (chess[i][j] == true) {
                    curr = curr + "Q";
                } else {
                    curr = curr + ".";
                }
            }
            ans.add(curr);
        }
        res.add(ans);
    }
    return;
}

for(int c=0; c<n; c++){
    if(isQueenSafe(r, c, ld, rd, col)){
        chess[r][c] = ld[r - c + n - 1] = rd[r + c] = col[c] = true;
        nqueen(r + 1, chess, ld, rd, col);
        chess[r][c] = ld[r - c + n - 1] = rd[r + c] = col[c] = false;
    }
}

```

$\mathcal{O}(n)$ extra space

Recursion call
Stack : $\rightarrow O(n)$

$\mathcal{O}(n)$ output

Time
Comp

$$(\mathcal{G})^n + (\mathcal{R})^n$$

\uparrow
isQueenSafe

$$\Rightarrow 4^n + n$$

N Queen Combinⁿ

Time Complexity

Brute force on level $\Rightarrow 2^{n^2} + n^3 \approx 120\text{ ms}$

Backtracking $\Rightarrow 4^n + n^2 \approx 4\text{ ms}$

Branch & Bound $\Rightarrow 4^n + n \approx 3\text{ ms}$

Space Complexity

Brute force on level $\Rightarrow O(n^2)$ Recursion call stack

Backtracking $\Rightarrow O(n)$ Recursion = ..

Branch & Bound $\Rightarrow O(N)$ extra Space

$O(N)$ recursion call stack

$O(n^2)$ output
(chess)

Bit Manipulation $\Rightarrow O(1)$ extra space

$O(N)$ Recursion call stack

④ Bit Manipulation

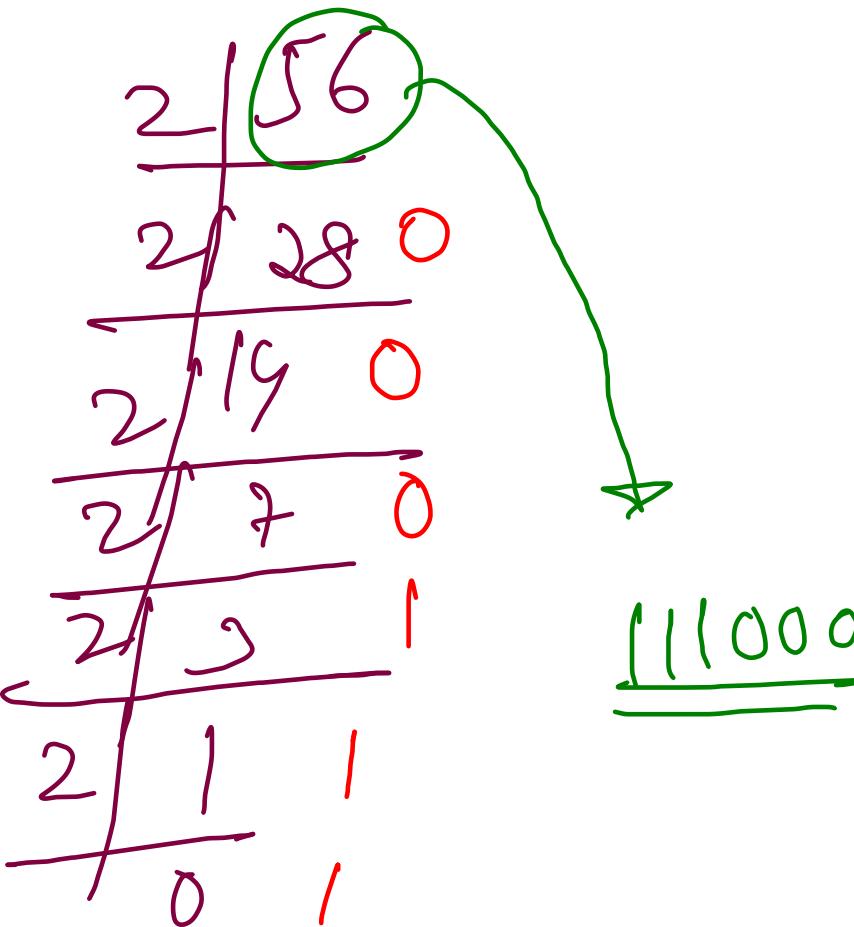
Replacing

boolean visited array by

\downarrow
 $O(n)$ bytes

T/F

bitset
 $\overbrace{\quad}$
 \downarrow
 n bits
 \uparrow {0/1}
 \downarrow
 $O(1)$ space



integer: 4 bytes : 32 bits
 $\overbrace{5 \text{ ld}}$
integer sd

integer^e
 $\overbrace{6 \text{ col}}$

```
public static boolean isQueenSafe(int r, int c, int n, BitSet ld, BitSet rd, BitSet col){
    return ((ld.get(r - c + n - 1) == true) || (rd.get(r + c) == true) || (col.get(c) == true)) ? false : true;
}

public static void nqueen(int r, boolean[][] chess, BitSet ld, BitSet rd, BitSet col) {
    int n = chess.length;
    if (r == n) {
        List<String> ans = new ArrayList<>();
        for (int i = 0; i < n; i++) {ans}
        res.add(ans);
        return;
    }

    for(int c=0; c<n; c++){
        if(isQueenSafe(r, c, n, ld, rd, col)){
            chess[r][c] = true;
            ld.set(r - c + n - 1, true);
            rd.set(r + c, true);
            col.set(c, true);

            nqueen(r + 1, chess, ld, rd, col);

            ld.set(r - c + n - 1, false);
            rd.set(r + c, false);
            col.set(c, false);
            chess[r][c] = false;
        }
    }
}
```

```
bitset<1000000> col,d1,d2;
void solve(int i,int n,int &ans)
{
    if(i == n) {ans++; return;}
    for(int j=0;j<n;j++)
    {
        if(!col[j] && !d1[i-j+n-1] && !d2[i+j])
        {
            col[j] = d1[i-j+n-1] = d2[i+j] = 1;
            solve(i+1,n,ans);
            col[j] = d1[i-j+n-1] = d2[i+j] = 0;
        }
    }
}
class Solution {
public:
    int totalNQueens(int n) {
        if(n <= 1) return n;
        int ans = 0;
        solve(0,n,ans);
        return ans;
    }
};
```

{ N Queen - 11 ?
using C++ }
Count

R&B - level-2 - Lecture 8

Sudoku & Other Puzzles

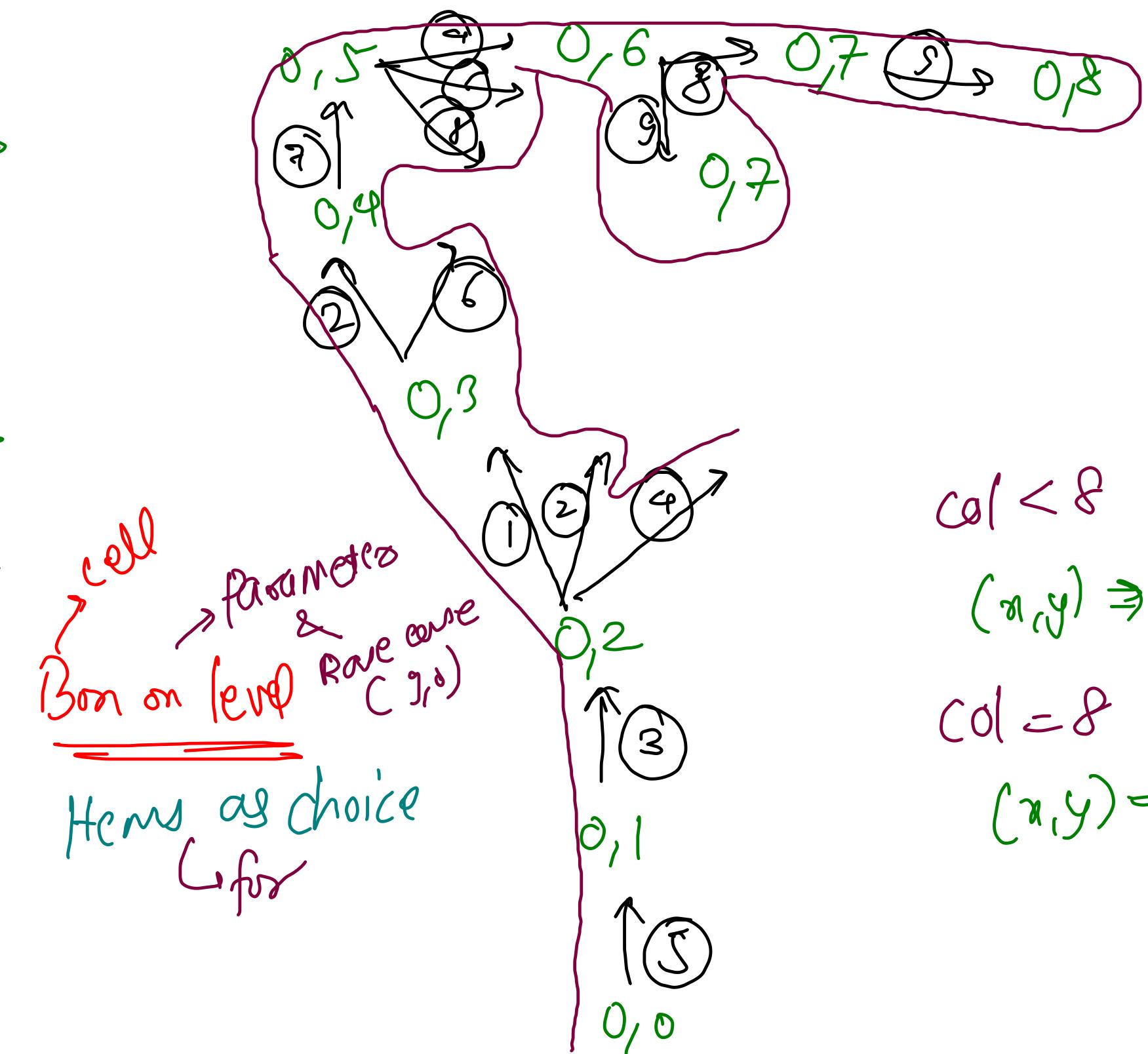
Puzzles	
1	<u>Valid Sudoku (IB)</u>
1	<u>Sudoku Solver (IB)</u>
1	<u>Using Backtracking</u>
1	<u>Using Bit Manipulation</u>
2	<u>Crossword Puzzle</u>
3	<u>Cryptarithmetic Puzzle</u>
3	<u>Cryptarithmetic Puzzle</u>
4	<u>Max Score</u>
4	<u>Max Score</u>
5	<u>N Knight</u> , <u>Knight's Tour</u>

	0	1	2	3	4	5	6	7	8
0	5	3	4	6	7	3	9	1	2
1	6		1	9	5				
2	9	8			6				
3	8		6			3			
4	4	8	3				1		
5	7		2		6				
6	6			2	8				
7		4	1	9			5		
8		8			7	9			

① In each row, we should have 1-9
 \hookrightarrow 9 columns

② In each column, we should have 1-9
 \hookrightarrow 9 rows

③ In each 3×3 matrix, we should have 1-9



	0	1	2	3	4	5	6	7	8
0	5	3		7		1			
1	6			1	9	5			
2		9	8				6		
3	8			6		1		3	
4			8		3				1
5	7			2				6	
6		6		1	2		8		
7				4	1	9			5
8							7	9	

f top-left

val
 0, 1, 2 $\Rightarrow 0 = 0 \times 3 - (\text{val}/3) \times 3$
 3, 4, 5 $\Rightarrow 3 = 1 \times 3 - (\text{val}/3) \times 3$
 6, 7, 8 $\Rightarrow 6 = 2 \times 3 - (\text{val}/3) \times 3$

8, 4 \Rightarrow 6, 3

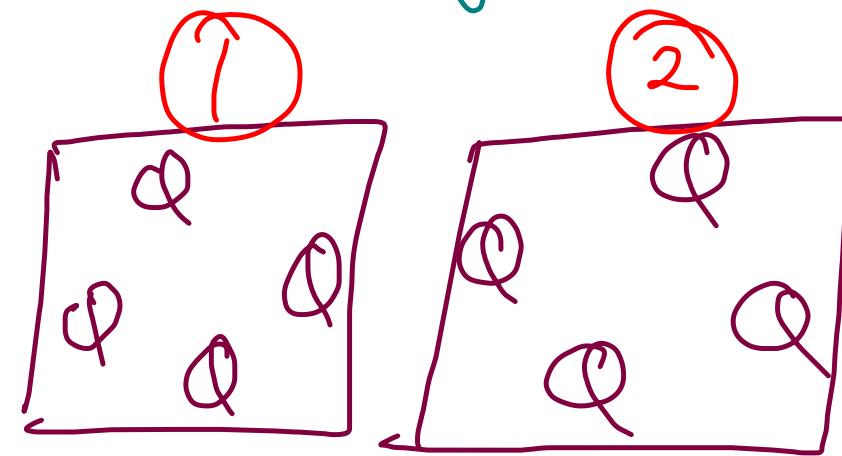
Diagram showing connections between cells:

- Cells (6,3), (6,4), (6,5) are connected to (6,3).
- Cells (7,3), (7,4), (7,5) are connected to (7,3).
- Cells (8,3), (8,4), (8,5) are connected to (8,3).
- Cell (7,3) is connected to (7,4) and (7,5).
- Cell (8,3) is connected to (8,4) and (8,5).
- Cell (6,3) is connected to (7,3) and (8,3).
- Cell (7,3) is connected to (6,3), (7,4), and (7,5).
- Cell (8,3) is connected to (6,3), (8,4), and (8,5).
- Cell (6,3) is connected to (6,4) and (6,5).
- Cell (7,3) is connected to (7,4) and (7,5).
- Cell (8,3) is connected to (8,4) and (8,5).
- Cell (6,3) is connected to (6,3).
- Cell (7,3) is connected to (7,3).
- Cell (8,3) is connected to (8,3).

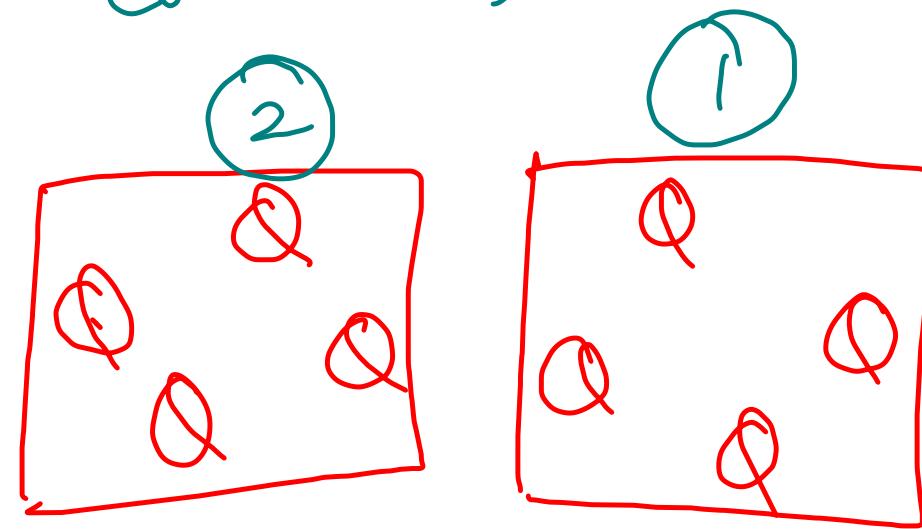
rw + 2

nqueen

row by row



column by column



Max Score

```
Input: words = ["dog", "cat", "dad", "good"], letters =  
["a", "a", "c", "d", "d", "g", "o", "o"], score =  
[1, 0, 9, 5, 0, 0, 3, 0, 0, 0, 0, 0, 2, 0, 0, 0, 0, 0, 0, 0, 0, 0]  
Output: 23
```

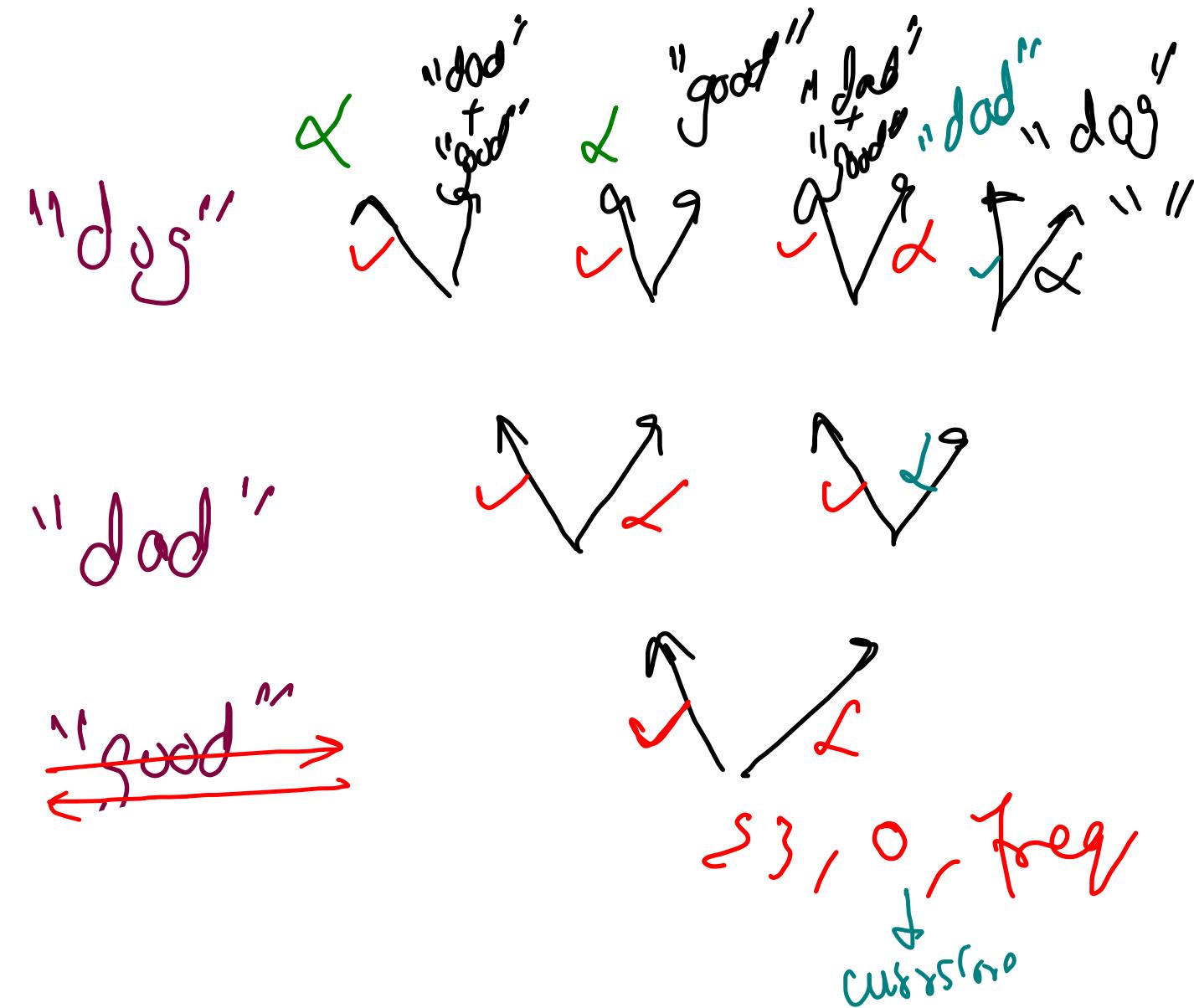
a-2 / 1 / 2

c-1

d-3 / 0 / 2 / 1 / 0 / 3

g-1 / 0 / 0 /

o-2 / 0 / 1 / 2 / 1



```
public int maxScore(int idx, String[] words, int[] freq, int[] score){
    if(idx == words.length) return 0;

    // yes
    int currScore = 0;
    boolean flag = true;
    for(char ch: words[idx].toCharArray()){
        freq[ch - 'a']--;
        if(freq[ch - 'a'] < 0){
            flag = false;
        }

        currScore += score[ch - 'a'];
    }

    if(flag == true){
        currScore += maxScore(idx + 1, words, freq, score);
    } else {
        currScore = 0;
    }

    // backtrack
    for(char ch: words[idx].toCharArray()){
        freq[ch - 'a']++;
    }

    // no
    return Math.max(currScore, maxScore(idx + 1, words, freq, score));
}
```

Move Zeros to End

{ 0, 1, 3, 12 }
3, 12 }

→ { 1, 3, 12, 0, 0 }

$f(0,0) \rightarrow f(1,0) \rightarrow f(2,1) \rightarrow f(3,1) \rightarrow f(4,2) \rightarrow f(5,3)$

```
public static void moveZeroes(int idx, int countOfNonZero, int[] nums) {
    if(idx == nums.length){
        int countOfZeros = nums.length - countOfNonZero;
        for(int i=0; i<countOfZeros; i++){
            nums[nums.length - 1 - i] = 0;
        }
        return;
    }

    if(nums[idx] != 0){
        nums[countOfNonZero] = nums[idx];
        moveZeroes(idx + 1, countOfNonZero + 1, nums);
    }
    else {
        moveZeroes(idx + 1, countOfNonZero, nums);
    }
}
```

Today's Questions

- ① Crossword
- ② Word Search - (I, II, III)
- ③ Cryptarithmefic
- ④ Josephus

Cryptarithm

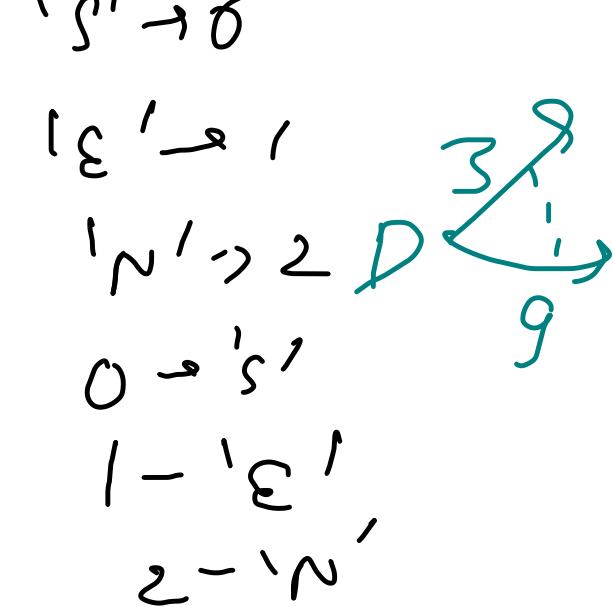
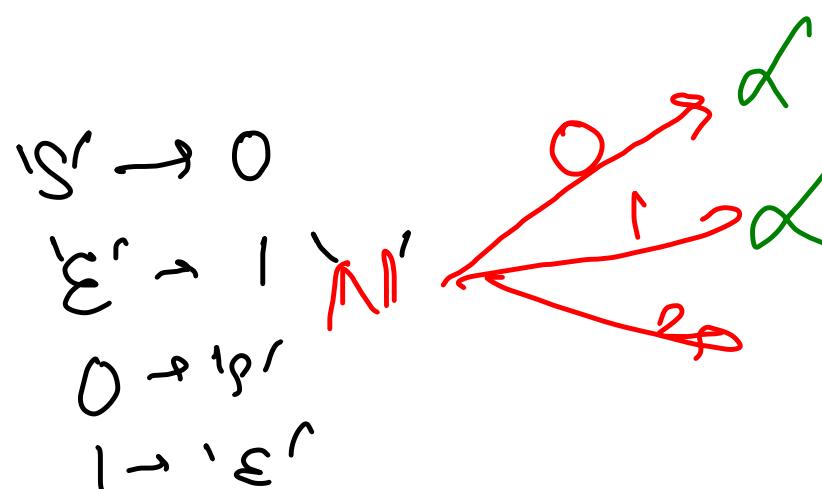
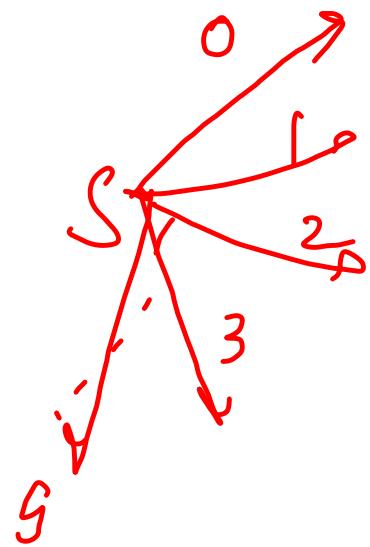
$$\text{"SEND"} + \text{"MORE"} = \text{"MONEY"}$$

1:1 mapping
 ↓ ↓
 char digit
 (A-Z) (0-9)

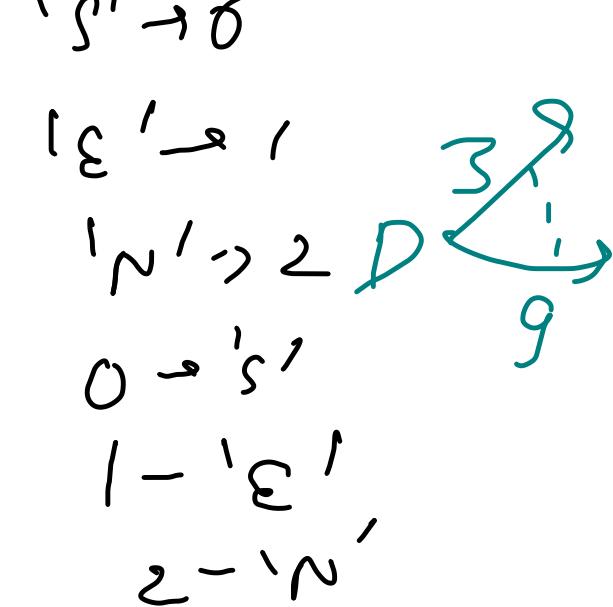
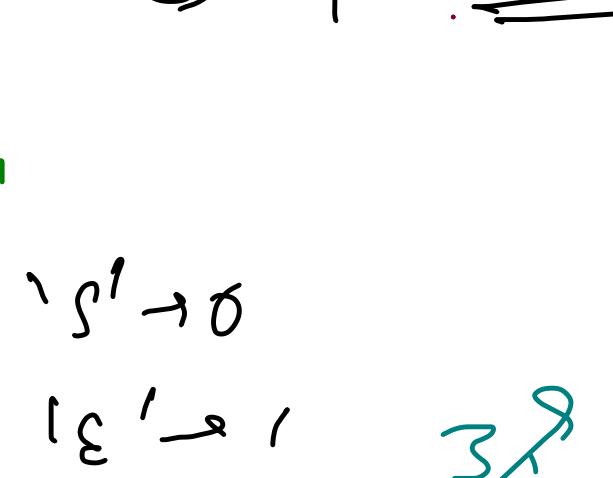
unique
char
on
level

{ 'S', 'E', 'N', 'D', 'M', 'O', 'R', 'Y' }
 base are

↳ equivalence check



{ 'S', 'E', 'N', 'D', 'M', 'O', 'R', 'Y' }
 base are

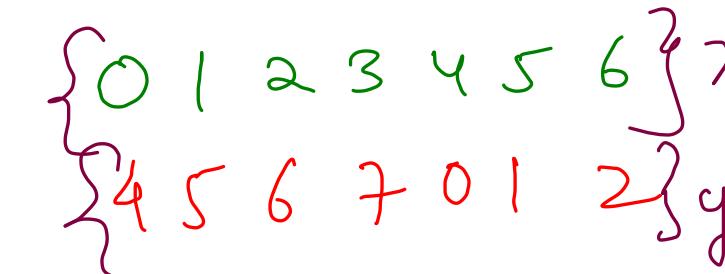
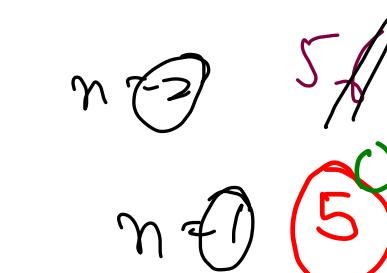
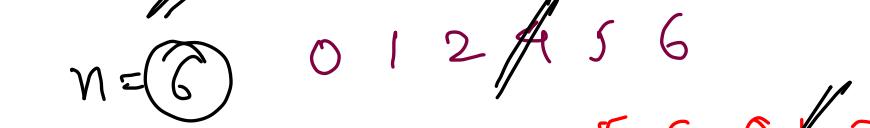
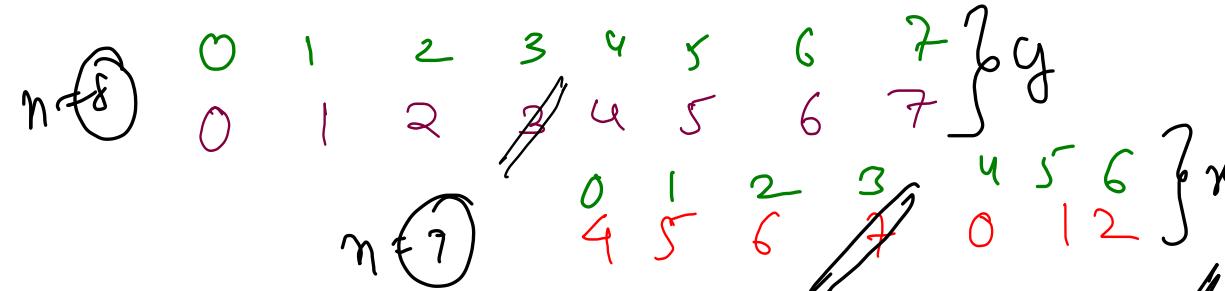


'S' → 0
 'E' → 1
 'O' → 2
 1 → 'E'

'S' → 0
 'E' → 1
 'N' → 2
 0 → 'S'
 1 → 'E'
 2 → 'N'

Josephus

$k=4$



$n=8$
 $k=4$

actual value
↓
 y index
↓
 x

$$= (x + k) / n$$

shifting on \odot
rotation on \odot

Mechan
Exp

$$(josephus(n-1, k) + k) \cdot f(n)$$

```

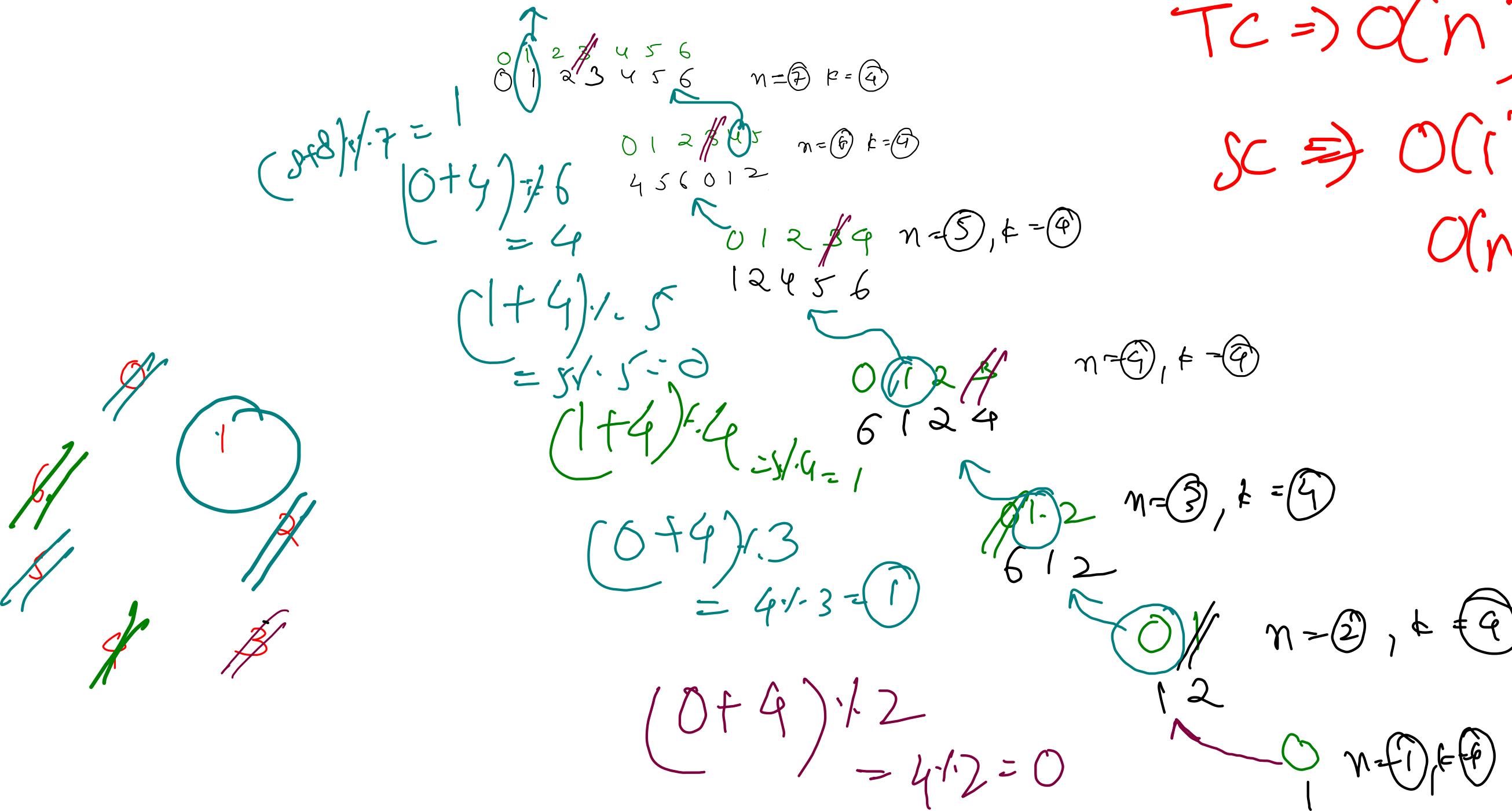
public static int solution(int n, int k){
    if(n == 1) return 0; // return index
    return (solution(n - 1, k) + k) % n;
}

```

$Tc \Rightarrow O(n)$

$Sc \Rightarrow O(1)$ extra space

$O(n)$ recursion



```
public static int solution(int n, int k){  
    Queue<Integer> q = new ArrayDeque<>();  
    for(int i=0; i<n; i++){  
        q.add(i);  
    }  
    int count = 0;  
    while(q.size() > 1){  
        int val = q.remove();  
        count++;  
  
        if(count == k){  
            count = 0;  
        } else {  
            q.add(val);  
        }  
    }  
    return q.peek();  
}
```

constructive
algorithm

Queue { Brute force }

$O(n \times R) \Rightarrow$ Time

$O(n)$ \Rightarrow Queue
Space

henicographical Nos

1 ↙

10 100 101 102 103 104 ... 109

11 110 111 112 113 ... 119

12 120 121 122 ... 129

13 191 192 ... 199

$$N = 999$$

2 ↙

20 200 201 202 ... 209

21 210 211 ... 219

22 220 221 ... 229

23 230 231 ... 239

24 240 241 ... 249

25 250 251 ... 259

26 260 261 ... 269

27 270 271 ... 279

28 280 281 ... 289

29 290 291 ... 299

3

30 → 301 ... 309

31 → 310 ... 319

32 -

⋮

39 → 390 ... 399

4

40, 41, 42, ... 49

5

50, ... 59

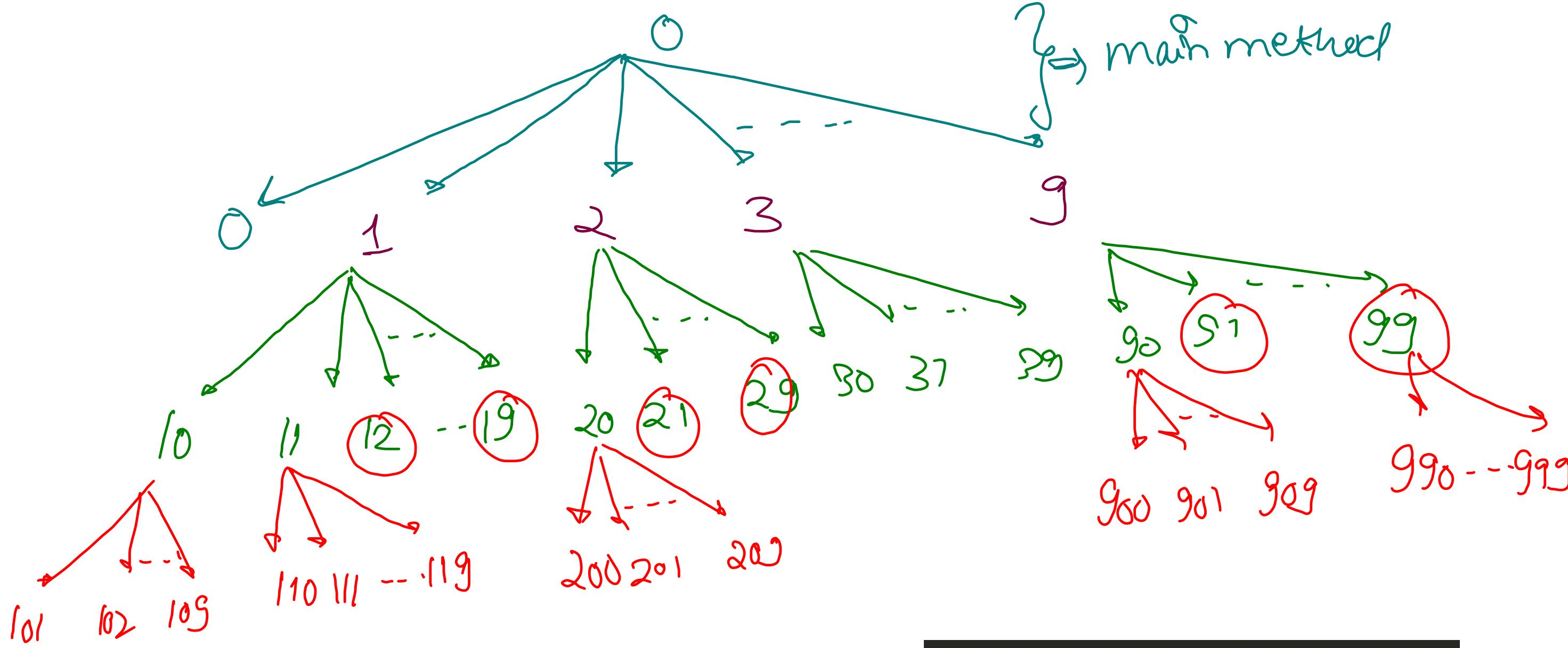
6

7

8

9

90, 91, 92, ... 99
900, 901, ... 909
990-999



```

public void lexicalOrder(int input, int n){
    if(input > n) return;

    res.add(input);
    for(int i=0; i<10; i++){
        lexicalOrder(input * 10 + i, n);
    }
}

```

R&B - Level 2 - Lecture 10

- ① Crossword
- ② Word search - II
- ③ Word Search - III { Homework }
- ④ Friends Pairings
- ⑤ Largest No in k swaps
- ⑥ Restore IP Address

More questions to practice

- Paranthesis Addition Ways
- Stepping Numbers
- Kth Symbol in Grammar
- Split into Fibonacci Sequence
- Distinct Subsets Permutations

N knights

Crossword

