

## # Hashmap

- » Hashmap - Introduction
- «/» Highest Frequency Character
- «/» Get Common Elements - 1
- «/» Get Common Elements - 2
- «/» Longest Consecutive Sequence Of Elements
- «/» Write Hashmap

## # Heap/Priority Queue

- » Heaps - Introduction And Usage
- «/» K Largest Elements
- » Efficient Heap Constructor
- «/» Write Priority Queue Using Heap
- «/» Sort K-sorted Array
- » Heap - Comparable V/s Comparator
- «/» Merge K Sorted Lists
- «/» Median Priority Queue

## #HashMap

Insert → put  $\Rightarrow O(1)$

Update

Delete  $\rightarrow$  remove  $\Rightarrow O(1)$

Read  $\rightarrow$  get  $\Rightarrow O(1)$

If key exist then  
return value  
else return null

Display

size  $\Rightarrow O(1)$

contains key  $\rightarrow$  find  $\Rightarrow O(1)$

keyset  $\rightarrow$  keys

Key  $\rightarrow$  value

e.g. IPL teams  $\rightarrow$  trophies count  
String  $\rightarrow$  int

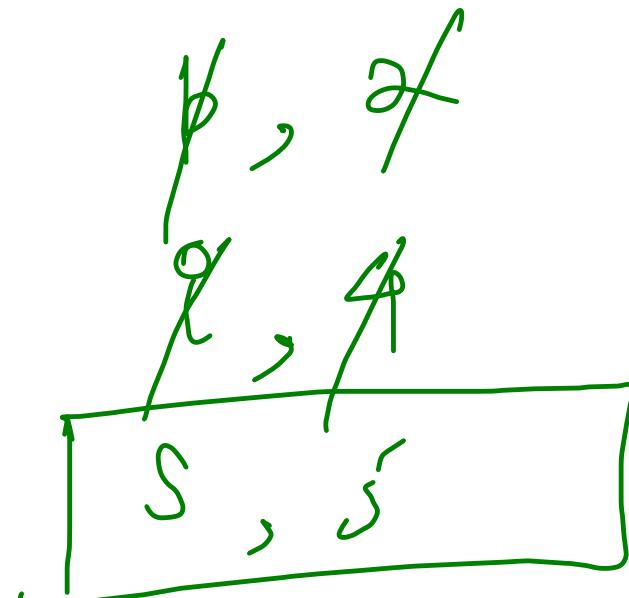
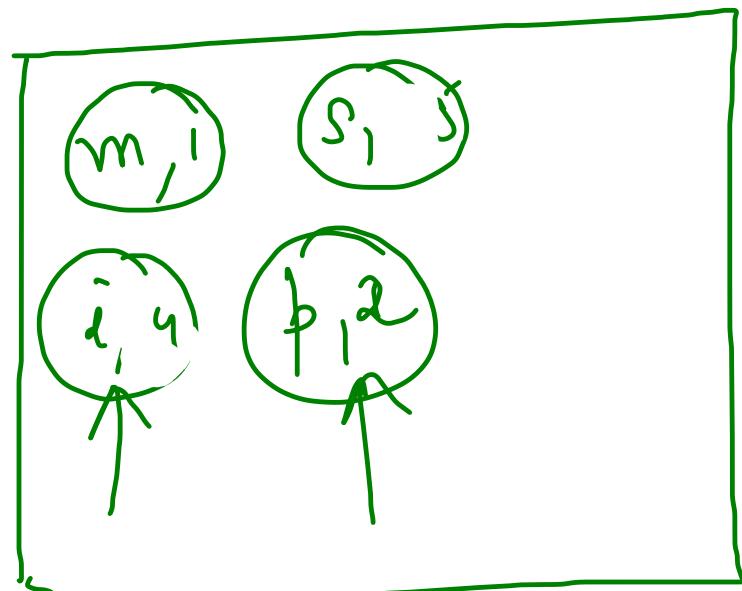
{ CSK  $\rightarrow$  4  
MI  $\rightarrow$  5  
SRH  $\rightarrow$  2  
DC  $\rightarrow$  0

e.g. countries  $\rightarrow$  population  
String  $\rightarrow$  integer

{ India  $\rightarrow$  130  
USA  $\rightarrow$  90  
China  $\rightarrow$  200

# Highest Frequency Character

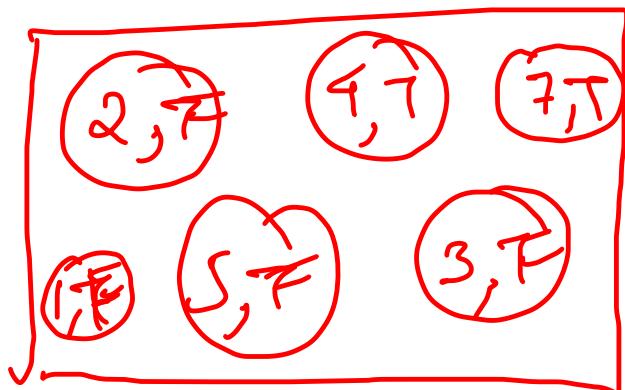
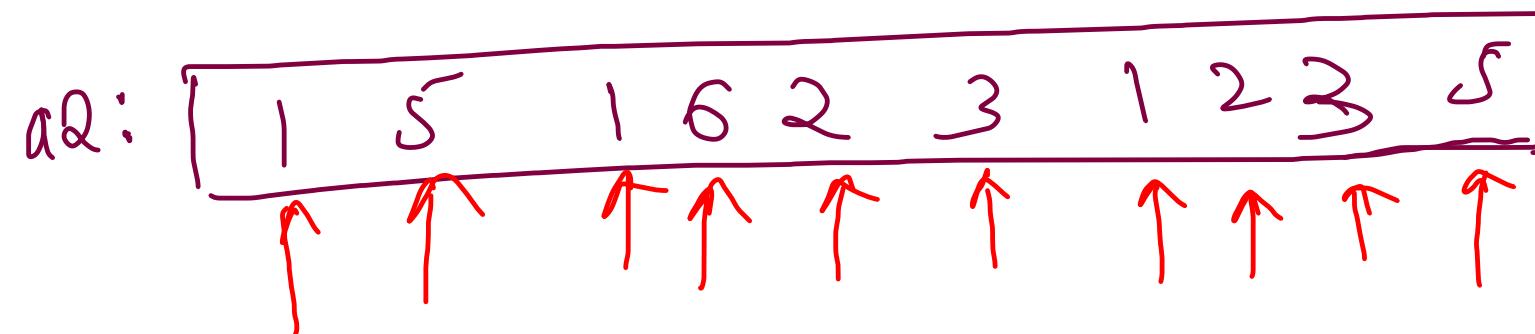
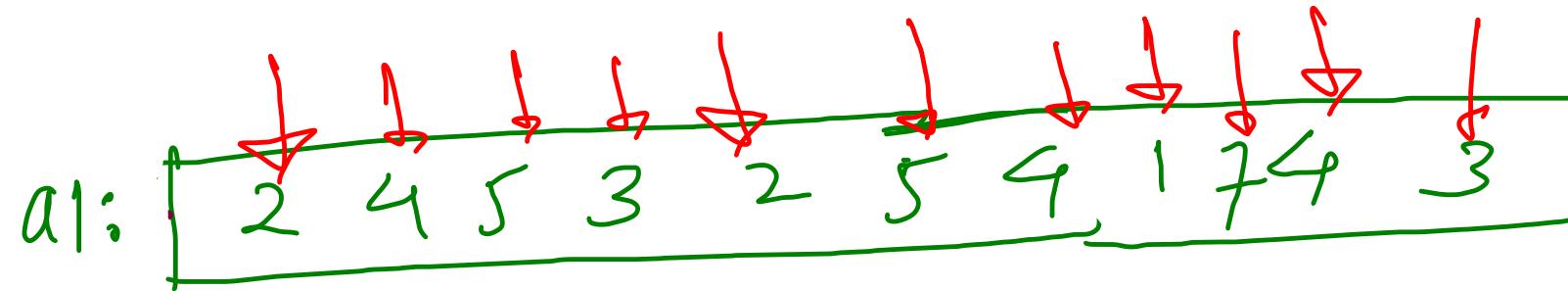
m i s s o s i p b o x f



```
HashMap<Character, Integer> freq = new HashMap<>();  
  
for(int i=0; i<str.length(); i++){  
    char ch = str.charAt(i);  
    if(freq.containsKey(ch)){  
        int oldFreq = freq.get(ch);  
        freq.put(ch, oldFreq + 1);  
    }  
    else {  
        freq.put(ch, 1);  
    }  
  
    char ch = str.charAt(0);  
    int maxFreq = freq.get(ch);  
  
    for(Character key: freq.keySet()){  
        int currFreq = freq.get(key);  
  
        if(currFreq > maxFreq){  
            ch = key;  
            maxFreq = currFreq;  
        }  
    }  
  
    System.out.println(ch);
```

The code is annotated with time complexities:  
- The first loop is labeled  $O(N)$ .  
- The nested loop is labeled  $O(\text{Keys})$ .  
- The final  $\maxFreq$  assignment is labeled  $= O(256)$ .

# # Get Common Elements ↗



$O(n)$

$\langle \text{Integer}, \text{Boolean} \rangle$

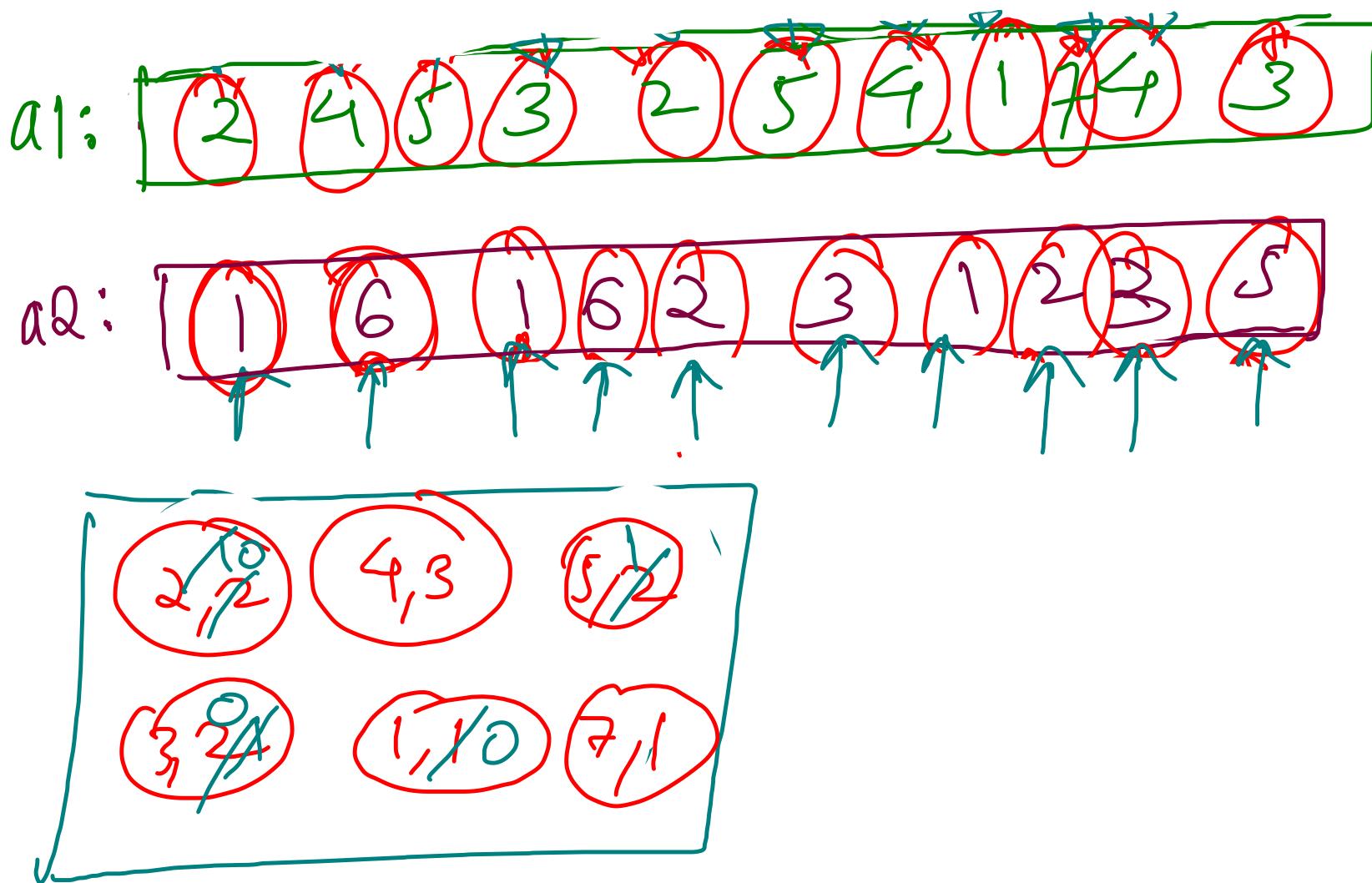
```
HashMap<Integer, Boolean> hm = new HashMap<>();
for(int i=0; i<n1; i++)
    hm.put(arr1[i], true);

for(int i=0; i<n2; i++){
    if(hm.containsKey(arr2[i]) && hm.get(arr2[i])){
        System.out.println(arr2[i]);
        hm.put(arr2[i], false);
    }
}
```

$\mathcal{O}(n_1 + n_2)$

1, 5, 2, 3

## # Get Common Elements - 2

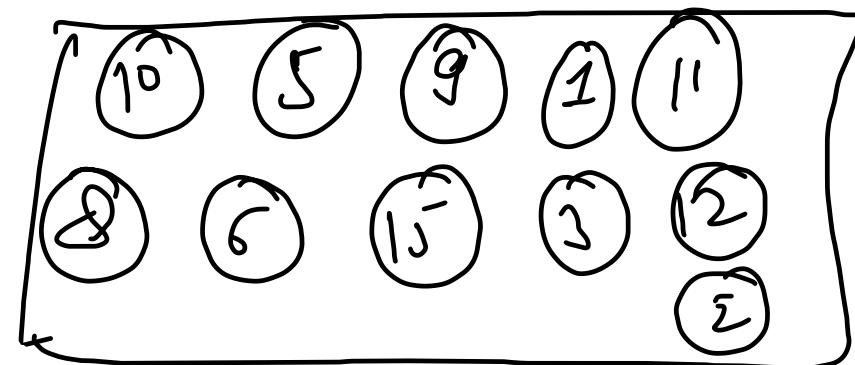
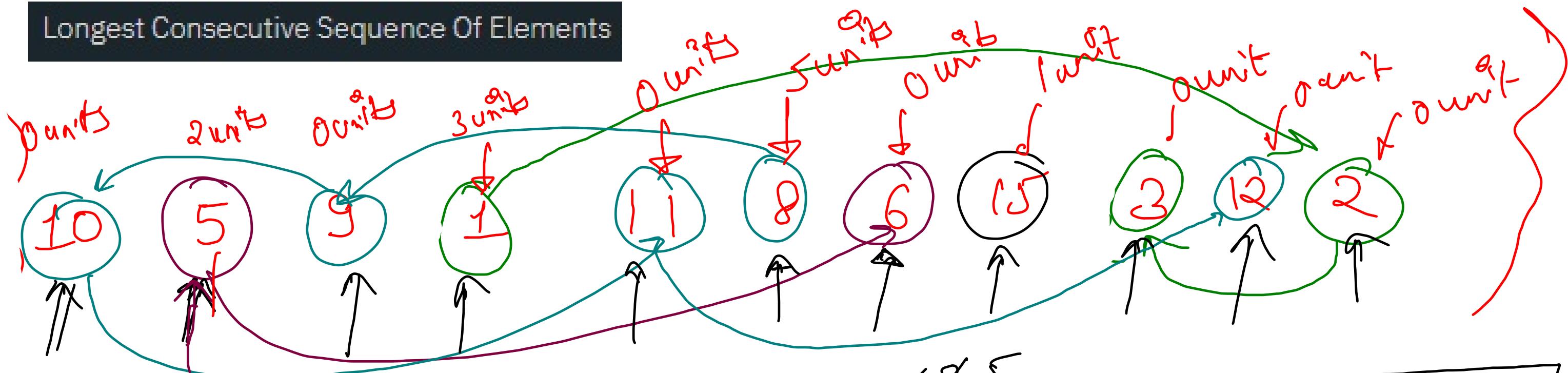


1, 2, 3, 2, 3, 5

```
HashMap<Integer, Integer> hm = new HashMap<>();
for(int i=0; i<n1; i++){
    if(hm.containsKey(arr1[i])){
        hm.put(arr1[i], hm.get(arr1[i]) + 1);
    } else {
        hm.put(arr1[i], 1);
    }
}

for(int i=0; i<n2; i++){
    if(hm.containsKey(arr2[i]) && hm.get(arr2[i]) > 0){
        System.out.println(arr2[i]);
        hm.put(arr2[i], hm.get(arr2[i]) - 1);
    }
}
```

## Longest Consecutive Sequence Of Elements



```

int maxChain = 0;
int startingPt = 0;

for(Integer key: hm.keySet()){
    if(hm.containsKey(key - 1) == false){
        // chain starting pt
        int length = 1;
        while(hm.containsKey(key + length) == true){
            length++;
        }
        if(length > maxChain){
            maxChain = length;
            startingPt = key;
        }
    }
}

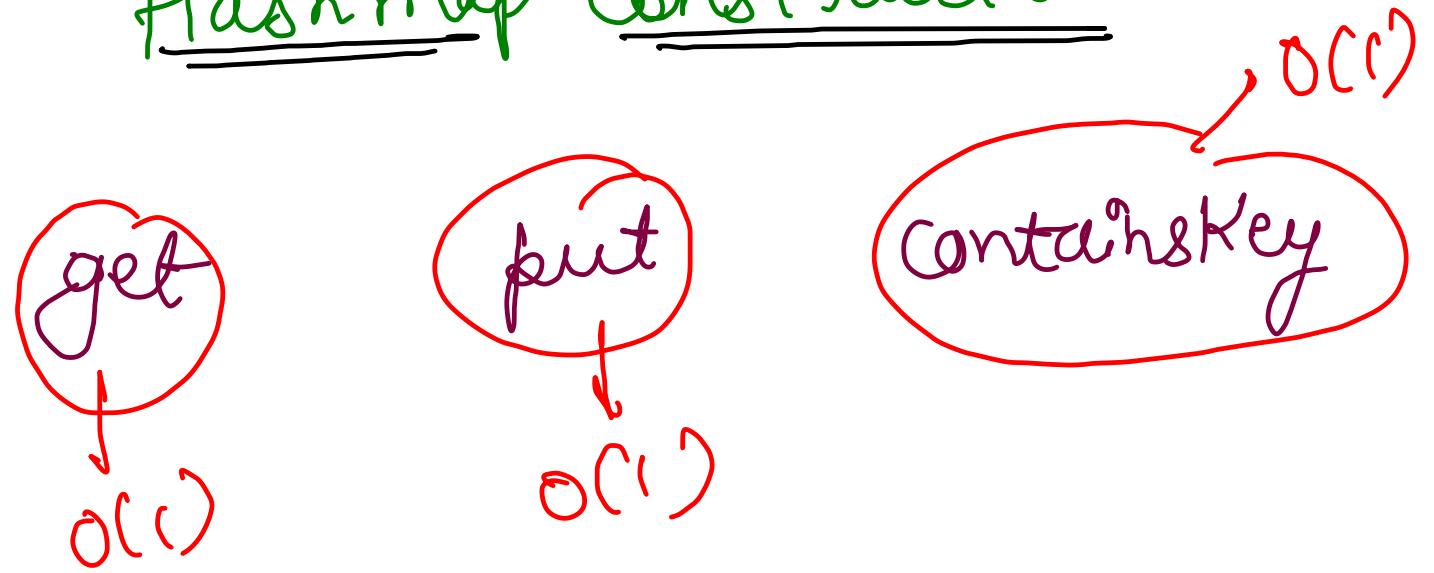
```

$2 \rightarrow 3$   
 $8 \rightarrow 9 \rightarrow 10 \rightarrow 11 \rightarrow 12$

$5 \rightarrow 6$   
 $15$

$O(N)$  Time

## HashMap Construction

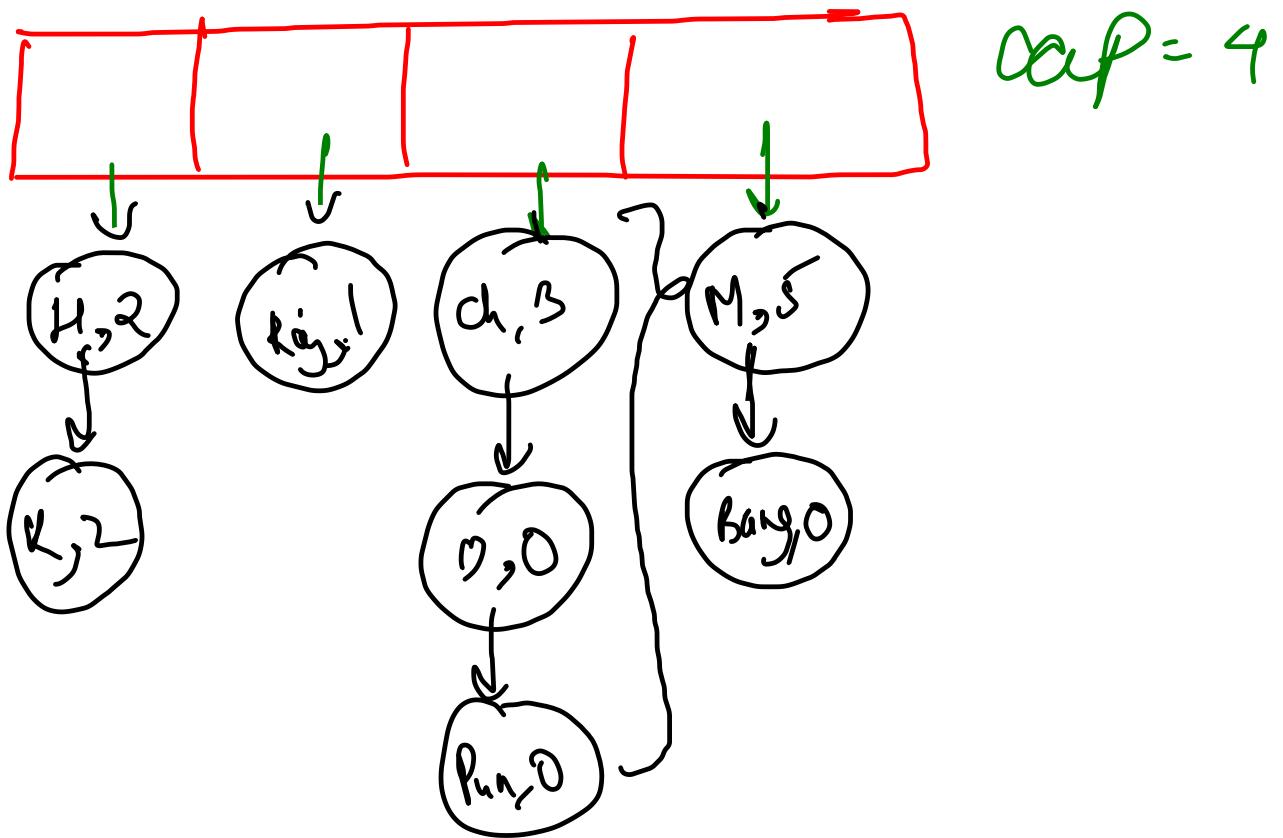


keySet

size

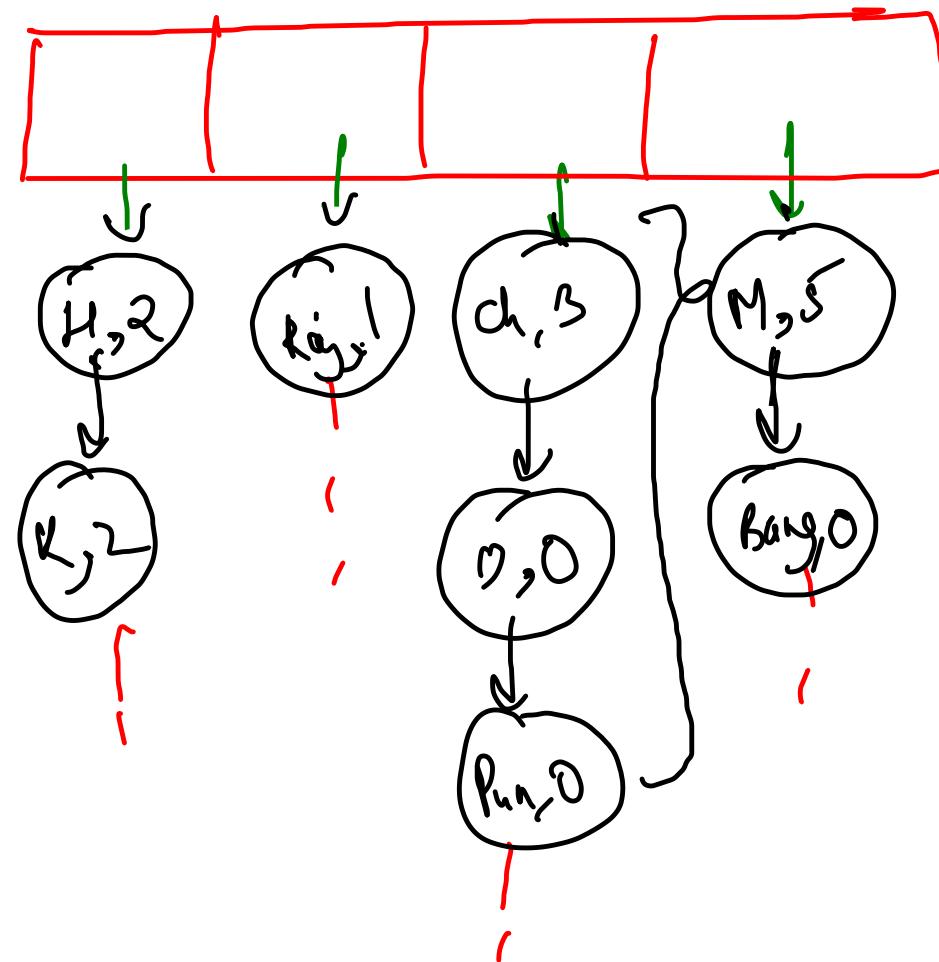
display

hashfunction

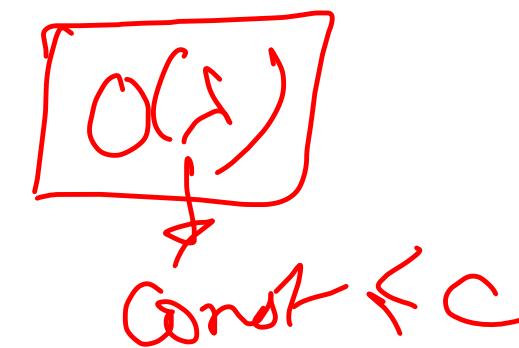


→ Two keys can have same hashfn

→ A given key will never have two hashfn values.



$\text{cap} = 4$

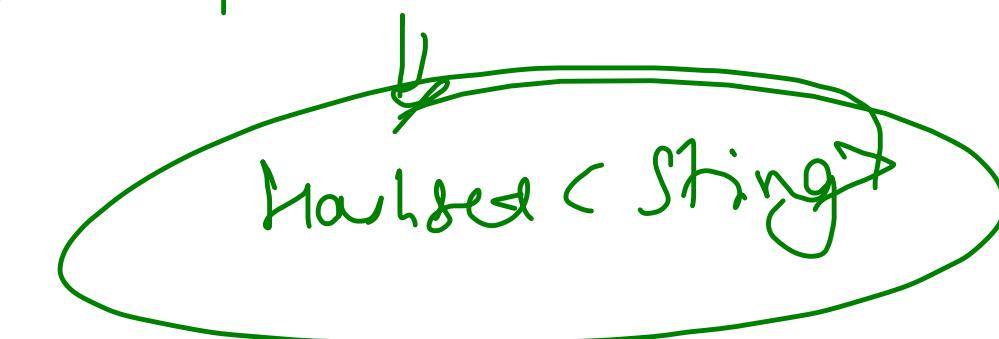


loading factor  $\frac{1}{4}$   
 $= \frac{\text{no of nodes}}{\text{in one bucket}}$   
 $= \frac{\text{total nodes}}{\text{no of buckets}}$

$$= \frac{8}{4}$$

$$= 2$$

HashMap < String, Boolean >



```

private LinkedList<HMNode>[] buckets;
private int noOfNodes;
private int noOfBuckets;

public HashMap(){
    noOfBuckets = 4;
    noOfNodes = 0;
    buckets = new LinkedList[noOfBuckets];

    for(int i=0; i<noOfBuckets; i++){
        buckets[i] = new LinkedList<>();
    }
}

```

```

public void put(K key, V value) throws Exception {
    // O(1)
    int bucketId = getBucketId(key);
    HMNode data = getData(bucketId, key);

    if(data == null){
        // Insertion
        HMNode node = new HMNode(key, value);
        buckets[bucketId].addLast(node);
        noOfNodes++;
    } else {
        // Updation
        data.value = value;
    }
}

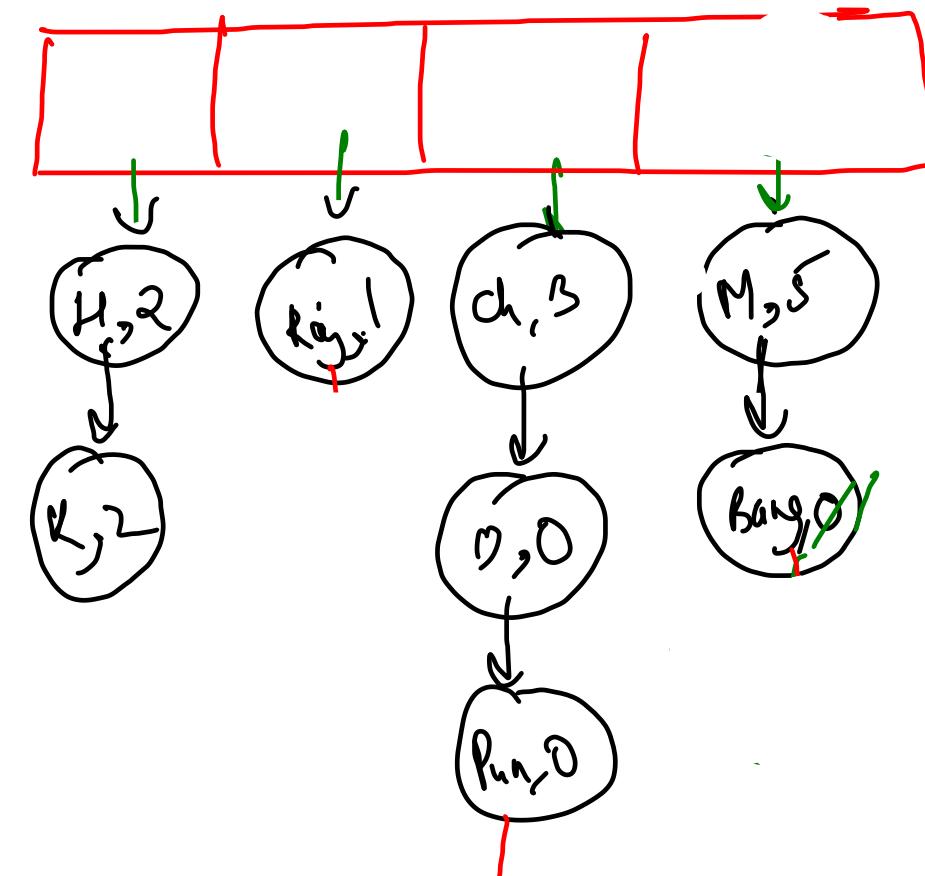
```

```

public int getBucketId(K key) throws Exception{
    // O(1)
    int hashCode = key.hashCode();
    int bucketId = (Math.abs(hashCode)) % noOfBuckets;
    return bucketId;
}

private HMNode getData(int bucketId, K key) throws Exception{
    for(HMNode node: buckets[bucketId]){
        if(node.key.equals(key) == true){
            return node; // data already exist
        }
    }
    return null; // data not exist
}

```



$\text{Cap} = 9$

$\text{Size} = 18$

$H = 815$

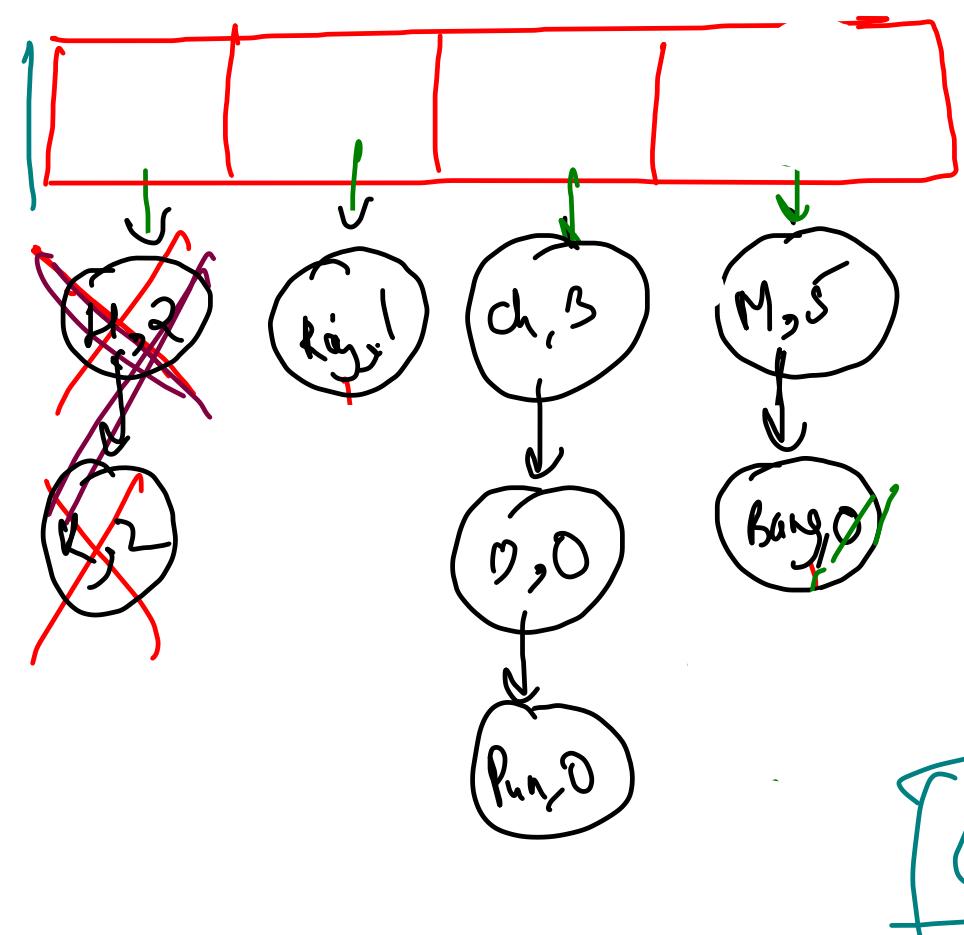
$\text{Cap}$

$O(2)$

$O(C)$

Ahm

$h \cdot 4 = O(1, 2, 3)$



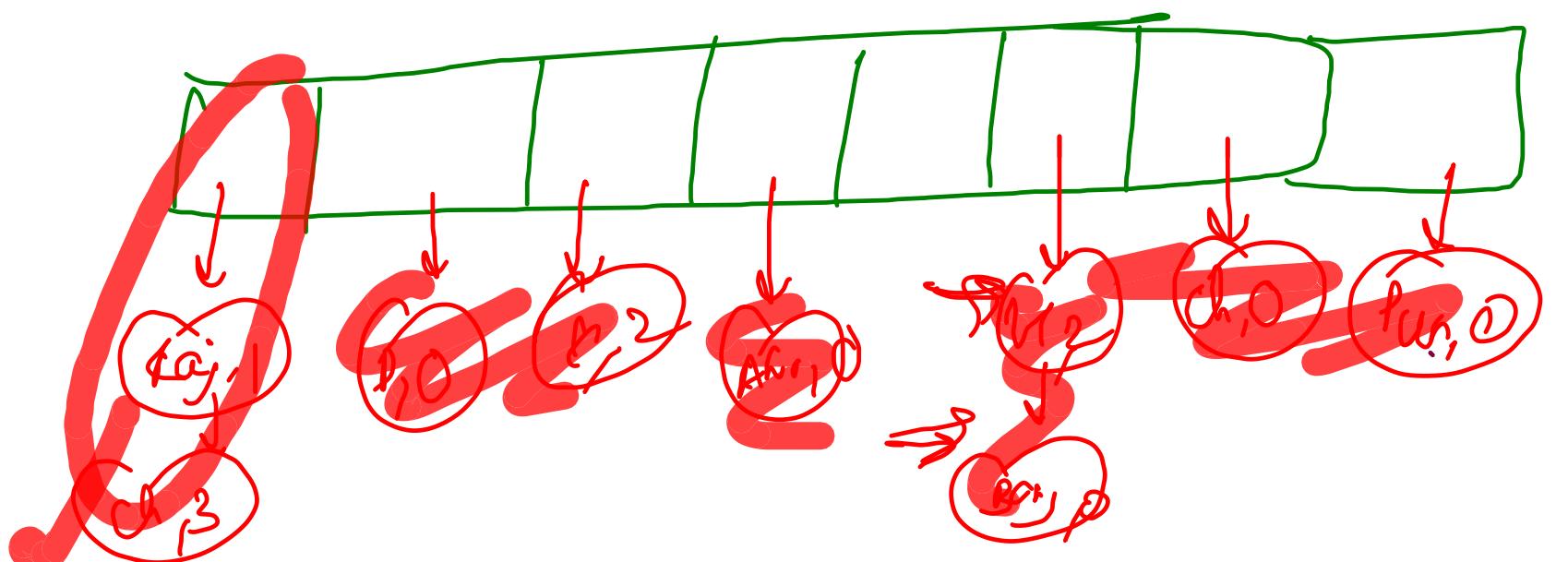
$\text{size} = 8$   
 $\text{cap} = 9$

$O(N)$

Rehashing

$$\lambda = \frac{8}{9} = 0.888\ldots$$

~~$O(c)$~~



hashcode = Hyp  
 $(h1) \mod 4$

$$\lambda = \frac{9}{8} = 1.125\ldots$$

$(h1) \leq 8$

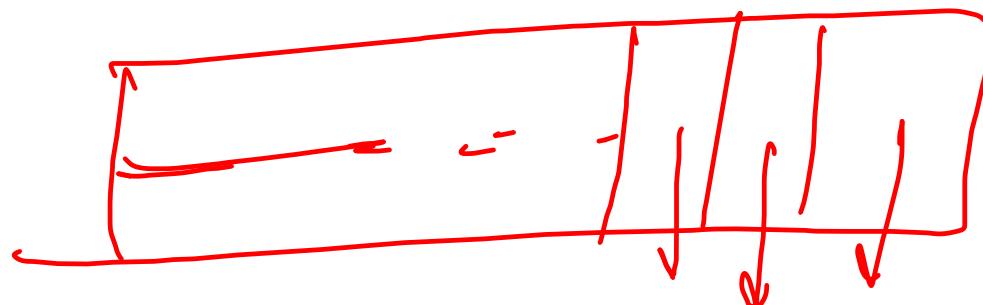
# Time Comp

→ Normal  
→ Rehashing

→ worst case

## Leading factors

Collision  $O(w)$



9<sup>..</sup>;  
931.  
92/1

C - 99°

Collisions of

100

93

1) 911 (889)

1 2 3 5 6 0 7

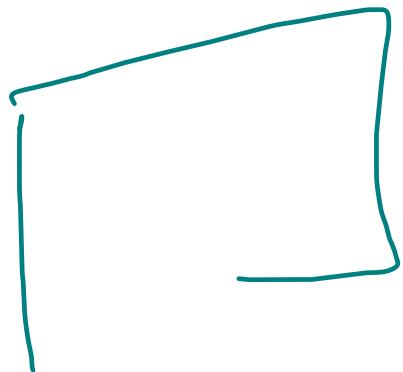
• 101

10 of 10

A hand-drawn red rectangle with a small circle inside the top-left corner.

19 'fro  
7, 8, 9

$\leftarrow$  OCI



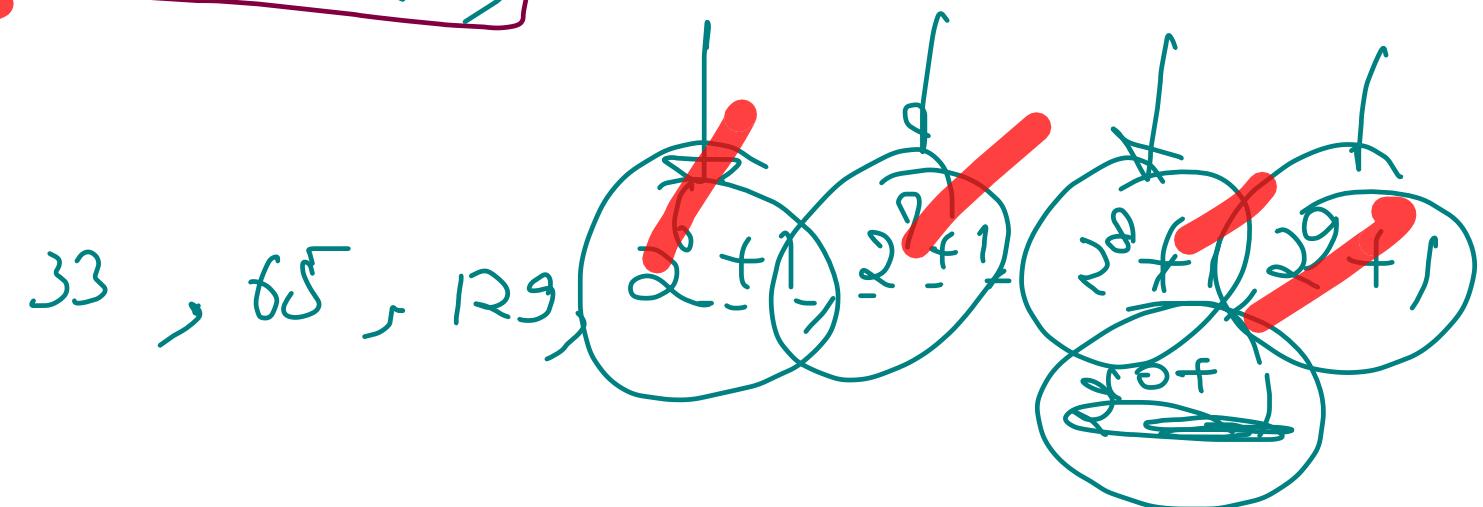
~~1st  $\rightarrow O(1)$~~   
~~2nd  $\rightarrow O(1)$~~   
~~3rd  $\rightarrow O(1)$~~   
~~4th  $\rightarrow O(1)$~~   
~~5th  $\rightarrow O(8)$~~   
~~6th  $\rightarrow O(1)$~~   
~~7th  $\rightarrow O(1)$~~

~~8  $\rightarrow O(1)$~~   
~~9  $\rightarrow O(16)$~~   
~~10  $\rightarrow O(1)$~~   
~~11  $\rightarrow O(32)$~~

Amortized

~~1 + 1 + 1 + 1 + 8 + 1 + ...~~  
~~+ 16 + ...~~  
~~+ 32 + ...~~

$$2^{10} + 1 = O(n)$$



✓ String  $\Rightarrow$

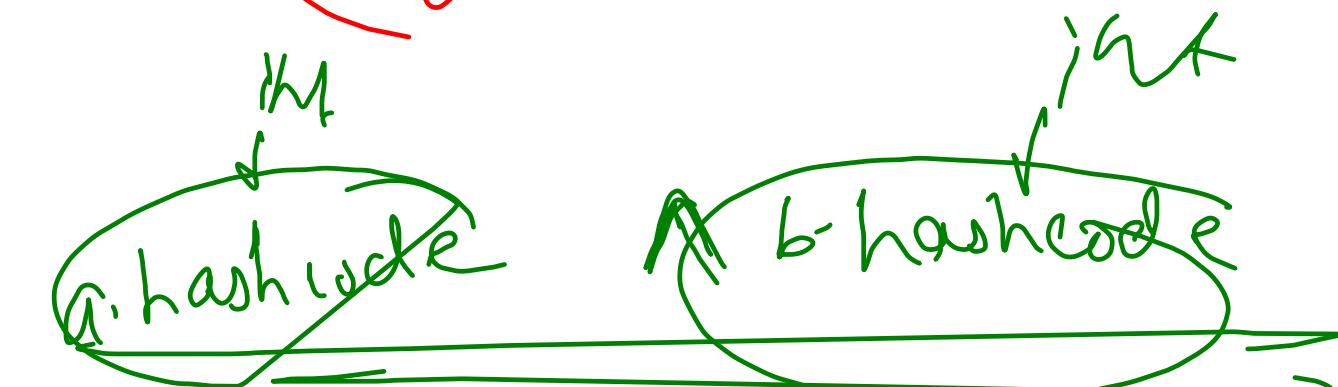
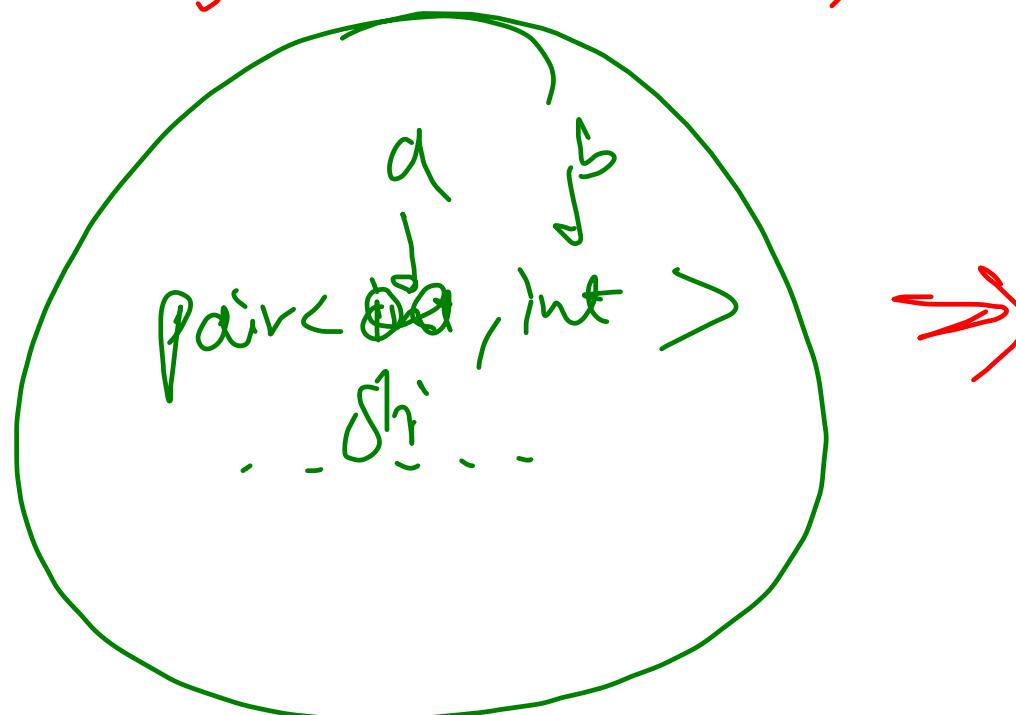
bepcoder  
0 1 2 3 4 5 6 7

$O \times P$   
 $+ 1 \times e$

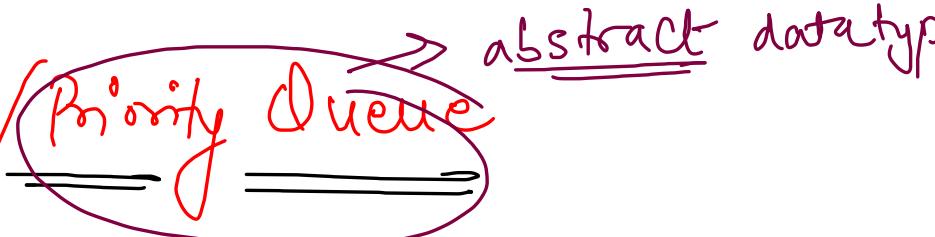
$+ 2 \times P + 3 \times C$   
 $+ \dots ) / f$

✓ int  $\Rightarrow R_1, R_2$

$\rightarrow (\text{key } \times P_1) \times P_2$



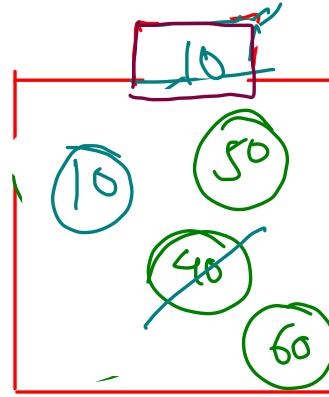
Data Structure



Priority Queue

- add →  $O(\log n)$
- removal →  $O(\log n)$
- peek →  $O(1)$

f+ Min Heap



- pq.remove()
- pq.remove()
- pq.remove()
- pq.remove()

$$n! = n^n$$

Student  
→ roll no  
→ marks  
→ height

- size { $O(1)$ }
- display { $O(n)$ }

{30, 50, 40, 20, 60}

(10)

Inserting  $n$  elements

$$\begin{aligned} & \log 1 + \log 2 + \log 3 \\ & + \dots \log n \\ & - \log (1 \times 2 \times 3 \times \dots \times n) \\ & = \log(n!) \\ & = \log(n^n) \end{aligned}$$

for inserting  
 $n$  elements  
one by one

$n \log n$

Deleting  $n$  elements

$$\begin{aligned} & \log n + \log(n-1) + \log(n-2) \\ & + \dots \log 1 \quad \text{for removing } n \text{ elements} \\ & = n \log n \end{aligned}$$

Binary Heap :- Array

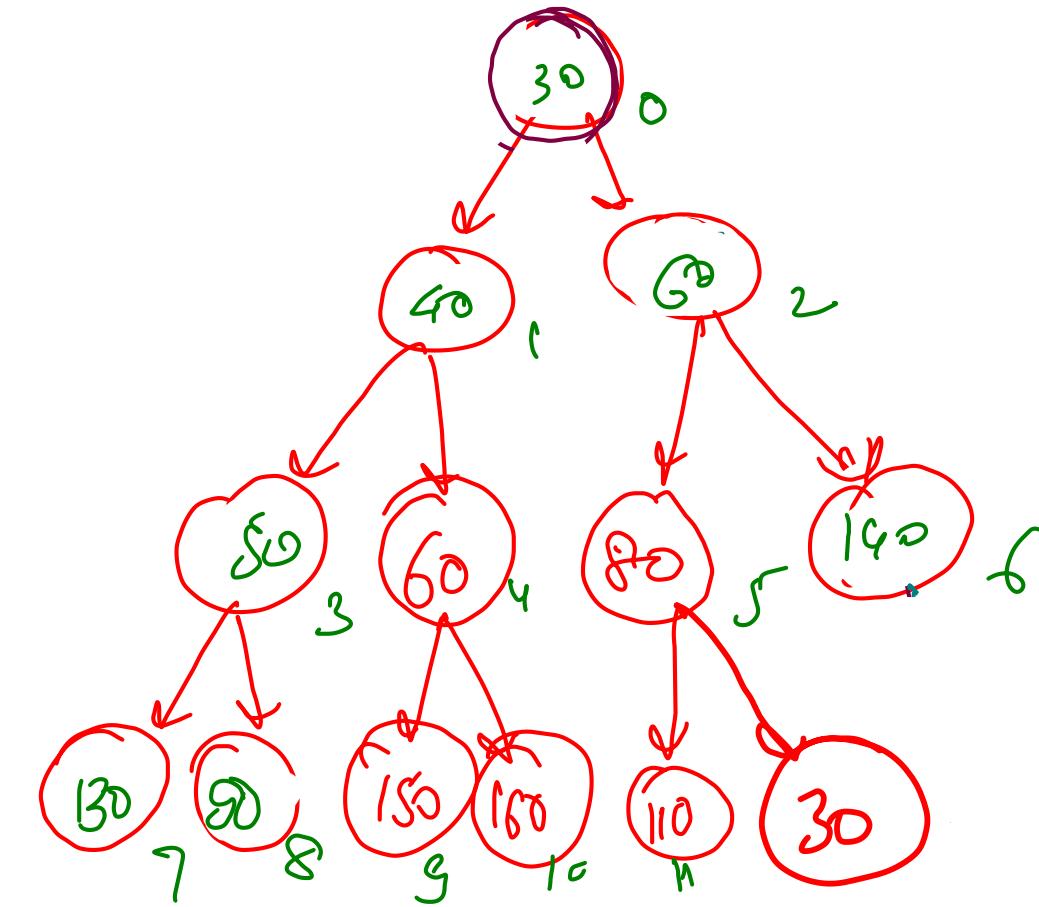
( $\Rightarrow$  complete Binary Tree)  $\Rightarrow$

$\checkmark$  {Heap Order Property}

→ left child :-  $2^i + 1$

→ right child :-  $2^i + 2$

→ parent :-  $(i-1)/2$



$$N \log(N) = \underline{N} \underline{\log N}$$

heapSort

5

10

20

30

40

60

Quick Sort :- ~~Stack~~

Merge Sort :- extra space

Heap Sort :- O(1) extra space

```

public void upheapify(int idx){
    int parIdx = (idx - 1) / 2;

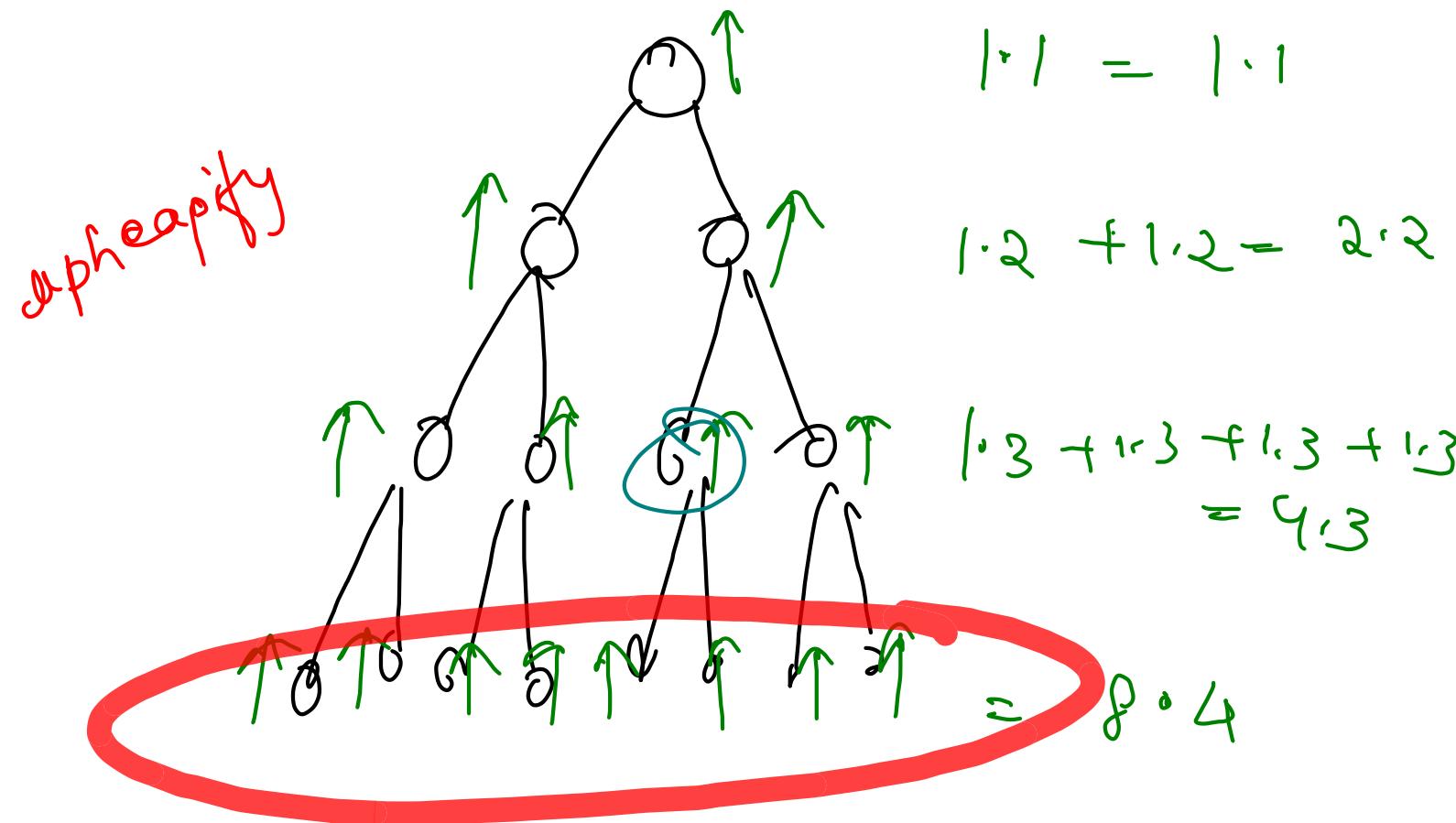
    if(isSmaller(idx, parIdx)){
        swap(idx, parIdx);
        upheapify(parIdx);
    }
}

```

```

public void add(int val) {
    // O(log n)
    data.add(val);
    upheapify(size() - 1);
}

```



insertion of  
N nodes

$$\left\{ 2^0 \cdot 1 + 2^1 \cdot 2 + 2^2 \cdot 3 + 2^3 \cdot 4 + 2^4 \cdot 5 + 2^5 \cdot 6 \dots \right\} = N \log N$$

3yoda  
nodes,  
3yada  
room  
↓  
 $O(n \log n)$   
3yada  
T.C.

```

public void downheapify(int idx){
    int min = idx;
    int left = 2 * idx + 1;
    int right = 2 * idx + 2;

    if(left < size() && isSmaller(left, min)){
        min = left;
    }

    if(right < size() && isSmaller(right, min)){
        min = right;
    }

    if(min != idx){
        swap(idx, min);
        downheapify(min);
    }
}

```

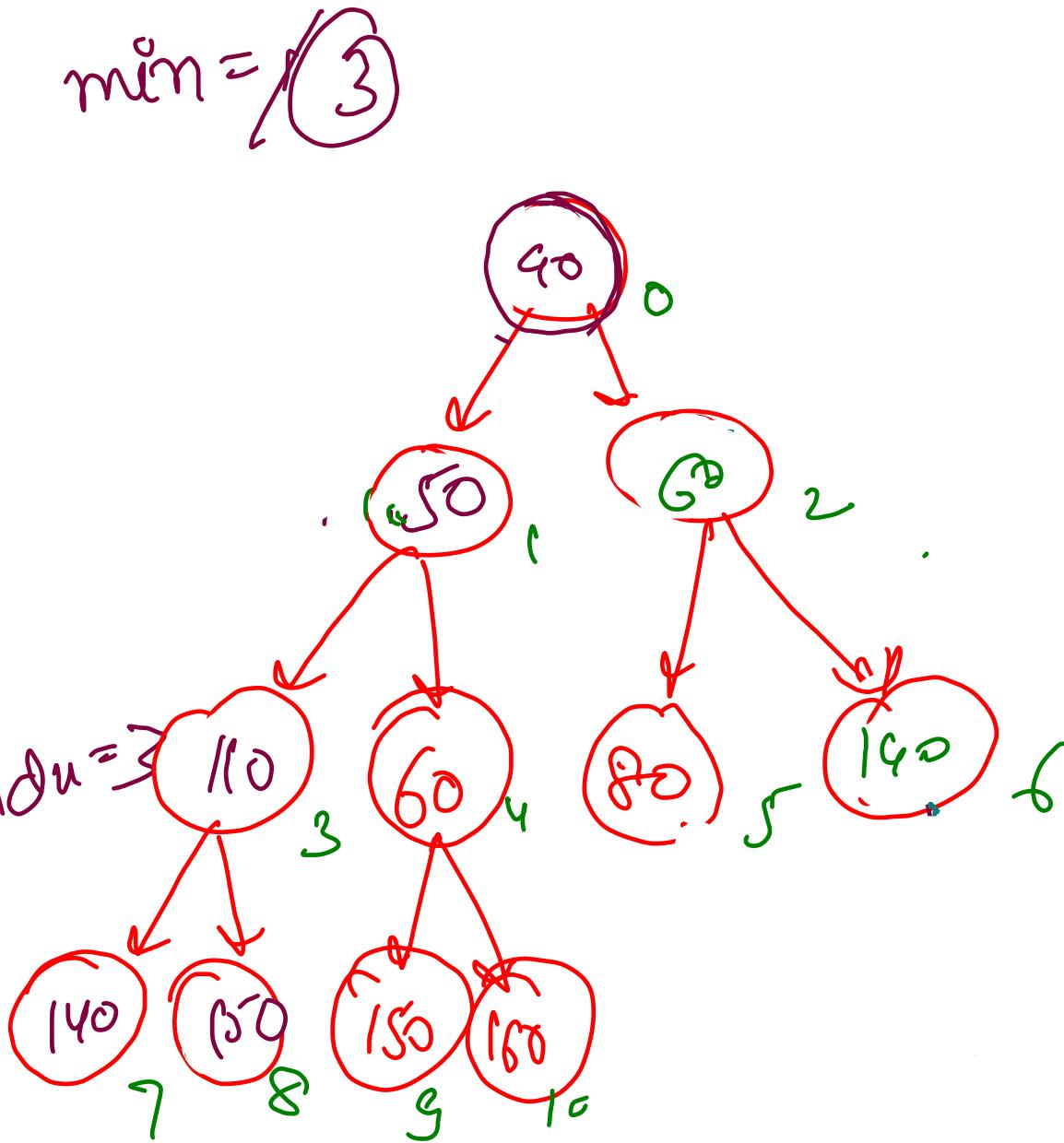
```

public int remove() {
    // O(log n)
    if(size() == 0){
        System.out.println("Underflow");
        return -1;
    }

    int val = peek();
    swap(0, size() - 1);
    data.remove(size() - 1);

    downheapify(0);
    return val;
}

```



# opti nodes, zyada kaam

0. 2<sup>0</sup> + 4

+ 2<sup>1</sup> + 3

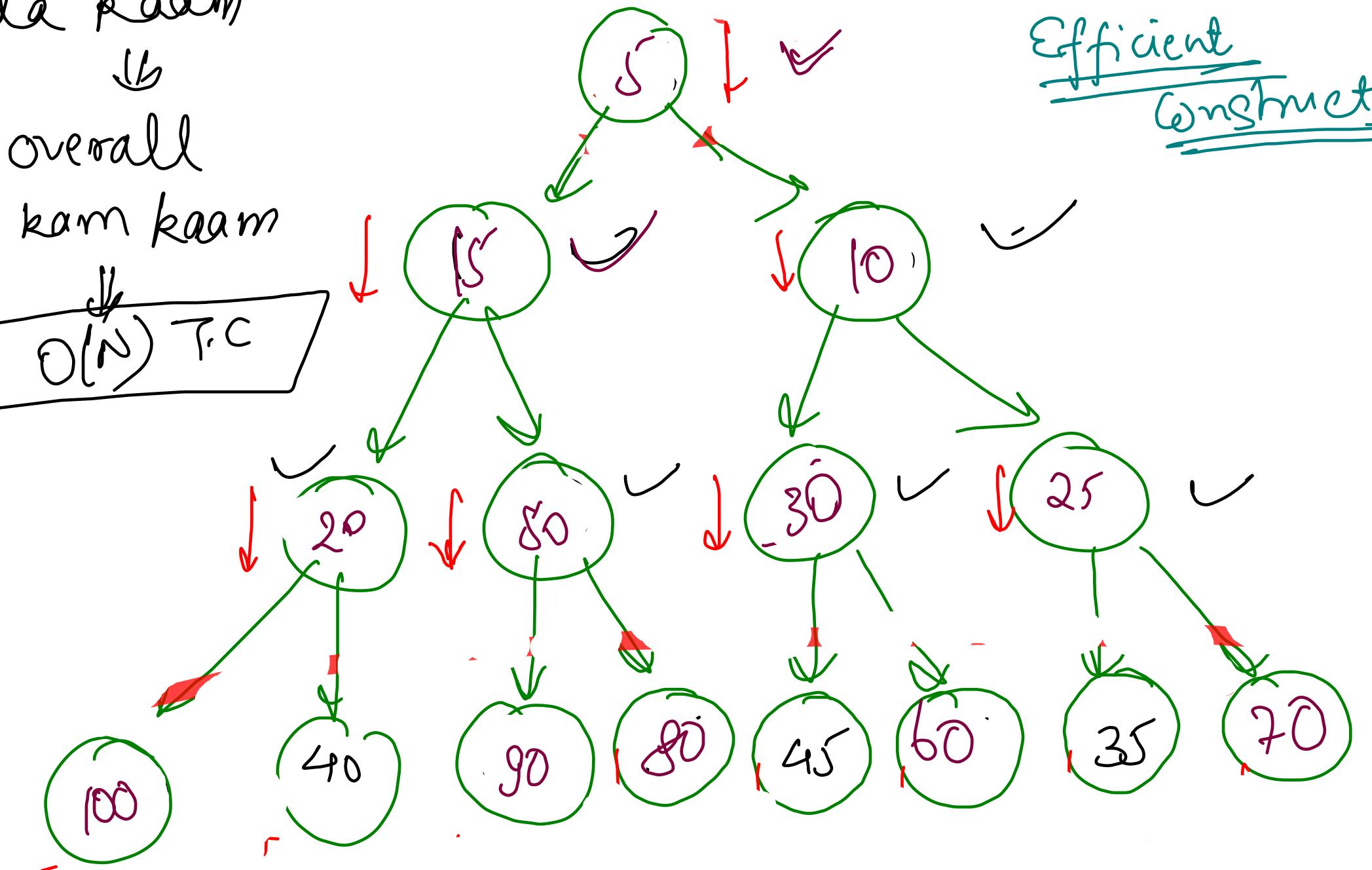
+ 2<sup>2</sup> + 2

+ 2<sup>3</sup> + 1

overall

kam kaam

$O(N)$  T.C



Efficient Constructor

AGP

$$T(n) = 2^0 \cdot h + 2^1 \cdot (h-1) + 2^2 \cdot (h-2) + 2^3 \cdot (h-3) + \dots + 2^{h-1} \cdot 1$$

$$+ 2^k T(n) = 2^0 \cdot h + 2^1 \cdot (h-1) + 2^2 \cdot (h-2) + \dots + 2^{h-1} \cdot 1$$

$$T(n) = -2^0 \cdot h + \{ 2^1 \cdot 1 + 2^2 \cdot 1 + 2^3 \cdot 1 + \dots + 2^{h-1} \cdot 1 \} + 2^h \cdot 1$$

$$T(n) = (2^h - 1) - 2^0 \cdot h$$

$$= 2^{\log_2 n} - 1 - \log_2 n$$

$$= n - \log n - 1$$

$$\boxed{h = \log_2 n}$$

$$\boxed{O(n)}$$

```
// O(n) for inserting n elements -> per element O(1)
public PriorityQueue(int[] arr){
    data = new ArrayList<>();
    for(int val: arr){
        data.add(val);
        size++;
    }
    for(int i=(size() - 1)/2; i>=0; i--){
        downheapify(i);
    }
}
```

Interview Prep

10 questions  $\Rightarrow$  3-4 Ques

$\Rightarrow$  3-4 Ques Algo's

$\Rightarrow$  1-2 complete

~~360~~  
~~660~~

3 monthly

5 classes \* 4 weekly

= 20 classes

$20 \times 3 = 60$  classes

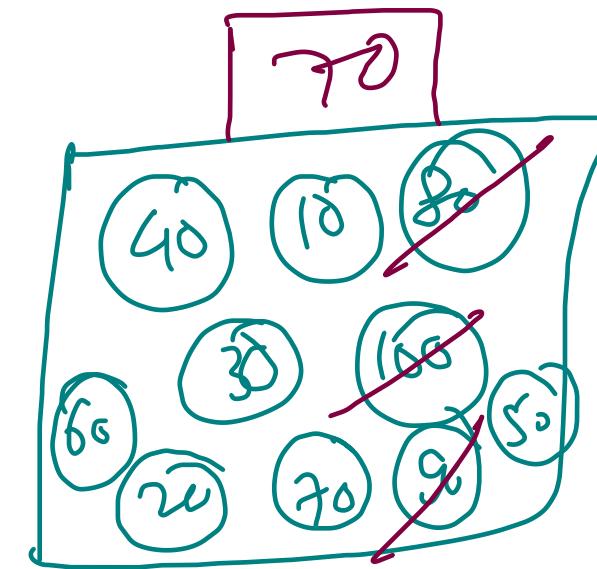
MTW

60 & 6 = 360

$\{ 40, 10, 80, 30, 100, 60, 20, 70, 90, 50 \}$   $k=3$

Approach 1  $\rightarrow$  Max Heap

$\{ 100, 30, 80 \}$

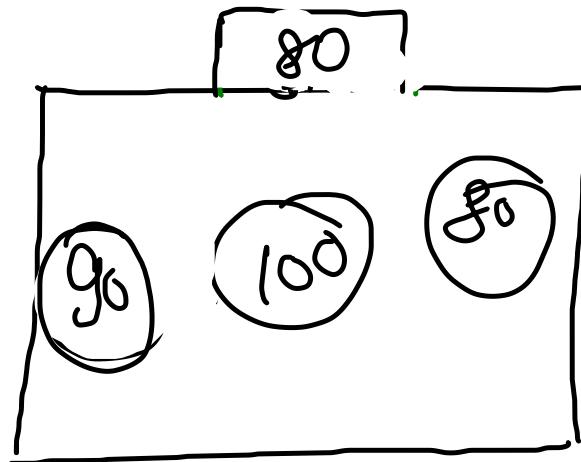


$O(N + K \log N)$

```
public static void approach1(ArrayList<Integer> arr, int k){  
    PriorityQueue<Integer> pq = new PriorityQueue<>(Collections.reverseOrder());  
  
    // O(n)  
    pq.addAll(arr);  
  
    // O(k log N)  
    ArrayList<Integer> res = new ArrayList<>();  
    while(k-- > 0)  
        res.add(pq.remove());  
  
    for(int i=res.size() - 1; i>=0; i--){  
        System.out.println(res.get(i));  
    }  
}
```

$\{ 40, 10, 80, 30, 100, 60, 20, 70, 90, 50 \}$   $\underline{k=3}$

f min heap



$$O(k \log k + (n-k) \log k) = O(n \log k)$$

```
// O(n log k)
public static void approach2(ArrayList<Integer> arr, int k){
    // Min heap
    PriorityQueue<Integer> pq = new PriorityQueue<>();

    // O(k log k)
    for(int i=0; i<k; i++){
        pq.add(arr.get(i));
    }

    // O((n-k) log k)
    for(int i=k; i<arr.size(); i++){

        if(arr.get(i) > pq.peek()){
            pq.remove();
            pq.add(arr.get(i));
        }
    }

    while(pq.size() > 0){
        System.out.println(pq.remove());
    }
}
```

## Today's Evening Questions

- </> Sort K-sorted Array
- ☒ Heap - Comparable V/s Comparator
- </> Merge K Sorted Lists
- </> Median Priority Queue

## level 2 & level 3

{ Level up (300 Questions) + Interview Prep (300 Ques) }

• Arrays & Strings { searching, sorting } (greedy, 2 pointers)  $\Rightarrow \pm 15$  classes

• Recursion & Backtracking  $\Rightarrow \pm 10$  classes

• Stack & Queue  $\Rightarrow \pm 5$  classes

• linked list  $\Rightarrow \pm 5$  classes

• Trees (Binary Tree, BST)  $\Rightarrow \pm 10$  classes

• Hashmap & Heap  $\Rightarrow \pm 5$  classes

• Graphs  $\Rightarrow \pm 8$  classes

• Dynamic Programming  $\Rightarrow \pm 15$  classes

80 classes

↓  
5 classes per week

↓  
4 months

• Bit Manipulation ② classes

• Tries ② classes

• Advanced Trees ① class

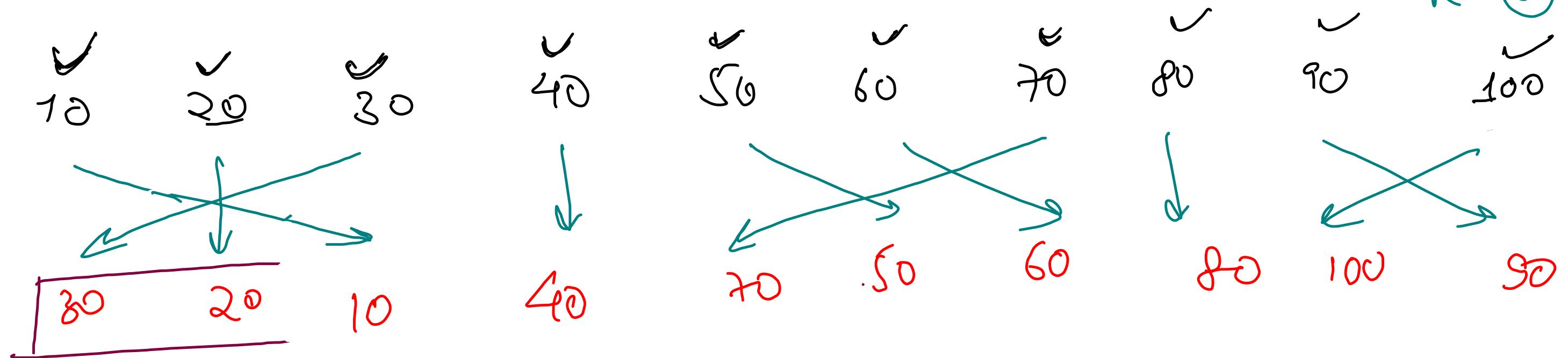
• Mathematics & Number Theory ② classes

• Range Query ② classes

• Geometry & Game Theory ① class

• String pattern matching ① class

Sort k Sorted Array [  $i_{dn-k}, i_{dn+k}$  ]



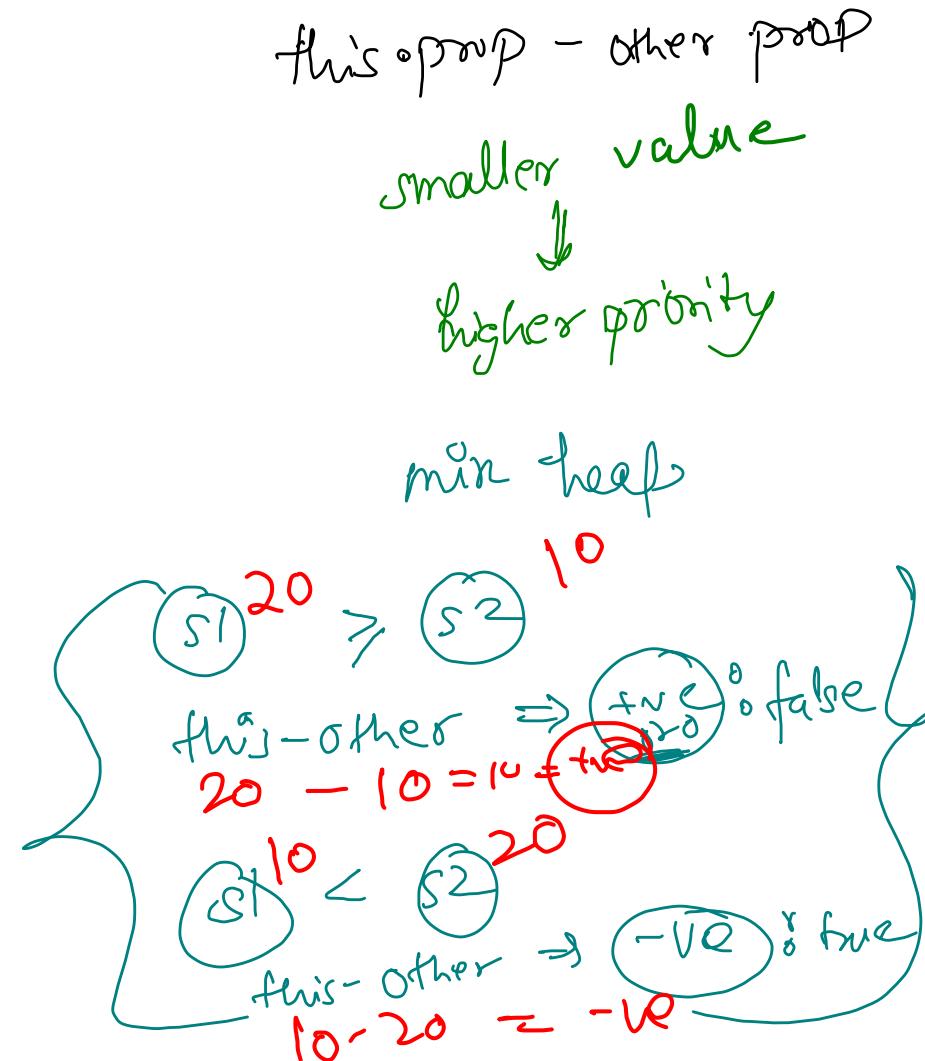
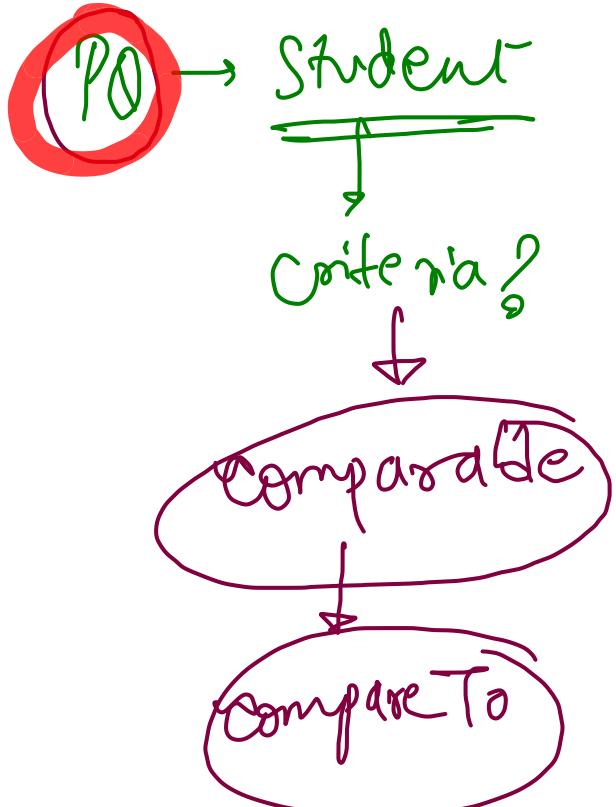
brute force :- Array sort :  $N \log N$

optimized code :-  $R \log k + (n-k) \log k = (k+n-k) \log k = n \log k$

→ Create Priority Queue of  $(k+1)$  elements { sliding window }

{ push 1 ele }  
pop 1 ele

## Comparable & Comparator



other.prop - this.prop

higher value  
↓

higher priority

max heap

$$S1 \leq S2$$

$$\text{other} - \text{this} = \text{the if false}$$

$$S2 - S1 = 20 - 10 = 10 \text{ +ve}$$

$$S1 > S2$$

$$\text{other} - \text{this} = \text{-ve: the}$$

$$10 - 20 = -10 \text{ -ve}$$

# Comparable & comparator

```
PriorityQueue<Student> pq = new PriorityQueue<>(new StudentMrksComparator());
```

```
public static class Student implements Comparable<Student>{
    int rollNo, marks, weight;

    Student(){}
    Student(int rollNo, int marks, int weight){}
    public String toString(){}

    public int compareTo(Student other){
        // Smaller Roll No -> Greater Priority
        return this.rollNo - other.rollNo;
    }
}

public static class StudentWtComparator implements Comparator<Student>{
    public int compare(Student s1, Student s2){
        // Higher Weight - Higher Priority
        return s2.weight - s1.weight;
    }
}

public static class StudentMrksComparator implements Comparator<Student>{
    public int compare(Student s1, Student s2){
        // Higher Marks - Higher Priority
        return s2.marks - s1.marks;
    }
}
```

- ① Generic
  - ② Comparable
  - ③ Obj. compare(s1, s2)
- is smaller

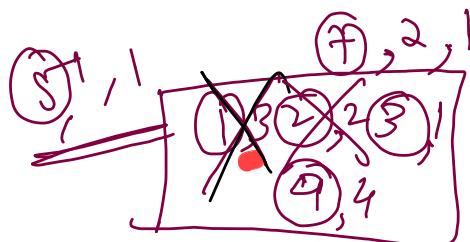
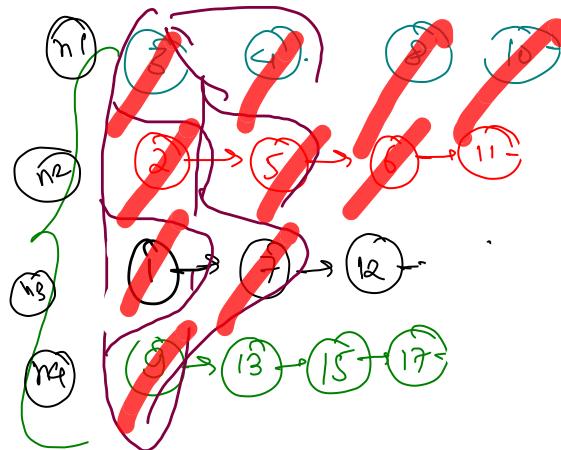
## # isSmaller

```
public boolean isSmaller(int i1, int i2){  
    K s1 = data.get(i1);  
    K s2 = data.get(i2);  
  
    if(comp == null){  
        // Comparable -> Compare To  
        Comparable s1c = (Comparable)s1;  
        Comparable s2c = (Comparable)s2;  
  
        // s1.compareTo(s2) will not work  
        if(s1c.compareTo(s2c) < 0)  
            return true;  
        return false;  
    } else {  
        // Comparator  
        if(comp.compare(s1, s2) < 0)  
            return true;  
        return false;  
    }  
}
```

## # Constructor

```
public ArrayList<K> data;  
private int size;  
private Comparator<K> comp;  
  
public Priorityqueue(){  
    data = new ArrayList<>();  
    size = 0;  
    this.comp = null;  
}  
  
public Priorityqueue(Comparator<K> comp){  
    data = new ArrayList<>();  
    size = 0;  
    this.comp = comp;  
}
```

## Merge k sorted arrays



$\log k$

(1, 2, 3, 4, 5, 7, 8, 9, 10)

d → {  
  data  
  listIdn  
  dataIdn}

brute force:

merge 2 at a time

$$O(n_1) +$$

$$O(n_1 + n_2)$$

$$O(n_1 + n_2 + n_3)$$

+

$$O(n_1 + n_2 + n_3 + n_4)$$

+

$$O(n_1 + n_2 + \dots + n_k)$$

$$O(n_1 * (k) + n_2 * (k-1))$$

$$+ n_3 * (k-2)$$

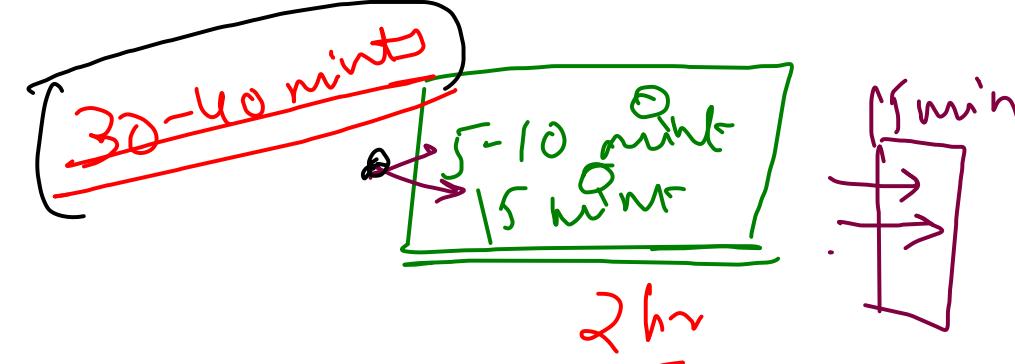
$$+ n_4 * (k-3) + \dots$$

$$+ n_k * 1)$$

$$N * k + N * (k-1) + N * (k-2) + \dots + N * 1$$

$$= N * \{ 1 + 2 + 3 + \dots + k \}$$

$$= O(N * k)$$

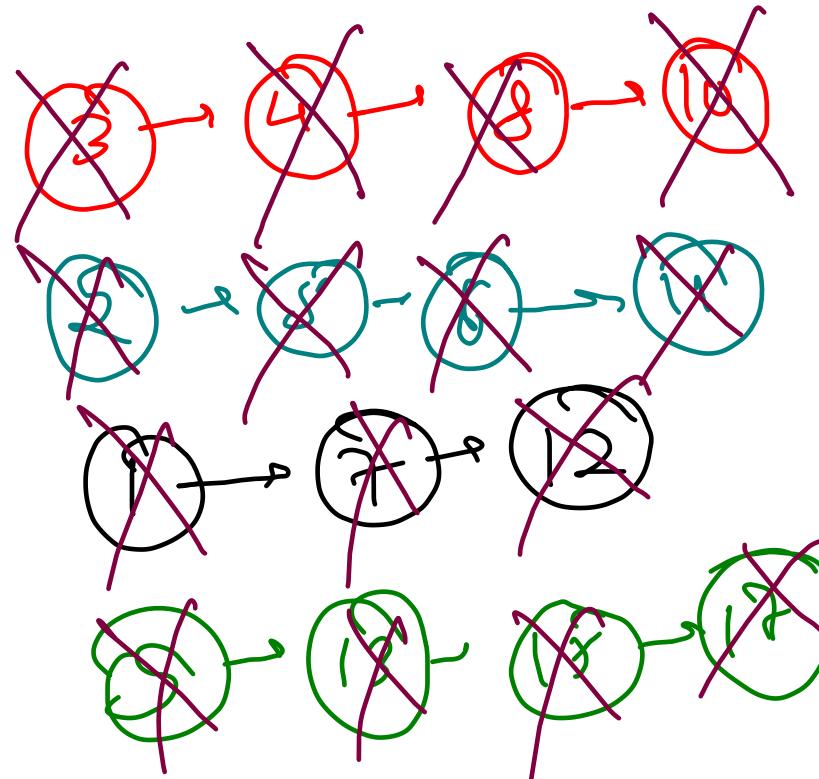


2 hr

g-10 hrs

⇒ push k elements :→ 0th element of each array(.)<sup>2</sup>

while (q.size()



```

public static ArrayList<Integer> mergeKSortedLists(ArrayList<ArrayList<Integer>> lists){
    PriorityQueue<Pair> q = new PriorityQueue<>();
    ArrayList<Integer> res = new ArrayList<>();

    // Insert 0th index element of each arraylist
    // O(k log k)
    for(int i=0; i<lists.size(); i++){
        q.add(new Pair(lists.get(i).get(0), i, 0));
    }

    // O(Nt * logk)
    while(q.size() > 0){
        // Pop one pair
        Pair top = q.remove();
        res.add(top.data);

        if(top.dataIdx + 1 < lists.get(top.listIdx).size()){
            // insert next pair
            q.add(new Pair(lists.get(top.listIdx).get(top.dataIdx + 1), top.listIdx, top.dataIdx + 1));
        }
    }

    // O(Nt * logk + k * logk) where Nt = n1 + n2 + n3 + ... nk
    return res;
}

```