

HashMap & Heap

DSA classes



Order of Content

- ① → Hashmap & Heap → level 1 {without constructor}

② → Graphs - lecture ① + ② # Chaining
 # Open Addressing
 # Good Hashfunc,
 Bad Hashfn

③ → Hashmap & Heaps + Arrays & Strings → level 2 ↑
 + constructor

④ → Bit Manipulation] → Competitive

⑤ → Number Theory]

⑥ → Trees - level 2 of remaining lectures ↗
 Self Balancing
 BST (AVL)

Data Structures

① add (insert)
 $\hookrightarrow O(1)$

② remove (delete)
 $\hookrightarrow O(1)$

③ update
 $\hookrightarrow O(1)$

④ search (find)
 $\hookrightarrow O(1)$

⑤ display / Traverse

⑥ length / size

⑦ read (get)
 $\hookrightarrow O(1)$

Arrays / Strings

$O(1)$ insert, $O(n)$ first

$O(1)$ last, $O(n)$ first

$O(1)$ {indexing}

$O(n)$ LS, $O(\log_2 n)$ BS

$O(n)$

$O(1)$

$O(1)$ {indexing}

Lists

$O(1)$ f, $O(1)$ last

$O(n)$ f, $O(1)$ last

$O(n)$ {indexing}

$O(n)$ LS

$O(n)$

$O(1)$

$O(n)$ worst case

Team (String)
↓
IPL Tournaments
Count
(integer)

Team (String)
↓
Years of winning
(ArrayList<Int>)

~~Java~~ Hashmap
key-value pairs
Read
Create
Update
Delete
Find
Objects

~~C++~~ Unordered_map
~~Python~~ Dictionary

~~Json~~ Key: value

(String) Key → Country
(long) Value → Population

HashMap

`HashMap<key, value> hm = new HashMap<>();`

① insert ↗ ~~key not found~~

② update ↗ ~~key found~~

③ delete ↗ remove

④ find/read → get ↗ $O(1)$

⑤ search → containsKey ↗ $O(1)$
 ↗ true/false

but ↗ $O(n)$

⑥ size() {
 ↗ $O(1)$
 ↗ k-v pairs } ↗ $O(n)$

⑦ display ↗ `System.out.println(hm);`

⑧ keySet ↗ $O(N)$

set & key ↗
Traverse

function, variable name



Java
getOrDefault()

populationOfCountry

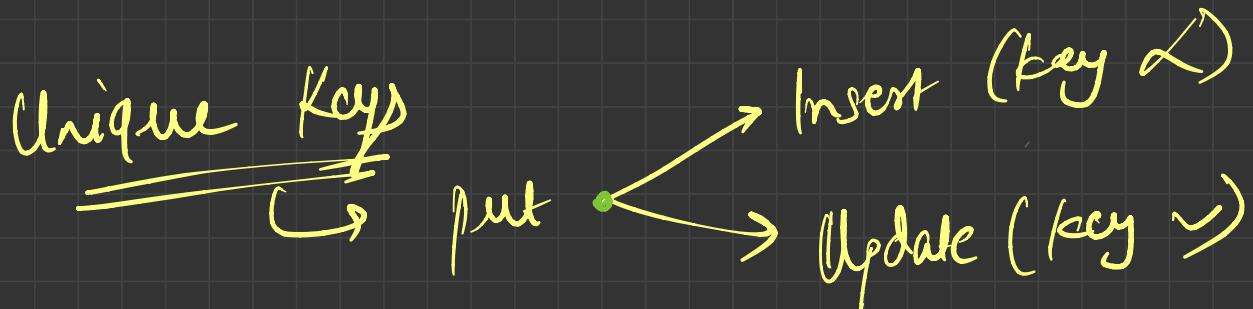
C++
get-or-default() population-of-country

Hashmap → { Order of insertion
 + Order of key-value pair }

Unordered
no indexing

+
sorted order based on key/value

Number of Key-value pairs
= Number of keys



1 key → 1 entry { No duplicates keys will
be present }

Common Elements / Intersection

{ 2, 1, 1, 3, 4, 2, 5, 1, 2, 5 }

{ 1, 1, 4, 4, 4, 5, 5, 3 }

LC 349

{ 1, 3, 4, 5 }

LC 350

{ 1, 1, 3, 4, 5, 5 }



Bonule force

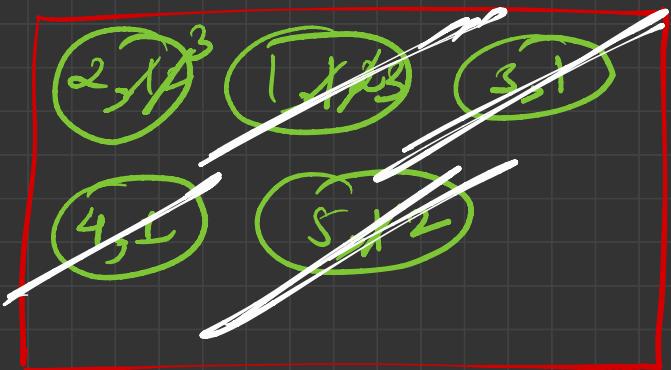
Sort Both Arrays + Two pointer

$$N_1 \log N_1 + N_2 \log N_2 + (N_1 + N_2)$$

~~Optimized~~ Flashmap

{ 2, 1, 1, 3, 4, 2, 5, 1, 2, 5 }

{ 1, 1, 4, 4, 4, 5, 5, 3 }



"1, 4, 5, 3"

HashMap < Integer, Integer >

Element → frequency

T,W,T → 9:45 PM - 11:45 PM (2hr)

S,S → 9 AM - 12 PM (3hr)

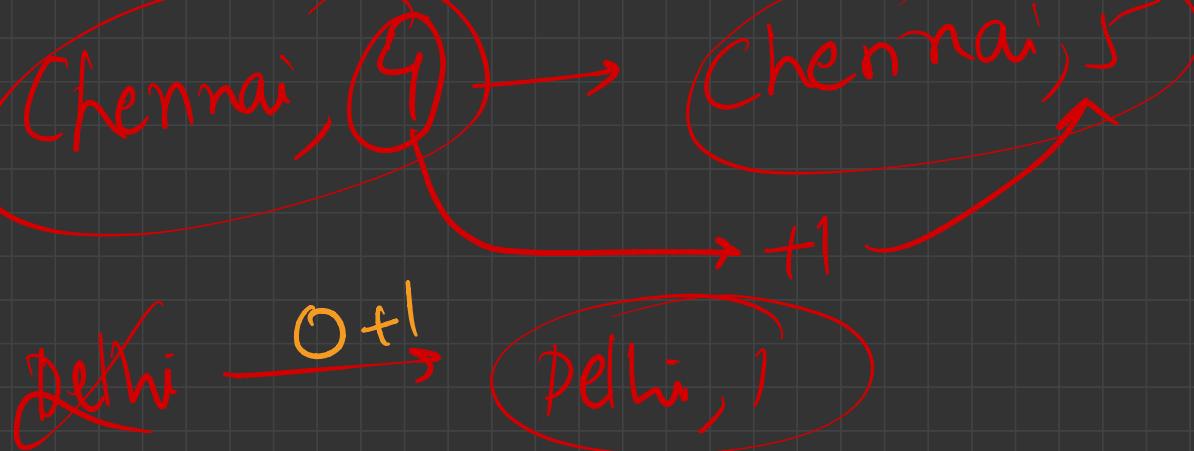


2-3 weeks

T,W,T → 9:15 PM - 11:15 PM (2hr)

S,S → 9 AM - 12 PM (3hr)

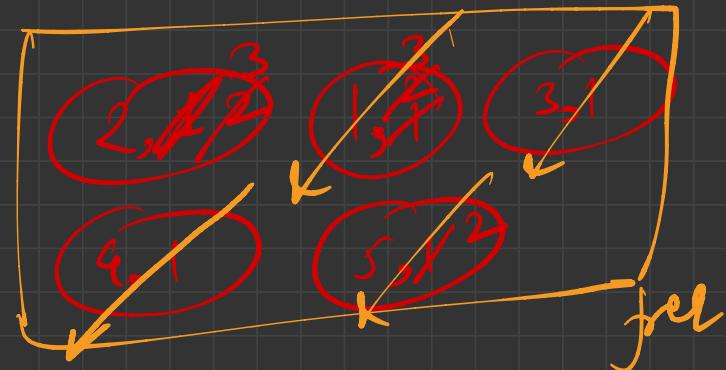
```
for (int i = 0, i < nums1.length, i++) {  
    if(freq.containsKey(nums1[i]) == true){  
        freq.put(nums1[i], freq.get(nums1[i]) + 1);  
    } else {  
        freq.put(nums1[i], 1);  
    }  
}  
  
freq.put(nums1[i], freq.getOrDefault(nums1[i], 0) + 1);
```



~~349 LC~~

```
HashMap<Integer, Integer> freq = new HashMap<>();  
  
for(int i=0; i<nums1.length; i++){  
    // if(freq.containsKey(nums1[i]) == true){  
    //     freq.put(nums1[i], freq.get(nums1[i]) + 1);  
    // } else {  
    //     freq.put(nums1[i], 1);  
    // }  
    freq.put(nums1[i], freq.getOrDefault(nums1[i], 0) + 1);  
}  
  
ArrayList<Integer> intersection = new ArrayList<>();  
  
for(int j=0; j<nums2.length; j++){  
    if(freq.containsKey(nums2[j]) == true){  
        intersection.add(nums2[j]);  
        freq.remove(nums2[j]);  
    }  
}  
  
int[] res = new int[intersection.size()];  
for(int i=0; i<res.length; i++)  
    res[i] = intersection.get(i);  
return res;
```

{ 2, 1, 1, 3, 4, 2, 5, 1, 2, 5 }
{ 1, 1, 4, 4, 4, 5, 5, 3, 6 }



{ 1, 4, 6, 3 }

$O(N_1 + N_2)$,
Time

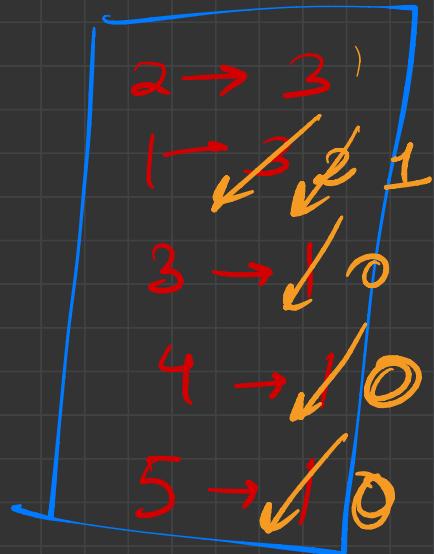
$O(N_1)$ extra
space
(HashMap)

$\{2, 1, 1, 3, 4, 2, 5, 1, 2, 5\}$

$\{1, 1, 4, 4, 4, 5, 3, 3, 6\}$

$\{1, 1, 4, 5, 3\}$

abcede
a⁻¹b⁻¹c⁻¹d⁻¹e⁻¹
(abde, c, bde)



for

aaaaa
a⁻⁵
 $\sum_{i=1}^5 \{a\}$

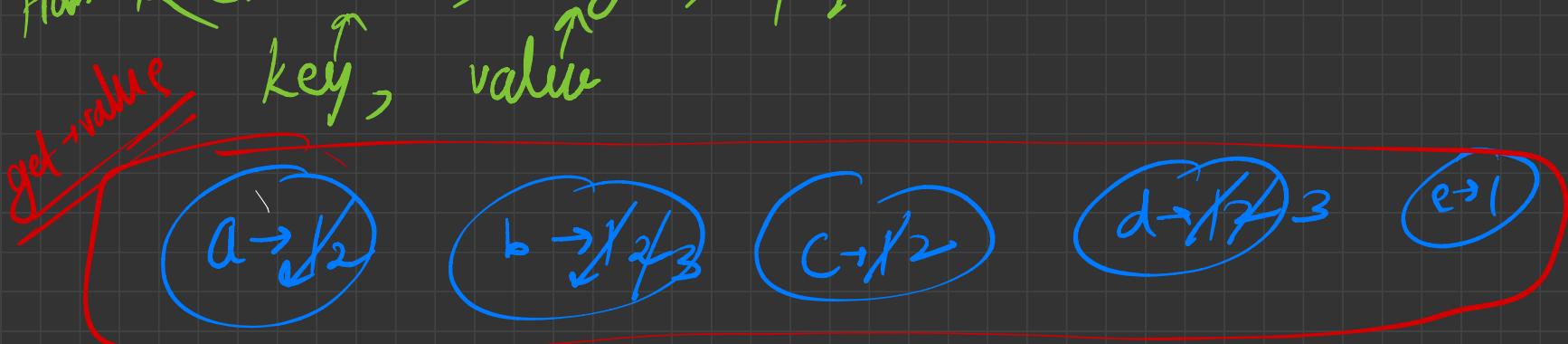
```
HashMap<Integer, Integer> freq = new HashMap<>();  
  
for(int i=0; i<nums1.length; i++)  
    freq.put(nums1[i], freq.getOrDefault(nums1[i], 0) + 1);  
  
ArrayList<Integer> intersection = new ArrayList<>();  
  
for(int j=0; j<nums2.length; j++){  
    if(freq.containsKey(nums2[j]) == true && freq.get(nums2[j]) > 0) {  
        intersection.add(nums2[j]);  
        freq.put(nums2[j], freq.getOrDefault(nums2[j], 0) - 1);  
    }  
}  
  
int[] res = new int[intersection.size()];  
for(int i=0; i<res.length; i++)  
    res[i] = intersection.get(i);  
return res;
```

Sort String by Frequency

abca**b**dd**c**d**b**e
0 1 2 3 4 5 6 7 8 9 10

$$N = 11$$

HashMap<Character, Integer> frey = new HashMap<>();
key, value



Expected output:

bbb ddd aacce
3 2 1

HashMap < Integer, ArrayList < Character >>

~~get k of~~

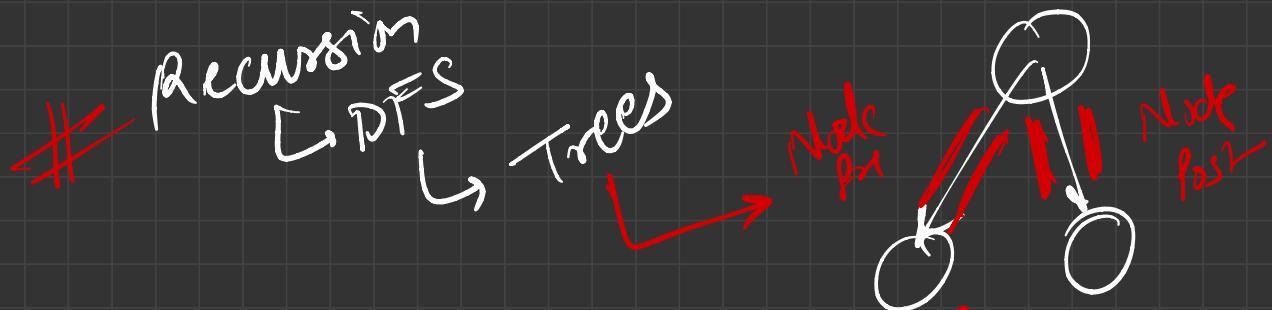
2 → {a,c}

3 → {b,d}

1 → c

bbbddd aace

	map[N]	
0	11	3 → {b,d} 2
0	10	2 → {a,c} 2
0	9	1 → {e} 1
0	8	
6	7	
0	6	
0	5	
0	4	



- Graph
- ① Hashmap / Map
 - ② Trees → DFS { record & visit }
 - ③ Arraylist / OOPS

```

    HashMap<Character, Integer> freq = new HashMap<>();
    for(int i=0; i<s.length(); i++){
        char ch = s.charAt(i);
        freq.put(ch, freq.getOrDefault(ch, 0) + 1);
    }
}

```

$O(N)$

```

    HashMap<Integer, ArrayList<Character>> rev = new HashMap<>();
    for(Character ch: freq.keySet()){
        int f = freq.get(ch);
        if(rev.containsKey(f) == false)
            rev.put(f, new ArrayList<>());
        rev.get(f).add(ch);
    }
}

```

$O(26+26)$

```

    StringBuilder res = new StringBuilder("");
    for(int f=s.length(); f>=1; f--){
        if(rev.containsKey(f) == true){
            for(Character ch: rev.get(f)){
                for(int i=0; i<f; i++)
                    res.append(ch);
            }
        }
    }
}

return res.toString();

```

$\hookrightarrow O(N)$

Time $\rightarrow O(N)$
Space $\rightarrow O(N+26)$

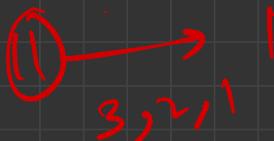
$C \rightarrow I$



Int \rightarrow Act?



eg = "bbdddaaccc"



$11 \rightarrow \{3\}_0$ $8 \rightarrow \{2\}_0$ $5 \rightarrow \{3\}_0$ $2 \rightarrow \{a, c\}_2$ $10 \rightarrow \{3\}_0$ $7 \rightarrow \{2\}_0$ $4 \rightarrow \{2\}_0$ $9 \rightarrow \{2\}_0$ $6 \rightarrow \{3\}_0$ $3 \rightarrow \{b, d\}_2$ $1 \rightarrow \{e, f\}_1$

$$N + 26^{26} + N = O(n)$$

$11 + \textcircled{5} + (8+3+2+2+1)$

 $b \rightarrow 3$ $a \rightarrow 2$ $e \rightarrow 1$ $d \rightarrow 3$ $c \rightarrow 2$ $\Rightarrow 11 \text{ total}$

HashMap < key, Boolean > hm \Rightarrow HashSet < key >



longest consecutive sequence

LC-128

{ 10, 2, 8, 15, 1, 5, 7, 6, 17, 16, 3 }

{ 1, 2, 3 } { 5, 6, 7, 8 } { 10 } { 15, 16, 17 }

① Sorting → Time $N \log N$, Space $\rightarrow O(1)$
 { 1, 2, 3, 5, 6, 7, 8, 10, 15, 16, 17 }

+
 Two pointer → N
 $O(N \log N + N) = O(N \log N)$

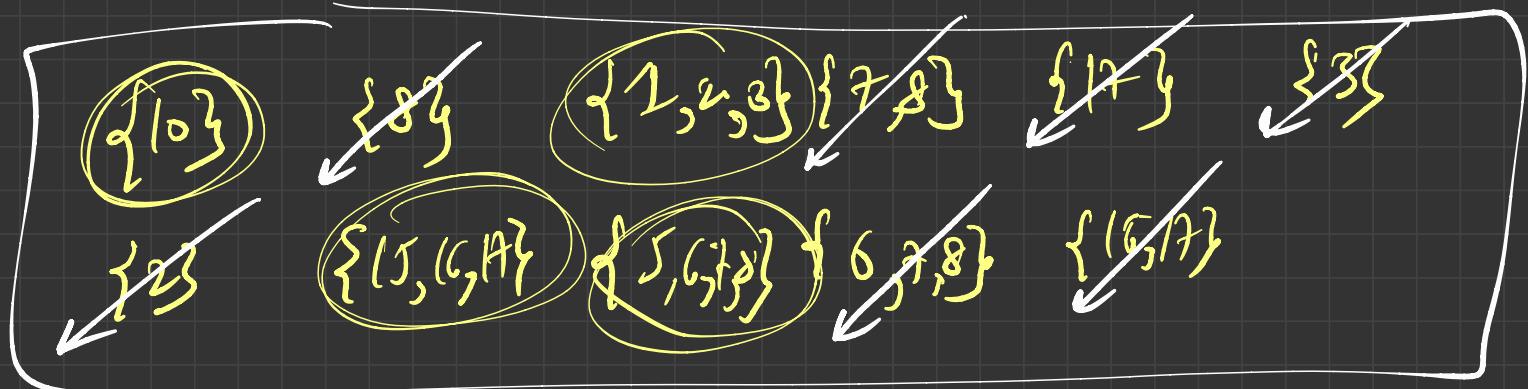
② Hashmap
 ↓
 Each elements if it
 starts the sequence
 (longest sequence)

③ Min → May
 No

$O(\text{Max}^n - \text{Min}^n)$
 $= O(10^9 - 10^9) =$


$\{ \overset{\checkmark}{9}, \overset{\checkmark}{1}, \overset{\checkmark}{7}, \overset{\checkmark}{14}, \overset{\checkmark}{0}, \overset{\checkmark}{4}, \overset{\checkmark}{6}, \overset{\checkmark}{8}, \overset{\checkmark}{16}, \overset{\checkmark}{15}, \overset{\checkmark}{2} \}$

$\{ 10, 2, 8, 15, 1, 5, 7, 6, 17, 16, 3 \}$



last set


```

public int longestConsecutive(int[] nums) {
    if(nums.length == 0) return 0;

    Arrays.sort(nums);

    int left = 0, max = 1;
    while(left < nums.length){

        int right = left + 1, count = 1;
        while(right < nums.length && nums[right] <= nums[right - 1] + 1){
            if(nums[right] == nums[right - 1] + 1)
                count++;
            right++;
        }

        max = Math.max(max, count);
        left = right;
    }

    return max;
}

```

① Sorting + Two pointer



count = 1 2 3 4 5 6 7 8 9 10

max = 3 4

Time $\rightarrow N \log N$, Space $\rightarrow O(1)$

```

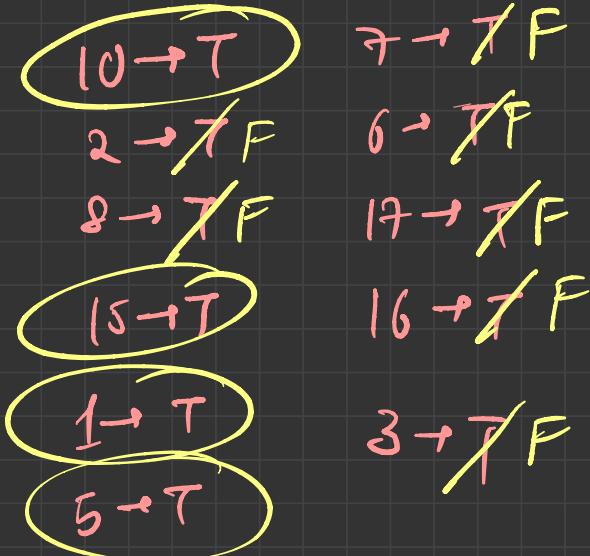
// Unique Key
HashMap<Integer, Boolean> isStarting = new HashMap<>();

// 1. Take All Elements as Potential Starting Sequence
for(int i=0; i<nums.length; i++){
    isStarting.put(nums[i], true);
}

// 2. Remove Elements which cannot start a new Sequence
for(int i=0; i<nums.length; i++){
    if(isStarting.containsKey(nums[i] - 1) == true){
        // We Will Join the Previous Sequence,
        // Hence, we cannot start the sequence
        isStarting.put(nums[i], false);
    }
}

int max = 0;

```



$$\begin{array}{cccc}
 & \downarrow & \downarrow & \downarrow \\
 3+2+1 & & 4+3+2+1 & \\
 \frac{3+9}{2} + \frac{4+5}{2} + \frac{1+2}{2} + \frac{3+4}{2} & = n^2
 \end{array}$$

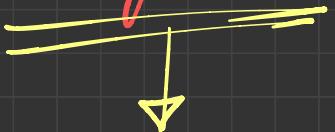
$$\begin{array}{cccc}
 3 & + & 4 & + \\
 & & 1 & + \\
 & & 3 & = n
 \end{array}$$

```

// 3. Finding the Maximum Length among all Disjoint Sets
for(int start: isStarting.keySet()){
    if(isStarting.get(start) == true){
        int curr = 1;
        for(int end=start+1; isStarting.containsKey(end) == true; end++){
            curr++;
        }
        max = Math.max(max, curr);
    }
}
return max;

```

Binary Heap



vs

Priority Queue



Concrete datastructure

(LIFO)

Stack

(FIFO)

Queue

(min-max)
BST

BinaryHeap

abstract datatype

↳ highest priority get
val

↳ remove(HP) - O(logn)

↳ insert(HP) - O(logn)

Order → Higher Priority
(LIFO)
Low - Priority
↳ First

get last → O(1)
remove last O(1)
insert last O(1)

Stack

Order → Least Priority
(FIFO)
Higher Priority → First
↳ Last (Insert Order)

get first
remove first
add last

Queue

10, ~, 20, 40, 50

min value → least priority

Sorted Array

value ↗
min value → Highest Priority

~~Self Balancing~~
BST

*
(AVL, RedBlack,
Splay)

get max → $O(\log_2 N)$ *

delete max → $O(\log_2 N)$

Insert → $O(\log_2 N)$

$O(D)$

→

$O(1)$

→

$O(N)$

Binary Heap

getHP() → O(1)

removeHP() → O(log₂n)

Insert () → O(log₂n)

Priority Queue (Binary Heap)

→ Upheaping, downheaping,
→ efficient construction, heap sort

→ Generic
→ Comparable
objects Comparator
Comparator

⇒ Priority Queue < Datatype > pq = new PriorityQueue<X>;

① Insert

② Remove

Highest Priority

③ Get

Highest Priority

④ Update

remove()
+
insert()

Insert-~~add~~ $\rightarrow O(\log_2 N)$

N = number of elements present in heap

Remove $\rightarrow O(\log_2 N)$

Get ~~peek~~ $\rightarrow O(1)$

Length $\rightarrow O(1)$
~~(size)~~

Min Heap

Comparison

Min element \rightarrow highest priority
value \rightarrow Wrapper class \rightarrow float, double
String (hierarchically)

Max Heap

Max element \rightarrow highest priority

Java → Min Heap
 (Default)



pair → Comparison?

C++ → Max Heap
 (Default)

pair<int, int>
 & default ✓

If n elements are inserted (1 at a time)

$$\begin{aligned}
 &\Rightarrow \log 1 + \log 2 + \log 3 + \dots + \log n \\
 &= \log(1+2+3+\dots+N) = \log N! \\
 &= \log N^N
 \end{aligned}$$

N^N

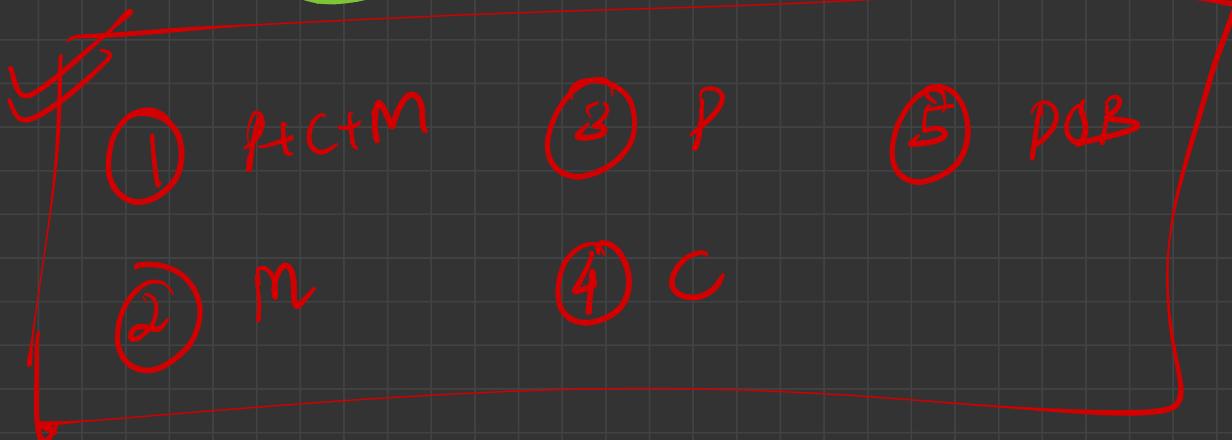
$$= \mathcal{O}(N \log n)$$

PCM

PCM

PCM

PCM



k largest Elements

{ 50, 20, 100, 30, 90, 10, 60, 70, 80, 40 }

$k=4$

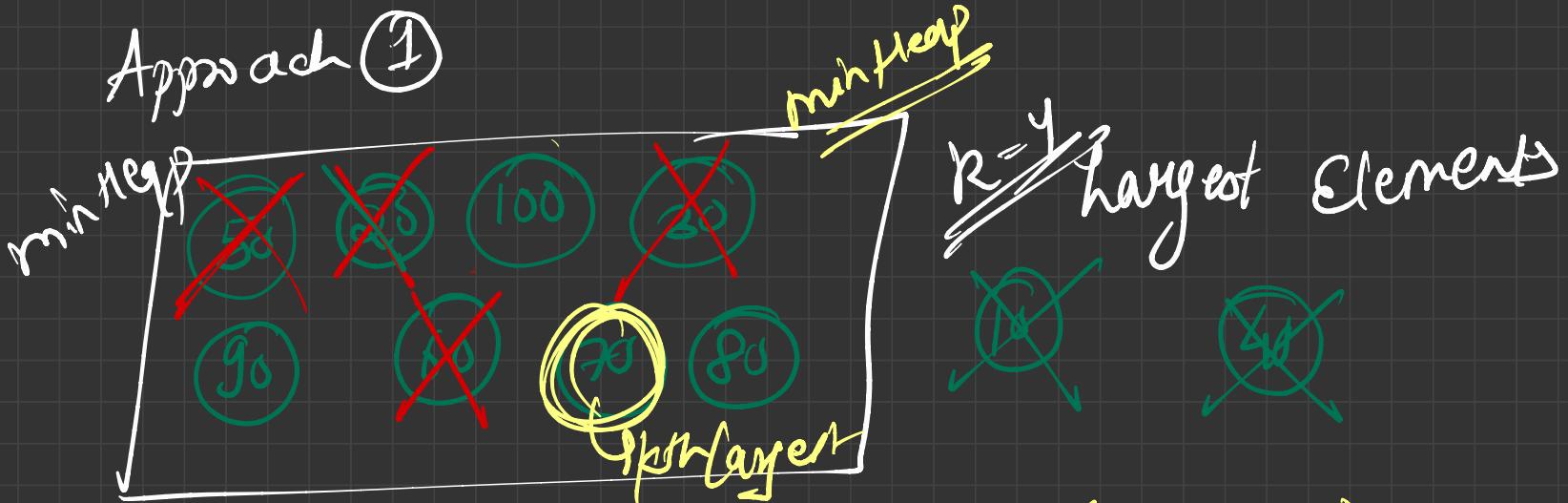
Brute force \rightarrow sorting \star

~~Time $\rightarrow O(N \log N)$~~

space $\rightarrow O(1)$

Priority Queue

Approach ①



→ First k elements in PQ (minHeap)

→ Each $(N-k)$ elements

↓

if $\text{curr} > \text{min}$
→ $\text{insert}(\text{curr})$
→ $\text{remove}()$

else do nothing

```

//Function to return k largest elements from an array.
public static ArrayList<Integer> kLargest(int arr[], int n, int k)
{
    PriorityQueue<Integer> q = new PriorityQueue<>();
    for(int i=0; i<n; i++){
        if(i < k){
            q.add(arr[i]);
        } else if(q.peek() < arr[i]) {
            q.remove();
            q.add(arr[i]);
        }
    }

    ArrayList<Integer> res = new ArrayList<>();
    while(q.size() > 0){
        res.add(q.remove());
    }
    Collections.reverse(res);
    return res;
}

```

$\left. \begin{array}{l} \\ \\ \\ \end{array} \right\} N \log K$

+

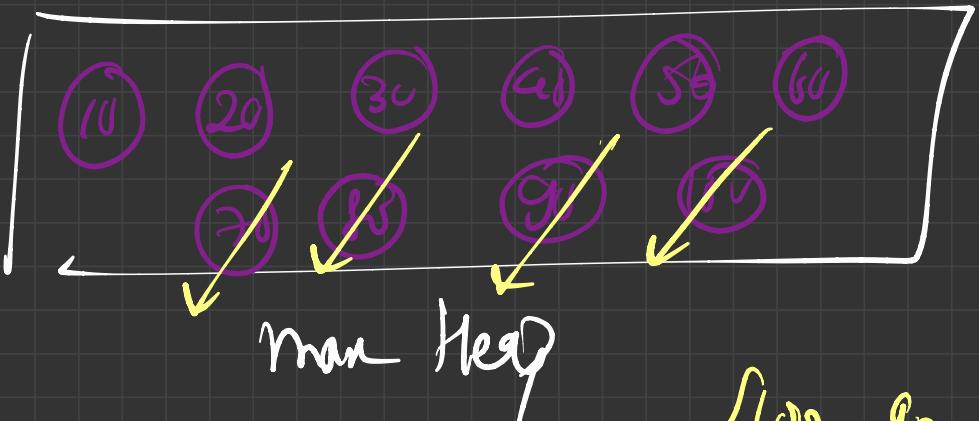
$\left. \begin{array}{l} \\ \\ \\ \end{array} \right\} K \log K$

Time
 $(N+k) \log K$

$\sim N \log K$

Space
 $O(k)$

Approach ②



If $\lceil \log n \rceil$ insert
+
Array is inserted once

$\{100, 90, 80, 70\}$
K elements $\rightarrow K \log N$

$$\text{Total time} = O(N + K \log N)$$

```
// Time - O(N + KlogN) MAX HEAP APPROACH
PriorityQueue<Integer> q = new PriorityQueue<>(Collections.reverseOrder());
for(int i=0; i<n; i++){
    q.add(arr[i]);
}
ArrayList<Integer> res = new ArrayList<>();
while(res.size() < k){
    res.add(q.remove());
}
return res;
```

Time \rightarrow $O(N + K \log N)$

Space \rightarrow $O(K)$

KC215

km largest Elements

Time $\rightarrow O(N \log N)$ insert 1 at a time
Space $\rightarrow O(N)$ removes k elements
Priority Queue

```
public int findKthLargest(int[] arr, int k) {  
    PriorityQueue<Integer> q = new PriorityQueue<>(Collections.reverseOrder());  
    for(int i=0; i<arr.length; i++){  
        q.add(arr[i]);  
    }  
  
    int kthLargest = q.peek();  
    while(k-- > 0){  
        kthLargest = q.remove();  
    }  
    return kthLargest;  
}
```

Priority Queue
 max heap

Time $\rightarrow O(N + K \log N)$
 space $\rightarrow O(N)$

(add all) \downarrow insert removal
~~8ms~~

```

public int findKthLargest(int[] arr, int k) {
  ArrayList<Integer> temp = new ArrayList<>();
  for(int val: arr) temp.add(val);

  PriorityQueue<Integer> q = new PriorityQueue<>(Collections.reverseOrder());
  q.addAll(temp);

  int kthLargest = q.peek();
  while(k-- > 0){
    kthLargest = q.remove();
  }
  return kthLargest;
}
    
```

App 3

Min Heap

Time → $N \log k$ $\approx m$

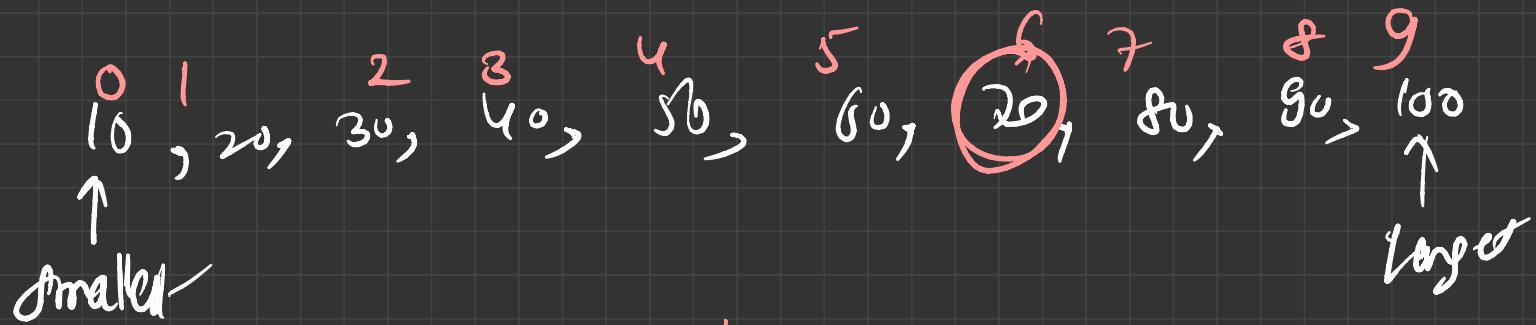
, space → $O(k)$

```
public int findKthLargest(int[] arr, int k) {
    PriorityQueue<Integer> q = new PriorityQueue<>();
    for(int i=0; i<arr.length; i++){
        if(i < k){
            q.add(arr[i]);
        } else if(q.peek() < arr[i]) {
            q.remove();
            q.add(arr[i]);
        }
    }
    return q.peek();
}
```

App 4 QuickSelect { Similar to Quick Sort }

$O(N)$ Partitioning

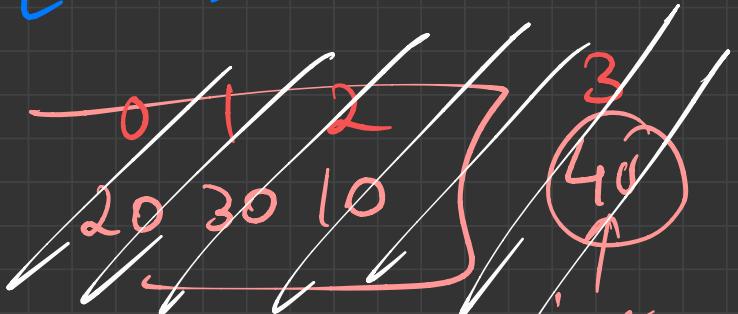
{ 50, 20, 100, 30, 90, 10, 60, 70, 80, 40 }



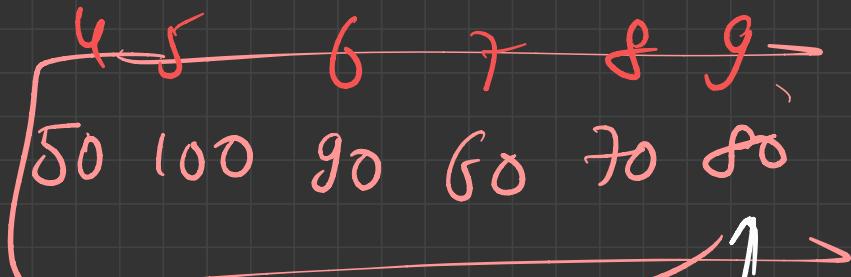
$$N = 10, k = 9, \text{ for Largest} = (N-k)$$

$\mathcal{O}(n^2)$

{ 50, 20, 100, 30, 90, 10, 60, 70, 80, 40 }

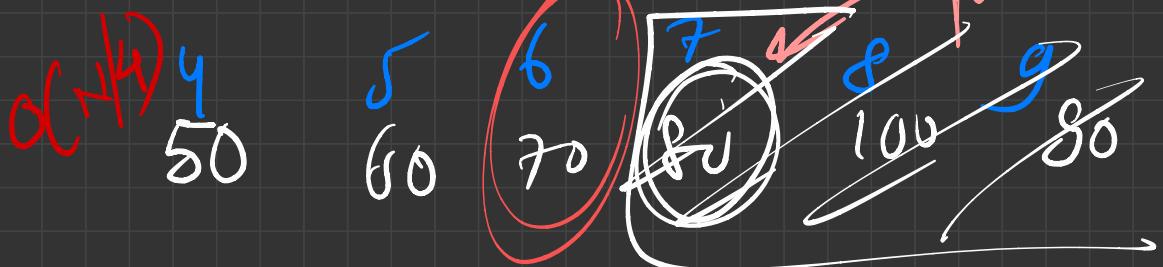


pivot



1
pivot

sorted partition



$\mathcal{O}(n^2)$

4

5

6
70

7
80

8
100

9
80

$$\text{Total time} = N + \frac{N}{2} + \frac{N}{4} + \dots 0$$

$$= N \left\{ 1 + \frac{1}{2} + \frac{1}{2^2} + \frac{1}{2^3} + \frac{1}{2^4} + \dots \right\}$$

worst case

$$N + (N^{1/2}) + (N^{1/4}) + \dots \approx O(N^2)$$

$$= N \left\{ 2 \right\} \Rightarrow O(N)$$

average case
Quick select

```
public int partition(int[] arr, int l, int r){  
    int j = l;  
    for(int i=l; i<=r; i++){  
        if(arr[i] <= arr[r]){  
            int temp = arr[j];  
            arr[j] = arr[i];  
            arr[i] = temp;  
            j++;  
        }  
    }  
    return j - 1;  
}  
  
public int quickSelect(int[] arr, int l, int r, int k){  
    if(l == r) return arr[l];  
  
    int pivot = partition(arr, l, r);  
    if(pivot == arr.length - k) return arr[pivot];  
    if(pivot > arr.length - k) return quickSelect(arr, l, pivot - 1, k);  
    return quickSelect(arr, pivot + 1, r, k);  
}  
  
public int findKthLargest(int[] arr, int k) {  
    return quickSelect(arr, 0, arr.length - 1, k);  
}
```

K closest Points to Origin

~~$$(x_2, y_2)$$~~

$$(6, 8) = \sqrt{6^2 + 8^2} = 10$$

$$(3, 2) = \sqrt{3^2 + 2^2} = \sqrt{13}$$

$$(-4, 1) = \sqrt{(-4)^2 + 1^2} = \sqrt{17}$$

$$(1, -1) = \sqrt{1^2 + (-1)^2} = \sqrt{2}$$

$$(1, 4) = \sqrt{1^2 + 4^2} = \sqrt{17}$$

$$\sqrt{8^2 + 6^2} = \sqrt{100} = 10$$

$$(12, 5) = \sqrt{12^2 + 5^2} = 13$$

$$(0, 0) = \sqrt{0^2 + 0^2} = 0$$

$$(3, 4) = \sqrt{3^2 + 4^2} = 5$$

$$(-5, 12) = \sqrt{(-5)^2 + 12^2} = 13$$

$$K = 4$$

(0, 0) 0 dist

(1, -1) $\sqrt{2}$ dist

(-4, 1) $\sqrt{17}$ dist

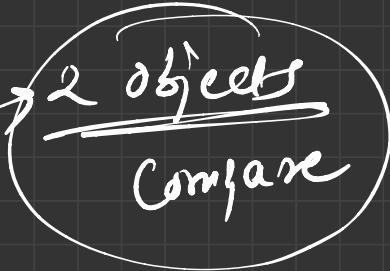
(1, 4) $\sqrt{17}$ dist

Euclidean distance = $\sqrt{(x_2 - x_1)^2 + (y_2 - y_1)^2}$

Manhattan distance = $|x_2 - x_1| + |y_2 - y_1|$

Comparator

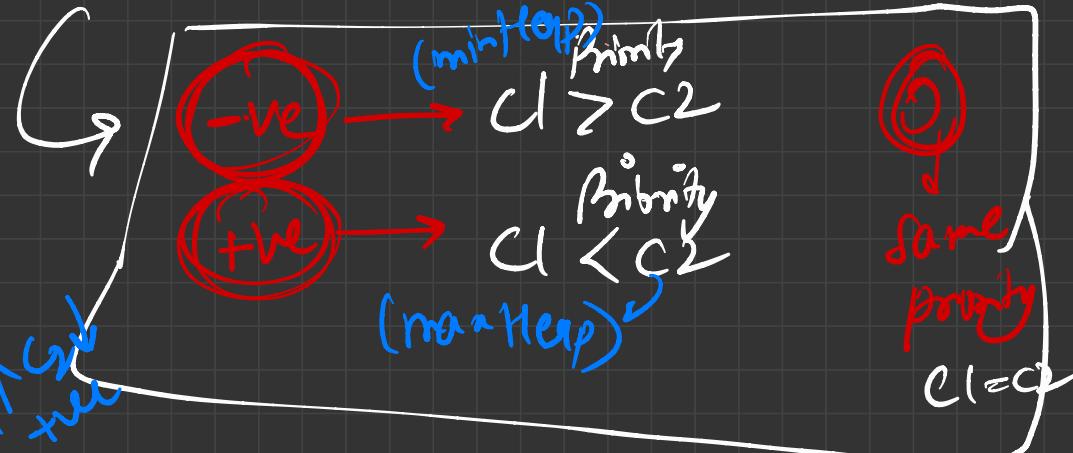
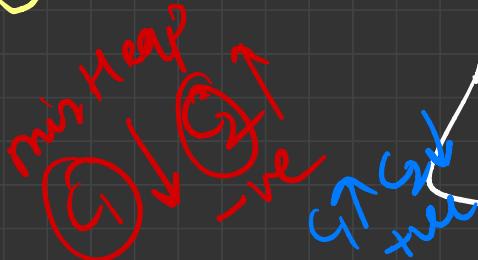
Function



class euclidean implements Comparator<int []> {

@Override
public int compare(int [] c1 , int [] c2) {

}



Closest Point to origin

↳ Euclidean Distance to origin

$$\sqrt{(x-0)^2 + (y-0)^2} = \sqrt{x^2 + y^2}$$

$$C_1(\overbrace{\sqrt{x^2+y^2}}^{HP}) < C_2(\overbrace{\sqrt{x^2+y^2}}^{HP})$$

return -ve;

min HP

$$C_1(\overbrace{\sqrt{x^2+y^2}}^{HP}) > C_2(\overbrace{\sqrt{x^2+y^2}}^{HP})$$

return +ve;

$$C_1(\overbrace{\sqrt{x^2+y^2}}^{HP}) = C_2(\overbrace{\sqrt{x^2+y^2}}^{HP})$$

return 0;

```

static class Pair{
    public int x, y;
}

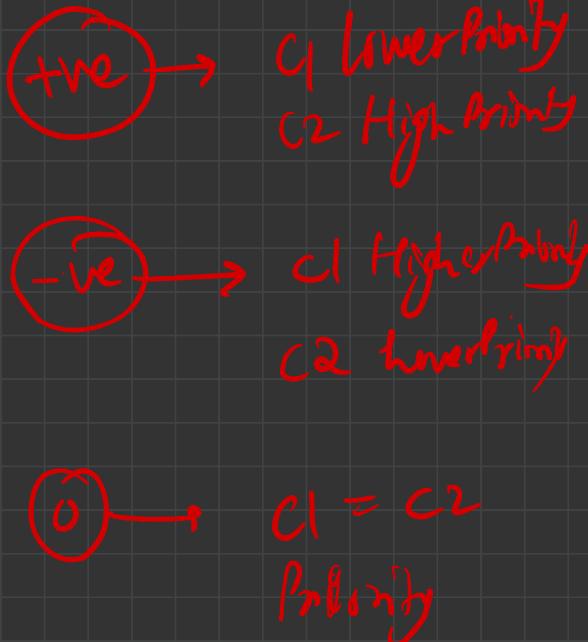
static class Euclidean implements Comparator<Pair> {
    public int compare(Pair c1, Pair c2){
        double d1 = Math.sqrt(c1.x * c1.x + c1.y * c1.y);
        double d2 = Math.sqrt(c2.x * c2.x + c2.y * c2.y);

        if(d1 < d2){ → c1 is closer to c2 → c1 Higher Priority
            // Smaller Value -> Higher Priority
            return -1;
        } else if(d1 > d2){ → c2 is closer than c1 → c2 Higher Priority
            // Higher Value -> Lower Priority
            return 1;
        } else return 0;
    }
}

```

euclidean { *overriden* *Interface*

c1 = c2 *Same Priority*



```

public int[][] kClosest(int[][] points, int k) {
    PriorityQueue<Pair> q = new PriorityQueue<>(new Euclidean());
    ↗ class → object → parameter →
    ↗ constructor

    for(int i=0; i<points.length; i++){
        Pair p = new Pair();
        p.x = points[i][0];
        p.y = points[i][1];
        q.add(p);
    }

    int[][] res = new int[k][2];
    int count = 0;
    while(count < k){
        Pair closest = q.remove();
        res[count][0] = closest.x;
        res[count][1] = closest.y;
        count++;
    }

    return res;
}

```

$\left. \begin{array}{l} \text{Insertion} \\ \text{NlogN} \end{array} \right\}$ $\left. \begin{array}{l} \text{Removal} \\ |k\log N| \end{array} \right\}$

\rightarrow object \rightarrow parameter \rightarrow \rightarrow constructor

$\xrightarrow{\text{Time}} O(N\log N)$

$+ O(k\log N)$

Removal

Auxiliary Space /
Extra Space

$O(N)$

Priority Queue

$$(6, 8) = \sqrt{6^2 + 8^2} = 10$$

$$(3, 2) = \sqrt{3^2 + 2^2} = \sqrt{13}$$

$$\begin{aligned} (-4, 1) &= \sqrt{(-4)^2 + 1^2} = \sqrt{17} \\ (1, -1) &= \sqrt{1^2 + (-1)^2} = \sqrt{2} \\ (1, 4) &= \sqrt{1^2 + 4^2} = \sqrt{17} \\ (8, 6) &= \sqrt{8^2 + 6^2} = \sqrt{100} = 10 \end{aligned}$$

$$(12, 5)$$

$$\sqrt{12^2 + 5^2} = 13$$

$$(0, 0)$$

$$\sqrt{0^2 + 0^2} = 0$$

①

$$(3, 4)$$

$$\sqrt{3^2 + 4^2} = 5$$

$$(-5, 12)$$

$$\sqrt{(-5)^2 + 12^2} = 13$$

$$\{(0, 0), (1, -1), (-1, 1), (1, 4)\}$$

Smaller Euclidean
Distance

more closer
to origin

Higher Probability

Pop First
from PQ

Marks of PCM

You are provided with marks of **N** students in Physics, Chemistry and Maths.

Perform the following 3 operations:

- Sort the students in Ascending order of their Physics marks.
- Once this is done, sort the students having same marks in Physics in the descending order of their Chemistry marks.
- Once this is also done, sort the students having same marks in Physics and Chemistry in the ascending order of their Maths marks.

Input:

N = 10

phy[] = {4 1 10 4 4 4 1 10 1 10}

chem[] = {5 2 9 6 3 10 2 9 14 10}

math[] = {12 3 6 5 2 10 16 32 10 4}

Output:

- ① Physics
↓ same
- ② Chemistry
↓ same
- ③ Maths

1 1 1 4 4 4 4 10 10
14 2 2 10 6 5 3 10 9
10 3 16 10 5 12 2 4 5

P.S.G.M
class MarksComp implements Comparator<Marks>{

```
public int compare(Marks s1, Marks s2){  
    if (s1.physics HP < s2.physics) {  
        return -1;  
    } else if (s1.physics > s2.physics) {  
        return +1;  
    } else {
```

```
if( s1.chem < s2.chem) {  
    return +1;  
}
```

```
else if( s1.chem > s2.chem) {  
    return -1;  
}
```

```
} else {  
    if( s1.math < s2.math)  
        return -1;  
    else  
        return +1;  
}
```

```
class MarksCmp implements Comparator<Marks>{
    public int compare(Marks s1, Marks s2){
        if(s1.phy < s2.phy) return -1; your key (increasing)
        else if(s1.phy > s2.phy) return +1;
        else {
            if(s1.chem < s2.chem) return +1; your key
            else if(s1.chem > s2.chem) return -1; (Decreasing)
            else {
                if(s1.maths < s2.maths) return -1;
                return +1;
            }
        }
    }
}
```

```
PriorityQueue<Marks> q = new PriorityQueue<>(new MarksCmp());  
  
for(int i=0; i<N; i++){  
    Marks stud = new Marks(phy[i], chem[i], math[i]);  
    q.add(stud);  
}  

```

Insertion
 $O(N \log N)$

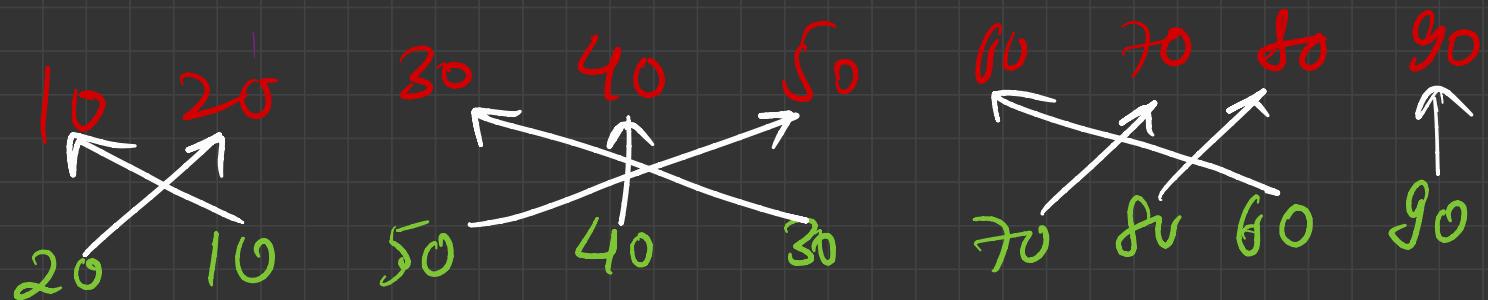
```
int idx = 0;  
while(q.size() > 0){  
    Marks stud = q.remove();  
    phy[idx] = stud.phy;  
    chem[idx] = stud.chem;  
    math[idx] = stud.maths;  
    idx++;  
}
```

Removal
 $O(N \log N)$

~~$O(N \log N)$~~
Custom
Priority(Heap Sort)

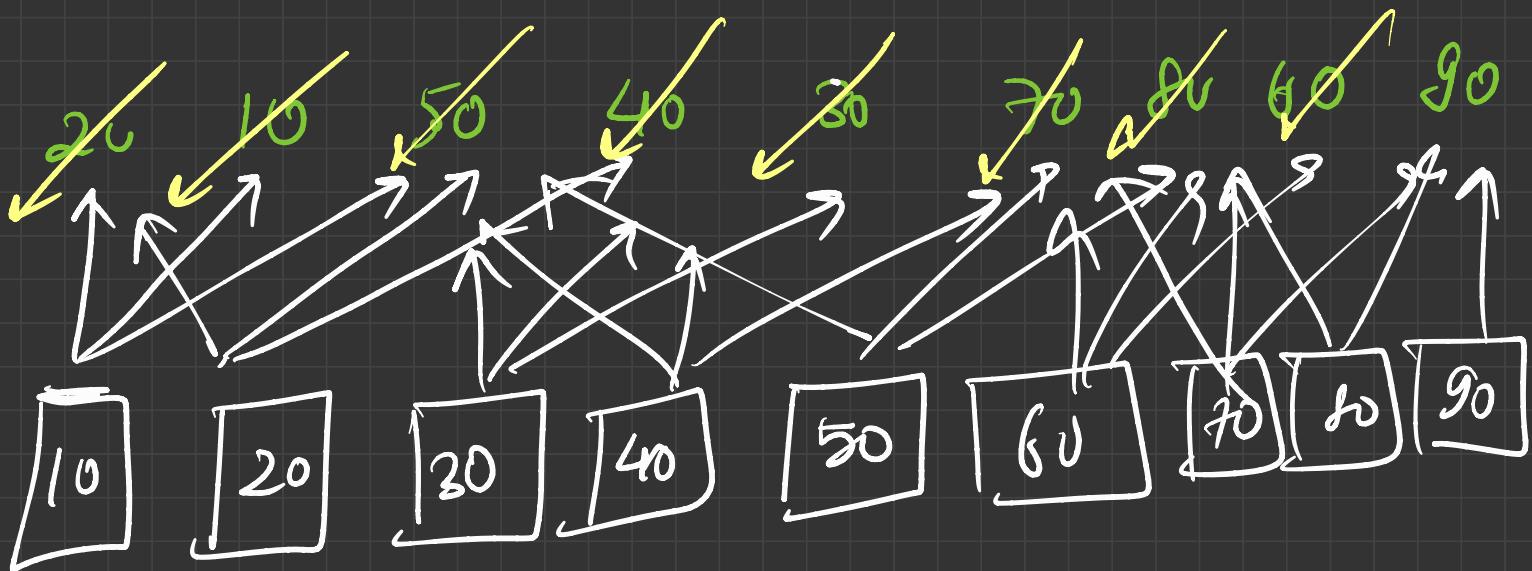
Nearly Sorted Array \rightarrow Sort

$k=2$ \rightarrow At most distance b/w actual position & sorted posn

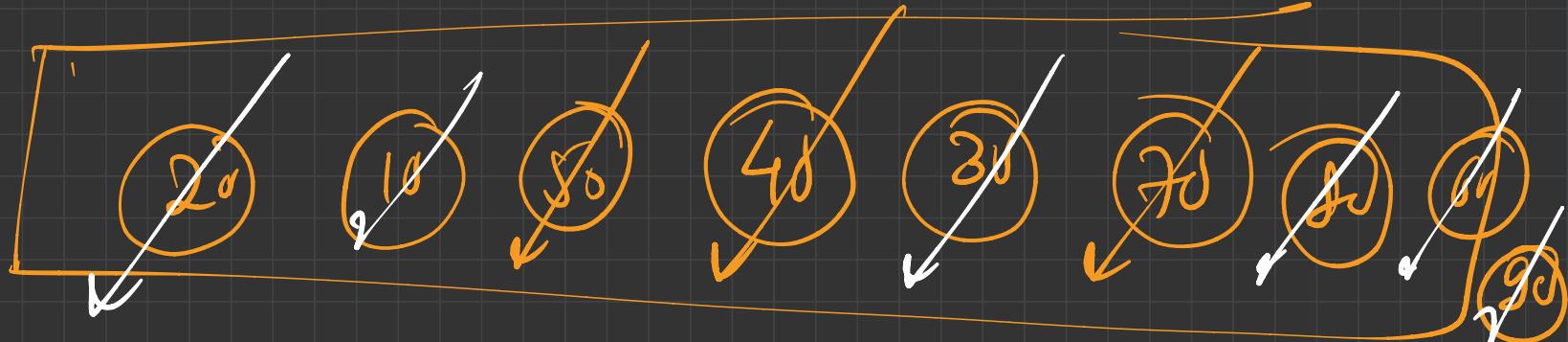
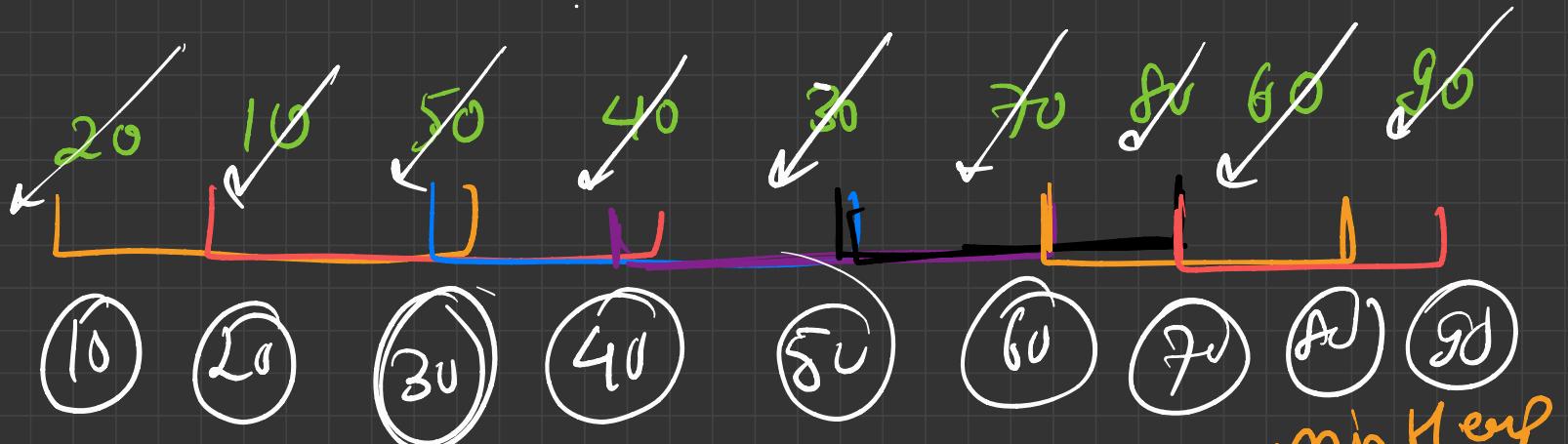


Bubbleforce \rightarrow Unsorted sorted \rightarrow Sort
 \rightarrow Time $\rightarrow O(n \log n)$

$k=2$



$k=2$

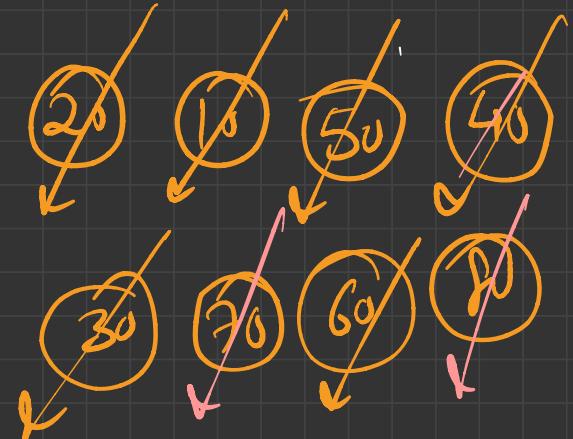
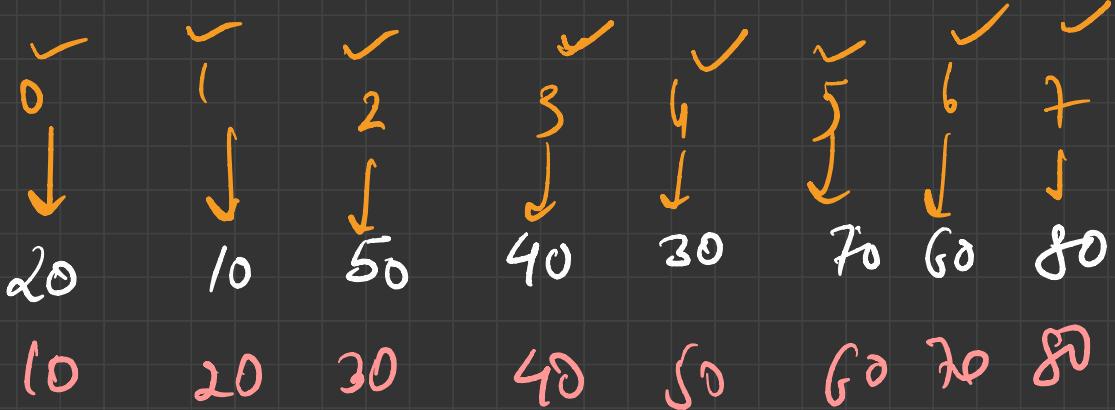


```

PriorityQueue<Integer> q = new PriorityQueue<>();
for(int i=0; i<n; i++){
    q.add(arr[i]);  $\rightarrow N \log K$ 
    if(i >= k){
        System.out.print(q.remove() + " ");
    }
}

while(q.size() > 0){  $\rightarrow N \log K$ 
    System.out.print(q.remove() + " ");
}

```



$O(N \log K)$

Time

$O(K)$ PQ

Space

$[K=2]$

Leetcode 295

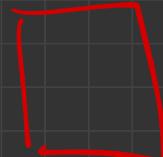
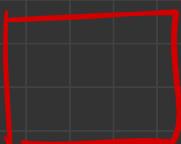
median - Stream

Online

Alg

Data

Alg



sol, Nosol

offline Data

Time series

Online data

Median in Data Stream

40
insert

median
40

70
insert

median
55

50
insert

Median
50

80
insert

median
60

10
insert

Median
50

30
insert

Median
45

Approach ①

ArrayList operation

[40] → 1 operation

median → $O(1)$ → 40

Time Complexity
worst-case
 $O(n)$

[40, 70] → 1 operation

median → $O(1)$ → 55

Insertion → $O(1)$
Get median → $O(1)$

[40, 50, 70] → 2 operations

median → $O(1)$ → 50

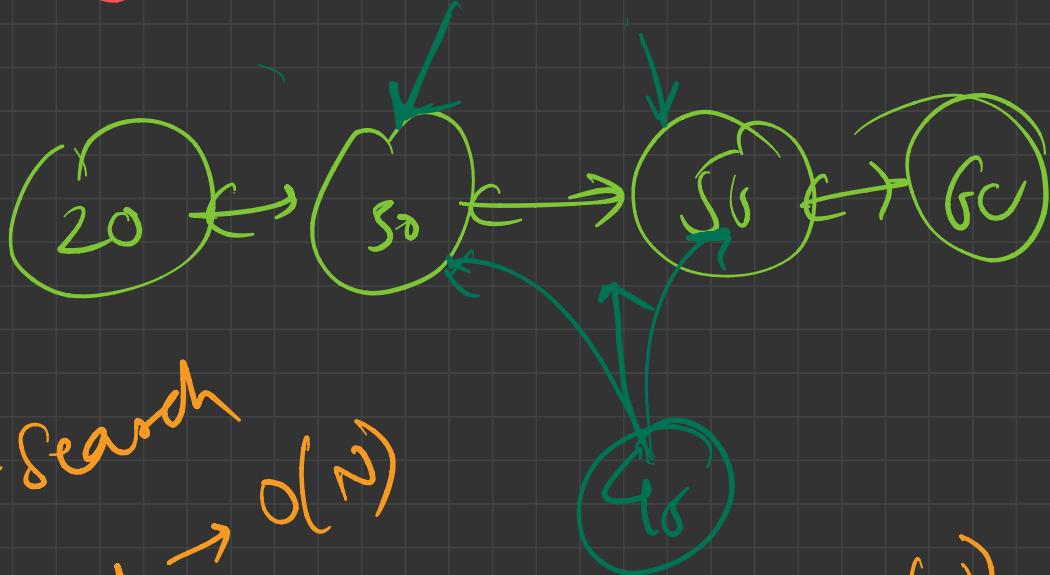
$[40, 50, 70, 80]$ → 1 operation

median $\rightarrow \mathcal{O}(1) \rightarrow 50$

$[10, 40, 50, 70, 80]$ → 5 operations

median $\rightarrow \mathcal{O}(4) \rightarrow 5$

Approach ② linked list (Doubly)



linear search
insert $\rightarrow O(N)$

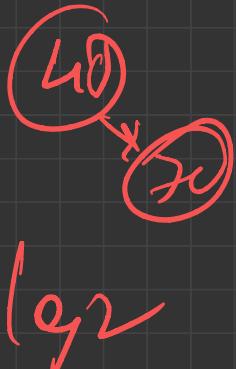
get median $\rightarrow O(N)$

Approach ③ ordered_set / TreeSet

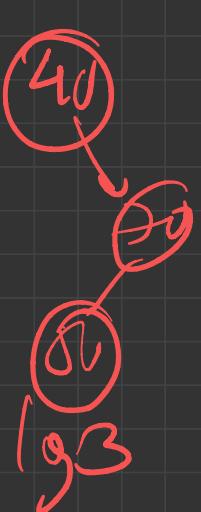
insert
log n

Get median
log n

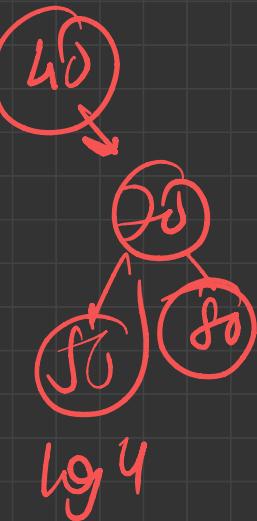
log 1



log 2



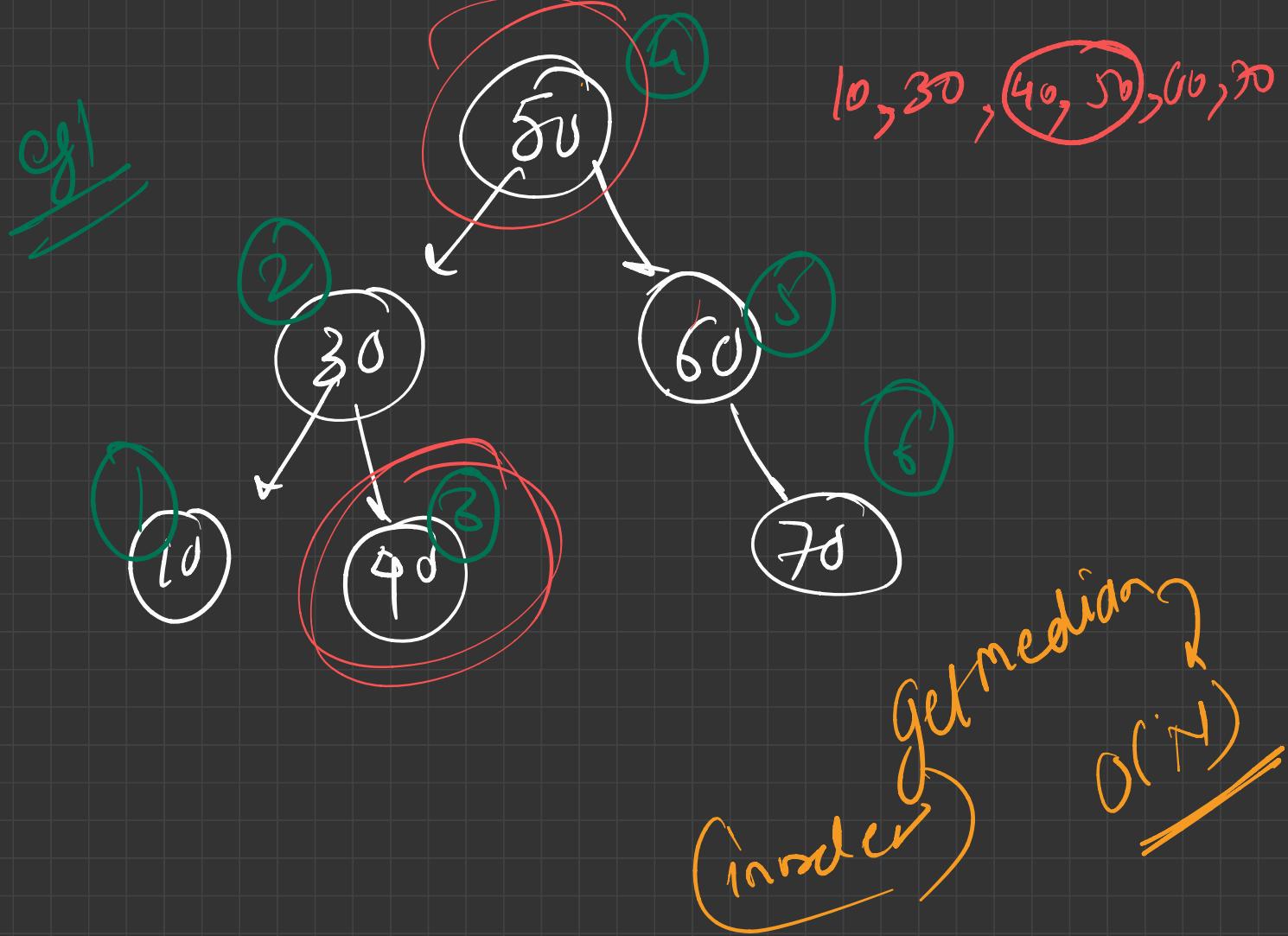
log 3



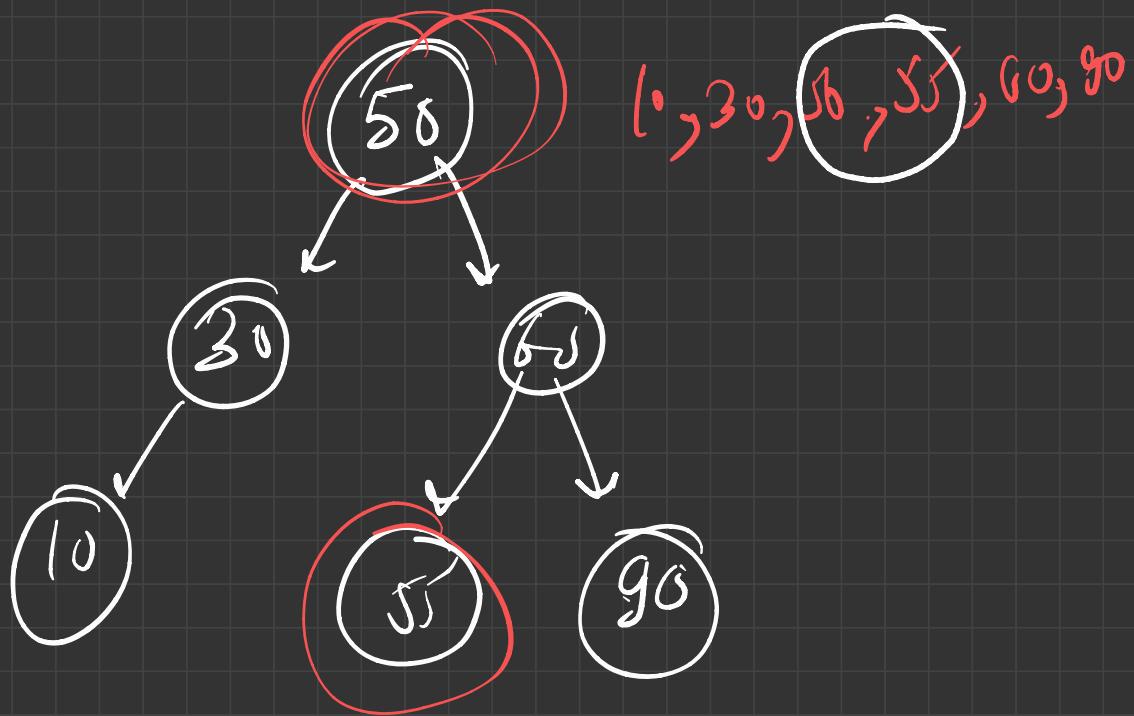
log 4

Binary search Tree

sorted data

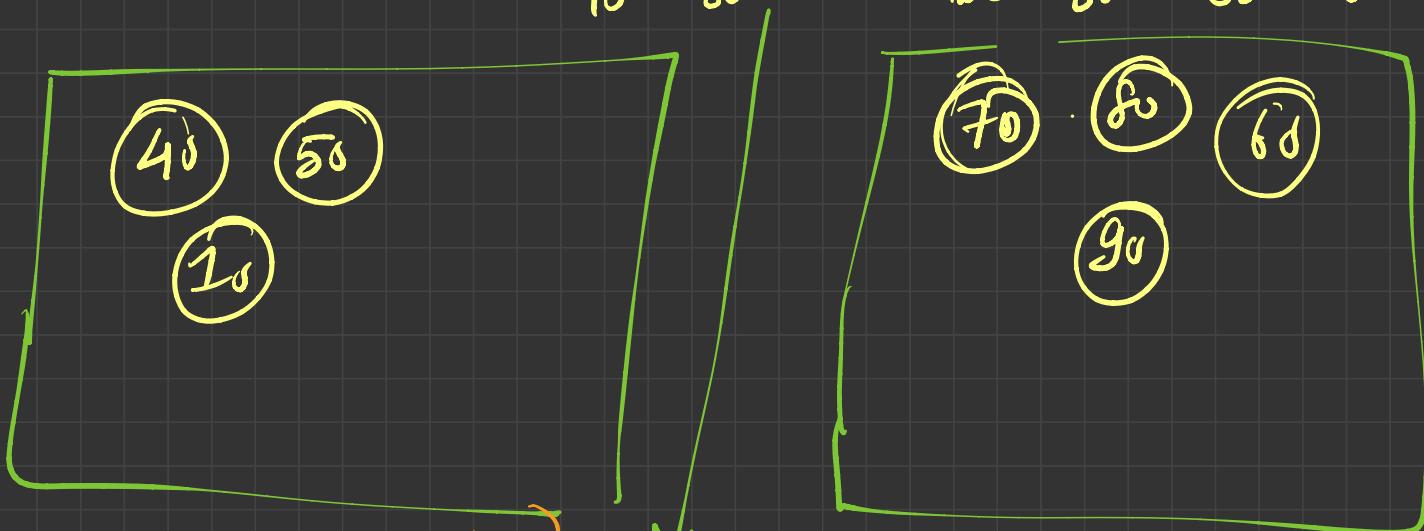


esl



Approach ④ Priority Queue

insert in order → [40, 70, 50, 80, 60, 10, 60, 55, 60, 90]



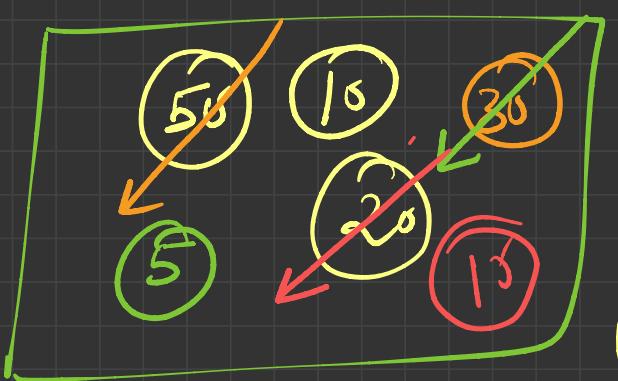
left (maxHeap)

median = 55

right (minHeap)

insert order \rightarrow

[50, 10, 30, 20, 5, 15]
60 30 30 $\frac{20+30}{2} = 25$ 20 (17.5)



minHeap

median = 20

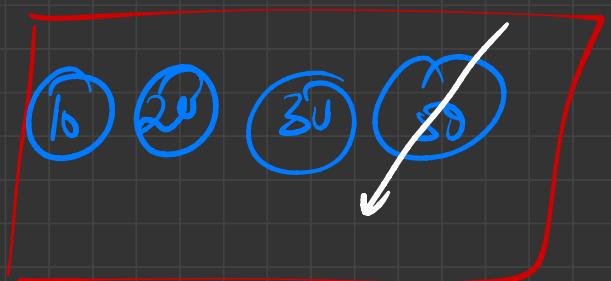


maxHeap

~~5, 10, 20~~, [30, 50]

insertion order \rightarrow [20, 50, 70, 60, 10, 30, 90]

20 35 50 55 60 70 90 10 30 50



min Heap
(left)

median



max Heap
(right)

10:55

```

PriorityQueue<Integer> left, right;

public MedianFinder() {
    left = new PriorityQueue<>(Collections.reverseOrder()); // Max Heap
    right = new PriorityQueue<>(); // Min Heap
}

```

```

public void addNum(int num) {
    // Insertion
    double median = findMedian();
    if(num <= median){
        left.add(num);
    } else {
        right.add(num);
    }

    // Balancing
    if(left.size() > right.size() + 1){
        right.add(left.remove());
    } else if(right.size() > left.size() + 1){
        left.add(right.remove());
    }
}

```

Insertion Removal

```

// Find Median -> Time = O(1)
public double findMedian() {
    if(left.size() + right.size() == 0) return -1;

    if(left.size() == right.size()){
        return (left.peek() + right.peek()) / 2.0;
    } else if(left.size() == right.size() + 1){
        return left.peek();
    } else {
        return right.peek();
    }
}

```

Insertion Time \rightarrow $(\log n + \log n)$

Get median Time \rightarrow $O(1)$

Recording → Backlog ↑

live classes

Reorganize String

"aaaaaa bbbcc"



a b a b a b a c a c a

No two adjacent characters are same

"aaaabbbccccdddo"



a b a b a b a b c d c d c d c

~~a → 10~~
~~b → 1000B0~~
~~c → 10~~
~~d → 10~~
~~e → 10~~

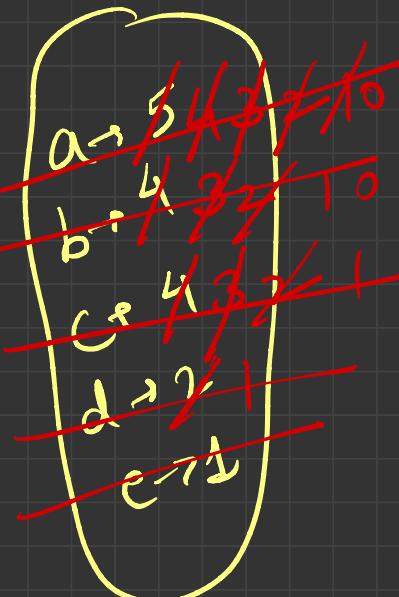
"aabbbbbbcdde"

Intuition
↓
Place
the
highest
frequency
character

a b a b — b — b — b

b a b a b c b d b e b

b a b a b c b d b e b


 " a a a a a b b b b c c c c d d e "

a → 5 → 4 → 3 → 2 → 1 → 0
 b → 4 → 3 → 2 → 1 → 0
 c → 4 → 3 → 2 → 1 → 0
 d → 2 → 1 → 0
 e → 1 → 0

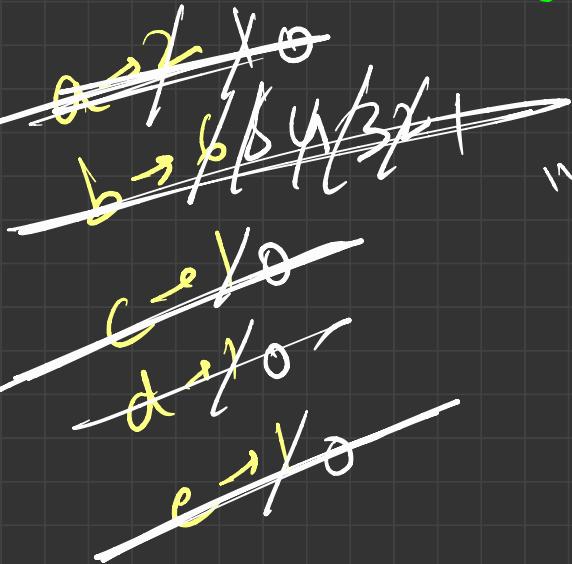
~~a b a c a b c a b c d~~
~~a b c d e~~
→ main heap of fair

① Pop/Place the highest Frequency character first

2 pairs pop

② Two pops cannot be of same character such that during 2nd pop, B↑C↑ is not in PQ

"aab bbb bb cde"



"babab cb d be b"

2pop → 2push alternatively

```
public static class Pair{  
    char ch;  
    int freq;  
  
    Pair(char ch, int freq){  
        this.ch = ch;  
        this.freq = freq;  
    }  
}  
  
public static class FreqComp implements Comparator<Pair>{  
    public int compare(Pair a, Pair b){  
        // Higher Frequency -> Higher Priority  
        return (b.freq - a.freq);  
    }  
}
```

max heap
S2-S1

min heap

S1-S2

```
HashMap<Character, Integer> freq = new HashMap<>();  
for(int i=0; i<s.length(); i++){  
    freq.put(s.charAt(i), freq.getOrDefault(s.charAt(i), 0) + 1);  
}  
  
PriorityQueue<Pair> q = new PriorityQueue<>(new FreqComp());  
for(Character ch: freq.keySet()){  
    q.add(new Pair(ch, freq.get(ch)));  
}  
  
StringBuilder str = new StringBuilder("");
```

```
StringBuilder str = new StringBuilder("");
|
while(q.size() > 0){
    Pair odd = q.remove();
    str.append(odd.ch);
    odd.freq--;
}

if(q.size() > 0){
    Pair even = q.remove();
    str.append(even.ch);
    even.freq--;
}

    if(even.freq > 0){
        q.add(even);
    }

    if(odd.freq > 0){
        q.add(odd);
    }

}
}

return str.toString();
```

aaaa bbbb

a \rightarrow y β γ δ γ δ α abababab
b \rightarrow y β γ δ γ δ α

"aaaaaa bb c"

~~a → 6 8 4 3~~

~~b → 2 X 0~~

~~c → 1 X 0~~

ababaca ada

q size = -1
for ?)

```
while(q.size() > 0){
    Pair odd = q.remove();
    str.append(odd.ch);
    odd.freq--;
}

if(q.size() == 0 && odd.freq > 0){
    return "";
}

if(q.size() > 0){
    Pair even = q.remove();
    str.append(even.ch);
    even.freq--;
}

if(even.freq > 0){
    q.add(even);
}
}

if(odd.freq > 0){
    q.add(odd);
}
}

return str.toString();
```