

Recursion & Backtracking

① Introduction

```
</> Print Increasing  
</> Print Decreasing  
</> Print Increasing Decreasing  
</> Power-linear  
</> Power-logarithmic  
</> Print Zigzag  
</> Tower Of Hanoi
```

② Recursion & Arrays

```
</> Display Array  
</> Display Array In Reverse  
</> Max Of An Array  
</> First Index  
</> Last Index  
</> All Indices Of Array
```

③ Recursion & ArrayList {Get}

```
</> Get Subsequence  
</> Get Kpc  
</> Get Stair Paths  
</> Get Maze Paths  
</> Get Maze Path With Jumps
```

⑤ Backtracking

```
</> Flood Fill  
</> Target Sum Subsets  
</> N Queens  
</> Knights Tour
```

④ Recursion on the way up (Print)

```
</> Print Subsequence  
</> Print Kpc  
</> Print Stair Paths  
</> Print Maze Paths  
</> Print Maze Paths With Jumps  
</> Print Encodings  
</> Print Permutations
```

Principle of Mathematical Induction (PMI)

① Trivial Case

$$\sum_{i=1}^0 i = 0$$

$$\sum_{i=1}^1 i = 1$$

$$\sum_{i=1}^n i = \frac{n*(n+1)}{2}$$

② Assume

$$\sum_{i=1}^k i = \frac{k*(k+1)}{2}$$

③ To prove

$$\begin{aligned}\sum_{i=1}^{k+1} i &= (\underbrace{1+2+\dots+k}_{\text{LHS}}) + (k+1) = \frac{k*(k+1)}{2} + (k+1) \\ &= \frac{(k+1)*(k+2)}{2} \end{aligned}$$

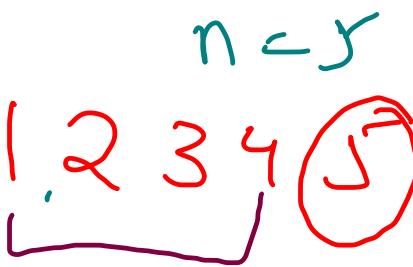
if $n=k+1$ LHS = RHS

Recursion

{ High-level thinking }

- ① Recursive function \rightarrow expectation

void printInc(int n) \rightarrow 1, 2, 3, - ... n



- ② Recursive fn \rightarrow smaller values \rightarrow faith (call)

void printInc(n-1) \rightarrow 1, 2, 3, - ... n-1 | 2 3 4

- ③ Meeting Expectation with faith

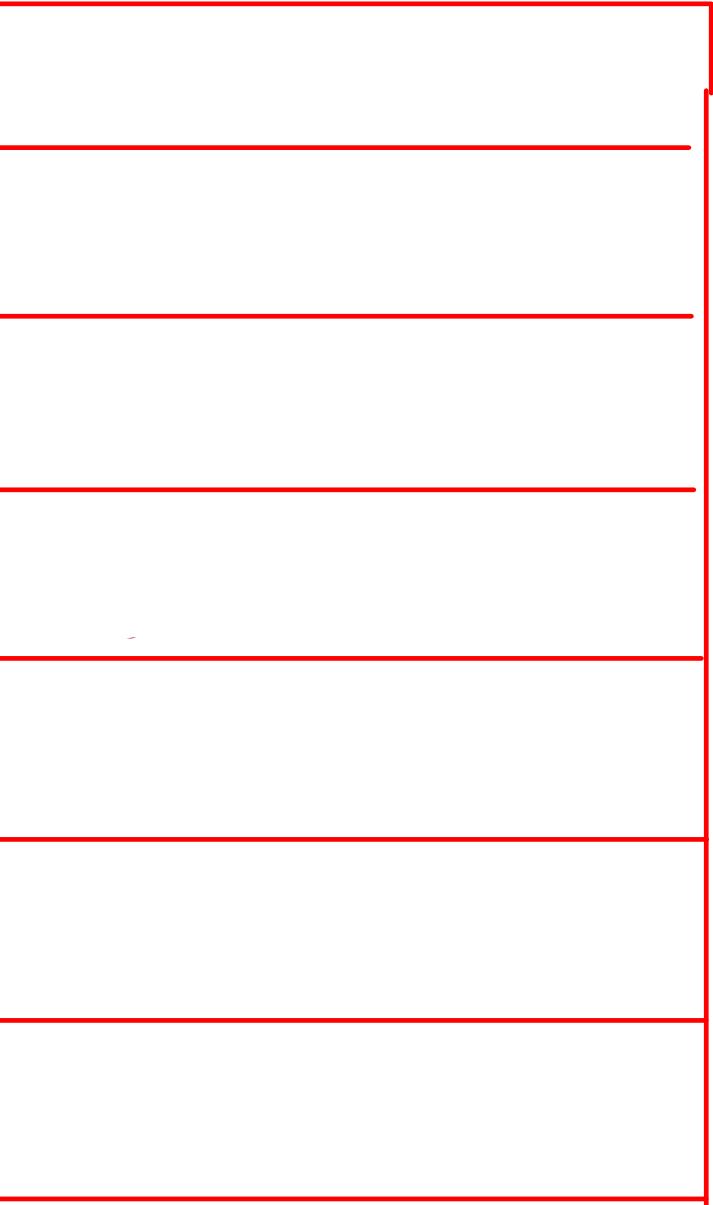
System.out.printn(n);

- ④ Base Case ↳ To stop recursion
- \swarrow Stack overflow / memory limit exceeded

low-level thinking

function call stack

```
void printInc(int n) {  
    Base case → ① if(n == 0) return;  
    Fact → ① printInc(n-1);  
    Meeting Expectation → ② System.out.println(n);  
  
    main() {  
        printInc(5);  
    }
```



```
public static void main(String[] args) throws Exception {
    Scanner scn = new Scanner(System.in);
    int n = scn.nextInt();
    printIncreasing(n);
}

public static void printIncreasing(int n){
    if(n == 0){
        // Base Case
        return;
    }

    // 1. Faith : 1, 2, .. n-1
    printIncreasing(n - 1);

    // 2. Meet Expectation with Faith
    System.out.println(n);
}
```

Tail Recursion

work is being done while returning

Q2) Point Decreasing

① pointDecreasing(int n) :->

Expectation

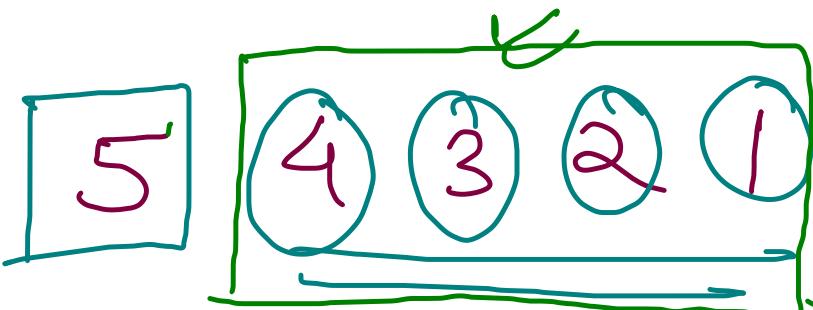
$n, n-1, n-2, \dots, 1$

② pointDecreasing(n-1) :->

$n-1, n-2, n-3, \dots, 1$

fail

③ System.out.println(b) :-> Meeting Expectation
Should be
before the
call.



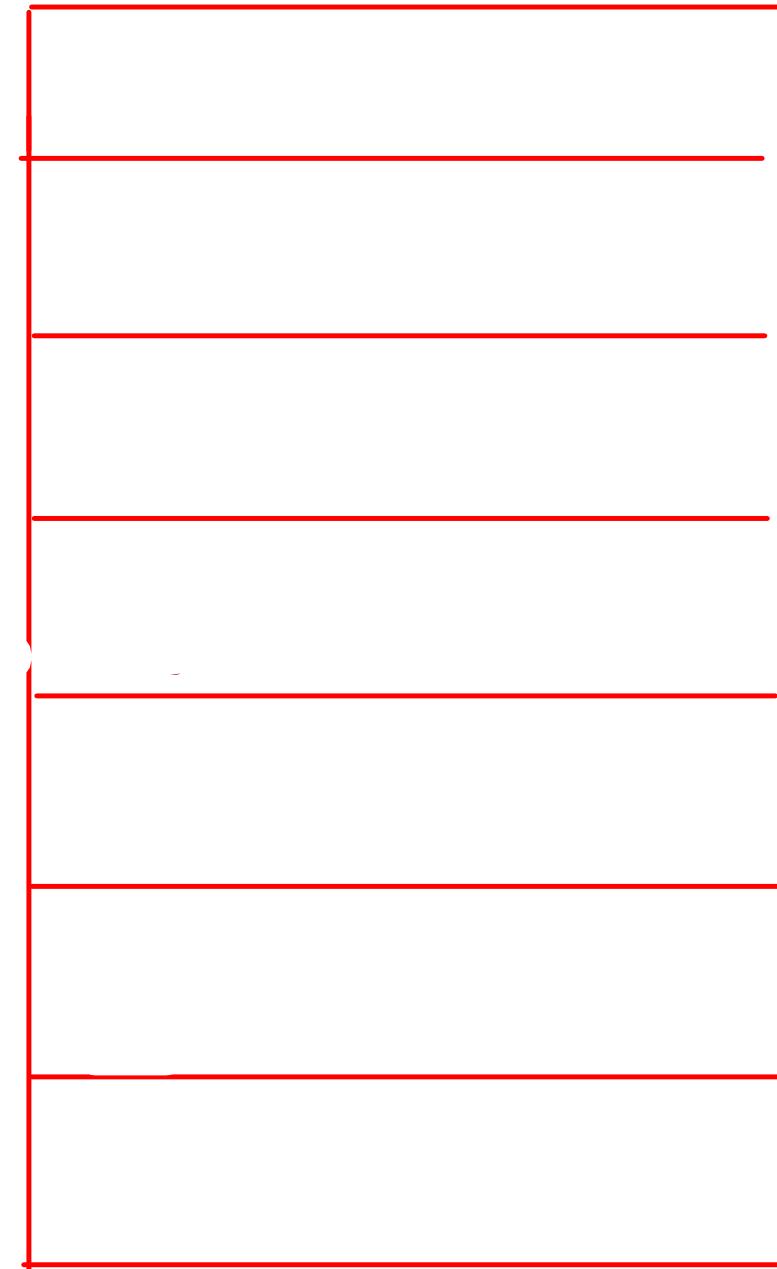
$n=5$ 5 4 3 2 1

void printDec (int n) {
 ① if ($n == 0$) return;
 System.out.println (n); } → Border

② printDec (n-1);

3

~~head recursion~~
work is done before
calling functions



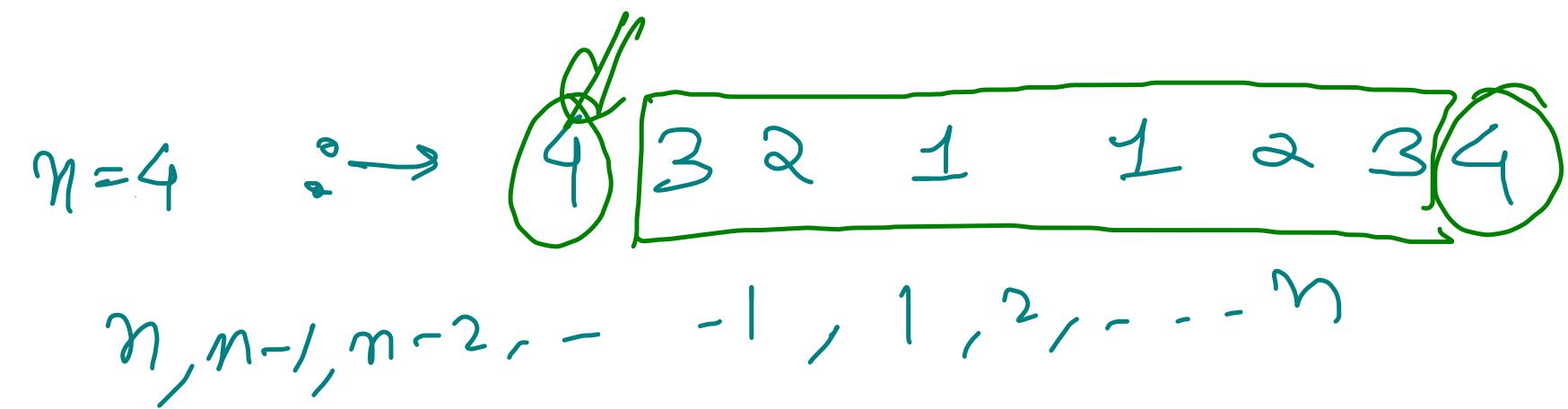
5, 4, 3, 2, 1

Q) Print Increasing Decreasing

High-Level Thinking

① Expectation

PDI(int n) : →



② Fault

PDI(n-1) : →

$n=3 \rightarrow$ 3 2 1 1 2 3

$n-1, n-2, n-3, \dots, 1, 1, 2, 3, \dots, n-2, n-1$

③ Meeting Expectation

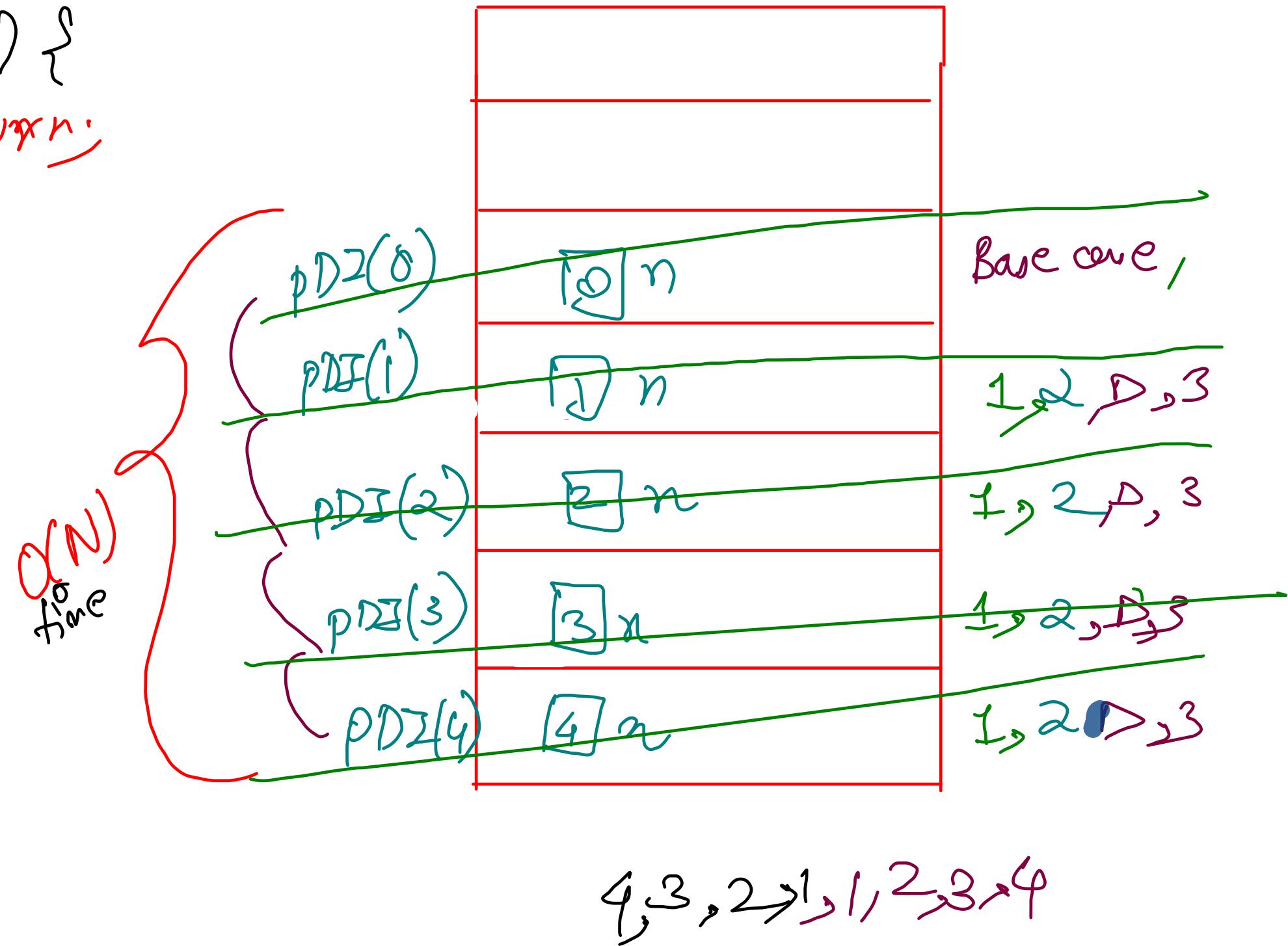
SPO(n) → Preorder
Postorder

$$N = 4 \quad 4, 3, 2, 1, 1, 2, 3, 4$$

```

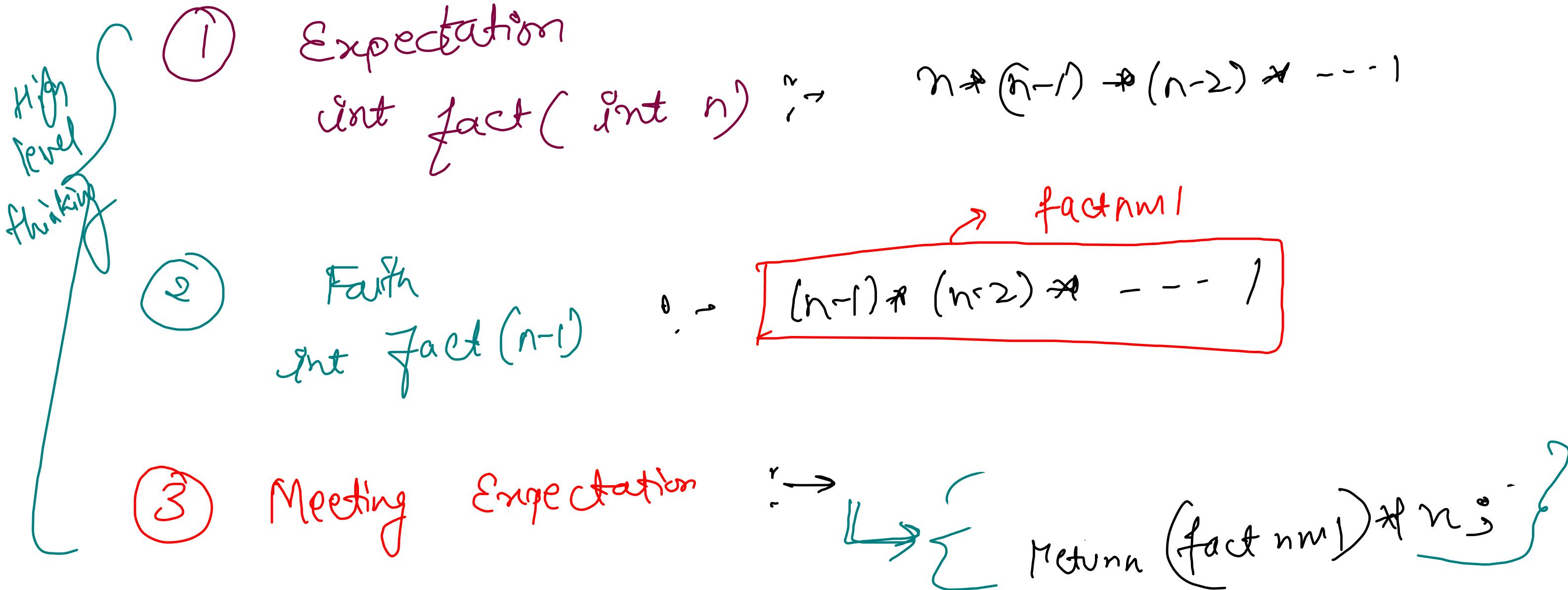
void PDI( int n ) {
    if ( n == 0 ) return;
    Preorder{ ① } Sys0(n);
    father{ ② } PDI( n-1 );
    Postorder{ ③ } Sys0(n);
}
main() {
    PDI( n );
}

```



Factorial

$$5! = 5 * 4 * 3 * 2 * 1$$



low-level thinking

$$N=4 \quad 4 \times 3 \times 2 \times 1$$

```
int fact (int n) {
```

Base case ① if ($n == 0$) return 1;

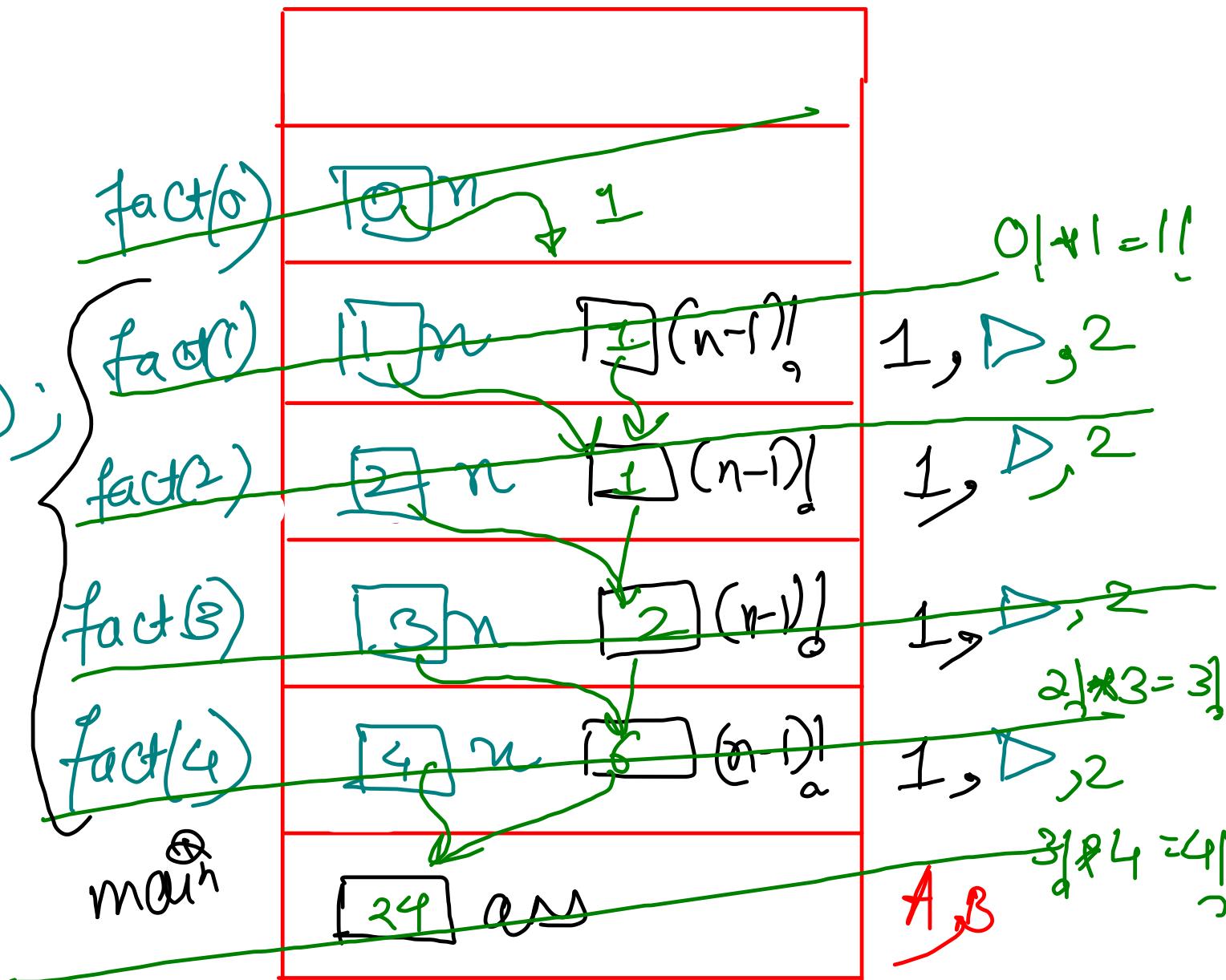
faith { ② int factm1 = fact(n-1);

Meeting
Expectation
3 }
return factm1 * n;

```
main() {
```

A) int ans = fact(4);
B) System.out.println(ans);

Call
 $n = N$
 $O(N)$



Power - Linear $\Rightarrow O(N)$

$$a^b = a * a * a * \dots b \text{ times}$$

$$3^5 = 3 * 3 * 3 * 3 * 3$$

① Expectation
power(x, n) $\xrightarrow{\text{constant variable}}$

$$\{ 3 * 3 * 3 * 3 * 3 \} \\ x * x * x * \dots n \text{ times}$$

② Faith
power($x, n-1$) $\xrightarrow{\text{ }}$

$$\{ 3 * 3 * 3 * 3 \} \\ x * x * \dots (n-1) \text{ times}$$

③ Meeting Expectation $\xrightarrow{\text{ }}$

return

$$\text{pxnum} * x^j \\ 3^4 * 3 = 3^5$$

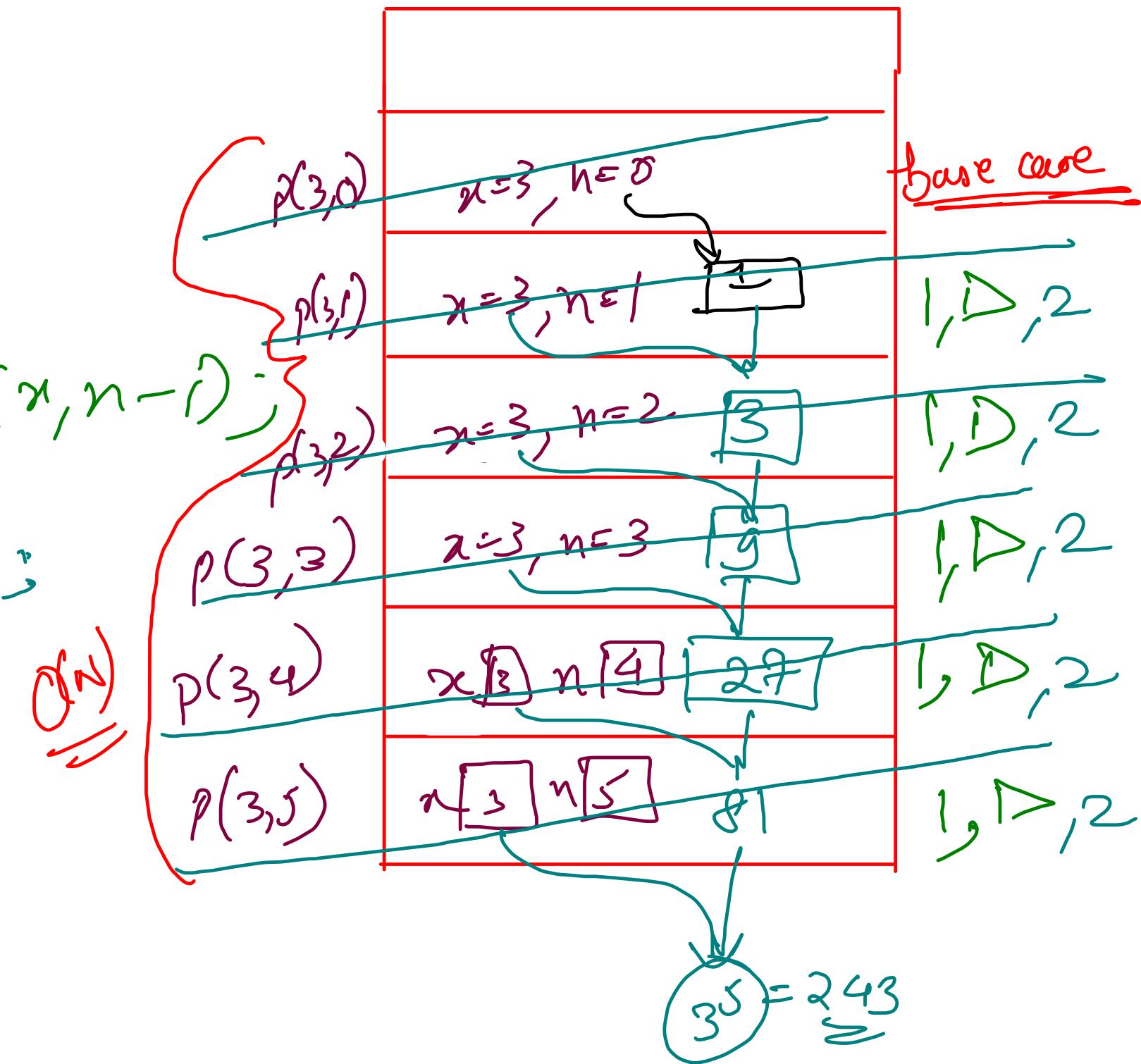
Low-level Thinking

```
int power(int x, int n) {
```

① if ($n == 0$) return 1;

② int pxnml = power(x, n-1);
return pxnml * x;

3



```
public static void main(String[] args) throws Exception {
    // write your code here
    Scanner scn = new Scanner(System.in);
    int x = scn.nextInt();
    int n = scn.nextInt();
    int ans = power(x, n);
    System.out.println(ans);
}

public static int power(int x, int n){
    if(n == 0) return 1; //  $x^0 = 1$ 
    int pxnm1 = power(x, n - 1); // faith
    return pxnm1 * x; // meeting expectation
}
```

Power - logarithmic

$$x = \cancel{x}^{\frac{n}{2}} + \cancel{x}^{\frac{n}{2}}$$

$$a^b = a * a * a * \dots \text{ b times}$$

$$3^5 = 3 * 3 * 3 * 3 * 3$$

$$3^8 = 3^4 * 3^4$$

$$\cancel{x}^{\frac{n}{2}} + \cancel{x}^{\frac{n}{2}}$$

$$3^9 = 3^4 * 3^4 * 3$$

~~if odd~~

① Expectation
 power(x, n) $\stackrel{\substack{\text{constant} \\ \text{variable}}}{\rightarrow}$

$\{ 3 * 3 * 3 * 3 * 3 \}$
 $x * x * x * \dots \text{ n times}$

② fair
 power($x, n/2$) $\stackrel{\substack{\text{?} \\ \text{?}}}{\rightarrow}$

$\overbrace{x * x * \dots}^{\text{n/2 times}} \xrightarrow{\text{pxnby2}}$

③ Meeting Expectation $\stackrel{\substack{\text{?} \\ \text{?}}}{\rightarrow}$

return $\boxed{\begin{array}{l} \text{pxnby2} \Rightarrow \text{pxnby2} \\ 0 * x \text{ if } (n/2 \text{ is } 1) \end{array}}$

Low-level Thinking

```
int power(int x, int n) {
```

① if ($n == 0$) return 1;

② int pxnby2 = power(x, $n/2$);

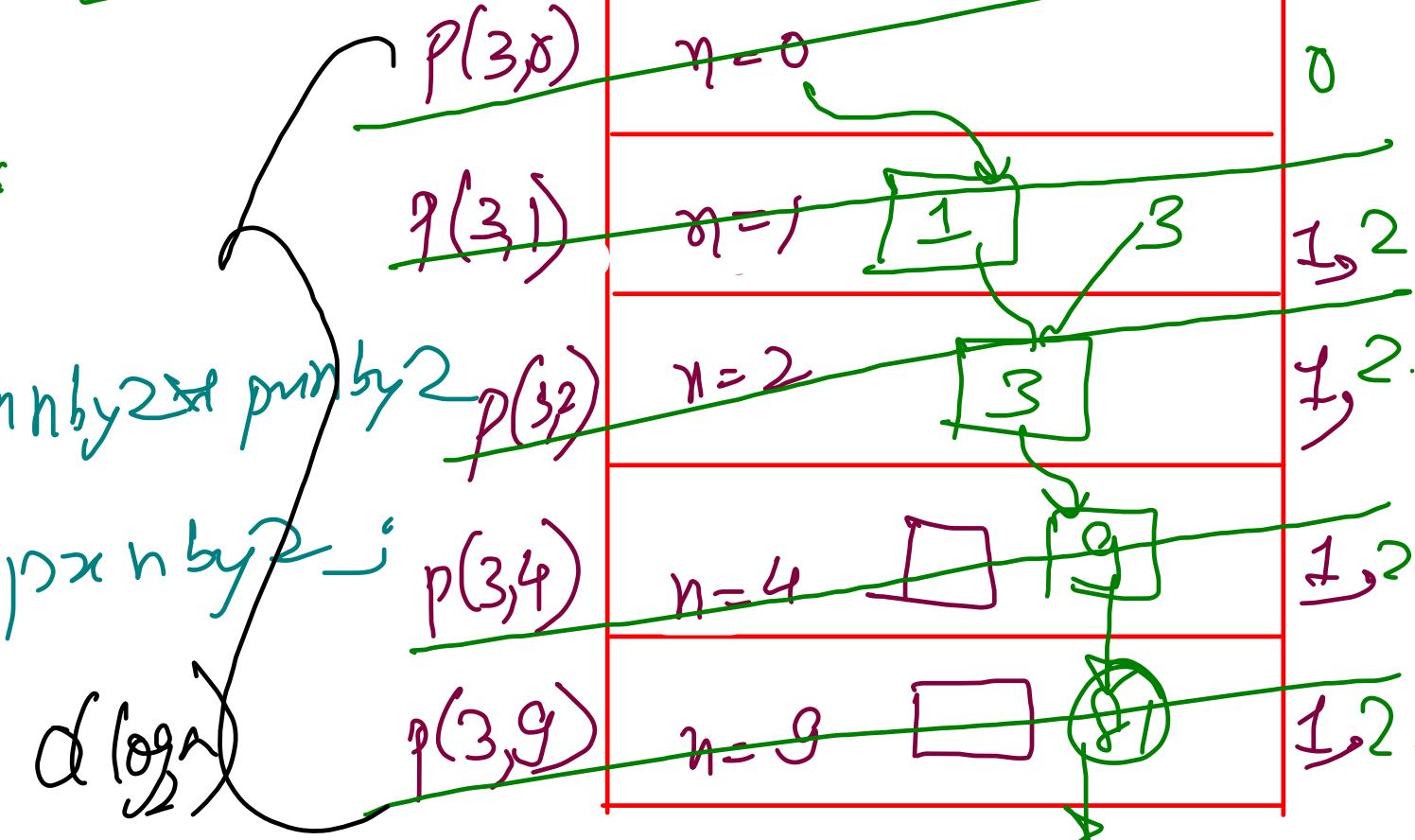
if ($n \% 2 == 1$)
return $x * pxnby2 * pxnby2$;
else
return $pxnby2 * pxnby2$;

3

$$3^2 = 3^1 * 3^1$$

$$3^4 = 3^2 * 3^2$$

$$3^9 = 3^4 * 3^4 * 3$$



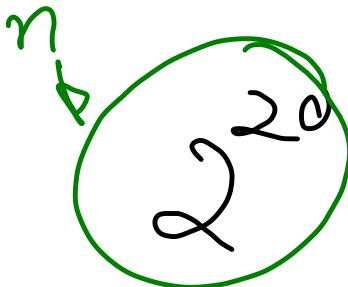
$n = 9$

$$3^9 = 81 * 81 * 3$$

```
public static void main(String[] args) throws Exception {  
    Scanner scn = new Scanner(System.in);  
    int x = scn.nextInt();  
    int n = scn.nextInt();  
    int ans = power(x, n);  
    System.out.println(ans);  
}
```

```
public static int power(int x, int n){  
    if(n == 0) return 1;  
  
    int pxnby2 = power(x, n/2);  
  
    if(n % 2 == 1) return pxnby2 * pxnby2 * x;  
    else return pxnby2 * pxnby2;  
}
```

$$n = 2^x \Rightarrow \text{no of calls} = x$$



$$\log n = \log(2^x)$$

$$O(n) \Rightarrow 1048576$$

$$\log n = x \Rightarrow \log_2 2$$

$$O(\log_2 n) \Rightarrow O(\log_2 2^{20}) \\ = O(20)$$

$$\log n = x \Rightarrow \text{no of call} \Rightarrow O(\log_2 n) \text{ constant}$$

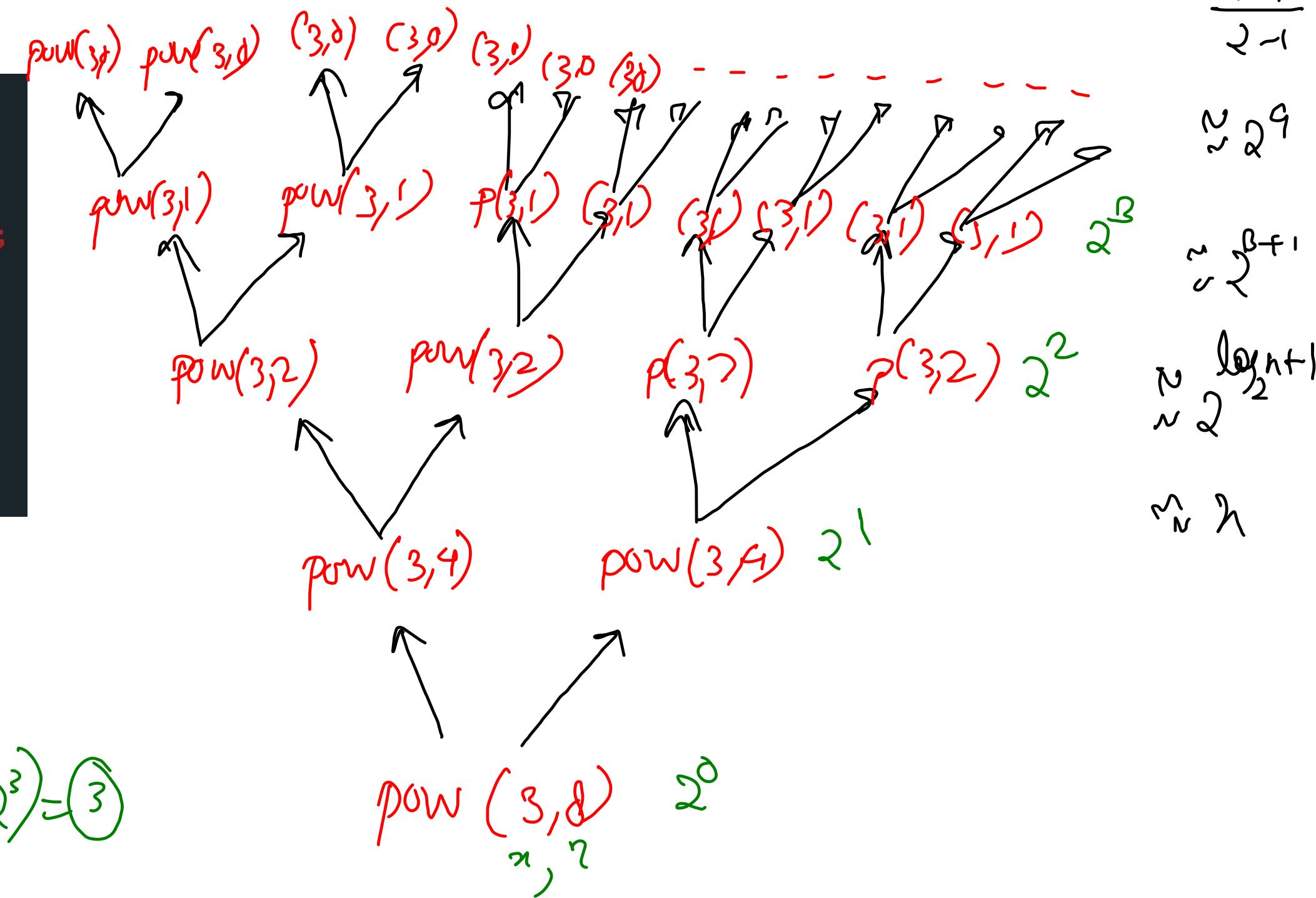
Today's Questions

- ① Power - 3rd method
- ② Print zigzag
- ③ Display Array
- ④ Display Array - Reverse of Home work }
- ⑤ Max of Array
- ⑥ First Index of Array
- ⑦ Last Index of Array

Power \rightarrow IIIrd method

$$2^0 + 2^1 + 2^2 + 2^3 = \frac{2^4 - 1}{2 - 1}$$

```
public static int power(int x, int n) {  
    if(n == 0){  
        return 1;  
    }  
    int xpn = power(x, n/2) * power(x, n/2);  
  
    if(n % 2 == 1){  
        xpn = xpn * x;  
    }  
  
    return xpn;
```



(Q) Print ZigZag

① Expectation :-

printZigzag(int n)

3 2 1 1 1 2 1 1 1 2 3 2 1 1 1 2 1 1 1 2 3

②

Fair :- printZigzag(n-1)

2 1 1 2 1 1 2

③ Meet Expectation :-

sys(n)

printZigzag(n-1)

sys(n)

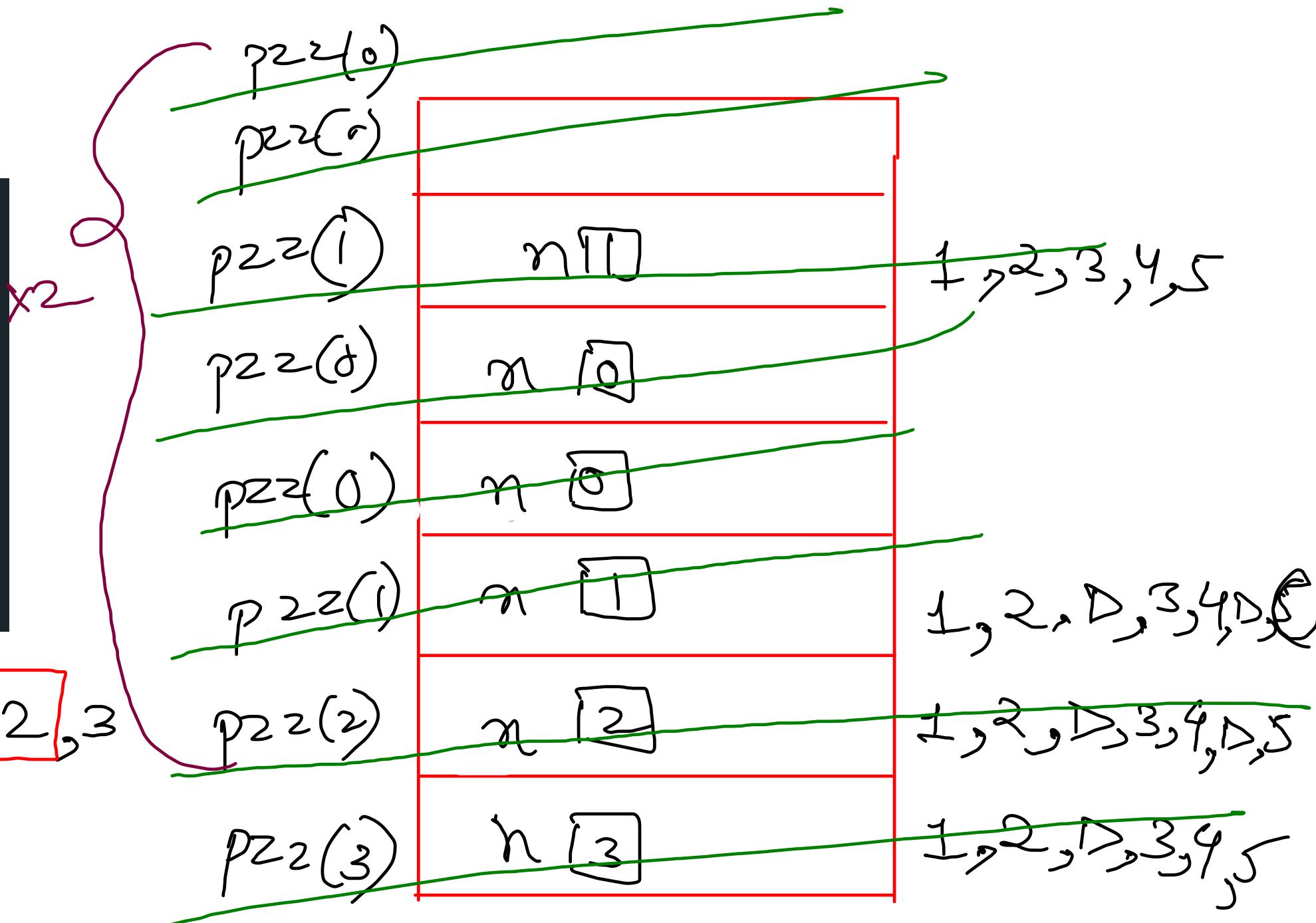
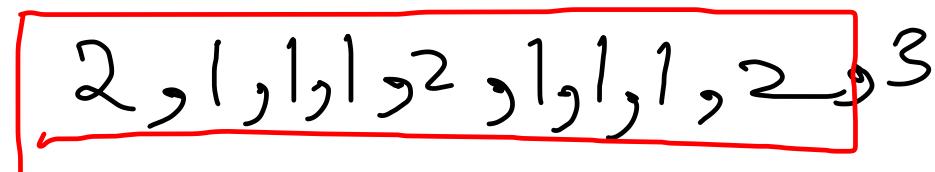
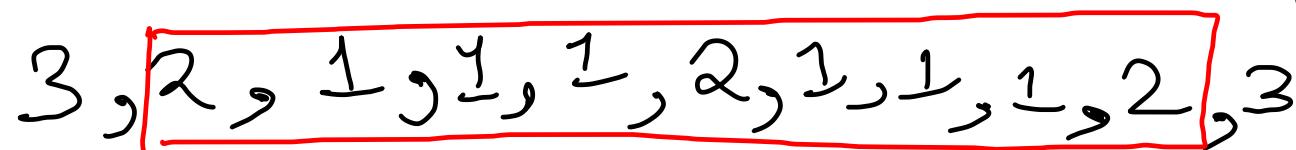
printZigzag(n-1)

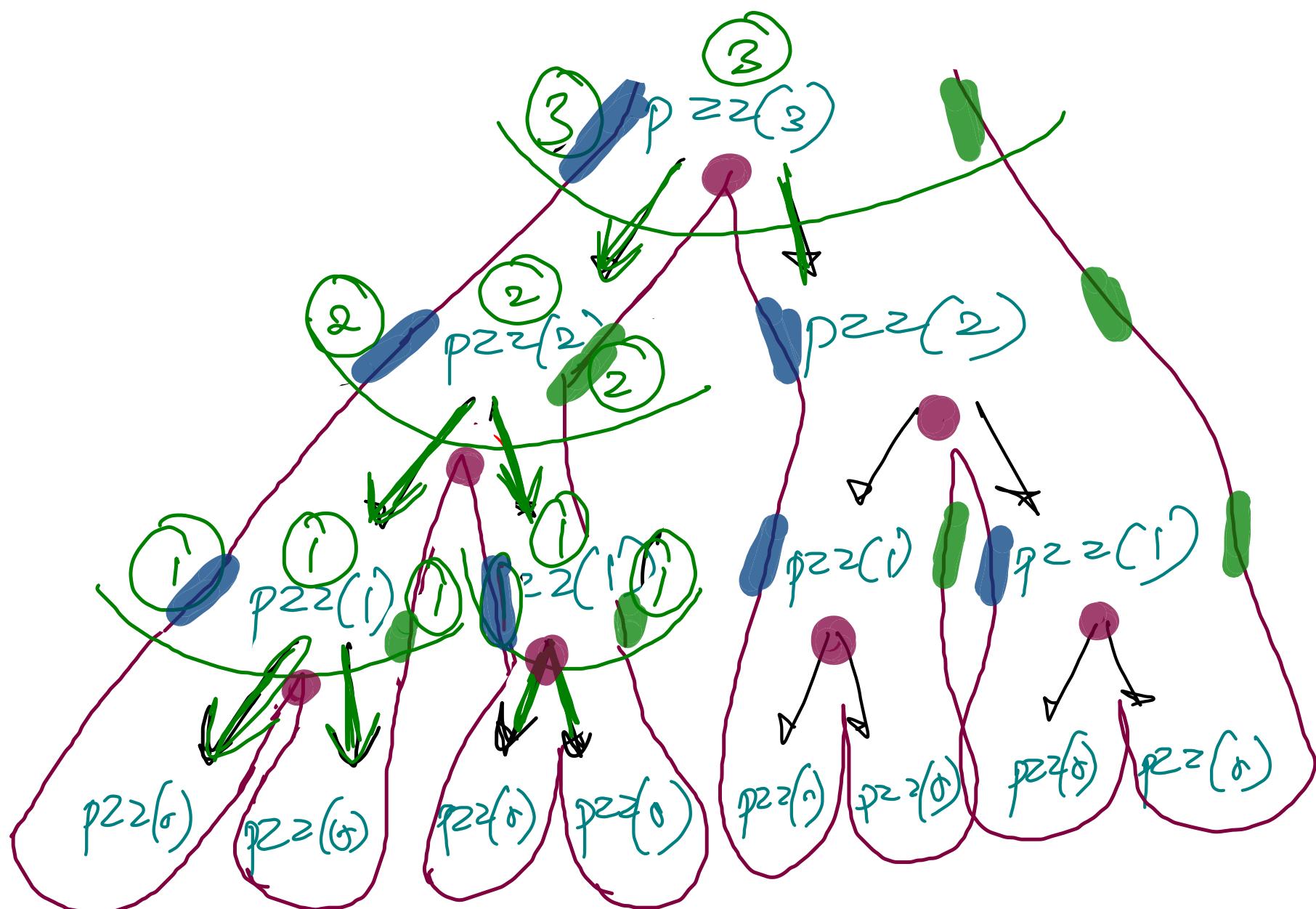
sys(n)

printZigzag(1)
1 1

printZigzag(0)
return ;

```
public static void pzz(int n){  
    if(n == 0) return;  
  
    ① System.out.print(n + " ");  
    ② pzz(n - 1);  
    ③ System.out.print(n + " ");  
    ④ pzz(n - 1);  
    ⑤ System.out.print(n + " ");  
}
```





```

public static void pzz(int n){
    if(n == 0) return;
    Preorder: System.out.print(n + " ");
    Left: pzz(n - 1);
    Middle: System.out.print(n + " ");
    Right: pzz(n - 1);
    Postorder: System.out.print(n + " ");
}

```

Depth first Search

↓
Recursive tree/
euler tree

Recursion

c ↗
(calls)^{height} + {preorder + inorder + postorder} * height

eg Point zigzag

$$2^n + \{k+k+k\} * n = 2^n + 3kn$$

$\boxed{O(2^n)}$

Recursion & Arrays

① Display Array

```
public static void displayArr(int[] arr, int idx){  
    ...  
}
```

③ Meeting Expectation

→ Preorder

sys(arr[idx])

display(arr, idx + 1)

① Expectation
{ 10, 20, 30, 40, 50 }

displayArr(arr, 0)

② Faith

display(arr, 1)

{ 20, 30, 40, 50 }

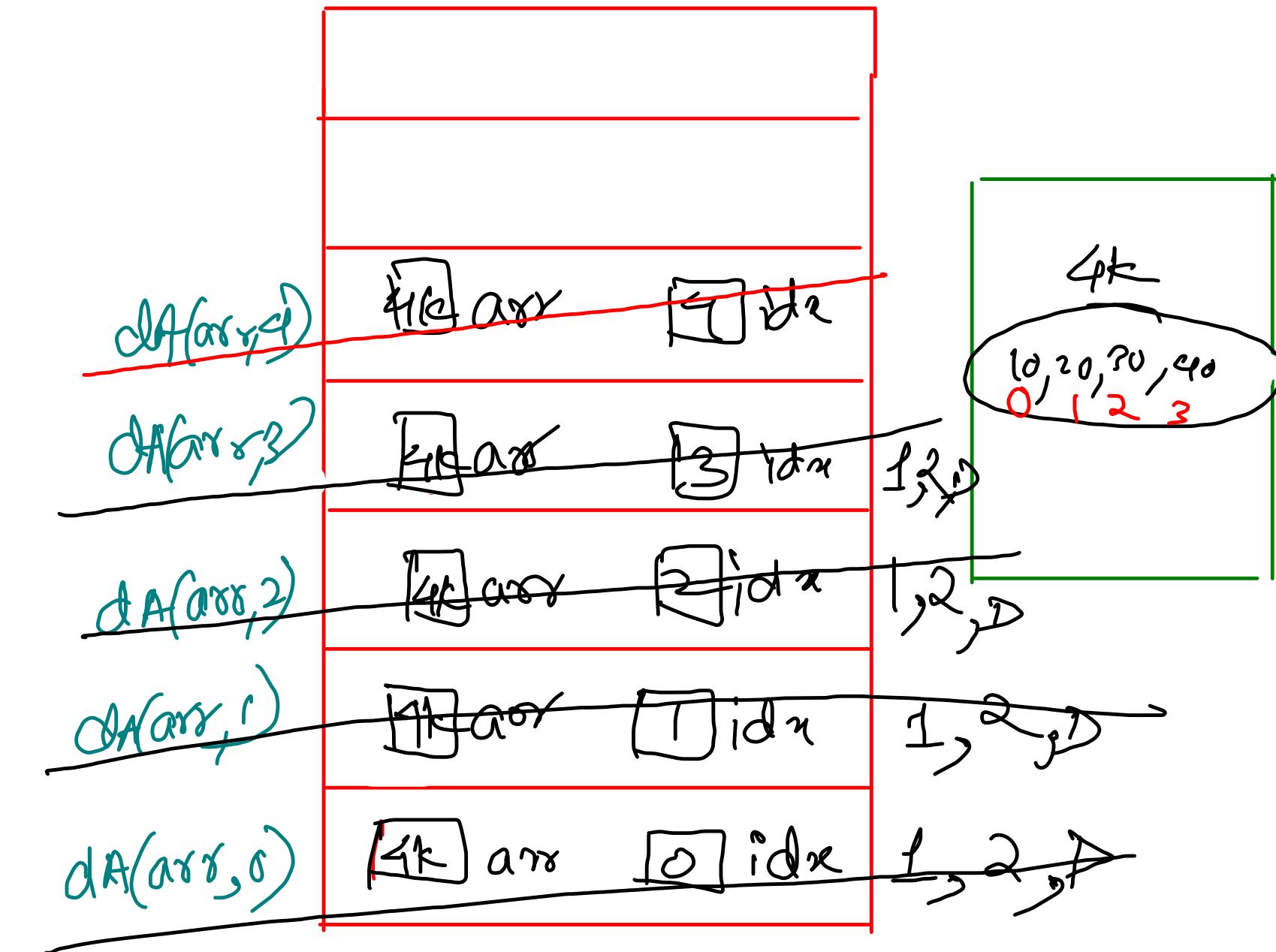
④ Base case
if(idx == arr.length)
 return

```

public static void displayArr(int[] arr, int idx){
    if(idx == arr.length) return;
    // preorder -> Meeting Expectation
    System.out.println(arr[idx]); {bordered}
    // faith
    displayArr(arr, idx + 1);
}

```

10, 20, 30, 40



Q) Maximum In an Array

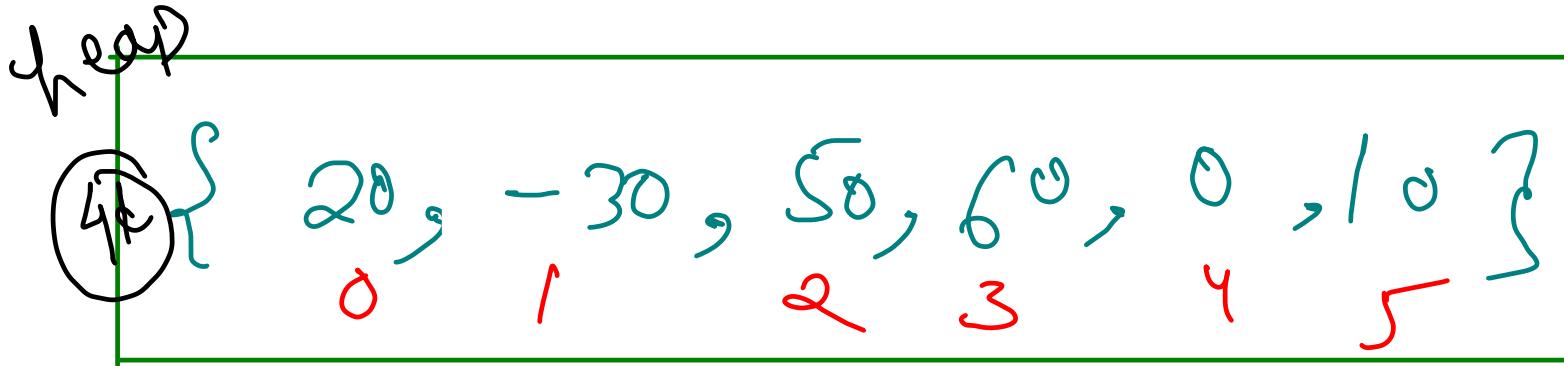
{ 20, -30, 50, 60, 0, 10, 40 }

① Expectation

int max of Array(arr, l) \Rightarrow max arr[0 : l-1] \Rightarrow 60

② Faith int max of Arr(arr, l) \Rightarrow max arr[1 : l-1]

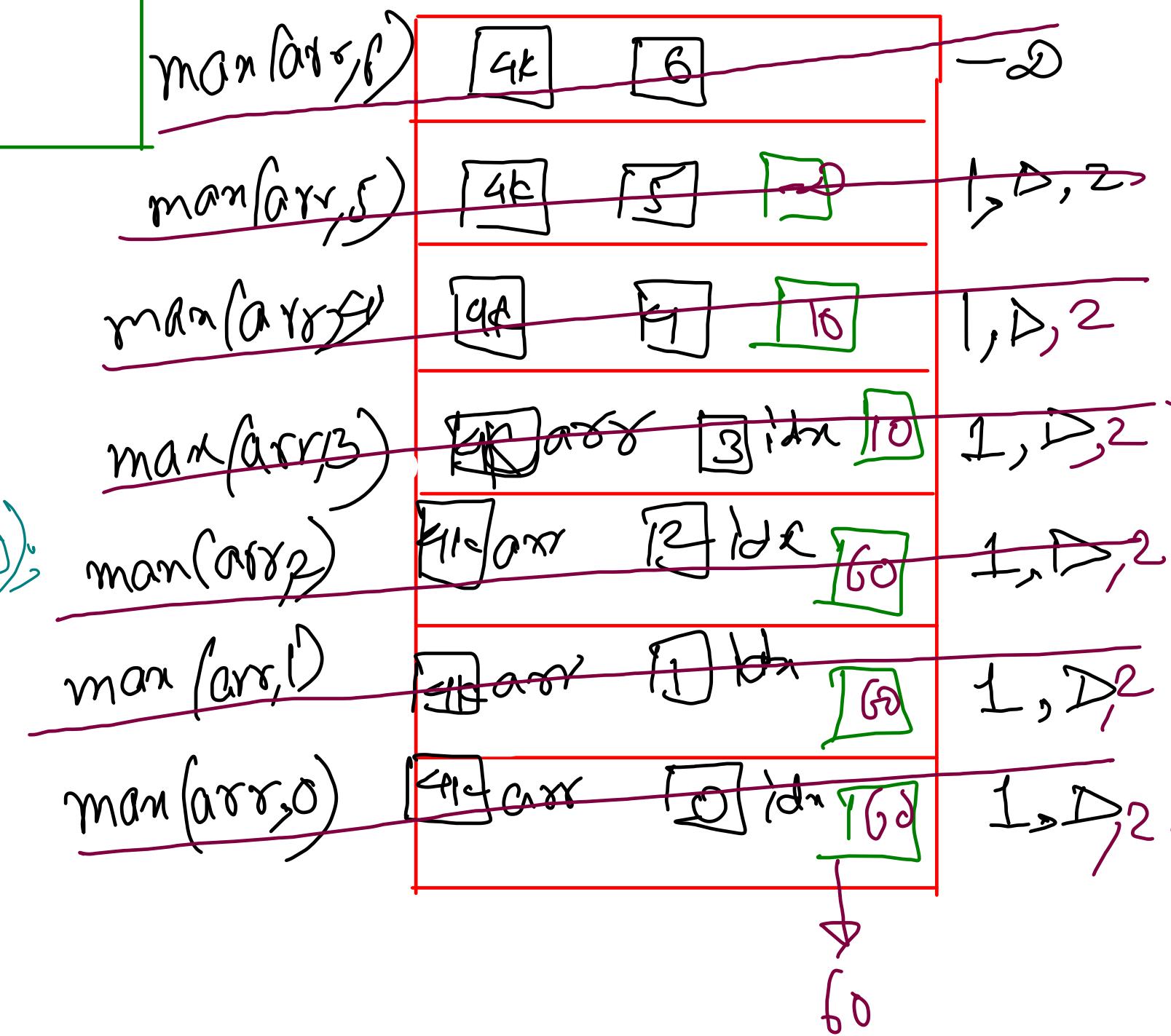
③ Meeting Expectation \Rightarrow return max(arr[i:l], maxall)



```

int maxOfArr(int [] arr, int idn) {
    ⑥ if( idn == arr.length ) return -∞;
    ① int maxTemp = max(arr, idn+1);
    ② return max(maxTemp, arr[idn]);
}

```



Addition

$$x + ? = x$$

Subtraction

$$x - ? = x$$

Multiplication

$$x * ? = x$$

Minimum

$$\min(x, +\infty) = x$$



Integer. MAX VALUE

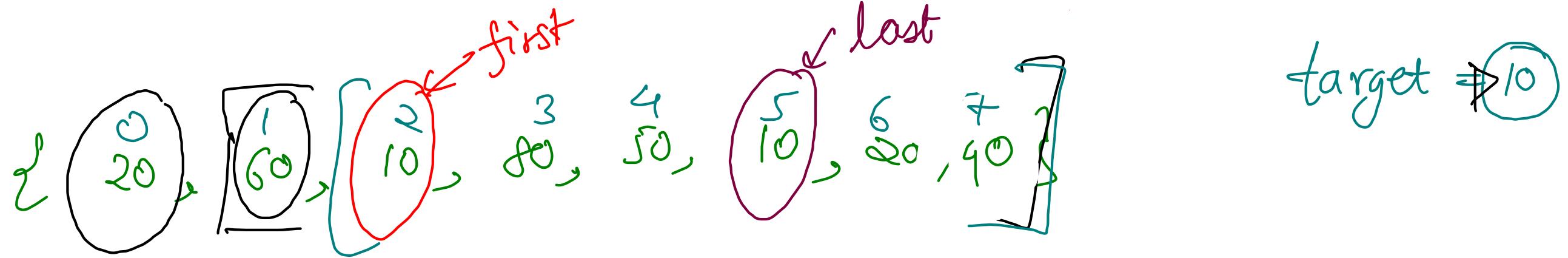
$$+2^{31} - 1$$

Maximum

$$\min(x, -\infty) = x$$

Integer. MIN VALUE

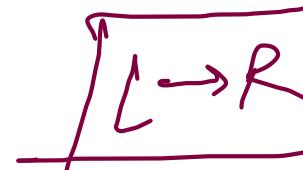
$$-2^{31}$$



First Index

$fI(\text{arr}, c) \rightarrow 2$

faith



temp = $fI(\text{arr}, \text{idn}+1)$

Meeting End

if ($\text{arr}[\text{idn}] == \text{target}$)
 T → return $\text{idn};$
 F → return temp;

lastIndex

$fI(\text{arr}, \text{a.l}-1)$

faith



temp = $fI(\text{arr}, \text{idn}-1)$

Meeting End

if ($\text{arr}[\text{idn}] == \text{target}$)
 T → return $\text{idn};$
 F → return temp;

```

public static int firstIndex(int[] arr, int idx, int x){
    if(idx == arr.length) return -1;

    ① if(arr[idx] == x){
        // meeting expectation
        return idx;
    } else {
        // faith
        return firstIndex(arr, idx + 1, x);
    }
}

```

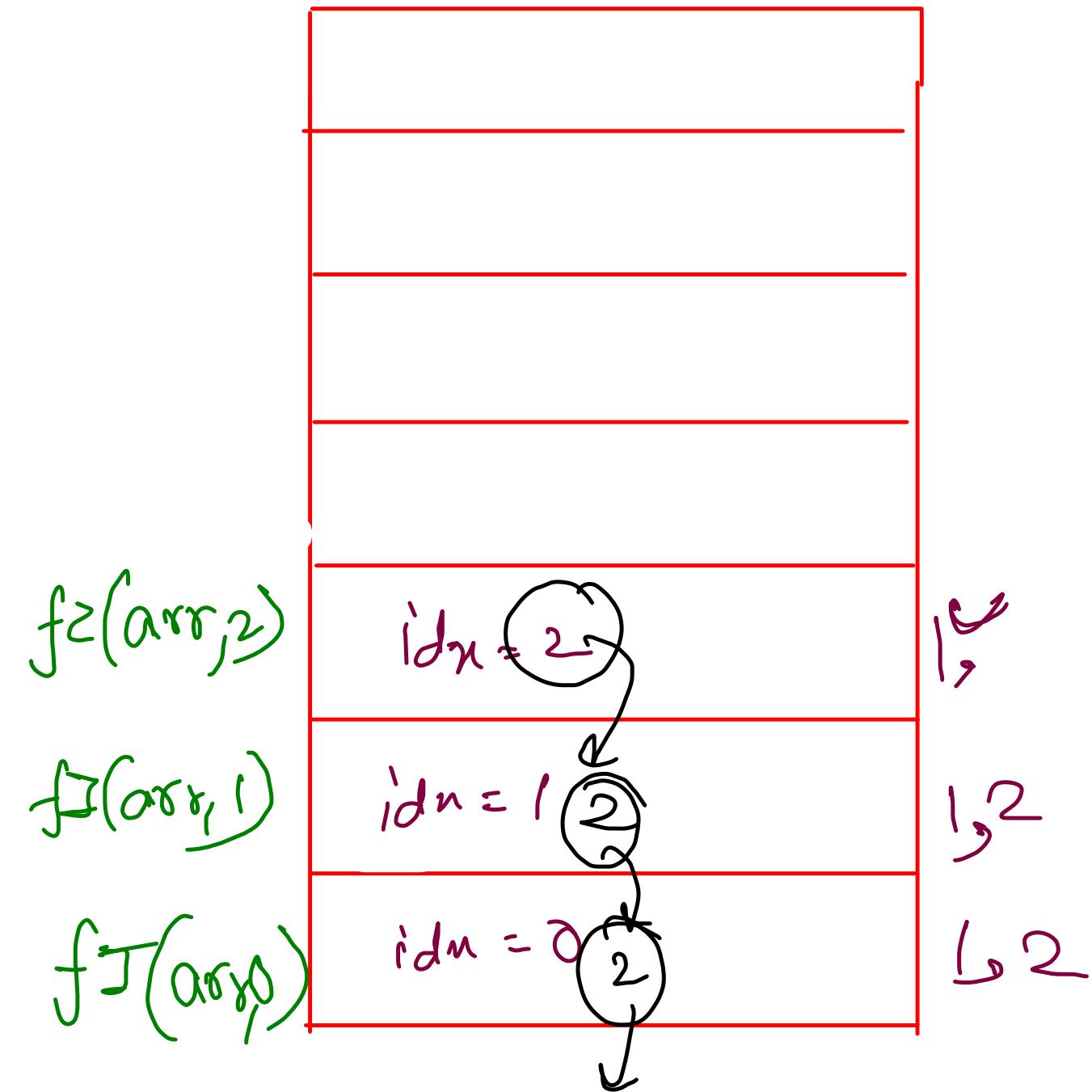
{ 20 50 10 80 50, 10, 20, 40 }

```

public static int lastIndex(int[] arr, int idx, int x){
    if(idx == -1) return -1; // search unsuccessful

    if(arr[idx] == x){
        return idx;
    } else {
        return lastIndex(arr, idx - 1, x);
    }
}

```



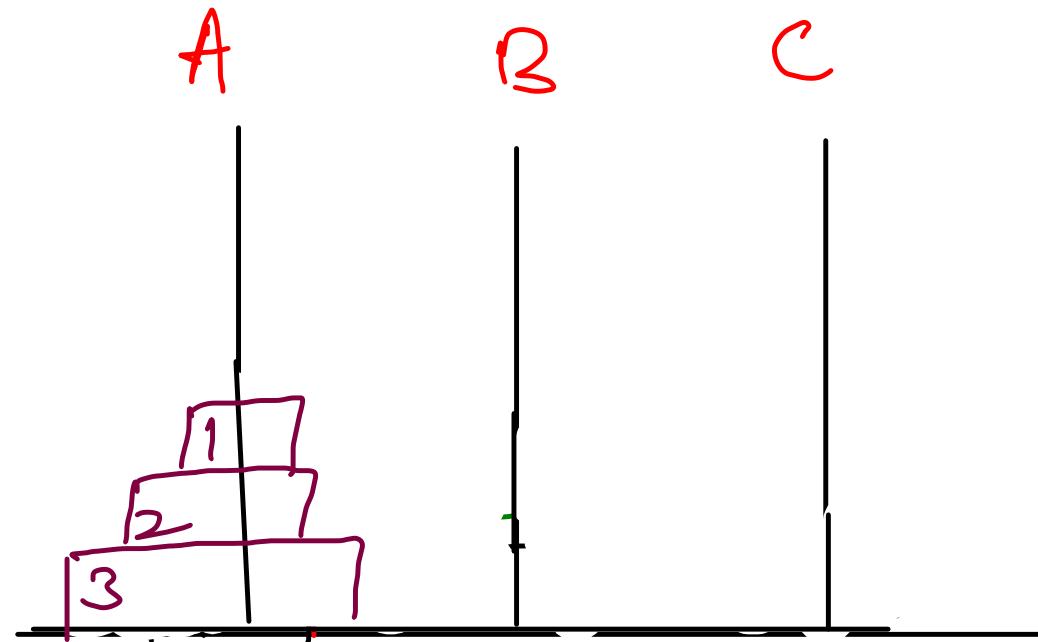
Today's Questions

- ① Tower of Hanoi^o
- ② All indices of Array

+ ~~Discussion~~ on DSA based projects

Tower of Hanoi

1. There are 3 towers. Tower 1 has n disks, where n is a positive number. Tower 2 and 3 are empty.
2. The disks are increasingly placed in terms of size such that the smallest disk is on top and largest disk is at bottom.
3. You are required to
 - 3.1. Print the instructions to move the disks.
 - 3.2. from tower 1 to tower 2 using tower 3
 - 3.3. following the rules
 - 3.3.1 move 1 disk at a time.
 - 3.3.2 never place a smaller disk under a larger disk.
 - 3.3.3 you can only move a disk at the top.



Constraints:

- ① only one disk can be moved at a time
- ② Larger disk cannot be placed on top of smaller disk

① Expectation

tower of hanoi(n , A, B, C)

(source) form
to
using



Instructions
to move 3 disk
from A to B
using C

② Faith

(2.1) tower of hanoi ($n-1$, A, C, B) \rightarrow Move $(n-1)$ disks
from A to C
using B

Meeting Expectation

(2.2) Move n^{th} disk from A to B

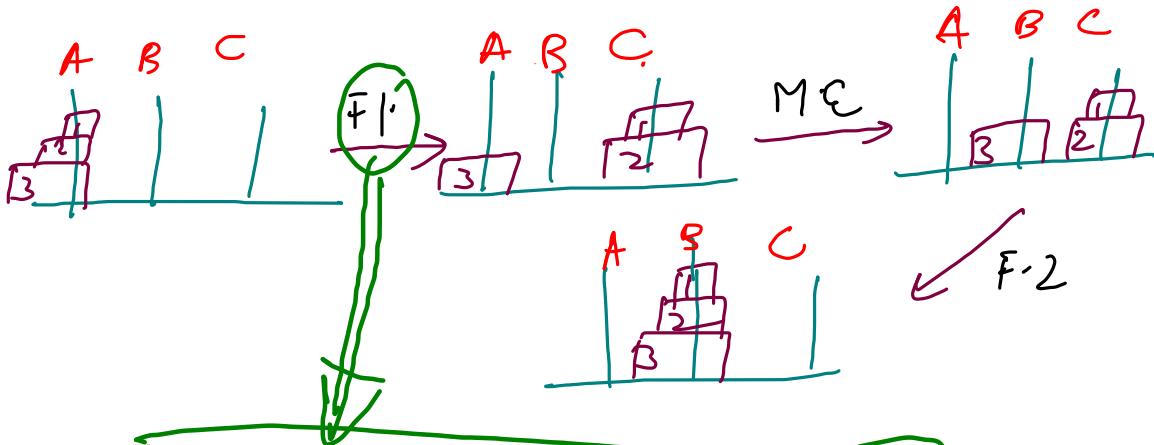
(2.3) tower of hanoi $n-1$, C, B, A) \rightarrow

move $(n-1)$ disk
from C to B
using A.

```

public static void toh(int n, int srcTower, int destTower, int auxTower){
    if(n == 0) return;
    ① toh(n - 1, srcTower, auxTower, destTower);
    // Move N - 1 disks from source to auxiliary
    ② System.out.println(n + "[" + srcTower + "-">+ destTower + "]");
    // Move nth disk from source to destination
    ③ toh(n - 1, auxTower, destTower, srcTower);
    // Move N - 1 disks from auxiliary to destination
}

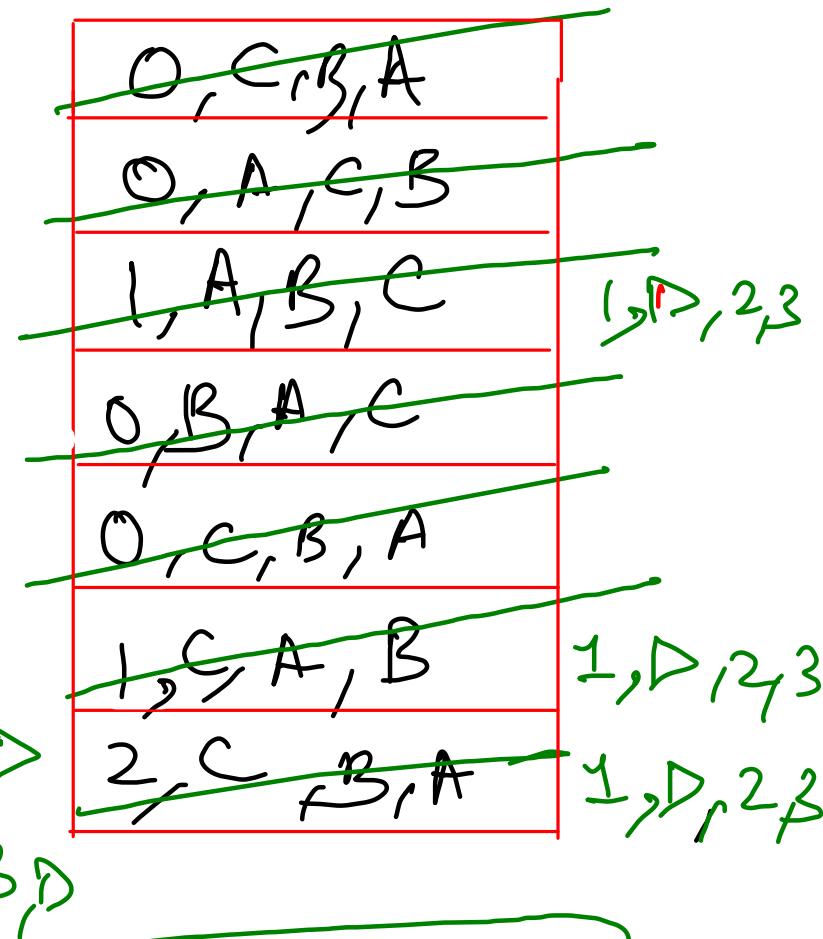
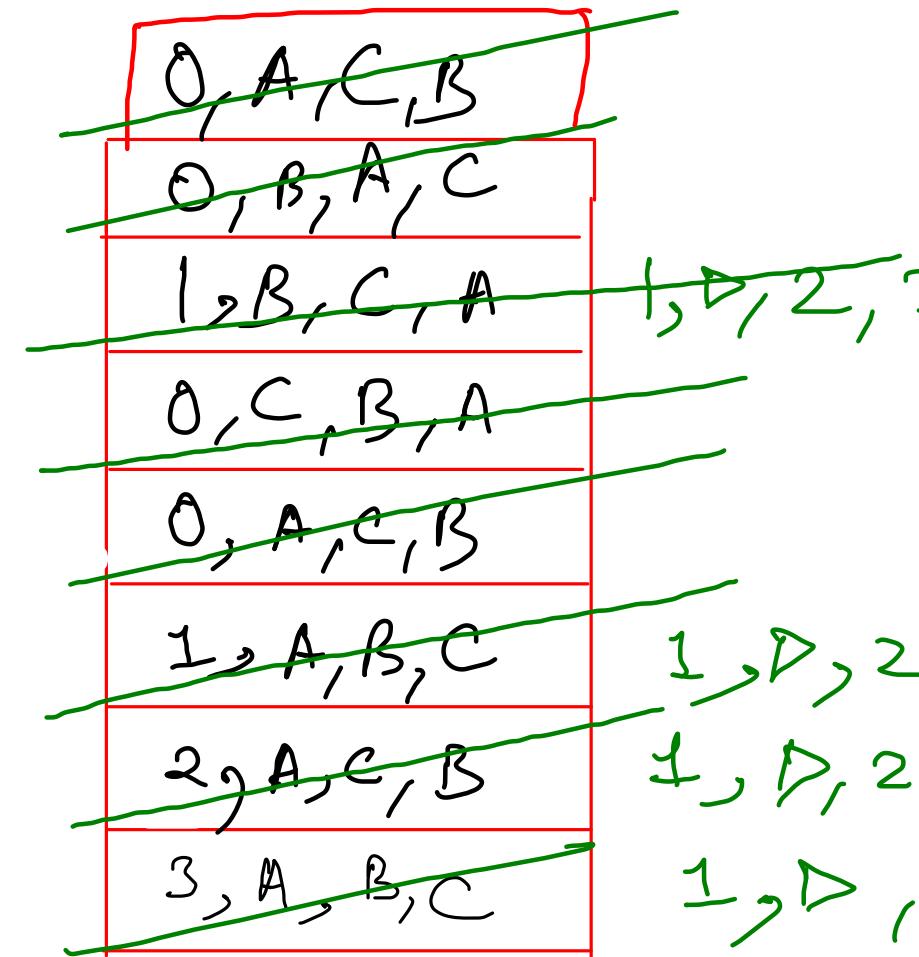
```



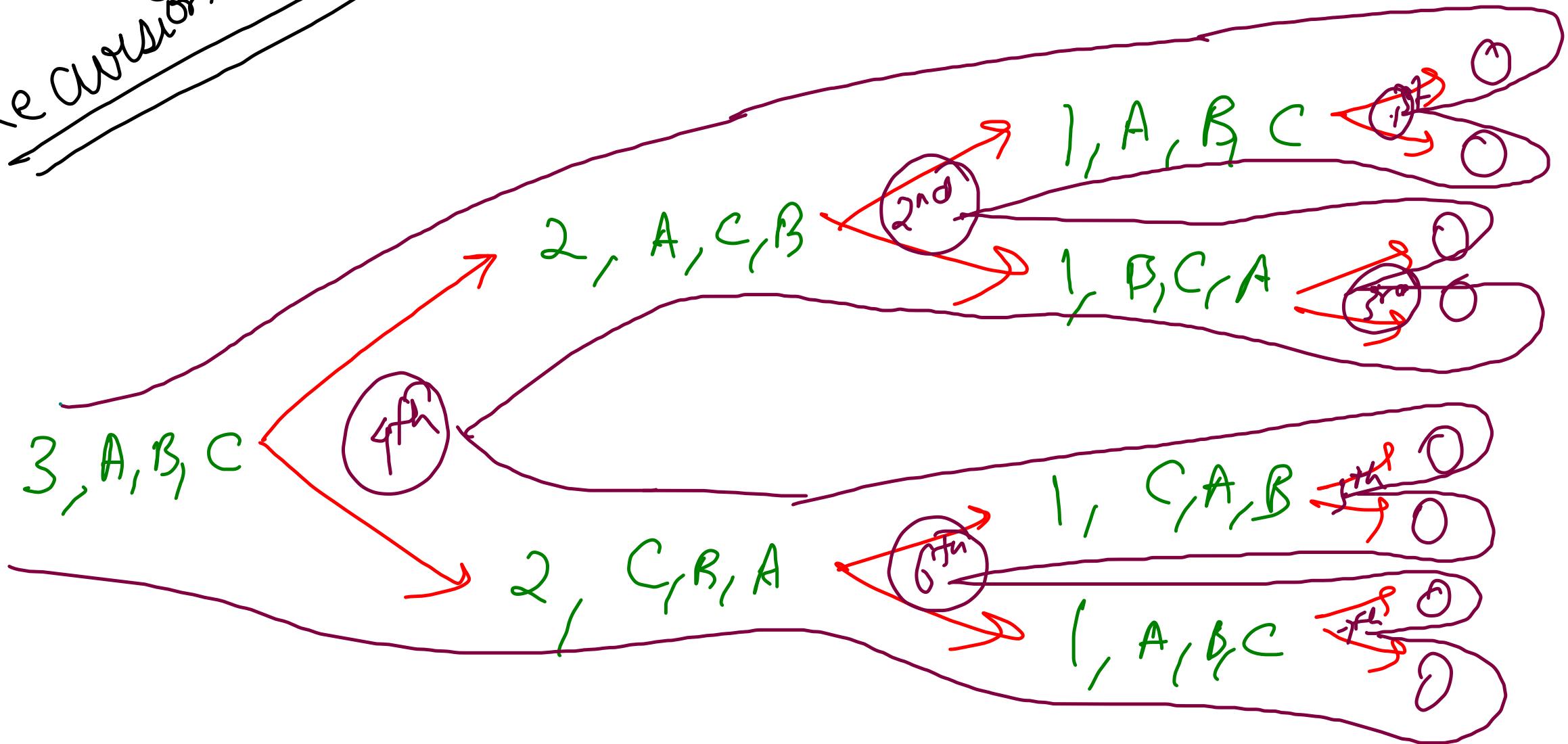
- ① 1st from A to B
- ② 2nd from A to C
- ③ 1st from B to C

- ④ 3rd from A to B

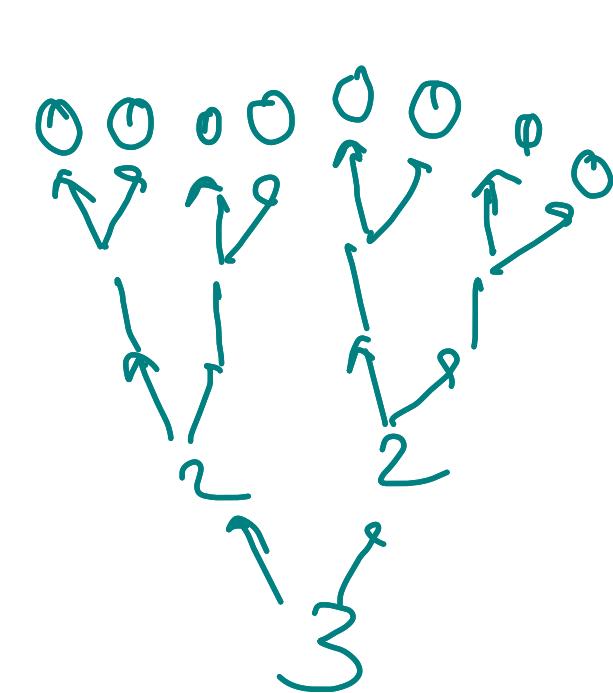
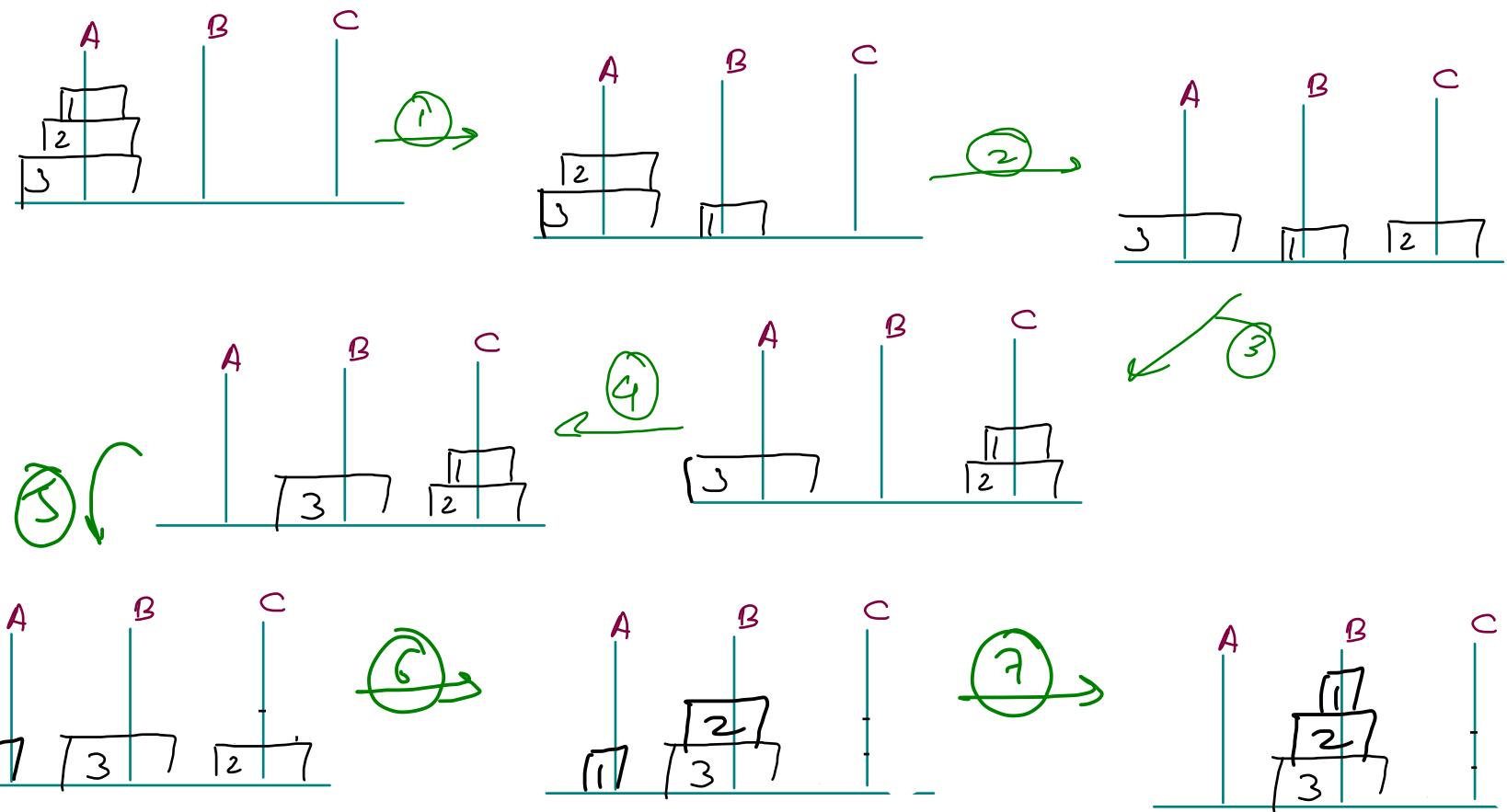
- ⑤ 1st from C to A
- ⑥ 2nd from C to B
- ⑦ 1st from A to B



Recurision Tree



- 2 discs from A to C using B
- ① 1st from A to B
 ② 2nd from A to C
 ③ 1st from B to C
- Meeting expectation
- ④ 3rd from A to B
- 2 discs from C to B
- ⑤ 1st from C to A
 ⑥ 2nd from C to B
 ⑦ 1st from A to B



DSA-based Projects {Unique Project Ideas}

MERN

- ① Writing own library/header file for
 - B Tree, B+ Tree & SQL + DSA}
 - Segment Tree, Fenwick Tree, Number Theory Algorithms,
DSU, Hashmaps, Priority Queue/Heaps, Graph Algorithms
- {Competitive (Advanced DSA) + Generic Programming
Templates + DSFS}

②

Games

~~Webdev~~
→ CSS
→ JS

→ {+ 2 players}

- ~~Canvas~~ → Snake & ladder → Graphs Algo {DFS}
- Sudoku → Backtracking {+ levels/themes}
- Crossword Puzzle → Backtracking {+ levels/themes}
- Chess → 2D matrix {+ AI → Bot}
- Jump Game → Geriody + Dynamic Programming
{Frogs/Cartoon} {Climb Stairs}
- Mario Game

③ WhatsApp/Telegram Clone

- LRU Cache Algo \Rightarrow Ranking/ ordering of chats
 - Trie Data Structure \Rightarrow Searching Contacts/ Messages
prefix tree
 - CRUD operations { Firebase database }
 - Chatbots { APIs }
 - Audios/ Video calls
 - 1 to 1
 - To many \oplus
- } Networking

④ Splitwise App Clone

→ GFG

nosql
↑
restore
↑
↑

↳ Minimizing Cash flow Algorithm

↳ CRUD applications { persistent data (database),
+ authentication }

↳ System Design

↳ Priority Queue, Hashmap → Multiset Data
Structures

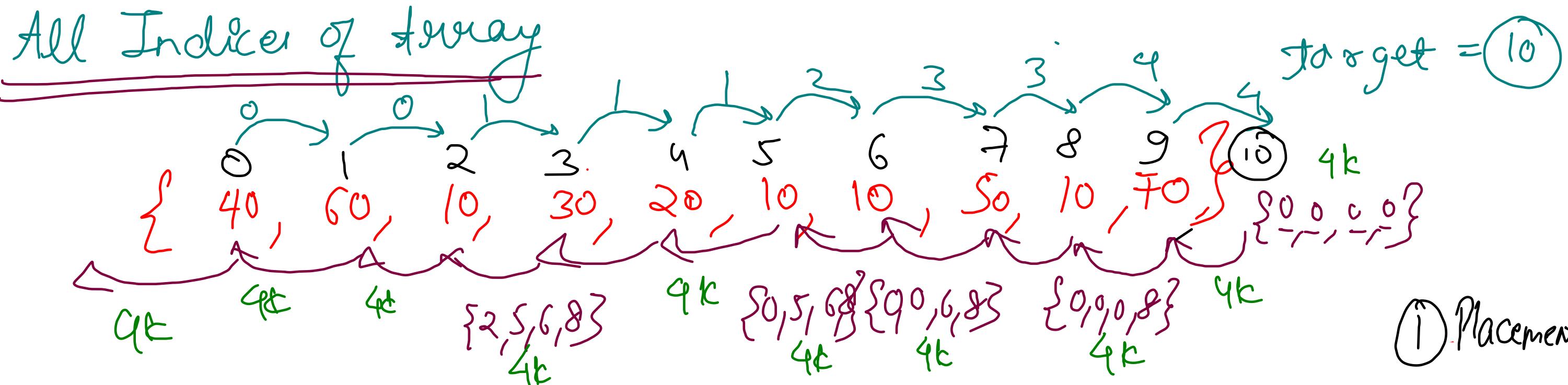
My Project :→ Flutter ⇒ LEN DEN

More Projects

{  } → Sorting visualizers
Path finding Algo visualizers

{  } → Text Editors : { Stack Data Structure }

{  } → File Zipping / Huffman Encoding
(TinyURL) URL Shortener



```

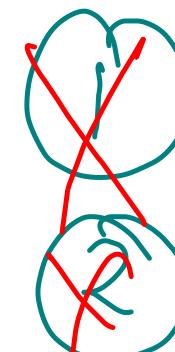
public static int[] allIndices(int[] arr, int target, int idx, int count) {
    if(idx == arr.length){
        int[] base = new int[count];
        return base;
    }
    if(arr[idx] == target){
        int[] res = allIndices(arr, target, idx + 1, count + 1);
        res[count] = idx;
        return res;
    } else {
        int[] res = allIndices(arr, target, idx + 1, count);
        return res;
    }
}

```

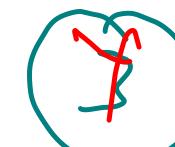
① Total occ
size array
return

② Occurrences
from left
part

Codechef & codeforces



Binary Tree



BST



Trie



Number Theory



Bit



DP

Graph



DSU

ST, f7

Leetcode Virtual Contest



BT



SDO



A&S



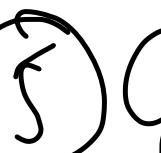
BST



SES



DP



Graph



R&B