

MULTIPARTY COMPUTATION USING SHAMIR'S SECRET SHARING



What does it even mean?

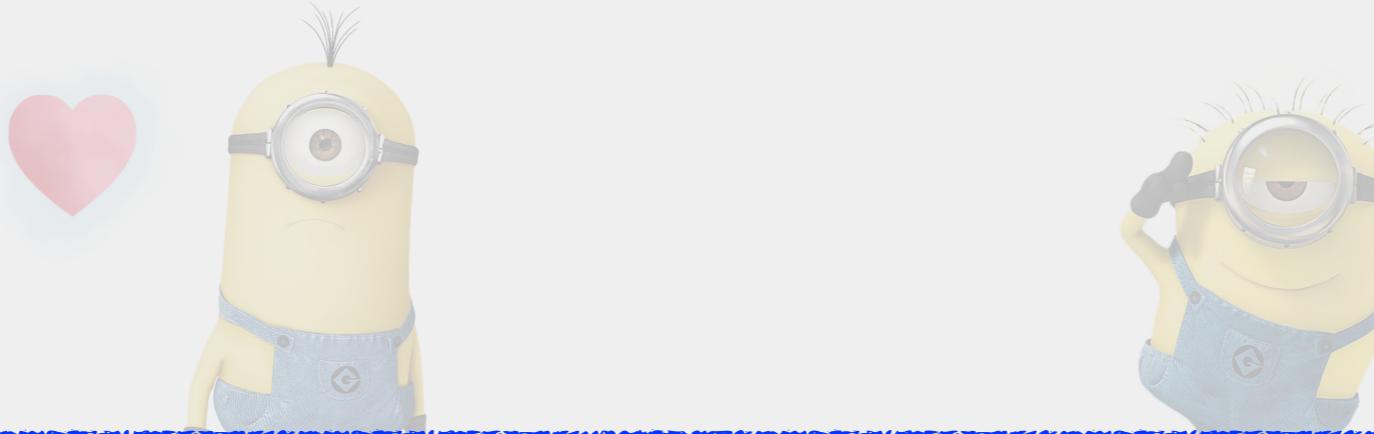
Evil Minions





Good Minions





What we want?

At least one good minion
participates in the division
process





Property of the locks :

Any 3 out of 5 keys can open
the cage ;-)



There's always a good guy in
the opening process!!

We control the minimum
number of people required
to open the cage!!



Given n parties

t are corrupted

$n-t$ are good

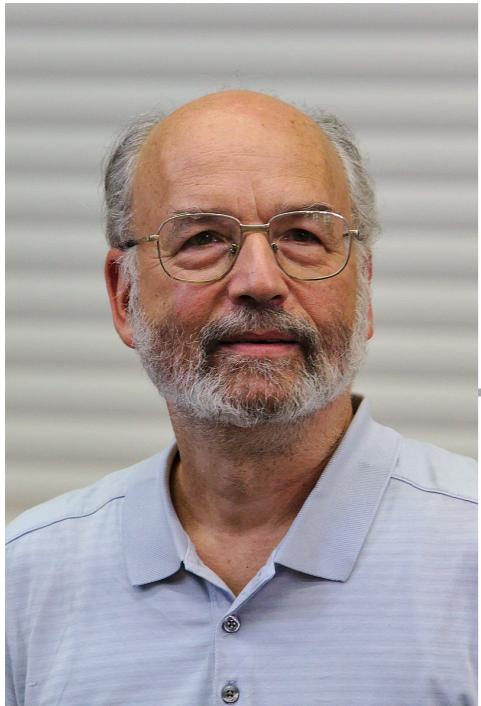
All interested in some computation such that

At least $t+1$ needed to do the computation

Minions
(good & bad)

Learn the code of lock

MULTIPARTY COMPUTATION USING SHAMIR'S SECRET SHARING



Private shares
known just to them

Shamir's Secret Sharing

POLYNOMIALS

POLYNOMIALS

degree

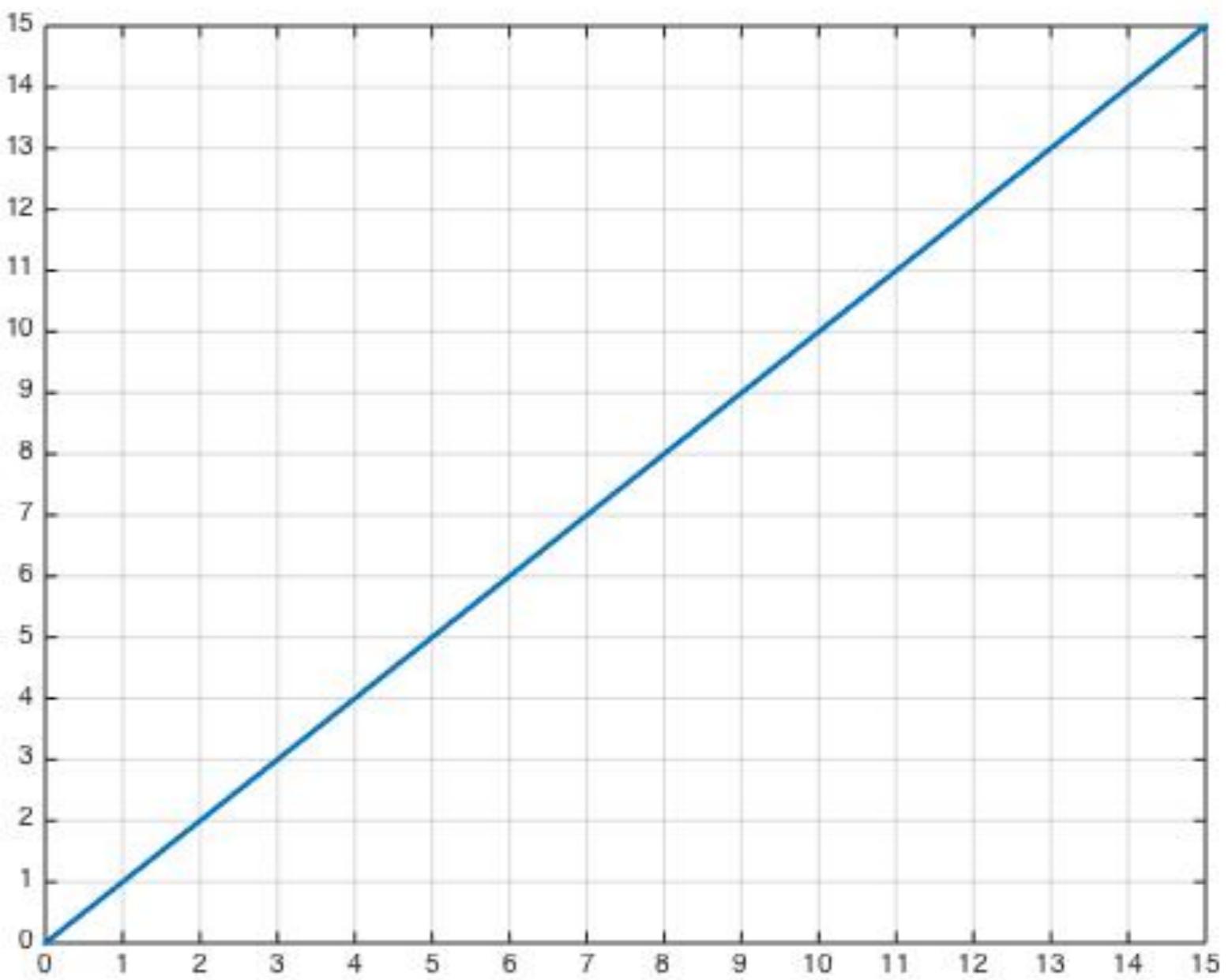
$$a_d x^d + a_{d-1} x^{d-1} + \dots + a_1 x + a_0$$

- ▶ $x - 3$ → 1 $a_1 x + a_0$
- ▶ $2x - 1.3$ → 1 $a_2 x^2 + a_1 x + a_0$
- ▶ $5x^2 - 0.5x + 1$ → 2 $a_2 x^2 + a_1 x + a_0$
- ▶ $10x^3 + 9x^2 + \overline{3}$ → 3 $a_3 x^3 + a_2 x^2 + a_1 x + a_0$
- ▶ 10 → 0 a_0

Degree 1 polynomials

12

$$f(x) = x$$

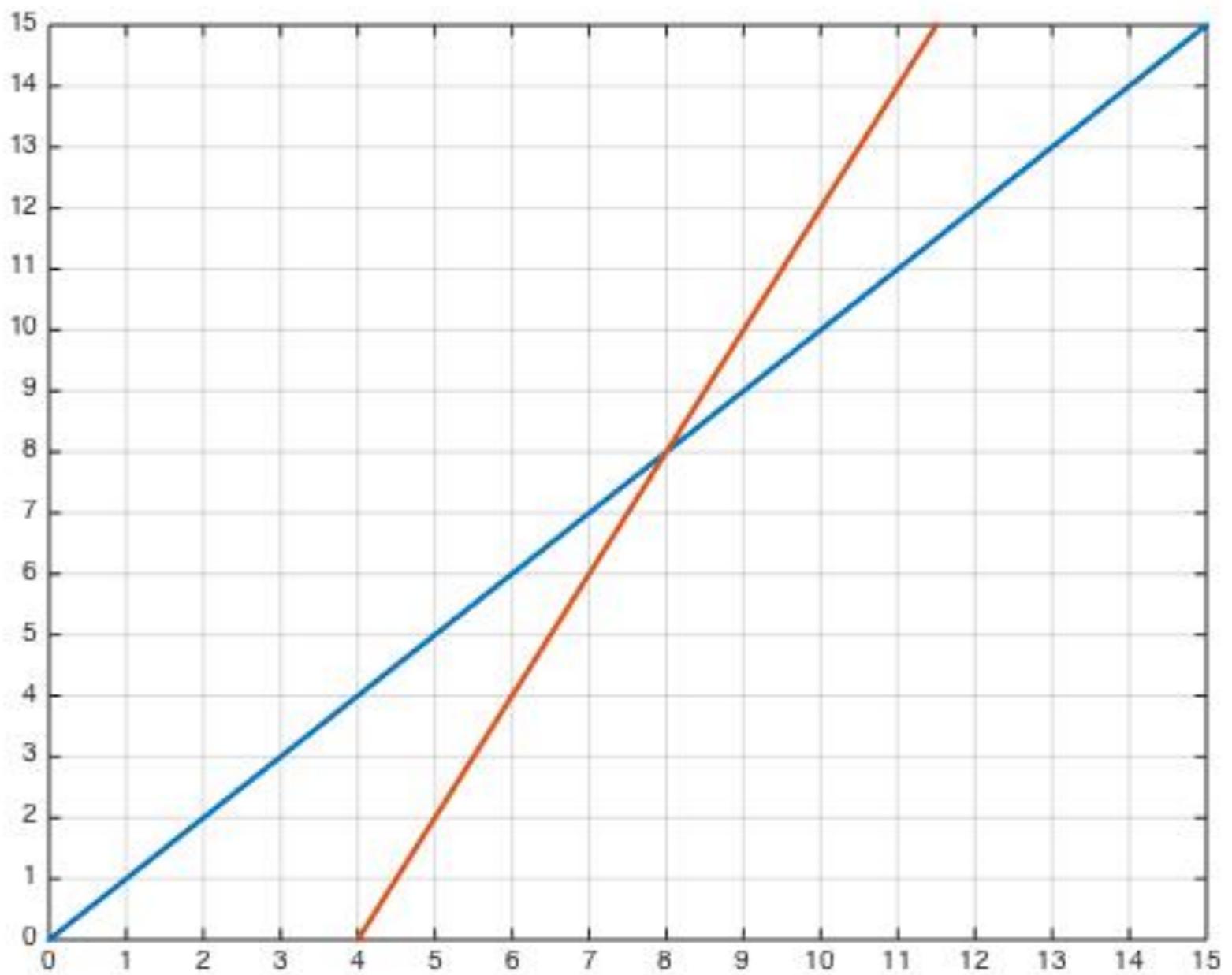


Degree 1 polynomials

13

$$f(x) = x$$

$$f(x) = 2x - 8$$



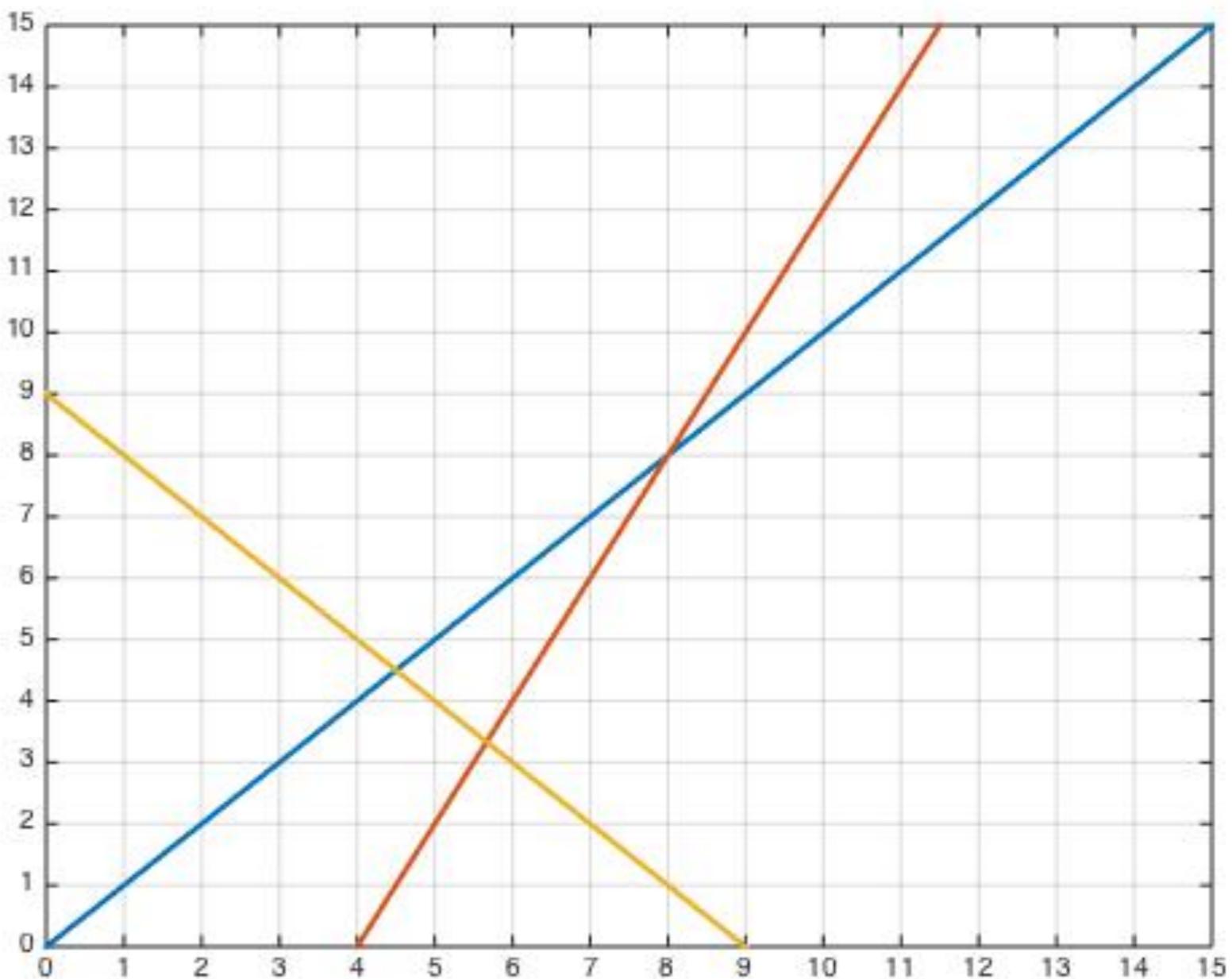
Degree 1 polynomials

14

$$f(x) = x$$

$$f(x) = 2x - 8$$

$$f(x) = -x + 9$$



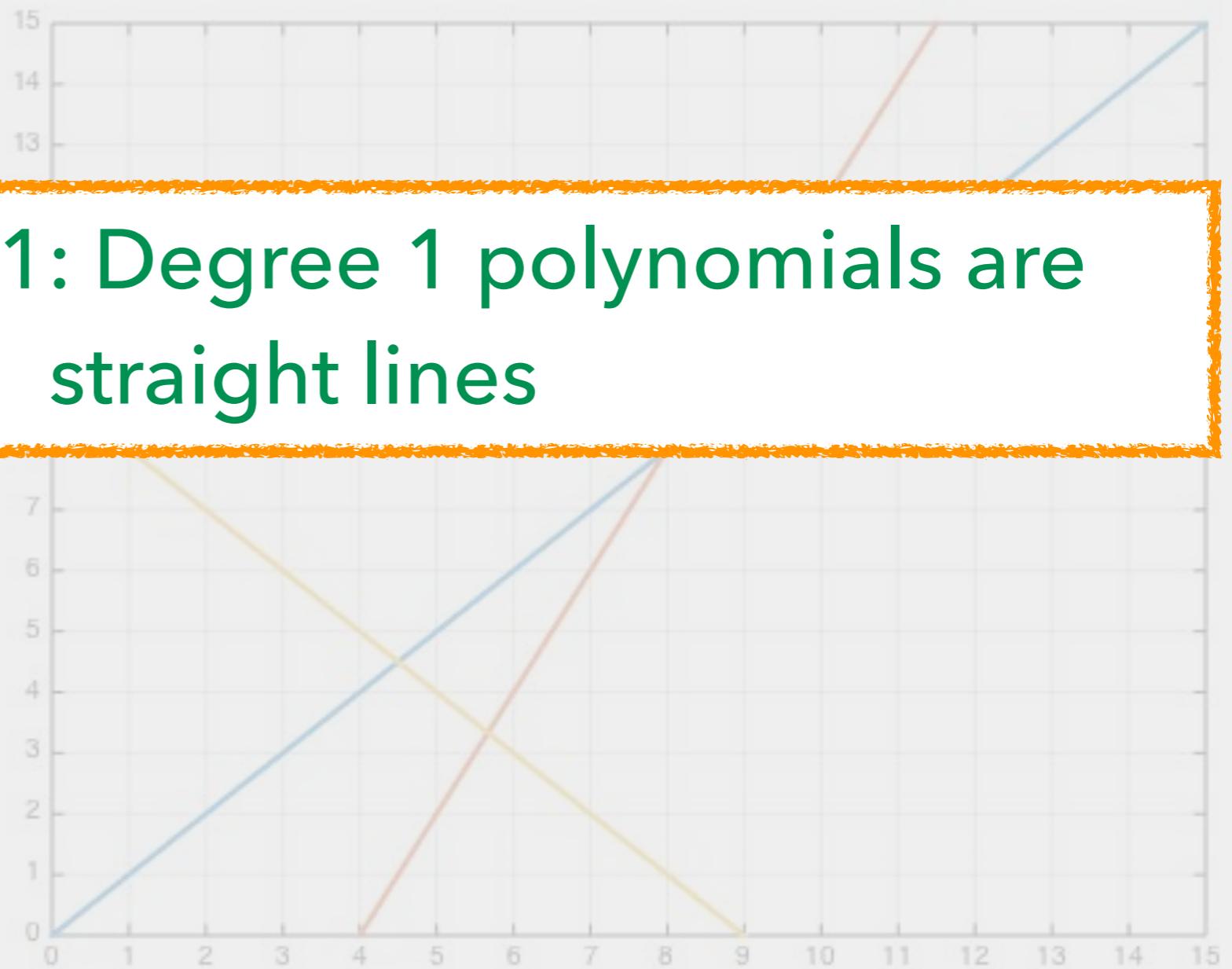
Degree 1 polynomials

15

$$f(x) = x$$

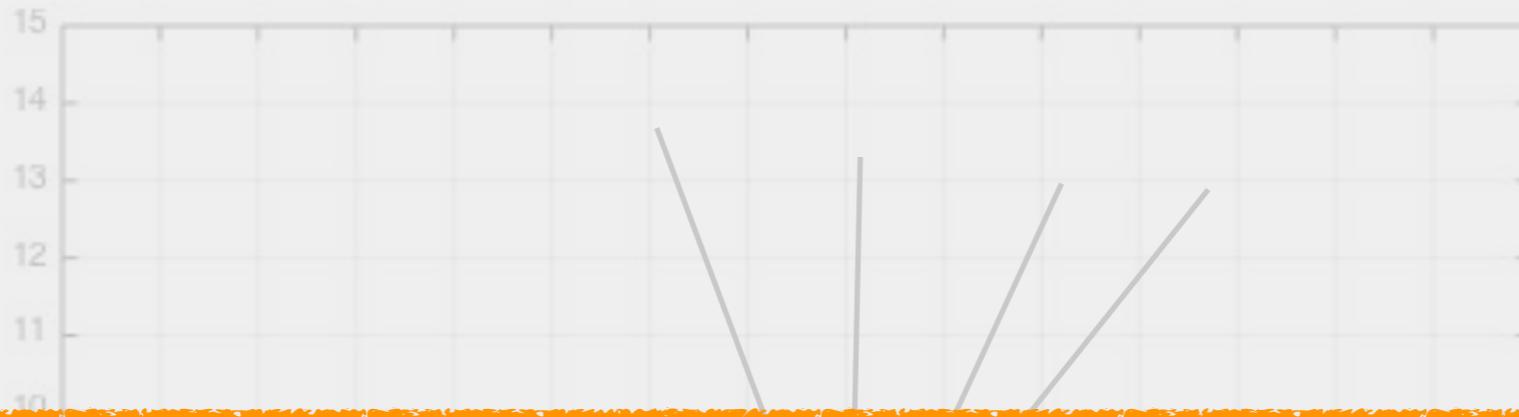
$$f(x) = 2x + 8$$

Observation 1: Degree 1 polynomials are straight lines

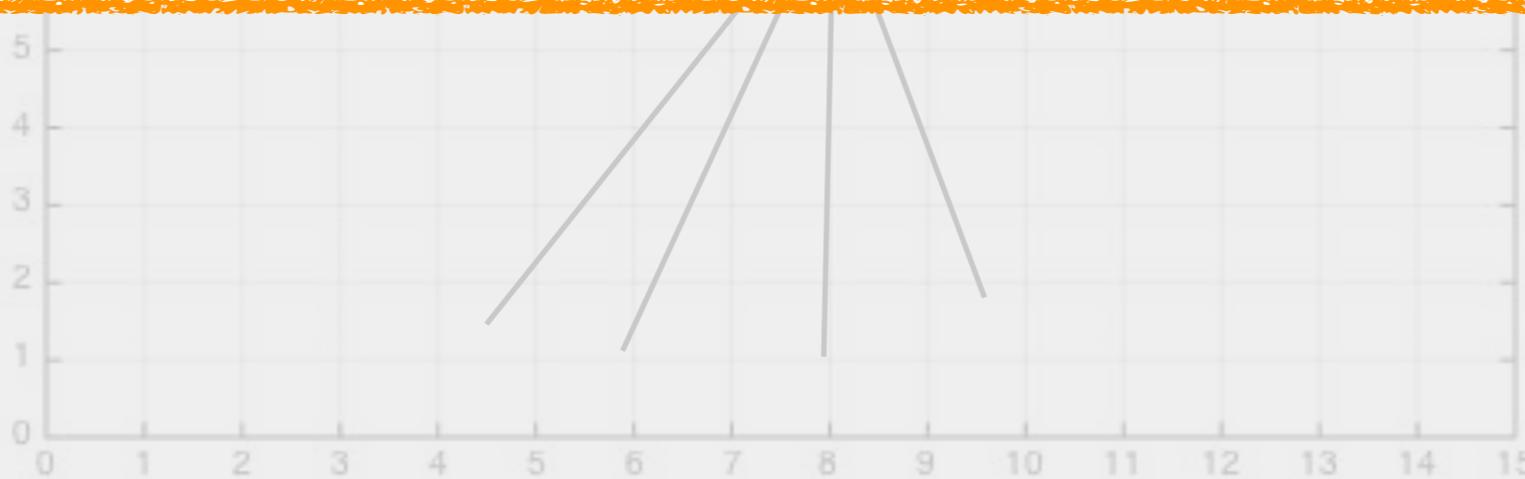


Degree 1 polynomials

16



Observation 2: Multiple degree 1 polynomials pass through a single point

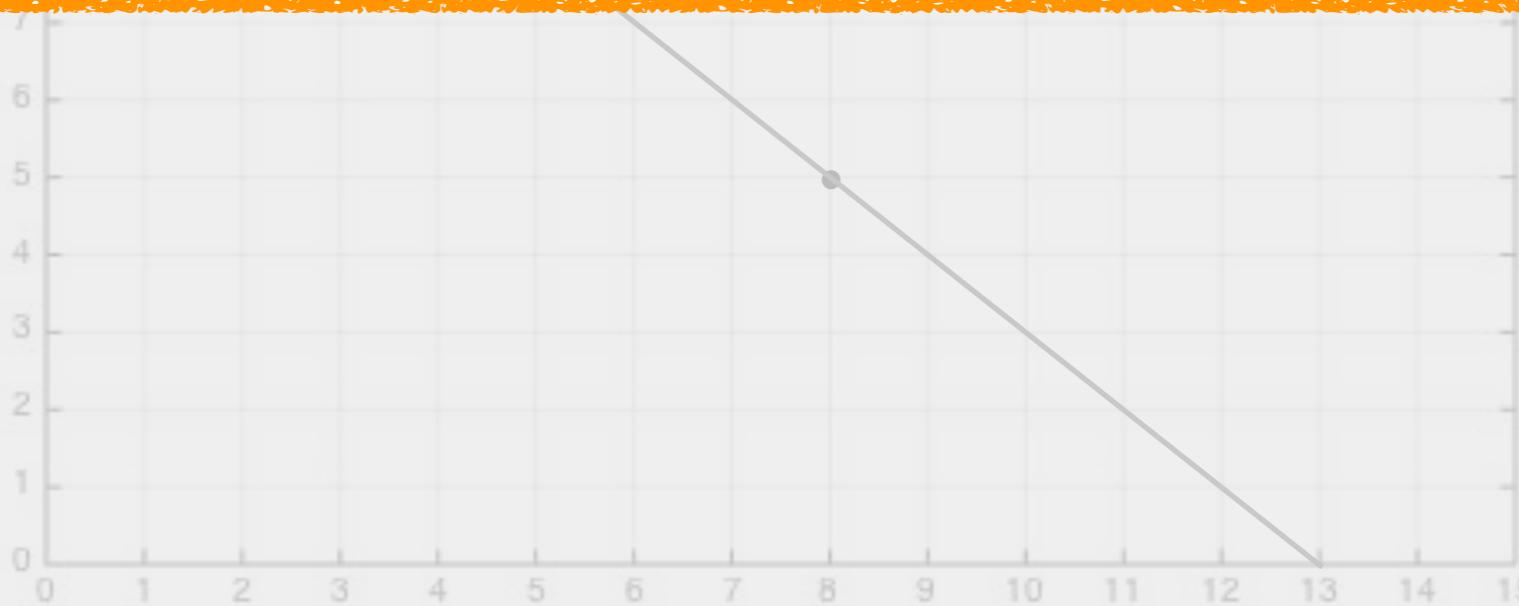


Degree 1 polynomials

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**Observation 3: Unique degree 1 polynomial
passes through 2 points**



Degree 2 polynomials

18

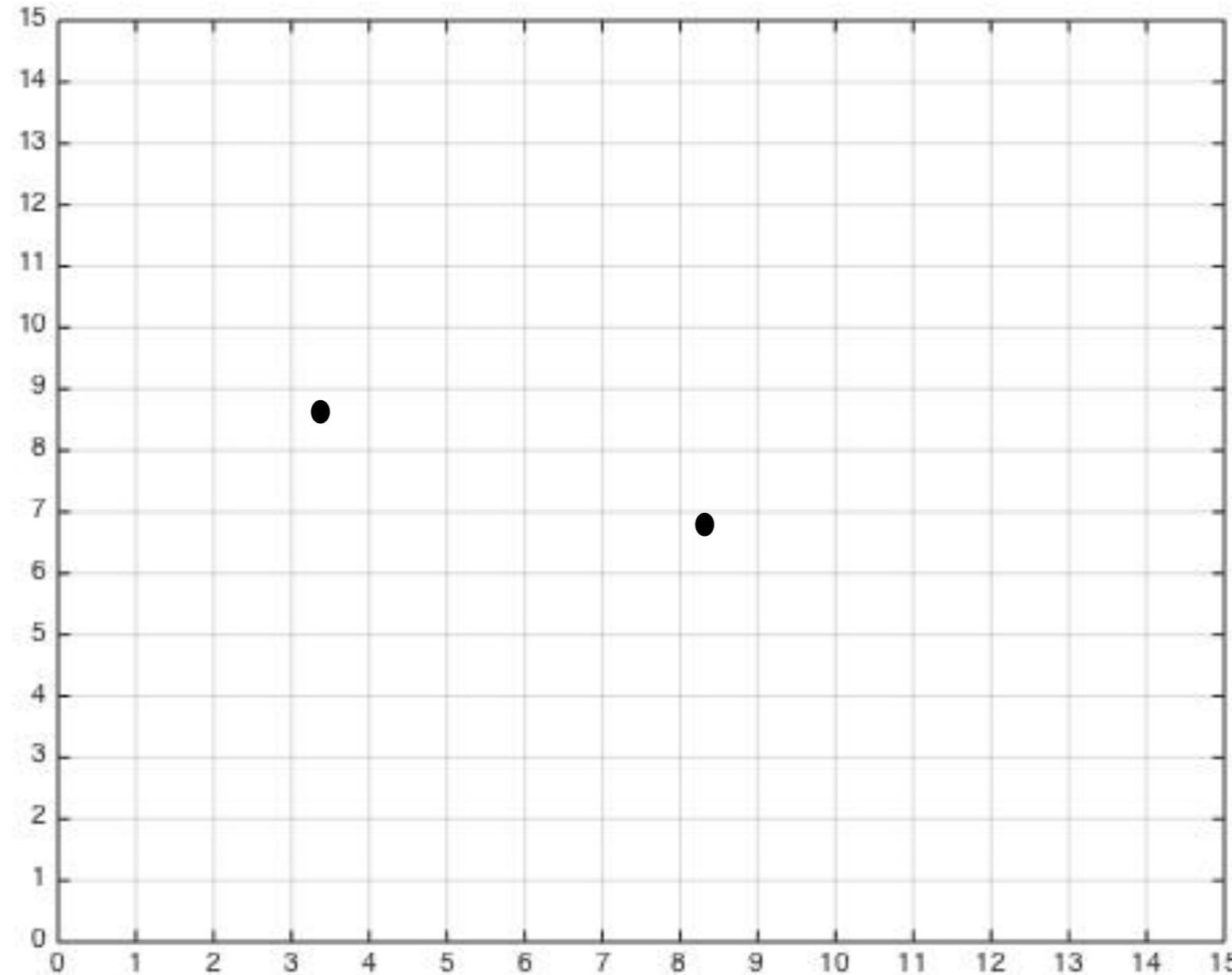
$$f(x) = x^2 - 6x + 12$$



Observation 1: Degree 2 polynomials are curved lines (parabolas)

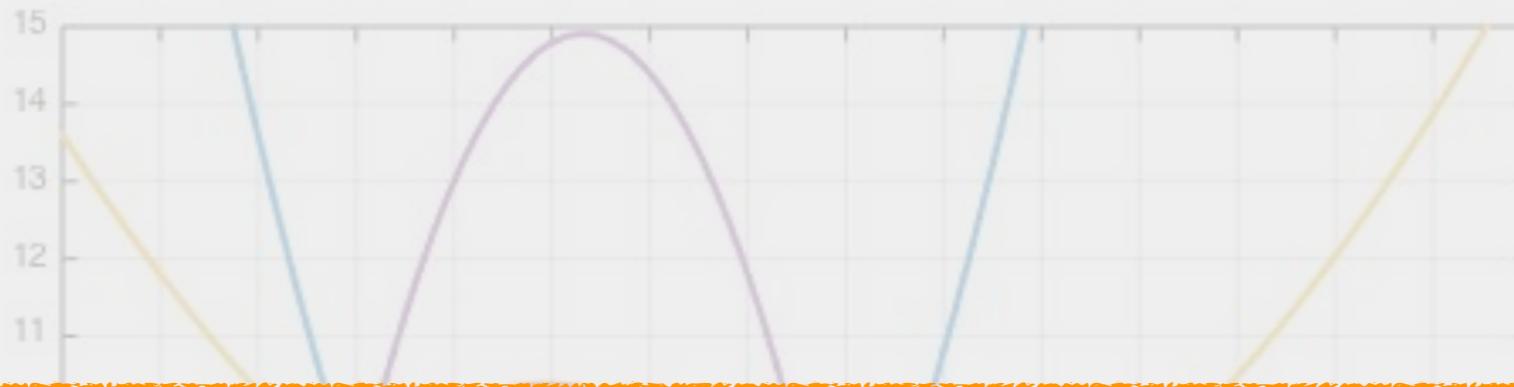


Degree 2 polynomials

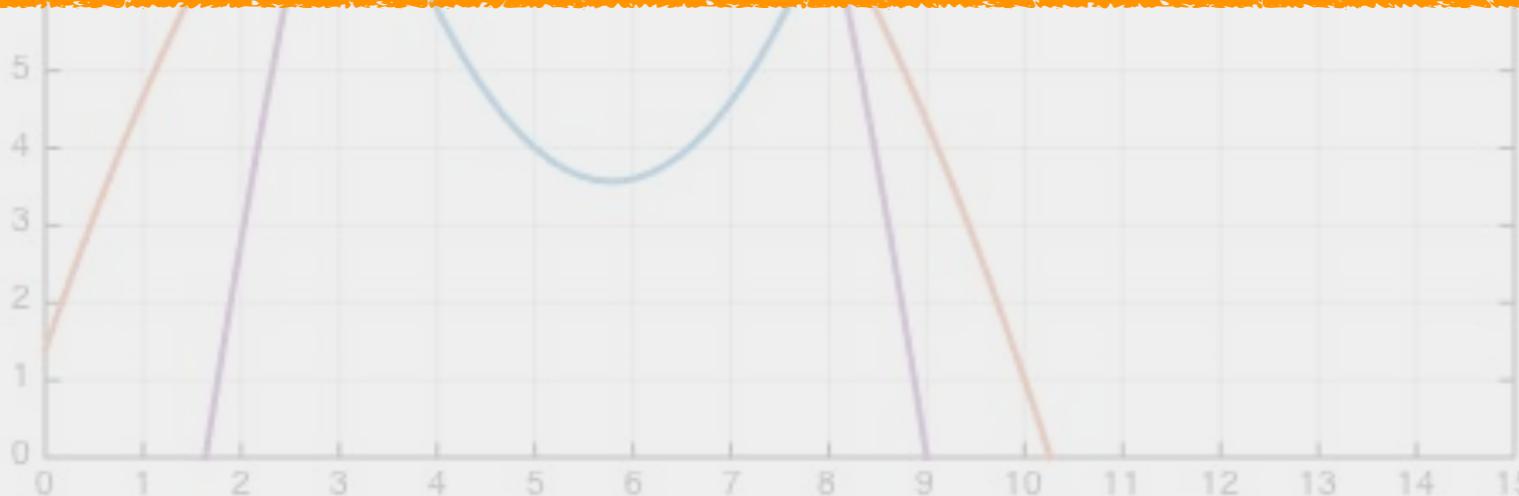


Degree 2 polynomials

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Observation 2: Multiple degree 2 polynomials pass through two points



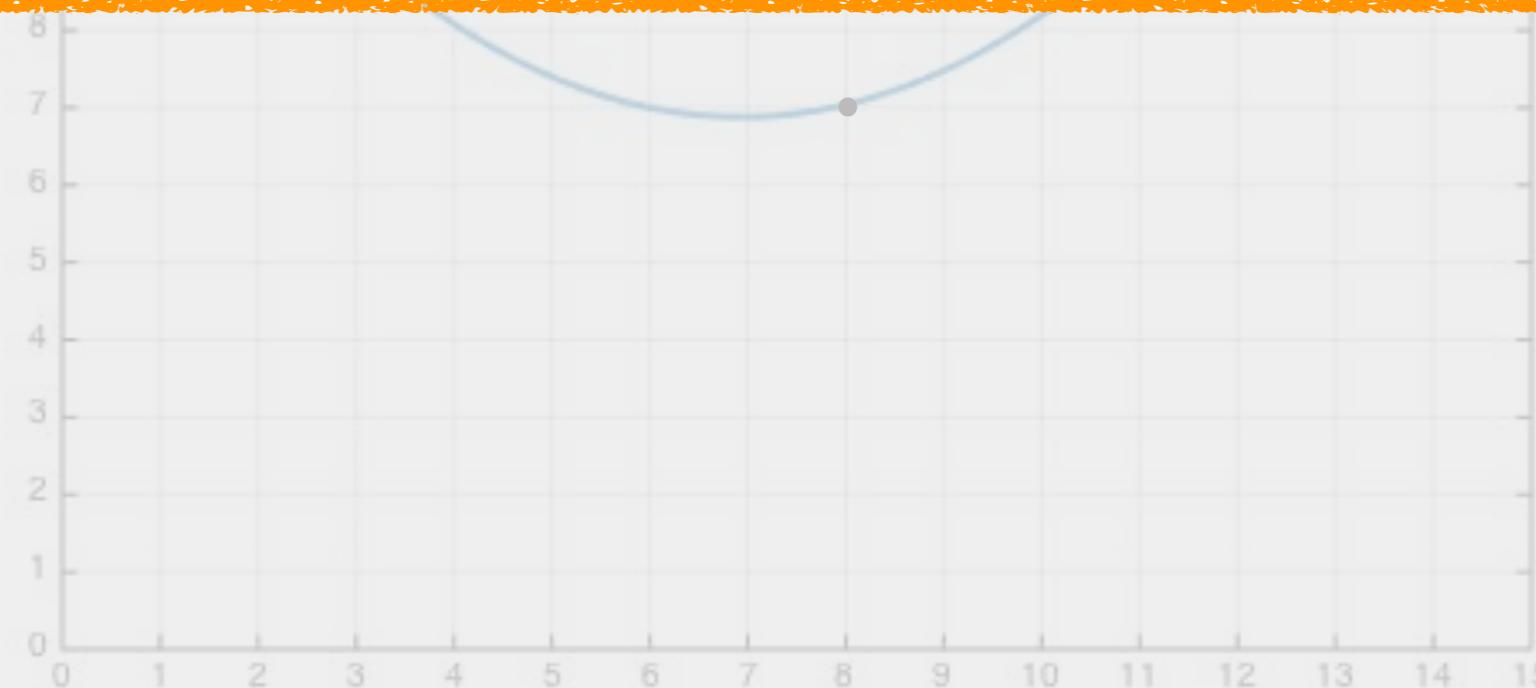
Degree 2 polynomials

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$$f(x) = 0.140x^2 - 1.94x + 13.6$$



**Observation 3: Unique degree 2 polynomial
passes through 3 points**



Property 1 of polynomials

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Infinite degree 1 polynomials pass through a single point

Infinite degree 2 polynomials pass through two points

:

Infinite degree d polynomials pass through d points

Property 2 of polynomials

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Degree 1 polynomial is uniquely determined
by 2 points

Degree 2 polynomial is uniquely determined
by 3 points

:

Degree d polynomial is uniquely determined
by $d+1$ points

But can we protect the
BANANA using polynomials

??



Shamir's secret sharing

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Alice



Bob



Charlie



Trusted
Gru

Shamir's secret sharing



$$f(x) = -x + 13$$

secret : $f(0) = 13$

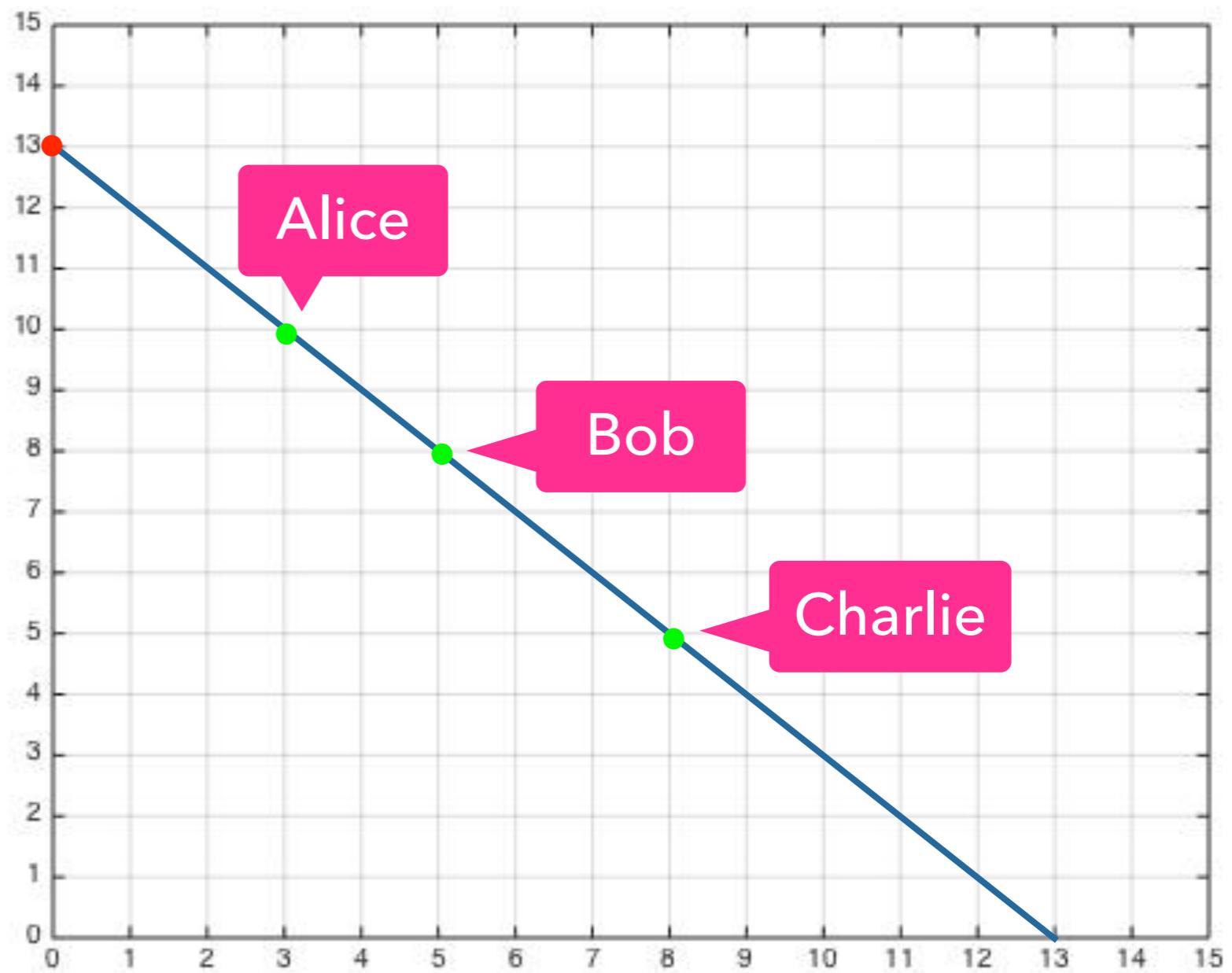
shares : $f(3), f(5), f(8)$

Choose a degree 1 polynomial

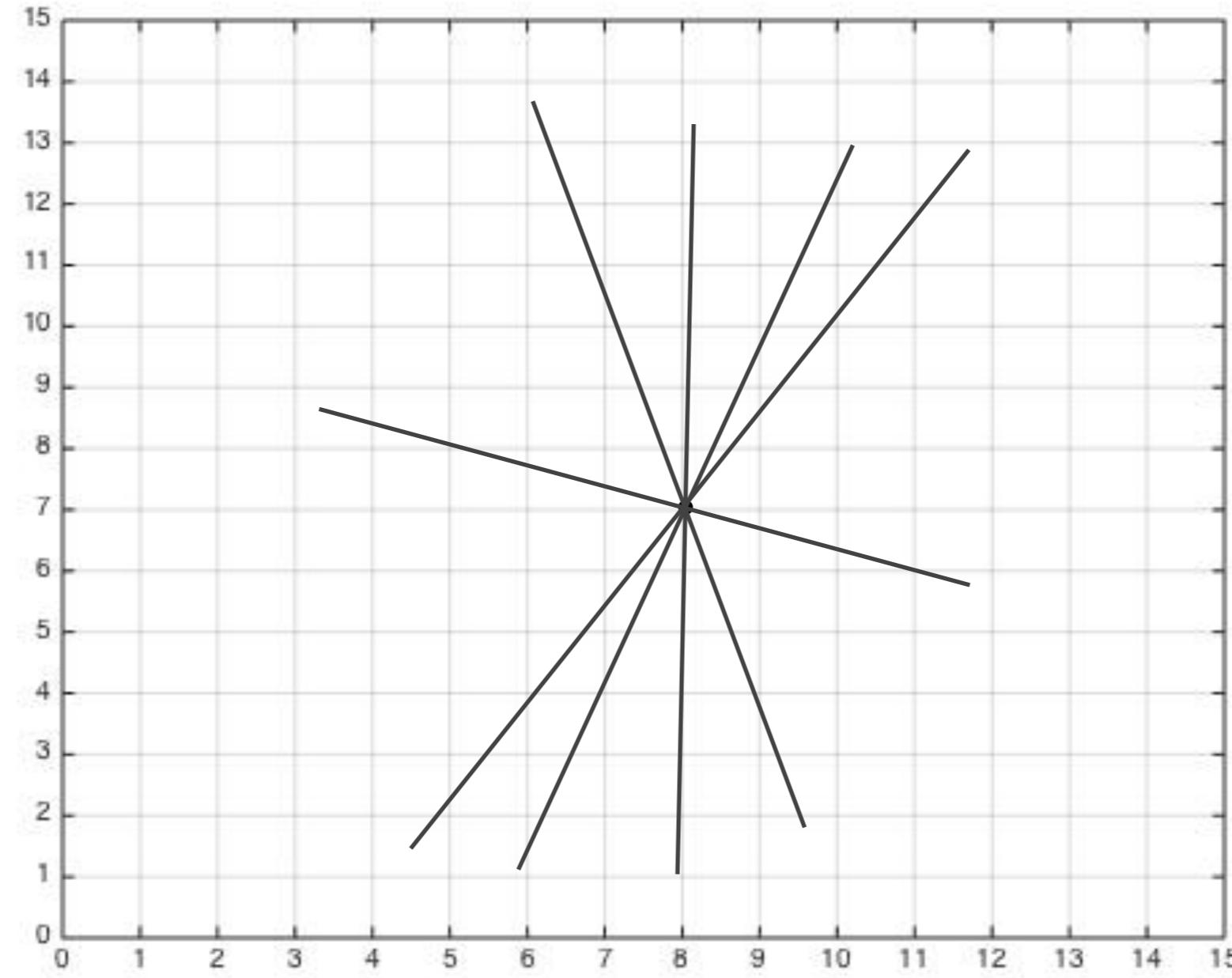
Secret in $f(0)$

Compute three shares by computing
 f at three more points

Distribute the shares to the parties



What is making it hard for Charlie to compute the secret using just his share?



Degree 1 polynomial is uniquely determined by 2 points

shares : $f(3), f(5), f(8)$

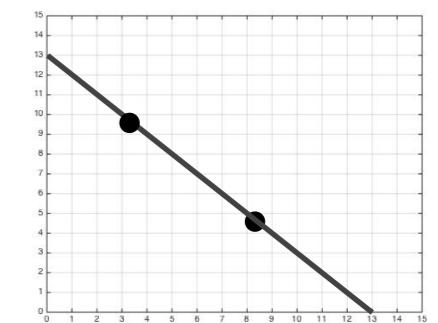
Choose a secret degree 1 polynomial

Hide the secret in $f(0)$

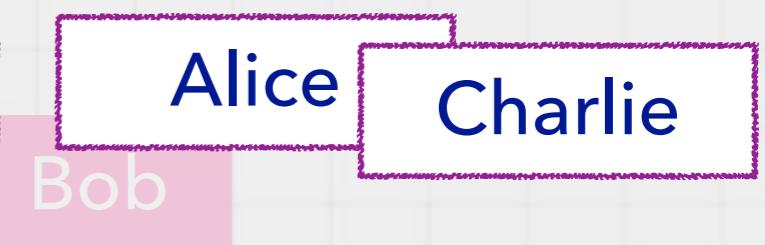
Com

f at three more points

Distribute the shares to the parties



At least 2 points needed to learn f



Once f is learnt, compute the secret $f(0)$

P L A Y

T I M E

G U E S S M Y S E C R E T

What if we need more than 2 people to be able to compute the secret ?

Prop 2: Degree d polynomial is uniquely determined by $d+1$ points

Ans: Use higher order (degree) polynomial

If t bad people, use a degree t polynomial

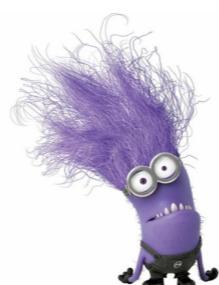
Hence, $t+1$ points needed to “learn” the polynomial

Secure secret sharing

vs

Insecure secret sharing

Anyone with $< t+1$ shares should not have more information than the one with 0 shares



2 - - -

- 0 - -

- - 1 -

- - - ?



Secure secret sharing

vs

Insecure secret sharing

Anyone with $< t+1$ shares should not have more information than the one with $t+1$

secret : 2017

INSECURE !!



0 S 33

$10 * 10 * 10 * 10$

- 0 -



3 Shares
2 0 1 -

10

-- 1 -



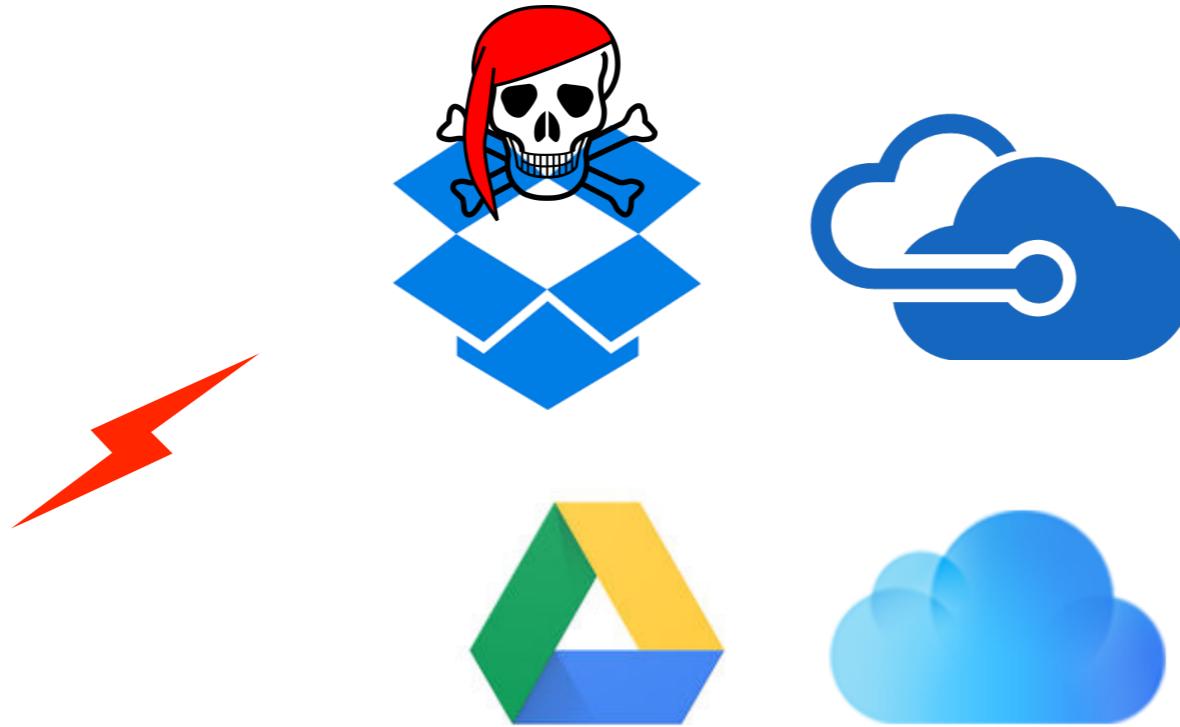
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Real World Applications

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Uploading personal data



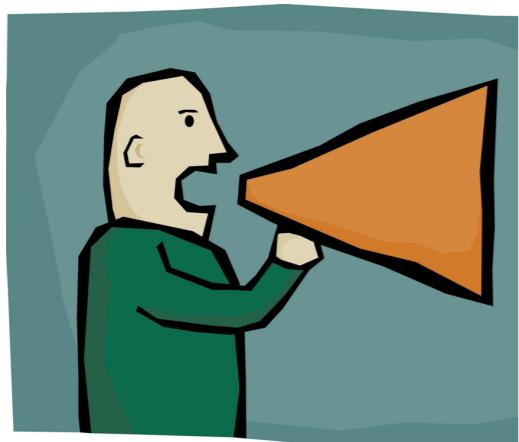
Real World Applications

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Uploading personal data



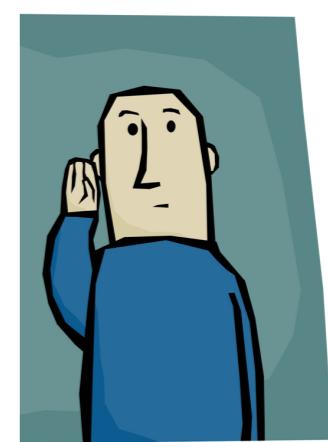
Error Correcting Codes



Sender

I will kill will kill you

Channel



Receiver

Real World Applications

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Uploading personal data



Error Correcting Codes



Add redundancy to the message to correct the errors introduced

I will kiss you



Any $t+1$ will recreate the message

A lot more corruptions needed to corrupt the message

Precisely, $n - t$

Real World Applications

Uploading personal data



Error Correcting Codes



Add redundancy to the message to correct the errors introduced

Storing highly sensitive data

Encryption keys

Missile launch codes

Numbered bank accounts

T H E M O R S E C O D E

Questions ?