

**0/9 Questions Answered**

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## Week 3 Wednesday Review Quiz

Student Name

**Q1 Set-wise concatenation****4 Points**

For alphabet  $\Sigma$ , given languages  $L_1$  and  $L_2$  over  $\Sigma$ , the set-wise concatenation is defined as  $L_1 \circ L_2 = \{w \in \Sigma^* \mid w = uv \text{ for some strings } u \in L_1 \text{ and } v \in L_2\}$

**Q1.1****1 Point**

Consider the alphabet  $\{a, b\}$ . How many strings are in the set  $\{\varepsilon, a, b\} \circ \{\varepsilon, a, b\}$ ?

0 (i.e. the set is empty)

3

6

9

Some other (finite) number

Infinitely many unique strings

Save Answer

## Q1.2

2 Points

Let  $N_1 = (Q_1, \Sigma, \delta_1, q_1, F_1)$  and  $N_2 = (Q_2, \Sigma, \delta_2, q_2, F_2)$  be NFAs.

When applying the construction in Theorem 1.47 to build the NFA  $N = (Q, \Sigma, \delta, q_1, F_2)$  that recognizes  $L(N_1) \circ L(N_2)$ , select all and only the statements below that are universally true.

☐  $|Q| > |Q_1|$

☐  $|Q| > |Q_2|$

☐  $|Q| > 2$

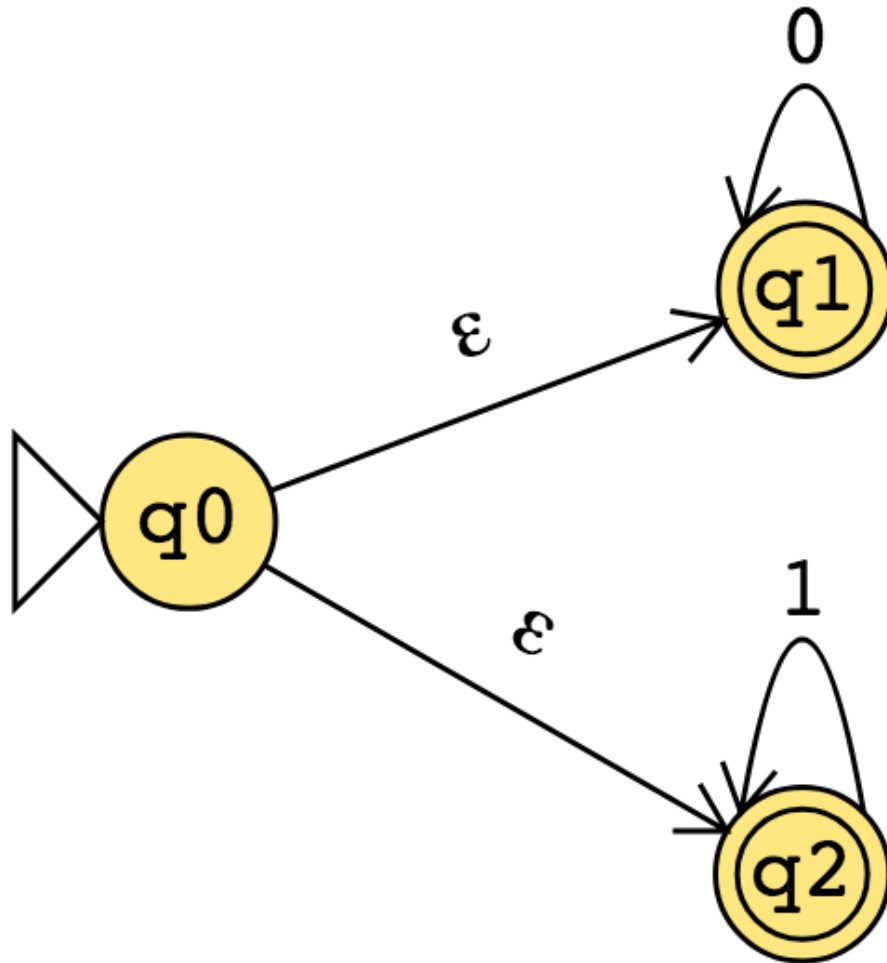
☐  $|Q| = |Q_1| + |Q_2|$

☐  $|Q| = |Q_1| \cdot |Q_2|$

Save Answer

**Q1.3**  
**1 Point**

The NFA whose state diagram below



is the result of applying the set-wise concatenation construction to obtain a machine that recognizes the language  $\{w \in \{0, 1\}^* \mid w \text{ has zeros followed by 1s}\}$

True

False, it is the result of applying the union construction instead to obtain the machine that recognizes the language  $\{w \in \{0, 1\}^* \mid w \text{ has all zeros or all 1s}\}$

False, it is the result of applying the intersection construction instead to obtain the machine that recognizes the language  $\{w \in \{0, 1\}^* \mid w \text{ has all zeros and all 1s}\}$

False, it is the result of applying the Kleene star construction instead to obtain the machine that recognizes the language  $\{0, 1\}^*$

Save Answer

**Q2 Kleene star****2 Points****Q2.1****1 Point**

Select all and only the languages below for which  $L^* = L$ .

☐  $\emptyset$ ☐  $\{\varepsilon\}$ ☐  $\{0\}$ ☐  $\{0, 1\}$ ☐  $\{0, 1\}^*$ Save Answer**Q2.2****1 Point**

True or False: The construction from Theorem 1.49 for an NFA that recognizes  $L^*$  from an NFA that recognizes  $L$  always gives the smallest number of states required in an NFA that recognizes  $L^*$ .

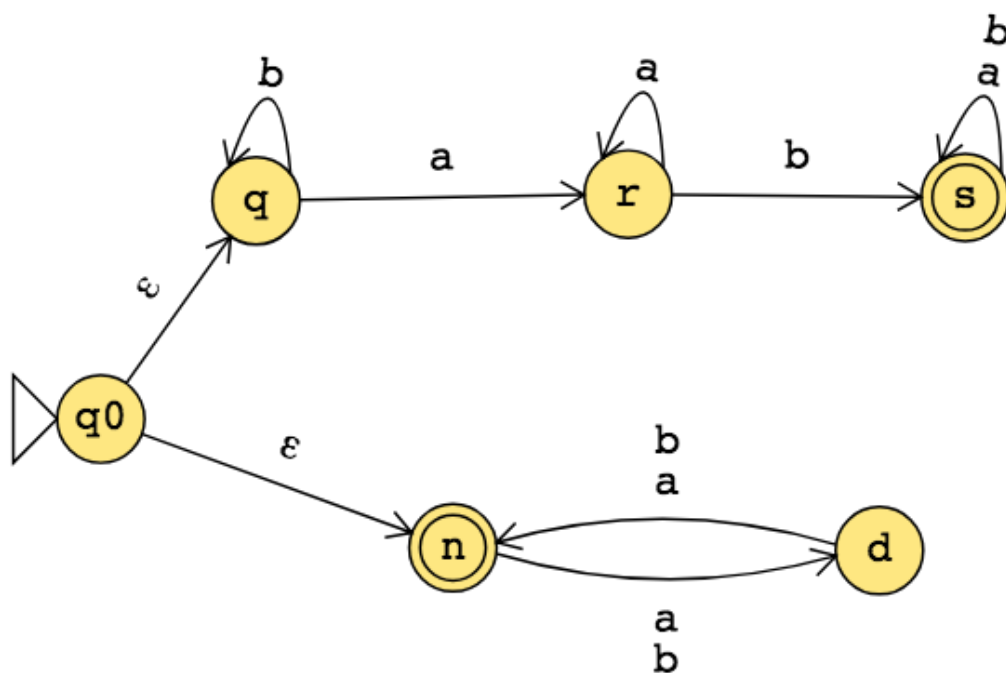
True

False

Save Answer

**Q3 NFA to DFA****2 Points**

Consider the following state diagram of a NFA over the alphabet  $\{a, b\}$ .



Answer the following questions about applying the construction for building an equivalent DFA from Theorem 1.39.

What is the start state of the equivalent DFA?

$q0$

$\{q0\}$

$\{q0, q, n\}$

What is the output of the transition function for the equivalent DFA from the start state on reading the character  $a$ ?



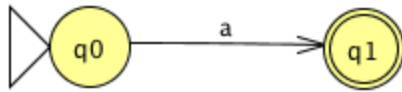
$\emptyset$  $\{q0\}$  $\{q\}$  $\{n\}$  $\{q, n\}$  $\{q, d\}$  $\{r, n\}$  $\{r, d\}$ 

None of the above, because DFA have a single state as the output of each transition function application, not a set of states.

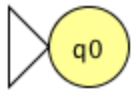
Save Answer

**Q4 Regular expression to NFA****1 Point**

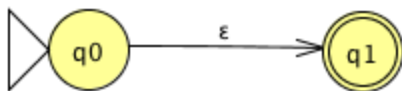
First diagram:



Second diagram:



Third diagram:



Which of the three diagrams above is a state diagram over the alphabet  $\{a, b\}$  for a NFA that recognizes the language  $L = \emptyset$ ?

first diagram

second diagram

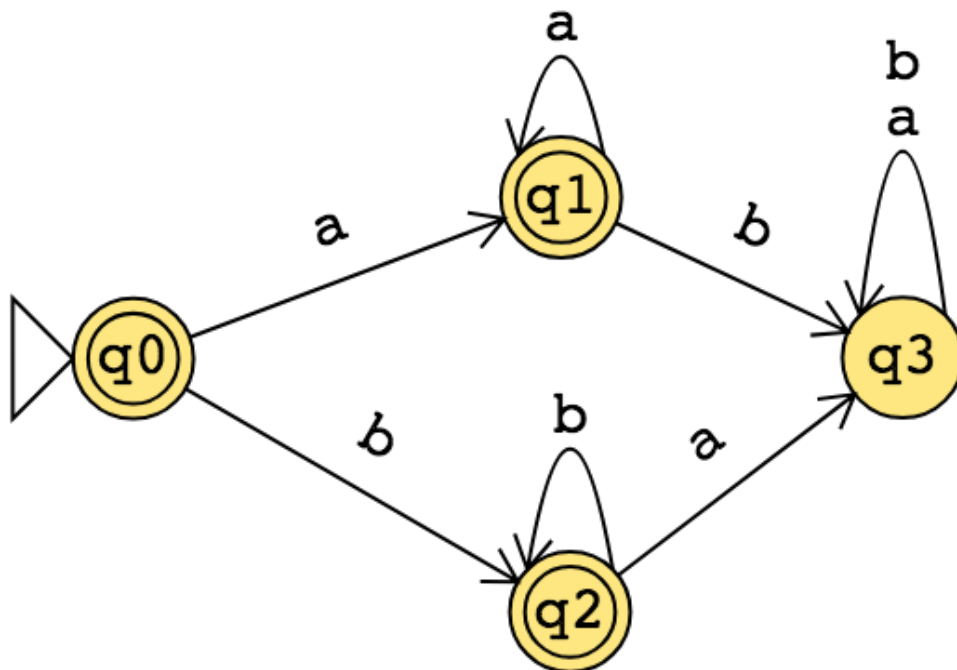
third diagram

None of the above

Save Answer

**Q5 DFA to regular expression**

1 Point



Which of the following regular expressions describe the language recognized by the DFA with state diagram above? (Select all and only that apply.)

☐  $a^+ \cup b^+$

☐  $\varepsilon \cup aa^* \cup bb^*$

☐  $a^*a \cup b^*b \cup \varepsilon$

☐  $aa^*b(a \cup b)^* \cup bb^*a(a \cup b)^*$

☐  $a^+b(a \cup b)^* \cup b^+a(a \cup b)^*$

Save Answer

### Q6 Feedback

0 Points

Any feedback about this week's material or comments you'd like to share?  
(Optional; not for credit)

Save Answer

Save All Answers

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