

Rethinking JavaScript Loops as Combinators

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Abstract—Loop combinators like *map* and *filter* have fallen out of use in performance critical sections of code in favor of *for* and *while* loops. This is tragic because not only are loop combinators more expressive than loop control structures but also they are simpler and more error resistant. Using inductive loop combinators, performance can be improved without sacrificing expressiveness. The resulting code is more modular and easier to maintain.

Index Terms—Functional Programming, JavaScript, Loop Combinators, Stateful Combinators, Method Chaining.

1 INTRODUCTION

LOOPS are one of the cornerstones of programming. Without loops every program would be uninteresting. Since loops are so important, it is paramount that they are efficient. Traditional loop control structures like *for* and *while* allow programmers to write efficient iteration constructs. However, they are neither modular nor composable.

Modularity and composability are the key to successful programming [1]. Together, they help combat software complexity — the biggest contributor to the “software crisis” identified in 1968 [2]. Functional programming provides new ways to decompose and combine programs, thereby mitigating software complexity [3].

In functional programming languages, loop combinators like *map* and *filter* are used instead of traditional loop control structures. Loop combinators are more expressive than loop control structures, but they are less efficient. For example, consider the following code which filters all the odd numbers in a list and then increments them by one:

```
var result = [1,2,3,4,5].filter(odd).map(inc);
```

Here, we are running two separate loops instead of one. Hence, the above functional program is less efficient than its equivalent imperative program:

```
var a = [1,2,3,4,5];
var l = a.length;
var i = 0, j = 0;
var result = [];

while (i < l) {
  var n = a[i++];

  if (odd(n)) {
    result[j++] = inc(n);
  }
}
```

```
}
}
```

The question is how to make loop combinators more efficient so that we don’t have to write loop control structures in performance critical sections of code? There are several existing solutions to this problem in JavaScript, inspired by transducers in Clojure [4]. However, we provide a solution that is on average thrice as fast as existing solutions:

- We define precisely what a loop combinator is and explain how they are composed (Section 2).
- We describe an alternative method of composing loop combinators using method chaining (Section 3).
- We describe how to create stateful combinators in Section 4, and how to exit from the loop early in Section 5.
- We describe an array specific optimization technique in Section 6, comparing the performance of our solution with existing solutions in Section 7.
- We conclude by showing how using loop combinators helps combat software complexity (Section 8).

2 BACKGROUND

In this section we lay the groundwork for the rest of the paper by providing a mathematical model of a loop. Furthermore, we use this model to give a precise definition of a loop combinator. To understand what a loop is, let’s take an example of a simple loop which computes the sum of a list of numbers:

```

var a = [1,2,3,4,5];
var i = 0, l = a.length;
var sum = 0;

while (i < l) {
  sum += a[i++];
}

```

This loop consists of three distinct elements: the initial result of the loop, the body of the loop which updates the result on every iteration and an input tape which produces one input element for every iteration of the loop body. The loop body may be executed several times. However, the loop always has an associated result even if the loop body is never executed.

For the purpose of abstraction, we won't concern ourselves as to how the input tape produces an input element for the loop body. The input tape could be a list which produces its own elements as input elements, or it could be a natural number which produces its predecessors as input elements, etc.

Let's assume that A is the smallest set that contains all the input elements produced by the input tape and that B is the set of all the possible results of the loop. Then the loop is defined by the 4-tuple $\langle A, B, B_0, F \rangle$ where $B_0 \in B$ is the initial result of the loop and $F : B \times A \rightarrow B$ represents the loop body.

Since, the loop is decoupled from the input tape we can use the same loop with any kind of input tape. For example, the aforementioned summation loop is defined by the 4-tuple $\langle \mathbb{N}, \mathbb{N}, 0, + \rangle$. Hence, it can be used with any kind of input tape that produces natural numbers.

Now that we provided a mathematical model for loops we can proceed to give a precise definition for loop combinators. A loop combinator is a function that transforms an input tape. For example, *map* and *filter* are loop combinators that transform lists. Loop combinators can be combined to form more complex loop combinators.

However, loop combinators need to be independent of the input tape. Since the loop body already abstracts away the input tape, an input tape independent loop combinator is simply a loop body transformer. For example, here are the definitions of the input tape independent *map* and *filter* loop combinators:

```

function map(func) {
  return function (body) {
    return function (result, value) {
      return body(result, func(value));
    };
  };
}

function filter(predicate) {
  return function (body) {
    return function (result, value) {
      return predicate(value) ?

```

```

    body(result, value) :
    result;
  };
}

```

A loop body is a partial function $F : (result : B) \times (value : A) \rightarrow B$. A loop combinator is a partial function $G : \forall R((R \times B \rightarrow R) \rightarrow R \times A \rightarrow R)$. Notice that the loop combinator knows nothing about the result of the loop. However, by returning the previous result instead of calling the output loop body it can continue with the next iteration. In Section 5 we define an alternative semantics of continuation.

Loop combinators can be composed using simple function composition. For example, the loop that computes the sum of all the odd numbers incremented by one can be defined as $\langle \mathbb{N}, \mathbb{N}, 0, (filter(odd) \circ map(inc))(+) \rangle$. This reads from left to right as "filter all the odd numbers, increment them by one and then compute their sum."

3 METHOD CHAINING

Object-oriented languages like JavaScript do not have a function composition operator. Although defining function composition is simple yet it's only efficient when composing two functions. When composing more than two functions, we create unnecessary intermediate functions:

```

function compose(f, g) {
  return function (x) {
    return f(g(x));
  };
}

```

Given, $(f : c \rightarrow d)$, $(g : b \rightarrow c)$ and $(h : a \rightarrow b)$, we get $(e : a \rightarrow d)$ as follows:

```

var e = compose(compose(f, g), h); // or
var e = compose(f, compose(g, h));

```

One way of getting rid of these unnecessary intermediate functions $(f \circ g)$ and $(g \circ h)$ is to create a generic function, *chain*, to compose a chain of functions:

```

var e = chain([f, g, h]);

function chain(xs) {
  return function (x) {
    var fs = xs, length = fs.length;
    while (length > 0) x = fs[--length](x);
    return x;
  };
}

```

However, a more object-oriented way would be to use method chaining. It is widely known that $(f \circ g \circ h)(x)$ in functional languages is equivalent to $x \cdot h() \cdot g() \cdot f()$ in object-oriented languages where x is an object, (\cdot) is the membership operator and $h()$, $g()$ and $f()$ are method calls. Hence, we can

write $(\text{filter}(\text{odd}) \circ \text{map}(\text{inc}))(+)$ in JavaScript as $\text{chain}(+) \cdot \text{map}(\text{inc}) \cdot \text{filter}(\text{odd}) \cdot \text{value}()$:

```
var LoopBody = defclass({
  constructor: function (body) {
    this.body = body;
  },
  map: function (func) {
    var body = this.body;

    return new LoopBody(mapBody);

    function mapBody(result, value) {
      return body(result, func(value));
    }
  },
  filter: function (predicate) {
    var body = this.body;

    return new LoopBody(filterBody);

    function filterBody(result, value) {
      return predicate(value) ?
        body(result, value) :
        result;
    }
  }
});

var body = new LoopBody(add)
  .map(inc)
  .filter(odd)
  .body;

// (1 + 1) + (3 + 1) + (5 + 1) = 12
console.log([1, 2, 3, 4, 5].reduce(body, 0));

function add(x, y) { return x + y; }
function inc(x)   { return x + 1; }
function odd(x)   { return x % 2 === 1; }

function defclass(prototype) {
  var constructor = prototype.constructor;
  constructor.prototype = prototype;
  return constructor;
}
```

This is commonly known as the chain-value pattern in JavaScript, which rose to prominence because of the popular Underscore.js library [5].

4 STATEFUL COMBINATORS

Until now, we've only created stateless loop combinators like *map* and *filter*. However, there are many useful loop combinators that require state such as the *drop* combinator which skips the first *n* elements of an input tape. The naïve solution is add a stateful loop combinator to the prototype of *LoopBody*:

```
LoopBody.prototype.drop = function (n) {
  var body = this.body;

  return new LoopBody(dropBody);

  function dropBody(result, value) {
    return n > 0 ?
      (n--, result) :
```

```
    body(result, value);
  }
};
```

However, this will only work once after which it will fail. Thus, this is not a good solution because it is not reusable. One way to make stateful loop combinators reusable is to create them anew for each loop. Indeed, this is what transducers in Clojure do [6]. However, creating a new chain of combinators for every loop is inefficient.

As an alternative, we propose to refresh the state once for every loop using an *init* function. In Section 6, we'll extend this function to implement an array specific optimization technique. The *init* function belongs to the same scope as the loop body and hence, it can modify the state of the loop body.

Although the *init* function is only used for its side effect yet it's still a function. Hence, the question arises as to what is its domain and codomain. Since we don't care about either of them, we get the partial function $\text{init} : \forall A(A \rightarrow A)$. Hence, the *init* chain of functions always returns its argument whether or not it performs any side effects.

```
var LoopBody = defclass({
  constructor: function (init, body) {
    this.init = init;
    this.body = body;
  },
  map: function (func) {
    var body = this.body;

    return new LoopBody(this.init, mapBody);

    function mapBody(result, value) {
      return body(result, func(value));
    }
  },
  filter: function (predicate) {
    var body = this.body;

    return new LoopBody(this.init,
      filterBody);

    function filterBody(result, value) {
      return predicate(value) ?
        body(result, value) :
        result;
    }
  },
  drop: function (_n) {
    var n;

    var init = this.init;
    var body = this.body;

    return new LoopBody(dropInit, dropBody);

    function dropInit(x) {
      n = _n;
      return init(x);
    }
  }
});
```

```

    function dropBody(result, value) {
      return n > 0 ?
        (n--, result) :
        body(result, value);
    }
  });

var loop = new LoopBody(id, add)
  .map(inc)
  .filter(odd)
  .drop(5);

loop.init();
// (7 + 1) + (9 + 1) = 18
console.log([1,2,3,4,5,6,7,8,9,10]
  .reduce(loop.body, 0));

function id(x)    { return x; }
function add(x, y) { return x + y; }
function inc(x)   { return x + 1; }
function odd(x)   { return x % 2 === 1; }

function defclass(prototype) {
  var constructor = prototype.constructor;
  constructor.prototype = prototype;
  return constructor;
}

```

Notice that stateless loop combinators like *map* and *filter* don't have their own *init* functions. They simply return the previous function. This makes refreshing the state a little more efficient than creating the entire loop combinator chain anew. However, there's also a downside to this. You can't use the same loop body concurrently because state would be shared. However, in practice this only happens in asynchronous loops and can be easily avoided.

5 EARLY TERMINATION

Traditional loop bodies have a certain amount of control over the flow of the loop itself. You can skip the rest of the iteration using *continue* and you can break out of the loop when you get the final answer without having to read the rest of the input tape using *break*. In fact, our current definition of a loop body captures the semantics of *continue*.

The loop body is a partial function $F : B \rightarrow A \rightarrow B$ where B is the set of all the possible results of the loop and A is the smallest set that contains all the input elements produced by the input tape. By returning the previous result, a loop combinator can effectively skip the rest of iteration. In fact, this is precisely what the *filter* and *drop* combinators do.

Unfortunately, our current definition of a loop body does not capture the semantics of *break*. Hence, we can't define certain loop combinators like *take*, which takes the first n elements of an input tape and skips the rest. Although we could keep skipping the rest of the elements using the semantics of *continue* yet that is neither efficient nor would it work on infinite input tapes.

Since our definition of a loop doesn't capture the semantics of *break* we need to revise our definition. This means that we won't be able to execute our loops with a simple *reduce* function like we've been doing all along. We will need something more powerful. In Clojure the *reduce* function can be explicitly terminated by applying *reduced* to the output of the loop body [7].

Although the semantics of *break* is explicit in Clojure yet the semantics of *continue* is implicit. However, we believe, in accordance with the Zen of Python, that explicit is better than implicit [8]. Therefore, in our new definition of a loop body we made both *break* and *continue* explicit $F : B \rightarrow A \rightarrow \langle B, \text{bool} \rangle$.

The set *bool* is a set of two elements, *true* and *false*. Hence, we may decide whether *break* corresponds to *true* or *false*. Although it doesn't really matter yet to follow a convention we choose *true* to denote *break* and *false* to denote *continue*. We also define a special *reduce* function for arrays which allows early termination:

```

function reduce(loop, a, xs) {
  loop.init();

  var f = loop.body;
  var l = xs.length;
  var i = 0;

  while (i < l) {
    var pair = f(a, xs[i++]);
    if (pair.snd) return pair.fst;
    a = pair.fst;
  }

  return a;
}

```

Now that we've defined the semantics of both *break* and *continue* we can define the *take* loop combinator. In addition, we must also redefine the previous loop bodies so that they explicitly *break* or *continue*. Since the *map* loop combinator doesn't *break* nor *continue*, it doesn't need to be modified.

```

var LoopBody = defclass({
  constructor: function (init, body) {
    this.init = init;
    this.body = body;
  },
  map: function (func) {
    var body = this.body;

    return new LoopBody(this.init, mapBody)
      ;

    function mapBody(result, value) {
      return body(result, func(value))
        ;
    }
  },
  filter: function (predicate) {
    var body = this.body;

    return new LoopBody(this.init,
      filterBody);
  }
});

```

```

    function filterBody(result, value) {
        return predicate(value) ?
            body(result, value) :
            { fst: result, snd: false };
    },
    drop: function (_n) {
        var n;

        var init = this.init;
        var body = this.body;

        return new LoopBody(dropInit, dropBody)
            ;

        function dropInit(x) {
            n = _n;
            return init(x);
        }

        function dropBody(result, value) {
            return n > 0 ?
                (n--, { fst: result, snd: false
                    }) :
                body(result, value);
        }
    },
    take: function (_n) {
        var n;

        var init = this.init;
        var body = this.body;

        return new LoopBody(takeInit, takeBody)
            ;

        function takeInit(x) {
            n = _n;
            return init(x);
        }

        function takeBody(result, value) {
            return n > 0 ?
                (n--, body(result, value)) :
                { fst: result, snd: true };
        }
    }
});

var loop = new LoopBody(id, add)
    .map(inc)
    .filter(odd)
    .take(5);

// (1 + 1) + (3 + 1) + (5 + 1) = 12
console.log(reduce(loop, 0,
    [1,2,3,4,5,6,7,8,9,10]));

function id(x)    { return x; }
function add(x, y) { return { fst: x + y,
                                snd: false }; }

function inc(x)   { return x + 1; }
function odd(x)   { return x % 2 === 1; }

function defclass(prototype) {
    var constructor = prototype.constructor;
    constructor.prototype = prototype;
    return constructor;
}

```

6 OPTIMIZATION FOR ARRAYS

Until now, we've only used the summation loop. Another common kind of loop is a generative loop. For example, we can create a loop that generates an array. The traditional way to do this would be to create an empty array for every loop and then push elements to the end of the array.

```

var toArray = new LoopBody(id, push);

// a new array [1,2,3,4,5]
console.log(reduce(toArray, [], [1,2,3,4,5]));

function push(result, value) {
    result.push(value);
    return { fst: result, snd: false };
}

```

However, if we knew the length of the array beforehand then we could allocate just enough space for the array, thereby making the loop faster. The question is, how do we determine the length of the array beforehand? As it turns out, the *init* function which we've only been using for its side effect of refreshing the state is the perfect solution to this problem.

Currently, we have defined $init : \forall A(A \rightarrow A)$ which is not very useful as a function. However, if we change its type to $init : \mathbb{N} \rightarrow R$ where R is the set of all the possible results of the loop then we can return the initial result of the loop, given the minimum expected number of loop iterations, in addition to refreshing the state of the loop.

This would mean that the entire loop would be represented as a single data structure. However, some stateless loop combinators like *filter* would also have to implement the *init* function. Nevertheless, the performance gain for generating arrays would make it worth implementing. Finally, we would have to modify the definition of our special *reduce* function for arrays.

```

function reduce(loop, xs) {
    var l = xs.length;
    var a = loop.init(l);
    var f = loop.body;
    var i = 0;

    while (i < l) {
        var pair = f(a, xs[i++]);
        if (pair.snd) return pair.fst;
        a = pair.fst;
    }

    return a;
}

var Loop = defclass({
    constructor: function (init, body) {
        this.init = init;
        this.body = body;
    },
    map: function (functor) {
        var body = this.body;

```



```

    return new Loop(this.init, mapBody);

    function mapBody(result, value) {
        return body(result, functor(value))
    }
},
filter: function (predicate) {
    var init = this.init;
    var body = this.body;

    return new Loop(filterInit, filterBody)

    function filterInit(length) {
        return init(0);
    }

    function filterBody(result, value) {
        return predicate(value) ?
            body(result, value) :
            { fst: result, snd: false };
    }
},
drop: function (_n) {
    var n;

    var init = this.init;
    var body = this.body;

    return new Loop(dropInit, dropBody);

    function dropInit(length) {
        n = _n;
        return init(Math.max(length - n, 0))
    }

    function dropBody(result, value) {
        return n > 0 ?
            (n--, { fst: result, snd: false }) :
            body(result, value);
    }
},
take: function (_n) {
    var n;

    var init = this.init;
    var body = this.body;

    return new Loop(takeInit, takeBody);

    function takeInit(length) {
        n = _n;
        return init(Math.min(length, n));
    }

    function takeBody(result, value) {
        return n > 0 ?
            (n--, body(result, value)) :
            { fst: result, snd: true };
    }
}
});

var toArray = (function () {
    var i;

    return new Loop(toArrayInit, toArrayBody);

```

```

    function toArrayInit(length) {
        i = 0;
        return length > 0 ?
            new Array(length) : [];
    }

    function toArrayBody(result, value) {
        result[i++] = value;
        return { fst: result, snd: false };
    }
})();

var loop = toArray
    .map(inc)
    .filter(odd)
    .take(5)
    .drop(2);

// prints [4, 6, 8]
console.log(reduce(loop,
    [1,2,3,4,5,6,7,8,9,10]));

function inc(x) { return x + 1; }
function odd(x) { return x % 2 === 1; }

function defclass(prototype) {
    var constructor = prototype.constructor;
    constructor.prototype = prototype;
    return constructor;
}

```

It's interesting to note that *dropInit* calls the next *init* function with the maximum of the *length - n* and zero while *takeInit* calls the next *init* function with the minimum of the *length* and *n*. It's even more interesting to note that *filterInit* calls the next *init* function with zero. This is because *filter* can't determine in advance how many elements it will filter. Hence, it returns the minimum number of possible elements (i.e. zero). Thus if you have a *filter* in your chain then this optimization won't work.

7 RELATED WORK AND BENCHMARKS

The loop combinators that we described in this paper are very similar to the Rich Hickey's transducers in Clojure [4]. In fact, our work on loop combinators has been inspired by Clojure transducers. However, there are several important differences between our loop combinators and Clojure transducers.

First, our loop combinators aren't as powerful as Clojure transducers because we don't have an output function that ties up loose ends after the loop body has completed its execution. This means that a few loop combinators, such as the *partition-by* combinator in Clojure, cannot be implemented. Nevertheless, adding the output function to the definition of the loop is straightforward.

Reducers in Clojure are just Moore machines [9] and transducers in Clojure are just Mealy machines [10] augmented with the ability to perform entry actions. Hence, our loop combinators are strictly less powerful than Moore machines and Mealy machines.

However, augmented with an output function they would be just as powerful as transducers in Clojure.

That being said, we are not the first people to import Clojure transducers to JavaScript. There are a lot of implementations of transducers in JavaScript. However, we believe that our approach to transducers have a lot of novel ideas that could be beneficial for everyone. We benchmarked our code against LoDash, the fastest functional programming library in JavaScript, and the results look quite promising.

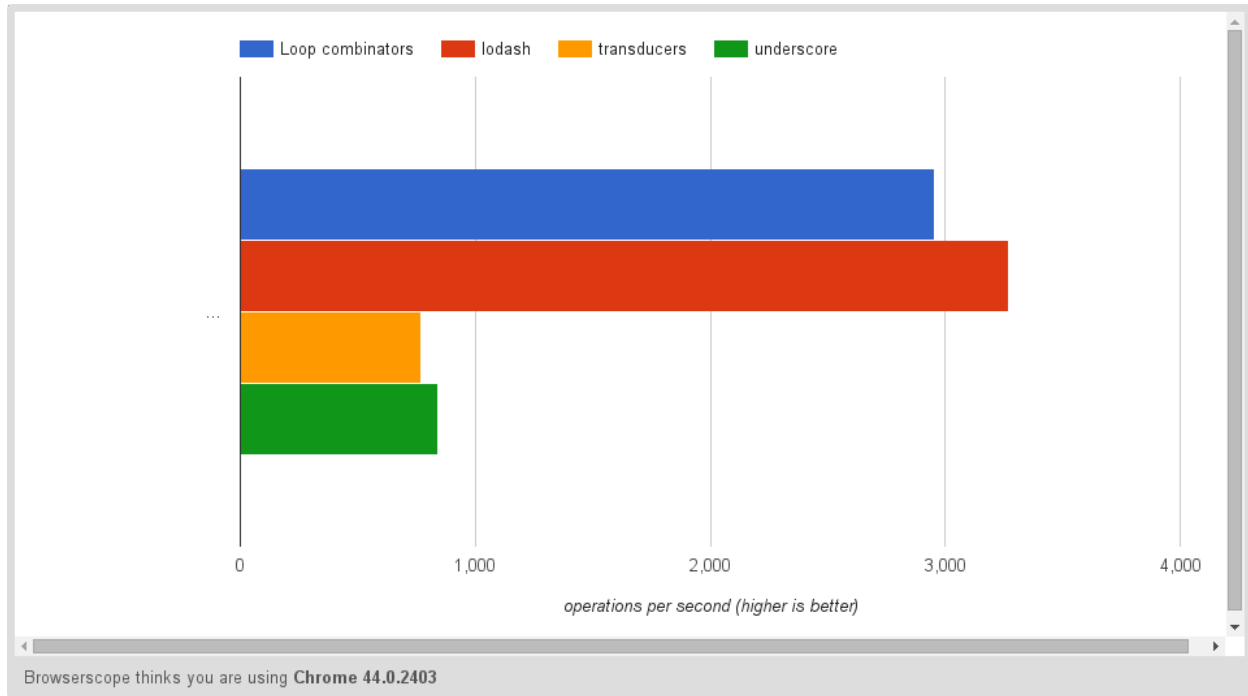


Fig. 1. Comparing `toArray · map(inc) · filter(odd)` across libraries.

8 CONCLUSION

Our goal was to promote the use of loop combinators over traditional loop control structures in order to reduce software complexity by making loop combinators more efficient. Although our loop combinators are quite efficient yet they are slower than the fastest combinator libraries. However, in the process of creating this library we discovered some insightful facts about method chaining and loop optimization which we hope will be helpful to others.

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