

Intelligence:

- AI: not copying exact human-level intelligence
: not a replica
- AI: creates env. to ease human functions
- Birth of AI: 1956 (workshop at Dartmouth College)
- Aim for general principles: Every aspect of learning
- Early success: Checker Board (1952)
: Problem Solving (1955) (Arithmetic & Typing)
- Overwhelming Expectation Vs Underwhelming Results
: m/c translation

Eg: translate 'The spirit is willing but the flesh is weak'

ERROR!!

↓ RUSSIAN (via AI)

'The vodka is good but the meat is rotten'

Reasons for failure:

- No generalised rule for all languages.
- Need to train the machine with every word & rule of languages involved.

Problems:

→ Limited computation

→ Limited information

• In 1956 - computation was expensive (expensive hard discs etc.)

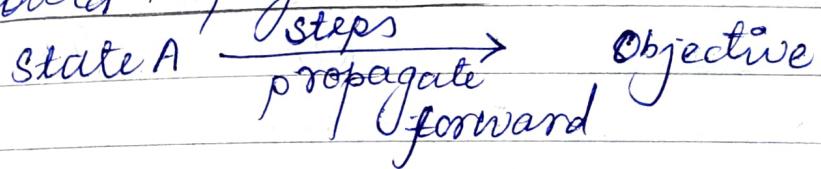
Contributions

- LISP (list processing) (logic-based prog. like if, 'else')
- Garbage collection

Knowledge Based Systems (KBS) (70-80's)

- Knowledge Based Systems (KBS) (70-80's)
- in Unit 2,
- Artificial Neural Netw. (ANN) (1943)
- thought: study human neurons to form artificial neural nets.

Front / Followed Propagation :



- Back Propagation : We know the src & dstn. & steps
- job is to go back & check if steps lead to source specified.

Eg: x is father of y, z is father of x.

Is z grandfather of y? → Back Prop.

Is y grandchild of z? → forward Prop.

- Convolution : to recognize handwritten digits (LeCun)

Deep Learning

- AlexNet (2012) : Object Recognition, computer vision commu
- AlphaGo (2016) : Deep reinforcement learning, defeat world champion (Lee Sedol)

Two Kineses of AI

Agents: how to create intelligence

Tools: how to benefit society?

AR: Augmented Reality
VR: Virtual Reality

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AR Vs VR

AR: In real world, using sensors, working IRL.
objects

VR: no connection with real world,

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wed

AGENTS:

Intelligent Agent (PEAS str.)

It will enact response to senses automatically.

P → Perceptors / performance measure

E → Environment

A → ~~Actions~~ Actuators (those which are take action)

S → Sensor

Ideal rational agents

- single Particular objective (Eg: cost minimization)
• ~~most of the time~~, we need a trade-off ^{to} time minimization

• Mostly, we need a trade-off ^{to} among all variables.

• optimization → single objective

multiple objective

(in collaboration, need to appreciate all needs & constraints)

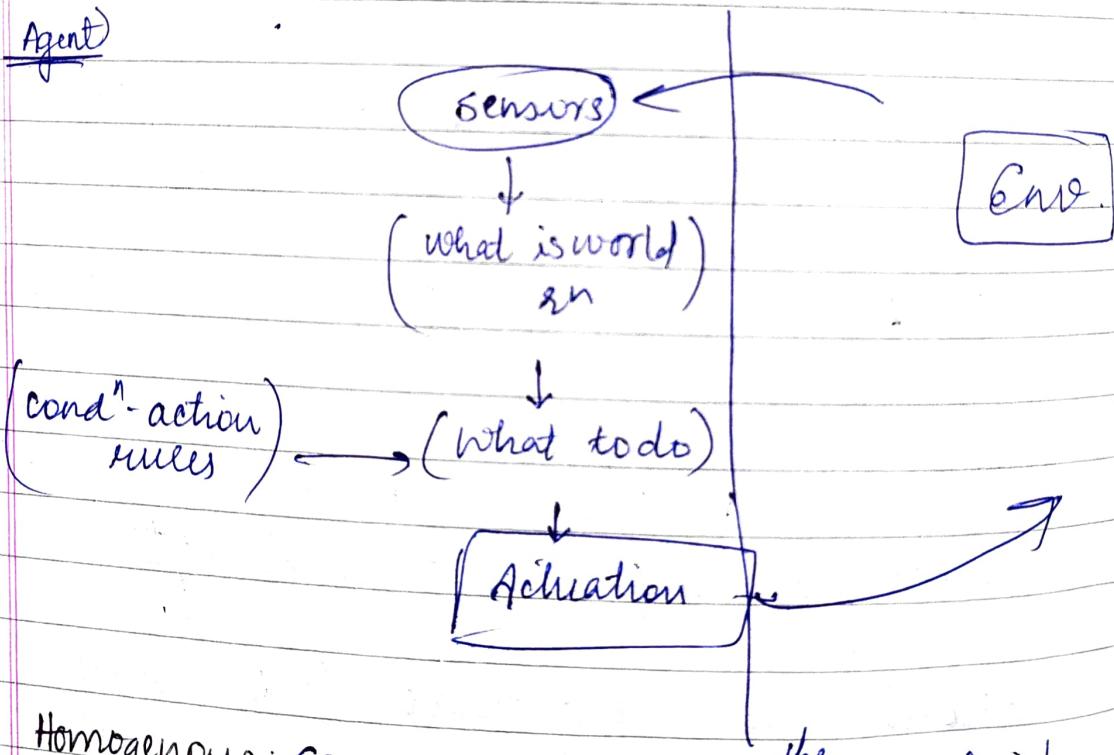
single obj → single constraint-satisfaction problem.

→ If more pple are concerned on 1 factor, we bias towards that single constraint-satisfaction prob.

- **Reflex:** catch the state of env. quickly & take action.
- only dependent on current state of env., \Rightarrow perceptible to you.
- on that state, action is taken.

A simple Reflex Agent

Agent)



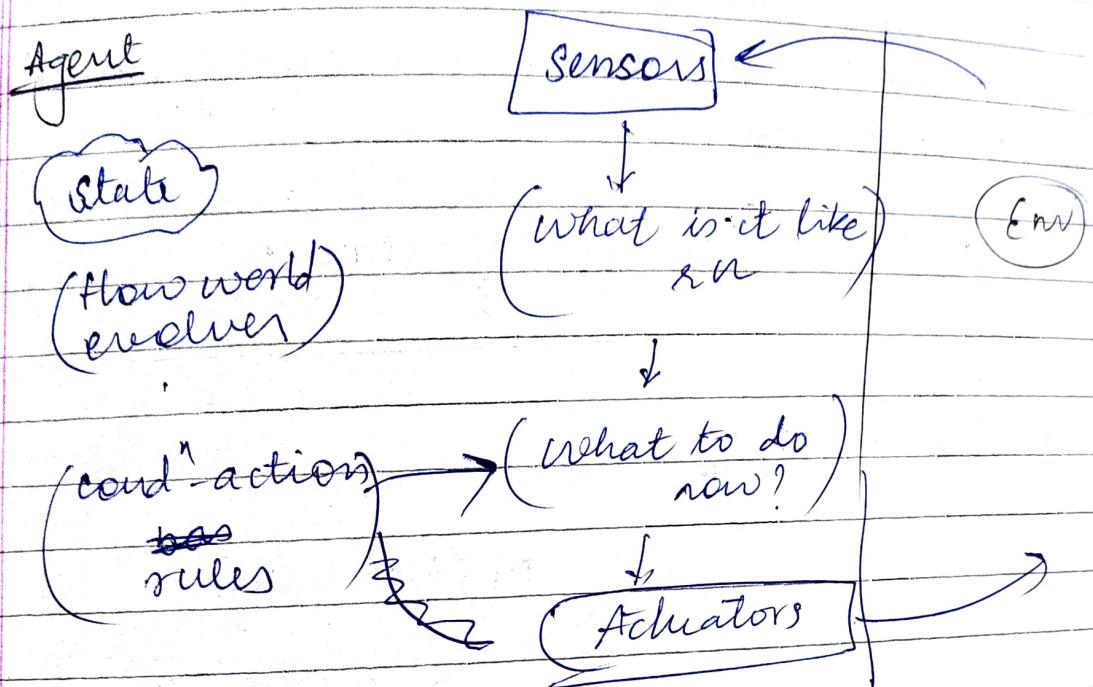
Homogeneous: Same rules for all of ^{the} ~~an~~ same kind.

Heterogeneous: Not same attributes, rules.



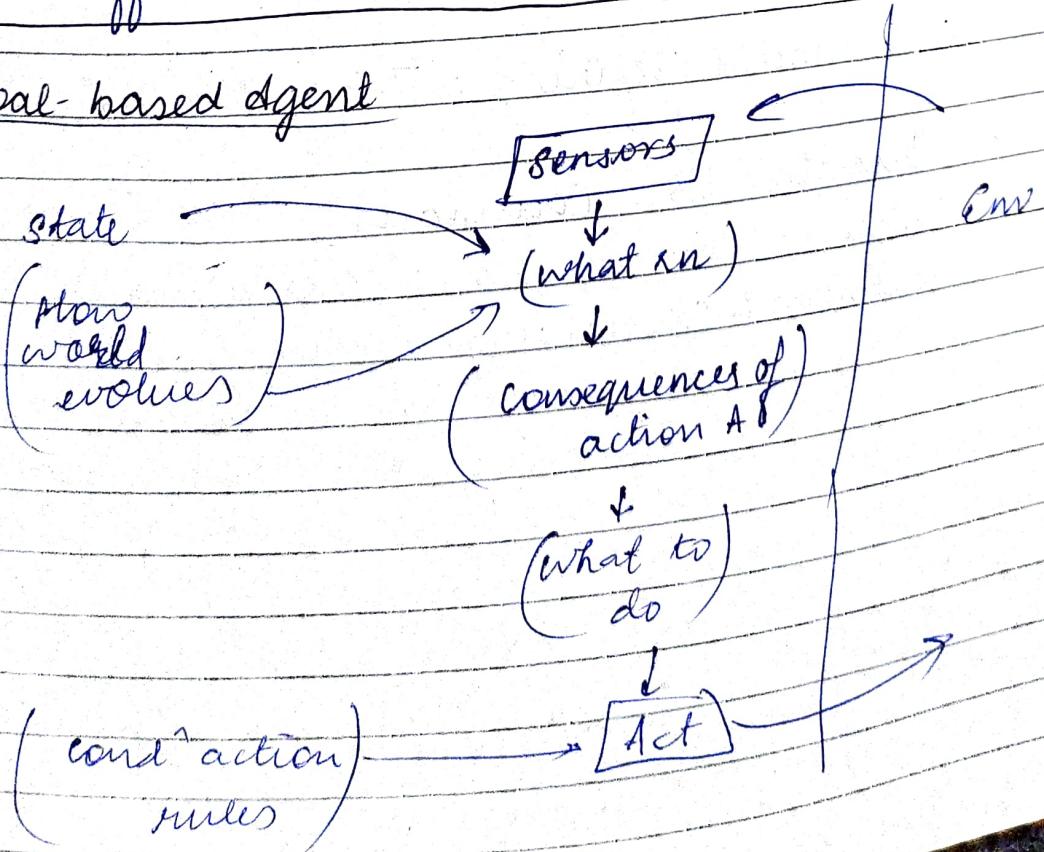
- Based on current state only
- Simple but limited
- works ~~if~~ only if env. is fully observable
- Eg: Vacuum Agent,
(Vacuum cleaner)
→ works if dust in env.

• Model-Based Reflex Agent



- considers history of environment.
- Eg: Rain detection (Weather forecast)
- somehow consider partially observable env.
- Eg: Traffic Detection.

• Goal-based agent



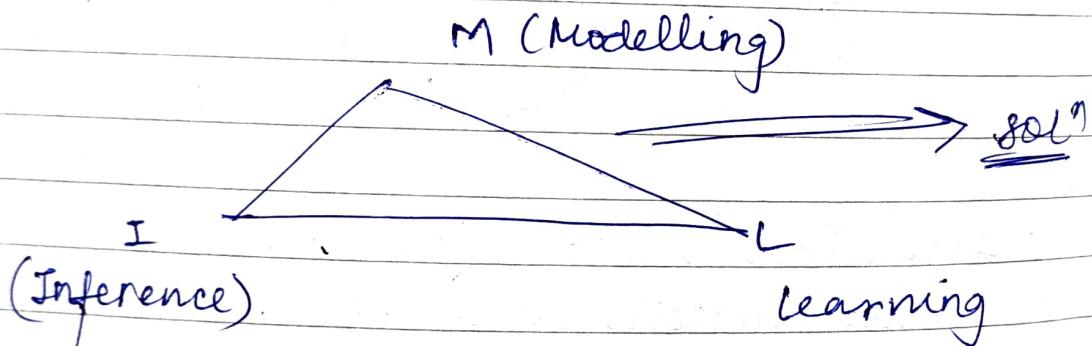
- Depends on consequences of action

Utility Based Agent

- Similar to goal-based
- A performance-measure which also gives a token of appreciation.

Utility → Benefit of action A.

• PEAS
Environment



Modelling : What is our env & constraints / req.

Environment Defining : 7-key Dimensions (clidy)

• Discrete Vs Continuous

Discrete : Playing chess: 1st step, wait, 2nd step

Continuous :- Car driving - acceleration

- All events are continuous, even if dependent of on env. factors.

- Complete Vs Incomplete

①: chess: Initially we know all rules all the time, response for all state.

②: Poker: Don't know what could happen after 2 chances.

- Fully observable: face recognition
Vs. Partially obs.: self-driven car.

- competitive vs collaborative

Eg: board games
- nin or loose

Eg: Alexa connected to music sys, listening music then ask alexa to play a song.

Eg: cricket is
- collab for a team
- compet for 2 team

- Static Vs Dynamic

③: Not changing on its own (voice/speech recog.)
- changes if action is taken (speech is recorded, no change)

④: self-driven car

- even if the car is stopped, there are changes in env.

- Deterministic

- We know the result of our action

Eg: speech recog.

(like finite automata machine)

vs

- stochastic

- self-driven car

- Applying brakes may or may not cause accident

- uncertain result for a specific action

- Single agent

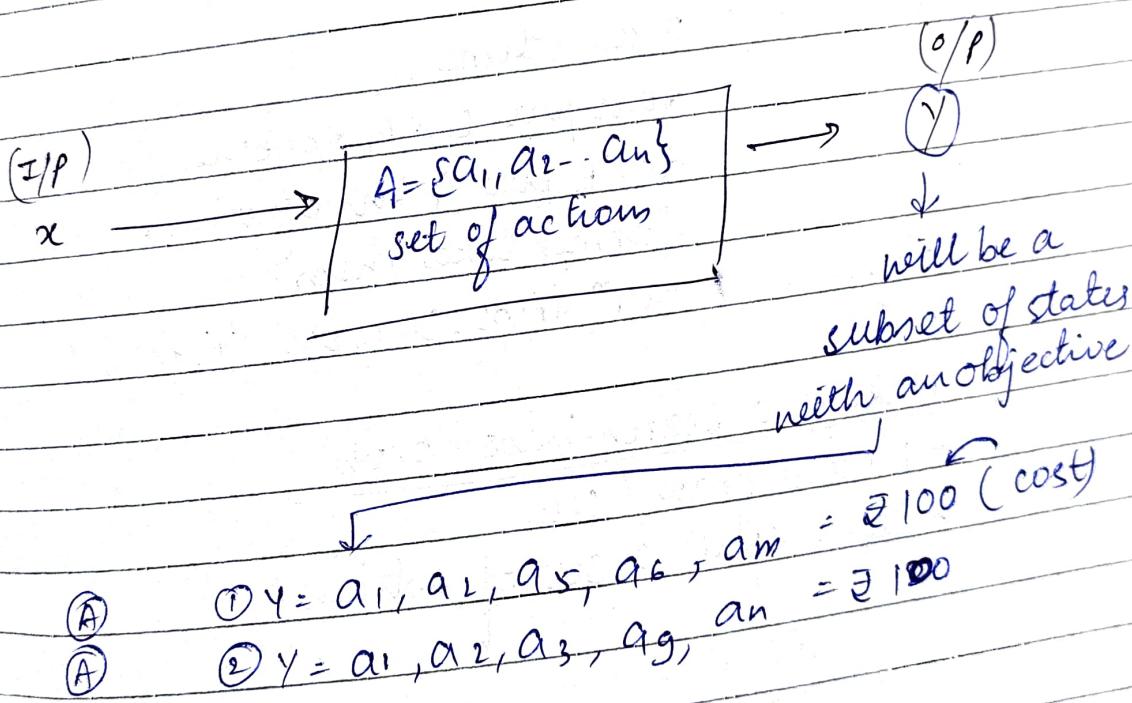
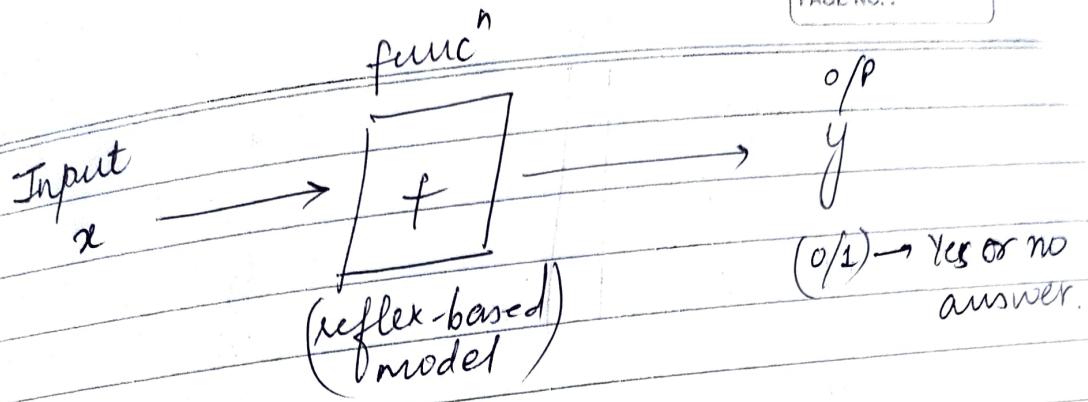
- Only 1 in the class attending

- Not competitive or collab

- maybe utility-based

- Multiagent

- either competitive or collaborative



Search
↓
search problem

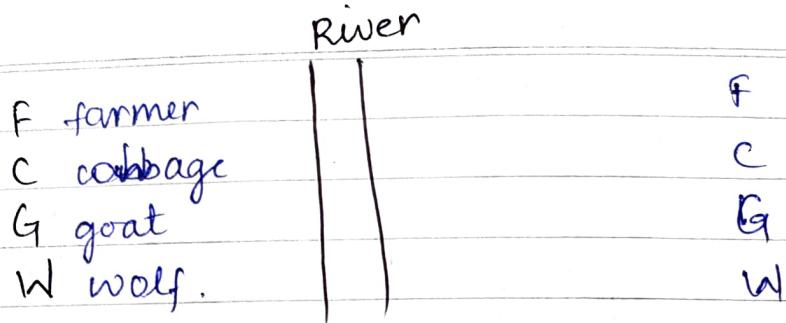
↑
Search algorithm

- Modelling real-world probs into a ~~set~~ subset of sol's.
- all possible set of actions
- Then comes constraints, and we go for search algos to achieve the objective

Eg: A real world prob.

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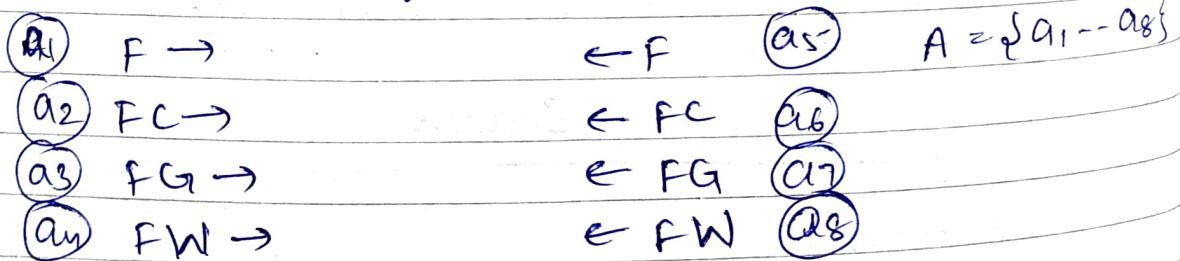
(y)

constraint : 2 at a time,
don't put C & G together,
don't put G & W together,
farmer takes them.

objective : Min^m trips. to cross river

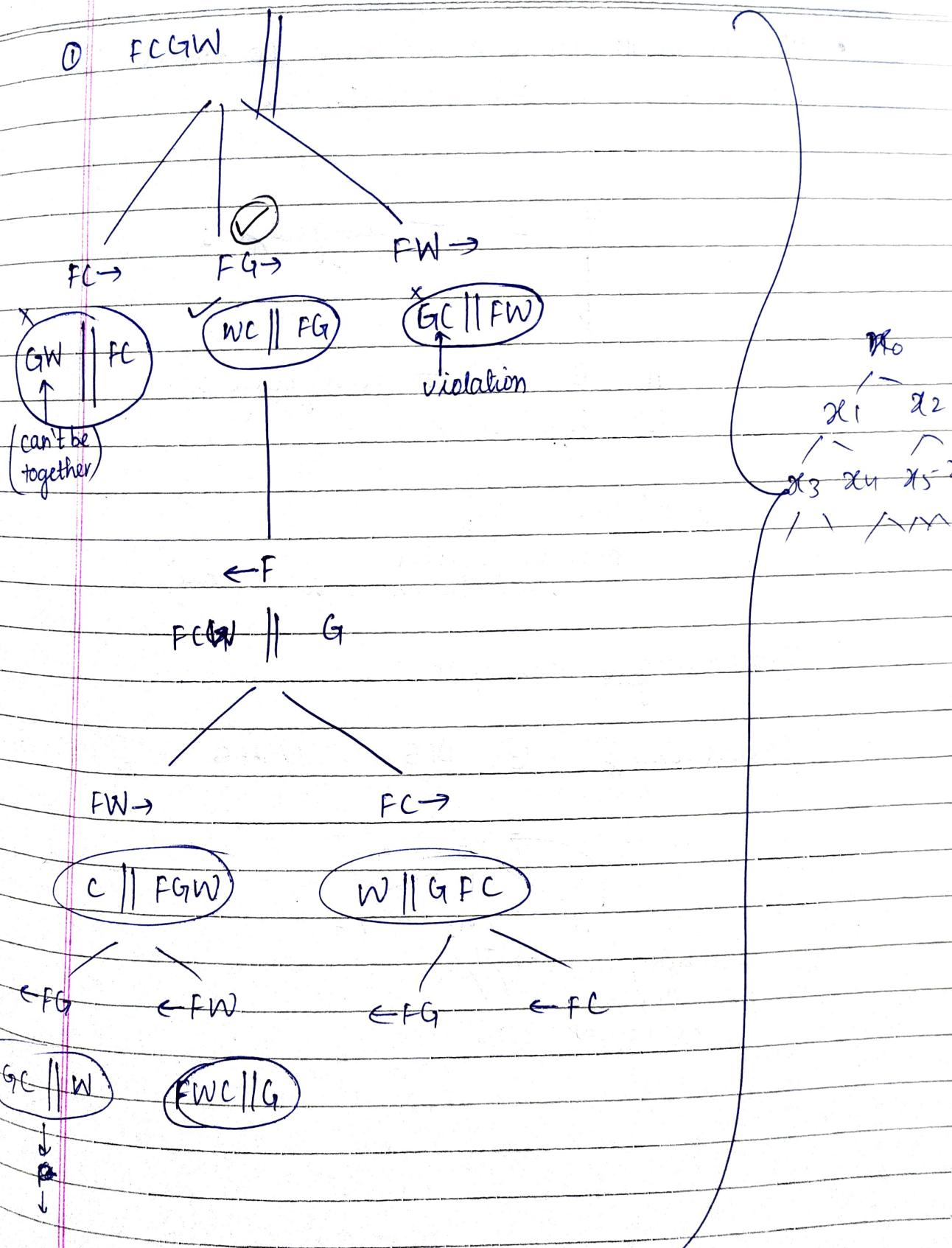
① Make it a search-based prob

- Make a graph/ tree
- Search space : set of Actions



Search Tree

→ Set of actions & possible states



- We need to train our machine to create these states & actions and then compute.

Eg:

X Y Z ← language 1
 ↓

A B C ← language 2
 (I) (have) (class)
 ↓

$A \{ a_1, \dots, a_n \}$
 pick the correct
 word (has/have/had/is/are)

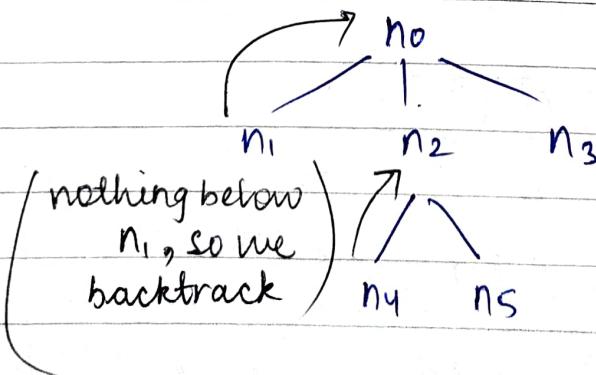
Traversing search space

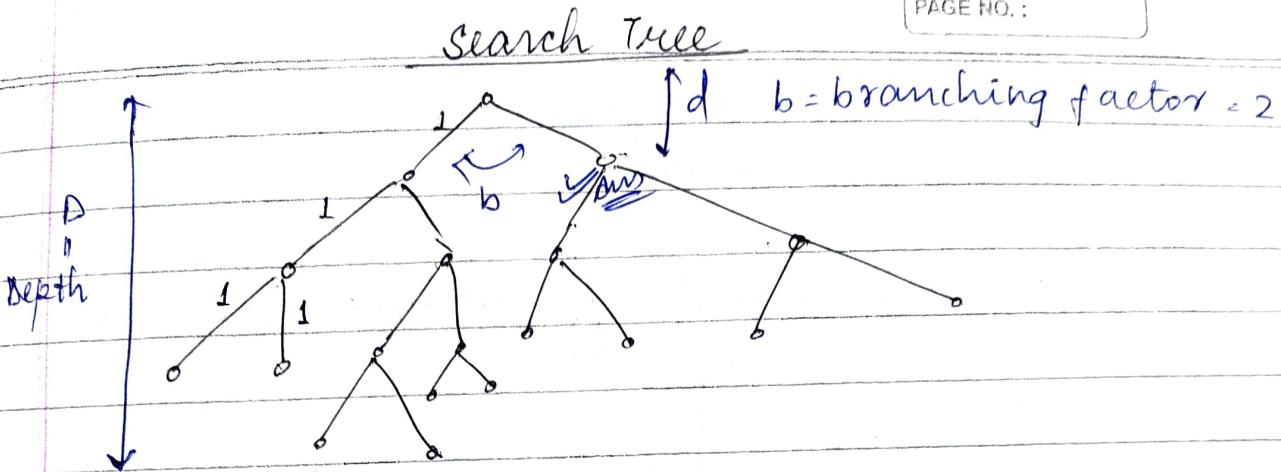
① Backtracking

② DFS

③ BFS

④ DFS-ID





Search Algo	cost	Time	Space
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① Backtrack	Homo/ Hetero (Any)	$O(b^D)$	$O(D)$
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Prob: traverse the whole tree in any case

- ② DFS : Backtracking but if we reach the colⁿ, we don't traverse the path further.
 ∴ Only in worst case it traverses whole tree

DFS	$c \geq 0$ (concerned by set of actions, not cost)	$O(b^D)$ (improved in avg/simple case)	$O(D)$
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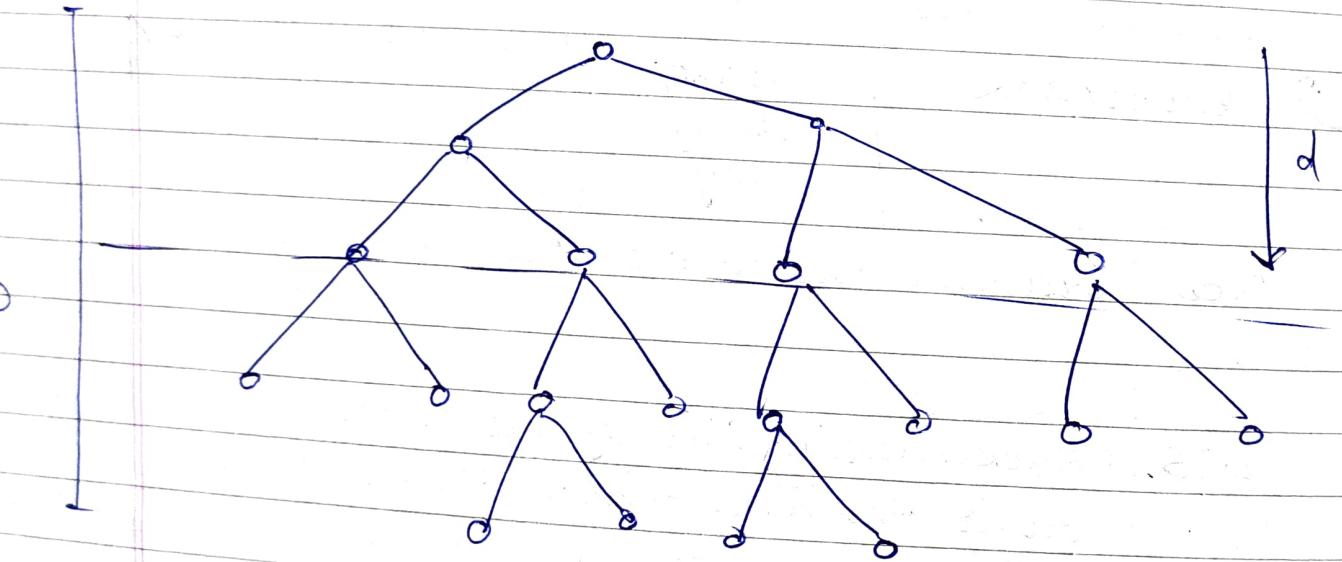
- ③ BFS : → level-by-level traversal.
 → stop if answer-node reached

BFS	$c \geq 0$	$O(b^d)$	$O(b^d)$ (keep track of every node)
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DFS → less time space
 BFS → less space

∴ we need a trade-off b/w time & space complexity

DFS - ID
 ↓
 Iterative Deepening

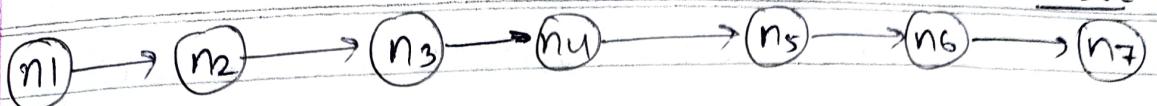


- First do DFS for d and hoping to get answer in this subtree
- If not found, we increment d by 1 and do BFS in the new layer and keep doing d and do BFS.

DFS-ID

$O(d)$

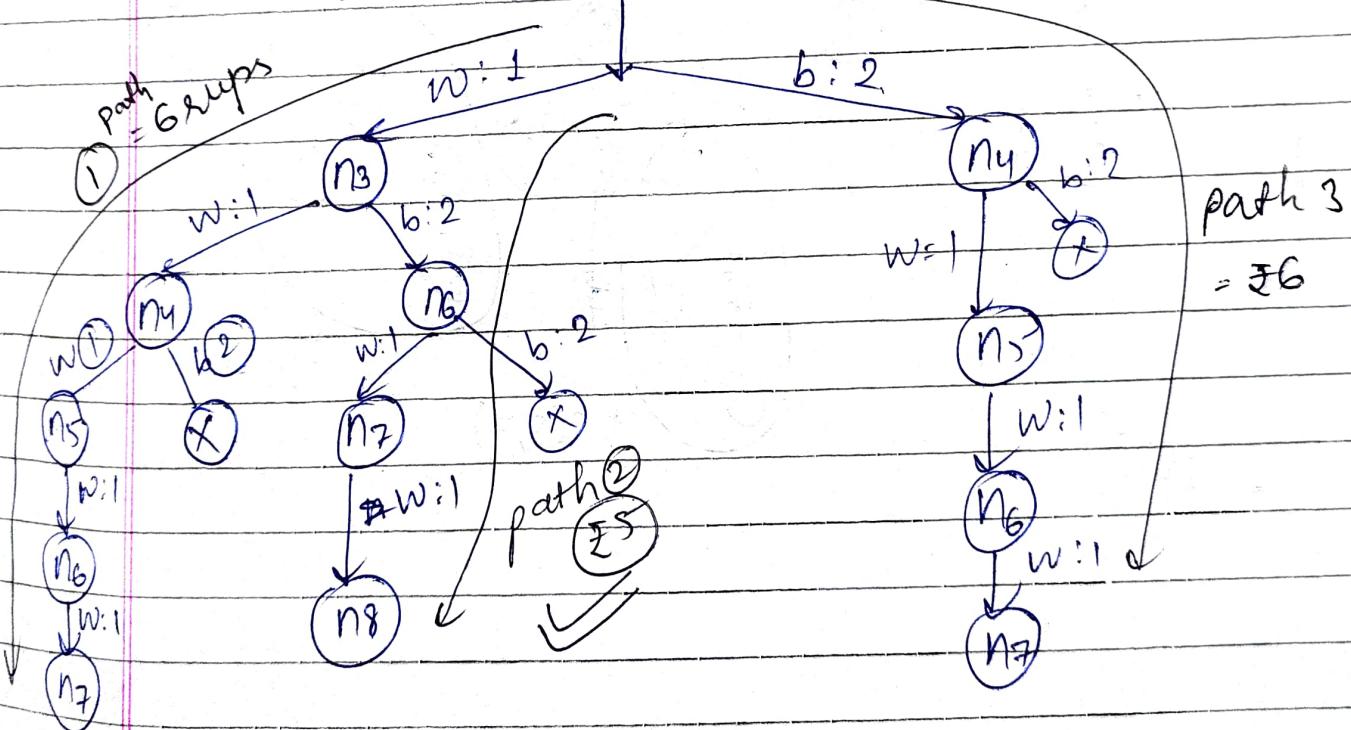
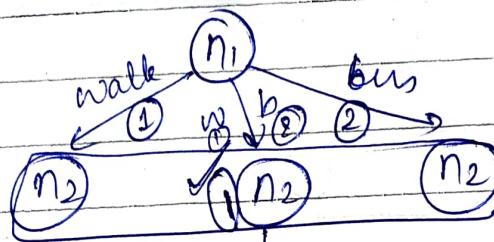
$O(d)$



walk : cost = £1 ($n \rightarrow n+1$)

bus : skips stoppages, cost = £2 ($n \rightarrow 2n$)

You can only move forward.



Ques: choose path with minimum cost.

path 2 is optimal

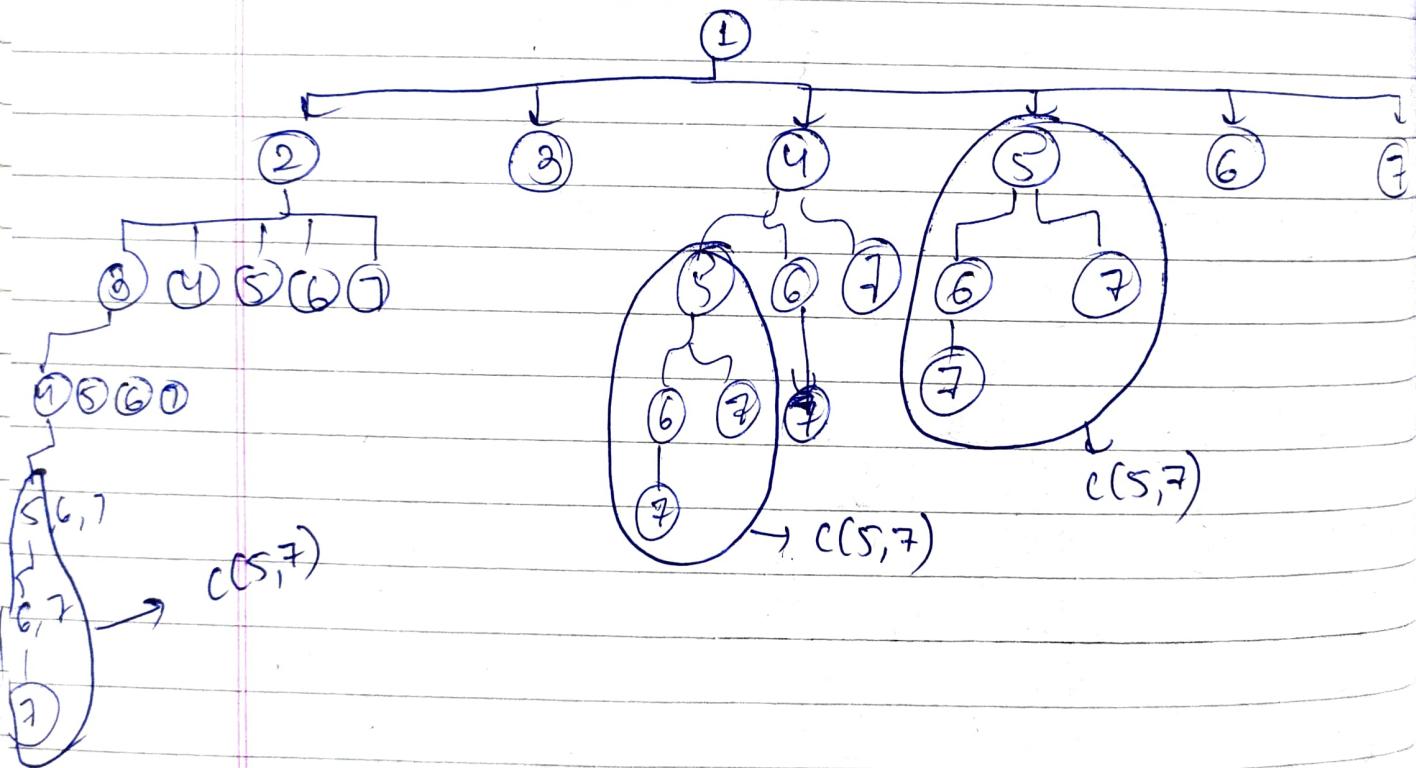
For exponential tree \rightarrow No improvement in complexity
coz we need to see all paths.

Soln: We can save cost from ny to n7 = £3
as it is used repeatedly.

We can save costs which are constant & reuse them

- This is Dynamic Programming

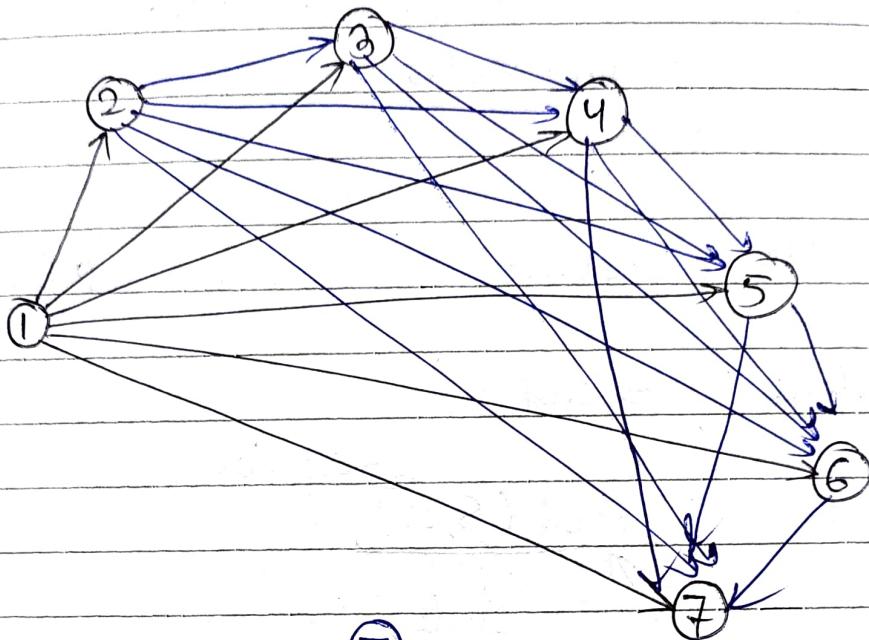
Eg: cost is $c(i, j) = \underline{\text{constant}}$
 $i \rightarrow j$



- Exponential search tree (2^7 nodes)

→ Reusing $c(5,7)$, and many more subtrees

Graph looks like :



No. of nodes here = 7

Time complexity = $O(n^2)$ (in worst case) $\frac{n(n+1)}{2}$

∴ complexity has gone down from exponential to polynomial

→ DP is not always n^2 , but definitely less than expo

Eg: city 1 → city 7

constraint → can't go to 3 consecutive odd no. city.

Concept of states

In search tree: states = all possibilities
graph: states

#

Start state: S_{start}

End state: S_{end}

cost: $\underbrace{\text{cost}(S, a)}_{\substack{\text{action} \\ \text{cost}}},$
 state

$a \in A$
 all possible set of actions

$\text{succ}(S, a) = S'$
 successor

{all states are subset of $e(S)$.}

Binary Func': IsEnd $\xrightarrow[=]{0} 1$

= true(1) or false(0)
 ending cond'.

Back to Example (odd consec.)

(prev city, current city) ← keep track.
 ↓
 1 or 2 3

if → (1, 3) → Now (1, 3), (3, 7) → Not possible
 (2, 3) → (2, 3), (3, 7) → possible

∴ current city = n options
 prev city = n-1 options

∴ complexity = $O(n^2)$ (for search space)
~~= $n(n+1)$~~

- But we don't really need $(n-1)$ options record.
- we only need to know if prev city is odd or even.

∴ prev city can have a flag for
odd & even

now complexity = $O(n)$ (search space
 $= 2n$)

Ques: Reach $1 \rightarrow 7$,

constraint: visit atleast 3 odd cities
before 7

Soln: States:

$\min(\text{cities visited}^{\text{odd}}, 3)$, current city

$O(n^2)$ complexity if keep track
of all prev. cities

But now $O(n)$

(coz $\min()$ returns
a constant)

Ques: Cost Minimization

$\min(\text{cost}(s, a), \text{future cost}(s))$

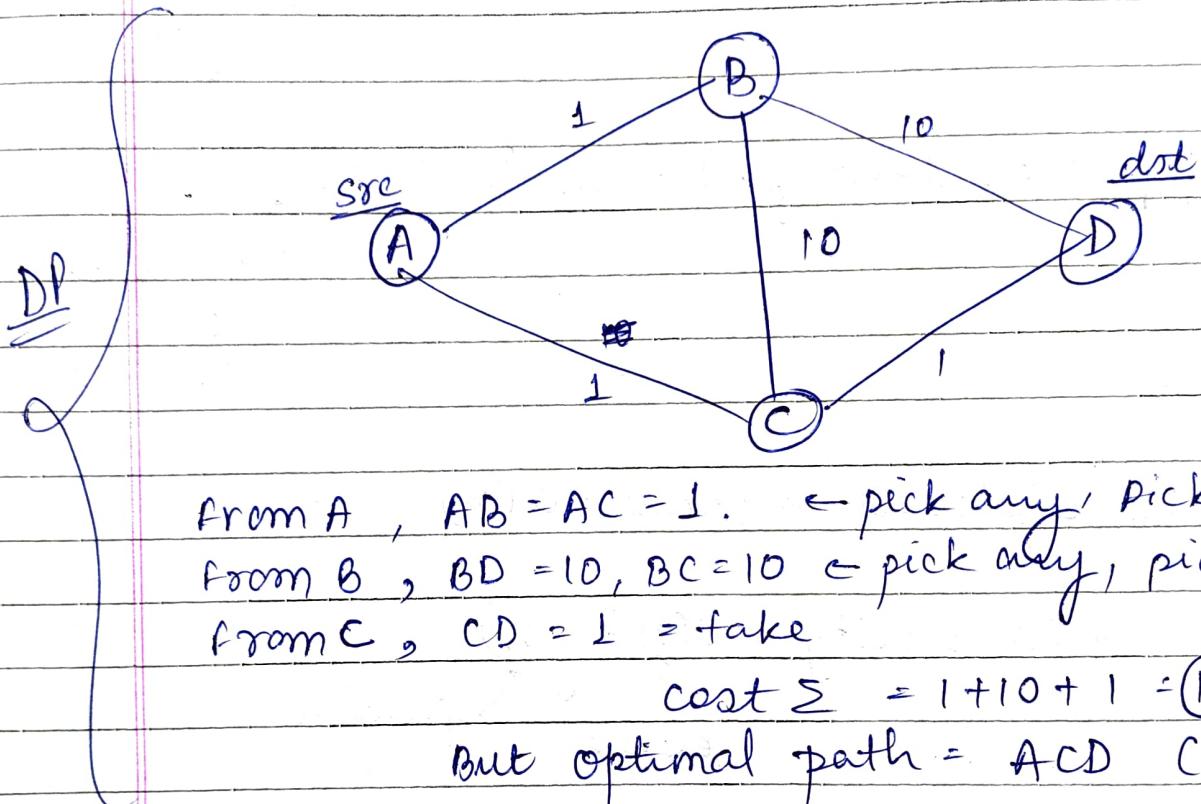
future cost =

O IsEnd.

- If graph is cyclic, ~~B~~ DP can't handle it
- For DP, graph should be acyclic, unidirectional
- For DP implementation, machine should have space for memorization. (In case space grows!)

Q: What if the graph is cyclic or bidirectional (can go back)

\therefore We use UCS. [Uniform Cost Search]



\therefore we need UCS

Keeps track of Explored, unexplored and frontier
 (priority list)

	unexplored	frontier priority	explored
A		(A, 0)	A, 0
B		(B, 1)	B, 1
C		(C, 10)	C, 10
D		(D, 10)	D, 10

C is already in ~~frontier~~ set, \therefore see B-D, if new value, update

$S_{start} = A$

- Update in frontier set,
- Don't change what is already in explored set.

Complexity

$$\frac{1}{n} + \frac{1}{n-1} + \frac{1}{n-2} + \dots - 1 = \log(n)$$

for frontier

If from A, (1 out of n), then (1 out of n-1)
 not check explored nodes

for explored \rightarrow space = $O(n)$

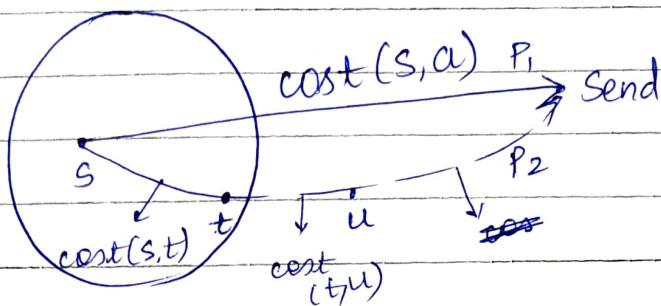
\therefore Complexity = $O(n \log n)$

explore

frontier

UCS

Unexplored \rightarrow Frontier \rightarrow Explored



$P_1 \rightarrow$ optimal
(Proof by contradiction)

Minimization Prob

Let P_2 be optimal (min. cost)

$$\begin{aligned}
 \text{cost}(s-t-u-Send) &\geq \text{cost}(s,t) + \text{cost}(t,u) \\
 &\geq \text{cost}(s,t) \quad \text{[coz } \text{Send} \text{ has more priority than } t \text{]} \\
 &\geq \text{cost}(s,a)
 \end{aligned}$$

* -ve cost is not allowed in UCS

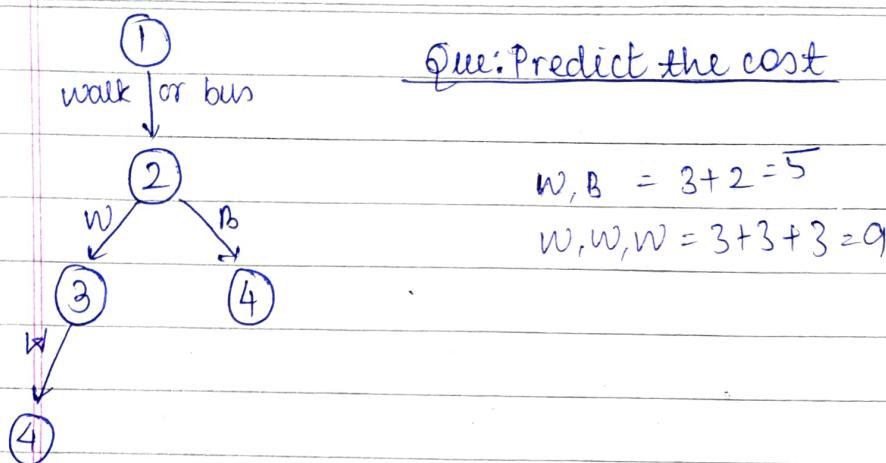
(b)

M (model)

I (inference)

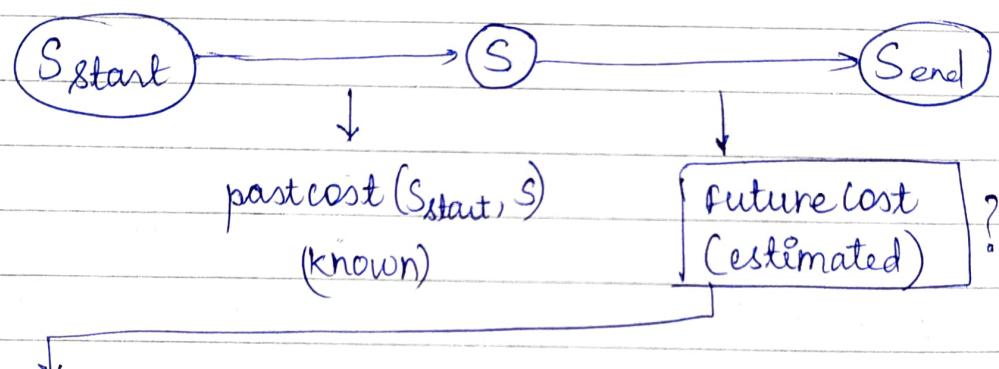
L (learning)

\rightarrow learning from previous data



$$W, B = 3 + 2 = 5$$

$$W, W, W = 3 + 3 + 3 = 9$$



Estimation func'

→ Heuristic \hat{h}

Prediction: cost (walk) = 3 (each step)

cost (bus) = 2 (each step)

Info provided / intuition / claim (from training data):

walk, walk, walk → optimal

My heuristic : walk, bus → optimal
(Prediction)

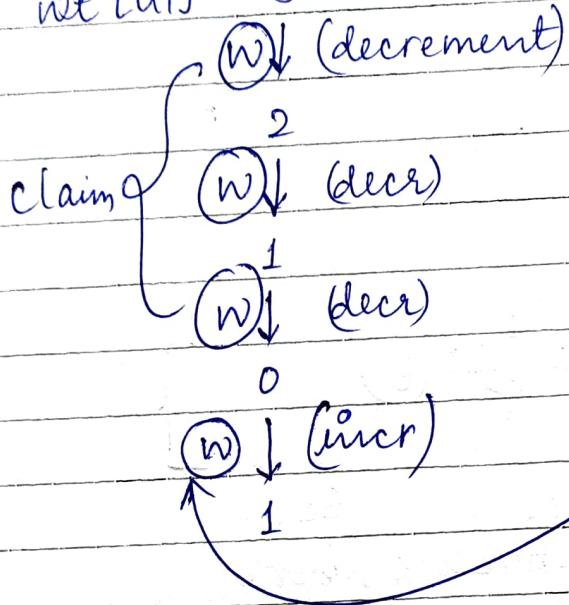
∴ with my prediction & prev. claim, if they match,
we get optimal.

~~weights~~ ~~action 1~~
~~wt [a₁], wt [a₂]~~

- We take a funcⁿ which ↑ or ↓ if we go by prediction or claim.

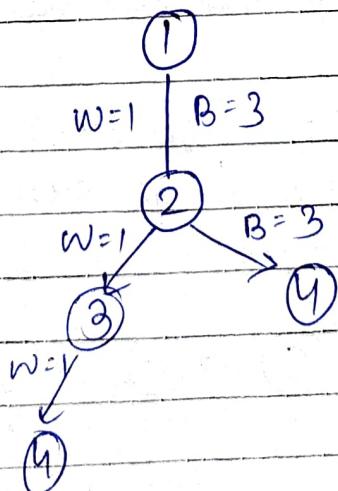
let $\text{wt } [a_1] = 3$ (walk)

$\text{wt } [a_2] = 2$ (bus)



$B \downarrow$ (incr)
 3

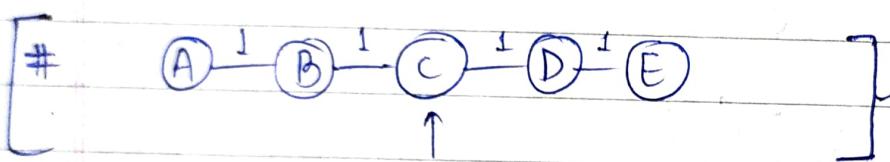
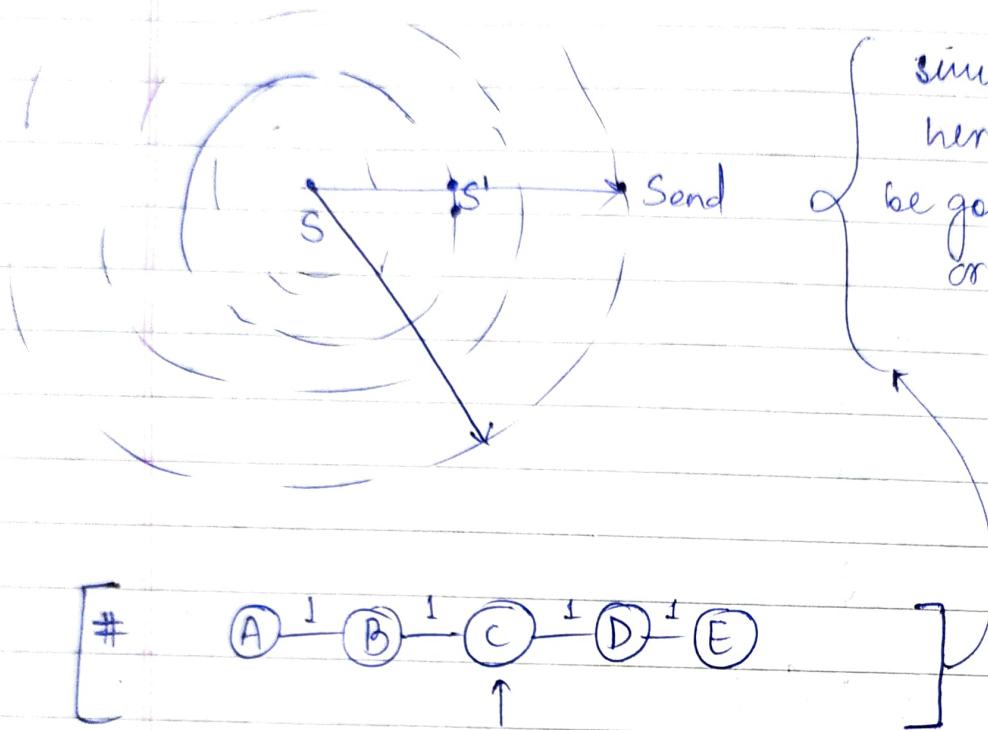
∴ Final values: $\text{wt } [\text{walk}] = 1$
 $\text{wt } [\text{bus}] = 3$



∴ Now optimal path = www ≠ prediction

∴ our prediction was wrong & we go
~~for~~ by training set.

∴ heuristic told us that we predicted wrong weights

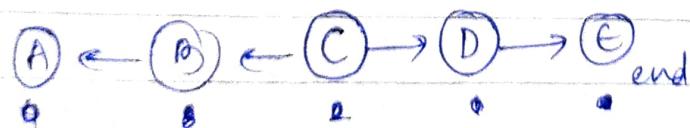
In UCS

From c, Send
idk where to go, B or D

* But what if we can predict that $C \rightarrow B$ is taking me away from Send and we need to go from C to D ??

Q: how do ensure we don't go back?

A: We define a heuristic (hop count)



$$h(S) = 4 \quad 3 \quad 2 \quad 1 \quad 0$$

$\text{cost}(S, a) \leftarrow$ not used

$$\text{cost}'(S, a) = \text{past cost}(c) + \underset{\substack{(known) \\ \downarrow}}{\text{heur.}(s')} - \text{heur}(s)$$

$\Rightarrow 1$

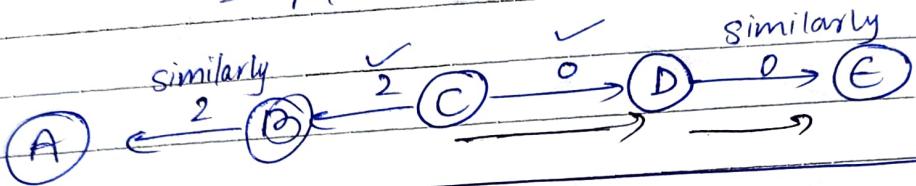
\uparrow
(given)

successor of s.
 $s' = \text{succ}(s, a)$

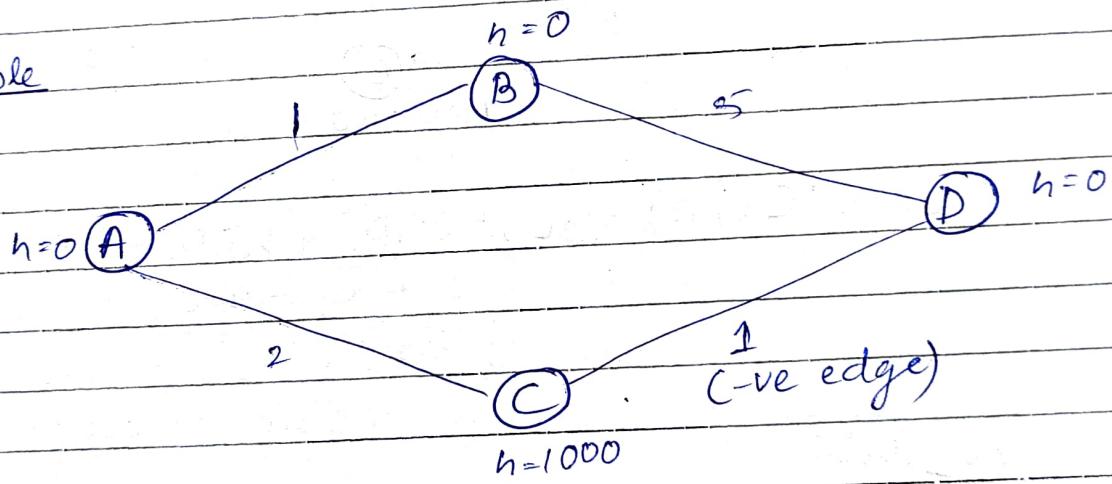
∴ here,

$$\begin{aligned}\text{cost}'(C, B) &= \text{past cost}(c) + h(B) - h(C) \\ &= 1 + 3 - 2 \\ &= 2\end{aligned}$$

$$\begin{aligned}\text{cost}'(C, D) &= \text{past cost}(c) + h(D) - h(C) \\ &= 1 + 1 - 2 = 0\end{aligned}$$



Example

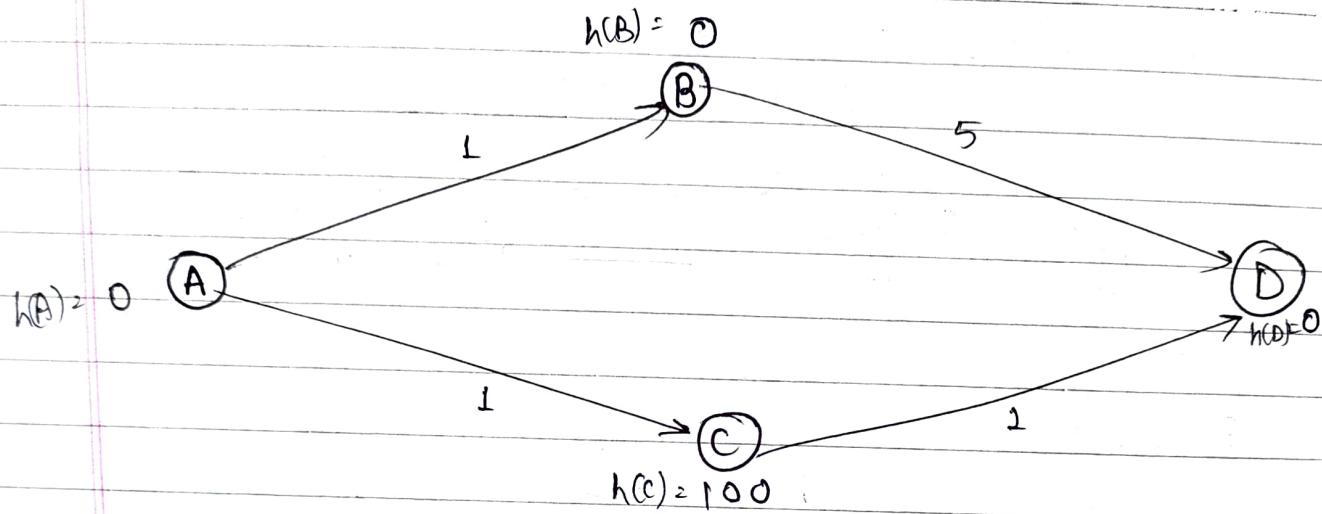


$$\text{cost}'(C, D) = 1 + (0 - 1000) = -999 \leftarrow \text{not allowed in UCS}$$

∴ can't apply 1 heuristic for all situations!

datasets, even if we are doing the same optimization in them.

- # • consistent h
- corrections
- greedy BFS
- CSP - COP



$$\begin{aligned}
 \text{cost}^*(C, C, D) &= \text{cost}(C, D) + h(D) - h(C) \\
 &= 1 + 0 - 100 \\
 &= -99
 \end{aligned}$$

-ve VCS not allowed (inconsistent)

heuristic

h is consistent iff it is always +ve.

∴ For it to be consistent,

$$\text{cost}^*(S_i, S_{i+1}) \geq 0$$

$$\therefore \text{cost}(C, D) + h(D) - h(C) \geq 0$$

$$\therefore \text{cost}(C, D) + h(D) \geq h(C)$$

$$\therefore \boxed{\text{past cost}(S) + h(S') \geq h(S)}$$

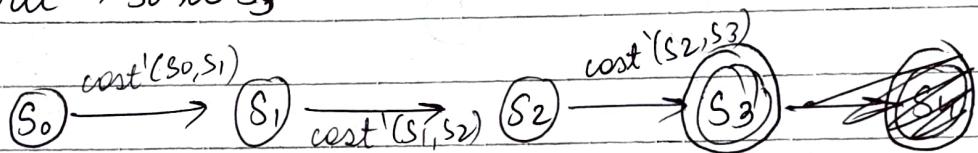
① heuristic value of end state = 0, always.

- when you reach end state / goal, no point in calculating h for further paths.

Q: How will you know path found is correct?

A:

Eg: Goal $\rightarrow S_0$ to S_g

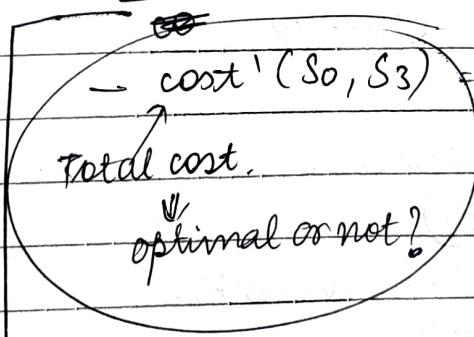


Path obtained: $S_0 - S_1 - S_2 - S_3 \xrightarrow{Sg}$

estimated cost (S_0, S_1)

or

$$\text{add } \left\{ \begin{array}{l} \text{cost}'(S_0, S_1) = \text{cost}(S_0, S_1) + h(S_1) - h(S_0) \\ \text{cost}'(S_1, S_2) = \text{cost}(S_1, S_2) + h(S_2) - h(S_1) \\ \text{cost}'(S_2, S_3) = \text{cost}(S_2, S_3) + h(S_3) - h(S_2) \end{array} \right.$$



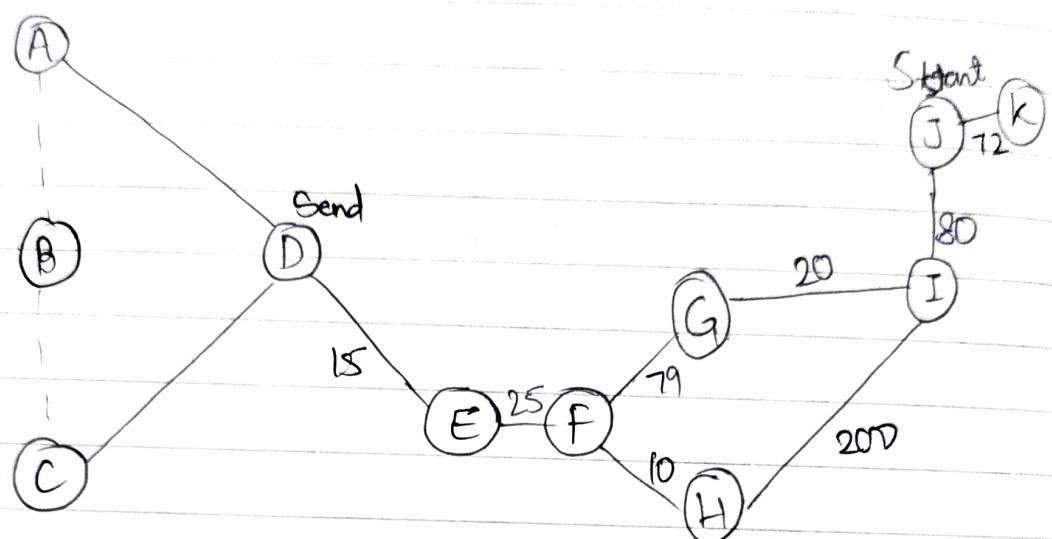
$$\sum_{i=1}^n \text{cost}'(S_{i-1}, S_i) = \sum_{i=1}^n \text{cost}(S_{i-1}, S_i) + h(S_n) - h(S_0)$$

↓ ↓ ↓
 - prev. path cost 0 constant
 - algorithmic consistency
 cost from UCS

- Since we already know that UCS gives optimal cost and we are decrementing some ^(the) constant from it,
 $\therefore \sum_{i=1}^n \text{cost}'(S_{i-1}, S_i)$ is the optimal one.

→ We can ~~only~~ do this: only estimate 'h' for next state.

Eg:



$$h(S) = \text{Euclidean dist of } (S, S')$$

heuristic defined as

(minimal)
(not using past cost)

Let,

$$\Rightarrow \text{Euc}(J-K) = d_1 \quad \& \quad \text{Euc}(J-I) = d_2$$

& $d_1 < d_2$ (known)

∴ We go from J to K.

⇒ From K, ~~there~~ only 1 option,

∴ We go from K to J

Now again heuristic takes us to K.

∴ We are stuck in $J \xrightarrow{\text{K}} J$ loop.

∴ Can't depend on heuristic alone

Reason: greedy of ~~but~~ taking simple heuristic
mode

This is the **Greedy BFS** approach

Greedy Best First Search

→ Best 1st search: coz we are prone to selecting the best ~~part~~ edge in current scenario.

∴ We modify it with $\underline{\text{pastcost}} + h(n)$

$$\text{Eq: } \underline{h(s^+)} - h(s) \\ \text{euc}(s_i \rightarrow s_{i+1})$$

- ① J → K → cost'(J, K) = 72 + h
- ② K → ~~J~~ J → cost'(K, J) = (72 + h) + 72 = 144 + h
- ③ J → K → cost'() = 144 + h X

$$\checkmark I \rightarrow \text{cost}'() = 80 + h \checkmark \text{(picked)}$$

→ Greedy BFS → not optimal

→ in tree, no option of loop

∴ Not complete for tree Greedy BFS
but complete for graph

Uninformed Search

→ ~~All~~ Brute force / blind search for goal / send

V8

Informed Search

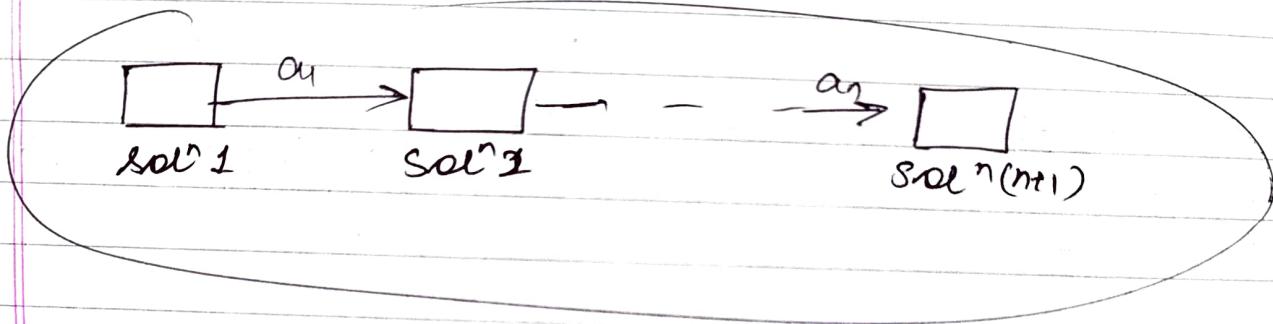
- Greedy BFS on heuristic
- extra info provided

Goal : 'a path' or 'a set of action' } so far
 Search space : states & actions }

Hill-climbing search

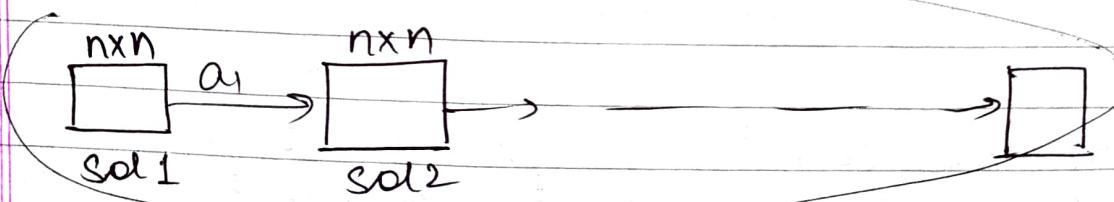
search space : states & actions

~~states~~ ~~actions~~ ~~states~~ ~~actions~~ ~~states~~ ~~actions~~ ~~states~~ ~~actions~~ ~~states~~ ~~actions~~

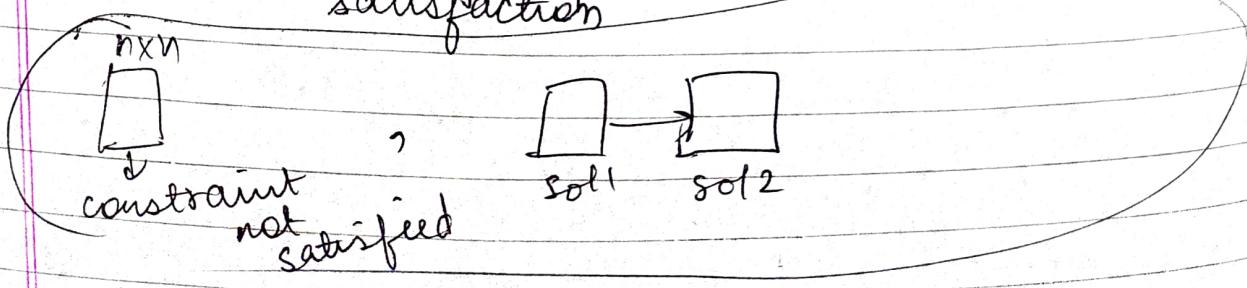


- Each sol^n is a state
- Goal states could be different but same sol space
- Each state is a sol space

N-queen problem



- different solution ⇒ different placements for N-queens
- No '1' goal state
- constraint-specific problems



(CSP)

flow?

(COP)

constraint-satisfaction $\xrightarrow{\text{to}}$ constraint - optimization
 for (N-queens) for (N-queens)

Q			
	Q	Q	
		Q	
S _{start}			

\leftarrow constraint not satisfied

- We have to find minimum moves to change this to ~~a~~ satisfactory placement
- That would be the optimal solution.

- ① Let after ' m ' no. of moves \rightarrow got ~~satisf~~ solⁿ
- ② If $m < m'$ \rightarrow got solⁿ

If $m < m' \rightarrow$ (ans 1 \rightarrow optimal)

Eg: CSP: ~~of~~ objective funcⁿ,
 subject to constraints $C_1, C_2 \dots C_n$.

Ques: If it satisfies 3 constraints, we reach our goal
 $\therefore nC_3$ options

COP - objective func^c & n constraints

Ques \Rightarrow Goal: the solⁿ which satisfies all n constraints

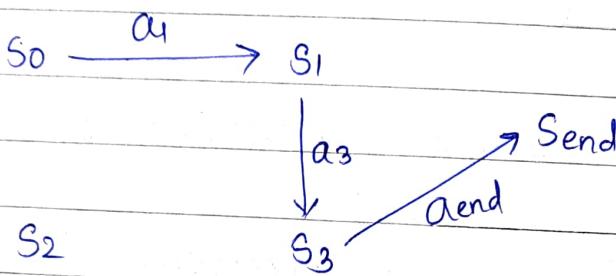
that will be the optimal solⁿ.

Q: COP is given, make it CSP.

A: ~~$\geq 3 \in \text{all sets}$~~

if optimality condⁿ = ≥ 3 constraints
= all solutions will
satisfy CSP.

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solution: a1a2a3aend.

(BUT)

Eg: In n-queen's prob : only 1 state is solⁿ, not a set of actions

In a local search,

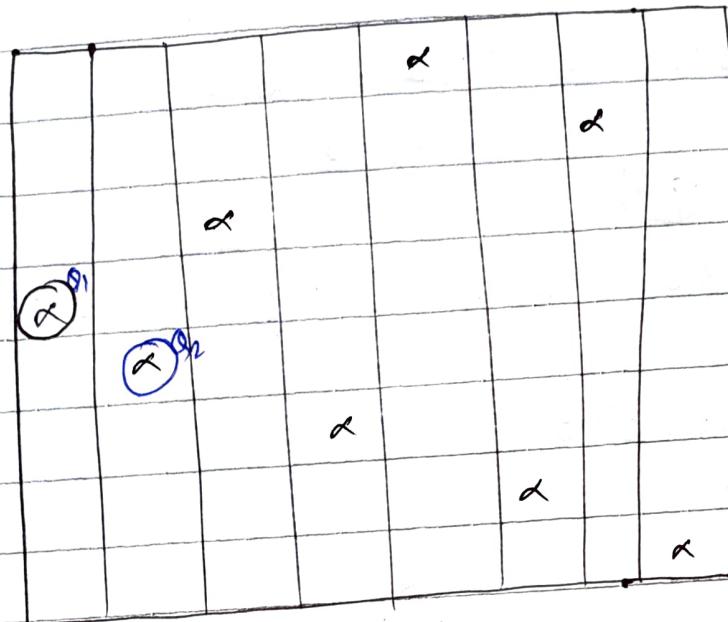
Eg: 8 queen prob.

→ check for neighbourhood nodes with some function.

Catch: How do you define neighbourhood func?

How many extra spaces you need to search?

determines if our neighbourhood funcⁿ is good or bad.

8-queens

Possible actions for all queens : a_1, a_2, a_3, a_4 .
 Neighbourhood funcⁿ? → how can you move taking every possible neighbour node.

case 1: can move anywhere in column, each column can have 1 queen atleast.

$$\rightarrow Q_1 = 7 \text{ moves } (\text{in its own column})$$

$$Q_2 = 7 \text{ moves}$$

$$\vdots$$

$$\vdots$$

$$Q_8 = 7 \text{ moves}$$

$$\left. \begin{array}{l} \\ \\ \\ \end{array} \right\} \therefore 8 \times 7 = 56 \text{ possible moves}$$

$$\rightarrow \text{If } n \text{ is increased } \rightarrow \boxed{\mathcal{O}(n^2)} \text{ moves.}$$

case 2: move only 1 step in adjacent node

Q_1 :

After every step, check if > 1 queen in a column
- Tiresome!!

Case 3: Move each queen anywhere on board.

: for $q_1 : (8 \times 8) - 8^8 + 1$ ~~huge no.~~

~~(8 × 8) - 1~~



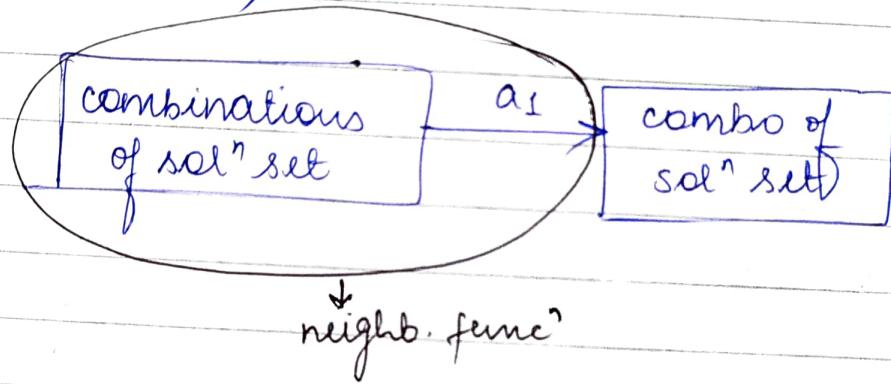
$$= (8 \times 8) - 1.$$

own place

: ~~$[8 \times (8 \times 8 - 1)]^8$~~ for all queen

Neighbourhood funcⁿ:

(NO successors)



→ local search : 1 action , possible states
 $\left\{ \begin{array}{l} \text{not a set} \\ \text{of actions} \end{array} \right\}$ (reachability)

How to choose neighbour?

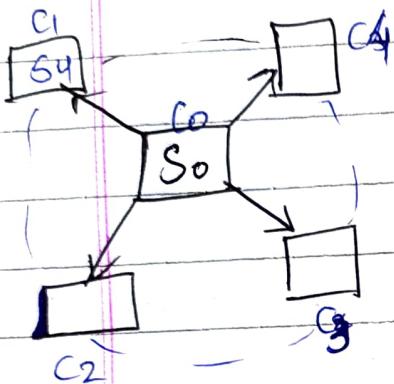
case 1

- Random sampling method: pick any 1 outta 56
- Random walk: maxⁿ similarity b/w neighbouring nodes.

: move towards most similar neighbour.

Hill Climbing Problem (HCP)

- 8-queens is also a HCP
- Greedy
- local, looking for best neighbour



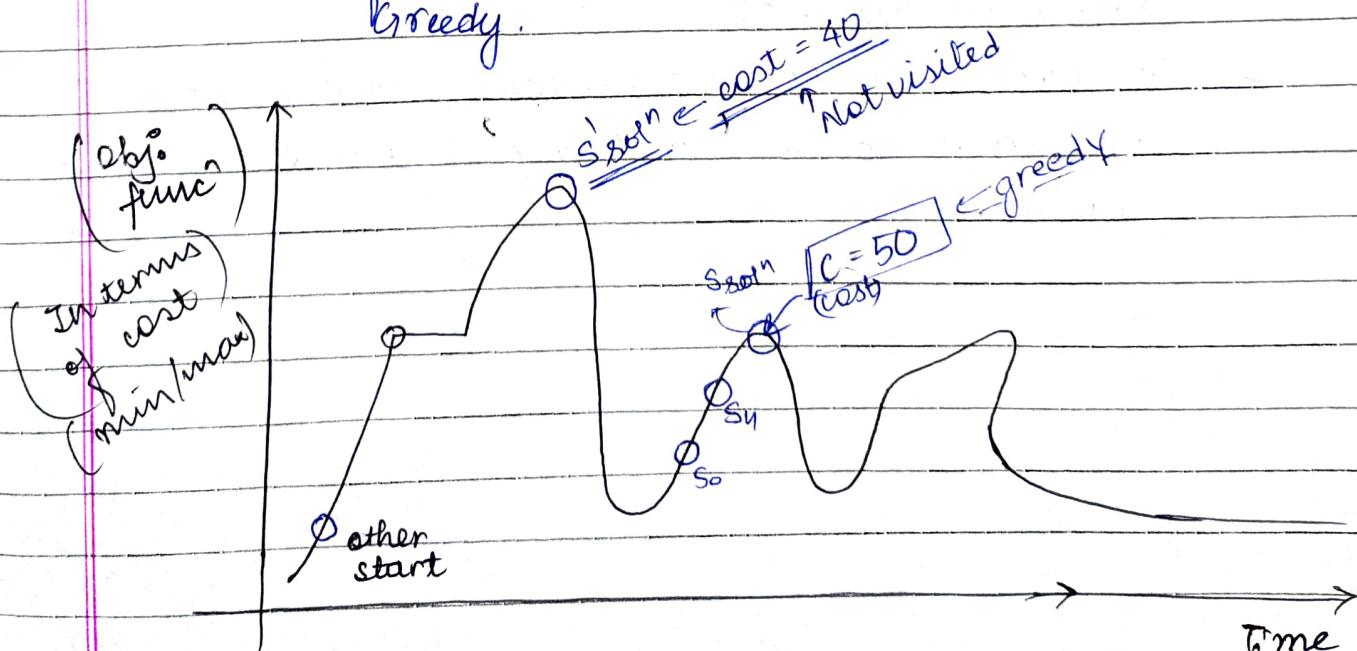
- Max/Min → optimisation
 - only 1 state will have max/min next move (assumed) ($>$) ($<$)
 - ∵ Others will be nullified.
- let it go from $S_0 \rightarrow S_4$

Is it the desired answer? optimal answer?

- $C_{sol^n} = 50$, say, but I haven't travelled a path with cost = 40.

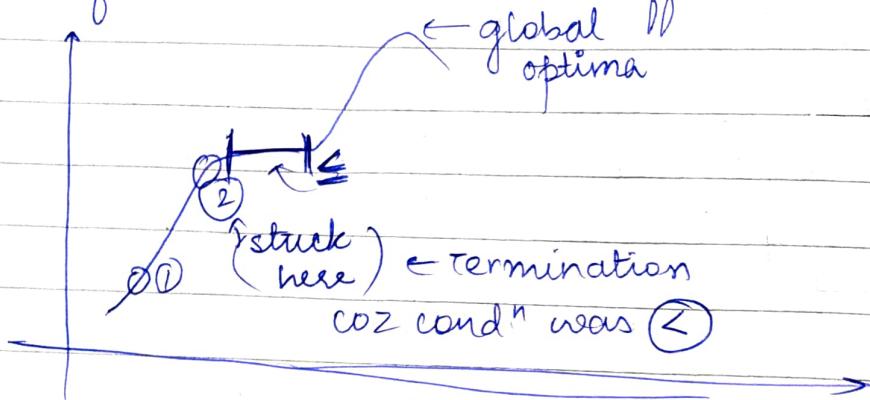
Drawback of HCP: Never reach global optimum.

Greedy.



We didn't visit cost 40 solⁿ coz neighbours of S0 were not optimal in the opposite dir.

But if the start state was different,



- But if cond^n is 0 ← now we can move on constant path.
- Now we are encouraging our machine to climb the hill.

If we put a stopping condⁿ, it drastically improves the success rate.

- * Read stats abt HCP.
- # To reach from local to global, you need to come back to the 'base camp' on the start point.