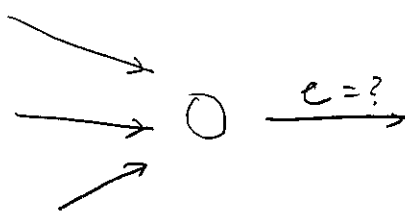


# Bufferless Statistical Multiplexing

Example:

1000 sources  
each one  
generates  
data at  
rate:

$\left. \begin{array}{l} 1 \text{ bps w.p } 0.5 \\ 0 \text{ bps w.p } 0.5 \end{array} \right\} \text{independent of each other}$



Goal: don't lose any packets.

$$\Rightarrow c = 1000 \text{ bps}$$

Goal: Allow small probability of overflow.

$X_i$  = rate at which source  $i$  generates data.

Find the maximum  $N$  such that

$$\mathbb{P}\left(\sum_{i=1}^N X_i > c\right) \leq \epsilon \quad \underline{X_i \text{ iid.}}$$

Chernoff bound  $c = Nc$  (ie. define  $c \triangleq c/N$ )

$$\begin{aligned}
 \mathbb{P}\left(\sum_{i=1}^N X_i \geq Nc\right) &= \mathbb{P}\left(e^{\theta \sum_{i=1}^N X_i} \geq e^{\theta Nc}\right) \quad \theta > 0 \\
 &\leq \frac{\mathbb{E}\left(e^{\theta \sum_{i=1}^N X_i}\right)}{e^{\theta Nc}}
 \end{aligned}$$

Define:

$$M_x(\theta) = \mathbb{E}(e^{\theta x}) \quad (\text{Moment generating function})$$

$$\Lambda_x(\theta) = \log M_x(\theta) \quad (\text{log moment generating fn.})$$

then we have,

$$\mathbb{P}\left(\sum_{i=1}^N X_i \geq Nc\right) \leq \frac{M_x^N(\theta)}{e^{\theta Nc}} \quad (\text{use independence})$$



$$= \frac{e^{N \log M_x(\theta)}}{e^{\theta Nc}}$$

$$= e^{-N(\theta c - \Lambda_x(\theta))}$$

Thus,

$$\mathbb{P}\left(\sum_{i=1}^N X_i \geq Nc\right) \leq e^{-N I(c)}$$

where  $I(c) = \sup_{\theta > 0} \{\theta c - \Lambda_x(\theta)\}.$



Lower Bound:

$$\max_{k \geq 1} P\left(\frac{1}{N} \sum_{i=1}^N A_i(k) \geq b + kc\right)$$

$$\stackrel{\sim}{\sim} \max_{k \geq 1} e^{-N I_A(b+kc)} = e^{-N \min_{k \geq 1} I_A(b+kc)}$$

where,

$I_A(\cdot)$  is the rate function of  $A_i(k)$

suppose  $k^* = \underset{k \geq 1}{\operatorname{argmin}} I_A(b+kc)$ . most likely time of overflow

Upper Bound:

$$\stackrel{\sim}{\sim} \sum_{k \geq 1} e^{-N I_A(b+kc)}$$

$$\stackrel{\sim}{\sim} e^{-N \min_{k \geq 1} I_A(b+kc)} \quad \text{for large } N$$

(one term dominates)

ie, under appropriate conditions,

$$\lim_{N \rightarrow \infty} \frac{1}{N} \ln(Q(0) \geq Nb) = \min_{k \geq 1} I_k(b+kc).$$

## Cramér's Theorem :

Let  $X_i$ 's iid; then

$$\lim_{N \rightarrow \infty} \frac{1}{N} \log \mathbb{P} \left( \frac{1}{N} \sum_{i=1}^N X_i \geq c \right) = - \sup_{\theta} (\theta c - \Lambda(\theta)) = - I(c)$$

## Assumptions :

- (i)  $c > \mathbb{E}(X_i) \triangleq \mu$
- (ii)  $\Lambda(\theta)$  is well defined in a neighborhood of  $\theta = 0$ .
- (iii)  $\exists \theta^*$  in this nbd such that
 
$$I(c) = \theta^* c - \Lambda(\theta^*)$$

Under these assumptions :

$$e^{-N(I(c) + \varepsilon)} \leq \mathbb{P} \left( \sum_{i=1}^N X_i \geq Nc \right) \leq e^{-N I(c)}$$

Given  $\varepsilon > 0$ ,  $\exists n : \forall N > n$  the above is true.

Upper bound :

$$\text{At } \theta = 0.$$

$$\theta c - \Lambda(\theta) = 0$$

If we show that

$$\theta c - \Lambda(\theta) \leq 0 \quad \forall \theta < 0,$$

then  $\theta \geq 0$  can be removed in the Chernoff bound.

Now, for  $\theta < 0$

$$e^{\theta(\mu - c)} \geq 1 \quad (\mu < c) \rightarrow \textcircled{1}$$

also,

$$\mathbb{E}(e^{\theta x_i}) \geq e^{\theta \mu} \quad (\mu = \mathbb{E}(x_i)) \rightarrow \textcircled{2}$$

$\uparrow$   
 Jensen's inequality.

$$\textcircled{1} \& \textcircled{2} \Rightarrow e^{\theta(\mu - c)} \mathbb{E}(e^{\theta x_i}) \geq e^{\theta \mu}$$

$$\Rightarrow -\theta c + \Lambda(\theta) \geq 0$$

$$\theta c - \Lambda(\theta) \leq 0$$

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Lower Bound :

$$e^{N I(c)} \mathbb{P}\left(\sum_{i=1}^N x_i \geq Nc\right) \geq e^{-N\varepsilon} \quad (\text{to show})$$

$$e^{N I(c)} \mathbb{P}\left(\sum_{i=1}^N x_i \geq Nc\right)$$

$$= \int_{\sum_{i=1}^N x_i \geq Nc} e^{N(\theta^* c - \Lambda(\theta^*))} d\mathbb{P}_1(x_1) d\mathbb{P}_2(x_2) \dots d\mathbb{P}_N(x_N)$$

$$\geq \int_{N(c+\delta) \geq \sum_{i=1}^N x_i \geq Nc} e^{N(\theta^* c - \Lambda(\theta^*))} d\mathbb{P}_1(x_1) d\mathbb{P}_2(x_2) \dots d\mathbb{P}_N(x_N)$$

$$= e^{-N\theta^* \delta} \int_{\substack{N(c+\delta) \geq \sum_{i=1}^N x_i \\ \geq Nc}} \frac{e^{N\theta^*(c+\delta)}}{(M(\theta^*))^N} d\mathbb{P}_1(x_1) \dots d\mathbb{P}_N(x_N)$$

Key step  
if  
 $\left(\begin{smallmatrix} c+\delta \\ \text{above is true} \end{smallmatrix}\right) \geq \sum x_i$

$$\geq e^{-N\theta^* \delta} \int \prod_{i=1}^N \frac{e^{\theta^* x_i}}{M(\theta^*)} d\mathbb{P}_i(x_i)$$

$$\text{Define } dQ_i = \frac{e^{\theta^* x_i}}{M(\theta^*)} d\mathbb{P}_i$$

$Y_i$  is a RV with cdf  $Q$ .

Aside: Show that  $E[Y_i] = c$

Proof.  $M_Y(\theta) = \int \frac{e^{\theta y} e^{\theta^* y}}{M(\theta^*)} dP$

$$= \frac{M(\theta + \theta^*)}{M(\theta^*)}$$

$$E(Y) = M'_Y(\theta) \Big|_{\theta=0} = \frac{M'(\theta^*)}{M(\theta^*)}$$

Since  $\theta^*$  maximizes  $\theta c - \log M(\theta)$

$$\Rightarrow c = \frac{M'(\theta^*)}{M(\theta^*)}$$

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$$e^{N\theta c} P\left(\sum_{i=1}^N X_i \geq Nc\right) \geq e^{-N\theta^* \delta} P\left(\sum_{i=1}^N (Y_i - c) \in [0, N\delta]\right)$$

$$= e^{-N\theta^* \delta} P\left(0 \leq \underbrace{\sum_{i=1}^N (Y_i - c)}_{\sqrt{N}} \leq \sqrt{N} \delta\right)$$

By central limit theorem, the above converges to  $\frac{1}{2}$

i.e., given  $\epsilon$ , choose  $N$  such that . such that .

$$\frac{1}{4} e^{-N\theta^* \delta} \geq e^{-N\epsilon}$$

# Loynes Formula

-  $a(k) = \#$  of packets arriving in time slot  $k$ .

$$- A(k) = \sum_{j=-k}^{-1} a(j)$$

Wish to estimate  $\mathbb{P}(Q(0) \geq b)$ .

Assume that  $\mathbb{E}(a(i)) < c$

We have,

$$\begin{aligned} Q(0) &= \max(Q(-1) + a(-1) - c, 0) \\ &= \max(Q(-1) + A(1) - c, 0) \end{aligned}$$

$$Q(-1) = \max(Q(-2) + a(-2) - c, 0)$$

$$\begin{aligned} \Rightarrow Q(0) &= \max(\max(Q(-2) + a(-2) - c, 0) + A(1) - c, 0) \\ &= \max(Q(-2) + A(2) - 2c, A(1) - c, 0) \end{aligned}$$

Assume that  $Q(\infty) = 0$ .

Then,

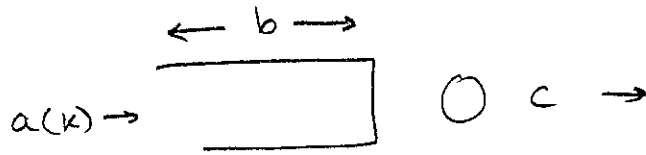
$$Q(0) = \max \left\{ \max_{k \geq 1} \{ A(k) - kc \}, 0 \right\}.$$

ie,

$$\mathbb{P}(Q(0) \geq b) = \mathbb{P}\left(\max_{k \geq 1} (A(k) - kc) \geq b\right).$$



SINGLE SOURCE LARGE BUFFER.



$$\lim_{b \rightarrow \infty} \frac{1}{b} \log \mathbb{P}(Q(0) \geq b) \leq -\min_{t \geq 0} t \mathbb{I}_{\infty}(c + 1/t)$$

Proof:

$$\mathbb{P}(Q(0) \geq b) = \mathbb{P} \left\{ \max_{k \geq 1} A(k) - kc \geq b \right\}$$

(Cayne's formula)

$$= \mathbb{P} \left\{ \max_{t \in \{1/b, 2/b, \dots\}} A(bt) \geq b + btc \right\}$$

$$= \mathbb{P} \left\{ \max_{t \in \{1/b, 2/b, \dots\}} A(bt) \geq b(1+tc) \right\}$$

$$\leq \sum_{t \in \{1/b, 2/b, \dots\}} \mathbb{P}(A(bt) \geq b(1+tc))$$

$$\leq \sum_{t \in \{1/b, 2/b, \dots\}} \frac{\mathbb{E}(e^{\theta A(bt)})}{e^{\theta b(1+tc)}}$$

Define,

$$\Lambda_{bt} = \frac{1}{bt} \log \mathbb{E} \left( e^{\theta \sum_{i=-bt}^{-1} a(i)} \right)$$

LHS

$$= \sum_{t \in \{1/b, 2/b, \dots\}} e^{-bt \left( \theta(c+1/t) + \Lambda_{bt}(\theta) \right)}$$

$$\leq \sum_{t \in \{1/b, \dots\}} e^{-bt \mathbb{I}_{bt}(c+1/t)}$$

$$\lim_{b \rightarrow \infty} \frac{1}{b} \log(\mathbb{P}(Q(0) > b)) \leq - \min_{t \geq 0} t \mathbb{I}_{\infty}(c+1/t)$$

□

Claim :

$$\min_{t \geq 0} t \, I_{\infty}(c + 1/t) \geq \varepsilon \left( \begin{array}{l} (\mathbb{P}(Q(0) > b) \\ \leq e^{-\varepsilon b}) \end{array} \right)$$



$$\frac{\Lambda(\varepsilon)}{\varepsilon} \leq c.$$

← effective bandwidths of the source.

Proof :

$$\inf_{t \geq 0} t \, I_{\infty}(c + 1/t) = \inf_{x \geq 0} \frac{1}{x} I(c+x)$$

$$\frac{1}{x} I_{\infty}(c+x) \geq \varepsilon \quad \forall x \quad (\text{by assumption})$$

$$\Rightarrow I_{\infty}(c+x) \geq \varepsilon x \quad \forall x$$

$$\Rightarrow \varepsilon x - I_{\infty}(c+x) \leq 0 \quad \forall x$$

$$\Rightarrow \varepsilon(c+x) - I_{\infty}(c+x) \leq c\varepsilon \quad \forall x$$

$$\Rightarrow \sup_{c+x} \varepsilon(c+x) - I_{\infty}(c+x) \leq c\varepsilon$$

$$\Rightarrow \frac{\Lambda(\varepsilon)}{\varepsilon} \leq c$$

Similarly in other direction.

# Properties of effective bandwidth.

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$$\frac{\Lambda(\varepsilon)}{\varepsilon} = \lim_{K \rightarrow \infty} \frac{1}{K\varepsilon} \ln \mathbb{E} \left( e^{\varepsilon \sum_{i=-K}^{-1} a(i)} \right)$$

$$\textcircled{1} \quad \lim_{\varepsilon \rightarrow 0} \frac{\Lambda(\varepsilon)}{\varepsilon} = \mathbb{E}(a(i)).$$

$$\text{"Proof"}: \quad \frac{1}{K\varepsilon} \log \mathbb{E} \left( e^{\varepsilon \sum_{i=1}^K a(i)} \right) = \frac{1}{K\varepsilon} \log \left( \mathbb{E} \left( 1 + \varepsilon \sum a(i) + \frac{\varepsilon^2 (\sum a(i))^2}{2!} + \dots \right) \right)$$

$$\sim \frac{1}{K\varepsilon} \log \mathbb{E} (1 + \varepsilon \sum a_i)$$

$$\sim \frac{\mathbb{E}(\sum a_i)}{K\varepsilon} = \mathbb{E}(a(i)).$$

② Suppose that

$$P = \inf \{ a : F(a) = 1 \} \quad \text{ie. } P = \text{peak rate}$$

$$\lim_{\theta \rightarrow \infty} \frac{\Lambda(\theta)}{\theta} = P.$$

Proof: Assume  $a = k_i$  w.p.  $\pi_i$

$$k_1 < k_2 < \dots < k_N = P$$

$$\frac{\mathbb{E}(e^{a\theta})}{e^{\theta P}} = \pi_N + \pi_{N-1} e^{\theta(\kappa_{N-1} - P)} + \dots + \pi_1 e^{\theta(\kappa_1 - P)} \quad (13)$$

For large  $\theta$ ,

$$\mathbb{E}[e^{a\theta}] \approx \pi_N e^{\theta P}$$

$$\Rightarrow \frac{\log \mathbb{E}[e^{a\theta}]}{\theta} = \frac{P\theta + \log \pi_N}{\theta}$$

$$\longrightarrow P$$

③  $\frac{\Lambda(\theta)}{\theta}$  is non-decreasing in  $\theta$ .

$$f(\theta) = \frac{\Lambda(\theta)}{\theta}, \quad a = \sum_{i=1}^K a_i$$

$$\frac{f(\theta_2)}{f(\theta_1)} = \frac{\ln \mathbb{E}[e^{\theta_2 a}] \theta_1}{\ln \mathbb{E}[e^{\theta_1 a}] \theta_2}$$

$$= \frac{\ln \left( \mathbb{E}[e^{\theta_2 a}] \right)^{\theta_1/\theta_2}}{\ln \mathbb{E}[e^{\theta_1 a}]}$$

If  $\theta_2 > \theta_1$ , we want to show that

$$\frac{f(\theta_2)}{f(\theta_1)} \geq 1$$

Now, if  $\theta_2 > \theta_1$ , then

$x^{\theta_1/\theta_2}$  is concave for  $x \in \mathbb{R}$ .

Let  $y = e^{\theta_2 a}$ , then

$$\begin{aligned} \frac{\ln [(\mathbb{E}[y])^{\theta_1/\theta_2}]}{\ln \mathbb{E}[e^{\theta_1 a}]} &\geq \frac{\ln \mathbb{E}[y^{\theta_1/\theta_2}]}{\ln \mathbb{E}[e^{\theta_1 a}]} \\ &= \frac{\ln \mathbb{E}[e^{\theta_2 a \theta_1/\theta_2}]}{\ln \mathbb{E}[e^{\theta_1 a}]} \\ &= 1 \end{aligned}$$

$$I(x) = \sup_{\theta} \theta x - \Lambda(\theta).$$

①  $I(x)$  is convex

Proof.

$$I(\alpha x_1 + (1-\alpha)x_2)$$

$$= \sup_{\theta} \theta (\alpha x_1 + (1-\alpha)x_2) - (\alpha \Lambda(\theta) + (1-\alpha) \Lambda(\theta))$$

$$\leq \alpha \sup_{\theta} \theta x_1 - \Lambda(\theta) + (1-\alpha) \sup_{\theta} \theta x_2 - \Lambda(\theta).$$

$$= \alpha I(x_1) + (1-\alpha) I(x_2)$$

②  $I(x) = \sup_{\theta} \theta x - \Lambda(\theta)$

$\Updownarrow$

$$\Lambda(\theta) = \sup_x \theta x - I(x)$$

$\Lambda(\theta)$  is a convex transform.

Proof:

$$I(x) = \sup_{\theta} \theta x - \Lambda(\theta)$$

$$\frac{dI(x)}{dx} = x - \Lambda'(\theta)$$

i.e.  $I(x) = \theta_x x - \Lambda(\theta_x) \longrightarrow \textcircled{1}$

∴  $\boxed{\Lambda'(\theta_x) = x} \longrightarrow \textcircled{2}$

$$\frac{dI(x)}{dx} = x \frac{d\theta_x}{dx} + \theta_x - \Lambda'(\theta_x) \frac{d\theta_x}{dx}$$

From (2),  $\Lambda'(\theta_x) = x$

$$\Rightarrow \frac{dI(x)}{dx} = \theta_x \quad \text{or} \quad \boxed{I'(x) = \theta_x} \rightarrow (2)$$

Now, we have to show,

$$\sup_x \theta x - I(x) = \Lambda(\theta)$$

LHS

To find sup. diff w.r.t  $x$ .

$$\theta - I'(x_0) = 0 \Rightarrow I'(x_0) = \theta \rightarrow (4)$$

and,

$$\text{LHS} = \theta x_0 - I(x_0) \rightarrow (5)$$

$$= \theta x_0 - (\theta x_0 - \Lambda(\theta_{x_0})) \quad (\text{from (1)})$$

$$\left. \begin{array}{l} \text{From (3), } \theta_{x_0} = I'(x_0) \\ \text{From (4) } I'(x_0) = \theta \end{array} \right\} \Rightarrow \theta_{x_0} = \theta.$$

$$\text{Thus, LHS } \theta x_0 - \theta x_0 - \Lambda(\theta) = -\Lambda(\theta).$$