

Problem Set 5

Out: February 21

Due: February 28

Reminder: you are encouraged to work in groups of two or three; however you must turn in your own write-up and note with whom you worked. You may consult the course notes and the text (Sipser). The full honor code guidelines can be found in the course syllabus.

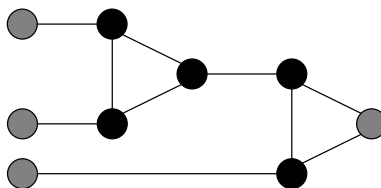
Please attempt all problems. **To facilitate grading, please turn in each problem on a separate sheet of paper and put your name on each sheet. Do not staple the separate sheets.**

- State whether each of the following statements is true or false, and briefly justify your answer. In this problem the reductions are always polynomial time many-one (or “mapping”) reductions.
 - Every language in EXP has algorithms that run in time $O(2^{n^{100}})$.
 - Either $P \neq NP$ or $NP \neq EXP$ (or both).
 - If for some fixed k , an NP-complete language has an $O(n^k)$ -time algorithm, then every language in NP has an $O(n^k)$ -time algorithm.
 - If $P = NP$, then every language in NP is NP-complete.
 - Every language in NP reduces to every EXP-complete language.
- Recall that a graph G is called k -colorable if there is a way to assign a color to each vertex so that no edge has both endpoints assigned the same color, using at most k distinct colors. On the last problem set, you showed that 2-COLORABLE was in P.

Show that the following language is NP-complete:

$$3\text{-COLORABLE} = \{G : G \text{ is 3-colorable}\}.$$

Hint: reduce from 3-SAT. Your graph will contain 1 vertex for each literal, and 3 special vertices connected in a triangle (which must then be colored with the three distinct colors). You may find this observation useful: in the following graph,



if each of the grey nodes are colored with one of two colors, then it is possible to extend this coloring to a 3-coloring if and only if at least one of the three grey nodes on the left has the same color as the one on the right.

3. Let $(3, 3)$ -SAT be the language consisting of satisfiable CNF formulas with at most 3 literals per clause, *and* at most 3 occurrences of any variable. Show that $(3, 3)$ -SAT is NP-complete.
4. MAX2SAT is the language consisting of all pairs (ϕ, k) where ϕ is a 2-CNF formula for which it is possible to simultaneously satisfy at least k clauses. Show that MAX2SAT is NP-complete. Hint: how many of the following clauses can be satisfied as a function of x, y and z ?

$$\begin{aligned} &(x \vee x), (y \vee y), (z \vee z), (w \vee w), \\ &(\neg x \vee \neg y), (\neg y \vee \neg z), (\neg z \vee \neg x), \\ &(x \vee \neg w), (y \vee \neg w), (z \vee \neg w) \end{aligned}$$