CS21
Decidability and Tractability

Lecture 12

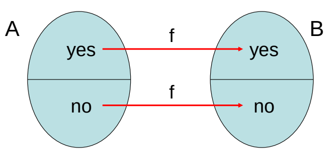
February 3, 2014

Outline

- undecidable problems
  - computation histories
  - surprising contrasts between decidable/undecidable
- Rice's Theorem
- Post Correspondence problem

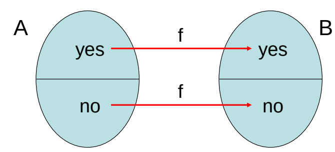
Definition of reduction

- More refined notion of reduction:
  - "many-one" reduction (commonly)
  - "mapping" reduction (book)



reduction from language A to language B

Definition of reduction



- function f should be computable

Definition: f : $\Sigma^* \rightarrow \Sigma^*$ is computable if there exists a TM M_f such that on every $w \in \Sigma^*$ M_f halts on w with f(w) written on its tape.

Definition of reduction

- Notation: "A many-one reduces to B" is written

$$A \leq_m B$$

  - "yes maps to yes and no maps to no" means:

$w \in A$ maps to $f(w) \in B$ & $w \notin A$ maps to $f(w) \notin B$

- B is at least as "hard" as A
  - more accurate: B at least as "expressive" as A

Using reductions

Definition: $A \leq_m B$ if there is a computable function f such that for all w

$$w \in A \Leftrightarrow f(w) \in B$$

Theorem: if $A \leq_m B$ and B is decidable then A is decidable

Proof:
- decider for A: on input w, compute f(w), run decider for B, do whatever it does.

Using reductions

- Main use: given language NEW, prove it is undecidable by showing OLD $\leq_m$ NEW, where OLD known to be undecidable
  - proof by contradiction
  - if NEW decidable, then OLD decidable
  - OLD undecidable. Contradiction.
- common to reduce in wrong direction.
- review this argument to check yourself.

Using reductions

Theorem: if A $\leq_m$ B and B is RE then A is RE

Proof:

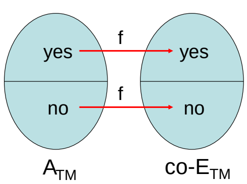
- TM for recognizing A: on input w, compute f(w), run TM that recognizes B, do whatever it does.

- Main use: given language NEW, prove it is not RE by showing OLD $\leq_m$ NEW, where OLD known to be not RE.

Many-one reduction example

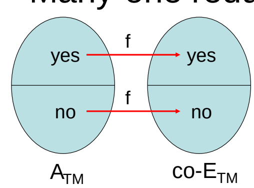
- Showed E_{TM} undecidable. Consider:

  co-E_{TM} = {<M> : L(M) $\neq \emptyset$}



- f(<M, w>) = <M'> where M' is TM that
  - on input x, if x $\neq$ w, then reject
  - else simulate M on x, and accept if M does
- f clearly computable

Many-one reduction example



- f(<M, w>) = <M'> where M' is TM that
  - on input x, if x $\neq$ w, then reject
  - else simulate M on x, and accept if M does
- f clearly computable

- yes maps to yes?
  - if <M, w> $\in A_{TM}$ then f(M, w) $\in$ co-E_{TM}
- no maps to no?
  - if <M, w> $\notin A_{TM}$ then f(M, w) $\notin$ co-E_{TM}

Undecidable problems

Theorem: The language

REGULAR = {<M>: M is a TM and L(M) is regular}

is undecidable.

Proof:

- reduce from A_{TM} (i.e. show $A_{TM} \leq_m$ REGULAR)
- what should f(<M, w>) produce?

Undecidable problems

Proof:

- f(<M, w>) = <M'> described below

on input x:
- if x has form 0^n1^n, accept
- else simulate M on w and accept x if M accepts

- is f computable?
- YES maps to YES?

  <M, w> $\in A_{TM} \Rightarrow$ f(M, w) $\in$ REGULAR
- NO maps to NO?

  <M, w> $\notin A_{TM} \Rightarrow$ f(M, w) $\notin$ REGULAR

Dec. and undec. problems

- the boundary between decidability and undecidability is often quite delicate
  - seemingly related problems
  - one decidable
  - other undecidable
- We will see two examples of this phenomenon next.

Computation histories

- Recall configuration of a TM: string uqv with $u,v \in \Gamma^*$, $q \in Q$
- The sequence of configurations M goes through on input w is a computation history of M on input w
  - may be *accepting*, or *rejecting*
  - reserve the term for halting computations
  - nondeterministic machines may have several computation histories for a given input.

Linear Bounded Automata

LBA definition: TM that is prohibited from moving head off right side of input.
- machine prevents such a move, just like a TM prevents a move off left of tape

- How many possible configurations for a LBA M on input w with |w| = n, m states, and p = |Γ| ?
  - counting gives: mnp^n

Dec. and undec. problems

- two problems we have seen with respect to TMs, now regarding LBAs:
  - LBA acceptance:
    A_{LBA} = {<M, w> : LBA M accepts input w}
  - LBA emptiness:
    E_{LBA} = {<M> : LBA M has L(M) = Ø}
- Both decidable? both undecidable? one decidable?

Dec. and undec. problems

<u>Theorem</u>: A_{LBA} is decidable.

Proof:
- input <M, w> where M is a LBA
- key: only mnp^n configurations
- if M hasn't halted after this many steps, it must be looping forever.
- simulate M for mnp^n steps
- if it halts, accept or reject accordingly,
- else reject since it must be looping

Dec. and undec. problems

<u>Theorem</u>: E_{LBA} is undecidable.

Proof:
- reduce from co-A_{TM} (i.e. show co-$A_{TM} \leq_m E_{LBA}$)
- what should f(<M, w>) produce?
- Idea:
  - produce LBA B that accepts exactly the accepting computation histories of M on input w

Dec. and undec. problems

Proof:

– f(<M, w>) = <B> described below

on input x, check if x has form

$\#C_1\#C_2\#C_3\#\ldots\#C_k\#$

- check that C_1 is the start configuration for M on input w
- check that $C_i \Rightarrow^1 C_{i+1}$
- check that C_k is an accepting configuration for M

- is B an LBA?
- is f computable?
- YES maps to YES?

  $\langle M, w\rangle \in \text{co-}A_{TM} \Rightarrow f(M, w) \in E_{LBA}$
- NO maps to NO?

  $\langle M, w\rangle \notin \text{co-}A_{TM} \Rightarrow f(M, w) \notin E_{LBA}$