



Page 60 and 63 ;

**2.1.2.** Describe informally the languages accepted by the deterministic finite automata shown in the next page.

**ANSWER:**

- (a)  $a(ba)^*$ . Strings with the prefix "a" followed by "ba" or "empty".
- (b)  $a^*b$ . Strings that can be with prefix "a" or "empty" and contain "b" as "a" substring.
- (c)  $(a(ab)^*b)^*$ . Empty or strings that can contain sequentially infinite "ab".
- (d)  $(ab \cup ba)^*$ . Strings containing ab or ba but not more than 1 consecutive same expression.
- (e)  $(aUb)^*(aabUbba)(aUb)^*$  All strings containing aab and bba as substrings.

**2.1.3.** Construct deterministic finite automata accepting each of the following languages.

- (a)  $\{w \in \{a, b\}^* : \text{each } a \text{ in } w \text{ is immediately preceded by a } b\}$ .
- (b)  $\{w \in \{a, b\}^* : w \text{ has } abab \text{ as a substring}\}$ .
- (c)  $\{w \in \{a, b\}^* : w \text{ has neither } aa \text{ nor } bb \text{ as a substring}\}$ .
- (d)  $\{w \in \{a, b\}^* : w \text{ has an odd number of } a\text{'s and an even number of } b\text{'s}\}$ .
- (e)  $\{w \in \{a, b\}^* : w \text{ has both } ab \text{ and } ba \text{ as substrings}\}$ .

**ANSWER:**

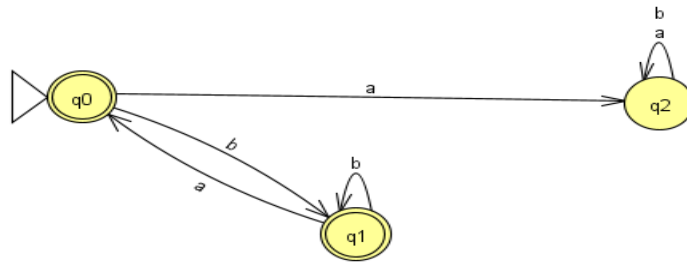
(a)

$$K = \{q_0, q_1, q_2\}$$

$$\Sigma = \{a, b\}$$

$$s = q_0$$

q	a	$\delta(q, a)$
$q_0$	a	$q_2$
$q_0$	b	$q_1$
$q_1$	a	$q_0$
$q_1$	b	$q_1$
$q_2$	a	$q_2$
$q_2$	b	$q_2$



$F = \{q_0, q_1\}$

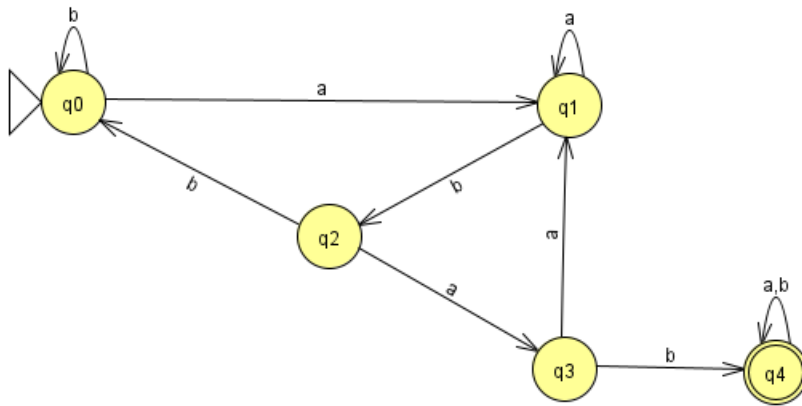
**(b)**

$K = \{q_0, q_1, q_2, q_3, q_4\}$

$\Sigma = \{a, b\}$

$s = q_0$

$F = q_4$



q	a	$\delta(q,a)$
q0	a	q1
q0	b	q0
q1	a	q1
q1	b	q2
q2	a	q3
q2	b	q0
q3	a	q1
q3	b	q4
q4	a	q4
q4	b	q4

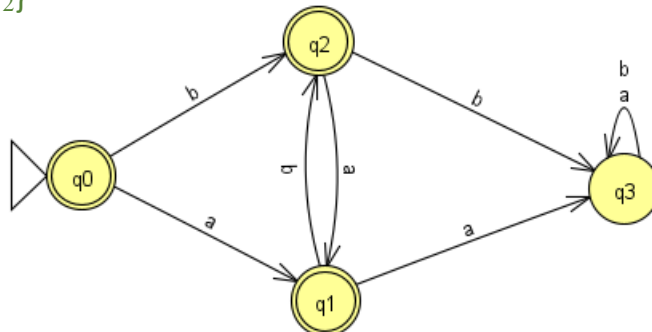
**(c)**

$K = \{q_0, q_1, q_2, q_3\}$

$\Sigma = \{a, b\}$

$s = q_0$

$F = \{q_0, q_1, q_2\}$



q	a	$\delta(q,a)$
q0	a	q1
q0	b	q2
q1	a	q3
q1	b	q2
q2	a	q1
q2	b	q3
q3	a	q3
q3	b	q3

q	a	$\delta(q,a)$
q0	a	q1
q0	b	q2

(d)

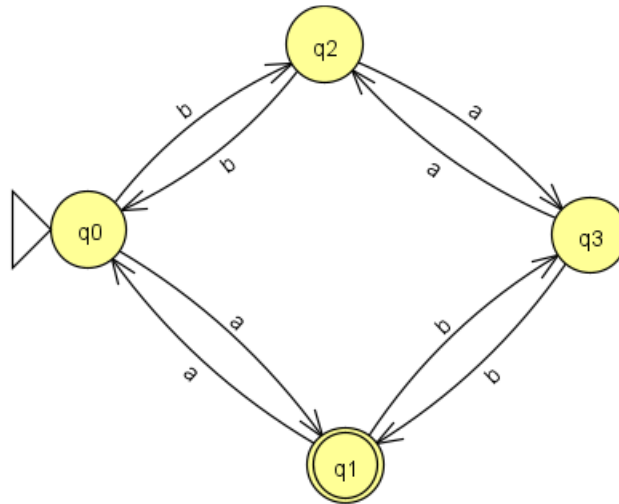
$K = \{q_0, q_1, q_2, q_3\}$

$\Sigma = \{a, b\}$

$s = q_0$

$F = q_1$

$q_1$	a	$q_0$
$q_1$	b	$q_3$
$q_2$	a	$q_3$
$q_2$	b	$q_0$
$q_3$	a	$q_2$
$q_3$	b	$q_1$



(e)

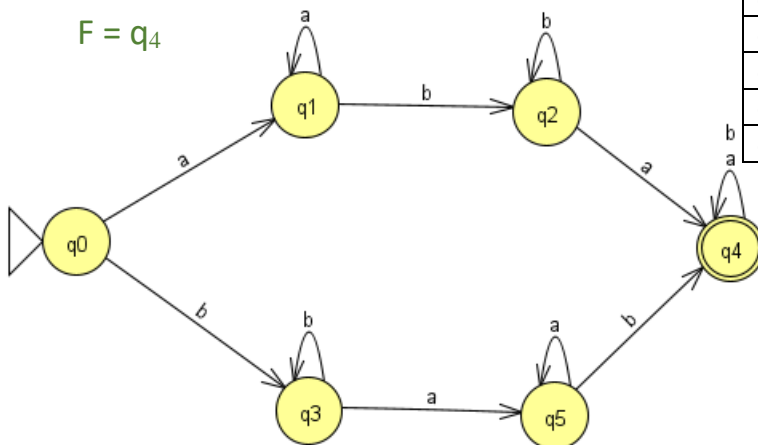
$K = \{q_0, q_1, q_2, q_3, q_4, q_5\}$

$\Sigma = \{a, b\}$

$s = q_0$

$F = q_4$

q	a	$\delta(q,a)$
$q_0$	a	$q_1$
$q_0$	b	$q_2$
$q_1$	a	$q_1$
$q_1$	b	$q_3$
$q_2$	a	$q_4$
$q_2$	b	$q_2$
$q_3$	a	$q_5$
$q_3$	b	$q_3$
$q_4$	a	$q_4$
$q_4$	b	$q_5$
$q_5$	a	$q_5$
$q_5$	b	$q_5$

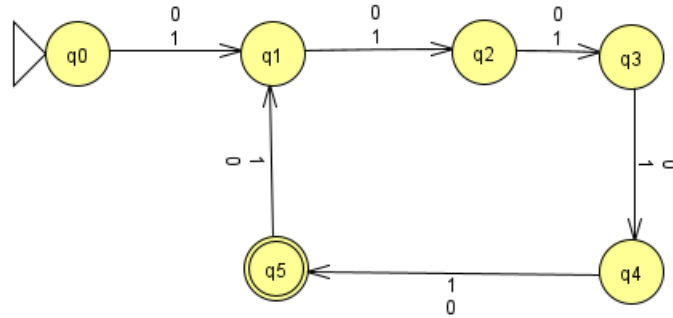


EXTRA:

Create a DFA that meets the rule “in  $\{0,1\}$  alphabet, for  $L(M)$   $w$  is 5 multiple of  $|w|$ .” by using JFLAP.

Add a short description that explains why each condition is included and why you accept each determining whether not. (in Turkish)

**ANSWER:**



**Regular expression :**  $[(0\cup 1)(0\cup 1)(0\cup 1)(0\cup 1)(0\cup 1)[(0\cup 1)(0\cup 1)(0\cup 1)(0\cup 1)(0\cup 1)]^*$

İstenilen kataşa eklenen elemanın ne olduđu önemli olmadığı için elimizdeki alfabaden herhangi bir eleman eklenmesi gerekmektedir. 5’e kadar olan sayıları (0 da dahil) ayrı bir durumda tutmalıyız, 5 durumda yapmaya çalışırsak boş katar inputu alındığında veya 4 elemanlı bir katarıda true output verir. Bu yüzden katarımızda 5’in katı kadar elemanımızın olduğunu temsil eden  $q_5$  durumu final state olur.