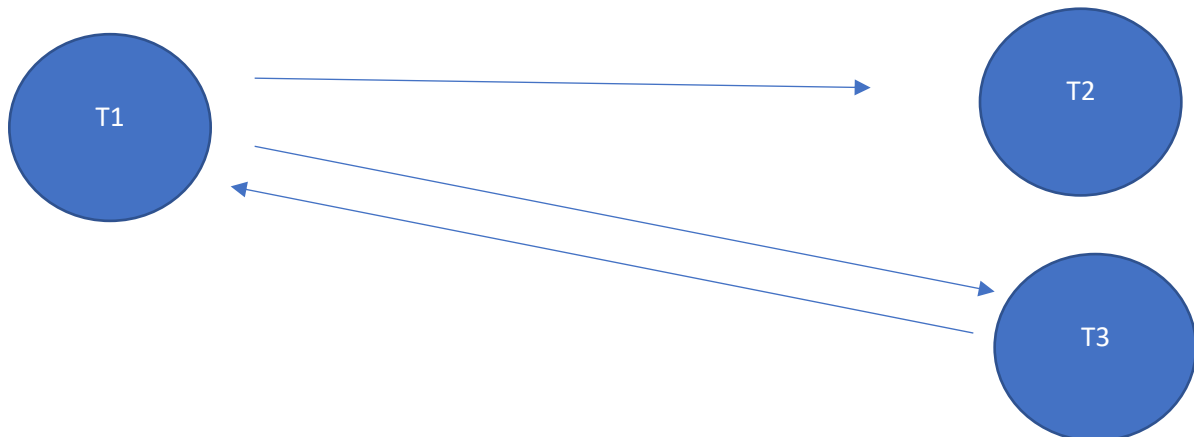


Consider the following three transactions and schedule (time goes from top to bottom). Is this schedule conflict-serializable? Show why or why not.

T1	T2	T3
R(A)		
W(A)		
		R(A)
		W(A)
	R(A)	
R(B)		
		R(B)
W(B)		
		W(B)
	R(B)	
	commit	
commit		
		commit

LOOP
ALERT!



This schedule is NOT conflict serializable because after drawing the precedence graph above, we can see that there is a loop between T1 and T3.