

$$7.1) \quad \tau_1 = (4, 11) \quad \tau_2 = (5, 16) \quad \tau_3 = (3, 31)$$

$$u = 4/11 + 5/16 + 3/31 = 0.77 < u(3) = 0.78$$

By UB test, they are schedulable

$$P_{m_1} = 1$$

$$\tau_1: (P(m_1), V(m_1))$$

$$\tau_3: (P(m_1), V(m_1))$$

No missed deadlines ✓

{See image attached}

$$\tau_1 \Rightarrow EEMM$$

$$\tau_2 \Rightarrow EEEEE$$

$$\tau_3 \Rightarrow MMM$$

$$2) \quad \tau_1 = (8, 7) \quad \tau_2 = (1, 9) \quad \tau_3 = (6, 26)$$

$$u = 8/7 + 1/9 + 6/26 = 0.77 < u(3) = 0.78$$

By UB test, they are schedulable

$$\tau_1: (P(m_1), V(m_1))$$

$$\tau_2: (P(m_2), V(m_2))$$

$$\tau_3: (P(m_1), V(m_1), P(m_2), V(m_2))$$

$$P_{m_1} = 1$$

$$P_{m_2} = 2$$

$$\tau_1 \Rightarrow EM_1M_1$$

$$\tau_2 \Rightarrow M_2$$

$$\tau_3 \Rightarrow M_1M_1M_{12}M_{12}M_2M_2$$

$\tau_2$  misses deadline

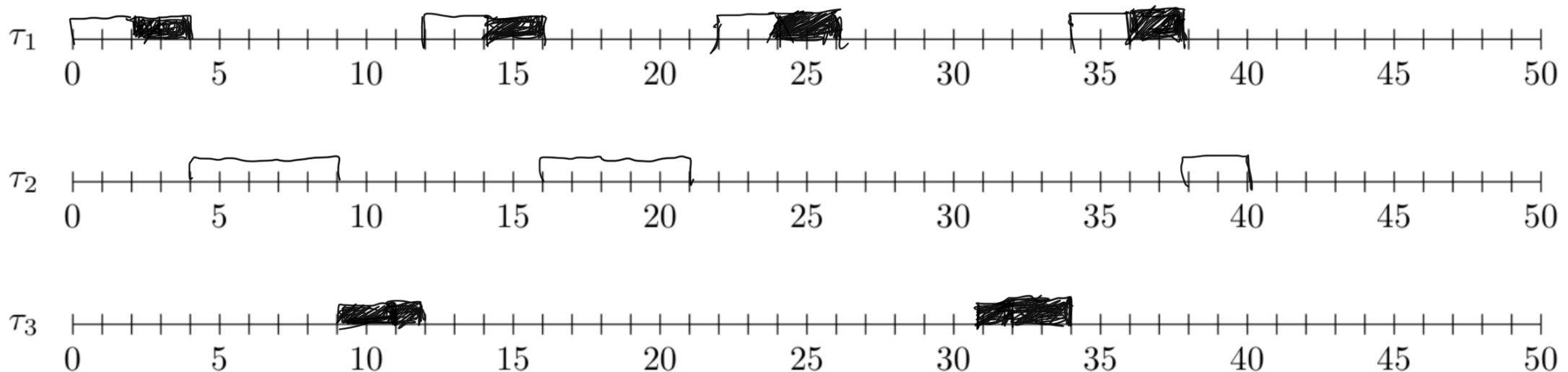
{see image attached}

Unfortunately, we can't always run our tasks in isolation. We need to share resources between each o protect this data from corruption with the following scheme:

- $\tau_1$  uses mutex  $m_1$  for the **last two ticks** of its computation time.
- $\tau_3$  uses mutex  $m_1$  for all three ticks of its computation time.

Please draw the expected scheduling using IPCP of this task set with the mutexes on the timelines  $t = 40$ . Assume that locks and unlocks can be done instantaneously (0 ticks). Indicate in which ticks holding a mutex by shading that tick in. Does any task miss its deadline?

No 



- $\tau_1$  uses mutex  $m_1$  for the **last two ticks** of its computation time.
- $\tau_2$  uses mutex  $m_2$  for all of its computation time (lock at the beginning of the tick, unlock at the end of the tick).
- $\tau_3$  uses mutex  $m_1$  for the **first four ticks** of its computation time, and mutex  $m_2$  for the **last four** of its computation time. (There will be two ticks when both mutexes are locked.)

Please draw the expected scheduling using IPCP of this task set with the mutexes on the timelines below. Assume that locks and unlocks can be done instantaneously (0 ticks). Indicate in which ticks each task uses each mutex by shading that tick in. Does any task miss its deadline? *Yes*

  $\Rightarrow m_1$   
  $\Rightarrow m_2$

