

# CLIQUEs

Refer Chapter 7

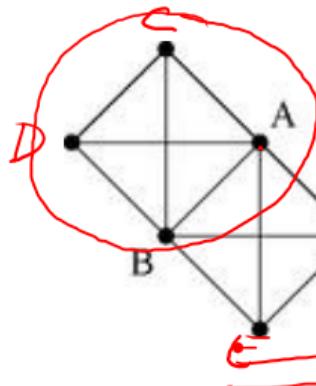
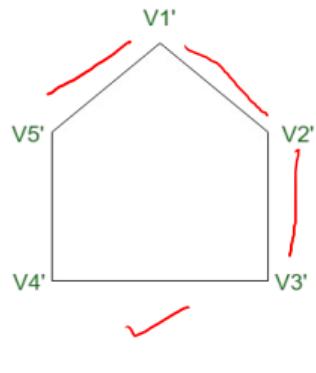
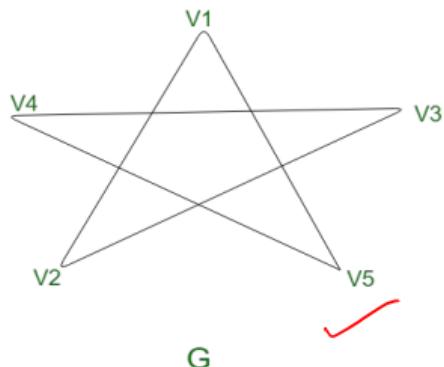
$$G : (V, E)$$

$$G(S) : V', E'$$

## CLIQUE

- A clique in a (an undirected) graph  $G$  is a subgraph of  $G$  isomorphic to a complete graph
- A Clique is a subset of vertices such that each vertex is connected to every other one. Vertices and edges in a clique makes that subgraph a complete graph.

## Isomorphism

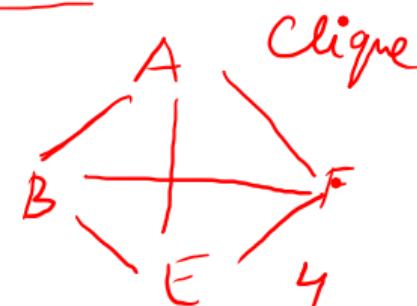


$$E - F \stackrel{?}{=}$$

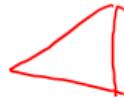
$F = D \setminus B \stackrel{?}{=}$

~~NP-Complete~~

A graph w/  $n$  vertices  
 $2^n$  subsets possible  
 $K_4$  clique



# CLIQUEs



## CLIQUEs ..

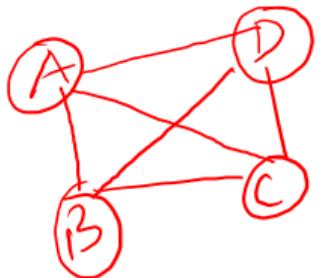
- Cliques of order 1 are just single vertices.
- Cliques of order 2 are just pairs of vertices connected by an edge, so for simple undirected graphs the number of cliques of order 2 is just the number of edges.
- The clique number of  $G$  is the maximum order of a clique in  $G$ .
  - For  $K_n$ , it is  $n$ .
  - For  $C_n$  it is 2 unless  $n = 3$  (for  $C_3 = 3$ ).

## CLIQUEs ..

- Cliques of order 3 are triangles in the graph

### Example: $K_4$

- This has one clique of order 4 and 4 of order 3, 6 of order 2 and 4 of order 1



K<sub>4</sub> → Complete graph w/ 4 vertices

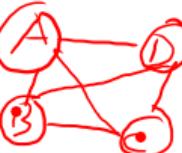
Cliques of order 1 → A    D    C    B  
A subgraph w/ 1 vertex

$\begin{matrix} G \\ \cong \\ H \end{matrix}$  is subgraph of H    How many cliques of order 1 = 4  
Cliques of order 2 → A - D  
= 6 cliques

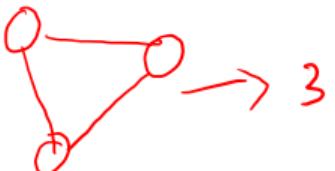
- Cliques of order 3 → Triangles = 4 cliques



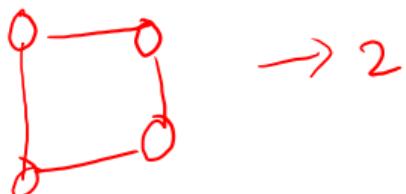
Cliques of order 4 →



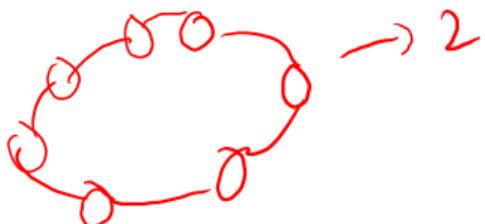
$C_n$  - circulant graph



$\rightarrow 3$  Maximal clique order



$\rightarrow 2$

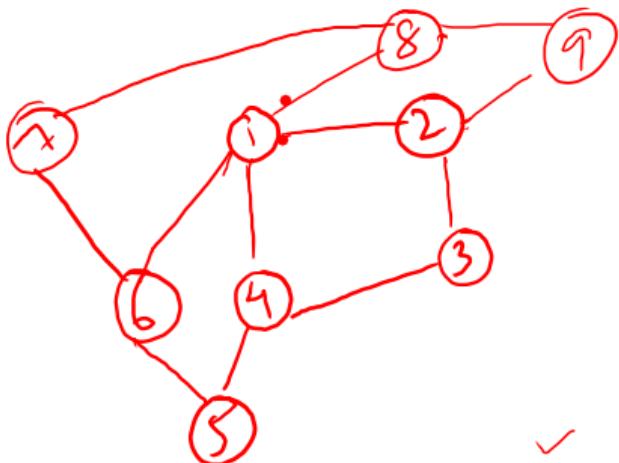


$\rightarrow 2$

$C_n = 2$  ; if  $n=3$ ,  
 $3$

$n = 9$

Algorithm to find clique



Deg(1) = 4      Deg(5) = ✓      Deg(9) = ✓ - X  
Deg(2) = 3      (6) = 3  
Deg(3) = 2 ✓      (7) = 2  
Deg(4) = 2 ✓      (8) = 3 → 2



1. Find a vertex  $v$  of the smallest possible degree in  $G$ . ✓
2. If the degree of  $v$  is  $n - 1$ , stop;  $G$  is a clique, so the largest clique in  $G$  has size  $n$ .
3. Otherwise, remove  $v$  and all of its edges from  $G$ .  
Find the largest clique in the smaller graph. Report that as the largest clique in  $G$ .

$n = 8$



$n = 7$

Clique of order  $\geq$

$n = 6$

$n = 5$

$n = 4$

$n = 3$

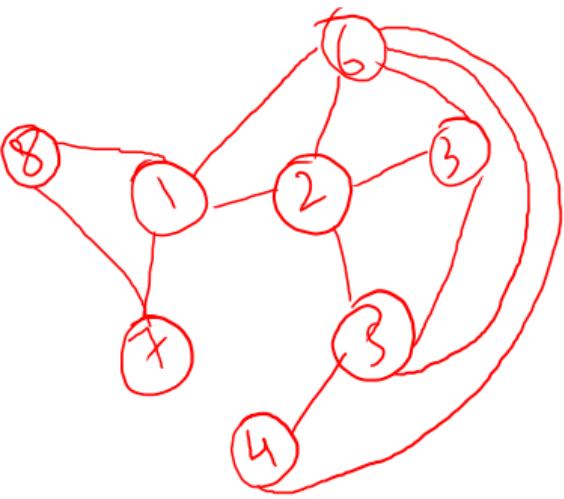
$n = 2$

$$\text{Deg}(1) = 4 - 3 \xrightarrow{v=1} \text{Deg}(5) = 2 - 1 \quad \cancel{\text{Deg}(9) = 2 - x}$$

$$(2) = 3 - 2 - 1 \times (6) = 3 \rightarrow 2 - x$$

$$(3) = 2 \vee x \quad (7) = 2 - x$$

$$(4) = 3 - 2 \quad (8) = 3 \rightarrow 2 \rightarrow 1 \times$$



$$\text{Deg}(1) = 4 \rightarrow 3 \rightarrow 2 - X$$

$$\text{Deg}(2) = 4 \rightarrow 3 \checkmark$$

$$\text{Deg}(3) = 3 \checkmark$$

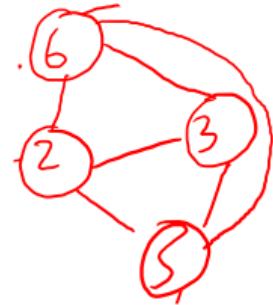
$$(4) = 2 \rightarrow X$$

$$(5) = 4 \rightarrow 3 \checkmark$$

$$(6) = 5 \rightarrow 4 \rightarrow 3 \checkmark$$

$$(7) = 2 \rightarrow 1 \rightarrow X$$

$$(8) = 2 \rightarrow X$$



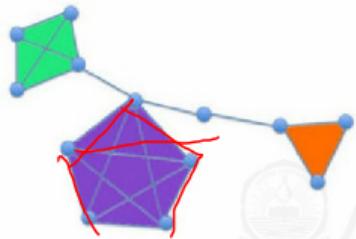
Clique of order 4

Find maximal clique.

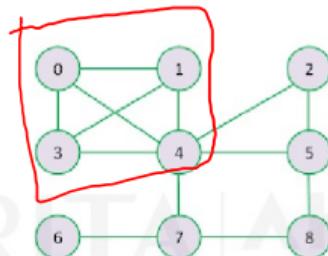
(1,2), (2,3), (3,1), (4,3), (4,1), (4,2) ✓

(1,2), (2,3), (3,1), (4,3), (4,5), (5,3) ✓

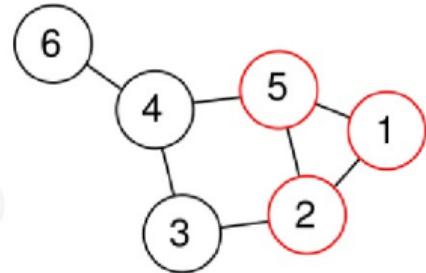
# CLIQUEs ..



- Cliques  $\rightarrow$  2, 3, 4 and ~~6~~<sup>5</sup>
- Clique Number  $\rightarrow$  ~~6~~<sup>5</sup>



- Cliques  $\rightarrow$  2, 3 and 4
  - Clique Number  $\rightarrow$  4
- Maximal order*



- Cliques  $\rightarrow$  2 and 3
- Clique Number  $\rightarrow$  3

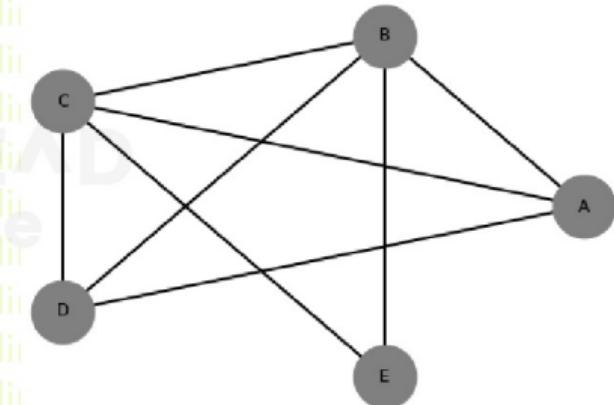
# FINDING CLIQUES ..

## "Cliques" in NetworkX

- `find_cliques(G[, nodes])`: Returns all maximal cliques in an undirected graph
- `graph_clique_number(G[, cliques])`: Returns the clique number of the graph
- `cliques_containing_node(G[, nodes, cliques])`: Returns a list of cliques containing the given node
- `node_clique_number(G[, nodes, cliques, ...])`: Returns the size of the largest maximal clique containing each given node

# CLIQUEs USING NETWORKX: EXAMPLES

```
breaklines
breaklines
list(nx.find_cliques(G))
[[ 'B' , 'C' , 'E' ] , [ 'B' , 'C' , 'A' ] ,
 [ 'B' , 'D' , 'E' ]]
nx.graph_clique_number(G)
4
nx.cliques_containing_node(G)
{ 'A' : [[ 'C' , 'B' , 'A' , 'D' ]],
 'B' : [[ 'C' , 'B' , 'E' ] ,
 [ 'C' , 'B' , 'A' , 'D' ]],
 'C' : [[ 'C' , 'B' , 'E' ] ,
 [ 'C' , 'B' , 'A' , 'D' ]],
 'D' : [[ 'C' , 'B' , 'A' , 'D' ]],
 'E' : [[ 'C' , 'B' , 'E' ]]}
nx.node_clique_number(G)
{ 'A' : 4, 'B' : 4, 'C' : 4, 'D' : 4, 'E' : 3}
```



## K-plex

A k-plex of size n is a maximal subset of n vertices within a network such that each vertex is connected to at least  $n - k$  of the others.

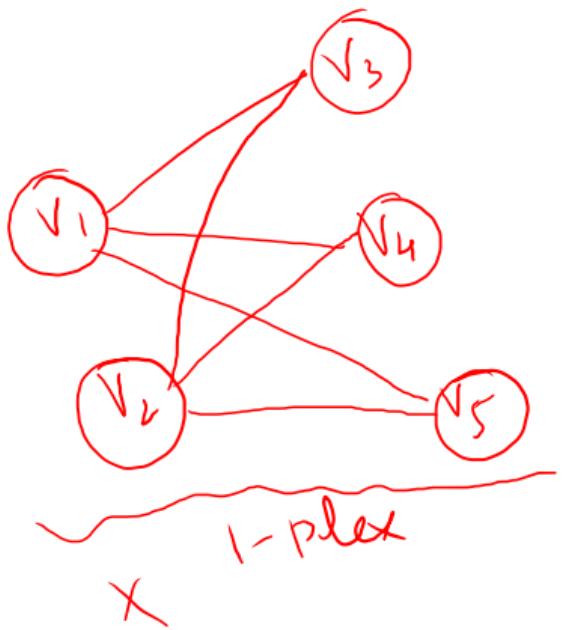
If  $k = 1$ ,  $\rightarrow$  ordinary clique

$$n - k = n - 1 \text{ of the vertices}$$

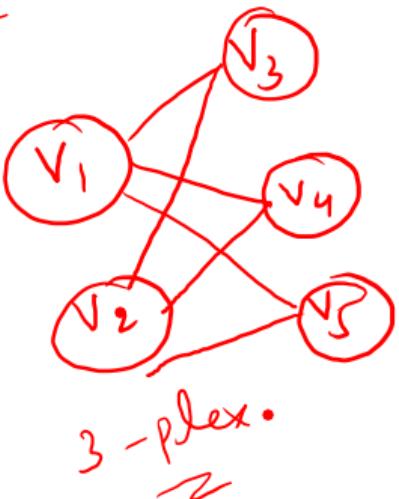
$\approx$   
Useful concept for discovering groups within networks: in real life many groups in social and other networks form k-plexes. There is no solid rule about what value k should take.

$$n - 3$$

K-core  
at least  $k$  vertices

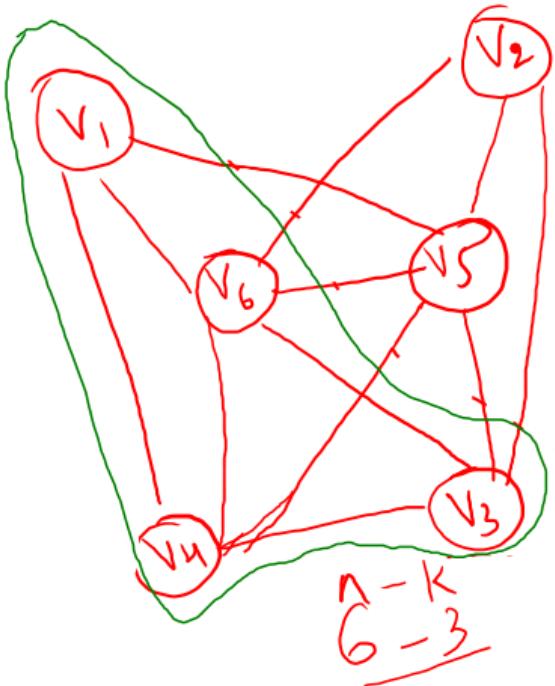


$k$ -plex?  
 $\approx$   
 $K = ?$   
 $\approx$



$5$  vertices  
 $\equiv$   
 $3 \rightarrow$   
 $n - k$   
 $\equiv$   
 $K = 2$   
 $3$  vertices.

$k$ -plex  
 $n - k = 2$   
 $\equiv$   
 $3$ -plex

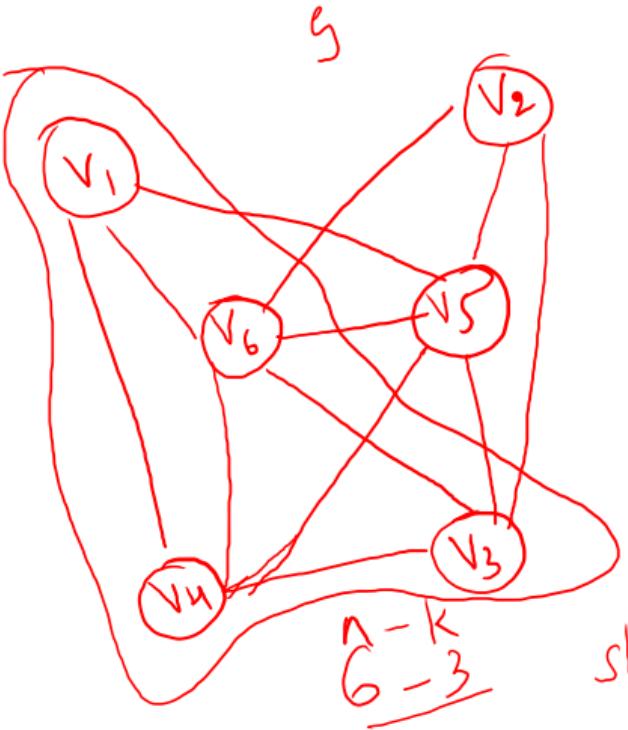


$\deg(v_1) = 3$  ✓  
 $(v_2) = 3$  ✓  
 $(v_3) = 4$   
 $(v_4) = 4$   
 $v_5 = 5^-$   
 $v_6 = 5^-$

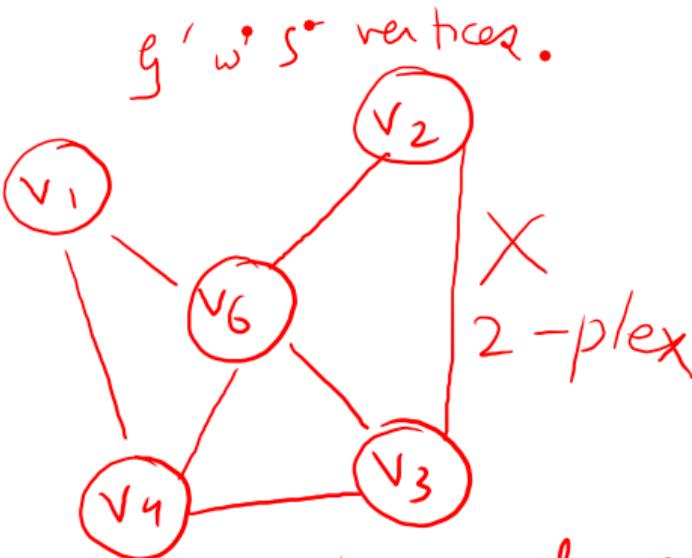
$\text{K-plex?}$   
 $\text{Min deg}(G) = ?$

$v_1 \rightarrow g'$   
 $v_1 \rightarrow v_5$   
 $v_1 \rightarrow v_6$   
 $v_1 \rightarrow v_4$   
 $v_1 \rightarrow v_3$   
 $v_6 \rightarrow v_5$   
 $v_6 \rightarrow v_4$   
 $v_6 \rightarrow v_3$   
 $v_4 \rightarrow v_5$   
 $v_4 \rightarrow v_3$   
 $v_3 \rightarrow v_5$   
 $\rightarrow 2\text{-plex.}$

$n = 5^-$   
 For  $g'$  to be 2-plex  
 $\deg(v^*) = n - 2 = 3$  ✓



should be  $n - k \Rightarrow n - 2$  vertices  
connected to  $\underline{\underline{3}}$  vertices



Is it 2-plex?

$n = s \cdot$  Each vertex

# OBJECTIVES

## Clustering in Networks

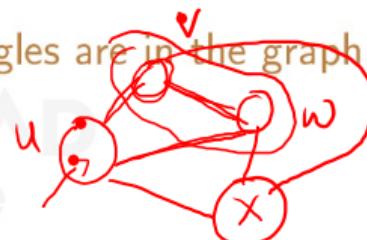
- Clustering Coefficient
- Transitivity

# CLUSTERING (1/3)

- Cliques are unstable under minor changes, and hence this measure is not considered as robust.

- *cohesiveness*: tightly-knit-groups
- identifying cohesive subgraphs: Clustering

\* Visually this is the notion of how many triangles are in the graph and how much they bunch together.



## DEFINITION (CLUSTERING COEFFICIENT)

Let  $G = (V, E)$  be a simple undirected graph. The (individual) clustering coefficient of a node  $u \in V$  is the probability  $\underline{Cl(u)}$  that two neighbors of  $u$  are themselves neighbors

### Clustering Coefficient

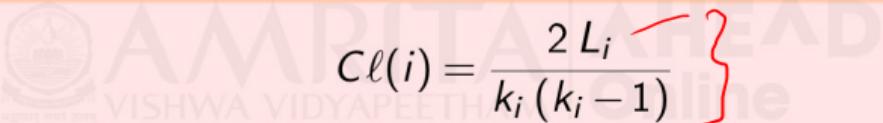
$$\underline{Cl(u)} = \frac{|\{(v, w) : v, w \in V; v \neq w; (u, v), (u, w), (v, w) \in E\}|}{|\{(v, w) : v, w \in V; v \neq w; (u, v), (u, w) \in E\}|}$$

## CLUSTERING (2/3)

Clustering Coefficient ..

$$C\ell(u) = \frac{|\{(v, w) : v, w \in V; v \neq w; (u, v), (u, w), (v, w) \in E\}|}{\deg(u)(\deg(u) - 1)}$$

Clustering Coefficient: Simplified Form

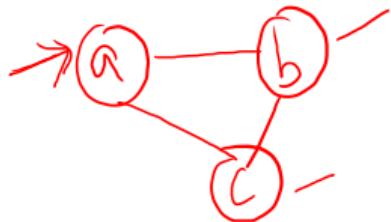

$$C\ell(i) = \frac{2 L_i}{k_i (k_i - 1)}$$

where,  $L_i$  is the number of edges between  $k_i$  neighbors, and  $k_i$  is the degree of node  $i$ .

EXAMPLE

Friendship networks: measures how many of your friends are friends with each other

- The average clustering coefficient is the average of  $C\ell(u)$  over all  $u \in V$



## CLUSTERING (3/3)



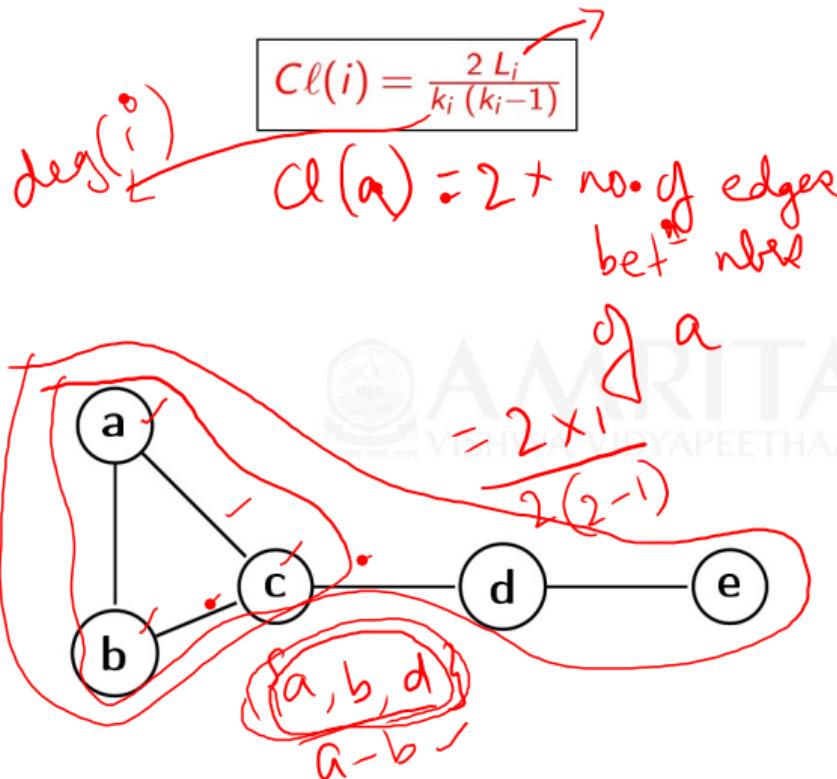
Overall Clustering Coefficient: Transitivity

$$C\ell(G) = \frac{\sum_{u \in V} |\{(v, w) : v, w \in V; v \neq w; (u, v), (u, w), (v, w) \in E\}|}{\sum_{u \in V} |\{(v, w) : v, w \in V; v \neq w; (u, v), (u, w) \in E\}|}$$

Transitivity Simplified

$$C\ell(G) = \frac{\text{Number of Triangles} \times 3}{\text{Number of Connected Triplets}}$$

# CLUSTERING COEFFICIENT: EXAMPLES (1/2)

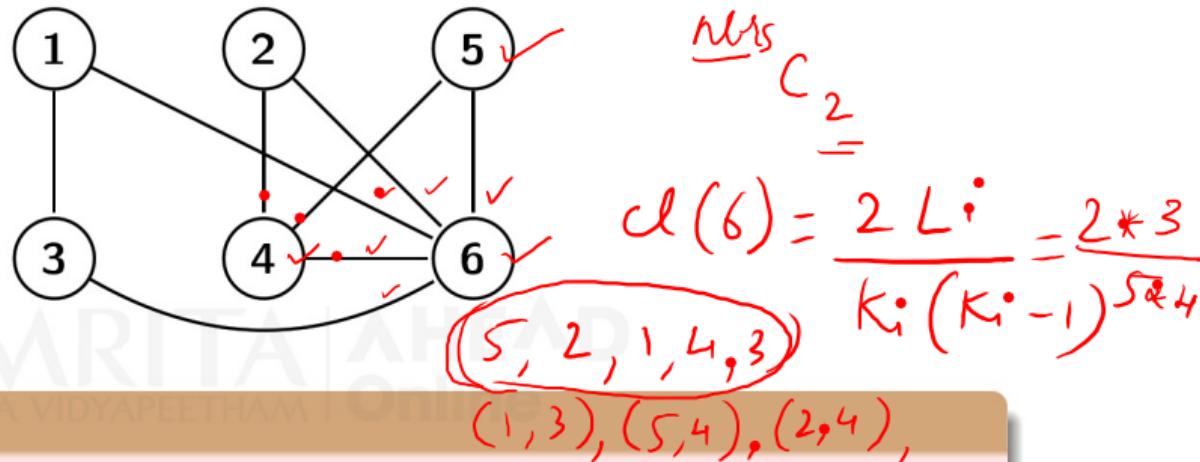


$$Cl(G) = \frac{\text{Number of Triangles} \times 3}{\text{Number of Connected Triplets}}$$

$Cl(i)$  &  $Cl(G)$

- $Cl(a) = \frac{2 \times 1}{2(2-1)} = 1$  ✓
- $Cl(b) = \frac{2 \times 1}{2(2-1)} = 1$  ✓
- $Cl(c) = \frac{2 \times 1}{3(3-1)} = 1/3$
- $Cl(d) = \frac{2 \times 0}{1(1-1)} = 0$
- $Cl(e) = \frac{2 \times 0}{2(2-1)} = 0$
- $Cl(G) = \frac{1 \times 3}{6} = 1/2$  ✓

## CLUSTERING COEFFICIENT: EXAMPLES (2/2)



$C\ell(G)$

- No. of Triangles: 3, No. of Triplets: 17

$$\text{--- Node 1: } \binom{2}{2} = 1 \checkmark$$

$$\text{--- Node 2: } \binom{2}{2} = 1$$

$$\text{--- Node 3: } \binom{2}{2} = 1$$

$$\text{--- Node 4: } \binom{3}{2} = 3$$

$$\text{--- Node 5: } \binom{2}{2} = 1$$

$$\text{--- Node 6: } \binom{5}{2} = 10$$

$$C\ell(G) = \frac{3 \times 3}{17} = 9/17$$

# OBJECTIVES

## Clustering using NetworkX

- Clustering Coefficient
- Transitivity
- Triangles

# FUNCTIONS FOR CLUSTERING

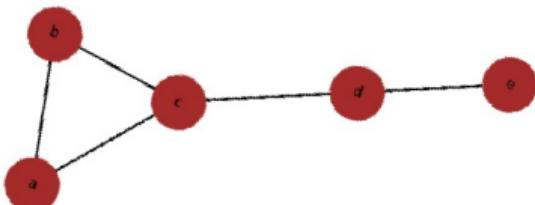
## In NetworkX

- `clustering(G[, nodes])`: Compute the clustering coefficient for nodes
- `average_clustering(G[, nodes])`: Compute the average clustering coefficient for the graph G
- `transitivity(G)`: Compute graph transitivity, the fraction of all possible triangles present in G.
- `triangles(G[, nodes])`: Compute the number of triangles.

# EXAMPLE 1

```
breaklines
import networkx as nx
import matplotlib . pyplot as plt
edge_widths = [2]
G2 = nx.Graph([( 'a' , 'b' ),( 'a' , 'c' ),( 'b' , 'c' ), ( 'c' , 'd' ), ( 'd' , 'e')])
nx . draw (G2 , with_labels = True , node_size=2000, node_color="brown",  
width=edge_widths)
```

breaklines

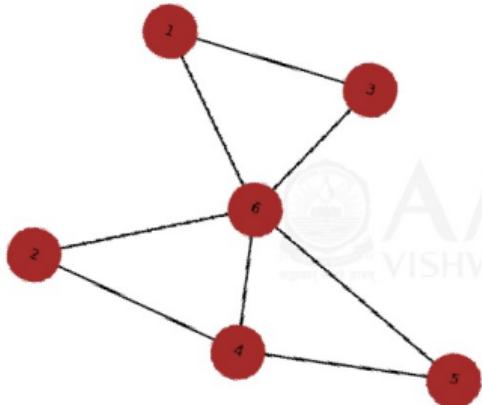


breaklines

```
print(nx.clustering(G2))
{'a': 1.0, 'b': 1.0,
 'c': 0.33, 'd': 0, 'e': 0}
nx.transitivity(G2)
0.5
```

breaklines

## EXAMPLE 2



breaklines

```
nx.transitivity(G2)
```

0.5294

```
print(nx.clustering(G2))
```

{1: 1.0, 3: 1.0, 6: 0.3,  
2: 1.0, 4: 0.66, 5: 1.0}

```
nx.average_clustering(G2)
```

0.8277

```
nx.triangles(G2)
```

{1: 1, 3: 1, 6: 3,  
2: 1, 4: 2, 5: 1}

breaklines

## EXAMPLE 3

breaklines

```
G2 = nx.karate_club_graph()  
G2 = nx.convert_node_labels_to_integers(G, first_label=1)
```

Breaklines



breaklines

```
nx.transitivity(G2)  
0.2556818181818182  
print(nx.clustering(G2))  
{1: 0.15, 2: 0.33, 3: 0.24, ...,  
17: 1.0, ..., 33: 0.196, 34: 0.110}  
nx.average_clustering(G2)  
0.5706384782076823  
nx.triangles(G2,1)  
18
```

Breaklines

# SUMMARY

## NETWORKX

- Clustering Coefficient
- Transitivity
- Triangles

## REFERENCES

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