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## Concise Representations of Frequent Itemsets

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## Do We Need to Know All Frequent Itemsets?

- The number of frequent itemsets is usually huge.
- Time of their discovery can be significant.
- There are cases in which one needs to know only a small subset of frequent itemsets! (*Representative* and minimal *non-redundant rules* can be derived directly from concise representations of frequent itemsets called *closed itemsets* and *generators*.)

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## Lossless Representations of Frequent Itemsets

- Itemsets representation is meant *lossless* if it allows derivation and support determination of all frequent itemsets without accessing the database.

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## Lossless Representations of Frequent Itemsets

Lossless representations of frequent itemsets are based on the following **sets subsuming other sets**:

- closed itemsets*
- (*key*) *generators*
- (*generalized*) *disjunctive sets*

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## Simple example of reasoning about supports of itemsets

- Let  $\text{sup}(\{ac\}) = 3$  and  $\text{sup}(\{abcde\}) = 3$ .
- This information is sufficient to determine the support of  $\{abce\}$  as follows:  

$$3 = \text{sup}(\{ac\}) \geq \text{sup}(\{abce\}) \geq \text{sup}(\{abcde\}) = 3.$$
Hence,  

$$\text{sup}(\{ac\}) = \text{sup}(\{abce\}) = \text{sup}(\{abcde\}) = 3.$$
- Property.** if  $X \subseteq Y$  and  $\text{sup}(X) = \text{sup}(Y) = k$ , then for each itemset  $Z$  such that:  $X \subseteq Z \subseteq Y$ , its supports also equals  $k$ .

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## Reasoning about supports of itemsets...

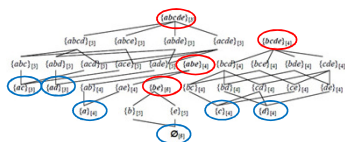
Id	Transaction
$T_1$	$\{abcde\}$
$T_2$	$\{abcdef\}$
$T_3$	$\{abcdehi\}$
$T_4$	$\{abe\}$
$T_5$	$\{bcdehi\}$

Example.  $\text{minSup}=2$ ;  $\text{minConf}=77\%$ .

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## Reasoning about supports of itemsets

Id	Transaction
$T_1$	$\{abcde\}$
$T_2$	$\{abcdef\}$
$T_3$	$\{abcdehi\}$
$T_4$	$\{abe\}$
$T_5$	$\{bcdehi\}$



Example.  $\minSup=2$ ;  $\minConf=77\%$ .

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## Supports and tid-lists of subsets/supersets

**Lemma.** Let  $X \subseteq Y$ . Then:

$$t(X) = t(Y) \Leftrightarrow sup(X) = sup(Y).$$

**Proof**

( $\Rightarrow$ ). Trivial.

( $\Leftarrow$ ). Let  $X \subseteq Y$  and  $sup(X) = sup(Y)$ .

Then,  $t(X) \supseteq t(Y)$  and  $|t(X)| = |t(Y)|$ .

Hence,  $t(X) = t(Y)$ .

Id	Transaction
$T_1$	$\{abcde\}$
$T_2$	$\{abcdef\}$
$T_3$	$\{abcdehi\}$
$T_4$	$\{abe\}$
$T_5$	$\{bcdehi\}$

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## Representation of Frequent Itemsets based on Closed Itemsets

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## Closures of Itemsets

- Let  $I$  be a set of all items that may occur in  $D$ .
- A closure of itemset  $X$  is denoted as  $\gamma(X)$  and defined as:

$$\gamma(X) = \bigcap \{T \in D \cup \{I\} \mid T \supseteq X\}.$$

- Note:** An itemset has exactly one itemset as its closure!
- Property.** The closure of itemset  $X$  is the greatest superset  $Y \supseteq X$  such that  $sup(Y) = sup(X)$ .

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## Supports of Proper Supersets of Closures

- Property.** The closure of itemset  $X$  is the greatest superset  $Y \supseteq X$  such that

$$sup(Y) = sup(X).$$



- Note:** Supports of all proper supersets of an itemset being a closure of any itemset is different from (less than) than the support of this closure.

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## Example: Closures of Itemsets

Id	Transaction
$T_1$	$\{abcde\}$
$T_2$	$\{abcdef\}$
$T_3$	$\{abcdehi\}$
$T_4$	$\{abe\}$
$T_5$	$\{bcdehi\}$

- $\gamma(\{abc\}_3) = \{abcde\}_3$
- $\gamma(\{abcde\}_3) = \{abcde\}_3$

**Question:** Can  $sup(\{abcdehi\})$  be equal to 3?

**Answer:** No.

$\{abcde\}_3$  is the closure of  $\{abc\}_3$ . Thus,  $sup(\{abcdehi\})$  must be less than  $sup(\{abcde\})$ . Indeed,  $sup(\{abcdehi\}) = 1 < sup(\{abcde\}) = 3$ .

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**Closed itemsets**

- An itemset is defined as *closed* if it is equal to its closure.
- Property.** Each closure is a closed itemset.

↓

- Property.** Supports of all proper supersets of a closed itemset are different from (less than) the support of this closed itemset.

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**Example: Closed Itemsets**

Id	Transaction
T1	{abcde}
T2	{abcdef}
T3	{abcdehi}
T4	{abe}
T5	{bcdehi}

- $\gamma(\{abc\}) = \{abcde\}$ .
- Hence,  $\{abc\}$  is not a closed itemset.
- $\gamma(\{abcde\}) = \{abcde\}$ .
- Hence,  $\{abcde\}$  is a closed itemset.

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**Important Property of Closed Itemsets**

- Important property of closed itemsets:** The set of all closed itemsets is sufficient to determine support of each itemset  $X$  in  $2^I$ , namely:  

$$\text{sup}(X) = \max\{\text{sup}(Y) \mid Y \text{ is closed} \wedge Y \supseteq X\}.$$

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**Closed Itemsets Representation**

- Closed itemsets representation (CR)* consists of all frequent closed itemsets and the information about their supports.

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**Example: Reasoning with CR**

- $\text{sup}(\{bef\})=?$ ,  $\text{sup}(\{bcd\})=?$

**CR**

- $\{bef\}$  has no superset in CR, so:  
 $\text{sup}(\{bef\}) \leq \text{minSup}.$
- $\text{sup}(\{bcd\}) = \max(\text{sup}(\{bcde\}_4), \text{sup}(\{abcde\}_3)) = 4.$

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**Calculating FCs with CHARM**

Id	Transaction
T1	{abcde}
T2	{abcdef}
T3	{abcdehi}
T4	{abe}
T5	{bcdehi}

- $t(X) = t(Y)$ :  
 – remove node  $Y$  from the tree;  
 – replace each  $X$  with  $X \cup Y$ .
- $t(X) \subset t(Y)$ :  
 – replace each  $X$  with  $X \cup Y$ .
- $t(X) \supset t(Y)$ :  
 – remove node  $Y$  from the tree;  
 – add  $X \cup Y$  as a child of node  $X$ .
- otherwise:  
 – add  $X \cup Y$  as a child of node  $X$ .

$$h(X \cup Y) = \sum_{i \in t(X \cup Y)} i$$

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Example: Calculating FCs with CHARM...

Id Transaction  
1 {abcde}  
2 {abcde f}  
3 {abcde h i}  
4 {abe}  
5 {bcde h i}

Assumption:  $\text{minSup} = 2$

$h(X \cup Y) = \sum_{i \in t(X \cup Y)} i$

1	...	6	...	10	11	...	15	19

Example: Calculating FCs with CHARM...

Id Transaction  
1 {abcde}  
2 {abcde f}  
3 {abcde h i}  
4 {abe}  
5 {bcde h i}

Assumption:  $\text{minSup} = 2$

$h(X \cup Y) = \sum_{i \in t(X \cup Y)} i$   $h(be) = 1+2+3+4+5 = 15$

1	...	6	...	10	11	...	15	20
							$be_5$	

Example: Calculating FCs with CHARM...

Id Transaction  
1 {abcde}  
2 {abcde f}  
3 {abcde h i}  
4 {abe}  
5 {bcde h i}

Assumption:  $\text{minSup} = 2$

$h(X \cup Y) = \sum_{i \in t(X \cup Y)} i$   $h(be) = 1+2+3+4+5 = 15$ ,  $h(abe) = 1+2+3+4 = 10$ ,

1	...	6	...	10	11	...	15	21
				$abe_4$			$be_5$	

Example: Calculating FCs with CHARM...

Id Transaction  
1 {abcde}  
2 {abcde f}  
3 {abcde h i}  
4 {abe}  
5 {bcde h i}

Assumption:  $\text{minSup} = 2$

$h(X \cup Y) = \sum_{i \in t(X \cup Y)} i$   $h(be) = 1+2+3+4+5 = 15$ ,  $h(abe) = 1+2+3+4 = 10$ ,  $h(abcde) = 1+2+3 = 6$

1	...	6	...	10	11	...	15	22
		$abcde_3$		$abe_4$			$be_5$	

Example: Calculating FCs with CHARM

Id Transaction  
1 {abcde}  
2 {abcde f}  
3 {abcde h i}  
4 {abe}  
5 {bcde h i}

Assumption:  $\text{minSup} = 2$

$h(X \cup Y) = \sum_{i \in t(X \cup Y)} i$   $h(be) = 1+2+3+4+5 = 15$ ,  $h(abe) = 1+2+3+4 = 10$ ,  $h(abcde) = 1+2+3 = 6$ ,  $h(bcde) = 1+2+3+5 = 11$

1	...	6	...	10	11	...	15	23
		$abcde_3$		$abe_4$	$bcde_4$		$be_5$	

Calculating FCs with dCHARM

Id Transaction  
T1 {abcde}  
T2 {abcde f}  
T3 {abcde h i}  
T4 {abe}  
T5 {bcde h i}

- $d(X) = d(Y)$ :
  - remove node Y from the tree;
  - replace each X with  $X \cup Y$ .
- $d(X) \supset d(Y)$ :
  - replace each X with  $X \cup Y$ .
- $d(X) \subset d(Y)$ :
  - remove node Y from the tree;
  - add  $X \cup Y$  as a child of node X.
- otherwise:
  - add  $X \cup Y$  as a child of node X.

$d(X \cup Y) = d(Y) \cdot d(X)$   
 $\text{sup}(X \cup Y) = \text{sup}(X) - |d(X \cup Y)|$   
 $h(X \cup Y) = h(X) - (\sum_{i \in t(X \cup Y)} i)$

Example: Calculating FCs with dCHARM...

Id Transaction  
1 {abcde}  
2 {abcdef}  
3 {abcdehi}  
4 {abe}  
5 {bcdehi}

Assumption:  $\text{minSup} = 2$

•  $h(X \cup Y) = h(X) - (\sum_{i \in d(X \cup Y)} i)$

1	...	6	...	10	11	...	15

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Example: Calculating FCs with dCHARM...

Id Transaction  
1 {abcde}  
2 {abcdef}  
3 {abcdehi}  
4 {abe}  
5 {bcdehi}

Assumption:  $\text{minSup} = 2$

•  $h(X \cup Y) = h(X) - (\sum_{i \in d(X \cup Y)} i)$

•  $h(be) = 1+2+3+4+5 = 15$

1	...	6	...	10	11	...	15
							be5

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Example: Calculating FCs with dCHARM...

Id Transaction  
1 {abcde}  
2 {abcdef}  
3 {abcdehi}  
4 {abe}  
5 {bcdehi}

Assumption:  $\text{minSup} = 2$

•  $h(X \cup Y) = h(X) - (\sum_{i \in d(X \cup Y)} i)$

•  $h(be) = 1+2+3+4+5 = 15$   
•  $h(abe) = h(be) - 5 = 15 - 5 = 10$

1	...	6	...	10	11	...	15
				abe4			be5

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Example: Calculating FCs with dCHARM...

Id Transaction  
1 {abcde}  
2 {abcdef}  
3 {abcdehi}  
4 {abe}  
5 {bcdehi}

Assumption:  $\text{minSup} = 2$

•  $h(X \cup Y) = h(X) - (\sum_{i \in d(X \cup Y)} i)$

•  $h(be) = 1+2+3+4+5 = 15$   
•  $h(abe) = h(be) - 5 = 15 - 5 = 10$   
•  $h(abcde) = h(be) - 4 = 6$

1	...	6	...	10	11	...	15
		abcde3		abe4			be5

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Example: Calculating FCs with dCHARM

Id Transaction  
1 {abcde}  
2 {abcdef}  
3 {abcdehi}  
4 {abe}  
5 {bcdehi}

Assumption:  $\text{minSup} = 2$

•  $h(X \cup Y) = h(X) - (\sum_{i \in d(X \cup Y)} i)$

•  $h(be) = 1+2+3+4+5 = 15$   
•  $h(abe) = h(be) - 5 = 15 - 5 = 10$   
•  $h(abcde) = h(be) - 4 = 6$   
•  $h(bcde) = h(be) - 4 = 15 - 4 = 11$

1	...	6	...	10	11	...	15
		abcde3		abe4	bcde4		be5

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Algorithm for Transforming FCs into Fs

Id Transaction  
1 {abcde}  
2 {abcdef}  
3 {abcdehi}  
4 {abe}  
5 {bcdehi}

• FC: {be}5, {abe}4, {abcde}3, {bcde}4  
• F: ?

• F5: {abcde}3  
• F4: {bcde}4, {abcd}3, {abce}3, {abde}3, {acde}3  
• F3: {abe}4, {bcd}4, {bce}4, {bde}4, {cde}4, {ace}3, {ade}3  
• F2: {be}5, {ab}4, {ae}4, {bc}4, {bd}4, {cd}4, {ce}4, {de}4, {ac}3, {ad}3  
• F1: {b}5, {e}5, {a}4, {c}4, {d}4  
• F0: ∅5

## Representation of Frequent Itemsets based on (Key) Generators

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## Generator of an Itemset

- $Y$  is defined a *generator* of itemset  $X$  if it is a minimal subset of  $X$  such that

$$\text{sup}(Y) = \text{sup}(X)$$

(equivalently,  $t(Y) = t(X)$ ).

- Note 1:** Supports of all proper subsets of a generator of any itemset are different from (greater than) than the support of this generator.
- Note 2:** An itemset may have more than one generator!

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## Example: Generators of Itemsets

Id Transaction  
T1 {abcde}  
T2 {abcdeh}  
T3 {abcdehi}  
T4 {abe}  
T5 {bcdehi}

- $G(\{abcd\}_3) = \{\{ac\}_3, \{ad\}_3\}$
- $G(\{ac\}_3) = \{\{ac\}_3\}$
- $G(\{ad\}_3) = \{\{ad\}_3\}$
- $\{ac\}_3$  is a generator of (among other)  $\{abcd\}_3$  and  $\{ac\}_3$  itself. Thus,  $\text{sup}(\{a\})$  must be greater than  $\text{sup}(\{ac\})$ . Indeed,  $\text{sup}(\{a\}) = 4 > \text{sup}(\{ac\}) = 3$ .

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## (Key) Generator

- An itemset  $X$  is defined as a (*key*) *generator* if  $X$  is a generator of  $X$ .
- (*Key*) *generator* has only one *generator* – itself.
- Property.** Supports of all proper subsets of a (*key*) generator are different from (greater than) than the support of this (*key*) generator.

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## Example: (Key) Generators

Id Transaction  
T1 {abcde}  
T2 {abcdeh}  
T3 {abcdehi}  
T4 {abe}  
T5 {bcdehi}

- $G(\{abcd\}_3) = \{\{ac\}_3, \{ad\}_3\}$ .
- Hence,  $\{abcd\}$  is not a (*key*) generator.
- $G(\{ac\}) = \{\{ac\}\}$ .
- Hence,  $\{ac\}$  is a (*key*) generator.
- $G(\{ad\}) = \{\{ad\}\}$ .
- Hence,  $\{ad\}$  is a (*key*) generator.

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## Supersets and Subsets of (Key) Generators

- Theorem.**
  - All supersets of a non-(*key*) generator are not (*key*) generators.
  - All subsets of a (*key*) generator are (*key*) generators.

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### Example: Supersets of non-(Key) Generators

Id Transaction  
T1 {abcde}  
T2 {abcdef}  
T3 {abcdehi}  
T4 {abe}  
T5 {bcdehi}

- $G(\{abcde\}_3) = \{\{ac\}_3, \{ad\}_3\}$ .
- Thus,  $\{abcde\}$  is not a (key) generator.
- Question: **Can  $\{abcde\}$  be a (key) generator?**
- Answer:  
 $\{ac\}$ , which is a generator of  $\{abcde\}$ , has the same support (3) as  $\{abcde\}$ , so  $t(\{ac\}) = t(\{abcde\})$ .  
 - Now,  $t(\{abcde\}) = t(\{abcde\}) \cap t(\{e\}) = t(\{ac\}) \cap t(\{e\}) = t(\{ace\})$ .  
 - So,  $sup(\{abcde\}) = sup(\{ace\})$  and  $\{abcde\} \supset \{ace\}$ .  
 - Therefore,  $\{abcde\}$  is not a key generator.

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### Supersets of Non-(Key) Generators

#### Theorem.

If  $X \notin G$ , then  $\forall Y \supset X, Y \notin G$ .

**Proof.** Let  $X \notin G, Y \supset X$ . Then:

$\exists X' \in G(X)$  such that  $X' \subset X$  and

$\exists Z \neq \emptyset$  such that  $Z = Y \setminus X$

$\Rightarrow t(X') = t(X)$  and

$t(Y) = t(X \cup Z) = t(X) \cap t(Z) = t(X') \cap t(Z) = t(X' \cup Z)$

$\Rightarrow sup(Y) = sup(X \cup Z) = sup(X' \cup Z)$  and

$Y = X \cup Z \supset X' \cup Z$

$\Rightarrow Y \notin G$ .

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### Subsets of (Key) Generators

**Theorem A.** If  $X \notin G$ , then  $\forall Y \supset X, Y \notin G$ .

**Theorem B.**

If  $X \in G$ , then  $\forall Y \subset X, Y \in G$ .

**Proof (by contradiction).**

Let  $X \in G, Y \subset X$  and  $Y \notin G$ . Then:

By Theorem A all proper supersets of  $Y$  (and thus also  $X$ ) are not (key) generators.

Hence,  $X \notin G$ , which contradicts the assumption.

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### Properties of (Key) Generators

#### • Important Property of Generators:

The set of all (key) generators is sufficient to determine the support of each itemset  $X$  in  $2^I$ , namely:

$$sup(X) = \min\{sup(Y) \mid Y \text{ is a (key) generator} \wedge Y \subseteq X\}.$$

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### Generators Representation (GR)

- The generators representation (GR) consists of:
  - 1) all frequent (key) generators and the information about their supports,
  - 2) the border of minimal infrequent (key) generators.

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### GR Border...

**Theorem.**  $X$  is a minimal infrequent (key) generator  $\Leftrightarrow X$  is a minimal infrequent itemset.

**Proof ( $\Rightarrow$ ).**  $X$  is a minimal infrequent (key) generator

$\Rightarrow X$  is a minimal infrequent (key) generator & **all proper subsets of  $X$  are (key) generators**

$\Rightarrow X$  is a minimal infrequent (key) generator & **all proper subsets of  $X$  are frequent (key) generators**

$\Rightarrow X$  is infrequent & all proper subsets of  $X$  are frequent

$\Rightarrow X$  is a minimal infrequent itemset.

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**GR Border**

**Theorem.**  $X$  is a minimal infrequent (key) generator  $\Leftrightarrow X$  is a minimal infrequent itemset.

**Proof ( $\Leftarrow$ ).**  $X$  is a minimal infrequent itemset

- $\Rightarrow X$  is infrequent & **proper subsets of  $X$  are frequent**
- $\Rightarrow X$  is infrequent & all proper subsets of  $X$  are frequent & **have supports different from  $sup(X)$**
- $\Rightarrow X$  is an infrequent (key) generator & all its proper subsets are frequent
- $\Rightarrow X$  is an infrequent (key) generator & **all proper subsets of  $X$  are frequent (key) generators**
- $\Rightarrow X$  is a minimal infrequent (key) generator.

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**Example: Reasoning with GR**

- $sup(\{bef\})=?$ ,  $sup(\{bcd\})=?$

**GR**

**Border:**  
f h i

**Frequent generators:**  
ac<sub>3</sub> ad<sub>3</sub>  
a<sub>4</sub> c<sub>4</sub> d<sub>4</sub>  
Ø<sub>5</sub>

- $\{bef\}$  is infrequent, as it is a superset of  $\{f\}$ , which is an infrequent (key) generator.
- $sup(\{bcd\}) = \min(sup(\{c\}_4), sup(\{d\}_4), sup(\{Ø\}_5)) = 4$ .

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**Calculating GR with GR-Apriori**

**Id Transaction**

- T1 {abcde}
- T2 {abcdef}
- T3 {abcdehi}
- T4 {abe}
- T5 {bcdehi}

- The GR-Apriori algorithm discovers GR-representation.
- GR-Apriori uses the fact that GR contains only frequent (key) generators and minimal infrequent (key) generators (that is, minimal infrequent itemsets).
- In consequence, GR-Apriori uses the following properties to generate GR candidates:
  - No proper superset of a non-(key) generator belongs to GR;
  - No proper superset of an infrequent itemset belongs to GR.

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**Calculating GR with GR-Apriori**

**Assumption: minSup = 2**

- Iteration 1:  
 $C_0 \rightarrow FG_0: \emptyset_5$   
 $C_1 \rightarrow FG_1: a_4 b_5 c_4 d_4 e_5 f_1 h_2 i_2$
- Iteration 2:  
 $C_2 \rightarrow FG_2: ac_3 ad_3 cd_4$
- Iteration 3:  
 $C_3: aed$

**Main component (FG)**  
ac<sub>3</sub> ad<sub>3</sub>  
a<sub>4</sub> c<sub>4</sub> d<sub>4</sub>  
Ø<sub>5</sub>

**Border**  
f h i

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