Gamma Language Reference Manual

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1 Introduction

Why GAMMA? - The Core Concept

We propose to implement an elegant yet secure general purpose object-oriented programming language. Interesting features have been selected from the history of object-oriented programming and will be combined with the familiar ideas and style of modern languages.

GAMMA combines three disparate but equally important tenants:

1. Purely object-oriented

GAMMA brings to the table a purely object oriented programming language where every type is modeled as an object–including the standard primitives. Integers, Strings, Arrays, and other types may be expressed in the standard fashion but are objects behind the scenes and can be treated as such.

2. Controllable

GAMMA provides innate security by choosing object level access control as opposed to class level access specifiers. Private members of one object are inaccessible to other objects of the same type. Overloading is not allowed. No subclass can turn your functionality on its head.

3. Versatile

GAMMA allows programmers to place "refinement methods" inside their code. Alone these methods do nothing, but may be defined by subclasses so as to extend functionality at certain important positions. Anonymous instantiation allows for extension of your classes in a quick easy fashion. Generic typing on method parameters allows for the same method to cover a variety of input types.

The Motivation Behind GAMMA

GAMMA is a reaction to the object-oriented languages before it. Obtuse syntax, flaws in security, and awkward implementations plague the average object-oriented language. GAMMA is intended as a step toward ease and comfort as an object-oriented programmer.

The first goal is to make an object-oriented language that is comfortable in its own skin. It should naturally lend itself to constructing API-layers and abstracting general models. It should serve the programmer towards their goal instead of exerting unnecessary effort through verbosity and awkwardness of structure.

The second goal is to make a language that is stable and controllable. The programmer in the lowest abstraction layer has control over how those higher

may procede. Unexpected runtime behavior should be reduced through firmness of semantic structure and debugging should be a straight-forward process due to pure object and method nature of GAMMA.

2 Grammar

```
• Classs may extend another class or default to extending Object
\langle class \rangle \Rightarrow
         class \langle class id \rangle \langle extend \rangle \{ \langle class section \rangle^* \}
\langle \text{extend} \rangle \Rightarrow
     | extends (class id)
• Sections – private protected public refinements and main
\langle class section \rangle \Rightarrow
         ⟨refinement⟩
        (access group)
        \langle main \rangle
• Refinements are named method dot refinement
\langle \text{refinement} \rangle \Rightarrow
        refinement { \langle refine \rangle^* }
\langle \text{refine} \rangle \Rightarrow
         \langle \text{return type} \rangle \langle \text{var id} \rangle \langle \text{var id} \rangle \langle \text{params} \rangle  { \langle \text{statement} \rangle^* }
• Access groups contain all the members of a class
\langle access group \rangle \Rightarrow
         \langle access type \rangle \{ \langle member \rangle^* \}
\langle access type \rangle \Rightarrow
        private
      protected
       public
\langle \text{member} \rangle \Rightarrow
         (var decl)
     |\langle \text{method} \rangle|
     |\langle init \rangle|
\langle \text{method} \rangle \Rightarrow
         \langle \text{return type} \rangle \langle \text{var id} \rangle \langle \text{params} \rangle  { \langle \text{statement} \rangle^* }
\langle \text{init} \rangle \Rightarrow
        init \langle params \rangle \{ \langle statement \rangle^* \}
```

• Main is special - not instance data starts execution

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\langle \text{main} \rangle \Rightarrow
        main (String[] (var id) ) { (statement)* }
• Finally the meat and potatoes
\langle statement \rangle \Rightarrow
         \langle \text{var decl} \rangle;
         \langle \text{super} \rangle;
         \langle \text{return} \rangle;
         \langle conditional \rangle
         \langle loop \rangle
         ⟨expression⟩;
• Super invocation is so we can do constructor chaining
\langle \text{super} \rangle \Rightarrow
        super ( \langle args \rangle )
• Methods need to be able to return something too
\langle {\rm return} \rangle \Rightarrow
        return (expression)
ullet Basic control structures
\langle conditional \rangle \Rightarrow
        if (\langle expression \rangle) { \langle statement \rangle^* } \langle else \rangle
\langle \text{else} \rangle \Rightarrow
     | \langle elseif \rangle  else \{ \langle statement \rangle^* \}
\langle elseif \rangle \Rightarrow
     |\langle elseif \rangle elsif (\langle expression \rangle) \{\langle statement \rangle^* \}
\langle loop \rangle \Rightarrow
         while ( \langle expression \rangle ) { \langle statement \rangle^* }
• Anything that can result in a value
\langle expression \rangle \Rightarrow
         \langle assignment \rangle
         (invocation)
         \langle field \rangle
         \langle deref \rangle
         \langle arithmetic \rangle
         \langle \text{test} \rangle
         ⟨instantiate⟩
        (refine expr)
```

```
\langle literal \rangle
          ( \( \text{\text{expression}} \) )
         null
• Assignment – putting one thing in another
\langle assignment \rangle \Rightarrow
          \langle expression \rangle := \langle expression \rangle
• Member / data access
\langle \text{invocation} \rangle \Rightarrow
          \langle expression \rangle . \langle invoke \rangle
      | (invoke)
\langle \mathrm{invoke} \rangle \Rightarrow
          \langle \text{var id} \rangle ()
      |\langle \text{var id} \rangle (\langle \text{args} \rangle)|
\langle \text{field} \rangle \Rightarrow
          \langle \text{expression} \rangle . \langle \text{var id} \rangle
\langle \mathrm{deref} \rangle \Rightarrow
          (expression) [ (expression) ]
• Basic arithmetic can and will be done!
\langle \operatorname{arithmetic} \rangle \Rightarrow
          \langle expression \rangle \langle bin op \rangle \langle expression \rangle
      |\langle unary op \rangle \langle expression \rangle
\langle \text{bin op} \rangle \Rightarrow
\langle \text{unary op} \rangle \Rightarrow
\bullet \ Common \ boolean \ predicates
\langle \text{test} \rangle \Rightarrow
          \langle expression \rangle \langle bin pred \rangle \langle expression \rangle
          ⟨unary pred⟩⟨expression⟩
      | refinable (\langle var id \rangle)
\langle \text{bin pred} \rangle \Rightarrow
          and
          \mathbf{or}
         \mathbf{xor}
```

```
nand
         nor
         <
         \leq =
         =
         !=
         >=
        >
\langle \text{unary pred} \rangle \Rightarrow
• Making something
\langle instantiate \rangle \Rightarrow
          \langle object\ instantiate \rangle
     | (array instantiate)
\langle \text{object instantiate} \rangle \Rightarrow
         \mathbf{new} \langle \text{class id} \rangle
      \mid new \langleclass id\rangle ( \langleargs\rangle )
\langle \text{array instantiate} \rangle \Rightarrow
         \mathbf{new} \langle \mathrm{type} \rangle [\langle \mathrm{expression} \rangle]
• Refinement takes a specific specialization and notes the required
return type
\langle \text{refine expr} \rangle \Rightarrow
         refine (specialize) to (type)
\langle {\rm specialize} \rangle \Rightarrow
          \langle \text{var id} \rangle ()
     |\langle \text{var id} \rangle (\langle \text{args} \rangle)|
• Literally necessary
\langle \text{literal} \rangle \Rightarrow
          \langle \text{int lit} \rangle
         ⟨bool lit⟩
         (float lit)
       (string lit)
\langle \text{float lit} \rangle \Rightarrow
          \langle digit \rangle + . \langle digit \rangle +
\langle \text{int lit} \rangle \Rightarrow
          \langle digits \rangle +
\langle \text{bool lit} \rangle \Rightarrow
         \mathbf{true}
      false
\langle \text{string lit} \rangle \Rightarrow
```

```
• Params and args are as expected
\langle params \rangle \Rightarrow
           ( \(\langle\) paramlist \(\rangle\) )
\langle \mathrm{paramlist} \rangle \Rightarrow
           \langle var decl \rangle
      \mid \langle paramlist \rangle, \langle var decl \rangle
\langle args \rangle \Rightarrow
           \langle {\rm expression} \rangle
      |\langle args \rangle, \langle expression \rangle
• All the basic stuff we've been saving up until now
\langle \text{var decl} \rangle \Rightarrow
            \langle {\rm type} \rangle \langle {\rm var~id} \rangle
\langle \text{return type} \rangle \Rightarrow
           unit
      |\langle type \rangle
\langle \text{type} \rangle \Rightarrow
            \langle class id \rangle
      |\langle \text{type} \rangle []
\langle \text{class id} \rangle \Rightarrow
           \langle upper \rangle \langle ualphanum \rangle^*
\langle \text{var id} \rangle \Rightarrow
            \langle lower \rangle \langle ualphanum \rangle^*
```

" $\langle \text{string escape seq} \rangle$ "