

We take the common nonterminals *digit*, *digits*, *lower*, *upper*, *alphanum*, and *alphanums* to be predefined with their obvious meaning (plurality such as *digits* implies at least one digit). Then the grammar for the language is:

- *Class may extend another class or default to extending Object*

$\langle \text{class} \rangle \Rightarrow$
 $\quad \mathbf{class} \langle \text{class id} \rangle \langle \text{extend} \rangle \langle \text{scope} \rangle \langle \text{class decls} \rangle$
 $\langle \text{extend} \rangle \Rightarrow$
 $\quad \epsilon$
 $\quad | \mathbf{extends} \langle \text{class id} \rangle$

- *A class can be split up into five sections – private protected public refinements and main*

$\langle \text{class decls} \rangle \Rightarrow$
 $\quad \langle \text{class decl} \rangle$
 $\quad | \langle \text{class decls} \rangle \langle \text{class decl} \rangle$
 $\langle \text{class decl} \rangle \Rightarrow$
 $\quad \langle \text{refinement} \rangle$
 $\quad | \langle \text{access group} \rangle$
 $\quad | \langle \text{main} \rangle$

- *Refinements are named method dot refinement*

$\langle \text{refinement} \rangle \Rightarrow$
 $\quad \mathbf{refinement} \langle \text{scope} \rangle \langle \text{refines} \rangle$
 $\langle \text{refines} \rangle \Rightarrow$
 $\quad \langle \text{refine} \rangle$
 $\quad | \langle \text{refines} \rangle \langle \text{refine} \rangle$
 $\langle \text{refine} \rangle \Rightarrow$
 $\quad \langle \text{refine sig} \rangle \langle \text{scope} \rangle \langle \text{statements} \rangle$
 $\langle \text{refine sig} \rangle \Rightarrow$
 $\quad \langle \text{return type} \rangle \langle \text{var id} \rangle . \langle \text{var id} \rangle \langle \text{params} \rangle$

- *Access groups contain all the members of a class*

$\langle \text{access group} \rangle \Rightarrow$
 $\quad \langle \text{access type} \rangle \langle \text{scope} \rangle \langle \text{members} \rangle$
 $\langle \text{access type} \rangle \Rightarrow$
 $\quad \mathbf{private}$
 $\quad | \mathbf{protected}$
 $\quad | \mathbf{public}$
 $\langle \text{members} \rangle \Rightarrow$
 $\quad \langle \text{member} \rangle$
 $\quad | \langle \text{members} \rangle \langle \text{member} \rangle$
 $\langle \text{member} \rangle \Rightarrow$

$$\begin{aligned}
& \langle \text{ivar} \rangle \\
& \quad | \quad \langle \text{method} \rangle \\
& \quad | \quad \langle \text{init} \rangle \\
\langle \text{ivar} \rangle & \Rightarrow \\
& \quad \langle \text{var decl} \rangle \\
\langle \text{method} \rangle & \Rightarrow \\
& \quad \langle \text{method sig} \rangle \langle \text{scope} \rangle \langle \text{statements} \rangle \\
\langle \text{method sig} \rangle & \Rightarrow \\
& \quad \langle \text{return type} \rangle \langle \text{var id} \rangle \\
& \quad | \quad \langle \text{params} \rangle \\
\langle \text{init} \rangle & \Rightarrow \\
& \quad \langle \text{init sig} \rangle \langle \text{scope} \rangle \langle \text{statements} \rangle \\
\langle \text{init sig} \rangle & \Rightarrow \\
& \quad \mathbf{init} \quad \langle \text{params} \rangle
\end{aligned}$$

- *The main is special – it belongs to no object and is what can be used to start an execution*

$$\begin{aligned}
\langle \text{main} \rangle & \Rightarrow \\
& \quad \langle \text{main sig} \rangle \langle \text{scope} \rangle \langle \text{statements} \rangle \\
\langle \text{main sig} \rangle & \Rightarrow \\
& \quad \mathbf{main} \\
& \quad | \quad \mathbf{main} \ (\ \mathbf{String}[] \ \langle \text{var id} \rangle \)
\end{aligned}$$

- *Finally the meat and potatoes*

$$\begin{aligned}
\langle \text{statements} \rangle & \Rightarrow \\
& \quad \langle \text{statement} \rangle \\
& \quad | \quad \langle \text{statements} \rangle \langle \text{statement} \rangle \\
\langle \text{statement} \rangle & \Rightarrow \\
& \quad \langle \text{local} \rangle \\
& \quad | \quad \langle \text{refine stmt} \rangle \\
& \quad | \quad \langle \text{assignment} \rangle \\
& \quad | \quad \langle \text{conditional} \rangle \\
& \quad | \quad \langle \text{loop} \rangle \\
& \quad | \quad \langle \text{expression} \rangle
\end{aligned}$$

- *Local variables can be uninitialized (end up null) or initialized*

$$\begin{aligned}
\langle \text{local} \rangle & \Rightarrow \\
& \quad \langle \text{var decl} \rangle \\
& \quad | \quad \langle \text{var decl} \rangle \mathbf{:} = \langle \text{expression} \rangle
\end{aligned}$$

- *Refine statements need not return type information*

$$\langle \text{refine stmt} \rangle \Rightarrow$$

refine $\langle \text{specialize} \rangle$

- *Assignment works as in most languages – lvalues receive the results of expressions*

$\langle \text{assignment} \rangle \Rightarrow$
 $\langle \text{lvalue} \rangle := \langle \text{expression} \rangle$
 $\langle \text{lvalue} \rangle \Rightarrow$
 $\langle \text{var id} \rangle$
 | **this** . $\langle \text{var id} \rangle$
 | $\langle \text{var id} \rangle$. $\langle \text{var id} \rangle$

- *If and else are the same as in every other language*

$\langle \text{conditional} \rangle \Rightarrow$
 $\langle \text{if} \rangle$
 | $\langle \text{if} \rangle \langle \text{else} \rangle$
 $\langle \text{if} \rangle \Rightarrow$
 if ($\langle \text{expression} \rangle$) $\langle \text{scope} \rangle \langle \text{statements} \rangle$
 $\langle \text{else} \rangle \Rightarrow$
 else $\langle \text{scope} \rangle$

- *C and java style loop constructs*

$\langle \text{loop} \rangle \Rightarrow$
 $\langle \text{loop head} \rangle \langle \text{scope} \rangle \langle \text{statements} \rangle$
 $\langle \text{loop head} \rangle \Rightarrow$
 for ($\langle \text{assignment} \rangle$, $\langle \text{expression} \rangle$, $\langle \text{assignment} \rangle$)
 | **while** ($\langle \text{expression} \rangle$)

- *An expression is anything that can result in a value – note that assignment is not an expression (should we change?)*

$\langle \text{expression} \rangle \Rightarrow$
 $\langle \text{invocation} \rangle$
 | $\langle \text{arithmetic} \rangle$
 | $\langle \text{array expression} \rangle$
 | $\langle \text{test} \rangle$
 | $\langle \text{instantiate} \rangle$
 | $\langle \text{refine expr} \rangle$
 | $\langle \text{literal} \rangle$
 | $\langle \text{var id} \rangle$
 | ($\langle \text{expression} \rangle$)
 | **null**

- *Method invocations always have a receiver*

$\langle \text{invocation} \rangle \Rightarrow$
 $\langle \text{receiver} \rangle . \langle \text{invoke} \rangle$
 $\langle \text{receiver} \rangle \Rightarrow$
 \mathbf{this}
 $| \langle \text{var id} \rangle$
 $\langle \text{invoke} \rangle \Rightarrow$
 $\langle \text{var id} \rangle ()$
 $| \langle \text{var id} \rangle (\langle \text{args} \rangle)$

- *Basic arithmetic can and will be done!*

$\langle \text{arithmetic} \rangle \Rightarrow$
 $\langle \text{expression} \rangle \langle \text{bin op} \rangle \langle \text{expression} \rangle$
 $| \langle \text{unary op} \rangle \langle \text{expression} \rangle$
 $\langle \text{bin op} \rangle \Rightarrow$
 $+$
 $| -$
 $| *$
 $| /$
 $| \%$
 $\langle \text{unary op} \rangle \Rightarrow$
 $-$

- *Build an array on the fly (should we leave this out? TODO – ARRAY ASSIGNMENT)*

$\langle \text{array expression} \rangle \Rightarrow$
 $[\langle \text{args} \rangle]$

- *Common boolean predicates*

$\langle \text{test} \rangle \Rightarrow$
 $\langle \text{expression} \rangle \langle \text{bin pred} \rangle \langle \text{expression} \rangle$
 $| \langle \text{unary pred} \rangle \langle \text{expression} \rangle$
 $| \mathbf{refinable} (\langle \text{var id} \rangle)$
 $\langle \text{bin pred} \rangle \Rightarrow$
 \mathbf{and}
 $| \mathbf{or}$
 $| \mathbf{xor}$
 $| <$
 $| <=$
 $| =$
 $| !=$
 $| >=$
 $| >$
 $\langle \text{unary pred} \rangle \Rightarrow$

!

- *Making something*

$\langle \text{instantiate} \rangle \Rightarrow$
 $\langle \text{object instantiate} \rangle$
 | $\langle \text{array instantiate} \rangle$
 $\langle \text{object instantiate} \rangle \Rightarrow$
 new $\langle \text{class id} \rangle$
 | **new** $\langle \text{class id} \rangle$ ($\langle \text{args} \rangle$)
 $\langle \text{array instantiate} \rangle \Rightarrow$
 new $\langle \text{type} \rangle$ [$\langle \text{digits} \rangle$]

- *Refinement takes a specific specialization and notes the required return type*

$\langle \text{refine expr} \rangle \Rightarrow$
 refine $\langle \text{specialize} \rangle$ **to** $\langle \text{type} \rangle$
 $\langle \text{specialize} \rangle \Rightarrow$
 $\langle \text{var id} \rangle$ ()
 | $\langle \text{var id} \rangle$ ($\langle \text{args} \rangle$)

- *Literally necessary*

$\langle \text{literal} \rangle \Rightarrow$
 $\langle \text{int lit} \rangle$
 | $\langle \text{bool lit} \rangle$
 | $\langle \text{float lit} \rangle$
 | $\langle \text{string lit} \rangle$
 $\langle \text{float lit} \rangle \Rightarrow$
 $\langle \text{digits} \rangle$. $\langle \text{digits} \rangle$
 $\langle \text{int lit} \rangle \Rightarrow$
 $\langle \text{sign} \rangle \langle \text{digits} \rangle$
 $\langle \text{bool lit} \rangle \Rightarrow$
 true
 | **false**
 $\langle \text{string lit} \rangle \Rightarrow$
 “ $\langle \text{string escape seq} \rangle$ ”

- *Params and args are as expected*

$\langle \text{params} \rangle \Rightarrow$
 ϵ
 | ($\langle \text{paramlist} \rangle$)
 $\langle \text{paramlist} \rangle \Rightarrow$
 $\langle \text{var decl} \rangle$
 | $\langle \text{paramlist} \rangle$, $\langle \text{var decl} \rangle$

$$\langle \text{args} \rangle \Rightarrow$$

$$\langle \text{expression} \rangle$$

$$| \langle \text{args} \rangle , \langle \text{expression} \rangle$$

- *All the basic stuff we've been saving up until now*

$$\langle \text{var decl} \rangle \Rightarrow$$

$$\langle \text{type} \rangle \langle \text{var id} \rangle$$

$$\langle \text{return type} \rangle \Rightarrow$$

$$\mathbf{unit}$$

$$| \langle \text{type} \rangle$$

$$\langle \text{type} \rangle \Rightarrow$$

$$\langle \text{class id} \rangle$$

$$| \langle \text{type} \rangle []$$

$$\langle \text{class id} \rangle \Rightarrow$$

$$\langle \text{upper} \rangle$$

$$| \langle \text{upper} \rangle \langle \text{alphanums} \rangle$$

$$\langle \text{var id} \rangle \Rightarrow$$

$$\langle \text{lower} \rangle$$

$$| \langle \text{lower} \rangle \langle \text{alphanums} \rangle$$