

CSC 225

Algorithms and Data Structures: I

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ECS 516

Building up a heap bottom-up

- n standard insert-operations for a heap result in $O(n \log(n))$ time.
- Can we build up a heap for n given elements faster? Is $O(n)$ possible?

Bottom-Up Heap

(pgs. 176-178)

Algorithm BottomUpHeap(S):

Input: A list S storing m keys

Output: A heap T storing the m keys

if S is empty **then**

return external node

remove the first key, k , from S

split S in half, lists S_1 and S_2

$T_1 \leftarrow \text{BottomUpHeap}(S_1)$

$T_2 \leftarrow \text{BottomUpHeap}(S_2)$

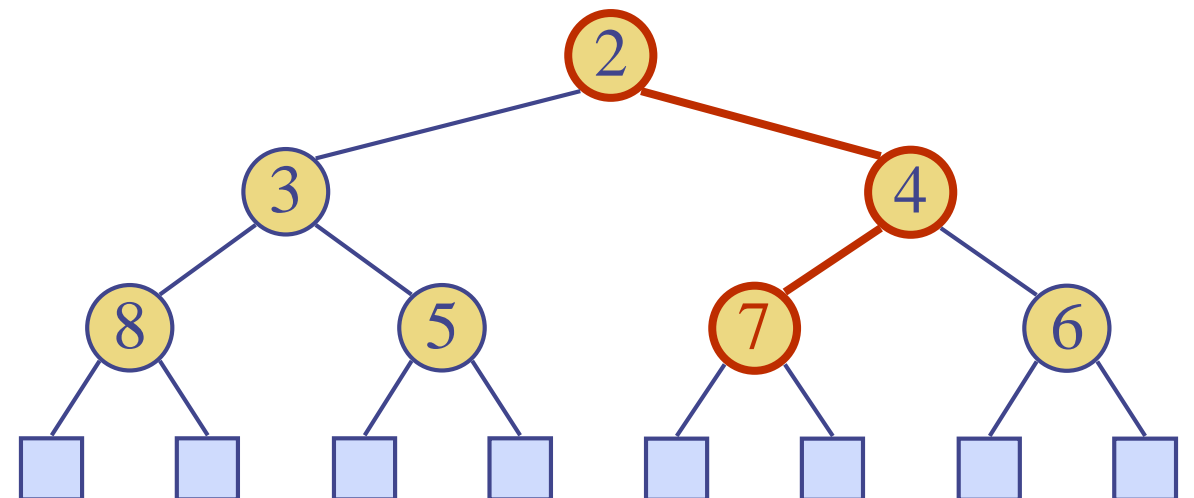
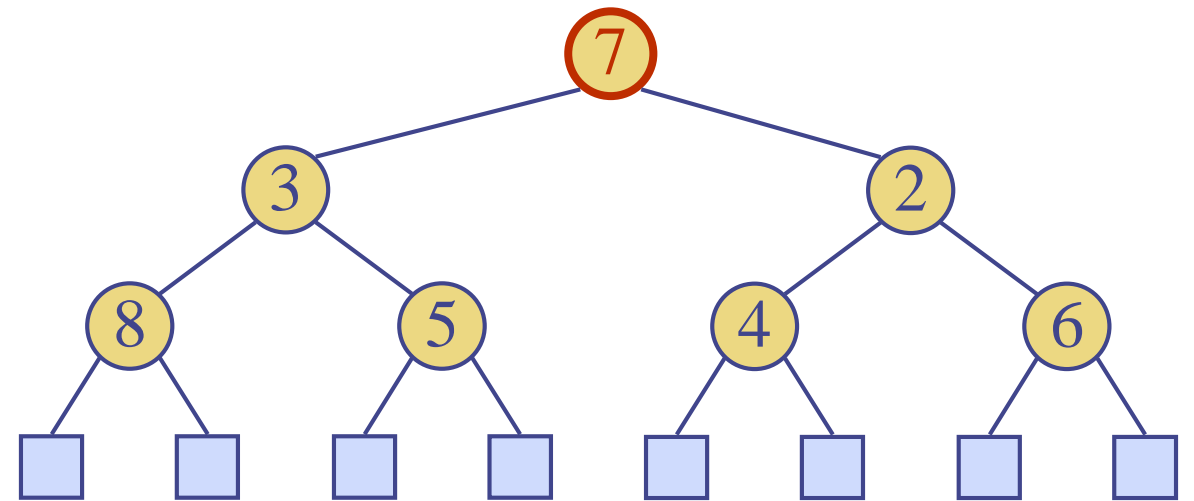
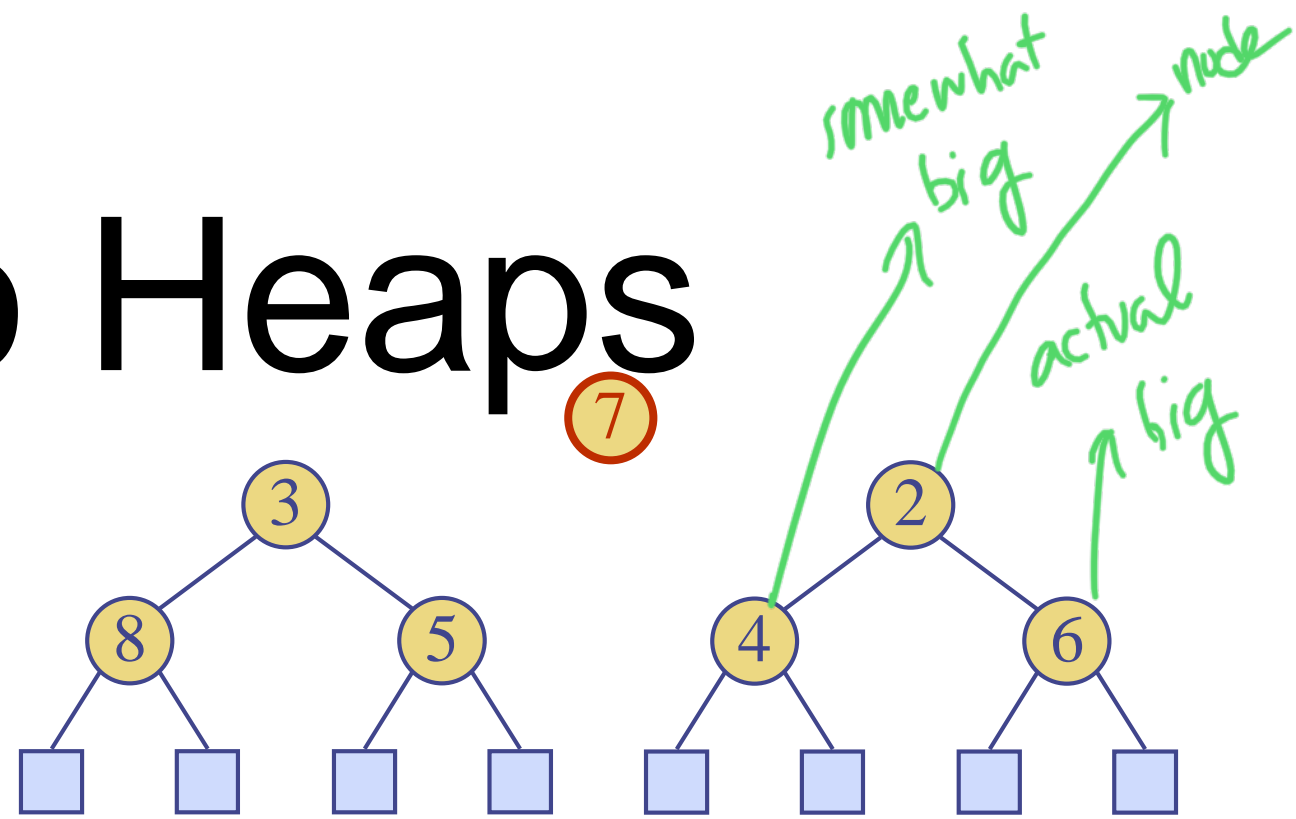
$T \leftarrow \text{merge}(k, T_1, T_2)$

DownHeap(T, root)

return T

Merging Two Heaps

- We are given two heaps and a key k
- We create a new heap with the root node storing k and with the two heaps as subtrees
- We perform downheap to restore the heap-order property

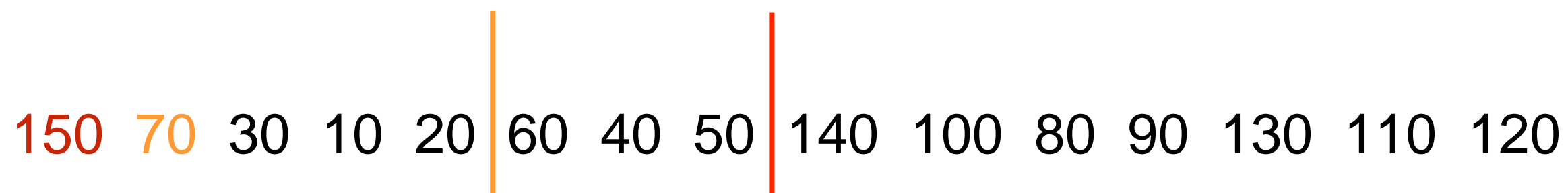


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150 70 30 10 20 60 40 50 | 140 100 80 90 130 110 120

150 70 30 10 20 60 40 50 | 140 100 80 90 130 110 120

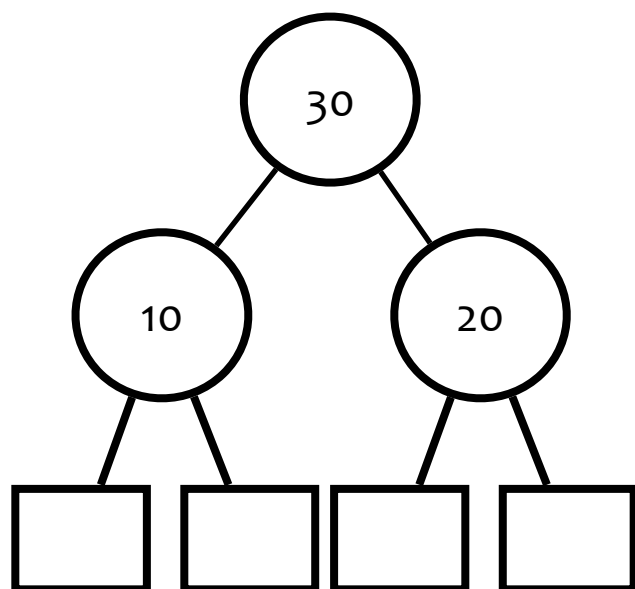
150 70 30 10 20 60 40 50 140 100 80 90 130 110 120



A horizontal number line with 15 tick marks. The numbers are: 150 (red), 70 (orange), 30, 10, 20, 60, 40, 50, 140, 100, 80, 90, 130, 110, 120. There are two vertical lines: an orange one between 20 and 60, and a red one between 50 and 140.

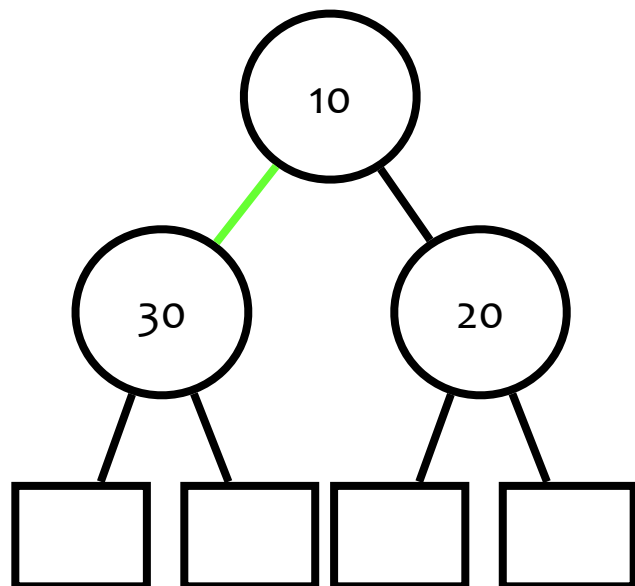
150 70 30 10 20 60 40 50 140 100 80 90 130 110 120

Three vertical lines are positioned between the numbers: a green line between 10 and 20, an orange line between 20 and 60, and a red line between 50 and 140.

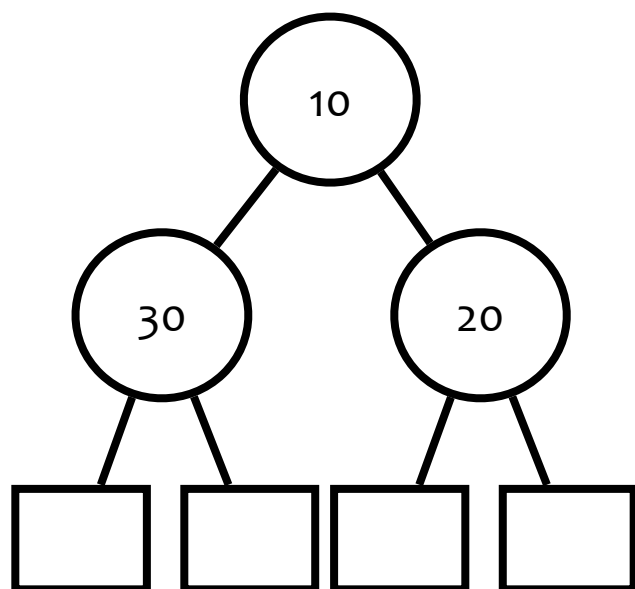


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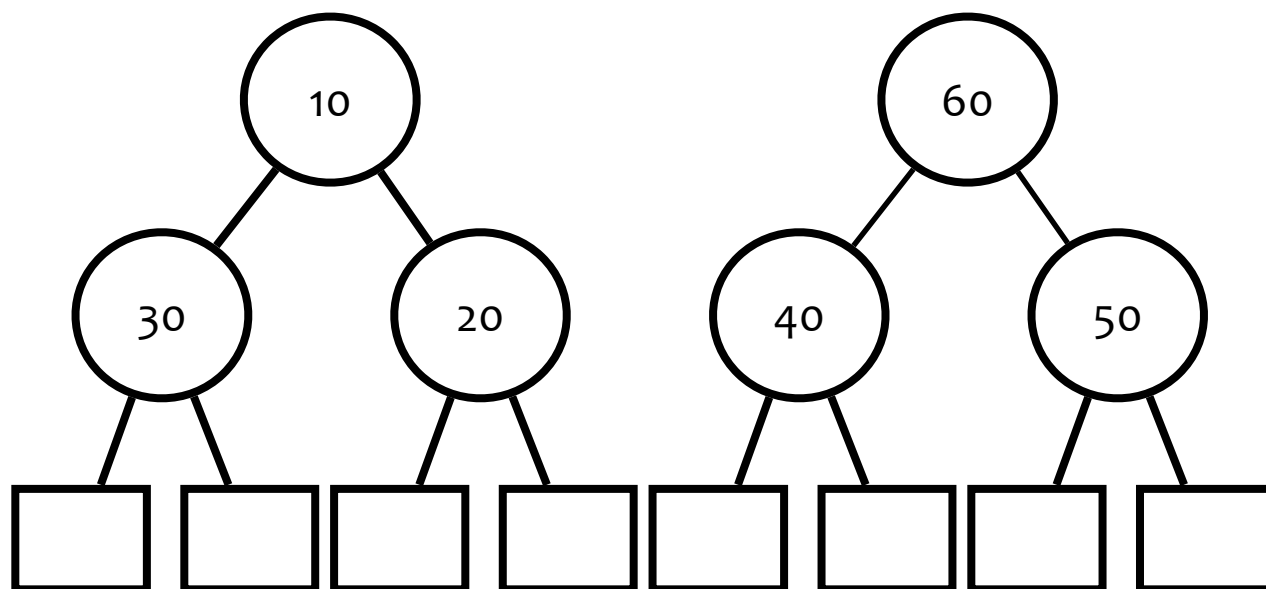
Three vertical lines are positioned between the numbers: a green line between 10 and 20, an orange line between 20 and 60, and a red line between 50 and 140.



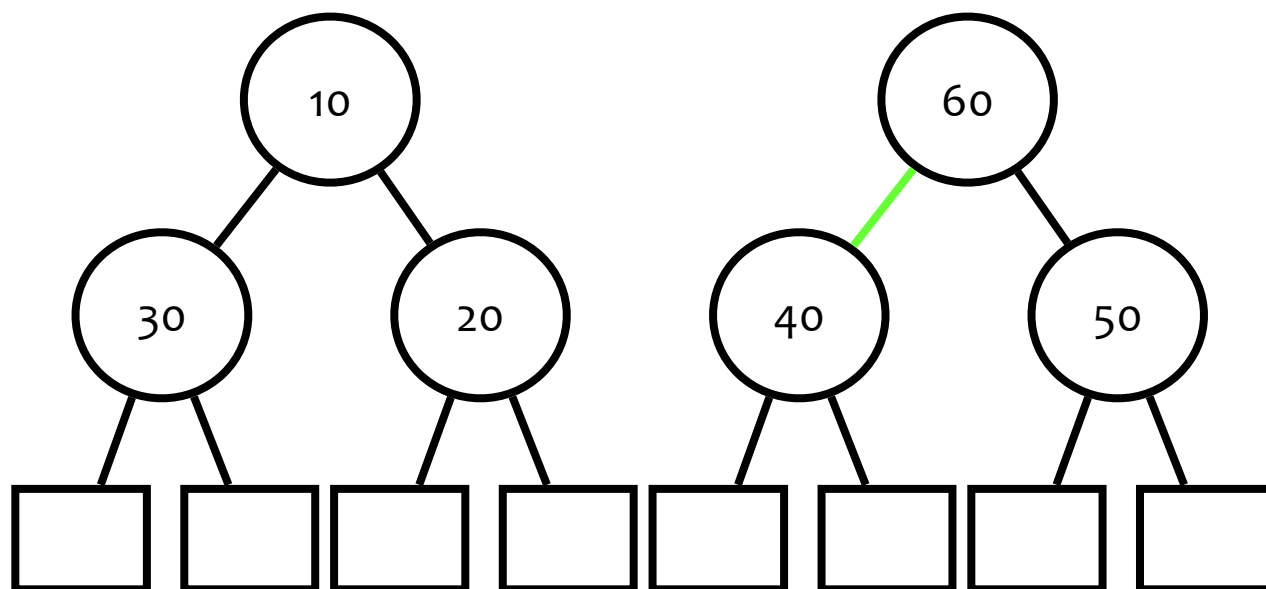
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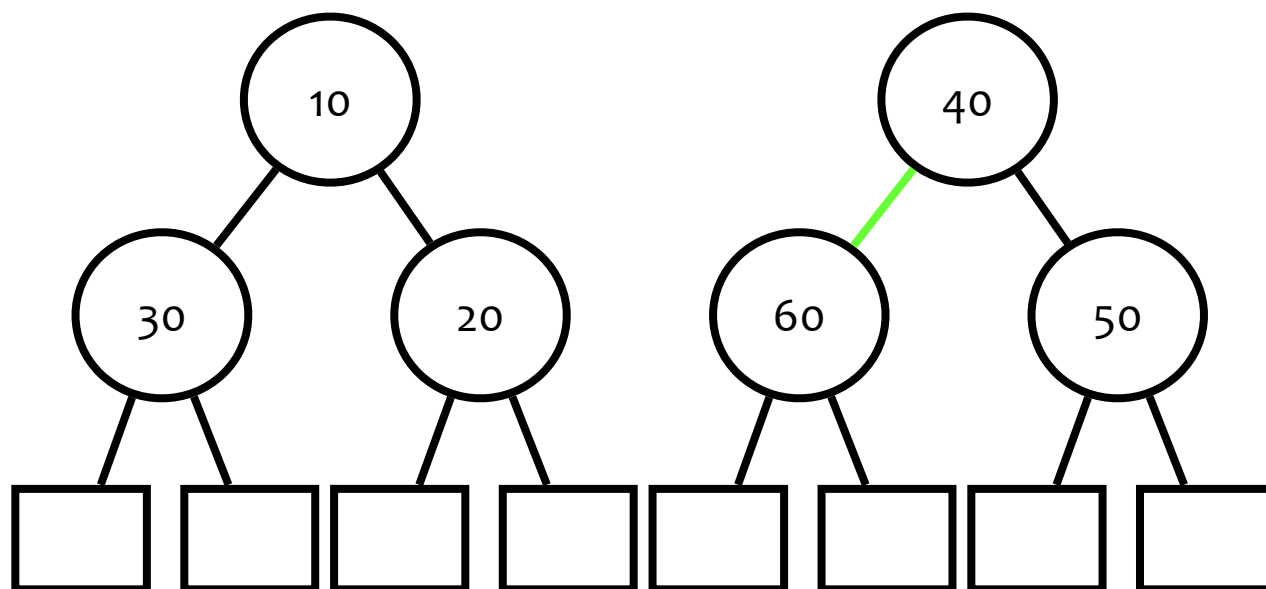
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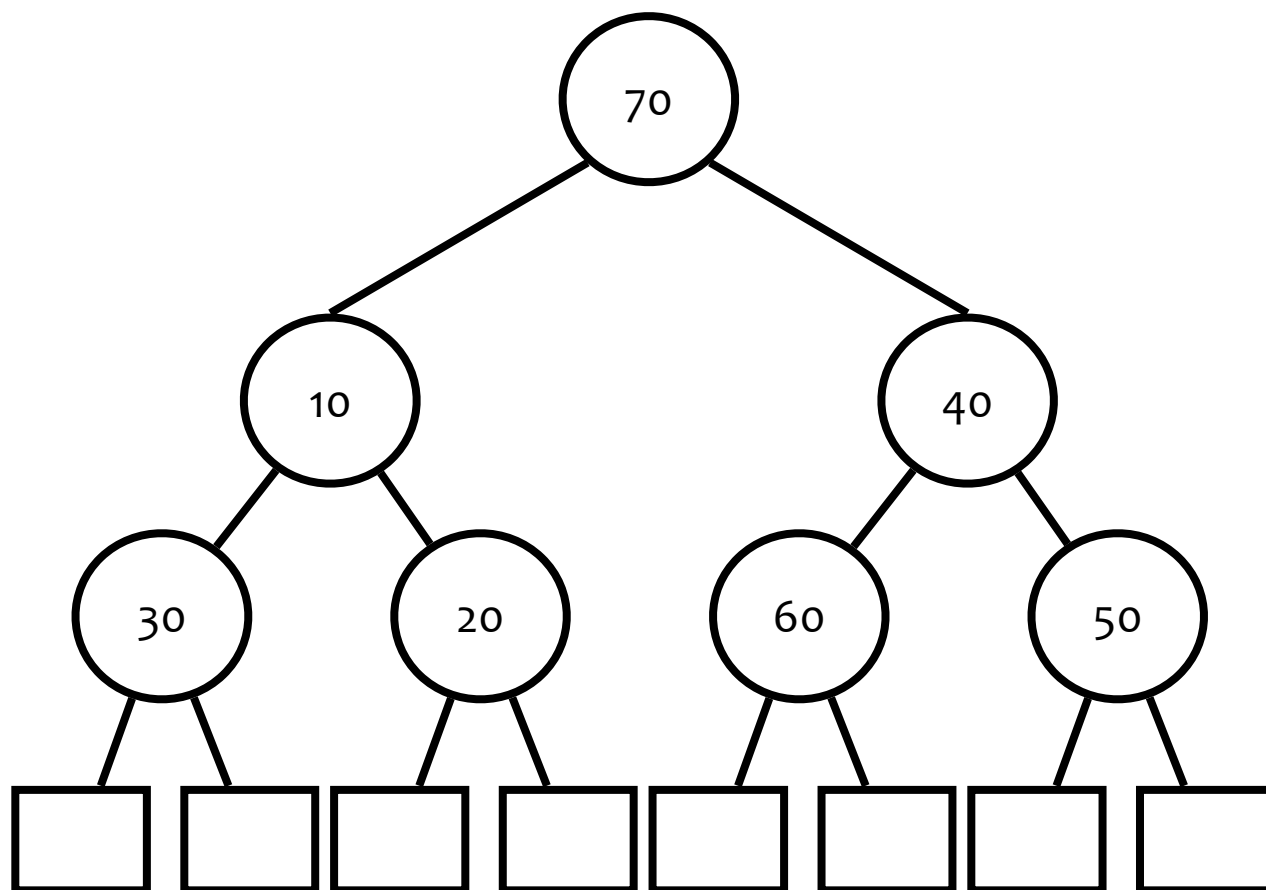
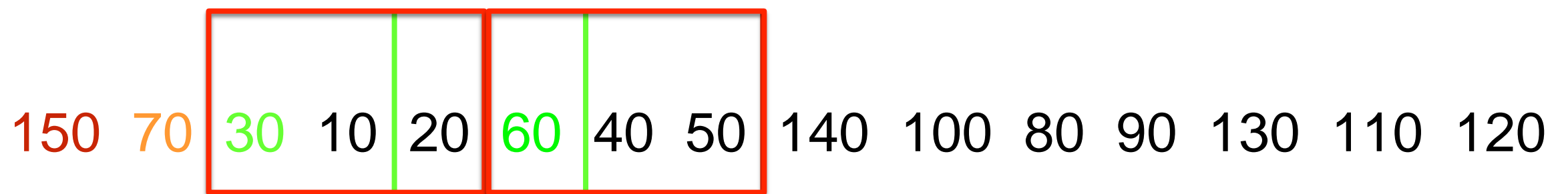


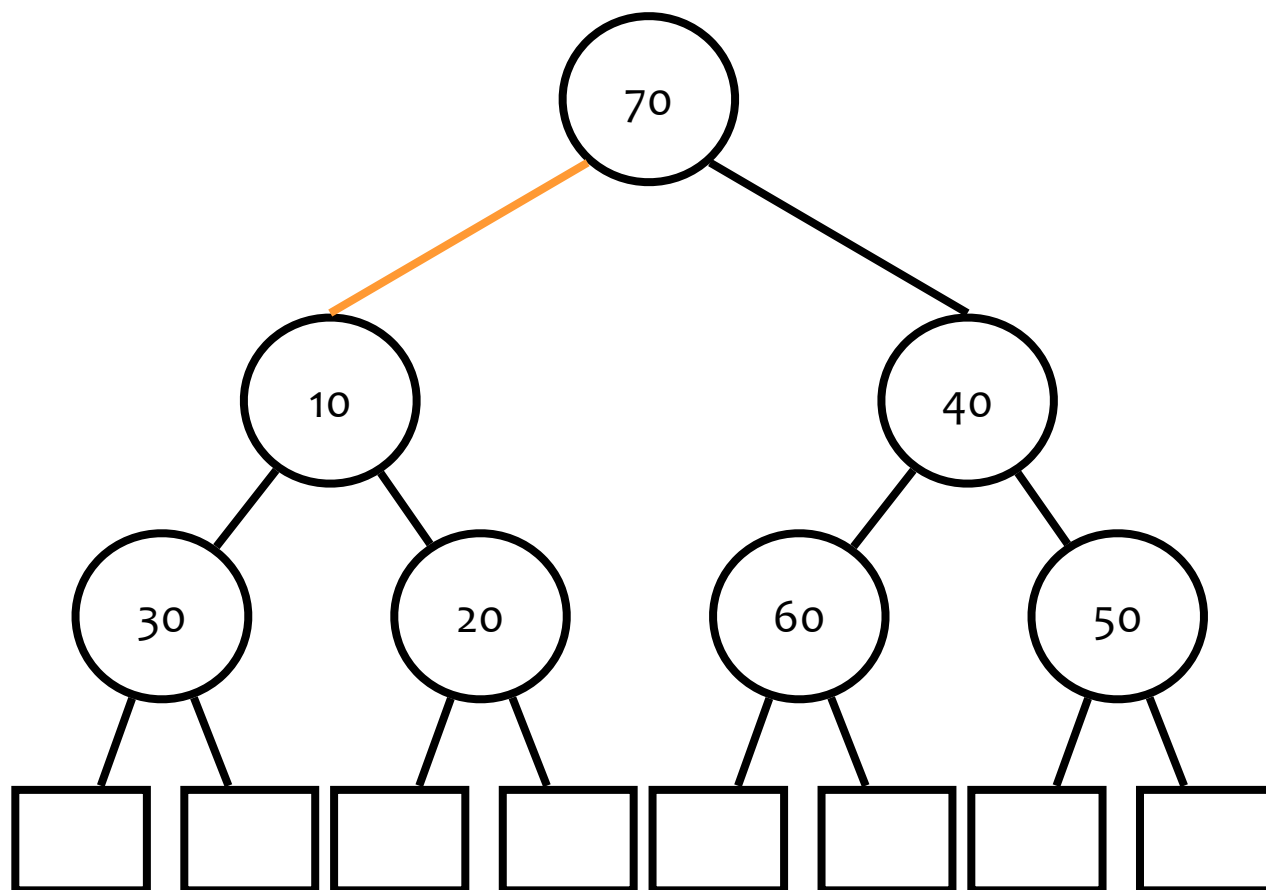
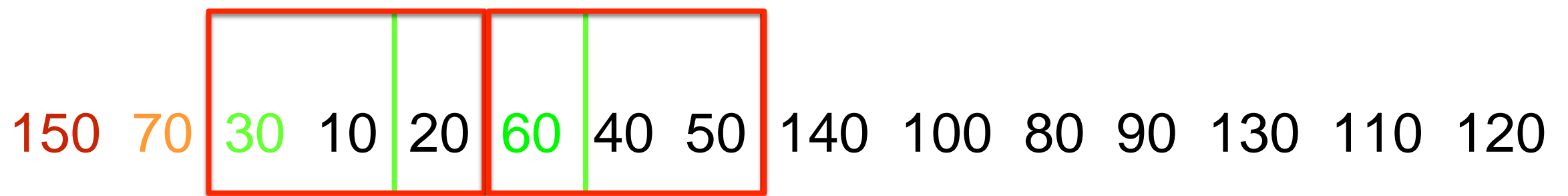
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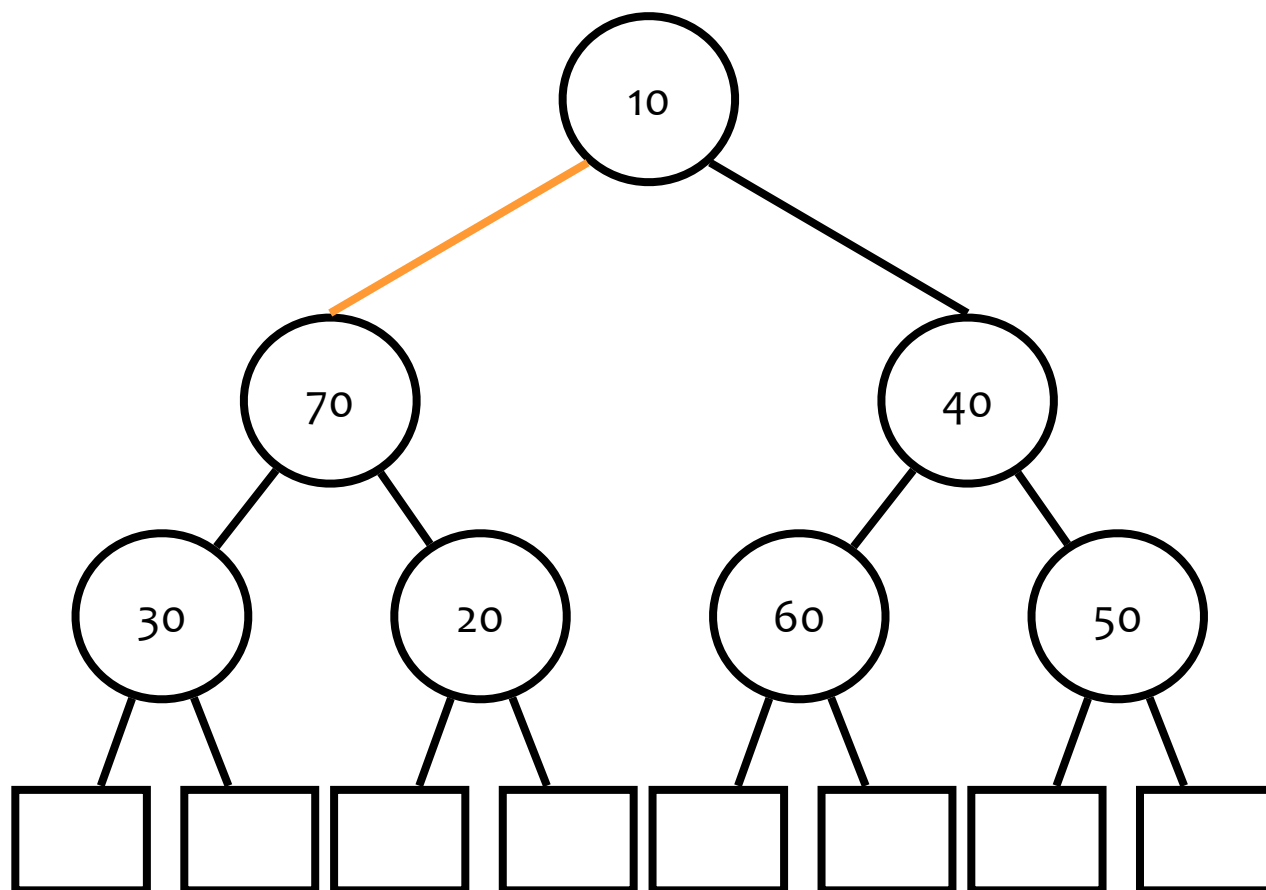
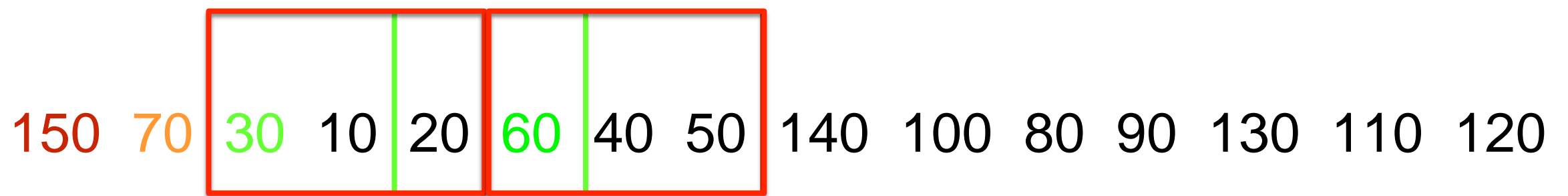


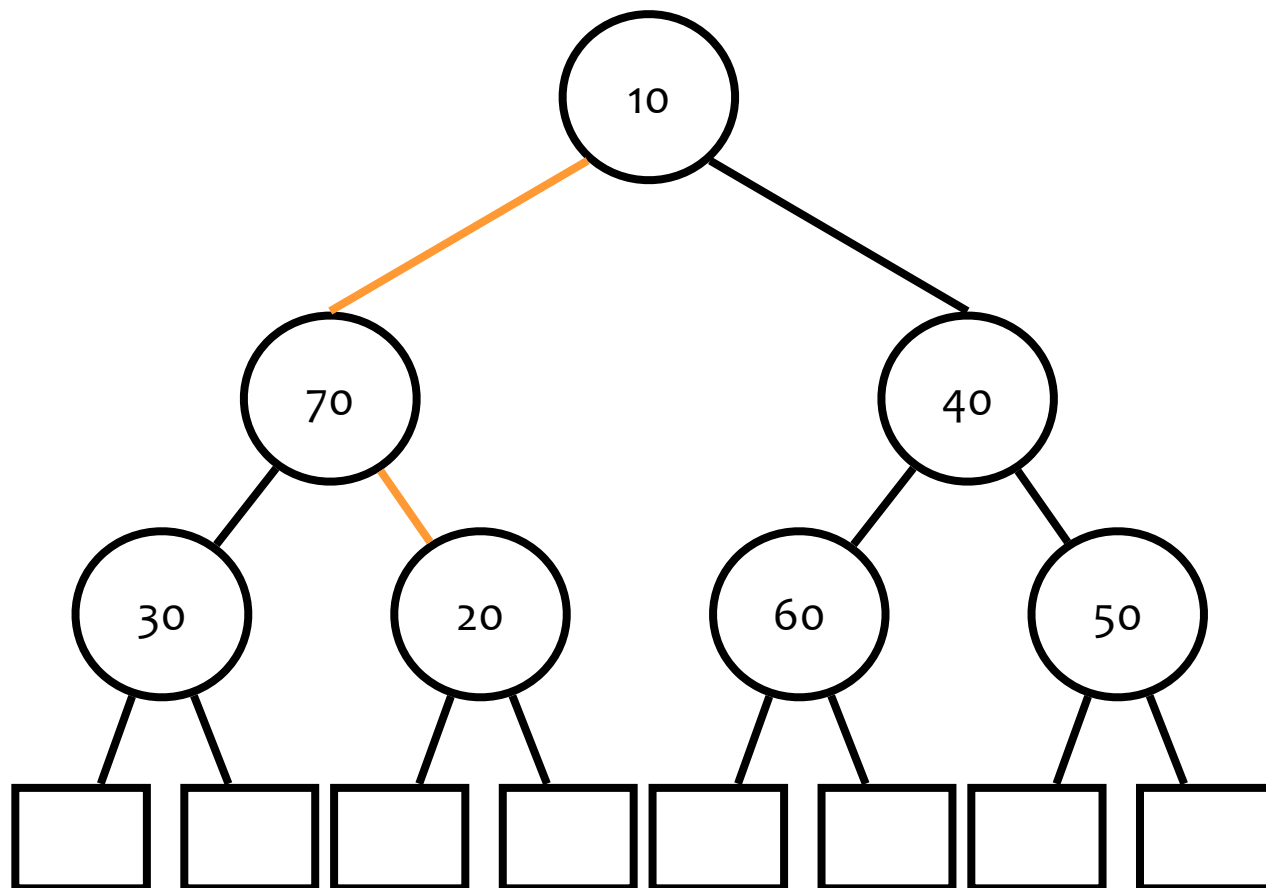
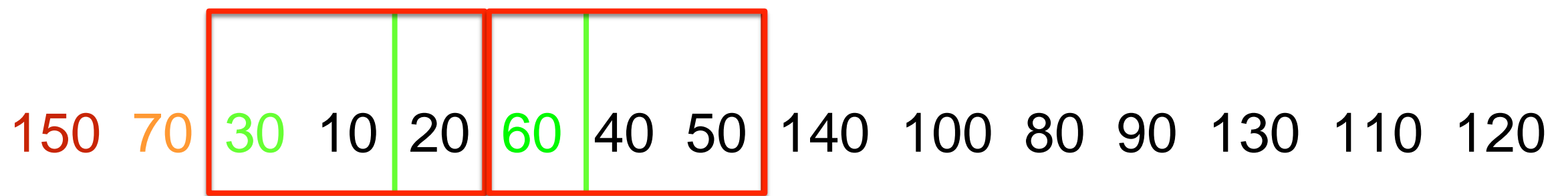
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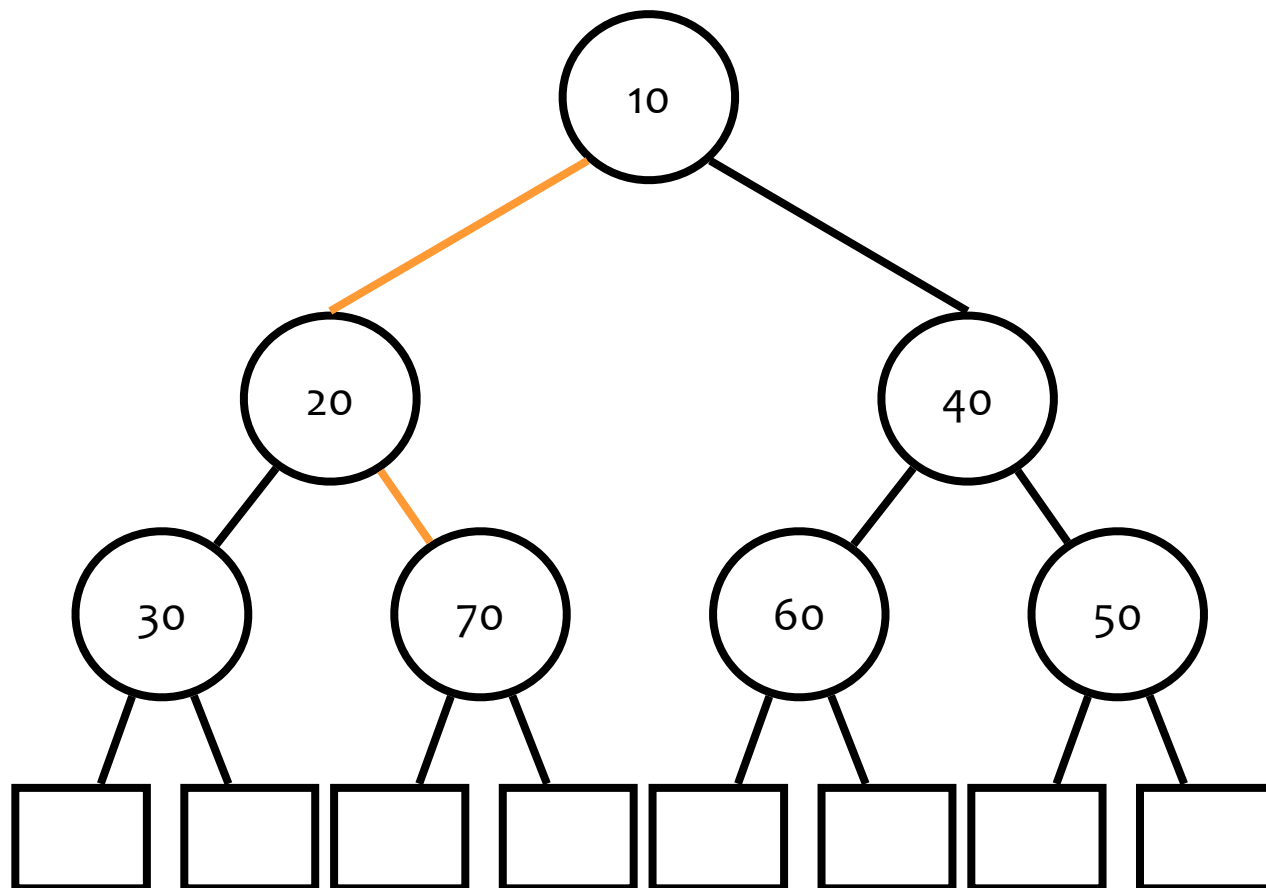
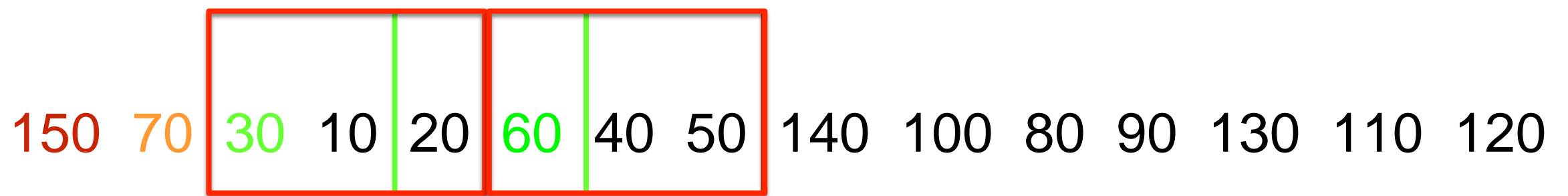


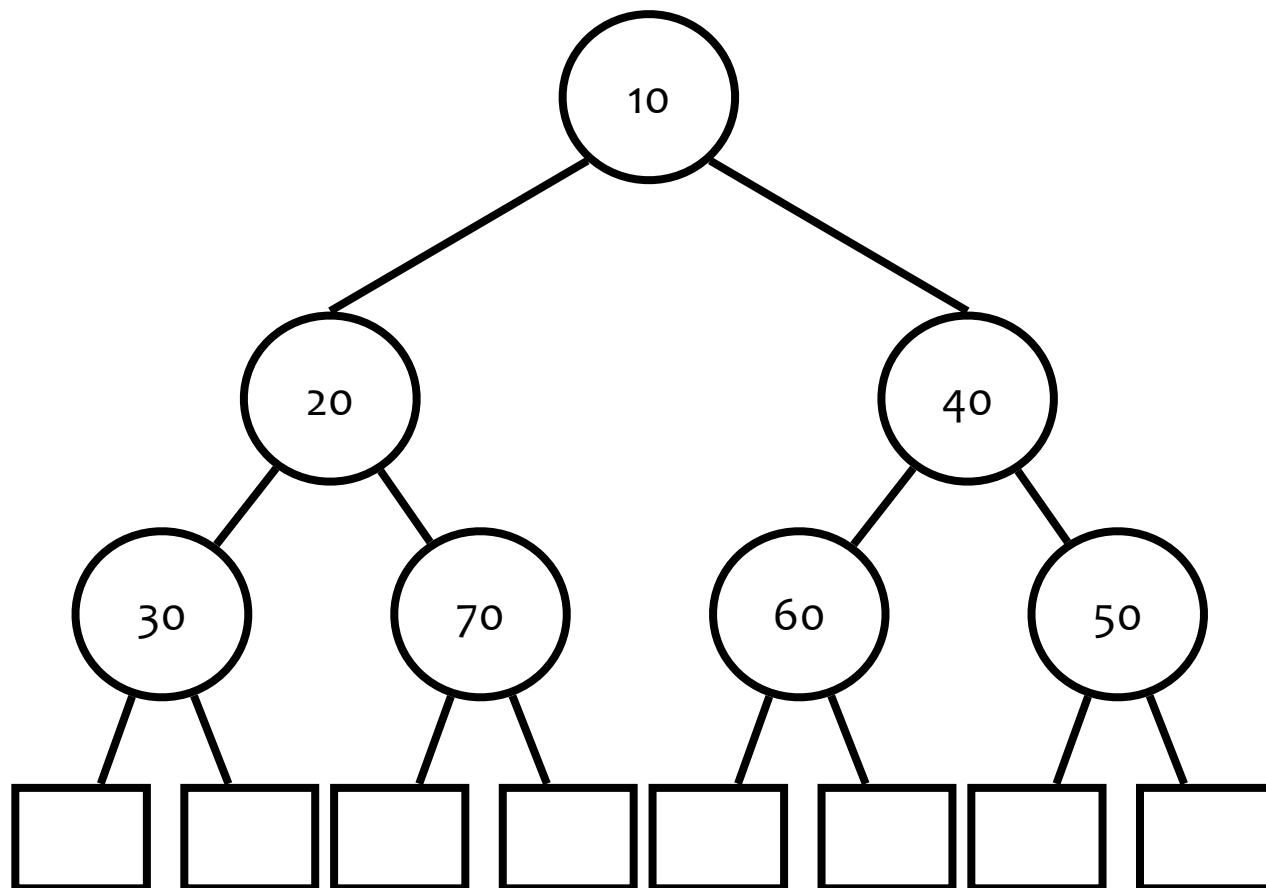
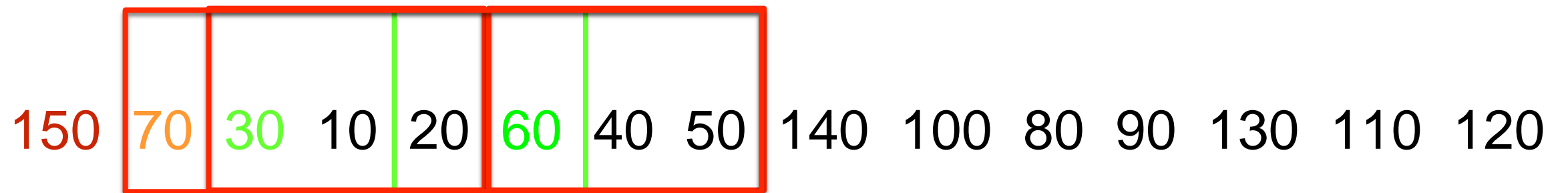


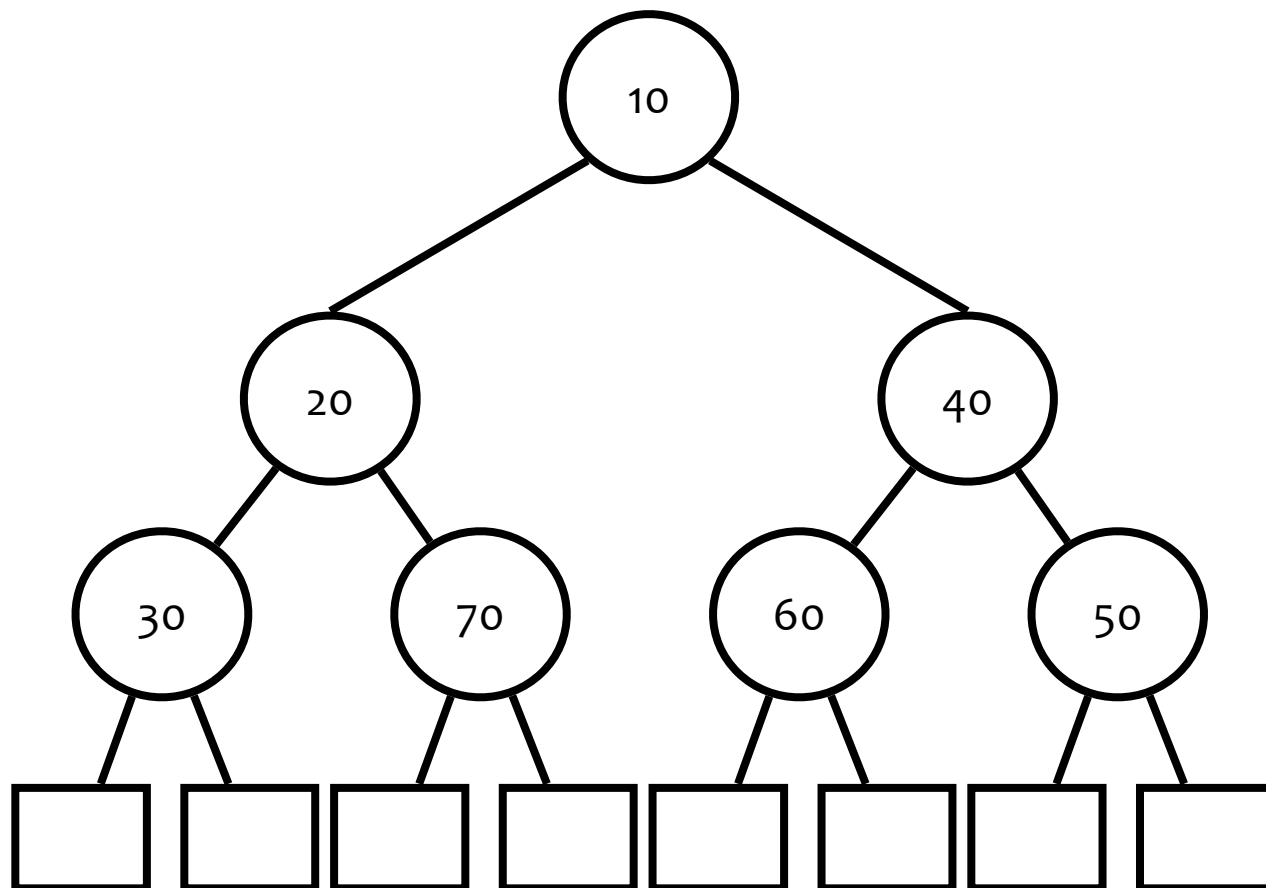
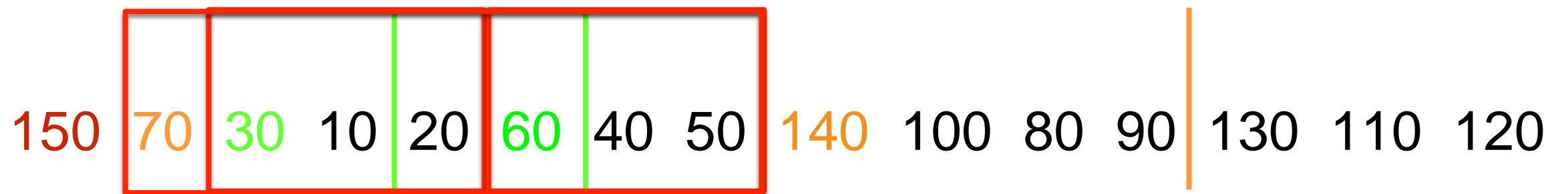


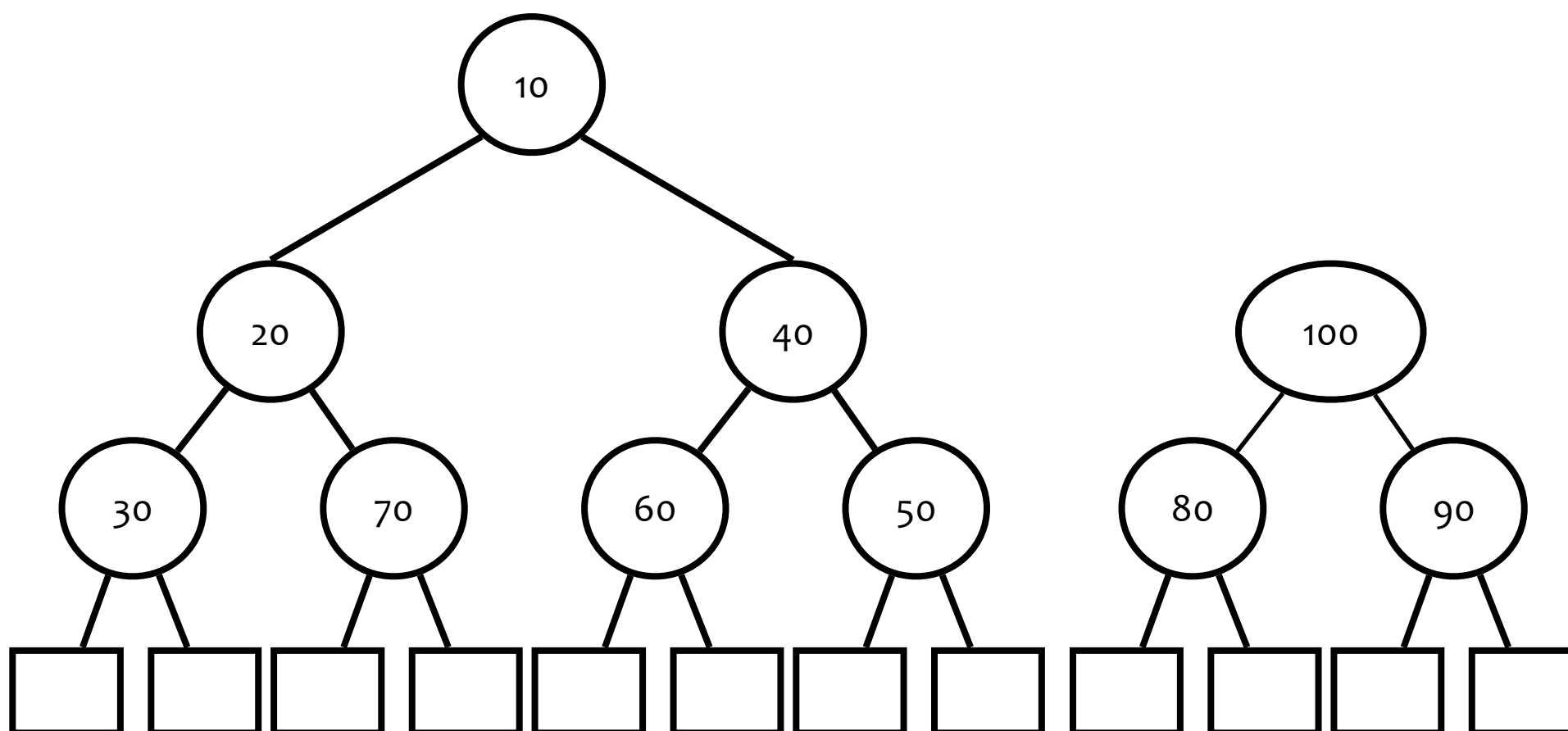
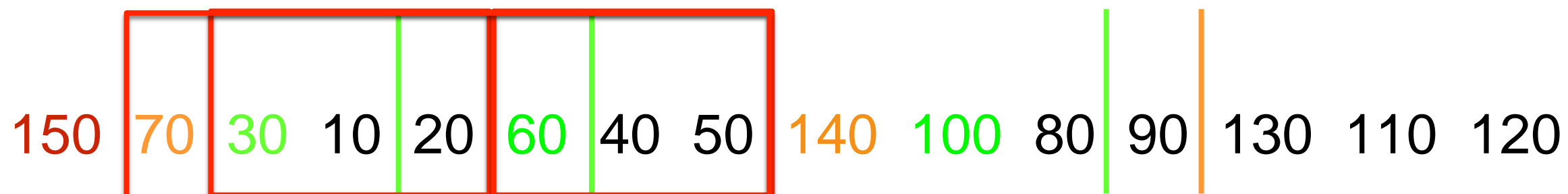


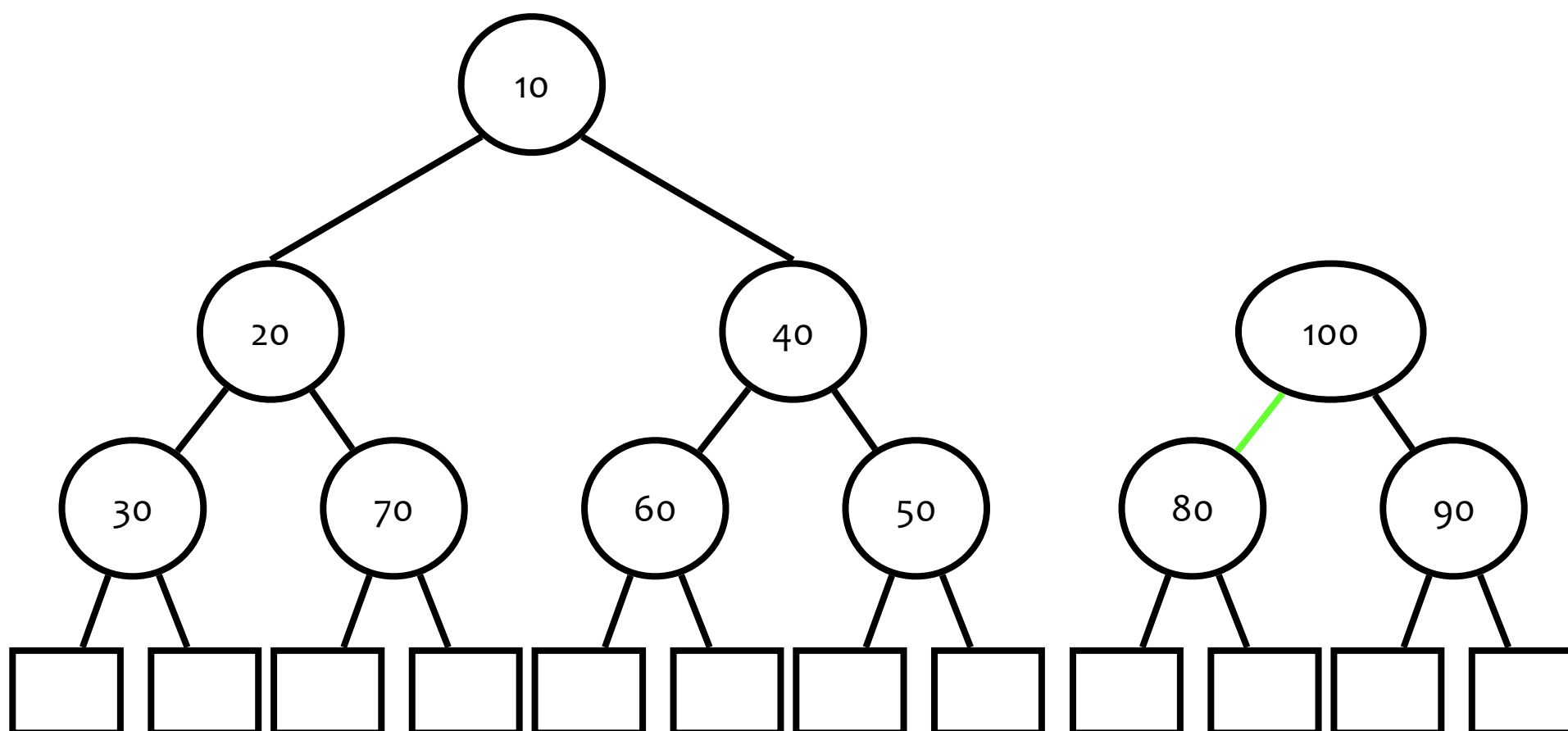
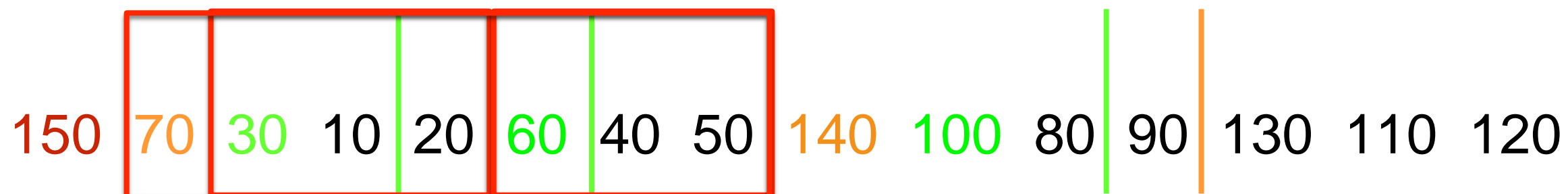


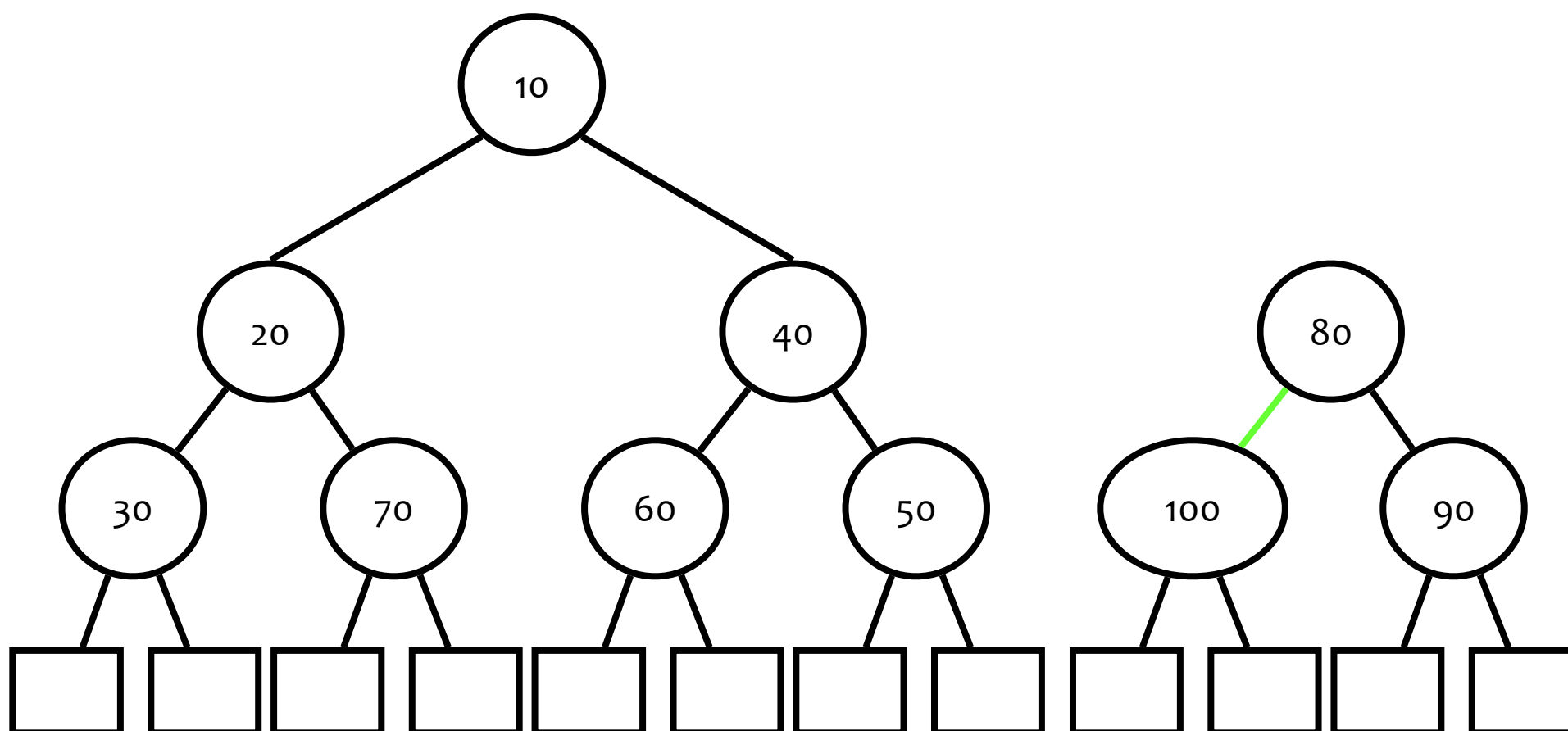
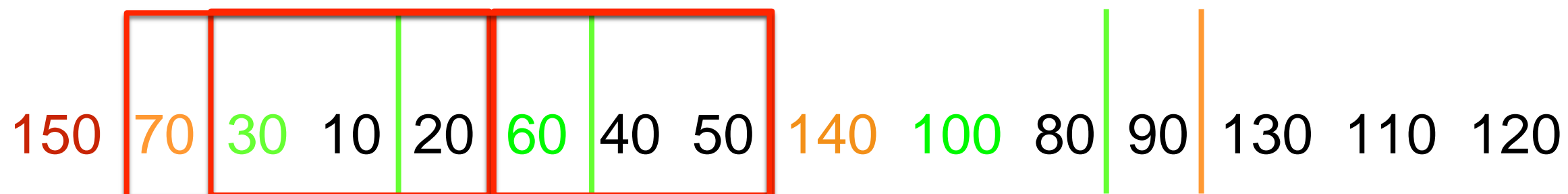


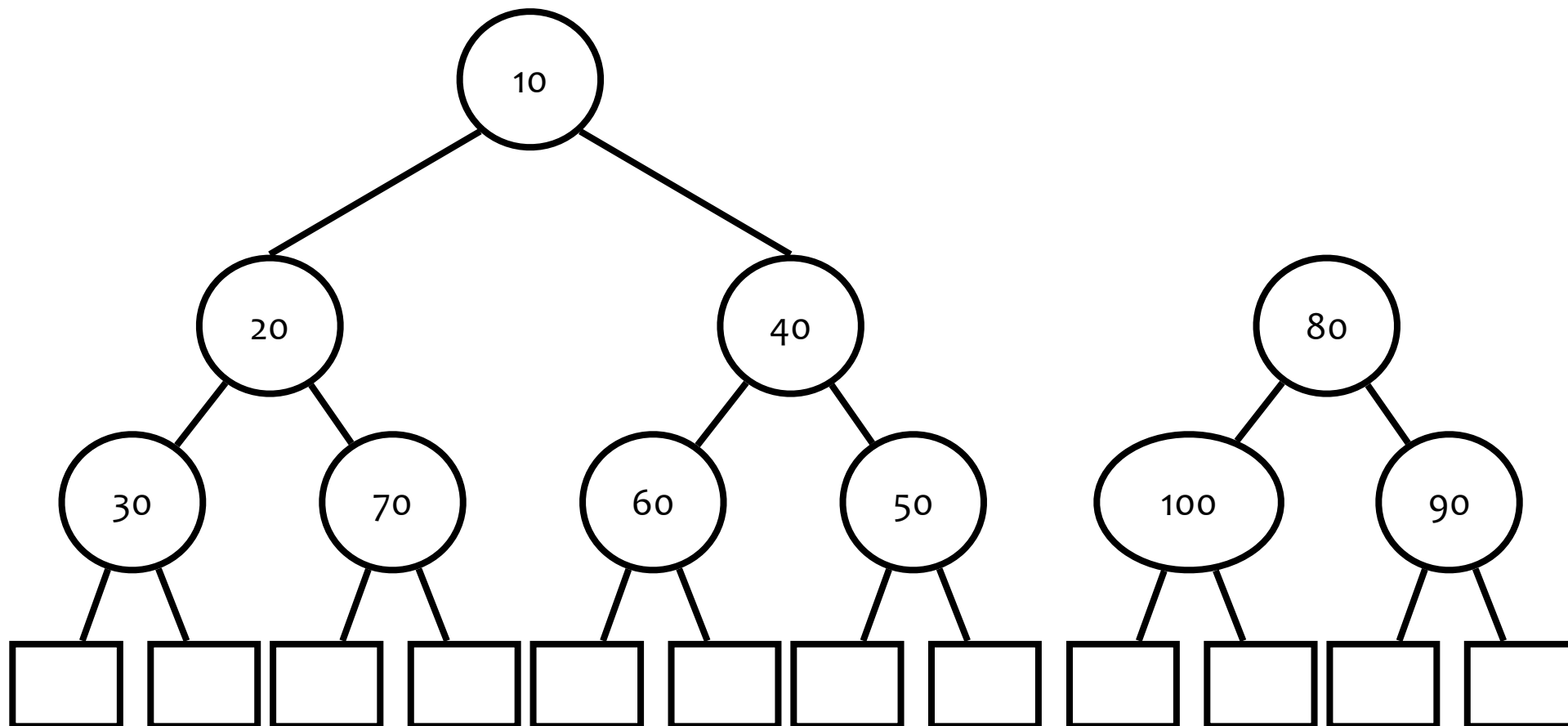


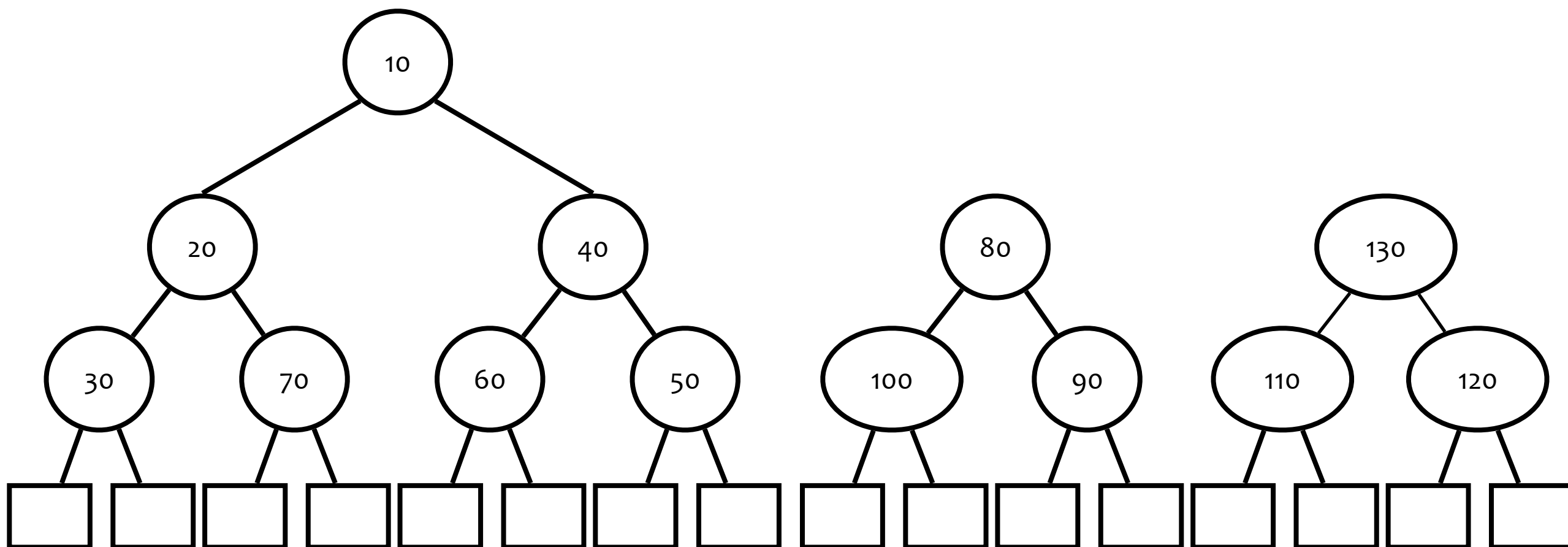
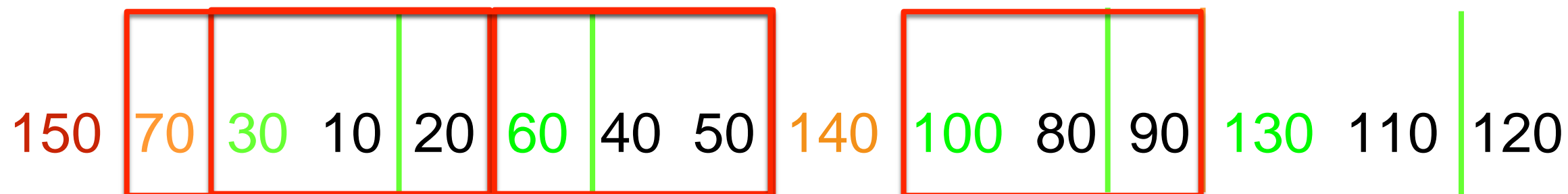


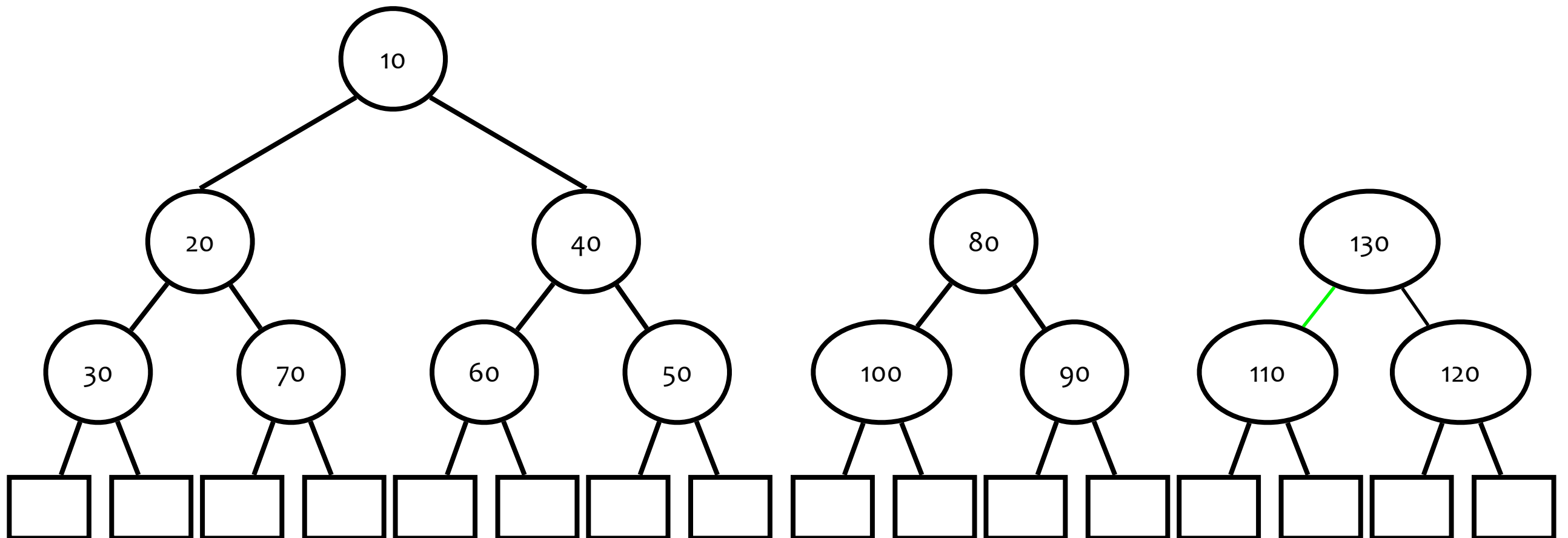


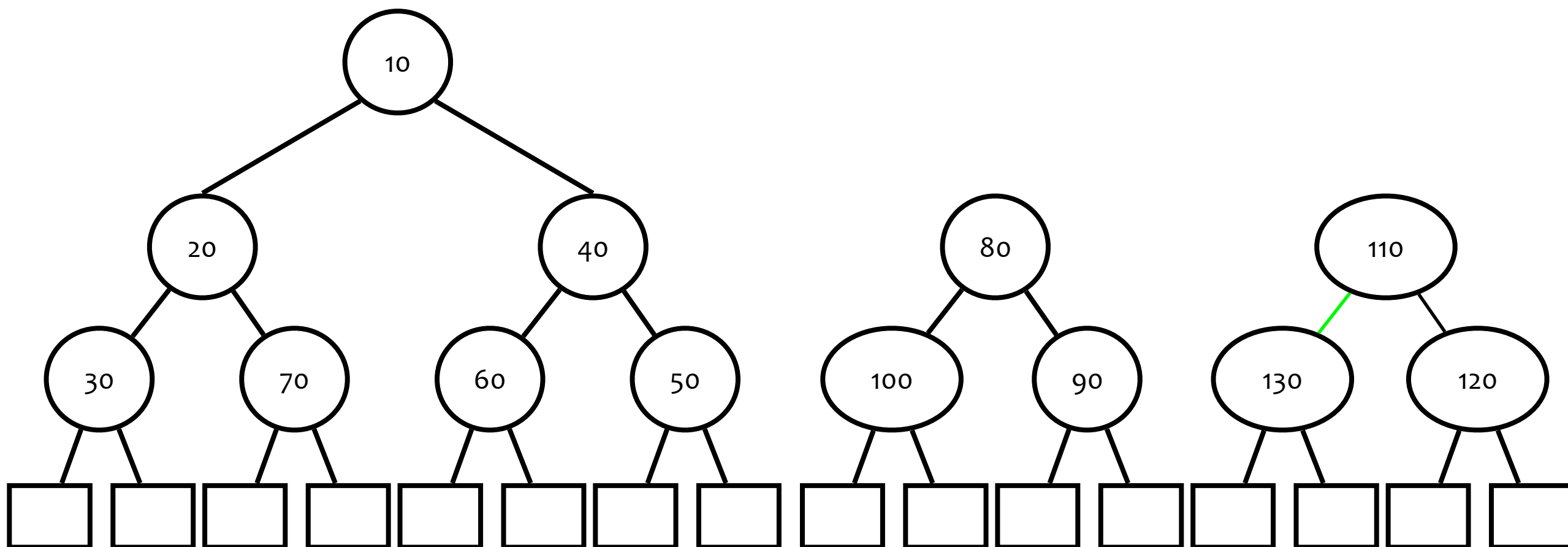
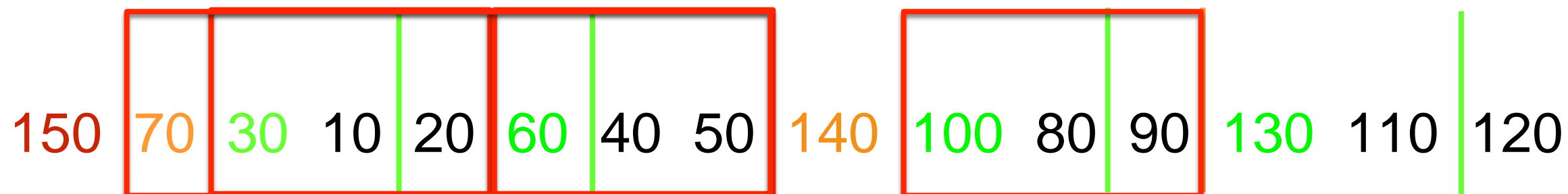


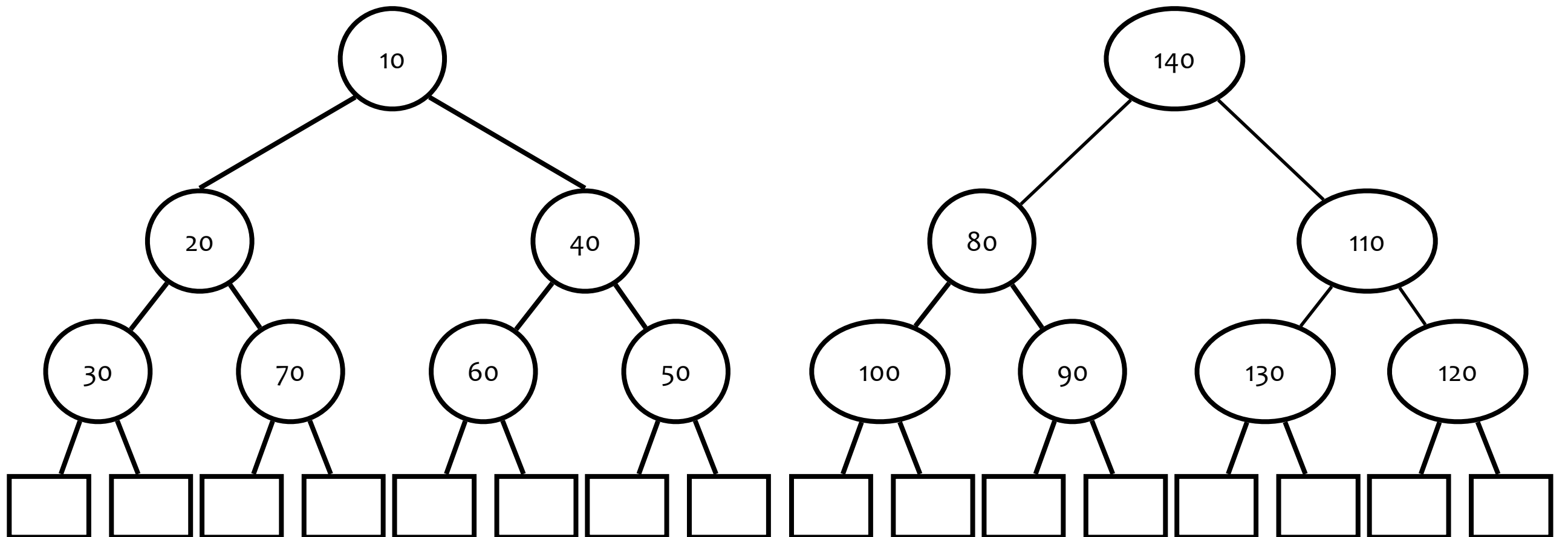
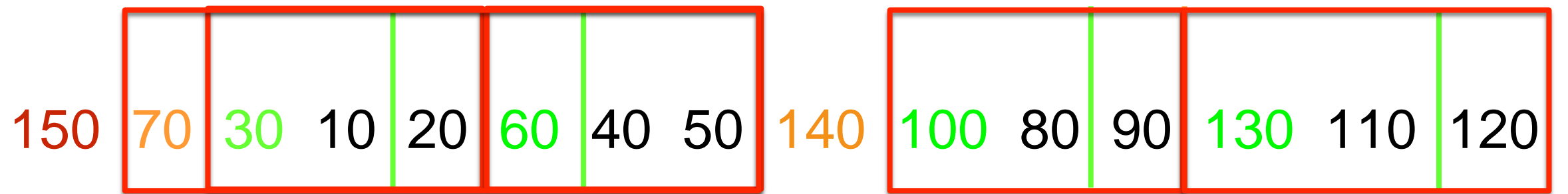


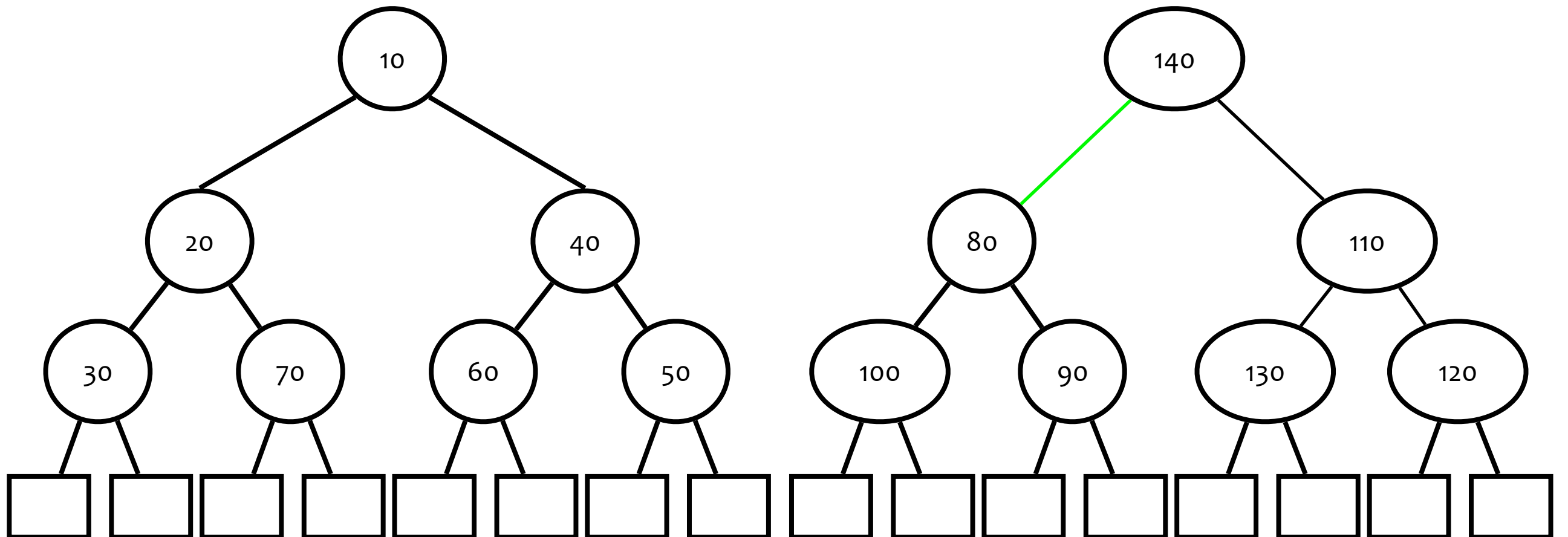
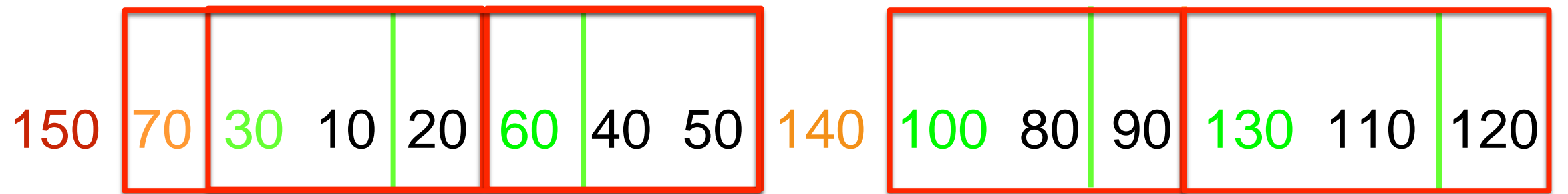


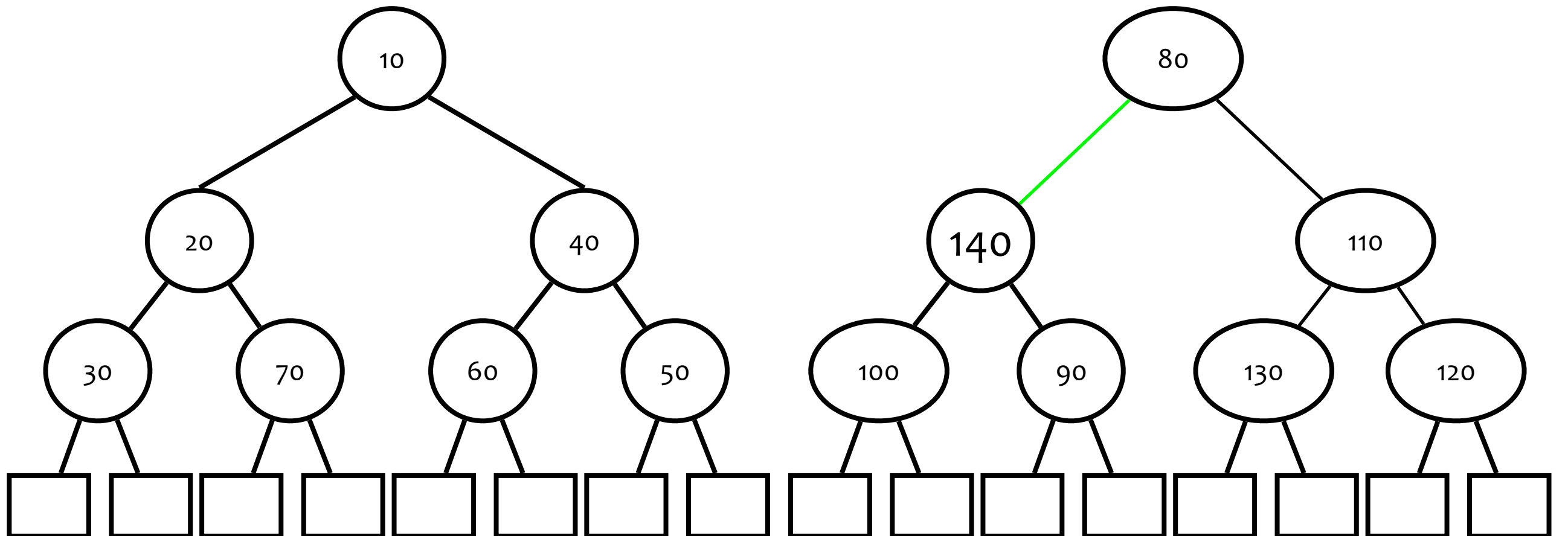


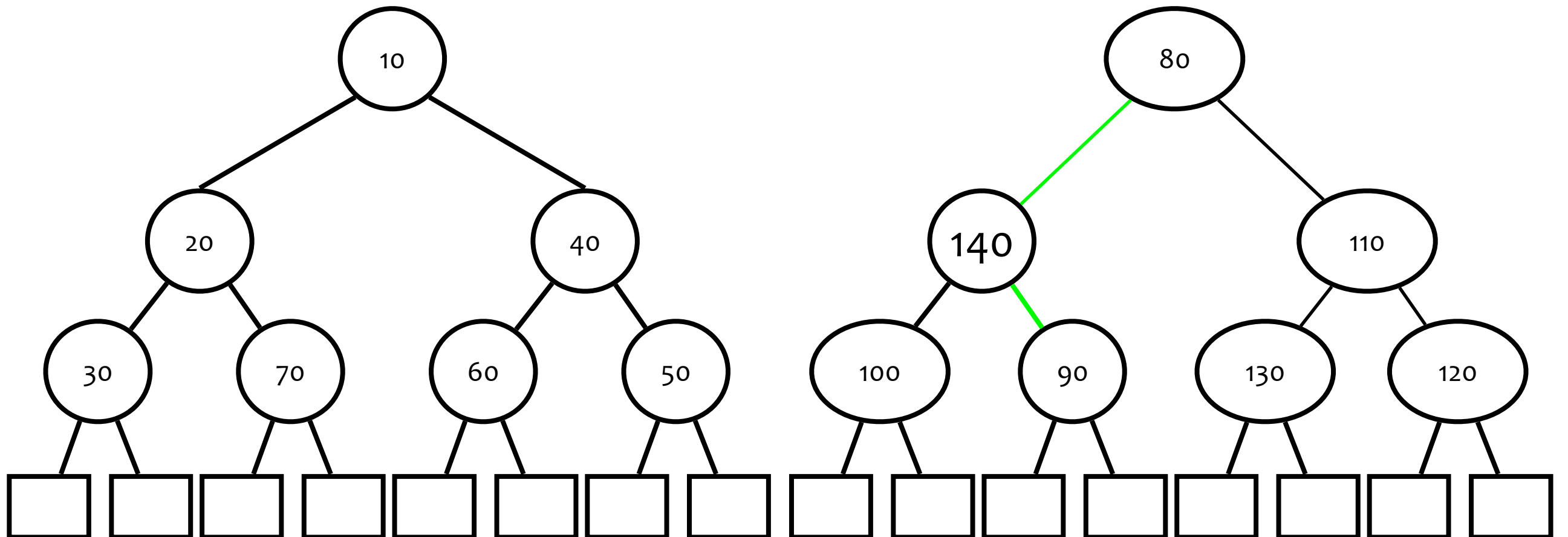
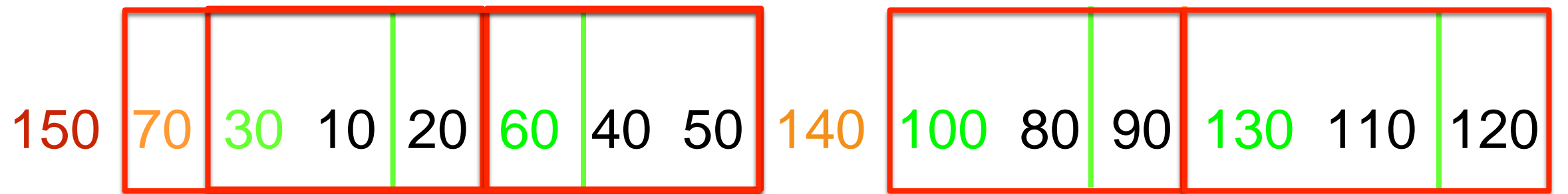


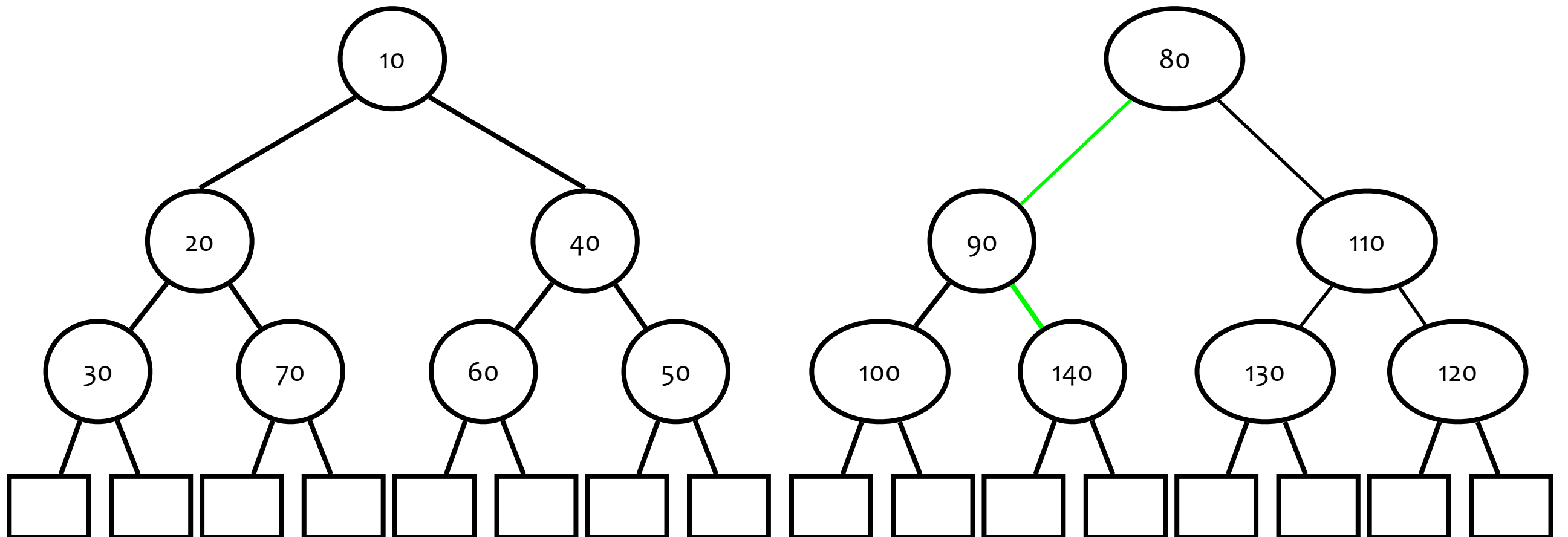
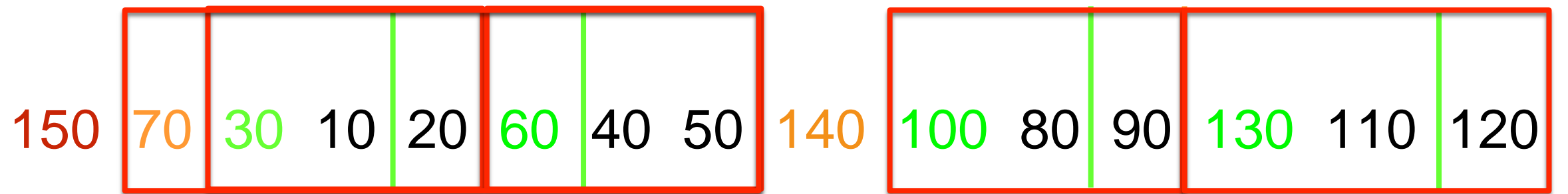




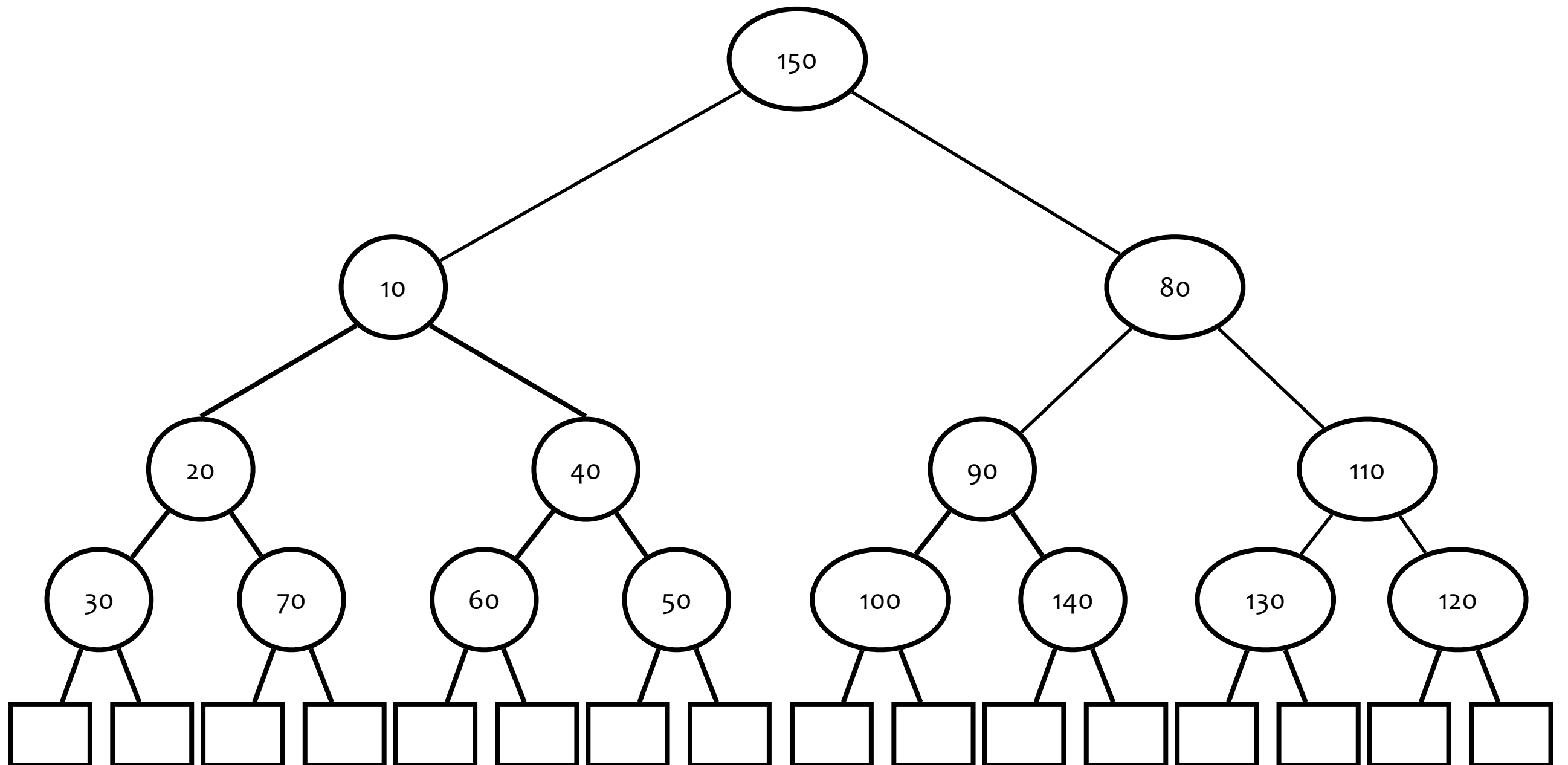




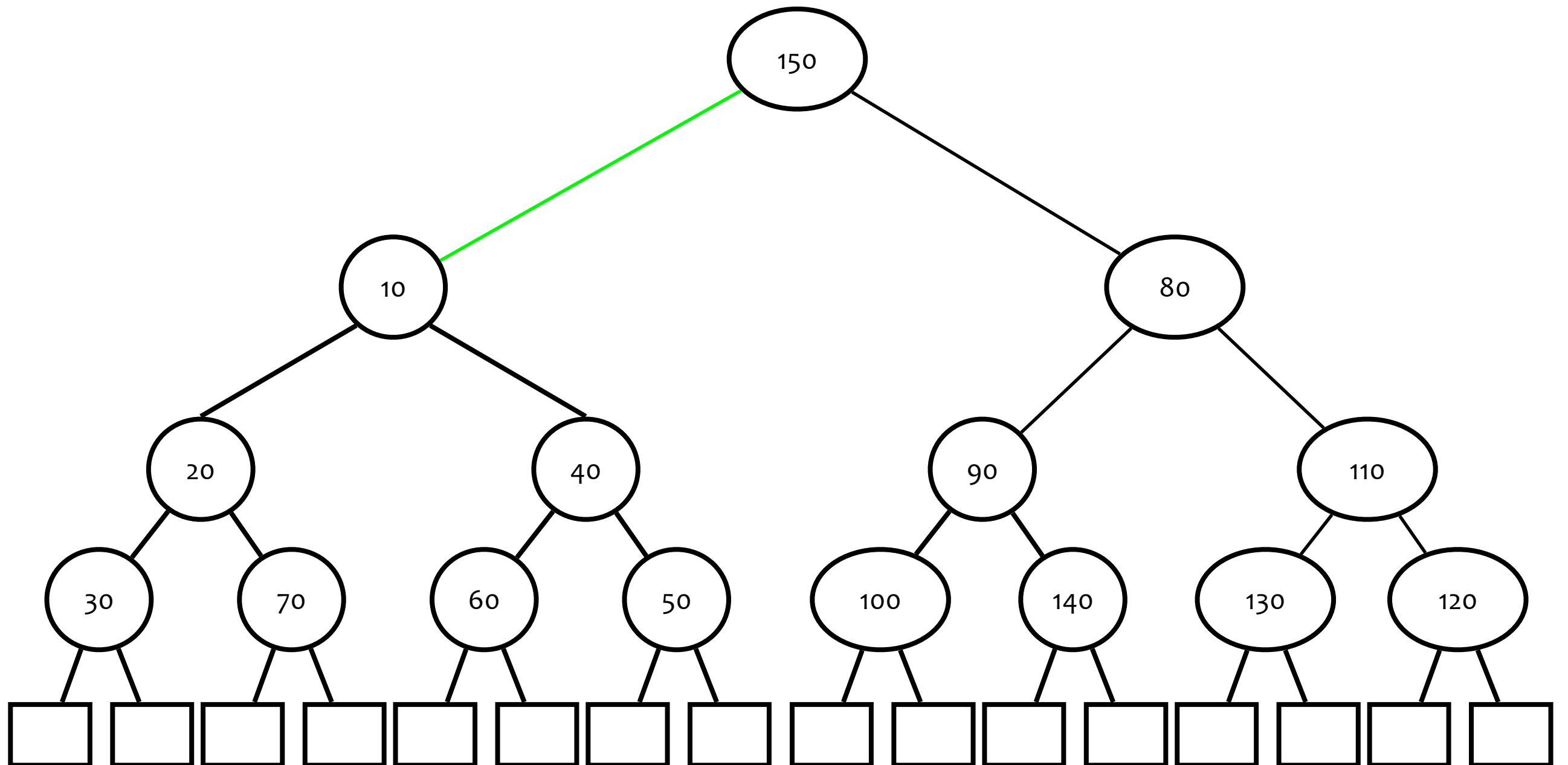




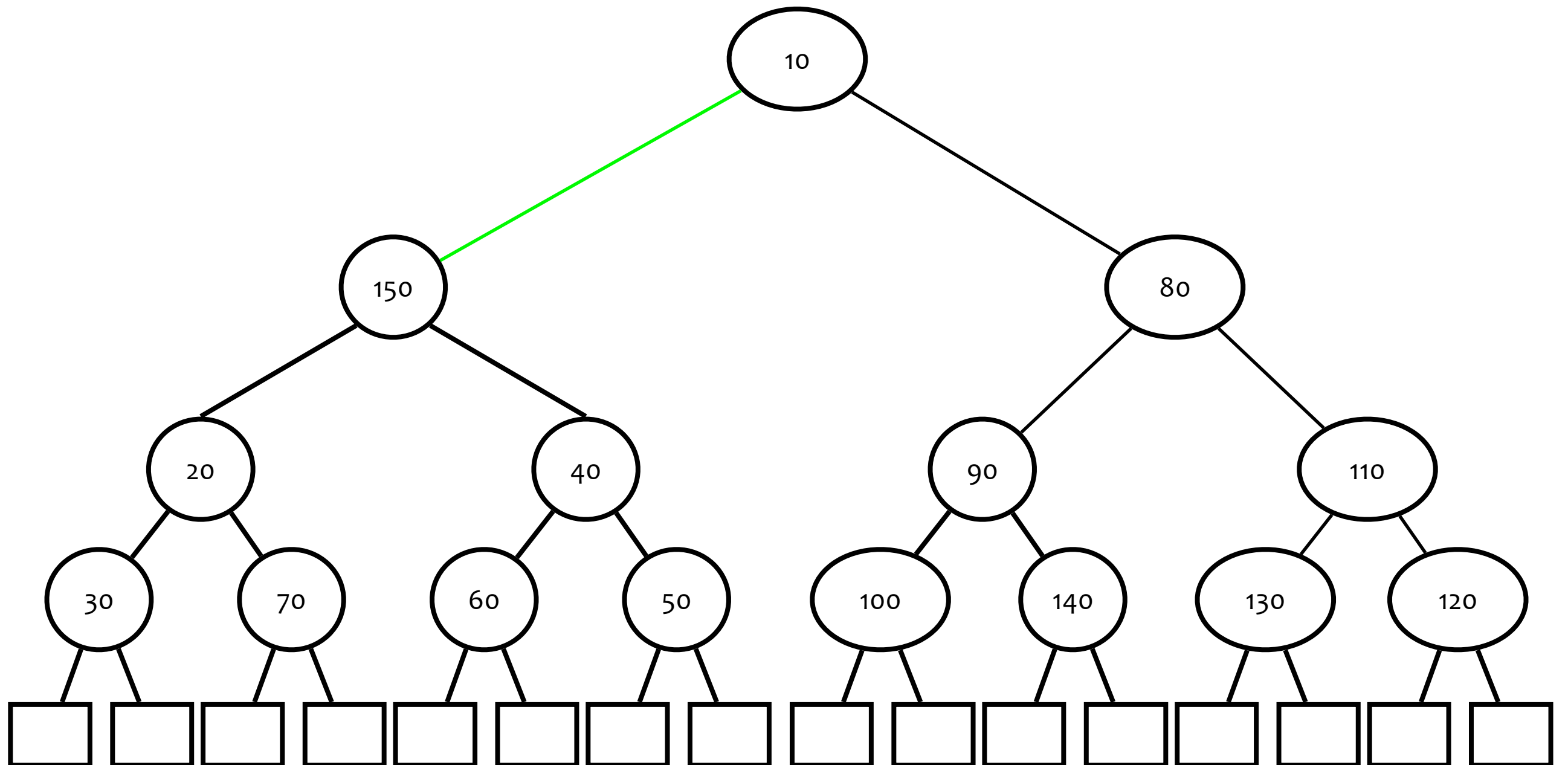
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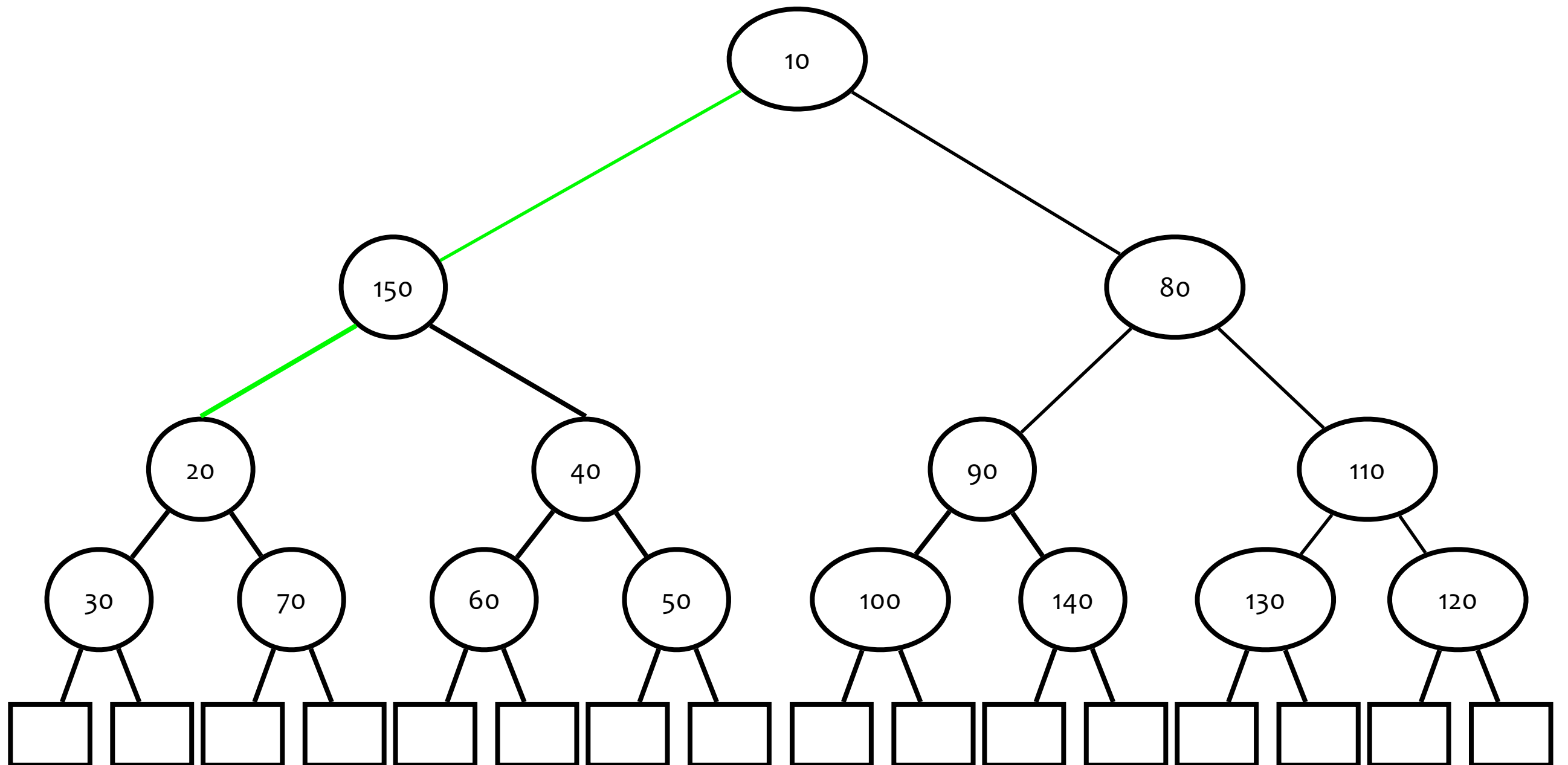
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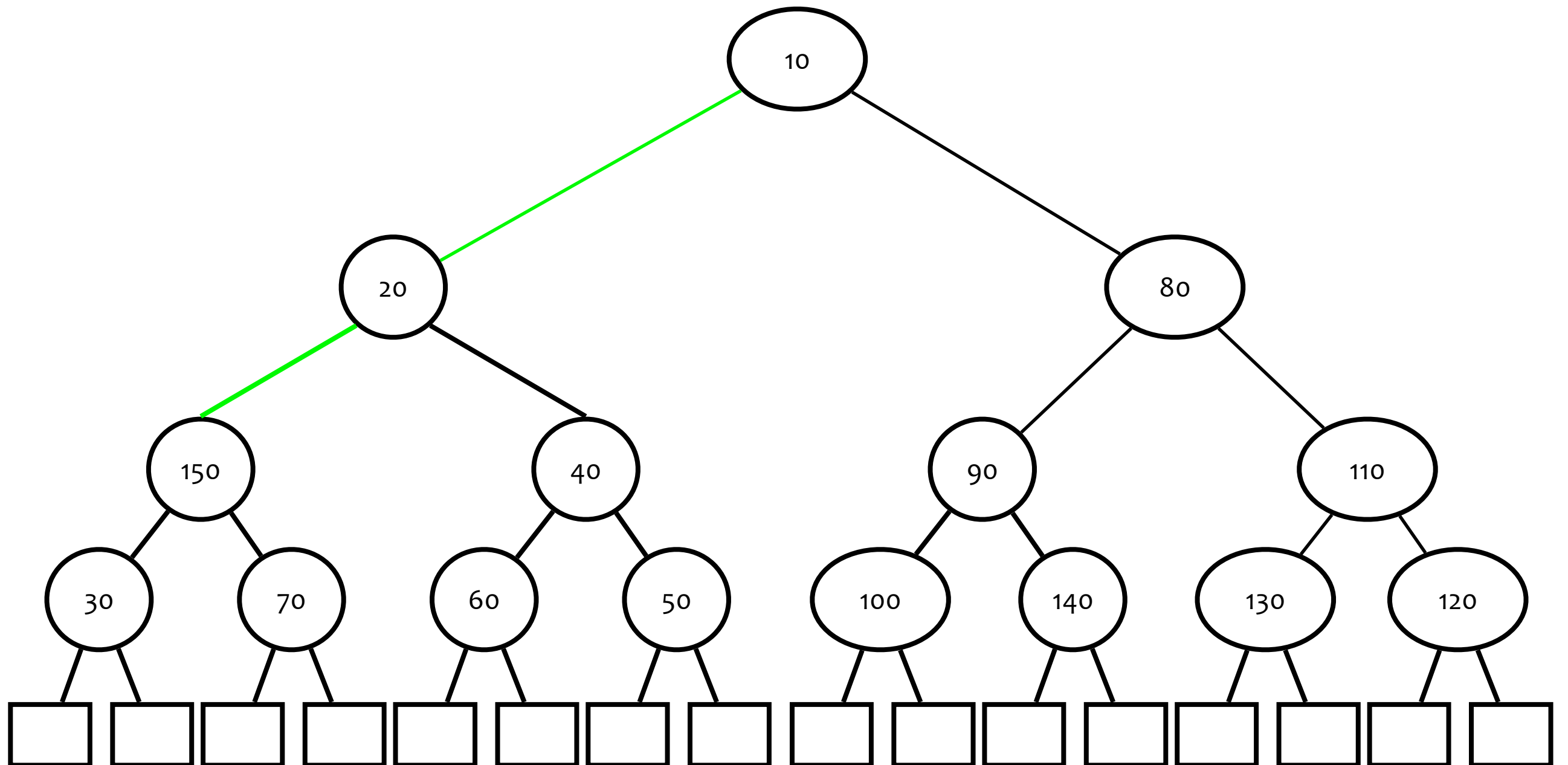
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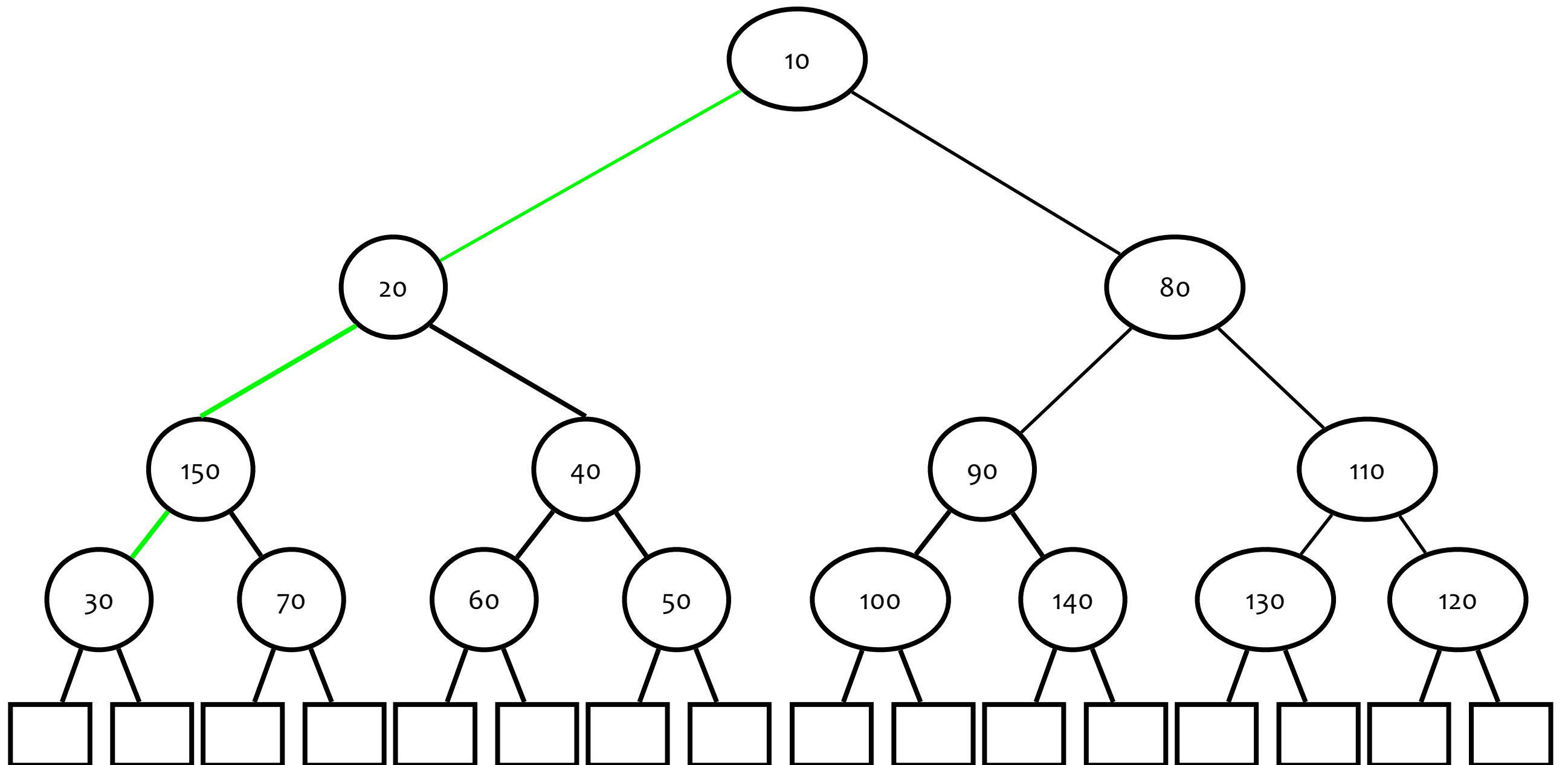
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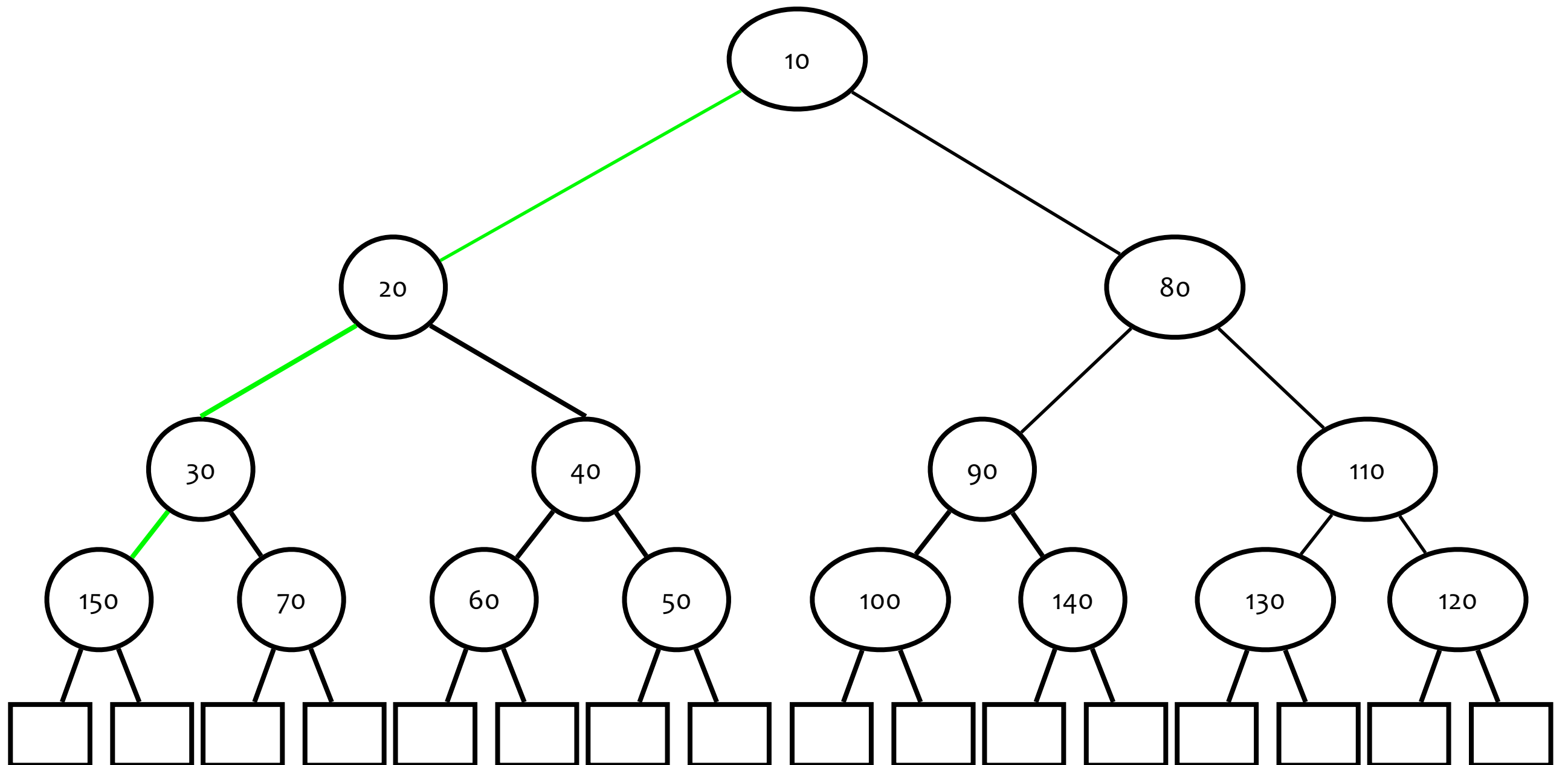
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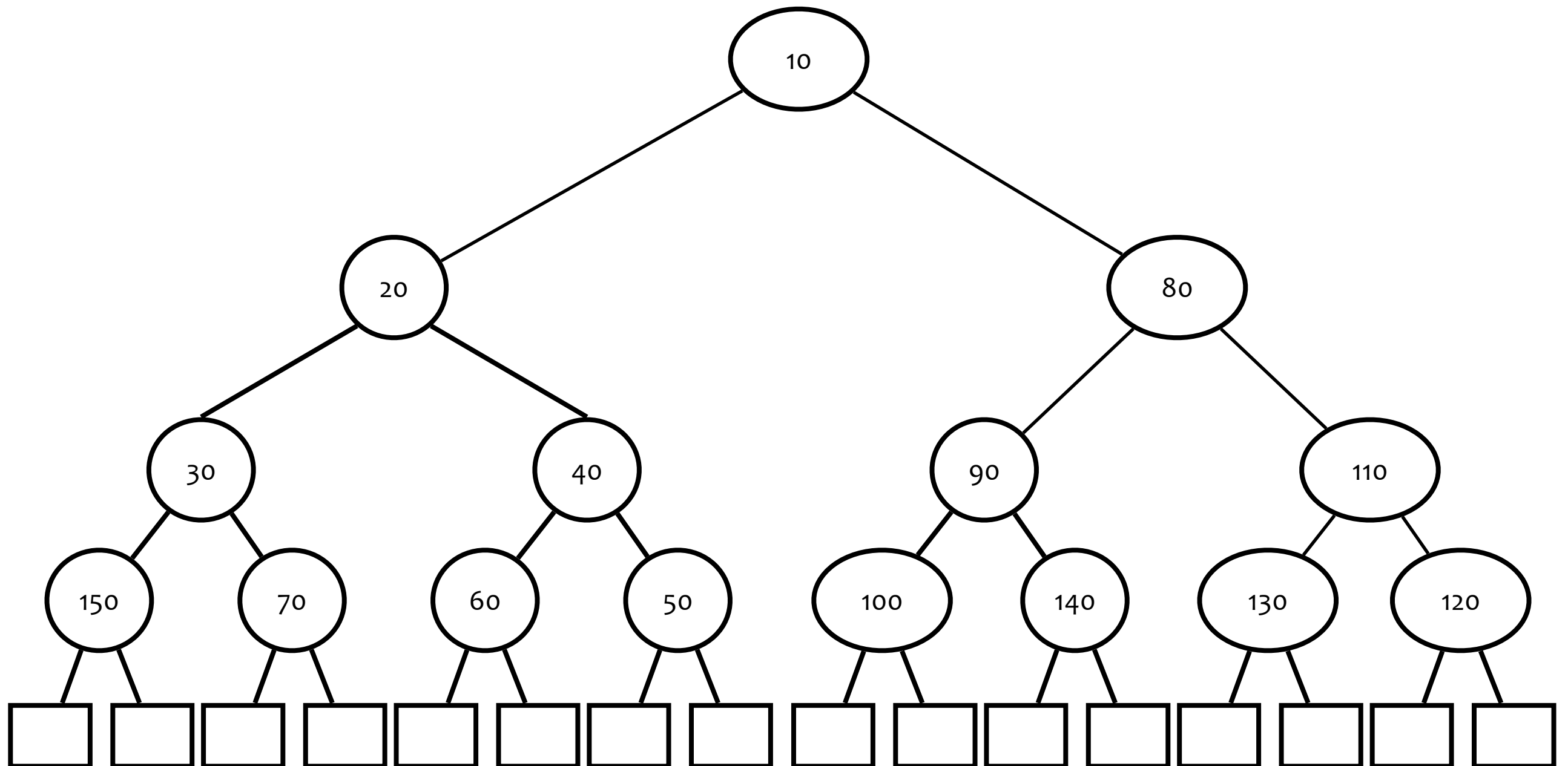
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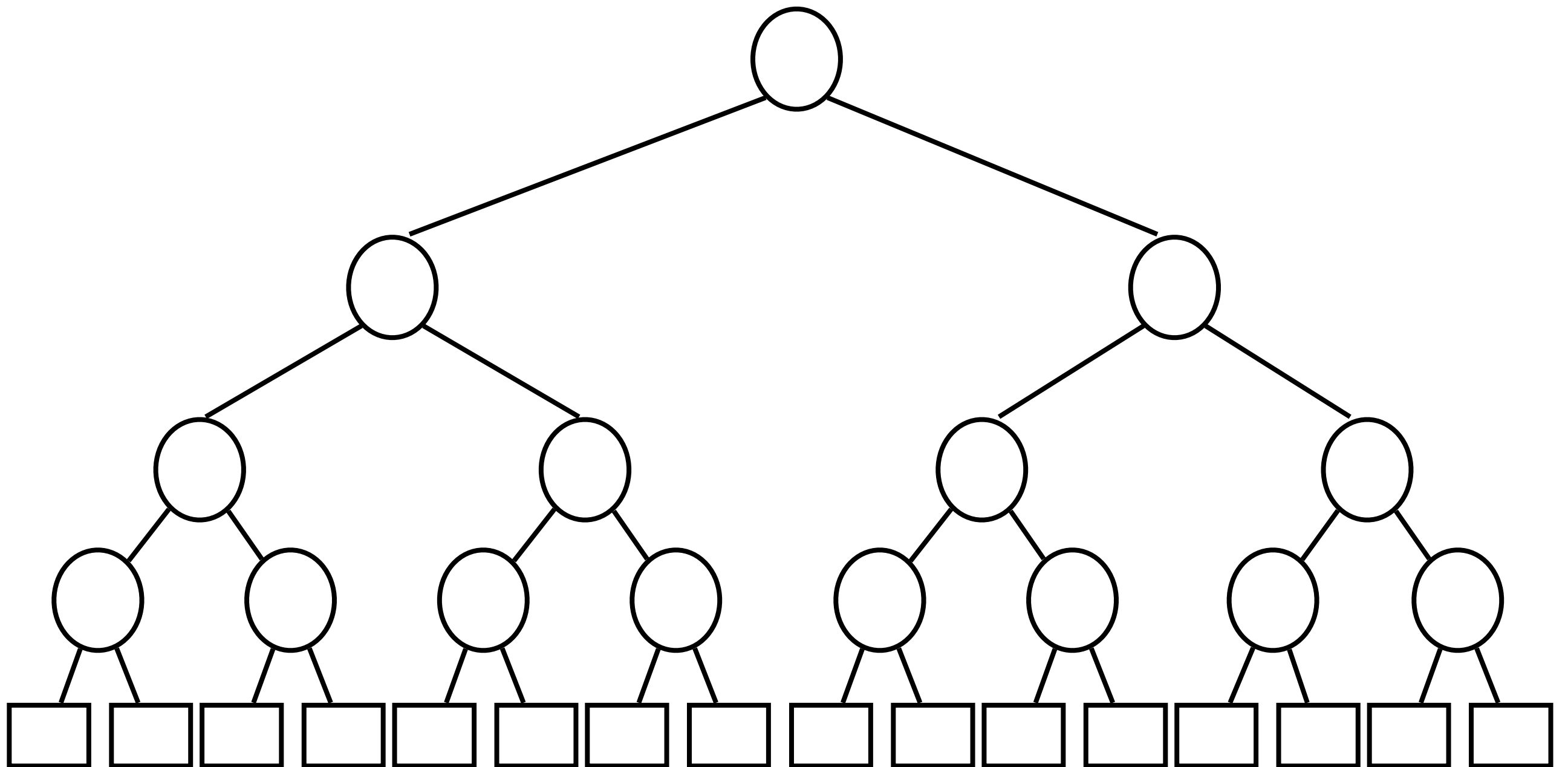
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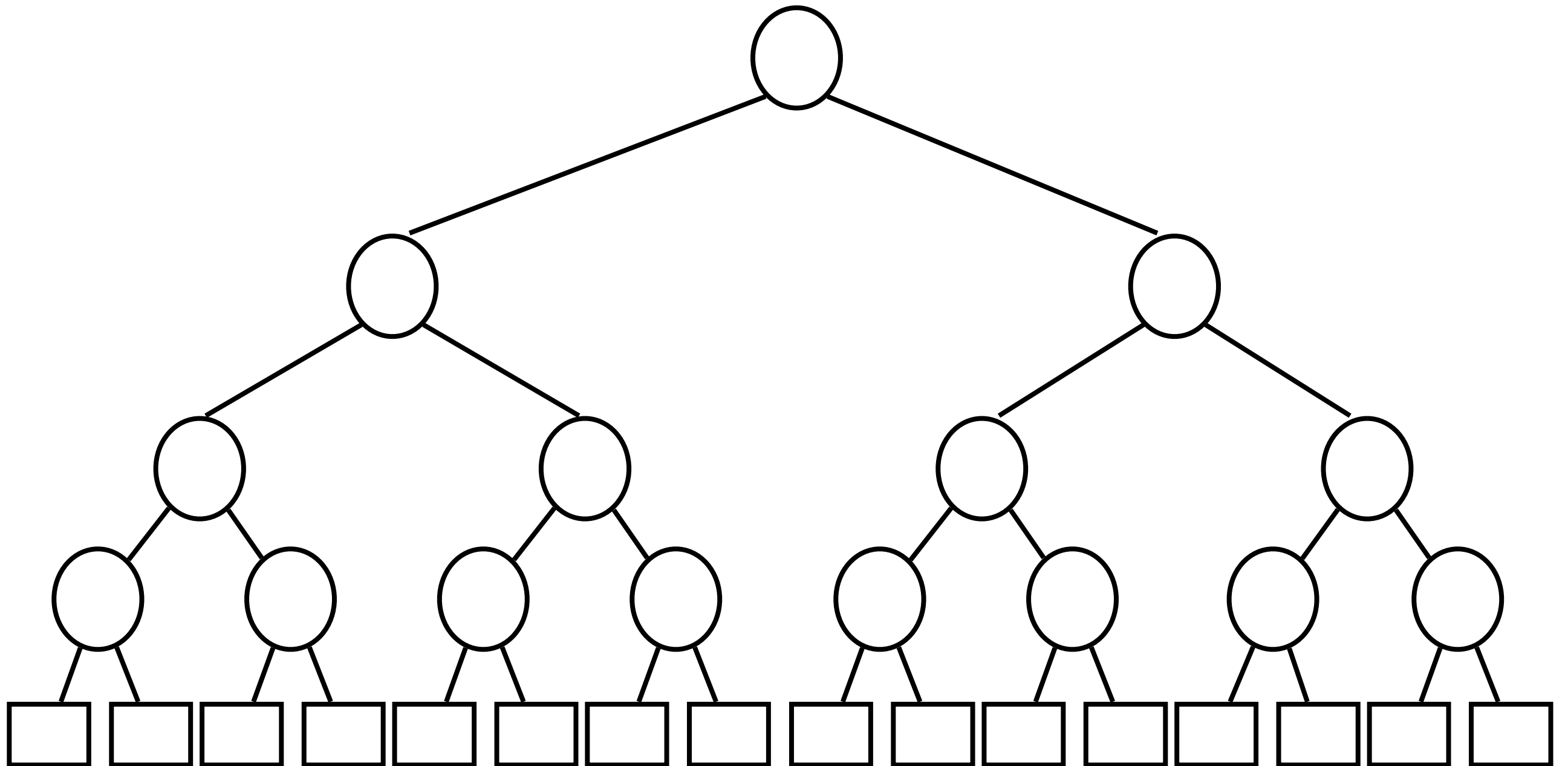
150	70	30	10	20	60	40	50	140	100	80	90	130	110	120
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Did we really insert all n
elements in $O(n)$ time??



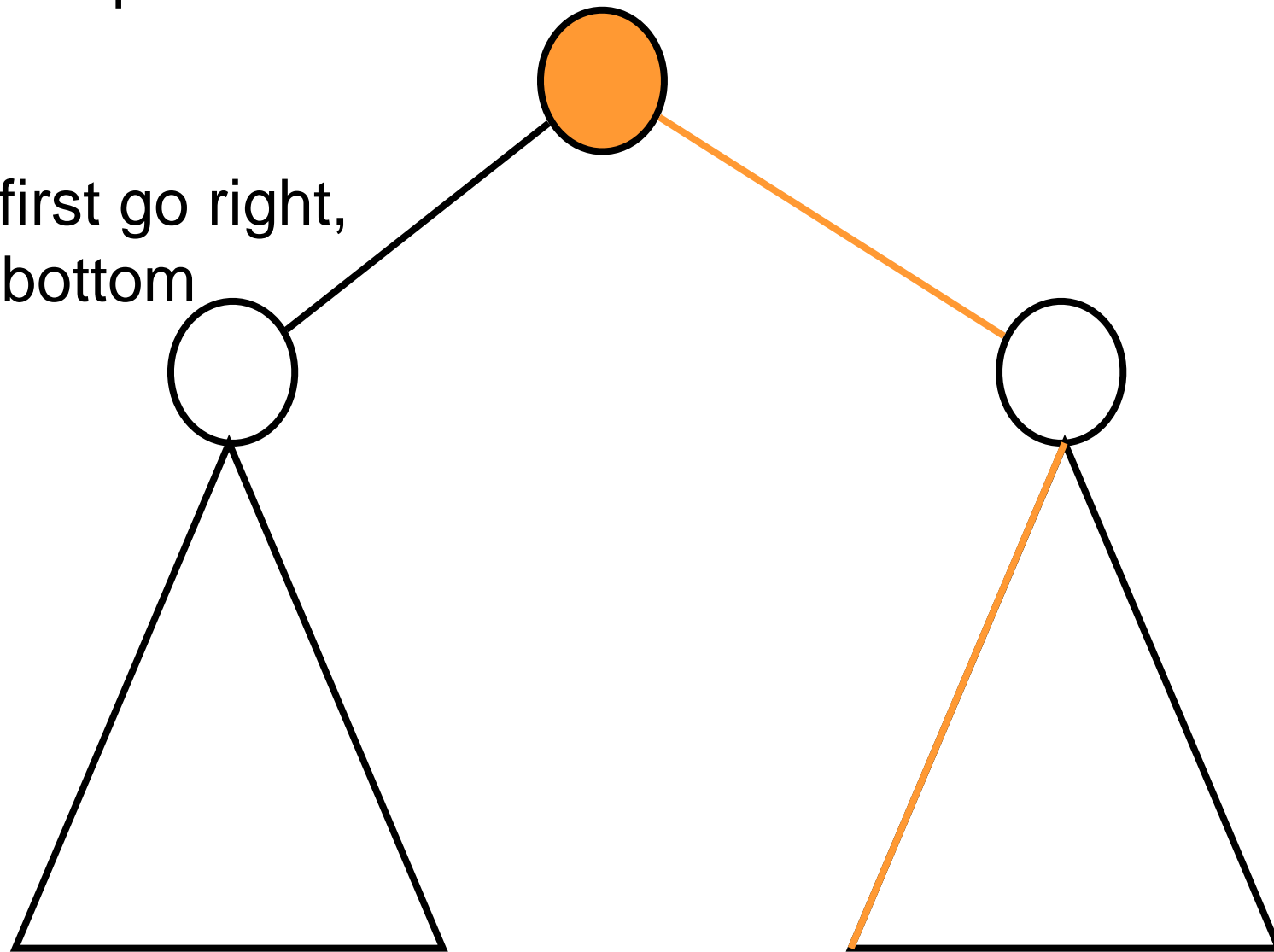
We show: $\max \# \text{ bubble down ops} < \# \text{ heap edges}$



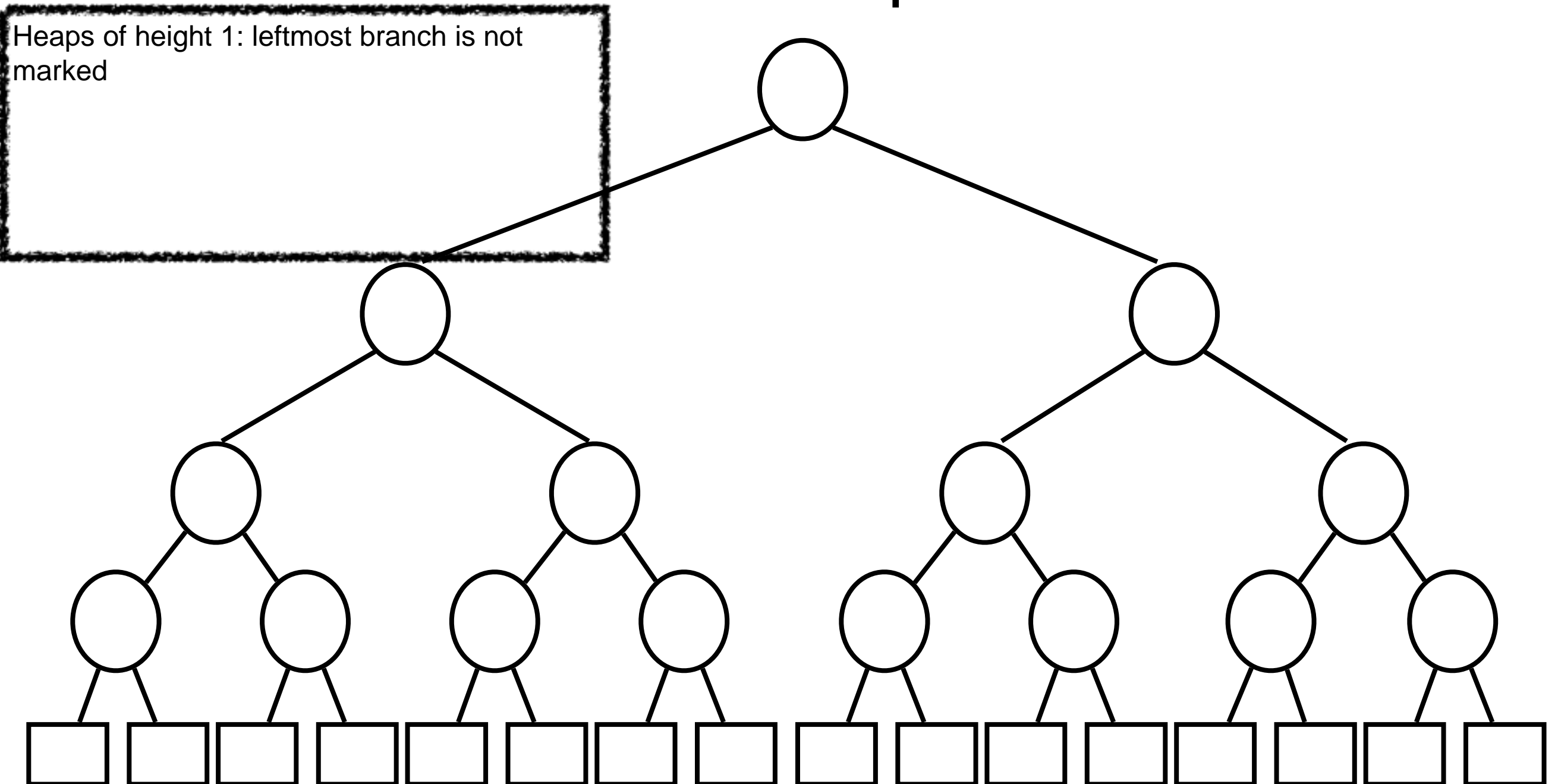
Proof idea

For each new node joining two heaps: mark path of maximum number of bubble-down operations

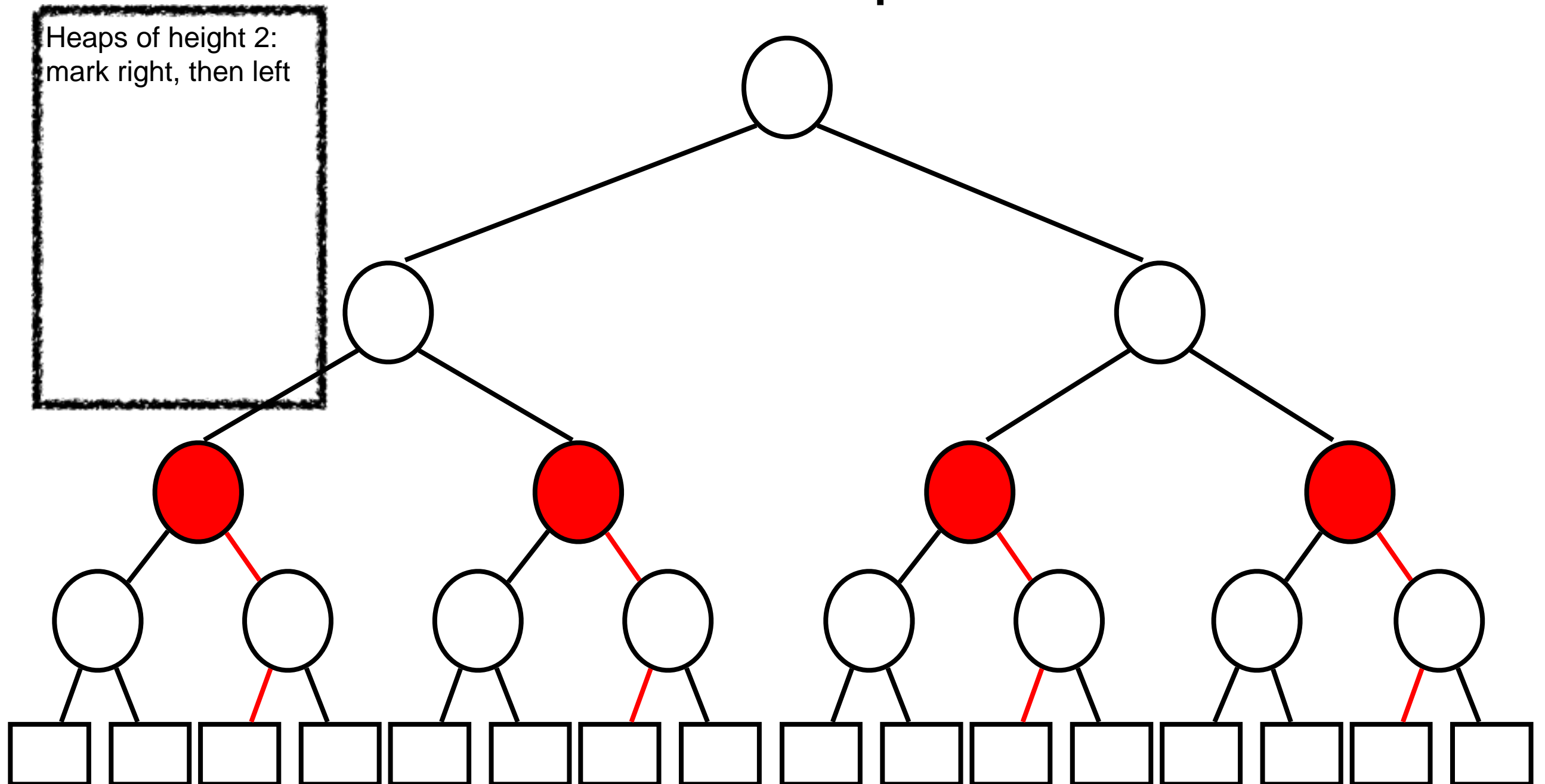
For marking: first go right,
then left until bottom



For each new node joining two heaps:
mark path with maximum number of
bubble-down operations



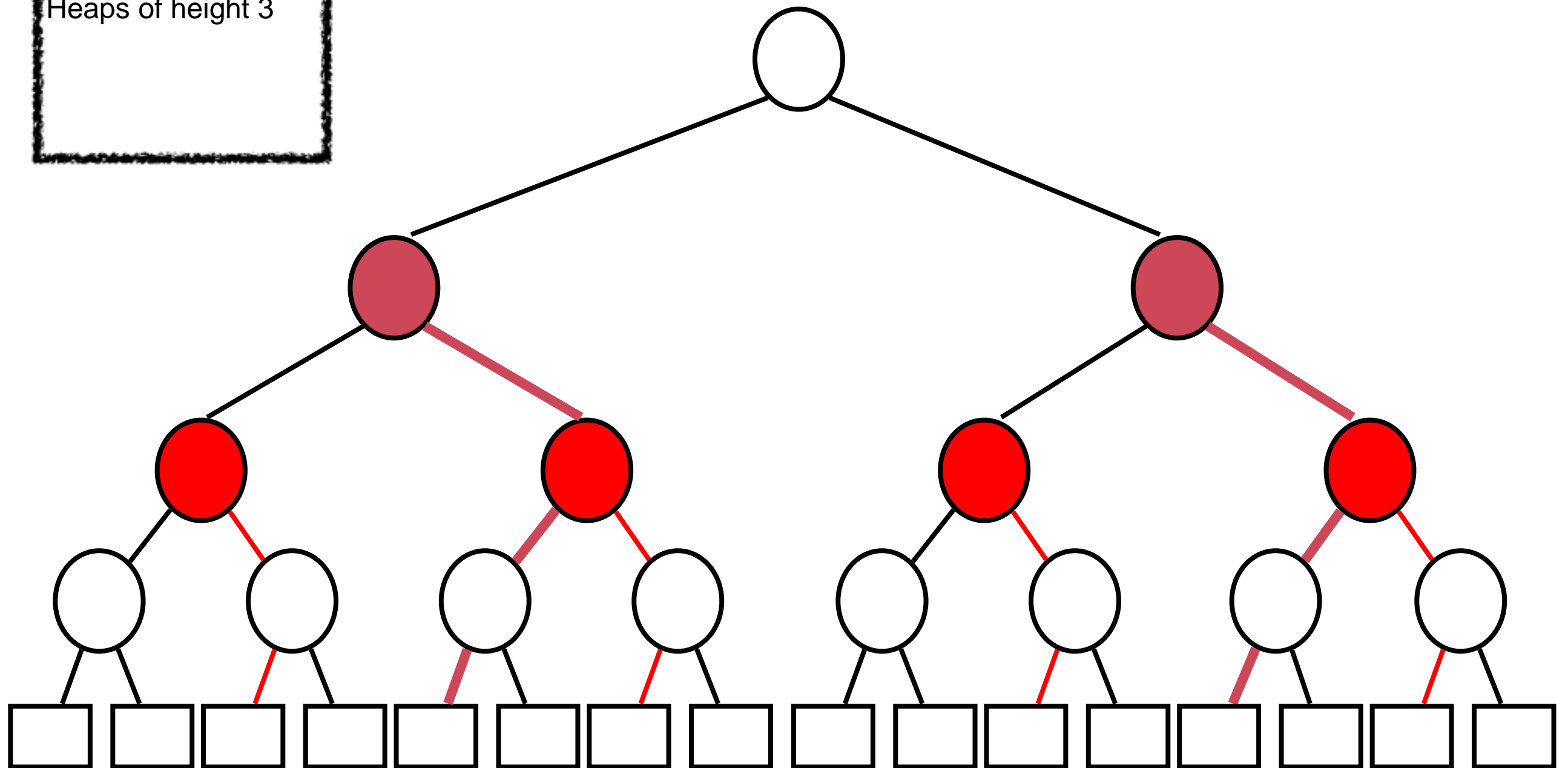
For each new node joining two heaps:
mark path with maximum number of
bubble-down operations



For each height-2 heap, leftmost branch not marked

For each new node joining two heaps:
mark path with maximum number of
bubble-down operations

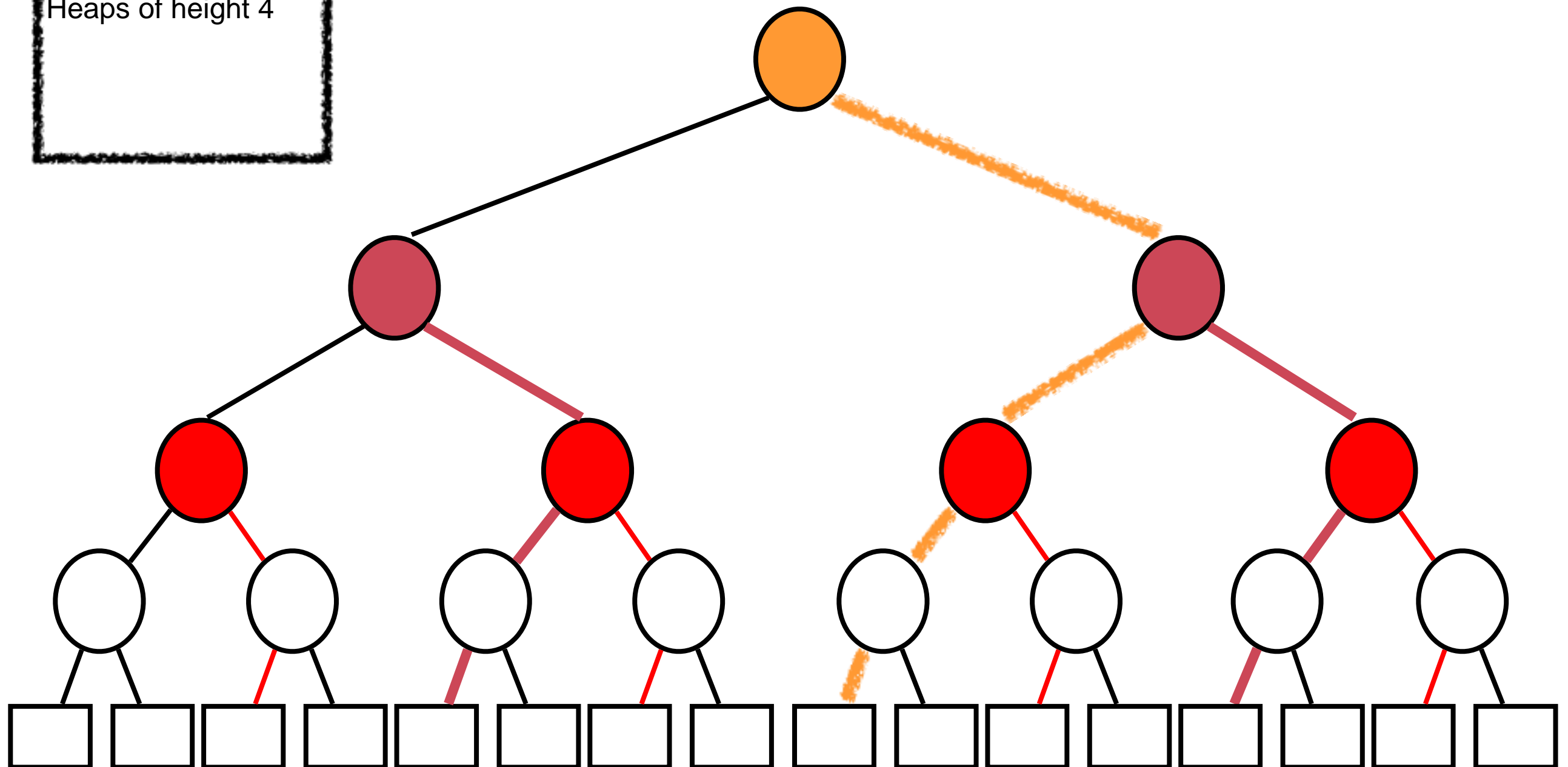
Heaps of height 3



For each height-3 heap, leftmost branch not marked

For each new node joining two heaps:
mark path with maximum number of
bubble-down operations

Heaps of height 4



For height-4 heap, leftmost branch not marked

Inductive argument: marking procedure will never mark all edges in heap, since the leftmost branch is never marked

- Note: leftmost branch in height- h heap: not marked
- When joining 2 heaps of height h to heap of height $h + 1$: new edges to be marked are
 - edge joining new node and right heap of height h , and
 - edges on left path in the right heap of height h
- We conclude: leftmost branch in height $(h+1)$ heap is not marked

Build Heap In-place

Algorithm buildHeap(A, n):

for $i \leftarrow \lfloor n/2 \rfloor$ **to** 1 **do**
 downHeap(A, i)

Algorithm downHeap(A, i):

$l \leftarrow 2i$

$r \leftarrow 2i + 1$

if $l \leq n \wedge A[l] < A[i]$ **then**

$min \leftarrow l$

else

$min \leftarrow i$

if $r \leq n \wedge A[r] < A[min]$ **then**

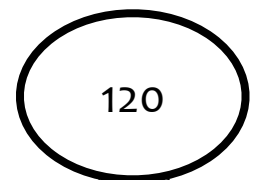
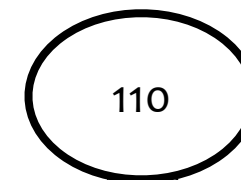
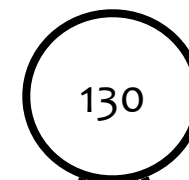
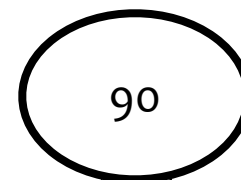
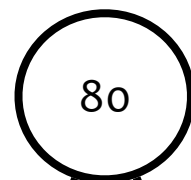
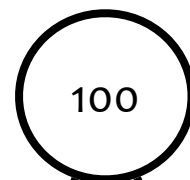
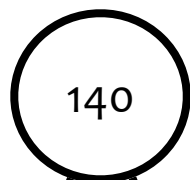
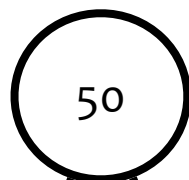
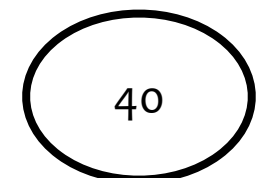
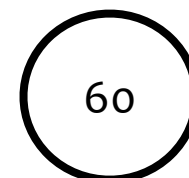
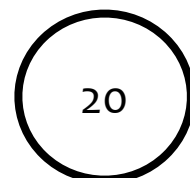
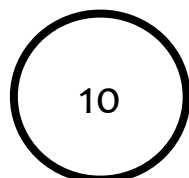
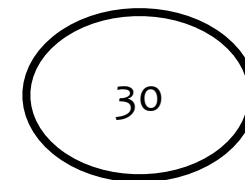
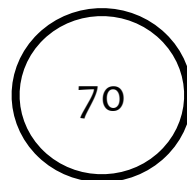
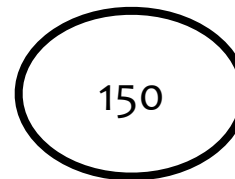
$min \leftarrow r$

if $i \neq min$ **then**

 swap(i, min)

 downHeap(A, min)

150 70 30 10 20 60 40 50 140 100 80 90 130 110 120



10 20 30 50 80 60 40 70 140 100 150 90 130 110 120

