

## Question 1

Assume for a contradiction that  $L$  is context-free. Let  $p$  be the pumping length.

Choose  $s = a^p b^{2p} c^{3p}$

Since  $s \in L$  and  $|s| \geq p$ , we can write  $s = uvxyz$  such that the properties of the PL hold.

By property 2,  $|vxy| \leq p$ , we have the following cases:

(Case 1)  $vxy$  consists of one type of symbol (a's, b's or c's)

(Case 2)  $vxy$  consists of two types of symbols ( $a \dots ab \dots b$  or  $b \dots bc \dots c$ )

Case 1) By property 1,  $v$  and/or  $y$  is a non-empty substring of the one type of symbol (a, b, or c).

The string  $uv^2xy^2z$  will result in increasing the number of the one type of symbol without changing the number of the other two types of symbols. So,  $uv^2xy^2z$  is not of form  $a^n b^{2n} c^{3n}$ .

Hence,  $uv^2xy^2z \notin L$

Case 2) By property 1,  $v$  and/or  $y$  is a nonempty substring of two symbols (eg.  $v=a \dots ab \dots b$ ) or  $v$  and/or  $y$  is one type of symbol (eg.  $v=a \dots a, y=b \dots b$ )

In the two symbol case, the string  $uv^2xy^2z$  will result in symbols out of order, ie not of form  $a \dots ab \dots bc \dots c$

In the one symbol case,  $uv^2xy^2z$  will increase the number of one or two types of symbols, but not the third type of symbol. So, the string will not be of form  $a^n b^{2n} c^{3n}$ .

Hence,  $uv^2xy^2z \notin L$

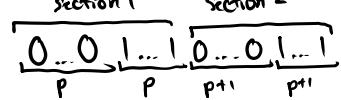
We have shown that we cannot write  $s = uvxyz$  satisfying all conditions of the PL. Therefore,  $L$  is not context-free.

## Question 2

Assume for a contradiction that  $L$  is context-free. let  $p$  be the pumping length.

Choose  $s = 0^p 1^p 0^{p+1} 1^{p+1}$

Since  $s \in L$  and  $|s| \geq p$ ,  $s$  can be rewritten as  $s = uvxyz$  such that the conditions of the PL hold.

By property 2,  $|vxy| \leq p$ , we have the following cases: 

(Case 1)  $vxy$  consists of one type of symbol

(Case 2)  $vxy$  consists of 0's followed by 1's ( $0 \dots 0 1 \dots 1$ )

(Case 3)  $vxy$  consists of 1's followed by 0's ( $1 \dots 1 0 \dots 0$ )

(Case 1)  $vxy$  is in the 0's of section 1, 1's of section 1, 0's of section 2, or 1's of section 2. By property 1,  $v$  and/or  $y$  is a non-empty substring of one type of symbol in the section.

The string  $uv^2zy^2z$  will increase the number of one type of symbol in the section without increasing the number of the other type of symbol in the same section. So,  $uv^2zy^2z$  will not be of form  $0^n1^m0^m1^n$ .

Hence  $uv^2zy^2z \notin L$ .

(Case 2) In this case, we need to consider the sections differently

a)  $vxy$  is of form 0...01...1 and is in section 1  
By property 1,  $v$  and/or  $y$  is non-empty.

If  $v$  or  $y$  consists of two symbols (e.g.  $v=0\dots01\dots1$ ),  $uv^2zy^2z$  will not be of form 0..01..10..01..1, so  $uv^2zy^2z \notin L$

If  $v$  and/or  $y$  consists of one symbol (e.g.  $v=0\dots0, y=1\dots1$ ),  $uv^2zy^2z$  will result in the section 1 having more (or equal) number of 0's or 1's than section 2, so  $uv^2zy^2z \notin L$ .

b)  $vxy$  is of form 0..01..1 and is in section 2.

Similar problem as a) but we take  $uv^0zy^0z$ , which will result in less (or equal) number of 0's or 1's in section 2 than section 1.

So,  $uv^0zy^0z \notin L$

(Case 3)  $vxy$  consists of 1..10..0 somewhere in the middle of s (between section 1 and 2)  
By property 1,  $v$  and/or  $y$  is non empty.

If  $v$  or  $y$  consists of two symbols,  $uv^2zy^2z$  will not be of form 0..01..10..01..1, so  $uv^2zy^2z \notin L$ .

If  $v$  and/or  $y$  consists of one symbol, then  $uv^2zy^2z$  will increase the one type of symbol in a section without increasing the other type of symbol in the same section, which will not be of form  $0^n1^m0^m1^n$ .

So,  $uv^2zy^2z \notin L$ .

We have shown that there is no way to rewrite  $s = uvxyz$  satisfying all properties of the PL.

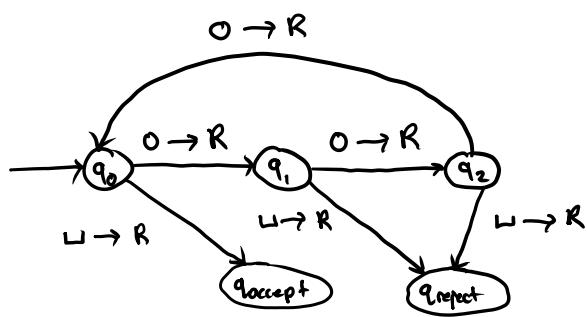
Therefore,  $L$  is not context-free.

### Question 3

- a)
  1. If the tape is empty, accept
  2. Scan the 0's on the tape from left to right in groups of 3  
If there are even less than 3 0's in a group, reject
  3. If we reach the end of the input string, accept

b)

M



$$M = (\{q_0, q_1, q_2, q_{\text{accept}}, q_{\text{reject}}\}, \{0\}, \{0, L\}, \delta, q_{\text{accept}}, q_{\text{reject}})$$

where  $\delta$  is described by the state diagram.