Introduction to AI: Useful Built-in predicates in Prolog (1/2)

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Outline

- Input/Output
- Controlling resolution
- 3 assert/retract

Input/Output on terms

- write/1 and writeq/1 : write(X) and writeq(X) display X on the output stream
- read/1: read(X) reads a term on the input stream and assigns it to X
- nl/0: displays a new line on the output stream

Outputting terms

write(T) causes the term T to be written to the current output stream, which by default is the user's screen

```
?- write(42).
42
true.
?- write(3.14).
3.14
true.
?- write('Artificial Intelligence').
Artificial Intelligence
true.
```

Outputting terms

write(T) causes the term T to be written to the current output stream, which by default is the user's screen

?- write(date(year(2020),month(december),day(1))).

```
date(year(2020),month(december),day(1))
true.
?- write([1,'a',date(year(2020),month(december),day(1)),[a,b]]).
```

```
[1,a,date(year(2020),month(december),day(1)),[a,b]] true.
```

```
?- write([1,'a',date(year(Y),month(M),day(D)),[a,b,c]]).
[1,a,date(year(_18088),month(_18092),day(_18096)),[a,b,c]]
true.
```

Outputting terms

 $\operatorname{writeq}(X)$ causes a quoted atom X to be written to the current output stream, as a quoted atom

```
?- writeq('Artificial Intelligence').
'Artificial Intelligence'
true.
?- writeq('student').
student
true.
?- writeq('Student').
'Student'
true.
?- writeq(42).
42
true.
```

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 $\operatorname{writeq}(X)$ causes a quoted atom term X to be written to the current output stream, as a quoted atom

```
?- writeq(date(year(2020),month(december),day(1))).
date(year(2020),month(december),day(1))
true.
?- writeq(date(year(2020),month('M'),day(1))).
date(year(2020),month('M'),day(1))
true.
?- write(date(year(2020),month('M'),day(1))).
date(year(2020),month(M),day(1))
true.
```

Outputting newline

nl/0 causes a new line to be written to the current output stream

```
go :-
   X=42, Y=73,
   write('Here is the value of X : '), write(X),
   nl,
   write('Here is the value of Y : '),write(Y).

?- go.
Here is the value of X : 42
Here is the value of Y : 73
true.
```

Inputting terms

 $\operatorname{read}(\mathsf{T})$ causes the term T to be read from the current input stream, which by default is the user's keyboard

```
?- read(T).
I: 42.
T = 42.
?- read(T).
l: 3.14.
T = 3.14.
?- read(T).
|: 'Artificial Intelligence'.
T = 'Artificial Intelligence'.
```

Note that you must enter a dot after the term entered.

Inputting terms

Input/Output

read(T) causes the term T to be read from the current input stream, which by default is the user's keyboard

```
?- read(T).
|: [1, 'a', date(year(2020), month(december), day(1)), [a,b]].
T = [1, a, date(year(2020), month(december), day(1)), [a, b]].
?- read(T).
|: [1,'a',date(year(Y),month(M),day(D)),[a,b,c]].
T = [1, a, date(year(21720), month(21724), day(21728)), [a, b, c]].
```

Inputting terms

 $\operatorname{read}(\mathsf{T})$ causes the term T to be read from the current input stream, which by default is the user's keyboard

T can be instantiated

```
?- T = mary, read(T).
|: mary.

T = mary.
?- ?- T = mary, read(T).
|: john.
false.
```

Input/Output on characters

- put/1: put(A) displays the character whose ASCII code is A on the output stream
- get0/1: get0(X) reads a character from the input stream and assigns it to X
- get/1: get(X) reads the next non-white-space character (i.e. character with an ASCII value less than or equal to 32) from the input stream

Input/Output on characters

ASCII TABLE

Decimal	Hex	Char	Decimal	Hex	Char	Decimal	Hex	Char	Decimal	Hex	Char
0	0	[NULL]	32	20	[SPACE]	64	40	@	96	60	
1	1	[START OF HEADING]	33	21	1	65	41	Α	97	61	a
2	2	[START OF TEXT]	34	22		66	42	В	98	62	b
3	3	[END OF TEXT]	35	23	#	67	43	C	99	63	C
4	4	[END OF TRANSMISSION]	36	24	\$	68	44	D	100	64	d
5	5	[ENQUIRY]	37	25	%	69	45	E	101	65	e
6	6	[ACKNOWLEDGE]	38	26	&	70	46	F	102	66	f
7	7	[BELL]	39	27	1	71	47	G	103	67	g
8	8	[BACKSPACE]	40	28	(72	48	H	104	68	h
9	9	[HORIZONTAL TAB]	41	29)	73	49	1	105	69	i
10	Α	[LINE FEED]	42	2A	*	74	4A	J	106	6A	j
11	В	[VERTICAL TAB]	43	2B	+	75	4B	K	107	6B	k
12	С	[FORM FEED]	44	2C	,	76	4C	L	108	6C	1
13	D	[CARRIAGE RETURN]	45	2D	-	77	4D	M	109	6D	m
14	E	[SHIFT OUT]	46	2E		78	4E	N	110	6E	n
15	F	[SHIFT IN]	47	2F	1	79	4F	0	111	6F	0
16	10	[DATA LINK ESCAPE]	48	30	0	80	50	P	112	70	р
17	11	[DEVICE CONTROL 1]	49	31	1	81	51	Q	113	71	q
18	12	[DEVICE CONTROL 2]	50	32	2	82	52	R	114	72	r
19	13	[DEVICE CONTROL 3]	51	33	3	83	53	S	115	73	S
20	14	[DEVICE CONTROL 4]	52	34	4	84	54	T	116	74	t
21	15	[NEGATIVE ACKNOWLEDGE]	53	35	5	85	55	U	117	75	u
22	16	[SYNCHRONOUS IDLE]	54	36	6	86	56	V	118	76	v
23	17	[ENG OF TRANS. BLOCK]	55	37	7	87	57	w	119	77	w
24	18	[CANCEL]	56	38	8	88	58	X	120	78	X
25	19	[END OF MEDIUM]	57	39	9	89	59	Υ	121	79	У
26	1A	[SUBSTITUTE]	58	3A	:	90	5A	Z	122	7A	Z
27	1B	[ESCAPE]	59	3B	;	91	5B	[123	7B	{
28	1C	[FILE SEPARATOR]	60	3C	<	92	5C	1	124	7C	1
29	1D	[GROUP SEPARATOR]	61	3D	=	93	5D	1	125	7D	}
30	1E	[RECORD SEPARATOR]	62	3E	>	94	5E	^	126	7E	~
31	1F	[UNIT SEPARATOR]	63	3F	?	95	5F	_	127	7F	[DEL]

Outputting characters

put(A) displays the character whose ASCII code is A on the current output stream, which by default is the user's screen

```
?- put(97).
a
true.
?- put(90).
Z
true.
?- put(35).
#
true.
```

get0(C) reads a character on the input stream (by default, the user's keyboard) and assigns it to C

```
?- get0(C).
1: a
C = 97.
?- get0(C).
1: 7.
C = 90.
?- get0(C).
1: #
C = 35.
```

get0(C) reads a character on the input stream (by default, the user's keyboard) and assigns it to C

T can be instantiated

```
?- C = 90, get0(C).
|: Z
C = 90.
?- C = 90, get0(C).
|: a
```

We can use get0 several times to input several characters

```
readline :-
   getO(X), continue(X).
continue(35).
continue(C):-
   C \== 35, put(C), readline.
?- readline.
|: this is a test#
this is a test
true .
```

We can extend this program, for example, to count the number of characters

```
readline(NumberChars):-
    readline(O,NumberChars).

readline(Count,NumberChars):-
    getO(C),continue(C,Count,NumberChars).

continue(35,Count,Count).
continue(C,Count,NumberChars):-
    C\==35,Count2 is Count+1,readline(Count2,NumberChars).
```

```
Let's run it
?- readline(N).
1: This is a test#
N = 14.
?- readline(N).
1: this is a test
1: and another one#
N = 30.
?- readline(N).
1: #
N = O.
```

Note we can write a simpler program (without using an accumulator):

```
readline(NumberChars):-
   getO(C),continue(C,NumberChars).

continue(35,0).
continue(C,NumberChars):-
   C\==35,
   readline(NumberChars2),
   NumberChars is NumberChars2 + 1.
```

We can extend this program to count the number of vowels

```
readline(NumberChars, NumberVowels):-
   readline(0, NumberChars, 0, NumberVowels).
readline(CountChars, NumberChars, CountVowels, NumberVowels):-
   getO(C),
   continue(C,CountChars,NumberChars,CountVowels,NumberVowels).
continue(35, CountChars, CountChars, CountVowels, CountVowels).
continue(C,CountChars,NumberChars,CountVowels,NumberVowels):-
   C\==35, CountChars2 is CountChars+1,
   isVowels(C, CountVowels, NumberVowels2),
   readline(CountChars2, NumberChars, NumberVowels2, NumberVowels).
isVowels(C,CountVowels,NumberVowels) :-
   vowel(C), NumberVowels is CountVowels + 1.
isVowels(_,CountVowels,CountVowels).
vowel(97). vowel(101). vowel(105). vowel(111). vowel(117). vowel(121).
```

Let's run it!

```
?- readline(N,V).
1: this is a test#
N = 14,
V = 4
?- readline(N,V).
|: zrtpml#
N = 6.
V = 0.
?- readline(N,V).
1: #
N = V, V = O.
```

Note we can write a simpler program (without using accumulators):

```
readline(NumberChars, NumberVowels):-
   get0(C),continue(C,NumberChars,NumberVowels).
continue(35,0,0).
continue(C, NumberChars, NumberVowels):-
   C = 35.
   vowel(C),
   readline(NumberChars2, NumberVowels2),
   NumberVowels is NumberVowels2 + 1,
   NumberChars is NumberChars2 + 1.
continue(C, NumberChars, NumberVowels):-
   C = 35.
   readline(NumberChars2, NumberVowels),
   NumberChars is NumberChars2 + 1.
vowel(97). vowel(101). vowel(105). vowel(111). vowel(117). vowel(121).
```

Defining the input stream

- The default current input stream is the user's keyboard
- The current input stream can be changed using the see/1 predicate
- We can restore the current input stream to be the user's keyboard by using the seen/0 predicate
- We can determine the name of the current input stream by solving seeing(F). F is unified to this name

Defining the input stream

For example we can run some previous predicates on a particular file foo.txt instead of the keyboard

```
readFile(NumberChars, NumberVowels) :-
    see('foo.txt'),
    readline(NumberChars, NumberVowels),
    seen.
```

Defining the output stream

- The default current output stream is the user's terminal
- The current output stream can be changed using the tell/1 predicate
- We can restore the current output stream to be the user's terminal by using the told/0 predicate
- We can determine the name of the current output stream by solving telling(F). F is unified to this name

Defining the output stream

For example we can copy a text from the keyboard to a file

```
copyText(Outfile):-
   tell(Outfile),
   copy_characters,
   told.

copy_characters:-
   getO(C),continue(C).

continue(35). /* ASCII value of # */
continue(N):-
   N\==35,
   put(N),copy_characters.
```

Defining the output stream

And then run:

```
?- copyText('fic1.txt').
|: this is a test
|: let's look at the content of
|: the file now#

true .

If we look at fic1.txt, it contains
this is a test
let's look at the content of
the file now
```

repeat

Recursivity may be good, but sometimes it's better to backtrack

- repeat/0
- it succeeds whenever it is called
- it also succeeds (as many times as necessary) on backtracking

This can be used to create a looping effect

repeat

An example

```
get_answer(A):-
   write('Enter an answer to the question'),nl,
   repeat,
   write('Your answer (yes or no): '),read(A),
   valid(A),
   write('Your answer is '),write(A),nl.

valid(yes).
valid(no).
```

repeat

Let's run it!

```
?- get_answer(A).
Enter an answer to the question
Your answer (yes or no): azer.
Your answer (yes or no): |: abc.
Your answer (yes or no): |: yes.
Your answer is yes
A = yes .
```

fail/0 is a predicate that always fail Why it may be useful?

```
student(mary).
student(john).
student(suzy).

allStudents :-
    student(S), write(S), write(' is a student'),nl,
    fail.
allStudents.
```

Let's run it!

?- allStudents.
mary is a student
john is a student
suzy is a student
true.

Consider another example

```
person(john,smith,45,lyon,teacher).
person(martin,williams,33,paris,teacher).
person(henry,smith,26,dijon,student).
person(jane,wilson,62,lyon,teacher).
person(mary,smith,29,paris,student).

allJobs(J):-
   person(FirstName,Name,_,_,J),
   write(FirstName),write(''),write(Name),nl,
   fail.
allJobs(_).
```

fail

Let's run it!

```
?- allJobs(teacher).
john smith
martin williams
jane wilson
true.
?- allJobs(student).
henry smith
mary smith
true.
?- allJobs(policeman).
true.
```

Exercise: Write a predicate findUnderAge(A) that list all the persons that are less than A years old.

A last example

What we call cut is the symbol!

It changes the way Prolog naturally backtracks

It prunes the derivation tree

Consider the following program and goals

```
max(A,B,B) :- B > A.

max(A,B,A).

?- max(3,4,M).

M = 4;

M = 3.

?- max(4,3,M).

M = 4.
```

This is not a good way to write max because the first goal gives one good solution but also one bad, due to the second clause

Obviously, we could write

```
\max(A,B,B) :- B > A.
\max(A,B,A) :- B =< A.
?- \max(3,4,M).
M = 4;
false
?- \max(4,3,M).
M = 4.
```

For this very basic example, finding the good condition to avoid unwanted backtracking is easy

BUT, this is not always the case

Consider a second example with some goals

```
sum(1,1).
sum(N,S):-
    N1 is N-1,
    sum(N1,S1),
    S is S1+N.

?- sum(6,S).
S = 21;
ERROR: Out of local stack
```

We get a solution, but if we ask for another one, the program loops

Obviously, we could write

```
sum(1,1).
sum(N,S):-
    N > 1,
    N1 is N-1,
    sum(N1,S1),
    S is S1+N.

?- sum(6,S).
S = 21;
false.
```

For this very basic example, finding the good condition to avoid unwanted backtracking is easy

BUT, this is not always the case

The cut may be used in such situations

```
max(A,B,B) :- B > A, !.
max(A,B,A).

?- max(3,4,M).
M = 4.

?- max(4,3,M).
M = 4.
```

The cut may be used in such situations

```
sum(1,1) :- !.
sum(N,S):-
   N1 is N-1,
   sum(N1,S1),
   S is S1+N.

?- sum(6,S).
S = 21.
```

More generally

```
p(X) :-
    condition(X),r(X,Y),s(Y).
p(X) :-
    \+ condition(X), t(X), u(X,Y).
```

BUT solving condition(X) may be expensive (time and memory)

Moreover, if condition(X) is false, we need to prove it again to prove the negation is true

It can be a big waste of time

Using the cut we can be more efficient

In fact it's an if then else...

```
p(X) :-
    condition(X),!, r(X,Y),s(Y).
p(X) :-
    t(X), u(X,Y).
```

Cut with failure

Consider this program

```
bird(eagle).
bird(duck).
bird(crow).
bird(ostrich).
bird(swan).

can_fly(ostrich) :-
   !,
   fail.
can_fly(X) :-
   bird(X).
```

Here the cut is useful to specify exceptions to general rules

Cut (recap)

Consider this program

```
p(a).
p(b):-!.
p(c).
```

What will produce the following goals:

```
?- p(X).

?- p(X), !.

?- p(X), p(Y).

?- !, p(X), p(Y).

?- p(X), !, p(Y).

?- p(X), p(Y), !.
```

Cut and derivation tree

Let's look at the impact of the cut on the derivation tree Consider the following program

```
p3(X,Y,Z) :- p2(X,Z), t(Y).

p2(X,Y) :- s(X), q(Y).
p2(X,Y) :- q(Z), r(X,Z,Y).

s(a). s(b). s(c).
q(a). q(e).
t(42). t(21).

r(f,a,g).
r(g,e,d).
```

Draw the derivation tree of the resolution of the goal: ?-p3(A,B,C).

Cut and derivation tree

Now consider the following program that contains a cut

```
p3(X,Y,Z) :- p2(X,Z), t(Y).

p2(X,Y) :- s(X), !, q(Y).

p2(X,Y) :- q(Z), r(X,Z,Y).

s(a). s(b). s(c).
q(a). q(e).
t(42). t(21).

r(f,a,g).
r(g,e,d).
```

Show on the previous derivation tree which branches are pruned

What happens if we run the goal: ?- p3(e,B,C).

assert/1 and retract/1

Allow dynamic change of the Prolog program

- asserta/1: asserta(C) adds clause C in the program (at the top of other clauses with the same predicate)
- assertz/1: assertz(C) adds clause C in the program (at the bottom of other clauses with the same predicate)
- retract/1: retract(C) deletes from the program the first clause that unifies with C (first from the top)
- retractall/1: retractall(H) deletes from the program all the clauses C whose head unifies with H

Those predicates are generally used to add new knowledge that has been learnt or delete knowledge that becomes outdated

asserta/assertz

To allow predicate p1 of arity n to be asserted or retracted it must be defined as dynamic: :- dynamic(p1/n). If you do not do that, you get an error while running assert or retract on p1

```
ERROR: No permission to modify static procedure 'p/1' ERROR: In:
```

ERROR: [8] retract(p(_8920))

ERROR: [7] <user>

?- retract(p1(X)).

With SWI Prolog assert(p1(a)) works even if you did not define p1/1 as dynamic (but retract fails)

Suppose we have no clause defining the predicate $\ensuremath{\text{p}}/1$

```
?- assertz(p(1)).
?- assertz(p(2)).
?- listing(p).
p(1).
p(2).
```

Let's continue to add facts

```
?- asserta(p(3)).
?- assertz(p(4)).
?- asserta(p(5)).
?- listing(p).
p(5).
p(3).
p(1).
p(2).
p(4).
```

We can also add clauses with head and body

```
?- asserta((grand_parent(X,Y) :- parent(X,Z), parent(Z,Y))).
?- listing(grand_parent).
:- dynamic grand_parent/2.
grand_parent(A, C) :-
   parent(A, B),
   parent(B, C).
```

Note the parentheses around the clause! and there is no dot at the end of the clause!

We can also add clauses with head and body

```
?- asserta((grand_parent(X,Y) :- parent(X,Y))).
?- listing(grand_parent).
:- dynamic grand_parent/2.
grand_parent(A, B) :-
   parent(A, B).
grand_parent(A, C) :-
   parent(A, B),
   parent(B, C).
```

retract(C) deletes from the program the first clause that unifies with C (first from the top)
Suppose we have those facts

```
:- dynamic(student/1).
student(mary).
student(john).
student(suzy).
```

Then look at what happens when we run various goals using retract/1

```
?- retract(student(john)).
true.
?- listing(student)
:- dynamic(student/1).
student(mary).
student(suzy).
true.
```

Take care:

```
?- retract(student(X)).
X = mary ;
X = suzy
?- listing(student).
:- dynamic student/1.
true.
```

```
But if we do not hit ";"
?- retract(student(X)).
X = mary
?- listing(student).
:- dynamic student/1.
student(suzy).
true.
```

Suppose we have the program:

```
p(X,Y) := q(X), r(X).

p(X,Y) := s(X), t(Y).

p(a,b).

p(c,d).
```

Then run retract/1 (and hit ";" each time we get a solution):

```
p(X,Y) := q(X), r(Y).
p(X,Y) := s(X), t(Y).
p(a,b).
p(c,d).
?- retract(p(A,B)).
A = a, B = b
A = c, B = d.
?- listing(p/2).
:- dynamic p/2.
p(A, B) := q(A), r(B).
p(A, B) := s(A), t(B).
```

true.

Continue to run retract/1 on the program:

```
p(X,Y) := q(X), r(Y).
p(X,Y) := s(X), t(Y).
and get:
?- retract((p(A,B) :- s(A),t(B))).
true.
?- listing(p/2).
:- dynamic p/2.
p(A, B) := q(A), r(B).
true.
```

Finally, on the program:

true.

Another example

Consider a file containing some information about persons

```
mary. smith. 22. 'Saint-Etienne'. student.
john. smith. 25. 'Lyon'. 'Computer Scientist'.
mike. portnoy. 51. 'New York'. musician.
stephen. hough. 58. 'London'. musician.
suzy. gallagher. 21. 'Paris'. student.
end.
```

Another example

We want to write a program that reads this file and add this information into the Prolog database

```
readfile :-
   see('people.txt'),
   read_data,
   write('Finished converting data.'), nl,
   seen.
read data:-
   read(FirstName), continue(FirstName).
continue(end).
continue(FirstName):-
   read(Name), read(Age), read(Town), read(Job),
   assertz(person(FirstName, Name, Age, Town, Job)),
   read data.
```

Another example

Running ?- readfile. we get the knowledge based defined by the person/5 predicate

```
?- listing(person).
:- dynamic person/5.
person(mary, smith, 22, 'Saint-Etienne', student).
person(john, smith, 25, 'Lyon', 'Computer Scientist').
person(mike, portnoy, 51, 'New York', musician).
person(stephen, hough, 58, 'London', musician).
person(suzy, gallagher, 21, 'Paris', student).
true.
```

assert/retract

Memorize previously computed values

Suppose we want to calculate the Fibonacci function of a number

The definition of this function is:

$$\begin{cases} fib(1) = 1\\ fib(2) = 1\\ fib(n) = fib(n-1) + fib(n-2) \text{ for } n > 2 \end{cases}$$

We can write the following Prolog program

```
fib(1, 1).
fib(2, 1).
fib(N, F) :-
   N > 2,
   N1 is N - 1, fib(N1, F1),
   N2 is N - 2, fib(N2, F2),
   F is F1 + F2.
```

assert/retract

Memorize previously computed values

Let's look at how it works on a small example

```
fib(5, F0)
   fib(4, F1)
       fib(3, F2)
          fib(2, F3) \rightarrow F3 = 1
          fib(1, F4) \rightarrow F4 = 1
          -> F2 = F3+F4 = 1+1 = 2
       fib(2, F5) \rightarrow F5 = 1
       -> F1 = F2+F5 = 2+1 = 3
   fib(3, F6)
       fib(2, F7) \rightarrow F7 = 1
       fib(1, F8) \rightarrow F8 = 1
       -> F6 = F7+F8 = 1+1 = 2
-> F0 = F1+F6 = 3+2 = 5
```

We can see that fib(3,X) is calculated twice

Memorize previously computed values

Using asserta/1 we can memorize some intermediate calculations

```
lazyFib(1, 1).
lazyFib(2, 1).
lazyFib(N, F) :-
   N > 2,
   N1 \text{ is } N - 1, \text{ lazyFib}(N1, F1),
   N2 is N - 2, lazyFib(N2, F2),
   F is F1 + F2.
   asserta(lazyFib(N,F)).
```

Memorize previously computed values

Let's look at how it works on a small example

```
lazyFib(5, F0)
   lazyFib(4, F1)
      lazyFib(3, F2)
          lazyFib(2, F3) \rightarrow F3 = 1
          lazyFib(1, F4) \rightarrow F4 = 1
          -> F2 = F3+F4 = 1+1 = 2
          asserta(lazyFib(3,2))
      lazyFib(2, F5) \rightarrow F5 = 1
       -> F1 = F2+F5 = 2+1 = 3
       asserta(lazyFib(4,3))
   lazyFib(3, F6)
       -> use the fact lazyFib(3,2)
      -> F6 = 2
-> F0 = F1+F6 = 3+2 = 5
asserta(lazyFib(5,5))
```

We can define a predicate that can remove some information about a particular person

```
remove(FirstName,Name):-
   retract(person(FirstName,Name,_,_,_)).
```

```
person(mary, smith, 22, 'Saint-Etienne', student).
person(john, smith, 25, 'Lyon', 'Computer Scientist').
person(mike, portnoy, 51, 'New York', musician).
person(stephen, hough, 58, 'London', musician).
person(suzy, gallagher, 21, 'Paris', student).
?- remove(john, smith).
true.
?- listing(person).
:- dynamic person/5.
person(mary, smith, 22, 'Saint-Etienne', student).
person(mike, portnoy, 51, 'New York', musician).
person(stephen, hough, 58, 'London', musician).
person(suzy, gallagher, 21, 'Paris', student).
true.
```

Finally we can define a predicate that can update some information about a particular person

```
update(FirstName,Name,NewJob):-
   person(FirstName,Name,Age,City,Job),!,
   retract(person(FirstName,Name,Age,City,Job)),
   assertz(person(FirstName,Name,Age,City,NewJob)).
update(FirstName,Name,NewJob):-
   write('Sorry but '),
   write(FirstName), write(' '),
   write(Name), write(' is not in the database...').
```

And run it:

```
?- update(mary,smith,doctor).
true.
?- listing(person).
:- dynamic person/5.
person(john, smith, 25, 'Lyon', 'Computer Scientist').
person(mike, portnoy, 51, 'New York', musician).
person(stephen, hough, 58, 'London', musician).
person(suzy, gallagher, 21, 'Paris', student).
person(mary, smith, 22, 'Saint-Etienne', doctor).
true.
```