



DP - Part IV

Complete Course on Algorithm for GATE - CS & IT

Searching

Linear Search

i/p: An array of n -ele, ele $= x$

o/p: return position of x if it found else return (-1)

ex

A $\begin{bmatrix} 20 & 50 & 20 & 10 & 80 & 90 & 100 & 5 & 6 & 200 & 7 & 400 \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 \end{bmatrix}$

$x=6$	$x=20$	$x=100$	$x=500$	$x=400$
9	1	7	-1	12

BC: 1 (minimum Time to solve)

WC: n (maximum Time to solve the problem)


$$Avg = 1 + 2 + 3 + 4 + 5 + 6 + \dots + n-1 + n$$

$$= \frac{\cancel{n(n+1)}}{2} = \frac{n+1}{2}$$

all $LS - TC - T(n)$

$$T(n) = O(n)$$

$$= \Omega(1)$$

$$= \cancel{\Theta(n)}$$


In general: $BC \leq AV \leq WC$

~~all $LS - WC - T(n)$~~

~~$T(n) = O(n)$~~

~~$= \Omega(n)$~~

~~$= \Theta(n)$~~

~~$= \cancel{\Theta(1)}$~~

~~all $LS - BC - TC - T(n)$~~

~~$T(n) = \Theta(1)$~~

$\frac{0+200}{2}$

100

BS

i/p: Sorted array of n -ele, ele- x

o/p: return position of x if found else (-1)

ex

A [$\begin{matrix} 10 \\ 1 \end{matrix}$ $\begin{matrix} 20 \\ 2 \end{matrix}$ $\begin{matrix} 30 \\ 3 \end{matrix}$ $\begin{matrix} 40 \\ 4 \end{matrix}$ $\begin{matrix} 50 \\ 5 \end{matrix}$ $\begin{matrix} 60 \\ 6 \end{matrix}$ $\begin{matrix} 70 \\ 7 \end{matrix}$]

$x = 10$

$$\text{mid} = \frac{1+7}{2}$$
$$= 4$$

$$= \frac{1+3}{2}$$

$$= 2$$

$BS(a, i, j, x) \Rightarrow T(n)$

$if(i == j) \left\{ \begin{array}{l} if(a[i] == x) \\ return(i) \\ else \\ return(-1) \end{array} \right\}$

$O(1)$

Partial
Application
Divide & Conquer

$else \left\{ \right.$

$mid = \lfloor (i+j)/2 \rfloor$
4

$if(a[mid] == x) return(mid) \Rightarrow C$

$else$

$if(x < a[mid])$

$BS(a, i, mid-1, x)$

$else$

$BS(a, mid+1, j, x)$

$\Rightarrow T(n/2)$

conquer

Divide

2

Let $T(n)$ be the TC of above algo.

RR-TC { Worst case
Avg case }

$$T(n) = \begin{cases} O(1) & \text{if } n=1 \\ c + T(n/2) & \text{if } n>1 \end{cases}$$

$$T(n) = T(n/2) + \underline{c}$$

$$= T(n/2^2) + c + c$$

$$= T(n/2^3) + c + c + c$$

$$\downarrow \underline{4n}$$

$$= T(1) + c \cdot 4n$$

$$= O(1) + c \cdot \log n \Rightarrow O(\log n)$$

Best case

$$T(n) = O(1)$$

Stack Space (wc)
Ac

$$\begin{array}{lcl} T(n) = O(\log n) & & \\ \swarrow \quad \searrow & & \\ \log n & 1 & = n(1) \\ & & = \underline{\underline{O(\log n)}} \\ & & = \underline{\underline{O(1)}} \end{array}$$

BS(a, i, j, x)

(10 20 30) (40) (50 60 70)
 1 2 3 4 5 6 7

while(i ≤ j)

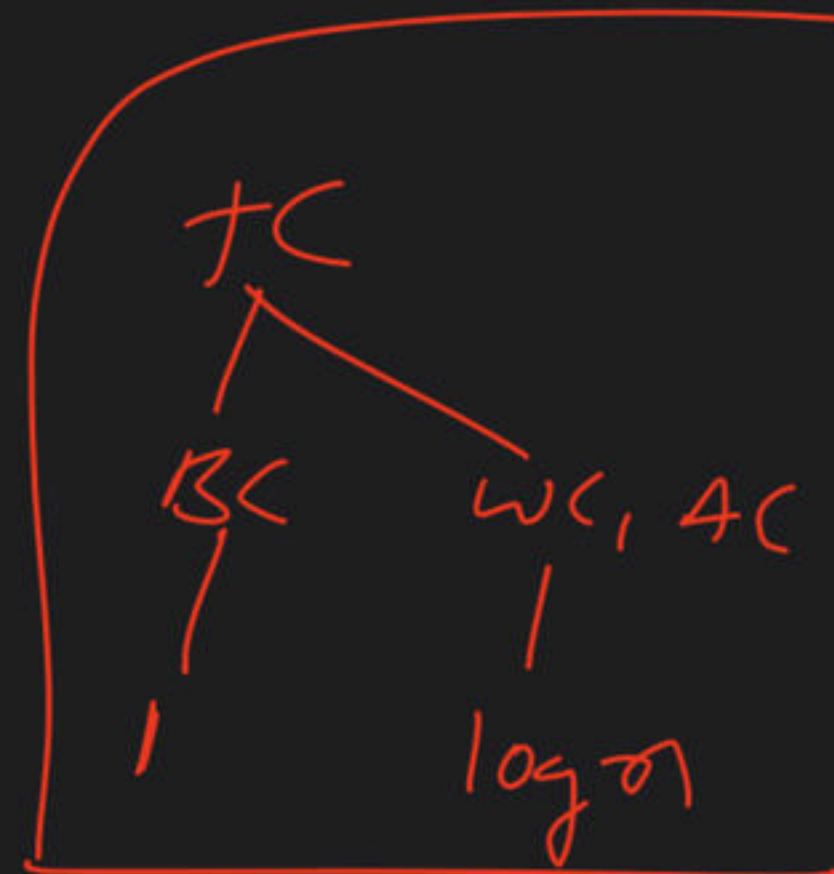
if(i == j) {
 if(a[i] == x)
 return(i)
 else
 return(-1)
}

else {
 mid = $\lfloor (i+j)/2 \rfloor$

if(a[mid] == x) return(mid)

else
 if(x < a[mid])
 else j = mid - 1

else i = mid + 1



Space
 ↓
 $\Theta(1)$