

Recursion - Part I

Complete Course on Algorithm for GATE - CS & IT

$$\textcircled{4} \quad f(n) = n + 5 \\ g(n) = n$$

$$n + 5 \leq \underbrace{c}_{=2} \cdot \underbrace{n}_{\geq 5}, \quad \forall n, n \geq n_0$$
$$n + 5 = O(n)$$

$$\textcircled{5} \quad f(n) = n, \quad g(n) = n + 5$$
$$n \leq \underbrace{c}_{=1} \cdot (n + 5), \quad n_0 = 1$$
$$n = O(n + 5)$$

ex

$$f(n) = n^2 + n + 10, \quad g(n) = n^2$$

$$f(n) = O(g(n))?$$



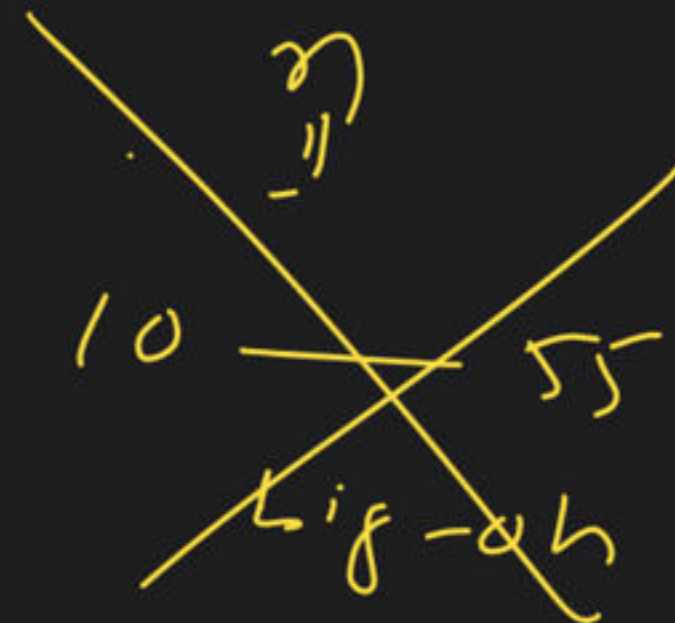
$$n^2 + n + 10 \leq c \cdot \underline{n^2}, \quad \forall n, \quad n \geq n_0$$



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$$n^2 + n + 10 = O(n^2)$$



ex

$$f(n) = n^2, \quad g(n) = n$$

$$f(n) = O(g(n)) \Leftrightarrow n^2 \leq c \cdot n, \quad \forall n, \quad n \geq n_0 \Leftrightarrow n^2 \neq O(n)$$



Theta - notation

$$f(n) = \Theta(g(n))$$

iff



① $f(n) \geq c_1 \cdot g(n)$

and

② $f(n) \leq c_2 \cdot g(n)$

$\forall n, n \geq n_0$

Omega-Notation (Ω, \geq)

$$f(n) = \Omega(g(n))$$

\Leftrightarrow

$$f(n) \geq C \cdot g(n), \forall n, n \geq n_0$$

Such that \exists 2 constants $C > 0, n_0 \geq 1$

ex

$f(n) = n, \quad g(n) = n+5$

$$f(n) = \Omega(g(n))$$

\Leftrightarrow

① $1 \cdot n \geq \underset{\downarrow 1/2}{C} \cdot (n+5), \forall n, n \geq \underset{\downarrow 5}{n_0} \iff \underline{\underline{n = \Omega(n+5)}}$

$$n \leq \underset{\downarrow 1}{C_1} \cdot \underset{\downarrow 2}{(n+5)}, \forall n, n \geq \underset{\downarrow 1}{n_0}$$

\checkmark

$$n = O(n+5)$$

$C_1 = 1/2, C_2 = 1$
 $n_0 = 5$

$n = O(n+5)$

ex

$$f(n) = n^2 + n + 10, \quad g(n) = n^2$$

$$f(n) = \Theta(g(n))?$$

Asymptotically
equal.

①

$$n^2 + n + 10 \geq c_1 \cdot n^2, \quad \forall n, \quad n \geq n_0$$

\Downarrow

$$n^2 + n + 10 = \Omega(n^2)$$

\Downarrow

②

$$n^2 + n + 10 \leq c_2 \cdot n^2, \quad \forall n, \quad n \geq n_0$$

\Downarrow
3

$$n^2 + n + 10 = O(n^2)$$

\Downarrow
3

$$c_1 = 1, c_2 = 3 \\ n_0 = 3$$

$$n^2 + n + 10 = \Theta(n^2)$$

ex

$$f(n) = n$$

$$g(n) = n$$

$$f(n) = \Theta(g(n)) \quad ?$$

$$n \geq c_1 \cdot n, \quad \forall n, \quad n \geq n_0$$

\Downarrow

$$n \leq c_2 \cdot n, \quad \forall n, \quad n \geq n_0$$

\Downarrow

$c_1 = 1, c_2 = 1$

$n_0 = 1$

$n = \Theta(n)$

ex

$$f(n) = n$$

$$g(n) = n^2$$

$$f(n) = \Theta(g(n)) ?$$

$$n = \Omega(n^2) \quad \text{No}$$

$$n \geq c_1 \cdot n^2, \quad \forall n, n \geq n_0$$

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$$n \leq c_2 \cdot n^2, \quad \forall n, n \geq n_0$$

$$n = O(n^2) \quad \checkmark$$

$$\boxed{n \neq \Theta(n^2)}$$

ex ②

$$f(n) = n + 5, \quad g(n) = n$$

$$f(n) = \Omega(g(n))$$



$$n + 5 \geq c \cdot n, \quad \forall n, n \geq n_0$$



$$n + 5 = \Omega(n)$$

ex $f(n) = n, \quad g(n) = n^2$

$$f(n) = \Omega(g(n))$$



$$n \geq c \cdot n^2, \quad \forall n, n \geq n_0$$



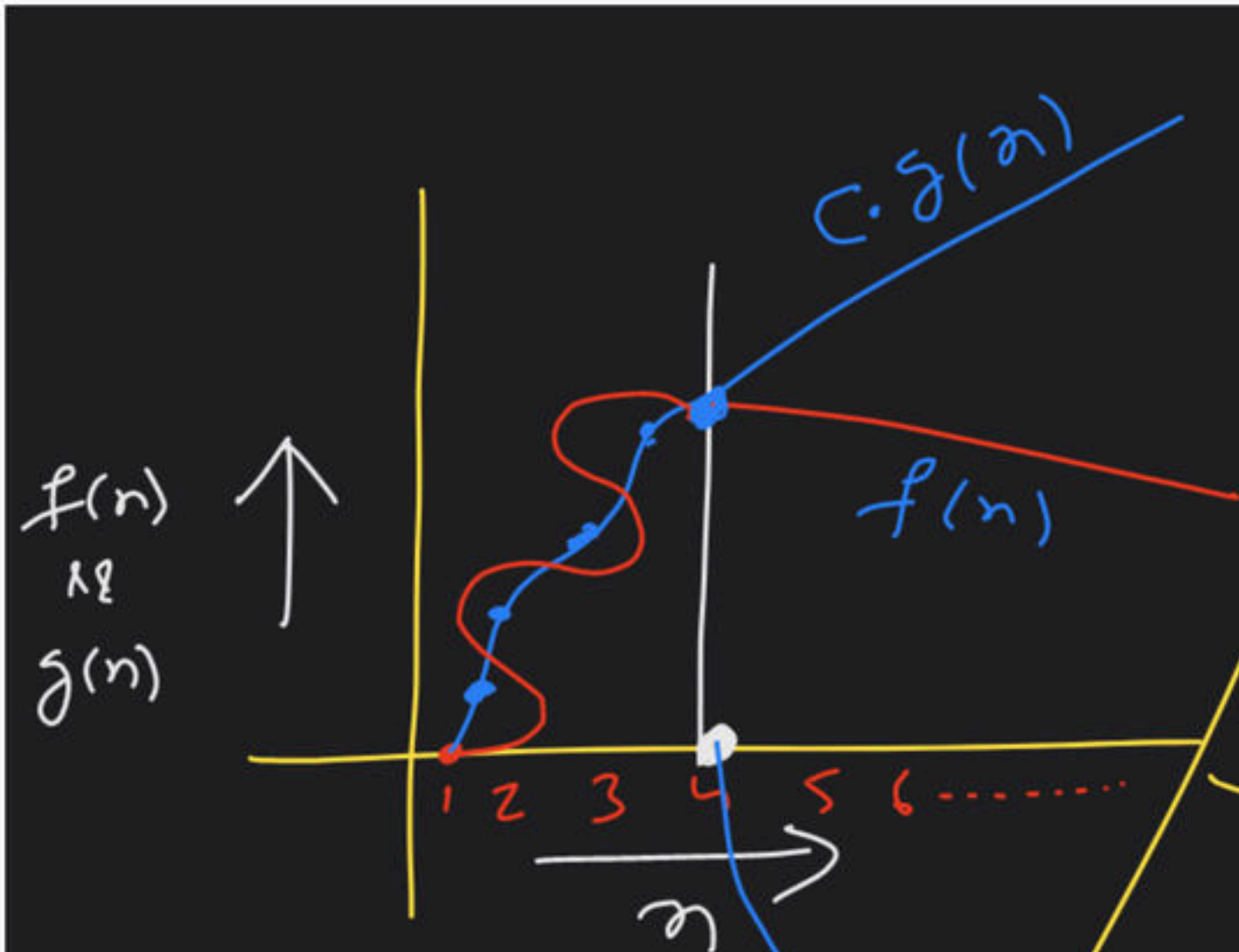
$$n \neq \Omega(n^2)$$

$-f(n)$

$g(n)$

$$1 \cdot n^2 \geq c(n^2 + n + 1) \quad \forall n, n \geq n_0$$

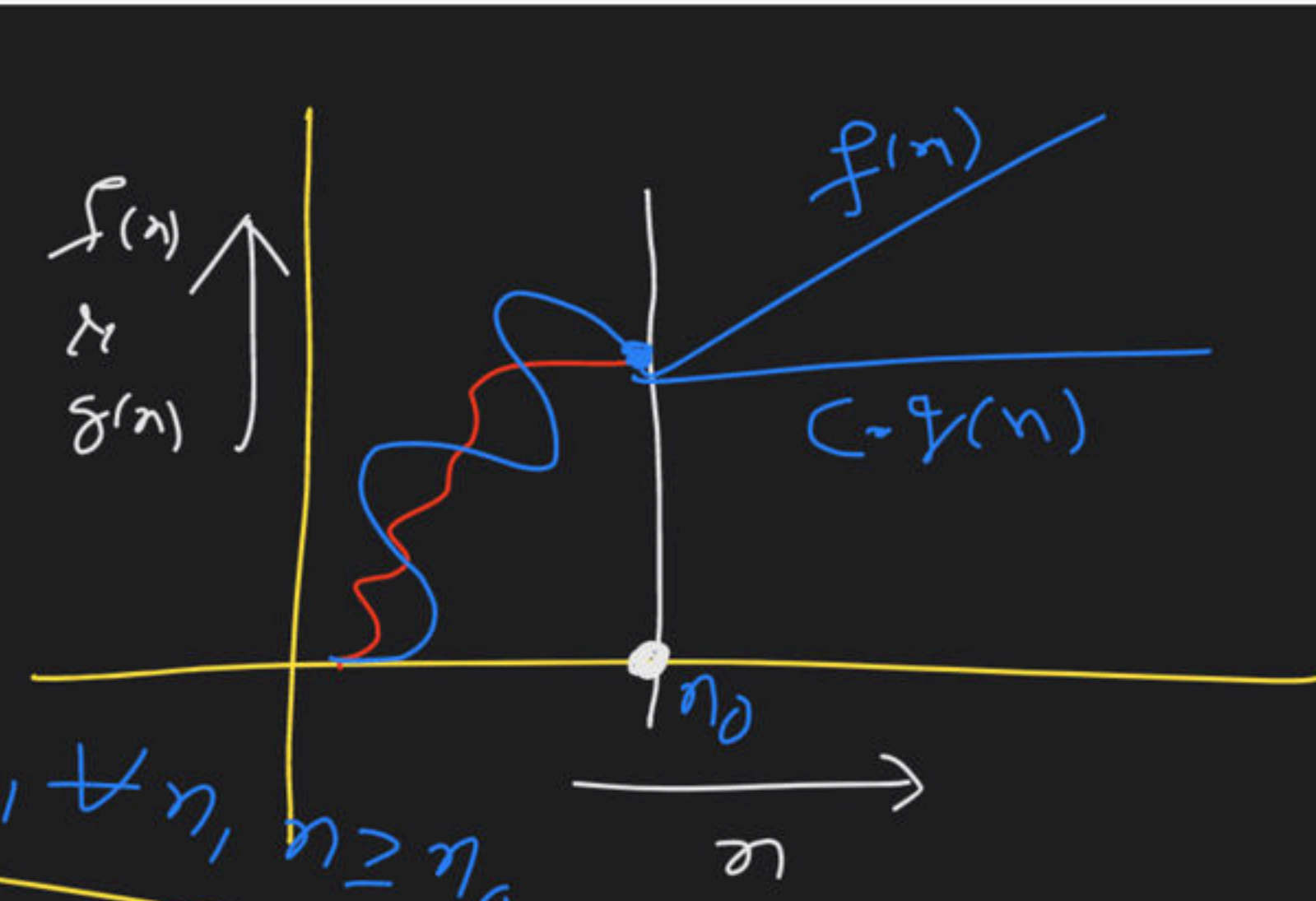
$$n^2 = \Omega(n^2 + n + 1)$$



$$f(n) = \Omega(g(n))$$

$$\Leftrightarrow$$

$$f(n) \geq c \cdot g(n), \forall n, n \geq n_0$$

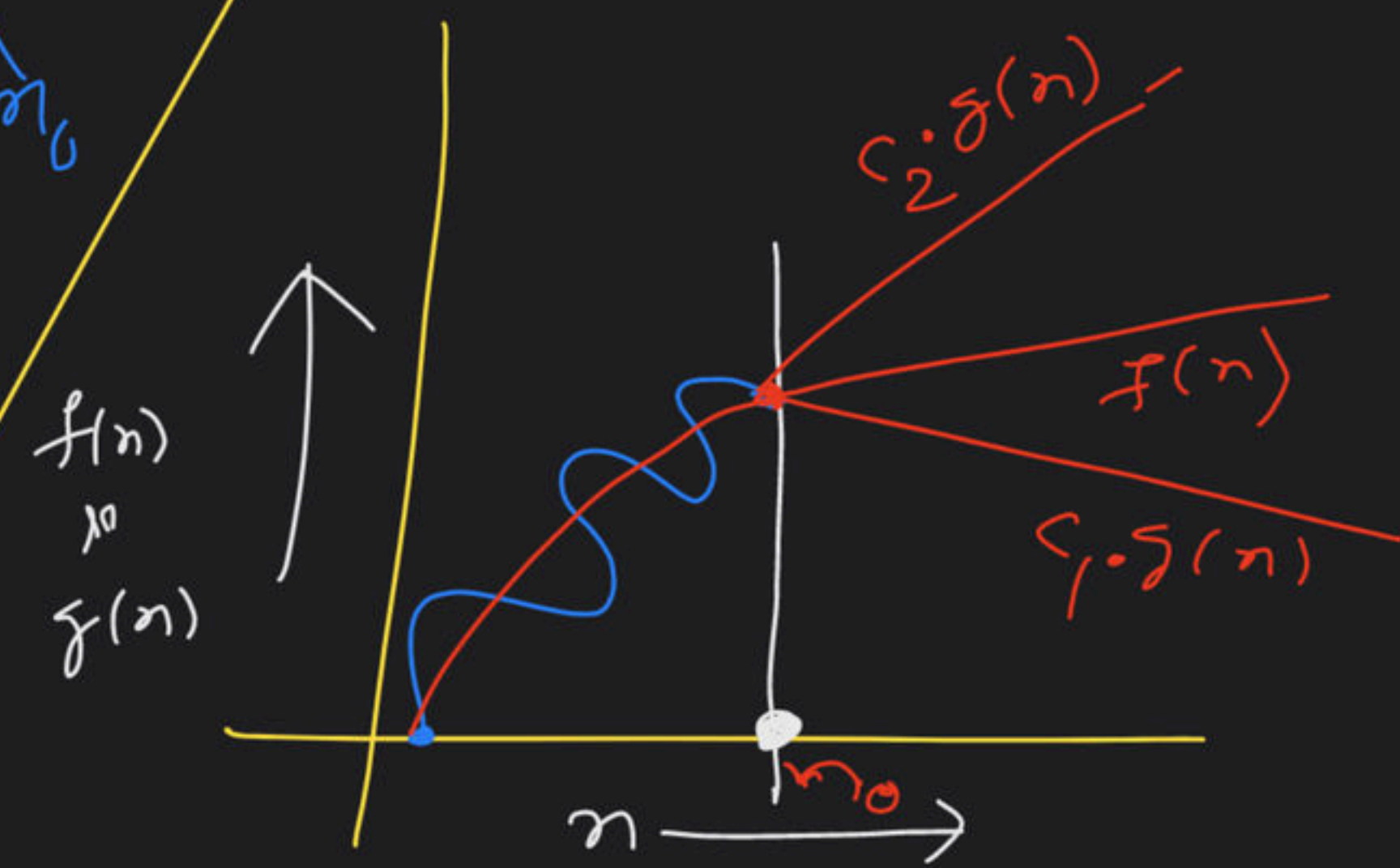


$$f(n) = O(g(n))$$

$$\Leftrightarrow$$

$$f(n) \leq c \cdot g(n)$$

$$\forall n, n \geq n_0$$



$$f(n) = \Theta(g(n))$$

$$\Leftrightarrow$$

$$(1) f(n) \geq c_1 \cdot g(n)$$

$$\text{and}$$

$$(2) f(n) \leq c_2 \cdot g(n)$$

