

Intermediate Code Generation

Part-2

Complete Course on Compiler Design

ex

$$E_1 \rightarrow E_2 + T \quad \{ E_1 \cdot val = E_2 \cdot val + T \cdot val \}$$

$$| T \quad \{ E_1 \cdot val = T \cdot val \}$$

$$T_1 \rightarrow T_2 * F \quad \{ T_1 \cdot val = T_2 \cdot val * F \cdot val \}$$

$$F \rightarrow (E) \quad \{ F \cdot val = 1; \quad \underline{E \cdot val = G \cdot val} \}$$

$$| num \quad \{ F \cdot val = num \}$$

$$G \rightarrow num \quad \{ G \cdot val = num \}$$

$$T \rightarrow F \quad \{ T \cdot val = F \cdot val \}$$

~~5. $E \cdot val$~~
~~6. $E \cdot val$~~
~~None~~

Intermediate Code Generation

ex

$$x = a + b * c$$

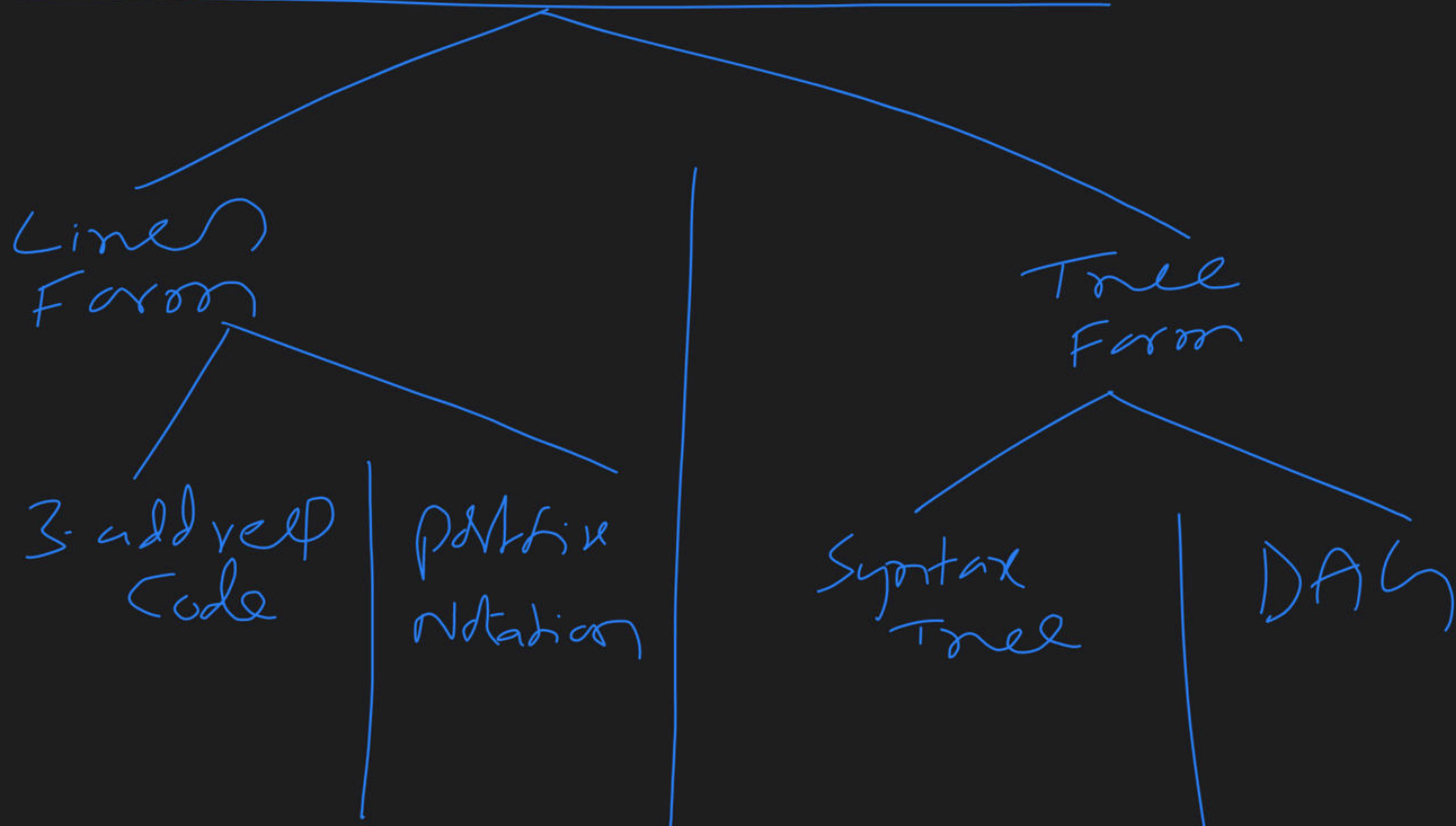
\Downarrow

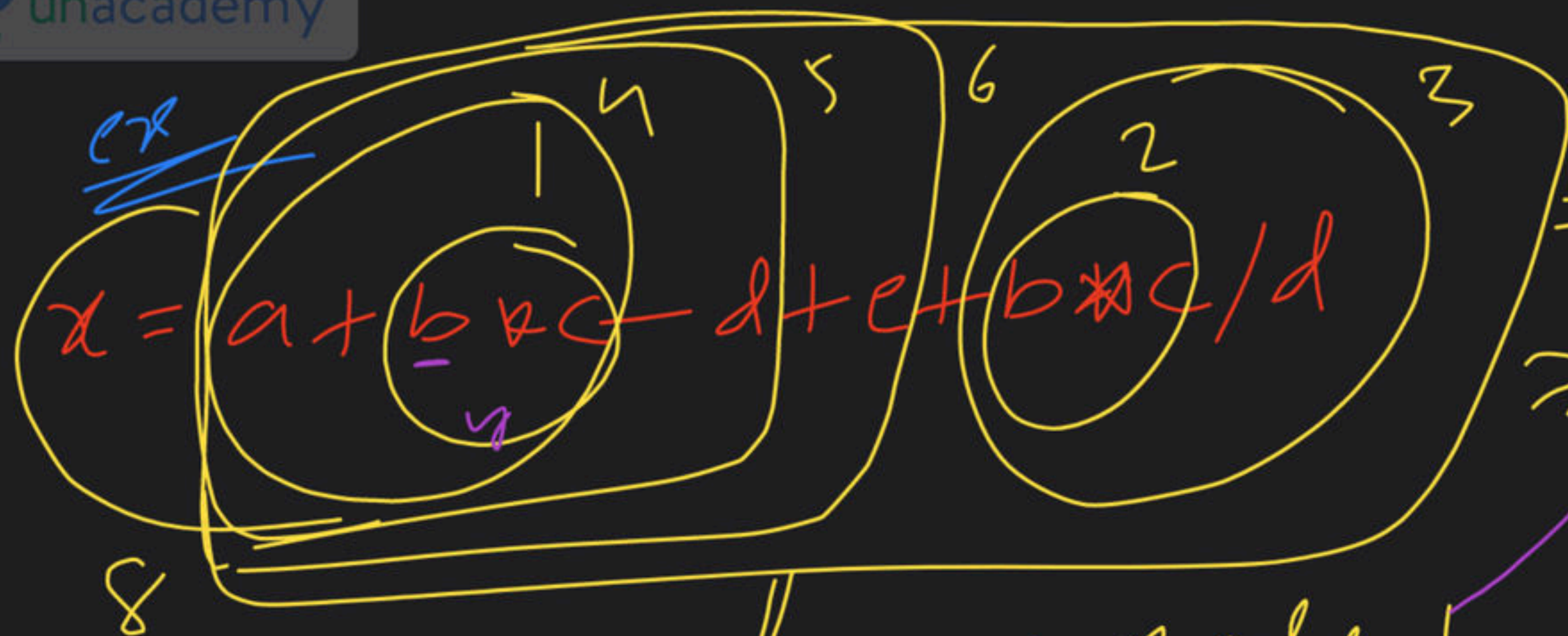
$$t_1 = b * c$$

$$t_2 = a + t_1$$

$$x = t_2$$

Representation of Intermediate Code



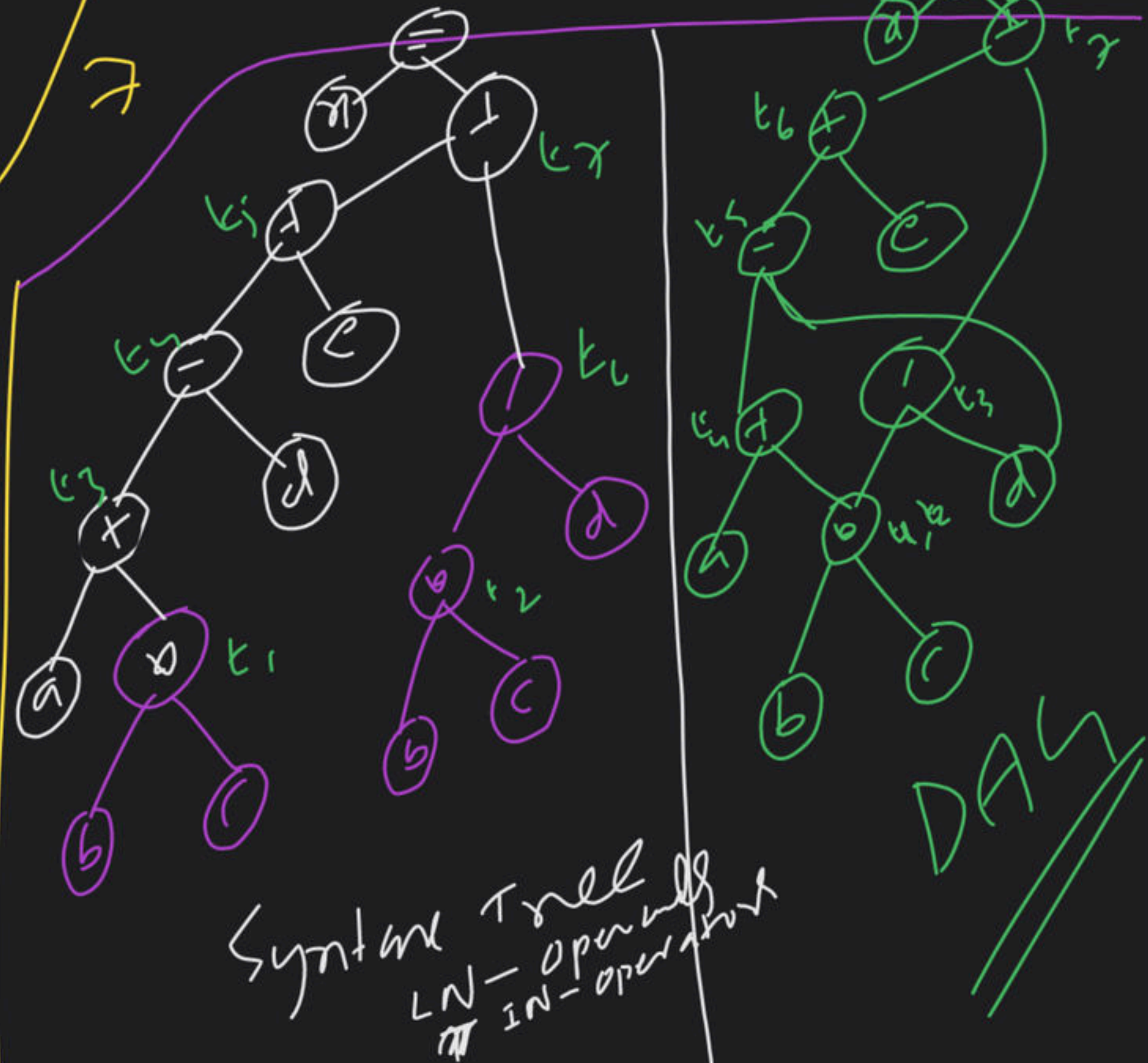


Praxis NALM

$$x = a + b * c - d + e + b * c / d$$

3-address code

$$\begin{aligned} t_1 &= b * c \\ t_2 &= b * c \\ t_3 &= t_2 / d \\ t_4 &= a + t_1 \\ t_5 &= t_4 - d \\ t_6 &= t_5 + e \\ t_7 &= t_6 + t_3 \\ x &= t_7 \end{aligned}$$



DAN

ex

$$X = \underbrace{(a+a)}_1 \vee \underbrace{(a+a)}_2 + \underbrace{\left(\underbrace{(a+a)}_4 \vee \underbrace{(a-a)}_5 \right)}_3$$

3 address code

$$t_1 = a + a$$

$$t_2 = a + a$$

$$t_3 = t_1 \vee t_2$$

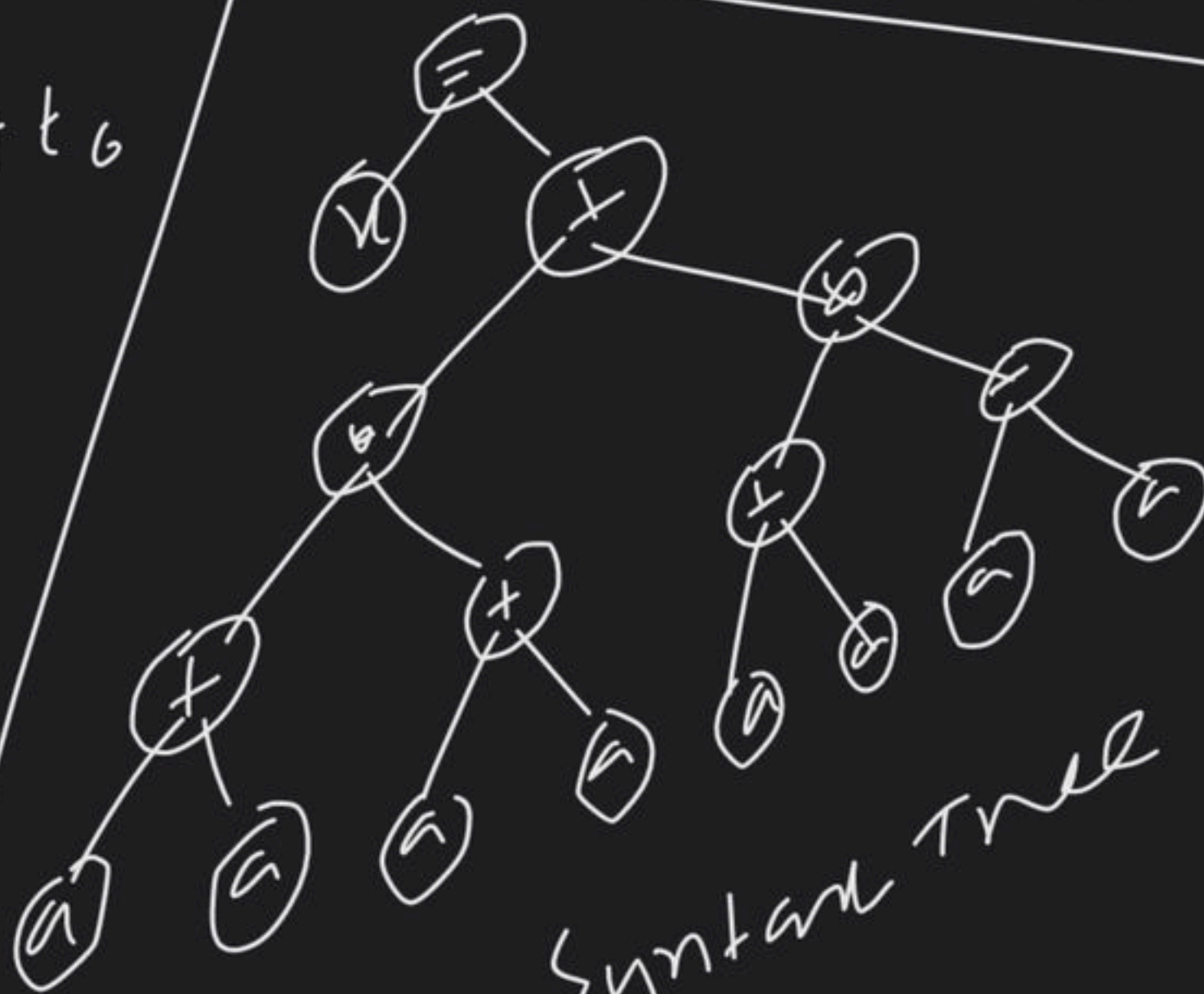
$$t_4 = a + a$$

$$t_5 = a - a$$

$$t_6 = t_4 \vee t_5$$

$$t_7 = t_3 + t_6$$

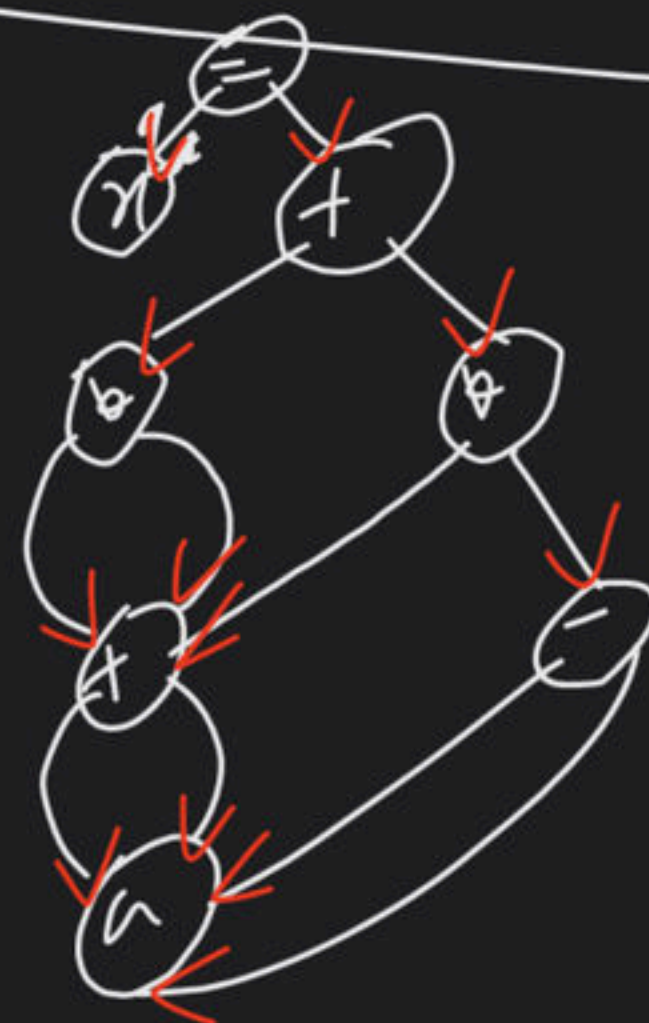
$$X = t_7$$



Syntax Tree

Prefix

$$X \quad \frac{a+a}{1} \quad \frac{a+a}{2} \quad \vee \quad \frac{a+a}{4} \quad \frac{a-a}{5}$$



3-address code (max 3-addresses)

① $x = y + z$

~~⑦ $x = a[i][i]$~~

② $x = y$

~~⑧ $x = f(a, b)$~~

③ $x = a[i]$

⑨ $\text{if}(a < b) \text{ goto } c$

④ $a[i] = x$

⑤ $x = *a$

⑥ $*a = x$

⑩ $\text{goto } c$

$\text{if } ((a == b) \text{ \& } (c == d)) \text{ then } e = f \text{ else } g = h$



1000) $\text{if } (a == b)$ goto ~~1001~~ 1003

1001) $g = h$

1002) goto ~~1001~~ 1006

1003) $\text{if } (c == d)$ goto ~~1004~~ 1005

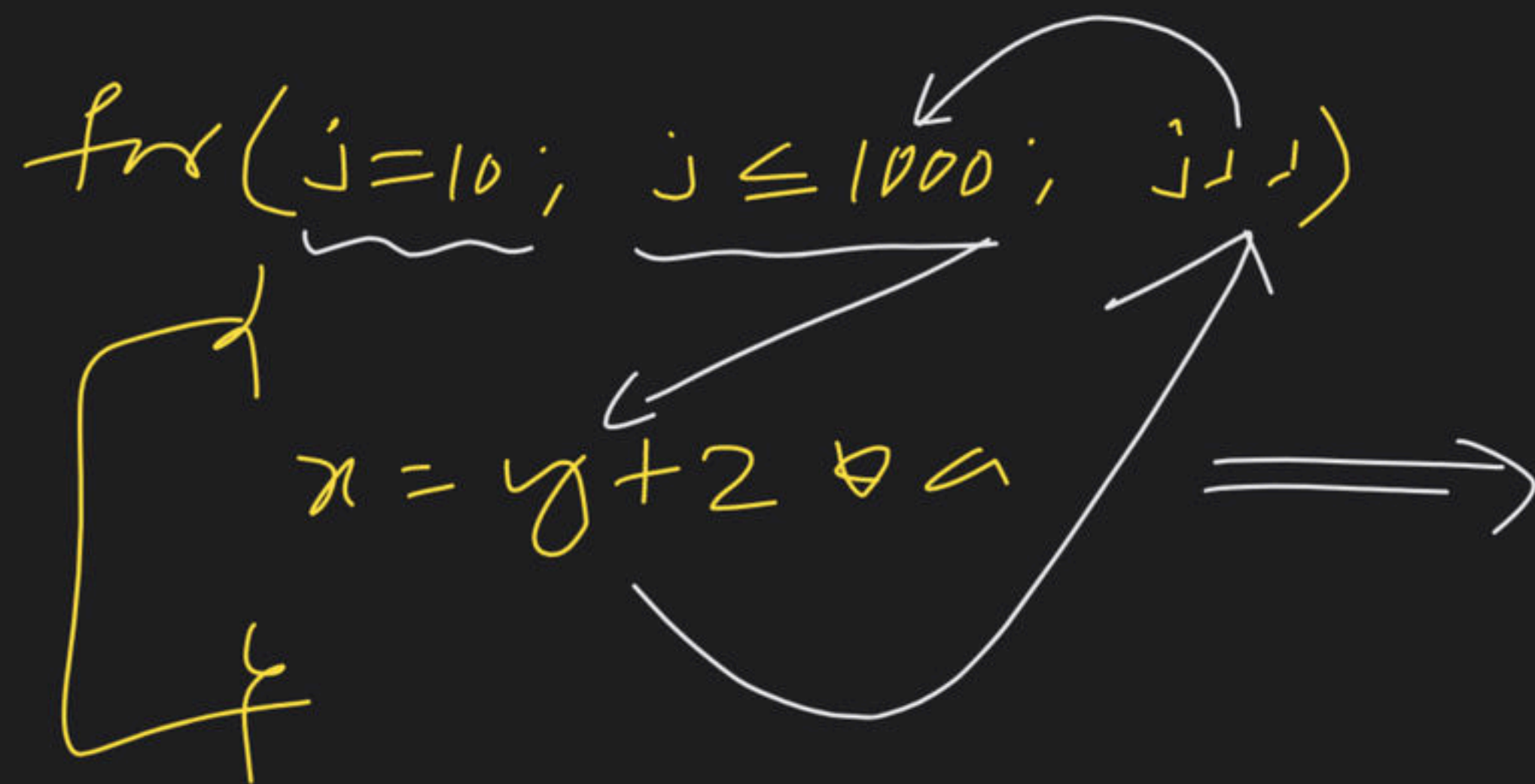
1004) goto ~~1001~~ 1001

1005) $e = f$

1006) _____

Backport

Write 3-address code for the following program //



q500) $j = 10$

q501) ~~$j \leq 1000$~~ goto q503 ^m

q502) ~~goto q508~~

q503) $t_1 = 2ba$

q504) $t_2 = y + t_1$

q505) $x = t_2$

q506) $j = j + 1$

q507) goto q501

ex

write 3-address code for the following program

2001) $t_3 = b_2 \wedge c_2$
 2002) $t_4 = a_2 \wedge t_3$
 2003) $x_2 = t_4$
 2004) goto 2006

$j = 25;$

Switch (j)

2007) $t_1 = b_1 \wedge c_1$

08) $t_2 = a_1 \wedge t_1$

09) $x_1 = t_2$

10) goto 2006

case 3: $x_1 = a_1 \wedge b_1 \wedge c_1$
 break;

case 10: $x_2 = a_2 \wedge b_2 \wedge c_2$
 break;

default: $x_3 = a_3 \wedge b_3 \wedge c_3$

2006

2000) $j = 25$

2001) if (j == 3) goto 2007

2002) if (j == 10) goto 2011

2003) $t_5 = b_3 \wedge c_3$

2004) $t_6 = a_3 \wedge t_5$

2005) $x_3 = t_6$

2006)

Write 3-address code for the following

Since N_5 size $\Rightarrow 27$

Stmt: int a[10][25]; \Rightarrow

$$BA = 1000$$

$$No. of rows = (0-9) \Rightarrow 10$$

$$No. of cols = (0-24) \Rightarrow 25$$

$x = a[i][j]$



$$t_1 = i \times 25$$

$$t_2 = t_1 + j$$

$$t_3 = t_2 \times 2$$

$$t_4 = 1000 + t_3$$

$$x = *t_4$$

$Loc(a[i][j])$

