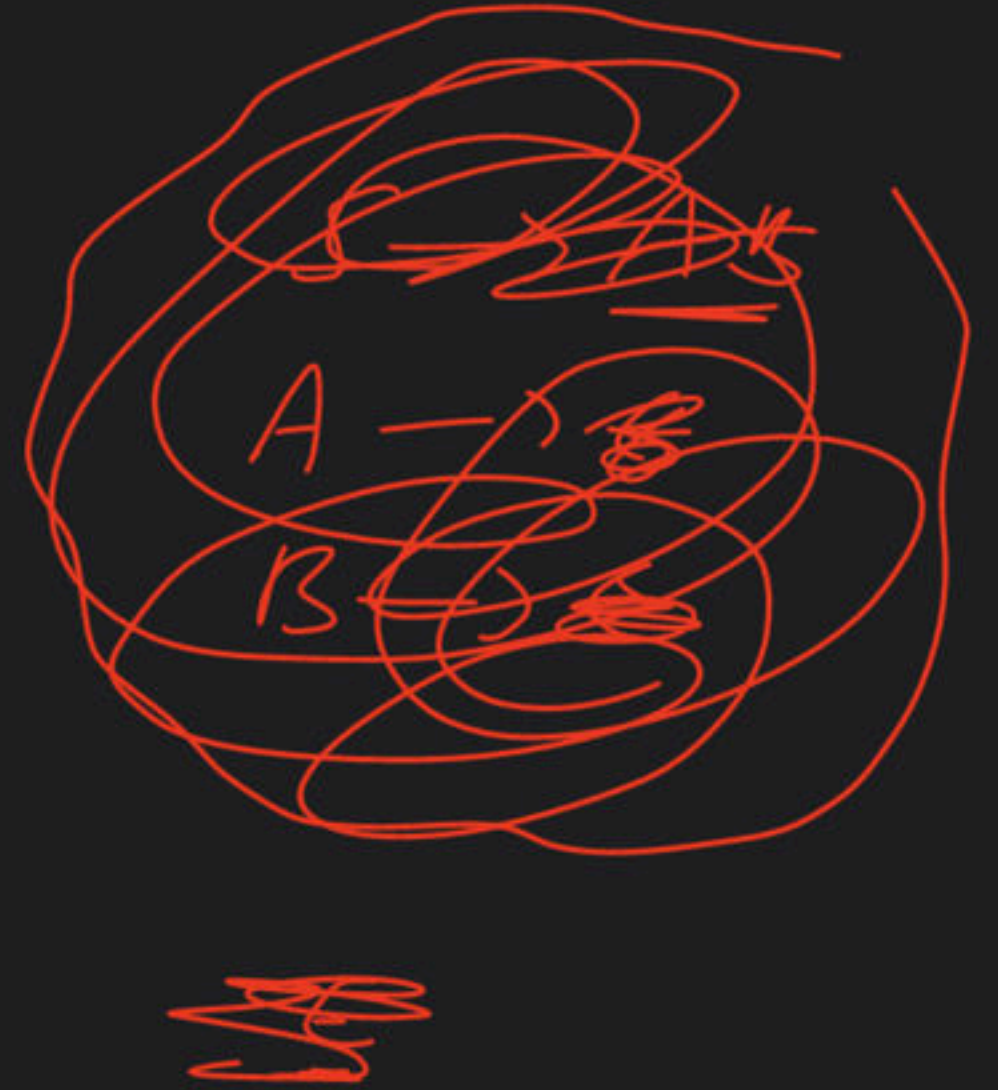
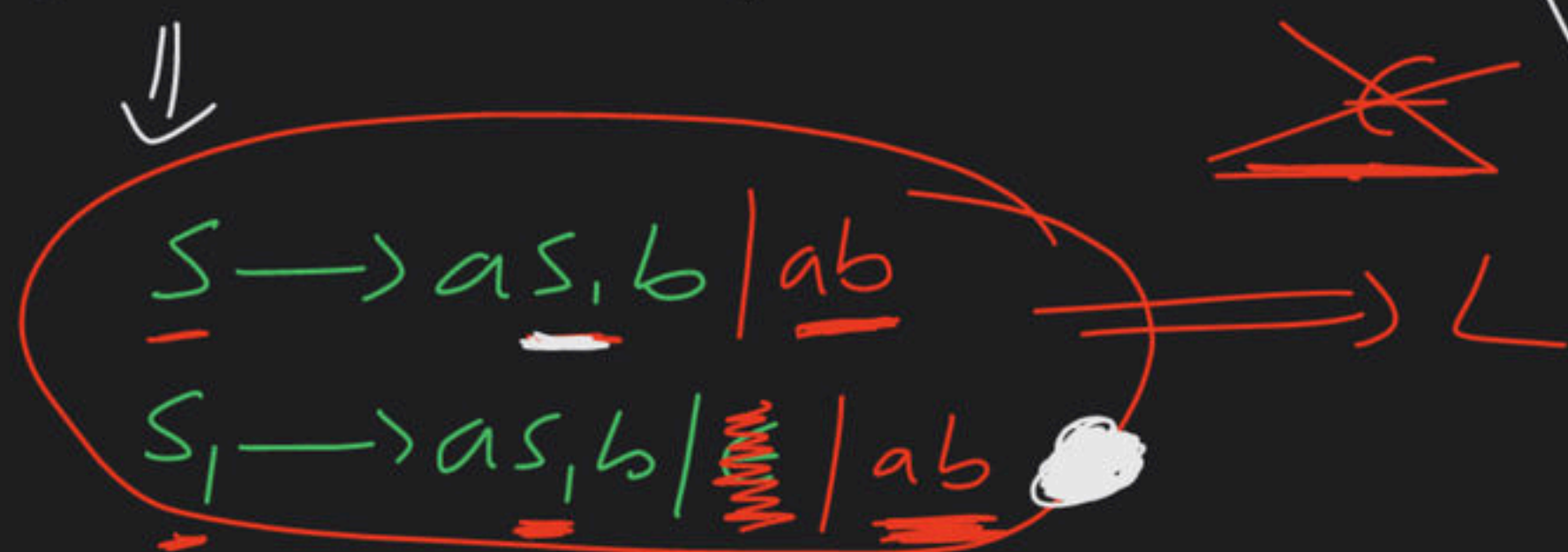
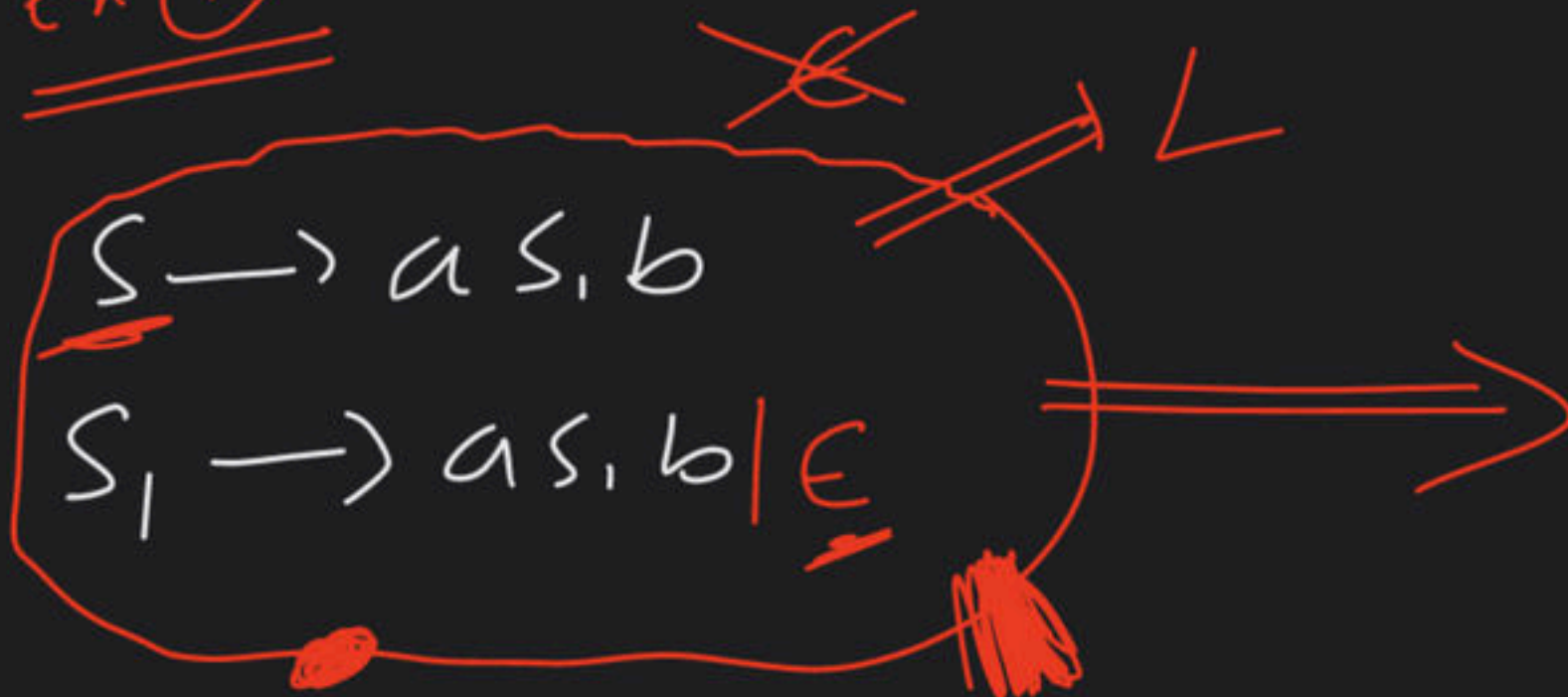


# Simplification of CFG

- ① Elimination of null productions
- ② " " Unit productions
- ③ " " useless symbols



## ex ① Elimination of null production



~~S -> aS, b~~  
~~S1 -> aS, b | epsilon~~  
S -> aS, b | ab



ex(2)

$S \rightarrow as, b \mid \underline{\epsilon}$   
 $S_1 \rightarrow as, b \mid \epsilon$

$\Downarrow L$

$ab \in L$

$aabb \in L$

$\epsilon \in L \checkmark$

$\Rightarrow$

$S \rightarrow \underline{as, b \mid \epsilon} \mid ab$   
 $S_1 \rightarrow \underline{as, b \mid \epsilon} \mid ab$

$\Downarrow L$

$ab \in L$

$aabb \in L$

$\epsilon \notin L$



In the given CFL  $\in \neq L(u)$  then we can get equivalent CFL after eliminating null production.

ex (3)

$S \rightarrow ABac$

$A \rightarrow Bc$

$B \rightarrow b|e$

$C \rightarrow \underline{D}|e$

$D \rightarrow d$

$S \rightarrow \underline{A} \underline{B} \underline{a} \underline{c} | ABa | Aac | Bac | Aa | \underline{a} | Ba | ac$

$A \rightarrow \underline{B} \underline{c} | \cancel{e} | \cancel{e} | \cancel{e} | b | d$

$B \rightarrow \underline{b} | \cancel{e}$

$C \rightarrow \cancel{D} | \cancel{e} | d$

$D \rightarrow \cancel{d}$

(2)

L



# Unit production

$\begin{matrix} \underline{A} \rightarrow \underline{B} \\ S.V & S.V \end{matrix} \left\{ \begin{array}{l} \text{Unit production} \end{array} \right.$

$A \rightarrow a$

S.V

S.T

~~U.P~~

ex ①

$\underline{S} \rightarrow \underline{A}$

$\underline{A} \rightarrow \underline{B}$

$B \rightarrow a$

$\Rightarrow$

$S \rightarrow \underline{A} a$

$\underline{A} \rightarrow \underline{B} a$

$B \rightarrow a$

$\underline{S} \rightarrow \underline{a}$

ex(2)

$S \rightarrow Aa|B$

$B \rightarrow \underline{A}|bb$

$A \rightarrow a|bc|\underline{B}$

E.U.P

$S \rightarrow Aa| \underline{B} | bb|a|bc$

$B \rightarrow \underline{A}| \underline{bb} | a|bc$

$A \rightarrow \underline{a}| \underline{bc} | \underline{B} | \underline{bb}$



# Useless Symbol Elimination

$S \rightarrow ABC \mid A$

$B \rightarrow bB$

$C \rightarrow a$

$D \rightarrow b$

$A \rightarrow a$

$\Rightarrow$

- ① Eliminate those variables & not reachable from start

~~$S \rightarrow ABC \mid A$~~

$A \rightarrow a$ ,  ~~$B \rightarrow bB$~~ ,  $C \rightarrow a$

- ② Eliminate those productions not producing any string

$S \rightarrow A$

$A \rightarrow a$

ex

$S \rightarrow as | AB | \epsilon$

$A \rightarrow bA$

$B \rightarrow AA$



~~①~~

②

~~$A \rightarrow bA$~~

~~$B \rightarrow AA$~~

~~$S \rightarrow as | \epsilon | AB$~~

$S \rightarrow as | \epsilon$



ex (2)

~~$S \rightarrow AB$~~

~~$A \rightarrow a$~~

~~$B \rightarrow b | \epsilon$~~

~~Not~~



~~$S \rightarrow AB | A$~~

~~$A \rightarrow a$~~

~~$B \rightarrow b$~~



$$G \Rightarrow SC$$



$$\cancel{L(G)} = \underline{L(SC)}$$

$$\cancel{\epsilon} \neq$$

For every  $\epsilon$ -free CFG we can write  
equivalent Simplified CFG

---



# Normal Form

1. CNF (Chomsky Normal Form)
2. GNF (Greibach " )



$$V \rightarrow VV / T$$

(or)

$$A \rightarrow \underline{\underline{BC}}$$

$$B \rightarrow \underline{b}$$

$$C \rightarrow \underline{c}$$

$$V \rightarrow \underline{a} \underline{V}^{\phi}$$

(or)

$$A \rightarrow b C D E$$

$$C \rightarrow a D \rightarrow e$$

$$E \rightarrow \epsilon //$$

CNF ( $V \rightarrow VV | T$ )

ex (1)

$S \rightarrow AA | a$

$A \rightarrow SA | b$

CNF ✓

ex (2)

$S \rightarrow AS | \overset{C}{\textcircled{AAS}} | \underline{a}$

$A \rightarrow SA | \underline{aa}$

~~CNF~~

① Simplified con  
no-chngl

② convert CNF

$S \rightarrow AS | \overset{\downarrow}{C}S | a$

$C \rightarrow AA$

$A \rightarrow SA | DD$

$D \rightarrow a$  ✓



$S \rightarrow A B a$

$A \rightarrow a b$

$B \rightarrow A c$

$\Rightarrow$  (1) Simply  
no-change

(2) Convert

$S \rightarrow A B a$

$S \rightarrow D E$

$E \rightarrow \underline{a}, D \rightarrow A B$

$A \rightarrow E F$

$F \rightarrow E G$

$G \rightarrow b$

$A \rightarrow a b$

$B \rightarrow A c$

$B \rightarrow A H$   
 $\underline{H} \rightarrow c$

Note

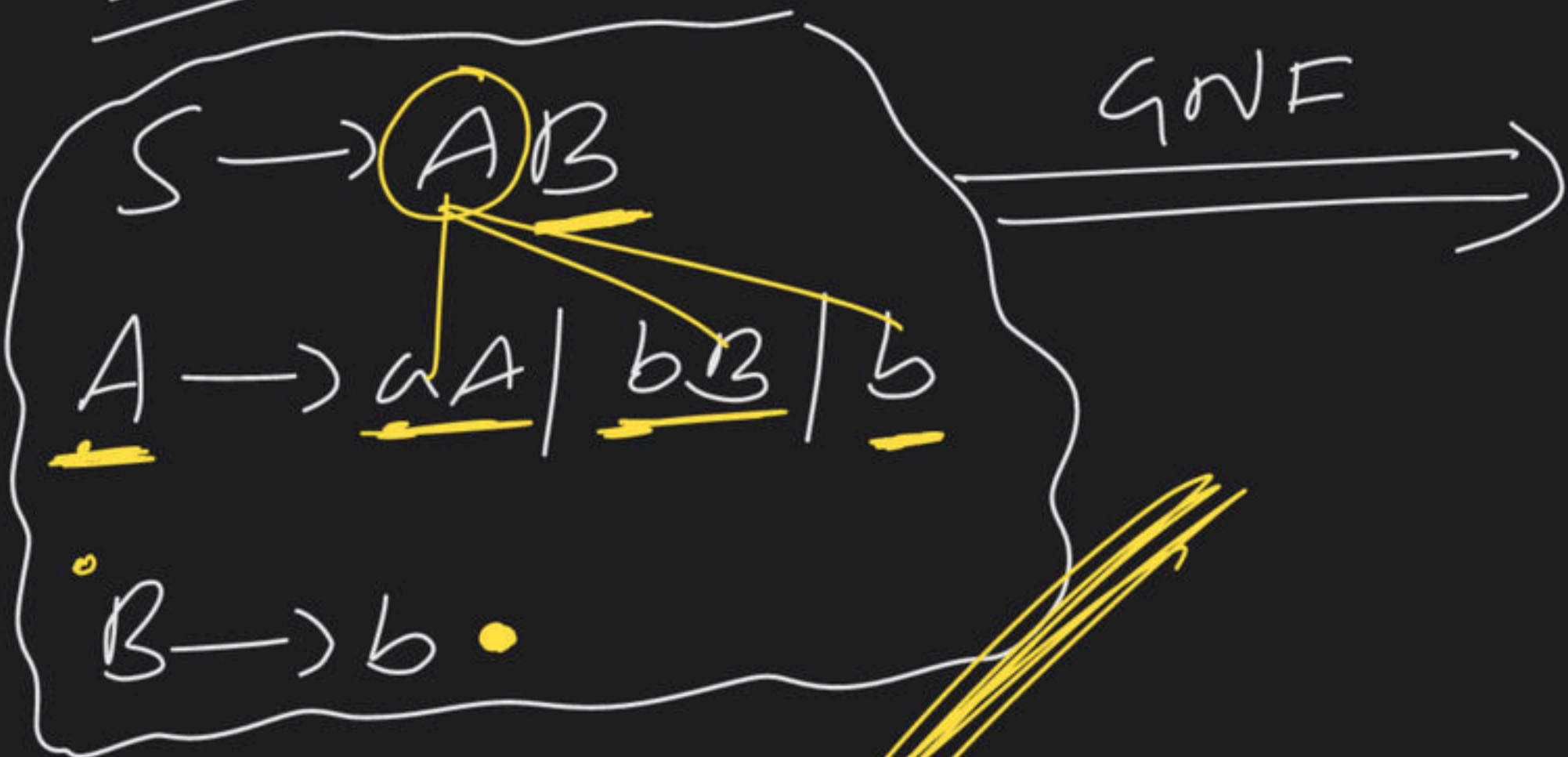
(1) From the given CNF grammar to generate  $n$ -length string we need  $2n-1$  productions

(2) For every E-free CFG we can construct equivalent CNF.



GNF ( $V \rightarrow TV^*$ )

ex



~~CNF~~  
~~GNF~~

① Simplifying  
no-simplification.

②

convert into  
GNF

$S \rightarrow aAB \mid bBB \mid bB$   
 $A \rightarrow aA \mid bB \mid b$   
 $B \rightarrow b$

ex

$S \rightarrow asb/ab$

~~GNF~~

① Simplify  
no-sim

② Conversion

$S \rightarrow \underline{a} s \underline{B} / \underline{a} \underline{B}$

$B \rightarrow \underline{b}$

Note

To generate  $n$ -length string from GNF  
we need  $n$ -productions.



Thank8 All

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D. H

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$S \rightarrow \underline{a} B c D E$

$B \rightarrow \underline{b}$

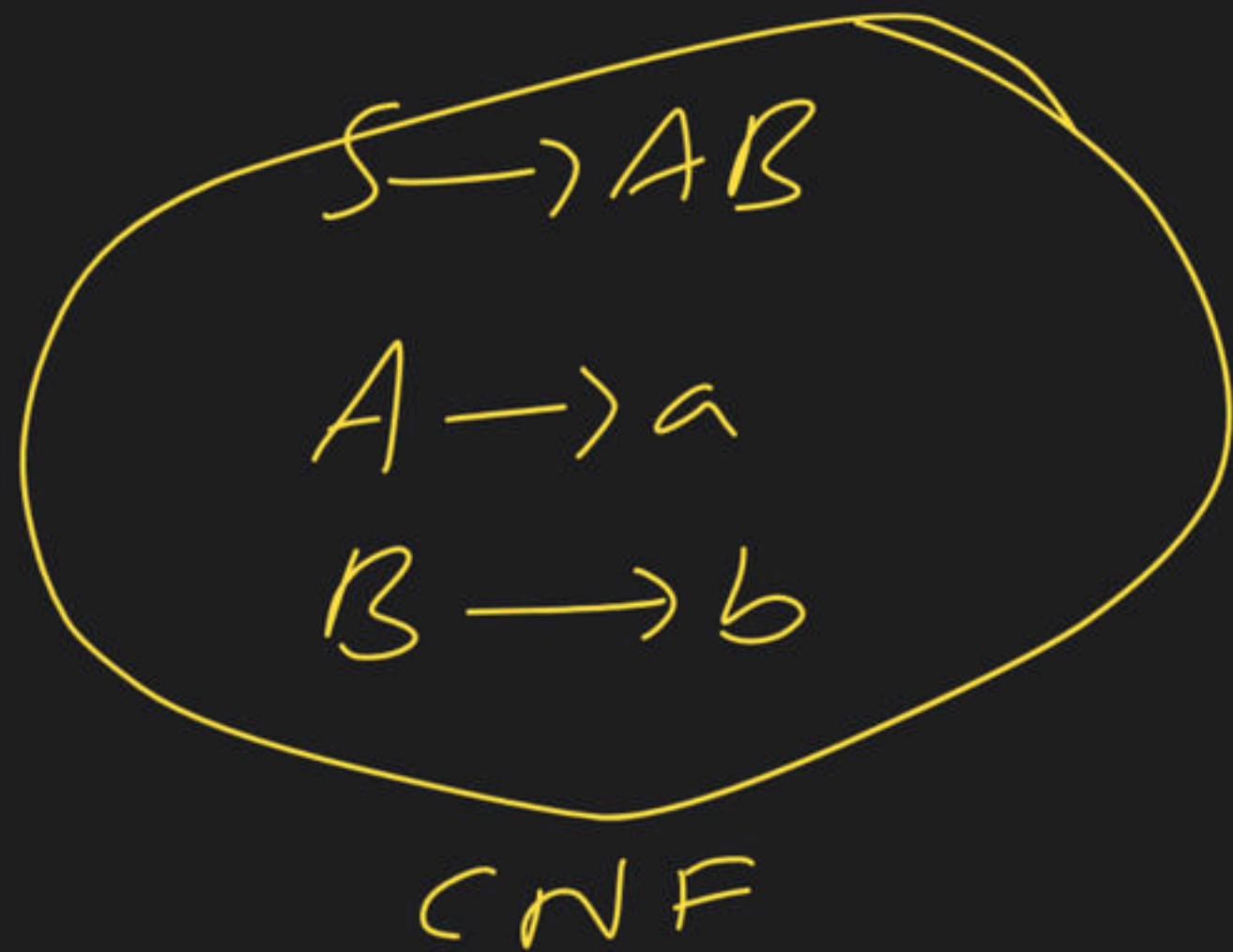
$C \rightarrow \underline{c}$

$D \rightarrow \underline{d}$

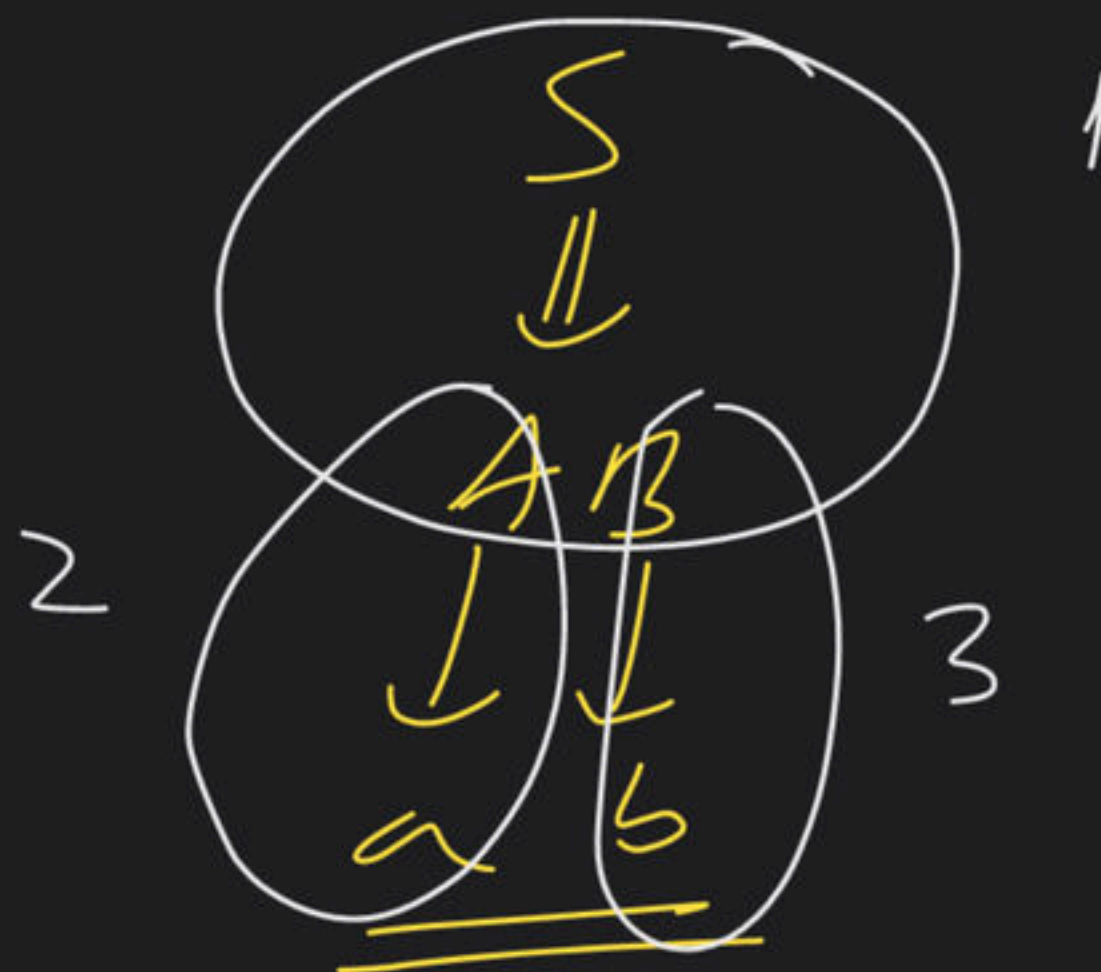
$E \rightarrow \underline{e}$







i/p: ab  
2



$$2n - 1$$

$$2 \times 2 - 1 \Rightarrow 3$$

$S \rightarrow ab$   


---

 $S \rightarrow A$   
 $A \rightarrow B$   
 $B \rightarrow C$   
 $C \rightarrow D$   
 $D \rightarrow ab$   


---



CNF

$S \rightarrow AB$

$A \rightarrow CD$

$B \rightarrow EF$

$C \rightarrow a$

$D \rightarrow b$

$E \rightarrow c$

$F \rightarrow d$

abcd (204-1)