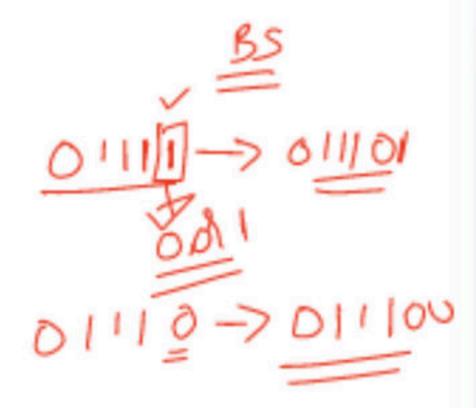
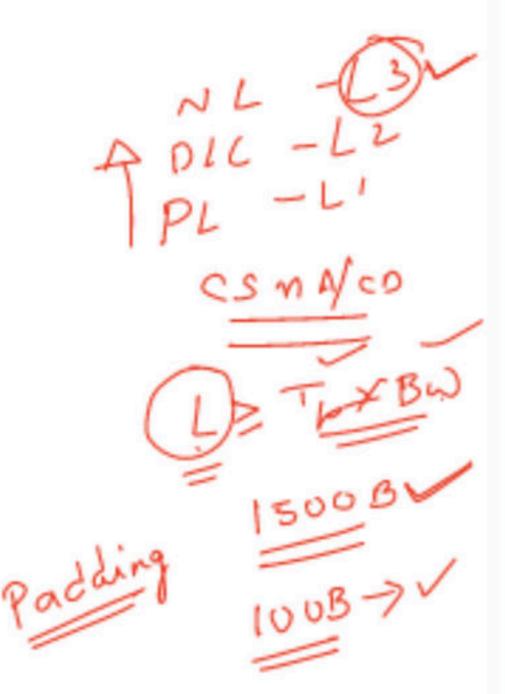


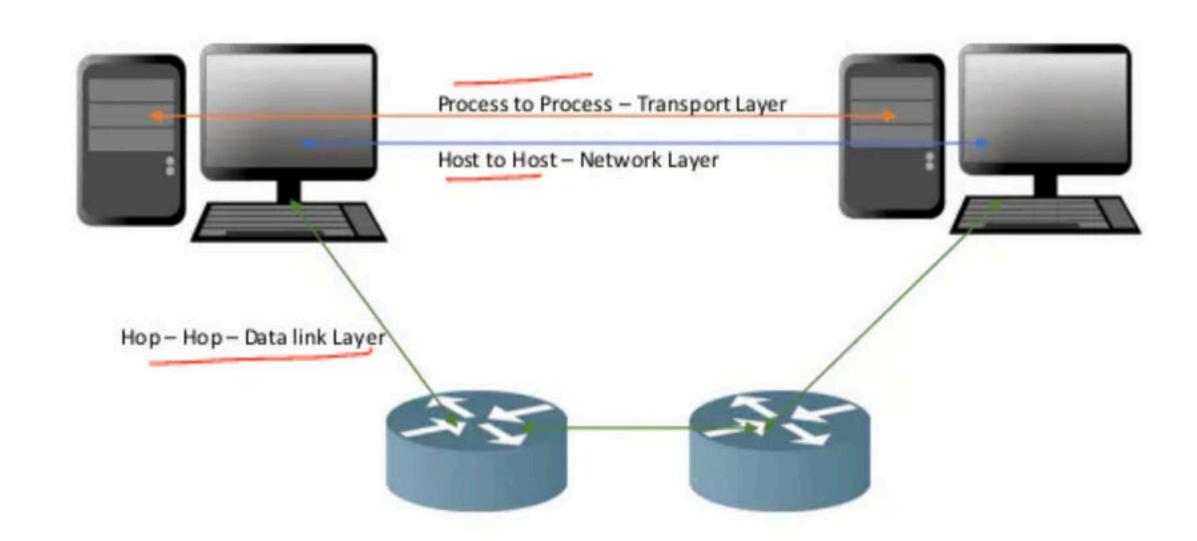
Complete Course on Computer Networks - Part II

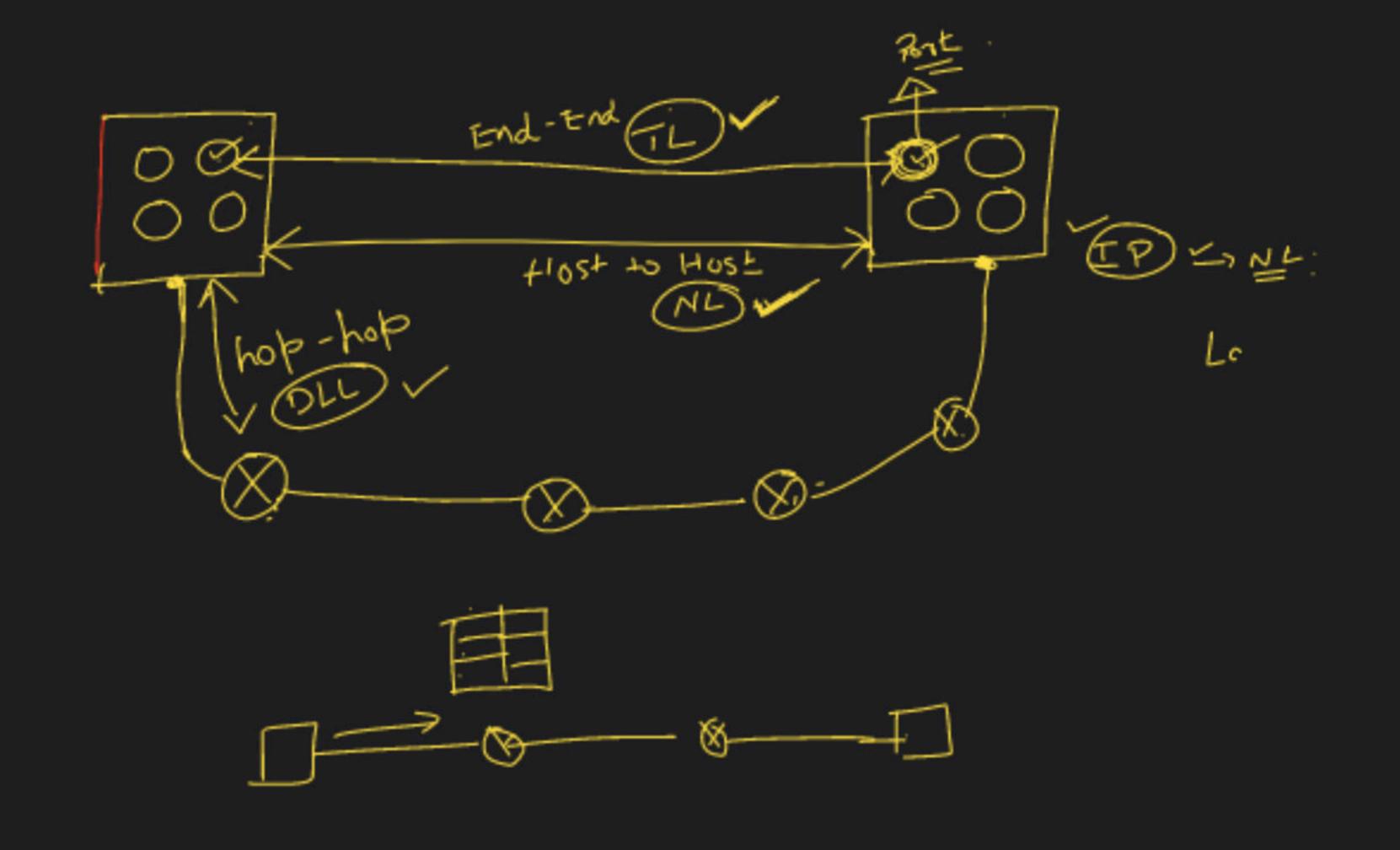


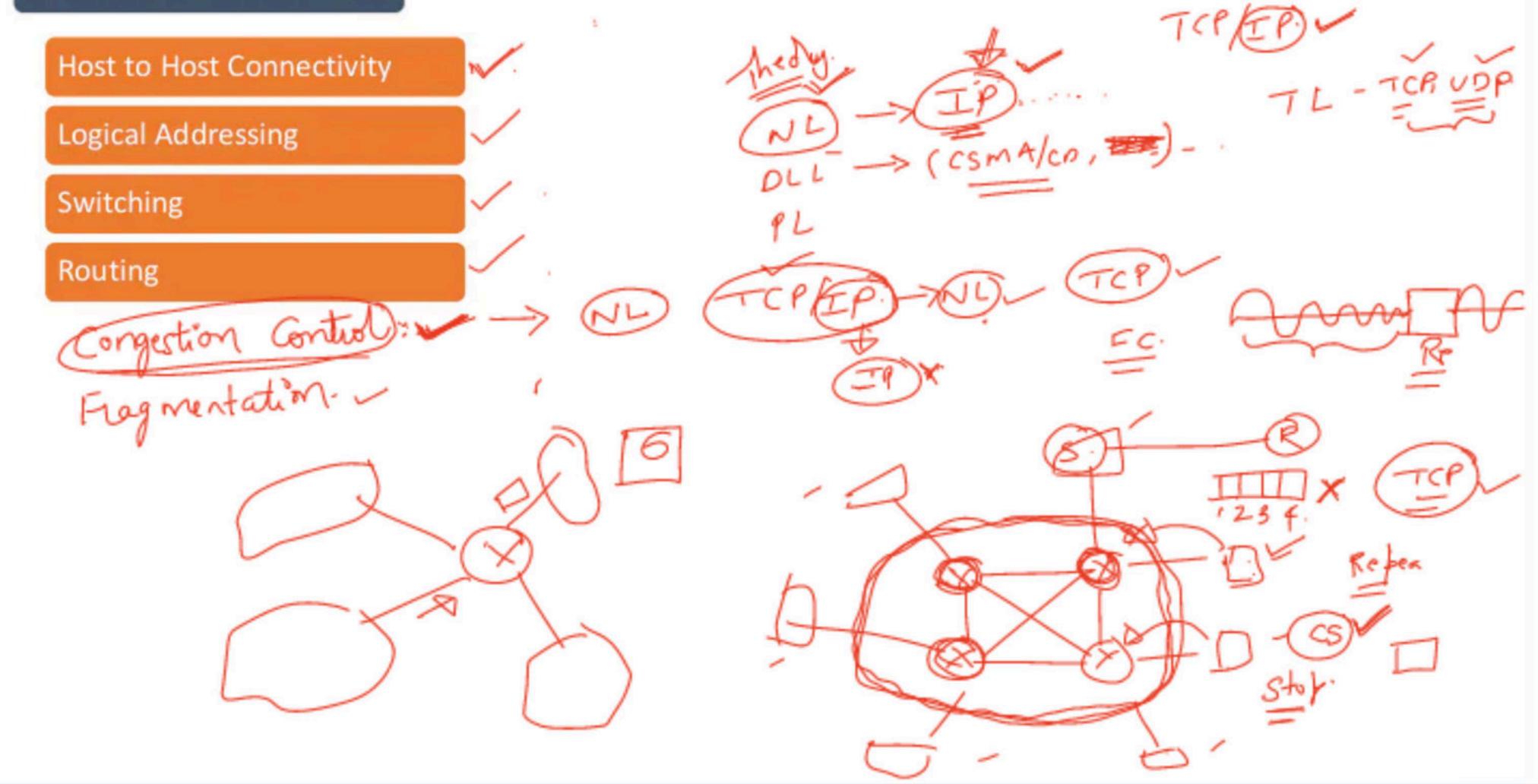
Computer Network

Network Layer 🗸









Host to Host Connectivity

Logical Addressing

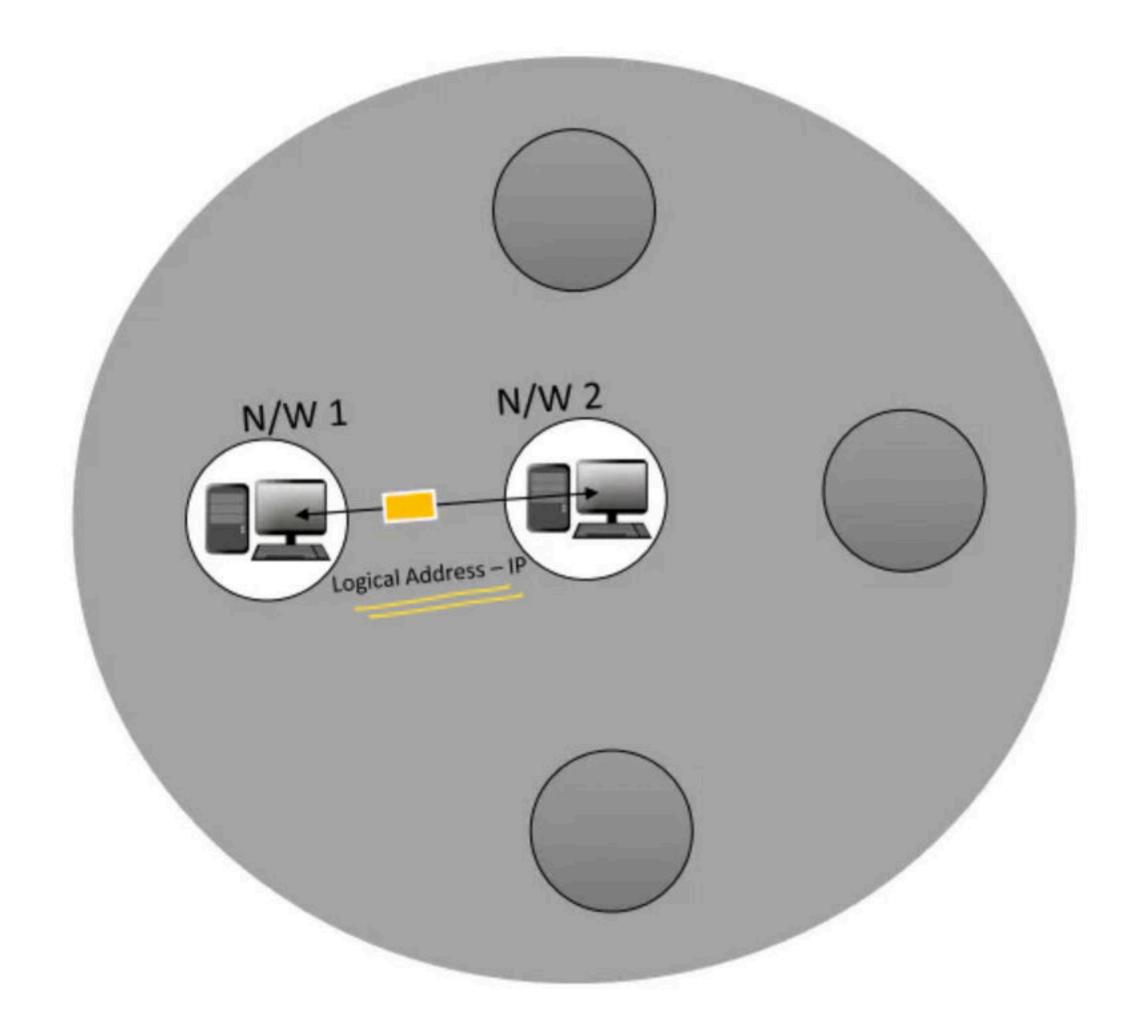
Switching



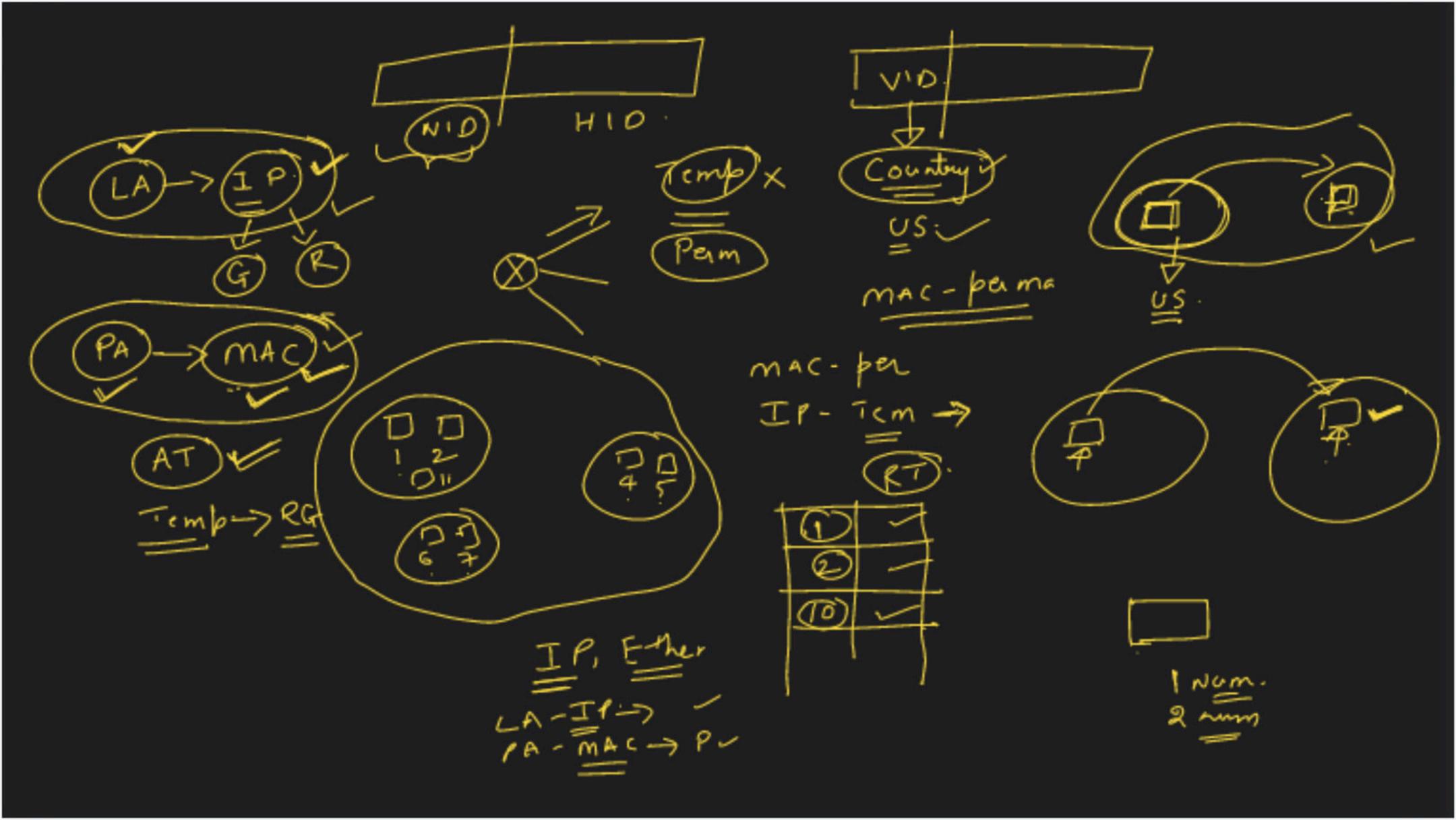
Host to Host Connectivity

Logical Addressing

Switching



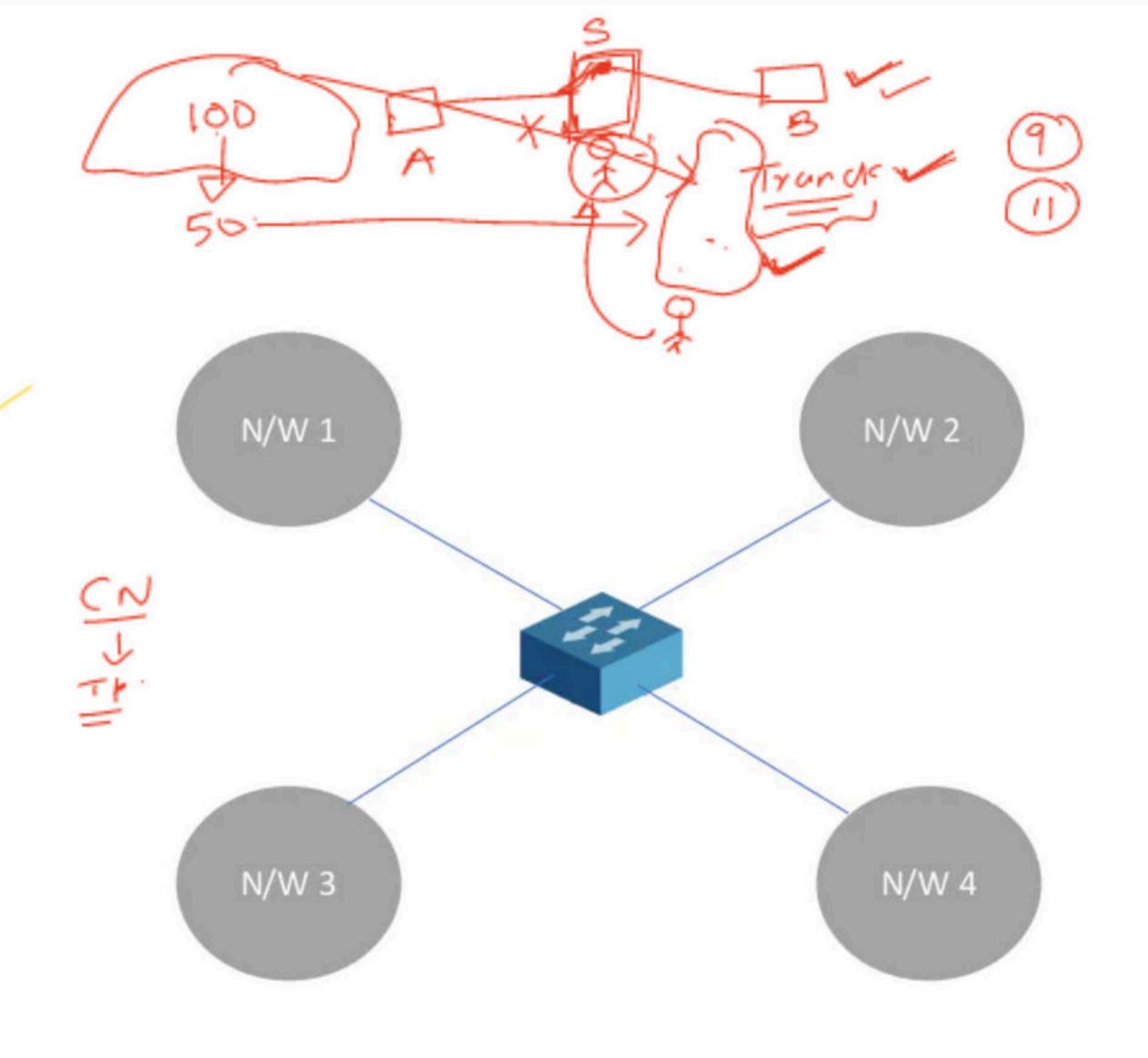




Host to Host Connectivity

Logical Addressing

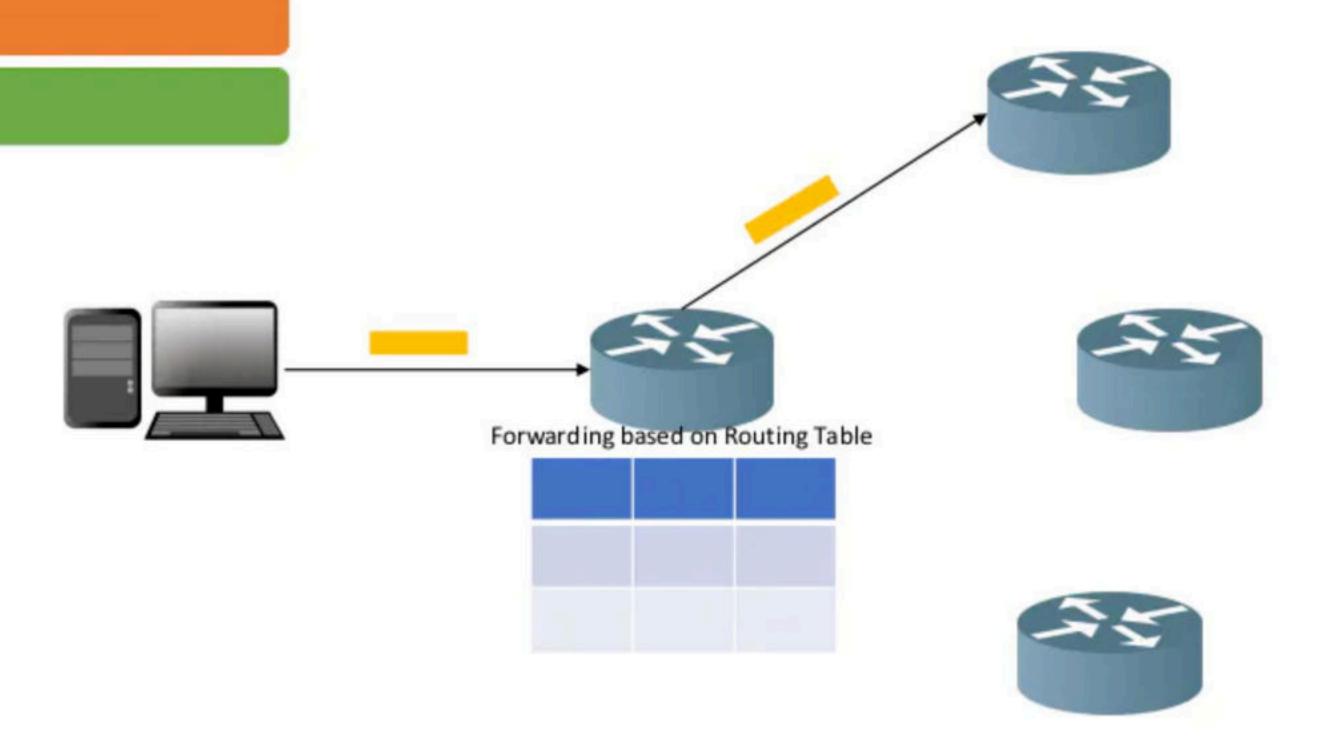
Switching



Host to Host Connectivity

Logical Addressing

Switching



Computer Networks

GATE QUESTIONS

1.) An organization requires a range of IP addresses to assign one to each of its 1500 computers. The organization has approached an Internet Service Provider (ISP) for this task. The ISP uses CIDR and serves the requests from the available IP address space 202.61.0.0/17. The ISP wants to assign an address space to the organization which will minimize the number of routing entries in the ISP's router using route aggregation. Which of the following address spaces are potential candidates from which the ISP can allot any one to the organization? [GATE 2020]

I. 202.61.84.0/21

II. 202.61.104.0/21

III. 202.61.64.0/21

IV. 202.61.144.0/21

1.) An organization requires a range of IP addresses to assign one to each of its 1500 computers. The organization has approached an Internet Service Provider (ISP) for this task. The ISP uses CIDR and serves the requests from the available IP address space 202.61.0.0/17. The ISP wants to assign an address space to the organization which will minimize the number of routing entries in the ISP's router using route aggregation. Which of the following address spaces are potential candidates from which the ISP can allot any one to the organization? (GATE 2020)

1. 202.61.84.0/21

II. 202.61.104.0/21

III. 202.61.64.0/21

IV. 202.61.144.0/21

Solution: Option C

Given CIDR IP is 202.61.0.0/17 and for HID 32 - 17 = 15 bits can be used.

And to Assign an IP address for 1500 computer, we require 11 bit from HID part.

So NID + SID = 17 + 4 = 21 bits and HID = 11 bits

NID HID

202.61.0 0000 000.00000000

So, from the given option, possible IP Address is

I. 84 - >0 1010 100 (Because in HID bit 1 is not possible)

II. 104 -> 0 1101 000

III. 64 -> 0 1000 000

IV. 144->1 0010 000 (Because in NID bit 1 is not possible)

2.) A computer network uses polynomials over GF(2) for error checking with 8 bits as information bits and uses x3 + x + 1 as the generator polynomial to generate the check bits. In this network, the message 01011011 is transmitted as [GATE 2017]

Α	01011011010
В	01011011011
С	01011011101
D	01011011100

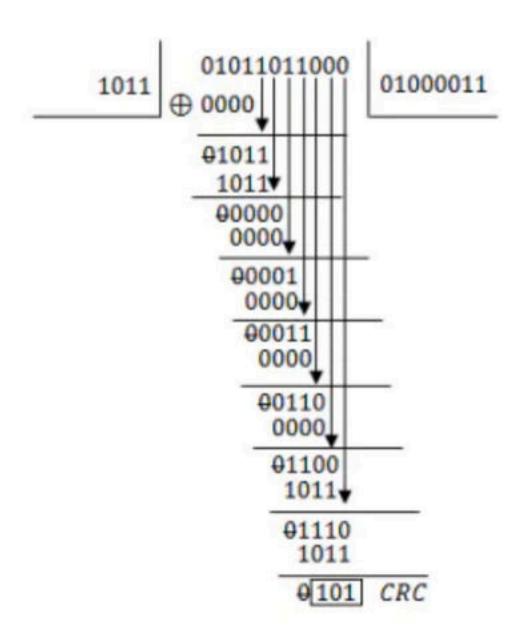
Solution:

Given CRC generator polynomial = x3+x+1

 $= 1 \cdot x3 + 0 \cdot x2 + 1 \cdot x1 + 1 \cdot x0$

=1011

Message =01011011



So, the message 01011011 is transmitted as 0101 1011 101

3.) The value of parameters for the Stop-and-Wait ARQ protocol are as given below:

Bit rate of the transmission channel = 1 Mbps.

Propagation delay from sender to receiver = 0.75 ms.

Time to process a frame = 0.25 ms.

Number of bytes in the information frame = 1980.

Number of bytes in the acknowledge frame = 20.

Number of overhead bytes in the information frame = 20.

Assume that there are no transmission errors. Then, the transmission efficiency (expressed in percentage) of the Stop-and-Wait ARQ protocol for the above parameters is ______ (correct to 2 decimal places). [GATE 2017]

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Number of overhead bytes in the information frame = 20.

Assume that there are no transmission errors. Then, the transmission efficiency (expressed in percentage) of the Stop-and-Wait ARQ protocol for the above parameters is ______ (correct to 2 decimal places). [GATE 2017]

Solution:

Given Data:

B = 1Mbps, L = 1980Bytes, Overhead = 20Bytes

TProc = 0.25ms, LAck = 20Bytes

Tp=0.75ms

Total Data size(L) = (L + overhead) = 1980+20 = 2000Bytes

Efficiency of Stop & Wait ARQ?

Tt = L/B = 2000Bytes/1Mbps = (2000×8bits)/(106 b/s) = 16msec

 $TAck = LAck/B = (20 \times 8bits)/(106 bits/sec) = 0.16msec$

: In Stop and Wait ARQ, efficiency

 $\eta = Tt/(Tt+TAck+2Tp+TProc) = 16ms/(16+0.16+2×0.75+0.25ms)$

= 16ms/17.91ms = 0.8933 ≈ 89.33%

4.)A sender uses the Stop-and-Wait ARQ protocol for reliable transmission of frames. Frames are of size 1000 bytes
and the transmission rate at the sender is 80 Kbps (1Kbps = 1000 bits/second). Size of an acknowledgement is 100
bytes and the transmission rate at the receiver is 8 Kbps. The one-way propagation delay is 100 milliseconds.

Assuming no frame is lost, the sender throughput is ______ bytes/second. [GATE 2016]

4.)A sender uses the Stop-and-Wait ARQ protocol for reliable transmission of frames. Frames are of size 1000 bytes and the transmission rate at the sender is 80 Kbps (1Kbps = 1000 bits/second). Size of an acknowledgement is 100 bytes and the transmission rate at the receiver is 8 Kbps. The one-way propagation delay is 100 milliseconds.

Assuming no frame is lost, the sender throughput is ______ bytes/second. [GATE 2016]

Solution:

Given,

Frame size (L) =1000 bytes

Sender side bandwidth (BS) = 80 kbps = 10 * 103 bytes/sec

Acknowledgement size (LA) =100 bytes

Receiver side bandwidth (BR) = 8 kbps = 1 * 103 bytes/sec

Propagation delay (Tp) =100 ms

By formula:

Transmission delay (Tt) = L/BS = 1000 bytes / 10 * 103 bytes/sec = 100 ms

Acknowledge delay (Tack) = LA / BR = 100 bytes / 1 * 103 bytes/sec = 100 ms

Total cycle time = Tt + 2 * Tp + Tack = 100 ms + 2 * 100 ms + 100 ms = 400 ms

Efficiency (η) = Tt / Total cycle time = 100 ms / 400 ms = 1 / 4 = 0.25

Throughput = Efficiency (η) * Bandwidth (BS) = 0.25 * 10 *103 bytes/s = 2500

bytes/second

5.) A network has a data transmission bandwidth of 20 × 106 bits per second. It uses CSMA/CD in the MAC layer. The maximum signal propagation time from one node to another node is 40 microseconds. The minimum size of a frame in the network is ______ bytes. [GATE 2016]

5.) A network has a data transmission bandwidth of 20 × 106 bits per second. It uses CSMA/CD in the MAC layer. The maximum signal propagation time from one node to another node is 40 microseconds. The minimum size of a frame in the network is ______ bytes. [GATE 2016]

Solution:

For frame size to be minimum, its transmission time should be equal to twice of one way propagation delay. i.e, $Tt = 2 \times TP$ Given,

Bandwidth (B) = 20×106 bps

TP = $40 \mu s \Rightarrow 40 \times 10-6 sec$

Suppose minimum frame size is L.

 $Tt = 2 \times TP \Rightarrow L/B = 2 \times TP$

 \Rightarrow L = 2 × TP × B = 2 × 40 × 10-6 × 20 × 106 = 1600 bits \Rightarrow 200

bytes

Therefore, L = 200 bytes

6.) Consider a 128 × 103 bits/ second satellite communication link with one way propagation delay of 150 milliseconds. Selective retransmission (repeat) protocol is used on this link to send data with a frame size of 1 kilobyte. Neglect the transmission time of acknowledgement. The minimum number of bits required for the sequence number field to achieve 100% utilization is _____ [GATE 2016]

6.) Consider a 128 × 103 bits/ second satellite communication link with one way propagation delay of 150 milliseconds. Selective retransmission (repeat) protocol is used on this link to send data with a frame size of 1 kilobyte. Neglect the transmission time of acknowledgement. The minimum number of bits required for the sequence number field to achieve 100% utilization is _____ [GATE 2016]

Solution:

```
To achieve 100% efficiency, the number of frames that we should send N = 1 + 2 * a a = Tp / Tt where Tp is propagation delay, and Tt is transmission delay. Given, B = 128 kbps, Tp = 150 msec, L = 1 KB = 1 * 8 * 210 bits

Tt = L / B \Rightarrow 1 * 8 * 210 bits / 128 * 103 bps \Rightarrow 0.064 sec = 64 msec

So, a = 150 msec / 64 msec = 2.343

Efficiency (\eta) = 100 % \Rightarrow 1 = N/ 1 + 2 * a

So, N = 1 + 2 * a \Rightarrow 1 + 2 * 2.343 = 5.686

No. of sequence numbers requires in SR is 2*N = 2 *5.686 = 11.375

Minimum No. of bits required in the sequence number = [log2 (11.375)] = 4
```