



# DAC - Part III

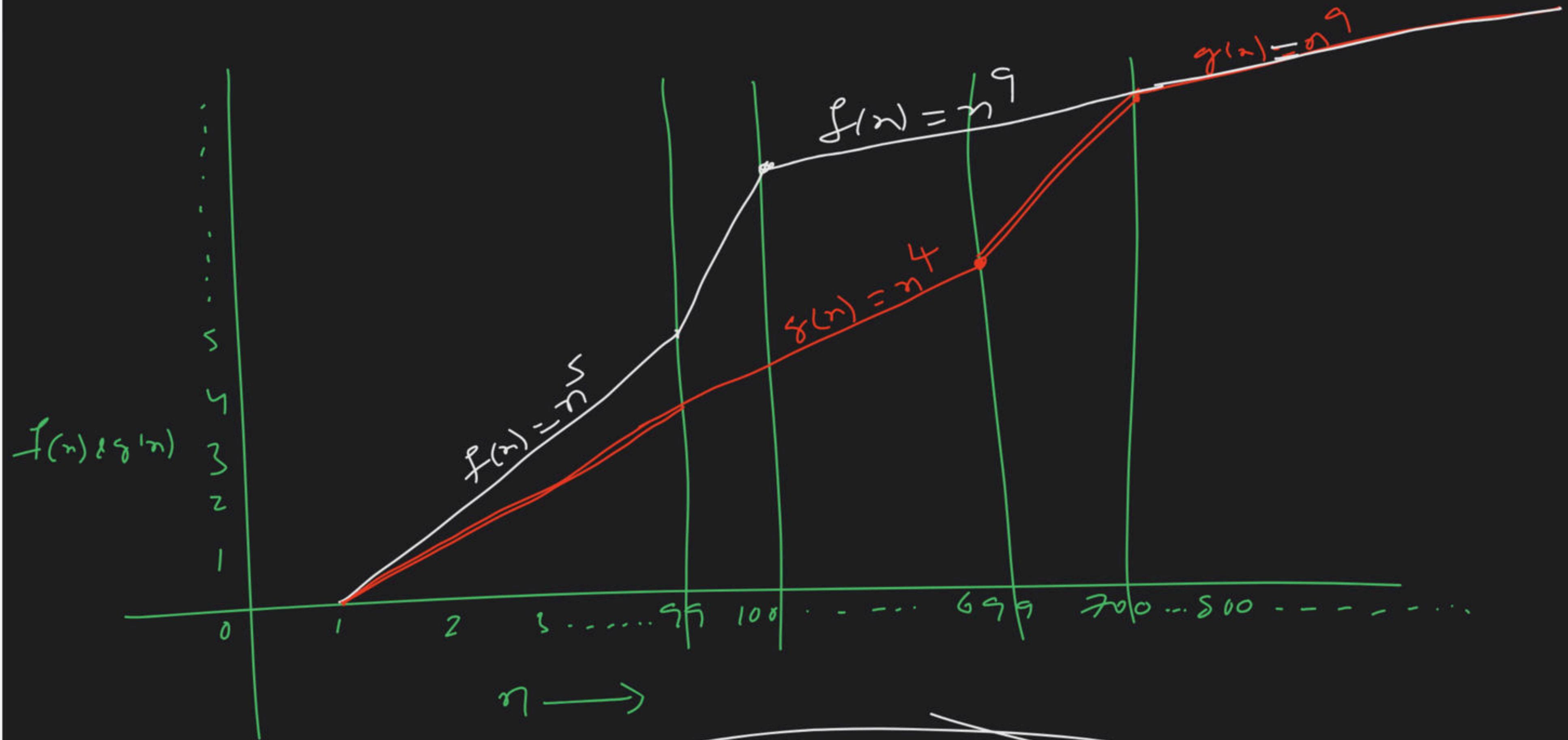
Complete Course on Algorithm for GATE - CS & IT

consider the following function

$$f(n) = \begin{cases} n^5 & \text{if } 0 < n < 100 \\ n^9 & \text{if } n \geq 100 \end{cases}$$

$$g(n) = \begin{cases} n^4 & \text{if } 0 < n < 700 \\ n^9 & \text{if } n \geq 700 \end{cases}$$

relate b/w  $f(n)$  &  $g(n)$ ?



$$f(n) = \Theta(g(n)) \quad n_0 = 70 \quad \Leftrightarrow \quad \left[ \begin{array}{l} f(n) = O(g(n)) \\ n_0 = 70 \end{array} \right] \quad \left| \quad \left[ \begin{array}{l} f(n) = \Omega(g(n)) \\ n_0 = 1 \end{array} \right] \right.$$



Which will satisfy reflexive property?

③

$O, \Omega, \Theta, o, \omega$   
 $\checkmark \quad \checkmark \quad \checkmark \quad \times \quad \times$

$$f(n) = O(f(n))$$

Which will satisfy symmetric property?

④

~~$O, \Omega, \Theta, o, \omega$~~   
 $\checkmark \quad \times \quad \times \quad \times \quad \times$

$$n^{20} = \Theta(n^5)$$

~~$$n^5 = \Theta(n^{20})$$~~

$$\text{If } f(n) = \Theta(g(n))$$

then

$$g(n) = \Theta(f(n))$$

~~$$n^5 = \Theta(n^5)$$~~

$$n^5 = \Theta(n^5)$$



# Transition proofs

$$\text{if } f(n) = O(g(n)) \text{ and } g(n) = O(h(n))$$

then

$$f(n) = O(h(n))$$

5

$$\Rightarrow O, \Omega, \Theta, \text{ or } \omega$$

~~$n^2 = n^2 = n^2$~~

$n^2 = n^3$

$n^3 < n^5 < n^2$

$n^2 < n^5 < n^9 \Rightarrow n^2 < n^9$

Q5

$$\overset{n^5}{f(n)} = \overset{n^7}{O(d(n))}$$

or

$$g(n) = O(e(n))$$

$n^9$

then

$n^{12}$

?

Note

①

$$f(n) + g(n)$$

~~$$= O(d(n))$$~~

~~$$O(e(n))$$~~



$$= O(d(n) + e(n))$$



$$O[\overset{(n)}{\max(d(n), e(n))}]$$





$$(2) \quad f(n) \cdot g(n) = O(d(n) \cdot e(n))$$

---

Note

$$\text{if } f(n) = O(g(n))$$

then

$$f(n) \cdot h(n) = O(g(n) \cdot h(n))$$







9/6  $f(n) = \cancel{O(g(n))}$  and  $g(n) \neq O(f(n))$   $\rightarrow f(n) < g(n)$

$g(n) = O(h(n))$  and  $h(n) = O(g(n))$

then

T/F

$g(n) = h(n)$

$f(n) < \boxed{g(n) = h(n)}$

(a)  $f(n) = O(h(n))$

$n^5 = O(n^9)$

(b)  $f(n) + g(n) = O(h(n))$

$n^7 = O(n^5)$

(c)  $f(n) \cdot g(n) = \Theta(g(n) \cdot h(n))$

$\boxed{n^5 = O(n^5)}$   $n^5 = O(n^5)$   
 $n^5 \neq O(n^5)$   $n^5 = O(n^5)$

(d)  $g(n) = \Theta(h(n))$

$2 \cdot n^7 = 2^7 \cdot n^7$



# Recursion

▲ 1 • Asked by Sai Teja

Please help me with this doubt

17 December 2022 at 11:31 PM

According Mathematical we can say about increment or decrement function but here we Min case is  $O(1)$  how can loop have  $1/n$  iteration.

- T/F
- $O, \omega, \theta, \cancel{<}, \cancel{>}$
- (a) ✓  $f(n) = O(f(n))$
- (b) ✗  $f(n) = O\left(\frac{1}{n}\right)^2$  [
- (c)  $f(n) = O\left(\frac{f(n)}{2}\right)$
- (d)  $f(n) = O\left(f\left(\frac{n}{2}\right)\right)$

✓