



# Parsing - Part I

Complete Course on Compiler Design

et

int main()

char a = 'b'; 10(5)

char c = "string"; 16(6)

float d = 20.5; 21(5)

char arr a = 'b'; 27(6)

char arr c = "string"; 34(7)

float d = 20.5; 40(6)

char\* com arr a = 'b'; 46(6)

char\* com arr c = "string"; 53(7)

49

51

52

54

~~ccc~~

55

void main()

int a = 10; (12)

while(1)

printf("%s", "a"); (24)

if(a)

a = a + 1;

}

36



$$\} \checkmark \implies \leq$$

`printf("%d", g, h, i);`  $\Rightarrow 16$

$c = d + e - \cancel{f} + g + h + k \text{ lön?}$

$$\underline{b = d + e;}$$

$c = g^h i;$

ex

$$\begin{array}{cccccccccccccccccccc} + & - & + & - & = & = & = & + & + & + & + & = & = & + & + & + & - & - & = & = & = \end{array}$$

~~Q~~  
~~Q~~  
13

Which one of the following is said to be  
Token without reading next i/p symbol.

(a) int / int

(b) float / float

(c) return / return

(d) ~~int~~ ; (e) , (f) (



Grammar

$G(V, T, P, S)$

terminal  $\{a, b\}$

start symbol

production  
variable  $\{S, A, B\}$

~~$S \rightarrow AB$~~

~~$A \rightarrow a$~~

~~$B \rightarrow b$~~

$T = \{i, j, g, h\}$   
 $V = \{S, A, B, C, D, E, F\}$

~~$S \rightarrow ABCD|i$~~

~~$A \rightarrow BCD|i$~~

~~$C \rightarrow DEF|h$~~

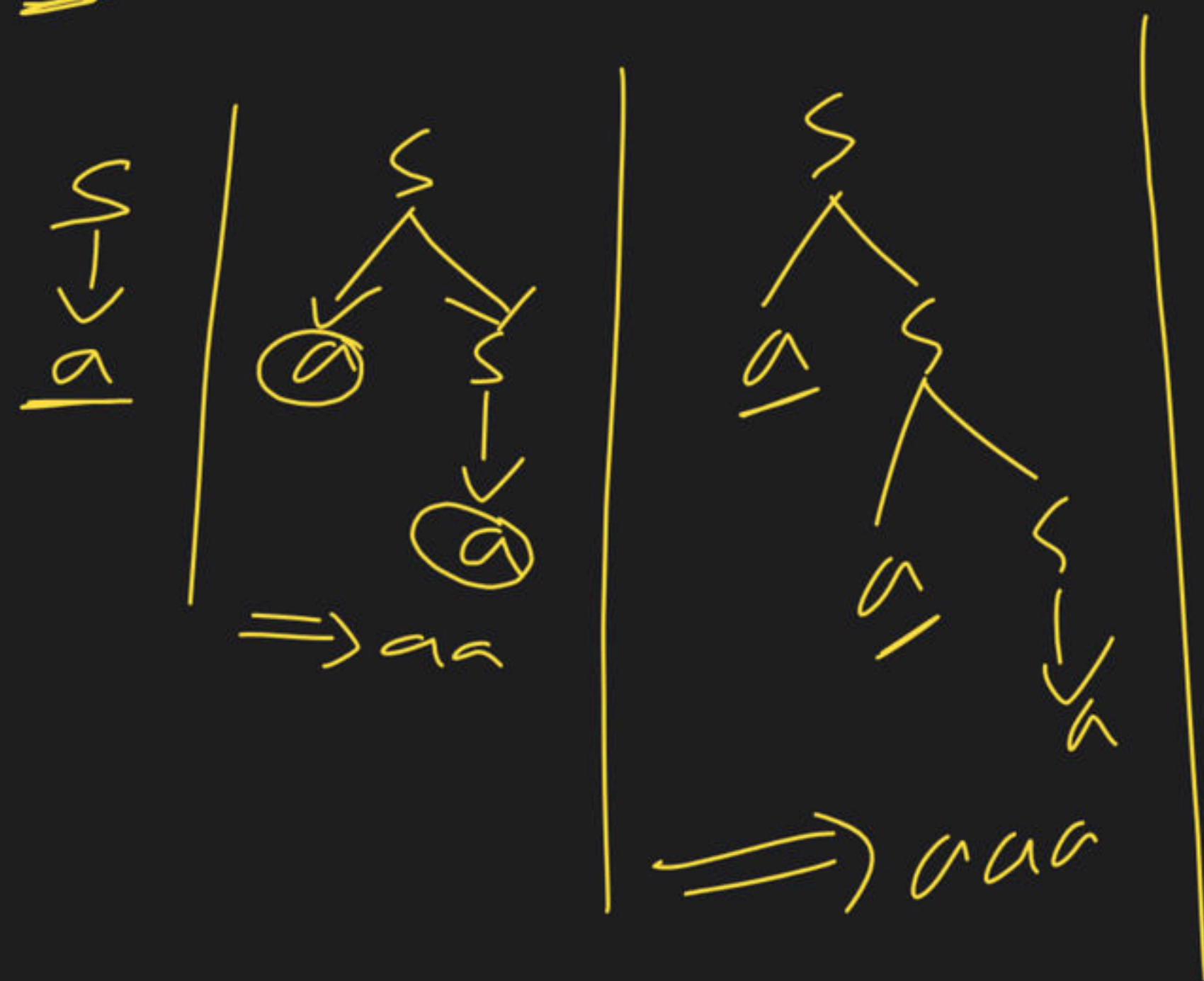
~~$E \rightarrow g$~~

P



ex

$s \rightarrow as/a$



$\Rightarrow \{a, aa, aaa, \dots\}$

$\Rightarrow L = \{a^n \mid n \geq 1\}$

ex

Context Free Gram

$S \rightarrow as | sa | a$

$a \Rightarrow 1$

string: ~~xy~~



2-possible parse tree  
 possible  
 Ambiguous

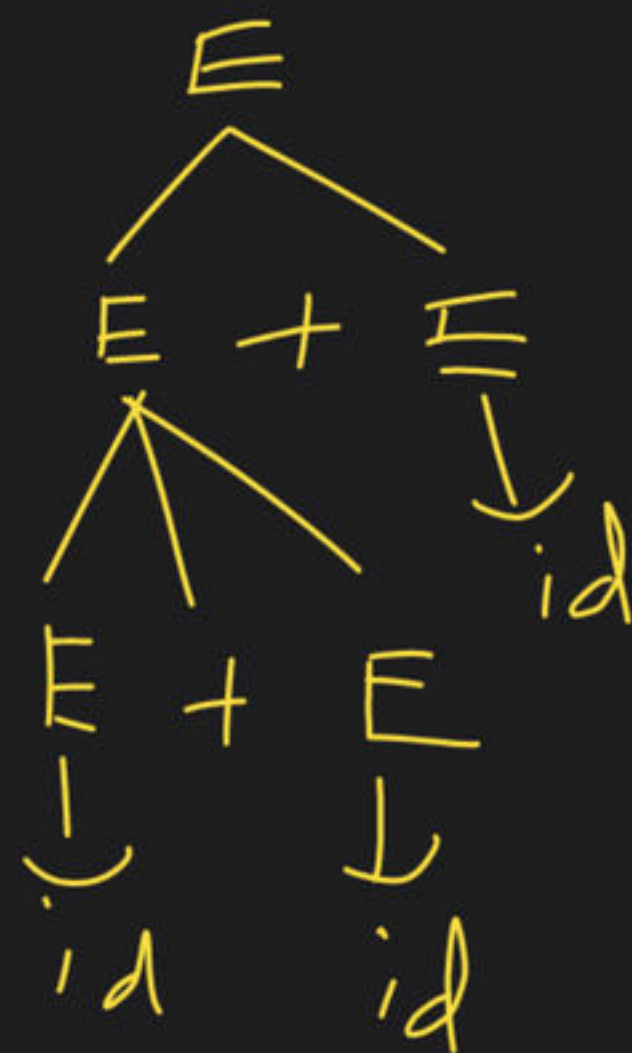
$aah \Rightarrow 4$



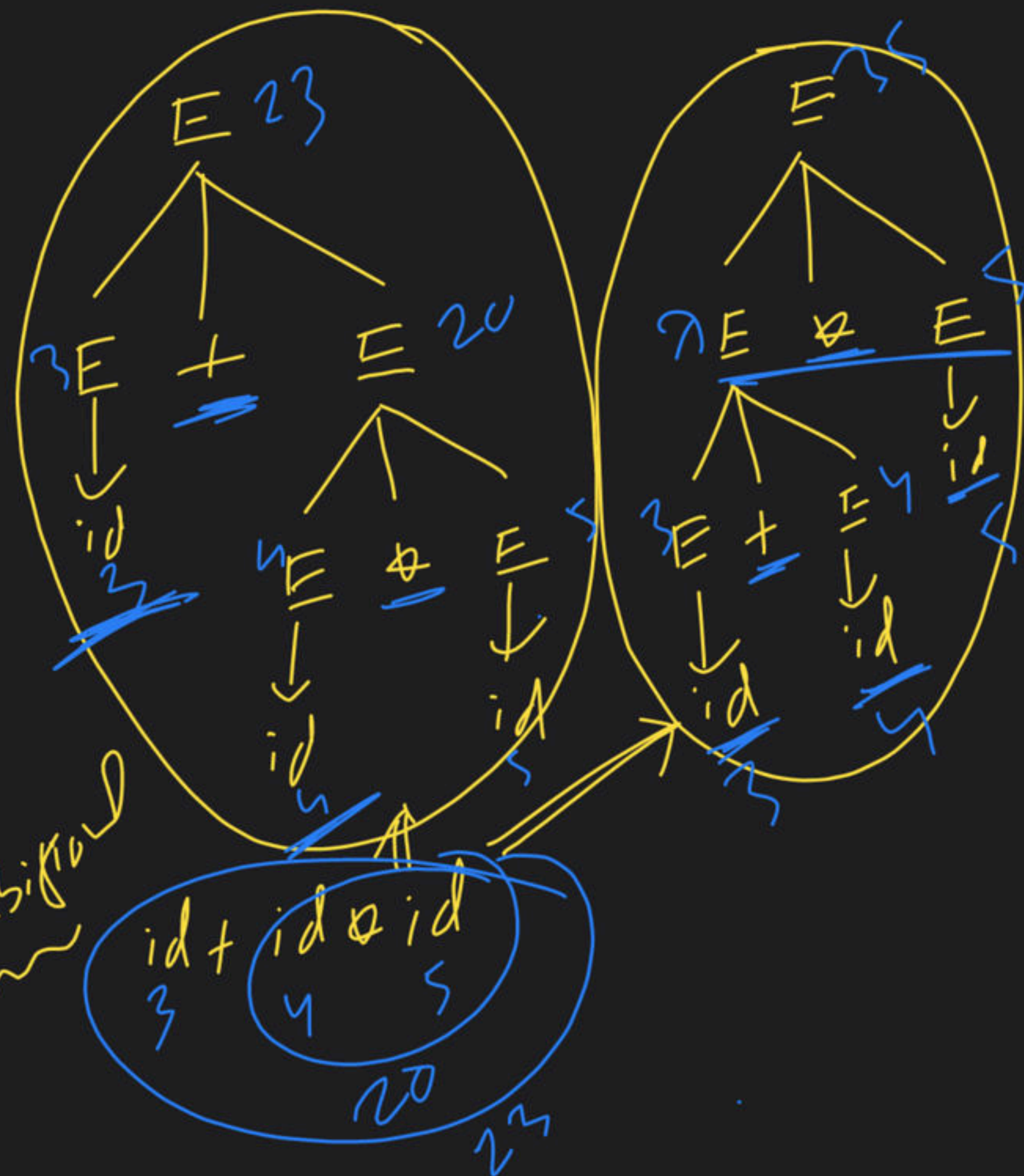


ex

$E \rightarrow E + E \mid E * E \mid id \mid E - E \mid E / E$

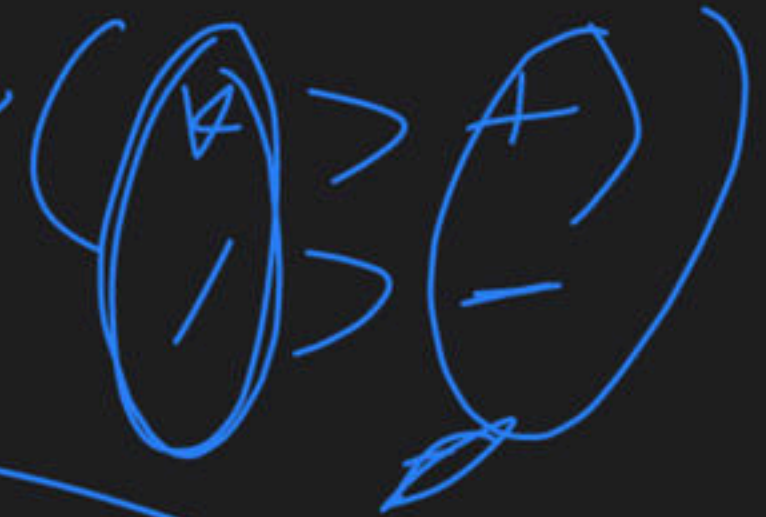


It is ambiguous grammar





conversion from Ambiguous to unambiguous  
according to C-Lang



$$\begin{aligned}
 E &\rightarrow E + T \mid T \mid E - T \\
 T &\rightarrow T * F \mid F \mid T / F \\
 F &\rightarrow id \mid (E)
 \end{aligned}$$