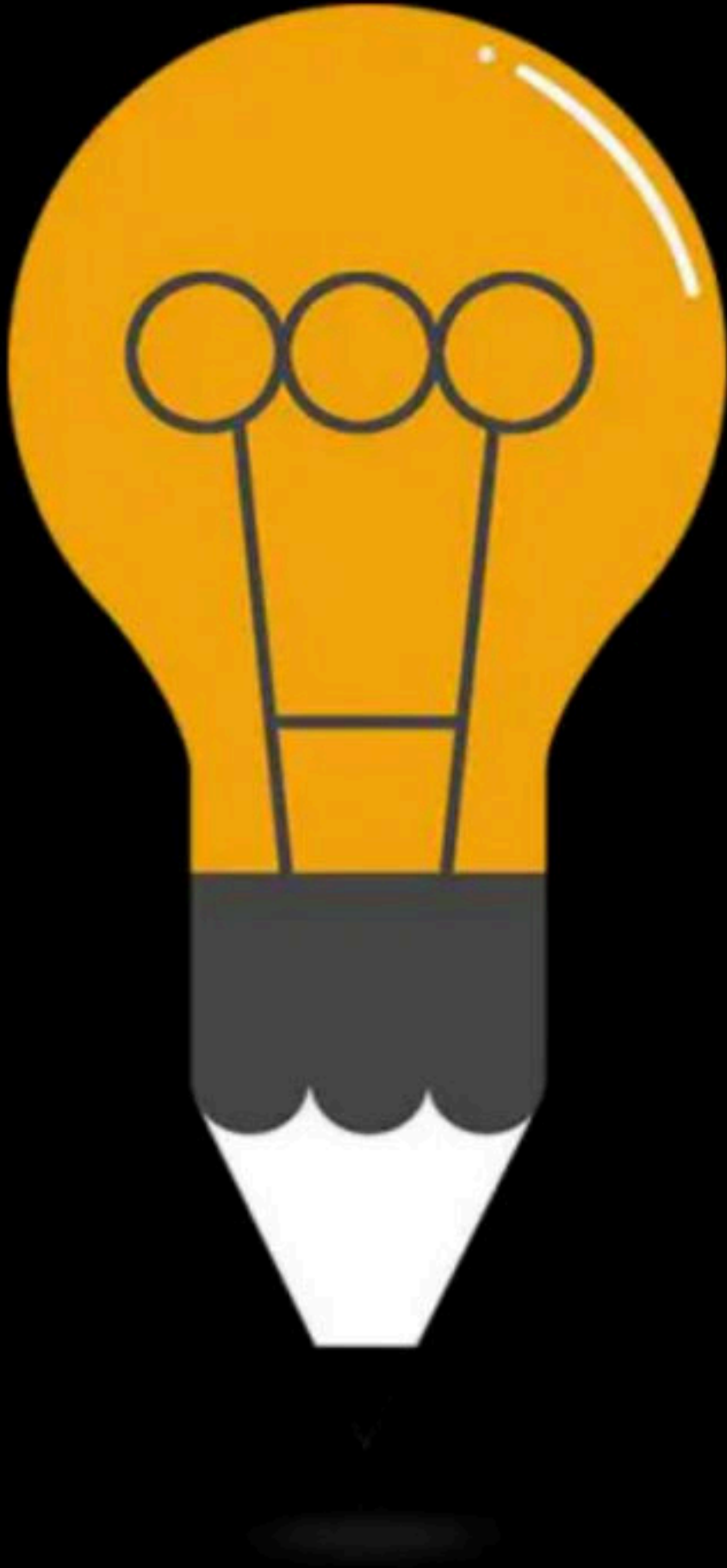




# Memory Management: Contiguous

Comprehensive Course on Operating System for GATE - 2024/25



# Operating System Memory Management

By: Vishvadeep Gothi

# Types of Locks

while (true);

1. Spinlock → busy waiting
2. Livelock ✓
3. Deadlock → 2 or more processes blocked because of each other
4. Semaphores
5. Reentrant Locks

Livelock :- 2 or more processes in busy waiting for each other forever.



ex:-

binary semaphores

$$S_x = 10$$

$$S_y = 10$$

Processes

$P_1$

$P_2$

$P(S_x)$

$P(S_y)$

$P(S_y)$

$P(S_x)$

⋮

⋮

Processes Run

$P_1$

$P_2$

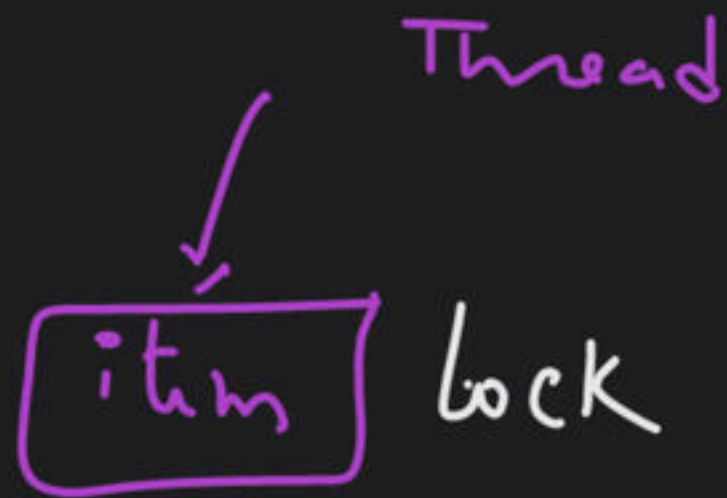
$P(S_x)$

$P(S_y)$

while ( $S_y \leq 0$ );  
while ( $S_x \leq 0$ );

lock:-

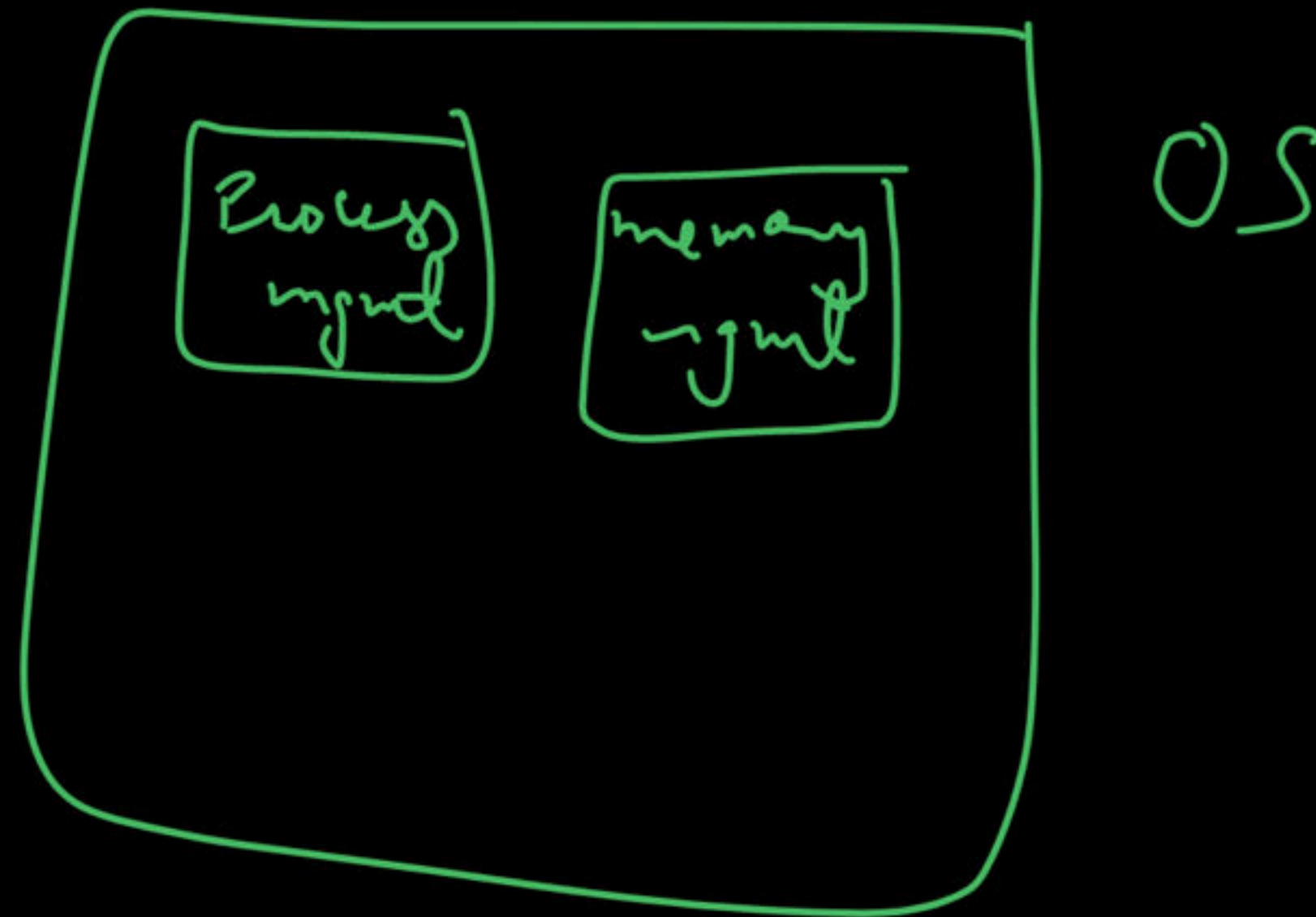
A thread can enter (acquire) same lock again.



count == 3

# Memory Management

© Module of OS

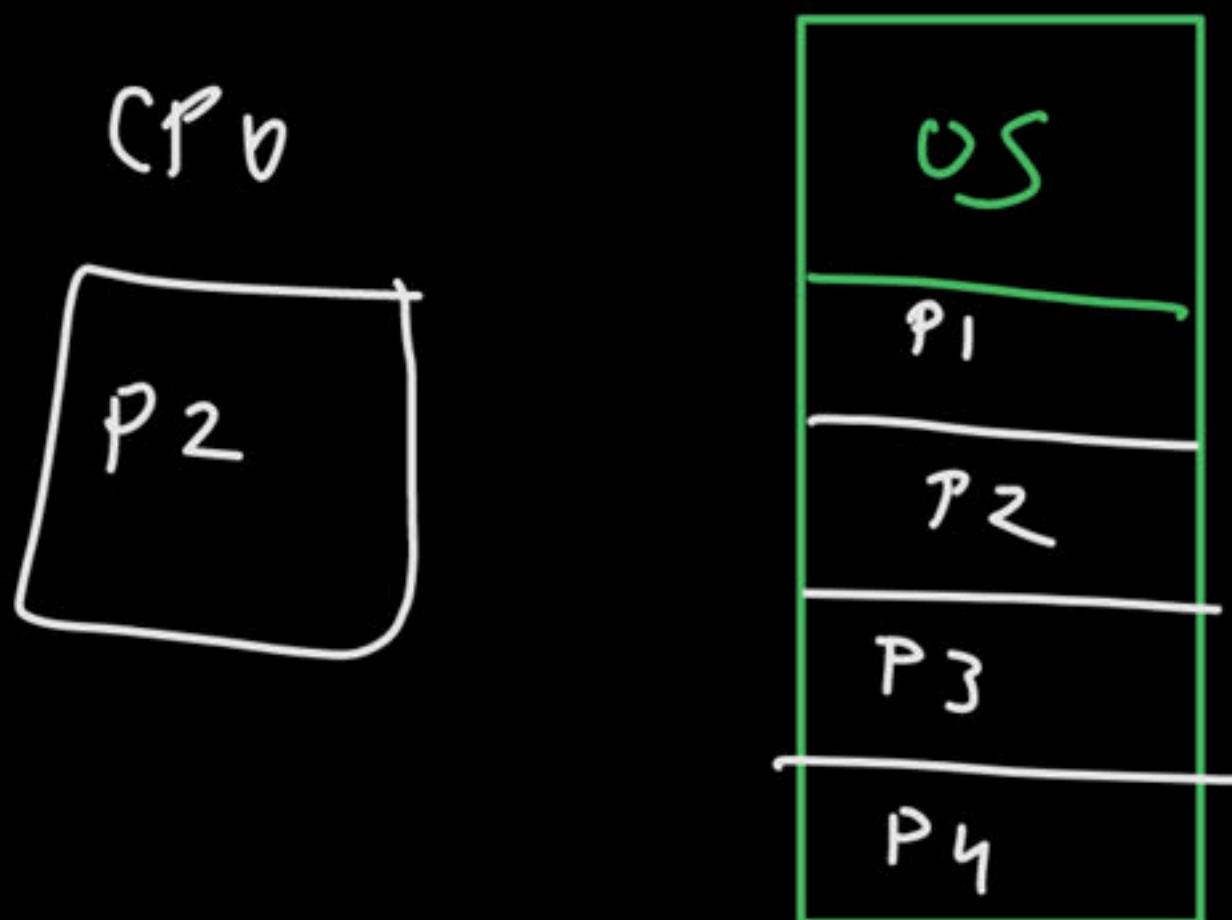




# Functions of Memory Management

1. Memory allocation :- Allocating space in mm to a newly admitted process.
2. Memory deallocation :- Deallocating space from mm for a terminated process.
3. Memory protection

A process can access only that part of memory which has been allocated to it by OS.



4. Free space management



# Goals of Memory Management

1. Maximum Utilization of space (min. space wastage)
2. Ability to run larger programs with limited space

↳ virtual memory

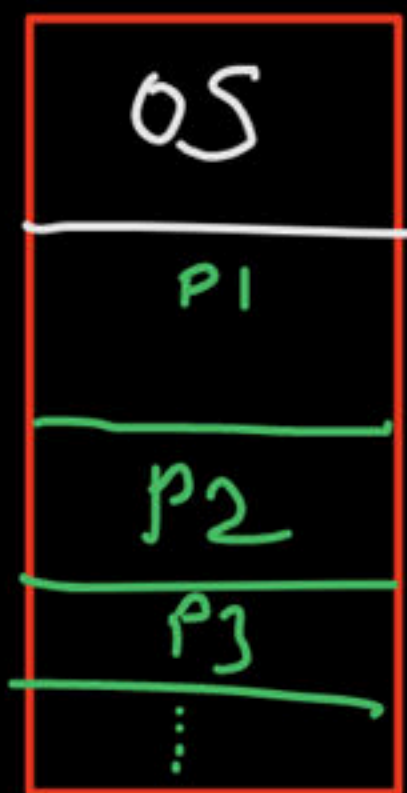
space wastage ⇒ fragmentation

↳ Internal  
↳ External

# Memory Management Techniques

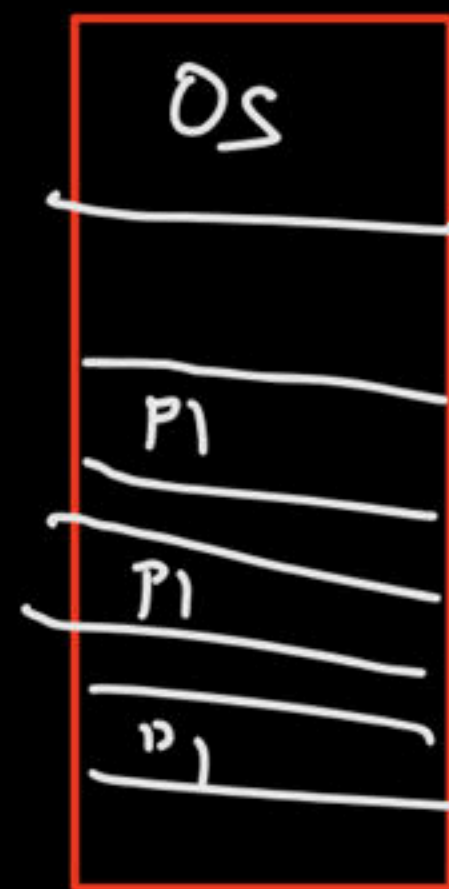
## Contiguous

Entire process is stored in memory on consecutive memory locations.



## Non-contiguous

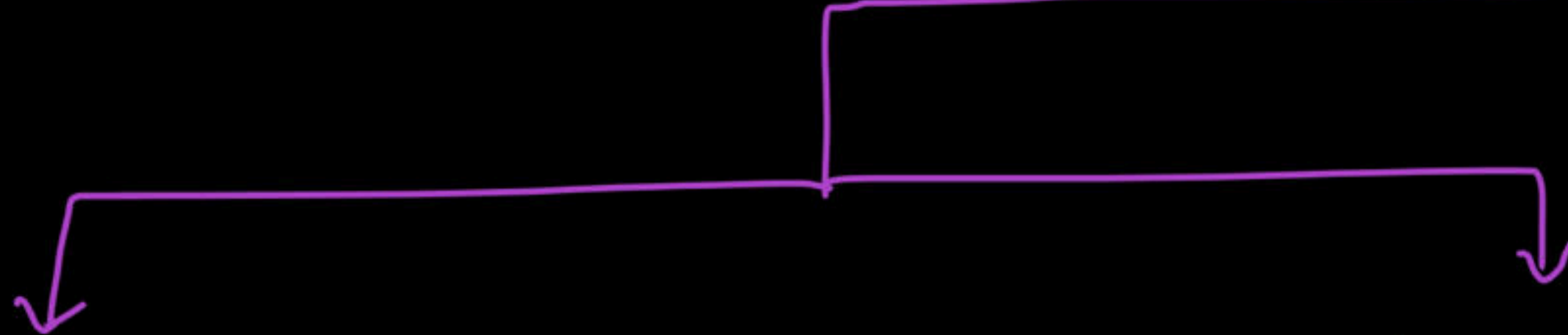
Entire process is divided into partitions and each partition can be stored on memory anywhere.





# Contiguous Memory Management

Entire process should be stored on consecutive memory locations



Fixed partition  
contiguous mem. mgmt  
Technique

=> Internal Fragmentation

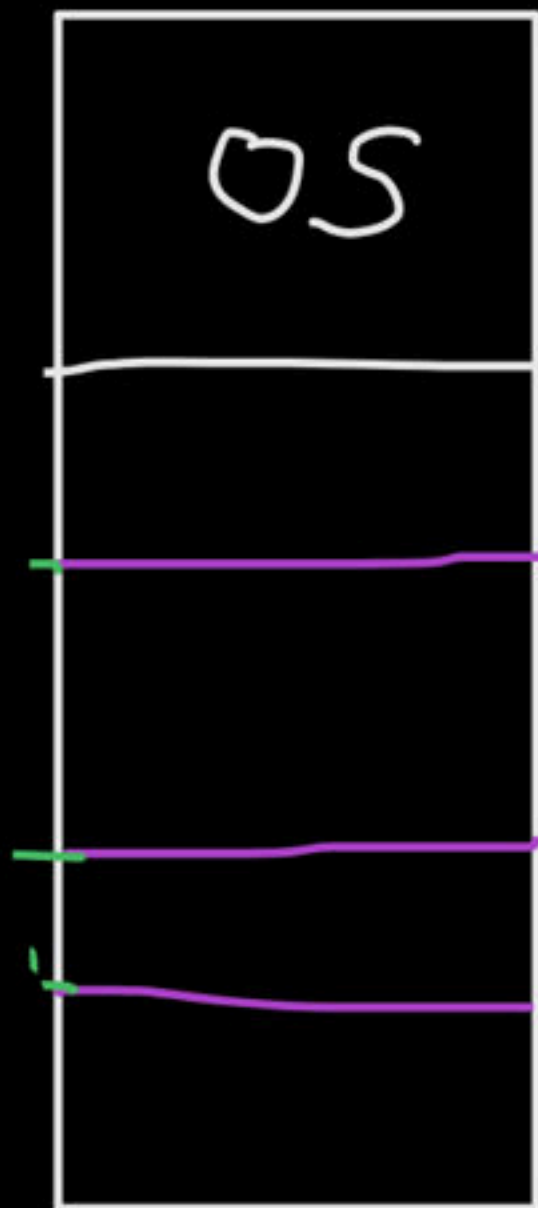
variable partition  
contiguous m.m. technique

=> External Fragmentation

# Fixed Partition Contiguous MMT

OS divides m.m. into fixed no. of partitions. Each partition can be used to accommodate <sup>only</sup> one process.

→ Each partition size can be different—diff.



---

Whenever a new process arrives mem. mgmt. provides it a free partition



# Partition Allocation Policy

	OS	Process P1 size = 120 KB	P2 size = 80 KB
100 KB			
150 KB	P1	150 KB	100 KB
300 KB			
125 KB			
200 KB			

First fit

150 KB

100 KB

Best fit

125 KB

worst fit

300 KB

next fit

150 KB

300 KB

# Partition Allocation Policy

First fit:- Find a partiti<sup>n</sup> which can accomodate the new process starting from first partiti<sup>n</sup>.

Best fit:- Find smallest possible partiti<sup>n</sup> which can accommodate the new process.

worst fit:- Find largest partiti<sup>n</sup>

Next fit:- Find next partiti<sup>n</sup> from current allocated partiti<sup>n</sup> in sequence.



Fixed partition contiguous mem. mgmt technique suffers from internal fragmentation.



The space allocated to a process is more than the required space. hence the extra unused space is known as internal fragmentation.

ex:- A process of size 120KB has been given a partition of size 130KB.  $\Rightarrow$  size of internal fragmentation = 10 KB.

# Variable Partition Contiguous MMT

no any partit<sup>n</sup> is created before hand.

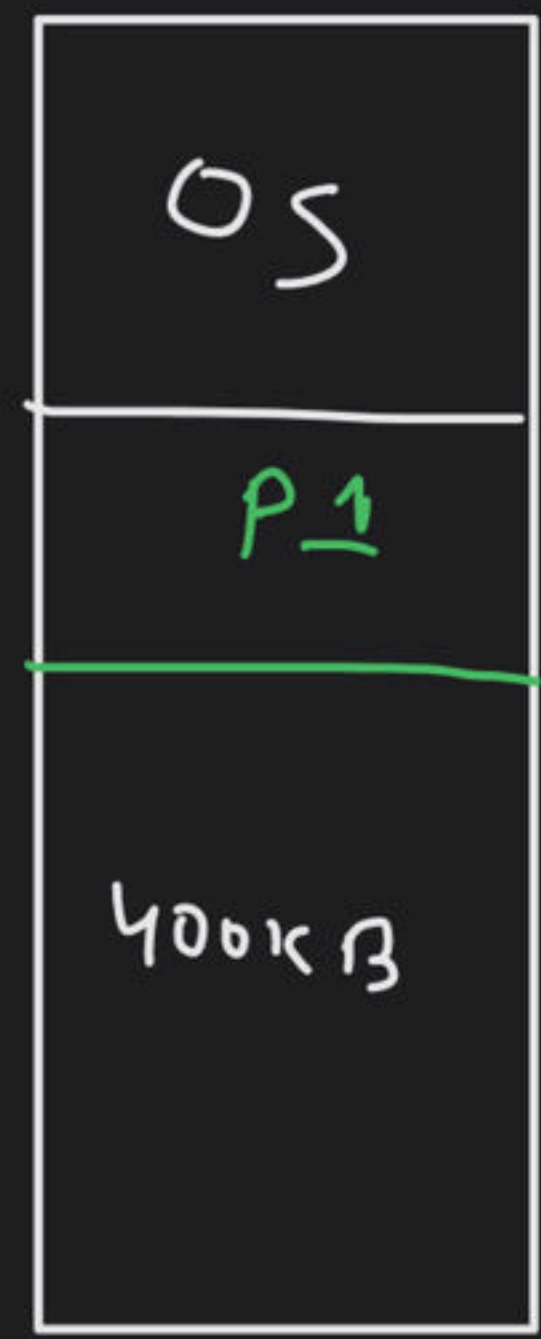
→ whenever a new process arrives then a new partition is created equal to the size of process. and the process is stored in that partition.

⇒ no any internal fragmentation

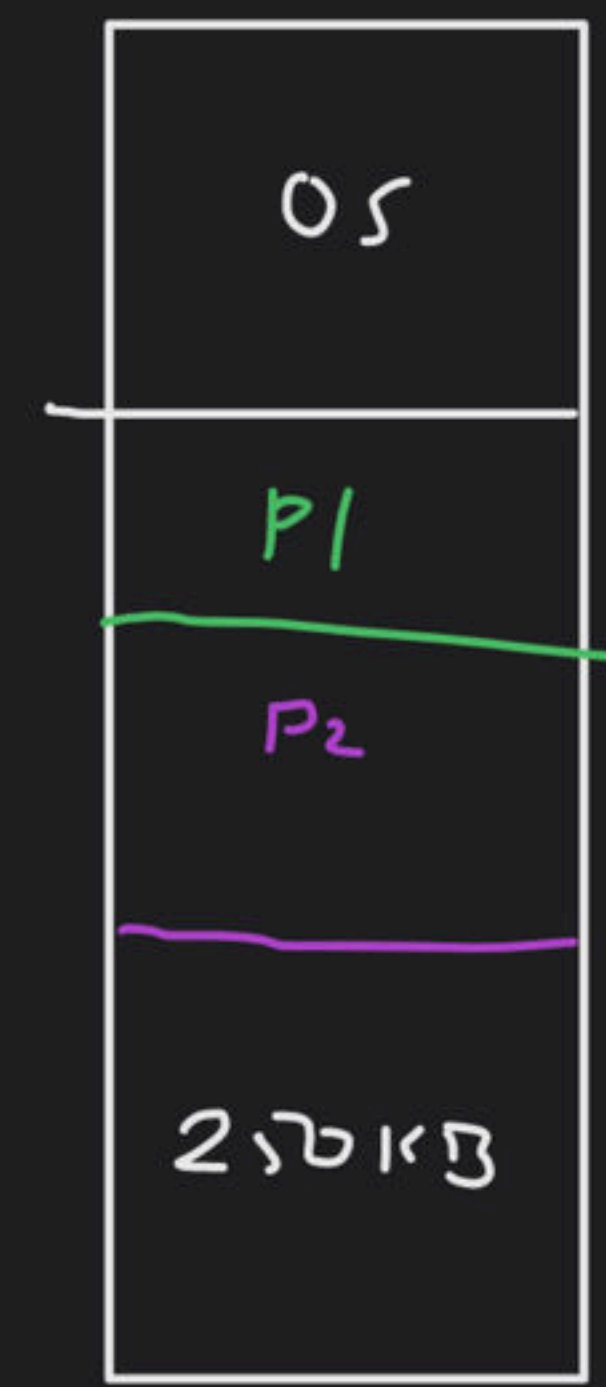




new process P<sub>1</sub>  
size = 200KB

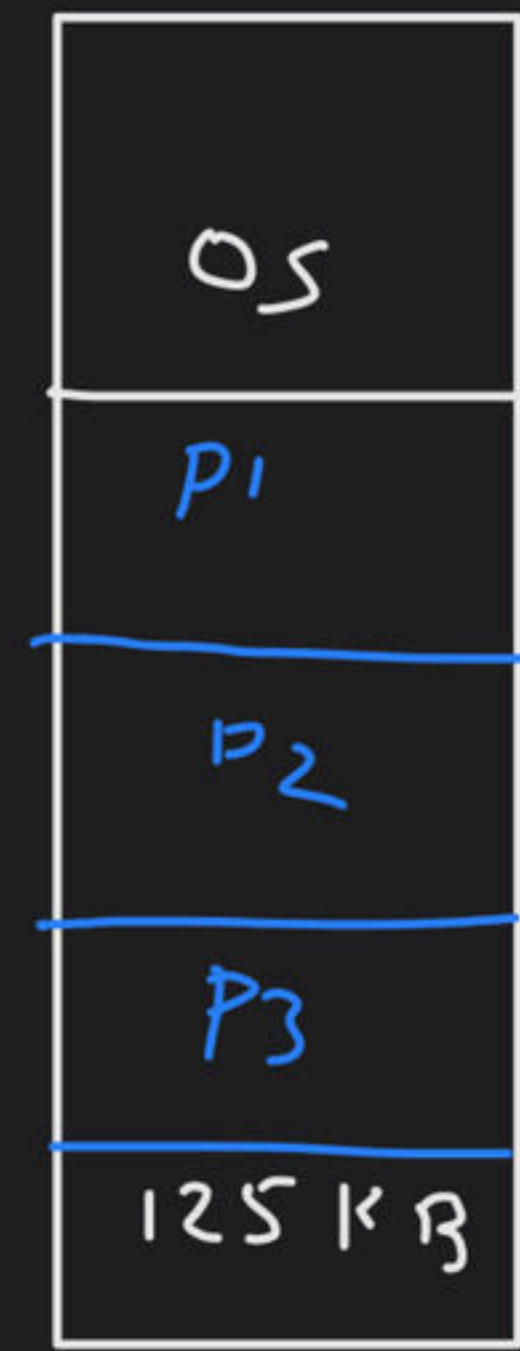


new process P<sub>2</sub>  
size = 150KB

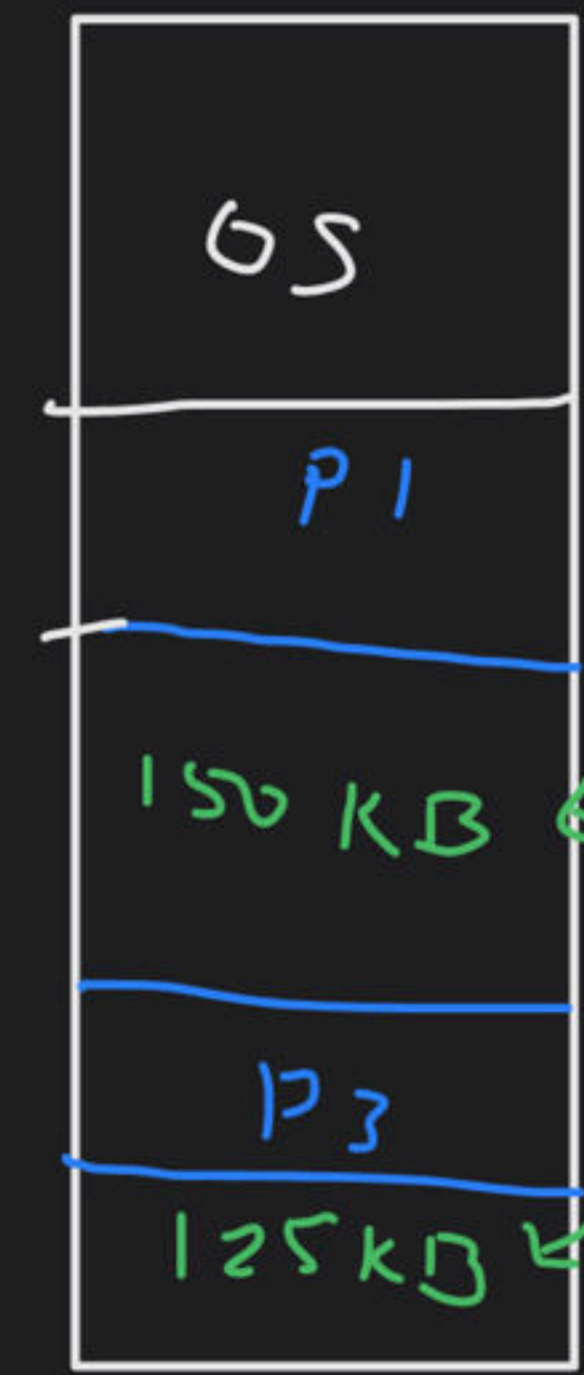


new process P3

size = 125 KB



Process P2  
Terminated



empty holes



new process P<sub>4</sub>

size = 170 KB

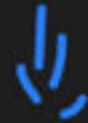


External  
fragmentation

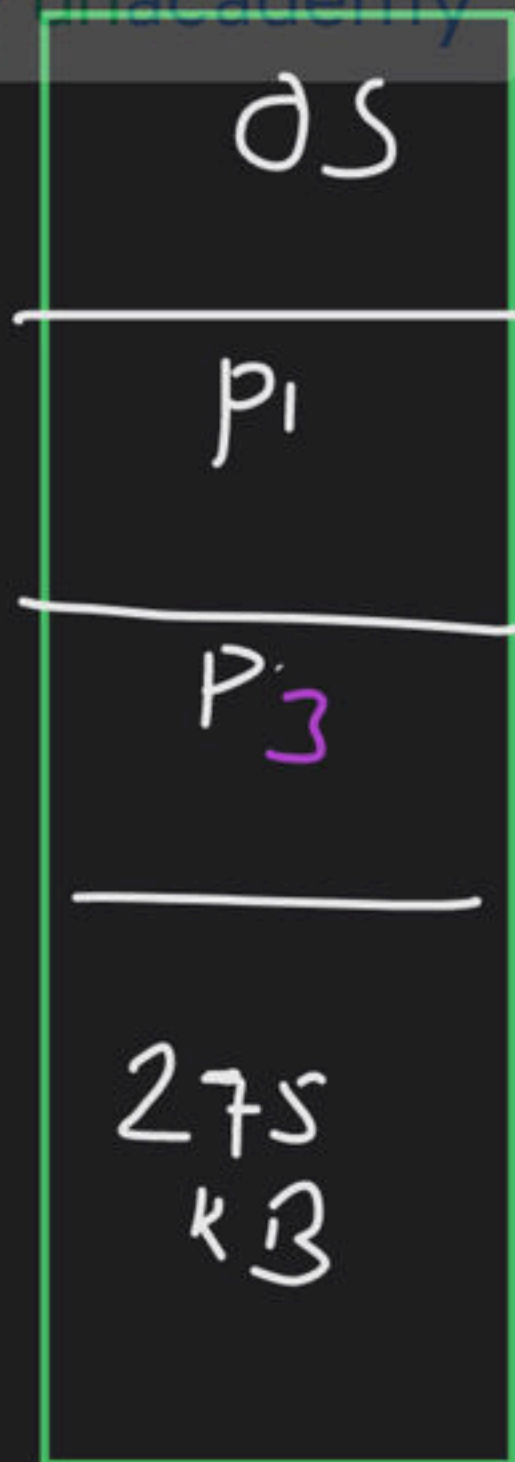
(Enough space is available  
in memory but not  
contiguous to store a  
process)



Sol<sup>n</sup> = compaction

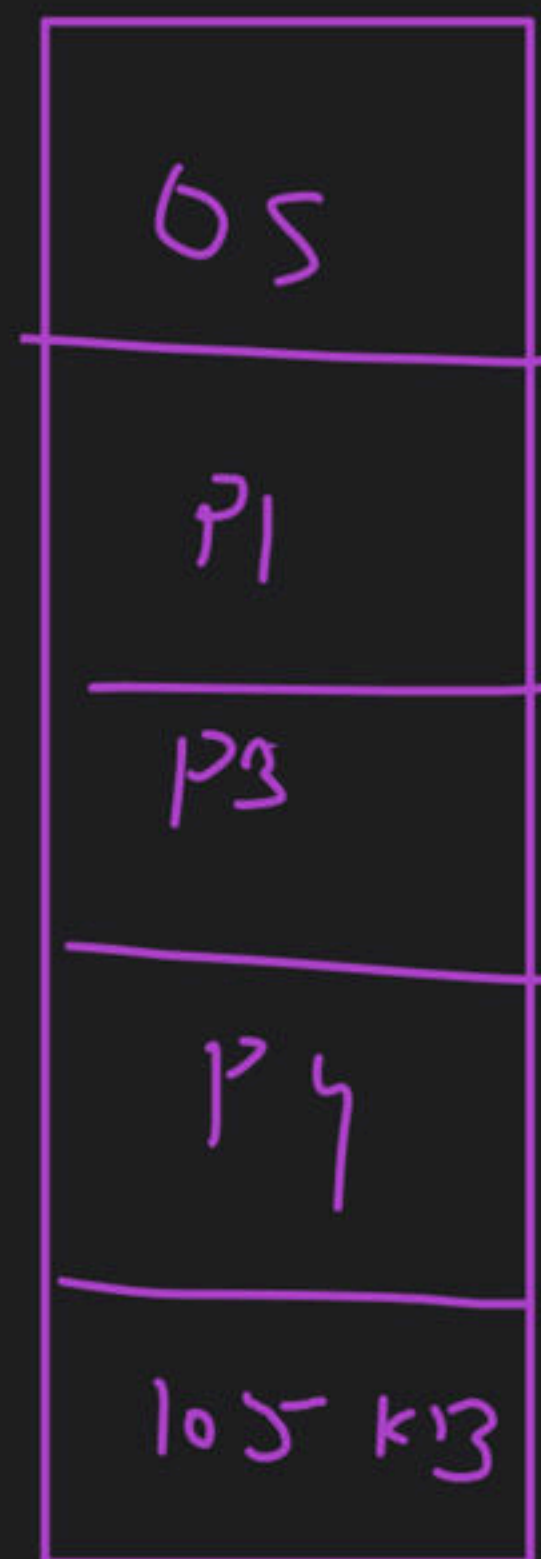


Gather all allocated spaces into one side  
of memory so that other side can have  
all empty spaces together.



P4  
170 KB

→





# Question

Consider the requests from processes in given order 300K, 25K, 125K, and 50K. Let there be two blocks of memory available of size 150K followed by a block size 350K. Which of the following partition allocation schemes can satisfy the above requests?

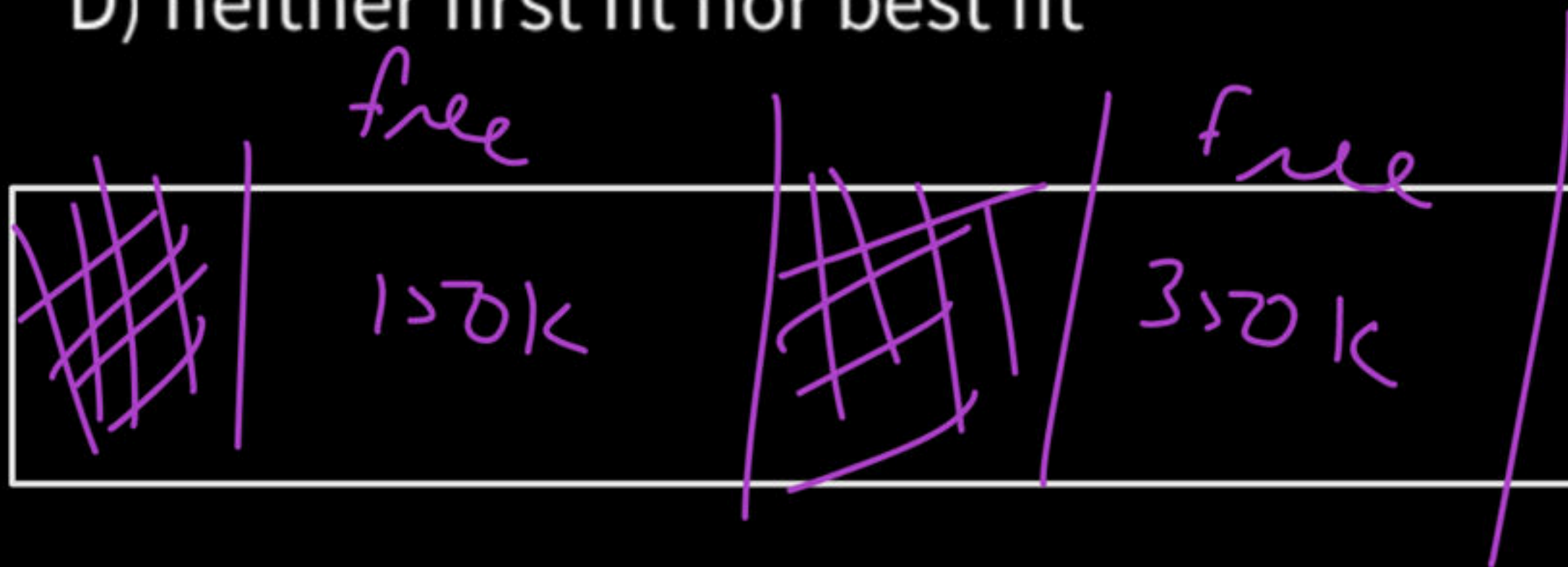
A) Best fit but not first fit

☒ B) First fit but not best fit

C) Both First fit & Best fit

D) neither first fit nor best fit

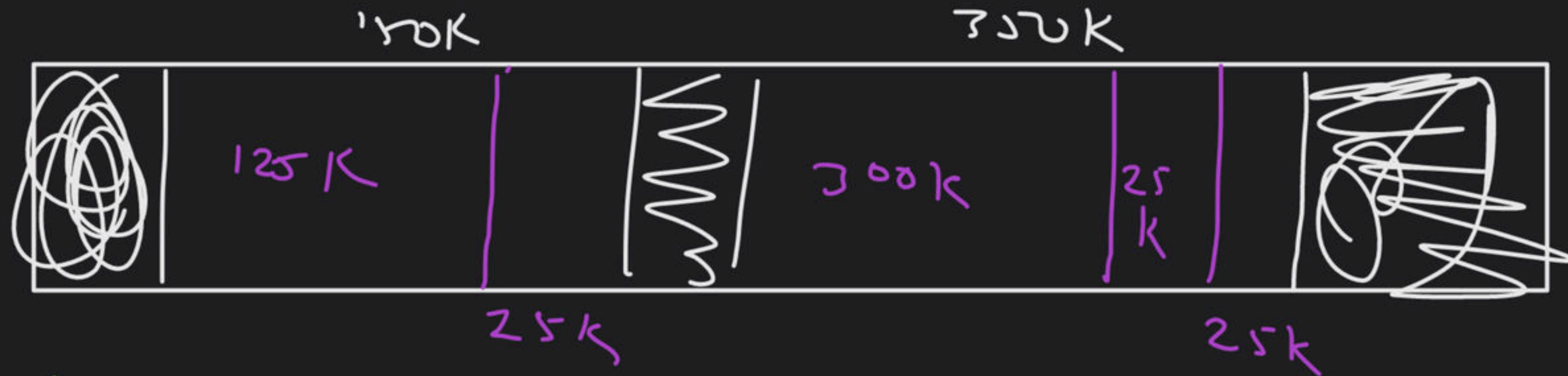
(variable partition contiguous memt)



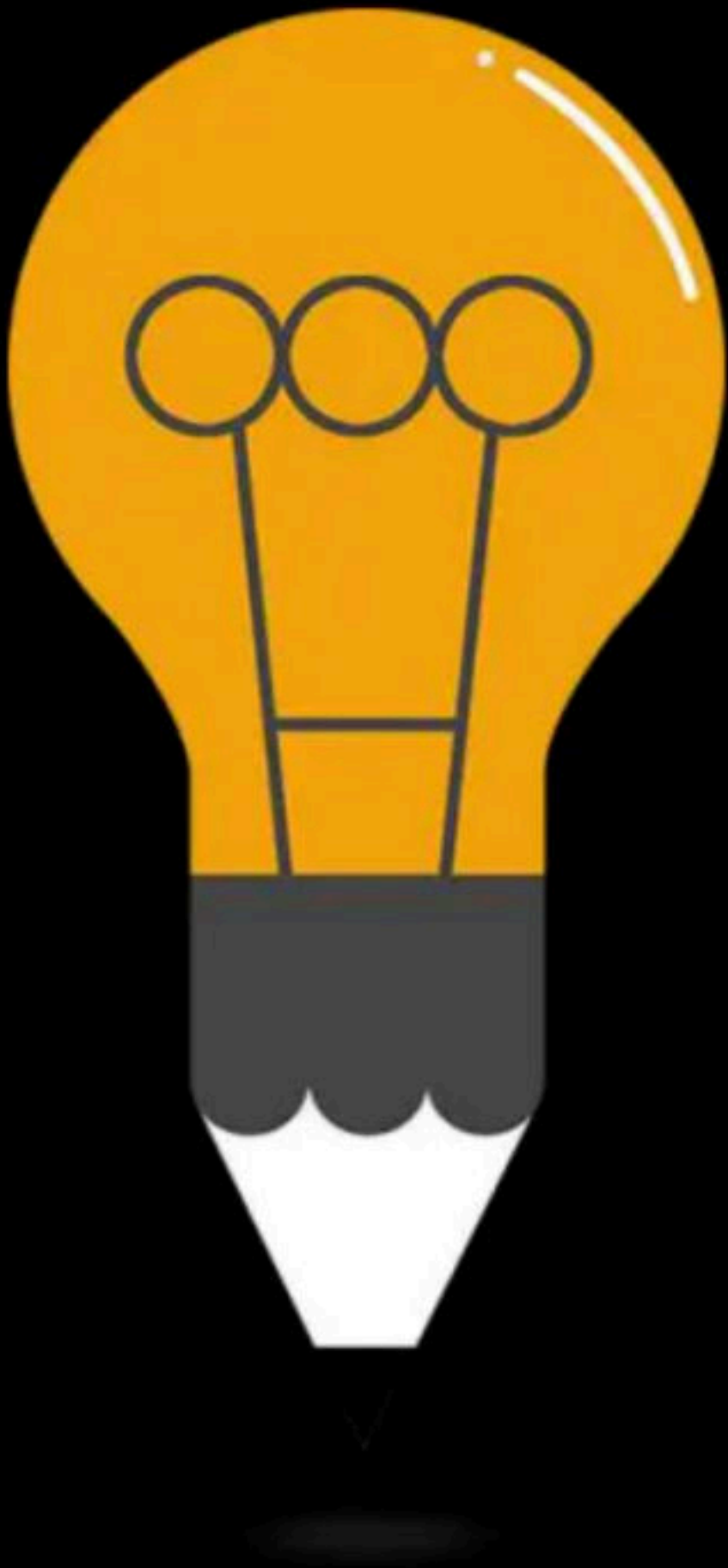
## First fit:-



## Best fit:-



last process not placed



# DPP

By: **Vishvadeep Gothi**



# Question 1

Consider a fixed partition MMT where there are 5 partitions of size 100MB, 250MB, 200MB, 500MB and 300MB. All Partitions are initially empty. The following process requests are made in the given order:

Process	Size
P1	150MB
P2	400MB
P3	270MB
P4	180MB
P5	80MB

Provide the following answers for First fit, Best fit and Worst Fit policies?

1. Maximum degree of multiprogramming?
2. What is the total internal fragmentation size?



# Question 2

Consider variable partition MMT where there are 4 partitions of size 250MB, 200MB, 500MB and 400MB. The following process requests are made in the given order:

Process	Size
P1	150MB
P2	400MB
P3	270MB
P4	180MB
P5	80MB
P6	50MB

Provide the processes are stored for First fit, Best fit and Worst Fit policies?

# Happy Learning.!

