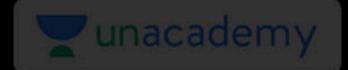
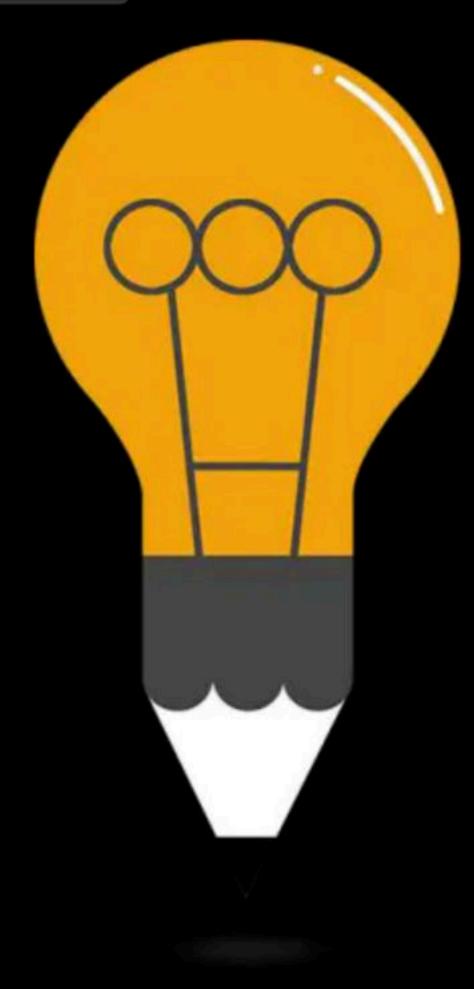




Relational Algebra PYQ Discussion: Part 1

Special class

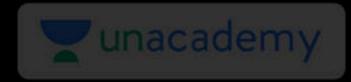




PYQ:

Relational DB

By: Vishvadeep Gothi



Suppose we have a database consisting of the following three relations:

FREQUENTS	(CUSTOMER, HOTEL)
SERVES	(HOTEL, SNACKS)
LIKES	(CUSTOMER, SNACKS)

The first indicates the hotels each customer visits, the second tells which snacks each hotel serves and last indicates which snacks are liked by each customer. Express the following query in relational algebra:

Print the hotels the serve the snack that customer Rama likes.



Consider the following relational schema:

- COURSES (cno, cname)
- STUDENTS (rollno, sname, age, year)
- REGISTERED_FOR (cno, rollno)

The underlined attributes indicate the primary keys for the relations. The 'year' attribute for the STUDENTS relation indicates the year in which the student is currently studying (First year, Second year etc.)

- a. Write a relational algebra query to print the roll number of students who have registered for cno 322.
 - b. Write a SQL query to print the age and year of the youngest student in each year.



Consider the relation scheme.

AUTHOR	(ANAME, INSTITUTION, ACITY, AGE)
PUBLISHER	(PNAME, PCITY)
BOOK	(TITLE, ANAME, PNAME)

Express the following queries using (one or more of) SELECT, PROJECT, JOIN and DIVIDE operations.

Express the following queries using (one or more of) SELECT, PROJECT, JOIN and DIVIDE operations.

- a. Get the names of all publishers.
- b. Get values of all attributes of all authors who have published a book for the publisher with PNAME='TECHNICAL PUBLISHERS'.
- c. Get the names of all authors who have published a book for any publisher located in Madras

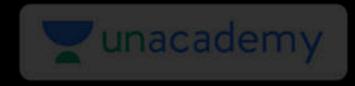


Tauthor. Anome, = norFlotion, (Sprane = ' Tech pub.' (Author & book))

acity, ge

<u>ری</u>

Traname (Prity = 'madras' (i'46 lisher & book)



A library relational database system uses the following schema

- USERS (User#, User Name, Home Town)
- BOOKS (Book#, Book Title, Author Name)
- ISSUED (Book#, User#, Date)

Explain in one English sentence, what each of the following relational algebra queries is designed to determine

- a. $\sigma_{\text{User}\#=6} \left(\pi_{\text{User}\#, \text{ Book Title}} \left((\text{USERS} \bowtie \text{ISSUED}) \bowtie \text{BOOKS} \right) \right)$
- b. $\pi_{\text{Author Name}}$ (BOOKS $\bowtie \sigma_{\text{Home Town=Delhi}}$ (USERS \bowtie ISSUED))
- Title of all books which have been issued to usere with user# 6.
 - have been issued to residents of delhi-

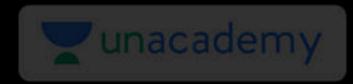
Consider the following relational database schema:

- EMP (eno name, age)
- PROJ (pno name)
- INVOLVED (eno, pno)

EMP contains information about employees. PROJ about projects and involved about which employees involved in which projects. The underlined attributes are the primary keys for the respective relations.

What is the relational algebra expression containing one or more of $\{\sigma, \pi, \times, \rho, -\}$ which is equivalent to SQL query.

select eno from EMP|INVOLVED where EMP.eno=INVOLVED.eno and INVOLVED.pno=3



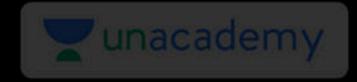
Given two union compatible relations $R_1(A, B)$ and $R_2(C, D)$, what is the result of the operation $R_1 \bowtie_{A=C \land B=D} R_2$?

A.
$$R_1 \cup R_2$$

B.
$$R_1 \times R_2$$

C.
$$R_1 - R_2$$

$$R_1 \cap R_2$$



Consider the following relational database schemes:

- COURSES (Cno, Name)
- PRE REQ(Cno, Pre Cno)
- COMPLETED (Student_no, Cno)

COURSES gives the number and name of all the available courses.

PRE_REQ gives the information about which courses are pre-requisites for a given course.

COMPLETED indicates what courses have been completed by students

Express the following using relational algebra:

List all the courses for which a student with Student_no 2310 has completed all the pre-requisites.



Consider the join of a relation R with a relation S. If R has m tuples and S has n tuples then the maximum and minimum sizes of the join respectively are

A. m+n and 0



C.
$$m+n$$
 and $|m-n|$ D. mn and $m+n$

Given the relations

- · employee (name, salary, dept-no), and
- department (dept-no, dept-name, address),

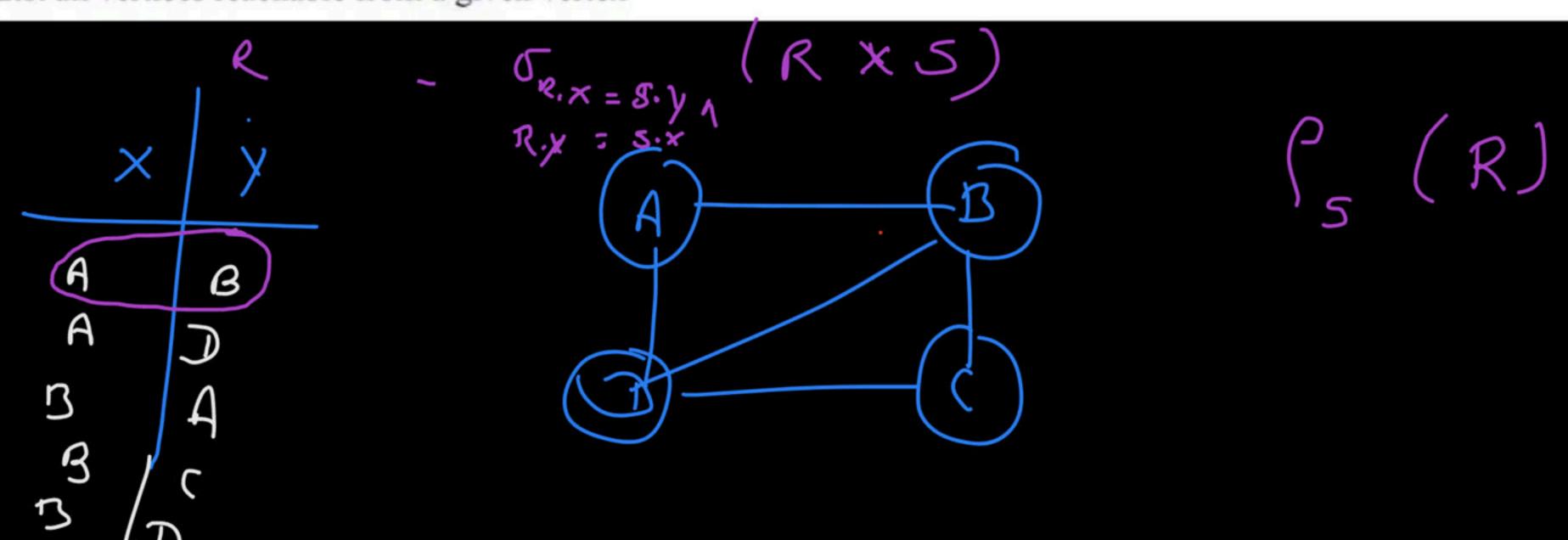
Which of the following queries cannot be expressed using the basic relational algebra operations $(\sigma, \pi, \times, \bowtie, \cup, \cap, -)$

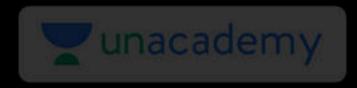
- A. Department address of every employee
- B. Employees whose name is the same as their department name
- The sum of all employees' salaries
- D. All employees of a given department



Suppose the adjacency relation of vertices in a graph is represented in a table Adj (X,Y). Which of the following queries cannot be expressed by a relational algebra expression of constant length?

- A. List all vertices adjacent to a given vertex
- B. List all vertices which have self loops
- C. List all vertices which belong to cycles of less than three vertices
- D. List all vertices reachable from a given vertex



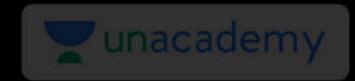


Let r and s be two relations over the relation schemes R and S respectively, and let A be an attribute in R. The relational algebra expression $\sigma_{A=a}(r\bowtie s)$ is always equal to

A.
$$\sigma_{A=a}(r)$$

C. $\sigma_{A=a}(r) \bowtie s$

B.
$$r$$
None of the above



A university placement center maintains a relational database of companies that interview students on campus and make job offers to those successful in the interview. The schema of the database is given below:

COMPANY(<u>cname</u>, clocation) STUDENT(<u>srollno</u>, sname, sdegree)
INTERVIEW(<u>cname</u>, srollno, idate) OFFER(<u>cname</u>, srollno, osalary)

The COMPANY relation gives the name and location of the company. The STUDENT relation gives the student's roll number, name and the degree program for which the student is registered in the university. The INTERVIEW relation gives the date on which a student is interviewed by a company. The OFFER relation gives the salary offered to a student who is successful in a company's interview. The key for each relation is indicated by the underlined attributes

- a. Write a relational algebra expressions (using only the operators \bowtie , σ , π , \cup , -) for the following queries.
 - i. List the rollnumbers and names of students who attended at least one interview but did not receive any job offer.
 - List the rollnumbers and names of students who went for interviews and received job offers from every company
 with which they interviewed.
- b. Write an SQL query to list, for each degree program in which more than five students were offered jobs, the name of the degree and the average offered salary of students in this degree program.

Consider the following SQL query

Select distinct a_1, a_2, \ldots, a_n

from r_1, r_2, \ldots, r_m

where P

For an arbitrary predicate P, this query is equivalent to which of the following relational algebra expressions?

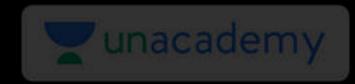
A.
$$\Pi_{a_1,a_2,\ldots a_n}\sigma_p\left(r_1\times r_2\times\cdots\times r_m\right)$$

- B. $\Pi_{a_1,a_2,\ldots a_n}\sigma_p\left(r_1\bowtie r_2\bowtie\cdots\bowtie r_m\right)$
- C. $\Pi_{a_1,a_2,\ldots a_n} \sigma_p (r_1 \cup r_2 \cup \cdots \cup r_m)$
- D. $\Pi_{a_1,a_2,\ldots a_n} \sigma_p (r_1 \cap r_2 \cap \cdots \cap r_m)$

Consider the relation Student (<u>name</u>, sex, marks), where the primary key is shown underlined, pertaining to students in a class that has at least one boy and one girl. What does the following relational algebra expression produce? (Note: ρ is the rename operator).

$$\pi_{name}\{\sigma_{sex=female}(Student)\} - \pi_{name}(Student \bowtie_{(sex=female \land x=male \land marks \le m)} \rho_{n,x,m}(Student))$$

- A. names of girl students with the highest marks
- B. names of girl students with more marks than some boy student
- C. names of girl students with marks not less than some boy student
- names of girl students with more marks than all the boy students



A table 'student' with schema (roll, name, hostel, marks), and another table 'hobby' with schema (roll, hobbyname) contains records as shown below:

Table: student

Roll	Name	Hostel	Marks
1798	Manoj Rathor	7	95
2154	Soumic Banerjee	5	68
2369	Gumma Reddy	7	86
2581	Pradeep pendse	6	92
2643	Suhas Kulkarni	5	78
2711	Nitin Kadam	8	72
2872	Kiran Vora	5	92
2926	Manoj Kunkalikar	5	94
2959	Hemant Karkhanis	7	88
3125	Rajesh Doshi	5	82

ren_			4.4	
13	ble	: D(3OC	YY .

Roll	Hobby Name
1798	chess
1798	music
2154	music
2369	swimming
2581	cricket
2643	chess
2643	hockey
2711	volleyball
2872	football
2926	cricket
2959	photography
3125	music
3125	chess

The following SQL query is executed on the above tables:

Am -

- L

select hostel from student natural join hobby where marks >= 75 and roll between 2000 and 3000;

Relations S and H with the same schema as those of these two tables respectively contain the same information as tuples. A new relation S' is obtained by the following relational algebra operation:

$$S' = \Pi_{\text{hostel}}((\sigma_{s.roll=H.roll}(\sigma_{marks>75 \text{ and } roll>2000 \text{ and } roll<3000}(S)) \times (H))$$

The difference between the number of rows output by the SQL statement and the number of tuples in S' is

Information about a collection of students is given by the relation studInfo(studId, name, sex). The relation enroll(studId, courseId) gives which student has enrolled for (or taken) what course(s). Assume that every course is taken by at least one male and at least one female student. What does the following relational algebra expression represent?

$$\Pi_{\text{courseId}} \left(\left(\Pi_{\text{studId}} \left(\sigma_{\text{sex="female"}} \left(\text{studInfo} \right) \right) \times \Pi_{\text{courseId}} \left(\text{enroll} \right) \right) - \text{enroll} \right)$$

- A. Courses in which all the female students are enrolled.
- Courses in which a proper subset of female students are enrolled.
 - C. Courses in which only male students are enrolled.
 - D. None of the above

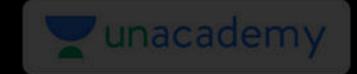
Let R and S be two relations with the following schema

where $\{P,Q\}$ is the key for both schemas. Which of the following queries are equivalent?

- I. $\Pi_P(R \bowtie S)$
- II. $\Pi_P(R)\bowtie\Pi_P(S)$
- III. $\Pi_{P}\left(\Pi_{P,Q}\left(R\right)\cap\Pi_{P,Q}\left(S\right)\right)$
- IV. $\Pi_{P}(\Pi_{P,Q}(R) (\Pi_{P,Q}(R) \Pi_{P,Q}(S)))$
- A. Only I and II

- B. Only I and III
- C. Only I, II and III





Suppose $R_1(\underline{A},\underline{B})$ and $R_2(\underline{C},D)$ are two relation schemas. Let r_1 and r_2 be the corresponding relation instances. B is a foreign key that refers to C in R_2 . If data in r_1 and r_2 satisfy referential integrity constraints, which of the following is **ALWAYS TRUE**?

A.
$$\prod_B(r_1)-\prod_C(r_2)=\varnothing$$

B.
$$\prod_{C}(r_2) - \prod_{B}(r_1) = \emptyset$$

C.
$$\prod_{B}(r_1) = \prod_{C}(r_2)$$

D.
$$\prod_B(r_1) - \prod_C(r_2) \neq \emptyset$$

In a relational database there are three relations:

- Customers = C(CName),
- Shops = S(SName),
- Buys = B(CName, SName).

Which of the following relational algebra expressions returns the names of shops that have no customers at all? projection operator.]

A.
$$\Pi_{SName}B$$

C.
$$S - \Pi_{SName}B$$

B.
$$S-B$$

D.
$$S - \Pi_{SName}((C \times S) - B)$$



What is the optimized version of the relation algebra expression $\pi_{A1}(\pi_{A2}(\sigma_{F1}(\sigma_{F2}(r))))$, where A1, A2 are sets of attributes in r with $A1 \subset A2$ and F1, F2 are Boolean expressions based on the attributes in r?

A.
$$\pi_{A1}(\sigma_{(F1\wedge F2)}(r))$$

C.
$$\pi_{A2}(\sigma_{(F1\wedge F2)}(r))$$

B.
$$\pi_{A1}(\sigma_{(F1\vee F2)}(r))$$

D.
$$\pi_{A2}(\sigma_{(F1\vee F2)}(r))$$



Consider the relational schema given below, where eld of the relation dependent is a foreign key referring to empld of the relation employee. Assume that every employee has at least one associated dependent in the dependent relation.

employee (empld, empName, empAge)

dependent (depId, eId, depName, depAge)

Consider the following relational algebra query:

 $\Pi_{empId} \ (employee) - \Pi_{empId} \ (employee \bowtie_{(empId=eID) \land (empAge \leq depAge)} \ dependent)$

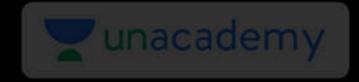
The above query evaluates to the set of emplds of employees whose age is greater than that of

A. some dependent.

C. some of his/her dependents.

B. all dependents.

D. all of his/her dependents.



SELECT operation in SQL is equivalent to

- A. The selection operation in relational algebra
- B. The selection operation in relational algebra, except that SELECT in SQL retains duplicates
- C. The projection operation in relational algebra
- D. The projection operation in relational algebra, except that SELECT in SQL retains duplicates

Consider a database that has the relation schema CR(StudentName, CourseName). An instance of the schema CR is as given below.

StudentName	CourseName
SA	CA
SA	CB
SA	CC
SB	CB
SB	CC
SC	CA
SC	CB
SC	CC
SD	CA
SD	CB
SD	CC
SD	CD
SE	CD
SE	CA
SE	CB
SF	CA
SF	CB
SF	CC

The following query is made on the database.

- $T1 \leftarrow \pi_{CourseName} \left(\sigma_{StudentName=SA}\left(CR\right)\right)$ $T2 \leftarrow CR \div T1$

The number of rows in T2 is

Consider the relations r(A, B) and s(B, C), where s. B is a primary key and r. B is a foreign key referencing s. B. Consider the query

$$Q:r\bowtie(\sigma_{B<5}(s))$$

Let LOJ denote the natural left outer-join operation. Assume that r and s contain no null values.

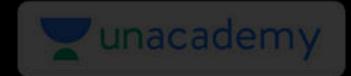
Which of the following is NOT equivalent to Q?

A.
$$\sigma_{B<5}(r\bowtie s)$$

C. $rLOJ(\sigma_{B<5}(s))$

B.
$$\sigma_{B<5}(r LOJ s)$$

D.
$$\sigma_{B<5}(r) LOJ s$$



Consider the following relations P(X, Y, Z), Q(X, Y, T) and R(Y, V).

Table: P			
X	Y	\mathbf{z}	
X1	Y1	Z1	
X1	Y1	Z2	
X2	Y2	Z2	
X2	Y4	Z4	

Table: Q			
X	Y	T	
X2	Y1	2	
X1	Y2	5	
X1	Y1	6	
X 3	Y3	1	

How many tuples will be returned by the following relational algebra query?

$$\Pi_x(\sigma_{(P.Y=R.Y\wedge R.V=V2))}(P\times R))-\Pi_x(\sigma_{(Q.Y=R.Y\wedge Q.T>2))}(Q\times R))$$



A relation r(A,B) in a relational database has 1200 tuples. The attribute A has integer values ranging from 6 to 20, and the attribute B has integer values ranging from 1 to 20. Assume that the attributes A and B are independently distributed.

The estimated number of tuples in the output of $\sigma_{(A>10)\vee(B=18)}(r)$ is ______.



The following relation records the age of 500 employees of a company, where empNo (indicating the employee number) is the key:

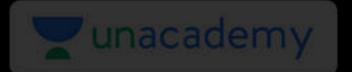
Consider the following relational algebra expression:

$$\Pi_{empNo}(empAge \bowtie_{(age>age1)} \rho_{empNo1,age1}(empAge))$$

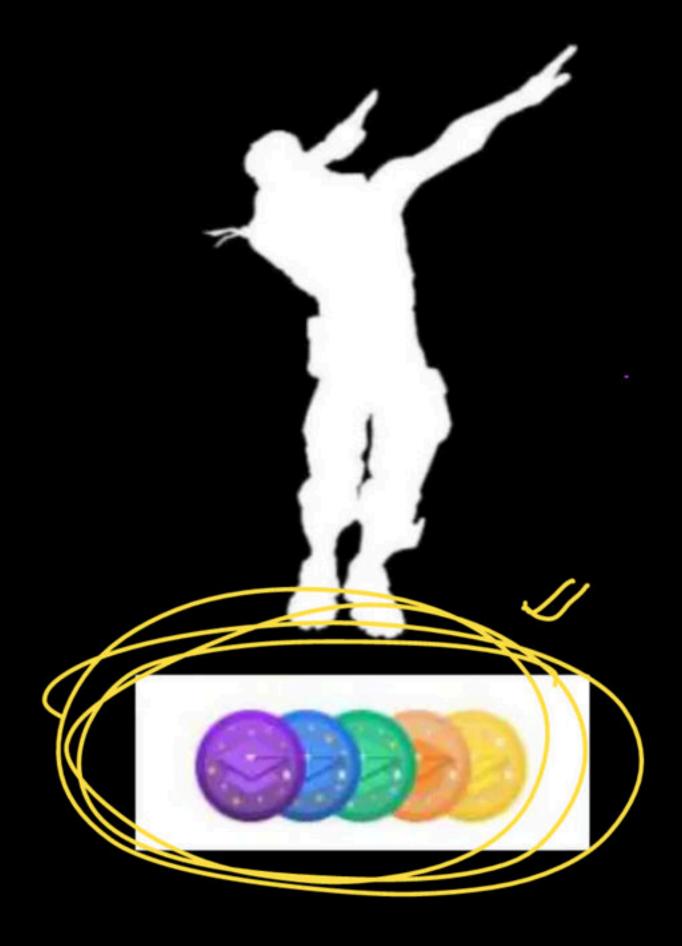
What does the above expression generate?

- A. Employee numbers of only those employees whose age is the maximum
- B. Employee numbers of only those employees whose age is more than the age of exactly one other employee
- C. Employee numbers of all employees whose age is not the minimum
- D. Employee numbers of all employees whose age is the minimum

Q.25	Consider the following three relations in a relational database.		
	Employee(eld, Name), Brand(bld, bName), Own(eld, bld)		
	Which of the following relational algebra expressions return the set of elds who own all the brands?		
(A)	$\Pi_{eld}(\Pi_{eld,bld}(Own)/\Pi_{bld}(Brand))$		
(B)	$\Pi_{eld}(Own) - \Pi_{eld}\left((\Pi_{eld}(Own) \times \Pi_{bld}(Brand)) - \Pi_{eld,bld}(Own)\right)$		
(C)	$\Pi_{eId}(\Pi_{eId,bId}(Own)/\Pi_{bId}(Own))$		
(D)	$\Pi_{\mathit{eId}}\left((\Pi_{\mathit{eId}}(\mathit{Own}) \times \Pi_{\mathit{bId}}(\mathit{Own})) / \Pi_{\mathit{bId}}(\mathit{Brand})\right)$		



Happy Learning.!





GATE & ESE

Win Big Rewards & Scholarship

Crack Unacademy Ultra Combat for

GATE & ESE Aspirants 2024/25

July 23 | 11 AM | Question = 40 MCQs | Time = 120 Minutes









Scholarship

Rank 1	90%
Rank 2-5	75%
Rank 6-10	60%
Rank 11-20	50%
Rank 21-50	40%
Rank 51-100	30%
Rank 100+	25%

Rank	
Rank 1	
Rank 2	
Rank 3	
Rank 4	
Rank 5	
Participat	ion Dr

ank		Reward
ınk 1	_	MacBook Air
ınk 2	-	■ Dell Laptop
ınk 3	-	Samsung Mobile
ınk 4	-	AGV - 10000 rs
ınk 5	-	AGV - 5000 rs
rticipation Prize	+	B AGV - 250 rs
ink 2 ink 3 ink 4 ink 5	-	Dell Laptop Samsung Mo AGV - 10000 AGV - 5000 r

500 participants on lucky draw

Enroll for Free

Use Code

*T&C apply, as available on the platform.