



# Doubt Clearing Session

Complete Course on Algorithm for GATE - CS & IT



$$T(n) = \begin{cases} 10 & \text{if } n=10 \\ T(n/2) + n & \text{if } n > 1 \end{cases}$$

$$\frac{n}{2^k} = 10, \quad n = 2^k \cdot 10, \quad k = \log_2 n$$

$$\begin{aligned} T(n) &= T(n/2^1) + \frac{n}{2^0} \\ &\Downarrow \\ T(n/2^2) + \frac{n}{2^1} + \frac{n}{2^0} \\ &\Downarrow \\ T(n/2^3) + \frac{n}{2^2} + \frac{n}{2^1} + \frac{n}{2^0} \end{aligned}$$

$$\begin{aligned} 2^2 \cdot 5n \\ 2^2 \cdot 5n \end{aligned}$$

$$\left\{ \begin{aligned} k &= \log_2 n \\ k &= \log_2 n \end{aligned} \right.$$

$$= T(n/2^k) + \frac{n}{2^{k-1}} + \frac{n}{2^{k-2}} + \dots + \frac{n}{2^1} + \frac{n}{2^0}$$

$$= T(1) + n \left[ \left(\frac{1}{2}\right)^0 + \left(\frac{1}{2}\right)^1 + \dots + \left(\frac{1}{2}\right)^{k-1} \right]$$

$$\begin{aligned} &= 10 + n \left[ \frac{1 - \left(\frac{1}{2}\right)^k}{1 - \frac{1}{2}} \right] \\ &= 10 + n \left[ \frac{(1-0)}{\frac{1}{2}} \right] \end{aligned}$$

$$\begin{aligned} 10 + n \left[ \frac{2}{1} \right] &\Rightarrow O(n) \\ &\Rightarrow \Theta(n) \\ &\Rightarrow O(n) \end{aligned}$$



$$T(n) = \begin{cases} 2 & \text{if } n=2 \\ \sqrt{n} T(\sqrt{n}) + n & \text{if } n > 2 \end{cases}$$

$$T(100) = 10T(10) + 100$$

$$T(n) = n^{1/2} T(n^{1/2}) + n$$

$$T(n) = n^{1/2} T(n^{1/2}) + \underbrace{n + n + n}_3$$

$$= \frac{1}{2} \left[ n^{1/2} T(n^{1/2}) + n \right] + n$$

$$= n^{1/2^k} T(n^{1/2^k}) + k \cdot n$$

$$= n^{1 - \frac{1}{2^k}} T(n^{1/2^k}) + k \cdot n$$

$$= \frac{n^{1/2^k}}{n^{1/2^k}} T\left(\frac{1}{n^{1/2^k}}\right) + k \cdot n$$

$$= \frac{n}{2} T(2) + n \cdot \frac{1}{2} \cdot 2$$

$$= n^{3/2} \left[ n^{1/2} T(n^{1/2}) + n \right] + n + n$$



$$= \frac{n}{2} \cdot 2 + n \cdot \log_2 n$$

$$= n + n \cdot \log_2 n = \Theta(n \log_2 n)$$

$$\frac{1}{n 2^k}$$

$$= 2$$

$$\frac{1}{2^k} \log_2 n = 1$$

$$\log_2 n = 2^k$$

$$\log_2 \log_2 n = k$$

$$2^k$$

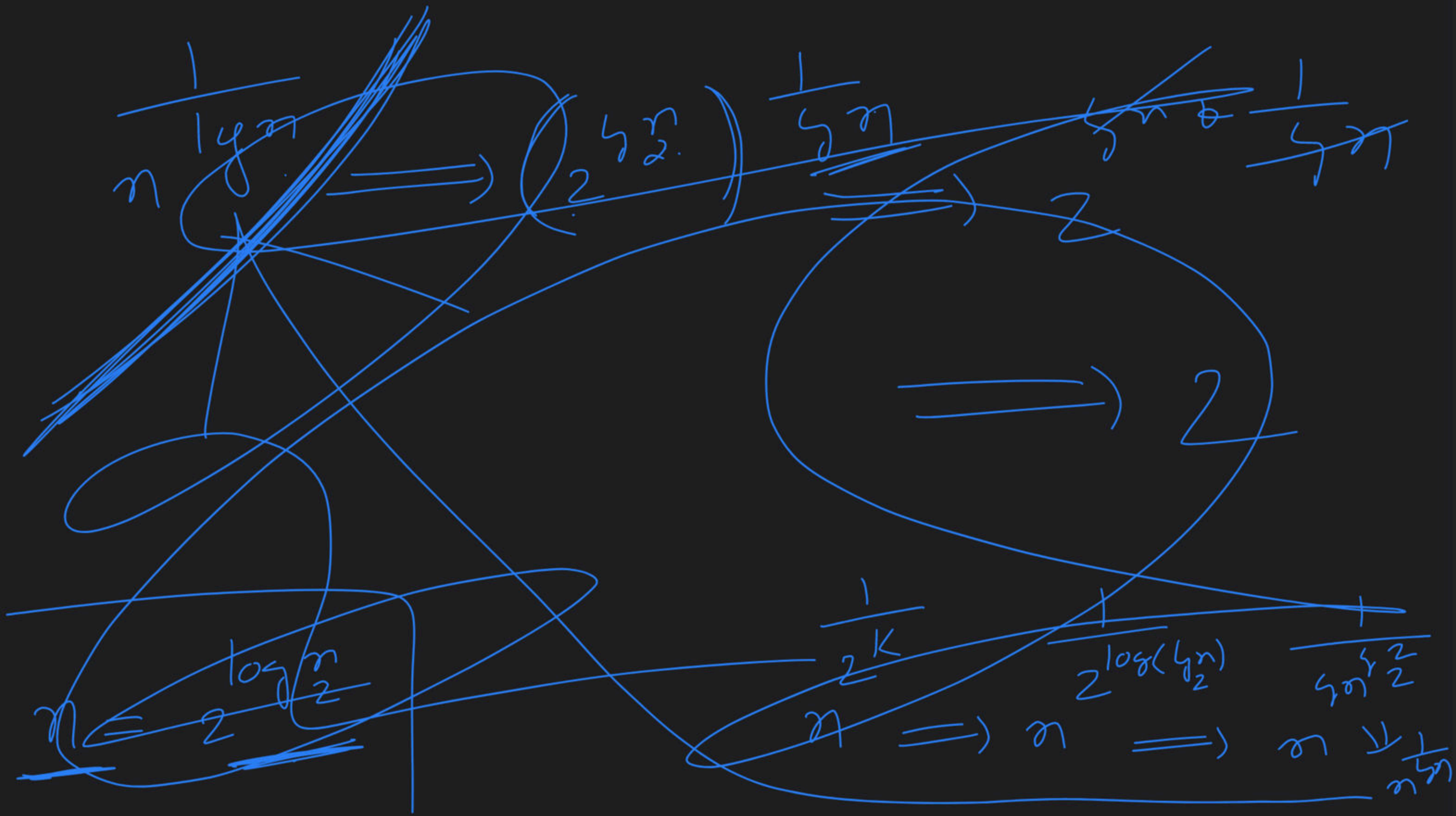
$$\log_2(\log_2 n)$$

$$\frac{1}{\log_2 n}$$

$$\log_2 n$$

$$\log_2 n$$







$$T(n) = \begin{cases} 1 & \text{if } n = \underline{1} \\ 2T(n-1) + n & \text{if } n > 1 \end{cases}$$

$$T(100) = 2T(99) + 100$$

$$\begin{aligned} T(n) &= 2T(n-1) + n \\ &= 2[2T(n-2) + (n-1)] + n \end{aligned}$$

$$= 2^2 T(n-2) + 2(n-1) + n$$

$$= 2^2 [2T(n-3) + (n-2)] + 2(n-1) + n$$

$$= 2^3 T(n-3) + 2^2(n-2) + 2^1(n-1) + 2^0(n-0)$$

$$= \frac{n-1}{2} T(n-(n-1)) + \overset{\substack{K^2 \\ \downarrow \\ 1}}{2^{n-2}}(n-(n-2)) + 2^{n-3}(n-(n-3)) + \dots + 2^1(n-1) + 2^0(n-0)$$

$$n - K = 1$$

$$n - 1 = K$$



$$\checkmark = 2^{n-1} \cdot 1 + 2^{n-2} \cdot 2 + 2^{n-3} \cdot 3 + \dots + 2^1 \cdot (n-1) + 2^0 \cdot (n-0)$$

$$T(n) = 2^0(n-0) + 2^1(n-1) + 2^2(n-2) + \dots + 2^{n-2} \cdot 2 + 2^{n-1} \cdot 1$$

$$T(n) = \frac{2^0(n-0)}{0} + \frac{2^1(n-1)}{2^1(n-0)} + \frac{2^2(n-2)}{2^2(n-1)} + \dots + \frac{2^{n-2} \cdot 2}{2^{n-2} \cdot 3} + \frac{2^{n-1} \cdot 1}{2^{n-1} \cdot 2} + 0$$

$$2T(n) = 0 + 2^1(n-0) + 2^2(n-1) + \dots + 2^{n-2} \cdot 3 + 2^{n-1} \cdot 2 + 2^n \cdot 1$$

$$T(n) - 2T(n) = 2^0(n-0) - 2^1 - 2^2 - 2^3 - 2^4 - \dots - 2^{n-2} - 2^{n-1} - 2^n$$

$$-T(n) = n - [2^1 + 2^2 + 2^3 + \dots + 2^n] \Rightarrow \frac{2(2^n - 1)}{2 - 1} - n$$

$$T(n) = -n + [2^1 + 2^2 + \dots + 2^n] \Rightarrow 2^{n+1} - 2 - n$$

$$T(n) = 2^{n+1} - 2 - n \Rightarrow \Theta(2^n)$$

$$T(n) = 2T(n/2) + n \log n$$



$$\begin{aligned}
 T(n) &= 10 \cdot 2^9 + 20 \cdot 2^{11} + 30 \cdot 2^{13} + 40 \cdot 2^{15} + \dots \\
 2 \cdot T(n) &= 0 + 10 \cdot 2^{11} + 20 \cdot 2^{13} + 30 \cdot 2^{15} + \dots
 \end{aligned}$$

$$T(n) - 2 \cdot T(n) = 10 \cdot 2^9 + 10 \cdot 2^{11} + 10 \cdot 2^{13} + 10 \cdot 2^{15}$$

$$-3T(n) = 10 \left[ 2^9 + 2^{11} + 2^{13} + 2^{15} + \dots \right]$$

$$T(n) = 10 \left[ \dots \right]$$

-3



$$\frac{a(r^n - 1)}{r - 1} \quad \text{if } r > 1$$

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$$\frac{a(1 - r^n)}{1 - r} \quad \text{if } r < 1$$

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$$\left(\frac{1}{2}\right)^2 = (0.25)^2$$

0.0625