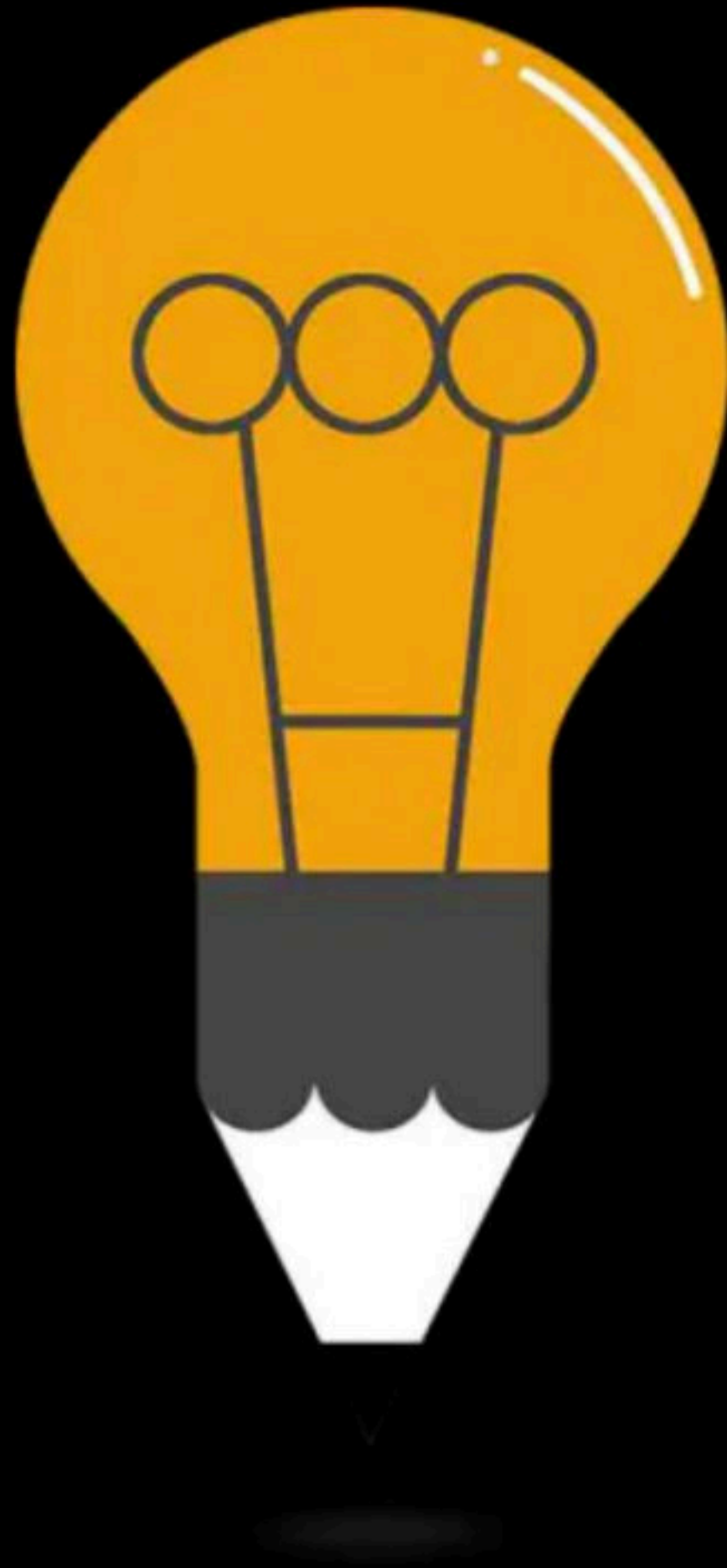




# Relational Algebra: Part I

Complete Course on Database Management System



# DBMS

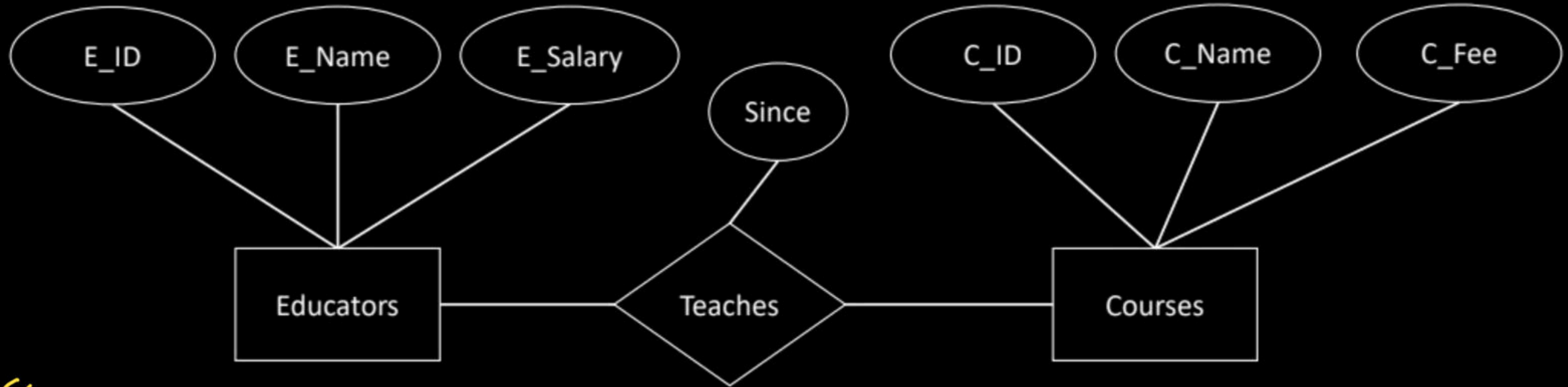
## ER to Relational Model & RA

By: **Vishvadeep Gothi**

# ER Diagram to Relational Model

1 Entity set  $\Rightarrow$  1 table

# Many to Many Relationship



✓  
option 1:-

3 tables →

- ① Educators
- ② Courses
- ③ Teachers

Best option



## Educators

<u>Eid</u>	Ename	Salary
------------	-------	--------

## Courses

<u>Cid</u>	Cname	Cfee
------------	-------	------

## Teaches

<u>Eid</u>	<u>Cid</u>	Since
1	C1	2010
1	C2	2015
1	C3	2012
2	C2	2015
3	C3	2020

Teaches.Eid  $\Rightarrow$  fk referring to Educators.Eid

Teaches.cid  $\Rightarrow$  fk referring to Courses.cid

 Educators

$Eid \rightarrow ename \text{ salary}$

Courses

$Cid \rightarrow cname \text{ cfee}$

Teaches

$Cid \text{ } Eid \rightarrow \text{since}$

all 3 relations are in BCNF

option 2:-

keep relationship info along with an entity set.

→ relationship towards Educators ⇒ 2 Tables

- Educators

<u>Eid</u>	Ename	Esalary	<u>Cid</u>	Since
<u>1</u>	—	—	C1	2010
<u>1</u>	—	—	C2	2015
<u>1</u>	—	—	C3	2012
2	—	—	C2	2015
3	—	—	C3	2020

Courses

<u>Cid</u>	Cname	Cfee
------------	-------	------



Primary key  $\Rightarrow$  Eid cid

but partial dependency  $\Rightarrow$  Eid  $\rightarrow$  Ename Esalary

after removing partial dependency  $\Rightarrow$  3 tables same  
as option 1

take relationship info<sup>n</sup> towards courses entity set  $\Rightarrow$  2 tables

Teachers

<u>Eid</u>	Ename	Esalary
------------	-------	---------

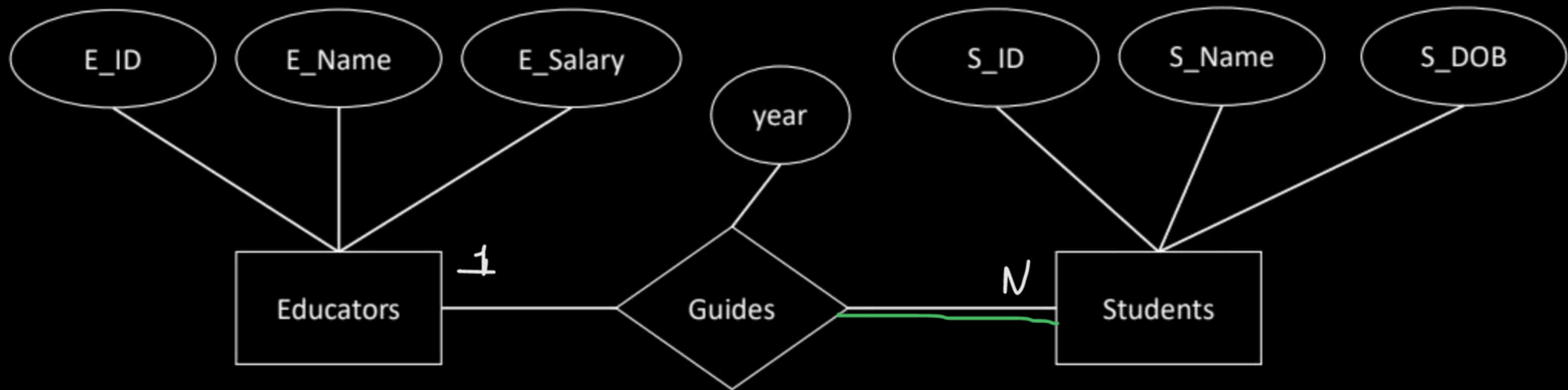
Courses

<u>Cid</u>	Cname	Cfee	<u>Sid</u>	since
------------	-------	------	------------	-------

Problem in Courses  $\Rightarrow$  partial dep.

$Cid \rightarrow Cname \ Cfee$

# One to Many Relationship



Best sol<sup>n</sup>:-

2 tables  $\Rightarrow$  Take relationship info<sup>n</sup> towards  
many side entity set



Educators

Eid Ename Esalary

Students

Sid sname slob guide eid year

not null

S1 Rohit 27 Jan E1 2022

S2 Govinda 1 Feb E1 2022

S3 Riya 23 Mar E2 2023

$\Rightarrow$

Eid  $\rightarrow$  Ename Esalary

$\Rightarrow$  Sid  $\rightarrow$  Name, Job, Guide, Year

options:-

Educators

Eid Ename Esalary Sid year

E1 ——— S1 2022

E1 ——— " ——— S2 2022

Eid not p.k. 4  
redundancy.



### option 3:-

3 tables

Editors

<u>Eid</u>	Ename	Salary
------------	-------	--------

Student

<u>Sid</u>	Sname	Sdob
------------	-------	------

Guides

<u>Sid</u>	Eid	year
------------	-----	------

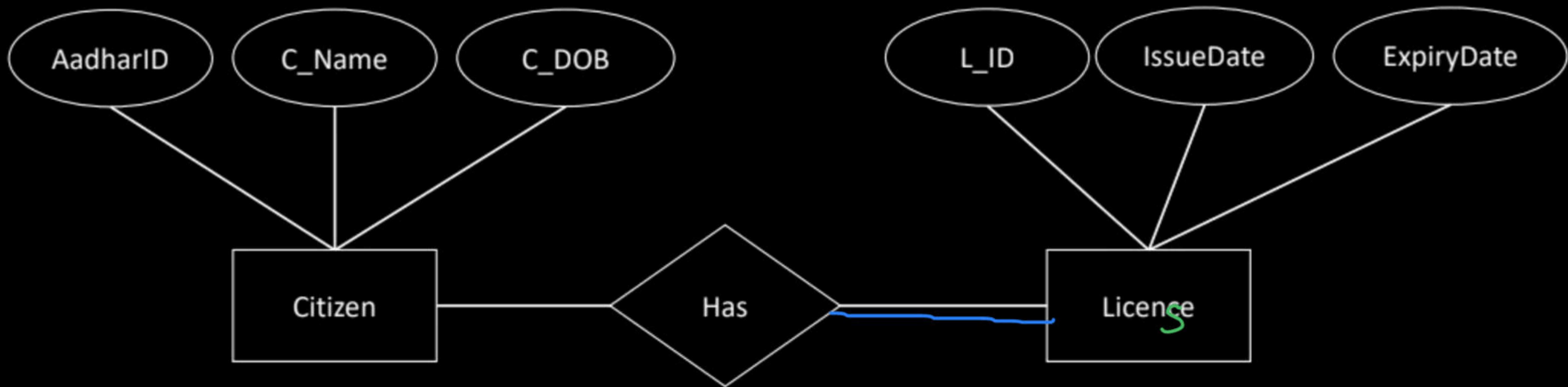
S1 E1

S2 E1



Sid is pk so better to  
keep it in same table

# One to One Relationship



option 1:- keep relationship info<sup>n</sup> towards citizen

Citizen

Aadharid C-name C-dob L-id

License

Lid issuedate Expirydate

Option 2: keep relationship info towards license side.

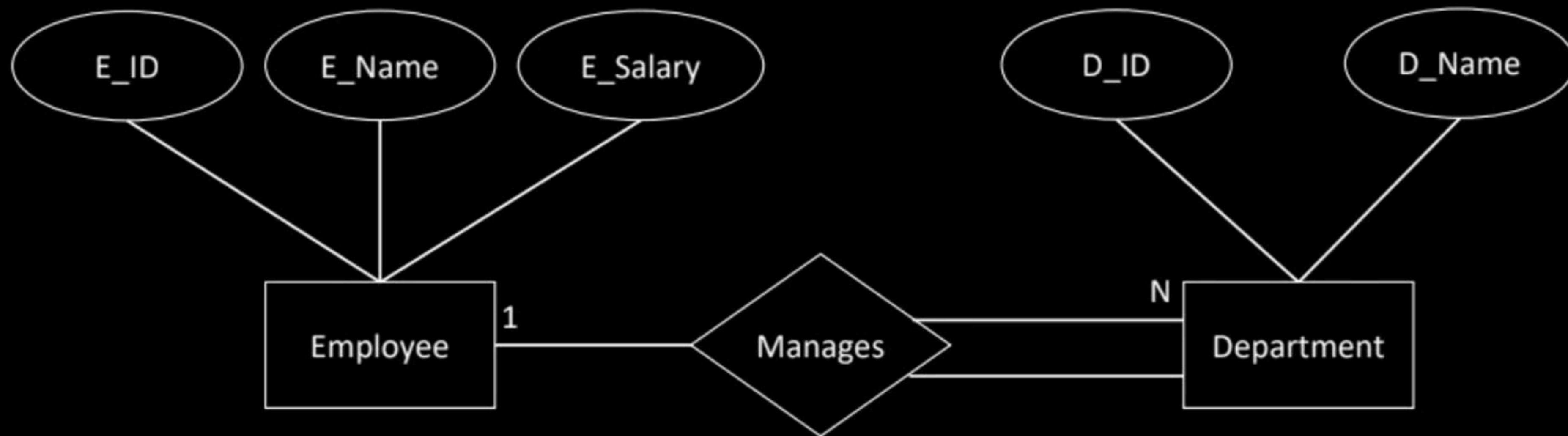
citizen			license		
Aadharid	C-name	C-dob	lid	issuere date	Expiry date

not NULL  
↓  
Aadharid

option 2 is better because license has total participation in relationship.



# Participation Constraint



Employee  
Eid Ename Esalary

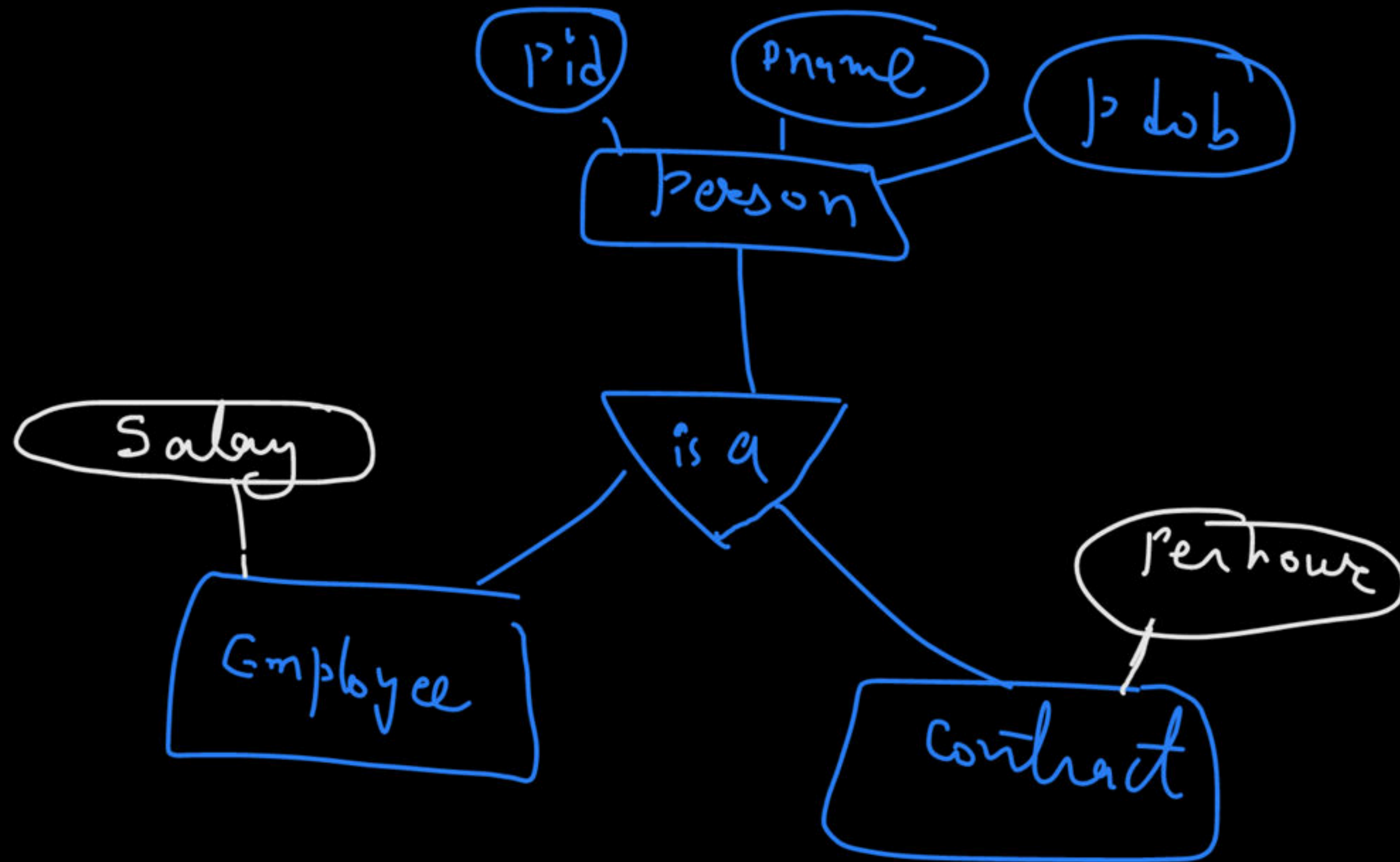
Department  
Did Dname Eid of manager

Not NULL



How to control participation in many-to-many  $\Rightarrow$  ?

# Generalization & Specialization



option 1:- 3 tables

Person (pid, pname, pdob)

Employee (pid, salary)

Contract (pid, per hour)



overlapping

option 2:-

2 tables

Employee (pid, pname, pdob, salary)

Contract (pid, pname, pdob, per hour)

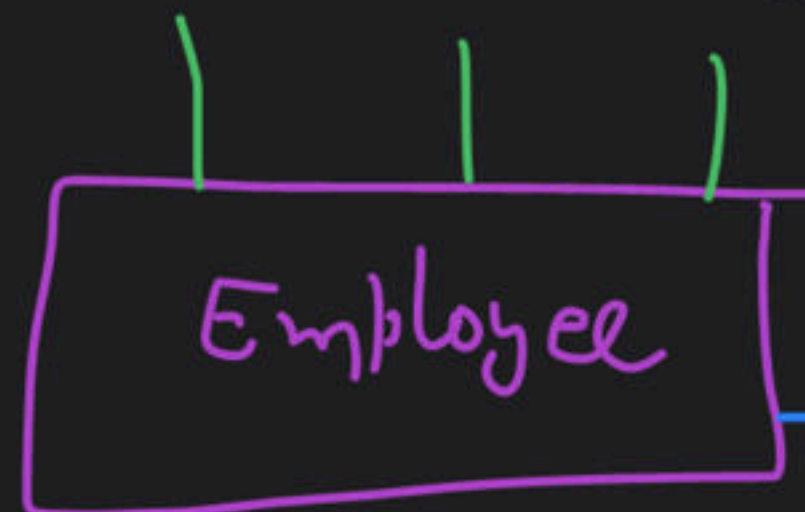


disjoint



# weak entity set

Eid Ename Esalary



dname djob



2 tables :-

Employee  
Eid Ename Esalary

dependents  
Eid dname djob



# Relational Algebra

Relational algebra is a procedural query language, which takes instances of relations as input and yields instances of relations as output

The main application of relational algebra is to provide a theoretical foundation for relational databases, particularly for query languages.

# Relational Algebra

## Basics Operators:

1. Select ( $\sigma$ )
2. Project ( $\Pi$ )
3. Union ( $\cup$ )
4. Set Difference ( $-$ )
5. Intersection ( $\cap$ )
6. Cartesian Product ( $\times$ )
7. Rename ( $\rho$ )
8. Division ( $\div$ )

9. natural join ( $\bowtie$ )
10. left outer join ( $\ltimes$ )
11. Right - " ( $\rtimes$ )

12. full outer join



select \* from table  $\Leftarrow$  SQL

R.A.  $\Rightarrow$   $\pi$  (table)

$\Downarrow$

(table)

# Select

The SELECT operator chooses those tuples in the output that satisfy the specified condition.

$\sigma_{\langle \text{selection-condition} \rangle} (\text{Relation name})$



sailors

## Select

Eid	Ename	Rating	Age
101	Richa	7	24
105	Rohan	9	20
120	Mahesh	8	26
145	Abhishek	10	29

select \* from sailors  
where rating > 8

To select those tuples of sailors relation  
where rating is greater than 8

$\sigma_{\text{rating} > 8}(\text{sailors})$

find all sailors who have rating greater than 8 but age less than 25.

$\sigma_{\text{rating} > 8 \wedge \text{age} < 25} (\text{sailors})$

$\wedge \Rightarrow$  And

$\vee \Rightarrow$  OR

# Select

Comparison operators that can be used in selection conditions are:

<, <=, =, ≠, >=, >

The select operation is commutative i.e.

$$\sigma_{\langle \text{cond } 1 \rangle} (\sigma_{\langle \text{cond } 2 \rangle} (R)) = \sigma_{\langle \text{cond } 2 \rangle} (\sigma_{\langle \text{cond } 1 \rangle} (R)) = \sigma_{\text{cond } 1 \wedge \text{cond } 2} (R)$$

$$\sigma_{\text{rating} > 8} (\sigma_{\text{age} < 25} (R)) = \sigma_{\text{age} < 25} (\sigma_{\text{rating} > 8} (R))$$



# Project

The project operation is used to choose certain columns from the table and trashes out the other attribute fields

A projection is a unary operation written as  $\Pi_{A_1, A_2, \dots, A_n}(r)$

# Project

Fetch only Ename and salary of all employees from employee relation

Eid	Ename	Salary	Dno	Sex
101	Raman	30,000	1	M
102	Sneha	20,000	1	F
103	Maya	20,000	2	F
104	Ranjith	20,000	2	M
105	Mahesh	15,000	3	M

# Project

The project operation results in a set of a distinct tuple as the Project operation removes duplicate tuples



# Project

Write a relation algebra statement to find name of all such employees from department no 2 whose salary is greater than 17000

Eid	Ename	Salary	Dno	Sex
101	Raman	30,000	1	M
102	Sneha	20,000	1	F
103	Maya	20,000	2	F
104	Ranjith	20,000	2	M
105	Mahesh	15,000	3	M

# Set Operations

- 1) Union
- 2) Intersection
- 3) Set-difference

# Set Operations

The two relations on which set operations are implemented upon must necessarily have similar data types of tuples. This condition is called type or union compatibility.



# Example

Eid	Ename	Age	Rating
20	Somya	24.0	7
30	Rahul	25.0	8
40	Ranjith	24.0	9
50	Yashvi	23.0	10
60	Sonam	27.0	8

Eid	Ename	Age	Rating
30	Rahul	25.0	8
35	Satyam	24.0	9
50	Yashvi	23.0	10
60	Sonam	27.0	8

# Question

Consider 2 relations Students(rno, sname, dob) and Employess(eld, ename, salary)

Write a relational algebra statement for corresponding SQL Query:

Select distinct sname from Students where dob='27-10-1988' Union Select distinct ename where salary>15000

# Set Operations

Union and intersection are commutative and associative



# Question

Consider 2 relations  $R1(x,y)$  and  $R2(x,y)$ .  $R1$  and  $R2$  contains all Not NULL values.  
Will the following 2 statements be equivalent or not?

$$\pi_x(R1 \cup R2)$$

$$\pi_x(R1) \cup \pi_x(R2)$$

# Question

Consider 2 relations  $R1(x,y)$  and  $R2(x,y)$ .  $R1$  and  $R2$  contains all Not NULL values.  
Will the following 2 statements be equivalent or not?

$$\pi_x(R1 \cap R2)$$

$$\pi_x(R1) \cap \pi_x(R2)$$

# Happy Learning.!



1 • Asked by Saloni

yeh ques nhi samaj aaya sir

$$z \leq w$$

$$w \leq x$$

GATE-2001

Consider a relation geq which represents "greater than or equal to", that is,  $(x, y) \in \text{geq}$  only if  $y \geq x$ .

```
create table geq
(
  ib integer not null,
  ub integer not null,
  primary key ib,
  foreign key (ub) references geq on delete cascade
);
```

Which of the following is possible if tuple  $(x, y)$  is deleted?

A. A tuple  $(z, w)$  with  $z > y$  is deleted  
 B. A tuple  $(z, w)$  with  $z > x$  is deleted  
 C. A tuple  $(z, w)$  with  $w < x$  is deleted  
 D. The deletion of  $(x, y)$  is prohibited

Handwritten notes on the video screen:

- A diagram showing a table with columns  $ib$  and  $ub$ . The  $ib$  column has values  $z, w, x$  and the  $ub$  column has values  $y, w$ . A red arrow points from  $y$  to  $w$  in the  $ub$  column.
- Handwritten inequalities:  $x \leq y$ ,  $w \leq x$ ,  $z \leq w$ ,  $z \leq y$ ,  $y \leq w$ .

Video player controls: 00:28:26 / 01:03:36

$$(ib) \geq (ub)$$



$$z$$

$$w$$

$$(z, w)$$

$$x < y$$



▲ 1 • Asked by Rishabh

If the relation is not in BCNF then all the attributes of every Dependency not having a superkey in LHS is kept in a separate relation.

**BCNF (Boyce Codd Normal Form)**

$R(A, B, C, D)$   
 $FDs = \{$   
 $AB \rightarrow CD,$   
 $D \rightarrow B$   
 $\}$

no any partial or transitive dependency. hence already in 3NF.

check for BCNF:-

$AB \rightarrow CD$  ✓ LHS AB is key

$D \rightarrow B$  ✗ LHS D is not key

↓

Not in BCNF

Remove B from original table & keep it in another table with D.

$R(A, C, D)$      $R_2(D, B)$   
 $A \rightarrow CD$      $D \rightarrow B$

keys  $\Rightarrow AB, AD$