

Parsing-VIII

Complete Course on Compiler Design

LR(0)-PT construction

1. Closure() 2. Goto()



ex
G: $S \rightarrow AA$
 $A \rightarrow aA/b$

$S \rightarrow S$
 $G' = S \rightarrow AA$
 $A \rightarrow aA/b$

Augmented grammar

Augmented Grammar

$S \rightarrow \cdot xy2$
 $S \rightarrow x \cdot y2$
 $S \rightarrow xy \cdot 2$
 $S \rightarrow xy2 \cdot$

LR(0)-item
Completed LR(0)-item
(or) Final LR(0)-item
"reduced"

Closure()

G: $S \rightarrow AA$
 $A \rightarrow aA/b$

ex

$Closure(S \rightarrow A \cdot A)$

① $S \rightarrow A \cdot \underline{A}$
② $A \rightarrow \cdot \underline{a}A$
 $\quad \quad \cdot \underline{b}$

ex
 $Closure(S \rightarrow \cdot S)$

① $S \rightarrow \cdot \underline{S}$
② $S \rightarrow \cdot \underline{A}A$
③ $A \rightarrow \cdot \underline{a}A$
 $\quad \quad \cdot \underline{b}$

ex

~~ex~~ $clone(A \rightarrow \cdot aA)$

\Downarrow

① $A \rightarrow \cdot \underline{a}A$

~~ex~~

ex $clone(S \rightarrow \cdot AA)$

\Downarrow

① $S \rightarrow \cdot AA$
② $A \rightarrow \cdot \underline{a}A$
No

Definition

$clone(I) =$ ① add - I

② if $A \rightarrow BC \cdot DE$ is in I & $D \rightarrow FH$ is in 4 item add

$D \rightarrow \cdot FH$ to $clone(I)$

③ repeat 2nd step for every newly added LR(0)-item.

goto(L)

G: $S \rightarrow AA$

$A \rightarrow aA/b$

\Downarrow

G': $S \rightarrow S$

$S \rightarrow AA$

$A \rightarrow aA/L$

ex

goto($S \rightarrow \cdot AA$, A)

\Downarrow

(1) $S \rightarrow A \cdot \underline{A}$

(2) $A \rightarrow \cdot aA$
 $\quad \quad \cdot b$

ex

goto($A \rightarrow \cdot a \cdot A$, A)

\Downarrow

(1) $A \rightarrow aA \cdot$

Def

goto(I_i , x)

(1) add- I by moving add after- x .

(2) apply closure (1st step)

ex

goto($A \rightarrow \cdot aA$, a)

\Downarrow

(1) $A \rightarrow a \cdot \underline{A}$

(2) $A \rightarrow \cdot aA$
 $\quad \quad \cdot b$

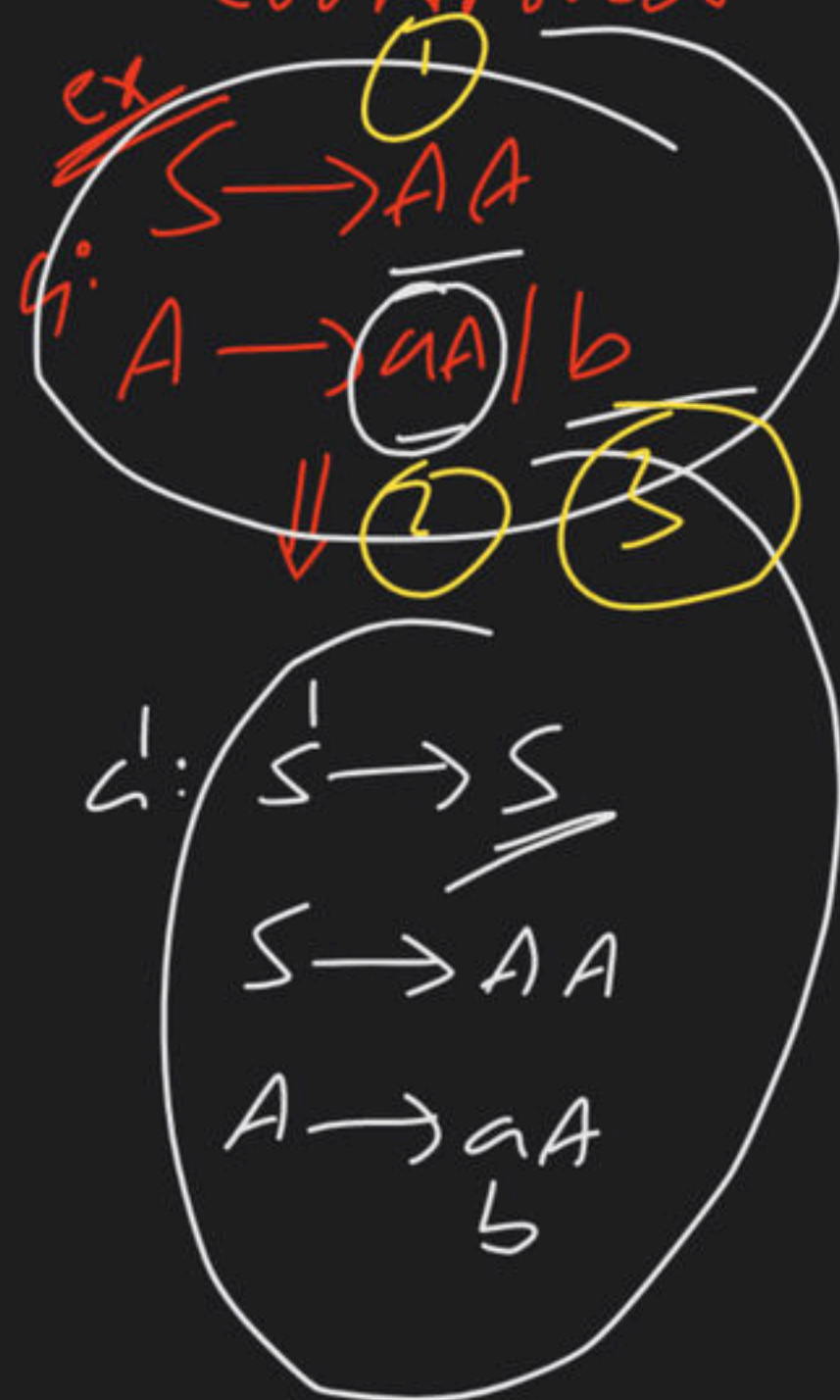
Algo

- ① $I_0 = \text{Closure (Augmented LR(0)-item)}$
- ② using I_0 construct DFA
- ③ convert DFA into LR(0)-PT .

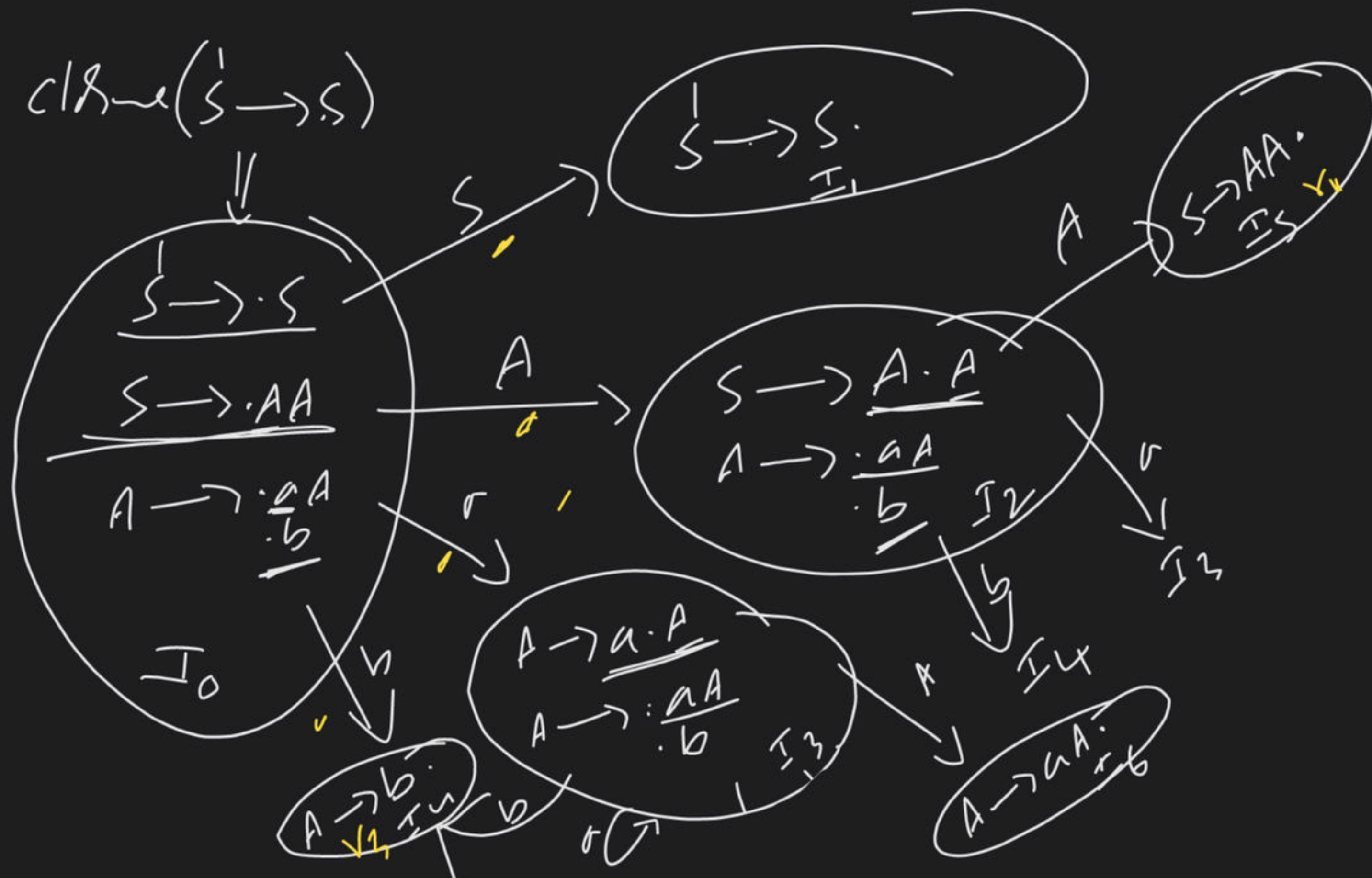
	action			goto	
	a	b	\$	S	A
0	s ₃	s ₄		1	2
1			acc		
2	s ₃	s ₄			5
3	s ₃	s ₄			6
4	r ₂	r ₇	r ₃		
5	r ₁	r ₁	r ₁		
6	r ₂	r ₂	r ₂		

LR(0)-PT

Construct LR(0)-PT for the following grammar



closure($S \rightarrow S$)



ex

$S \rightarrow (L) / a$

~~LLR~~ Left LR ✓

$LL(1) \times$

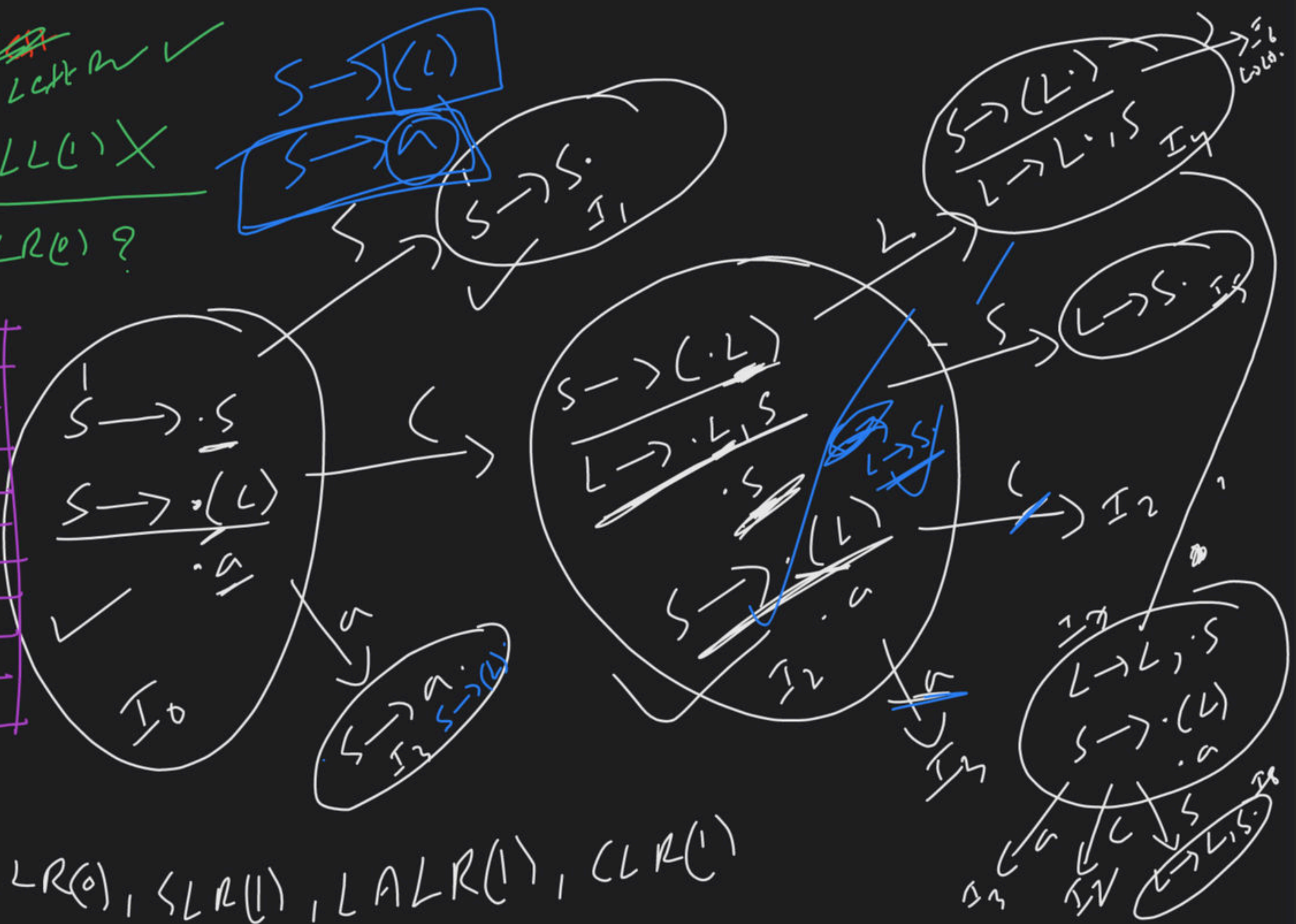
$LR(0) ?$

$L \rightarrow L, S / S$

	action						state	
	1	2	3	4	5	6	S	L
0			$S_3 S_2$				1	
1					a			
2	δ_1	S_2	S_2	γ_1	γ_1		5	4
3	γ_2	γ_2	γ_2	γ_1	γ_2			
4	S_3				S_6			
5	γ_1	γ_1	γ_1	γ_1	γ_1			
6	γ_1	γ_1	γ_1	γ_1	γ_1			
7		S_3	S_2				8	
8	γ_1	γ_1	γ_1	γ_1	γ_1			

no-conflict

$\Rightarrow LR(0), SLR(1), LALR(1), CLR(1)$



RR
SR
SR