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# Parsing-V

Complete Course on Compiler Design

**Subbarao Lingamgunta** • Lesson 8 • Aug 24, 2023

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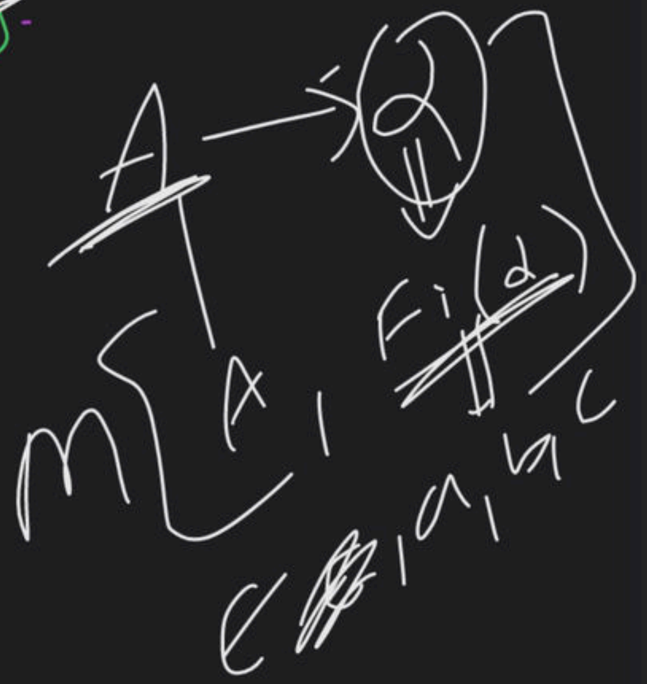
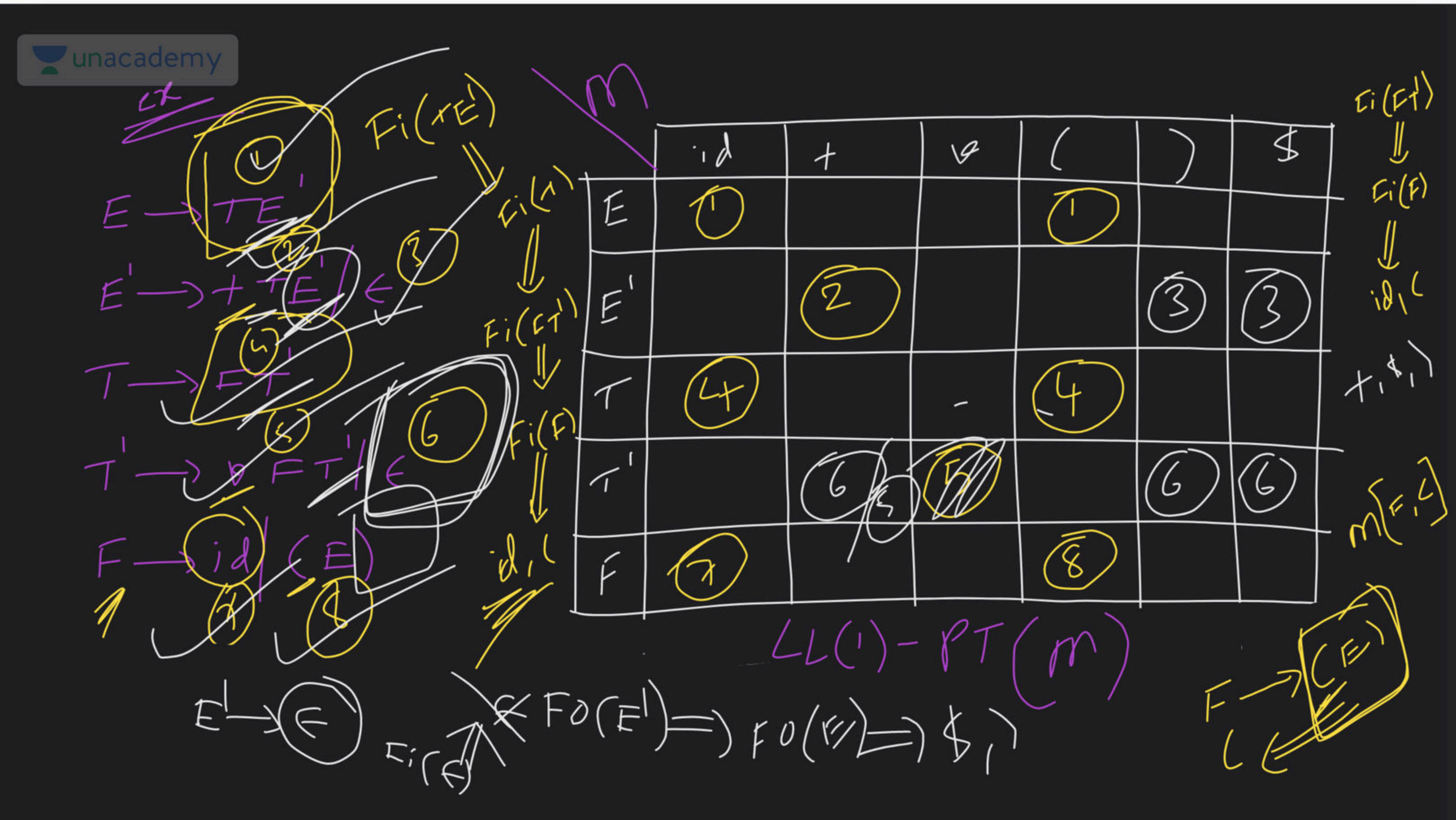
Algo

for each production $A \to \alpha$
repeate following steps:

(1) add $A \to \alpha$ under $M[A, d]$.
$\forall d \in First(\alpha)$

(2) if $First(\alpha)$ contain $\epsilon$ then
add $A \to \alpha$ under $M[A, e]$
$\forall e \in Follow(A)$.

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CFG
E → TE'
E' → +TE' | ∈
T → FT'
T' → *FT' | ∈
F → id | (E)
Fi(TE')
Fi(T)
Fi(FT')
Fi(F)
id, (
M
id
+
*
(
)
$
E
E'
T
T'
F
LL(1) - PT (M)
Fi(FT')
Fi(F)
id, (
+, $, )
M[F, (]
F → (E)
(
E' → ∈
Fi(∈)
Fo(E')
Fo(∈)
$, )

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ex

$S \rightarrow (L) \mid a$ ① ②

$L \rightarrow SL'$ ③

$L' \rightarrow ,SL' \mid \epsilon$ ④ ⑤

| | , | a | ( | ) | $ |
|---|---|---|---|---|---|
| S | | ② | ① | | |
| L | | ③ | ③ | | |
| L' | ④ | | | ⑤ | |

$Fi(\epsilon) = \{\epsilon\}$

$Fo(L') \Rightarrow$ LL(1) - PT

$Fi(SL') = Fi(S)$

$L, a$

$\alpha_1 \rightarrow \beta$

$Fi(\beta_1) = Fo(\alpha_1)$

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Ex

S → a B D h

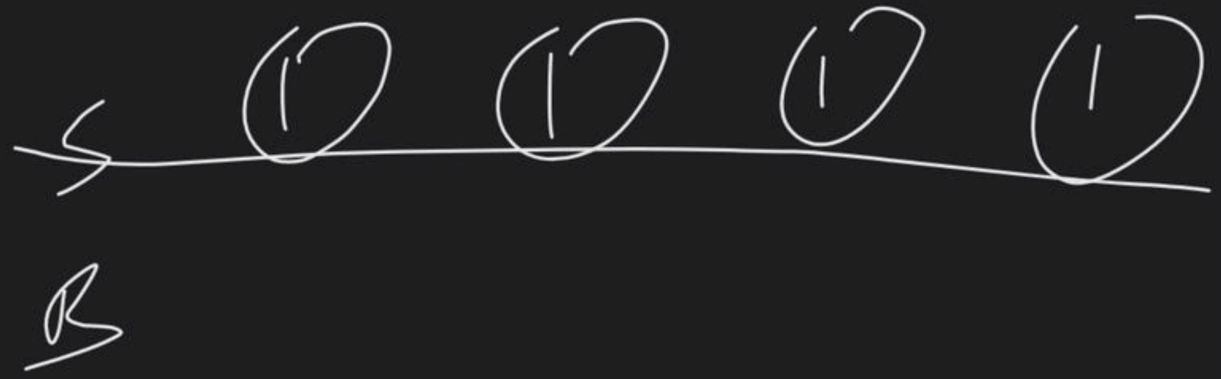
B → c C | ∈

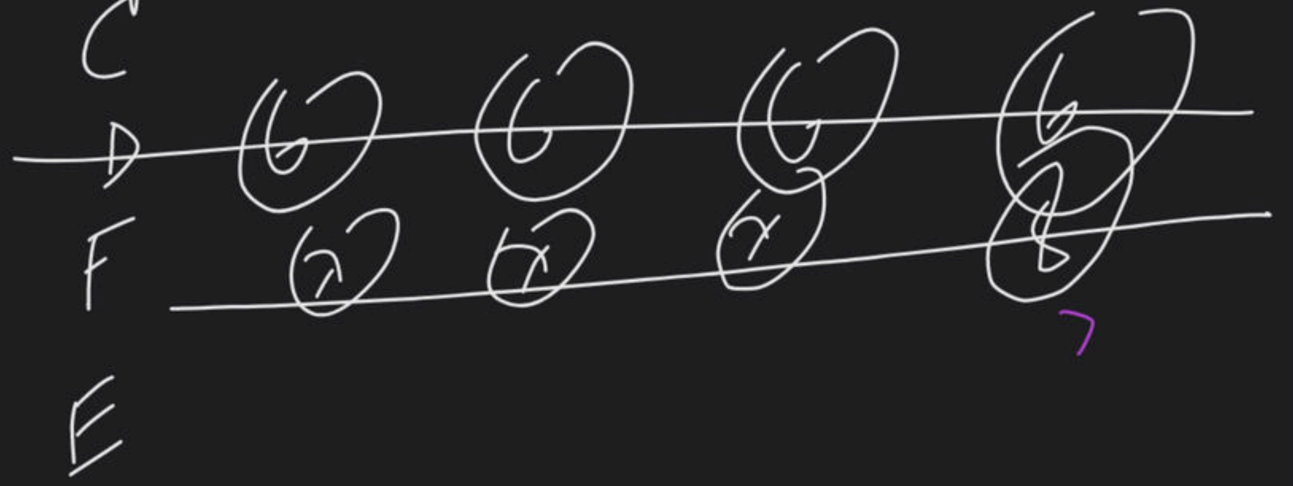
C → b C | ∈

D → E E

F → f | ∈

E → f | ∈



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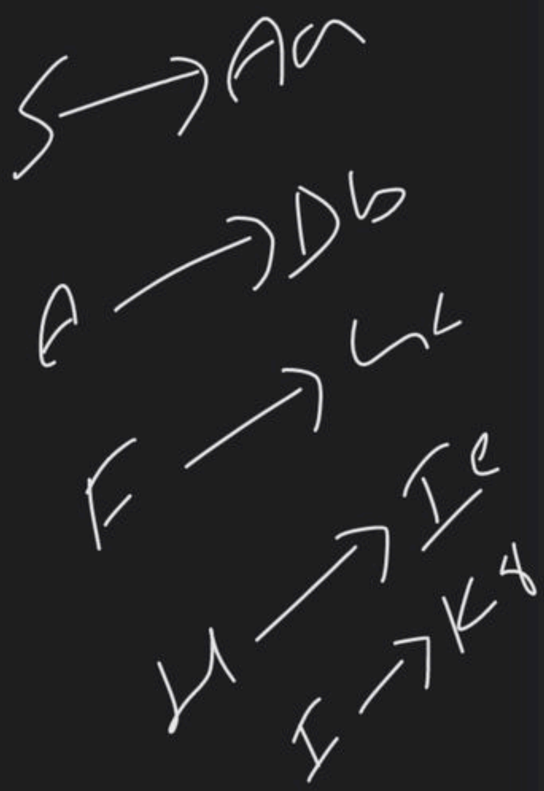

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$S \rightarrow aSbS \mid bSaS \mid \epsilon$

| | Fi() | Fo() |
|---|---|---|
| S | a, b, ε | $, b, a |

$S \rightarrow aAb \mid BbBa$

$A \rightarrow \epsilon$

$B \rightarrow \epsilon$

| | Fi() | Fo() |
|---|---|---|
| S | a, b | $ |
| A | ε | a, b |
| B | ε | b, a |