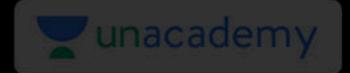
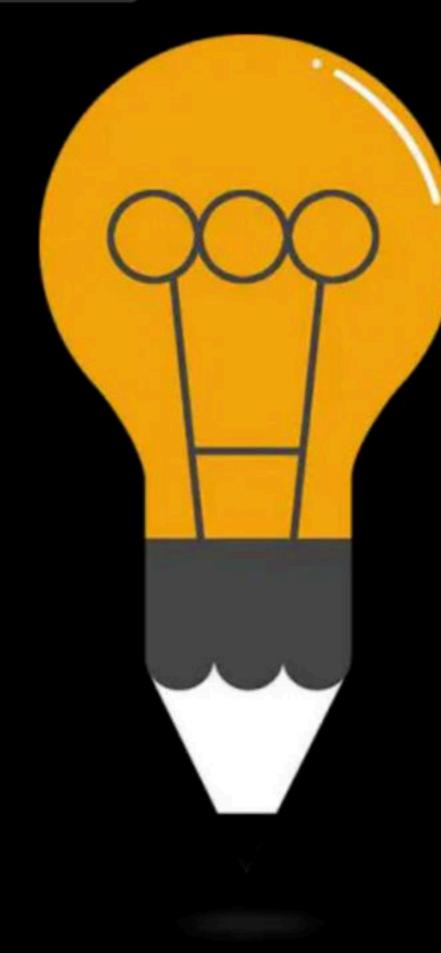




#### CPU Scheduling Algo: FCFS & SJF

Comprehensive Course on Operating System for GATE - 2024/25





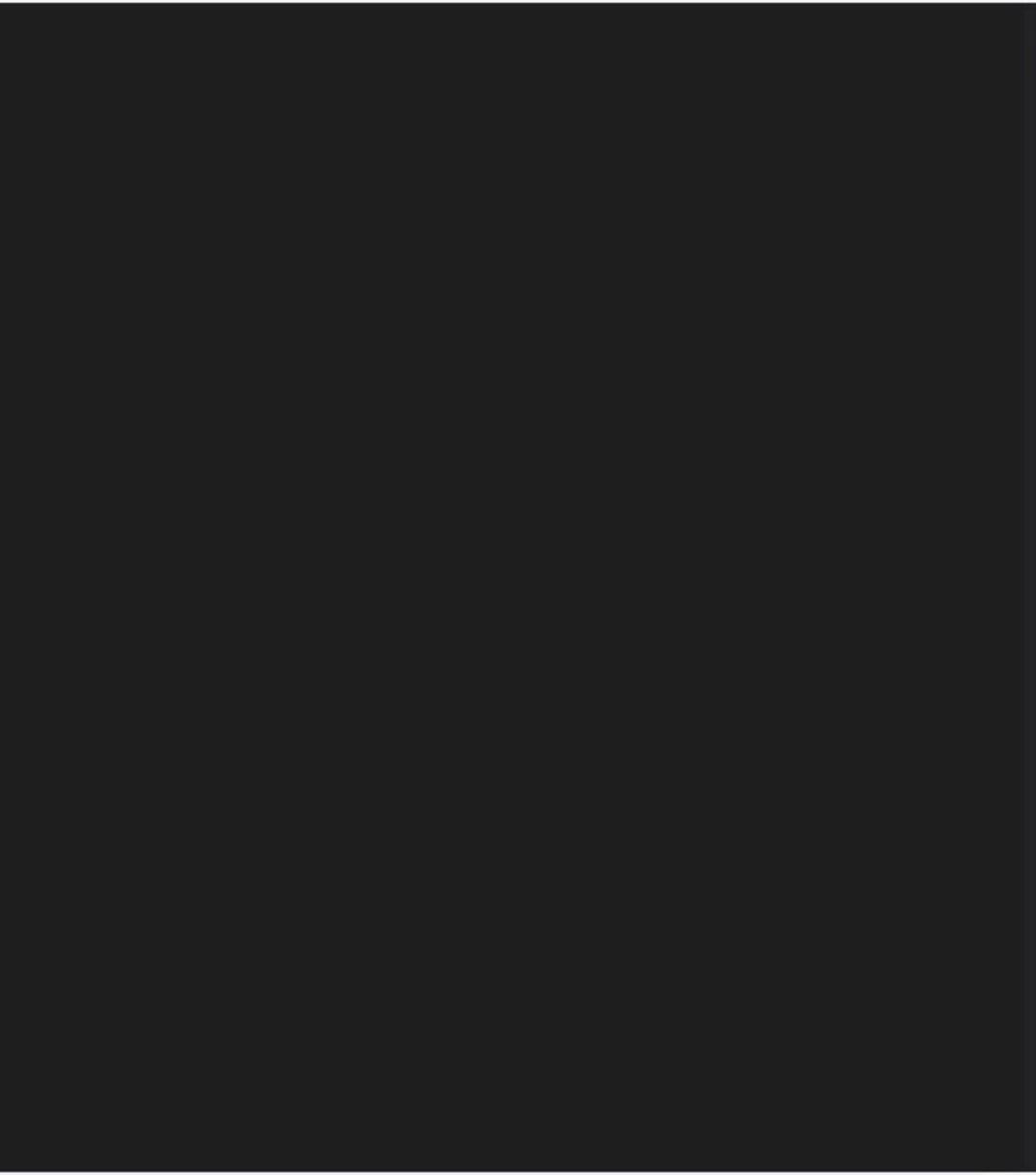
# Operating System CPU Scheduling FCFS SJF SRTF

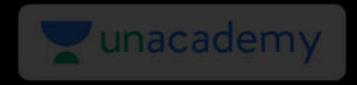
By: Vishvadeep Gothi



#### ▲ 1 • Asked by Bivas

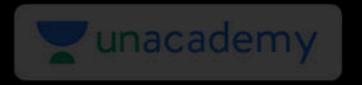
can mutiple process present in running state? if yes then what is the flow of a process running state to suspened ready state?





#### Scheduling Algorithms

- 1. FCFS
  - 2. SJF
  - 3. SRTF
  - HRRN
  - Priority Based
  - Round Robin
  - 7. Multilevel Queue Scheduling
  - Multilevel Feedback Queue Scheduling



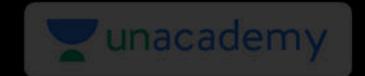
Criteria: Lesser Ar Circt Tie breaker = smaller\_id first

Mode: Non-preemptire

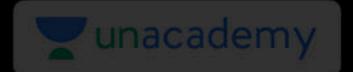


any	TAT	36
•	_	

Process	Arrival Time	<b>Burst Time</b>	C 1	T f
P1	0	30	30	30
P2	0	6	٦(	24
P3	0	6	して	42



Process	Arrival Time	Burst Time	Completion Time	Turnaround Time	Waiting Time
P1	0	30			
P2	0	6			
Р3	0	6			



Process	Arrival Time	Burst Time
P1	0	5
P2	1	1
P3	2	2
P4	3	4
P5	4	5
P6	5	3



Process	Arrival Time	Burst Time	CT	TAT	1 WT
P1	5	4	೨	Ч	0
P2	8	2	2	13	11
P3	6	3	12	6	3
P4	3	1	7	2	,
P5	2	2	4	2_	6
P6	7	7	19	12	5

Cru idle

P5 Py P1 D3 P6 P2

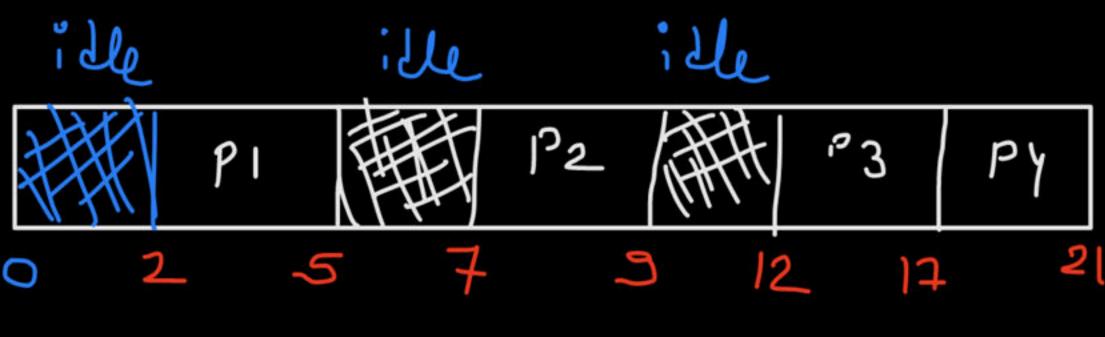
awy 
$$\omega T = \frac{39}{6}$$

Scheduling length (L) =  $21-2=19$ 

Thoughput =  $\frac{6}{19}$ 



Process	Arrival Time	Burst Time	CT	TAT	- W r
P1	2	3	5	3	1-
P2	7	2	9	2	0
P3	12	5	17	5	6
P4	13	4	21	8	4

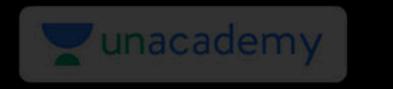


Imacadethie - like from 0 to 21,

(PU hoet = in line with cpu idle = 7 time unit

1. of time CPU utilized = 14 + 100% = 66.67%

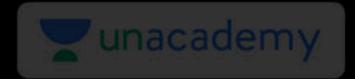
1. of time CPU ible =  $\frac{7}{21}$  \$ 100%. = 33.33 %.



### Convoy Effect



If a bij process is scheduled ahead of small processes, then w.T. of small processes will increase; and it will eventually slow down the system.

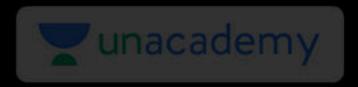


#### Advantages:

- Easy to implement
- No complex logic
- No starvation

#### Disadvantages:

- No option of Preemption
- 2. Convoy effect makes the system slow

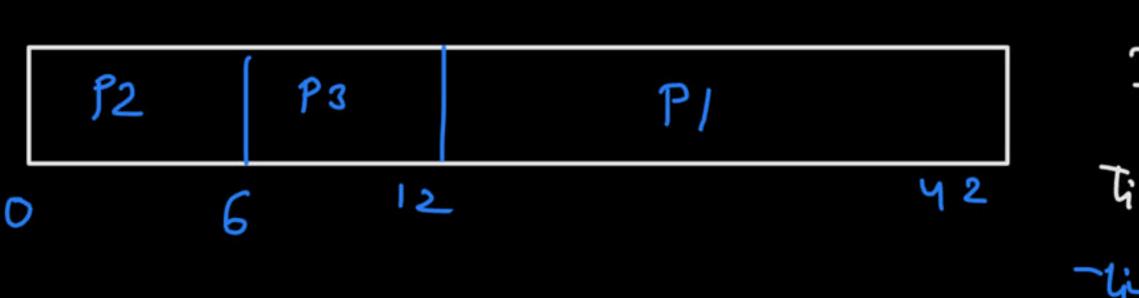


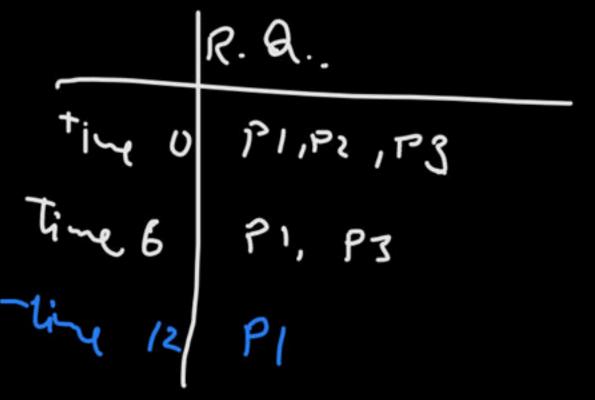
Criteria: Small Burst Time Tie breaker => FCF5

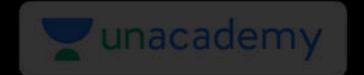
Mode: Non-preemptive



Process	Arrival Time	Burst Time	Completion Time	Turnaround Time	Waiting Time
P1	0	30	42	42	12
P2	0	6	6	6	O
Р3	0	6	2	12	6

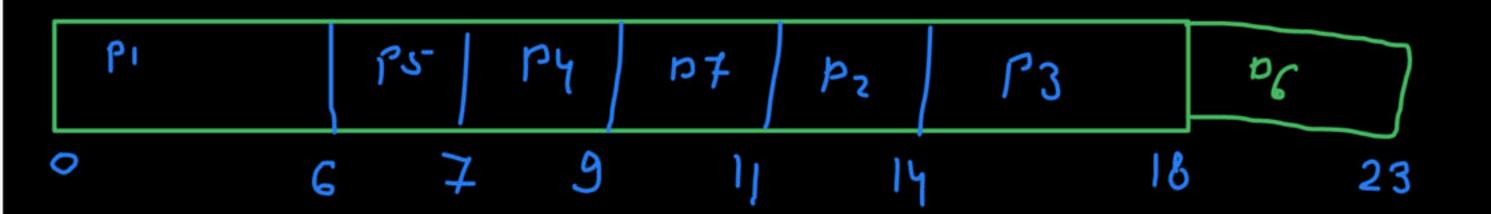






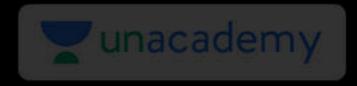
1136		17.9.
	0	PI
	6	P2, P3, P4 X5, P6
		P2, P3, P6, P7

Process	Arrival Time	Burst Time	Completion Time	Turnaround Time	Waiting Time
P1	0	6	6	6	0
P2	1	3	- ነ ካ	13	٥
Р3	2	4	18	16	12
P4	4	2	عا	5	3
P5	5	1	1	م	1
P6	6	5	23	17	12
P7	8	2		73	



Pro cassen	AT	<b>B</b>	Γ
PI	0	Ч	
rz	2	3	
P3	3	6	
84	4	3	

PI		r2	Р Ч	<b>†</b> 3
0	Ŋ	7	. [	०  ४

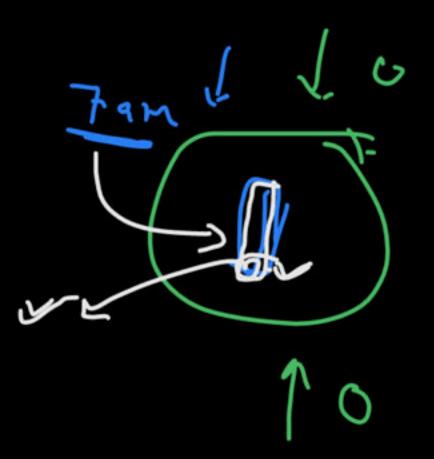


#### Advantages:

- Minimum average waiting time among non-preemptive scheduling
- Better throughput in continuous execution

#### Disadvantages:

- 1. No practical implementation because Burst time is not known in advance
- No option of Preemption
- Longer Processes may suffer from starvation

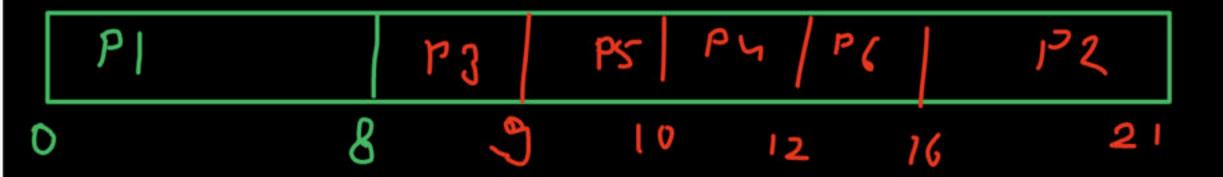


Criteria: Smaller 37 Leust Tie breaker => FCF5

Mode: Ereenstine

enemption happiers when a new process (Reemptive SJF) arrives, and it has beson BT then current running process's remaining time.

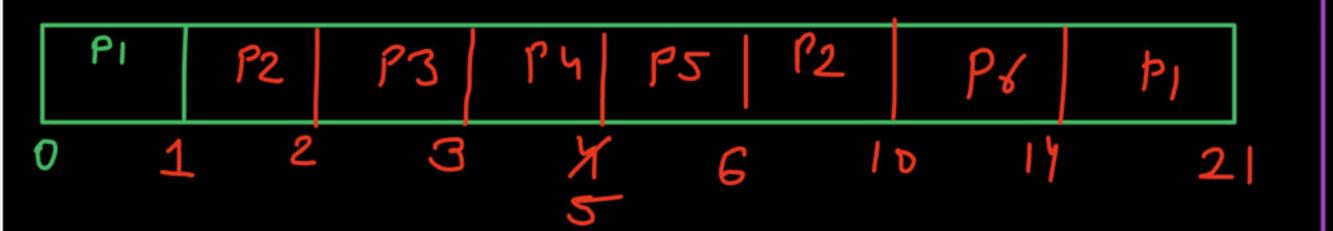
Process	Arrival Time	Burst Time	Completion Time	Turnaround Time	Waiting Time
P1	0	8			
P2	1	5			
Р3	2	1 .			
P4	3	2			
P5	4	1			
P6	5	4			



Process	Arrival Time	Burst Time	Completion Time	Turnaround Time	Waiting Time	7
P1	0	8	21	21	'3	0
P2	1	5	10	3	4	0
Р3	2	1	3	1	0	O
P4	3	2	5	2	U	0
P5	4	1	6	2	-1	1
P6	5	4	14	3	5	5

Time	13 Q
Ö	اع
1	P1, P2
2	P1, 02, BB
3	P1, P2, P4
7	P1. P2 , 24, P5
5	P1, 82. 26 P6

	Br
PZ	87
P	X
PS	24
126	<b>*</b>



Process	Arrival Time	Burst Time	Completion Time	Turnaround Time	Waiting Time	RT
P1	4	7	20	16	ಲ	೨
P2	5	5	13	8	3	2
Р3	3	1	5	2	1	1
P4	1	2	3	2	0	6
P5	2	1	4	2	1	1
P6	0	4	8	8	7	0

	RA
0	P(
1	P6, 134
2	P6, P4, P5
3	P61P5, P3
4	P6, 00, P1
5	P6, P1, P2

	16	PY	P5	Р3	P6	P2	þį
כ	1	23	4	5	- 8	17	3 20

#### unacademy

$$\sim 38$$

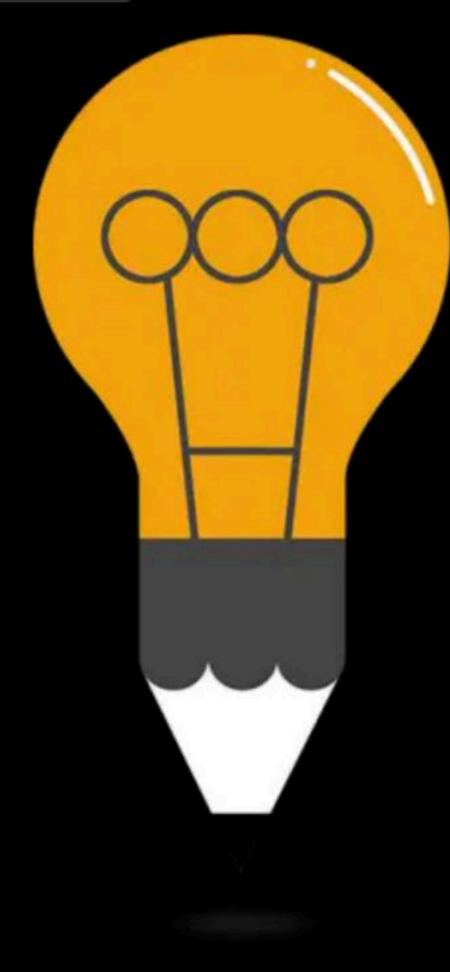
#### Advantages:

- Minimum average waiting time among all scheduling algorithm
- Better throughput in continue run

#### Disadvantages:

- 1. No practical implementation because Burst time is not known in advance
- 2. Longer Processes may suffer from starvation





## DPP

By: Vishvadeep Gothi



Response time of processes in non-preemptive scheduling algorithms are equal to waiting time of processes?

Frue or False

Justify your answer with appropriate explanation.

After waiting for CPU initially, process runs and completes.

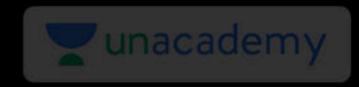
Hence afterwards no any waiting of process.



Consider the following process scenario.

Calculate the average waiting time and average TAT for processes for FCFS?

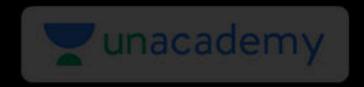
Process	Arrival Time	Burst Time
P1	5	6
P2	5	3
P3	8	4
P4	7	2
P5	3	1
P6	0	2
P7	7	2



Consider the following process scenario.

Calculate the average waiting time and average TAT for processes for SJF algo?

Process	Arrival Time	Burst Time
P1	5	6
P2	3	3
P3	1	4
P4	2	2
P5	4	1
P6	0	5
P7	1	2



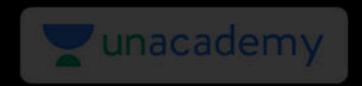
Consider a CPU performance metric throughput which is calculated as:

$$Throughput = \frac{Number\ of\ processes\ executed}{Total\ scheduling\ duration}$$

For the following process scenario calculate the throughput if process execution is done using:

- FCFS algorithm
- 2. SJ algorithm

Process	Arrival Time	Burst Time
P1	0	4
P2	0	3
P3	0	1
P4	0	5



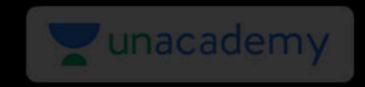
Consider a CPU performance metric throughput which is calculated as:

$$Throughput = \frac{Number\ of\ processes\ executed}{Total\ scheduling\ duration}$$

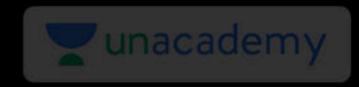
For the following process scenario calculate the throughput if process execution is done using:

- FCFS algorithm
- 2. SJ algorithm

Process	Arrival Time	Burst Time
P1	0	3
P2	3	6
P3	4	4
P4	5	2
P5	6	1
P6	3	3



Consider 4 processes A, B, C and D. All arrived at time 0 in the given order. The processes needed 5ns, 3ns, 9ns and 10ns respectively for their CPU burst to complete. What is the average waiting time of processes if executed in FCFS order?



Consider the following process scenario.

How much average waiting time can be saved using SRTF execution as compared to SJF execution?

Process	Arrival Time	Burst Time
P1	6	1
P2	2	4
P3	0	9
P4	5	6
P5	3	2
P6	1	8
P7	4	7



# Happy Learning.!



