

Graph Theory

Lecture 13

Tutte's 1-factor thm. (contd.)

G-graph; $g(G) = \# \text{odd comp. of } G$. (^{conn.} comp. with odd no. vertices)
Thm. $\Rightarrow G$ has perfect matching iff $g(G-S) \leq |S| \forall S \subseteq V(G)$

Pf. : ② Assumed the condition. Assumed that
 G has no 1-factor.

We added edges to G to reach a stage where we can assume that $G+e$ has 1-factor, G has no 1-factor & G satisfies $g(G-S) \leq |S|$.

If $U = \{v \mid \deg_G(v) = n-1 ; n = |V(G)|\}$

\rightarrow ④ $\langle G-U \rangle = \text{disjunct union of complete graphs. } \checkmark$

⑤ \exists a conn. comp. of $\langle G-U \rangle$ which is not a complete graph.

$\rightarrow \exists$ vertices x, y, z, w s.t.

$xy, yz \in E(G) \quad xz \notin E(G)$
 ie. $d_G(x, z) = 2$.

AND $\exists w \in V(G)$.

s.t. $yw \notin E(G)$.

By assumption on G , we know $G+xz$ has a perfect matching say M_1 . (clearly $xz \in M_1$)
 Ily $G+yw$ has 1-factor say M_2 ($yw \in M_2$).

Let H be the graph formed by $M_1 \Delta M_2$.

ie. $V(H) = V(G)$; & $E(H) = E(M_1) \cup E(M_2) - E(M_1 \cap M_2)$

Conn. comp. of H are either isolated vertices
or 2-regular.

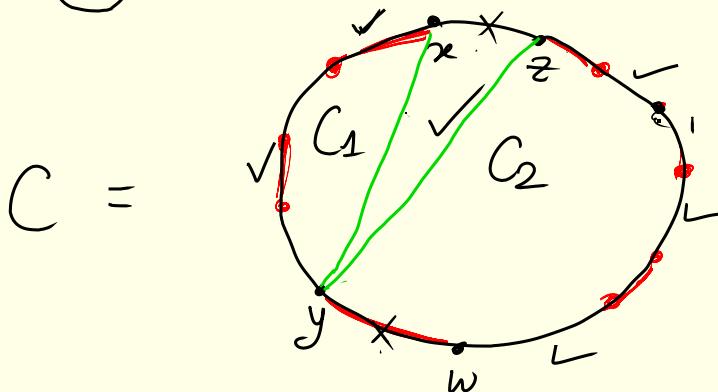
Since xz & yw remain in $M_1 \Delta M_2$,

$$\deg_H(x) = \deg_H z = \deg_H y = \deg_H w = 2.$$

This is where we stopped last time.

\Rightarrow the edge $e_1 = xz \in$ 2-regular conn. comp. of H
i.e. $e_2 = yw \in$ 2-regular conn. comp. of H

(a) e_1, e_2 are in same connected comp. of H , say C .

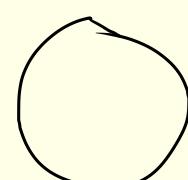
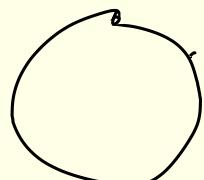
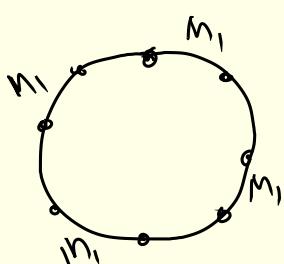


red edges are from M_2
black edges are from M_1 .
green edges are not in H
but they exist in $E(G)$.

Form M by taking edges from M_2 in $C_1 + yz$.
& edges from M_1 in C_2

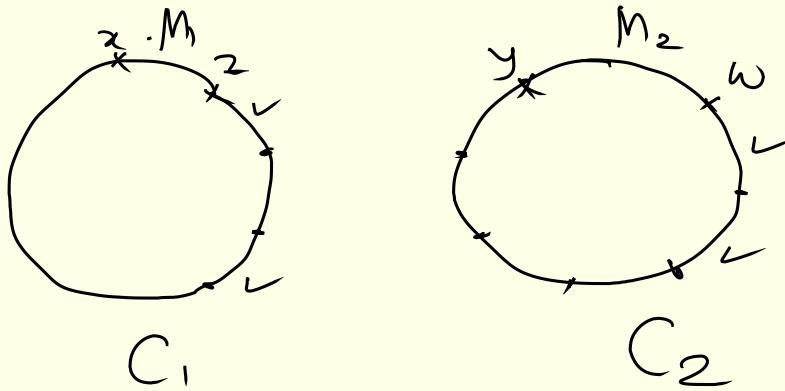
This M constitute a 1-factor of C .

which can be extended to a 1-factor of G by
adding all common edges in $M_1 \cap M_2$, & choosing
edges from M_1 from all other 2-reg. conn. comp.
of H .



This M is a 1-factor of G . a contradiction.

(b) e_1, e_2 belong to distinct connected components of H .



Construct a 1-factor for G by taking all M_2 -edges from C_1 , all M_1 -edges from C_2 , all common edges from $M_1 \cap M_2$ and in all other 2-reg. conn. comp. of H take edges from M_1 .

which again contradicts the construction of G being a graph without 1-factor.

QED.

As a consequence we reprove an old theorem of Petersen (1891)

Theorem (Petersen) Every 3-regular graph without any cut-edge has perfect matching.

Pf. \rightarrow Since $\sum_{v \in V(G)} d(v) = 2|E(G)|$, any $2k+1$ -reg. graph must have even no. of vertices.

We will show that $H \subseteq V(G)$, $|G - S| \leq |S|$.

Let $S \subseteq V(G)$. Let H be an odd comp. of $G - S$.

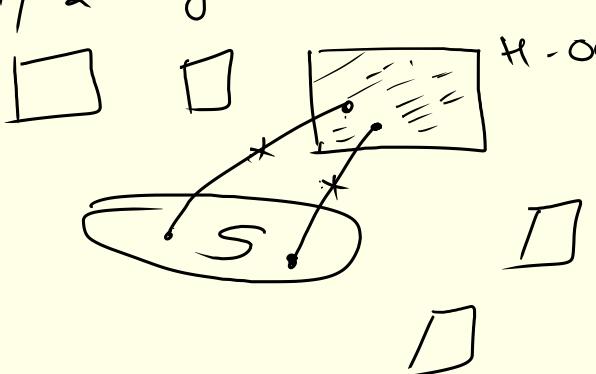
(First note that if we prove that every conn. comp. of G has a 1-factor, then G will have a 1-factor. \therefore assume G is connected)

Since G is connected, \Rightarrow at least one edge from H to S . (ie $\exists e \in E(G)$ s.t. one end pt of $e \in H$, other end pt. $\in S$)

If \exists an H s.t. $\exists!$ e betⁿ $H \leftarrow S$ then e will be an cut-edge of G . \Rightarrow every H must have at least 2-edges joining $H \leftarrow S$.



If H is odd comp. of $G-S$ then we can't have exactly 2-edges betⁿ $H \leftarrow S$.



H -odd comp. of $G-S$.

$$\text{then } \sum_{v \in V(H)} d(v) = \frac{3|V(H)| - 2}{2} = \text{odd.}$$

contradiction!

\Rightarrow every odd. comp. of $G-S$ must have at least 3-edges "joining it" to S .

3-edges $\leq \sum_{v \in V(H)} \deg(v)$

* edges with exactly one end pt. in $S \leq 3|S|$.

But \Rightarrow atleast $3 \cdot q(G-S)$ such edges

$$\Rightarrow 3 \cdot q(G-S) \leq 3|S| \Rightarrow q(G-S) \leq |S| \quad \forall S \subseteq V(G).$$

Now use Tutte's 1-factor theorem!

QED.

Defect version of Tutte's theorem.

$$q(G-S) \leq |S| \text{ but for some } S,$$

$$q(G-S) > |S|. \quad ; \quad \text{let } d(S) = \max_{S \subseteq V(G)} \{ q(G-S) - |S| \}.$$

Note. $d(S)$ & $V(G)$ have same parity.

$$\text{i.e. } d(S) \equiv V(G) \pmod{2}$$

Pf.

$$\textcircled{a} \quad |V(G)| - \text{odd} \begin{cases} \xrightarrow{|S| - \text{even}} |G-S| = \text{odd} \Rightarrow q(G-S) \text{ is odd} \\ \xrightarrow{|S| - \text{odd}} |G-S| = \text{even} \Rightarrow q(G-S) \text{ even} \end{cases} \Rightarrow d(S) \text{ odd}$$

$$\textcircled{b} \quad |V(G)| - \text{even} \begin{cases} \xrightarrow{|S| - \text{even}} |G-S| = \text{even} \\ \xrightarrow{|S| - \text{odd}} |G-S| = \text{odd} \end{cases} \begin{array}{l} \text{Prove that} \\ d(S) \equiv V(G) \pmod{2}. \\ (\text{Exercise}) \end{array}$$

Defect version of Tutte's thm.

(Berge 1958) The largest matching in an n -vertex graph G covers at least $n - \max_{S \subseteq V(G)} \{d(S)\}$ vertices.

Pf :- Clearly $\max_{S \subseteq V(G)} d(S) \geq 0$ take $S = \emptyset$.

Let $d = \max d(S)$. Define G' by

$$V(G') = V(G) \cup \{x_1, \dots, x_d\}$$

$$\& E(G') = E(G) \cup \{x_i x_j \mid i, j \in V(G)\}$$

We prove that Tutte's condition is satisfied by G' . (Exercise) $\Rightarrow G'$ has 1-factor

Then removing at most d edges gives a matching of required size for G .

QED.