

# Graph Theory

## Lecture 2

On MOODLE, New ref. book has been added.

→ AINS Open Notes Series  
Benny Sudakov - Graph Theory.

+ D. West - Introduction to G.T.

Connectedness in Graphs.

Path ; If any two vertices can be "joined by a path"  
then  $G$  - (graph) is connected.

Actually the relation  $\sim$  on  $V(G)$  is an equivalence  
relation where  $v \sim w$  iff  $\exists$  a  $v-w$  path.  
Equivalence classes are called connected components  
of a graph.

Thm: If a graph  $G$  has  $n$  vertices &  $k$  edges, then  
it must have at least  $n-k$  connected components.

Pf:- Note that this theorem needs to be proved only  
for  $k \leq n-1$ .

Prove this by induction. ①  $k=0$ .

Then  $G = \begin{matrix} \bullet \\ \vdots \\ \bullet \end{matrix} \rightarrow n$ -points.

&  $G$  has  $n = n-k$  conn. comp.

②  $k=1$ .  $G = \{c\} \cup \{v_1, v_2, \dots, v_n\}$

Thm. is true in this case too.

③ Induction hypothesis : If  $G$  has  $n$ -vertices  
&  $k-1$ -edges

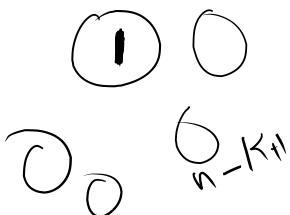
then  $G$  has  $\geq n+1-k$   
conn. comp.

④ If  $|V(G)| = n$  &  $|E(G)| = k$

Then remove one of the edges, use induction hypothesis & then observe that at most two conn. comp. of  $G-e$  can become into one compo. with rest being undisturbed.

$$\begin{aligned} \text{# comp of } G &\geq (\text{# comp of } G-e) - 1 \\ &\geq (n+1-k)-1 = n-k \end{aligned}$$

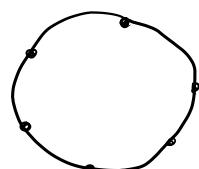
QED.



Tree :- A Tree is a graph with  $n$  vertices  $n-1$  edges & is connected.

(Extremal case in the above theorem!)

Defn :- A Cycle of length  $n$  is a closed path with  $n$  edges.  
 $\hookrightarrow$  (0-0 path)



= cycle of length 6.

Exercise : ① Prove that a tree can not have a cycle as a subgraph.

② Prove that given any two vertices  $v$  &  $w$  in a tree,  $\exists$  a unique path betw  $v$  &  $w$ .

③ A connected graph without cycle is a tree.

Forest :- Any graph without cycle is called a forest.

Exercise :- Let  $G$  be a graph, that is not connected.

If  $|V(G)| = 10$ , then prove that  $|E(G)| \leq 36$ .

Does equality hold?

Sol :- If  $G_1$  &  $G_2$  are two conn. comp. of  $G$ , then  $G$  can not have any edge with one end pt in  $V(G_1)$  & other in  $V(G_2)$ .

If  $|V(G_1)| = m$  &  $|V(G_2)| = n$  then from the possible  $45$  edges ( ${}^{10}C_2 = 45$ ),  $mn$  edges must be missing.

To find upper bound on  $|E(G)|$ , we need consider only the case of 2 conn. comp.

$$\Rightarrow m+n = 10.$$

$$\min \{mn \mid m+n=10\}.$$

45 - mn $\geq  E(G) $	
1+9	9
2+8	16
3+7	21
4+6	24
5+5	25

$$\Rightarrow |E(G)| \leq 36.$$

$$\Rightarrow G = K_9 \cup \{\cdot\}. \Leftrightarrow \text{equality is achieved!}$$

Exercise :- Let  $A_1, \dots, A_n$  be distinct subsets of  $\{1, 2, \dots, n\}$ . P.T. there exist a number  $1 \leq x \leq n$  s.t.  $A_i \setminus \{x\}$  are also distinct.

Pf :- When  $A_i \neq A_j$  becomes  $A_i \setminus \{x\} = A_j \setminus \{x\}$  ?

iff  $A_i = A_j \cup \{x\}$  OR  $A_j = A_i \cup \{x\}$ .

iff  $|A_i \Delta A_j| = 1$ .

Aside : S, T two subsets of a set  $X$ , then  $S \Delta T = S \setminus T \cup T \setminus S$   
 $= (S \setminus T) \cup (T \setminus S)$ .

on  $\wp(X)$ ,  $\Delta$  is a binary relation &  $\phi$  is identity  
 $A^2 = \text{id}$ .

$(\wp(X), \Delta)$  is an abelian group.

We form  $G$  whose vertices are labelled  $A_1, \dots, A_n$   
&  $\{A_i, A_j\} \in E(G)$  iff  $|A_i \Delta A_j| = 1$ . Let the edge  
 $A_i, A_j$  has colour  $x$  if  $A_i \Delta A_j = \{x\}$

→ We need to prove that there is one colour missing from  
these edges. If  $G$  was a tree,  $\exists$  only  $n-1$  edges & we  
are done.

If  $G$  was a forest, then it has  $k$  trees  $k \geq 2$ .

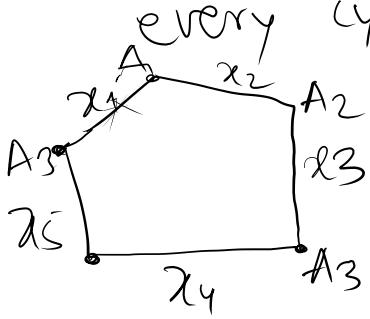
& each one have  $n_i - 1$  edges  $\sum n_i = N$ .

$$\Rightarrow \star = \sum n_i - k < n.$$

Hence we are done.

What happens if  $G$  has a cycle?

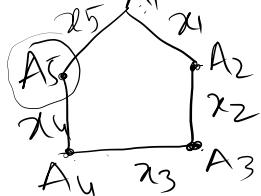
claim : In the graph  $G$  (as described above),  
every cycle must have a repeated colour.



- Assuming this claim, we can drop an edge from  $G$  s.t. that a given cycle is broken & no. of colours are still same!

Repeating this argument (if needed), we arrive at  
a graph  $\tilde{G} = G - \{e_1, \dots, e_k\}$  s.t.  $\tilde{G}$  is a forest and  
colours in  $\tilde{G}$  = colours in  $G$ . QED.

(proof of the claim)



Assume WLOG that  
 $|A_1| = \max_i \{|A_i| \mid 1 \leq i \leq 5\}$   
 $\Rightarrow A_2 = A_1 - \{x_1\}$ ;  $A_5 = A_1 - \{x_5\}$

$$\{x_2\} = A_2 \Delta A_3. \quad |A_3| = |A_1|$$

or  $|A_3| = |A_1| - 2 \leftarrow |A_4| = |A_3| - 1$   
 $|A_4| = |A_1| - 3.$

we know that  $|A_5| = |A_1| - 1.$

$\Rightarrow \exists i \text{ s.t. } |A_i| = |A_1| - 1$   
 $\Leftrightarrow |A_{i+1}| = |A_1| + 1.$

Also  $A_5 \subset A_1 \supset A_2$

- $\Rightarrow$  No new element remains by the time we reach  $A_5$
- $\Rightarrow$  If new element is added then it has to be removed also!
- $\Rightarrow$  One colour is repeated.
- $\Rightarrow$  only elements from  $A_1$  are added or subtracted (one at a time).
- $\Rightarrow$  A removed element must be added.
- $\Rightarrow$  A colour is repeated. Thereby proving the claim!  $\rightarrow \text{QED}$

Exercise :- Complete the proof of the claim by considering a cycle of length  $m$  (instead of 5).

— x — x — x —

Generate more exercises out of this.

$\rightarrow$  What happens if we take  $A_1, \dots, A_m$  distinct subsets of  $\{A_1, \dots, A_n\}$   $m \leq n$   $\rightarrow$  add  $B_{m+1}, \dots, B_n$

$\rightarrow$  What happens if you remove  $x \& y$ ?