

Combinatorics

Lecture 5

Conditions on parameters

{ Grading will be 60% - home assignments
40% - final exam.
NO mid-term. (like last sem) }

→ Incidence Structure. (P, B, I) .

P - set. elts are called points

B - set. elts are called blocks

$I \subset P \times B$ elts are called flats. ←

Linear space

B is a collection of subsets of P s.t.
any two points are in a unique block
& $|B| \geq 2 \quad \forall B \in B$.

(Erdős + De Bruijn) Either $|B| = 1$
or $|B| \geq |P|$.

Conway's proof.

t -design

Defⁿ:- Let v, k, t, λ be integers s.t. $v \geq k \geq t > 0$

A t -design is an incidence structure $\mathcal{D} = (P, B, I)$

s.t. ① $|P| = v$,

② $|B| = k \quad \forall B \in B$

③ Any subset of t points is contained in exactly λ blocks.

Ex:- $\mathcal{P} = \mathbb{P}^1(\text{GF}(q)) = \{ \text{set of all 1-diml subspaces of } \text{GF}(q)^{n+1} \}$

$$\mathcal{B} = \{ L_W \mid W - 2 \text{ diml subspace of } \text{GF}(q)^{n+1} \}$$

$$\& L_W = \{ v \in \mathcal{P} \mid v \subseteq W \}$$

$$|L_W| = \frac{(q^2-1)}{q-1} = q+1 = k, \quad v = \frac{q^{n+1}-1}{q-1} = 1+q+\dots+q^n$$

$t=2, \lambda=1$. Any two distinct one-dimensional subspaces of a v-space span a unique 2-diml subspace.

Aim of today:-

To give conditions that v, k, t, λ must satisfy if a $S_\lambda(t, k, v)$ -design exists.

Theorem 1. The number of blocks (i.e. $|\mathcal{B}|$) in an $S_\lambda(t, k, v)$ equals $\frac{\lambda \binom{v}{t}}{\binom{k}{t}} =: b$

(we denote this number by b).

($b = |\mathcal{B}|$, \mathcal{B} - a block.
 \mathcal{B} - set of all blocks)

Proof:- • If one knows how to double count, then one has a very good understanding of combinatorics

① Typically one constructs a set of tuples $S = \{ (a, b) \mid a \in \dots, b \in \dots \}$ - cleverness

② Count $|S|$ by fixing 1st coordinate } technical.

③ Count $|S|$ by fixing 2nd coordinate

② = ③ gives the required result.

$$① \quad S = \{ (T, B) \mid B \in \mathcal{B}, T \subset B \text{ s.t. } |T| = t \}$$

② Fixing T first we get $\binom{v}{t} - T$'s.
 & for each such T , $\exists \lambda - B$'s

$$\Rightarrow |S| = \lambda \cdot \binom{v}{t}.$$

③ Fixing B - first we see that $\exists |B|$ such B 's. & each one contains exactly $\binom{k}{t}$ subsets of size t .

$$\therefore |S| = |B| \cdot \binom{k}{t}.$$

$$\Rightarrow \lambda \binom{v}{t} = |B| \binom{k}{t} \Rightarrow |B| = \frac{\lambda \binom{v}{t}}{\binom{k}{t}}.$$

Q.E.D.

→ If one knows the result, it's easy to find a proof ←

Theorem 2. (Every t -design is also an i -design for every $0 \leq i \leq t$.)

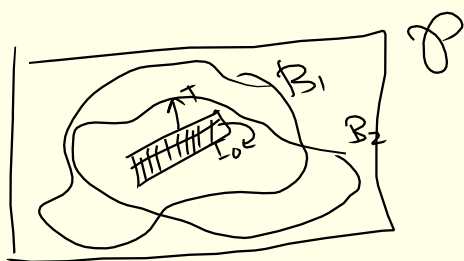
Given $0 \leq i \leq t$, the number of blocks that contain a subset of size i is :

$$b_i := \frac{\lambda \binom{v-i}{t-i}}{\binom{k-i}{t-i}} \quad 0 \leq i \leq t$$

$$\left(\Rightarrow S_{\lambda}(t, k, v) = S_{\frac{\lambda \binom{v-i}{t-i}}{\binom{k-i}{t-i}}}(i, k, v) \text{ design.} \right)$$

proof Fix $I_0 \subseteq \mathcal{P}$ s.t. $|I_0| = i$

$$\textcircled{1} \quad S_1 = \{ (T, B) \mid B \in \mathcal{B}, T \subseteq P; |T| = t \text{ \& } T \supseteq I_0 \}$$



\textcircled{2} Fixing T first :

$\exists \binom{v-i}{t-i}$ such T 's & for each such $T \exists \lambda$ B 's.

$$\therefore |S_1| = \lambda \binom{v-i}{t-i}.$$

\textcircled{3} On the other hand, Fixing B first we get

$\#$ blocks that contain I_0 say b_{I_0}

& for each such block, $\exists \binom{k-i}{t-i}$ T 's of size t inside it that contain I_0 .

$$\therefore |S_1| = b_{I_0} \cdot \binom{k-i}{t-i}$$

$$\textcircled{2} = \textcircled{3} \Rightarrow b_{I_0} = \frac{\lambda \binom{v-i}{t-i}}{\binom{k-i}{t-i}}. \quad \therefore b_{I_0} = b_{I_1} \text{ whenever } |I_0| = |I_1| = i \text{ } i \leq t$$

QED

Corollary (Prob. 19B in van lint + Wilson)

If $S_1(3, 6, v)$ exists, then $v \equiv 2$ or $6 \pmod{20}$.

Nomenclature: $S_\lambda(t, k, v)$ will be denoted as $S(t, k, v)$ if $\lambda = 1$.
↙
Steiner.

$S(3, 6, v)$ = collection of 6-subsets of a set of size v such that every 3-subset occurs in exactly one of these chosen 6-subsets.

\Rightarrow simplest non-trivial design must have at least 22 elts!

pf: $b_2 \Rightarrow v \equiv 2 \pmod{4}$; $b_1 \Rightarrow 20 \mid (v-1)(v-2)$.
 \Downarrow
 $5 \mid \frac{(v-1) \cdot (v-2)}{4}$

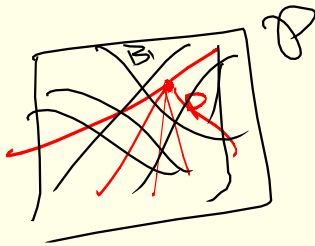
$$\begin{aligned} \text{If } 5 \mid v-1 &\Rightarrow 5 \mid v-6 \Rightarrow 5 \& 4 \text{ divides } v-6. \\ &\Rightarrow 20 \text{ divides } v-6 \\ &\Rightarrow v \equiv 6 \pmod{20} - (*) \end{aligned}$$

$$\text{If } 5 \mid \left(\frac{v-2}{4}\right) \Rightarrow v \equiv 2 \pmod{20}. - (**)$$

— x — x — x —
Remark: Since any t -design (for $t \geq 2$) is also a 2-design, theorems regarding 2-designs are applicable for all designs too. For 2-design.

$$b_1 = \frac{\lambda \binom{v-1}{1}_{21}}{\binom{k-1}{1}} = b_1 = \frac{\lambda(v-1)}{(k-1)}$$

"
 no. of blocks containing a subset of size 1. i.e. ~~*~~ blocks containing an element.



It is generally accepted to write $b_1 = r$ (the no. of times an elt. repeats in \mathcal{B})

$$\therefore \boxed{r(k-1) = \lambda(v-1)} \rightarrow \text{true in a 2-design.}$$

Count in two ways the set $S_x = \{(x, B) \mid x \in \mathcal{P}, B \in \mathcal{B}, x \in B\}$
 to get $\boxed{v \cdot r = b \cdot k}.$