



Introducing Core Solidity

Argot Collective





Solidity, today

1. Most used smart contract language
2. Successfully secures billions of dollars
3. However, it has some limitations:
 - Type system lacks expressiveness
 - Does not support compile time eval.
4. Extending the current impl. is **hard**.



What is Core Solidity?

1. A complete rebuild of the language type system and compiler pipeline.
2. Introduce new features
3. A formal semantics to enable analysis and verification.



Core Solidity New features

1. Generics and type classes (traits)
2. Algebraic data types and pattern matching
3. Type inference
4. High-order and anonymous functions.
5. Compile time evaluation.



Algebraic data types

```
data wad = wad(uint256)
```

- Data type for representing a 18 decimal fixed point



Algebraic data types

```
data wad = wad(uint256)
```

Type name

Constructor name

The diagram illustrates the components of the Solidity code. The variable name 'wad' is annotated with two labels: 'Type name' pointing to the first occurrence and 'Constructor name' pointing to the second occurrence. This visual distinction highlights that in Solidity, type names and constructor names share the same identifier.

- Data type for representing a 18 decimal fixed point
- Constructors and type names live in different namespaces



Algebraic data types

```
data AuctionState =  
    NotStarted(uint256)  
    | Active(uint256, address)  
    | Ended(uint256, address)  
    | Cancelled(uint256, address);
```

- Allows data modelling using sum and product types.



Algebraic data types

```
data AuctionState =  
    NotStarted(uint256)  
    | Active(uint256, address)  
    | Ended(uint256, address)  
    | Cancelled(uint256, address);
```

Type name

Sum types/Alternatives

```
graph LR; A["Type name"] --> B["data AuctionState =  
    NotStarted(uint256)  
    | Active(uint256, address)  
    | Ended(uint256, address)  
    | Cancelled(uint256, address);"]; C["Sum types/Alternatives"] --- B;
```

- Sums allows the definition of exclusive alternatives in a type.



Algebraic data types

Type name

```
data AuctionState =  
    NotStarted(uint256)  
  | Active(uint256, address)  
  | Ended(uint256, address)  
  | Cancelled(uint256, address);
```

Product type

- Products combine values into structured tuples.



Pattern Matching

```
data wad = wad(uint256)

let WAD = 10 ** 18;

function wmul(lhs : wad, rhs : wad) -> wad {
    match (lhs, rhs) {
        | (wad(l), wad(r)) => return wad((l * r) / WAD);
    }
}
```

- Pattern matching allows changing control flow based on value shape.



Pattern Matching

```
let WAD = 10 ** 18;                                Pattern match block  
  
function wmul(lhs : wad, rhs : wad) -> wad {  
    match (lhs, rhs) {  
        | (wad(l), wad(r)) => return wad((l * r) / WAD);  
    }  
}
```



Pattern Matching

```
let WAD = 10 ** 18;  
  
function wmul(lhs : wad, rhs : wad) -> wad {  
    match (lhs, rhs) {  
        | (wad(l), wad(r)) => return wad((l * r) / WAD);  
    }  
}
```

A white arrow points from the word "lhs" in the first line of the code to the word "lhs" in the pattern matching expression "match (lhs, rhs) {". Above the arrow, the word "Discriminee" is written in white text.



Pattern Matching

```
let WAD = 10 ** 18;

function wmul(lhs : wad, rhs : wad) -> wad {
    match (lhs, rhs) {
        | (wad(l), wad(r)) => return wad((l * r) / WAD);
    }
}
```

Wrapped uint256 values



Pattern Matching

```
let WAD = 10 ** 18;

function wmul(lhs : wad, rhs : wad) -> wad {
    match (lhs, rhs) {
        | (wad(l), wad(r)) => return wad((l * r) / WAD);
    }
}
```



Wrapping result in
data constructor



Generics

```
forall T . function identity(x : T) -> T {  
    return x;  
}
```

- Generics allows functions that work in a uniform way for all types.
- Example: polymorphic identity function.



Generics

Type variable definition.

```
forall T . function identity(x : T) -> T {  
    return x;  
}
```

Parameter type.

- Generics allows functions that work in a uniform way for all types.
- Example: polymorphic identity function.



Generics

Type variable definition.

```
forall T . function identity(x : T) -> T {  
    return x;  
}
```

Function return type.

- Generics allows functions that work in a uniform way for all types.
- Example: polymorphic identity function.



Generics

```
data Result(T) = Ok | Err(T);
```

- Core supports the definition of generic types
 - Type parameter for the payload used in Err constructor



Generics

Type parameter definition

```
data Result(T) = Ok | Err(T);
```

- Core supports the definition of generic types
 - Type parameter for the payload used in Err constructor



Generics

Type parameter use

```
data Result(T) = Ok | Err(T);
```

- Core supports the definition of generic types
 - Type parameter for the payload used in Err constructor



Generics and Type classes

1. Generics are a bit limited.
 - Few functionality works **for all** types.
2. We need functionality which works for **some** types.
3. Type classes are the solution!
 - Allow the systematic combination of generics and overloading.



Type classes

```
forall T . class T:Mul {  
    function mul(lhs : T, rhs : T) -> T;  
}
```

- Type classes define a set of operations with their types.
 - Example: class for defining a multiplication operation



Type classes

```
forall T . class T:Mul {  
    function mul(lhs : T, rhs : T) -> T;  
}
```



The word 'T' in the code is highlighted in purple. A white arrow points from the text 'Class name' to the word 'T', indicating it is a placeholder for the class name.

- Type classes define a set of operations with their types.
 - Example: class for defining a multiplication operation



Type classes

```
forall T . class T:Mul {  
    function mul(lhs : T, rhs : T) -> T;  
}
```

Type argument

Function signature

The diagram consists of two annotations. A curly brace is positioned under the entire function definition (the class declaration and the function line). An arrow originates from the word 'Type' in the label 'Type argument' and points to the first 'T' in 'T:Mul'.



Type classes

```
instance wad:Mul {  
    function mul(lhs : wad, rhs : wad) -> wad {  
        return wmul(lhs, rhs);  
    }  
}
```

Type argument

```
graph TD; A[Type argument] --> B[wad]; A --> C[wad]; A --> D[wad]; A --> E[wad];
```

- Instances provide an implementation for a specific type.
- Function types should “replace” the class type argument for the corresponding concrete type.



Type classes

Type class constraint

```
forall T . T:Mul => function square(val : T) -> T {  
    return Mul.mul(val, val);  
}
```

- Type class constraints allow us to restrict polymorphic types.
 - Function square can be called on all types which are instances of Mul.



High-order and anonymous functions

1. Functions are first-class in Core Solidity.
2. This means that they can be used:
 - Function parameters
 - Return values



Anonymous functions

```
forall T U . function count_calls(fn : (T) -> U) -> (memory(word), (T) -> U) {  
    let counter : memory(word) = allocate(32);  
    return (counter, lam (a : T) -> {  
        counter += 1;  
        return fn(a);  
    } );  
}
```

- Anonymous functions are defined using the **lam** keyword.
- Can capture values of their definition scope.
- Example: count the number of times which an argument function is called.



Anonymous functions

```
forall T U . function count_calls(fn : (T) -> U) -> (memory(word), (T) -> U) {  
    let counter : memory(word) = allocate(32);  
    return (counter, lam (a : T) -> {  
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    });  
}
```



**Anonymous function
definition**

- Anonymous functions are defined using the **lam** keyword.



Anonymous functions

```
forall T U . function count_calls(fn : (T) -> U) -> (memory(word), (T) -> U) {  
    let counter : memory(word) = allocate(32);  
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        counter += 1;  
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    });  
}
```



Captured value

- Anonymous functions are defined using the **lam** keyword.
- Can capture values of their definition scope.



High-order functions

```
forall T U . function count_calls(fn : (T) -> U) -> (memory(word), (T) -> U) {  
    let counter : memory(word) = allocate(32);  
    return (counter, lam (a : T) -> {  
        counter += 1;  
        return fn(a);  
    });  
}
```



Function type

- Type $(T) \rightarrow U$ represents a function
 - That takes a value of type T as argument.
 - Returns a value of type U .



Type inference

1. Core Solidity infers types in functions and local variables.
2. Annotations are required only to solve ambiguities.
 - Occur in very rare situations.
3. Can help readability by omitting unnecessary type annotations.



Type inference

```
uint256[3] memory a = [1, 2, 3];
```

- Consider this simple Classic Solidity definition.



Type inference

```
uint256[3] memory a = [1, 2, 3];
```

- It is rejected with the following error message:

Error: Type uint8[3] is not implicitly convertible to type uint256[3].



Type inference

```
uint256[3] memory a = [1, 2, 3];
```

- Classic Solidity has a limited support for inferring the type of array literals.
 - Elements have the type of the first expression such that all other elements can be casted to it.
- Inferred type for the first element: uint8



Type inference

```
uint256[3] memory a = [uint256(1), 2, 3];
```

- Solution: use a type coercion on the first array element.



Type inference

```
uint256[3] memory a = [1, 2, 3];
```

- In Core Solidity, this definition is accepted
 - Thanks to type inference, which uses a more general strategy which eliminates the need of the type coercion.



Extended example

1. Core Solidity introduces several new features.
2. How those features help?



Console.log

```
function logUint(uint256 p0) internal pure {
    _sendLogPayload(abi.encodeWithSignature("log(uint256)", p0));
}

function logString(string memory p0) internal pure {
    _sendLogPayload(abi.encodeWithSignature("log(string)", p0));
}
```

- Library forge-std implementation has a combinatorial explosion of overloaded functions.
- More than 30 different functions, one for each argument type.



Console.log

```
function _sendLogPayloadImplementation(bytes memory payload) internal view {
    address consoleAddress = CONSOLE_ADDRESS;
    /// @solidity memory-safe-assembly
    assembly {
        pop(
            staticcall(
                gas(), consoleAddress,
                add(payload, 32),
                mload(payload), 0, 0
            )
        )
    }
}
```

- Log functions just redirect to `_sendLogPayloadImplementation`.



Console.log - Core Solidity

```
forall T . T:ABIEncode => function log(val : T) {
    let CONSOLE_ADDRESS : word = 0x00000000000000000000000000000000636F6e736F6c652e6c6f67;
    let payload = abi_encode(val);
    // extract the underlying word representation of the payload
    let ptr = TypeDef.rep(payload);
    assembly {
        pop(
            staticcall(
                gas(), CONSOLE_ADDRESS,
                add(ptr, 32),
                mload(ptr), 0, 0
            )
        )
    }
}
```

- Generics and type classes eliminate all repetitive definitions!



Console.log - Core Solidity

```
forall T . T:ABIEncode => function log(val : T) {
    let CONSOLE_ADDRESS : word = 0x00000000000000000000000000000000636F6e736F6c652e6c6f67;
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                gas(), CONSOLE_ADDRESS,
                add(ptr, 32),
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            )
        )
    }
}
```

ABI encode constraint

ABI encode argument

- Generics and type classes eliminate all repetitive definitions!



Console.log - Core Solidity

```
forall T . T:ABIEncode => function log(val : T) {
    let CONSOLE_ADDRESS : word = 0x00000000000000000000000000000000636F6e736F6c652e6c6f67;
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            )
        )
    }
}
```

Convert to assembly
level representation.

- Generics and type classes eliminate all repetitive definitions!



SAIL , desugaring and standard library

- SAIL is a new mid-level IR for Core.
 - Solidity Algebraic Intermediate Language.
- Features
 - Functions and contracts
 - Algebraic types and pattern matching
 - Assembly blocks
 - Type classes and generics
 - Variable introduction and assignment.
 - Short-circuiting if-then-else expression.



SAIL, desugaring and standard library

- Core Solidity high-level constructs will be implemented by a combination of desugaring steps and std-lib definitions.
- Similar approach used in proof assistants like Lean.
- Objective: RFC style for language and std-lib changes



Compatibility and interoperability

- Introducing major reviews to a language is challenging.
- We plan a smooth transition to avoid language split.
 - ABI compatibility will be maintained.
 - We'll try to minimize syntax changes.
 - Investigate tooling for automated migration.



Simple Contract - Classic Solidity

```
pragma solidity >=0.4.16 <0.9.0;

contract SimpleStorage {
    uint storedData;

    function set(uint x) public {
        storedData = x;
    }

    function get() public view returns (uint) {
        return storedData;
    }
}
```



Simple Contract - Core Solidity

```
import std;
contract SimpleStorage {
    storedData : uint256;

    function set(x : uint256) -> () {
        storedData = x;
    }

    function get() -> uint256 {
        return storedData;
    }
}
```

- Minimal changes: postfix types, import of standard library.



The road to production

- Prototype implementation available.
- Standard library has implementations of:
 - ABI compatible contracts
 - ABI encoding / decoding
 - Dispatch
 - Storage access.
- Current prototype is able to compile a ERC20 contract.



The road to production

- Next steps (finish type system):
 - Define compile time evaluation.
 - Module system.
- Prototype stable:
 - Start a production implementation.
 - Mechanize the meta-theory using a proof assistant.



Beyond 1.0

- Our focus: deliver the language described so far.
- Possible future iterations:
 - Linear types.
 - Refinement types.
 - Macros
 - Theorem proving



Conclusion

- Core Solidity is a foundational re-imagining of the language.
- Our objective is to build a language:
 - More expressive.
 - More secure.
 - Mathematically sound.



Thanks!

Thanks !



**Core Solidity
repository**



**Core Solidity
feedback
forum thread**