The Poseidon Hash Function formalisation in Lean 4

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1 Introduction

In this document, we describe the Poseidon formalisation in Lean 4. Initially, Poseidon has been introduced as a family of hash functions based on the permutations $Poseidon^{\pi}$, see [1]. This approach utilises the Hades strategy [2] according to which (TODO: explain the Hades strategy with a couple of sentences).

2 Formalisation itself

We need some preliminaries first. Let p, t be natural numbers and p is prime. As usual, \mathbb{Z}_p stands for prime field of order p. In the original text, it is also assumed that $\lceil \log_2(p) \rceil = n$ for some n > 0.

We introduce these requirements in Lean 4 by declaring the following variables.

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variable (p t : \mathbb{N}) [Fact p.Prime] [Field (Zmod p)] [Fintype (Fin_x t)]
\operatorname{def} ARC (c a : Fin<sub>x</sub> t \rightarrow Zmod p) (i : Fin<sub>x</sub> t) : Zmod p := (a i) + (c i)
{\tt def} \ {\tt R\_f\_round} \ ({\tt S\_box'} \ : \ {\tt Zmod} \ {\tt p} \ \to \ {\tt Zmod} \ {\tt p}) \ ({\tt c} \ : \ {\tt Fin}_x \ {\tt t} \ \to \ {\tt Zmod} \ {\tt p})
   (MDS': Matrix (Fin_x t) (Fin_x t) (Zmod p)) (a : Fin_x t 	o Zmod p) : Fin_x t 	o Zmod p :=
   Matrix.mulVec<sub>x</sub> MDS' (\lambda i => S_box' (ARC p t c a i))
      {\tt def} \ {\tt R\_p\_round} \ ({\tt S\_box'} \ : \ {\tt Zmod} \ {\tt p} \ \to \ {\tt Zmod} \ {\tt p}) \ ({\tt c} \ : \ {\tt Fin}_x \ {\tt t} \ \to \ {\tt Zmod} \ {\tt p})
      (MDS' : Matrix (Fin_x t) (Fin_x t) (Zmod p)) (a : Fin_x t 	o Zmod p) : Fin_x t 	o Zmod p :=
      Matrix.mulVec_x MDS
          (\lambda i => dite ((i : \mathbb{N}) = 0) (\lambda => S_box' (ARC p t c a i)) (\lambda => ARC p t c a i))
\texttt{def} \ \texttt{P\_perm} \ (\texttt{R\_f} \ \texttt{R\_p} : \mathbb{N}) \ (\texttt{S\_box'} : \texttt{Zmod} \ \texttt{p} \to \texttt{Zmod} \ \texttt{p}) \ (\texttt{c} \ \texttt{a} : \texttt{Fin}_x \ \texttt{t} \to \texttt{Zmod} \ \texttt{p})
   (MDS' : Matrix (Fin_x t) (Fin_x t) (Zmod p)) : Fin_x t 	o Zmod p :=
   (R_f_round p t S_box' c MDS')^[R_f] ((R_p_round p t S_box' c MDS')^[R_p]
   ((R_f_round p t S_box' c MDS')^[R_f] a))
{\tt def} add_to_state (r cap : {\mathbb N}) (m : {\sf Fin}_x r 	o Zmod p)
   (a : \operatorname{Fin}_x t 	o Zmod p) (h : t = r + cap) : \operatorname{Fin}_x t 	o Zmod p :=
   \lambda i => dite ((i : \mathbb{N}) < r) (\lambda h => a i + m (Fin<sub>x</sub>.castLt i h)) (\lambda h => a i)
\textcolor{red}{\texttt{def}} \ \texttt{P\_hash} \ (\texttt{R\_f} \ \texttt{R\_p} \ \texttt{r} \ \texttt{o} \ \texttt{cap} \ : \ \mathbb{N}) \ (\texttt{hr} \ : \ \texttt{1} \ \leq \ \texttt{r}) \ (\texttt{S\_box'} \ : \ \texttt{Zmod} \ \texttt{p} \ \to \ \texttt{Zmod} \ \texttt{p})
   (c : Fin_x (r + cap) \rightarrow Zmod p)
   (MDS': Matrix (Fin_x (r + cap)) (Fin_x (r + cap)) (Zmod p)) (ho : o \leq r + cap)
   (k : \mathbb{N}) (a : \mathrm{Fin}_x (k * r + (r + cap)) 	o Zmod p) : \mathrm{Fin}_x o 	o Zmod p
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^{*}The Lean 3 formalisation is by Ashvni Narayanan. Daniel Rogozin prepared the Lean 4 version and its text description.

References

- [1] Lorenzo Grassi, Dmitry Khovratovich, Christian Rechberger, Arnab Roy, and Markus Schofnegger. Poseidon: A new hash function for {Zero-Knowledge} proof systems. In 30th USENIX Security Symposium (USENIX Security 21), pages 519–535, 2021.
- [2] Lorenzo Grassi, Reinhard Lüftenegger, Christian Rechberger, Dragos Rotaru, and Markus Schofnegger. On a generalization of substitution-permutation networks: The hades design strategy. In *Annual International Conference on the Theory and Applications of Cryptographic Techniques*, pages 674–704. Springer, 2020.