

# Formalizing ZkSNARKs

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# Introduction

In this blueprint we formalize the knowledge soundness proof for the BabySNARK protocol.

The protocol is a toy example of a Zk SNARK protocol defined and implemented in the repository .

The outline of the proof used in this blueprint is a port of the work done by Bolton Bailey in Lean 3 found [here](#) .

We will eventually include proofs of knowledge soundness for other protocols as well, and proofs of completeness and zero knowledge.

# Chapter 1

## Supporting Lemmas

**Definition 1.1.** Fix a finite field  $F$ . The polynomial  $t \in F[X]$  is then defined as  $t = \prod_{i=0}^{m-1} (X - r_i)$  for  $r_i \in F$  for  $0 \leq i \leq m-1$  where  $m > 0$ .

**Lemma 1.2.** The polynomial  $t$  is monic of positive degree  $0 < m$ .

*Proof.*

□

## Chapter 2

# Knowledge Soundness for BabySNARK

Fix natural numbers  $n, l$  with  $l < n$ , and a sequence  $a_s = (a_0, \dots, a_{l-1})$  and  $a_w = (a_l, \dots, a_n)$  (the *statement* and *witness*).

Fix also a collection of polynomials  $u_i(X) \in F[X]$  for  $0 \leq i < n$  split up into the first  $l$  denoted  $u_s$  and the last  $n-l$  denoted  $u_w$ .

Finally, fix strings  $(b_0, \dots, b_m), (h_0, \dots, h_m), (v_0, \dots, v_m) \in F^m$  where  $m = \deg t$ . Similarly fix  $(b'_i)_{i=l}^{n-l-1}, (v'_i)_{i=l}^{n-l-1}$  and  $(h'_i)_{i=l}^{n-l-1} \in F^{n-l}$ , and  $b_\gamma, v_\gamma, h_\gamma, b_{\gamma\beta}, v_{\gamma\beta}$ .

**Definition 2.1.** Define the polynomials  $V_s = \sum_{i=0}^{n-1} a_i u_i(X) = V_{ss} + V_{sw}$ , and

$$\begin{aligned} B_w &= \sum_{i=0}^{m-1} b_i X^i + b_\gamma Z + b_{\gamma\beta} YZ + \sum_{i=l}^{n-1} b'_i Y u_i(X) \\ V_w &= \sum_{i=0}^{m-1} v_i X^i + v_\gamma Z + v_{\gamma\beta} YZ + \sum_{i=l}^{n-1} v'_i Y u_i(X) \\ H &= \sum_{i=0}^{m-1} h_i X^i + h_\gamma Z + h_{\gamma\beta} YZ + \sum_{i=l}^{n-1} h'_i Y u_i(X) \end{aligned}$$

**Definition 2.2.** Call a sequence  $(a_i)_{i=0}^{n-1}$  satisfying if

$$\sum_{i=0}^{l-1} a_i u_i(X) + \sum_{i=l}^{n-1} a_i u_i(X) \equiv 1 \pmod{t}$$

**Lemma 2.3.**  $\forall 0 \leq i < m$ , the coefficient of  $X^i$  in  $B_w$  (or  $B_{wit}$ ) is  $b_i$ .

*Proof.*

□

**Lemma 2.4.** The coefficient of  $Z^2$  in  $Ht + 1$  is 0.

**Lemma 2.5.** Given  $(a_i)_{i=0}^{l-1}$ , the coefficient of  $Z^2$  in  $(b_{\gamma\beta} \cdot Z + \sum_{i=0}^{l-1} a_i u_i(X) + \sum_{j=l}^{n-1} b'_j u_j(X))^2$  is  $b_{\gamma\beta}^2$ .

For the following lemmas assume we are in the setting of the proof of knowledge soundness. In particular, assume:

$$B_w = YV_w \quad (2.1)$$

$$Ht = V^2 - 1 \quad (2.2)$$

Knowledge soundness for BabySNARK follows from the following lemmas:

**Lemma 2.6.** *Then given a monomial  $m$  not having a  $Y$ -term, the coefficient of  $m$  in  $B_w$  is 0.*

*Proof.* □

**Lemma 2.7.**  $\forall 0 \leq i \leq m-1, b_i = 0.$

*Proof.* □

**Lemma 2.8.**  $b_\gamma = 0$

*Proof.* □

**Lemma 2.9.**  $B_w = b_{\gamma\beta}ZY + \sum_{i=l}^{n-1} b'_i Y u_i(X)$

*Proof.* □

**Lemma 2.10.**  $V_w = b_{\gamma\beta}Z + \sum_{i=l}^{n-1} b'_i u_i(X)$

*Proof.* □

**Lemma 2.11.**  $V(a_i)_{i=0}^l = b_{\gamma\beta}Z + \sum_{i=0}^{l-1} a_i u_i(X) + \sum_{i=l}^{n-1} b'_i u_i(X)$

*Proof.* □

**Lemma 2.12.**  $b_{\gamma\beta} = 0$

*Proof.* □

**Lemma 2.13.**  $V(a_i)_{i=0}^l = \sum_{i=0}^{l-1} a_i u_i(X) + \sum_{i=l}^{n-1} b'_i u_i(X)$

*Proof.* □

**Lemma 2.14.**  $(Ht + 1) \equiv (V(a_i)_{i=0}^l)^2 \pmod{t}$

**Lemma 2.15.**  $\text{singlify}(Ht + 1)/t = \text{singlify}H$

*Proof.* □

**Theorem 2.16.** *If an adversary produces polynomials  $B(X, Y, Z), V(X, Y, Z), H(X, Y, Z)$  which satisfy  $B_w = YV_w$  and  $Ht = V^2 - 1$ , then the adversary can extract a satisfying witness.*

*Proof.* □