Convex Optimization II

Lecture 18: Sequential Convex Programming: A Method for Non-convex Optimization Problems

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1400-2

MOTIVATIONS

For convex optimization problems, in general, we can *easily* obtain a local optima which is the same as the *global* solution.

For non-convex optimization problems, we have to give up one. We can consider two approaches:

- Local optimization methods: Fast but sub-optimal
- Global optimization methods: Often slow, but attain the global solution.
- [1] John Duchi, "Lecture Notes on Sequential Convex Programming", [Online] https:
- //web.stanford.edu/class/ee364b/lectures/seq_notes.pdf

Thanks to Professor John Duchi, Stanford University, for the slides.

Sequential Convex Programming (SCP): A local optimization method for non–convex problems that leverages convex optimization.

- Handle the convex portions of the problem exactly and efficiently
- Model the non–convex portions of the problem by convex functions that are (at least locally) accurate.

SCP is a heuristic method.

- may fail to find an optimal (or even feasible) point,
- may depend on the starting point,

but often finds good solutions.

An iterative approach

- At iteration k, for any non-convex function f, form a model \hat{f} that is good enough near the current iterate $x^{(k)}$.
- Then, minimize that model or a regularized version of it, and repeat.

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Consider the (potentially non-convex) problem

minimize
$$f_0(x)$$

subject to $f_i(x) \leq 0$, for $i = 1, \dots, m$
 $h_j(x) = 0$, for $j = 1, \dots, p$
 $x \in \mathbb{R}^n$.

Here, the functions f_0 and f_i are (possibly) non–convex, and the functions h_j may be non-affine.

SCP: iterate by maintaining an estimate of the solution $x^{(k)}$ and a convex trust region, denoted by $\mathcal{T}^{(k)} \subset \mathbb{R}^n$, over which we trust our solutions and models.

The generic SCP strategy then forms a

- convex approximation $\widehat{f_i}$ of the functions f_i over the trust region $\mathcal{T}^{(k)}$
- affine approximation \widehat{h}_j of the functions h_j over the trust region $\mathcal{T}^{(k)}$

SCP iterations:

$$\begin{array}{ll} x^{(k+1)} = & \arg\min & \widehat{f}_0(x) \\ & \text{subject to} & \widehat{f}_i(x) \leq 0, \text{for } i = 1, \dots, m \\ & \widehat{h}_j(x) = 0, \text{for } j = 1, \dots, p \\ & x \in \mathcal{T}^{(k)}. \end{array}$$

TRUST REGION

 ℓ_2 -norm Ball

$$\mathcal{T}^{(k)} = \{ x \in \mathcal{R}^n \mid ||x - x^{(k)}||_2 \le \rho \}.$$

Box

$$\mathcal{T}^{(k)} = \{ x \in \mathcal{R}^n \mid |x_i - x_i^{(k)}| \le \rho_i, \ i = 1, \dots, n \}.$$

In the latter, if x_i appears only in convex objectives and constraints, we can set $\rho_i = \infty$.

AFFINE AND CONVEX APPROXIMATIONS

 \widehat{f} either takes an affine (first–order) Taylor approximation

$$\widehat{f}(x) = f(x^{(k)}) + \Delta f(x^{(k)})^T (x - x^{(k)})$$

or the convex part of the second order Taylor expansion

$$\widehat{f}(x) = f(x^{(k)}) + \Delta f(x^{(k)})^T (x - x^{(k)}) + \frac{1}{2} (x - x^{(k)})^T P(x - x^{(k)})$$

where P is the positive semidefinite part of the Hessian $\Delta^2 f(x^{(k)})$. That is, if

$$\Delta^2 f(x^{(k)}) = U \Lambda U^T \quad \Longrightarrow \quad P = U[\Lambda]_+ U^T.$$

A NUMERICAL EXAMPLE

Non-convex Quadratic Program

where *P* is symmetric but not positive semindefinite.

Second-order Taylor (convex) approximation

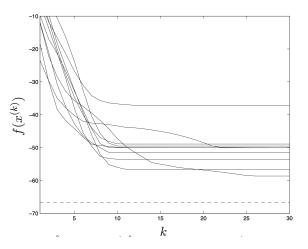
$$\widehat{f}(x) = f(x^{(k)}) + (Px^{(k)} + q)^{T}(x - x^{(k)}) + \frac{1}{2}(x - x^{(k)})^{T}P_{+}(x - x^{(k)})$$

Trust Region

$$\mathcal{T}^{(k)} = \{ x \in \mathcal{R}^n \mid ||x - x^{(k)}||_{\infty} \le \rho \}.$$

A NUMERICAL EXAMPLE

Example with n=20 and $\rho=0.2$, and different initial points.



SEQUENTIAL CONVEX PROGRAMMING ISSUES

Each iteration is, in principle, easy to solve as it is a convex problem. However, numerous issues arise:

Infeasibility: How to decide when to accept a step by taking into account:

- constraint violations,
- trade-off between feasibility of constraints and quality of the objective.

Trust region size: How large trust region should be.

- \bullet ρ too large, poor approximations
- \bullet $\,\rho$ too small, more accurate approximation, but slow.

INFEASIBILITY ISSUE VIA PENALIZATION

Workaround: Assign penalties to constraint violations rather than to directly enforce the constraints.

Replace the original problem with a penalized approximation

minimize
$$\phi(x) = f_0(x) + \lambda \left(\sum_{i=1}^m [f_i(x)]_+ + \sum_{j=1}^p |h_j(x)| \right)$$

This approach can often allow progresses that would otherwise be impossible.

For λ large enough, this turns into exact penalization. The solutions will satisfy the constraints.

INFEASIBILITY ISSUE VIA PENALIZATION

So, we will replace the update in SCP by:

$$\begin{split} \tilde{x} = & \text{ arg min } & \hat{\phi}(x) = \hat{f}_0(x) + \lambda \left(\sum_{i=1}^m [\hat{f}_i(x)]_+ + \sum_{j=1}^p |\hat{h}_j(x)| \right) \\ & \text{ subject to } & x \in \mathcal{T}^{(k)}. \end{split}$$

If this provides a good solution, we accept it and set

$$x^{(k+1)} = \tilde{x}$$

TRUST REGION SIZE ISSUE

Adjust the radius of the trust region considering how good is the approximation.

The trust region is:

$$\mathcal{T}^{(k)} = \{ x \in \mathcal{R}^n \mid ||x - x^{(k)}||_2 \le \rho \}.$$

Predicted Decrease:

$$\widehat{\delta} = \phi\left(x^{(k)}\right) - \widehat{\phi}(\widetilde{x}).$$

Actual Decrease:

$$\delta = \phi\left(x^{(k)}\right) - \phi(\tilde{x}).$$

TRUST REGION SIZE ISSUE

Let $\alpha \in (0,\frac{1}{2})$ and $\beta^{\rm succ} > 1$ and $\beta^{\rm fail} < 1.$

Algorithm

- If $\delta \geq \alpha \widehat{\delta}$ (i.e., sufficient decrease),
 - accept the solution,
 - ② set $x^{(k+1)} = \tilde{x}$,
 - **1** enlarge the trust region by $\rho^{(k+1)} = \beta^{\text{succ}} \rho^{(k)}$.
- Otherwise,
 - reject the step,
 - $oldsymbol{0}$ shrink the trust region by $ho^{(k+1)}=eta^{\mathrm{fail}}
 ho^{(k)},$
 - \odot obtain \tilde{x} again.