

Introduction to Stochastic Processes

Assignment 3
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3.1 Exercise

Given an undirected graph $G = (V, E)$, what are the requirements for the simple random walk on it to

- (1) be irreducible?
- (2) contain a transient state?
- (3) be 2-periodic?
- (4) be 3-periodic?
- (5) be aperiodic?

- (1) A Markov chain is irreducible if and only if it is connected (as S is by definition closed). Since for every $a, b \in V$, $a \rightarrow b$ if and only if a and b are connected in G (as if there exists a path of length n then $P^n(a \rightarrow b) > 0$). Thus in order for the walk to be irreducible, G must be connected (these are equivalent).
- (2) Suppose G contains a transient state $a \in V$. Then define $C = \{b \in V \mid a \rightarrow b\}$, and since $C \subseteq V$, C is finite. And C must be closed: if $b \in C$ and $b \rightarrow b'$ then $a \rightarrow b \rightarrow b'$ and so $a \rightarrow b'$ meaning $b' \in C$. Therefore this means C contains a recurrent state as a finite closed set. But if $b \in C$ is recurrent, since $a \rightarrow b$ and therefore $b \rightarrow a$ since G is undirected, we must have that a is recurrent, in contradiction. So G cannot have a transient state.
- (3) We claim that G is 2-periodic if and only if it contains no odd cycles. Any cycle containing a of length n means that $n \in \tau(a) = \{t \geq 1 \mid P^t(a \rightarrow a) > 0\}$, and so $d(a) \mid n$. So if G is 2-periodic, it cannot contain odd cycles as 2 must divide the length of any cycle. And if G contains no odd cycles, then for every $a \in V$, $\{a, a\}$ cannot be an edge as it is a 1-cycle, and so a must be connected to some $b \neq a$. But then a, b, a is a 2-cycle and so $P^2(a \rightarrow a) > 0$ and therefore $2 \in \tau(a)$. And since all cycles are even, this means that $\tau(a)$ contains only even numbers, so $d(a) = 2$.
- (4) G must contain two vertices a, b such that $\{a, b\} \in E$ (a and b need not be distinct) then a, b, a is a path and so $2 \in \tau(a)$. Thus $d(a) \neq 3$ as it doesn't divide 2.
- (5) We already showed that $2 \in d(a)$ for every $a \in V$. So for every $a \in V$, either $d(a) = 2$ or $d(a) = 1$. So G is aperiodic if and only if there exists a vertex a whose period is 1 (for the converse, if $d(a) = 1$ then either all vertices have degree one or some vertices have different periods, and in either case the graph is then aperiodic). This is if and only if there exists a path of odd length (since $2 \in \tau(a)$, $d(a) = 2$ if and only if there exists an odd number in $\tau(a)$, which corresponds to an odd cycle). So G is aperiodic if and only if there exists an odd cycle.

3.2 Exercise

Cookie Monster is collecting cookies. At every step, he either buys a new cookie or eats all the cookies he has collected, each with the same probability (when he has no cookies, he buys one with probability 1).

- (1) What is the probability that within a finite amount of time, he will have half a million cookies?
- (2) If Cookie Monster starts with 3 cookies, what is the expected time that it will take him to return to having 3 cookies?
- (3) One day Cookie Monster decides that he will limit himself to only a million cookies, and once he gets to a million cookies, he eats them all. How will your previous answers change?

- (1) Let us define X_n to be the number of cookies Cookie Monster has on the n th step. So

$$\mathbb{P}(X_n = i + 1 \mid X_{n-1} = i) = \frac{1}{2}, \quad \mathbb{P}(X_n = 0 \mid X_{n-1} = i) = \frac{1}{2}, \quad \mathbb{P}(X_n = 1 \mid X_{n-1} = 0) = 1 \quad (i > 0)$$

This is asking for $\mathbb{P}\left(T_{\frac{1}{2}\text{mil}} \mid X_0 = 0\right)$. We claim that all states (which are $\mathbb{N}_{\geq 0}$) are recurrent, meaning that this probability is 1. Since 0 is connected to all the states, it is sufficient to show that 0 is recurrent. Notice that the probability that $T_0 > n$ is $\frac{1}{2}^n$, and by the continuity of probability, $\mathbb{P}(T_0 = \infty \mid X_0 = 0) = 0$. Thus $\mathbb{P}(T_0 < \infty \mid X_0 = 0) = 1$ and so 0 is recurrent, therefore so is half a million, so the probability is one.

- (2) Let us define $e_k = \mathbb{E}[T_k \mid X_0 = 0]$. Let us denote $t := T_k$, and so $(\mathbb{E}_t[T_{k+1}] := \mathbb{E}[T_{k+1} \mid T_k])$, and all probability is implicitly under the assumption that $X_0 = 0$,

$$\mathbb{E}_t[T_{k+1}] = \frac{1}{2} \mathbb{E}_t[T_{k+1} \mid X_{t+1} = k_1] + \frac{1}{2} \mathbb{E}_t[T_{k+1} \mid X_{t+1} = 0] = \frac{1}{2}(t+1) + \frac{1}{2} \mathbb{E}_t[T_{k+1} \mid X_{t+1} = 0]$$

Now, $\mathbb{E}_t[T_{k+1} \mid X_{t+1} = 0] = \mathbb{E}_t[T_{k+1}] + t + 1$ since this is like starting over but having “wasted” $t + 1$ steps (due to homogeneity). Thus

$$\mathbb{E}_t[T_{k+1}] = \frac{1}{2} \mathbb{E}_t[T_{k+1}] + t + 1 \implies \mathbb{E}_t[T_{k+1}] = 2t + 2$$

By the law of total expectation, we then get $\mathbb{E}[T_{k+1}] = 2\mathbb{E}[T_k] + 2$, so $e_{k+1} = 2e_k + 2$. Solving this recurrence we get $e_k = 2^k + 2^{k-1} - 2$.

Now, let us define $\tilde{e}_k = \mathbb{E}[T_k \mid X_0 = 3]$, so we are attempting to compute \tilde{e}_3 . Then we get that, now with probability being conditioned under $X_0 = 3$,

$$\mathbb{E}_t[T_{k+1}] = \frac{1}{2} \mathbb{E}_t[T_{k+1} \mid X_{t+1} = k_1] + \frac{1}{2} \mathbb{E}_t[T_{k+1} \mid X_{t+1} = 0] = \frac{1}{2}(t+1) + \frac{1}{2}(e_{k+1} + t + 1) = t + 1 + \frac{1}{2}e_{k+1}$$

Thus again by the law of total expectation, $\tilde{e}_{k+1} = \tilde{e}_k + 2^k + 2^{k-1}$. Now, notice that $\mathbb{P}_3(T_0 = k) = \frac{1}{2^k}$ since in order for $T_0 = k$, Cookie Monster must buy $k - 1$ cookies and then eat them all, so $\frac{1}{2^{k-1}} \cdot \frac{1}{2}$. Thus $T_0 \mid X_0 = 3 \sim \text{Geo}(\frac{1}{2})$ and so $\tilde{e}_0 = 2$. And so we get from our recurrence that $\tilde{e}_1 = 3, \tilde{e}_2 = 6, \tilde{e}_3 = 12$. So the answer is 12.

- (3) For the first subquestion, the answer will not change as Cookie Monster will reach half a million cookies before one million, and so all the probabilities will remain the same. For the second subquestion, all the computations remain valid until computing $\mathbb{P}_3(T_0 = k)$. While for $k \leq 1\text{mil} - 2$ this remains the same, for $k = 1\text{mil} - 1$ this is just $\mathbb{P}_3(T_0 = k) = \frac{1}{2^{k-1}}$, as the probability of returning to zero is one. And for larger k it is zero. Thus the difference between this new $\mathbb{E}_3[T_0]$ by

$$\frac{a}{2^{a-1}} + \sum_{k=a+1}^{\infty} \frac{k}{2^k} \quad (a = 1\text{mil} - 1)$$

this is very very small, so the change in the answer is negligible.

3.3 Exercise

Let $p \in [0, 1]$ and X_n be a random walk on \mathbb{Z} where the transition probabilities are

$$P(i, j) = \begin{cases} p & j = i + 1, i > 0 \\ 1 - p & j = i - 1, i > 0 \\ \frac{1}{2} & j = i + 1, i \leq 0 \\ \frac{1}{2} & j = i - 1, i \leq 0 \\ 0 & |i - j| \neq 1 \end{cases}$$

Determine, for each p , whether or not 0 is transient or recurrent.

If $p > \frac{1}{2}$, then 0 must be transient. This is as $0 \rightarrow 1$ and 1 is transient:

$$\mathbb{P}(T_1 < \infty \mid X_0 = 1) = p \mathbb{P}(T_1 < \infty \mid X_1 = 2) + (1 - p) \mathbb{P}(T_1 < \infty \mid X_1 = 0)$$

The probability being summed on the left is the same if we replaced X_n with a random walk on \mathbb{Z} of probability p : all the probabilities until it hits 1 again are the same. Now since $p > \frac{1}{2}$, if we were to replace this with a random walk of probability p , $\mathbb{P}(T_1 < \infty \mid X_1 = 2) \leq \mathbb{P}(T_1 < \infty \mid X_1 = 1)$. And since all states in a random walk of probability p are transient, this means that this is less than 1, and so $\mathbb{P}(T_1 < \infty \mid X_0 = 1) < 1$ so 1 is transient and therefore so is 0.

If $p \leq \frac{1}{2}$, then the probability of returning to zero if we go to the right is more than it would be on a fair walk. And the left is a fair walk. We showed last assignment that in a fair walk, all states are recurrent, so 0 must be a recurrent state.