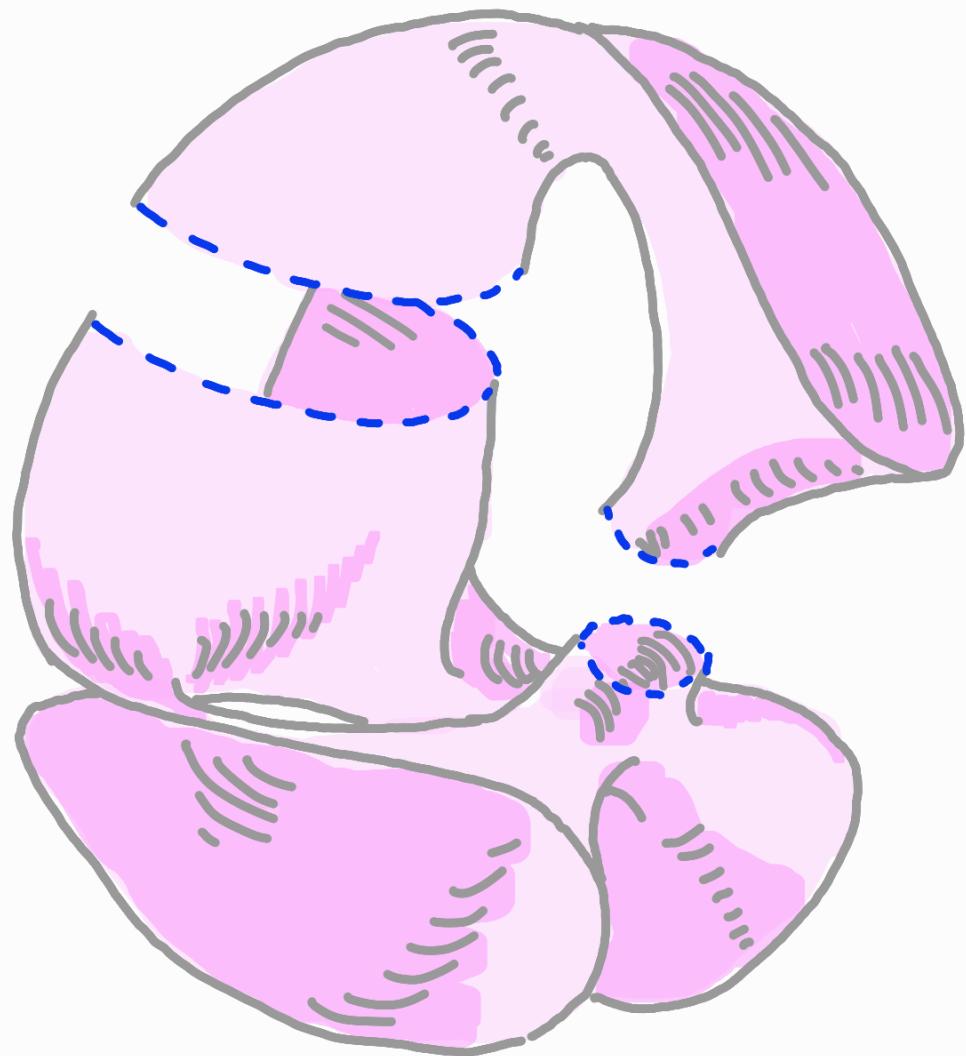


ON CATEGORIES  
OF ENRICHED  
SHEAVES



ARI ROSENFIELD  
UNIV. OF CALIFORNIA,  
IRVINE

## OVERVIEW

TWO PROJECTS :

I . ENRICHED SHEAVES UNDER  
CHANGE OF ENRICHING BASE

II . (J/W ANA TENDRIO) CLASSIFYING  
CATEGORIES OF ENRICHED SHEAVES

# OVERVIEW

UNIFYING THEME:

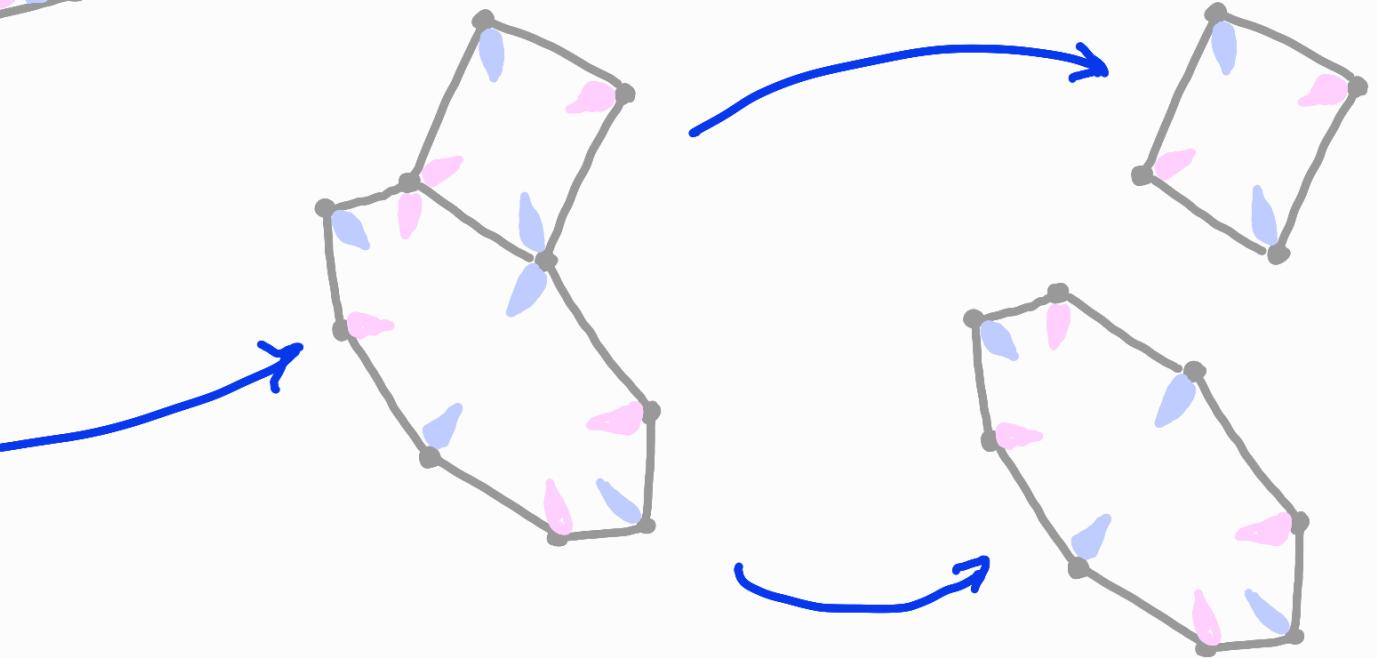
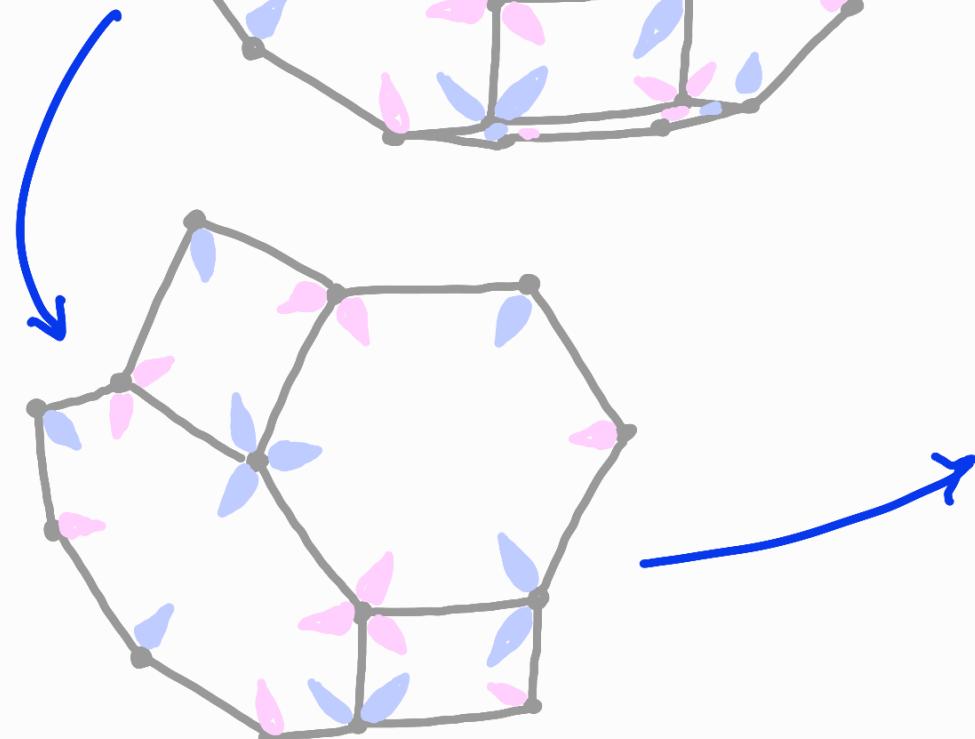
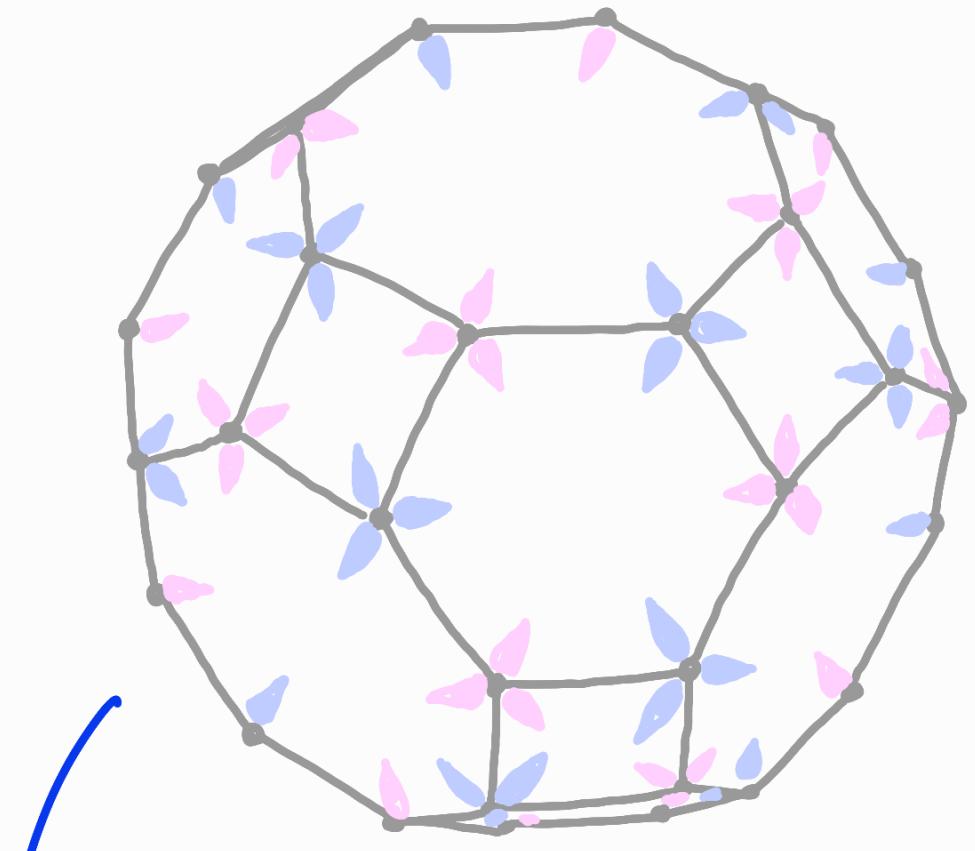
ENRICHED GROTHENDIECK TOPOSES.

## OUR PERSPECTIVE

A GROTHENDIECK TOPOS IS A CATEGORY EQUIVALENT TO SOME CATEGORY OF SHEAVES.

AN ENRICHED GROTHENDIECK TOPOS IS AN ENRICHED CATEGORY EQUIVALENT TO SOME CATEGORY OF ENRICHED SHEAVES.

# I. ENRICHED SHEAVES AND COVERAGES UNDER CHANGE OF BASE



arXiv : 2405.19529

# I. A. INTRODUCTION

## NOTATION AND CONVENTIONS

$(\mathcal{V}, \otimes, *_\mathcal{V})$  AND  $(\mathcal{U}, \otimes, *_\mathcal{U})$  ARE  
ALWAYS SYMMETRIC MONOIDAL CLOSED  
AND LOCALLY SMALL.

$\mathcal{C}$  IS ALWAYS A SMALL (ESSENTIALLY  
SMALL)  $\mathcal{V}$ -ENRICHED CATEGORY.

## ENRICHMENT

THINK OF AN ENRICHED CATEGORY AS A  
CATEGORY WHOSE HOMS HAVE MORE  
GENERAL STRUCTURE THAN THAT OF A  
SET.

ENRICHED FUNCTORS, TRANSFORMATIONS  
RESPECT THIS STRUCTURE.

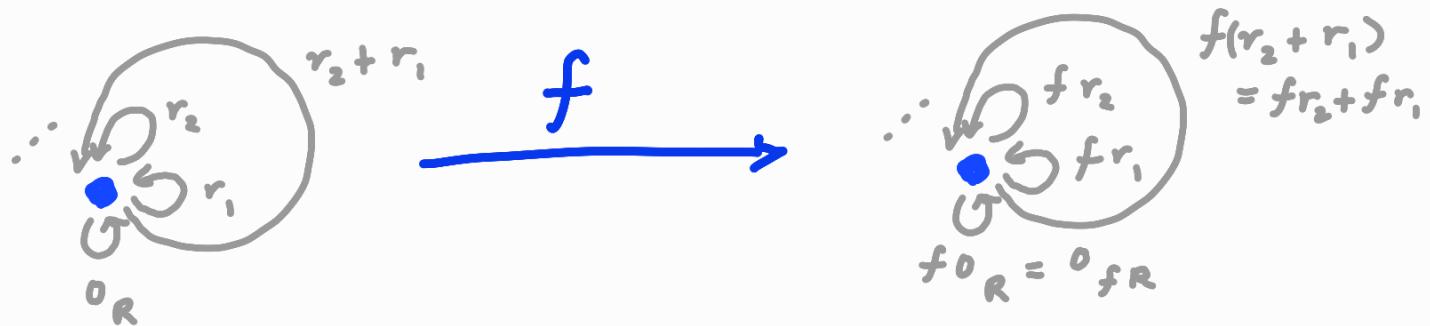
## ENRICHMENT

e.g.

A RING IS A ONE-OBJECT Ab - CATEGORY.

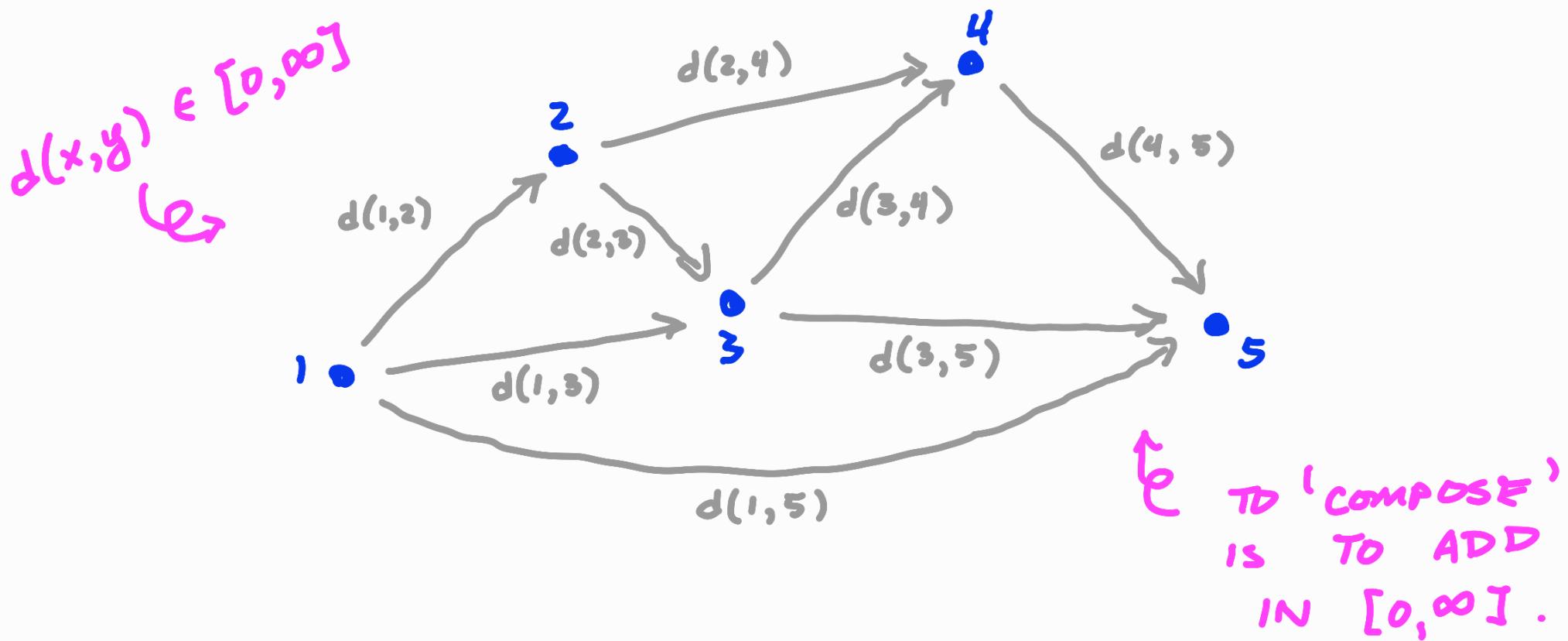
A RING HOMOMORPHISM IS AN

Ab - FUNCTOR.



## ENRICHMENT

e.g.) WE CAN ENRICH OVER MONOIDAL PREORDERS! A  $[0, \infty]$  - CATEGORY IS A LAWVERE METRIC SPACE.



## SHEAVES

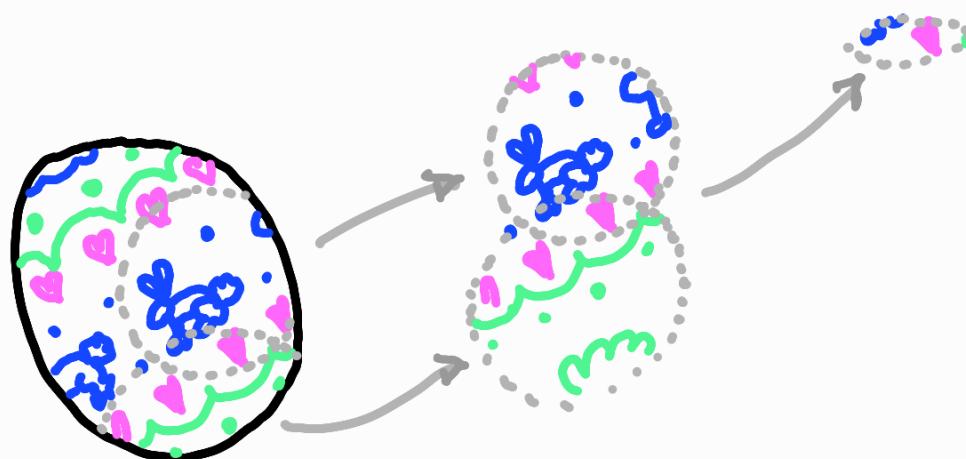
A GROTHENDIECK TOPOLOGY IS A WAY TO ENDOW ANY CATEGORY (NOT JUST  $\mathcal{O}(X)$ ) WITH A NOTION OF WHEN A GIVEN OBJECT IS "COVERED" BY OTHERS.

A COVERAGE ON A CATEGORY IS A "WEAK" GROTHENDIECK TOPOLOGY.



SHEAVES

WE WANT TO KNOW WHEN OBJECTS COVER ONE  
ANOTHER BECAUSE THEN WE CAN DEFINE  
SHEAVES : "COVER - RESPECTING" ASSIGNMENTS  
OF DATA TO THE OBJECTS OF OUR CAT'S .



# ENRICHED SHEAVES!

WE HAVE ENRICHED GROTHENDIECK TOPOLOGIES  
AND ENRICHED SHEAVES, TOO (SAY, OVER  $\mathcal{V}$ ).

e.g., A GABRIEL LOCALIZING SYSTEM  
ON A RING IS AN AB - TOPOLOGY.  
AN AB - SHEAF ON A RING w/r/t  
SUCH A LOCALIZING SYSTEM IS ITS  
GABRIEL LOCALIZATION.

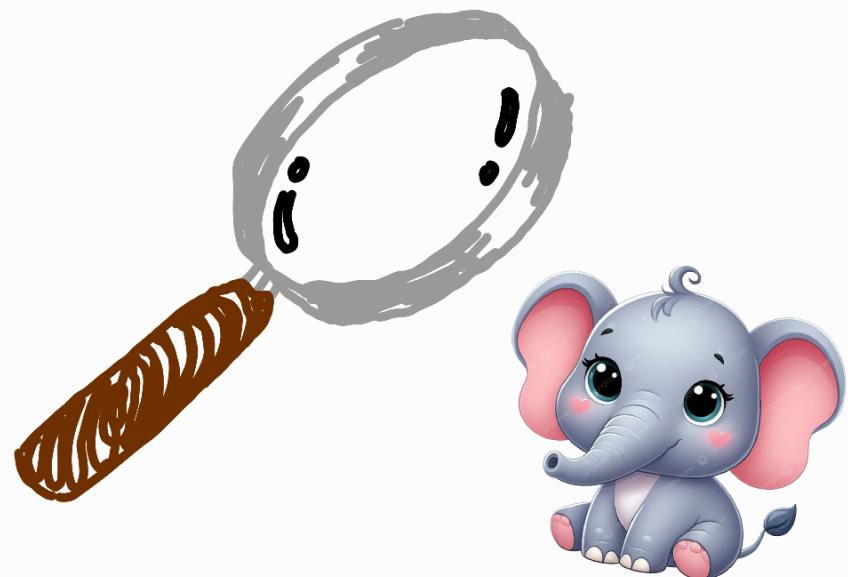


LEARN MORE ON  
THIS LATER.

## HISTORY

CAT'S OF  
3 WAYS TO STUDY  $\mathcal{V}$ -SHEAVES ON  $\mathcal{C}$ :

- (1) "NICE"  $\mathcal{V}$ -SUBCATEGORIES (LOCALIZATIONS) OF  $[\mathcal{C}^{\text{op}}, \mathcal{V}]$ ;
- (2)  $\mathcal{V}$ -GROTHENDIECK TOPOLOGIES ON  $\mathcal{C}$ ;
- (3) UNIVERSAL CLOSURE OPERATIONS ON  $[\mathcal{C}^{\text{op}}, \mathcal{V}]$ .



## HISTORY

BORCEUX - QUINTEIRO, 1996 -

FOR "NICE ENOUGH"  $\nu$ , THERE ARE

BIVIJECTIONS

$$(1) \leftrightarrow (2)$$

$$(2) \leftrightarrow (3)$$

$$(3) \leftrightarrow (1).$$

## HISTORY

BORCEUX - QUINTEIRO, 1996 -

$\mathcal{V}$  = Set is "NICE ENOUGH", so

WE RECOVER THE CLASSICAL  
THEORY IN THAT CASE.

## HISTORY

WE'LL FOCUS ON THE  
"NICE" SUBCATEGORY AND  
 $\mathcal{V}$ -TOPOLOGY PERSPECTIVES.

## ARCHAEOLOGY?

BY 1996, ENRICHED SHEAF THEORY HAD  
ALREADY BEEN DEVELOPED FOR  $\mathcal{V} = \text{Ab}$   
AND  $\mathcal{V} = \text{gr Mod}_K$ . (STENSTRÖM'S "RINGS  
OF QUOTIENTS" IS A TEXTBOOK ON  $\mathcal{V} = \text{Ab}!$ )

BORCEUX - QUINTERO'S WORK HASN'T  
BEEN WIDELY USED TO DATE OUTSIDE  
OF THESE CASES.

ARCHAEOLOGY?

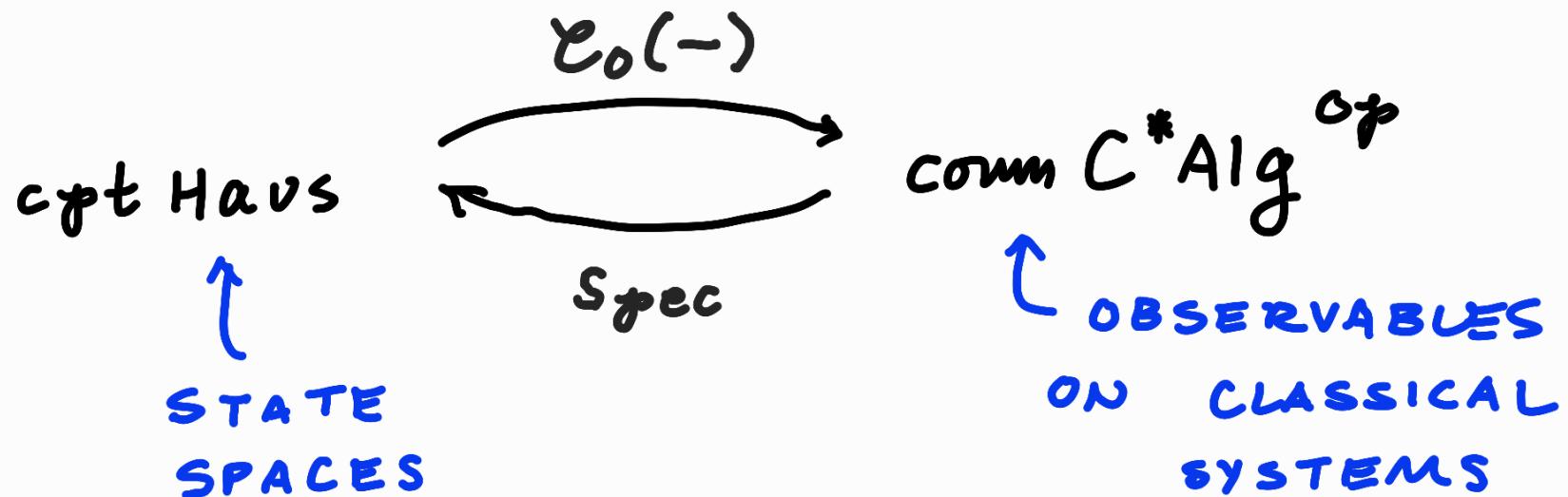
OK, SO WHY DO YOU  
CARE ABOUT IT THEN?



WHY  
REVIVE  
THIS  
TOPIC ?

# NONCOMMUTATIVE SPECTRAL THEORY

CLASSICAL GEL'FAND DUALITY: AN EQUIVALENCE OF CATEGORIES



## NONCOMMUTATIVE SPECTRAL THEORY

WHAT IF I WANT GEL'FAND DUALITY  
FOR NON COMMUTATIVE  $C^*$ -ALGEBRAS?  
(i.e., OBSERVABLES ON QUANTUM SYSTEMS?)

THAT WOULD BE SO COOL...

THEN I COULD DO  
ACTUAL GEOMETRY...

I COULD TRY JUST APPLYING THE  $\text{Spec}$   
FUNCTOR I ALREADY HAVE ??

# NONCOMMUTATIVE SPECTRAL THEORY

PROBLEM! (REYES, 2012)

THEOREM: ANY FUNCTOR  $F: \text{C}^*\text{Alg}^{\text{op}} \rightarrow \text{Top}$   
FOR WHICH



$$\begin{array}{ccc} \text{comm } \text{C}^*\text{Alg}^{\text{op}} & \xrightarrow{\text{Spec}} & \text{Top} \\ \text{inc} \downarrow & & \\ \text{C}^*\text{Alg}^{\text{op}} & \xrightarrow{F} & \end{array}$$



COMMUTES MUST HAVE  $F(M_n(\mathbb{C})) \cong \emptyset$   
FOR  $n \geq 3$ .

## NONCOMMUTATIVE SPECTRAL THEORY

IN PHYSICAL TERMS: GEL'FAND DUALITY SAYS  
ANY CLASSICAL SYSTEM (conn.  $C^*$ -Alg) IS  
DETERMINED BY A STATE SPACE (cpt. top. space)  
IN A WAY THAT RESPECTS OPERATIONS ON  
THE SYSTEM (FUNCTORIALLY).

THIS THEOREM SAYS THAT NO COMMUTATIVE  
 $C^*$ -ALG DETERMINES A QUANTUM SYSTEM IN  
THIS WAY.

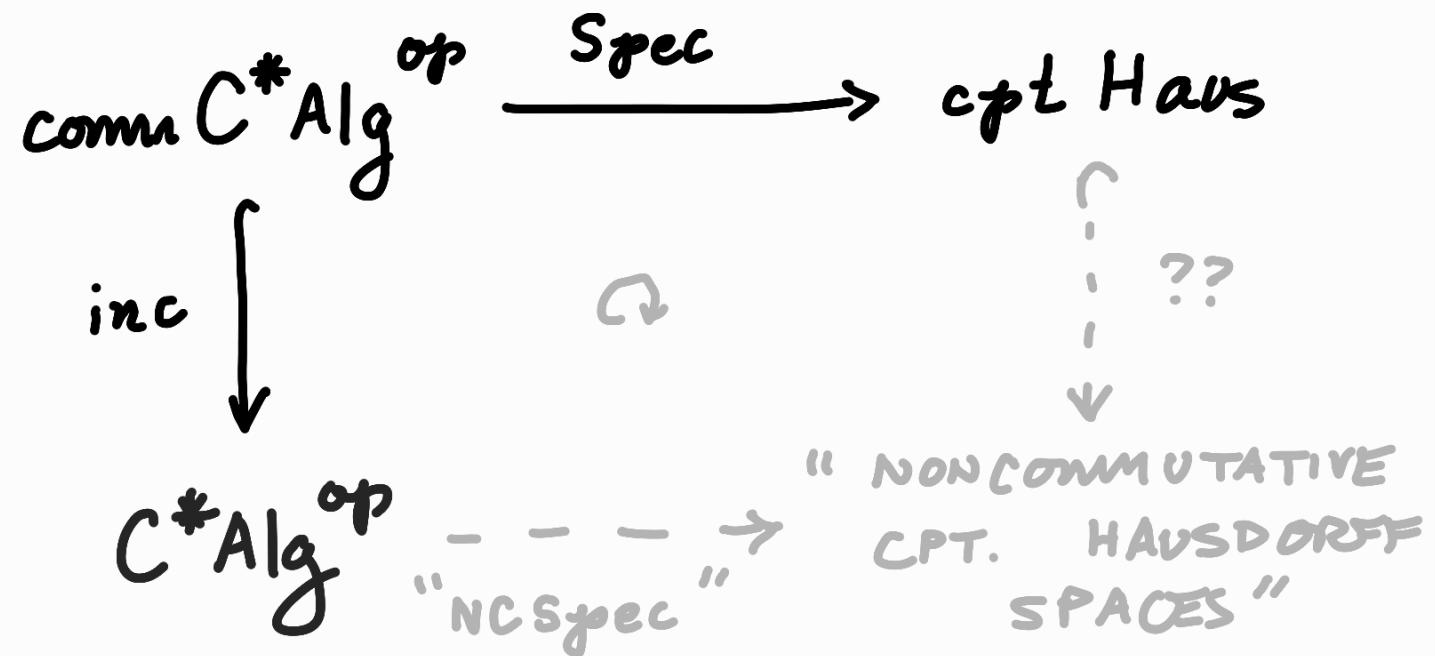
# NONCOMMUTATIVE SPECTRAL THEORY

RELATED OBSTRUCTIONS HOLD FOR

- (v. d. BERG-HEUNEN, 2014) ZARISKI, STONE, AND PIERCE SPECTRA AS FUNCTORS INTO  $\text{Set}$ ,  $\text{Loc}$ ,  $\text{Top}$ , Scheme, Quantale, Topos;
- (REYES, 2014) SUBCANONICAL COVERAGES ON THE BIG ZARISKI SITE FOR  $\text{cRing}$ .

# NONCOMMUTATIVE SPECTRAL THEORY

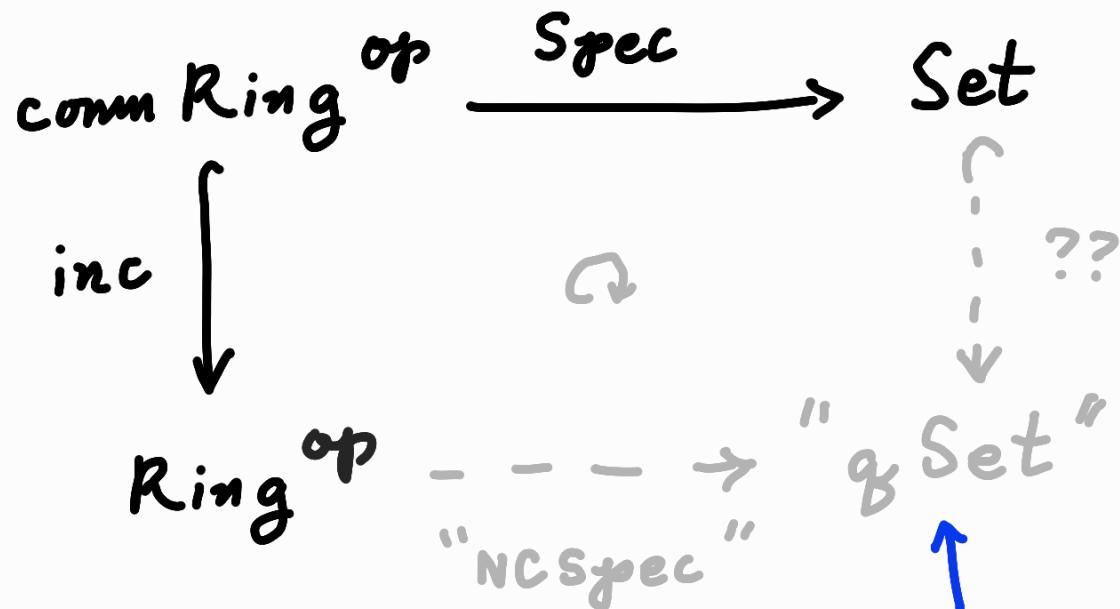
THE DREAM OF NONCOMMUTATIVE GEOMETRY :



# NONCOMMUTATIVE SPECTRAL THEORY

TOPOLOGICAL SPACES ARE SETS OF POINTS, SO I  
MAY AS WELL THINK ABOUT A SLIGHTLY SIMPLER

PROBLEM:



FINDING A CAT. OF  
"QUANTUM / N.C." POINT - SETS.

## NONCOMMUTATIVE SPECTRAL THEORY

WE EXPECT  $\mathcal{B}\text{Set}$  TO BE A CLOSED  
SYMMETRIC MONOIDAL CATEGORY, AND WE  
EXPECT A FULLY FAITHFUL EMBEDDING

$$\text{Set} \longrightarrow \mathcal{B}\text{Set}.$$

HOW TO TELL IF A GIVEN CANDIDATE FOR  
 $\mathcal{B}\text{Set}$ , SAY  $\mathcal{V}\text{e MonCat}$ , IS THE RIGHT ONE?

COOL, WHAT DOES IT HAVE TO DO WITH QUANTUM?

ONE WAY TO TEST WHETHER WE'VE FOUND A  
GOOD CANDIDATE FOR  $\mathbf{qSet}$  IS TO MAKE  
SURE THERE ARE  $\mathbf{qSet}$ -ENRICHED GROTH-  
ENDIECK TOPOLOGIES ON A GIVEN  $\mathbf{qSet}$ -CATEGORY  
WHICH DON'T ARISE FROM UNENRICHED ONES.

COOL, WHAT DOES IT HAVE TO DO WITH QUANTUM?

SO, NEED TOOLS TO DETECT WHETHER  
A GIVEN ENRICHED GROTHENDIECK TOPOLOGY  
ARISES FROM AN UNENRICHED ONE.

THAT'S WHERE BASE CHANGE COMES IN...

## BASE CHANGE

GIVEN A "NICE ENOUGH" FUNCTOR  $G: \mathcal{V} \rightarrow \mathcal{U}$ ,  
WE CAN CHANGE THE BASE OF  $\mathcal{C}$  VIA  $G$ ,  
OBTAINING A  $\mathcal{U}$ -CATEGORY  $G_* \mathcal{C}$ .

THINK OF  $\mathcal{U} : \mathcal{V}$  AS DIFFERENT TYPES  
OF STRUCTURE THE HOMS IN  $\mathcal{C}$  CAN  
HAVE.

## BASE CHANGE

FOR EXAMPLE, WE CAN USE THE FUNCTOR

$$\text{Hom}_{\text{Ab}}(\mathbb{Z}, -) : \text{Ab} \longrightarrow \text{Set}$$

TO TURN A 1-OBJECT Ab-CATEGORY  
(A RING) INTO A 1-OBJECT  
Set-CATEGORY (A MONOID).

## NEW CONTRIBUTIONS

HOW BASE CHANGE INTERACTS  
WITH ENRICHED TOPOLOGIES  
OR ENRICHED LOCALIZATIONS  
HASN'T YET BEEN  
ADDRESSED , UNTIL NOW!

I . B . A      CHANGE - OF - BASE  
CONSTRUCTION FOR  
PRE SHEAVES

## HYPOTHESIS

ENRICHING CATEGORIES  $\mathcal{U}, \mathcal{V}$  ARE ASSUMED  
TO BE:

- SYMMETRIC MONOIDAL CLOSED
- LOCALLY FINITELY PRESENTABLE
- BARR REGULAR ( $[\mathcal{C}^{\text{op}}, \mathcal{V}]_0$  IS REGULAR)
- FINITELY COMPLETE

WE'LL CALL SUCH A CATEGORY A COSMOS.

## LFP CATEGORIES

"LFP"

A LOCALLY FINITELY PRESENTABLE <sup>V</sup> CATEGORY IS  
ONE IN WHICH ANY OBJECT IS REALIZABLE AS A  
"WELL-BEHAVED" COLIMIT OF "FINITE" OBJECTS.

e.g., for  $x \in \text{Set}$ ,  $x \cong \bigcup_{\substack{S \subseteq x \\ \text{FINITE}}} S$ .

WE SAY  $\text{Ob}(\text{FinSet})$  IS A DENSE GENERATING  
FAMILY FOR  $\text{Set}$ .

## LFP CATEGORIES

FACT: ANY LFP CATEGORY W...

- ADMITS A DENSE GENERATING FAMILY  
 $\mathcal{G}_\omega$  MADE UP OF FINITELY PRESENTABLE OBJECTS.
- IS COCOMPLETE AND WELL-Powered.

## LFP CATEGORIES

WHO CARES ABOUT A GENERATING FAMILY?

IDEA. IF  $\mathcal{C} \in \mathcal{V}\text{-Cat}$ , so  $\mathcal{C}(x,y) \in \text{Ob}(\mathcal{V})$ ,

AND  $\mathcal{V} \neq \text{Set}$ , so  $*_{\mathcal{V}} \neq \{\cdot\}$ , WHAT'S AN

"ELEMENT" OF  $\mathcal{C}(x,y)$ ? WE HAVE TO

CONSIDER GENERALIZED ELEMENTS INSTEAD:

MOE'S  $g \rightarrow \mathcal{C}(x,y)$  FOR  $g \in \mathcal{G}_{\mathcal{V}}$ .

"cosmoi"

EXAMPLES OF COSMOI:

Set, Ab, RMod (FOR COMM. RINGS  $R$ ),  
grMod $_k$  (FOR FIELDS  $k$ ), ANY  
PRESHEAF TOPOS.

"ALMOST EXAMPLES": THE MONOIDAL PREORDERS  
 $[0,1]$  AND  $[0,\infty]$  (NOT LFP),  
Cat (NOT REGULAR).

"cosmo!"

WHY CONSIDER THESE "ALMOST EXAMPLES?" IT  
WAS HARD TO FIND EXAMPLES OF BASE CHANGE  
THAT WEREN'T BASICALLY "FORGETFUL" FUNCTORS.  
  
PLUS, NEITHER LFP NOR REGULARITY ARE NEEDED  
FOR THE DEFINITION OF ENRICHED TOPOLOGY.

GOAL! WANT TO FIGURE OUT HOW TO  
CHANGE THE BASE OF A  $\mathcal{V}$ -TOPOLOGY.

WHAT ARE  $\mathcal{V}$ -TOPOLOGIES : BASE CHANGE,  
REALLY?

# WHAT IS A GROTHENDIECK TOPOLOGY?

BORCEUX - QUINTEIRO, 1996 :

$\mathcal{V}$ -  $\mathcal{V}$ -  
**Definition 1.2** Let  $\mathcal{C}$  be a small category. A Grothendieck topology on  $\mathcal{C}$  is the choice, for every object  $C \in \mathcal{C}$ , of a family  $T(C)$  of subobjects of the representable  $\mathcal{V}$ -presheaf  $\mathcal{C}(-, C)$ . Those data must satisfy the following axioms:

$f: G \rightarrow \mathcal{C}(D, C)$  FOR  
 $G \in \mathcal{G}_\mathcal{V}$

(T1)  $\mathcal{C}(-, C) \in T(C)$  for every object  $C \in \mathcal{C}$ ;

(T2) given  $R \in T(C)$  and  $f \in_G \mathcal{C}(D, C)$ , one has  $f^{-1}(R) \in T(D)$ , where  $f^{-1}(R)$  is defined by the following pullback:

$$\begin{array}{ccc} f^{-1}(R) & \longrightarrow & \{G, R\} \\ \downarrow & & \downarrow \\ \mathcal{C}(-, D) & \xrightarrow{f} & \{G, \mathcal{C}(-, C)\}; \end{array}$$

$\leftarrow$  COTENSOR IN  
 $[\mathcal{C}^{\text{op}}, \mathcal{V}]$

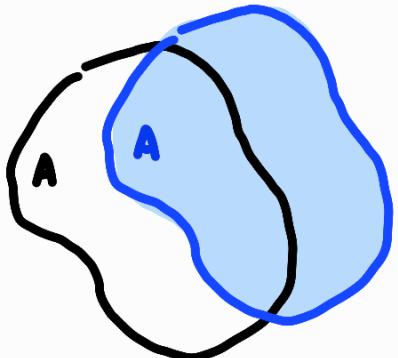
(T3) given  $S \in T(C)$  and a subobject  $R \rightarrowtail \mathcal{C}(-, C)$  such that  $f^{-1}(R) \in T(D)$  for all  $f \in_G S(D)$ , one has  $R \in T(C)$ .

# WHAT IS A GROTHENDIECK TOPOLOGY?

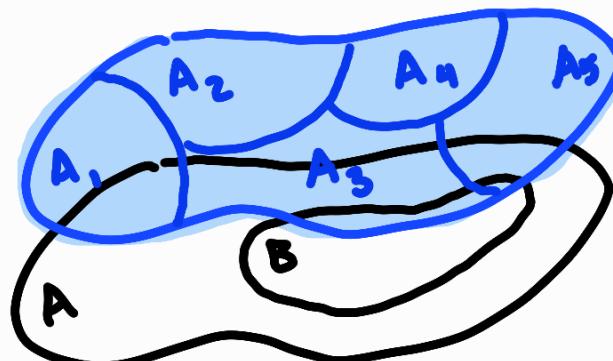
(T<sub>1</sub>) - (T<sub>3</sub>) ENCODE SOME FAMILIAR PROPERTIES  
OF OPEN COVERS IN TOPOLOGICAL SPACES X:

# WHAT IS A GROTHENDIECK TOPOLOGY?

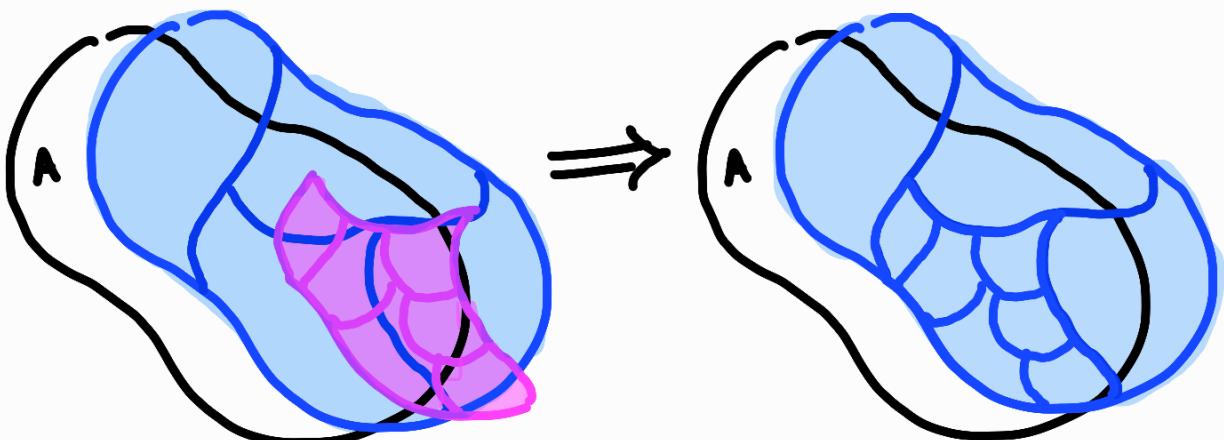
(T1)  $A \in \mathcal{O}(X)$  is an open cover of itself:



(T2) If  $A = \bigcup_i A_i$  and  $B \subset A$ , then  $B = \bigcup_i (B \cap A_i)$ :



(T3) If  $A = \bigcup_i A_i$  and  $A_i = \bigcup_j B_{ij} \quad \forall i$ , then  $A = \bigcup_{i,j} B_{ij}$ .



# WHAT IS A GROTHENDIECK TOPOLOGY?

A SIEVE ON  $x \in \text{Ob}(\mathcal{C})$  IS A SUB-  
OBJECT  $R \rightarrowtail \mathcal{C}(-, x)$  — i.e., AN EQUIV-  
ALENCE CLASS OF MONOMORPHISMS INTO  
 $\mathcal{C}(-, x)$ . THESE FORM A COMPLETE  
LATTICE,  
ABUSE!  
 $\text{Sub}(\mathcal{C}(-, x))$ .

"AN OPEN COVER OF  $x$ ."

## WHAT IS A GROTHENDIECK TOPOLOGY?

SAY THAT A FAMILY  $\mathcal{J}$  OF SIEVES IS A  
COVERAGE ON  $\mathcal{C}$  IF IT SATISFIES  
(T1) AND (T2).

DENOTE

$$\Sigma(\mathcal{C}, \mathcal{V}) = \{ \text{ALL } \mathcal{V}\text{-COVERAGES ON } \mathcal{C} \}$$

$$\tau(\mathcal{C}, \mathcal{V}) = \{ \text{ALL } \mathcal{V}\text{-TOPOLOGIES ON } \mathcal{C} \}$$

## COVERAGES VS. TOPOLOGIES

CONDITION ( $T_3$ ) WAS TECHNICALLY  
DIFFICULT TO WORK WITH UNDER BASE  
CHANGE.

CONJECTURE. ANY  $\mathcal{V}$ -COVERAGE UNIQUELY  
DETERMINES A  $\mathcal{V}$ -TOPOLOGY WITH THE SAME  
SHEAVES.

## WHAT IS A GROTHENDIECK TOPOLOGY?

e.g., FOR A RING  $A$ , AN AB-SIEVE ON  $A$ 'S ONE OBJECT IS A RIGHT IDEAL

OF  $A$ . AN AB-COVERAGE  $J$  ON  $A$  IS A FAMILY OF RIGHT IDEALS

S.T.

$$(T1) A \in J$$

"CLOSURE UNDER  
IDEAL QUOTIENT"

$$(T2) \forall a \in A, \forall I \in J, \quad \checkmark$$

$$(I : a) := \{x \in A : xa \in I\} \in J$$

## WHAT IS A GROTHENDIECK TOPOLOGY?

e.g., let  $\mathcal{X}$  be a Lawvere metric space with distance function  $d(-, -)$ .

A  $[0, \infty]$ -sieve on  $x \in \mathcal{X}$  is a Lipschitz function  $f: \mathcal{X} \rightarrow [0, \infty]$  such that  $\forall y, z,$

$$(i) f_z \geq d(z, x)$$

$$(ii) d(y, z) \geq \max \{0, d(y, x) - f_z\}.$$

These say that  $f$  is  $[0, \infty]$ -natural.

$$[0, \infty] \xrightarrow{f} d(-, x)$$

## WHAT IS A GROTHENDIECK TOPOLOGY?

A  $[0, \infty]$ -COVERAGE ON  $X$  IS,  $\forall x$ , A  
COLLECTION  $J(x)$  OF SIEVES S.T.

$$(T1) d(-, x) \in J(x)$$

(T2)  $\forall y \in X$ ,  $\forall q \in [0, \infty]$  WITH  $q \geq \max\{0, y-x\}$ ,  
AND  $\forall f \in J(x)$ , THE FUNCTION

$$f_q(z) := \max\{fz, \max\{0, d(z, y) - q\}\}$$

IS IN  $J(y)$ .

WHAT IS A GROTHENDIECK TOPOLOGY?

IN SHORT :

"A FAMILY OF SIEVES  
SATISFYING SOME CLOSURE  
CONDITIONS."

## BASE CHANGE

Fix  $\mathcal{U}, \mathcal{V}$ , AND  $\overset{\text{LAX MONOIDAL}}{\Downarrow} G: \mathcal{V} \rightarrow \mathcal{U}$ .

$G$  DETERMINES A 2-FUNCTOR

$$\mathcal{V}\text{-Cat} \xrightarrow{G_*} \mathcal{U}\text{-Cat}.$$

GIVEN A  $\mathcal{V}$ -CATEGORY  $\mathcal{C}$ , THE  $\mathcal{U}$ -CATEGORY

$G_* \mathcal{C}$  HAS

$$\text{Ob}(G_* \mathcal{C}) := \text{Ob}(\mathcal{C}),$$

$$G_* \mathcal{C}(x, y) := G(\mathcal{C}(x, y)).$$



## BASE CHANGE

GIVEN A  $\mathcal{V}$ -FUNCTOR  $A: \mathcal{X} \rightarrow \mathcal{Y}$ , DENOTE  
THE CORR.  $\mathcal{U}$ -FUNCTOR BY

$$G_* A : G_* \mathcal{X} \longrightarrow G_* \mathcal{Y}$$

( $G_*$  ALSO ACTS ON  $\mathcal{V}$ -NATURAL TRANSFORMATIONS).

## HYPOTHESIS

WE CONSIDER LAX MONOIDAL FUNCTORS

$G: \mathcal{V} \rightarrow \mathcal{U}$  WHICH ARE

- i. FAITHFUL
- ii. CONSERVATIVE
- iii. RIGHT ADJOINTS (IN  $\text{MonCat}$ )

}  $\Delta$

# BASE CHANGE

EXAMPLES!

FORGETFUL FUNCTORS ON  
 $(G \circ F)$ -Mod FOR MONADS  
 $G \circ F$

(1) MONADIC FUNCTORS. AMONG THEM:

- FOR  $(f: R \rightarrow S) \in \text{cRing}$ , RESTRICTION OF SCALARS  $S\text{Mod} \xrightarrow{f^*} R\text{Mod}$ .
- FOR  $W \in \text{Cosmos}$ , THE FORGETFUL FTR.  $\text{Mon}(W) \rightarrow W$
- REFLECTIVE SUBCATEGORY INCLUSIONS.

## BASE CHANGE

ALMOST EXAMPLES!

(2) THE INCLUSION  $\{0,1\} \hookrightarrow [0,1]$ . A

$\{0,1\}$ -CATEGORY IS JUST A PREORDER.  
ENRICHING OVER  $[0,1]$  IS SIMILAR TO  $[0,\infty]$ .

(3) THE FUNCTIONS

-  $\log(*): [0,1] \hookrightarrow [0,\infty]: e^{-(*)}$ .

(IN FACT, THIS IS AN ISOMORPHISM  
OF CATEGORIES.)

## BASE CHANGE

### ALMOST EXAMPLES!

(4) THE NERVE CONSTRUCTION  $\text{Cat} \rightarrow \text{sSet}$ .  
IT'S MONADIC, BUT THESE AREN'T  
COSMOI.

## BASE CHANGE FOR SIEVES

GOAL! WANT TO FIGURE OUT HOW TO  
CHANGE THE BASE OF A COVERAGE.

A COVERAGE IS MADE UP OF SIEVES,  
SO LET'S FIGURE OUT BASE CHANGE  
FOR SIEVES.

## BASE CHANGE FOR SIEVES

IF  $R: \mathcal{C}^{\text{op}} \rightarrow \mathcal{V}$  IS A  $\mathcal{V}$ -PRESHEAF, BASE  
CHANGE GIVES US A  $\mathcal{U}$ -FUNCTOR

$G_* R: G_* \mathcal{C}^{\text{op}} \rightarrow G_* \mathcal{V}$  ←  
- NOT QUITE A  $\mathcal{U}$ -PRESHEAF!

IT DOESN'T  
TAKE VALUES IN  
 $\mathcal{U}$ .

## BASE CHANGE FOR SIEVES

IN CASE  $G$  IS A <sup>V</sup> MONOIDAL RIGHT ADJOINT, WE HAVE AN INDUCED ADJUNCTION

$$G_* \mathcal{V} \begin{array}{c} \xrightarrow{G^u} \\[-1ex] \xleftarrow{T} \\[-1ex] \xrightarrow{F^u} \end{array} \mathcal{U}$$

IN  $\mathcal{U}\text{-Cat}$ :

## BASE CHANGE FOR SIEVES

SAY  $(F \dashv G, \epsilon, \eta)$  IS OUR ADJUNCTION AND  
 X, Y, Z ARE OBJECTS:

$$\text{Hom}_{\mathcal{U}}(x, \mathcal{U}(y, Gz)) \cong \text{Hom}_{\mathcal{U}}(x \otimes y, Gz)$$

HOM-TENSOR  
IN  $\mathcal{U}$

$$\cong \text{Hom}_{\mathcal{V}}(Fx \otimes Fy, z)$$

$F \dashv G$ , STRENGTH  
OF  $F$

$$\cong \text{Hom}_{\mathcal{U}}(x, G\mathcal{V}(Fy, z))$$

HOM-TENSOR  
IN  $\mathcal{V}$ ,  $F \dashv G$

$\dagger$  LEMMA  
 $\implies$

$$\phi: \mathcal{U}(y, Gz) \cong G\mathcal{V}(Fy, z) \quad \text{IN } \mathcal{U}.$$

## BASE CHANGE FOR SIEVES

SETTING

$$G^u_{yz} := GV(y, z) \xrightarrow{G(\epsilon^*)} GV(FGy, z) \xrightarrow{\phi} U(Gy, Gz)$$

DEFINES A  $U$ -FUNCTION WITH  $G^u(x) = Gx$  FOR  
 $x \in Ob(V)$ .

## BASE CHANGE FOR SIEVES

WE CAN THEN DEFINE

$$\tilde{G}R := G^u \circ G_* R : G_* \mathcal{C}^{op} \longrightarrow \mathcal{U}.$$

A  $\overset{\uparrow}{\text{PERFECTLY GOOD}}$   
 $\mathcal{U}$ -PRESHEAF.

## BASE CHANGE FOR SIEVES

PROPOSITION. SUPPOSE  $G$  SATISFIES ( $\Delta$ ).

THE OPERATION  $R \mapsto \tilde{G}R$  EXTENDS TO  
A FAITHFUL, CONSERVATIVE FUNCTOR

$$[e^{\text{op}}, \mathcal{V}] \xrightarrow{\tilde{G}} [G_* e^{\text{op}}, \mathcal{U}]$$

WHICH PRESERVES AND  
REFLECTS LIMITS.

↑  
LIKE 6 TECHNICAL LEMMATA  
REQUIRED TO PROVE  
THIS



## BASE CHANGE FOR SIEVES

IN PARTICULAR, SINCE  $\hat{G}$  PRESERVES LIMITS, IT  
INDUCES A FUNCTION

$$\text{Sub}_v(\mathcal{C}(-, \times)) \xrightarrow{\hat{G}} \text{Sub}_u(\tilde{G}\mathcal{C}(-, \times)).$$

## BASE CHANGE FOR SIEVES

MORE IS TRUE WHEN  $G$  SATISFIES ( $\alpha$ ):

THEOREM: THE FUNCTION

$$\text{Sub}(\mathcal{C}(-, x)) \xrightarrow{\tilde{G}} \text{Sub}(\tilde{G}\mathcal{C}(-, x))$$

IS AN INJECTIVE LATTICE MORPHISM.

" $G$  PRESERVES COARSENESS"

BASE CHANGE FOR COVERAGES

WE CAN EXTEND THIS RESULT TO  
COVERAGES AFTER ADDRESSING  
SOME TECHNICALITIES :

## BASE CHANGE FOR COVERAGES

THEOREM.  $\Sigma(\mathcal{C}, \mathcal{V})$  AND  $\mathcal{T}(\mathcal{C}, \mathcal{V})$  ARE  
COMPLETE LATTICES. ← SAME FOR  $G_* \mathcal{C}$ ,  
 $\mathcal{U}$ , ETC.

PROPOSITION. GIVEN A  $\mathcal{V}$ -COVERAGE  $J$  ON  $\mathcal{C}$ ,  
G INDUCES A  $\mathcal{U}$ -COVERAGE  $\tilde{G}J$  ON  $G_* \mathcal{C}$   
BY SETTING  $\tilde{G}J(x) := \{\tilde{G}R : R \in J(x)\}$ .

## BASE CHANGE FOR COVERAGES

THEOREM. SUPPOSE  $G$  SATISFIES  $(\Delta)$ .

THE ASSIGNMENT

$$\Sigma(e, v) \xrightarrow{\tilde{G}} \Sigma(G_* e, u)$$

IS AN INJECTIVE LATTICE MORPHISM.

## BASE CHANGE FOR COVERAGES

IT APPEARS THAT FAITHFULNESS IS NECESSARY!

THEOREM. FOR A FIELD  $K$ ,  $\mathcal{V} = \text{grMod}_K$ ,

$\mathcal{U} = \text{Set}$ , AND

NOT FAITHFUL!

$$G = \text{Hom}_{\mathcal{V}}(K, -) : \mathcal{V} \rightarrow \mathcal{U},$$

THERE EXIST DISTINCT  $\mathcal{V}$ -COVERAGES ON

$K[x, y]$  WHICH GIVE RISE TO THE SAME

$\mathcal{U}$ -COVERAGE.

BASE CHANGE FOR TOPOLOGIES??

CAN WE STRENGTHEN THAT TO HOLD  
FOR TOPOLOGIES? IDK, MAYBE!

PROPOSITION. GIVEN A  $\mathcal{V}$ -TOPOLOGY  $\mathcal{J}$  ON  $\mathcal{C}$ ,  
 $G$  INDUCES A  $\mathcal{U}$ -TOPOLOGY  $\overline{G\mathcal{J}}$  ON  $G_*\mathcal{C}$ .  
↑ SMALLEST TOPOLOGY  
CONTAINING  $\tilde{G}\mathcal{J}$ .

IDEA: FORM  $\tilde{G}\mathcal{J}$ , THEN "CLOSE UNDER (T3)"  
USING A TRANSFINITE CONSTRUCTION.

BASE CHANGE FOR TOPOLOGIES??

IT'S STILL ONLY A CONJECTURE.

CONJECTURE. SUPPOSE  $G$  SATISFIES  $(\blacktriangle)$ .

THE ASSIGNMENT

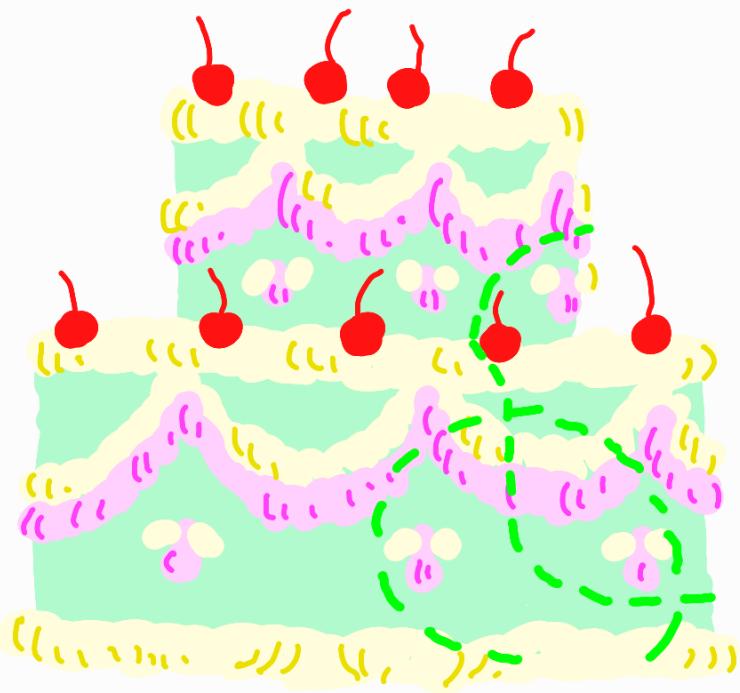
$$\tau(e, v) \xrightarrow{\overline{G(-)}} \tau(G_* e, u)$$

IS AN INJECTIVE LATTICE MORPHISM.

BASE CHANGE FOR TOPOLOGIES??

THIS WOULD LET US MAKE USE OF BORCEUX -  
QUINTEIRO'S CORRESPONDENCE THEOREM,  
SINCE IT APPLIES TO TOPOLOGIES, NOT  
COVERAGES.

PROBLEM:  $\overline{G(-)}$  IS HARD TO WORK WITH -  
NEED AN EASIER WAY TO CONSTRUCT.



I . C.

ENRICHED SHEAVES

AND BASE

CHANGE



## ENRICHED SHEAVES + BASE CHANGE

BORCEUX - QUINTEIRO ALSO GIVE DEFINITIONS FOR

(i) ENRICHED SHEAF

(ii) ENRICHED SHEAIFICATION.

WHEN  $G$  IS A FULLY FAITHFUL RIGHT  
ADJOINT, IT INTERACTS NICELY WITH

i.e.,  $V$  is THESE.  
A REFLECTIVE SUBCAT.  
OF  $U$

QUICKLY, DOCTOR! THE DEFINITIONS!

SAY  $\mathcal{W} \in \text{Cosmos}$  AND  $\mathcal{X}$  IS A SMALL  $\mathcal{W}$ -CATEGORY.

**Definition 5.1.** A  $\mathcal{W}$ -functor  $P \in [\mathcal{X}^{\text{op}}, \mathcal{W}]$  is a **sheaf** for a  $\mathcal{W}$ -coverage  $J$  when, given any object  $x \in \mathcal{X}$ , any  $R \in J(x)$ , any  $g \in \mathcal{G}_{\mathcal{W}}$ , and  $\alpha$  such that

$$\begin{array}{ccc} R & \xrightarrow{r} & \mathcal{X}(-, x) \\ \alpha \downarrow & \swarrow \exists! \beta & \\ \{g, P\} & & \end{array},$$

there exists a unique  $\beta$  for which the diagram commutes.

QUICKLY, DOCTOR! THE DEFINITIONS!

THIS DEFINITION ENCODES THE MORE FAMILIAR  
ONE FROM GEOMETRY:

"GIVEN A FAMILY OF COMPATIBLE LOCAL  
SECTIONS,  $\exists!$  GLOBAL SECTION WHICH  
RESTRICTS TO THEM."

QUICKLY, DOCTOR! THE DEFINITIONS!

FOR A COVERAGE  $J$ , DENOTE THE CAT.  
OF SHEAVES ON  $X$  BY

$$Sh(X, J).$$

QUICKLY, DOCTOR! THE DEFINITIONS!

FACT. WE CAN DEFINE A  $\mathcal{W}$ -FUNCTOR

$$[\mathcal{X}^{\text{op}}, \mathcal{W}] \xrightarrow{\Sigma} [\mathcal{X}^{\text{op}}, \mathcal{W}]$$

WHICH SENDS PRESHEAVES TO "NICER" ONES

AND "NICER" PRESHEAVES TO SHEAVES

w/r/t A COVERAGE  $\mathcal{J}$ .

QUICKLY, DOCTOR! THE DEFINITIONS!

FACT. THE COMPOSITE  $\Sigma\Sigma$  EXHIBITS  
 $\text{Sh}(X, J)$  AS A REFLECTIVE SUBCAT.  
OF  $[X^{\text{op}}, W]$ .

WE CALL

$$[X^{\text{op}}, W] \xrightarrow{\Sigma\Sigma_W} \text{Sh}(X, J)$$

SHEAFIFICATION (W/R/T J).

## ENRICHED SHEAVES + BASE CHANGE

FOR A FULLY FAITHFUL RIGHT ADJOINT  $G$ :

PROPOSITION. FOR  $J \in \Sigma(\mathcal{C}, \mathcal{D})$  AND A SHEAF  $P$  ON  $\mathcal{C}$  WITH RESPECT TO  $J$ ,  $\tilde{G}P$  IS A SHEAF FOR  $\tilde{G}J$ .

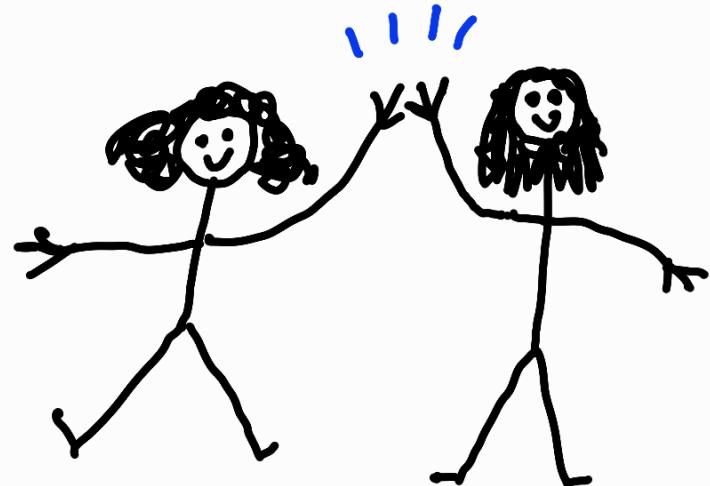
THEOREM. FOR  $P \in [\mathcal{C}^\text{op}, \mathcal{V}]$ , WE HAVE

$$\tilde{G}(\Sigma\Sigma, P) \cong \Sigma\Sigma_u (\tilde{G}P).$$

END OF I.

II. TOWARD A CLASSIFICATION  
OF ENRICHED  
GROTHENDIECK TOPOSES

(j/w ANA TENÓRIO)



## THE UNENRICHED SETTING

WHAT IF I WANT TO KNOW WHETHER A  
GIVEN CATEGORY  $\mathcal{E}$  (ENRICHED OR NOT) IS  
A SHEAF CATEGORY?

IN THE UNENRICHED CASE, WE CAN CHECK

GIRAUD'S AXIOMS.



THE UNENRICHED SETTING

UNFORTUNATELY, GENERALIZING THEM TO  $V$   
IS HARD:

## THE UNENRICHED SETTING

(GIRAUD, c. 1960)

A CATEGORY  $\mathcal{E}$  IS EQUIVALENT TO  $Sh(X, J)$   
FOR SOME (SET-) SITE  $(X, J)$  EXACTLY WHEN IT

- ① IS LOCALLY PRESENTABLE
- ② HAS PULLBACK-STABLE COLIMITS
- ③ HAS DISJOINT COPRODUCTS
- ④ HAS EFFECTIVE QUOTIENTS

## THE UNENRICHED SETTING

THE CLASSIFICATION AS REFLECTIVE SUBCAT'S  
OF PRESHEAF CATEGORIES HOLDS BOTH  
FOR SET AND GENERAL  $V$ .

BUT THERE'S ANOTHER CLASSIFICATION FOR  
SET WHICH WE CAN TRY TO EXTEND  
TO  $V$ ...

## THE UNENRICHED SETTING

THEOREM (STREET, 1981) A LOCALLY SMALL  
CATEGORY  $\mathcal{E}$  IS A GROTHENDIECK  
TOPOS IFF

$$\mathcal{L}_{\mathcal{E}} : \mathcal{E} \longrightarrow [\mathcal{E}^{\text{op}}, \text{Set}]$$

HAS A LEFT EXACT LEFT ADJOINT.

"TOPOS  $\Leftrightarrow$  LEX-TOTAL"



## THE ENRICHED SETTING

WELL, THAT'S NICE! I KNOW WHAT A LEFT  
 $\mathcal{V}$ -ADJOINT IS, I HAVE A GOOD NOTION  
OF ENRICHED YONEDA EMBEDDING, AND  
(LACK, TENDAS 2019) ALREADY DEFINED  
 $\mathcal{V}$ -LEFT EXACTNESS OF  $\mathcal{V}$ -FUNCTORS.

(PRESERVE FINITE WEIGHTED  
 $\mathcal{V}$ -LIMITS)

## THE ENRICHED SETTING

PROVED SO FAR:

THEOREM. IF  $\mathcal{E}$  IS A  $\mathcal{V}$ -GROTHENDIECK  
TOPOS, THEN

$$\mathfrak{L}_{\mathcal{E}} : \mathcal{E} \longrightarrow [\mathcal{E}^{\text{op}}, \mathcal{V}]$$

HAS A  $\mathcal{V}$ -LEFT EXACT  $\mathcal{V}$ -LEFT ADJOINT.

" $\mathcal{V}$ -TOPOS  $\Rightarrow$   $\mathcal{V}$ -LEX TOTAL."

## THE ENRICHED SETTING

IDEA. SHOW THAT THE "STANDARD" YONEDA  
STRUCTURE ON  $\mathcal{V}\text{-Cat}$  RESTRICTS  
TO ONE ON  $\mathcal{V}\text{-Lex}$ , THEN APPLY  
ANOTHER RESULT OF STREET THAT  
ALLOWS US TO DETECT "TOTAL" 0-CELLS.



THANKS FOR LISTENING!